

# PP Lecture 3

## 2020

# Learning Goals

- Simulation of OOP with functions
- Handling of imperative language mechanisms in relation to functional programming
- Understanding and use of continuations
  - Including continuation passing style
- Delayed evaluation with trampolining

# The plan for this live stream

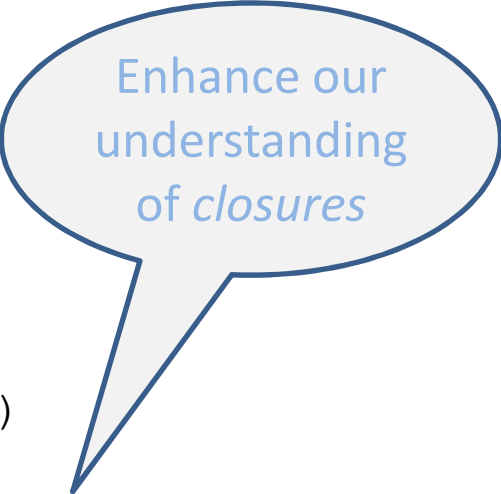
- Something about OOP simulation
- Working with objects in Scheme
- Flexible parameters
- Something about continuations
- Programming as language development
- Exercise intro

Questions and answers – mediated by Jonas Hansen

OOP simulation

```
(define (send message obj . par)
  (let ((method (obj message)))
    (apply method par)))
```

```
(define (point x y)
  (letrec ((getx (lambda () x))
            (gety (lambda () y))
            (add (lambda (p)
                    (point
                     (+ x (send 'getx p))
                     (+ y (send 'gety p))))))
    (type-of (lambda () 'point))
  )
  (lambda (message)
    (cond ((eq? message 'getx) getx)
          ((eq? message 'gety) gety)
          ((eq? message 'add) add)
          ((eq? message 'type-of) type-of)
          (else (error "Message not understood"))))))
```



