

# 1 Exercise 1

1.

$D^0$ ,  $D^1$ , and  $D^2$  are as follows:

$$\begin{pmatrix} 0 & 80 & +\infty & 10 \\ +\infty & 0 & 10 & +\infty \\ 40 & +\infty & 0 & +\infty \\ +\infty & 20 & 90 & 0 \end{pmatrix}$$
$$\begin{pmatrix} 0 & 80 & +\infty & 10 \\ +\infty & 0 & 10 & +\infty \\ 40 & \mathbf{120} & 0 & \mathbf{50} \\ +\infty & 20 & 90 & 0 \end{pmatrix}$$
$$\begin{pmatrix} 0 & 80 & \mathbf{90} & 10 \\ +\infty & 0 & 10 & +\infty \\ 40 & 120 & 0 & 50 \\ +\infty & 20 & \mathbf{30} & 0 \end{pmatrix}$$

**2.1**  $\Theta(n)$  **2.2**  $\Theta(1)$  **2.3**  $\Theta(1)$

**3.1** 6, F. Visited event points: ABCDEF

**3.2** 8, H. Visited event points: BFCDH

**4.1** 4

**4.2** 4

**4.3** 3

Range tree:  $\Theta(\lg^2 n + k)$

kd tree:  $\Theta(\sqrt{n} + k)$

2 BSTs:  $\Theta(\lg n + k_x + k_y)$

workload 1:  $k = \frac{n}{4}$ , workload 2:  $k = \frac{n}{100}$ , workload 3:  $k$  is a constant.

5

**5.1** 12000 or 20000

**5.2** 4000

**5.3** 4000

1 TB file means  $\frac{1TB}{1GB} = 1000$  disk pages.

**5.1** 50 GB memory means 50 disk pages. 2-way merge sort.  $\lceil \log_2(1000) \rceil = 10$ .  
10 iterations. each iteration reads in 1000 disk and writes 1000 disk. in total, each iteration 2000 I/O operations. So  $2000 * 10 = 20000$ .

Or: Depends on how do we define a 2-way merge sort. Initial main memory sortings (resulting in 20 runs) + ceil ( $\log_2 20$ ) = 1 + 5 = 6 iterations. each iteration 2000 I/O of the file so, 12000.

**5.2** 50 GB memory means 50 disk pages. multi-way merge sort.  $\lceil \log_{49}(1000) \rceil = 2$ . 2 iterations. each iteration reads in 1000 disk and writes 1000 disk. in total, each iteration 2000 I/O operations. So  $2000 * 2 = 4000$ .

Or: Initial main memory sortings (resulting in 20 runs) + ceil ( $\log_{49} 20$ ) = 1 + 1 = 2 iterations. each iteration 2000 I/O of the file so, 4000.

**5.3** 50 GB memory means 500 disk pages. multi-way merge sort.  $\lceil \log_{499}(1000) \rceil = 2$ . 2 iterations. each iteration reads in 1000 disk and writes 1000 disk. in total, each iteration 2000 I/O operations. So  $2000 * 2 = 4000$ .

Or: Initial main memory sortings (resulting in 20 runs) + ceil ( $\log_{499} 20$ ) = 1 + 1 = 2 iterations. each iteration 2000 I/O of the file so, 4000.

**6.3**

## 2 Exercise 2

1

Label activities by finishing time such that:  $f_1 \leq f_2 \leq f_3 \leq \dots \leq f_n$

Define  $i_j = \text{largest index } i < j \text{ such that}$   
 $\text{activity } i \text{ is compatible with activity } j.$

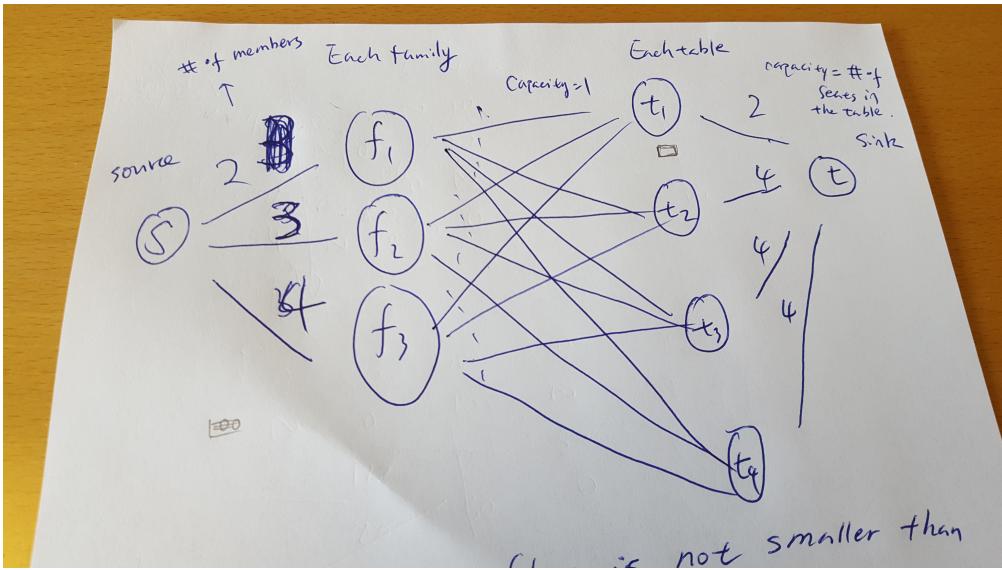
$$OPT(j) = \begin{cases} 0 & \text{if } j = 0 \\ \max(w_j + OPT(i_j)), OPT(j-1) & \text{otherwise} \end{cases}$$

$\Downarrow$  choose activity  $j$       not choosing activity  $j$

**2** We have a source and a sink. We have X vertices, one for each family. We have Y more vertices, one for each table.

Connect the source to each family vertex with an edge, where the capacity is the number of the family members.

Connect the sink to each table vertex with an edge, where the capacity is the number of tables.



Connect each family vertex with each table vertex with an edge, where the capacity is 1.

If the maximum flow is not smaller than the total number of family members, it is possible to arrange the seats according to the rule. Otherwise, it is impossible.

For an edge that connect family vertex  $f_i$  and table vertex  $t_j$ : 1 means that a family member from  $f_i$  should sit table  $t_j$ .