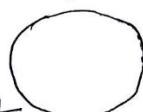


SYNCHRONIZATION

~~Synchronized~~ is a modifier applicable only for methods / blocks but not for classes & variables.

- > If multiple threads are trying to operate simultaneously on ^{the} same object then there may be a chance of data inconsistency problem.
- > To overcome this problem we should use synchronized keyword.
- > If a method / block declared as synchronized then at a time only one thread is allowed to execute that method/block on the given object so that data inconsistency problem will be resolved.
- > The main advantage of synchronized keyword is  we can resolve data inconsistency problems but the main disadvantage of synchronized keyword is, it increases waiting time of threads & creates performance problems. Hence if there is no specific requirement then it is not recommended to use synchronized keyword.
- > Internally synchronization ^{concept} is implemented by using lock. Every object in java has a unique lock. Whenever we are using synchronized keyword then only lock concept will come into the picture.
 - . If a thread wants to execute synchronized method on the given object first it has to get lock of that object. Once thread gets the lock then it is allowed to execute any synchronized method on that object.
 - Once method execution completes automatically thread releases lock.
- > Acquiring & releasing lock internally takes care by JVM. Programmer is not responsible for this activity.

class X

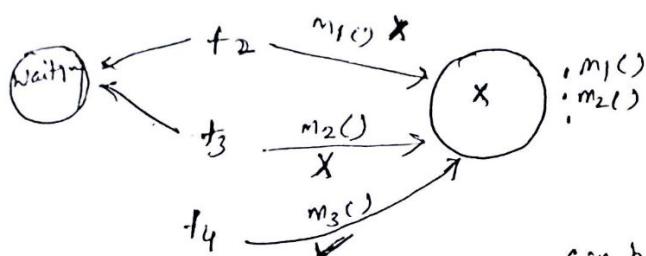
{

 synchronized m1();
 synchronized m2();

 m3();

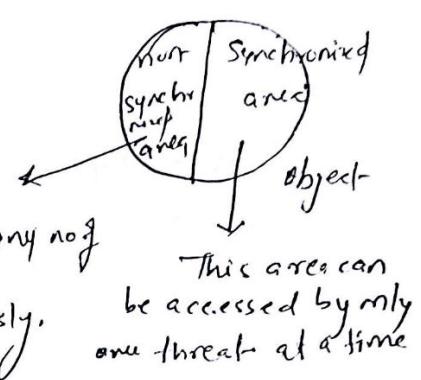
}

→ While a thread executing synchronized method on the given object, the remaining threads are not allowed to execute any synchronized method simultaneously on the same object. But, remaining threads are allowed to execute non synchronized methods simultaneously.



$t_1 \rightarrow l(X)$

can be
accessed by any no. of
threads
simultaneously.



This area can
be accessed by only
one thread at a time

class X

{
 synchronization {

 wherever we are performing update operation [add/delete/
 replace/remove] where state of object changing

}

Non-synchronized

&

whenever object state won't be changed [read() operation]

3

ex

 class ReservationSystem

 {
 checkAvailability()

 // read operation

3

 synchronized bookTicket()

