Bubble Sort

(A Taste of Algorithms: Definition, Design, and Analysis)

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Bubble Sort

1. The Sorting Problem
2. Bubble Sort
3. Analysis of Bubble Sort

Algorithms

What is an algorithm? What is computation?

$$\xrightarrow{\text{Input}} \boxed{\begin{array}{c}\text{Algorithm}\\ \text{(sequence of operations)}\end{array}} \xrightarrow{\text{Output}}$$

Correctness!

Definiteness: precisely defined operations

Finiteness: termination

Effectiveness: a reasonable model; basic operations

- for sorting: compare, swap

Sorting

The sorting problem:

Input: A sequence of n integers A: $a_1\ a_2\ \cdots\ a_n$.

Output: A permutation $a'_1\ a'_2\ \ldots\ a'_n$ of A *s.t.* $a'_1 \le a'_2 \le \cdots \le a'_n$ (non-decreasing order).

$$3 \quad 1 \quad 4 \quad 2 \quad \Longrightarrow 1 \quad 2 \quad 3 \quad 4$$

Inversions

$$A = a_1 \quad a_2 \quad \ldots \quad a_n.$$

If $i < j$ and $a_i > a_j$, then (a_i, a_j) is an **inversion**.

Adjacent inversion: (a_i, a_{i+1})

3 1 4 2

#inversions = 3
#adjacent inversions = 2

4 3 2 1

#inversions = 3 + 2 + 1 = 6
#adjacent inversions = 3

1 2 3 4

#inversions = 0
#adjacent inversions = 0

Inversions

Theorem: A is sorted $\iff$ A has no adjacent inversions.

$$\begin{aligned} A \text{ is sorted } &\implies A \text{ has no inversions} \\ &\implies A \text{ has no adjacent inversions.} \end{aligned}$$

$$\begin{aligned} A \text{ has no adjacent inversions } &\implies \forall i \in [1, n-1] : a_i \le a_{i+1} \\ &\implies A \text{ is sorted.} \end{aligned}$$

Bubble Sort

Bubble Sort: Basic Idea

Basic idea: to eliminate all adjacent inversions

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repeat
  pick any i $\triangleright$ **Definiteness!**
  if $a_i > a_{i+1}$ **then**
    $\textsc{Swap}(a_i, a_{i+1})$
until no adjacent inversions $\triangleright$ **Finiteness! Definiteness!**

Bubble Sort: Definiteness

procedure BubbleSort($A : a_1\ a_2\ \cdots\ a_n$)

  repeat

    $\boxed{\text{swapped} \leftarrow \text{false}}$

    for $\boxed{i \leftarrow 1 : n-1}$ **do** ▷ Pick i

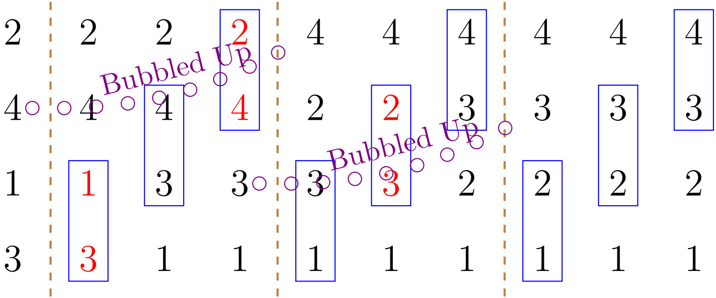
      if $a_i > a_{i+1}$ **then**

        Swap(a_i, a_{i+1})

        $\boxed{\text{swapped} \leftarrow \text{true}}$

  until $\boxed{\text{no adjacent inversionsswapped} = \text{false}}$

Bubble Sort: Example



After each "**for**" loop, one more element is bubbled up to its final position.

Bubble Sort: Finiteness

```
procedure BubbleSort(A : a_1 a_2 ··· a_n)
    repeat
        swapped ← false
        for i ← 1 : n − 1 do                ▷ Pick i
            if a_i > a_{i+1} then
                Swap(a_i, a_{i+1})
                swapped ← true
    until swapped = false                   ▷ No swaps
```

The inner "**for**" loops:

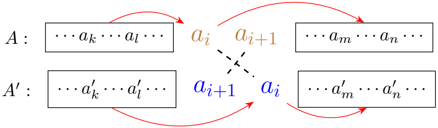
1) $\exists$ loop : no swaps $\implies$ swapped = false $\implies$ terminates

2) $\forall$ loop : has swaps **Impossible!**

Bubble Sort: Finiteness

$$\boxed{\text{Fact: total \#inversions is finite.}}$$

Effects of $\textsc{Swap}(a_i, a_{i+1})$ on #inversions[1]:



$$-1 : (a_i, a_{i+1})$$

$+0$: relative order between any other two elements does not change!

$$(a_k, a_l), (a_m, a_n), (a_k, a_i), (a_i, a_m)$$

$$\boxed{\textsc{Swap}(a_i, a_{i+1}) \implies -1 \text{ inversion}}$$

[1] Not on #adjacent inversions! Think about it.

