

PROBLEM 3 Counting Shortest Paths [13 points]

Given a graph $G = (V, E)$, and a starting node s , let $\ell(s, t)$ be the length of the shortest path in terms of number of edges between s and t . Design an algorithm that computes the number of distinct paths from s to t that have length exactly $\ell(s, t)$.

Answer

Graph $G = (V, E)$

Starting node = s

Ending node = t

Length of shortest path from s to $t = \ell(s, t)$

For this problem, we just need to use the Breadth First Search Algorithm and modify it a bit. While running BFS, we can keep track of all the nodes [Distance from the start and number of paths from the start].

To do this, we can make use of some extra arrays/maps:

- Distance to store the current distance from the start.
- Paths to store the number of paths from the start.

First, we can initialize the arrays/maps, Distance with infinity (start with 0) and Paths with 0 (start with 1 as there is one path from start to itself).

Then we can update the arrays/maps in the following manner:

- If the distance of some neighbor of node n is bigger than the distance of current node $n + 1$, then we set the distance of the neighbor = to the current node and the $\text{Paths}[\text{neighbor}] = \text{Path}[\text{current}]$.
- If the distance of some neighbor of node n is equal to the distance of the current node $n + 1$, then we add the $\text{Paths}[\text{current}]$ to $\text{Paths}[\text{neighbor}]$.

In the $\text{Paths}[x]$, where x is some node, is stored the number of unique paths from the start to node x . Looking at how this algorithm operates, it goes through each node only once, thus the time complexity will be the same as BFS **$O(n+m)$** .

Some possible algorithm...

- Start performing the BFS algorithm.
- For every neighbor Y of each vertex X , after pushing or not to the queue do:
 - If $\text{distance}[Y] > \text{distance}[X] + 1$, then decrease the $\text{distance}[Y]$ to $\text{distance}[X] + 1$ and assign the number of paths of vertex X to the number of paths of vertex Y .
 - Else if $\text{distance}[Y] = \text{distance}[X] + 1$, then add the number of paths of vertex X to the number of paths of vertex Y .
 - When the BFS queue is empty, then stop.

In the Paths array/map, we have stored all the number of distinct paths from start to x – node.