

**Course: Data Structures and Algorithms** 

# Fundamental Data Structures Arrays and Linked Lists





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#### Arrays vs. Linked Lists: The Eternal Struggle





Arrays are great... until they aren't!



#### Why study arrays and linked lists



- Arrays store data in order for fast access
- Linked Lists handle dynamic data without fixed size
- Used in web applications, databases, operating systems, games, and so many more applications

They help **store** and **manage data** efficiently in programs and realworld applications

**Array Fundamentals** 

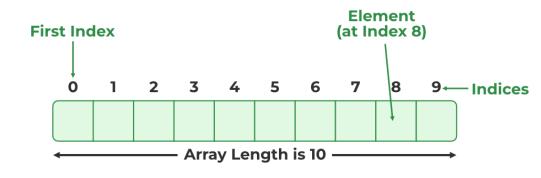
### What is an array: definition and properties





An array is a fixed-size collection of elements stored in continuous memory locations

- Stores multiple values of the same type in order
- Elements are accessed using an index starting from 0
- Memory is allocated in a single, continuous block

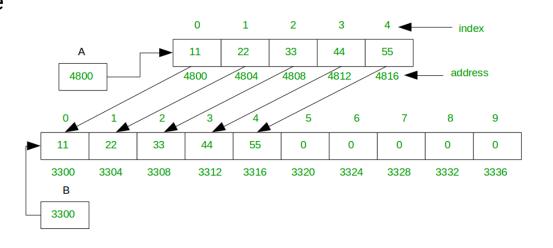


Efficient for fast lookups but fixed in size

#### Static vs dynamic arrays



- Static arrays have a fixed size set at creation
- Dynamic arrays resize automatically when needed
- Static arrays are faster but waste memory if too large
- Dynamic arrays are flexible but need extra resizing time

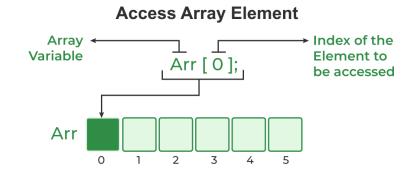


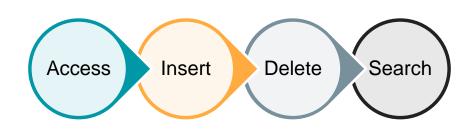
In Java, ArrayLists act as dynamic arrays by automatically resizing when needed

#### Array operations overview



- Access: retrieve any element instantly using its index
- Insert: add at index, may require shifting elements
- Delete: remove elements, shifting needed to fill the gap
- Search: finding an element using linear or binary search





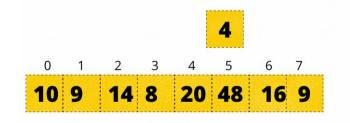
#### Insertion in arrays



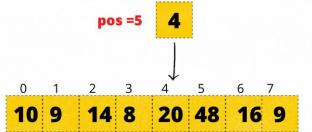


### Inserting elements in an array may involve shifting elements for space

- Insert at the end: O(1) time complexity, no shifts needed
- Insert at middle: O(n) time complexity, elements need shifting
- Shifting: Involves moving elements to create space for the new element



#### Add an array element at specific position



Click to see interactively how an element is inserted

#### Search in arrays: linear search





### Search for elements using linear or binary search based on array order

- Linear search: check each element sequentially until the target is found or array ends
- Simple but slow for large arrays
- O(1) for best case; O(n) for average and worst case

Linear Search



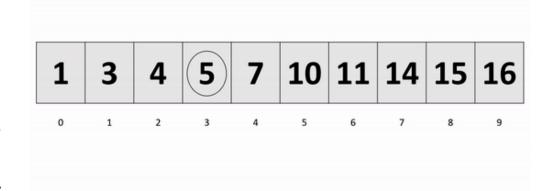
#### Search in arrays: binary search





## Search for elements using linear or binary search based on array order

- Binary search: divides a <u>sorted</u> <u>array</u> in half, repeatedly narrowing down search range
- Faster but requires sorted data
- O(1) for best case;  $O(\log n)$  for average and worst case



