

An Introduction to Statecharts Modelling and Simulation

Simon Van Mierlo

simon.vanmierlo@uantwerpen.be

Hans Vangheluwe

hans.vangheluwe@uantwerpen.be



McGill

INTRODUCTION

STATECHARTS BASICS

YAKINDU IN DEPTH

ADVANCED CONCEPTS

Reactive Systems

- Completely reactive
- In control of events
- Events trigger output



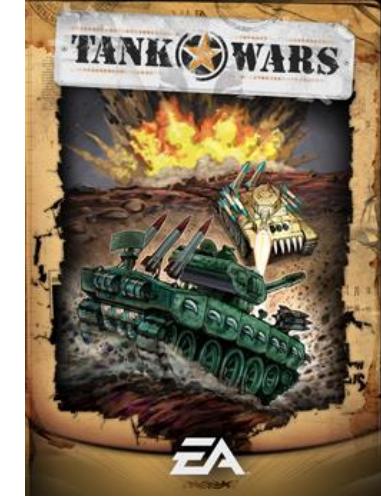
Modelling Reactive Systems

- Interaction with the environment: reactive to events
- Autonomous behaviour: *timeouts + spontaneous transitions*
- System behaviour: *modes (hierarchical) vs concurrent units*
- Use programming language (and OS) is too low-level
-> most appropriate formalism: "what" vs. "how"
*"nontrivial software written with threads, semaphores, and mutexes are incomprehensible to humans"*¹

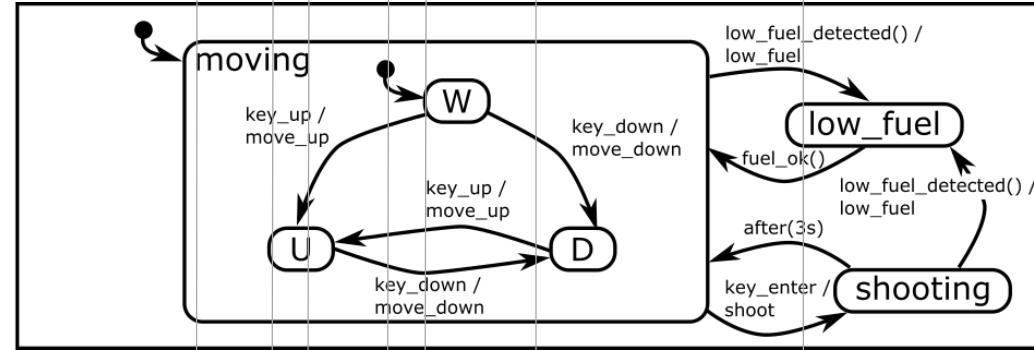


¹E. A. Lee, "The problem with threads," in *Computer*, vol. 39, no. 5, pp. 33-42, May 2006.

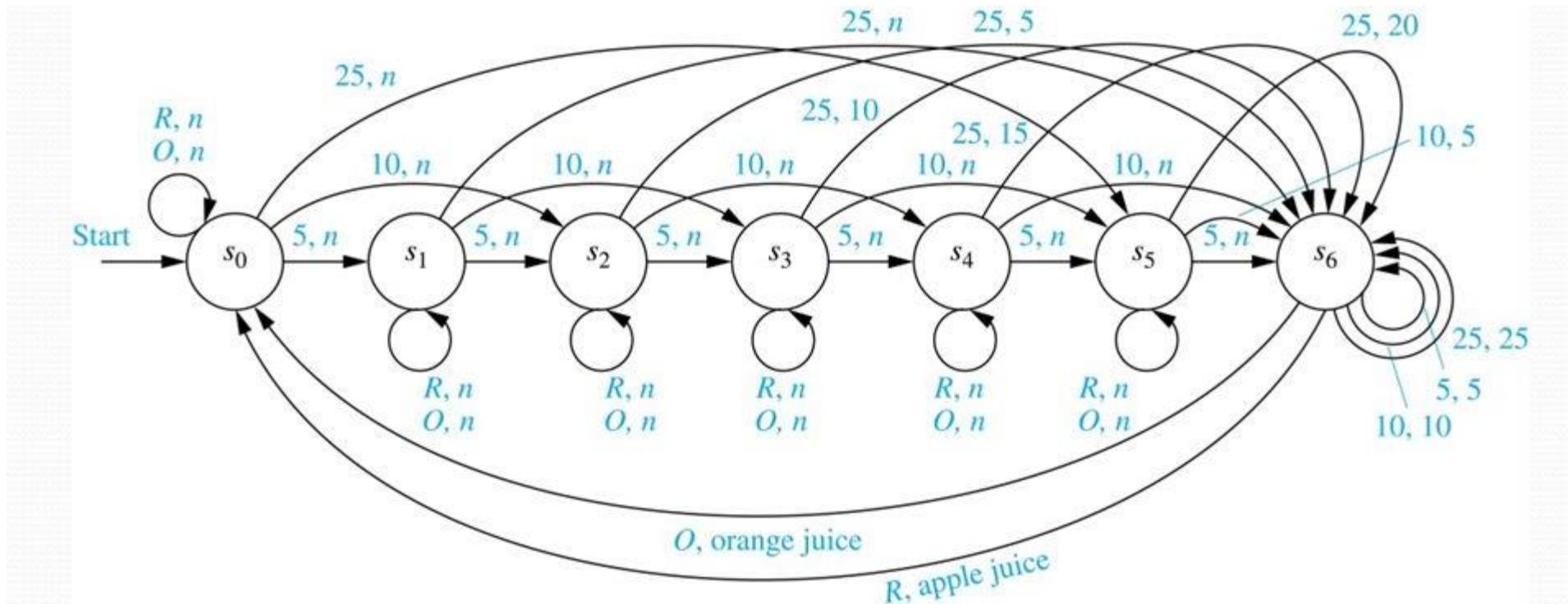
Discrete-Event Abstraction



behavioural
model

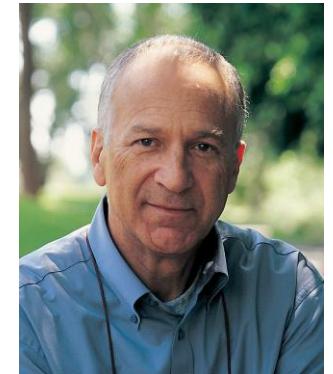


State Diagrams



- Non-modular: hierarchical decomposition (orthogonal/depth) not possible
- State space limited (positive: analysability, negative: expressitivity)
- Becomes too large too quickly to be usable

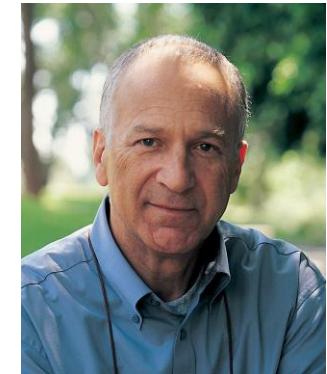
Statecharts History



- Introduced by David Harel in 1987¹
- Notation based on higraphs = hypergraphs + Euler diagrams
- Semantics extend deterministic finite state automata with:
 - Depth (Hierarchy)
 - Orthogonality
 - Broadcast Communication
 - Time
 - History
 - Syntactic sugar, such as enter/exit actions

¹David Harel, Statecharts: a visual formalism for complex systems, Science of Computer Programming, Volume 8, Issue 3, 1987, Pages 231-274

Statecharts History



Science of Computer Programming 8 (1987) 231-274
North-Holland

STATECHARTS: A VISUAL FORMALISM FOR COMPLEX SYSTEMS*

David HAREL
Department of Applied Mathematics, The Weizmann Institute of Science, Rehovot, Israel

Communicated by A. Pnueli
Received December 1984
Revised July 1986

Abstract. We present a broad extension of the conventional formalism of state machines and state diagrams, that is relevant to the specification and design of complex discrete-event systems, such as multi-computer real-time systems, communication protocols and digital control units. Our diagrams, which we call *statecharts*, extend conventional state-transition diagrams with essentially three elements, dealing, respectively, with the notions of hierarchy, concurrency and communication. These transform the language of state diagrams into a highly structured and economical description language. Statecharts are thus compact and expressive—small diagrams can express complex behavior—as well as compositional and modular. When coupled with the capabilities of computerized graphics, statecharts enable viewing the description at different levels of detail, and make even very large specifications manageable and comprehensible. In fact, we intend to demonstrate here that statecharts counter many of the objections raised against conventional state diagrams, and thus appear to render specification by diagrams an attractive and plausible approach. Statecharts can be used either as a stand-alone behavioral description or as part of a more general design methodology that deals also with the system's other aspects, such as functional decomposition and data-flow specification. We also discuss some practical experience that was gained over the last three years in applying the statechart formalism to the specification of a particularly complex system.

1. Introduction

The literature on software and systems engineering is almost unanimous in recognizing the existence of a major problem in the specification and design of large and complex reactive systems. A reactive system (see [14]), in contrast with a *transformational system*, is characterized by being, to a large extent, event-driven, continuously having to react to external and internal stimuli. Examples include telephones, automobiles, communication networks, computer operating systems, missile and avionics systems, and the man-machine interface of many kinds of ordinary software. The problem is rooted in the difficulty of describing reactive behavior in ways that are clear and realistic, and at the same time formal and

* The initial part of this research was carried out while the author was consulting for the Research and Development Division of the Israel Aircraft Industries (IAI), Lod, Israel. Later stages were supported in part by grants from IAI and AD CAD, Ltd.

© 1987/\$3.50 © 1987, Elsevier Science Publishers B.V. (North-Holland)

ARTICLES

ON VISUAL FORMALISMS

The higraph, a general kind of diagramming object, forms a visual formalism of topological nature. Higraphs are suited for a wide array of applications to databases, knowledge representation, and, most notably, the behavioral specification of complex concurrent systems using the higraph-based language of statecharts.

DAVID HAREL

Visualizing information, especially information of complex and intricate nature, has for many years been the subject of considerable work by many people. The information that interests us here is nonquantitative, but rather, of a structural, set-theoretical, and relational nature. This should be contrasted with the kinds of quantitative information discussed at length in [43] and [46]. Consequently, we shall be interested in diagrammatic paradigms that are essentially topological in nature, not geometric, *terming them topvisual in the sequel*.

To represent a (single) set of elements S and some binary relation R on them, The precise meaning of the relation R is part of the application and has little to do with the mathematical properties of the graph itself. Certain restrictions on the relation R yield special classes of graphs that are of particular interest, such as ones that are connected, directed, acyclic, planar, or bipartite. There is no need to elaborate on the use of graphs in computer science—they are used extensively in virtually all branches of the field. The elements represented by the nodes in these applications range from the most concrete (e.g., physical gates in a circuit diagram) to the most abstract (e.g., complexity classes in a classification schema), and the edges have been used to represent almost any conceivable kind of relation, including ones of temporal, causal, functional, or epistemological nature. Obviously, graphs can be modified to support a number of different kinds of nodes and edges, representing different kinds of elements and relationships.

A somewhat less widely used extension of graphs is the formalism of *hypographs* (see, e.g., [1]), though these are also finding applications in computer science, mainly in database theory (see [14], [15], and [31]). A hypograph is a graph in which the relation being specified is not necessarily binary; in fact, it may not even be a fixed arity. Formally, an edge connects two sets of nodes, but rather, a subset thereof. This makes hypographs somewhat less amenable to visual representation, but various ways of overcoming this difficulty can be conceived (see Figure 1). In analogy with graphs, several special kinds of hypographs are of particular interest, such as directed or acyclic.

It is important to emphasize that the information

* Interestingly, both these topo-visual achievements of Euler were carried out during the period in which he could see with one eye only. (Euler lost sight in his right eye in 1735, and in the left around 1766.) It is tempting to attribute this in part to the fact that the lack of stereoscopic vision reduces one's ability to estimate size and distance, possibly causing a sharper awareness of topological features.

Part of this work was carried out while the author was at the Computer Science Department of Carnegie-Mellon University, Pittsburgh, Pennsylvania.

© 1988 ACM 0001-0782/88/0500-0514 \$1.50

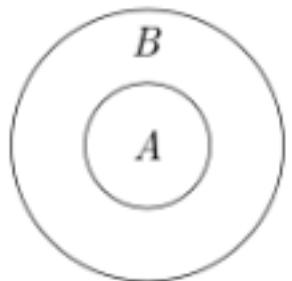
communications of the ACM

May 1988 Volume 31 Number 5

Higraphs

Euler Diagrams

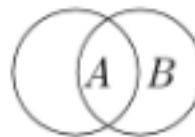
All A are B .



No A is B .

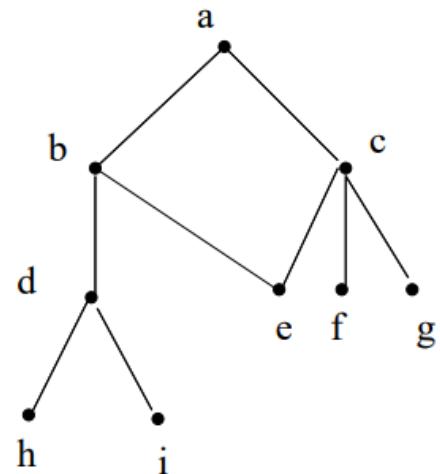


Some A is in B . Some A is not in B .

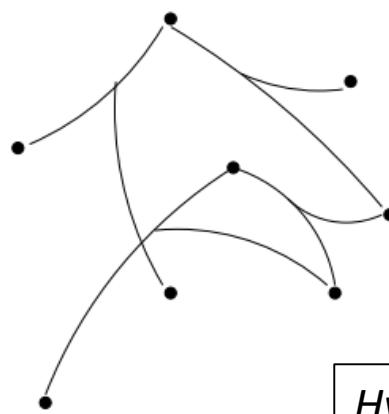


topological notions for set union, difference, intersection

Hypergraphs



a graph

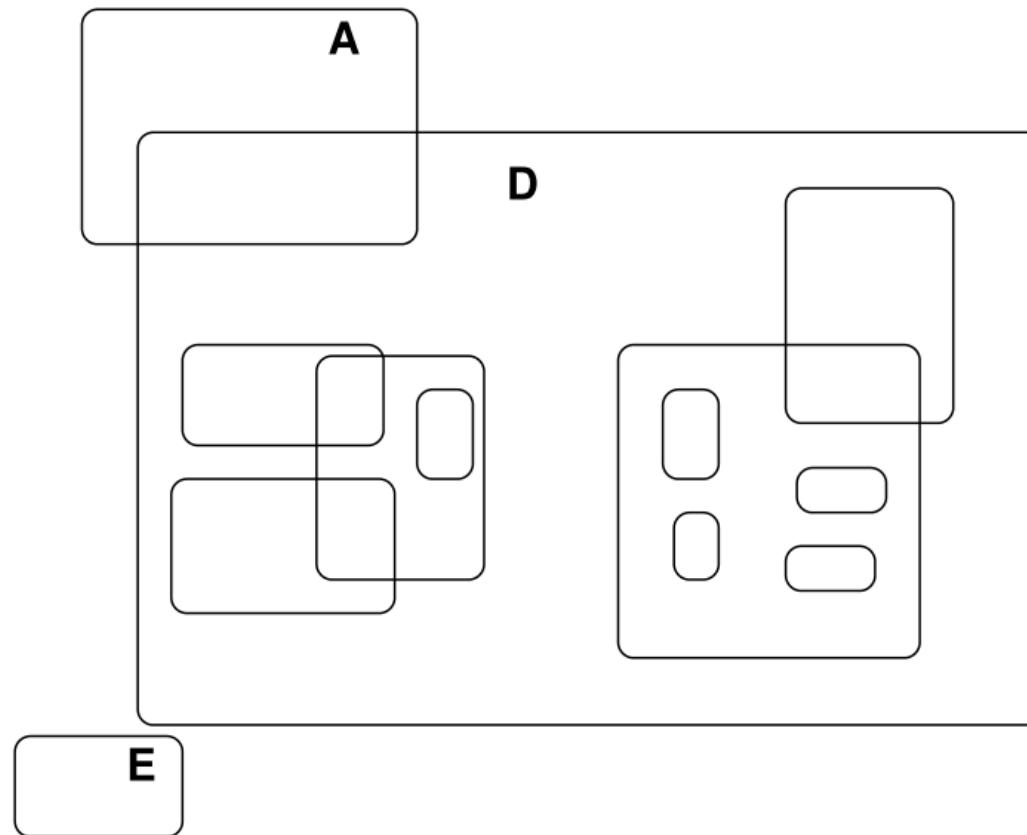


topological notion (syntax): connectedness

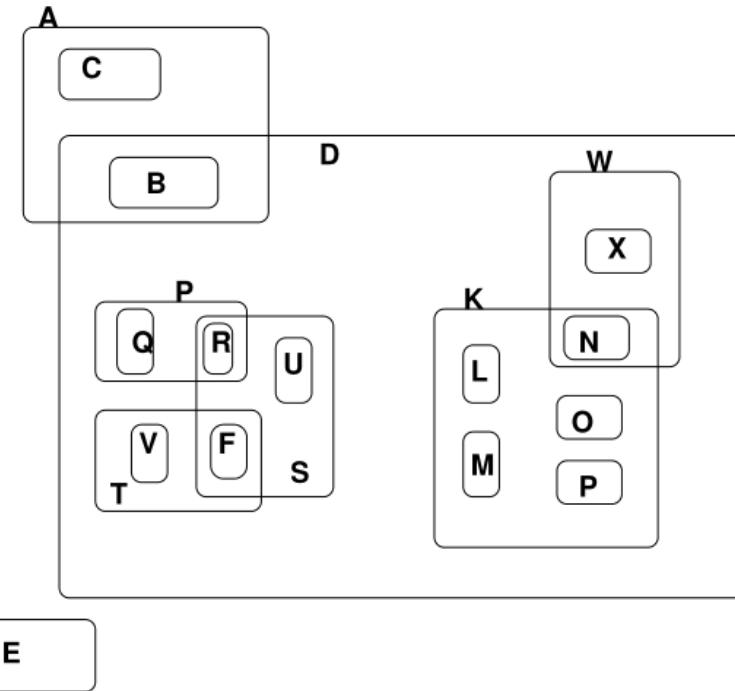
Hyperedges: $\subseteq 2^X$ (undirected), $\subseteq 2^X \times 2^X$ (directed).

a hypergraph

Blobs: set *inclusion*, not *membership*

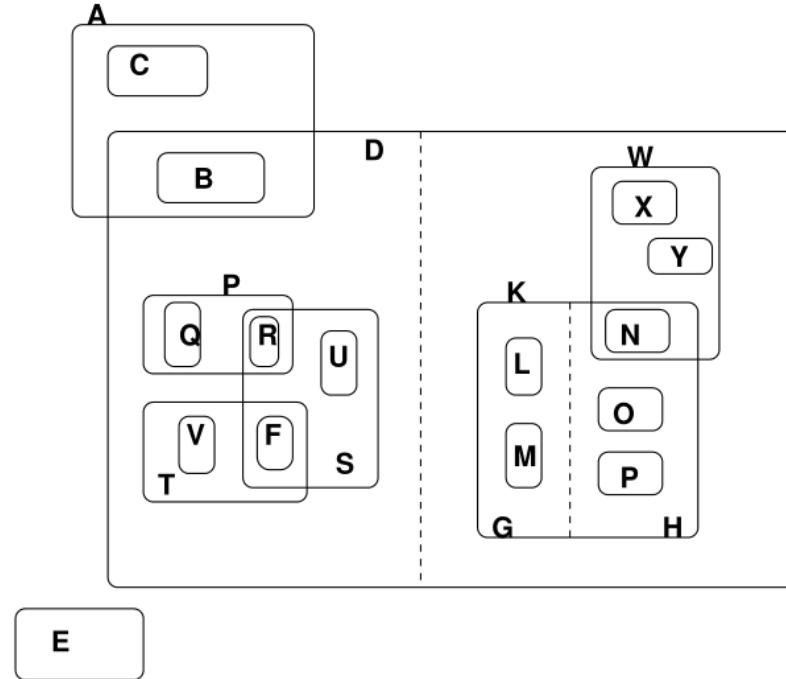


Unique Blobs (atomic sets, no *intersection*)



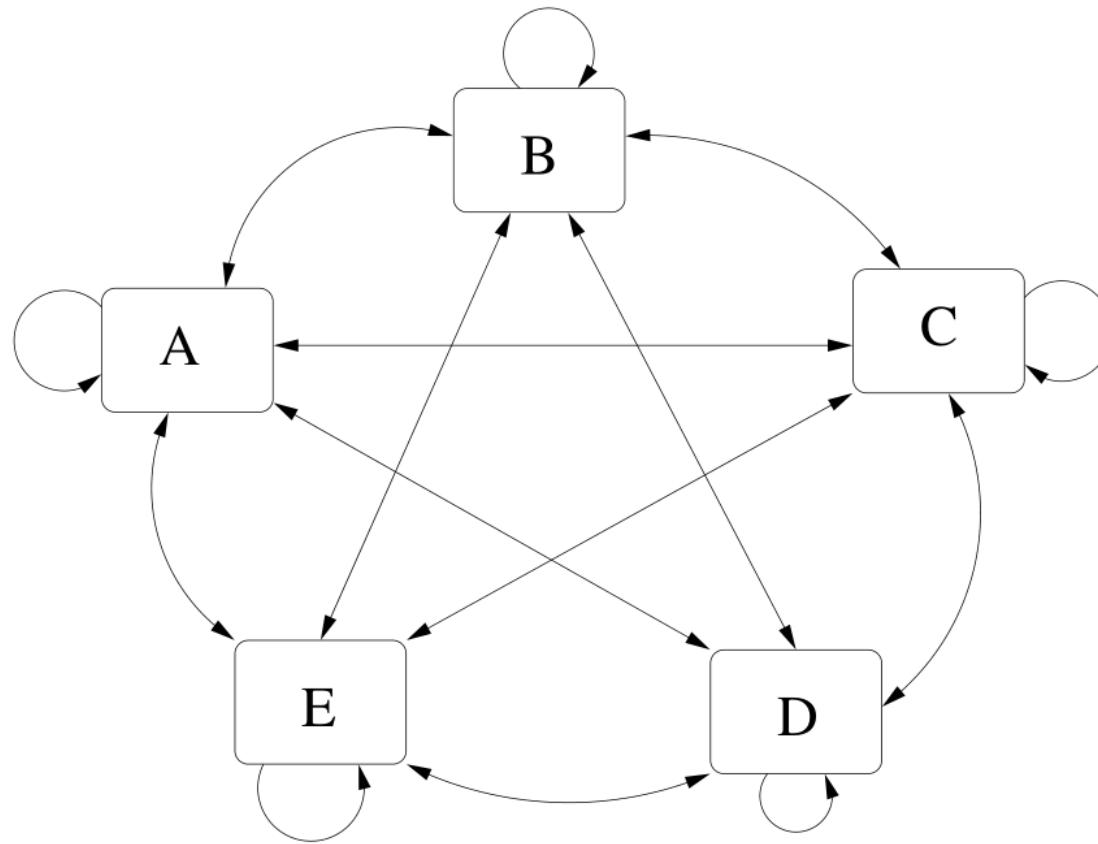
- Atomic blobs are identifiable sets
- Other blobs are union of enclosed set (e.g., $K = L \cup M \cup N \cup O \cup P$)

