

CSE 5526: Introduction to Neural Networks

Hopfield Network for Associative Memory

The next few units cover unsupervised models

- Goal: learn the distribution of a set of observations
- Some observations are a better “fit” than others
- Hopfield networks store a set of observations
 - Deterministic, non-linear dynamical system
- Boltzmann machines can behave similarly
 - Stochastic, non-linear dynamical system
- Boltzmann machines with hidden units have a much greater capacity for learning the data distribution

Content-addressable memory basic task

- Store a set of “fundamental memories”
 $\{\xi_1, \xi_2, \dots, \xi_M\}$
- So that when presented with a new pattern \mathbf{x}
- The system outputs the stored memory that is most similar to \mathbf{x}
- The first content-addressable memory we will consider is the Hopfield network
 - Introduced in the influential (14,000 citations) paper Hopfield (1982). “Neural networks and physical systems with emergent collective computational abilities.” PNAS.

Is this possible? How good can it be?

- Is this possible to implement as a neural network?
 - For a single pattern?
- Does it work equally well for any pattern?
- How many patterns can such a system store?
 - How do its storage requirements compare to other sys's?
- How much corruption can it tolerate?
 - And still retrieve the correct pattern?
 - Corruption of noise or of partial information

Hopfield (1982) describes the problem

- “Any physical system whose dynamics in phase space is dominated by a substantial number of locally stable states to which it is attracted can therefore be regarded as a general content-addressable memory. The physical system will be a potentially **useful** memory if, in addition, **any** prescribed set of states can readily be made the stable states of the system.”

One associative memory: the Hopfield network

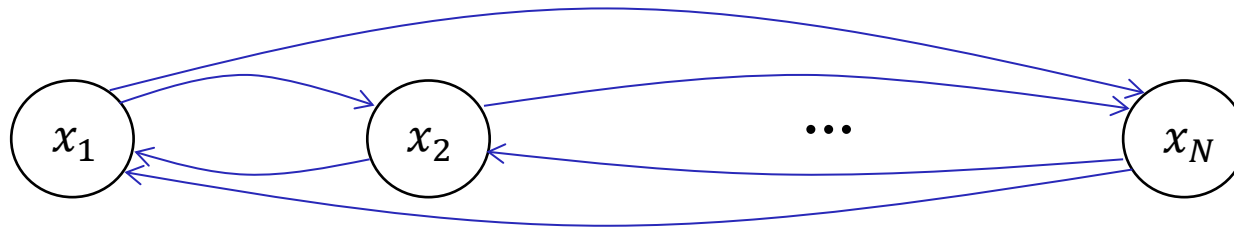
- The Hopfield net consists of N McCulloch-Pitts neurons, recurrently connected among themselves



- The network is initialized with a (corrupted) pattern

One associative memory: the Hopfield network

- The Hopfield net consists of N McCulloch-Pitts neurons, recurrently connected among themselves



- Then runs recurrently until it reaches a fixed point

State of each neuron defines the “state space”

- The network is in state \mathbf{x}_t at time t
- The state of the network evolves according to
$$\mathbf{x}_{t+1} = \varphi(W\mathbf{x}_t + \mathbf{b})$$
 - Where we set $\mathbf{b} = 0$ without loss of generality
 - Meaning that each state leads to at most one next state
- $\{\mathbf{x}_1, \mathbf{x}_2, \dots, \mathbf{x}_t\}$ is called a state trajectory
- Goal: set W so that state trajectory of corrupted memory $\xi_i + \Delta$ converges to true memory ξ_i

One-shot storage phase uses Hebbian learning

- Hopfield nets set W using the outer-product rule, one choice for doing so.

$$W = \frac{1}{N} \sum_{\mu=1}^M \xi_{\mu} \xi_{\mu}^T - I$$

Where N is the number of bits. Or, equivalently

$$w_{ji} = \frac{1}{N} \sum_{\mu=1}^M \xi_{\mu,j} \xi_{\mu,i} - \delta_{ij}$$

- The $-I$ and $-\delta_{ij}$ terms enforce $W_{ii} = 0$
 - no self-feedback

Hebbian learning

- “Neurons that fire together, wire together”
- In the Hopfield network, increase the weights of neurons that receive correlated inputs
- This notion is symmetric between neurons
 - And since $w_{ji} = w_{ij}$, the weight matrix is symmetric

Retrieval phase

- Play out dynamics $\mathbf{x}_{t+1} = \varphi(W\mathbf{x}_t)$
 - Until reaching a stable state $\mathbf{x}_{t+1} = \mathbf{x}_t$
 - If argument to $\varphi(\cdot)$ is 0, neuron stays in previous state
 - Leads to symmetric flow diagrams
- Can also use “asynchronous” updates
 - Pick one neuron at random
 - Update it based on the others
 - Repeat

With one memory, that memory is stable

- Let the input \mathbf{x}_0 be the same as the single memory ξ

$$\begin{aligned}\mathbf{x}_1 &= \varphi(W\mathbf{x}_0) = \varphi\left(\frac{1}{N}(\xi\xi^T - I)\xi\right) \\ &= \varphi\left(\frac{1}{N}\xi(\xi^T\xi - 1)\right) \\ &= \varphi\left(\frac{\|\xi\|^2 - 1}{N}\xi\right) \\ &= \varphi\left(\frac{N-1}{N}\xi\right) = \xi\end{aligned}$$