 {
 update operation

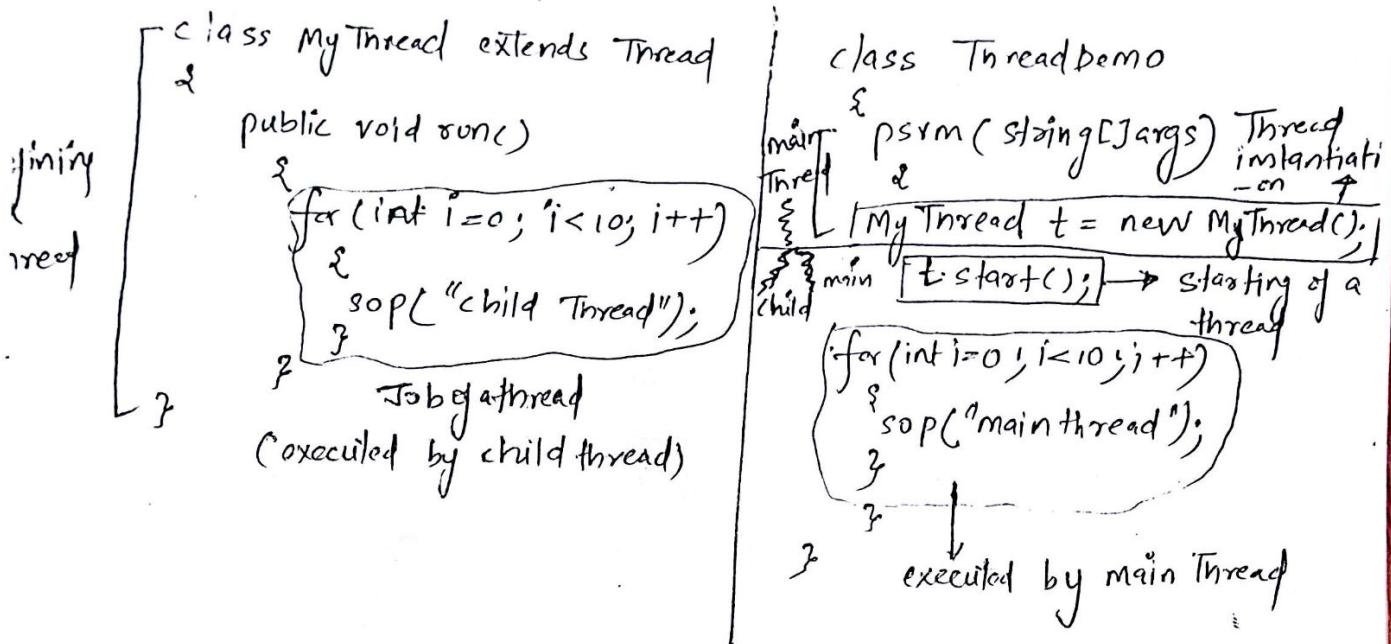
3

3

By extending Thread class

`t.start()` → calls `start()` method of Thread class which internally calls `run()` method of MyThread class.

(2)



output → mixed output (no dependency)

<any order threads can follow>

Case 1 Thread scheduler

- 1) It is the part of JVM
- 2) It is responsible to schedule threads i.e. if multiple threads are waiting to get ~~the~~ chance of execution then in which order threads will be executed is decided by Thread scheduler.
- 3) We can't expect exact algorithm followed by thread scheduler, it is varied from JVM to JVM. Hence we can't expect threads execution order & exact output.

Hence whenever situation comes to multithreading, there is no guarantee for exact output but we can provide several possible outputs

POSSIBLE OUTPUTS :-

- 1) main Thread
main Thread
:
10 times
child Thread
child Thread
:
10 times
- 2) child Thread
child Thread
:
10 times
main Thread
main Thread
:
10 times
- 3) main Thread
child Thread
main Thread
:
10 times
- 4) child Thread
main Thread
child Thread
:
10 times

~~case 1~~
case 2 Difference between `t.start()` & `t.run();`

```
class MyThread extends Thread
{
    public void run()
    {
        for (int i=0; i<10; i++)
        {
            System.out.println("child Thread");
        }
    }
}
```

```
Child Thread
Child Thread
:  
10 times
main Thread
10 times
```

Thread class
method

```
class ThreadDemo
{
    public static void main (String [] args)
    {
        MyThread t = new MyThread();
        t.run();
        for (int i=0; i<10; i++)
        {
            System.out.println("main Thread");
        }
    }
}
```

In the case of `t.start()`, a new thread will be created which is responsible for the execution of `run()` method.

