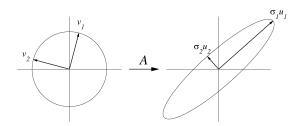
# Data mining: lecture 8

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We will see that any matrix  $A \in \mathbb{R}^{m \times n}$  can be written as  $A = U \Sigma V^T$  such that  $U \in \mathbb{R}^{m \times m}$  is unitary,  $V \in \mathbb{R}^{n \times n}$  is unitary, and  $\Sigma \in \mathbb{R}^{m \times n}$  is a nonnegative real diagonal matrix.  $\Sigma(i,i)$ , denoted  $\sigma_i$ , are unique. If A the singular values are distinct, then the singular vectors are unique up to a multiplication by  $z \in \mathbb{C}$  with |z| = 1.

Remark 0.1. Note the difference in notation from what we saw in class. The matrices V and U are what we denoted by  $[V; \overline{V}]$  and  $[U; \overline{U}]$  respectively. This makes the proofs a little cleaner and hopefully more easy to follow. Note also that  $\Sigma$ , unlike the matrix we denoted by S, is not square. The non square matrix  $\Sigma$  is still diagonal though, i.e.  $\Sigma(i,j) = 0$  for all  $i \neq j$ .

## 1 The geometry of SVD



#### 2 Proof of existence

Set  $\sigma_1 = ||A||_2$ . Let  $u_1 \in \mathbb{R}^n$  and  $v_1 \in \mathbb{R}^m$  be unit 2-norm vectors such that  $Av_1 = \sigma_1 u_1$ . To find these vectors, find the unit vector  $v_1$  that brings to maximum the expression

$$\max_{\|x\|=1} \|Ax\|.$$

Then  $Av_1 = \mu u_1$  for some  $\mu$  and a unit vector  $u_1$ . Since  $||Av_1|| = \sigma_1$ , we get that  $\sigma_1 = ||Av_1|| = |\mu| ||u_1|| = |\mu|$ . Set  $\mu = \sigma_1$  to be positive, by flipping the sign of  $u_1$  if needed.

Complete  $v_1$  into an orthonormal basis of  $\mathbb{C}^n$ , denote  $V_1$ . Complete  $u_1$  into

an orthonormal basis of  $\mathbb{C}^m$ , denoted  $U_1$ .

$$S = U_1^T A V_1 = U_1^T \left[ \sigma_1 u_1, A v_2, \dots, A v_n \right] = \begin{pmatrix} \sigma_1 & w^T \\ 0 & B \end{pmatrix}.$$

We will show that  $w^T = 0$ .

$$\left\| S \begin{pmatrix} \sigma_1 \\ w \end{pmatrix} \right\|_2 = \left\| \begin{pmatrix} \sigma_1 & w^T \\ 0 & B \end{pmatrix} \begin{pmatrix} \sigma_1 \\ w \end{pmatrix}_2^T \right\| \ge \sigma_1^2 + w^T w = \sqrt{\sigma_1^2 + w^T w} \left\| \begin{pmatrix} \sigma_1 \\ w \end{pmatrix} \right\|_2.$$

That is  $||S|| \ge \sqrt{\sigma_1^2 + w^T w}$ . But  $||S||_2 = ||A||_2 = \sigma_1$  and so w = 0. By induction,  $B = U_2 \Sigma_2 V_2^T$  and

$$A = U_1 S V_1^T = U_1 \begin{pmatrix} 1 & 0 \\ 0 & U_2 \end{pmatrix} \begin{pmatrix} \sigma_1 & 0 \\ 0 & \Sigma_2 \end{pmatrix} \begin{pmatrix} 1 & 0 \\ 0 & V_2^T \end{pmatrix} V_1^T.$$

The matrices

$$U = U_1 \begin{pmatrix} 1 & 0 \\ 0 & U_2 \end{pmatrix}, \quad V = V_1 \begin{pmatrix} 1 & 0 \\ 0 & V_2^T \end{pmatrix}$$

are unitary and the proof is complete.

## 3 More properties of SVD

**Lemma 3.1.** The rank of A equals the number of nonzero singular values.

*Proof.* Recall that if B is  $n \times k$  with rank n than rank(AB) = rank(A), and if C is  $l \times m$  with rank m then rank(CA) = rank(A). Thus,

$$rank(A) = rank(U\Sigma V^T) = rank(\Sigma V^T) = rank(\Sigma).$$

Since  $\Sigma$  is diagonal, its rank it the number its nonzero elements.