Unordered Cartesian Product: *Orthogonal* Components

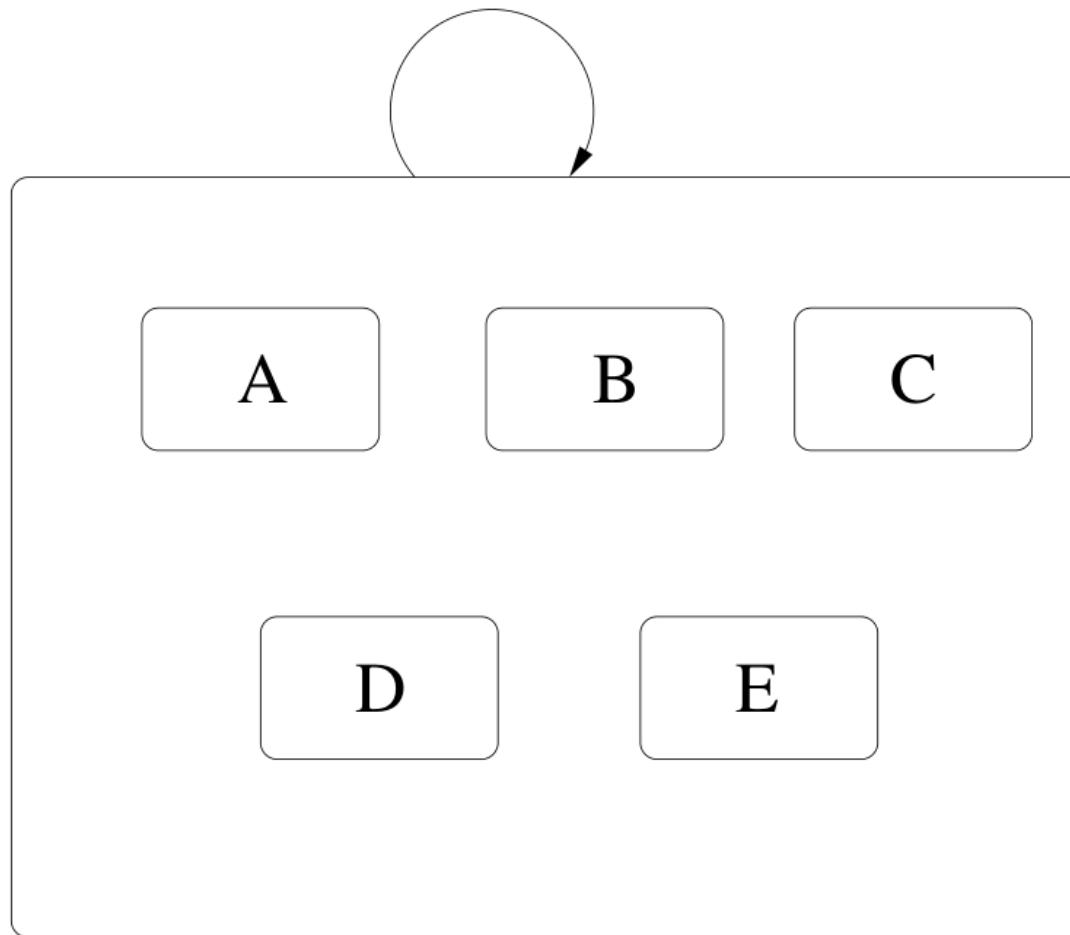


- $K = G \times H = (L \cup M) \times (N \cup O \cup P)$

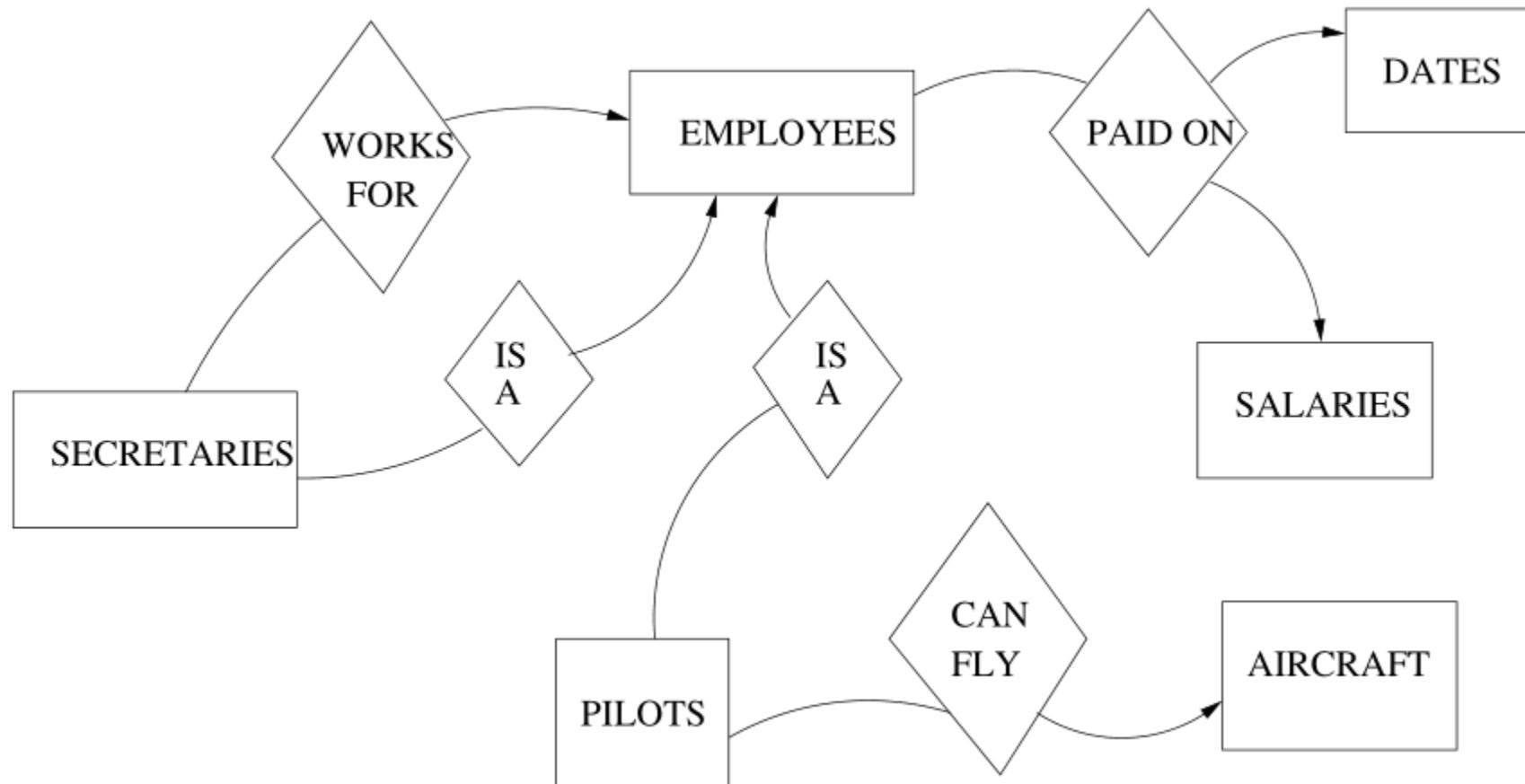
Clique Example



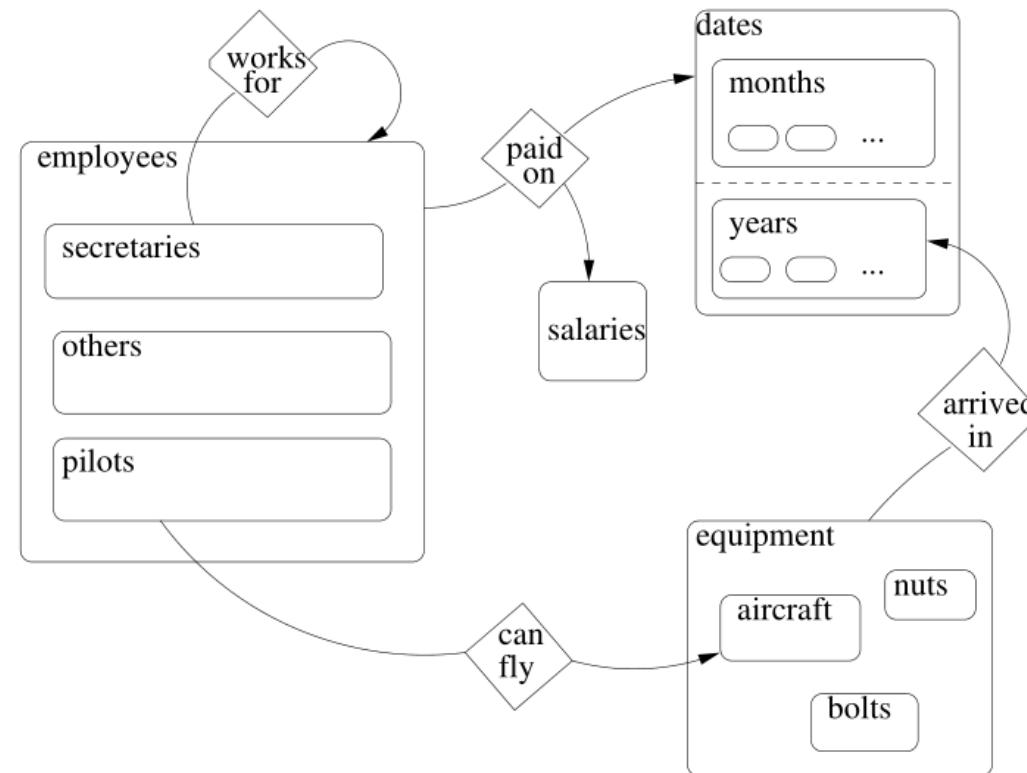
Clique: fully connected semantics



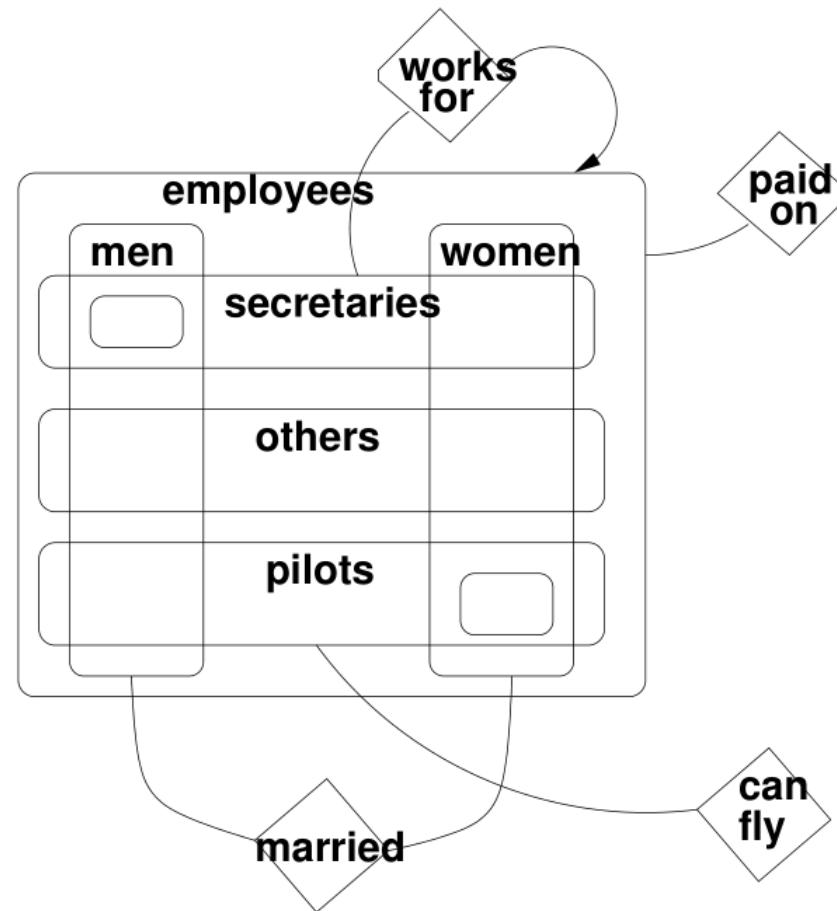
Entity Relationship Diagram



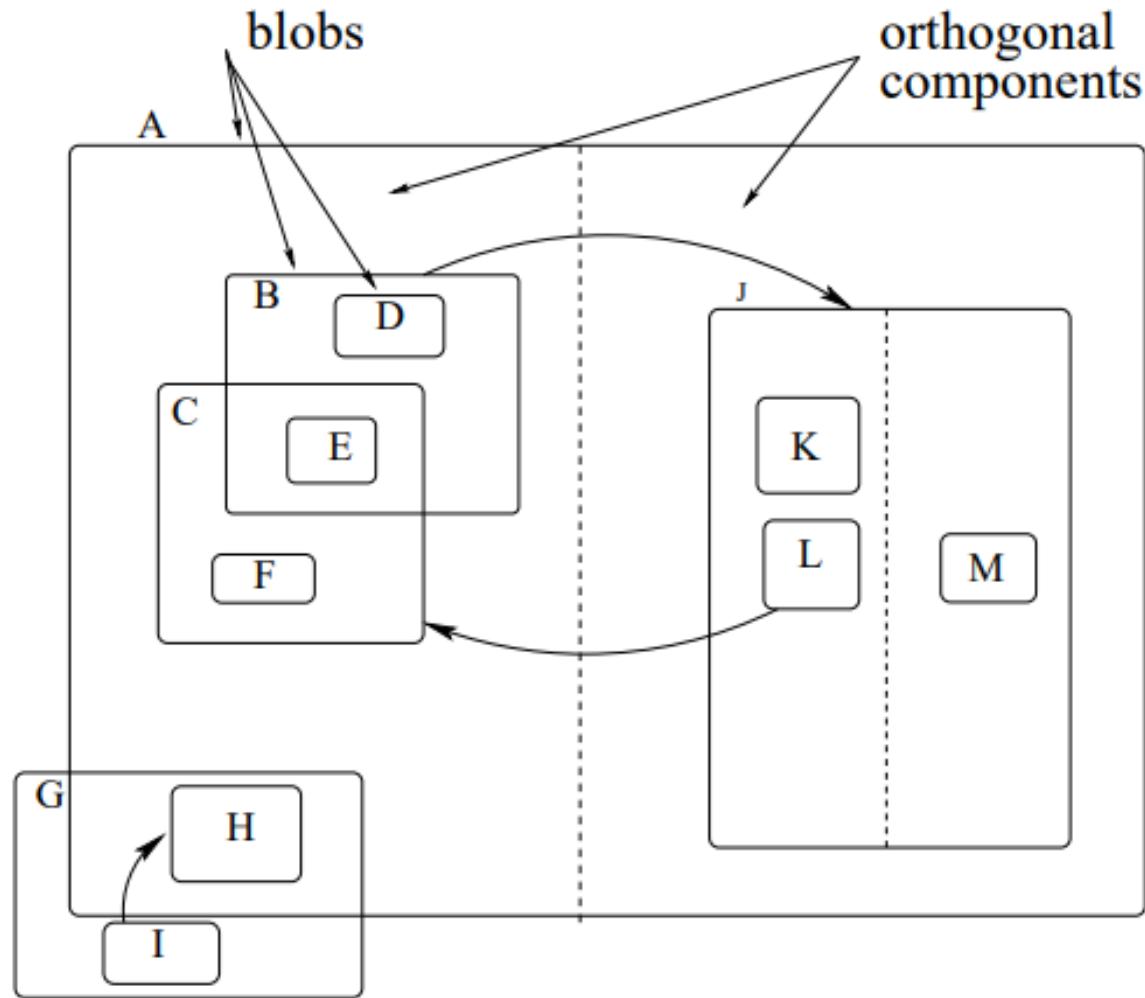
Higraph version of E-R diagram



Extending E-R Diagram



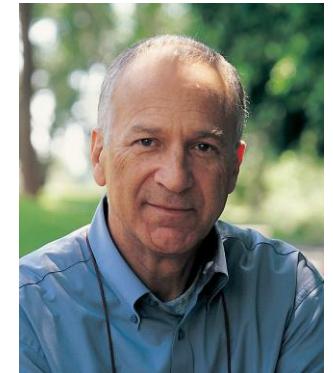
Simple Higraph



Statecharts

- Visual (topological, not geometric) formalism
- Precisely defined syntax and semantics
- Many uses:
 - Documentation (for human communication)
 - ~~Analysis (of behavioural properties)~~
 - Simulation
 - Code synthesis
 - ... and derived, such as testing, optimization, ...

Statecharts History



- Introduced by David Harel in 1987¹
- Notation based on higraphs = hypergraphs + Euler diagrams
- Semantics extend deterministic finite state automata with:
 - Depth (Hierarchy)
 - Orthogonality
 - Broadcast Communication
 - Time
 - History
 - Syntactic sugar, such as enter/exit actions

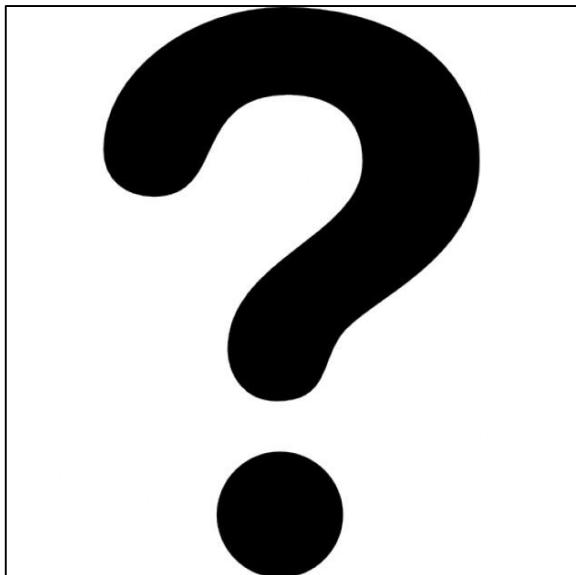
¹David Harel, Statecharts: a visual formalism for complex systems, Science of Computer Programming, Volume 8, Issue 3, 1987, Pages 231-274

Statecharts History: Website

Lectures
The theory exam will cover the highlighted papers/presentations below.
Blackboard scribbles [pdf].
Overview presentation [pdf]
Modelling and Simulation to Tackle Complexity presentation [pdf] exploring the causes of complexity. abstraction video
Formalisms: Use Cases, Sequence Diagrams, Regular Expressions and Finite State Automata presentation [pdf] discussing these formalisms in the context of checking the requirements of a system.
Formalisms: Causal Block Diagrams (CBDs) Analog computers and CSMP [pdf] CSMP: Robert D. Brennan: Digital simulation for control system design. DAC. New Orleans, Louisiana, USA, May 16-19, 1966. [pdf] (old) Blackboard Scribbles . Topological Sorting and Strong Component algorithms. Lecture on Algebraic and Discrete-Time CBDs [video]. Lecture on Continuous-Time CBDs [video]. Note: the above are not recordings of this year's class, but rather of an older version of the course, with the same content however. Lecture on (PID) controllers [pdf]
Formalisms: Petri Nets presentation[pdf] Christos G. Cassandras. Discrete Event Systems. Irwin, 1993. Chapters 4, 5. [pdf (MoSIS access only)]. Carl Adam Petri. Kommunikation mit Automaten . 1962. (this is Petri's doctoral dissertation). Tadao Murata. Petri nets: Properties, analysis and applications. Proceedings of the IEEE, 77(4):541-580, April 1989. James L. Peterson. Petri Net Theory and the Modeling of Systems. Prentice Hall, 1981.
Formalisms: Statecharts Higraphs presentation[pdf] . Statecharts presentation[pdf] . David Harel. Statecharts: A Visual Formalism for Complex Systems. Science of Computer Programming. Volume 8. 1987. pp. 231 - 274. [pdf]. David Harel. On Visual Formalisms. Communications of the ACM. Volume 31, No. 5. 1988. pp. 514 - 530. [pdf] [pdf (MoSIS access only)]. David Harel and Amnon Naamad, The STATEMATE semantics of statecharts. ACM Transactions on Software Engineering and Methodology (TOSEM) Volume 5 , Issue 4 (October 1996) pp.293 - 333. [pdf] [pdf (MoSIS access only)]. D. Harel and M. Politi. Modeling Reactive Systems with Statecharts: The STATEMATE Approach. McGraw-Hill, 1998. (available online). David Harel and Hillel Kugler. The Rhapsody Semantics of Statecharts (or, On the Executable Core of the UML). Springer, Lecture Notes in Computer Science 3147. 2004. pp. 325 - 354. [pdf] Michael von der Beeck. A structured operational semantics for UML-statecharts. Software and Systems Modeling. Volume 1, No. 2 pp.130 - 141. December 2002. [pdf]. The digital watch assignment (not an assignment this year).

