1 General

1.1 Interviewer Considerations

Notes:

- How did the candidate **analyze** the problem?
- Did the candidate miss any special or **edge** cases?
- Did the candidate approach the problem **methodically** and logically?
- Does the candidate have a strong foundation in basic computer science **concepts**?
- Did the candidate produce working code? Did the candidate test the code?
- Is the candidate's code clean and easy to read and maintain?
- Can the candidate **explain** their ideas clearly?

1.2 Steps for Success During the Technical Interview

Summary:

1. Clarify the question

- (a) Understand what the question is asking and gather example inputs and outputs.
- (b) Clarify constraints such as:
 - i. Can numbers be negative or repeated?
 - ii. Are values sorted or do we need to sort them?
 - iii. Can we assume input validity?
- (c) Asking clarifying questions shows communication skills and prevents missteps.

2. Design a solution

- (a) Avoid immediate coding; propose an initial approach and refine it.
- (b) Analyze the algorithm's time and space complexity.
- (c) Consider and address edge cases.
- (d) Think aloud to demonstrate logical reasoning and collaboration.
- (e) Discuss non-optimal ideas to show your thought process.

3. Write your code

- (a) Structure the solution using helper functions.
- (b) Confirm API details when uncertain.
- (c) Use your strongest programming language and full syntax.
- (d) Write complete, working code—not pseudocode.

4. Test your code

- (a) Validate your solution with 1–2 example test cases.
- (b) Walk through each line using inputs.
- (c) Do not assume correctness—prove it through testing.
- (d) Discuss any further optimizations and their trade-offs.

1.3 Common Mistakes to Avoid

Warning:

- 1. Starting to code without clarifying the problem.
- 2. Failing to write or discuss sample inputs and outputs.
- 3. Using pseudocode instead of fully functional code.
- 4. Misunderstanding the problem or optimizing prematurely.

1.4 Syntax

Summary:

- 1. dict.items()
 - Returns a view object that displays a list of a dictionary's key-value tuple pairs.
- 2. sorted(iterable, key=..., reverse=...)
 - iterable: The sequence or collection (e.g., list, dictionary view) to be sorted.
 - key=...: A function that extracts a comparison key from each element. Sorting is performed based on the result of this function.
 - key=lambda x: x[0]: Sort by the first element of each tuple.
 - key=lambda x: x[1]: Sort by the second element of each tuple.
 - reverse=...: A boolean value. If True, sorted in descending order; otherwise, sorted in ascending order (default is False).

2 Arrays and Hashing

2.1 When to Use?

Summary:

•

2.2 Hashing

Algorithm:

2.3 Common Problems

Summary:

Problem Description:

217. Contains Duplicate Given an integer array nums, return true if any value appears at least twice.

- Use a set to store the elements. If an element is already in the set, return True.
- Otherwise, add it to the set.

242. Valid Anagram Given two strings s and t, return true if t is an anagram of s and false otherwise.

- Use a hashMap to count the frequency of each character in s and t.
- If the frequency maps are equal, return True. Otherwise, return False.
- 1. Two Sum Given an array of integers, return indices of the two numbers s.t. they add up to a specific target.
 - Tricks:
 - Use a hashMap to store the indices of the elements, prevMap[nums[i]] = i
 - For each element, check if the target nums[i] is in the map.
 - If it is, return the index of the target nums[i] (from prevMap) and i. Otherwise, add target nums[i].
- **49. Group Anagrams Given an array of strings, group the anagrams together.
 - Use a hashMap to store a tuple of count of each char as the key and the list of words as the value.
 - For each word, create a tuple of count of each char and add the word to the list in the map.
 - Finally, return the values of the map.

3 Two Pointers

3.1 When to Use?

Summary:

- If we need to find a pair of elements that satisfy a condition.
- If we need to find a subarray that satisfies a condition.

3.2 Slow and Fast Pointers

Algorithm:

1.

3.2.1 Common Problems

Problem	Description:
15. 3Sum	Given an array of integers, return all the triplets $[nums[i], nums[j], nums[k]]$ s.t. $i != j$, $i != k$, and $j != k$.
• Tricks:	
125. Valid Palindrome	Given a string, determine if it is a palindrome, considering only alphanumeric characters and ignoring cases.
• s_new = ".join(char.lower() for lowercase.	or char in s if char.isalnum()) to remove non-alphanumeric as
• Use front and back pointers. If they not equal, return False. If equal move both pointers.	
167. Two Sum II - Input array is sorted	Given an array of integers that is already sorted in ascending order, find two numbers such that they add up to a target.

3.3 Left and Right Pointers

Algorithm:

- 1. Initialize two pointers. Some common choices:
 - \bullet One at the front and one at the back of the array.
 - Both at the front of the array.
 - Both at the back of the array.

3.3.1 Common Problems

Problem	Description:
15. 3Sum	Given an array of integers, return all the triplets $[nums[i], nums[j], nums[k]]$ s.t. $i != j$, $i != k$, and $j != k$.
• Tricks:	
125. Valid Palindrome	Given a string, determine if it is a palindrome, considering only alphanumeric characters and ignoring cases.
• s_new = ".join(char.lower() for lowercase.	or char in s if char.isalnum()) to remove non-alphanumeric an
• Use front and back pointers. If they	y not equal, return False. If equal move both pointers.
167. Two Sum II - Input array is sorted	Given an array of integers that is already sorted in ascending order, find two numbers such that they add up to a target.
 Use front and back pointers. If > to 	arget, move back pointer left. If < target, move front pointer right.

