



Compiler Optimization and Runtime SystEms



# Understanding, Scripting and Extending GDB

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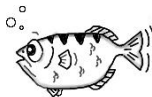
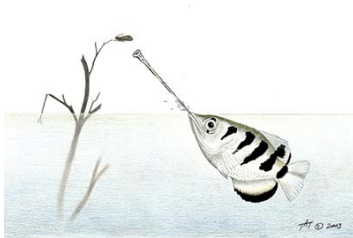


**What is a debugger?**



## What is a debugger?

It's not a tool to remove bugs!



(not even to shoot them like the Archerfish of GDB's logo ;-)



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Tools like GDB have the ability to ...

- access **the program state**
  - ▶ read and write memory cells and CPU registers ...
  - ▶ in the language's type system
- control the **execution execution**
  - ▶ execute internal code on specific events
  - ▶ execute code in the process' address-space



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Like



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Like   **QEMU**  ?

Nope:

- the execution is **100% native**
- everything done through **collaboration** between ...
  - ▶ the **OS**, the **compiler**, the **CPU** ... and old hackers' tricks!



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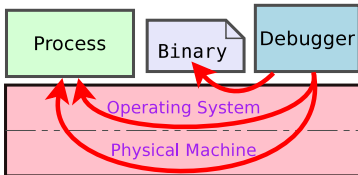
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## Help from the compiler

- DWARF debug info: type system and calling conventions

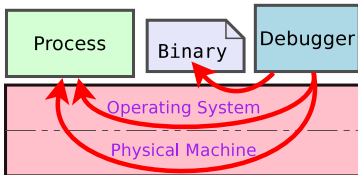
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- not much (mainly watchpoint and instruction-level step-by-step)

## Help from the OS

- ... the rest (access to the memory/registers + scheduler)





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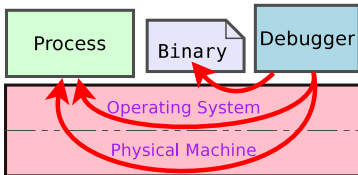
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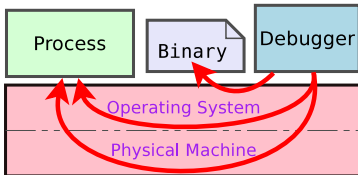
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- **Stopping** the execution ...

- breakpoint on an address execution

- watchpoint on an address access (read or write)

- catchpoints on particular events (signals, syscalls, fork/exec, ...)

- Controlling the execution:

- next/i go to next line/instruction

- step/i step into the current line's function call (if any)

- finish continue until the end of the current function

- return abort the current function call

- Stopping the execution ...

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- 1 GDB Under the Hood
- 2 Programming GDB in Python
- 3 New GDB Fonctionnalités



## 1 GDB Under the Hood

- Help from the Compiler
- Help from the OS
- Help from the CPU
- Internal algorithms

## 2 Programming GDB in Python

- Python Interface Capabilities
- Ex. 1: (re)discovering gdb-cli and gdb.py
- Ex. 2: gdb simple scripting

## 3 New GDB Functionnalités

- Section breakpoint
- Return true breakpoint
- Register watchpoint
- Step into next call
- Faking function execution



# Agenda

Compiler Optimization and Runtime Systems



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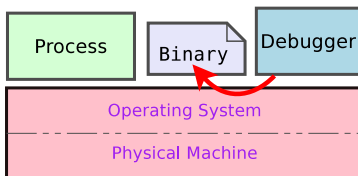


# Under the Hood: Help from the Compiler


Compiler Optimization and Runtime Systems

Everything GDB knows about the **language** (DWARF)

- the type system
- the calling conventions and local variables
- the address-to-line mapping



```
docker@[host]/dwarf $ dwarfdump prodconsum
```




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Compiler Optimization and Runtime Systems