Running Example

Controller



system
input

system
output

(Physical) Plant

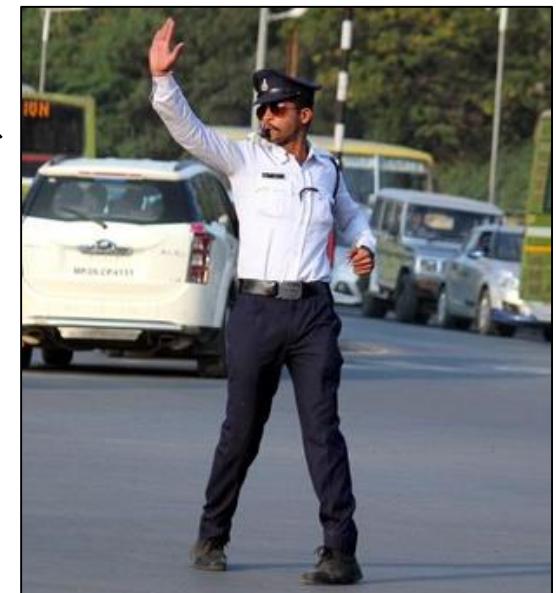


<<observe>>

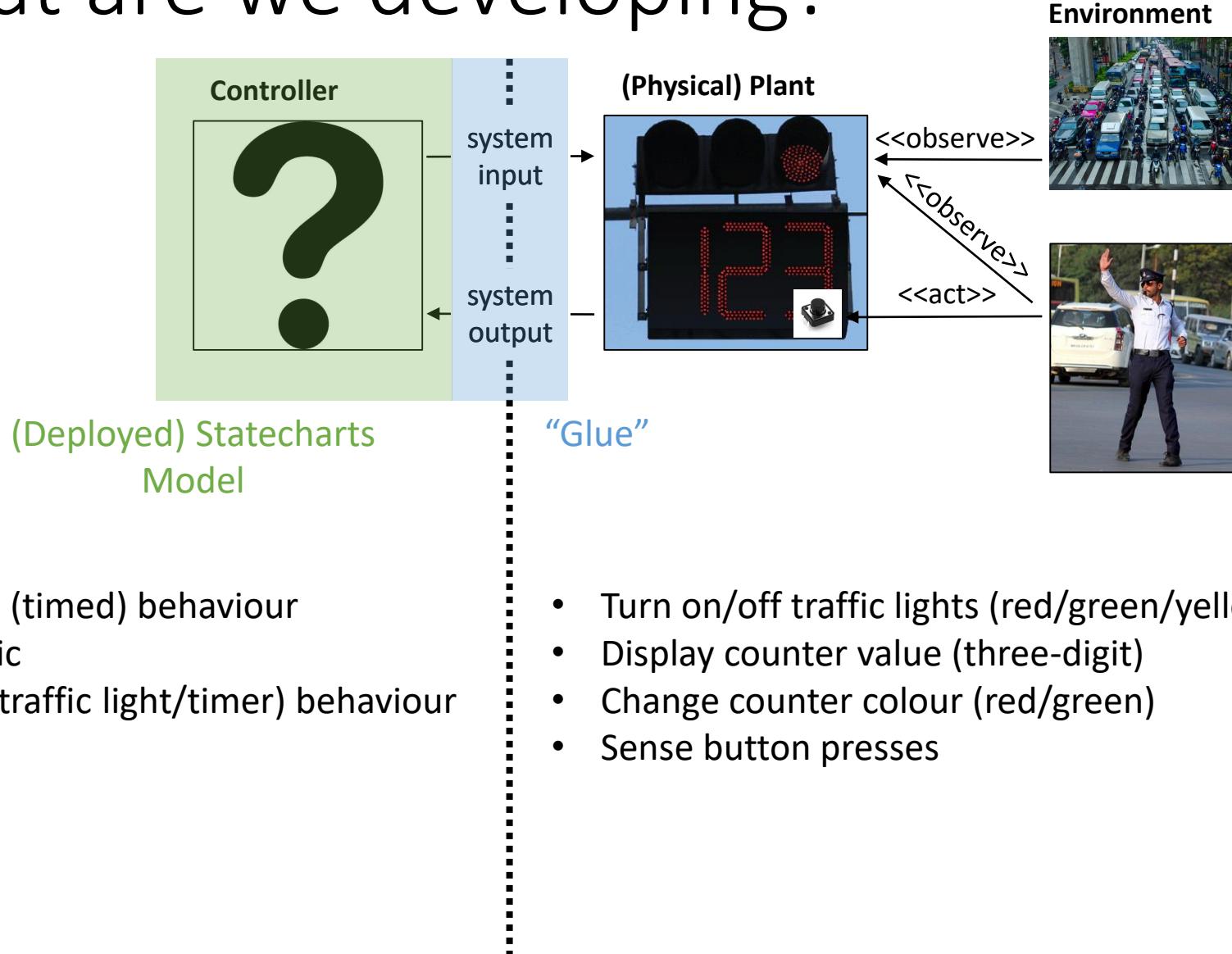
<<observe>>

<<act>>

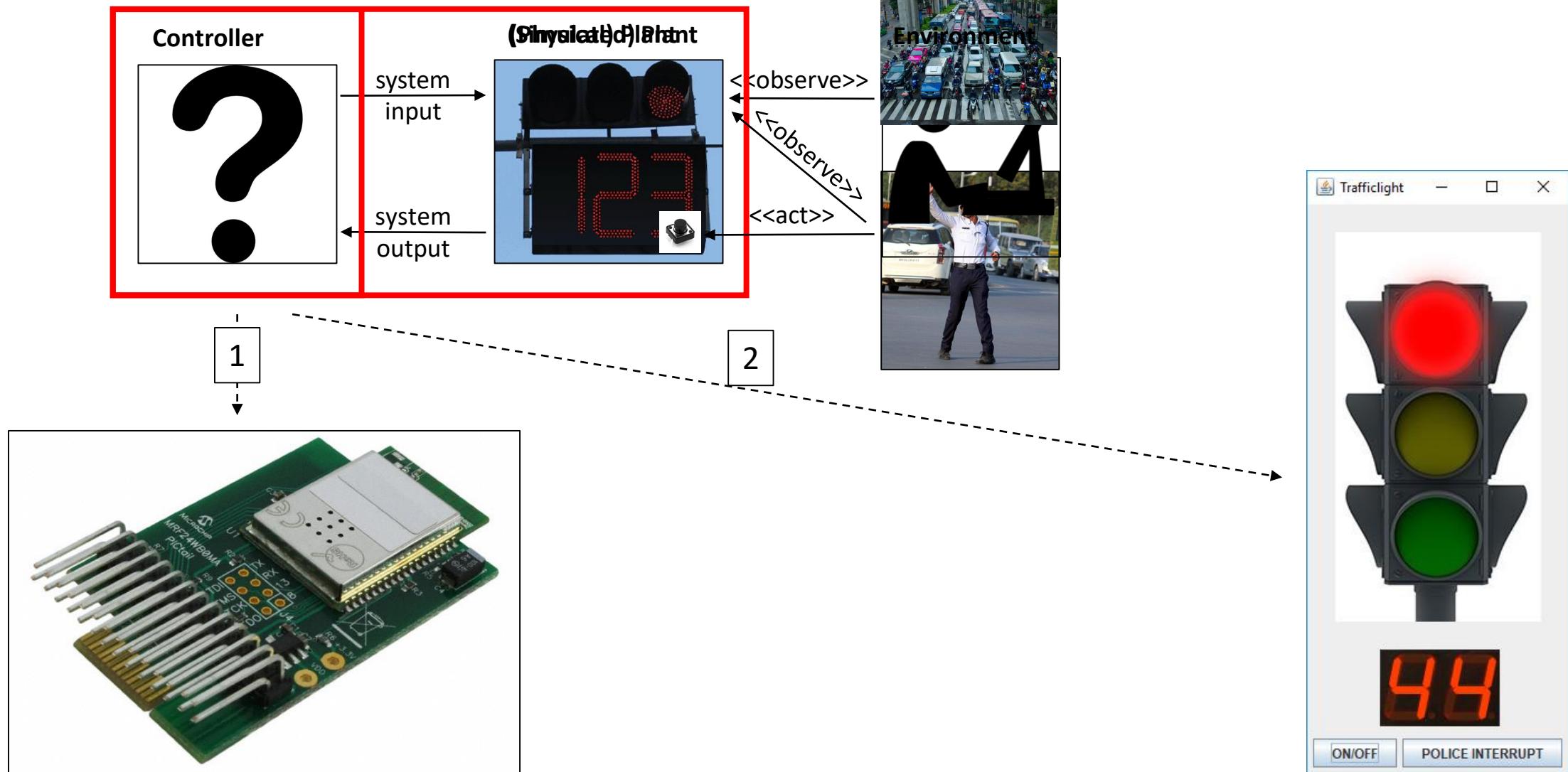
Environment



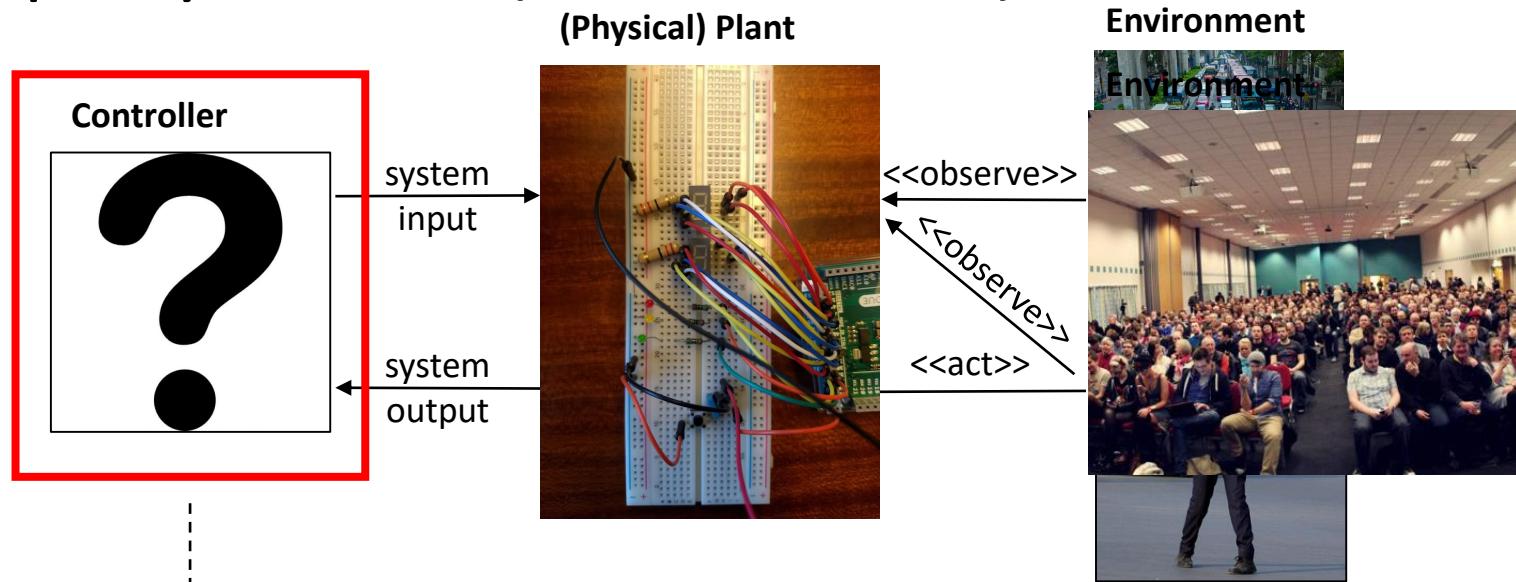
What are we developing?



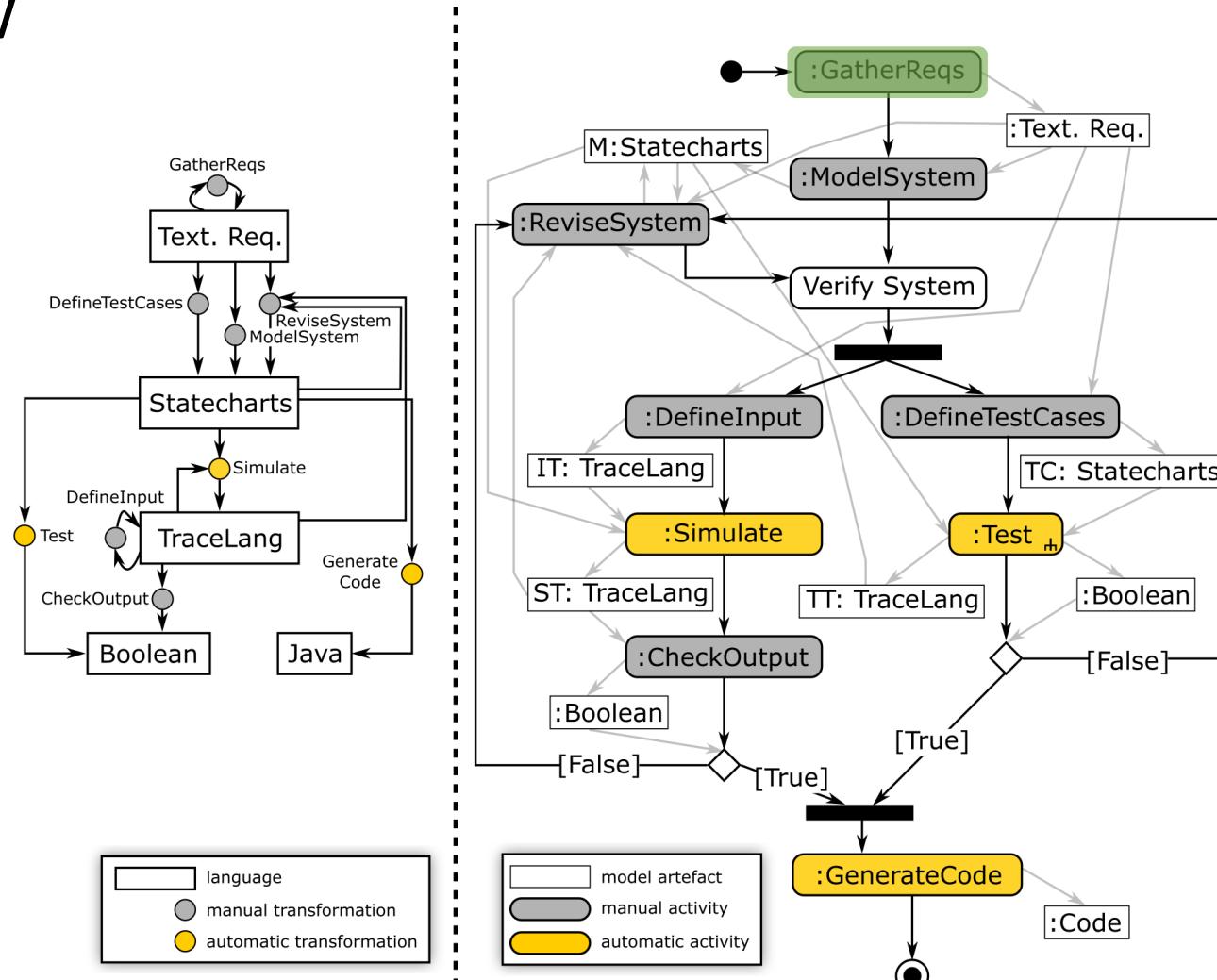
Deployment (Simulation)



Deployment (Hardware)



Workflow

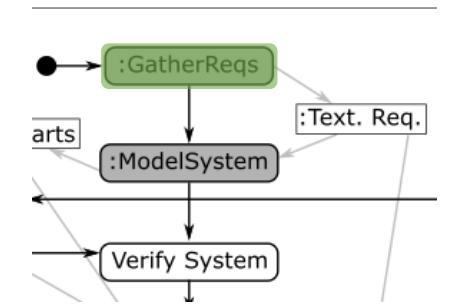


³ Hans Vangheluwe and Ghislain C. Vansteenkiste. A multi-paradigm modeling and simulation methodology: Formalisms and languages. In European Simulation Symposium (ESS), pages 168-172. Society for Computer Simulation International (SCS), October 1996. Genoa, Italy.

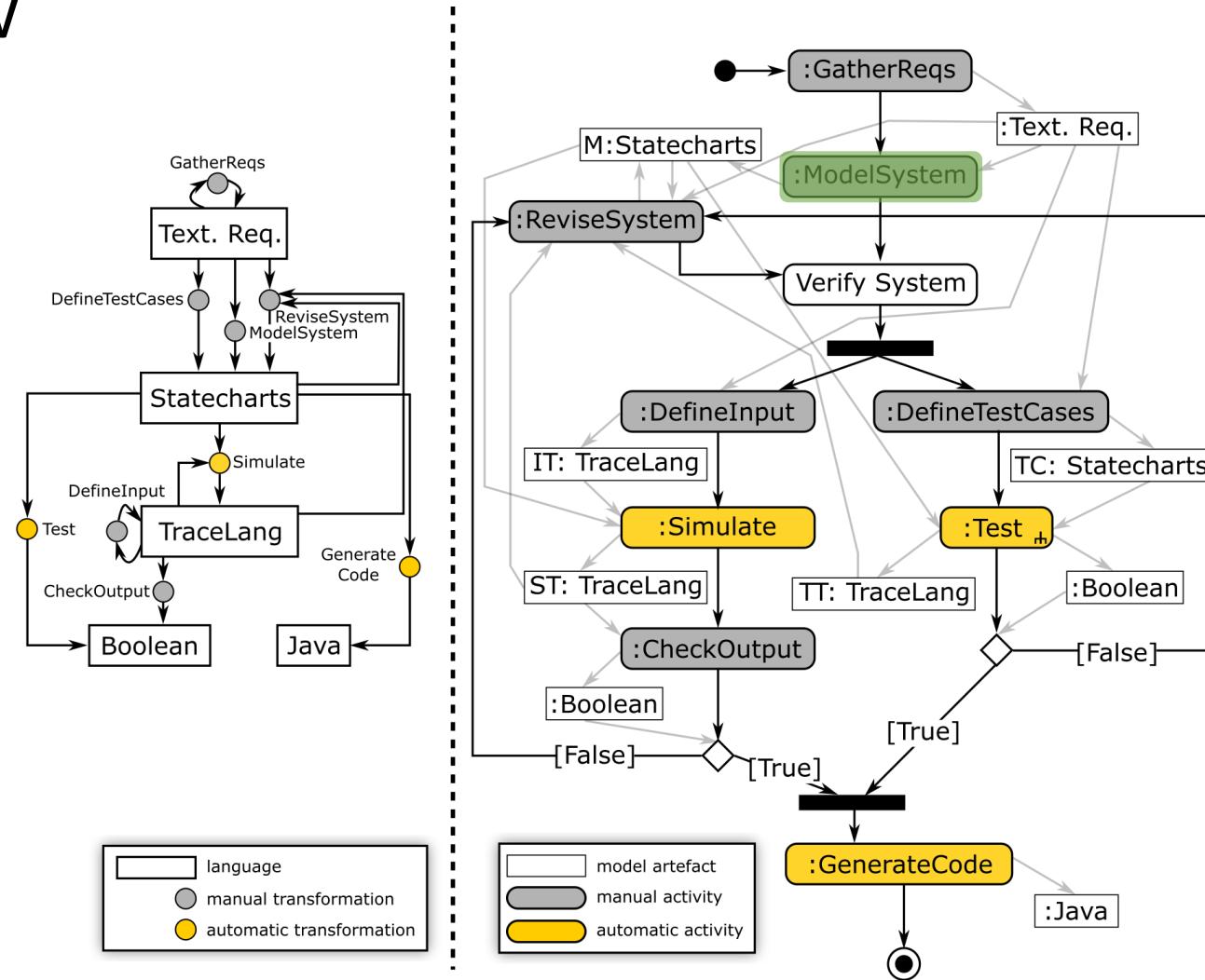
⁴ FTG+PM: An Integrated Framework for Investigating Model Transformation Chains, Levi Lúcio, Sadaf Mustafiz, Joachim Denil, Hans Vangheluwe, Maris Jukss. System Design Languages Forum (SDL) 2013, Montreal, Quebec. LNCS Volume 7916, pp 182-202, 2013.

Requirements

- R1: three differently coloured lights: red, green, yellow
- R2: at most one light is on at any point in time
- R3: at system start-up, the red light is on
- R4: cycles through red on, green on, and yellow on
- R5: red is on for 60s, green is on for 55s, yellow is on for 5s
- R6: police can interrupt autonomous operation
 - Result = blinking yellow light (on -> 1s, off -> 1s)
- R7: police can resume an interrupted traffic light
 - Result = light which was on at time of interrupt is turned on again
- R8: a timer displays the remaining time while the light is red or green; this timer decreases and displays its value every second. The colour of the timer reflects the colour of the traffic light.



Workflow



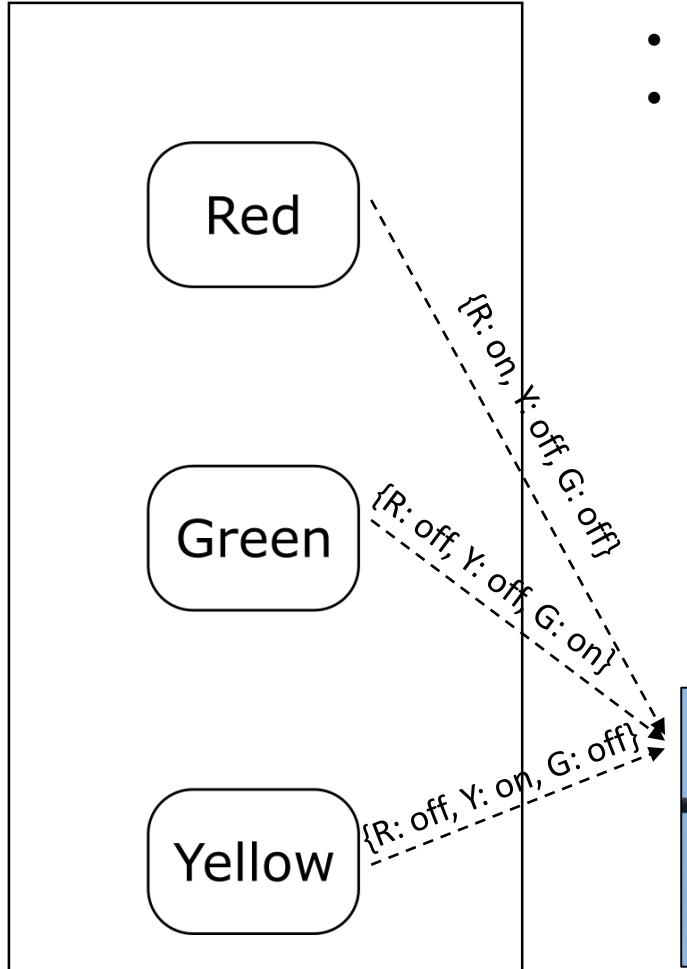
INTRODUCTION

STATECHARTS BASICS

YAKINDU IN DEPTH

ADVANCED CONCEPTS

States



- R1: three differently coloured lights: red (R), green (G), yellow (Y)
- R2: at most one light is on at any point in time

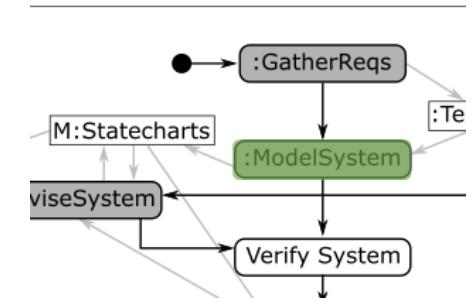
(Simulated) Plant



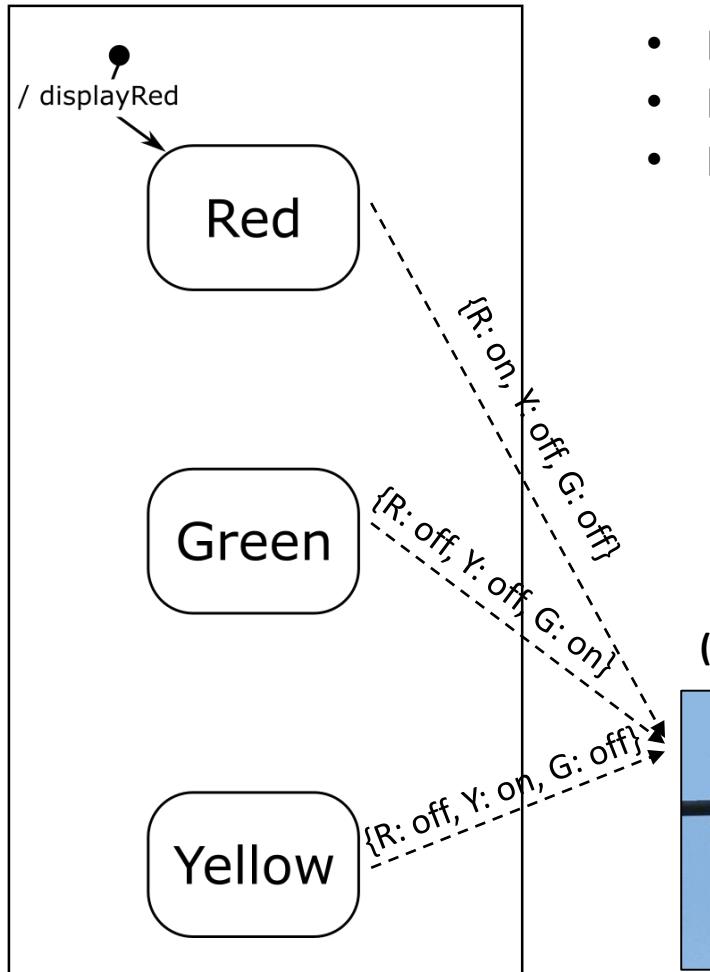
Environment



<<observe>>



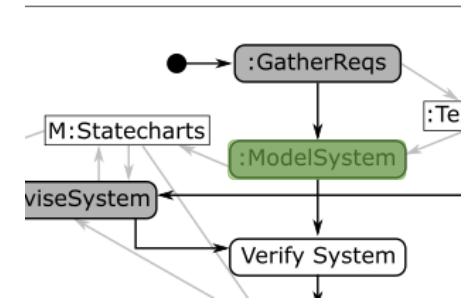
Default State



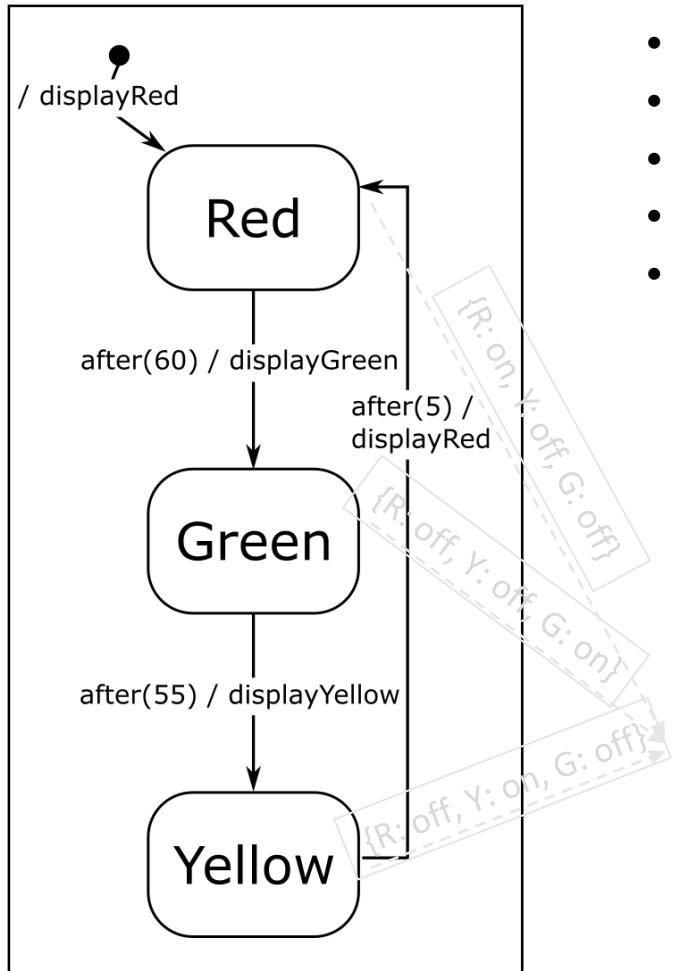
- R1: three differently coloured lights: red (R), green (G), yellow (Y)
- R2: at most one light is on at any point in time
- R3: at system start-up, the red light is on



<<observe>>

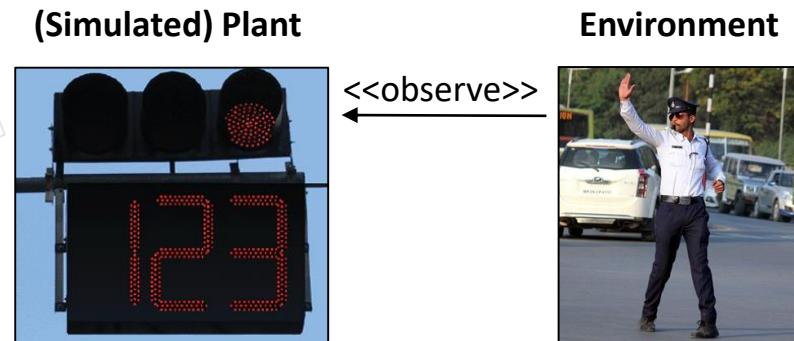


Transitions



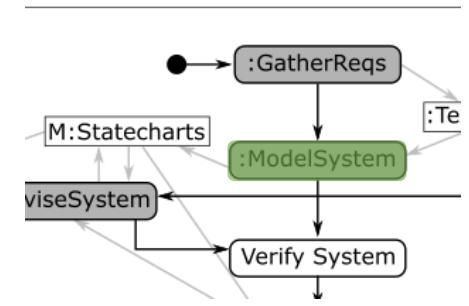
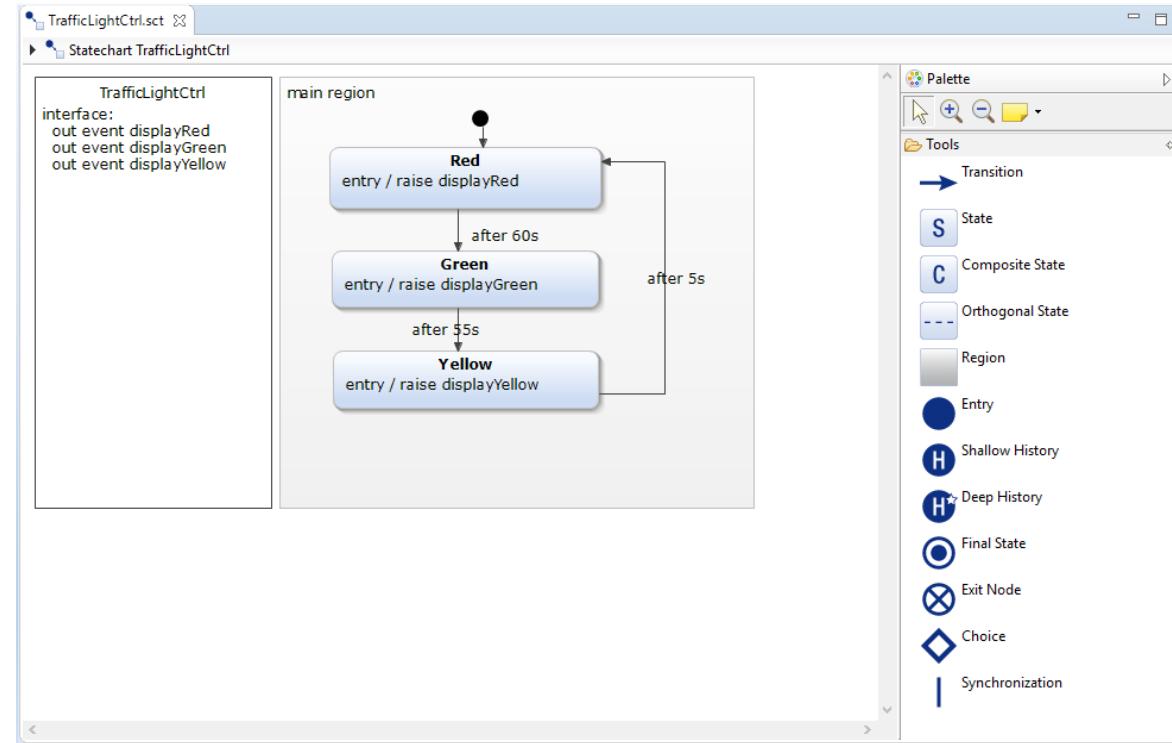
- R1: three differently coloured lights: red (R), green (G), yellow (Y)
- R2: at most one light is on at any point in time
- R3: at system start-up, the red light is on
- R4: cycles through red on, green on, and yellow on
- R5: red is on for 60s, green is on for 55s, yellow is on for 5s

event(params) [guard] / output_action(params)



