

Computer Graphics -Subdivision & Simplification

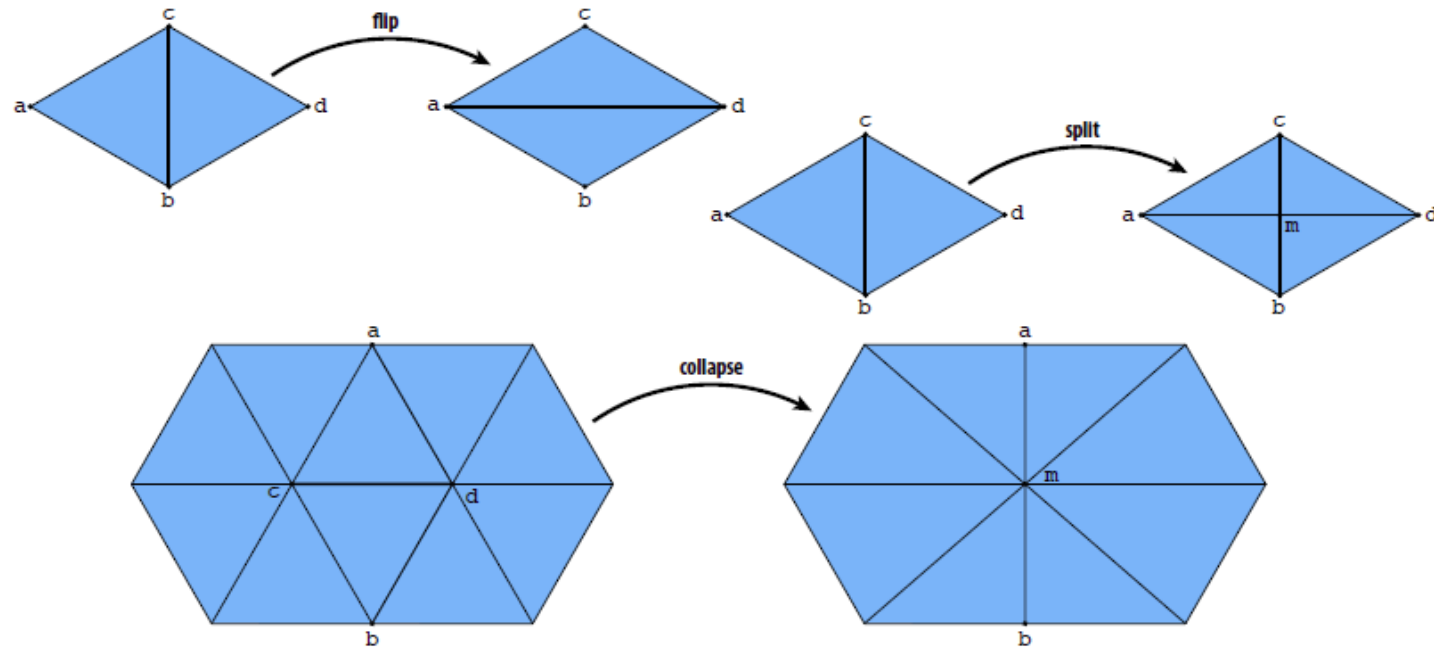
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Spring 2017

<http://jjcao.github.io/ComputerGraphics/>

Processing geometry with halfedges

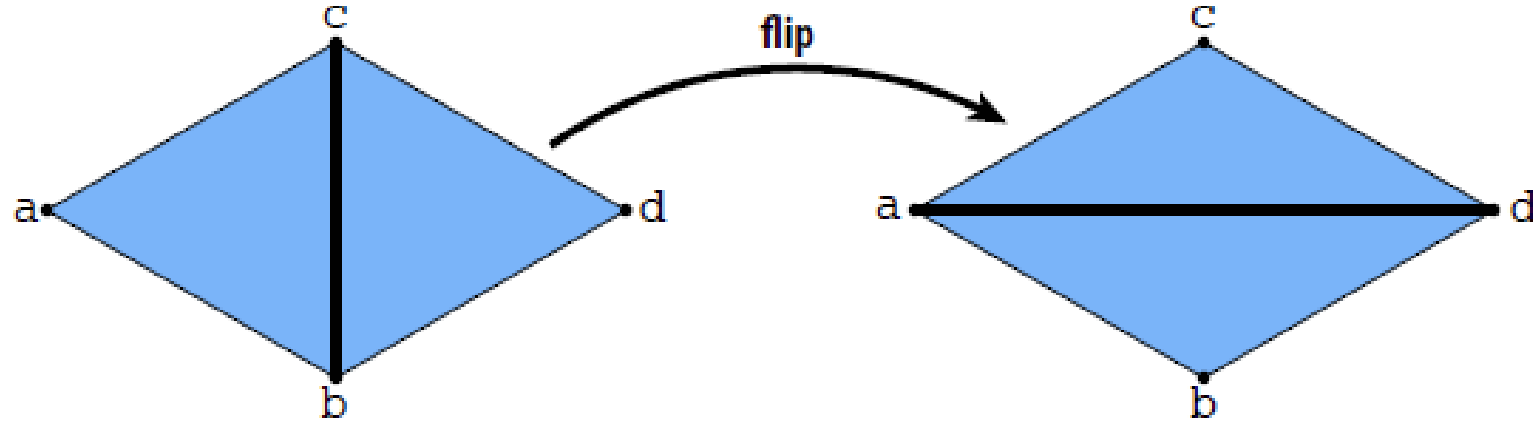
- Remember key feature of linked list: insert/delete elements
- Same story with halfedge mesh (“linked list on steroids”)
- Several atomic operations for triangle meshes:



- How? Allocate/delete elements; reassigning pointers.
- (Should be careful to preserve manifoldness!)

Edge Flip

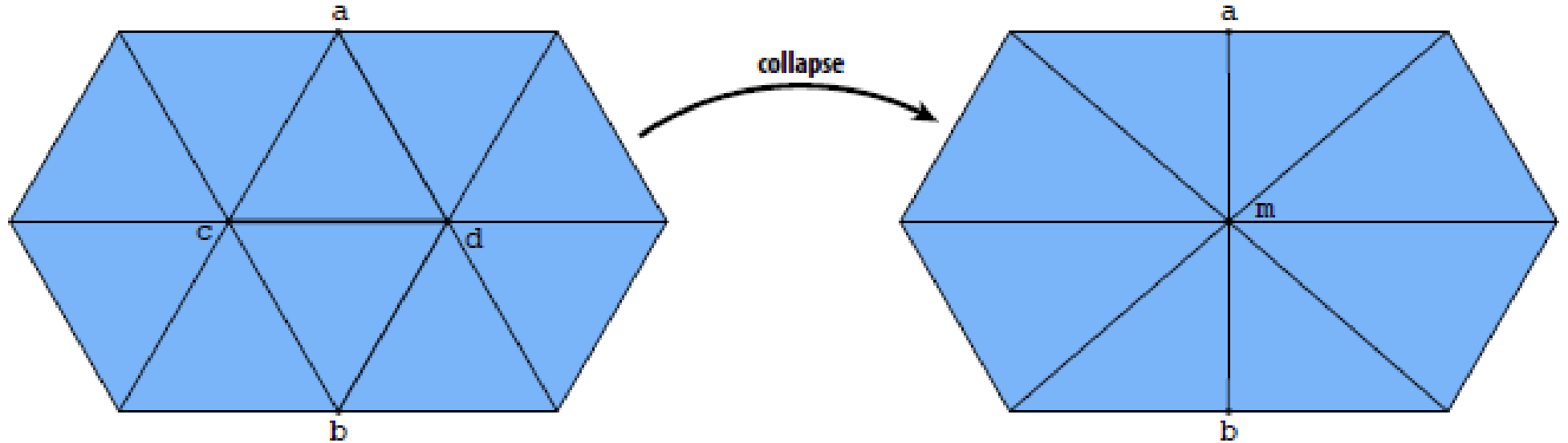
- Triangles (a,b,c) , (b,d,c) become (a,d,c) , (a,b,d) :



- Long list of pointer reassignments (edge \rightarrow halfedge = ...)
- However, no elements created/destroyed.
- Q: What happens if we flip twice?
- (Challenge: can you implement edge flip such that pointers are unchanged after two flips?)

