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Not Sealed: Practical Attacks on Nostr, a Decentralized Censorship-Resistant Protocol

Keywords: Distributed SNS, signature verification bypass, CBC mode malleability, cache poisoning, plaintext recovery

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Contributors: Ryoma Ito, Kazuhiko Minematsu, Shogo Shiraki and Takanori Isobe

(Also, IEEE EuroS&P2025)

Our Team



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The dawn of the Distributed SNS

FORBES DIGITAL ASSETS

Jack Dorsey Backs Ocean In Shifting Toward Decentralized Bitcoin Mining

By [Susie Violet Ward](#), Contributor. © Bitcoin journalist and financial analyst b...

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Social media platform has become a 'refuge' from the far-right activism on X, experts say, after Elon Musk teamed up with Donald Trump

BITCOIN

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May 16, 2025 — 04:56 pm EDT

Written by [Frank Corva](#) for Bitcoin Magazine →

SOCIAL

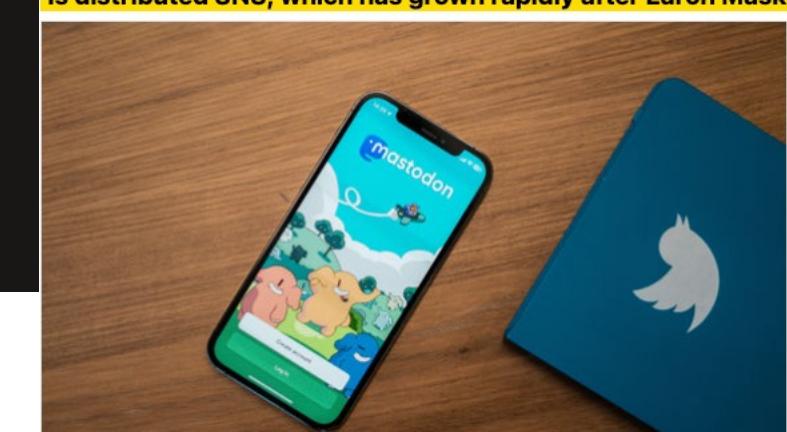
Jack Dorsey pumps \$10M into a nonprofit focused on open source social media

Sarah Perez — 9:01 PM PDT · July 16, 2025

Threads Surpasses 350M Monthly Users: A Growth Milestone with Mixed Monetization Prospects

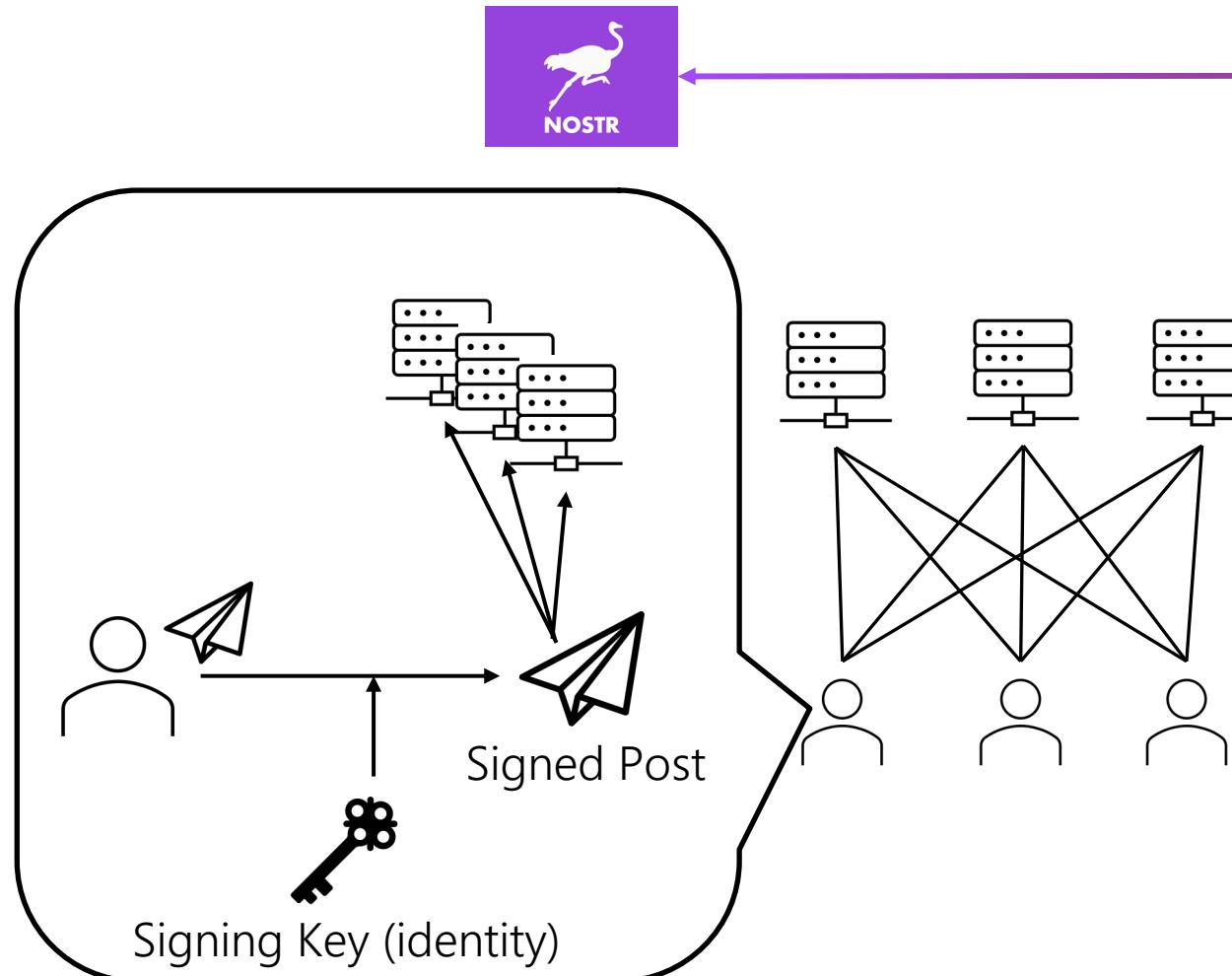
Nathaniel Stone • Wednesday, Apr 30, 2025 6:54 pm ET