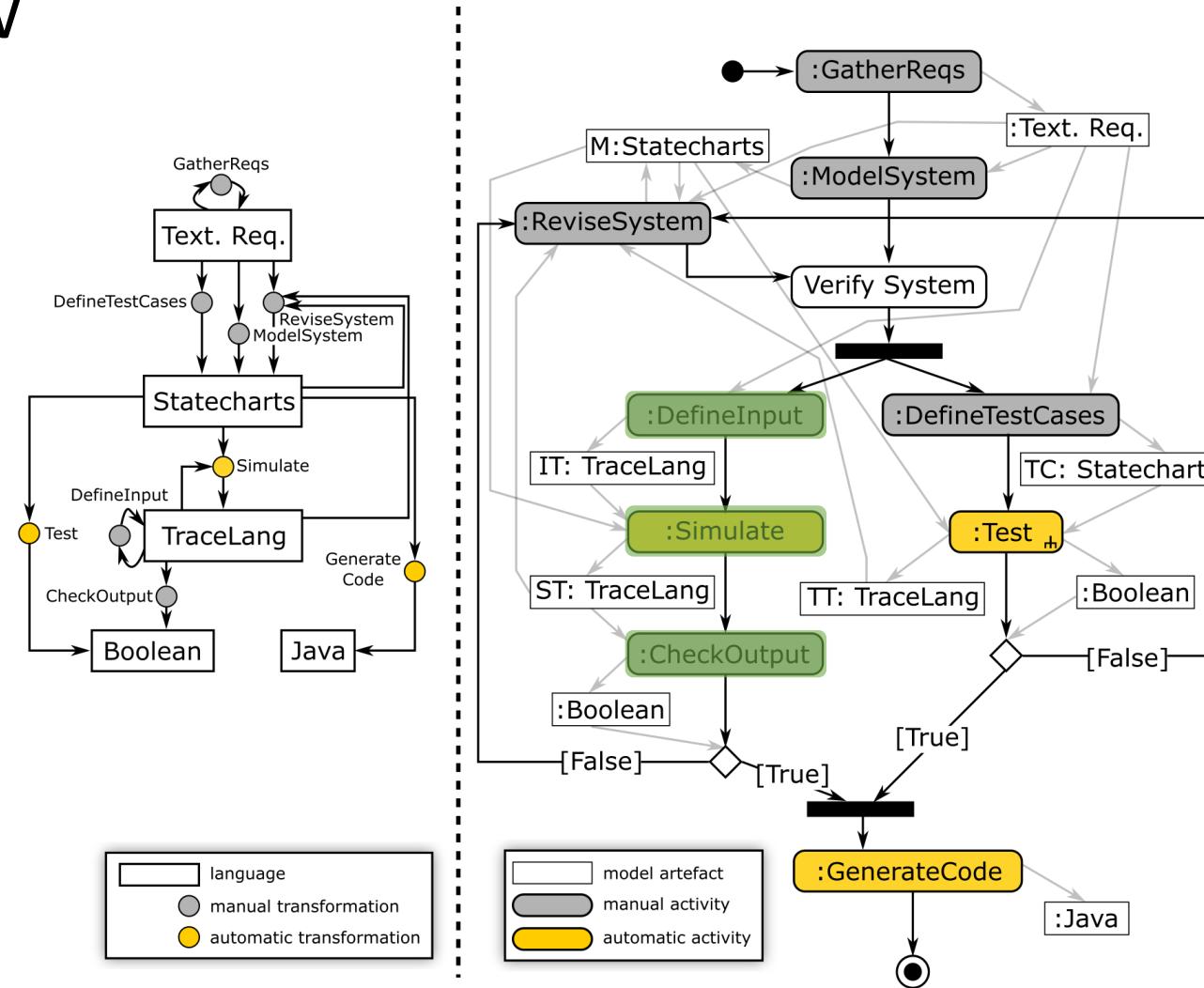
Yakindu⁵: Modelling

(introducing syntactic sugar: **enter actions**)

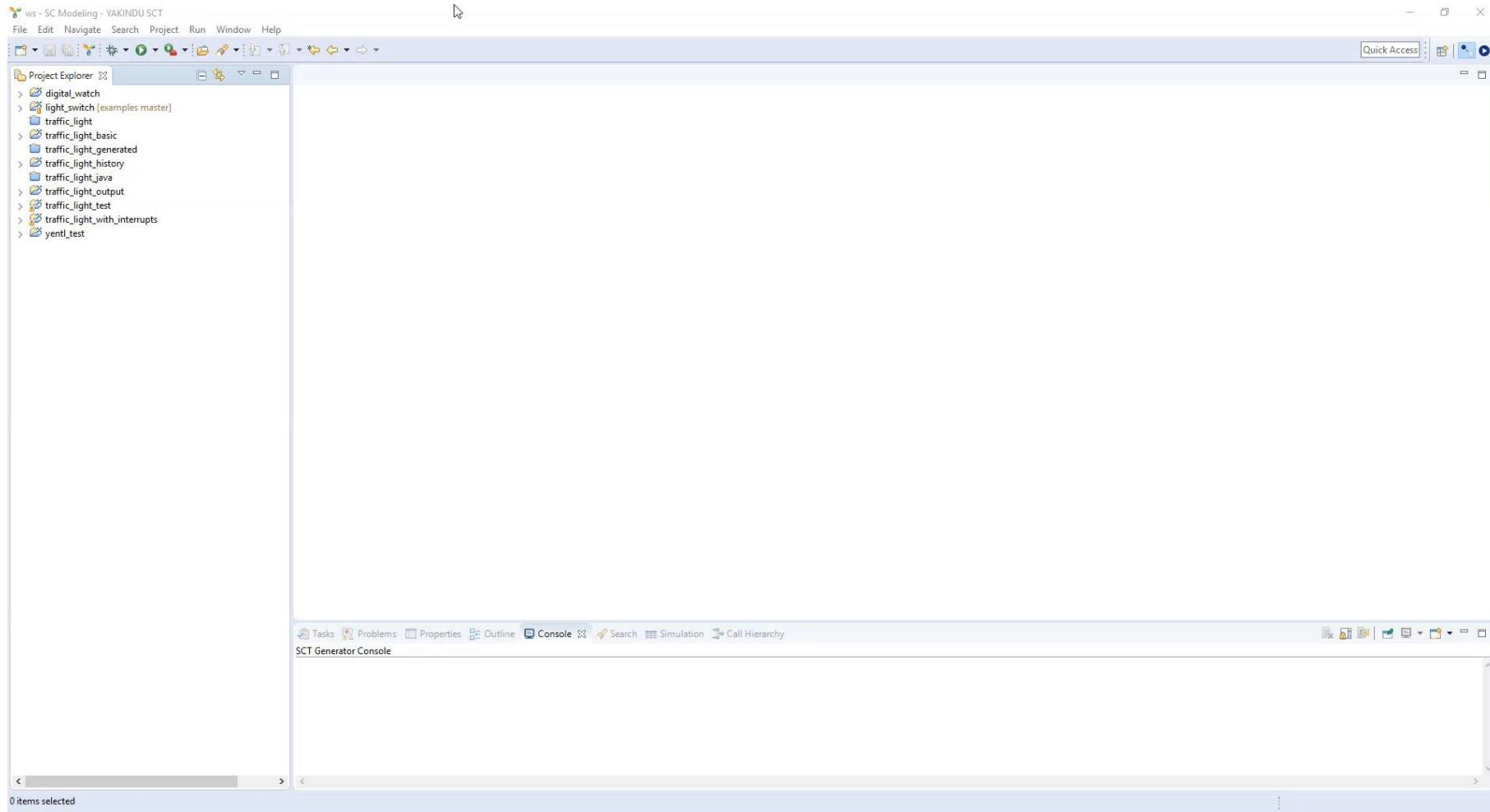
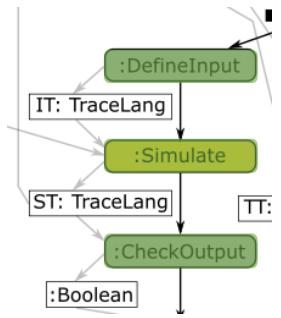


⁵ <https://www.itemis.com/en/yakindu/state-machine/>

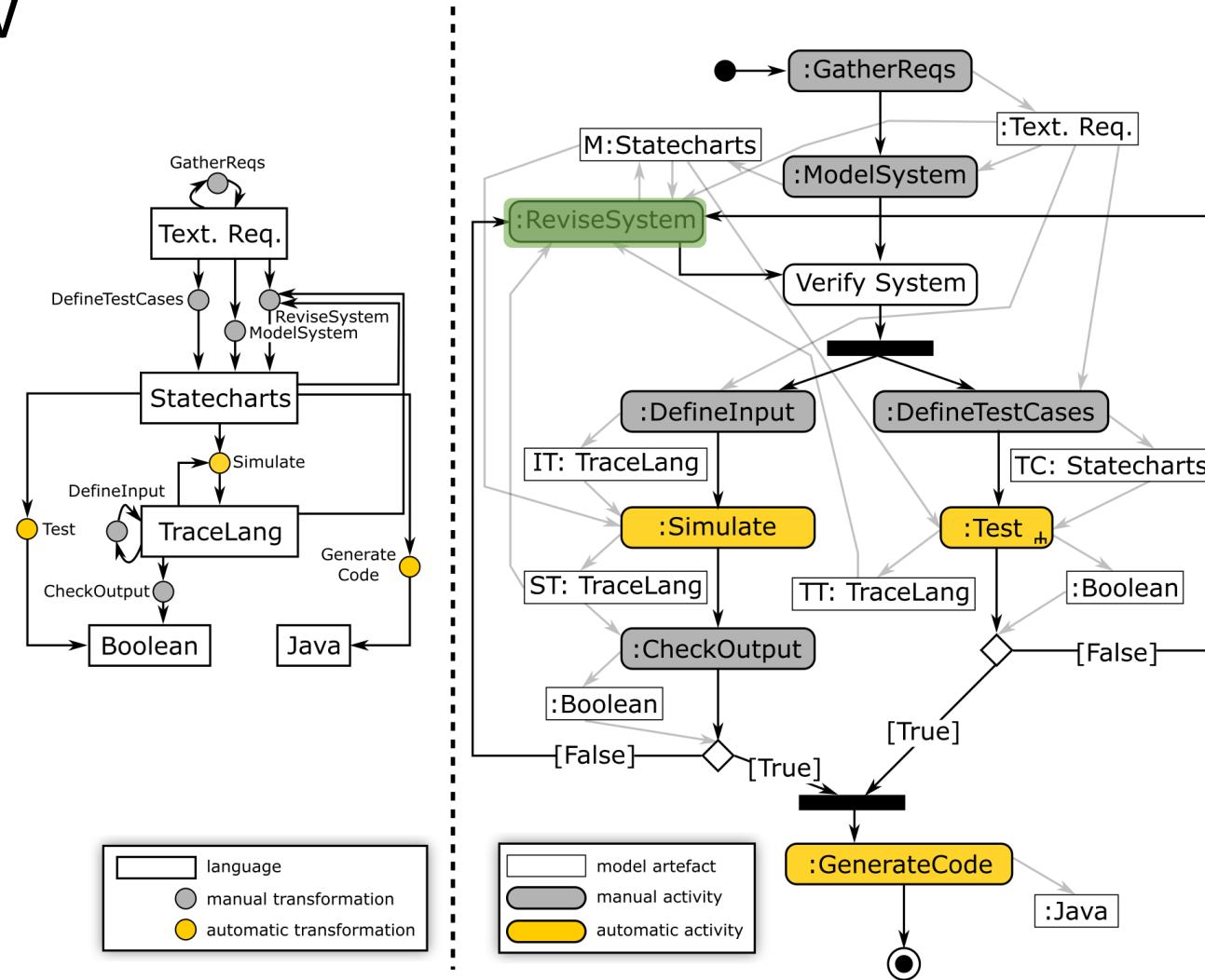
Workflow



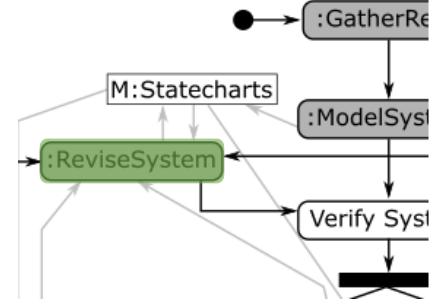
Yakindu: Simulation (Scaled Real-Time)



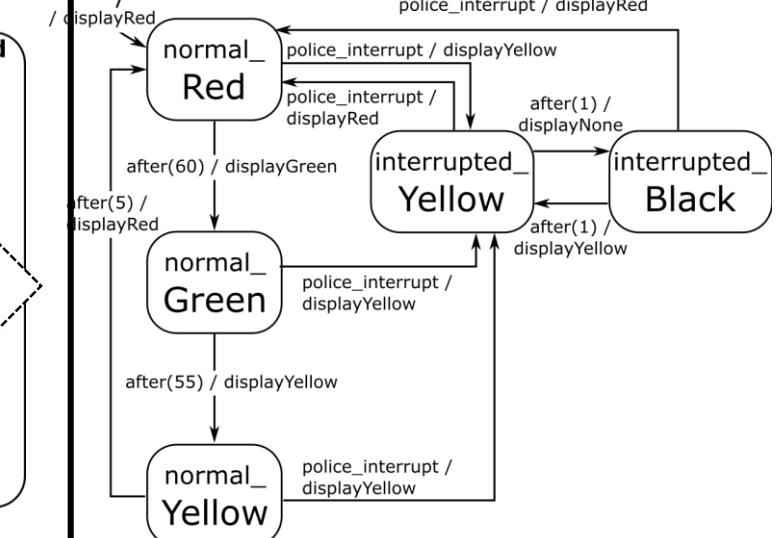
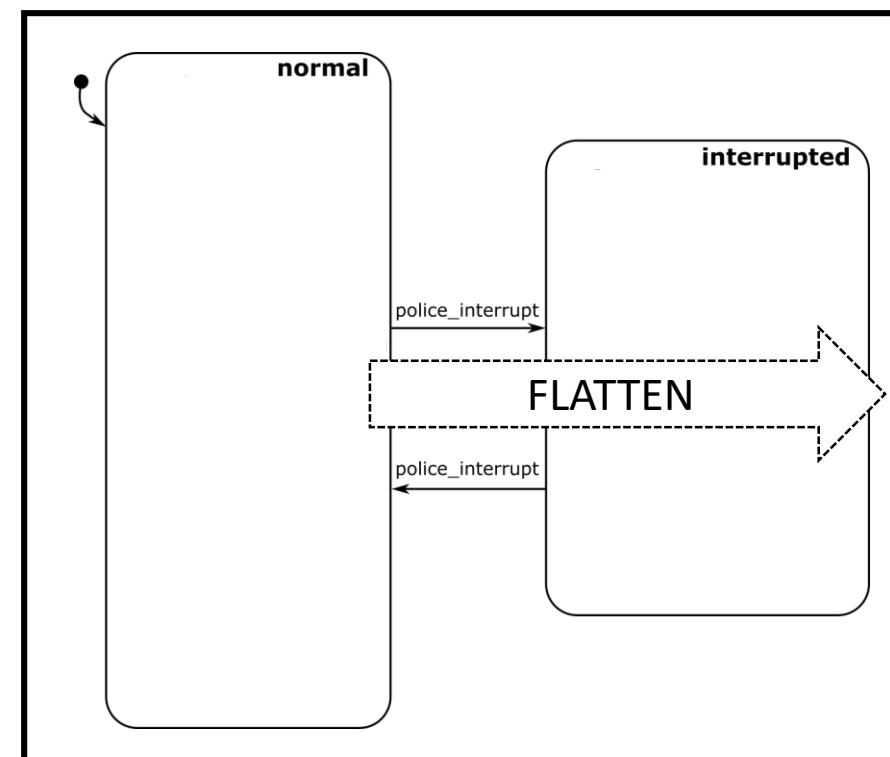
Workflow



Hierarchy

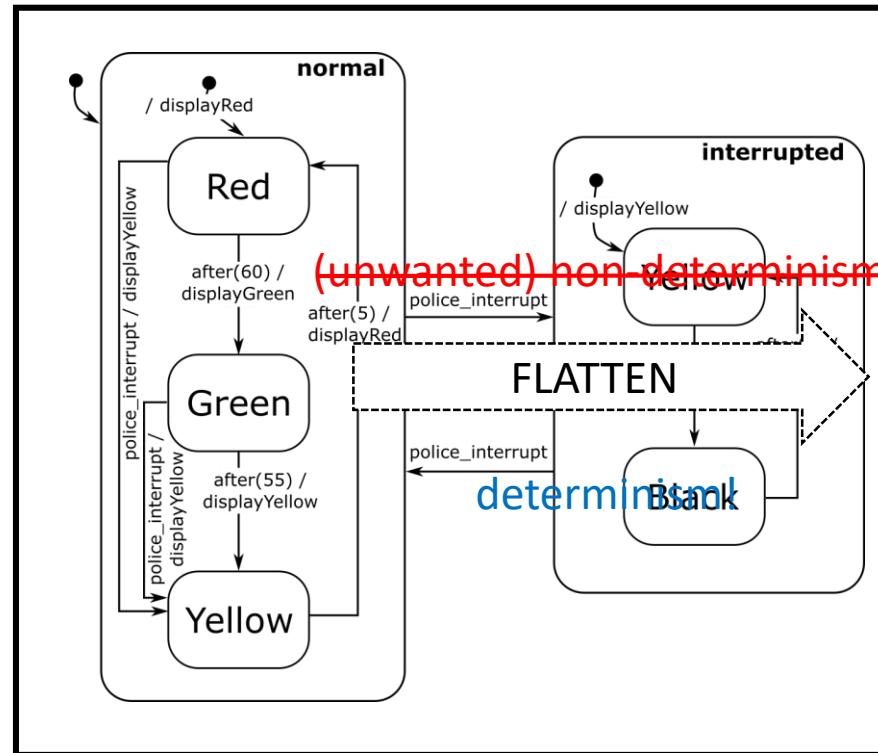


- R6: police can interrupt autonomous operation
 - Result = blinking yellow light (on -> 1s, off -> 1s)
- R7: police can resume an interrupted traffic light



Hierarchy: Modified Example

Statemate, Yakindu, ...

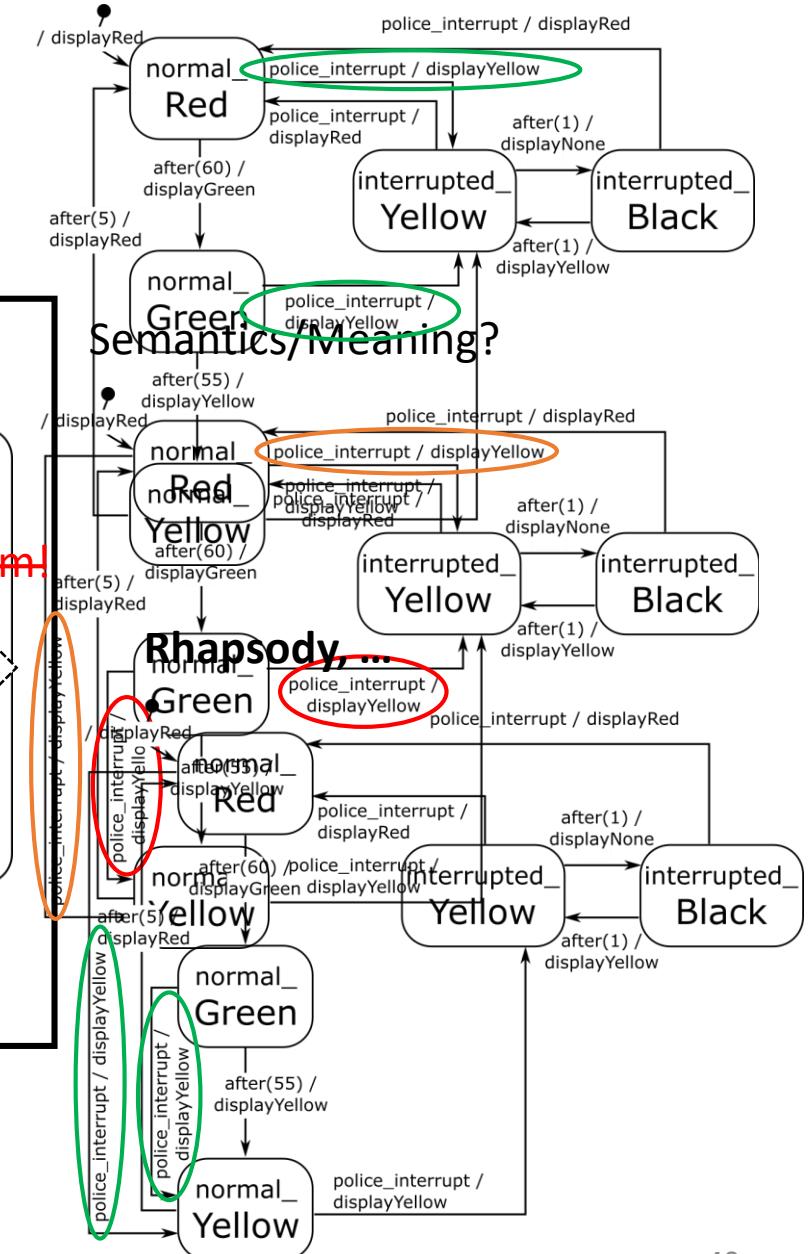


~~(unwanted) non-determinism~~

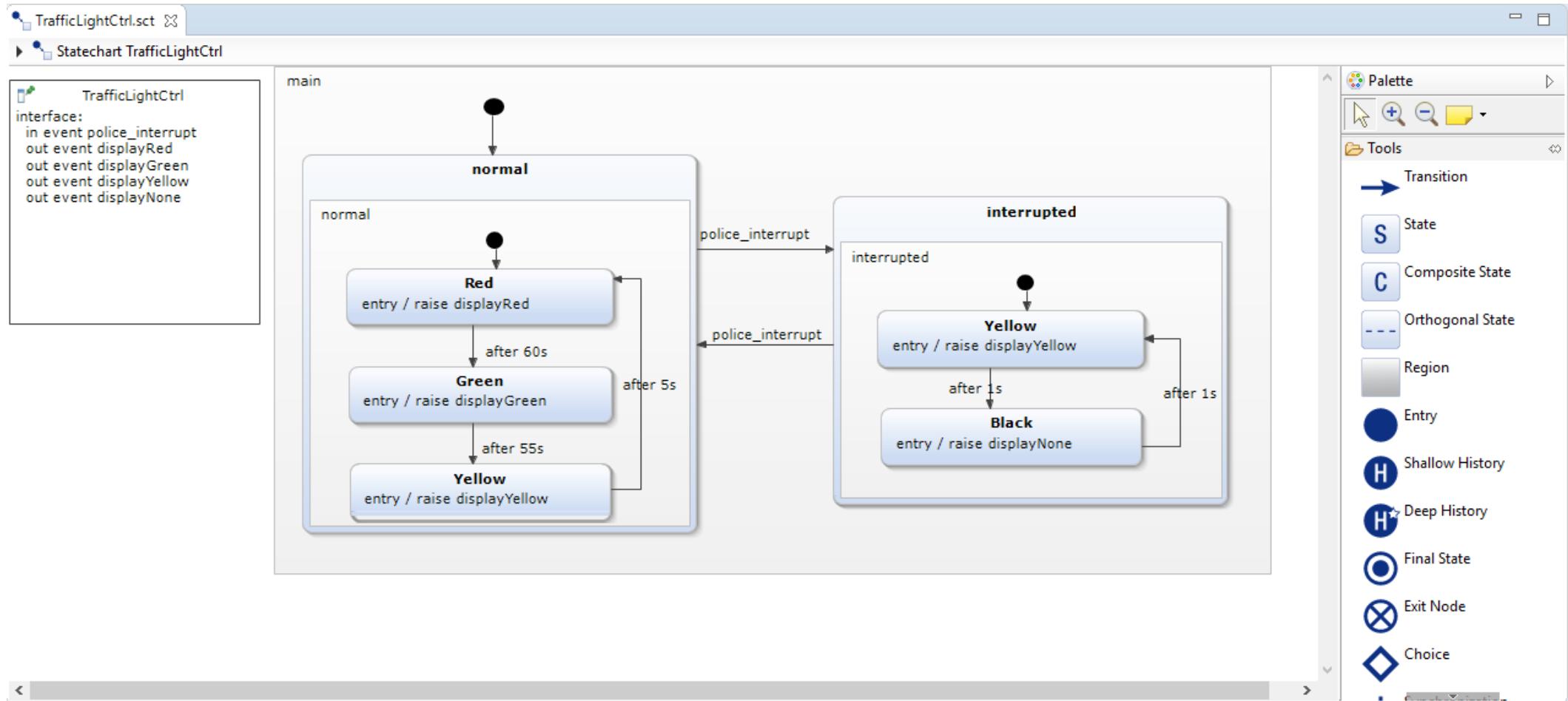
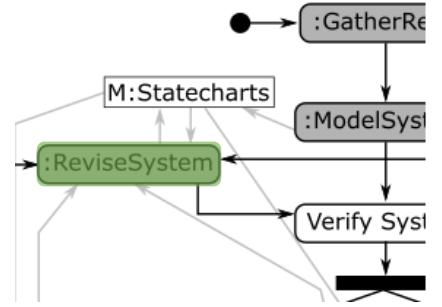
FLATTEN

determinism

Semantics/Meaning?