**Lemma 3.2.** Let rank(A) = r. Then,

range(A) = span(
$$u_1, ..., u_r$$
),  
null(A) = span( $v_{r+1}, ..., v_n$ ).

Proof.

$$y \in \text{range}(A) \iff \exists x \text{ such that } y = Ax$$

$$\iff y = U\Sigma V^T x$$

$$\iff y = U\Sigma z, \text{ where } z = V^T x$$

$$\iff y = U\left(\sigma_1 z_1, \dots, \sigma_r z_r, 0, \dots, 0\right)^T$$

$$\iff y = \sum_{i=1}^r (\sigma_i z_i) u_i$$

$$\iff y \in \text{span}(u_1, \dots, u_r).$$

$$x \in \text{null}(A) \iff ||Ax||_2 = 0 \iff ||U\Sigma V^T x||_2 = 0$$

$$\iff ||\Sigma V^T x||_2 = 0 \iff ||\Sigma y||_2 = 0 \text{ where } y = V^T x$$

$$\iff y = (0, \dots, 0, y_{r+1}, \dots, y_n)^T \text{ where } y = V^T x$$

$$\iff x = Vy, \ y = (0, \dots, 0, y_{r+1}, \dots, y_n)^T$$

$$\iff x = \sum_{i=r+1}^n y_i v_i$$

$$\iff x \in \text{span}(v_{r+1}, \dots, v_n).$$

**Lemma 3.3.**  $||A||_2 = \sigma_1$  (even if you don't know the above proof).

*Proof.* Immediate from the invariance of  $\|\cdot\|_2$  under unitary transformations.

Similarly, 
$$||A||_F = (\sigma_1^2 + \dots + \sigma_r^2)^{1/2}$$
.

# 4 Relation between singular values and eigenvalues

**Lemma 4.1.** The singular values of A are the square roots of the nonzero eigenvalues of  $A^TA$  and  $AA^T$ .

*Proof.* If  $A = U\Sigma V^T$ , then  $A^T = V\Sigma U^T$  and

$$AA^T = \left(U\Sigma V^T\right)\left(V\Sigma U^T\right) = U\Sigma\Sigma U^T = U\Sigma^2 U^{-1}.$$

 $AA^T$  is positive semi-definite and therefore all eigenvalues are non-negative and there is no problem with the square root.

Do not use this observation to compute the SVD! Reason: Assume for simplicity that we have a  $2 \times 2$  matrix A (not diagonal) whose SVD is given by  $A = \sigma_1 u_1 v_1^T + \sigma_2 u_2 v_2^T$  (See 5.1 below). If  $\sigma_2/\sigma_1 < 10^{-15}$ , then the second term will disappear due to round-off, that is, we cannot represent such a matrix in double precision. Now, if the matrix A has  $\sigma_2/\sigma_1 < 10^{-8}$ , then  $A^TA$  and  $AA^T$  have ratio of singular values that is smaller than  $10^{-15}$ , and so those matrices cannot be represented, and will be approximated as rank-1 matrices with the second singular value being due to round-off. In other words, although A is not terribly conditioned, we loose the small eigenvalues if we try to compute the SVD by computing the eigenvalues of  $A^TA$  or  $AA^T$ .

**Lemma 4.2.** If A is hermitian, then the singular values of A are the absolute values of its eigenvalues.

*Proof.* A hermitian matrix is diagonalized by a unitary matrix with real eigenvalues. That is,

$$A = Q\Lambda Q^T = Q|\Lambda|\operatorname{sign}(\Lambda)Q^T.$$

Now set U = Q,  $\Sigma = |\Lambda|$ ,  $V^T = \text{sign}(\Lambda)Q^T$ .