Therefore the memory is stable

Aside: Hamming distance is the number of differing bits between two patterns

- Hamming distance of 1 from $\{+1, +1, +1\}$
 - $\{+1, +1, -1\}, \{+1, -1, +1\}, \{-1, +1, +1\}$
- Hamming distance of 2 from $\{+1, +1, +1\}$
 - $\{+1, -1, -1\}, \{-1, +1, -1\}, \{-1, -1, +1\}$
- Hamming distance of 3 from $\{+1, +1, +1\}$
 - $\{-1, -1, -1\}$
- For $x_1, x_2 \in \{\pm 1\}^N$, $x_1^T x_2 = N - 2d_H(x_1, x_2)$
 - So $-N \leq x_1^T x_2 \leq N$

With one memory, Hopfield net converges to the closer of ξ or $-\xi$

- For input of \mathbf{x}_0

$$\begin{aligned}\mathbf{x}_1 &= \varphi \left(\frac{1}{N} W \mathbf{x}_0 \right) = \varphi \left(\frac{1}{N} (\xi \xi^T - I) \mathbf{x}_0 \right) \\ &= \varphi \left(\frac{1}{N} (\xi \xi^T \mathbf{x}_0 - \mathbf{x}_0) \right) \\ &= \pm \xi\end{aligned}$$

- Assuming that $|\xi^T \mathbf{x}_0| > 1$
- Closer is measured by inner product
 - Or equivalently in this case, by Hamming distance

Example: Hopfield net with one memory

- Let's use $\xi = [-1, +1, -1]^T$, then

$$W = \frac{1}{3} \begin{bmatrix} 0 & -1 & +1 \\ -1 & 0 & -1 \\ +1 & -1 & 0 \end{bmatrix}$$

- Test memory stability

$$W\xi = \frac{1}{3} \begin{bmatrix} 0 & -1 & +1 \\ -1 & 0 & -1 \\ +1 & -1 & 0 \end{bmatrix} \begin{bmatrix} -1 \\ +1 \\ -1 \end{bmatrix} = \frac{1}{3} \begin{bmatrix} -2 \\ +2 \\ -2 \end{bmatrix}$$

- So $\varphi(W\xi) = \xi$

Example: Hopfield net with one memory

- Follow state trajectory from $\mathbf{x}_1 = [-1, -1, +1]^T$

$$W\mathbf{x}_1 = \frac{1}{3} \begin{bmatrix} 0 & -1 & +1 \\ -1 & 0 & -1 \\ +1 & -1 & 0 \end{bmatrix} \begin{bmatrix} -1 \\ -1 \\ +1 \end{bmatrix} = \frac{1}{3} \begin{bmatrix} 2 \\ 0 \\ 0 \end{bmatrix}$$

- So $\varphi(W\mathbf{x}_1) = [+1, -1, +1]^T = -\xi$

- Follow state trajectory from $\mathbf{x}_2 = [+1, +1, -1]^T$

$$W\mathbf{x}_2 = \frac{1}{3} \begin{bmatrix} 0 & -1 & +1 \\ -1 & 0 & -1 \\ +1 & -1 & 0 \end{bmatrix} \begin{bmatrix} +1 \\ +1 \\ -1 \end{bmatrix} = \frac{1}{3} \begin{bmatrix} -1 \\ 0 \\ 0 \end{bmatrix}$$

- So $\varphi(W\mathbf{x}_2) = [-1, +1, -1]^T = \xi$

For a Hopfield net with multiple memories

- The stability condition for any memory ξ_{ϑ} is

$$\begin{aligned}\xi_{\vartheta} &= \varphi(W \xi_{\vartheta}) \\ &= \varphi \left(\left(\frac{1}{N} \sum_{\mu} \xi_{\mu} \xi_{\mu}^T - I \right) \xi_{\vartheta} \right) \\ &= \varphi \left(\frac{N-1}{N} \xi_{\vartheta} + \underbrace{\frac{1}{N} \sum_{\mu \neq \vartheta} \xi_{\mu} (\xi_{\mu}^T \xi_{\vartheta} - 1)}_{\text{crosstalk}} \right)\end{aligned}$$

Multiple memories can be stored if $M \ll N$

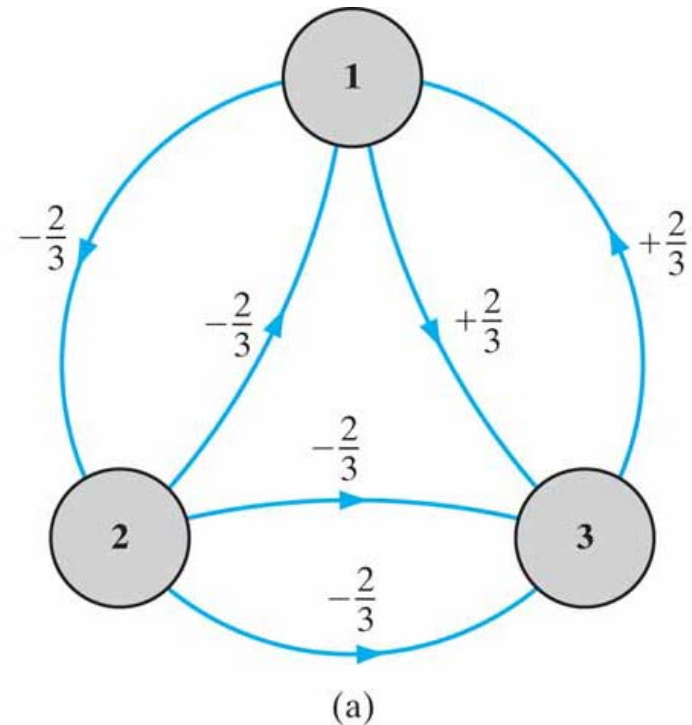
- Crosstalk is a weighted sum of the memories
- If memories are random variables (i.e., uncorrelated with each other)
 - Then this is a sum of $N(M - 1)$ random ± 1 variables
 - Which is asymptotically Gaussian
- If the crosstalk is small, compared to the ξ_{θ} term
 - Then the memory system is stable
 - In general, fewer memories are more likely stable
- More on this shortly