Edge Collapse

- Replace edge (c,d) with a single vertex m:

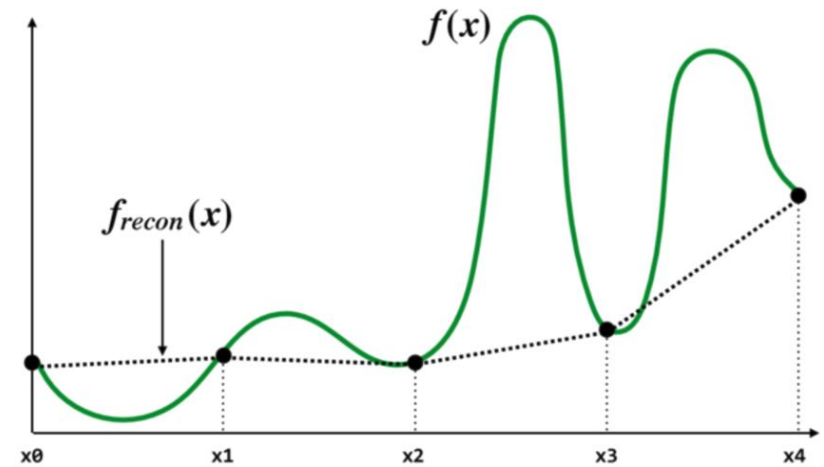


- Now have to *delete* elements.
- Still lots of pointer assignments!
- Q: How would we implement this with a polygon soup?
- Any other good way to do it? (E.g., different data structure?)

**Ok, but what can we actually *do* with
these operations?**

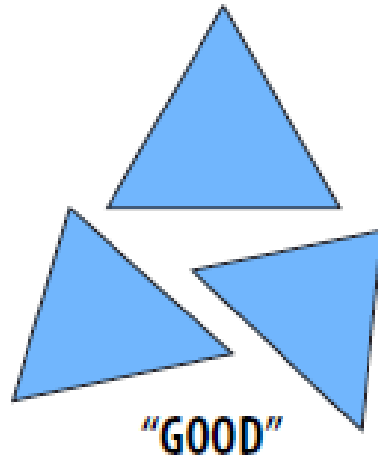
Remeshing as resampling

- Remember our discussion of *aliasing*
- Bad sampling makes signal appear different than it really is
- E.g., undersampled curve looks flat
- Geometry is no different!
 - undersampling destroys features
 - oversampling destroys performance

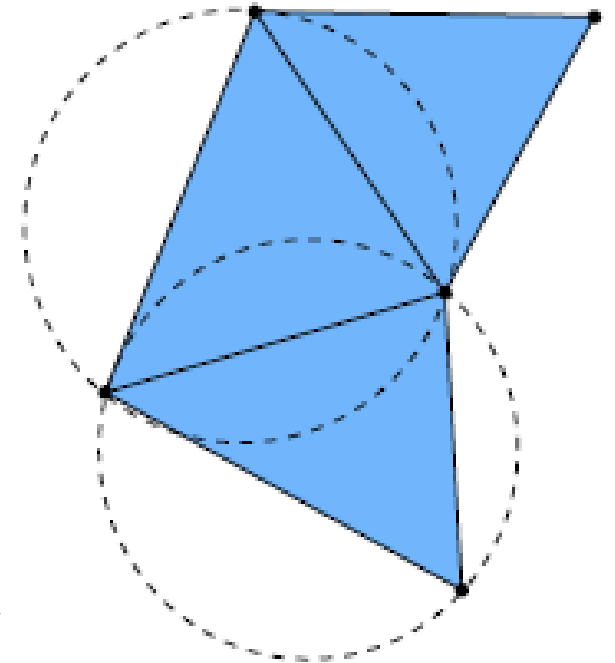


What makes a “good” triangle mesh?

- One rule of thumb: *triangle shape*

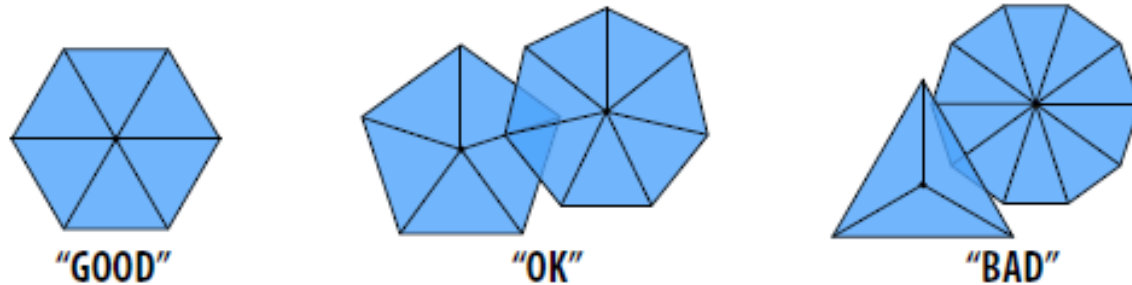


- More specific condition: *Delaunay*
- “Circumcircle interiors contain no vertices.”
- Not *always* a good condition, but often*
- especially important for simulation
- for approximation, long triangles may be better

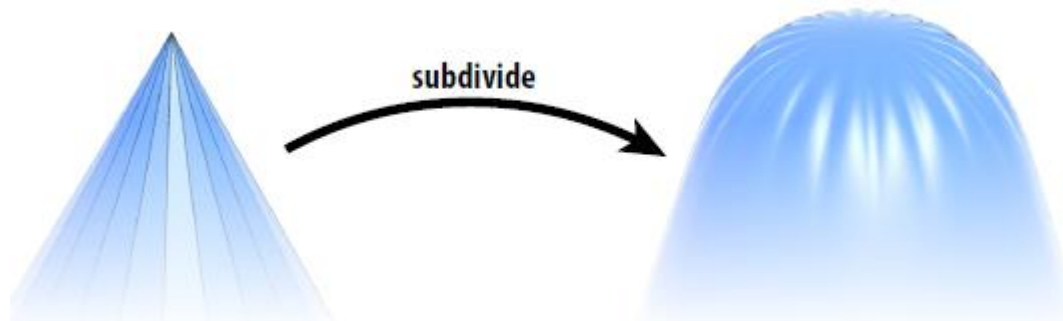


What else constitutes a good mesh?

