Improvements on a Time Slot Allocation Algorithm

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- Maximum General Assignment Problem
  - Still misses some parts of the model

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- Hard to encode the last preference without making the problem exponential in size

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- This will run at most k times

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- Some matchings infeasible
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- Iterations allow us to change weights in between

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- Thus, we have a k-approximation (in our case, $k = 2$)