Example 2 from textbook

- Consider the Hopfield network with

$$W = \frac{1}{3} \begin{bmatrix} 0 & -2 & +2 \\ -2 & 0 & -2 \\ +2 & -2 & 0 \end{bmatrix}$$

- 8 possible states
 - See where each goes



Two states are stable

- Two states are stable

- $\xi_1 = [+1, -1, +1]^T$ and $\xi_2 = [-1, +1, -1]^T = -\xi_1$

$$W\xi_1 = \frac{1}{3} \begin{bmatrix} 0 & -2 & +2 \\ -2 & 0 & -2 \\ +2 & -2 & 0 \end{bmatrix} \begin{bmatrix} +1 \\ -1 \\ +1 \end{bmatrix} = \frac{1}{3} \begin{bmatrix} +4 \\ -4 \\ +4 \end{bmatrix}$$

- So $\varphi(W\xi_1) = \xi_1$

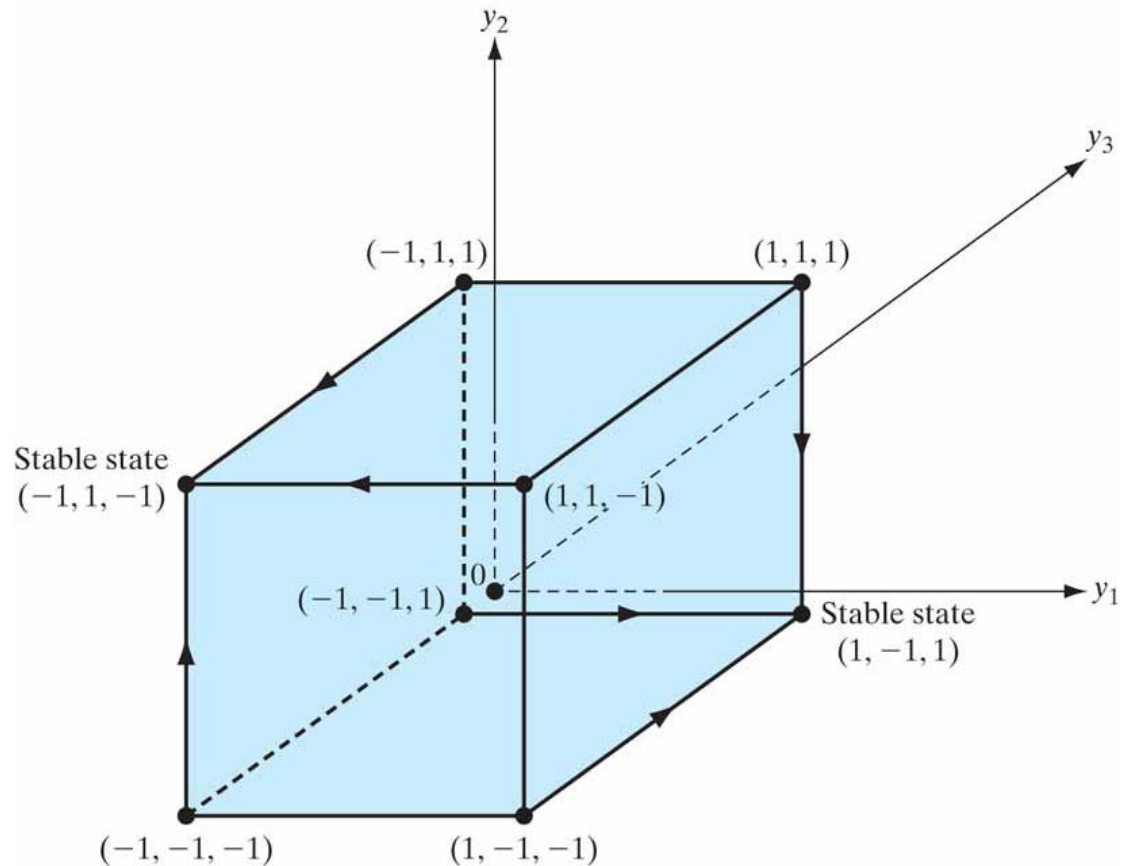
$$W\xi_2 = \frac{1}{3} \begin{bmatrix} 0 & -2 & +2 \\ -2 & 0 & -2 \\ +2 & -2 & 0 \end{bmatrix} \begin{bmatrix} -1 \\ +1 \\ -1 \end{bmatrix} = \frac{1}{3} \begin{bmatrix} -4 \\ +4 \\ -4 \end{bmatrix}$$