- Another rule of thumb: *regular vertex degree*
- Ideal for triangle meshes: make every vertex valence 6:



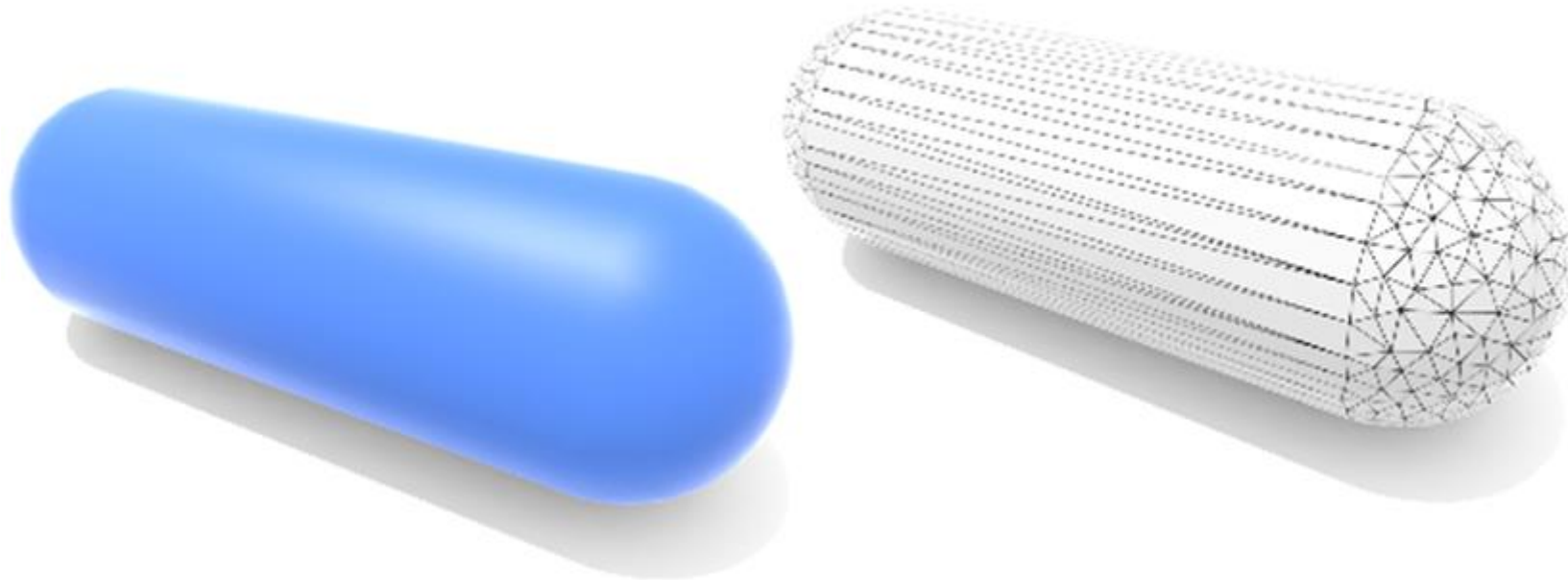
- Why? Better triangle shape, important for (e.g.) subdivision:



- **FACT:** Can't have perfect valence everywhere! (except on torus)

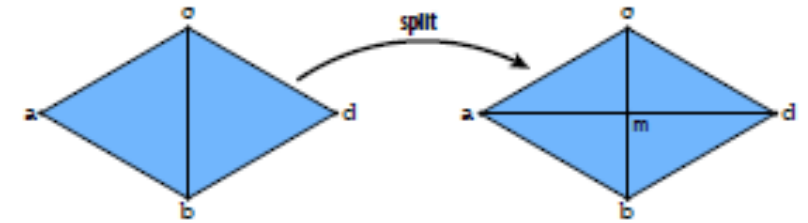
What *else* makes a “good” mesh?

- Keep only elements that contribute *information* about shape
- Add additional information where, e.g., *curvature* is large
- Balance with element quality
 - Delaunay, “round” triangles, regular degree, etc.

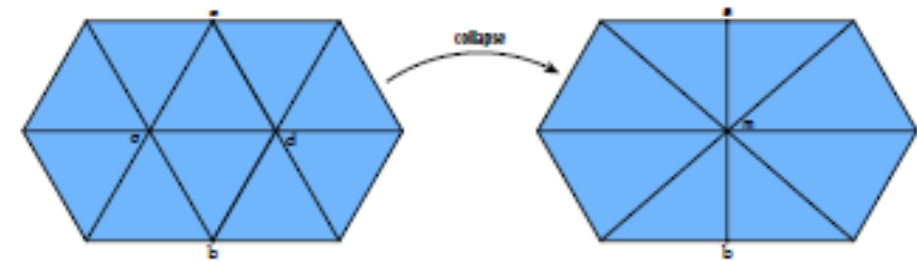


How do we resample? Already know how!

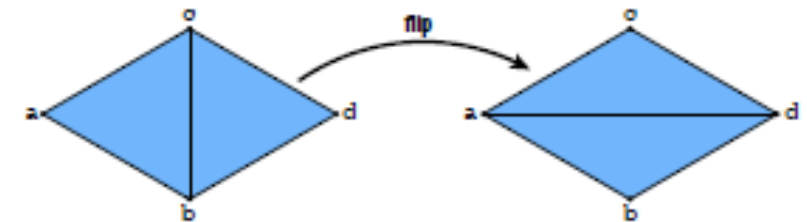
- Edge split is (local) upsampling:



- Edge collapse is (local) downsampling:



- Edge flip is (local) resampling:

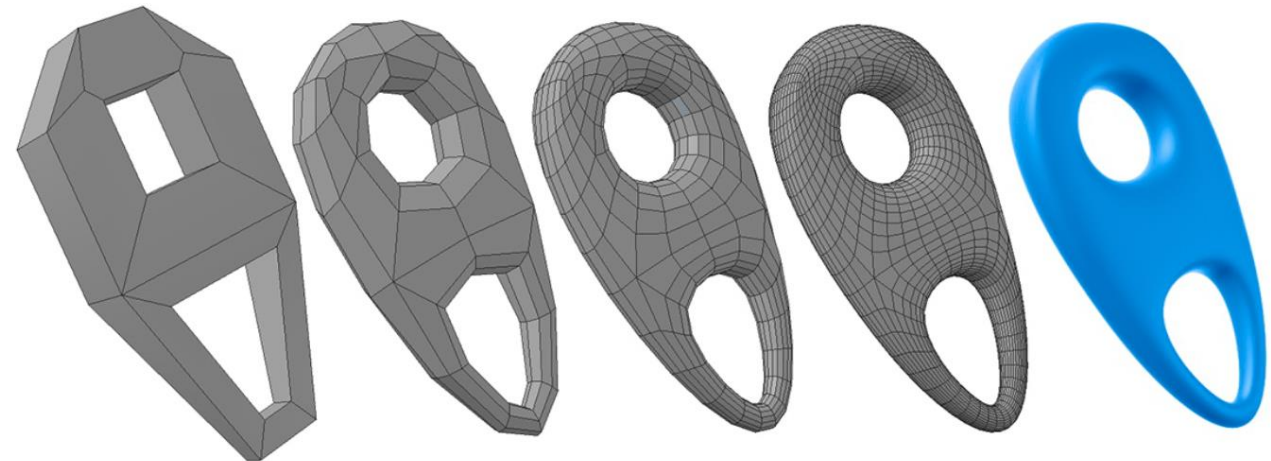
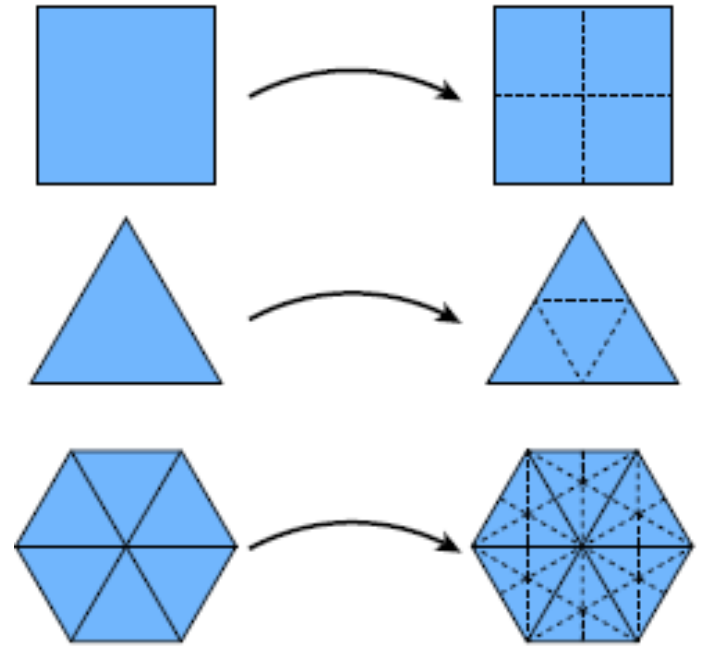


- Still need to intelligently decide *which* edges to modify!

