

WiSeDB: A Learning-based Workload Management Advisor for Cloud Databases

NTHU DB / AI Bootcamp

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Basic information of this paper

- Issue : 2016. 6. 10
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Outline

- Cloud DBMS
- About the problem
 - workload management
 - notations
 - problem definition
- Approach
 - overview
 - graph construction
 - feature extraction

Recalling the Definition of Cloud DBMS

- A ***cloud DBMS*** is a DBMS designed to run in the cloud and that could ensure ***SAE***:
 - ***high Scalability***
 - ***high Availability***
 - ***Elasticity***
- **Workload Management**



Workload Management

- What is workload management?
 - Provisioning resource
 - Assigning incoming queries to provisioned VMs
 - Query scheduling within a VM
- Specific by application workload characteristics and performance goals

Notations Definition

- Query templates $\mathbf{T} = \{ T1, T2, \dots \}$
- Workload $\mathbf{Q} = \{ q_1^x, q_2^y, \dots \}$
- $\mathbf{VM}_i = \{ q_1^x, q_2^y, \dots \}$
- Schedule $\mathbf{S} = \{ \mathbf{VM}_i, \mathbf{VM}_j, \dots \}$
- Fix VM start-up cost : f_s ; Running cost per unit of time : f_r
- Latency of a query q_j^x run on $\mathbf{VM}_i = l(q_j^x, i)$.

Problem Definition

- Given performance goal **R** and query templates **T**, we try to find a S that minimize the cost:

$$cost(R, S) = \sum_{vm_j^i \in S} \left[f_s^i + \sum_{q_k^m \in vm_j^i} f_r^i \times l(q_k^m, i) \right] + p(R, S) \quad (1)$$

- Where **p(R,S)** is penalty function for queries didn't meet up with given deadline.
- NP-Hard problem

How?

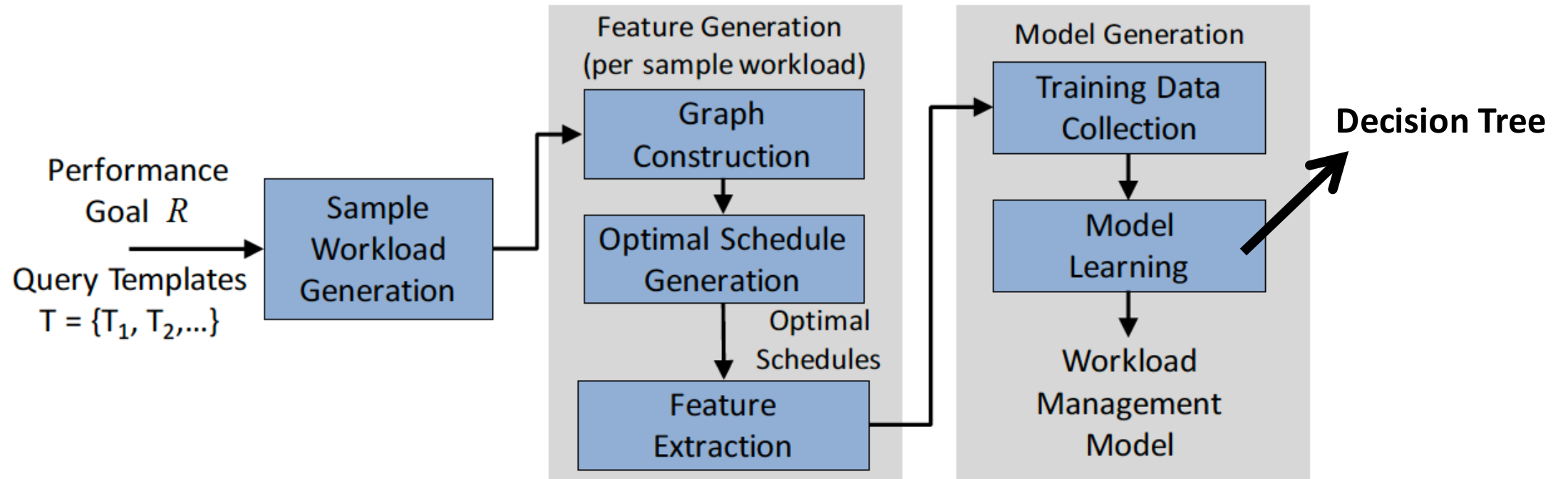
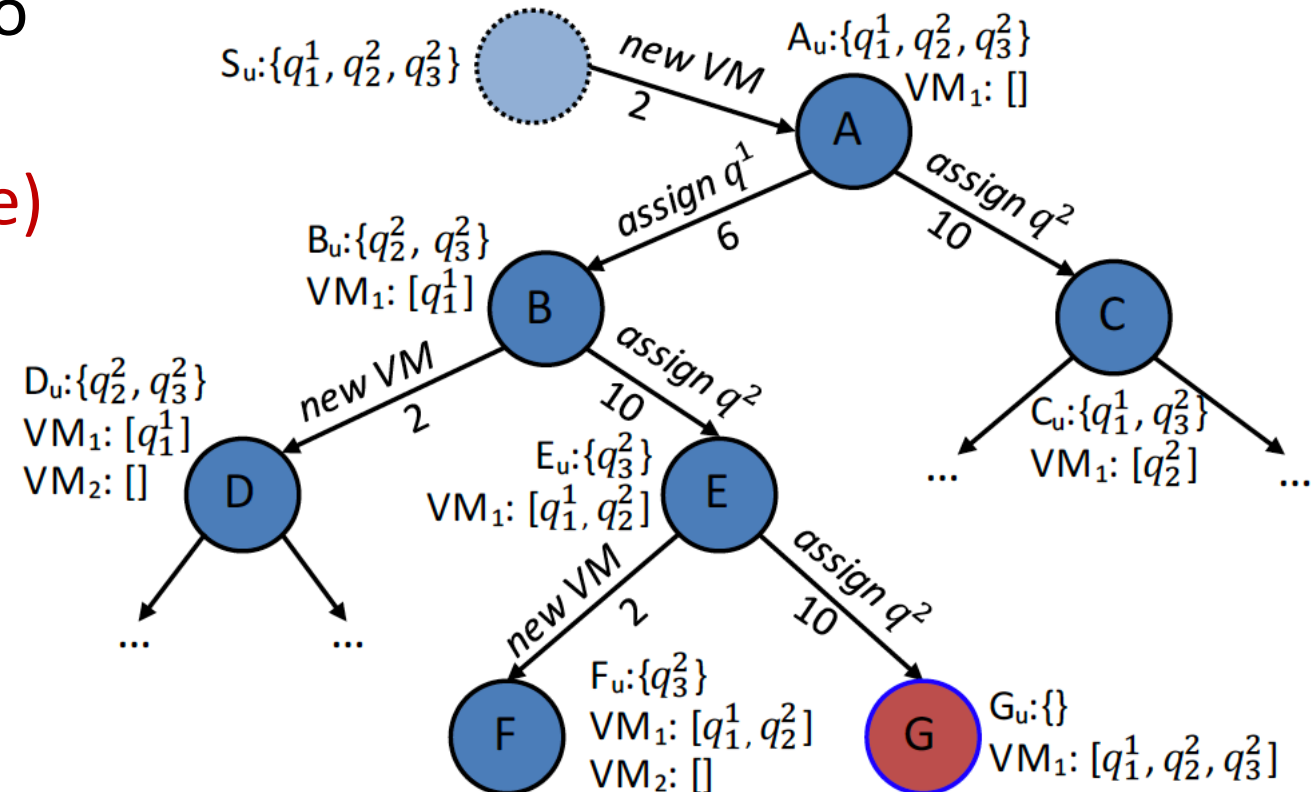