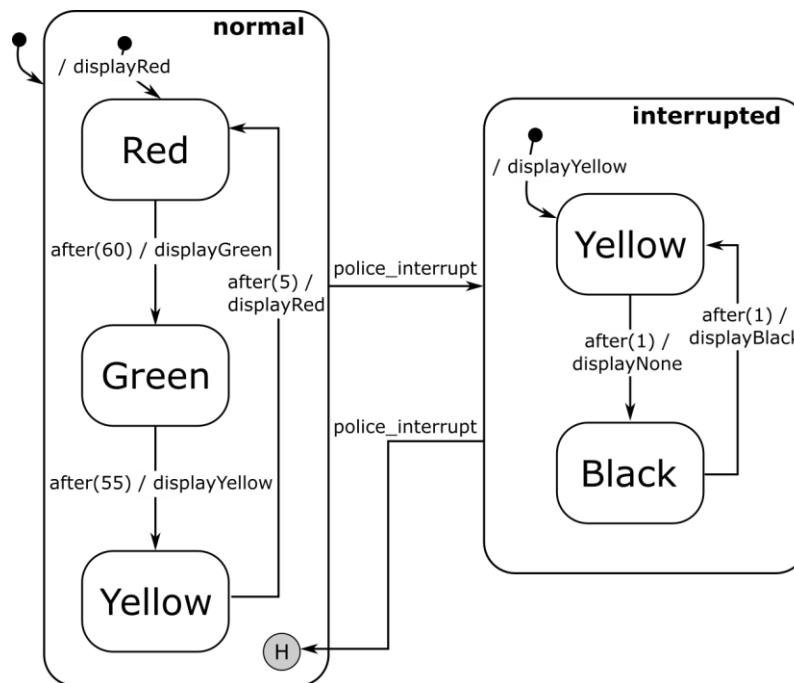
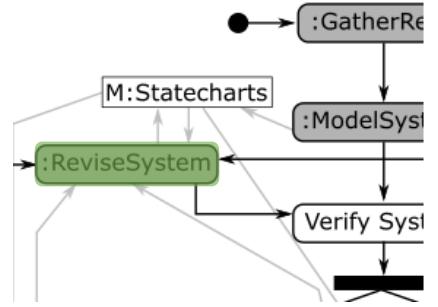


Yakindu: Hierarchy



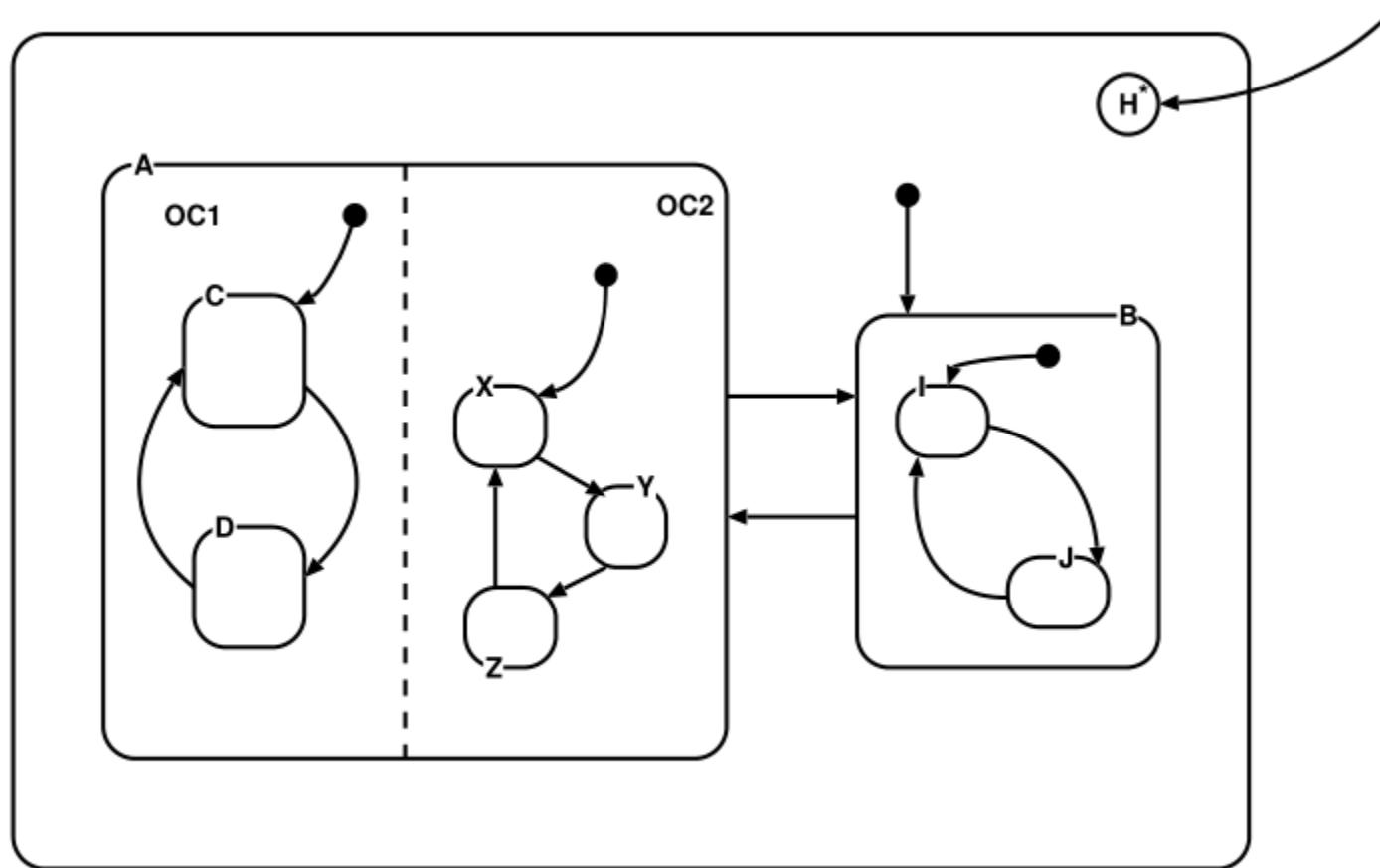
History

- R7: police can resume an interrupted traffic light
 - Result = light which was on at time of interrupt is turned on again

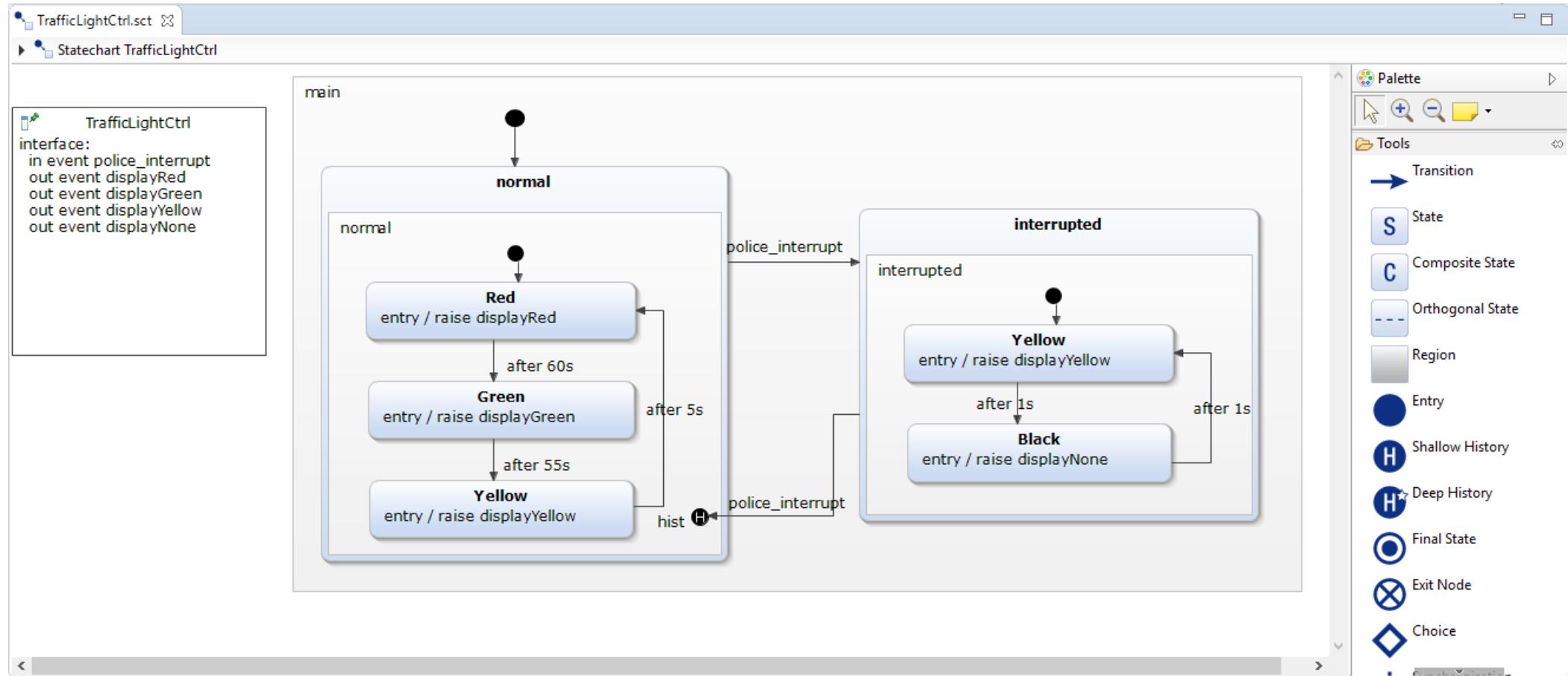
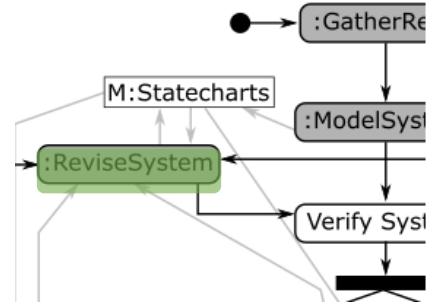


H shallow history
 H^* deep history

Deep History



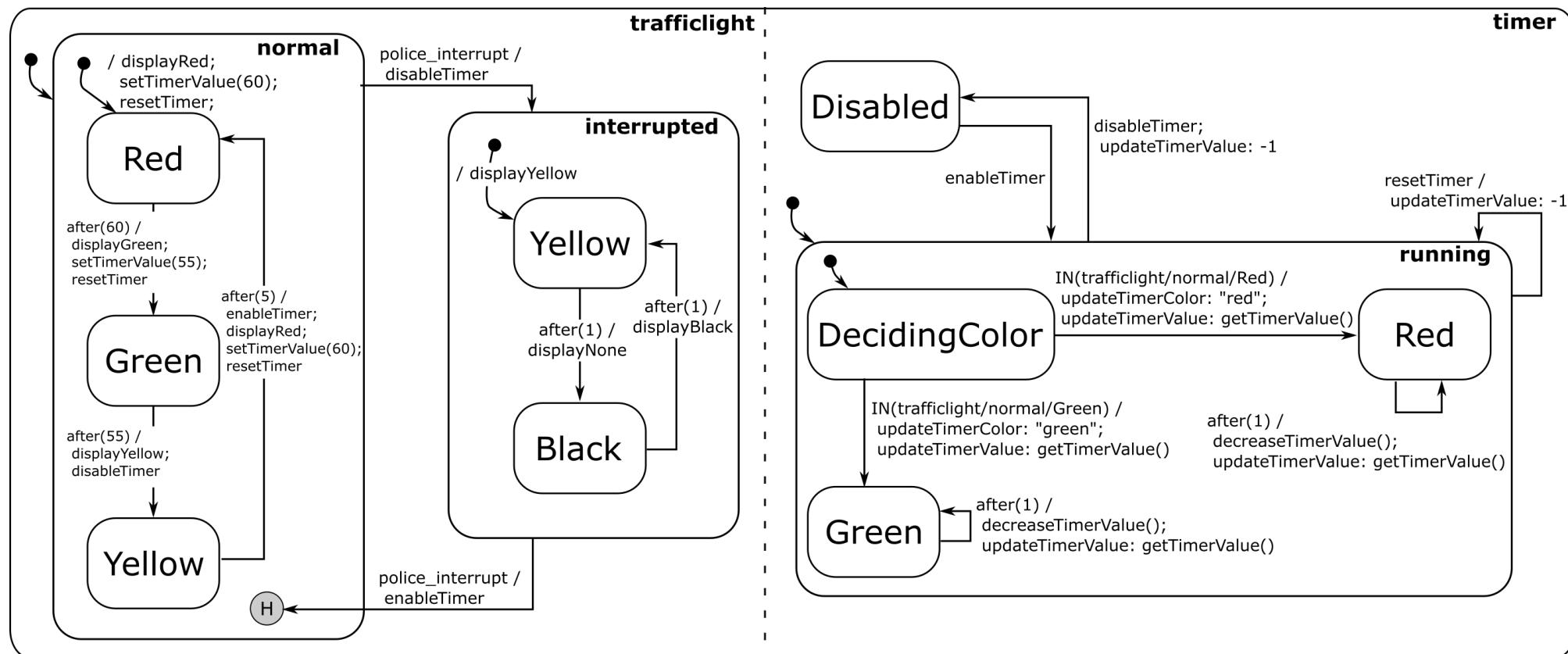
Yakindu: History



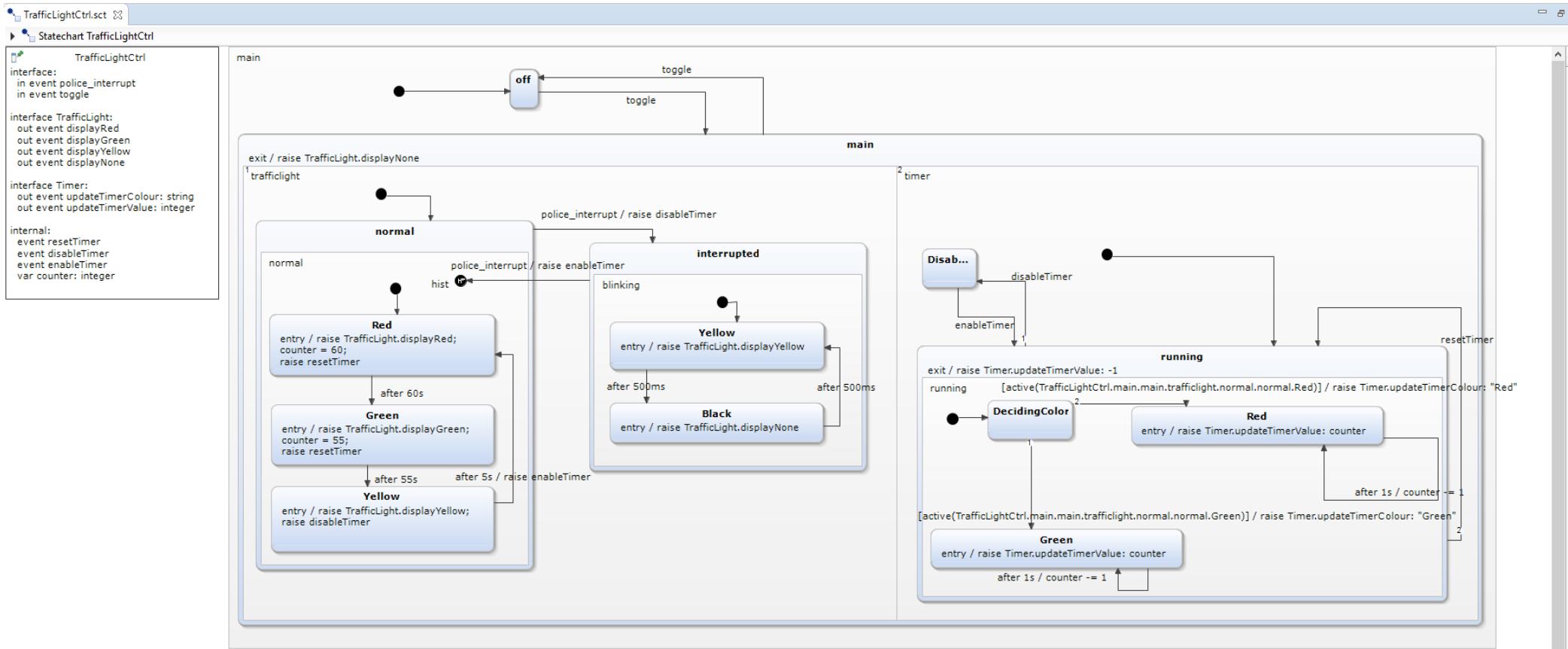
Concurrency

TrafficLight
- timer: int

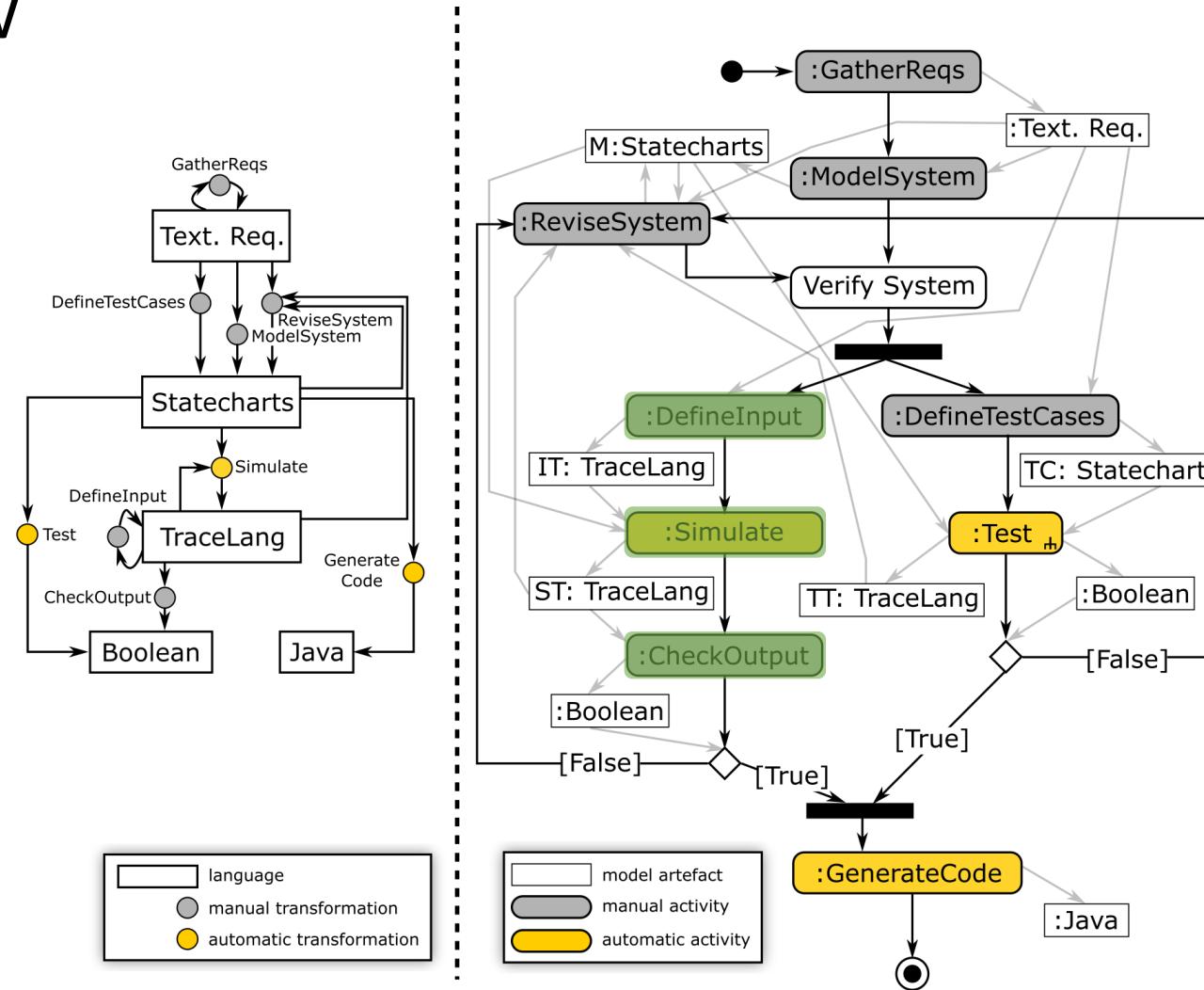
- R8: a timer displays the remaining time while the light is red or green; this timer decreases and displays its value every second. The colour of the timer reflects the colour of the traffic light.