## Everything GDB knows about the language (DWARF)

- the type system
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```
struct Context {  
    pthread_cond_t *cond;  
    ...  
};  
  
void *consumer(void *_context){  
    struct Context *context = ...;  
    ...  
}
```



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Compiler Optimization and Runtime Systems

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DW\_TAG\_subprogram

DW_AT_name	<b>consumer</b>
DW_AT_decl_file	prodconsum.c
DW_AT_type	<0x00000094> # void *
DW_AT_low_pc	0x00400d47
DW_AT_high_pc	<offset-from-lowpc>237
DW_AT_frame_base	len 0x0001: 9c: DW_OP_call_frame_cfa
...	



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- the **type system**
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- the address-to-line mapping

```
DW_TAG_subprogram
```

```
DW_AT_name          consumer
```

```
...
```

```
DW_TAG_formal_parameter
```

```
DW_AT_name          _context
```


```
DW_AT_decl_file     0x00000001 prodconsum.c
```

```
DW_AT_decl_line     0x0000007b  # 123
```

```
DW_AT_type          <0x00000094> # void *
```

```
DW_AT_location      len 0x0002: 9158: DW_OP_fbreg -40
```

```
...
```



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## Everything GDB knows about the language (DWARF)

- the **type system**
- the calling conventions and **local variables**
- the address-to-line mapping

```
DW_TAG_subprogram
```

```
DW_AT_name          consumer
```

```
...
```

```
DW_TAG_variable
```

```
DW_AT_name          context
```

```
DW_AT_decl_file     0x00000001 prodconsum.c
```

```
DW_AT_decl_line     0x0000007d    # 125
```

```
DW_AT_type           <0x00000596> # struct Context *
```

```
DW_AT_location       len 0x0002: 9168: DW_OP_fbreg -24
```

```
...
```


# Under the Hood: Help from the Compiler

Compiler Optimization and Runtime Systems

## Everything GDB knows about the language (DWARF)

- the **type system**
- the calling conventions and local variables
- the address-to-line mapping

```
DW_TAG_pointer_type      # <0x00000596> struct Context*
  DW_AT_byte_size        0x00000008
  DW_AT_type              <0x0000050a>
DW_TAG_structure_type    # <0x0000050a> struct Context
  DW_AT_name              Context
  DW_AT_byte_size        0x00000018
  DW_TAG_member
    DW_AT_name            cond
    DW_AT_type            <0x0000054c> # pthread_cond_t *
  DW_AT_data_member_location 0
```



# Under the Hood: Help from the Compiler

Compiler Optimization and Runtime Systems

## Everything GDB knows about the language (DWARF)

- the **type system**
- the calling conventions and local variables
- the address-to-line mapping

DW_TAG_pointer_type	# 0x00000094 void *
DW_AT_byte_size	0x00000008
DW_TAG_base_type	# 0x0000003f int
DW_AT_name	int
DW_AT_byte_size	0x00000004
DW_AT_encoding	DW_ATE_signed

# Under the Hood: Help from the Compiler

Compiler Optimization and Runtime Systems

## Everything GDB knows about the language (DWARF)

- the type system
- the calling conventions and local variables
- the **address-to-line mapping**