- So $\varphi(W\xi_2) = \xi_2$

Weight matrix agrees with the one
calculated from the two stable states

$$\begin{aligned} W &= \frac{1}{3} \xi_1 \xi_1^T + \frac{1}{3} \xi_2 \xi_2^T - \frac{2}{3} I \\ &= \frac{1}{3} \begin{bmatrix} +1 \\ -1 \\ +1 \end{bmatrix} [+1, -1, +1] + \frac{1}{3} \begin{bmatrix} -1 \\ +1 \\ -1 \end{bmatrix} [-1, +1, -1] - \frac{2}{3} I \\ &= \frac{1}{3} \begin{bmatrix} 3 & -1 & +1 \\ -1 & 3 & -1 \\ +1 & -1 & 3 \end{bmatrix} + \frac{1}{3} \begin{bmatrix} 3 & -1 & +1 \\ -1 & 3 & -1 \\ +1 & -1 & 3 \end{bmatrix} - \frac{2}{3} I \\ &= \frac{1}{3} \begin{bmatrix} 0 & -2 & +2 \\ -2 & 0 & -2 \\ +2 & -2 & 0 \end{bmatrix} \end{aligned}$$

Asynchronous updates follow this flow diagram



Memory capacity for a single bit:

Prob of error is defined by amount of cross-talk

- Define

$$C_j^{\vartheta} = -\xi_{\vartheta,j} \sum_i \sum_{\mu \neq \vartheta} \xi_{\mu,j} \xi_{\mu,i} \xi_{\vartheta,i}$$

- Amount cross-talk pushes bit j in the wrong direction

$$\begin{aligned} C_j^{\vartheta} < 0 &\Rightarrow \text{stable} \\ 0 \leq C_j^{\vartheta} < N &\Rightarrow \text{stable} \\ C_j^{\vartheta} > N &\Rightarrow \text{unstable} \end{aligned}$$

Capacity: Crosstalk is approximately Gaussian

- Consider random memories where each element takes $+1$ or -1 with equal probability.
- For random patterns, C_j^{ϑ} is proportional to a sum of $N(M - 1)$ random numbers of $+1$ or -1
- For large NM , it can be approximated by a Gaussian distribution (central limit theorem)
 - With zero mean and variance $\sigma^2 = NM$
- Capacity M_{max} is defined by an error criterion
 - Acceptable level of $P_{error} = \text{Prob}(C_j^{\vartheta} > N)$

Capacity: Prob of error is a function of N/M

- So

$$P_{\text{error}} = \frac{1}{\sqrt{2\pi}\sigma} \int_N^{\infty} \exp\left(-\frac{x^2}{2\sigma^2}\right) dx$$

$$= \frac{1}{2} - \frac{1}{\sqrt{2\pi}\sigma} \int_0^N \exp\left(-\frac{x^2}{2\sigma^2}\right) dx$$

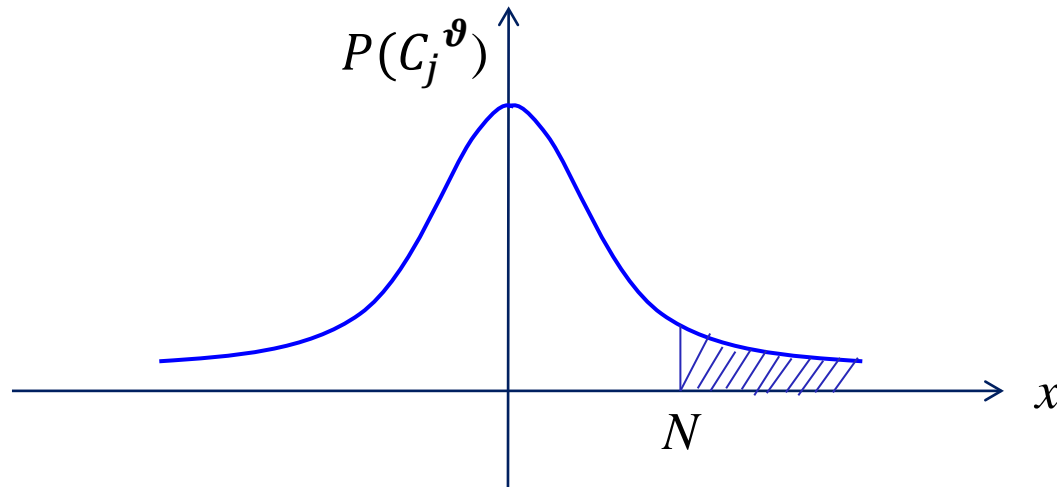
define $\mu = \frac{x}{\sigma\sqrt{2}}$

$$= \frac{1}{2} \left(1 - \underbrace{\frac{2}{\sqrt{\pi}} \int_0^{\sqrt{N/(2M)}} \exp(-\mu^2) d\mu}_{\text{error function}} \right)$$

error function

Capacity: Visualizing prob of error

- So $P_{\text{error}} = \frac{1}{2} \left(1 - \text{erf} \left(\sqrt{\frac{N}{2M}} \right) \right)$



Capacity: Lower error prob requires smaller M

P_{error}	M_{max}/N
0.001	0.105
0.0036	0.138
0.01	0.185
0.05	0.37
0.1	0.61

- So $P_{\text{error}} < 0.01 \Rightarrow M_{\text{max}} = 0.185N$, an upper bound
- Or $0.138N$ just to be safe