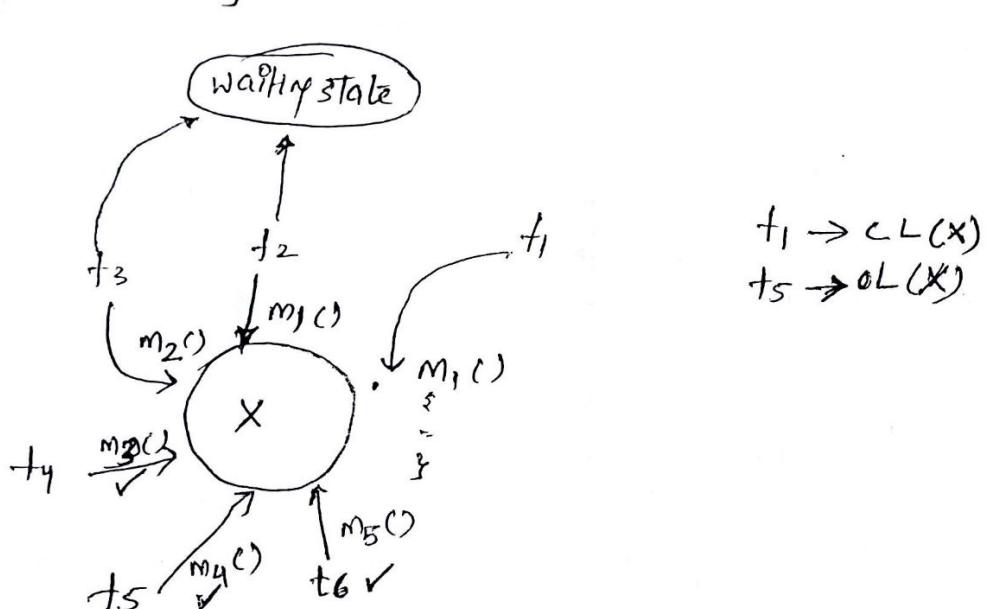
But, In the case of `t.run()`, a new thread won't be created & `run()` method will be executed just like a normal method call by main thread. Hence in the above program, output is fixed

Child Thread (10 times) followed by main Thread (10 times). This total output produced by only main Thread.

Once method execution completes, automatically thread releases the lock.

class X

```
{  
    static synchronized m1() {}  
    static synchronized m2() {}  
    static m3() {}  
    synchronized m4() {}  
    m5() {}  
}
```



$t_1 \rightarrow CL(X)$
 $t_5 \rightarrow OL(X)$

While a thread executing static synchronized method, the remaining threads are not allowed to execute any static synchronized method of that class simultaneously. But, remaining threads are allowed to execute ~~and~~ the following methods simultaneously :-

- normal static method
- instance synchronized method
- normal instance method

```

class Display
{
    public void display()
    {
        for (int i=1; i<=10; i++)
        {
            sop(i);
            try
            {
                Thread.sleep(2000);
            }
            catch (InterruptedException e)
            {
                e.printStackTrace();
            }
        }
    }
}

```

```

public void displayC()
{
    for (int i=65; i<=75; i++)
    {
        sop((char)i);
        try
        {
            Thread.sleep(2000);
        }
        catch (IE e)
        {
        }
    }
}

```

class MyThread1 extends Thread

```

    Display d;
    MyThread1 (Display d)

```

```

        {
            this.d = d;
        }

```

```

    public void run()

```

```

        {
            d.display();
        }
    }
}

```

class MyThread2 extends Thread

```

    Display d;
    MyThread2 (Display d)

```

```

        {
            this.d = d;
        }

```

```

    public void run()

```

```

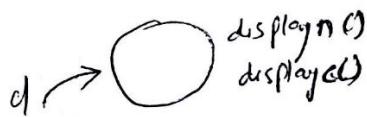
        {
            d.display();
        }
    }
}

```

```

class Test
{
    JVM (String[] args)
    {
        Display d = new Display();
        MyThread t1 = new MyThread1(d);
        MyThread t2 = new MyThread2(d);
        t1.start();
        t2.start();
    }
}

```

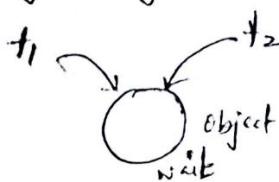


Object communication

(A)

Two threads can communicate with each other by using wait, notify & notifyAll methods.