## 5 Approximation properties

#### 5.1 Rank-k approximation in the spectral norm

**Lemma 5.1.** A can be written as a sum of rank-1 matrices. Explicitly,

$$A = \sum_{j=1}^{r} \sigma_j u_j v_j^T.$$

Theorem 5.1. Set

$$A_k = \sum_{j=1}^k \sigma_j u_j v_j^T.$$

Then,

$$\min_{\substack{B \in \mathbb{C}^{m \times n} \\ \operatorname{rank}(B) \le k}} \|A - B\|_2 = \|A - A_k\|_2 = \sigma_{k+1}.$$

Proof.

$$A - A_k = \sum_{j=1}^{r} \sigma_j u_j v_j^T - \sum_{j=1}^{k} \sigma_j u_j v_j^T = \sum_{j=k+1}^{r} \sigma_j u_j v_j^T$$

and thus  $\sigma_{k+1}$  is the largest singular value of  $A-A_k$ . Alternatively, look at  $U^TA_kV=\mathrm{diag}(\sigma_1,\ldots,\sigma_k,0,\ldots,0)$ , which means that  $\mathrm{rank}(A_k)=k$ , and that

$$||A - A_k||_2 = ||U^T (A - A_k)V||_2 = ||\operatorname{diag}(0, \dots, 0, \sigma_{k+1}, \dots, \sigma_r)||_2 = \sigma_{k+1}.$$

Let B be an arbitrary matrix with  $rank(B_k) = k$ . Then, it has a null space of dimension n - k, that is,

$$\operatorname{null}(B) = \operatorname{span}(w_1, \dots, w_{n-k}).$$

A dimension argument shows that

$$\operatorname{span}(w_1, \dots, w_{n-k}) \cap \operatorname{span}(v_1, \dots, v_{k+1}) \neq \{0\}.$$

Let w be a unit vector from the intersection. Since

$$Aw = \sum_{j=1}^{k+1} \sigma_j(v_j^T w) u_j,$$

we have

$$||A - B||_2^2 \ge ||(A - B)w||_2^2 = ||Aw||_2^2 = \sum_{j=1}^{k+1} \sigma_j^2 |v_j^T w|^2 \ge \sigma_{k+1}^2 \sum_{j=1}^{k+1} |v_j^T w|^2 = \sigma_{k+1}^2,$$

since  $w \in \text{span}\{v_1, \dots, v_{n+1}\}$ , and the  $v_j$  are orthogonal.

#### 5.2 Rank-k approximation in the Frobenius norm

The same theorem holds with the Frobenius norm.

Theorem 5.2. Set

$$A_k = \sum_{j=1}^k \sigma_j u_j v_j^T.$$

Then,

$$\min_{\substack{B \in \mathbb{C}^{m \times n} \\ \operatorname{rank}(B) \leq k}} \|A - B\|_F = \|A - A_k\|_F = \sqrt{\sum_{i=k+1}^n \sigma_i^2}.$$

*Proof.* Suppose  $A = U\Sigma V^T$ . Then

$$\min_{\operatorname{rank}(B) \leq k} \left\| A - B \right\|_F^2 = \min_{\operatorname{rank}(B) \leq k} \left\| U \Sigma V^T - U U^T B V V^T \right\|_F^2 = \min_{\operatorname{rank}(B) \leq k} \left\| \Sigma - U^T B V \right\|_F^2.$$

Now,

$$\|\Sigma - U^T B V\|_F^2 = \sum_{i=1}^n (\Sigma_{ii} - (U^T B V)_{ii}))^2 + \text{off-diagonal terms.}$$

If B is the best approximation matrix and  $U^TBV$  is not diagonal, then write  $U^TBV = D + O$ , where D is diagonal and O contains the off-diagonal elements. Then the matrix  $B = UDV^T$  is a better approximation, which is a contradiction. Thus,  $U^TBV$  must be diagonal. Hence,

$$\|\Sigma - D\|_F^2 = \sum_{i=1}^n (\sigma_i - d_i)^2 = \sum_{i=1}^k (\sigma_i - d_i)^2 + \sum_{i=k+1}^n \sigma_i^2,$$

and this is minimal when  $d_i = \sigma_i$ , i = 1, ..., k. The best approximating matrix is  $A_k = UDV^T$ , and the approximation error is  $\sqrt{\sum_{i=k+1}^n \sigma_i^2}$ .