**Which edges should we split to
upsample the whole mesh?**

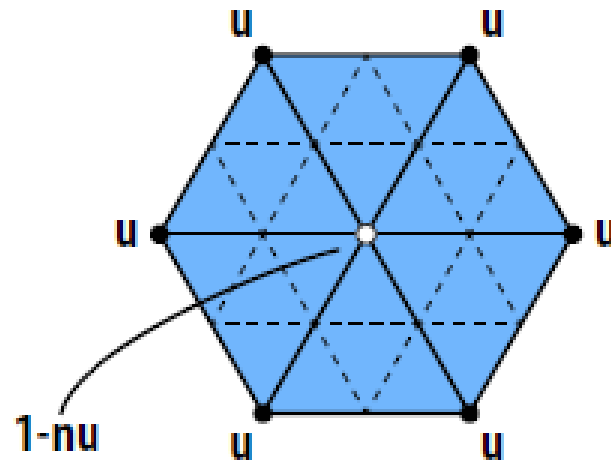
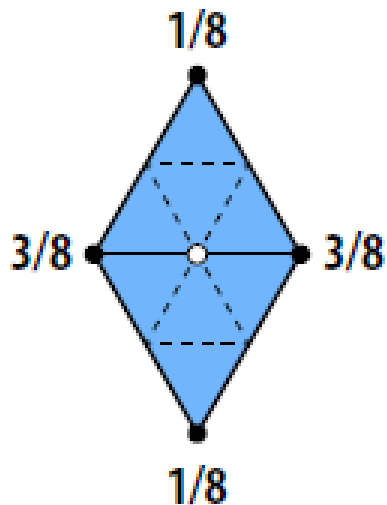
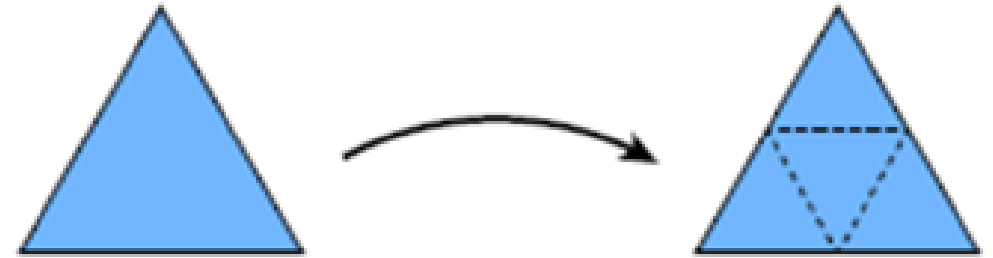
Subdivision as Upsampling

- Repeatedly split each element into smaller pieces
- Replace vertex positions with weighted average of neighbors
- Main considerations:
 - interpolating vs. approximating
 - limit surface continuity ($C1$, $C2$, ...)
 - behavior at irregular vertices
- Many options:
 - Quad: Catmull-Clark
 - Triangle: Loop, Butterfly, Sqrt(3)



Loop Subdivision

- Fairly common subdivision rule for triangles
- Curvature is continuous away from irregular vertices (“C2”)
- *Approximating*, not interpolating
- Algorithm:
 - Split each triangle into four
 - Assign new vertex positions according to weights:

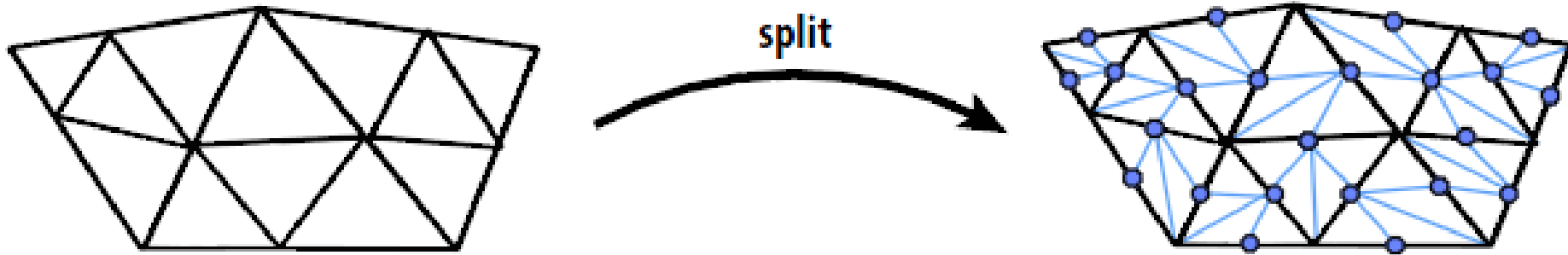


n : vertex degree

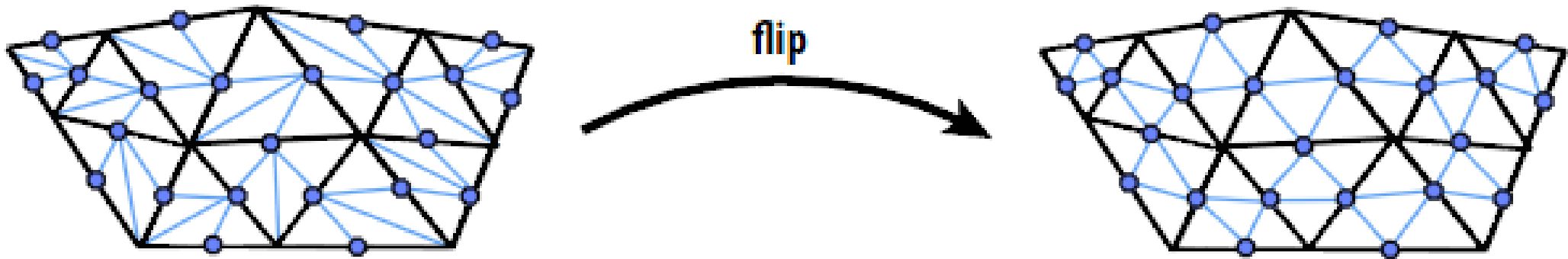
u : $3/16$ if $n=3$, $3/(8n)$ otherwise

Loop Subdivision via Edge Operations

- First, split edges of original mesh in *any* order:



- Next, flip new edges that touch a new & old vertex:



What if we want fewer triangles?