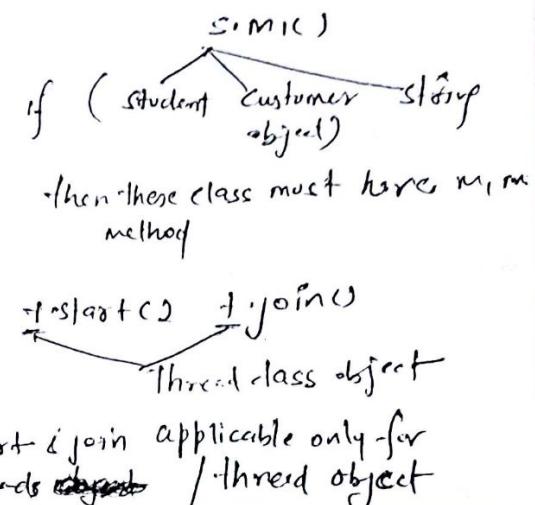
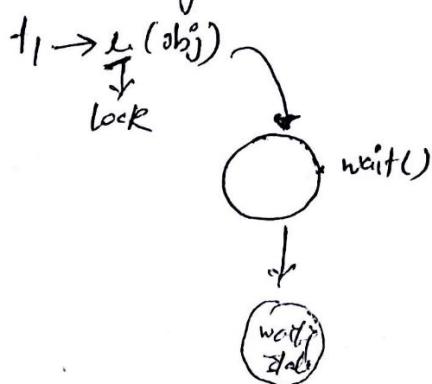
- The thread which is expecting update is responsible to call wait method. The immediately the thread will enter into waiting state.
 - The thread which is responsible to perform update, after performing update, it is responsible to call notify method then waiting thread will get that notification & continue its execution with those updated items.
- But* wait, notify & notifyAll methods present in Object class but not in Thread class because Thread can call these methods on any Java object.



To call wait, notify & notifyAll methods on any object, thread should be owner of that object (must acquire lock on that object) i.e. thread should be inside synchronized area.

Hence we can call wait, notify & notifyAll methods only from synchronized areas otherwise we will get Runtime Exception "Illegal Monitor State Exception"

- If a thread calls wait method on any object, it immediately releases the lock of that particular object & enter into waiting state.



- If a thread calls notify method on any object, it releases the lock of that object but may not immediately (postpone)
- * Except wait, notify & notifyAll methods, there is no other method where thread releases the lock.

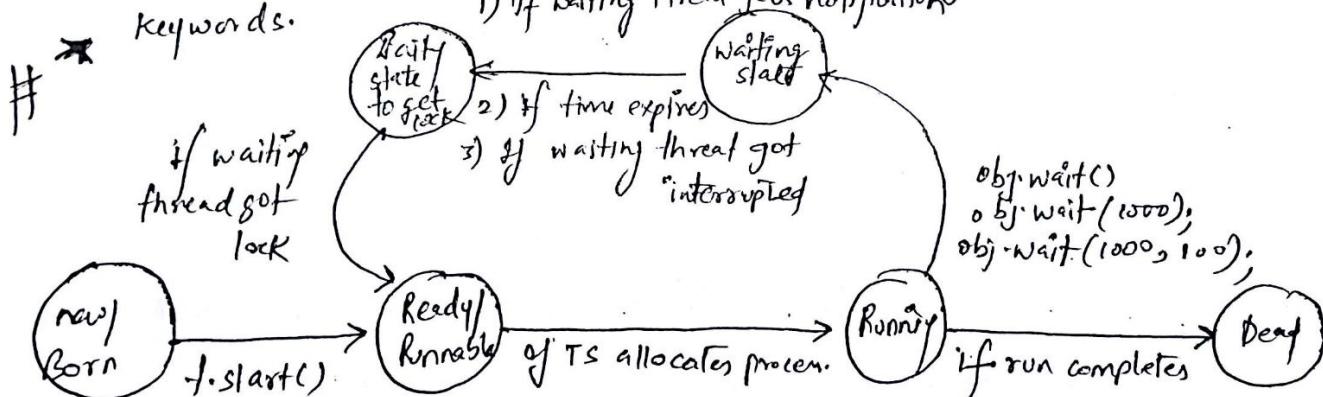
D) wait() method →
public final void wait() throws InterruptedException

public final native void wait(long ms) "

public final void wait(long ms, int ns) "