To get all N bits correct requires smaller M

- The above analysis is for one bit
- If we want perfect retrieval for ξ^v with prob 0.99
$$(1 - P_{\text{error}})^N > 0.99$$
 - Approximately $P_{\text{error}} < \frac{0.01}{N}$
- For this case $M_{\text{max}} = \frac{N}{2 \log N}$
 - See (McEliece, Posner, Rodemich, and Venkatesh, 1987)
- This is a bit disappointing compared to various error correction codes

Non-random memories modify capacity

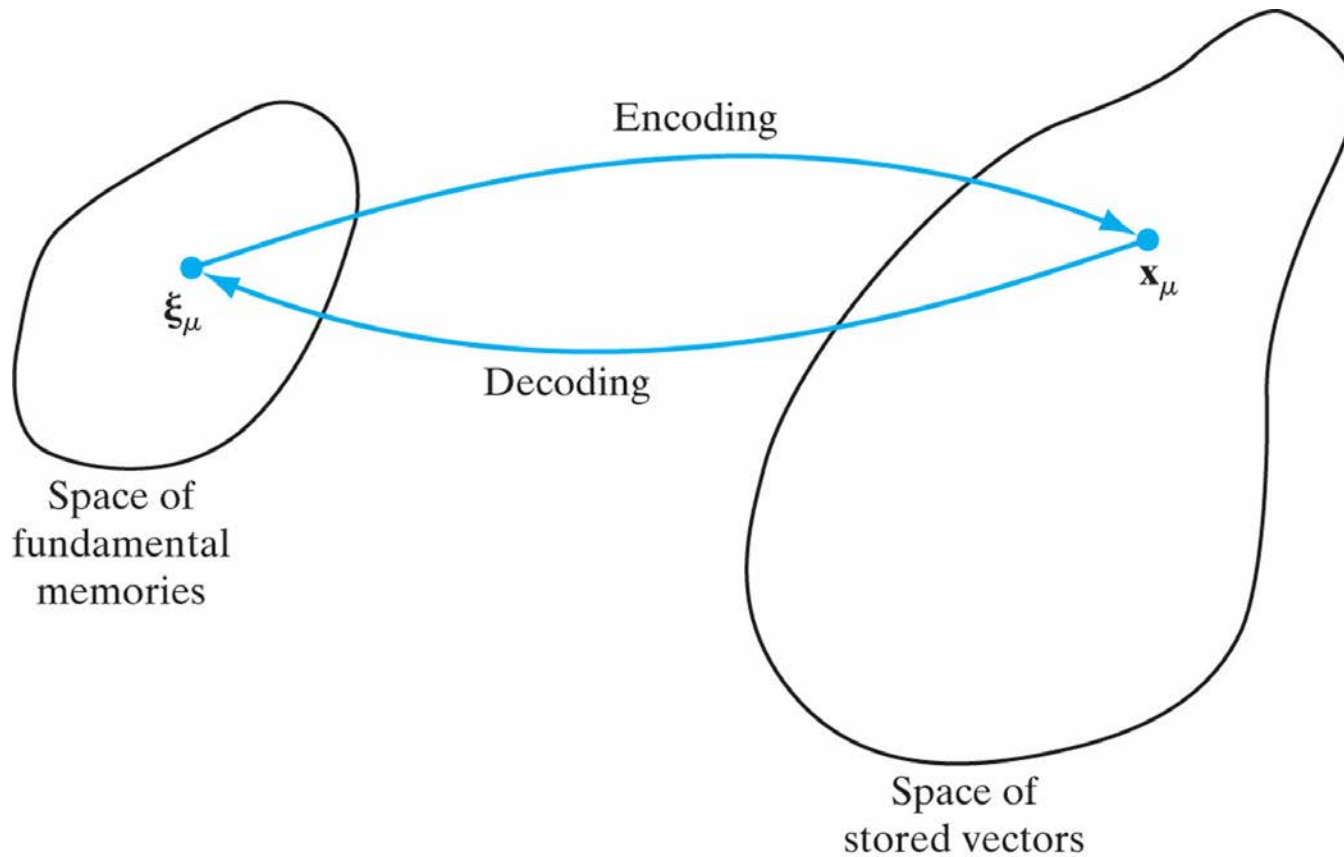
- Real patterns are not random
 - Although they could be encoded to be almost random
 - The capacity is worse for correlated patterns
- At the favorable extreme, for orthogonal memories

$$\sum_i \xi_{\mu,i} \xi_{\vartheta,i} = 0 \quad \text{for } \vartheta \neq \mu$$

then $C_j^{\vartheta} = 0$ and $M_{\max} = N$

- This is the maximum number of orthogonal patterns
- Use fewer memories to allow some evolution, otherwise, why bother?

Coding illustration



Energy function (Lyapunov function)

- The existence of an energy (Lyapunov) function for a dynamical system ensures its stability
- The energy function for the Hopfield net is

$$E(\mathbf{x}) = -\frac{1}{2} \sum_i \sum_j w_{ji} x_i x_j = -\frac{1}{2} \mathbf{x}^T W \mathbf{x}$$

- **Theorem:** Given symmetric weights, $w_{ji} = w_{ij}$, the energy function does not increase as the Hopfield net evolves asynchronously

Energy function (cont.)

- Let x_j' be the new value of x_j after an update

$$x_j' = \varphi \left(\sum_i w_{ji} x_i \right)$$

- If $x_j' = x_j$, E remains the same

Energy function (cont.)