#### Time and space complexity of arrays





### Arrays offer **fast access** but have limitations in insertion, deletion, and resizing

Operation	Complexity
Access	0(1)
Insert (end)	0(1)
Insert (middle)	O(n)
Delete (end)	0(1)
Delete (middle)	O(n)
Linear search	O(n)
Binary search	$O(\log n)$

Use arrays for fast access and fixed-size data, but avoid frequent manipulations

### Linked List Fundamentals

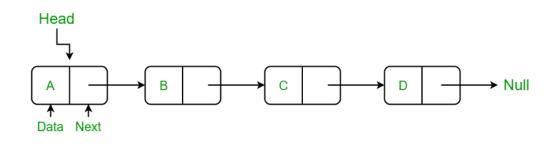
#### What is a list: concept and node structure





## A linked list is a collection of nodes, where each node points to the next

- Nodes contain data and a reference to the next node
- Head points to the first node, last node points to null
- Dynamic size allows efficient insertions and deletions



Used when data is frequently modified or added/removed

#### Singly linked list (SLL)



A singly linked list has nodes that point only to the next node

```
class Node {
   int data;
   Node next;

   Node(int data) {
      this.data = data;
      this.next = null;
   }
}
```

```
class SinglyLinkedList {
   Node head;
    SinglyLinkedList() {
        head = null;
   void add(int data) {
        // To be implemented
   void remove(int data) {
       // To be implemented
   void display() {
     // To be implemented
```

#### Key operations in lists: traversal

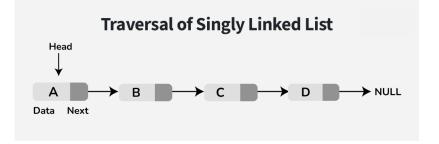




### Traversal involves visiting each node from head to tail to access all data

- Start at the head and visit each node sequentially
- Continue until the last node (next is null)
- Used for printing or searching data in the list

```
void traverse() {
   Node current = head;
   while (current != null) {
       System.out.print(current.data + " ");
       current = current.next;
   }
}
```



#### Key operations in lists: insertion





### Insertion can be done at the beginning, middle, or end of the list

- Insert at the beginning: new node becomes head
- Insert at the end: traverse and add after the last node
- Insert in the middle: traverse to the position and link the new node

```
void insertAtEnd(int data) {
   Node newNode = new Node(data);
    if (head == null) {
        head = newNode:
   } else {
        Node current = head;
        while (current.next != null) {
            current = current.next;
        current.next = newNode;
```

Click here for an interactive visualization

#### Key operations in lists: deletion

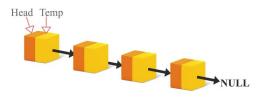




### Deletion can be done at the beginning, middle, or end of the list

- Delete the first node: move the head to the next node
- Delete the last node: traverse to the second-to-last node and unlink
- Delete a specific node: traverse to the node before and unlink it

```
void deleteFirst() {
   if (head != null) {
     head = head.next;
   }
}
```



Never change or remove the head without ensuring the list stays intact

Click here for an interactive visualization

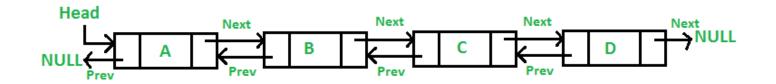
#### Other types of lists: doubly and circular



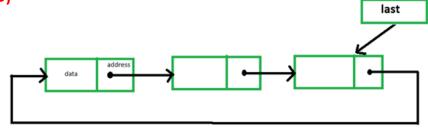


### Doubly and circular linked lists offer more flexibility for complex operations





Singly linked lists are like one-way streets; sometimes you need a U-turn!