Enhance our  
understanding  
of *closures*

Define `send` before `point` in Racket

```
(define (point x y)
  (letrec ((getx      (lambda () x))
            (gety      (lambda () y))
            (add        (lambda (p)
                           (point
                            (+ x (send 'getx p))
                            (+ y (send 'gety p))))))
    (type-of (lambda () 'point)))
  )
  (lambda (message)
    (cond ((eq? message 'getx) getx)
          ((eq? message 'gety) gety)
          ((eq? message 'add)  add)
          ((eq? message 'type-of) type-of)
          (else (error "Message not understood")))))
```

```
(define (point x y)
  (let ((x x)
        (y y))

    (define (getx) x)

    (define (gety) y)

    (define (add p)
      (point
       (+ x (send 'getx p))
       (+ y (send 'gety p)))))

    (define (type-of) 'point)

    (define (self message)
      (cond ((eqv? message 'getx) getx)
            ((eqv? message 'gety) gety)
            ((eqv? message 'add) add)
            ((eqv? message 'type-of) type-of)
            (else (error "Message not understood" message))))

    self))
```

```
(define (new-instance class . parameters)
  (apply class parameters))

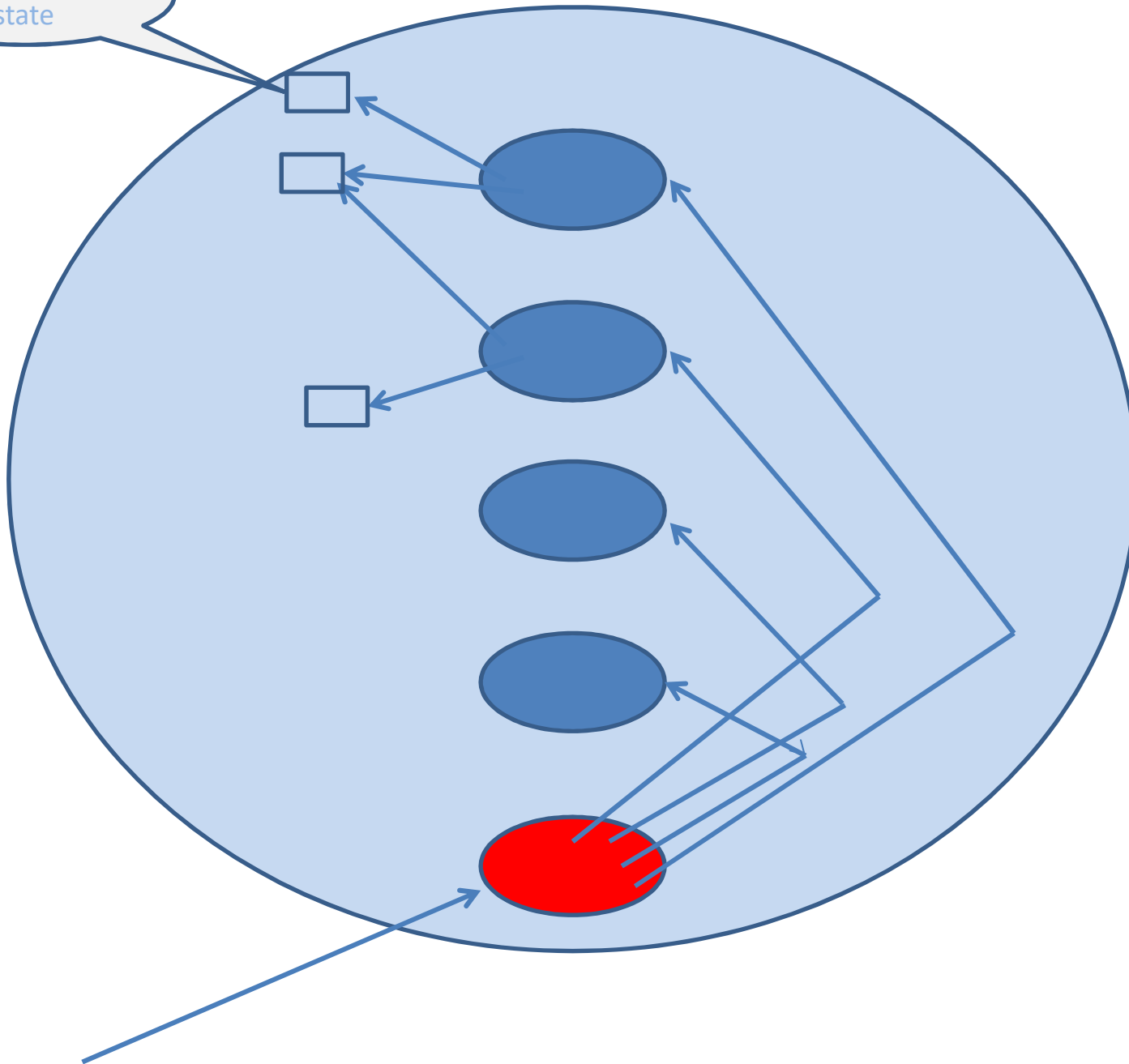
(define (send message object . args)
  (let ((method (method-lookup object message)))
    (cond ((procedure? method) (apply method args))
          (else (error "Error in method lookup " method)))))

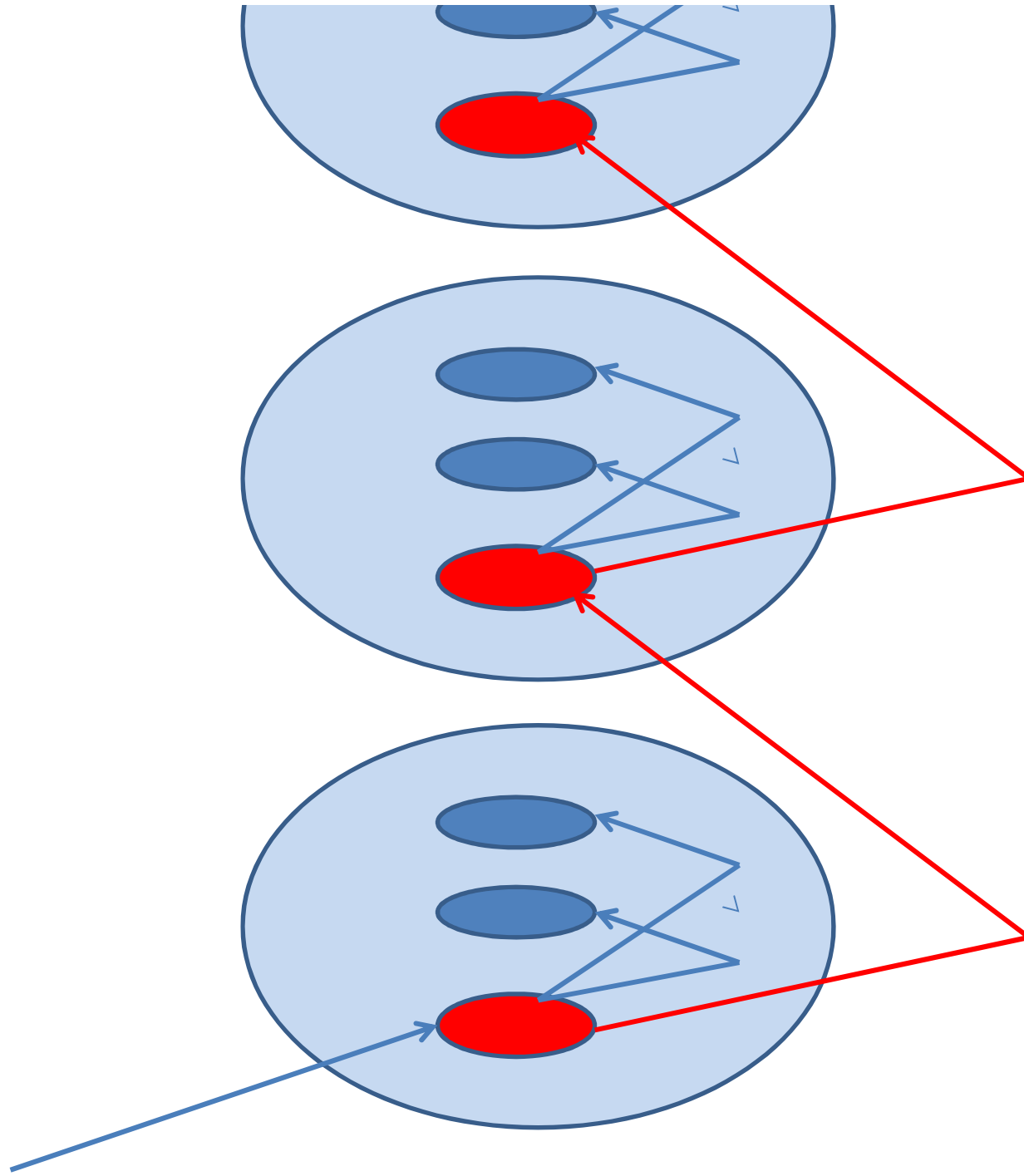
(define (method-lookup object selector)
  (cond ((procedure? object) (object selector))
        (else
         (error "Inappropriate object in method-lookup: "
                object)))))
```