```
<pc>          [lno,col] NS BB ET PE EB IS= DI= uri: "filepath"
0x00400aa6     [ 44, 0] NS uri: "prodconsum.c"
0x00400aae     [ 46, 0] NS
0x00400abc     [ 47, 0] NS
0x00400aca     [ 48, 0] NS
0x00400ad1     [ 50, 0] NS
0x00400ae2     [ 51, 0] NS
0x00400af3     [ 56, 0] NS
0x00400afd     [ 57, 0] NS
```





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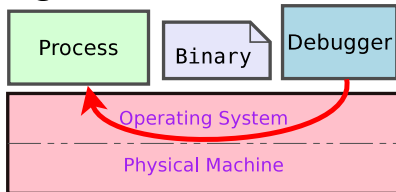
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# Under the Hood: Help from the OS

Compiler Optimization and Runtime Systems

Everything GDB knows about the **execution**



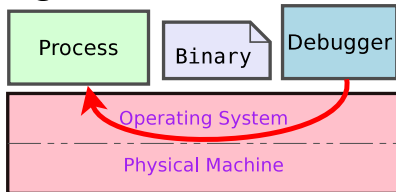
In LINUX: the `ptrace` API  
(link: `kernel/ptrace.c`)

- read/write access to memory addresses
- read/write access to CPU registers
- (re)start and stop the process
- a few more notifications...

# Under the Hood: Help from the OS

Compiler Optimization and Runtime Systems

## Everything GDB knows about the execution

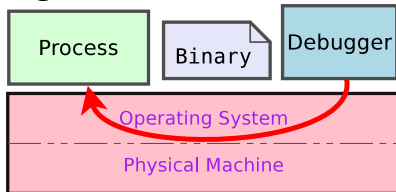


- read/write **access to memory addresses**
  - ▶ `PTRACE_PEEKTEXT`, `PTRACE_PEEKUSER`, `PTRACE_POKE...`
  - ▶ `copy_to_user()`, `copy_from_user()`
- read/write access to CPU registers
- (re)start and stop the process
- a few more notifications...
  - ▶ catching syscalls
  - ▶ handling signals

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Compiler Optimization and Runtime Systems

## Everything GDB knows about the execution

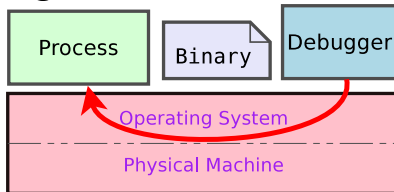


- read/write access to memory addresses
  - ▶ `PTRACE_PEEKTEXT`, `PTRACE_PEEKUSER`, `PTRACE_POKE...`
  - ▶ `copy_to_user()`, `copy_from_user()`
- read/write **access to CPU registers**
  - ▶ registers are saved in the scheduler's `struct task_struct`
  - ▶ `copy_regset_to`, `copy_regset_from_user`
- (re)start and stop the process
- a few more notifications

# Under the Hood: Help from the OS

Compiler Optimization and Runtime Systems

## Everything GDB knows about the execution

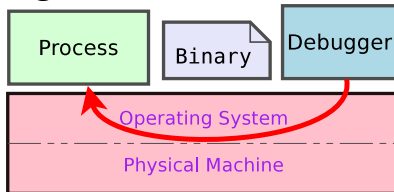


- read/write access to memory addresses
- read/write access to CPU registers
- (re)start and stop the process
  - ▶ basic scheduler operations
  - ▶ ie: put it on the run-queue, send a signal-like interruption request, ...
- a few more notifications...
  - ▶ catching syscalls
  - ▶ handling signals

# Under the Hood: Help from the OS

Compiler Optimization and Runtime Systems

## Everything GDB knows about the execution



- read/write access to memory addresses
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  - ▶ ...



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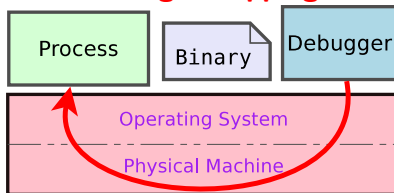
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# GDB Under the Hood: Help from the CPU

Compiler Optimization and Runtime Systems

Everything GDB ... **Single-stepping** and Watchpoints



Single-stepping execute **one CPU instruction**

Watchpoint stop on **memory-address reads and writes**

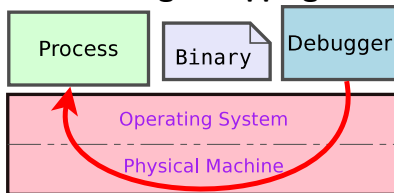
- it's inefficient to implement in software
- main CPUs only have 4 debug registers



# GDB Under the Hood: Help from the CPU

Compiler Optimization and Runtime Systems

~~Everything GDB ...~~ Single-stepping and Watchpoints



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Compiler Optimization and Runtime Systems



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- Callstack
  - newest frame based on CPU registers (IP, FP, BP)
  - older frames based on calling conventions  
(=where registers are stored)
- Finish
  - set temporary breakpoint on the upper-frame PC  
(+ exception handlers / setjumps)
- Step
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Compiler Optimization and Runtime Systems

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# GDB Under the Hood: Internal algorithms

Compiler Optimization and Runtime Systems

Catchpoint

- Kernel notification (via ptrace)

Watchpoint

- CPU notification to the kernel (trap)
- Kernel notification to GDB (ptrace)

or

- Instruction-by-instruction execution
- Instruction parsing to figure out reads and writes  
⇒ very slow!

Breakpoint

- it's a bit more complicated ...





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# GDB Under the Hood: Internal algorithms

Compiler Optimization and Runtime Systems

## The algorithm behind breakpoints

- `original_insn = *&to_breakpoint`
- `*&to_breakpoint = <special instruction>`
- `continue && wait(signal)`
  - ▶ SIGTRAP if ISA has a breakpoint insn (0xcc in x86)
  - ▶ SIGILL if illegal instruction
- `if PC  $\neq$  set(bpts): deliver(signal); done;`
- `otherwise: # breakpoint hit`
  - ▶ `cancel(signal)`
  - ▶ `stop if bpt.cli_condition() || bpt.py.stop()`
  - ▶ `*&to_breakpoint = original_insn`
  - ▶ `cpu(single_step)`
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Compiler Optimization and Runtime Systems


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# GDB Python interface

Compiler Optimization and Runtime SysEms

## Extending

(not for today)

- pretty-printers custom variable printing based on its type
- frame decorators custom display of the callstack
- frame unwinders tell GDB how your callstacks are structured
- more to come (one day):
  - ▶ thread management and process abstractions
    - ★ bypass existing process access mechanisms
    - ★ access to embedded systems, virtual machines, core files ...
    - ★ already possible but in C !

## Scripting

(for today)



## Extending

(not for today)

## Scripting

(for today)

- values and types manipulation
- access the callstack and local variables, registers, ...
- create new commands
- action on breakpoints
- action on events (exec. stop/cont/exit, library loading, ...)
- ...
- for the rest: `gdb.execute("command", to_string=True)`



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- ...
- for the rest: `gdb.execute("command", to_string=True)`



## Interactive part!

- `docker run -it`
  - ▶ `-v $HOME/gdb.py_debug:/home/gdb.py/host`
  - ▶ `-e GROUPEID=$(id -g) -e USERID=$(id -u)`
  - ▶ `--cap-add sys_ptrace # or --privileged`
  - ▶ `pouget/gdb-tuto`
- edit in `host@$HOME/gdb.py_debug` or `docker@~/host`
- consider adding this line in your `$HOME/.gdbinit`
  - ▶ `source $HOME/gdb.py_debug/gdbinit`



## Exercise 1: (re)discovering gdb-cli and gdb.py

- print a variable