- Otherwise, $x'_j = -x_j$:

- Let s be a vector of 1s except for $s_j = -1$

$$\begin{aligned} E(x') - E(x) &= -\frac{1}{2} \sum_i \sum_k w_{ki} x_i x_k s_i s_k + \frac{1}{2} \sum_i \sum_k w_{ki} x_i x_k \\ &= -\sum_{i \neq j} w_{ji} x_i x_j s_j + \sum_{k \neq j} w_{kj} x_j x_k s_j \end{aligned}$$

since $W = W^T$

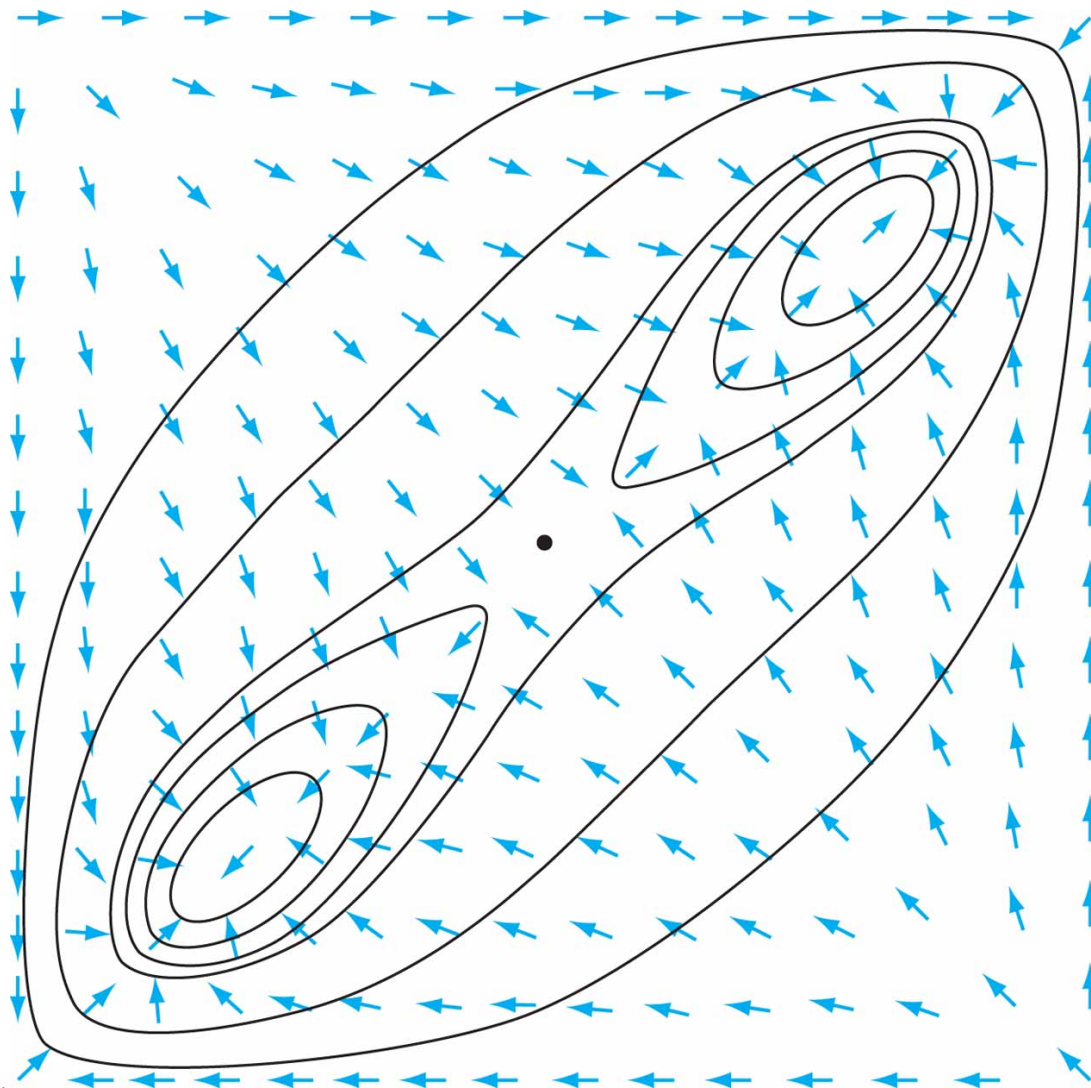
$$\begin{aligned} &= -2x_j s_j \sum_{i \neq j} w_{ji} x_i \\ &= 2x_j \sum_{i \neq j} w_{ji} x_i < 0 \end{aligned}$$