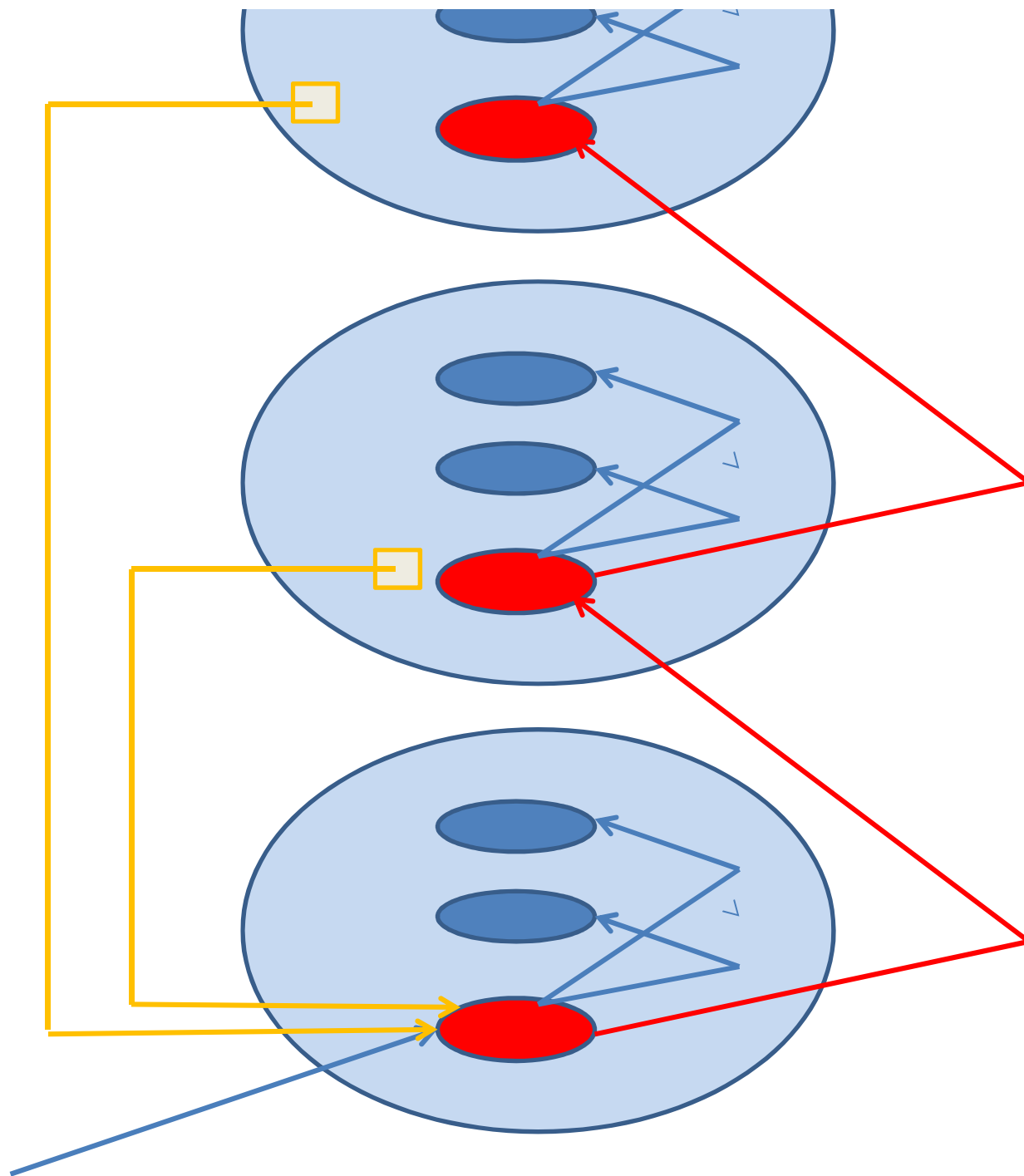
REPL SESSIONS:  
naked-point.scm and point.scm



Encapsulated  
state







# Working with *objects* in Scheme

Objects in a broad sense

# Objects as lists

(point 5 -10)

- Just lists, not alists
  - First element is a tag
- OOP inspiration:
  - Constructor, predicate, selectors

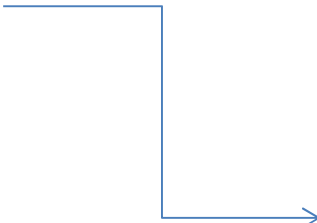
List-based representations are flexible

It is easy to add additional elements, for the purpose of future program modifications.

# Objects as association lists

- Association lists
  - A flexible approach
  - Updates by shadowing
  - Constructor functions
  - Predicate
  - Selector functions

```
(  
  (type . point)  
  (x . 5)  
  (y . -10)  
)
```



```
(  
  (x . 17)  
  (type . point)  
  (x . 5)  
  (y . -10)  
)
```

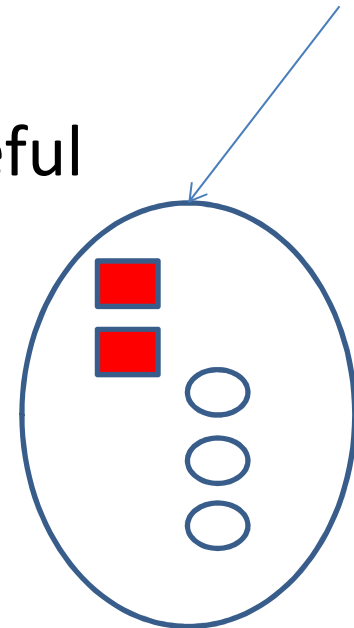
A kind of dictionary – mapping keys to values

# Avoid car/cdr programming

- With constructors, predicates and selector functions:
  - Only these functions know about the list representation.
  - Only these functions use car/cdr
- Without:
  - Hard to change your mind wrt. Representation
  - All parts of your program will be polluted with car/cdr.

