

# Assignment 3: Searching

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## Assignment 3: Searching

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## 1. Find Peak in a Array

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### Description

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A peak element is an element that is greater than its neighbors. Given an input array where  $\text{num}[i] \neq \text{num}[i+1]$ , find a peak element and return its index.

The array may contain multiple peaks, in that case return the index to any one of the peaks is fine.

You may imagine that  $\text{num}[-1] = \text{num}[n] = -\infty$ .

### Examples

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#### Example 1

input: [1,8,9,5]

output:9

Explanation: The number 9 is the peak because it is greater than its neighbors. So return the index of 9.

#### Example 2

input: [9,8,5,1]

output: 9

Explanation: The number 9 is the peak because it is greater than its neighbors (its left is  $-\infty$ ).

#### Example 3

input: [1,5,8,9]

output: 9

Explanation: The number 9 is the peak because it is greater than its neighbors (its right is  $-\infty$ ).

## Solutions

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### Solution 1: Use intuition

#### Algorithm

The idea is simple: go through the array and find the max number in flight.

#### Implementation

(This is very straightforward so I would rather spend my time on other high efficiency solutions.)

#### Complexity

Time:  $O(n)$

Space:  $O(1)$

### Solution 2: Use 'binary search'

#### Algorithm

This could be improved to a complexity of  $O(\log(n))$  by using 'binary search' strategy.

- use two pointers low and high, one of which from the start and the other one from the end
- compare the middle element between the two pointers with its next
- if the next is bigger, set the low pointer to mid+1
- otherwise if it is smaller, set the high pointer to mid
- continue to do when low < high
- when exit, low equals high and either of nums[low] or nums[high] is the result

## Implementation

### Java

```
public static int findPeakElement(int[] nums) {  
    int low = 0, high = nums.length - 1;  
  
    while (low < high) {  
        int mid = low + (high - low) / 2;  
        if (nums[mid] < nums[mid + 1]) {  
            low = mid + 1;  
        } else if (nums[mid] > nums[mid + 1]) {  
            high = mid; // trick  
        }  
    }  
    return nums[low];  
}
```

### Complexity

**Time:**  $O(\log(n))$ . Each time the array is divided to two, one of which is dropped.

**Space:**  $O(1)$ . Constant space is used.

## References

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[Find Minimum in Rotated Sorted Array](#)

## 2. Find Pair Where Sum is Closest to 0

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### Description

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Given an integer array, you need to find the two elements such that their sum is closest to zero and print them in ascending order.

Note: If there are more than two sums equal close to 0, output any pair if fine.

### Examples

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*Example 1*

input: -8 -66 -60

output: -60 -8

*Example 2*

input: -21 -67 -37 -18 4 -65

output: -18 4

Example 3

input: -24 -73

output: -73 -24

## Algorithm

Sort this array in ascending order and use two pointers and move following this rule:

- if the sum of the low index and high index number is 0, return these two numbers directly
- if the sum is greater than 0, move the high pointer left if the new sum is closer to 0
- otherwise, move the low right if the new sum is closer to 0
- if the new sum is further from 0, then stop anytime in the previous two steps

## Solution: Two pointers

Java

```
public static int[] findClosestToZero(int[] nums) {  
    if (nums == null || nums.length < 2)  
        throw new IllegalArgumentException();  
    Arrays.sort(nums);  
  
    int low = 0, high = nums.length - 1;  
    while (low < high - 1) {  
        int sum = nums[low] + nums[high];  
        if (sum == 0)  
            break;  
        else if (sum < 0) {  
            if (Math.abs(nums[low + 1] + nums[high] - 0) > Math.abs(sum  
- 0))  
                break;  
            low++;  
        } else {  
            if (Math.abs(nums[low] + nums[high - 1] - 0) > Math.abs(sum  
- 0))  
                break;  
            high--;  
        }  
    }  
    return new int[] { nums[low], nums[high] };  
}
```

## Complexity

**Time:**  $O(n)$ . In the worst case (such as the middle two elements are the result), all the number in this array will be visited.

**Space:**  $O(1)$ . The sorting is in space and the used sapce is constant.

## References

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[Two numbers with sum closest to zero](#)