4 Sliding Window

4.1 Fixed Sliding Window

Summary:

- Find a subarray/substring of a fixed size that satisfies a condition.
- Find the maximum or minimum of a subarray of a fixed size.

Algorithm:

```
initialize window_sum = 0
initialize max_result (or other required value)

# Set up initial window
for i in range(0, k):
    window_sum += arr[i]

max_result = window_sum # Initialize result

# Slide the window
for i in range(k, n):
    window_sum += arr[i] - arr[i - k] # Add new element and remove 1st element of prev window
    max_result = max(max_result, window_sum) (or other computation)

return max_result (or other required value)
```

4.1.1 Common Problems

Summary:

Problem	Description:
643. Maximum Average Subarray I	Given an integer array nums and an integer k, return the maximum average value of a subarray of length k.
• Follow template.	
567. Permutation in String	Given two strings s1 and s2, return true if s2 contains a permutation of s1, or false otherwise.
 than sum, get freq of chars. Special Case: If len(s1) > len For: Since contiguous, slide the char (make sure to del key if free free free free free free free f	arough s2 and update freq Map_window by adding new char and removing old

4.2 Dynamic Sliding Window

Summary:

• Find longest or shortest subarray/substring that satisfies a condition.

Algorithm:

```
initialize left = 0
initialize window_state (sum, count, frequency map, etc.)
initialize min_or_max_result

for right in range(n):
    update window_state to include arr[right] # Expand the window

while window_state violates the condition:
    update min_or_max_result (if needed)
    update window_state to exclude arr[left] # Shrink the window
    move left pointer forward

return min_or_max_result
```

4.2.1 Common Problems

Summary	
Summary	७:

Problem	Description:
121. Best Time to Buy and Sell Stock	Given an array where the ith element is the price of a given stock on day i, find the maximum profit you can achieve. You may not engage in multiple transactions.

- Buy low, sell high principle
 - Use left = buy and right = sell, initialized at 0, 1.
 - If price[right] >= price[left], update max profit. Move right pointer since we can still sell for a profit.
 - If price[right] < price[left], move left pointer since we need to find a lower price to buy.
 - Continue until right pointer reaches the end of the array.
- 3. Longest Substring W/O Repeating Characters Given a string s, find the length of the longest substring without repeating characters.
 - Init: Follow template and use frequency map of chars for window state.
 - While: If a char is repeated, move left pointer to right by 1 and adjust freqMap until current char is unique.
 - Change: Compare substring length outside of while with max_res = max(max_res, right left + 1).
- 424. Longest Repeating Character Replacement

 Given a string s that consists of only uppercase English letters, you can replace any letter with another letter.

 Find the length of the longest substr containing the same letter after performing at most k replacements.
 - Init: Follow template and use freqMap of chars for window state.
 - While: If the number of replacements needed exceeds k, i.e. $(r 1 + 1) max_freq > k$
 - Move left pointer to right by 1 and adjust freqMap until the condition is satisfied.
 - Change: Compare substring length outside of while with max_res = max(max_res, right left + 1).
- **76. Minimum Window Substring

 Given two strings s and t, return the minimum window substr of s such that every character in t (including duplicates) is included in the window. If there is no such substring, return ""
 - Init: Set left = 0. Initialize count_t as frequency map of t, count_s for current window, and variables have = 0, required = len(count_t), res = [-1, -1], and resLen = \infty.
 - For right in range(n): Expand window by adding s[right] to count_s. If frequency matches count_t, increment have.
 - While have == required:
 - Update result if current window is smaller.
 - Shrink window by decrementing count_s[s[left]]; if below count_t, decrement have; increment left.
 - Return: s[res[0]:res[1]+1] if valid window found, else empty string.

239. Sliding Window Maximum Given an integer array nums and an integer k, return the maximum value in each sliding window of size k.

- **Init:** Use deque to store indices of elements in the current window.
- For right in range(n):
 - Remove indices that are out of the current window.
 - Remove indices from the back of the deque while the current element is greater than the element at those indices.
 - Append the current index to the deque.
 - If the window size is reached, append the maximum (element at the front of the deque) to the result list.

5 Binary Search

5.1 When to Use?

Summary:

• If array is sorted.

6 Linked List

Summary: Data structure for storing objects in linear order.

• Object: Data and a pointer to the next object.

6.1 When to Use?

Summary:

- Implement other DS: stacks, queues, hash tables.
- Dynamic memory allocation.

6.2 Operations

Summary	σ.

Operation	Time Complexity
Search	O(n)
Insert	O(1)
Delete	O(1)
Access	O(n)

6.3 Singly Linked List

Algorithm:

6.4 Doubly Linked List

Algorithm:

6.5 Circular Linked List

Algorithm:

6.6 Common Problems

Summary:

Problem Description:

206. Reverse Linked List Given the head of a singly linked list, reverse the list and return the reversed list.

• Iterative:

- Init: None
$$\rightarrow$$
 0 \rightarrow 1 \rightarrow 2

- While loop until curr is None. curr will point to prev, then curr will get updated to a temp that has curr.next and prev will be updated to curr.

$$\begin{array}{c} * \ \, \underbrace{\mathrm{None}}_{\mathrm{prev}} \leftarrow \underbrace{0}_{\mathrm{curr}} \rightarrow \underbrace{1}_{\mathrm{temp}} \rightarrow 2 \\ * \ \, \mathrm{None} \leftarrow \underbrace{0}_{\mathrm{prev=curr}} \rightarrow \underbrace{1}_{\mathrm{curr=temp}} \rightarrow 2 \end{array}$$

- prev will be the new head.