# Objects as closures

- Representation of objects as closures
  - Local variables in let-bindings inside a closure
- Interesting in principle
  - But it can be used for real-life purposes
  - Troublesome in several respects
  - Better language support would be useful
    - Syntactically
    - Implementation-wise





# Objects as structs

- Structs are not part of R5RS
  - But they exist in R6RS and Racket
  - SRFI 9 defines a portable record type
    - May rely on R5RS vectors
    - Vectors are heterogeneous in Scheme
      - Of fixed lengths – cannot easily be enlarged
      - Mapping of field names to indexes

type → 0  
x → 1  
y → 2

0	point
	10
	17
	...
4	...

# Objects in the PP1 Mini Project

- Representation of student, group and grouping
- Be sure to encapsulate your representation decisions
- List-based representations are natural in a list-based language...

You may start on this part of the mini project  
already today

# Flexible parameters

- the use of rest parameters

REPL session: flexible-parameters

# The LAML libraries

**LAML & Scheme  
Browser:**  
*All defined names*  
R5RS Scheme Manual  
LAML Root Library  
General LAML Library  
XML-in-LAML Library  
LAML Compatibility  
XHTML Transitional  
XHTML Convenience

General LAML Library:  
**Top**  
**Abstract**  
**Table of Content**  
**Definition Index**  
CR  
absolute-file-path?  
absolute-url?  
accumulate-right  
alist-from-keys-and-values  
alist-to-propertylist  
as-01-boolean  
as-boolean  
as-char  
as-list  
as-number  
as-quoted-string  
as-string  
as-symbol  
as-two-complement-signed-number  
binary-search-in-vector  
binary-to-hex-string  
bit-list-to-byte-string  
bite-of-length  
bite-of-varied-length  
bite-while-accumulate

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**Reference Manual of the General LAML library**  
Kurt Nørmark © normark@cs.aau.dk    Department of Computer Science, Aalborg University, Denmark.  
LAML Source file: lib/general.scm

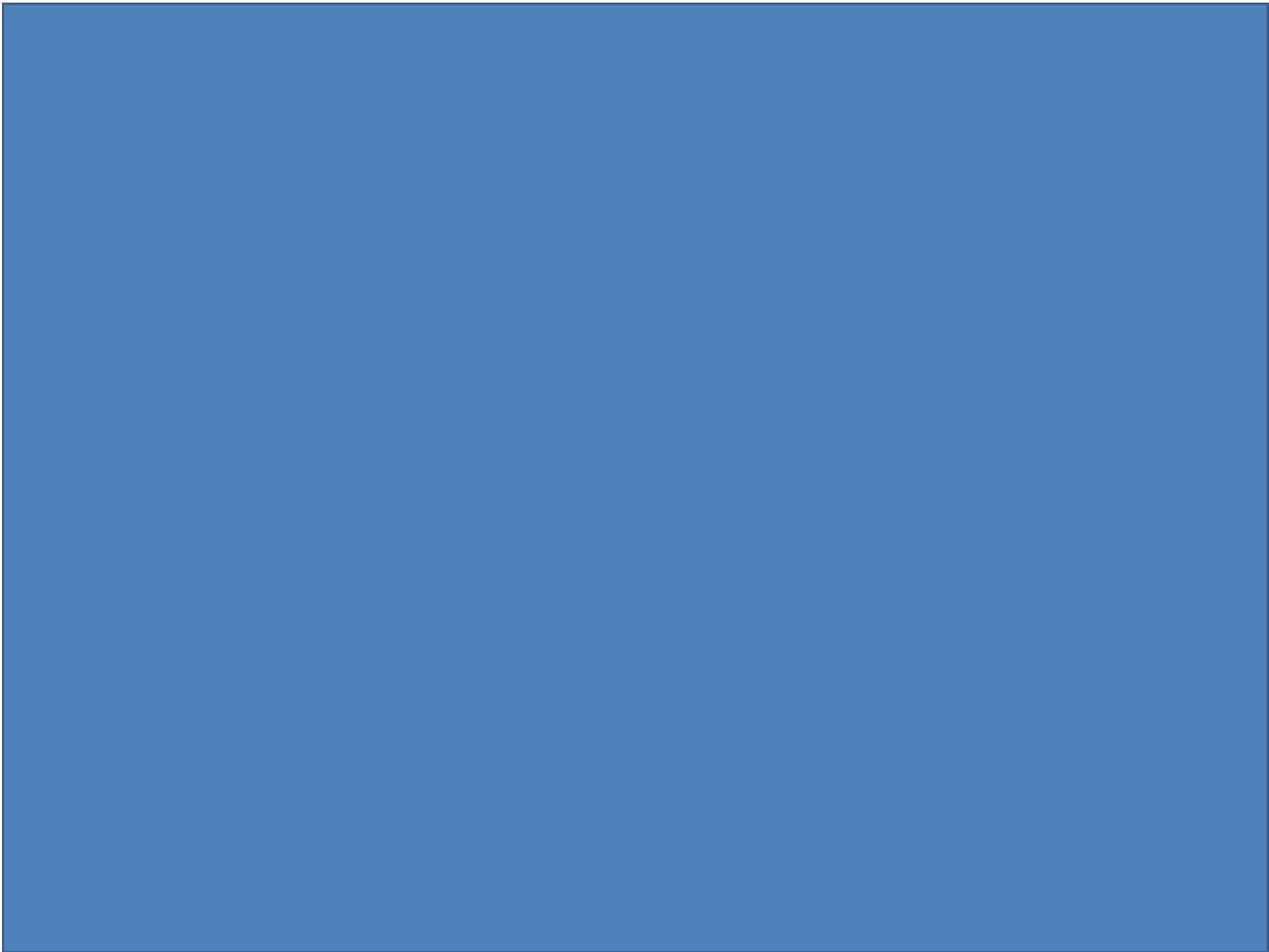
This is a library of common and generally useful Scheme functions, which are used in other LAML libraries, in LAML styles, and in LAML tools. Far the majority of the functions can also be used outside LAML.