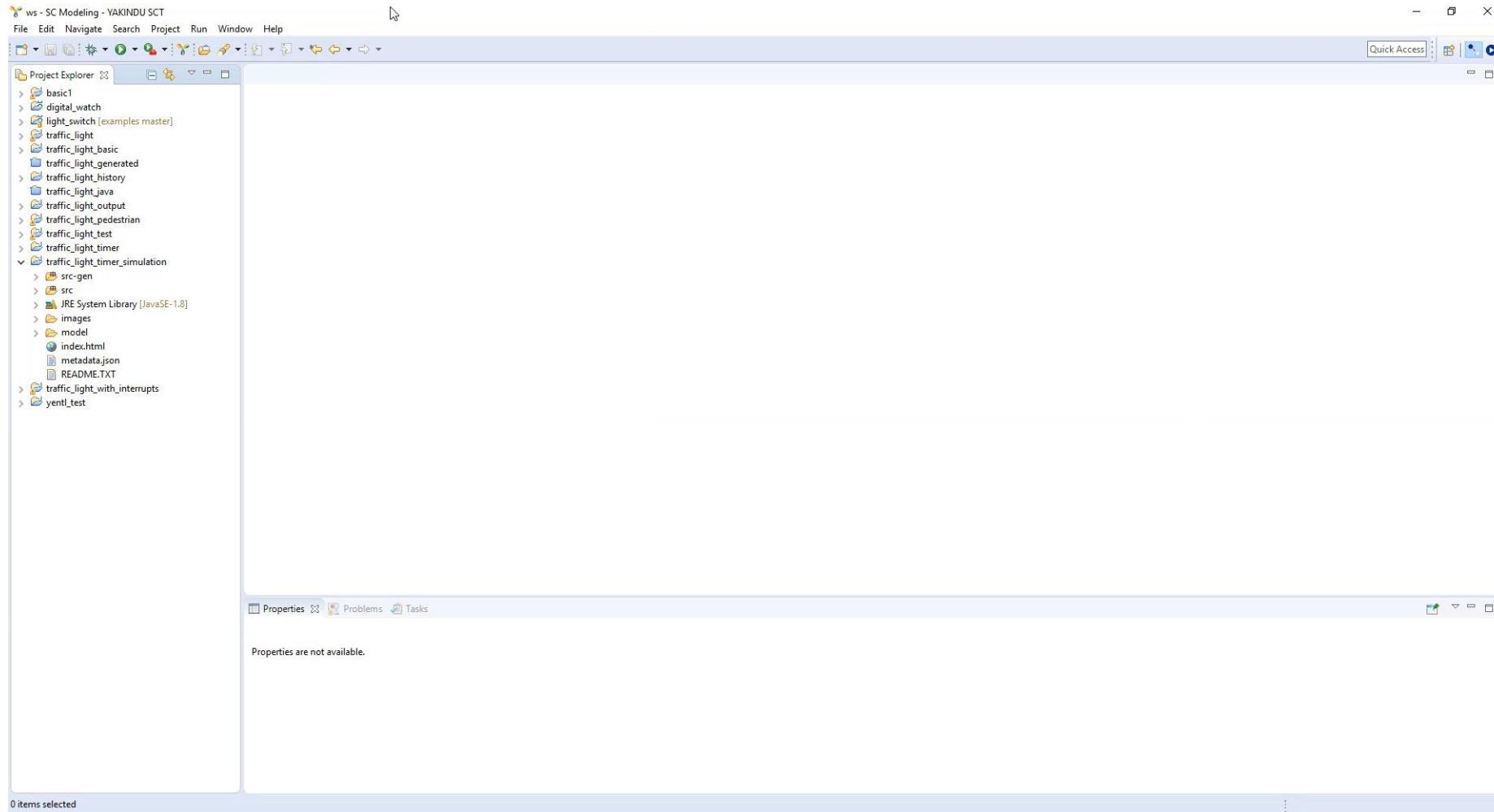
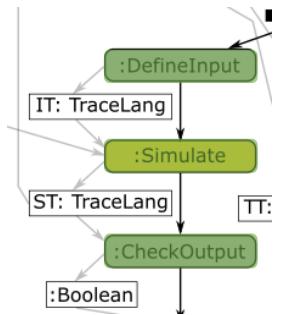
Yakindu: Concurrency



Workflow



Yakindu: Simulation (Scaled Real-Time)



Statechart Semantics: Initialization

```
init(sc):
    targetStates =
        getEffectiveTargetStates(getDefaultState(sc))
    for target in targetStates:
        enter(target)
```

Statecharts Semantics: “Main Loop”

while True:

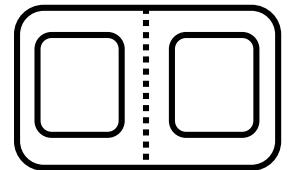
 for all concurrent regions:

 candidates =

 findEnabledTransitions(getEnabledEvents(),
 getCurrentState())

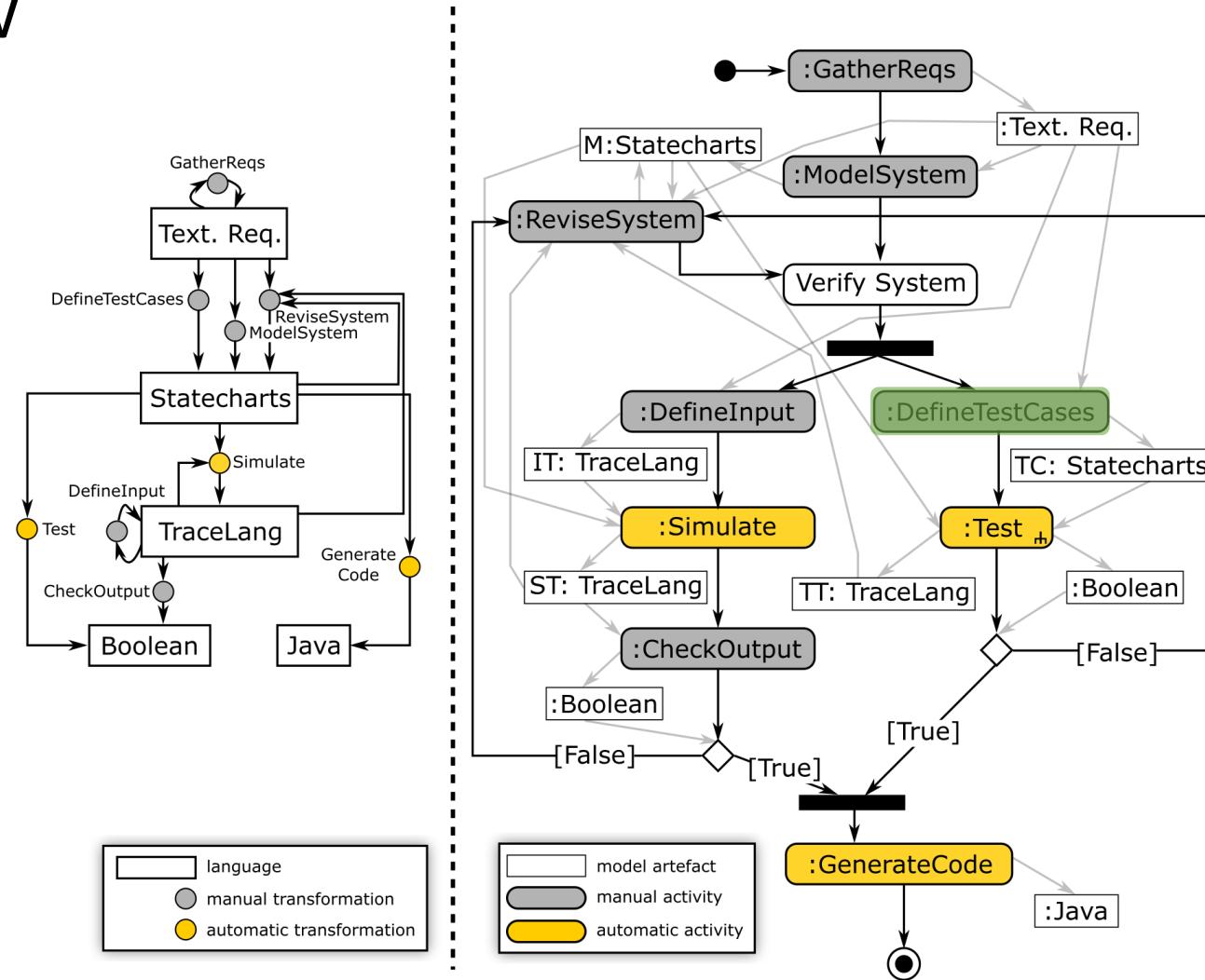
 removeConflicts(candidates)

 execute(chooseOne(candidates))

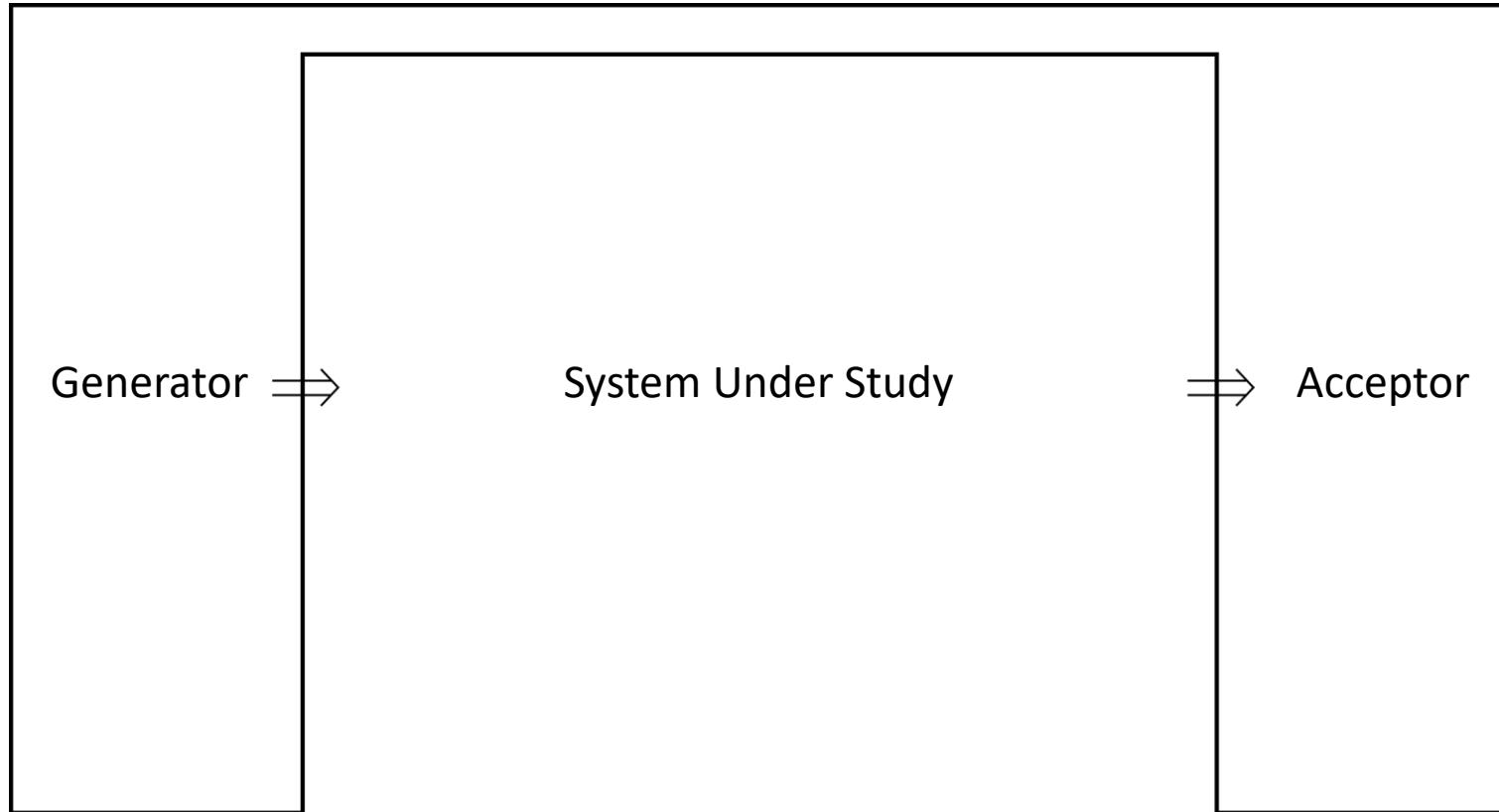
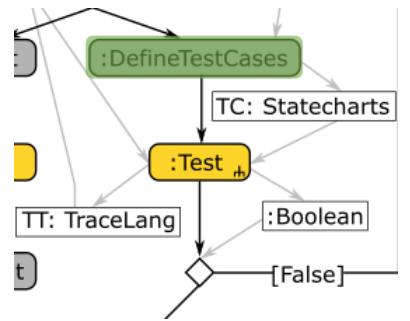


1. Find LCA
2. Leave states up the hierarchy
3. Execute action a
4. Enter states down the hierarchy
(getEffectiveTargetStates())

Workflow



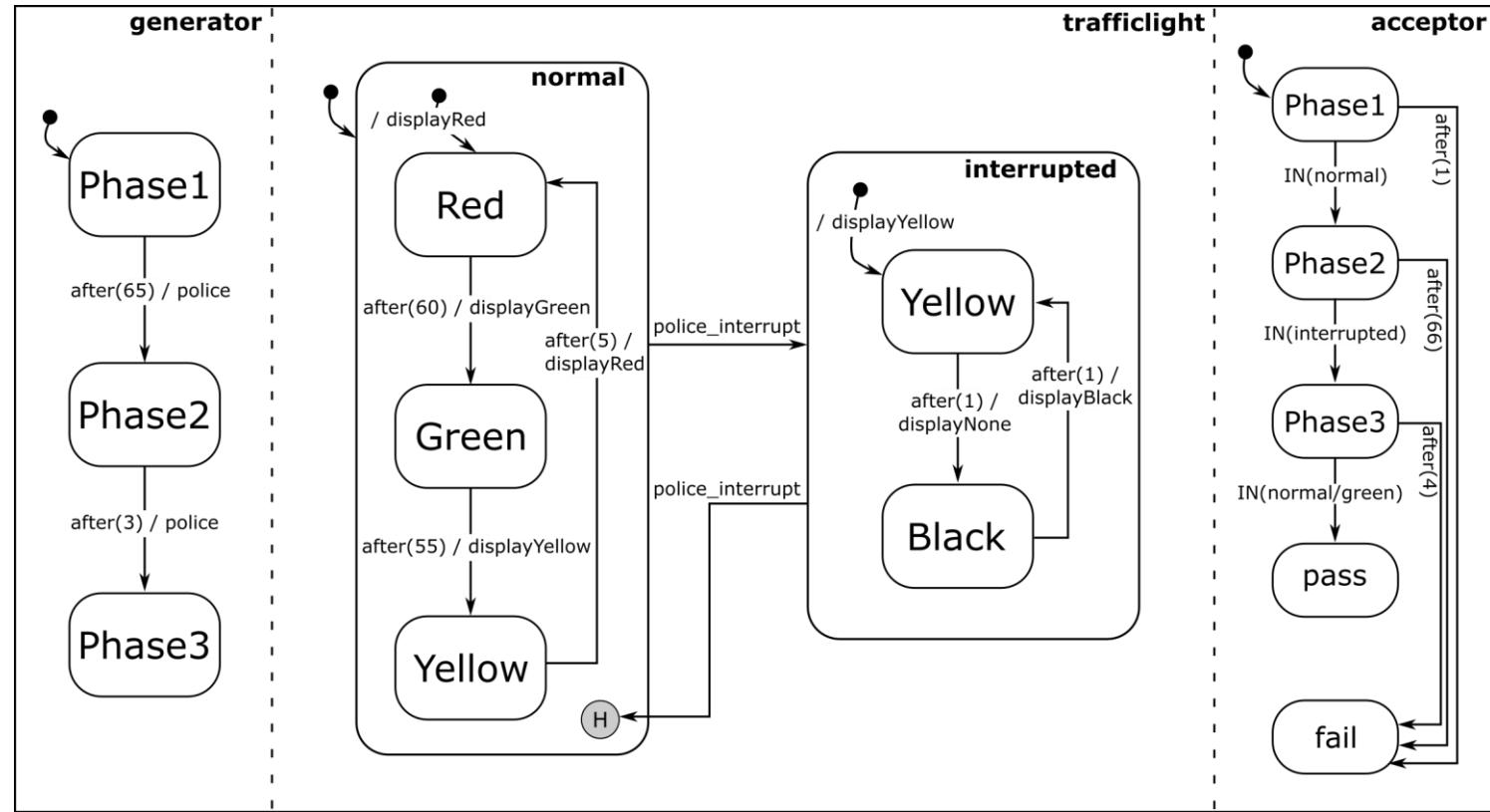
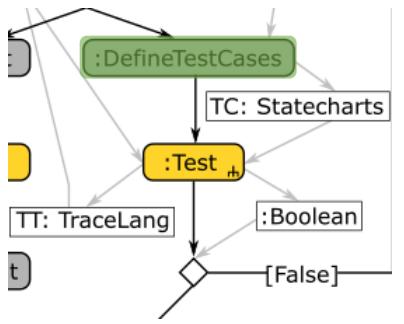
Statecharts Testing



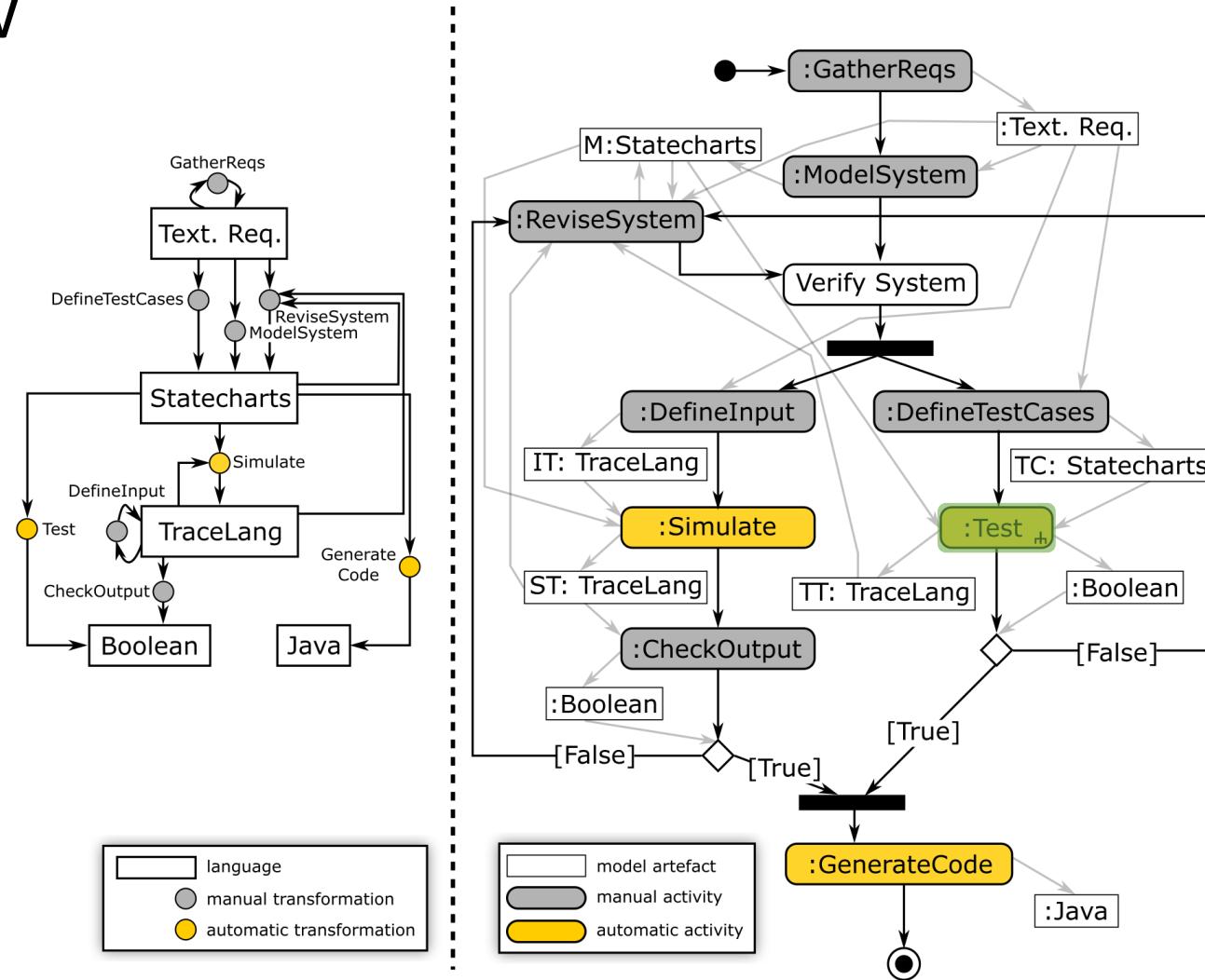
⁶ Zeigler BP. Theory of modelling and simulation. New York: Wiley-Interscience, 1976.

⁷ Mamadou K. Traoré, Alexandre Muzy, Capturing the dual relationship between simulation models and their context, Simulation Modelling Practice and Theory, Volume 14, Issue 2, February 2006, Pages 126-142

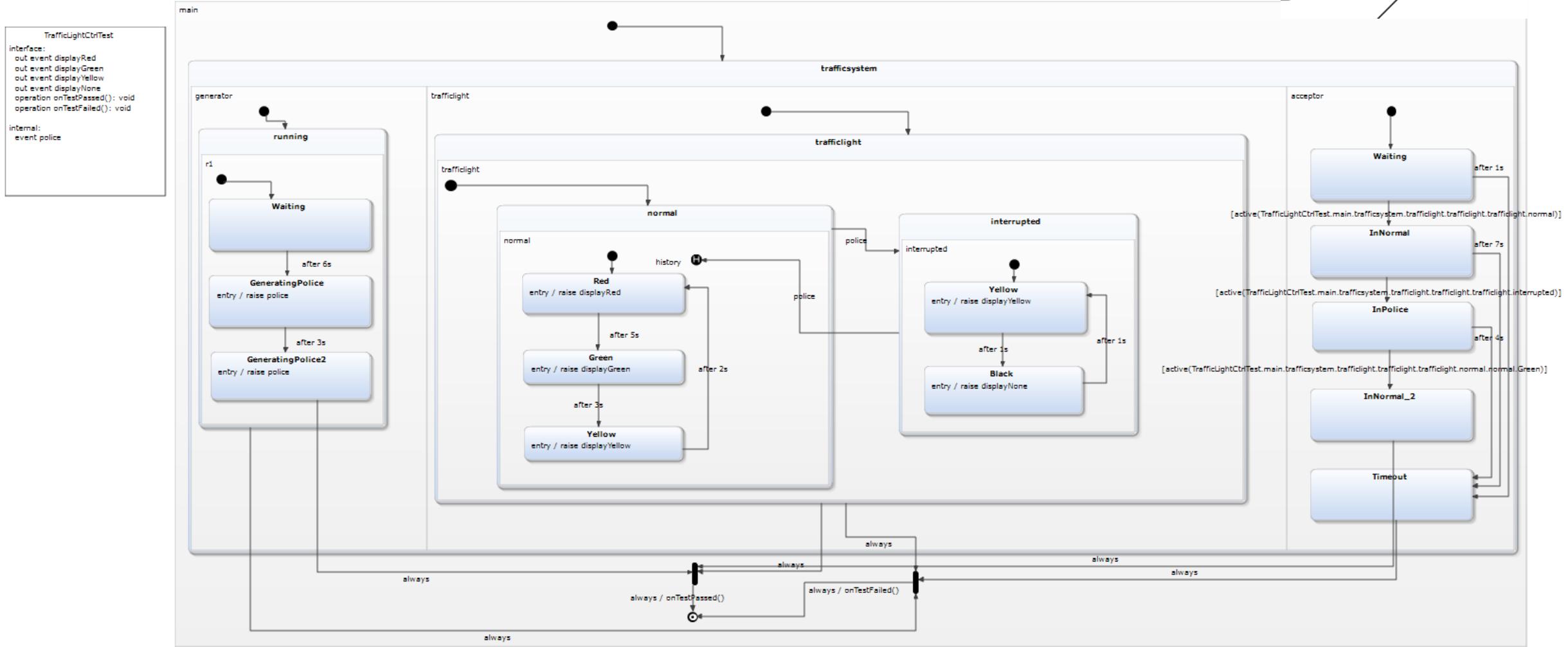
Orthogonal Components (White-Box)



Workflow



Yakindu: Testing



Yakindu: Testing

Interface

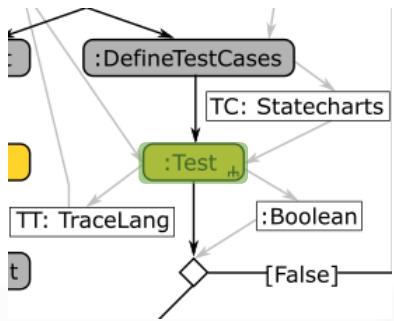
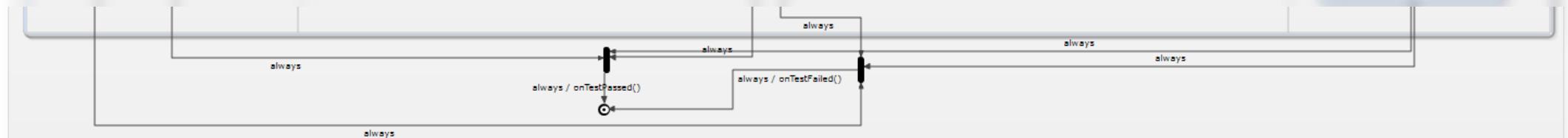
```
TrafficLightCtrlTest
interface:
out event displayRed
out event displayGreen
out event displayYellow
out event displayNone
operation onTestPassed(): void
operation onTestFailed(): void

internal:
event police
```

