

Data Structures and Algorithms
with applications in Machine Learning
- MCQ 2 -

NAME: _____

GROUP: _____

Each Question: 1 Mark

Duration: 20 Minutes

Completely fill the circles as shown: ○○●○

Answer sheet

Q1. ○ a.
 ○ b.
 ○ c.
 ○ d.

Q2. ○ a.
 ○ b.
 ○ c.
 ○ d.

Q3. ○ a.
 ○ b.
 ○ c.
 ○ d.

Q4. ○ a.
 ○ b.
 ○ c.
 ○ d.

Q5. ○ a.
 ○ b.
 ○ c.
 ○ d.

Q6. ○ a.
 ○ b.
 ○ c.
 ○ d.

Q7. ○ a.
 ○ b.
 ○ c.
 ○ d.

Q8. ○ a.
 ○ b.
 ○ c.
 ○ d.

Q9. ○ a.
 ○ b.
 ○ c.
 ○ d.

Q10. ○ a.
 ○ b.
 ○ c.
 ○ d.

The Quiz

Q. 1 According to the definition of an instability in a matching:

- ☐ a. A pair (m, w) is unstable if the man m prefers another woman w' over w , but w' does not prefer m over her current match.
- ☐ b. A matching is unstable if there are participants left unmatched.
- ☐ c. A matching is unstable if every pair (m, w) prefers their current partners over all others.
- ☐ d. A pair (m, w) is unstable if both m and w prefer each other over their current matches.

Q. 2 A matching $\phi : \mathcal{M} \rightarrow \mathcal{W}$ is considered stable if:

- ☐ a. Every participant is matched to exactly one partner, regardless of preferences.
- ☐ b. There is no instability, meaning no pair (m, w) exists such that m and w prefer each other over their current matches.
- ☐ c. No man m prefers any woman w' over his current match, regardless of w' 's preference.
- ☐ d. The matching ensures that every participant is paired with their least preferred valid partner.

Q. 3 Regarding the *laissez-faire* approach to resolving instabilities:

- ☐ a. It guarantees that the process will always terminate with a stable matching.
- ☐ b. It resolves all instabilities simultaneously, ensuring optimal efficiency.
- ☐ c. It resolves instabilities iteratively but does not guarantee convergence to a stable matching.
- ☐ d. It avoids resolving instabilities and instead focuses on random pairings.

Q. 4 Which of the following statements about the Gale-Shapley algorithm is true?

- ☐ a. The Gale-Shapley algorithm always results in the most optimal matching for both men and women simultaneously.
- ☐ b. The Gale-Shapley algorithm terminates only if there are no instabilities but does not guarantee a stable matching.
- ☐ c. The Gale-Shapley algorithm guarantees a stable matching for any set of preferences.
- ☐ d. The Gale-Shapley algorithm ensures a matching, but stability depends on the initial preferences.

Q. 5 The following pseudo-code implements the Gale-Shapley algorithm:

Algorithm 1 Gale-Shapley Algorithm

Require: Lists of preferences for men \mathcal{M} and women \mathcal{W} **Ensure:** Stable matching ϕ

```
1: All men and women start as free
2: while  $\exists$  free man  $m$  who has not proposed to every woman on his list do
3:   Pick such a man  $m$ 
4:   Let  $w$  be the next woman on  $m$ 's preference list
5:   if _____ then                                     ▷ Fill in the blank
6:     Engage  $m$  and  $w$ 
7:   else
8:      $w$  rejects  $m$ 
9:   end if
10: end while
```

What should replace the blank in the condition?

- ☐ a. w is free or prefers m to her current match
- ☐ b. w is free and prefers m to her current match
- ☐ c. w prefers her current match to m
- ☐ d. w is free but has already rejected m

Q. 6 The following function implements the Gale-Shapley stable matching algorithm.

```
def gale_shapley(men_prefs, women_prefs):
    """
    Implementation of Gale-Shapley stable matching algorithm.
    """
    n = len(men_prefs)
    free_men = list(range(n))
    proposals = np.zeros(n, dtype=int)
    women_partners = [-1] * n

    while free_men:
        m = free_men[0]
        w = men_prefs[m][proposals[m]]
        proposals[m] += 1

        if women_partners[w] == -1:
            women_partners[w] = m
            free_men.pop(0)
        else:
            current_m = women_partners[w]
            if _____: # Fill in the blank
                women_partners[w] = m
                free_men.pop(0)
                free_men.append(current_m)

    matches = {(m, w) for w, m in enumerate(women_partners)}
    return matches
```

What should replace the blank?