**Table of Contents:**

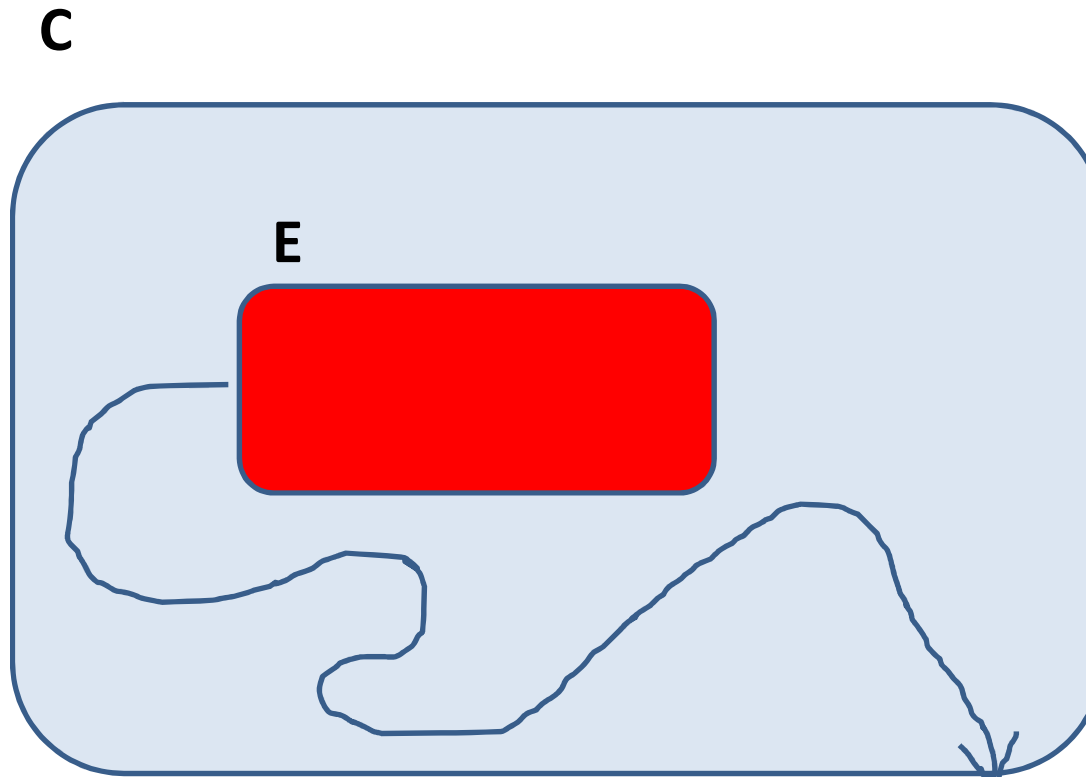
<a href="#">1. Optional parameter handling.</a>	<a href="#">7. Other higher-order functions.</a>	<a href="#">13. Functions that change letter case in string.</a>
<a href="#">2. List selection functions and their generators.</a>	<a href="#">8. List and Sexpr functions.</a>	<a href="#">14. Byte string functions.</a>
<a href="#">3. Association and property list functions.</a>	<a href="#">9. Vector functions.</a>	<a href="#">15. Message displaying and error handling procedures.</a>
<a href="#">4. Filter and accumulation functions.</a>	<a href="#">10. Conversion functions.</a>	<a href="#">16. File name, file path and URL functions.</a>
<a href="#">5. Mapping functions.</a>	<a href="#">11. String predicates.</a>	<a href="#">17. Other functions.</a>
<a href="#">6. Higher-order bite functions.</a>	<a href="#">12. Other string functions.</a>	<a href="#">18. Bite Generators.</a>