#### 5.3 Closest orthogonal matrix

The SVD also allows to find the orthogonal matrix that is closest to a given matrix. Again, suppose that  $A = U\Sigma V^T$  and W is an orthogonal matrix that minimizes  $\|A - W\|_F^2$  among all orthogonal matrices. Now,

$$\left\|U\Sigma V^T - W\right\|_F^2 = \left\|U\Sigma V^T - UU^T W V V^T\right\| = \left\|\Sigma - \tilde{W}\right\|,$$

where  $\tilde{W} = U^T W V$  is another orthogonal matrix. We need to find the orthogonal matrix  $\tilde{W}$  that is closest to  $\Sigma$ . Alternatively, we need to minimize  $\left\|\tilde{W}^T \Sigma - I\right\|_F^2$ .

If U is orthogonal and D is diagonal and positive, then

$$\operatorname{trace}(UD) = \sum_{i,k} u_{ik} d_{ki} \le \sum_{i} \left( \left( \sum_{k} u_{ik}^{2} \right)^{1/2} \left( \sum_{k} d_{ik}^{2} \right)^{1/2} \right)$$

$$= \sum_{i} \left( \sum_{k} d_{ki}^{2} \right)^{1/2} = \sum_{i} \left( d_{ii}^{2} \right)^{1/2} = \sum_{i} d_{ii} = \operatorname{trace}(D).$$
(1)

Now

$$\begin{split} \left\| \tilde{W}^T \Sigma - I \right\|_F^2 &= \operatorname{trace} \left( \left( \tilde{W}^T \Sigma - I \right) \left( \tilde{W}^T \Sigma - I \right)^T \right) \\ &= \operatorname{trace} \left( \left( \tilde{W}^T \Sigma - I \right) \left( \Sigma \tilde{W} - I \right) \right) \\ &= \operatorname{trace} \left( \tilde{W}^T \Sigma^2 \tilde{W} \right) - \operatorname{trace} \left( \tilde{W}^T \Sigma \right) - \operatorname{trace} \left( \Sigma \tilde{W} \right) + n \\ &= \operatorname{trace} \left( \left( \Sigma \tilde{W} \right)^T \left( \Sigma \tilde{W} \right) \right) - 2 \operatorname{trace} \left( \Sigma \tilde{W} \right) + n \\ &= \left\| \Sigma \tilde{W} \right\|_F^2 - 2 \operatorname{trace} \left( \Sigma \tilde{W} \right) + n \\ &= \left\| \Sigma \right\|_F^2 - 2 \operatorname{trace} \left( \Sigma \tilde{W} \right) + n. \end{split}$$

Thus, we need to maximize trace  $(\Sigma \tilde{W})$ . But this is maximized by  $\tilde{W} = I$  by (1). Thus, the best approximating matrix is  $W = UV^T$ .

#### 6 The "Thin" SVD

Also called "economy size" SVD. If  $A \in \mathbb{C}^{m \times n}$ ,  $A = U \Sigma V^T$ , and  $m \geq n$ , then the "thin" SVD is  $A = U_1 \Sigma_1 V^T$  where

$$U_1 = [u_1, \dots, u_n] \in \mathbb{C}^{m \times n}$$

and

$$\Sigma_1 = \operatorname{diag}(\sigma_1, \dots, \sigma_n) \in \mathbb{R}^{n \times n}.$$

# 7 Applications of the SVD

- 1. Determining range, null space and rank (also numerical rank).
- 2. Matrix approximation.
- 3. Inverse and Pseudo-inverse: If  $A = U\Sigma V^T$  and  $\Sigma$  is full rank, then  $A^{-1} = V\Sigma^{-1}U^T$ . If  $\Sigma$  is singular, then its pseudo-inverse is given by  $A^{\dagger} = V\Sigma^{\dagger}U^T$ , where  $\Sigma^{\dagger}$  is formed by replacing every nonzero entry by its reciprocal.

- 4. Least squares: If we need to solve Ax=b in the least-squares sense, then  $x_{LS}=V\Sigma^{\dagger}U^Tb$ .
- 5. Denoising Small singular values typically correspond to noise. Take the matrix whose columns are the signals, compute SVD, zero small singular values, and reconstruct.
- 6. Compression We have signals as the columns of the matrix S, that is, the i signal is given by

$$S_i = \sum_{i=1}^r \left(\sigma_j v_{ij}\right) u_j.$$

If some of the  $\sigma_i$  are small, we can discard them with small error, thus obtaining a compressed representation of each signal. We have to keep the coefficients  $\sigma_j v_{ij}$  for each signal and the dictionary, that is, the vectors  $u_i$  that correspond to the retained coefficients.