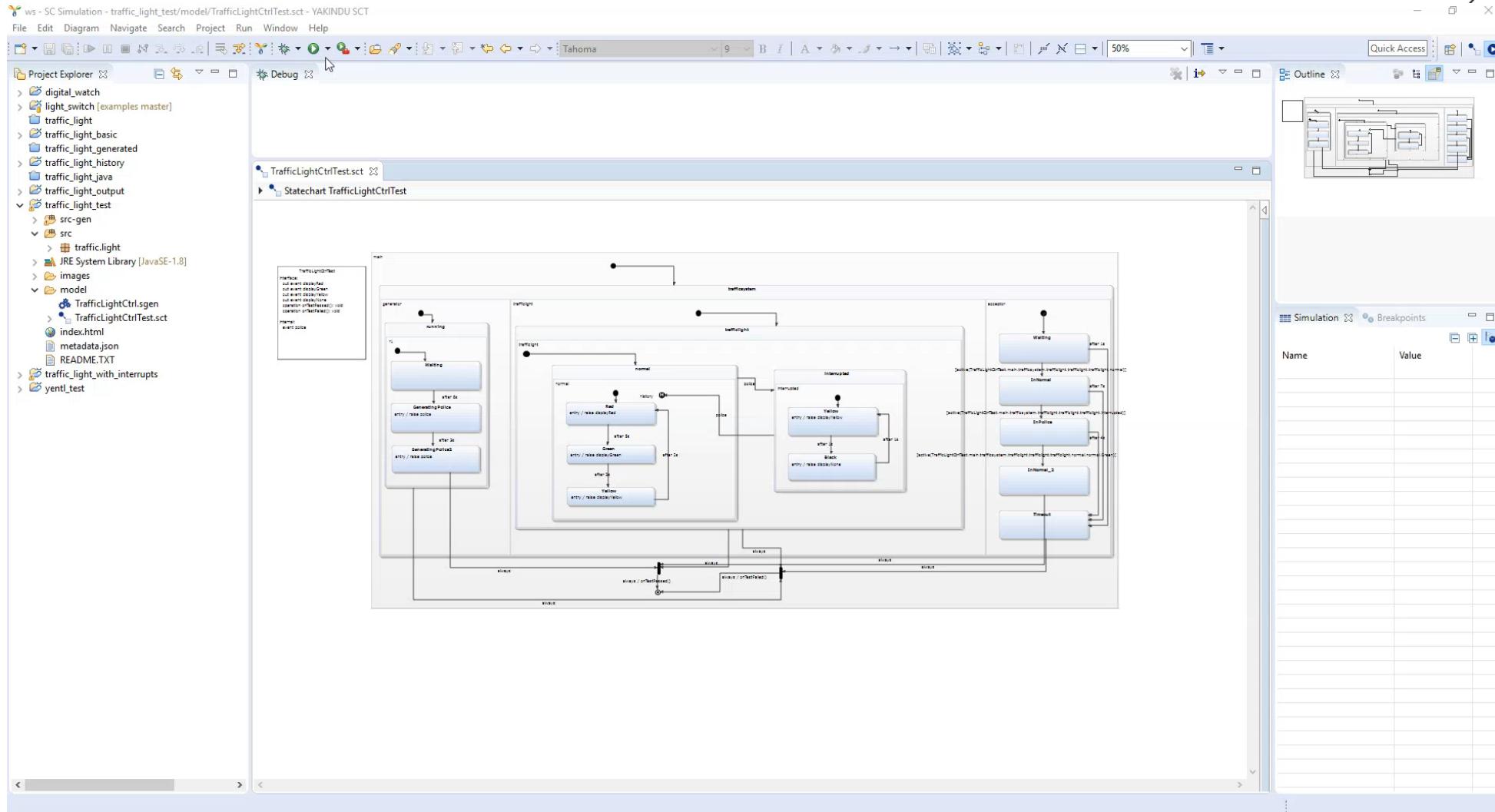
Callback

```
statemachine = new SynchronizedTrafficLightCtrlTestStatemachine();
timer = new TimerService();
statemachine.setTimer(timer);
statemachine.getSCIface().setSCIfaceOperationCallback(new ITrafficLightCtrlTestStatemachine.SCIfaceOperationCallback() {
    @Override
    public void onTestPassed() {
        System.out.println("Test passed!");
        System.exit(0);
    }
    @Override
    public void onTestFailed() {
        System.out.println("Test failed!");
        System.exit(0);
    }
});
```

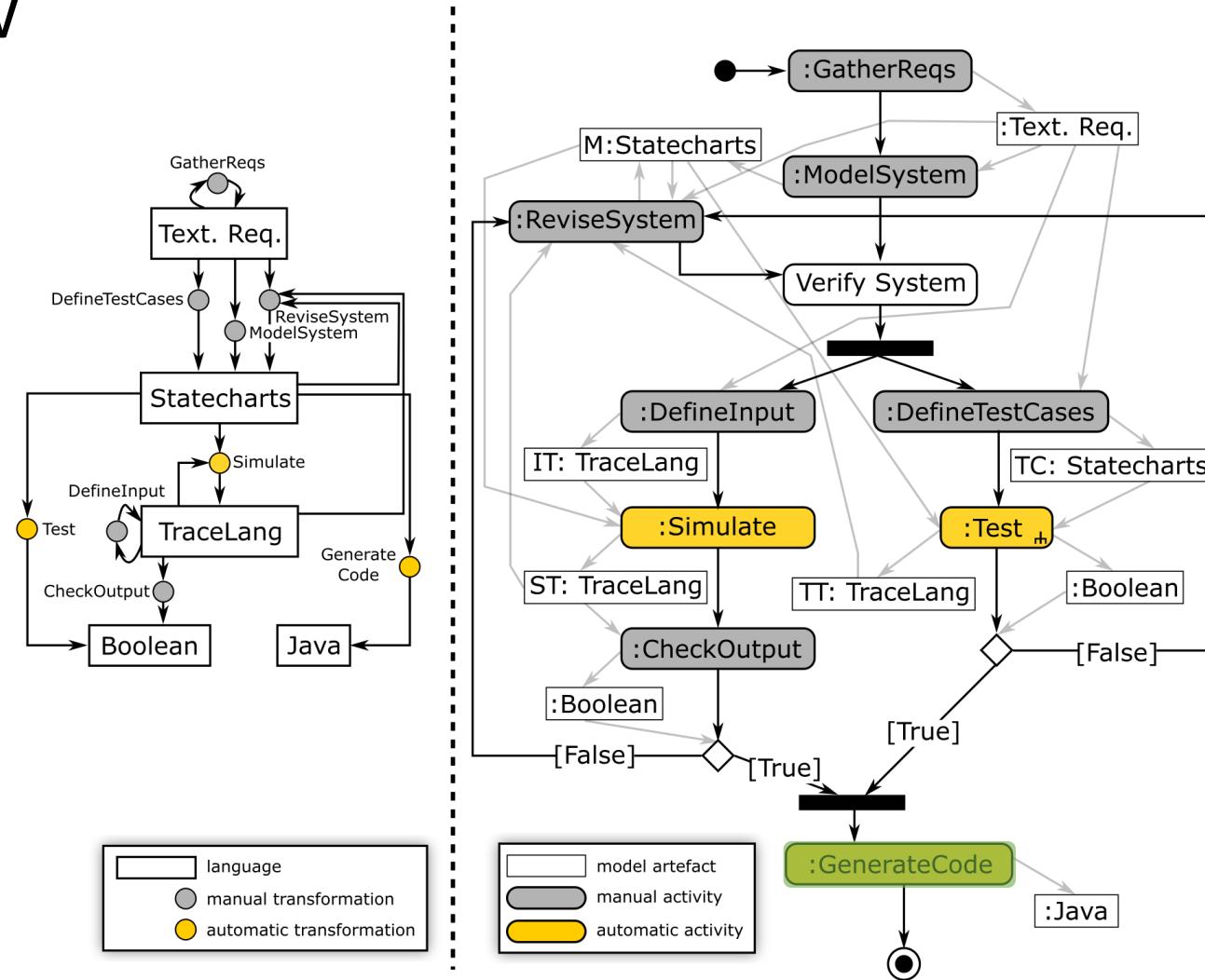
Synchronization



Yakindu: Testing



Workflow



Code Generation



interrupts

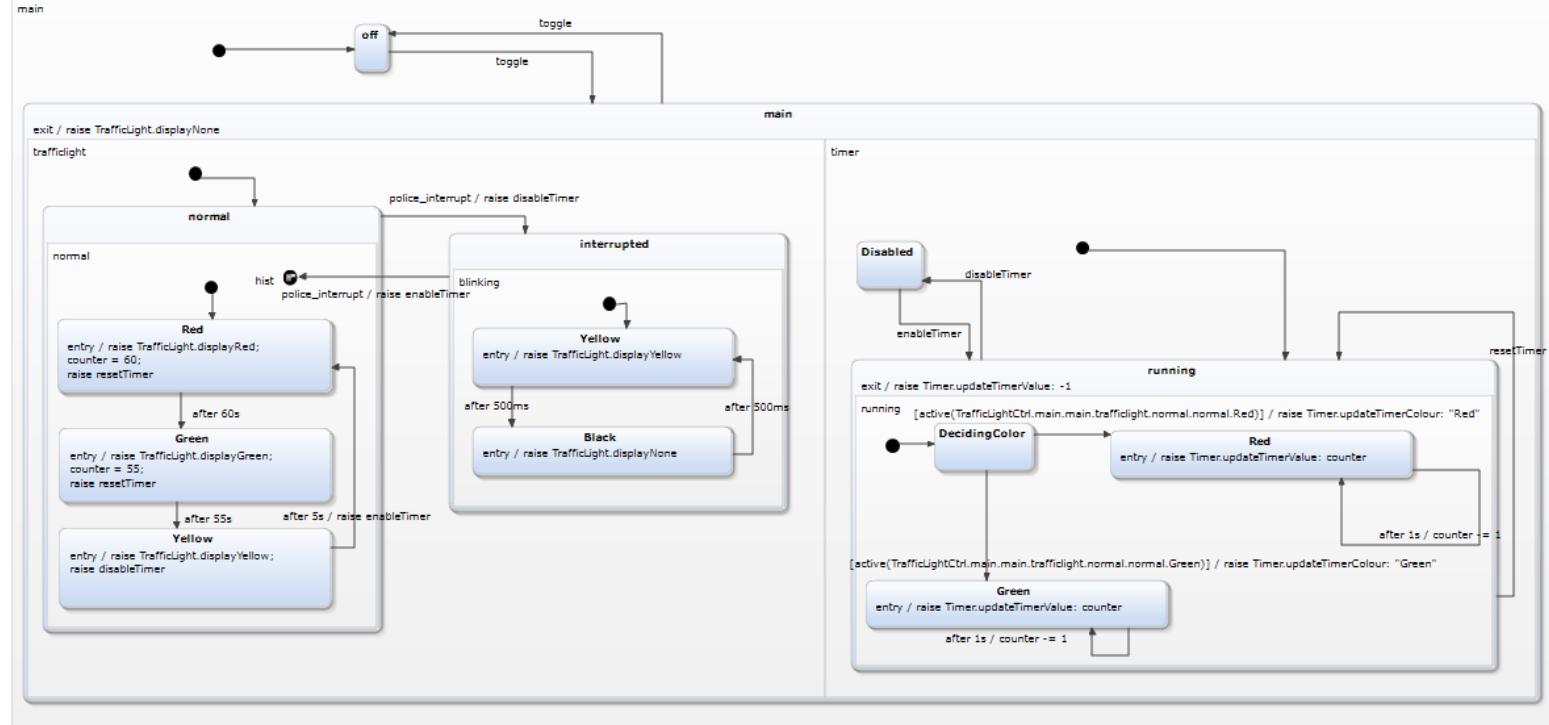
events

```
TrafficLightCtrl
interface:
in event police_interrupt
in event toggle

interface Trafficlight:
out event displayRed
out event displayGreen
out event displayYellow
out event displayNone

interface Timer:
out event updateTimerColour: string
out event updateTimerValue: integer

internal:
event resetTimer
event disableTimer
event enableTimer
var counter: integer
```

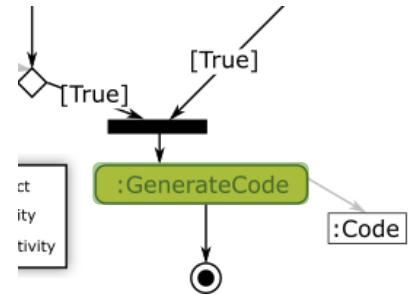


Interface:

- setRed(boolean)
- setGreen(boolean)
- setYellow(boolean)
- setTimerValue(int)
- setTimerColour(string)

Interface:

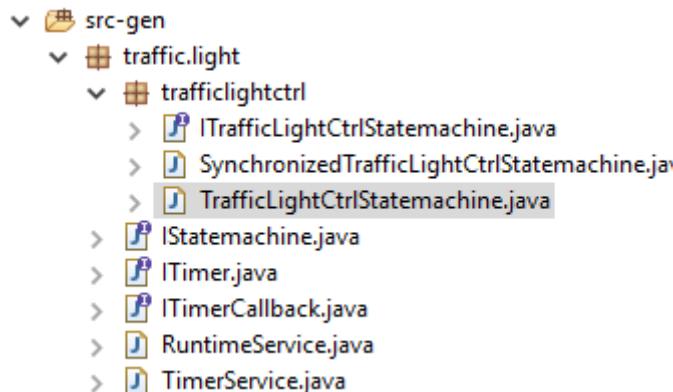
- in event *police_interrupt*
- in event *toggle*
- out event *updateTimerColour: string*
- out event *updateTimerValue: int*
- out event *displayRed, displayYellow, displayGreen, displayNone*



Generated Code

Sample

Files



- 8 files
- 1311 lines of code
- 302 manual (UI) code

The screenshot shows a code editor with two tabs: 'TrafficLightCtrl.sct' and 'TrafficLightCtrlStatemachine.java'. The code in the editor is as follows:

```
break;
case main_main_trafficlight_interrupted_blinking_Yellow:
    exitSequence_main_main_trafficlight_interrupted_blinking_Yellow();
    break;
case main_main_trafficlight_normal_normal_Red:
    exitSequence_main_main_trafficlight_normal_normal_Red();
    break;
case main_main_trafficlight_normal_normal_Yellow:
    exitSequence_main_main_trafficlight_normal_normal_Yellow();
    break;
case main_main_trafficlight_normal_normal_Green:
    exitSequence_main_main_trafficlight_normal_normal_Green();
    break;
default:
    break;
}

/* Default exit sequence for region blinking */
private void exitSequence_main_main_trafficlight_interrupted_blinking() {
    switch (stateVector[0]) {
        case main_main_trafficlight_interrupted_blinking_Black:
            exitSequence_main_main_trafficlight_interrupted_blinking_Black();
            break;
        case main_main_trafficlight_interrupted_blinking_Yellow:
            exitSequence_main_main_trafficlight_interrupted_blinking_Yellow();
            break;
        default:
            break;
    }
}

/* Default exit sequence for region normal */
private void exitSequence_main_main_trafficlight_normal_normal() {
    switch (stateVector[0]) {
        case main_main_trafficlight_normal_normal_Red:
            exitSequence_main_main_trafficlight_normal_normal_Red();
            break;
        case main_main_trafficlight_normal_normal_Yellow:
            exitSequence_main_main_trafficlight_normal_normal_Yellow();
            break;
        case main_main_trafficlight_normal_normal_Green:
            exitSequence_main_main_trafficlight_normal_normal_Green();
            break;
        default:
            break;
    }
}

/* Default exit sequence for region timer */
private void exitSequence_main_main_timer() {
    ... (truncated)
}
```

Interface

```
TrafficLightCtrl
interface:
in event police_interrupt
in event toggle

interface TrafficLight:
out event displayRed
out event displayGreen
out event displayYellow
out event displayNone

interface Timer:
out event updateTimerColour: string
out event updateTimerValue: integer

internal:
event resetTimer
event disableTimer
event enableTimer
var counter: integer
```

Setup Code (Excerpt)

```
protected void setupStatemachine() {
    statemachine = new SynchronizedTrafficLightCtrlStatemachine();
    timer = new MyTimerService(10.0);
    statemachine.setTimer(timer);

    statemachine.getSCITrafficLight().getListeners().add(new ITrafficLightCtrlStatemachine.SCITrafficLightListener() {
        @Override
        public void onDisplayYellowRaised() {
            setLights(false, true, false);
        }

        public void onDisplayRedRaised() { }

        public void onDisplayNoneRaised() { }

        public void onDisplayGreenRaised() { }

    });
    statemachine.getSCITimer().getListeners().add(new ITrafficLightCtrlStatemachine.SCITimerListener() {
        @Override
        public void onUpdateTimerValueRaised(long value) {
            crossing.getCounterVis().setCounterValue(value);
            repaint();
        }

        @Override
        public void onUpdateTimerColourRaised(String value) {
            crossing.getCounterVis().setColor(value == "Red" ? Color.RED : Color.GREEN);
        }
    });

    buttonPanel.getPoliceInterrupt()
        .addActionListener(e -> statemachine.getSCIface().raisePolice_interrupt());

    buttonPanel.getSwitchOnOff()
        .addActionListener(e -> statemachine.getSCIface().raiseToggle());

    statemachine.init();
}

private void setLights(boolean red, boolean yellow, boolean green) {
    crossing.getTrafficLightVis().setRed(red);
    crossing.getTrafficLightVis().setYellow(yellow);
    crossing.getTrafficLightVis().setGreen(green);
    repaint();
}

protected void run() {
    statemachine.enter();
    RuntimeService.getInstance().registerStatemachine(statemachine, 100);
}
```

Generator

GeneratorModel for yakindu::java {

```
statechart TrafficLightCtrl {

    feature Outlet {
        targetProject = "traffic_light_history"
        targetFolder = "src-gen"
    }

    feature Naming {
        basePackage = "traffic.light"
        implementationSuffix =""
    }

    feature GeneralFeatures {
        RuntimeService = true
        TimerService = true
        InterfaceObserverSupport = true
    }

    feature SynchronizedWrapper {
        namePrefix = "Synchronized"
        nameSuffix = ""
    }
}
```

Runner

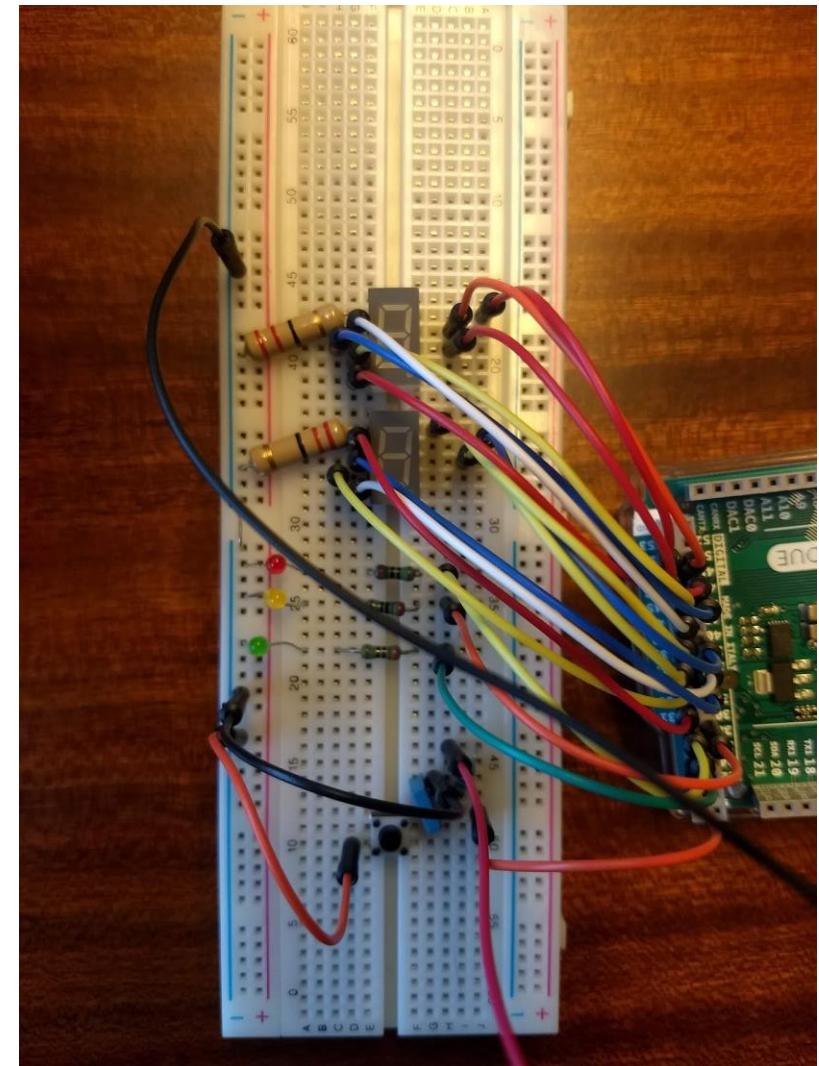
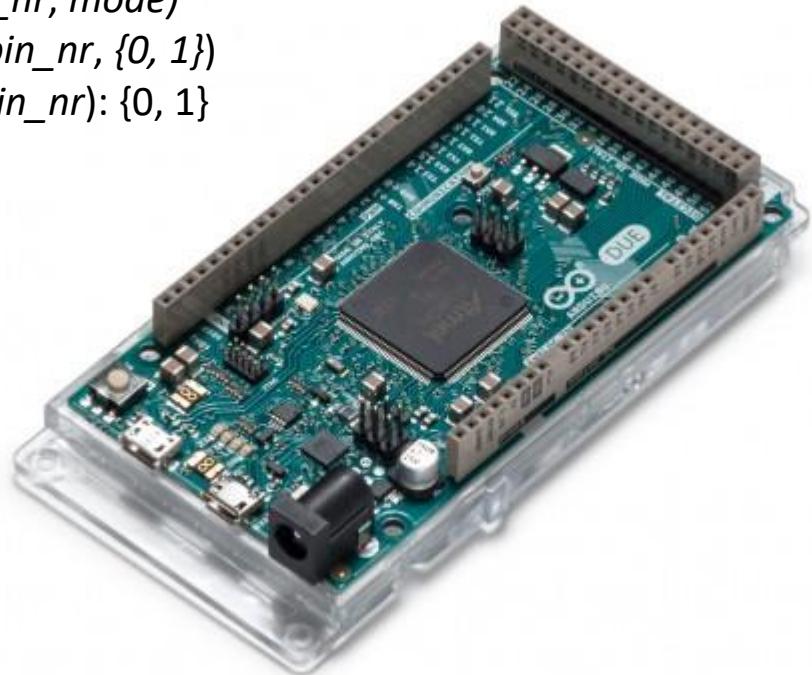
Deployed Application (Scaled Real-Time)



Deploying onto Hardware

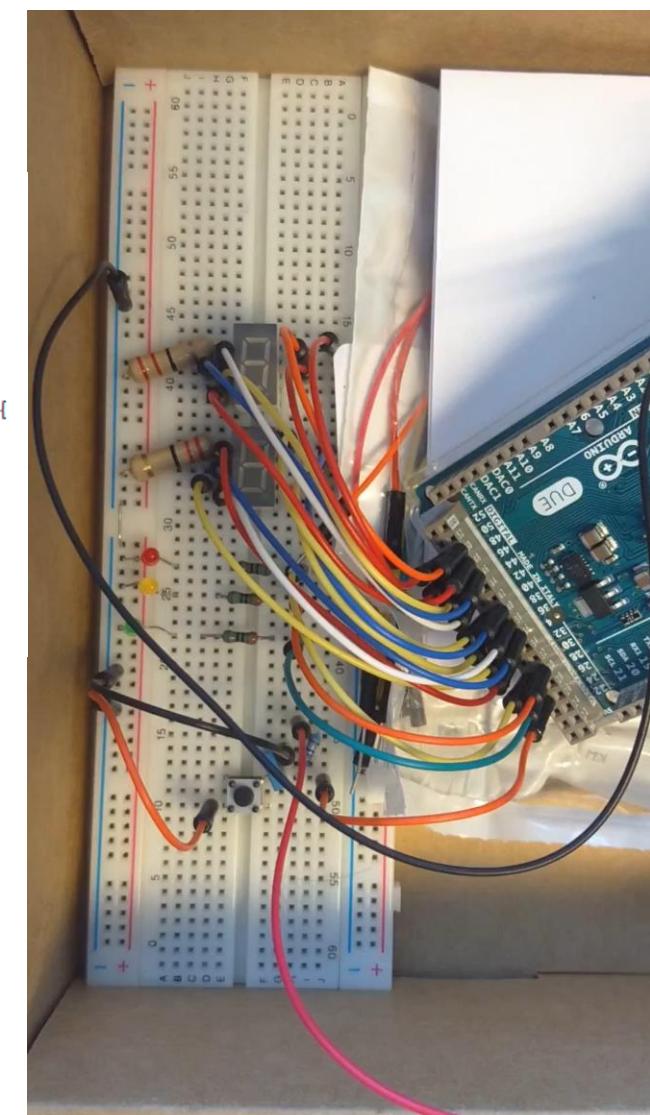
Interface:

- `pinMode(pin_nr, mode)`
- `digitalWrite(pin_nr, {0, 1})`
- `digitalRead(pin_nr): {0, 1}`



Deploying onto Hardware

Deployed Application



Generator

```
GeneratorModel for yakindu::c {

    statechart TrafficLightCtrl {

        feature Outlet {
            targetProject = "traffic_light_arduino"
            targetFolder = "src-gen"
            libraryTargetFolder = "src-gen"
        }

        feature FunctionInlining {
            inlineReactions = true
            inlineEntryActions = true
            inlineExitActions = true
            inlineEnterSequences = true
            inlineExitSequences = true
            inlineChoices = true
            inlineEnterRegion = true
            inlineExitRegion = true
            inlineEntries = true
        }

    }
}
```

Runner

```
#define CYCLE_PERIOD (10)
static unsigned long cycle_count = 0L;
static unsigned long last_cycle_time = 0L;

void loop() {
    unsigned long current_millis = millis();
    read_pushbutton(&pushbutton);
    if (cycle_count == 0L || (current_millis >= last_cycle_time + CYCLE_PERIOD) ) {
        sc_timer_service_proceed(&timer_service, current_millis - last_cycle_time);
        synchronize(&trafficLight);
        trafficLightCtrl_runCycle(&trafficLight);
        last_cycle_time = current_millis;
        cycle_count++;
    }
}
```