Bubble Sort: Correctness

$$
\begin{aligned}
\textbf{Finiteness} \quad &\Longrightarrow \exists \text{ loop} : \text{ no swaps} \\
&\Longrightarrow A \text{ has no adjacent inversions any more} \\
&\Longrightarrow A \text{ is already sorted.}
\end{aligned}
$$

Optimizing Bubble Sort (I)[2]

After each "**for**" loop, one more element is bubbled up to its final position.

procedure BubbleSort($A : a_1 \; a_2 \; \cdots \; a_n$)
    $n \leftarrow \text{len}(A)$
    repeat
        swapped $\leftarrow$ false
        for $i \leftarrow 1 : n - 1$ **do**
            if $a_i > a_{i+1}$ **then**
                Swap(a_i, a_{i+1})
                swapped $\leftarrow$ true
        $n \leftarrow n - 1$ $\triangleright$ One maximal bubbles up
    until swapped $=$ false

[2] See Appendix for "Optimizing Bubble Sort (II)".

Bubble Sort

Time Complexity of Bubble Sort

- Finiteness is NOT enough $\implies$ Quantitative finiteness
- Time on real computers varies $\implies$ #Ops on our model:

$|C|$: #Comparisons (**if** $a_i > a_{i+1}$)

$|S|$: #Swaps $(\textsc{Swap}(a_i, a_{i+1}))$

$$|C| \geq |S|$$

- Different inputs $\implies$ $|C|$ and $|S|$ vary:
  - Best-case, worst-case, and average-case analysis

Best-case and Worst-case Analysis

Best-case: $1 \quad 2 \cdots n$

$$
\begin{aligned}
& \qquad\qquad \underline{\text{Best-case:}} \qquad\qquad\qquad \underline{\text{Worst-case:}} \\
& \qquad \text{non-decreasingly sorted} \qquad \text{non-increasingly sorted} \\
|C| &= (\min : n-1, \qquad\qquad \max : \tfrac{n^2-n}{2}); \\
|S| &= (\min : 0, \qquad\qquad\qquad \max : \tfrac{n^2-n}{2}).
\end{aligned}
$$

Worst-case: $n \quad n-1 \cdots 1$

$$\#\text{inversions} = (n-1) + (n-2) + \cdots + 1 = \frac{n^2-n}{2}.$$

$|S|$: #Swaps (Average-case Analysis)[3]

Assumptions on inputs:

1. The input is a random permutation ("average input")
2. All numbers are all different (for simplicity)

$$\boxed{\text{SWAP}(a_i, a_{i+1}) \implies -1 \text{ inversion}}$$

$$|S| = \mathbb{E}(\#\text{inversions})$$

[3] An exercise: what is $|C|$ (#Comparisions) in average?

$|S|$: #Swaps (Average Analysis)

$$I_{ij} = \begin{cases} 1 & (a_i, a_j) \text{ is an inversion} \\ 0 & \text{o.w.} \end{cases}$$

$$X = \sum_{1 \le i < n} \sum_{i < j \le n} I_{ij} \qquad (\#\text{inversions})$$

$$\mathbb{E}(X) = \mathbb{E}(\sum_i \sum_{j>i} I_{ij}) = \sum_i \sum_{j>i} \mathbb{E}(I_{ij}) \quad (\text{linearity of expectation})$$

$$\mathbb{E}(I_{ij}) = \mathbb{P}\{I_{ij} = 1\} = \frac{1}{2} \qquad (a_i \neq a_j; \text{half: } a_i < a_j, \text{half: } a_i > a_j)$$

$$\mathbb{E}(X) = \sum_i \sum_{i<j} \frac{1}{2} = \binom{n}{2} \cdot \frac{1}{2} = \frac{n(n-1)}{4} = O(n^2)$$

Faster Algorithms

It took a good deal of work to analyze the bubble sort; and although [. . .], the results are disappointing since they tell us that the bubble sort isn't really very good at all.

— Donald E. Knuth

faster: $O(n^2) \to O(n \lg n)$?

. . . and faster: $O(n \lg n) \to O(n)$?

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4. Appendix

Bubble Sort: Correctness

procedure BUBBLESORT($A : a_1\ a_2\ \cdots\ a_n$)
  repeat
    swapped $\leftarrow$ false
    for $i \leftarrow 1 : n-1$ **do** ▷ **Loop invariant?**
      if $a_i > a_{i+1}$ **then** ▷ CAS
        SWAP(a_i, a_{i+1})
        swapped $\leftarrow$ true
  until swapped = false

Loop invariant:

Before the k-th ($k \geq 1$) "**for**" loop, $a_{n-(k-1)} \cdots a_n$
(1) consists of the largest $(k-1)$ elements
(2) in sorted order.

Correctness: Initialization + Maintenance + Termination

Bubble Sort: Finiteness

Idea: well-founded relation over $\mathbf{N}$

Effects of $\textsc{Swap}(a_i, a_{i+1})$ on adjacent inversions:

$$5 \quad 3 \quad 1 \quad 6 \quad 7 \quad 2 \quad 4 \quad 8$$

$$5 \quad 3 \quad 1 \quad 6 \quad 2 \quad 7 \quad 4 \quad 8$$

$-1 : (7, 2)$ $\qquad$ $+2 : (6, 2), (7, 4)$

Optimizing Bubble Sort (II)

Idea: After each "**for**" loop, all elements after "lsp" are settled.

```
procedure BubbleSort(A : a_1 a_2 ··· a_n)
    repeat
        swapped ← false
        lsp ← 0                                  ▷ lsp: the last swap position
        for i ← 1 : n − 1 do
            if a_i > a_{i+1} then
                Swap(a_i, a_{i+1})
                swapped ← true
                lsp ← i                          ▷ Update lsp
        n ← lsp                                  ▷ Elements after lsp are sorted
    until swapped = falselsp = 0
```