**Alphabetic index:**


<a href="#">absolute-file-path?</a>	(absolute-file-path? x)	Return whether x represents an absolute path to a file.
<a href="#">absolute-url?</a>	(absolute-url? x)	Does the string x represent an absolute URL.
<a href="#">accumulate-right</a>	(accumulate-right f init lst)	A higher-order function which right accumulates the list lst by means of the binary function f, using init as the initial value of the accumulation.
<a href="#">alist-from-keys-and-values</a>	(alist-from-keys-and-values key-list val-list)	Make an alist from a key-list and a val-list.
<a href="#">alist-to-propertylist</a>	(alist-to-propertylist alist)	Make and return a property list from an association list.
<a href="#">as-01-boolean</a>	(as-01-boolean x)	Convert x to C-style boolean values, 0 or 1.
<a href="#">as-boolean</a>	(as-boolean x)	Convert x to a boolean.
<a href="#">as-char</a>	(as-char x)	Convert x to a character.
<a href="#">as-list</a>	(as-list x)	Convert x to a list.
<a href="#">as-number</a>	(as-number x)	Convert x to a number.
<a href="#">as-quoted-string</a>	(as-quoted-string x)	Convert x to a string, in which string constituents themselves are quoted.
<a href="#">as-string</a>	(as-string x)	Convert x to a string.
<a href="#">as-symbol</a>	(as-symbol x)	Convert x to a symbol.
<a href="#">as-two-complement-signed-number</a>	(as-two-complement-signed-number i n)	Re-interpret the positive integer i as an n bit two's complement number.
<a href="#">binary-search-in-vector</a>	(binary-search-in-vector v el sel el-eq? el-leq?)	Search for an element el in the sorted vector v.
<a href="#">binary-to-hex-string</a>	(binary-to-hex-string byte-string)	Given byte-string, which is binary data.



# Continuations



A continuation of an expression  $E$  in a contextual expression  $C$  is the future/rest of the computation  $C$ , which waits for (and depends on) the value of  $E$

(lambda (E) )

```
(let ((x 5)
      (y 4))
  (if (even? x)
      (+ (* x x) (* y y))
      (+ (* x x) (call/cc
                  (lambda (e)
                    (+ (* y y) (e y) x)))))))
```

*The value is 29.*

*Why?*



```
(define (list-length l)
  (call-with-current-continuation
    (lambda (do-exit)
      (letrec ((list-length1
                  (lambda (l)
                    (cond ((null? l) 0)
                          ((pair? l)
                           (+ 1 (list-length1 (cdr l))))
                          (else (do-exit 'improper-list))))))
        (list-length1 l)))  ))
```

# Continuation Passing Style

```
(define (f a b c d k)
  (k (+ a b c d)))
```

```
(define (g a k)
  (k (* a a)))
```

```
(define (h a k)
  (k (* a a a)))
```

```
(define (w a b k0)
  (g a (lambda (v1)
    (h b (lambda (v2)
      (g b (lambda (v3)
        (h a (lambda (v4)
          (f v1 v2 v3 v4 k0))))))))))
```

# Direct Style

```
(define (f a b c d)  
  (+ a b c d))
```

```
(define (g a)  
  (* a a))
```

```
(define (h a)  
  (* a a a))
```

```
(define (w a b)  
  (f (g a) (h b) (g b) (h a)))
```

# Continuation Passing Style

## Interesting Observations

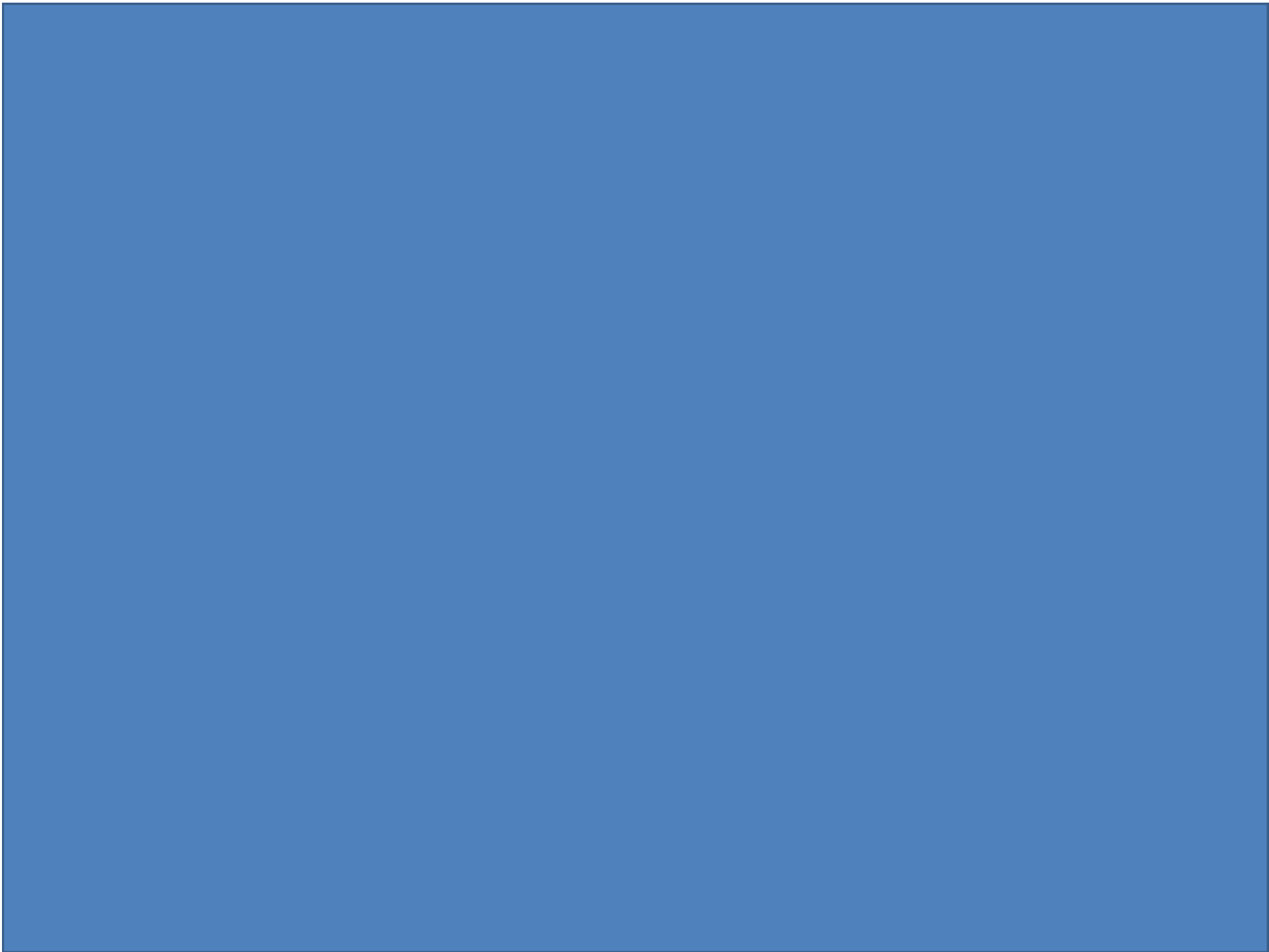
- Functions written in CPS are always tail recursive
- Tail recursive in CPS functions do not need to construct new continuations functions
- Functions written in CPS do not need 'the magic primitive' call-with-current-continuation
- Functions written in CPS are specific about the evaluation order of sub-expressions
- Functions in CPS are typically automatically translated from functions in direct style
- Functions written in CPS are subject to trampolining