different signs
by assumption

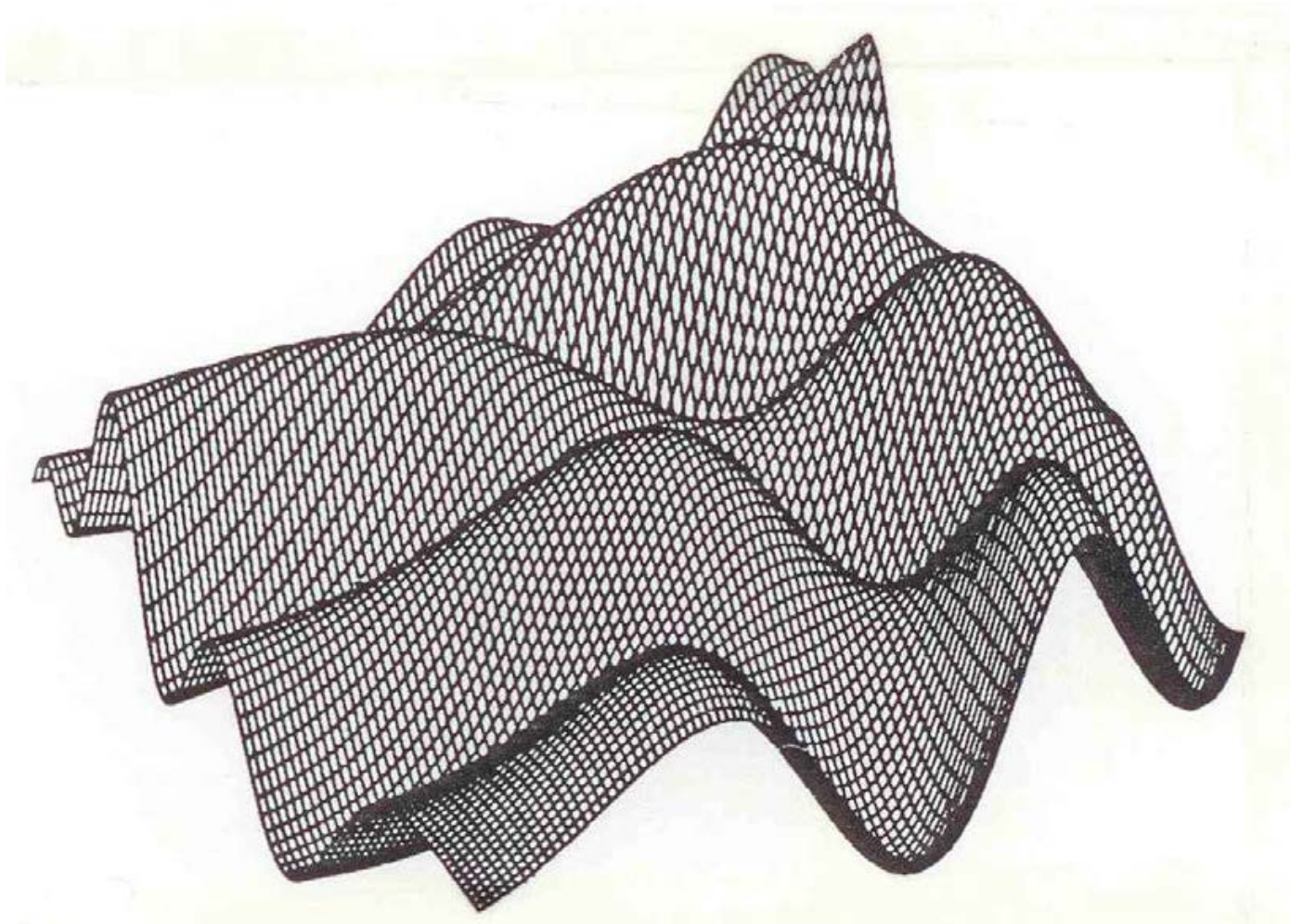
Energy function (cont.)

- Thus, $E(\mathbf{x})$ decreases every time x_j flips. Since E is bounded, the Hopfield net is always stable
- **Remarks:**
 - Useful concepts from dynamical systems: attractors, basins of attraction, energy (Lyapunov) surface or landscape