4 min read



Distributed SNS

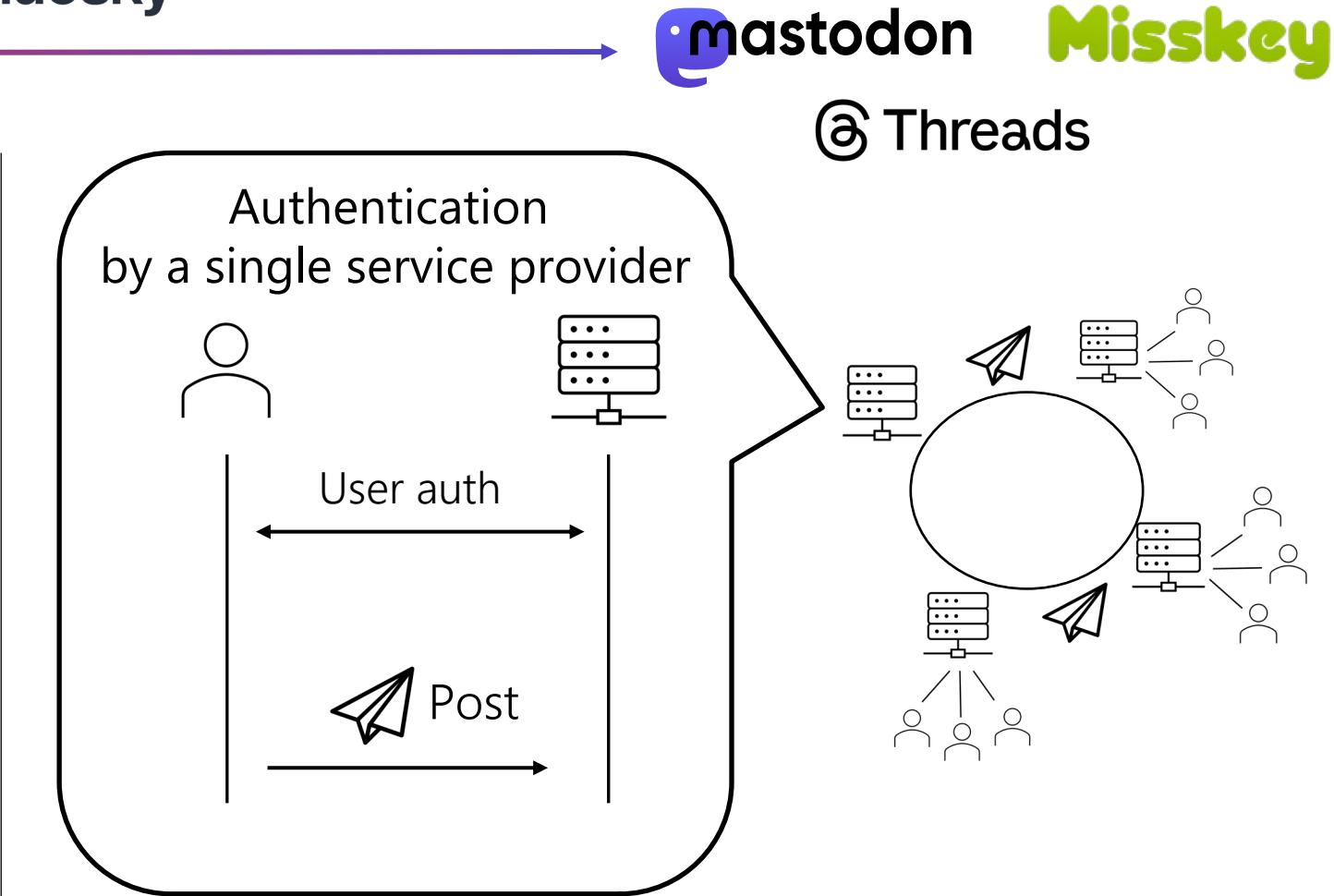
Self-sovereign



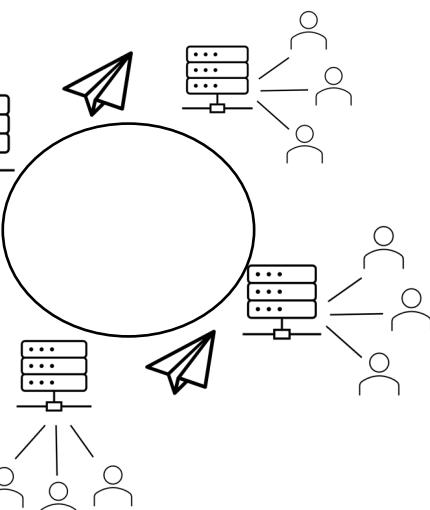
Service providers are independent
User's identity is managed by user



Federated

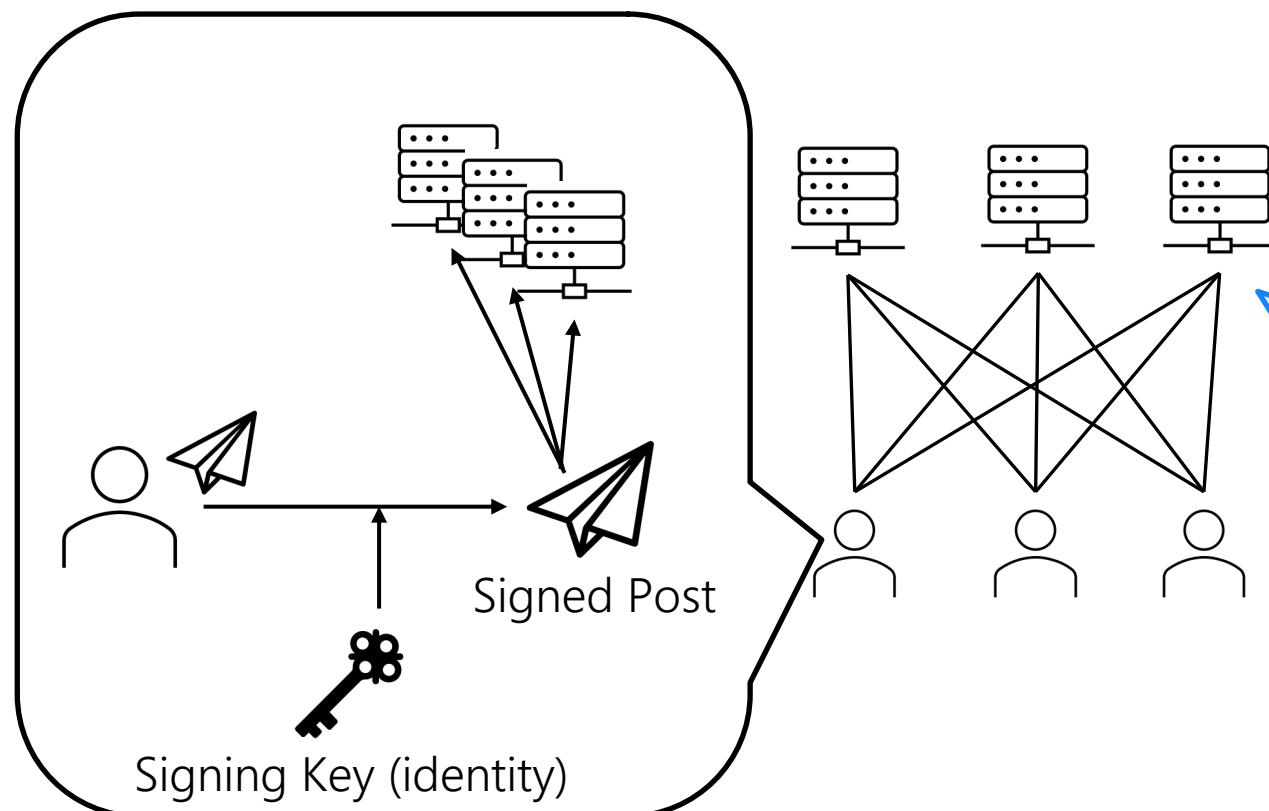


Service providers are interconnected
But identity managed like a **centralized SNS**



Distributed SNS

Self-sovereign



Service providers are independent
User's identity is managed by user

Quite different architecture from
traditional centralized SNS / messaging

Research Questions

- How to trust public keys?
- New architecture, new attack surface?