```
print i
```

```
(gdb) p context
```

```
$1 = {  
  cond = 0x400e40 <__libc_csu_init>,  
  mutex = 0x4009b0 <_start>,  
  holder = -128,  
  error = 32767  
}
```

- print its type

```
ptype i
```

- print it as another type

```
print (unsigned int) i
```

- print its address / target

```
print &i; print *i
```



## Exercise 1: (re)discovering gdb-cli and gdb.py

- print a variable

```
print i
```

- print its type

```
ptype i
```

*(gdb)* ptype context

```
type = volatile struct Context {  
    pthread_cond_t *cond;  
    thread_mutex_t *mutex;  
    char holder;  
    int error;  
}
```

- print it as another type

```
print (unsigned int) i
```

- print its address / target

```
print &i; print *i
```



## Exercise 1: (re)discovering gdb-cli and gdb.py

- print a variable

```
print i
```

- print its type

```
ptype i
```

- print it as another type

```
print (unsigned int) i
```

```
(gdb) print (unsigned int) context.holder
```

```
$3 = 4294967168
```

- print its address / target

```
print &i; print *i
```

- evaluate C expression

```
i + 1; i & 0x4
```

- evaluate functions

```
f(i)
```





## Exercise 1: (re)discovering gdb-cli and gdb.py

- print a variable
- print its type
- print it as another type
- print its address / target

```
print i
```

```
ptype i
```

```
print (unsigned int) i
```

```
print &i; print *i
```

```
(gdb) p &context.mutex  
$5 = (pthread_mutex_t **) 0x7fffffff588
```

```
(gdb) p *context.mutex  
$6 = {  
  __data = {  
    __lock = -1991643855,  
    ...  
  }  
}
```



## Exercise 1: (re)discovering gdb-cli and gdb.py

- print a variable
- print its type
- print it as another type
- print its address / target