Button Code

```
void read_pushbutton(pushbutton_t *button) {
    int pin_value = digitalRead(button->pin);
    if (pin_value != button->debounce_state) {
        button->last_debounce_time = millis();
    }
    if ((millis() - button->last_debounce_time) > button->debounce_delay) {
        if (pin_value != button->state) {
            button->state = pin_value;
            button->callback(button);
        }
    }
    button->debounce_state = pin_value;
}
```

INTRODUCTION

STATECHARTS BASICS

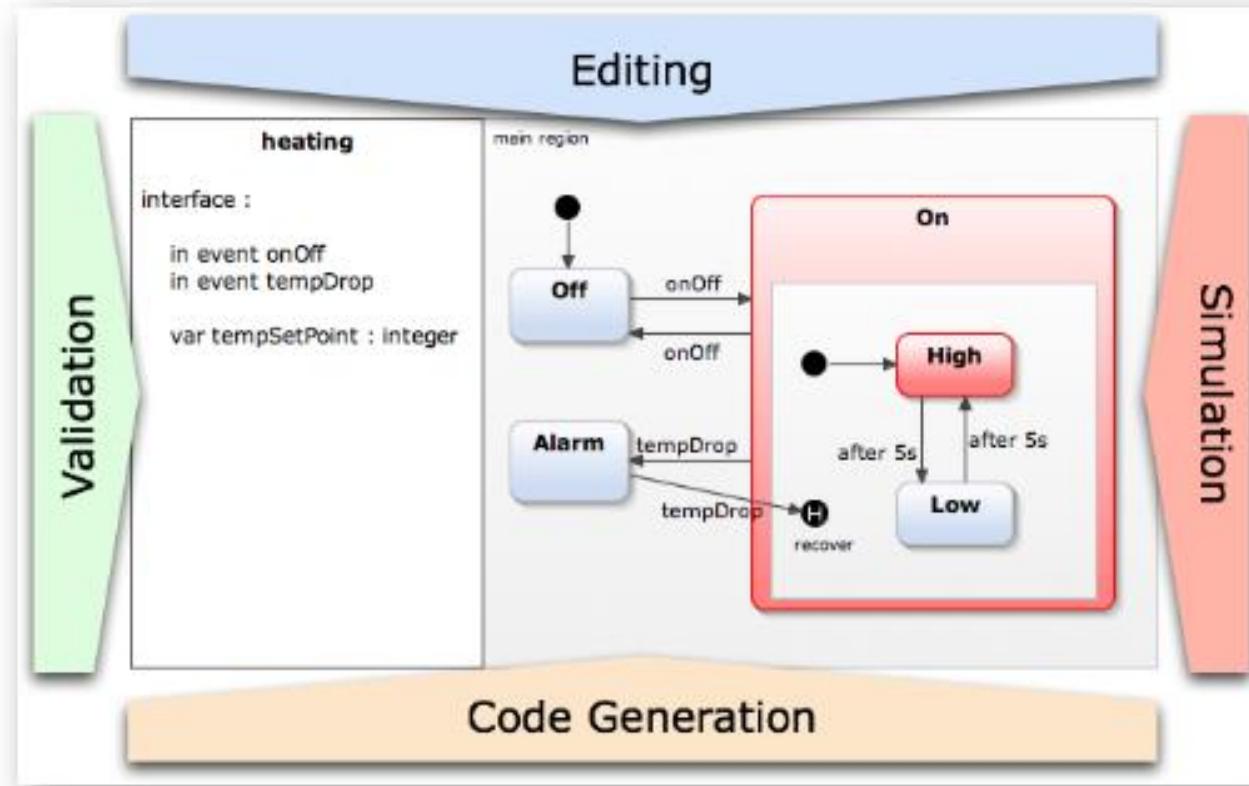
YAKINDU IN DEPTH

ADVANCED CONCEPTS

Overview

graphically create and edit Statecharts

validate syntax and
(static) semantics



code generators for Java, C, and C++
+ custom code generators

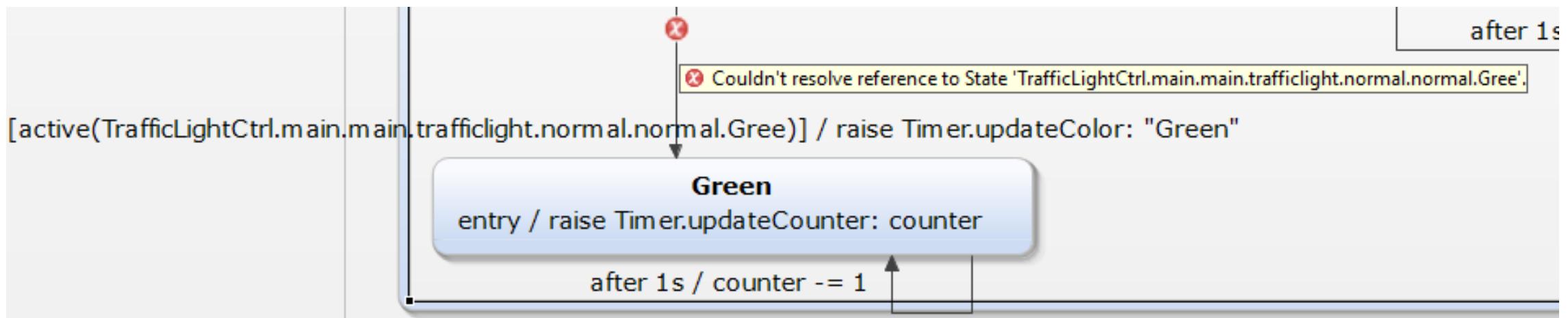
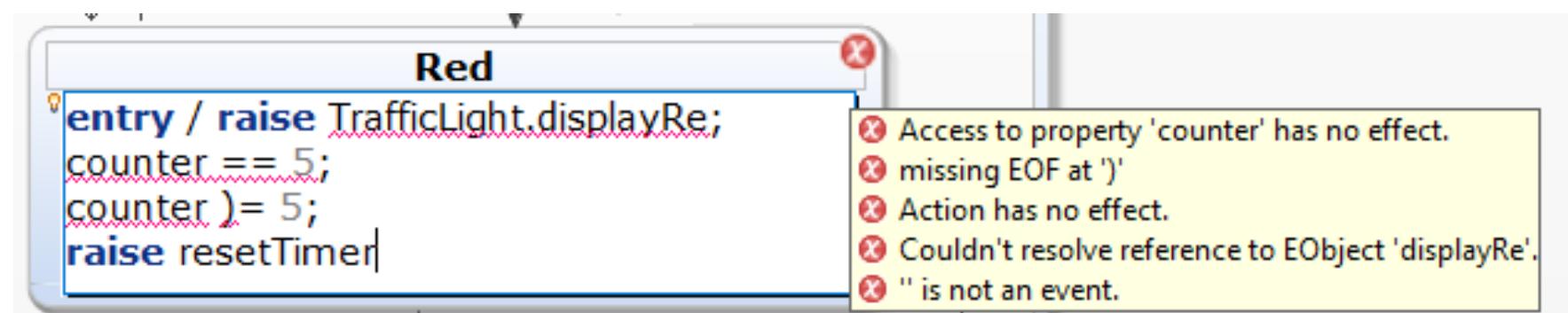
simulate the behavior
of Statecharts

Validator

```
interface TrafficLight:  
    out event displayRed  
    out event displayGreen  
    out event displayYellow  
    out event displayNone
```

```
interface Timer:  
    out event updateColor: string  
    out event updateCounter: integer
```

```
internal:  
    event resetTimer  
    event disableTimer  
    event enableTimer  
    var counter: integer
```



Simulation and Debugging

The screenshot shows a statechart simulation interface for a traffic light control system. The main window displays a statechart diagram with several states: off, normal, blinking, interrupted, and running. Transitions between states are triggered by events like 'police_interrupt' or time intervals like 'after 5s'. The 'normal' state has sub-states Red, Green, and Yellow, each with its own entry actions and timer triggers. A 'DecidingColor' state is used to switch between Red and Green. The 'running' state has a 'Disable...' transition that leads back to the 'normal' state.

Current State Information

Play/Pause/Stop/Step

Current Time

Events (Raise)

Local Variable Values (Inspect + Modify)

Debug X
main.py [C:\Python37\python3.exe]
TrafficLightCtrl.sct (7) [Statechart Simulation]
TrafficLightCtrl

TrafficLightCtrl.sct X
Statechart TrafficLightCtrl

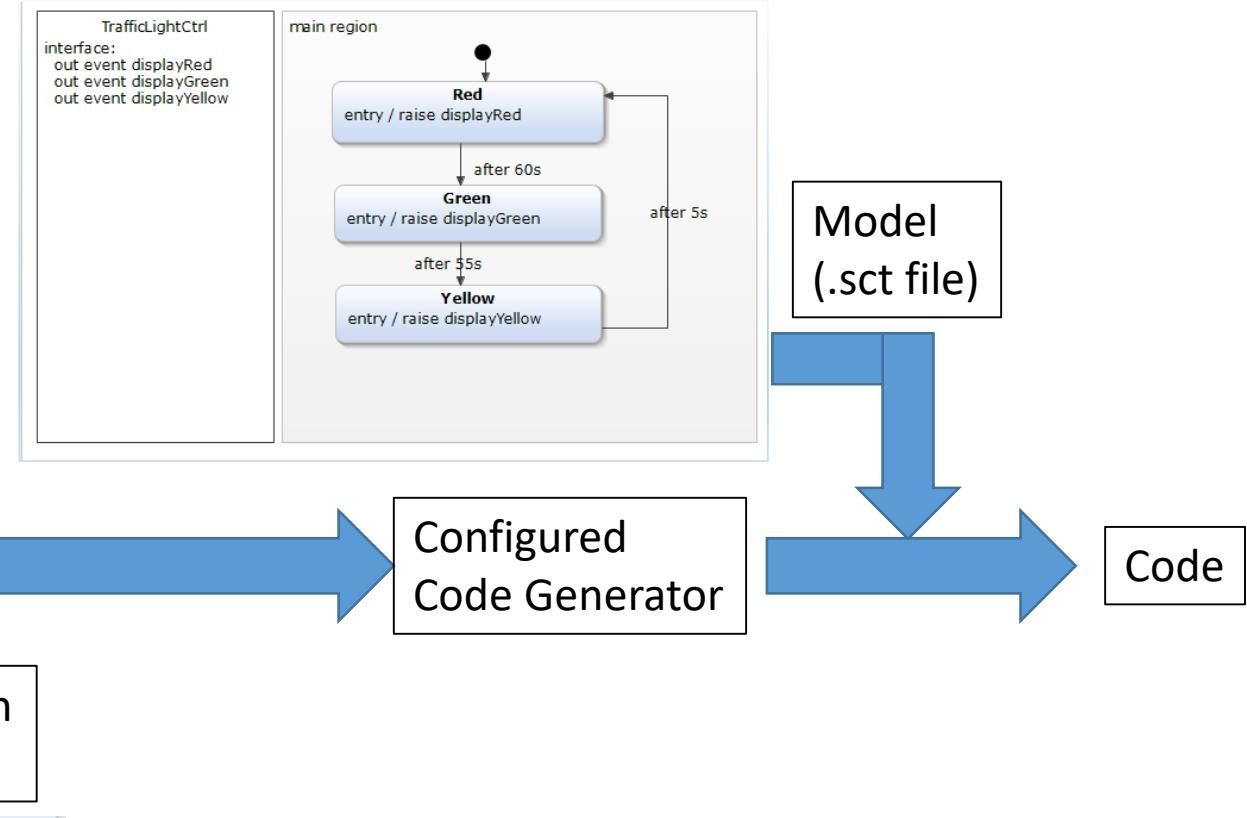
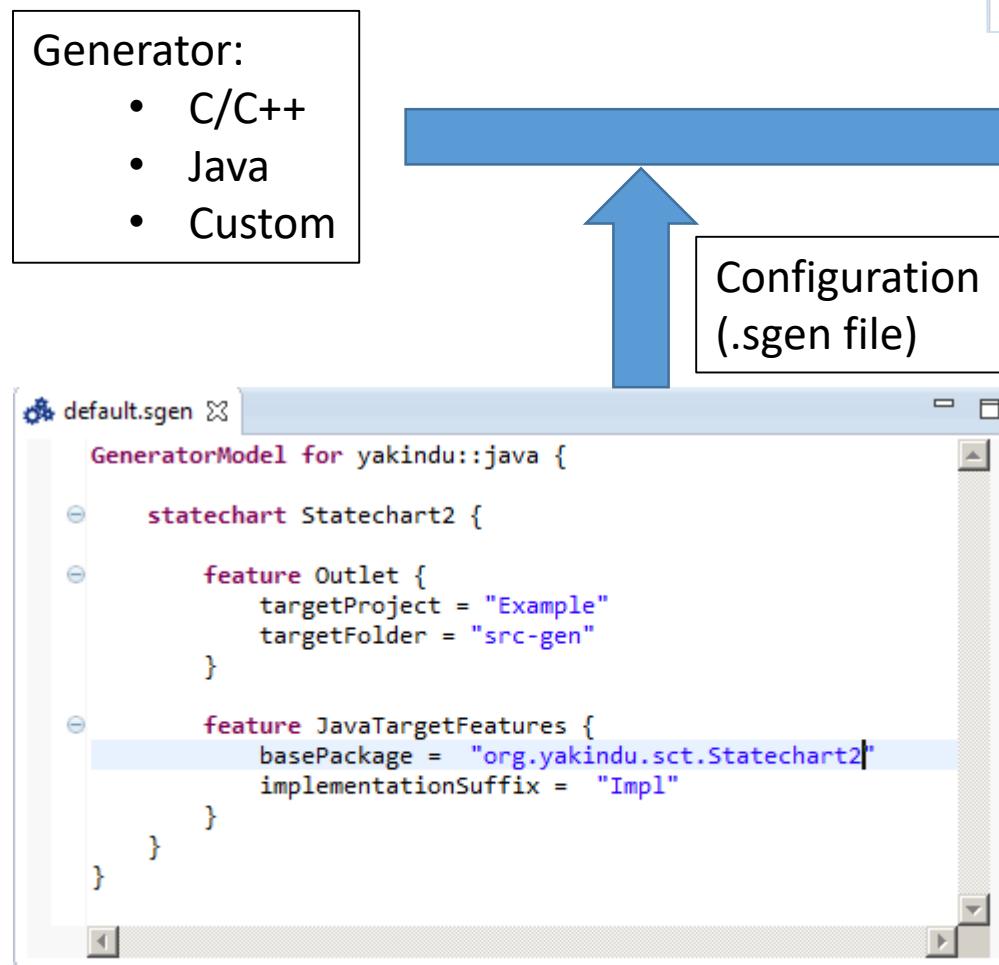
Outline X

Transition
S State
C Composite State
Orthogonal State
Region
Entry
H Shallow History
Deep History
Final State
X Exit Node
Diamond Choice
Synchronization

TrafficLightCtrl 00:00:06.800

Name	Value
default	
police_interrupt	
toggle	
TrafficLight	
Timer	
internal	
time events	

Code Generators



```
GeneratorModel for [GeneratorId] {
    statechart [StatechartReference] {
        feature [Feature] {
            [ParameterName] = [ParameterValue]
        }
    }
}
```

Examples

File -> New... -> Example... ->
Yakindu Statechart Examples

YAKINDU Examples

Select an example

Choose an example project to import it.

- [Pro] Autonomous Robot
- [Pro] Coffee Machine with C++ integration
- [Pro] Coffee Machine with C integration
- [Pro] Headless: Ant
- [Pro] Headless: Gradle
- [Pro] Headless: Make
- [Pro] Headless: Maven
- [Pro] MSP430 Blinky LED
- [Pro] Sensor example with Arrays
- Basic Finite State Machine for Arduino
- Custom Generator Example for SCXML
- Digital Watch
- Example Template
- Four Iterations of a Coffee Machine
- Light Switch Example
- Light Switch Example Series
- MSP430 Blinky LED
- Raspberry Pi Hello World
- Swift Example
- Testing Elevator Model with SCTUnit
- Traffic Light (C++)
- Traffic Light (Java)
- Traffic Light (Java) with SCTUnit
- Traffic Light (Python)
- Traffic Light (Python) for Raspberry Pi
- Traffic Light Ported to Arduino
- Webbased YCar App

1. Four Iterations of a Coffee Machine

This example introduces composite states with three consecutive statecharts.

- The first statechart is the base for all other statecharts. It models a very basic automated coffee machine. The statechart does not contain any features but states and transitions. The user can switch the machine on and off, and order a cappuchino when the machine is turned on. He can also cancel the operation anytime.

main region

```
graph TD
    Start(( )) -->|User.on| Idle1[Idle]
    Idle1 -->|User.cappuchino| Milling[\"Milling Bea...\"]
    Milling -->|after 9s| Pouring[\"Pouring Milk\"]
    Pouring -->|User.cancel| Idle1
    Milling -->|User.cancel| Start
    Idle1 -->|User.off| Start
    Idle1 -->|User.cancel| Start
```

?

< Back Next > Finish Cancel

Neutral Action Language

- Types: integer/real/boolean/string/void
- Statements:
 - Assignment
 - Event Raising
 - Operation Call
- Expressions: arithmetic/condition/logical
- Built-in Functions:
 - *valueOf(event)*: <value>
 - <var> as <type>
 - *active(state)*: boolean

Neutral Action Language

trigger [guard] / effect

- Trigger:
 - after: execute after a given time (s, ms, us, ns)
 - every: execute periodically after a given time (s, ms, us, ns)
 - always: execute always
 - oncycle: same as always
 - else: useful for choice states
 - default: same as else
 - entry: execute upon entering the state
 - exit: execute upon exiting the state
- Guard:
 - Expression (boolean!)
- Effect:
 - Statement
 - Event Raise

Operation Callbacks

Model (Excerpt)

```
TrafficLightCtrl

interface TrafficLight:
    var red:boolean
    var yellow:boolean
    var green:boolean

interface Pedestrian:
    var request:boolean
    var red:boolean
    var green:boolean

interface:
    in event pedestrianRequest
    in event onOff
    operation synchronize() : void

internal:
    every 200ms / synchronize
    event blinkOff
    event blinkOn
```

Generated Code (Excerpt)

```
public interface SCInterface {
    public void raisePedestrianRequest();
    public void raiseOnOff();
    public void setSCInterfaceOperationCallback(SCInterfaceOperationCallback operationCallback);
}

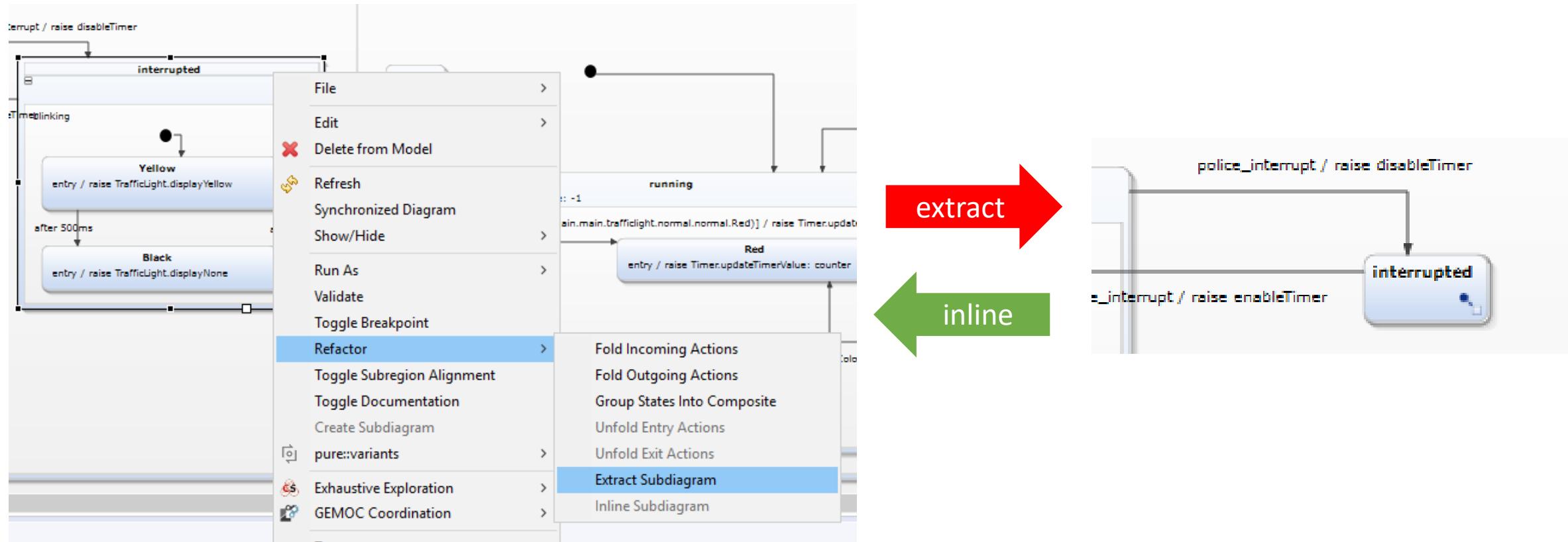
public interface SCInterfaceOperationCallback {
    public void synchronize();
}
```

Runner (Excerpt)

```
statemachine.getSCInterface().setSCInterfaceOperationCallback(
    new ITrafficLightCtrlStatemachine.SCInterfaceOperationCallback() {
        @Override
        public void synchronize() {
            checkTrafficLightStates();
            repaint();
        }
    });
}
```

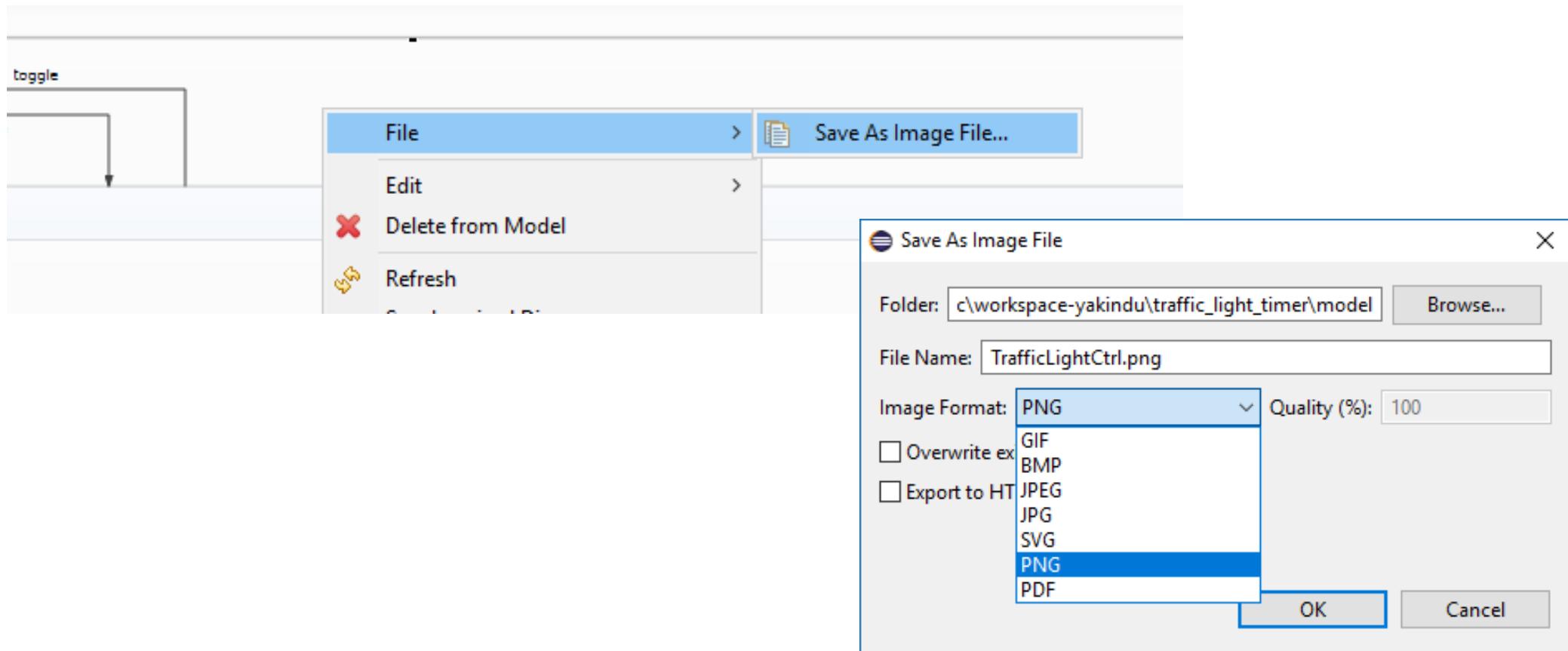
Editor Tricks

Subdiagrams



Editor Tricks

Export Model as Image



INTRODUCTION

STATECHARTS BASICS

YAKINDU IN DEPTH

ADVANCED CONCEPTS

Recap

- Model the behaviour of complex, timed, reactive, autonomous systems
 - “What” instead of “How” (= implemented by Statecharts compiler)
- Abstractions:
 - States (composite, orthogonal)
 - Transitions
 - Timeouts
 - Events
- Tool support:
 - Yakindu
 - SCCD