# The use of continuations in Scheme Programming

- Tricky one-liners
- Exit from deep expressions
  - list-length of improper lists
  - find-in-tree programmed with for-each
- Advanced control flow
  - Producer consumer functions
  - Simultaneous traversal of two binary trees
  - Coroutines

# Programming as language development

- Your own functions are similar to the native functions in the language
- Your own syntactic extensions are similar to the special forms
- As part of your solution, you may implement an interpreter for a little language
- Ultimately, you may play with a meta circular interpreter
  - Scheme in Scheme



# About the exercises

- *There are intros and outros to all exercises apart from the last one*
- Two exercises about simulated OOP with closures
- One about call-with-current-continuation
- One about continuation passing style
- One about trampolining



# Exercise 3.1

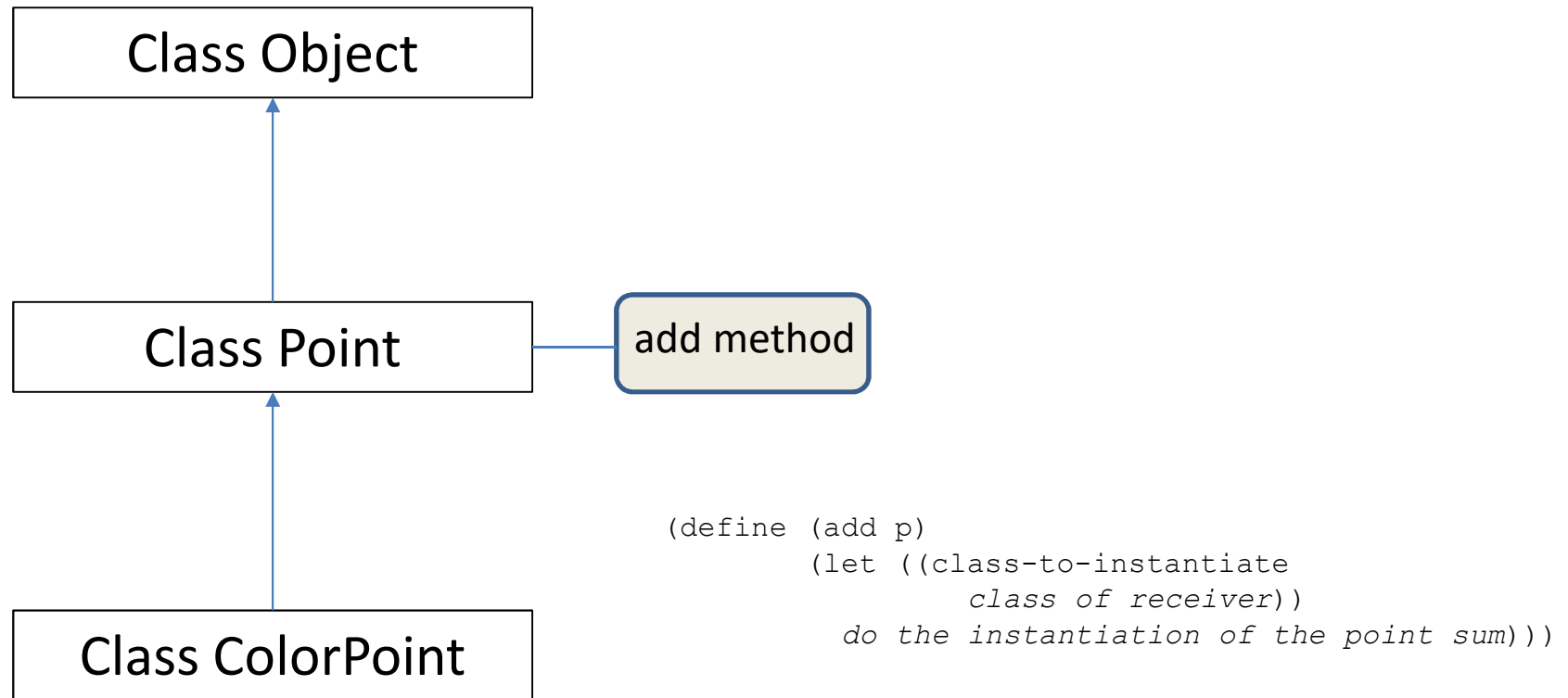
Class Point and class Rectangle

- A move method in class point
- A new class rectangle
  - Represented by points
  - move and area method

