(What are Algorithms and How to Analyze Algorithms)

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- Sorting
- 2 Bubble Sort
- 3 Analysis

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# Sorting

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# Sorting

#### The sorting problem:

Given a sequence A of sortable elements, arrange them in ascending/descending order.

$$3 \ 1 \ 4 \ 2 \Longrightarrow 1 \ 2 \ 3 \ 4$$

A little more formalism: ordering relation "<" on A.

 $\forall a, b, c \in A$ ,

Trichotomy: a < b, a = b, a > b

Transitivity:  $a < b \land b < c \implies a < c$ 



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Input: an integer array A

Ouput: A' sorted

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Roughly, what is computation?

$$A = a_1, a_2, \dots, a_n.$$

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$$\iff: \forall i \in [1, n-1] : a_i \le a_{i+1}.$$



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#### Basic idea: to eliminate all adjacent inversions

```
repeat
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until no adjacent inversion
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#### Definiteness!

```
1 repeat
2 swapped = false
3 for i = 1 to n-1
4 if a_{i-1} > a_i
5 swap(a_{i-1}, a_i)
6 swapped = true
7 until not swapped
```

 $5 \ 3 \ 1 \ 6 \ 7 \ 2 \ 4 \ 8$ 

```
repeat

swapped = false

for i = 1 to n-1

if a_{i-1} > a_i

swap(a_{i-1}, a_i)

swapped = true

until not swapped
```

#### Finiteness!

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#### Proof.

The inner "for" loop:

- I. not swapped  $\implies$  terminates
- II. counting #inversions
  - swap $(a_{i-1}, a_i) \implies -1$  inversion; + 0 inversion
  - total #inversions is finite

# Optimizing Bubble Sort

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# Time Complexity of Bubble Sort

Finiteness



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#### Finiteness

```
|P|: #Passes (the "for" loops)

|C|: #Comparisons (if a_{i-1} > a_i)

|S|: #Swaps (swap(a_{i-1}, a_i))
```

# Time Complexity of Bubble Sort

#### Finiteness

```
|P|: \mbox{\#Passes} \qquad \qquad (\mbox{the "for" loops}) \\ |C|: \mbox{\#Comparisons} \qquad \qquad (\mbox{if } a_{i-1} > a_i) \\ |S|: \mbox{\#Swaps} \qquad \qquad (\mbox{swap}(a_{i-1}, a_i))
```

Different inputs  $\iff$  different execution time:

Best-case, Worst-case, and Average-case Analysis

	Best-case:	Worst-case:			
P	= (	);			
C	= (	);			
S	= (	).			



Best-case: 1 2 3 4 5 6 7 8

	Best-case:	Worst-case:			
	ascendingly sorted				
P	= (	);			
C	= (	);			
S	= (	).			



Best-case: 1 2 3 4 5 6 7 8

	Best-case:	Worst-case:			
	ascendingly sorted				
P	$= (\min: 1,$	);			
C	$= (\min: n-1,$	);			
S	$= (\min: 0,$	).			



Best-case: 1 2 3 4 5 6 7 8

	$\frac{\text{Best-case:}}{\text{ascendingly sorted}}$	Worst-case: descendingly sorted					
P	$= (\min: 1,$		);				
C	$= (\min: n-1,$		);				
S	$= (\min: 0,$		).				
Wo	rst-case: 8 7 6	5	4	3	2	1	



Best-case: 1 2 3 4 5 6 7 8

 $\frac{\text{Best-case:}}{\text{ascendingly sorted}} \quad \frac{\text{Worst-case:}}{\text{descendingly sorted}}$   $|P| = (\min : 1, \quad \max : n);$   $|C| = (\min : n - 1, \quad \max : \frac{n^2 - n}{2});$   $|S| = (\min : 0, \quad \max : \frac{n^2 - n}{2}).$   $\text{Worst-case: } 8 \quad 7 \quad 6 \quad 5 \quad 4 \quad 3 \quad 2 \quad 1$ 



|S|: #Swaps