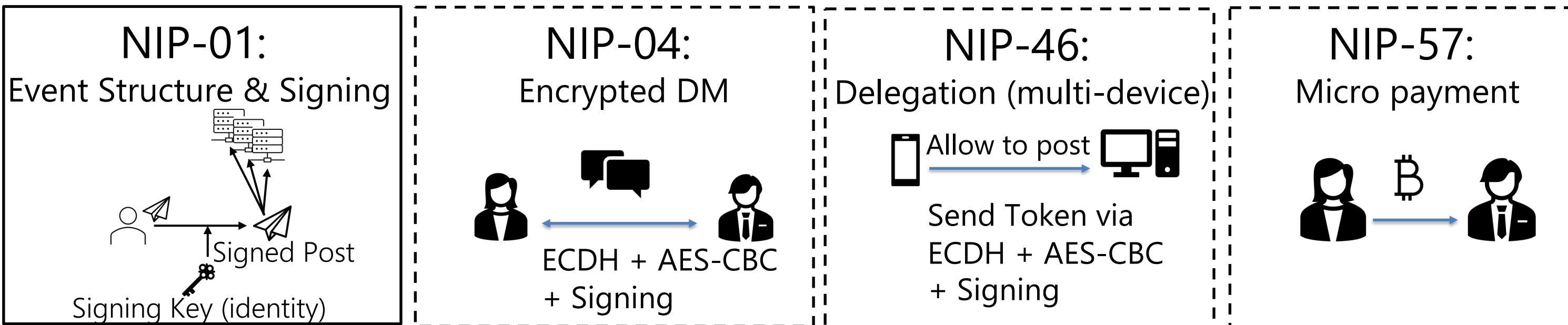
Service providers are interconnected
But identity managed like a centralized SNS

What is Nostr?

- **Open, censorship-resistant social-network**
- **1.1 million registration users**
- **No centralized authority, users must manage Public-key-based identities**
 - A secp256k1 key pair defines who you are; every post carries a signature
- **Zero barriers to participation**
 - Anyone can run a relay server or client
 - Covers most of the attractive features of centralized SNS
 - E.g., Post, Profile, Encrypted DMs, Micro payment, Multiple device sign-in

Cryptography in Nostr Specs

- NIP = Nostr implementation possibilities
- 56+ specifications
- 1 mandatory protocol & 55+ optional protocols
- 4 key feature protocols



Our Contributions

- Analyze 56 specs
- Analyze 9 implementations
- Find 7 vulnerabilities on 4 key features
- Implement 8 attacks
- Breaking confidentiality, integrity, availability
- Propose mitigation
- Two years of persistent disclosure process



First Comprehensive Analysis



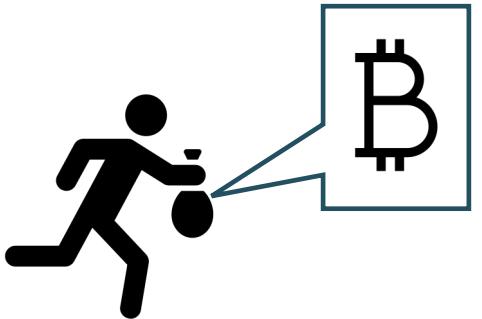
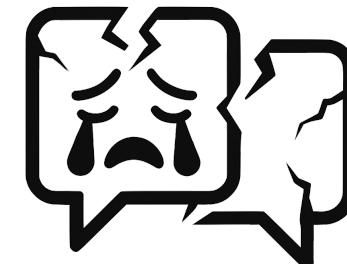
Practical Attacks
& PoCs



Mitigation
& Responsible Disclosure

Our findings

- Breaking integrity on All items (e.g., Profile, Contact List Encrypted DMs...)
- Impersonating to another user
- Breaking confidentiality on Encrypted DMs
- Hijacking micro payment (subset of impersonating)