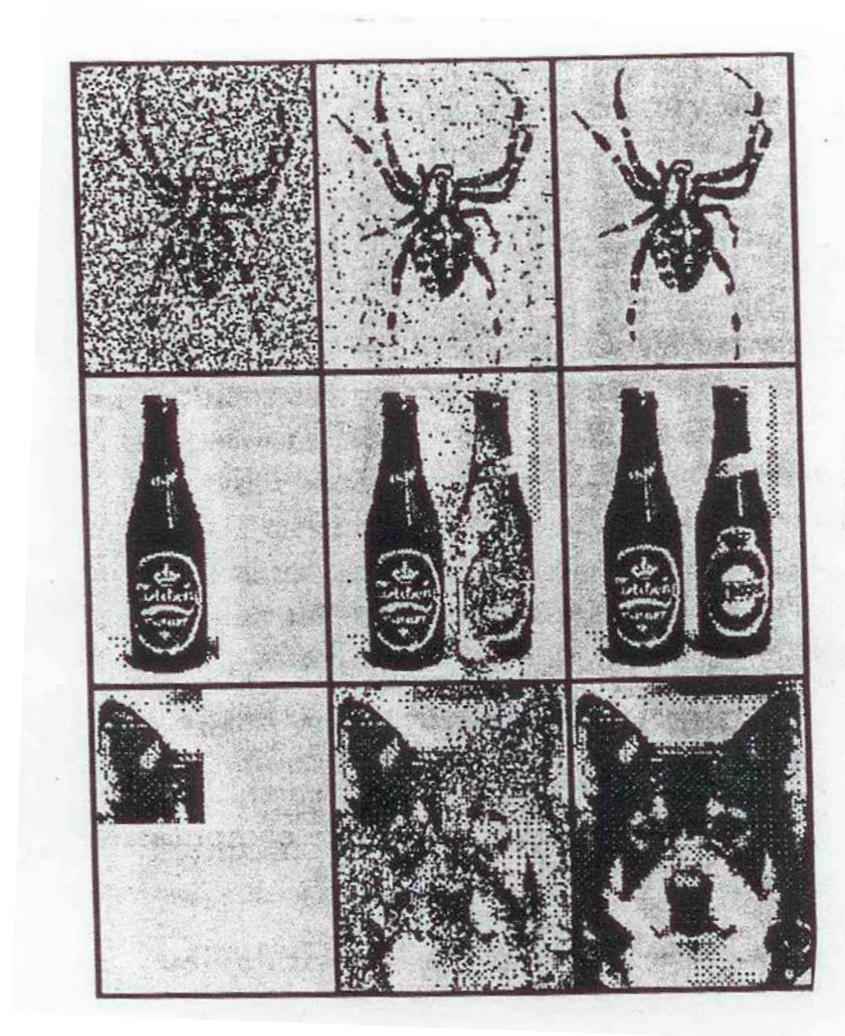
Energy contour map



2-D energy surface



Memory recall illustration



Hertz, Krogh, and
Palmer (1991), Ch 2

Remarks (cont.)

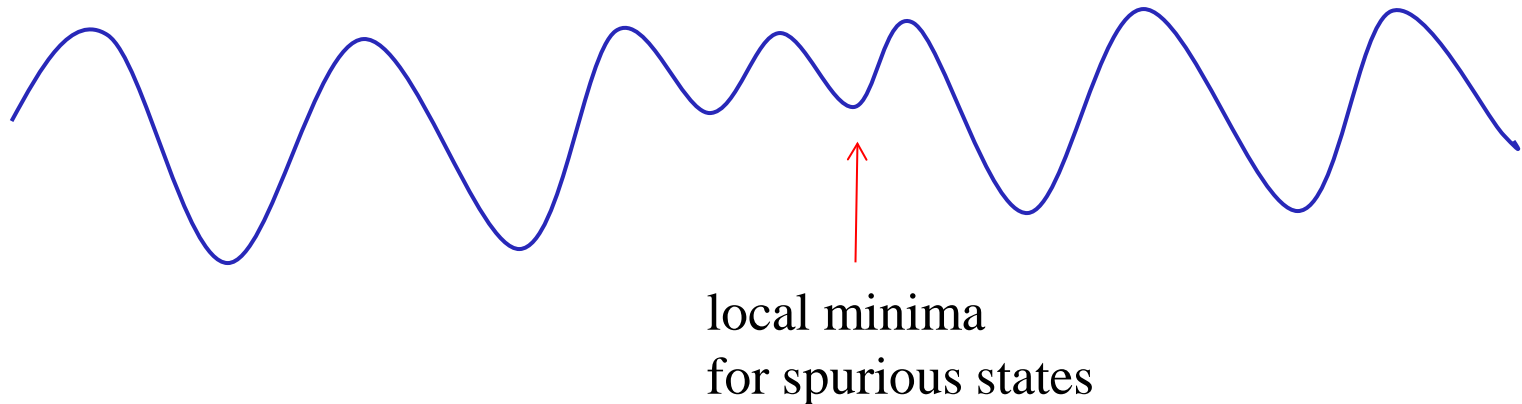
- Bipolar neurons can be extended to continuous-valued neurons by using hyperbolic tangent activation function, and discrete update can be extended to continuous-time dynamics (good for analog VLSI implementation)
- The concept of energy minimization has been applied to optimization problems (neural optimization)

Spurious states

- Not all local minima (stable states) correspond to fundamental memories.
- Other attractors:
 - $-\xi_{\mu}$
 - linear combination of odd number of memories
 - other uncorrelated patterns
- Such attractors are called spurious states

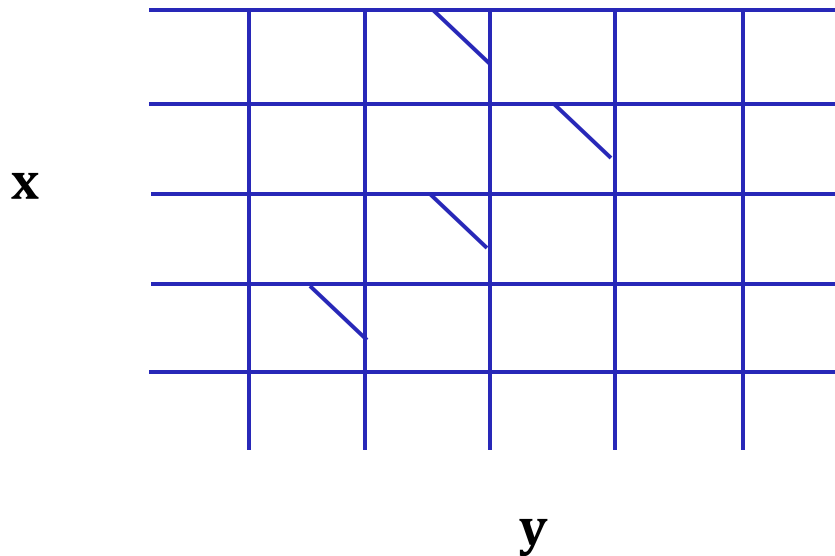
Spurious states (cont.)

- Spurious states tend to have smaller basins and occur higher on the energy surface



Kinds of associative memory

- { Autoassociative (e.g. Hopfield net)
Heteroassociative: store pairs $\langle x_\mu, y_\mu \rangle$ explicitly



matrix memory
(Anderson 1972)

holographic memory
(van Heerden, 1963)