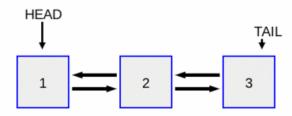
#### Doubly linked lists (DLL)





### DLLs allow navigation in both directions, making insertions and deletions more flexible

- Each node has two references: next and prev
- Traversal from head to tail and tail to head
- Used in complex data structures like deques and browser history



Click here for an interactive visualization

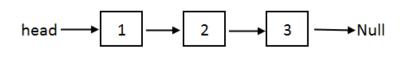
#### Circular linked lists (CLL)



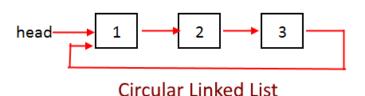


### CLLs have a circular structure, making them ideal for continuous cycles and buffers

- Tail node points to the head, forming a loop
- Traversal starts from head and loops back after reaching the tail
- Used in applications needing repetitive cycles (e.g., buffers, games)



Singly Linked List



Click here for an interactive visualization

## Key Characteristics and Applications of Linked List Types



#### Singly Linked Lists

- One-direction traversal, simple structure
- Basic data storage, simple queues, and stacks
- Simple operations, space efficiency, and low overhead

#### **Doubly Linked Lists**

- Two-way traversal, more complex node structure
- Deques, browser history, undo functionality
- Bidirectional navigation, efficient insertions/deletions at both ends

#### Circular Linked Lists

- Circular structure, continuous looping
- Circular queues,
   round-robin
   scheduling, buffers
- Continuous cycles and circular data management

#### Complexity of key operations in lists



- Head operations: Insertion and deletion at the head are always O(1)
- Tail/Middle operations: Insertion and deletion from the tail or middle require O(n) due to linear traversal
- Search: Searching requires O(n)
   as each node must be visited
   one by one

Operation	Complexity
Insertion (at head)	0(1)
Insertion (at tail/middle)	O(n)
Deletion (from head)	0(1)
Deletion (from tail/middle)	0(n)
Search	O(n)

Understanding complexity
helps optimize performance
and resource usage

#### Real world applications of linked lists





Web browsers

Manage the history of visited pages going back and forth



Handle playlists to add and remove files



Maintain and manipulate history of user actions



Polynomial arithmetic

Represent polynomials and conduct operations efficiently



Graph representations

<u>Used as helper data structures</u> <u>to build graphs</u>

Arrays versus Linked Lists

#### Arrays versus linked lists: the overgoing battle



Feature	Arrays	Linked List
Memory allocation	Contiguous block, fixed size	Dynamic allocation, flexible size
Access time	O(1) for random access	O(n) for random access
Insertion/Deletion	O(n) (shifting needed)	O(1) at head/tail, O(n) elsewhere
Search	O(n) (sequential search)	O(n) (sequential search)
Memory efficiency	Inefficient for dynamic size	Efficient for dynamic size
Resizing	Fixed size (resize expensive)	Dynamic resizing (efficient)

Arrays are best when you need quick access and a fixed size





Linked lists work well for data that changes often, but access is slower and less efficient

#### Common mistakes and pitfalls



Assuming fixed size for arrays: arrays cannot grow after creation;
 use ArrayList for dynamic sizing



 Not handling null in linked lists: always check for null pointers when traversing or modifying linked lists



 Inefficient insertion in arrays: inserting in the middle of an array requires shifting elements, which can be slow



 Overusing linked lists for simple data: linked lists are overkill when simple arrays or arraylists suffice



### Key Takeaways



- Arrays and linked lists are fundamental data structures with distinct strengths and use cases
- Arrays provide fast access and efficient memory for fixed-size data
- Linked Lists are ideal for dynamic data that changes frequently
- Arrays are used in matrices, buffers, and lookup tables;
   Linked Lists in dynamic data like playlists, browser history,
   and undo functionality

#### Helpful resources on arrays and linked lists



- VisualAlgo: an interactive platform that visualizes data structures and algorithms
- <u>Responsible</u> BroCode video to learn linked lists in a concise and interactive manner
- Programiz: beginner-friendly tutorials on data structures and algorithms with clear explanations and examples
- <u>FreeCodeCamp</u> course on algorithms and data structures

#### Quote of the Week



Arrays provide speed in simplicity, but linked lists offer flexibility in complexity



