

MORE PROGRAMMING IN CUDA C: PERFORMANCE MEASUREMENT AND ATOMICS

Will Landau, Prof. Jarad Niemi

EVENTS: MEASURING PERFORMANCE ON THE GPU

Event: a time stamp for the GPU.

Use events to measure the amount of time the GPU spends on a task.

```

int main(){
    float    elapsedTime;
    cudaEvent_t start, stop;

    cudaEventCreate(&start);
    cudaEventCreate(&stop);
    cudaEventRecord( start, 0 );

    // SOME GPU KERNEL YOU WANT TIMED HERE

    cudaEventRecord( stop, 0 );
    cudaEventSynchronize( stop );
    cudaEventElapsedTime( &elapsedTime, start, stop );

    cudaEventDestroy( start );
    cudaEventDestroy( stop );

    printf("GPU Time elapsed: %f\n", elapsedTime);
}

```

The variable, `elapsedTime`, is the GPU time spent on the task. You can now print it any way you like.

Note: only GPU elapsed time is measured, not CPU time.

GPU time and CPU time must be measured separately.

MEASURING CPU TIME

```
#include <stdio.h>
#include <time.h>

int main(){

    clock_t start = clock();

    // SOME CPU CODE YOU WANT TIMED HERE

    float elapsedTime = ((double)clock() - start) /
                        CLOCKS_PER_SEC;

    printf("CPU Time elapsed: %f\n", elapsedTime);
}
```

ATOMICS

Consider the following GPU operation on integer **x**:

x++;

which tells the GPU to do 3 things...

`x++;`

1. Read the value stored in x.
2. Add 1 to the value read in step 1.
3. Write the result back to x.

Say we need threads A and B to increment x. We want:

| STEP | EXAMPLE |
|--|-------------------|
| 1. Thread A reads the value in x. | A reads 7 from x. |
| 2. Thread A adds 1 to the value it read. | A computes 8. |
| 3. Thread A writes the result back to x. | x <- 8. |
| 4. Thread B reads the value in x. | B reads 8 from x. |
| 5. Thread B adds 1 to the value it read. | B computes 9. |
| 6. Thread B writes the result back to x. | x <- 9. |

but we might get:

| STEP | EXAMPLE |
|---------------------------------------|-------------------------|
| Thread A reads the value in x. | A reads 7 from x. |
| Thread B reads the value in x. | B reads 7 from x. |
| Thread A adds 1 to the value it read. | A computes 8. |
| Thread B adds 1 to the value it read. | B computes 8. |
| Thread A writes the result back to x. | <code>x <- 8.</code> |
| Thread B writes the result back to x. | <code>x <- 8.</code> |

Race Condition: A computational hazard that arises when the results of a program depend on the timing of uncontrollable events, such as threads.

Atomic Operation: A command that is executed one thread at a time, thus avoiding a race condition.

To avoid the race condition in our example, we atomically add 1 to x:

```
atomicAdd( &x, 1 );
```

instead of using `x++`;

LIST OF CUDA C BUILT-IN ATOMICS

atomicAdd()

atomicSub()

atomicMin()

atomicMax()

atomicInc()

atomicDec()

atomicExch()

atomicCAS()

```
int atomicAdd(int* address, int val);
unsigned int atomicAdd(unsigned int* address,
                       unsigned int val);
unsigned long long int atomicAdd(unsigned long long int* address,
                                unsigned long long int val);
float atomicAdd(float* address, float val);
```

reads the 32-bit or 64-bit word **old** located at the address **address** in global or shared memory, computes **(old + val)**, and stores the result back to memory at the same address. These three operations are performed in one atomic transaction. The function returns **old**.

```
int atomicSub(int* address, int val);
```

```
unsigned int atomicSub(unsigned int* address,  
                        unsigned int val);
```

reads the 32-bit word **old** located at the address **address** in global or shared memory, computes **(old - val)**, and stores the result back to memory at the same address. These three operations are performed in one atomic transaction. The function returns **old**.

```
int atomicExch(int* address, int val);
unsigned int atomicExch(unsigned int* address,
                        unsigned int val);
unsigned long long int atomicExch(unsigned long long int* address,
                                unsigned long long int val);
float atomicExch(float* address, float val);
```

reads the 32-bit or 64-bit word **old** located at the address **address** in global or shared memory and stores **val** back to memory at the same address. These two operations are performed in one atomic transaction. The function returns **old**.

```
--  
int atomicMin(int* address, int val);  
unsigned int atomicMin(unsigned int* address,  
                        unsigned int val);
```