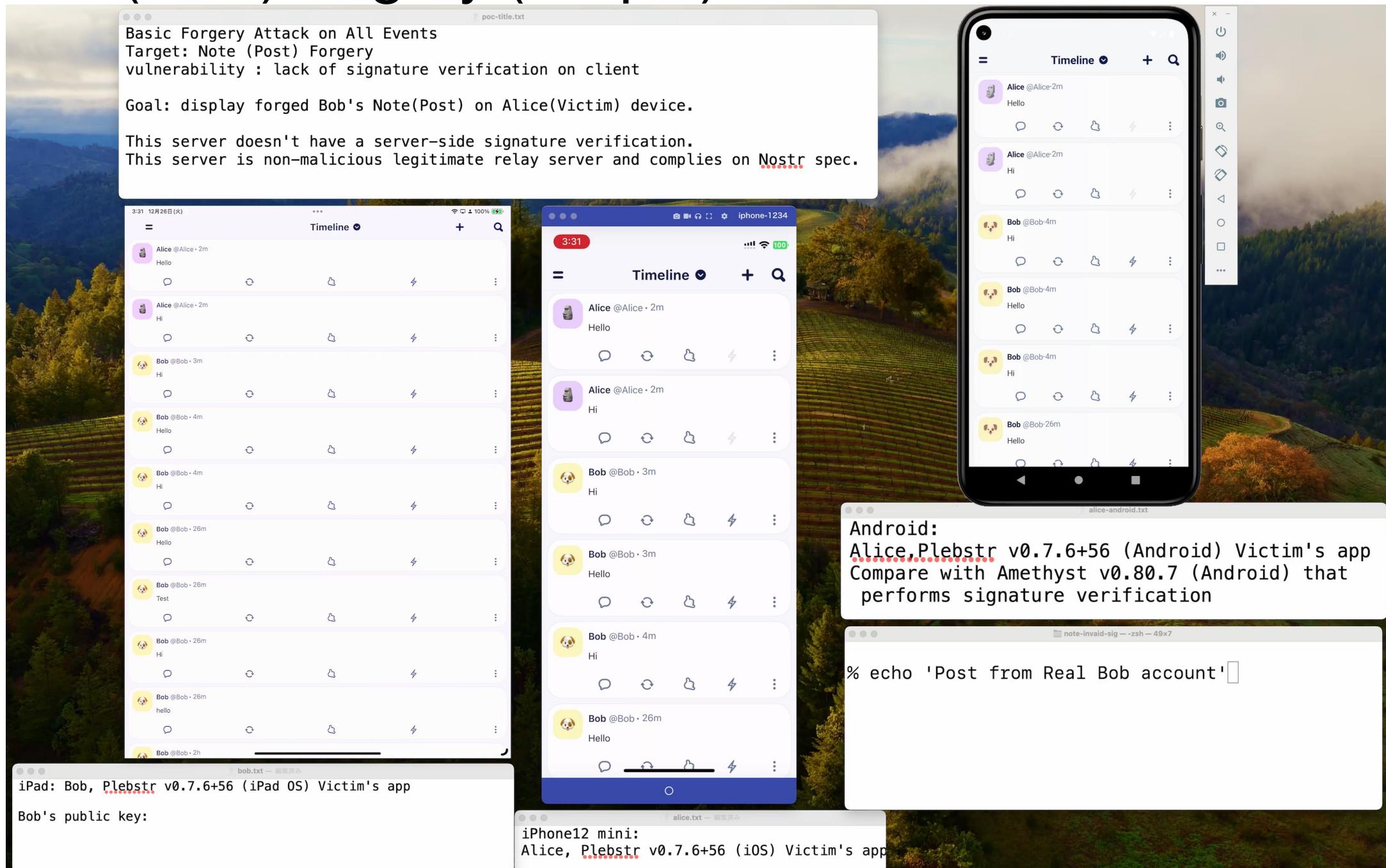


These are not theoretical flaws—they enable practical exploitation

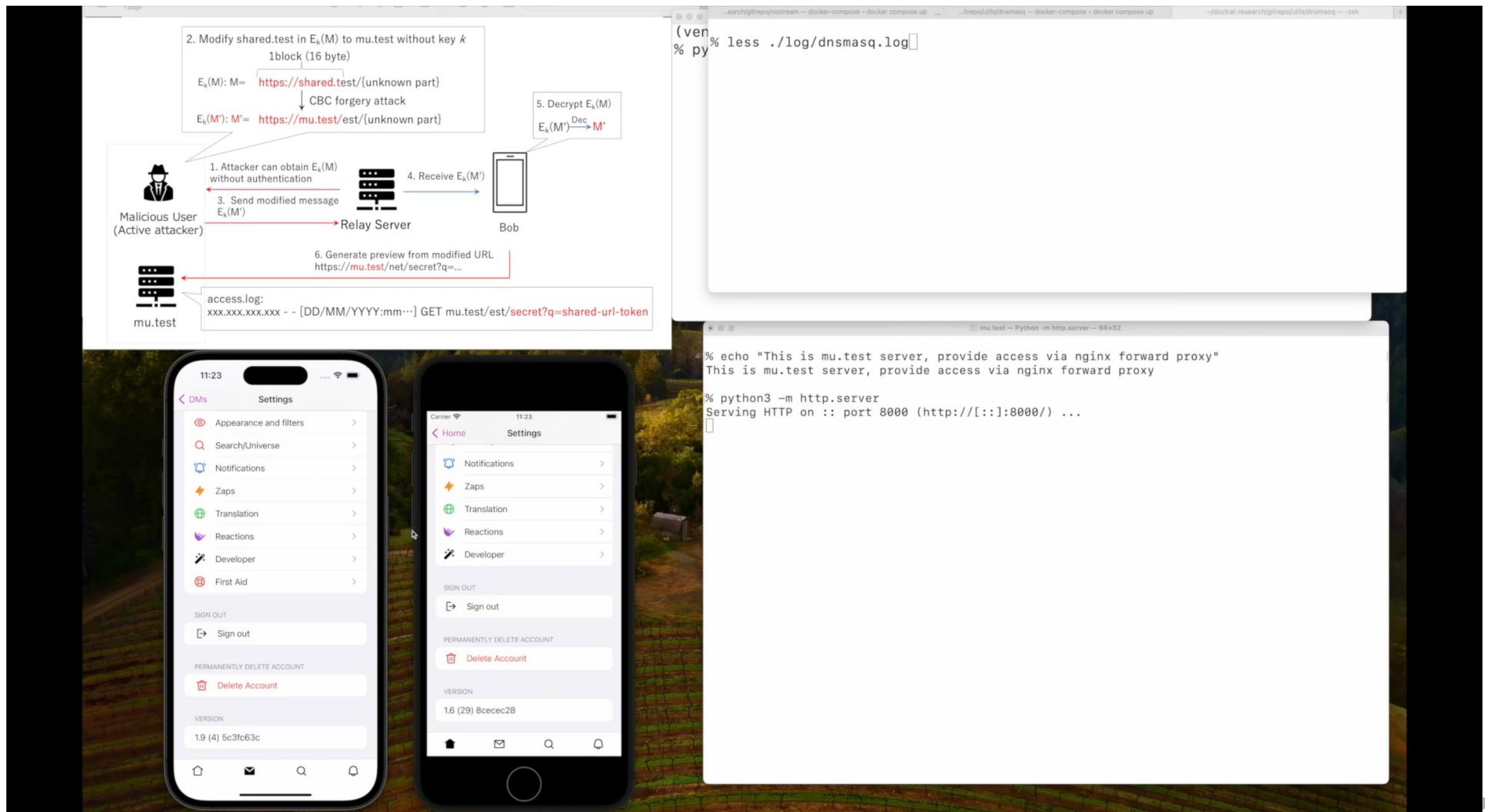
The required threat model varies