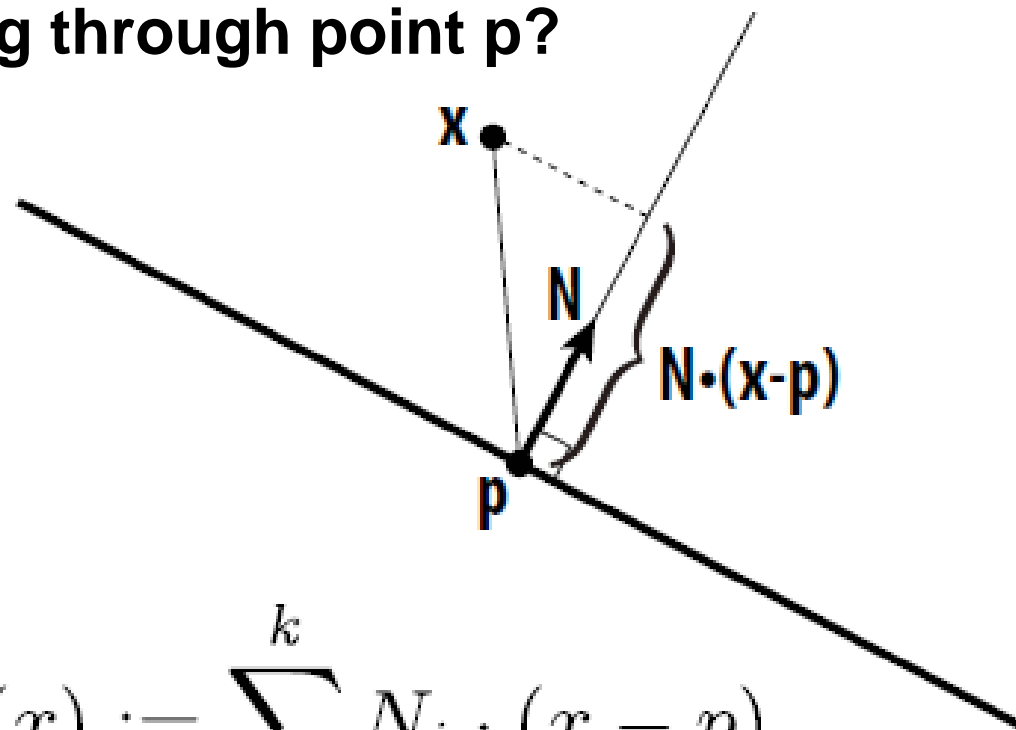
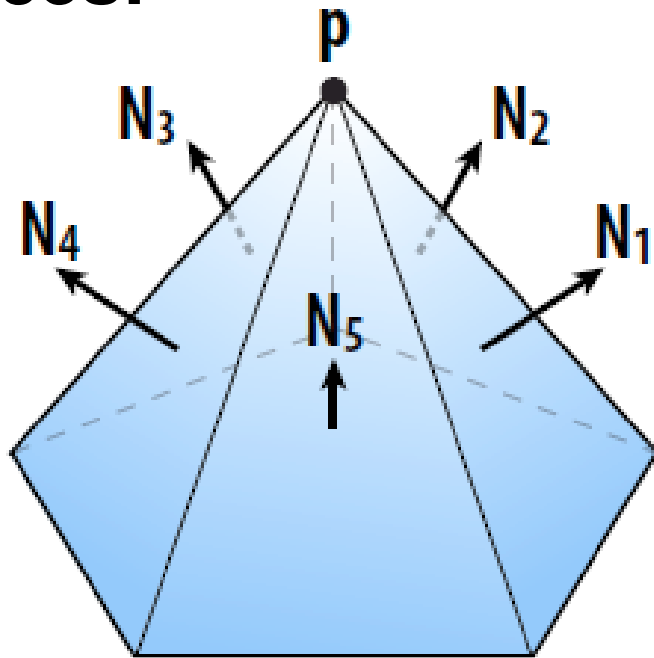
Simplification via Quadric Error Metric

- One popular scheme: iteratively collapse edges
- Which edges? Assign score with *quadric error metric**
 - approximate distance to surface as sum of distance to aggregated triangles
 - iteratively collapse edge with smallest score
 - greedy algorithm... great results!



Quadric Error Metric

- Approximate distance to a collection of triangles
- Distance is sum of point-to-plane distances
 - Q: Distance to plane w/ normal N passing through point p ?
 - A: $d(x) = N \cdot (x - p)$
- Sum of distances:



$$d(x) := \sum_{i=1}^k N_i \cdot (x - p)$$

Quadric Error - Homogeneous Coordinates

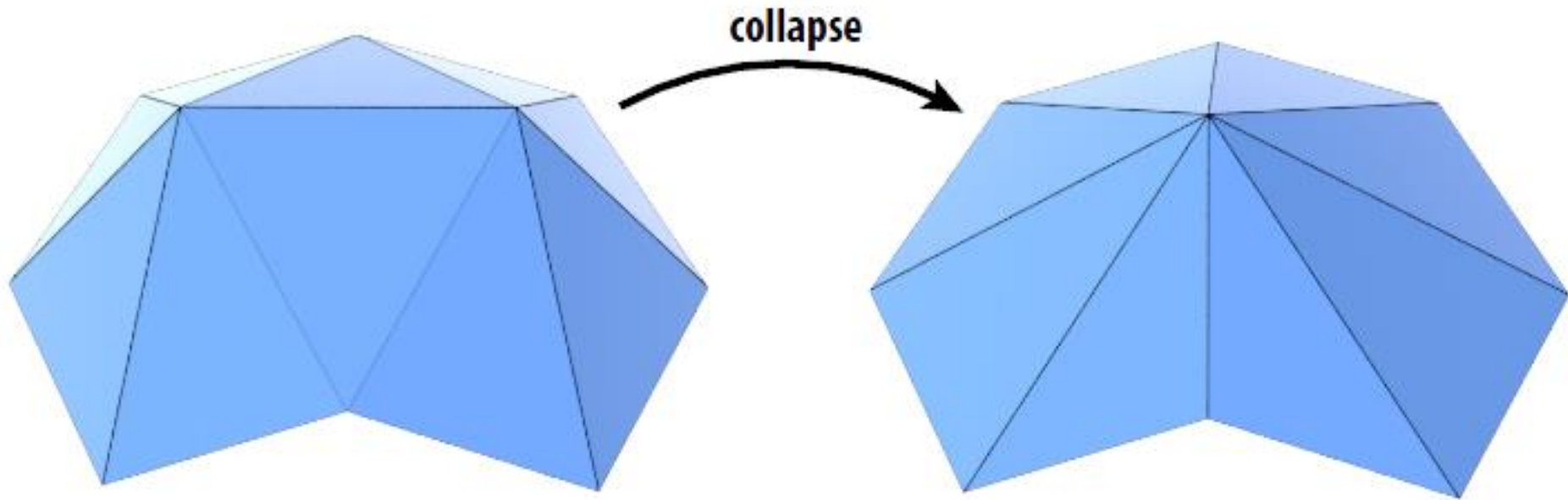
- Suppose in coordinates we have
 - a query point (x,y,z)
 - a normal (a,b,c)
 - an offset $d := -(x,y,z) \cdot (a,b,c)$
- Then in homogeneous coordinates, let
 - $u := (x,y,z,1)$
 - $v := (a,b,c,d)$

$$Q = \begin{bmatrix} a^2 & ab & ac & ad \\ ab & b^2 & bc & bd \\ ac & bc & c^2 & cd \\ ad & bd & cd & d^2 \end{bmatrix}$$

- *Signed* distance to plane is then just $u \cdot v = ax+by+cz+d$
- *Squared* distance is $(u'v)^2 = u'(vv')u =: u'Qu$
- Key idea: *matrix Q encodes distance to plane*
- Q is symmetric, contains 10 unique coefficients (small storage)

Quadric Error of Edge Collapse

- How much does it cost to collapse an edge?
- Idea: compute edge midpoint, measure quadric error



- Better idea: use point that minimizes quadric error as new point!
- (More details in assignment; see also Garland & Heckbert 1997.)

Quadric Error Simplification

- Compute Q for each triangle
- Set Q at each vertex to sum of Q s from incident triangles
- Until we reach target # of triangles:
 - collapse edge (i,j) with smallest cost to get new vertex k
 - add Q_i and Q_j to get new quadric Q_k
 - update cost of any edge touching new vertex k
- Store edges in *priority queue* to keep track of minimum cost



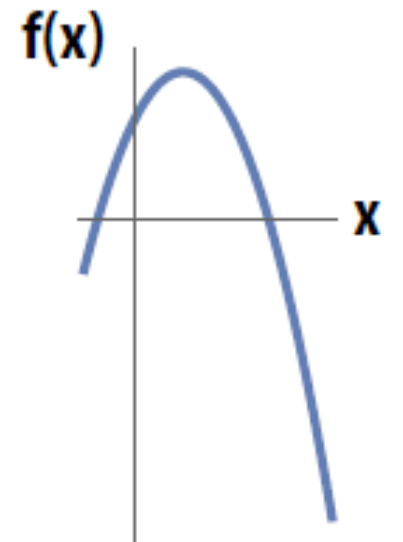
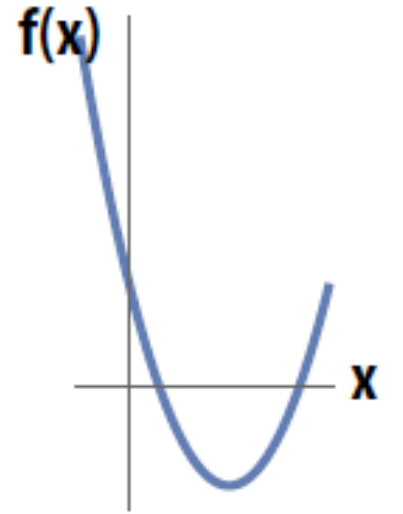
Review: Minimizing a Quadratic Function

