

# COMP4500 Assignment 1

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## Part A (30 marks total)

### Question 1: Constructing SNI and directed graph

a) My origin SNI: 9845 0048 3052

1 for  $i = 2$  to 12

2 if  $d[i] == d[i - 1]$

3  $d[i] = (d[i] + 3) \bmod 10$

My new SNI: 9845 0348 3052

b) 0→3 0→5

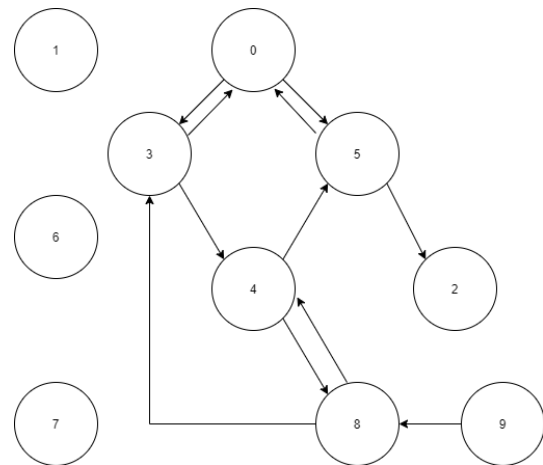
3→4 3→0

4→5 4→8

5→0 5→2

8→4 8→3

9→8



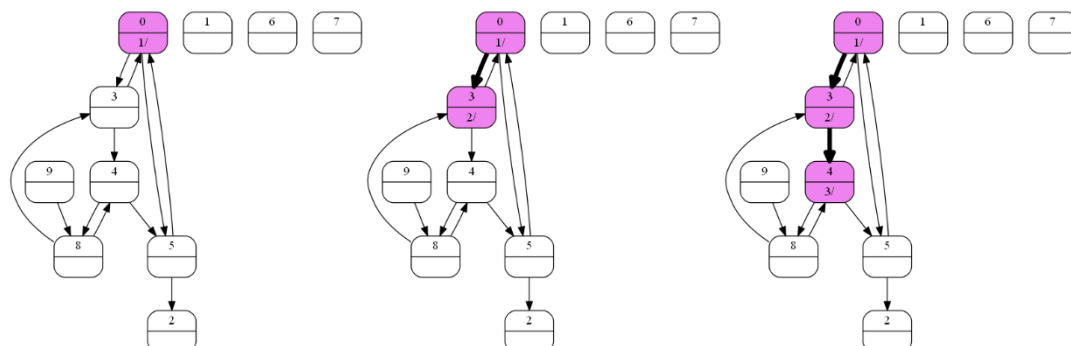
### Question 2: Strongly connected components

(a) Perform step 1 of the SCC algorithm using S as input.

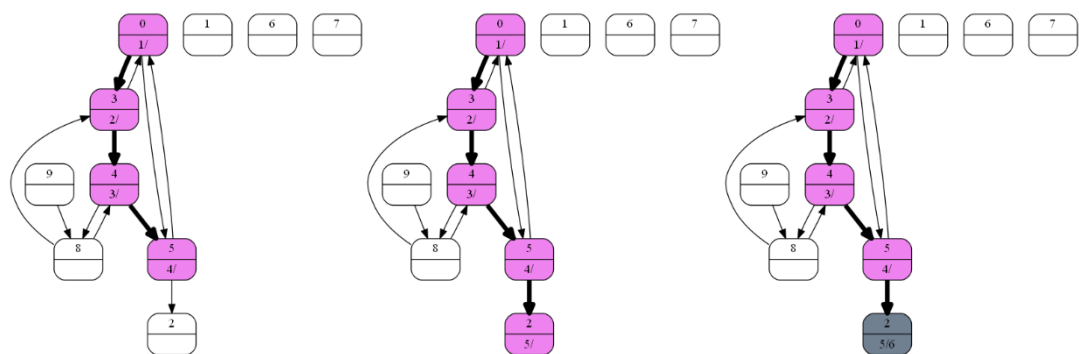
Note: the purple node means on visit, gray node means finish visited, white node means still not visited.