Some attacks assume a malicious user; others work under a malicious relay server

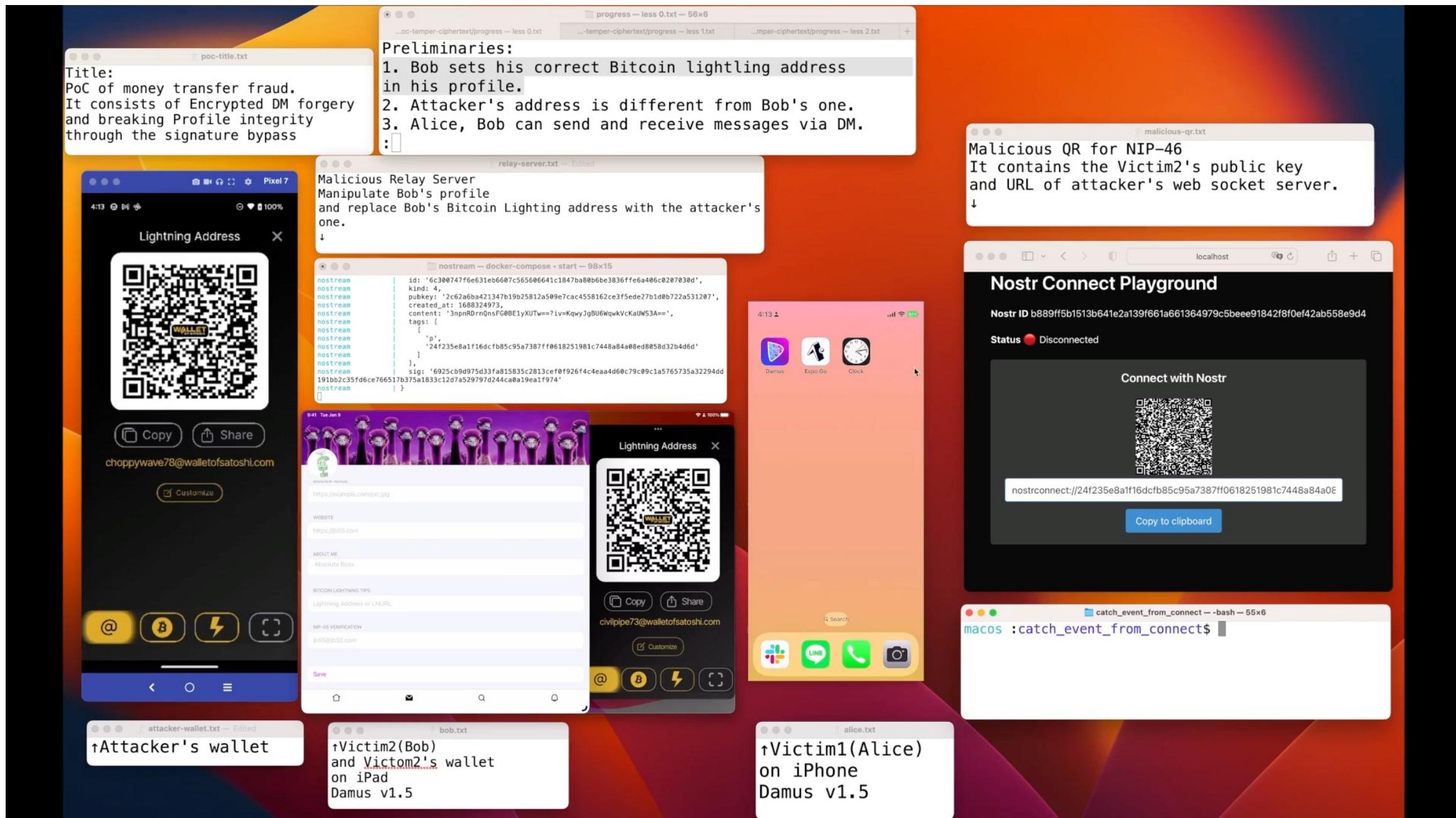
PoC: Note (Post) forgery (simple)