D) public final native void notify()
public final native void notifyAll()

Note → Every wait method throws InterruptedException which is checked exception. Hence, whenever we are using wait method, compulsory we should handle the exception either by using try-catch or throws keywords.



CHANGES IN TRANSITION STATE DIAGRAM

- notification
- get the lock

class ThreadA

```
    {
        public (String[] args)
    }
```

```
ThreadB = new ThreadB();
→ t.start();
sop("H");
sop(t.total);
}
}
```

class ThreadB extends Thread

```
(5)
int total = 0;
public void run()
{
    for (int i=0; i<1000; i++)
    {
        total = total + i;
    }
}
```

500500

* Using wait() & notify()

class ThreadA

```
    {
        public (String[] args) throws InterruptedException
    }
```

```
ThreadB = new ThreadB();
t.start(); synchronized (t) { t.wait(); }
// break sleep(10000);
sop(t.total);
}
}
```

class ThreadB

```
    {
        int total = 0;
        public void run()
        {
            synchronized (this)
            {
                for (int i=0; i<1000; i++)
                {
                    total = total + i;
                    this.notify();
                }
            }
        }
    }
```

3) program (file)

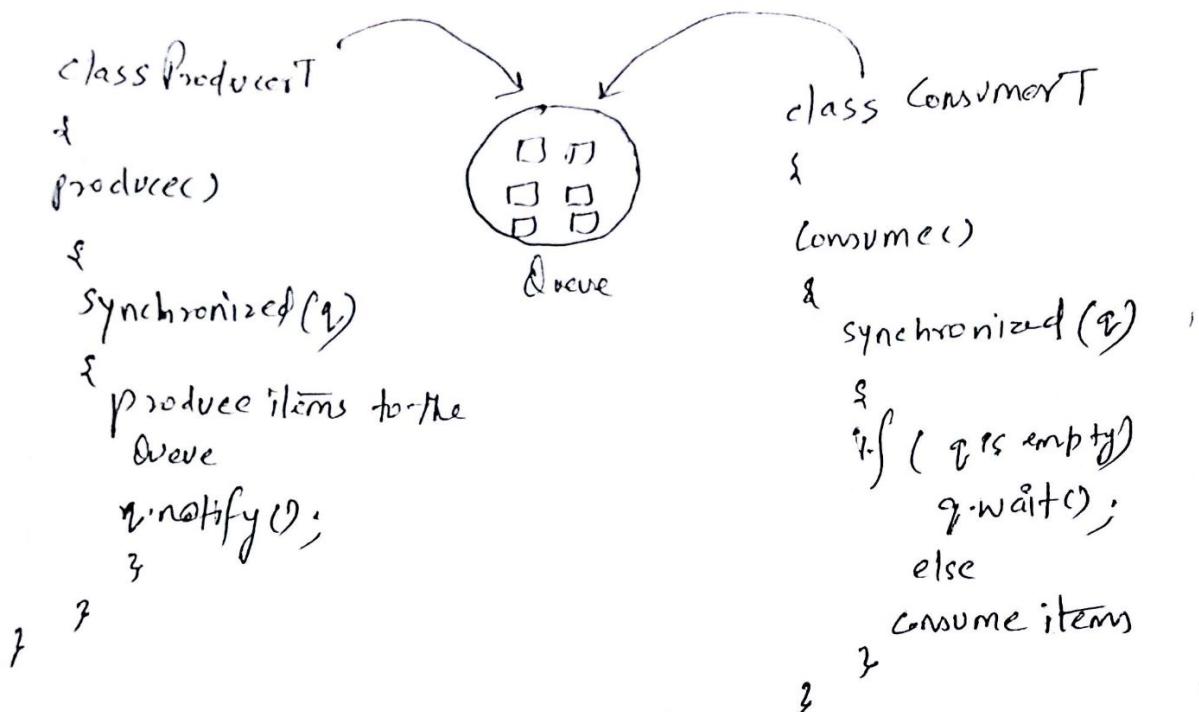
Producer-consumer Problem

→ Producer thread is responsible to produce items to the Queue
& consumer thread is responsible to consume items from the Queue

→ If queue is empty then consumer thread will call wait method
& enter into waiting state

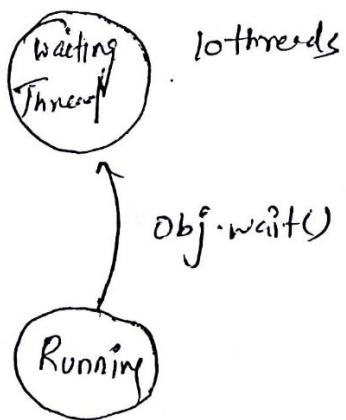
→ After producing items to the Queue, producer thread is responsible to call notify method then waiting consumer will get that

notification & continue its execution with updated items



Difference b/w notify & notify All

- We can use notify method to give the notification for only one waiting thread. If multiple threads are waiting then only one thread will be notified & the remaining threads have to wait for further notifications.
 - Which thread will be notified, we can't expect, it depends on JVM



We can use notifyAll to give the notification for all waiting threads
of a particular object then only one thread (6)

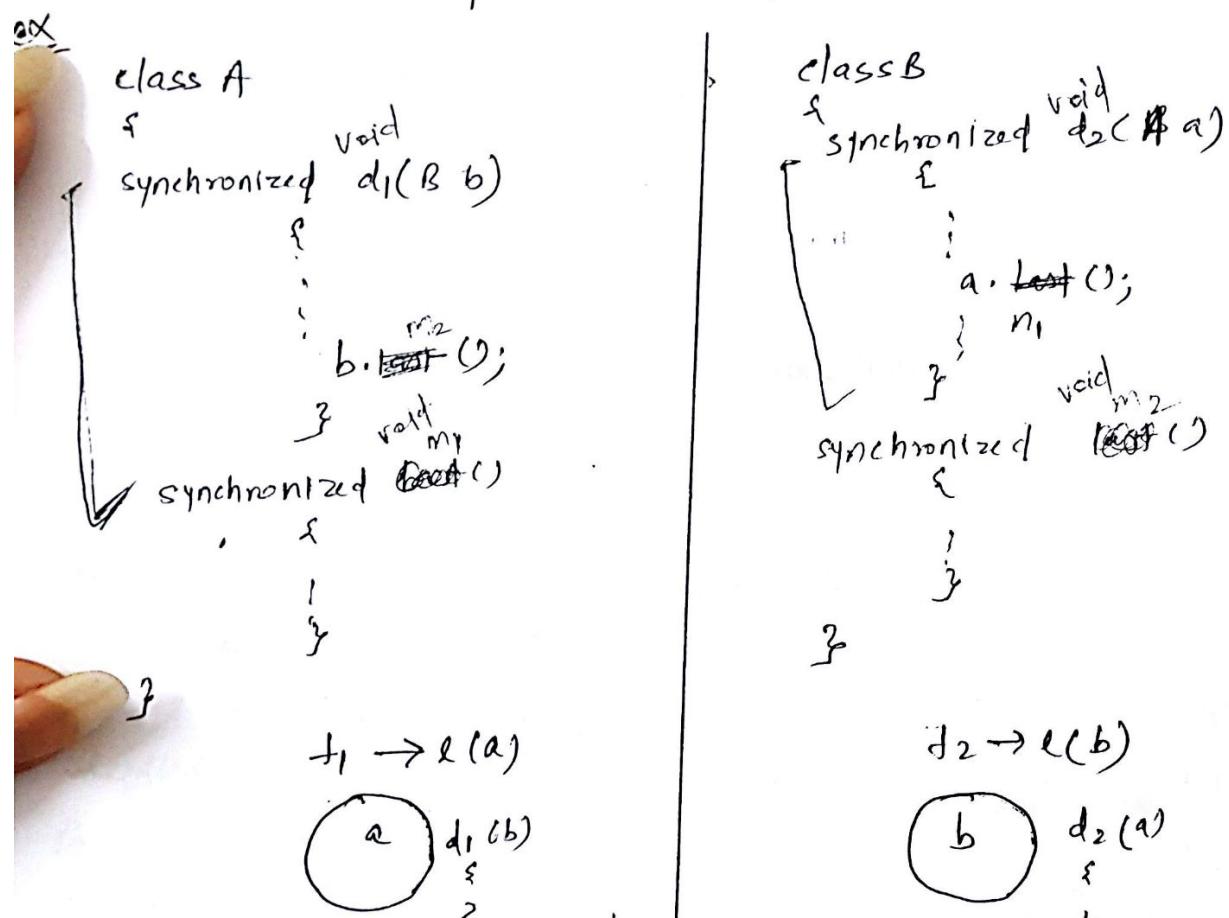
Even though multiple threads notified, execution will be performed
one by one because threads require lock & only one lock is
available.