The bold edge is the path of performing depth-first search

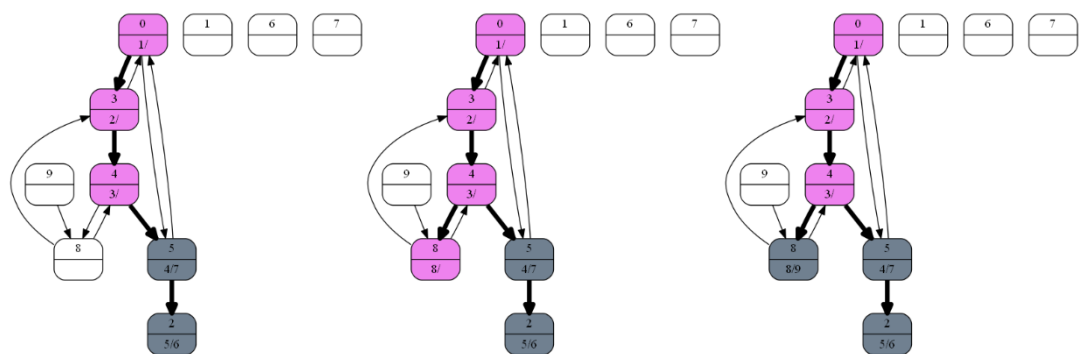
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4	white			4	white			4	purple	3
5	white			5	white			5	white	
6	white			6	white			6	white	
7	white			7	white			7	white	
8	white			8	white			8	white	
9	white			9	white			9	white	



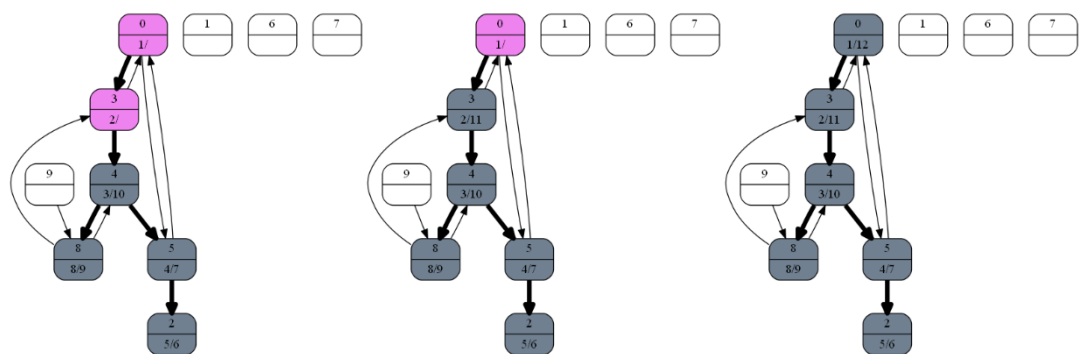
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2	white			2	purple	5		2	grey	5
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4	purple	3		4	purple	3		4	purple	3
5	purple	4		5	purple	4		5	purple	4
6	white			6	white			6	white	
7	white			7	white			7	white	
8	white			8	white			8	white	
9	white			9	white			9	white	



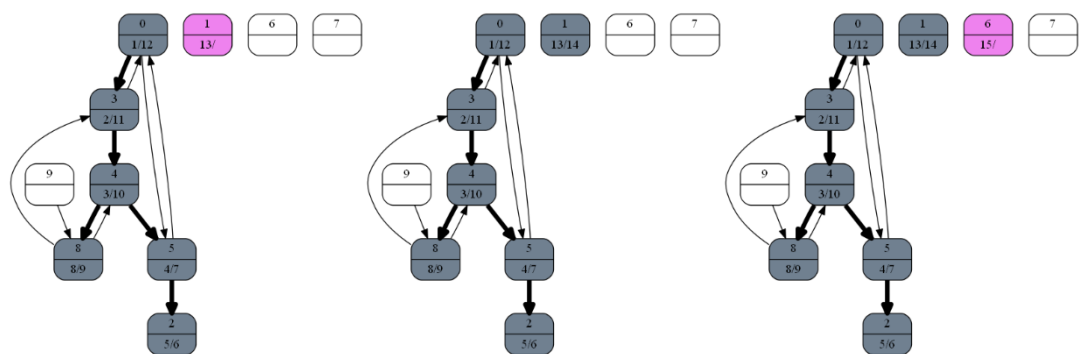
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2	grey	5		2	grey	5		2	grey	5
3	purple	0		3	purple	0		3	purple	0
4	purple	3		4	purple	3		4	purple	3
5	grey	4		5	grey	4		5	grey	4
6	white			6	white			6	white	
7	white			7	white			7	white	
8	white			8	purple	4		8	grey	4
9	white			9	white			9	white	