Figure 3: Generation of the decision model

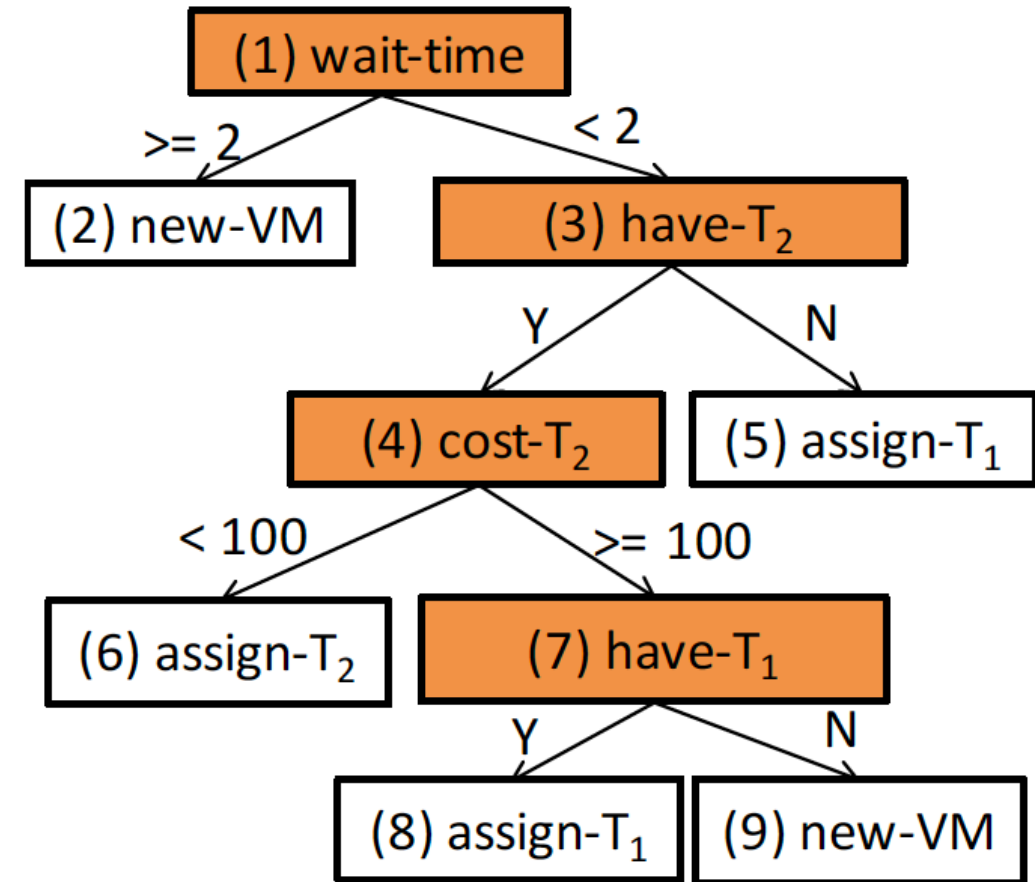
Graph Construction

- Every node of the graph represents a scheduling step.
- To simplify question, we have two choices in every step:
 - Renting a new VM (start-up edge)
 - Assign a query to the newest VM (replacement edge)



Feature Extraction

- Edge selection for a vertex u only depends on the unassigned queries u_u and the scheduled u_s so far.
- Features :
 - wait – time
 - proportion – of – x
 - cost – of – x
 - have - x



Wrap up

- input :
 - query templates T
 - Performance goal R
- output : a schedule S that could minimize cost function

$$cost(R, S) = \sum_{vm_j^i \in S} \left[f_s^i + \sum_{q_k^m \in vm_j^i} f_r^i \times l(q_k^m, i) \right] + p(R, S) \quad (1)$$