reads the 32-bit word **old** located at the address **address** in global or shared memory, computes the minimum of **old** and **val**, and stores the result back to memory at the same address. These three operations are performed in one atomic transaction. The function returns **old**.

```
int atomicMax(int* address, int val);  
unsigned int atomicMax(unsigned int* address,  
                        unsigned int val);
```

reads the 32-bit word **old** located at the address **address** in global or shared memory, computes the maximum of **old** and **val**, and stores the result back to memory at the same address. These three operations are performed in one atomic transaction. The function returns **old**.


```
--  
unsigned int atomicInc(unsigned int* address,  
                        unsigned int val);
```

reads the 32-bit word **old** located at the address **address** in global or shared memory, computes $((\text{old} \geq \text{val}) ? 0 : (\text{old} + 1))$, and stores the result back to memory at the same address. These three operations are performed in one atomic transaction. The function returns **old**.

```
unsigned int atomicDec(unsigned int* address,  
                        unsigned int val);
```

reads the 32-bit word **old** located at the address **address** in global or shared memory, computes $((\text{old} == 0) \mid (\text{old} > \text{val})) ? \text{val} : (\text{old}-1)$, and stores the result back to memory at the same address. These three operations are performed in one atomic transaction. The function returns **old**.

```
int atomicCAS(int* address, int compare, int val);
unsigned int atomicCAS(unsigned int* address,
                      unsigned int compare,
                      unsigned int val);
unsigned long long int atomicCAS(unsigned long long int* address,
                                unsigned long long int compare,
                                unsigned long long int val);
```

reads the 32-bit or 64-bit word **old** located at the address **address** in global or shared memory, computes (**old == compare ? val : old**), and stores the result back to memory at the same address. These three operations are performed in one atomic transaction. The function returns **old** (Compare And Swap).

```
int atomicAnd(int* address, int val);  
unsigned int atomicAnd(unsigned int* address,  
                        unsigned int val);
```

reads the 32-bit word **old** located at the address **address** in global or shared memory, computes **(old & val)**, and stores the result back to memory at the same address. These three operations are performed in one atomic transaction. The function returns **old**.

```
int atomicOr(int* address, int val);  
unsigned int atomicOr(unsigned int* address,  
                      unsigned int val);
```

reads the 32-bit word **old** located at the address **address** in global or shared memory, computes **(old | val)**, and stores the result back to memory at the same address. These three operations are performed in one atomic transaction. The function returns **old**.

```
int atomicXor(int* address, int val);  
unsigned int atomicXor(unsigned int* address,  
                        unsigned int val);
```

reads the 32-bit word **old** located at the address **address** in global or shared memory, computes **(old ^ val)**, and stores the result back to memory at the same address. These three operations are performed in one atomic transaction. The function returns **old**.

LOCKS

Lock: a mechanism in parallel computing that forces a block of code to be executed atomically.

mutex: short for “mutual exclusion”, the idea behind locks: while a thread is running code inside a lock, it blocks all other threads from running the code.

lock.h:

```
struct Lock {
    int *mutex;
    Lock( void ) {
        int state = 0;
        HANDLE_ERROR( cudaMalloc( (void**)& mutex,
                                   sizeof(int) ) );
        HANDLE_ERROR( cudaMemcpy( mutex, &state, sizeof(int),
                                   cudaMemcpyHostToDevice ) );
    }

    ~Lock( void ) {
        cudaFree( mutex );
    }

    __device__ void lock( void ) {
        while( atomicCAS( mutex, 0, 1 ) != 0 );
    }

    __device__ void unlock( void ) {
        atomicExch( mutex, 0 );
    }
};
```


NEW DOT PRODUCT CODE

```
__global__ void dot( Lock lock, float *a,
                    float *b, float *c ) {
    __shared__ float cache[threadsPerBlock];
    int tid = threadIdx.x + blockIdx.x * blockDim.x;
    int cacheIndex = threadIdx.x;

    float temp = 0;
    while (tid < N) {
        temp += a[tid] * b[tid];
        tid += blockDim.x * gridDim.x;
    }

    // set the cache values
    cache[cacheIndex] = temp;

    // synchronize threads in this block
    __syncthreads();
}
```

```

// for reductions, threadsPerBlock must be a power of 2
// because of the following code
int i = blockDim.x/2;
while (i != 0) {
    if (cacheIndex < i)
        cache[cacheIndex] += cache[cacheIndex + i];
    __syncthreads();
    i /= 2;
}

if (cacheIndex == 0) {
    lock.lock();
    *c += cache[0];
    lock.unlock();
}
}

```

REFERENCES

NVIDIA CUDA C Programming Guide. Version 3.2.
2010.