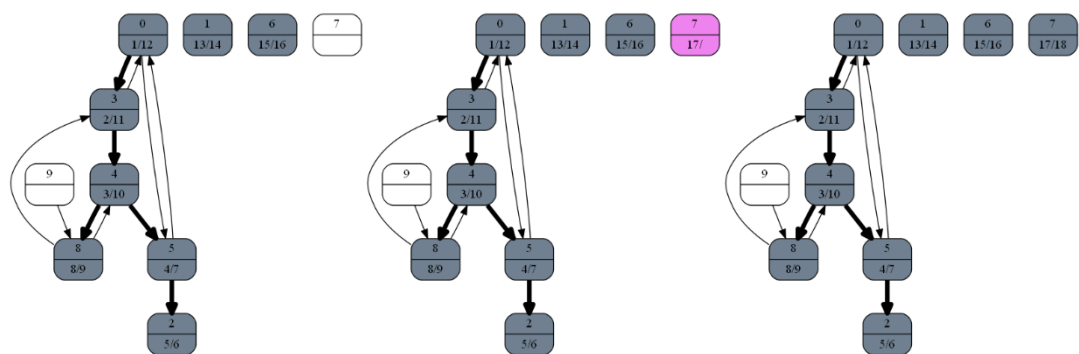
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3	purple	0		3	grey	0		3	grey	0
4	grey	3		4	grey	3		4	grey	3
5	grey	4		5	grey	4		5	grey	4
6	white			6	white			6	white	
7	white			7	white			7	white	
8	grey	4		8	grey	4		8	grey	4
9	white			9	white			9	white	



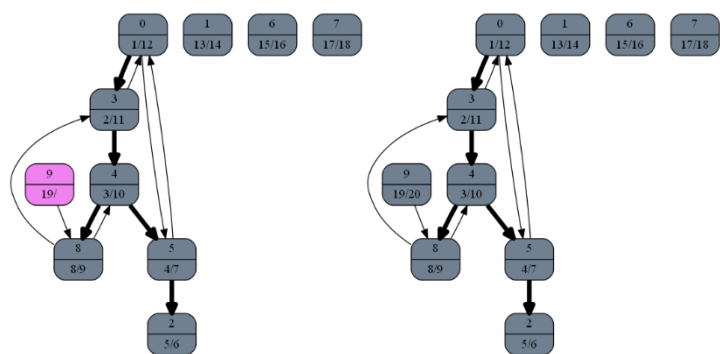
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2	grey	5		2	grey	5		2	grey	5
3	grey	0		3	grey	0		3	grey	0
4	grey	3		4	grey	3		4	grey	3
5	grey	4		5	grey	4		5	grey	4
6	white			6	white			6	purple	undef
7	white			7	white			7	white	
8	grey	4		8	grey	4		8	grey	4
9	white			9	white			9	white	



x	Color[x]	Pi[x]		x	Color[x]	Pi[x]		x	Color[x]	Pi[x]
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1	grey	undef		1	grey	undef		1	grey	undef
2	grey	5		2	grey	5		2	grey	5
3	grey	0		3	grey	0		3	grey	0
4	grey	3		4	grey	3		4	grey	3
5	grey	4		5	grey	4		5	grey	4
6	grey	undef		6	grey	undef		6	grey	undef
7	white			7	purple	undef		7	grey	undef
8	grey	4		8	grey	4		8	grey	4
9	white			9	white			9	white	



x	Color[x]	Pi[x]		x	Color[x]	Pi[x]
0	grey	undef		0	grey	undef
1	grey	undef		1	grey	undef
2	grey	5		2	grey	5
3	grey	0		3	grey	0
4	grey	3		4	grey	3
5	grey	4		5	grey	4
6	grey	undef		6	grey	undef
7	grey	undef		7	grey	undef
8	grey	4		8	grey	4
9	purple	undef		9	grey	undef