PoC: Encrypted DMs forgery & URL recovery



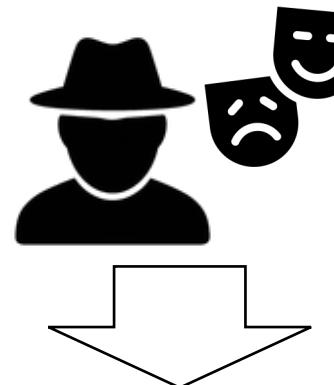
PoC: Hijacking micro payment ← Profile forgery (cache) & DMs forgery



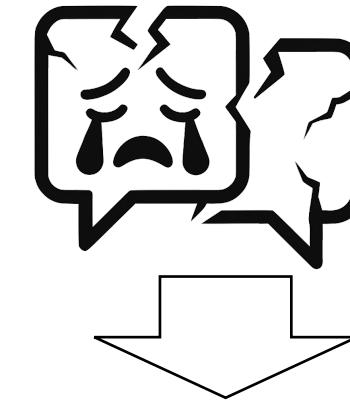
Why does it happen?

Cryptographic protocol design flaw + Implementation flaw

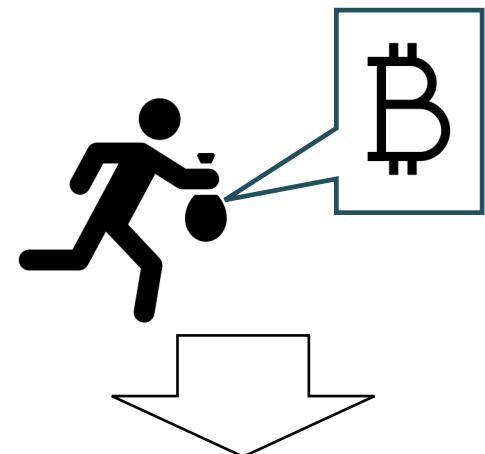
- Breaking integrity on All items
- Impersonating to another user
- Breaking confidentiality on Encrypted DMs
- Hijacking micro payment (subset of impersonating)



- Signature verification Bypass



- Lack of key separation
- Receiver-side preview generation



- Verification Bypass

Why does it happen?

Cryptographic protocol design flaw + Implementation flaw

Step by step attack tracing

Breaking...

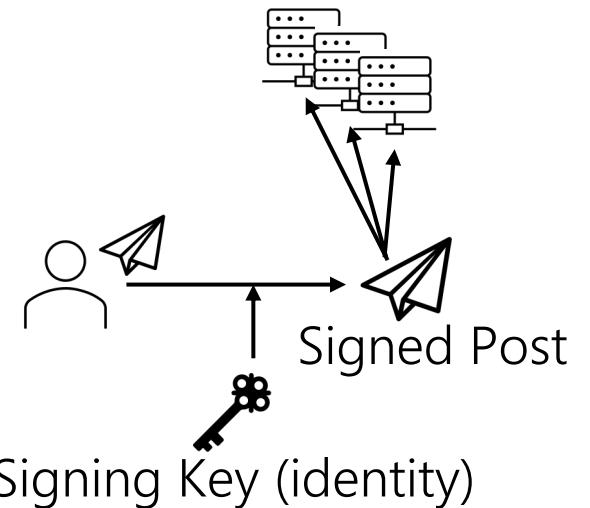
Plaintext integrity
(simple / cache poisoning)