- ☐ a. `list(women_prefs[m]).index(w) < list(women_prefs[m]).index(current_m)`
- ☐ b. `list(women_prefs[w]).index(m) < list(women_prefs[w]).index(current_m)`
- ☐ c. `m > current_m`
- ☐ d. `current_m == -1`

Q. 7 Consider the following preference lists for 4 men and 4 women:

```
men_prefs = [  
    [0, 1, 2, 3], # Man 0's preferences  
    [2, 0, 1, 3], # Man 1's preferences  
    [1, 3, 0, 2], # Man 2's preferences  
    [0, 2, 3, 1]  # Man 3's preferences  
]  
  
women_prefs = [  
    [2, 1, 3, 0], # Woman 0's preferences  
    [0, 3, 1, 2], # Woman 1's preferences  
    [1, 0, 3, 2], # Woman 2's preferences  
    [0, 1, 2, 3]  # Woman 3's preferences  
]
```

What is the output of the Gale-Shapley algorithm (as a set of stable matches)?

- ☐ a. $\{(0, 0), (1, 2), (2, 1), (3, 3)\}$
- ☐ b. $\{(0, 3), (1, 0), (2, 2), (3, 1)\}$
- ☐ c. $\{(0, 2), (1, 1), (2, 0), (3, 3)\}$
- ☐ d. $\{(0, 1), (1, 2), (2, 3), (3, 0)\}$

Q. 8 In the context of stable matchings, a valid partner for a man m is defined as:

- ☐ a. A woman w such that m and w are matched in at least one stable matching.
- ☐ b. The woman w at the top of m 's preference list.
- ☐ c. Any woman w who accepts a proposal from m during the Gale-Shapley algorithm.
- ☐ d. A woman w such that w and m are unmatched in all stable matchings.

Q. 9 Consider the following preferences of 5 men over 5 women, given as indices:

```
0 : [2, 0, 1, 3, 4]  
1 : [1, 0, 2, 3, 4]  
2 : [3, 0, 1, 4, 2]  
3 : [3, 1, 2, 0, 4]  
4 : [4, 3, 2, 1, 0]
```

The preferences shown in **bold** indicate the valid partners for each man. Using the Gale-Shapley algorithm (where men propose), the resulting stable matches are:

- ☐ a. $\{(0, 2), (1, 0), (2, 1), (3, 3), (4, 4)\}$
- ☐ b. $\{(0, 3), (1, 1), (2, 0), (3, 2), (4, 4)\}$
- ☐ c. $\{(0, 0), (1, 2), (2, 1), (3, 3), (4, 4)\}$
- ☐ d. $\{(0, 0), (1, 2), (2, 1), (3, 0), (4, 3)\}$

Q. 10 Consider the following preferences of 5 women over 5 men, given as indices:

0 : $[2, \mathbf{0}, 1, \mathbf{3}, 4]$

1 : $[1, \mathbf{0}, 2, 3, \mathbf{4}]$

2 : $[\mathbf{3}, 0, \mathbf{1}, 4, 2]$

3 : $[\mathbf{3}, 1, \mathbf{2}, 0, 4]$

4 : $[\mathbf{4}, 3, \mathbf{2}, 1, \mathbf{0}]$

The preferences shown in **bold** indicate the valid partners for each woman. Using the Gale-Shapley algorithm (where men propose), the resulting stable matches are:

- ☐ a. $\{(0, 3), (1, 4), (2, 1), (3, 2), (4, 0)\}$
- ☐ b. $\{(0, 2), (1, 0), (2, 1), (3, 3), (4, 4)\}$
- ☐ c. $\{(0, 3), (1, 1), (2, 0), (3, 2), (4, 4)\}$
- ☐ d. $\{(0, 0), (1, 2), (2, 1), (3, 0), (4, 3)\}$