### hhuOS

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### Heinrich Heine University Facts about hhuOS

- A small operating system for teaching and learning purposes
- written for x86 32-bit architecture
- written in C++ and x86-Assembler using gcc and nasm
- Open-Source, published under the GPL v3 license
- $\sim 25000$  lines of code



- Processes & Threads
  - Round-Robin based preemptive scheduling
  - Binary files are executed as processes
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  - Round-Robin based preemptive scheduling
  - Binary files are executed as processes
  - Each process has its own address space
- Adress spaces and memory management
  - Using x86-Paging mechanism (virtual/physical memory)
  - Lazy mapping implementation
  - Kernel address space is always mapped at 3 GiB (not accessible from user space)
  - Different memory managing algorithms are implemented (Free List, Bitmap, Table)



- Unified Library
  - Single codebase for user- and kernel-space library
  - Functions requiring kernel access are outsourced into an interface (implemented two times - for kernel and user space)
  - Kernel access from user space via system calls (software interrupts)



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- Further features
  - Virtual Filesystem (can mount physical filesystems)
  - Hardware support (Keyboard, PCI, Floppy AHCI, VESA & CGA Graphics)
  - Multiboot compatible: Bootable on UEFI and BIOS systems (BIOS calls are supported)
  - Own UEFI bootloader (towboot, developed by Niklas Sombert)
  - Experimental support for networking, thanks to several bachelors thesis (2 drivers and a UDP/IP Stack)



# Overview: Memory & Paging

- Paging: Abstract physical memory from virtual address spaces
- New pages can be mapped in/out dynamically
- Use of different address spaces for process separation
- $lue{}$  Kernel is mapped at 3 GiB ightarrow Higher-Half-Kernel (always visible)
- Adresses below 3 GiB are used for user space memory



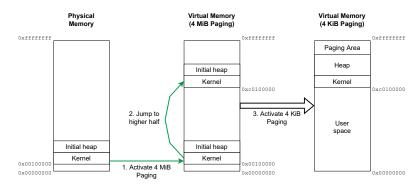
### Paging: The boot process

- How to allocate memory when no memory manager is available?
- How to map the Kernel-code at 3 GiB without losing the EIP?
- Solution: Activate paging in several steps



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# Allocating heap memory

- Invoke Free list manager to find free memory block in heap
- 2 Slice free block (with respect to alignment)
- 3 Return pointer to found block
- If the block was not previously mapped, the first access will generate a pagefault
- 5 Fault handler is invoked during interrupt handling
- 6 Search unused pageframe and map it to the virtual fault address
- Return to program



# Storage Devices

- StorageDevice as interface for block devices
- Only 4 methods: getSectorCount(), getSectorSize(), read(), write()
- Implemented for AHCI (hard drives), Floppy, Virtual Drives, and Partitions

# hhu Heinrich Heine University Disselder Virtual File System

- Overlay over physical file systems (e.g. FAT)
- Storage devices can be mounted to any folder
- Files and directories are represented by Filesystem::Node (similar to INode)
- Interface Filesystem::Driver for physical file systems
- Integrate FatFs¹ to support all FAT variants

<sup>1</sup> http://elm-chan.org/fsw/ff/00index\_e.html



# **Process Scheduling**

- Global scheduler manages processes in a queue
- Timer invokes scheduler every 10ms to switch to next process (and address space)
- Each process has a separate thread scheduler, switching the current thread each time the process becomes active
- Threads can be blocked and unblocked by other threads
- Threads can sleep for a defined time
- It is possible to join other threads (or processes)

### **Future Work**

- Implement more device drivers (sound, graphics, network, etc.)
- Implement new library functions (libc, graphics/game, etc.)
- Enhance network stack
- Better scheduling (priorities, I/O management)
- New (custom) filesystems
- Multicore support
- Support modern x86-Features (Physical Address Extensions, Long mode, APIC, HPET, etc.)