- Suppose I give you a function $f(x) = ax^2 + bx + c$
- Q: What does the graph of this function look like?
- Could also look like this!
- Q: How do we find the *minimum*?
- A: Look for the point where the function isn't changing (if we look “up close”)
- I.e., find the point where the *derivative* vanishes

$$f'(x) = 0$$

$$2ax + b = 0$$

$$x = -b/2a$$



Minimizing a Quadratic Form

- A *quadratic form* is just a generalization of our quadratic
- polynomial from 1D to nD
- E.g., in 2D: $f(x,y) = ax^2 + bxy + cy^2 + dx + ey + g$
- Can always (always!) write quadratic polynomial using a
- *symmetric* matrix (and a vector, and a constant):

$$f(x, y) = \begin{bmatrix} x & y \end{bmatrix} \begin{bmatrix} a & b/2 \\ b/2 & c \end{bmatrix} \begin{bmatrix} x \\ y \end{bmatrix} + \begin{bmatrix} d & e \end{bmatrix} \begin{bmatrix} x \\ y \end{bmatrix} + g$$

$$= \mathbf{x}^T A \mathbf{x} + \mathbf{u}^T \mathbf{x} + g \quad \text{(this expression works for any } n\text{)}$$

- Q: How do we find a critical point (min/max/saddle)?
- A: Set derivative to zero!

$$2A\mathbf{x} + \mathbf{u} = 0$$

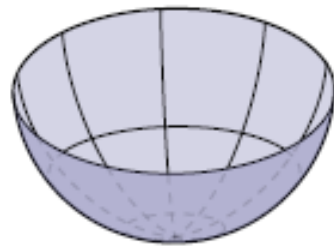
$$\mathbf{x} = -\frac{1}{2}A^{-1}\mathbf{u}$$

Positive Definite Quadratic Form

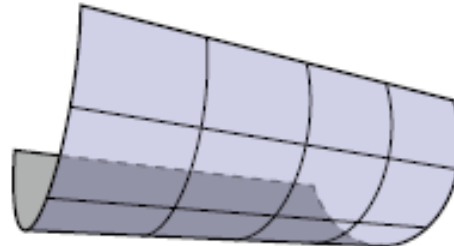
- Just like our 1D parabola, critical point is *not* always a min!
- Q: In 2D, 3D, nD, when do we get a *minimum*?
- A: When matrix A is *positive-definite*:

$$\mathbf{x}^T A \mathbf{x} > 0 \quad \forall \mathbf{x}$$

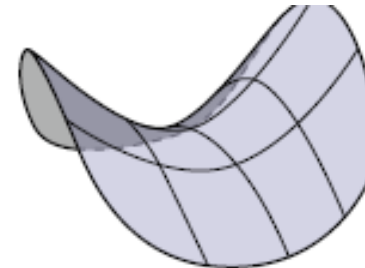
- 1D: Must have $xax = ax^2 > 0$. In other words: a is positive!
- 2D: Graph of function looks like a “bowl”:



positive definite



positive semidefinite



indefinite

- Positive-definiteness is *extremely important* in CG: it means we can find a minimum by solving linear equations. Basis of many, many modern algorithms (geometry processing, simulation, ...).

Minimizing Quadratic Error

- Find “best” point for edge collapse by minimizing quad. form

$$\min_u \mathbf{u}^\top K \mathbf{u}$$

- Already know fourth (homogeneous) coordinate is 1!
- So, break up our quadratic function into two pieces:

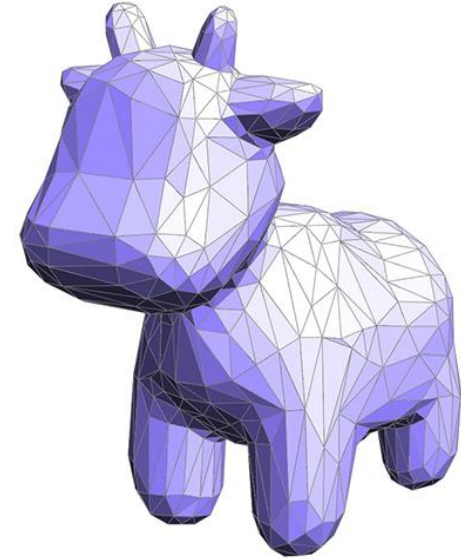
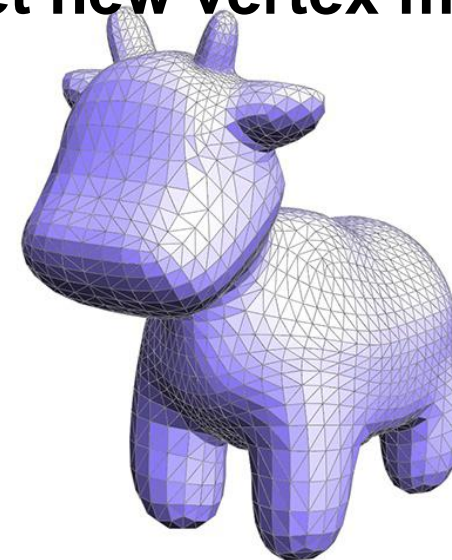
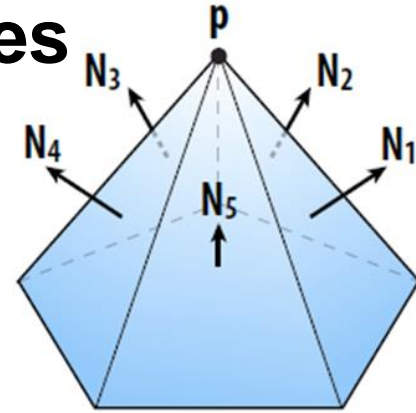
$$\begin{bmatrix} \mathbf{x}^\top & 1 \end{bmatrix} \begin{bmatrix} B & \mathbf{w} \\ \mathbf{w} & d^2 \end{bmatrix} \begin{bmatrix} \mathbf{x} \\ 1 \end{bmatrix} \\ = \mathbf{x}^\top B \mathbf{x} + 2\mathbf{w}^\top \mathbf{x} + d^2$$

- Now we have a quadratic form in the 3D position \mathbf{x} .
- Can minimize as before:

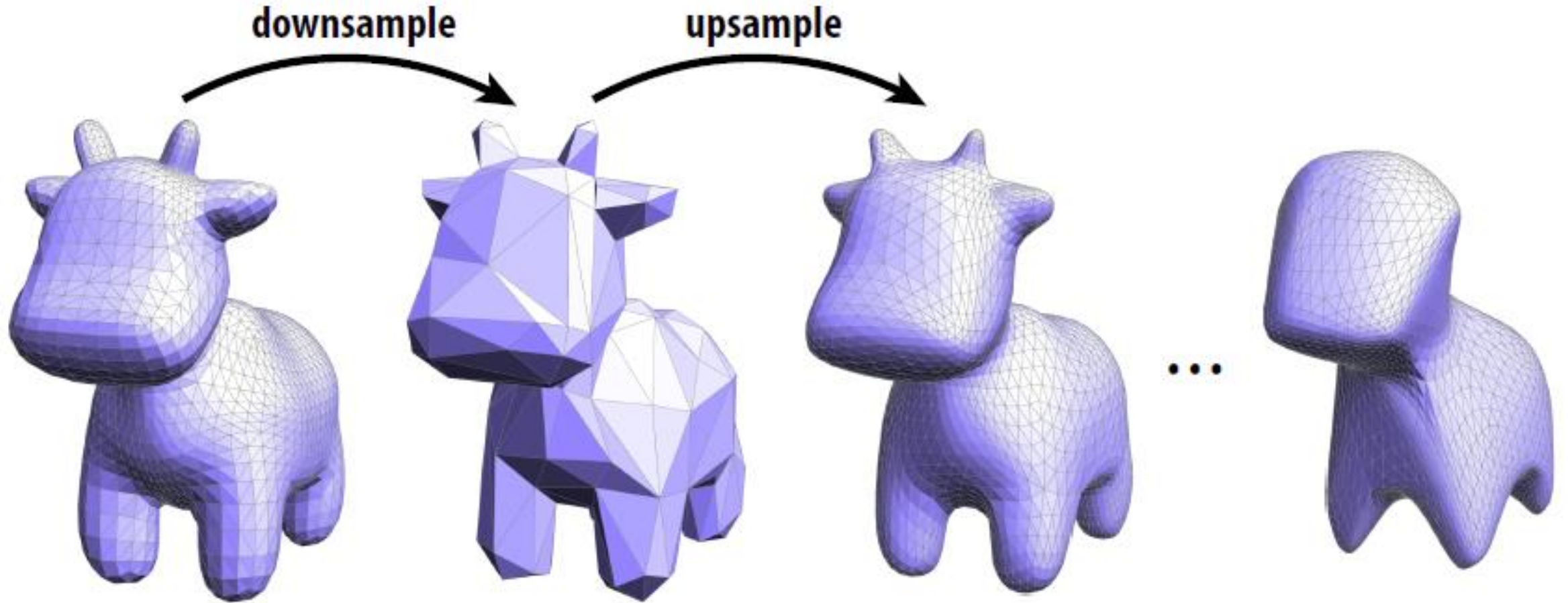
$$2B\mathbf{x} + 2\mathbf{w} = 0 \quad \Longleftrightarrow \quad \mathbf{x} = -B^{-1}\mathbf{w}$$

Quadric Error Simplification: Final Algorithm

- Compute K for each triangle (distance to plane)
- Set K at each vertex to sum of K s from incident triangles
- Set K at each edge to sum of K s at endpoints
- Find point at each edge minimizing quadric error
- Until we reach target # of triangles:
 - collapse edge (i,j) with smallest cost to get new vertex m
 - add K_i and K_j to get quadric K_m at m
 - update cost of edges touching m
- *More details in assignment writeup!*



Demo: Danger of Resampling

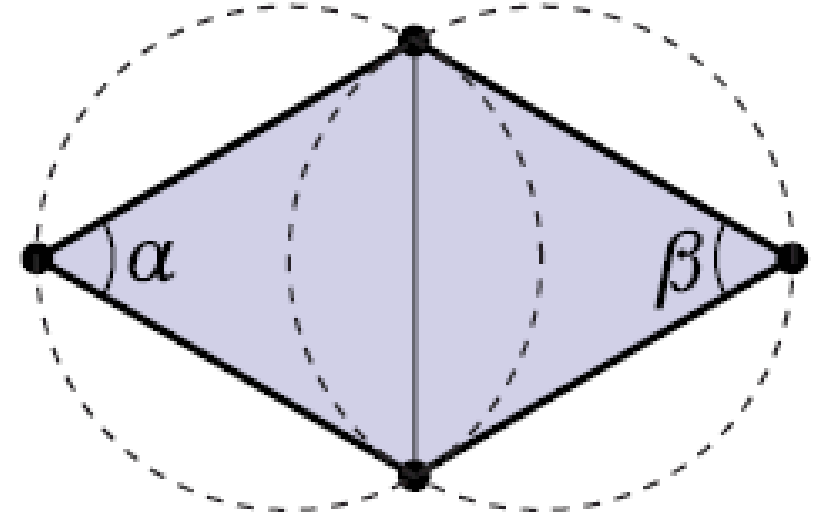
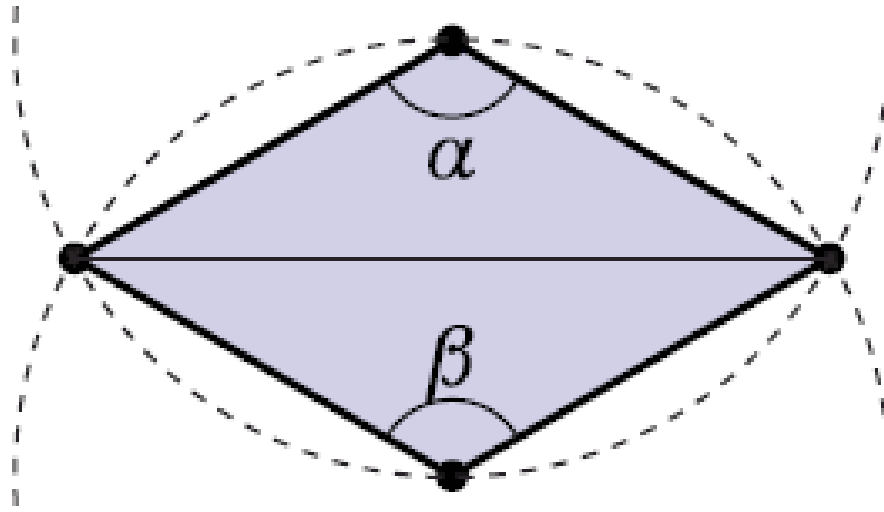


(Q: What happens with an image?)

What if we're happy with the number of triangles, but want to improve quality?

How do we make a mesh “more Delaunay”?

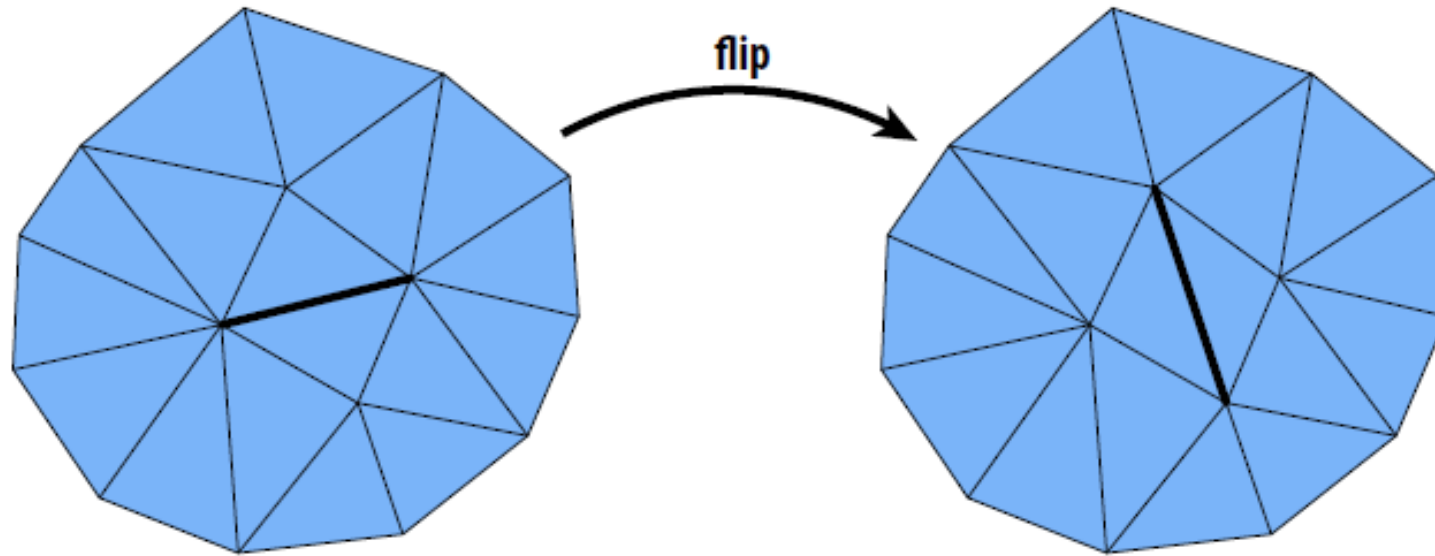
- Already have a good tool: edge flips!
- If $\alpha + \beta > \pi$, flip it!



