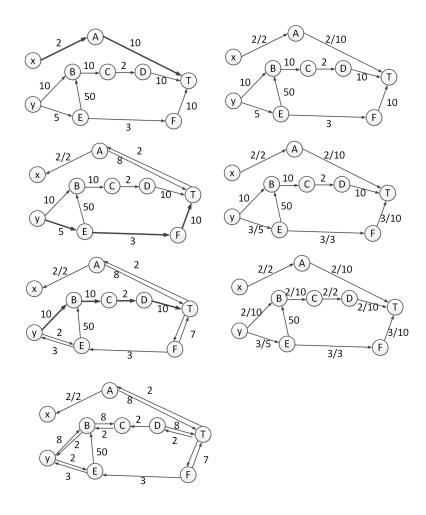
## 1 Exercise 1

1.



Critical edges: XA, EF, and CD.

Maximum flow: 7.

## **2.** A:111 D:10

**3** J

Sweepline status:

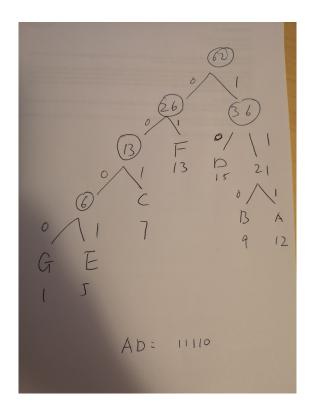
Α

H, A

Η

H, D

J, H, D



**4** a

Range tree:  $\Theta(\lg^2 n + k)$ 

kd tree:  $\Theta(\sqrt{n} + k)$ 

2 BSTs:  $\Theta(\lg n + k_x + k_y)$ k is zero in this specific case

## 2 Exercise 2

1

<i>i</i> -th Operation	1	2	3	4	5	6	7	8	9	10
Cost	1	1	3	1	1	1	1	1	9	1
Cost	1	1	$1+(3^1-1)$	1	1	1	1	1	$1+(3^2-1)$	1

Aggregate analysis: When i is a power of 3, we write the cost into  $1 + (3^x - 1)$ , so that it also has a cost of 1 and the remaining cost  $(3^x - 1)$ . x is from 1 to  $log_3 n$ .

Then, since we have n operations, each has at least a cost of 1, so the sum is n. Next, the sum of the remaining cost is  $\sum_{x=1}^{\log_3 n} (3^x - 1) = \sum_{x=1}^{\log_3 n} 3^x - \log_3 n = \frac{3}{2}n - \frac{3}{2} - \log_3 n \le 1.5n$ .

So all together the cost is less than 2.5n. The amortized cost is at most 2.5. **2** Configuration: (X,Y): X contains the vertices that are already in the cycle and Y contains the vertices to be considered.

From the last vertex in X, we check whether there exist an adjacent vertex that is not in X. If we find such a vertex, we add the vertex as part of the solution. If we do not find a vertex and Y is not empty, then it is a deadend. If we do not find a vertex and Y is empty, and there exists an edge connecting the last vertex in X and the first vertex in X, then it is a solution.