```
print i
```

```
ptype i
```

```
print (unsigned int) i
```

```
print &i; print *i
```

```
# access to variables
i = gdb.parse_and_eval("i")           <gdb.Value(int)>
i.type                                <gdb.Type(int)>
uint = gdb.lookup_type("unsigned int") <gdb.Type(uint)>
i.cast(uint)                          <gdb.Value(uint)>
gdb.newest_frame().read_var("i")
```



## Exercise 1: (re)discovering gdb-cli and gdb.py

- print a variable
- print its type
- print it as another type
- print its address / target

```
print i
```

```
ptype i
```

```
print (unsigned int) i
```

```
print &i; print *i
```

- evaluate C expression
- evaluate functions

```
i + 1; i & 0x4
```

```
f(i)
```



## Exercise 1: (re)discovering gdb-cli and gdb.py

- print a variable
- print its type
- print it as another type
- print its address / target

```
print i
```

```
ptype i
```

```
print (unsigned int) i
```

```
print &i; print *i
```

- evaluate C expression
- evaluate functions

```
i + 1; i & 0x4
```

```
f(i)
```

```
(gdb) p puts("creating first thread") # print or call  
creating first thread  
$10 = 23
```



Your turn! print, evaluate, access, ...

Compiler Optimization and Runtime Systems

```
# frame register access
gdb.newest_frame().older().read_reg("pc")
# function call
ret = gdb.parse_and_eval("puts")("text")    <gdb.Value()>
text
```

- disassemble a specified section of memory `disassemble main`
- in Python: `gdb.execute("disa fct", to_string=True)` or

```
frm = gdb.selected_frame()
frm.architecture().disassemble(frm.read_register("pc"))
[{'addr': 4595344, 'asm': 'sub $0x28,%rsp', 'length': 4}]
```



Exercise 1: (re)discovering gdb-cli and gdb.py

```
docker@~$ cat exercices.md # Discovering gdb-cli and gdb.py
```


Time to work!



# Your turn! commands, breakpoints and events

Compiler Optimization and Runtime Systems

## Exercise 2: GDB Simple Scripting



# Your turn! commands, breakpoints and events

Compiler Optimization and Runtime Systems

## Exercise 2: GDB Simple Scripting

### Defining new commands

CLI

```
define cmd
...
...
end
```

Python

```
class MyCommand(gdb.Command):
    def __init__(self):
        gdb.Command.__init__(self, "cmd", gdb.COM)

    def invoke(self, args, from_tty):
        ...
```



# Your turn! commands, breakpoints and events

Compiler Optimization and Runtime SysEms

## Exercise 2: GDB Simple Scripting

### ■ Conditional breakpoints

```
break <loc> if f(i) == &j
```

- ▶ internally, the **breakpoint is hit all the time**
- ▶ but GDB only notifies the user if the condition is met

CLI

```
break fct
command
  silent
  print i
  cont
end
```

# Your turn! commands, breakpoints and events

Compiler Optimization and Runtime Systems

## Exercise 2: GDB Simple Scripting

### ■ Conditional breakpoints

```
break <loc> if f(i) == &j
```

- ▶ internally, the **breakpoint is hit all the time**
- ▶ but GDB only notifies the user if the condition is met

CLI

```
break fct
command
  silent
  print i
cont
end
```

Python

```
class MyBreakpoint(gdb.Breakpoint):
    def __init__(self):
        gdb.Breakpoint.__init__(self, "fct",
                                internal=True)

        self.silent = True
    def stop(self):
        print(gdb.parse_and_eval("i"))
        return True or False
```



## Executing code on events

```
def say_hello(evt): print("hello")

gdb.events.stop.connect(say_hello) # then disconnect
gdb.events.cont
gdb.events.exited

gdb.events.new_objfile # shared library loads, mainly
gdb.events.clear_objfiles

gdb.events.inferior_call_pre/post
gdb.events.memory/register_changed # user-made changes

gdb.events.breakpoint_created/modified/deleted
```



# Your turn! commands, breakpoints and events

Compiler Optimization and Runtime Systems

## Exercise 2: GDB Simple Scripting