Second last graph: Finishing times for the original graph G

Last graph: Strongly Connected Components

(b) 0 -> 5 0 -> 3

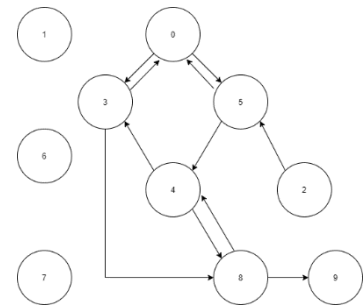
2 -> 5

3 -> 0 3 -> 8

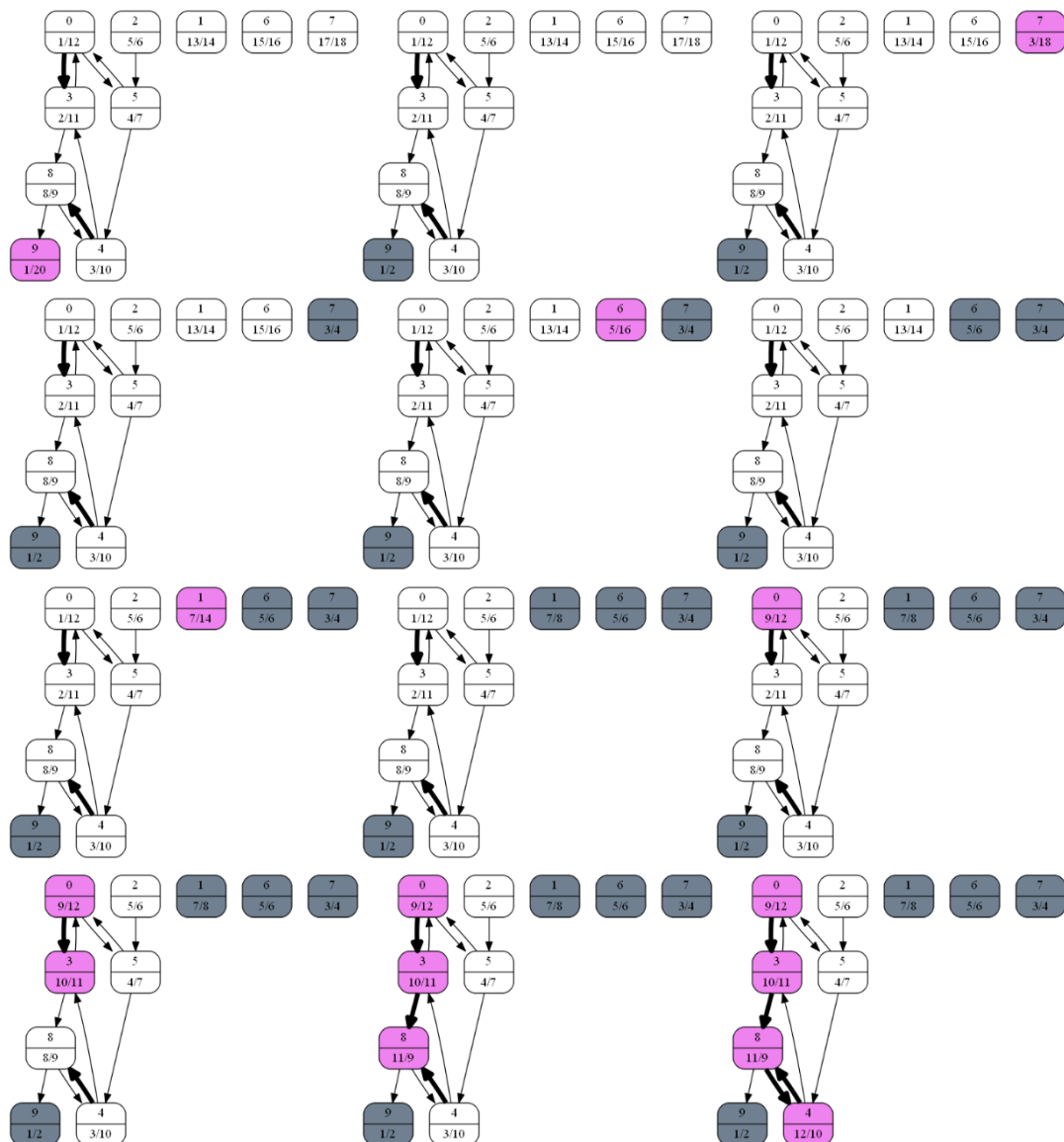
4 -> 8 4 -> 3

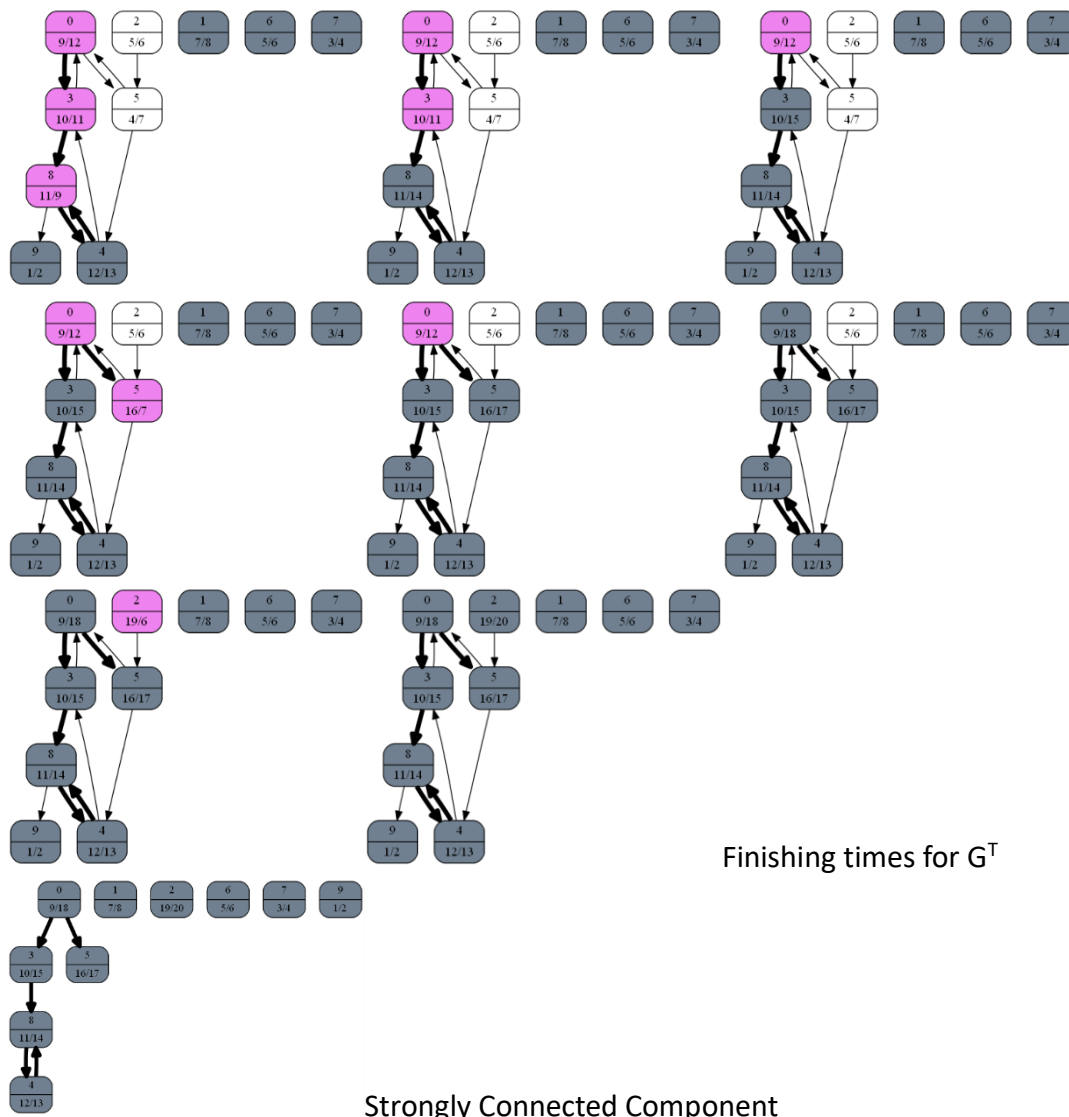
5 -> 4 5 -> 0

8 -> 9 8 -> 4



(c) Perform steps 3, 4 of the SCC algorithms.