## 8 Differences between SVD and eigen-decomposition

- 1. Not every matrix has an eigen-decomposition (not even any square matrix). Any matrix (even rectangular) has an SVD.
- 2. In eigen-decomposition  $A = X\Lambda X^{-1}$ , that is, the eigen-basis is not always orthogonal. The basis of singular vectors is always orthogonal.
- 3. In SVD we have two singular-bases (right and left).
- 4. SVD tells everything on a matrix.
- 5. SVD as no numerical problems.
- 6. Relation to condition number; the numerical problems with eigen-decomposition; multiplication by an orthogonal matrix is perfectly conditioned.

# 9 Linear regression in the least-squared loss

In Linear regression we aim to find the best linear approximation to a set of observed data. For the m data points  $\{x_1, \ldots, x_m\}$ ,  $x_i \in \mathbb{R}^n$ , each receiving the value  $y_i$ , we look for the weight vector w that minimizes:

$$\sum_{i=1}^{n} (x_i^T w - y_i)^2 = ||Aw - y||_2^2$$

Where A is a matrix that holds the data points as rows  $A_i = x_i^T$ .

**Proposition 9.1.** The vector w that minimizes  $||Aw - y||_2^2$  is  $w = A^{\dagger}y = V\Sigma^{\dagger}U^Ty$  for  $A = U\Sigma V^T$  and  $\Sigma_{ii}^{\dagger} = 1/\Sigma_{ii}$  if  $\Sigma_{ii} > 0$  and 0 else.

Let us define  $U_{\parallel}$  and  $U_{\perp}$  as the parts of U corresponding to positive and zero singular values of A respectively. Also let  $y_{\parallel}=0$  and  $y_{\perp}$  be two vectors such that  $y=y_{\parallel}+y_{\perp}$  and  $U_{\parallel}y_{\perp}=0$  and  $U_{\perp}y_{\parallel}=0$ .

Since  $y_{\parallel}$  and  $y_{\perp}$  are orthogonal we have that  $\|Aw - y\|_2^2 = \|Aw - y_{\parallel} - y_{\perp}\|_2^2 = \|Aw - y_{\parallel}\|_2^2 + \|y_{\perp}\|_2^2$ . Now, since  $y_{\parallel}$  is in the range of A there is a solution w for which  $\|Aw - y_{\parallel}\|_2^2 = 0$ . Namely,  $w = A^{\dagger}y = V\Sigma^{\dagger}U^Ty$  for  $A = U\Sigma V^T$ . This is because  $U\Sigma V^TV\Sigma^{\dagger}U^Ty = y_{\parallel}$ . Moreover, we get that the minimal cost is exactly  $\|y_{\perp}\|_2^2$  which is independent of w.

## 10 Optimal squared loss dimension reduction

Given a set of n vectors  $x_1, \ldots, x_n$  in  $\mathbb{R}^m$ . We look for a rank k projection matrix  $P \in \mathbb{R}^{m \times m}$  that minimizes:

$$\sum_{i=1}^{n} ||Px_i - x_i||_2^2$$

If we denote by A the matrix whose i'th column is  $x_i$  then this is equivalent to minimizing  $||PA - A||^2_{Fro}$  Since the best possible rank k approximation to the matrix A is  $A_k = \sum_{i=1}^k \sigma_i u_i v_i^T$  the best possible solution would be a projection P for which  $PA = A_k$ . This is achieved by  $P = U_k U_k^T$  where  $U_k$  is the matrix corresponding to the first k left singular vectors of A.

If we define  $y_i = U_k^T x_i$  we see that the values of  $y_i \in \mathbb{R}^k$  are optimally fitted to the set of points  $x_i$  in the sense that they minimize:

$$\min_{y_1, ..., y_n} \min_{\Psi \in \mathbb{R}^{k \times m}} \sum_{i=1} ||\Psi y_i - x_i||_2^2$$

The mapping of  $x_i \to U_k^T x_i = y_i$  thus reduces the dimension of any set of points  $x_1, \ldots, x_n$  in  $\mathbb{R}^m$  to a set of points  $y_1, \ldots, y_n$  in  $\mathbb{R}^k$  optimally in the squared loss sense. This is commonly referred to as Principal Component Analysis (PCA).