```
docker@~$ cat exercices.md # Hooking into gdb.py
```

Time to work!



- 1 GDB Under the Hood
  - Help from the Compiler
  - Help from the OS
  - Help from the CPU
  - Internal algorithms
- 2 Programming GDB in Python
  - Python Interface Capabilities
  - Ex. 1: (re)discovering gdb-cli and gdb.py
  - Ex. 2: gdb simple scripting
- 3 New GDB Functionnalities
  - Section breakpoint
  - Return true breakpoint
  - Register watchpoint
  - Step into next call
  - Faking function execution



## Adding new functionalities to GDB

- 1 Section breakpoint
  - ▶ `break_section start_profiling stop_profiling run`
- 2 Break when returned `true`
  - ▶ `break_return run 1`
- 3 Register watchpoint
  - ▶ `reg_watch eax main void *`
- 4 Step-to-next-call
  - ▶ `step-before-next-call`
  - ▶ `step-to-next-call`
- 5 Faking function execution
  - ▶ `skip_function run`
  - ▶ `fake_run_function`

<https://sourceware.org/gdb/current/onlinedocs/gdb/Python-API.html>



- `make all; make help`
- `make run_{section|return|watch|step|fake} DEMO={y|n}`
  - ▶ `DEMO=y` to run my code, `DEMO=n` for yours (default)



## Your turn: section.c (1/2)

Compiler Optimization and Runtime Systems

```
int main() {  
    int i;  
  
    srand(time(NULL));  
    int bad = rand() % NB_ITER;  
  
    for(i = 0; i < NB_ITER; i++) {  
        if (i != bad) start_profiling();  
        run(i); # calls bugs(i) if not profiling  
        if (i != bad) stop_profiling();  
    }  
}
```





## Your turn: section.c (2/2)


Compiler Optimization and Runtime Systems

```
void start_profiling(void) {
    assert(!is_profiling);
    is_profiling = 1;
}

void stop_profiling(void) {
    assert(is_profiling);
    is_profiling = 0;
}

int run(int i) {
    if (!is_profiling) bug(i);

    return is_profiling;
}
```



## Section breakpoint

Compiler Optimization and Runtime Systems

### Context

- We want to profile the function `run()`.
  - ▶ profiling starts with function `start_profiling()`
  - ▶ and stops with function `stop_profiling()`.

### Problem

- `run()` is sometimes called outside of the profiling region.  
⇒ we want to stop the debugger there.

```
(gdb) break_section start_profiling stop_profiling run
Section bpt set on start_profiling/run/stop_profiling
(gdb) run
Section breakpoint hit outside of section
15         if (!is_profiling) bug(i);
```



## Section breakpoint


Compiler Optimization and Runtime Systems

### Idea:

- breakpoint on `start_profiling()` that sets a flag,
- breakpoint on `stop_profiling()` that unsets a flag,
- breakpoint on `run()` that checks the flag

### Better:

- `start()` / `stop()` breakpoints enable/disable the bpt on `run()`



## Return true breakpoint

Compiler Optimization and Runtime Systems

### Context

- I want to stop the execution whenever function `run()` has returned `true`.

### Problem (kind of :)


- Function `run()` has many return statements
- I don't want to breakpoint all of them.

```
(gdb) break_return run 1
```

```
(gdb) run
```

Stopped after finding 'run' return value = 1 in \$rax.

```
#0 0x00000000004006f7 in main () at section.c:36
```



## Return true breakpoint

Compiler Optimization and Runtime SysEms

```
(gdb) break_return <fct> <expected value>
```

### Idea:

- BreakReturn\_cmd.invoke
  - ▶ parse and cast the expected value:  
gdb.parse\_and\_eval(<expected value>)
  - ▶ Function breakpoint on target function:  
FunctionReturnBreakpoint(<fct>, <expected value>)
- FunctionReturnBreakpoint.prepare\_before()
  - ▶ before the function call: nothing to do
- FunctionReturnBreakpoint.prepare\_after()
  - ▶ after the call: read register `eax`  
my\_gdb.my\_archi.return\_value(<expected value>.type)



# Register watchpoint

Compiler Optimization and Runtime Systems

## Context

- Inside a function, we want to see all the accesses to a register.

## Problem

- GDB only supports memory watchpoints

