T1
```
lock - X(B);
read (B);
B = B-50;
write (B);
unlock (B);
lock - X(A);
read (A);
A := A + 50;
write (A);
unlock (A);
```

T2:
```
lock-S(A);
read (A);
unlock (A);
lock - S(B);
read (B);
unlock (B);
display (A+B);
```

T3:
```
lock - X(B);
read (B);
b := B-50;
write (B);
lock - X(A);
read (A);
A := A +50;
write (A);
unlock (B);
unlock (A);
```

T4:
```
lock-S(A);
read (A);
lock - S(B);
read (B);
display (A+B);
unlock (A);
unlock (B);
```

T3	T4
lock - X(B)	
read (B)	
B := B-50	
write (B)	
	lock -S(A)
	read (A)
	lock - S(B)
lock - X(A)	

Schedule 2

Granting of Locks

$T_3 \rightarrow X(Q)$

$T_2 \rightarrow S(Q)$

$T_4 \rightarrow S(Q)$

$T_5 \rightarrow S(Q)$

when a transaction T_i requests a lock on a data item Q, in a particular mode M, the concurrency control manager runs the lock provided that:

1) there is no other transaction holding a lock on Q in a mode that conflicts with M.

2) there is no other transaction that is waiting for a lock on Q, and that made its lock request before T_i.

2-Phase locking Protocol (2PL)

1. Growing Phase
2. shrinking phase

Lock point - last of the locks being acquired

$T_1 \rightarrow T_2 \rightarrow T_3$

Precedence graph