# Exercise 3.2

Class Point and subclass ColorPoint.  
The add method should get the class of the point  
... and instantiate a new point of that class

- Almost trivial methods class-of in point and color-point
- Modifying the add method, such that it instantiates the most class
  - With a single problem, perhaps



## Exercise 3.7

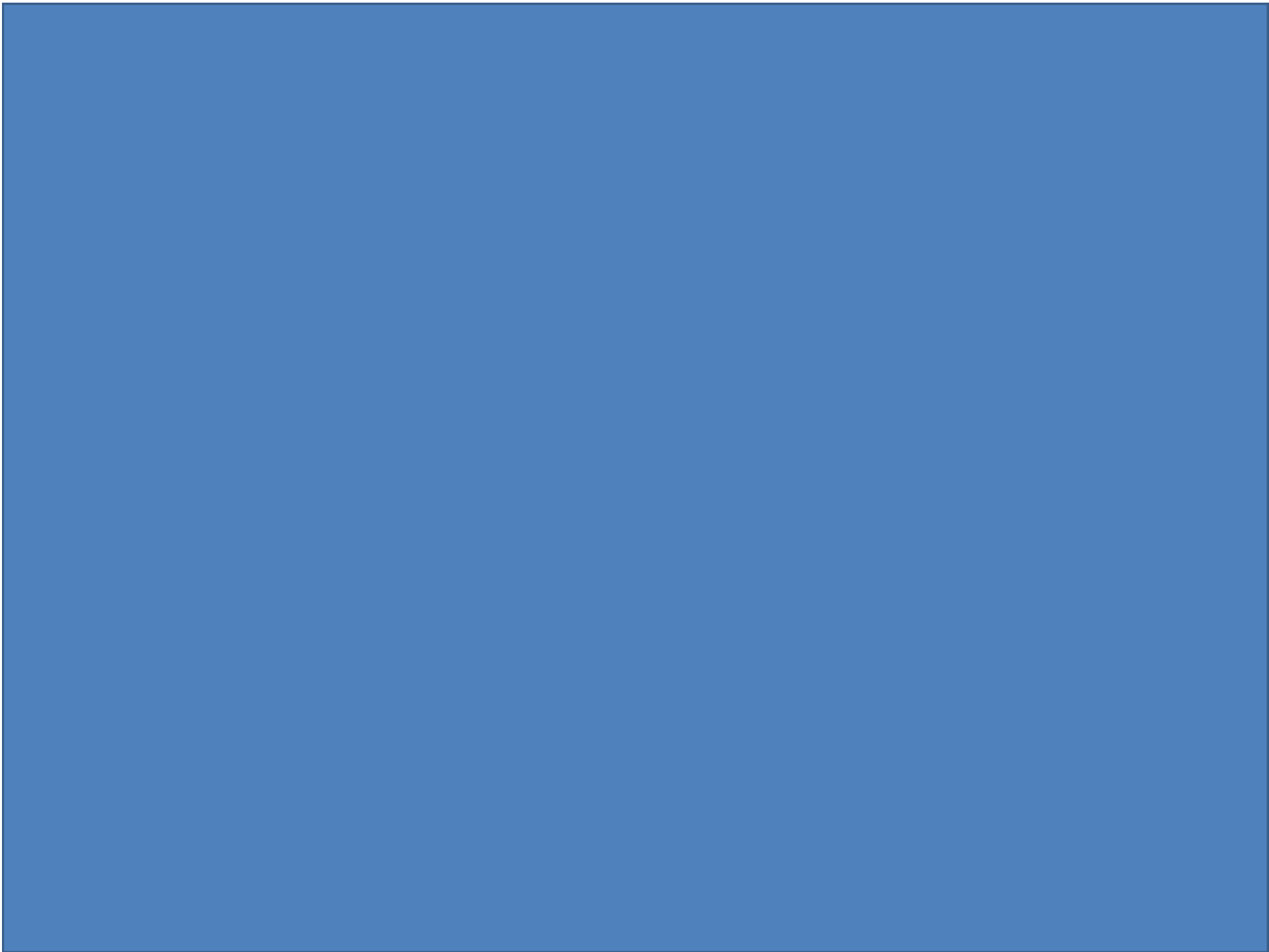
- An exercises about
  - call-with-current-continuation
- Understanding an expression that
  - captures a continuation
  - uses the continuation

## Exercise 3.4

- An expression that trains you to rewrite simple expressions to *continuation passing style*

## Exercise 3.6

- Not all of you get to this exercise
- About a refined sessaw trampolining function



# Start to think about your work on the PP1 miniproject

- Representation of students, groups and groupings
- Constructors
  - Students
  - Groups
  - Groupings
- Selectors and predicates



# PP from 10:15 – 12:00

- Uninterrupted exercises in groups
- Enter your group channel in the Programming Paradigms team
  - **Start a meeting** in the channel
  - We will drop into the meeting and discuss PP topics
- Ask for help in the general channel:
  - Group @<group-id> needs help about ...