1. 9
2. 7
3. 6
4. 1
5. 0 3 8 4 5
6. 2

## Part B (70 marks total)

Question 4: Worst-case time complexity analysis

a) In my method, I treat each interaction as a node in the graph, and the edges represent the two element sub paths of possible route of transmission from personFrom to personTo.

```
List<Interaction> FindTransmissionPath(int start, int end, List<Interaction> interactions) {
1    personToInteractions: HashMap<Integer, HashSet<Interaction>> = new
HashMap().onLookupFail(new HashSet())
2    endInteractions: HashSet<Interaction> = new HashSet()
3    // I (number of interactions) iterations with O(1) loop body
4    // with O(P) (number of person) Hashset construction cost
5    // Overall: O(I + P)
6    // Worst-case: |P| = 2 * |I|, if each interaction transmit between two unique person
7    // 3 * |I| give us Overall O(I)
8    for interaction in interactions:
9        startInteractions: HashSet<Interaction> = personToInteractions[interaction.PersonFrom]
10       startInteractions.add(interaction) // O(1)
11       if interaction.personTo == end -> endInteractions.add(interaction)
12       // O(I) due to implementation limitations
13       sources: HashSet<Interaction> = personToInteractions[start]
14       adjacency: HashMap<Interaction, HashSet<Interaction>> = new
HashMap().onLookupFail(new HashSet())
15       // I iterations with worst-case: O(I) loop body, if successors' time >= predecessors' time
16       // and successors' personFrom don't equal to predecessors' personTo
17       // Overall: O(I^2)
18       for interaction in interactions:
19           neighbors: HashSet<Interaction> = personToInteractions[interaction.personTo]
20           .filter(i -> interaction.time <= i.time && interaction.PersonFrom != i.PersonTo)//O(I)
21           adjacency[interaction] = neighbors // O(1)
22       // Running time for Dijkstra's algorithm using a Java Heap as a priority queue is
23       // O((|E| + |V|) * log|V|).
24       // Worst-case: |E| = |V|^2, we get O((|V|^2 + |V|) * log|V|).
25       // as we use Interaction I as our Vertex
26       // so O((|V|^2 + |V|) * log|V|) give us Overall O(I^2 * logI)
27       Dijkstra(adjacency, sources);
28       maximumProb: Double = personToInteractions[end].
29       .filter(v -> v.prob != Double.MAX_VALUE).maxBy(v -> v.prob).prob
30       finalInteractionList: List<Interaction> = new ArrayList<>();
31       finalInteractionList.add(lowestDDestinationVertex.get().element);
32       while True // O(I)
33           if head is not null
34               finalInteractionList.add(0, head.element);
35               head = head.predecessor;
36           else // head is null (finish)
37               return finalInteractionList;
```



b) The first part of my algorithm to prepare the parameters of Dijkstra's algorithm, which are a HashMap of one interaction to its valid neighbors and a HashSet of source interactions.

We assume that when HashSet and HashMap execute put, add, and get, its worst-case time complexity will be  $O(1)$  instead of  $O(p)$ .

For doing a graph search, I first convert all interaction to a vertex and build an endInteraction HashSet where the path finish, which is Overall  $O(i)$ ;

For creating the HashSet of sources, it is  $O(i)$  in worst-case;

For creating the HashMap, I run through all interactions and filter its time elapsed and check if there is a loop in interactions, which is overall  $O(i^2)$ .

Then the Running time for Dijkstra's algorithm using a Java Heap as a priority queue is Overall  $O(i^2 * \log i)$  in worst-case. We use a Binary Heap in the implementation of Dijkstra's, with  $O(\lg V)$  Extract-Mins and Decrease-Keys.

In the last part, my algorithm is to find the path with the highest probability through all endInteraction HashSets, and then add it to an ArrayList, both of which are  $O(i)$ .

Thus, the time complexity is  $O(i^2 * \log i)$  which describes an asymptotic upper bound on the worst-case time complexity of this algorithm.