- **FACT:** in 2D, flipping edges eventually yields Delaunay mesh
- **Theory:** worst case $O(n^2)$; no longer true for surfaces in 3D.
- **Practice:** simple, effective way to improve mesh quality

Alternatively: how do we improve degree?

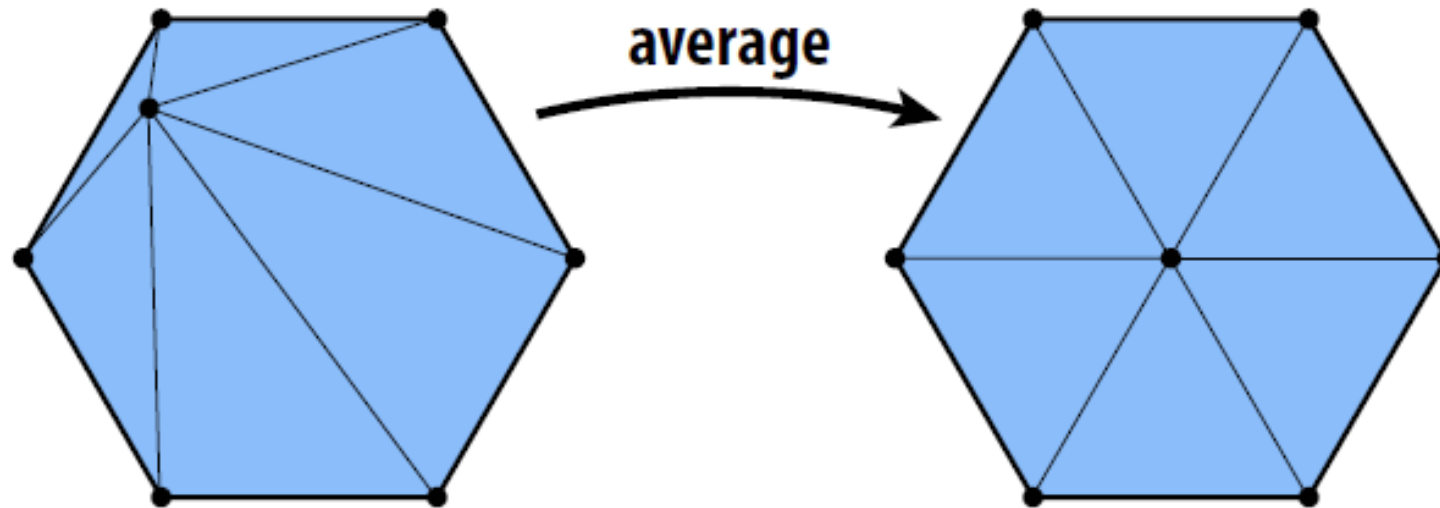
- Same tool: edge flips!
- If total deviation from degree-6 gets smaller, flip it!



- **FACT:** average valence of any triangle mesh is 6
- Iterative edge flipping acts like “discrete diffusion” of degree
- Again, no (known) guarantees; works well in practice

How do we make a triangles “more round”?

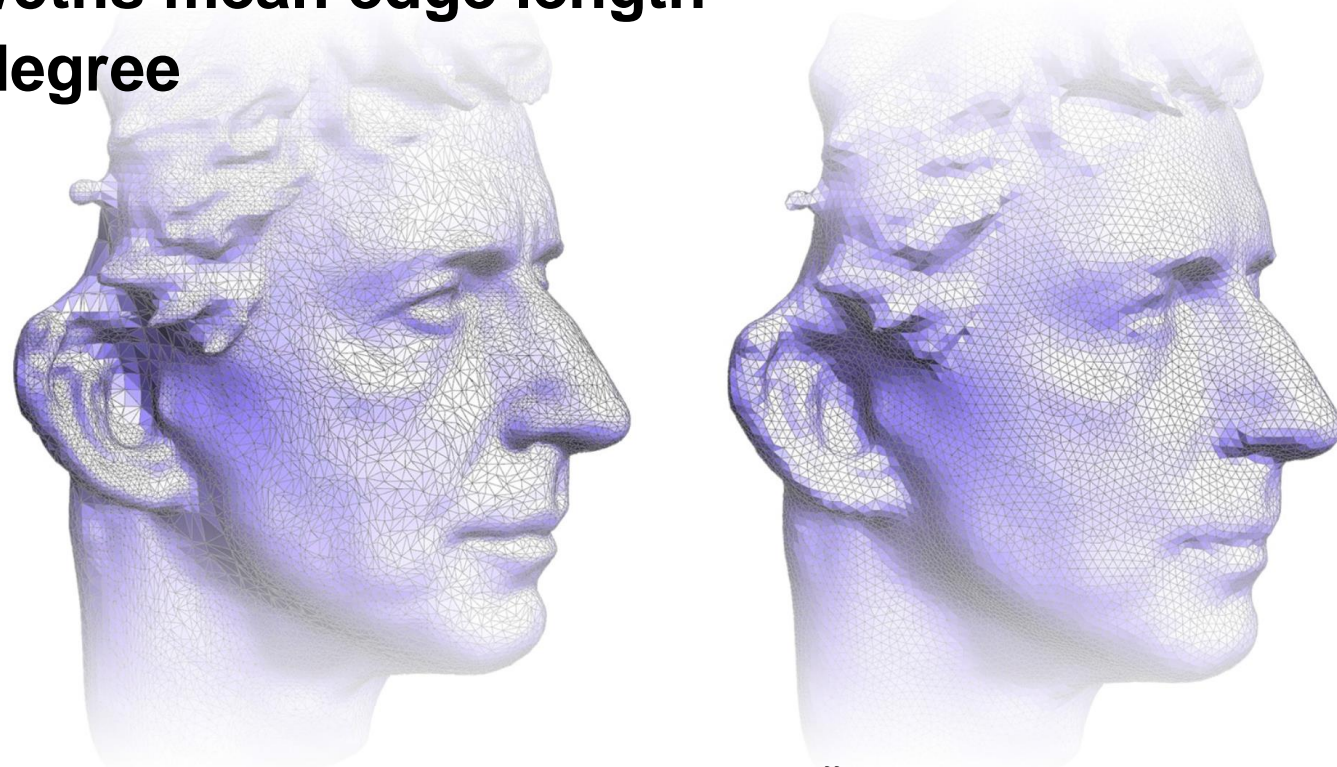
- Delaunay doesn't mean triangles are “round” (angles near 60°)
- Can often improve shape by centering vertices:



- Simple version of technique called “Laplacian smoothing”.*
- On surface: move only in *tangent* direction
- How? Remove normal component from update vector.

Isotropic Remeshing Algorithm*

- Try to make triangles uniform shape & size
- Repeat four steps:
 - Split any edge over $\frac{4}{3}$ ths mean edge length
 - Collapse any edge less than $\frac{4}{5}$ ths mean edge length
 - Flip edges to improve vertex degree
 - Center vertices tangentially



*Based on Botsch & Kobbelt, “A Remeshing Approach to Multiresolution Modeling”

Summary

- **Extend signal processing to curved shapes**
 - encounter familiar issues (sampling, aliasing, etc.)
 - some new challenges (irregular sampling, no FFT, etc.)
- **Focused on resampling triangle meshes**
 - local: edge flip, split, collapse
 - global: subdivision, quadric error, isotropic remeshing



Thank you