DEADLOCK

If two threads are waiting for each other forever, such type of infinite waiting is called deadlock.

Synchronized keyword is the only reason for deadlock situation. Hence, while using synchronized keyword we have to take special care.

There are no resolution techniques for deadlock but several prevention techniques are available.



$t_1 \rightarrow d_1(b)$
 $t_2 \rightarrow d_2(a)$

t_1 will enter into waiting state to gets b 's lock & currently t_1 holds a 's lock

$t_2 \rightarrow d_2(a)$
 t_2 will enter into waiting state to gets a 's lock & currently holds b 's lock.

class A

```
{  
    public synchronized void d1(B b)  
    {  
        sop ("Thread-1 starts execution of  
              d1() method");  
        try {  
            Thread.sleep (5000);  
        } catch (IE e) {}  
        sop ("Thread-1 trying to call b's last method");  
        b.last(  
            m2);  
    }  
}  
public synchronized void lastm1()
```

class B

```
{  
    public synchronized void d2(A a)  
    {  
        sop ("Thread-2 starts execution of d2() method");  
        try {  
            Thread.sleep (5000);  
        } catch (IE e) {}  
        sop ("Thread-2 trying to call A's last() method");  
        a.last(  
            m1);  
    }  
}  
public synchronized void lastm2()  
{  
    sop ("Inside-B, last() method");  
}
```

(8)

```

class MyThread extends Thread
{
    A a = new A();
    B b = new B();
    public void m1()
    {
        this.start(); // main starts child thread
        a.d1(b);      → main thread
    }
    public void run()
    {
        b.d2(a);      ] → child thread
    }
}
public class Main
{
    public static void main(String[] args)
    {
        MyThread t = new MyThread();
        t.m1();
    }
}

```

In the above program if we remove at least one synchronized keyword then the program won't enter into deadlock. Hence synchronized keyword is the only reason for deadlock situation. Due to this, while using synchronized keyword we have to take special care.

a.d1(b) }
 sync? }

~~Starvation~~ Deadlock vs

long waiting of a thread where waiting never ends is called deadlock

whereas, long waiting of a thread where waiting ends at certain point is called starvation.

ex low priority thread has to wait until completing all high priority threads. This waiting can be long waiting but ends at certain point, is called starvation.

$\boxed{\begin{array}{l} \downarrow \text{wait} \\ \begin{array}{l} \text{1 thread - P(1)} \\ \text{1000 threads P(10)} \end{array} \end{array}}$