```
(gdb) reg_watch eax main void *  
20 watchpoints added in function main  
(gdb) cont  
before: (void *) 0xffffffffffffd256  
0x0000000004006a4 <+18>:      mov    %eax,%edi  
after: <unchanged>  
(gdb) cont  
before: (void *) 0xffffffffffffd256  
0x0000000004006be <+44>:      mov    %ecx,%eax
```



# Register watchpoint

Compiler Optimization and Runtime Systems

```
(gdb) reg_watch <reg name> <fct> [<fmt>]
```

## Idea:

- ensure that target function exists

```
if not gdb.lookup_symbol(fct)[0]:...
```

- ▶ may throw a `gdb.error` if there is no frame selected

- examine the function binary instructions

- ▶ `gdb.execute("disassemble {fct}", to_string=True)`

- for all of them,

- ▶ check if `<reg name>` appears
- ▶ if yes, breakpoint its address (`*addr`)

- ...



# Register watchpoint

Compiler Optimization and Runtime Systems


```
(gdb) reg_watch <reg name> <fct> [<fmt>]
```

## Idea:

- on breakpoint hit:

- ▶ read and print the current value of the register  
`gdb.parse_and_eval("({fmt}) ${regname}")`
- ▶ print the line to be executed (from disassembly)
- ▶ in `my_gdb.before_prompt`:
  - ★ execute instruction (`nexti`)
  - ★ re-read the register value
  - ★ print it if different
- ▶ mandatory stop here  
(GDB cannot `nexti` from a `Breakpoint.stop` callback)





## Step into next call

Compiler Optimization and Runtime Systems

### Context

- I want to step into the next function call, even if far away.

- ▶ stop right **before**

step-before-next-call

- ▶ stop right **after**

step-into-next-call

```
(gdb) step-before-next-call
```

```
step-before-next-call: next instruction is a call.
```

```
0x4006ed: callq 0x40062f <start_profiling>
```

```
(gdb) step-into-next-call
```

```
Stepped into function start_profiling
```

```
#0 start_profiling () at section.c:21
```

```
21         assert(!is_profiling);
```

```
#1 0x00000000004006f2 in main () at section.c:37
```

```
37         if (i != bad) start_profiling();
```



## Idea:

### ■ step-before-next-call:

- ▶ run instruction by instruction

```
gdb.execute("stepi")
```

- ▶ until the current instruction contains a call

```
gdb.selected_frame().read_register("pc")
```

```
arch = gdb.selected_frame().architecture()
```

```
"call" in arch.disassemble(current_pc)[0]["asm"]
```

### ■ step-into-next-call:

- ▶ run step by step: `gdb.execute("stepi")`

- ▶ stop when the stack depth increases

```
def callstack_depth():
```

```
    depth = 1; frame = gdb.newest_frame()
```

```
    while frame: frame = frame.older(); depth += 1
```

```
    return depth
```



# Faking function execution

Compiler Optimization and Runtime Systems

## Context

- I don't want function `run()` code to execute,
- Instead I want to control its side effects from the debugger.

```
(gdb) run
```

```
BUG BUG BUG (i=<random>)
```

```
(gdb) skip_function run; run
```

```
[nothing]
```

```
(gdb) fake_run_function # calls bug(i) if not i % 10
```

```
BUG BUG BUG (i=0)
```

```
BUG BUG BUG (i=10)
```

```
BUG BUG BUG (i=20)...
```



# Faking function execution

Compiler Optimization and Runtime Systems

## Idea:

- `skip_function <fct>`:

- ▶ Breakpoint on `<fct>`, then call return:  
`gdb.execute("return")`

- `fake_run_function`:

- ▶ as above, but run code before return:  
`i = int(gdb.newest_frame().read_var("i"))`  
`if not i % 10: gdb.execute("call bug({})".format(i))`



# Understanding, Scripting and Extending GDB

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