Ciphertext integrity

Ciphertext confidentiality

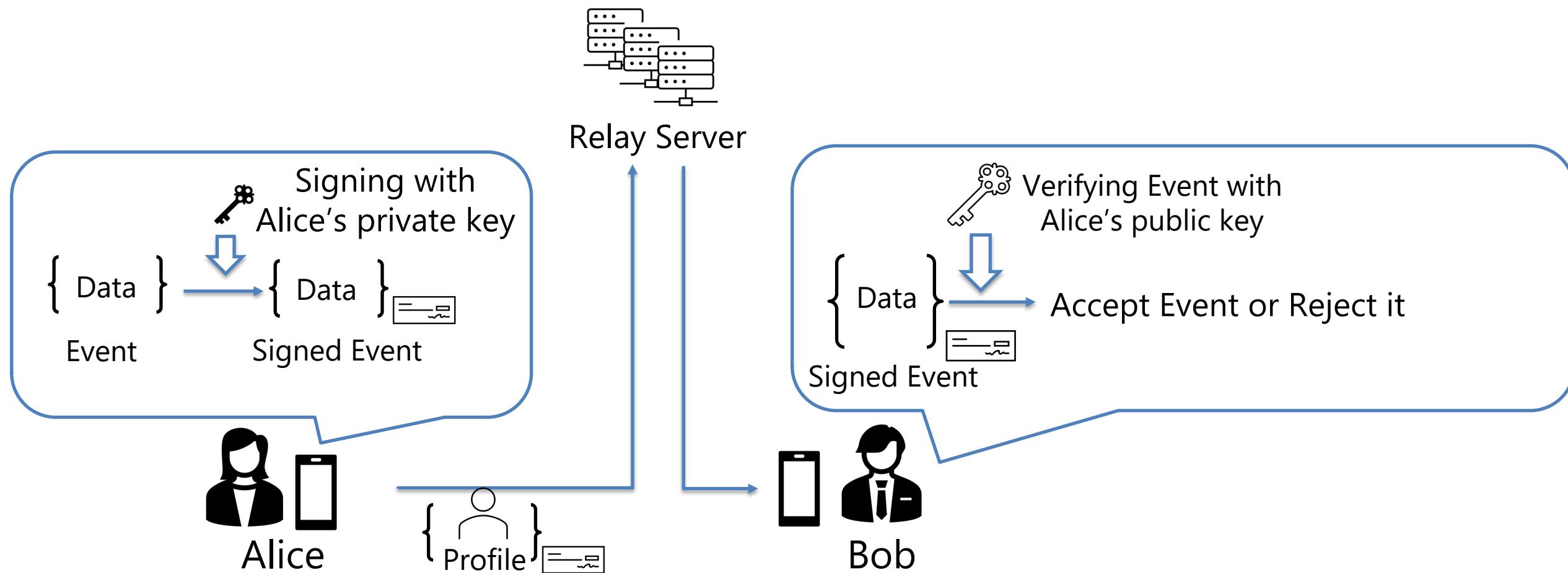
Remark: mandatory signing specification
(simplified)

NIP-01:
Event Structure & Signing



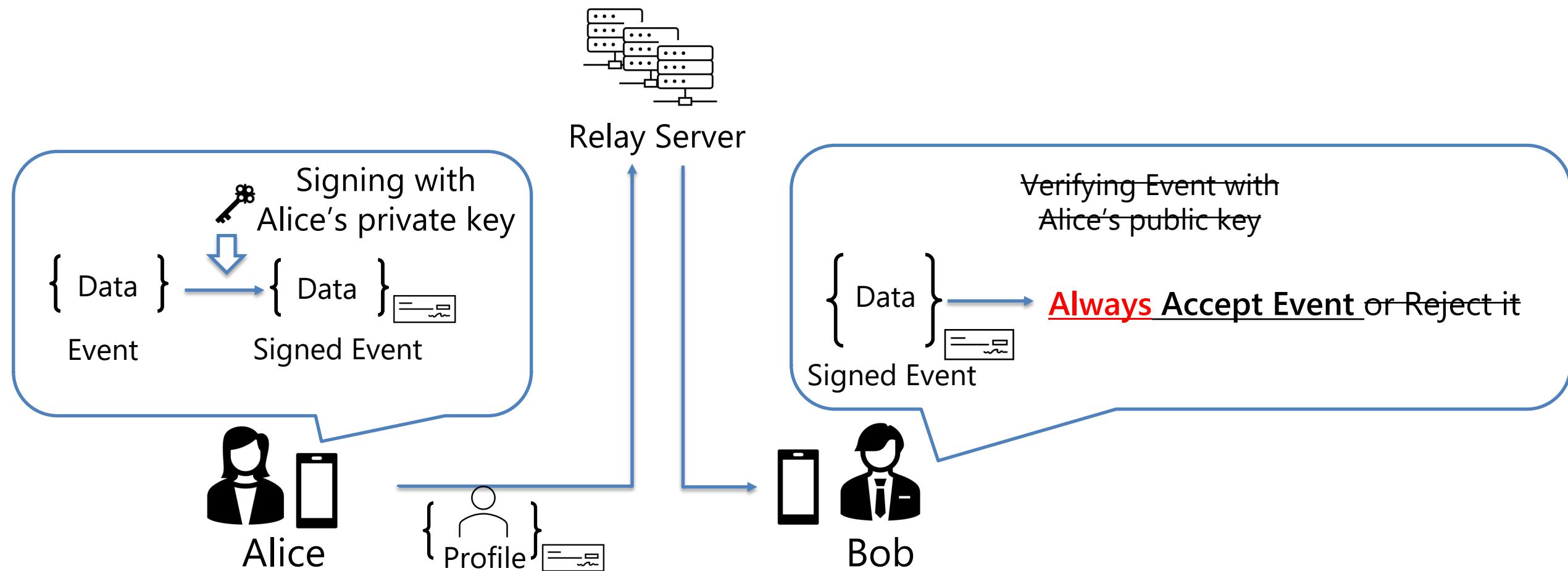
Breaking...
Plaintext integrity
(simple)

Remark: mandatory signing specification
(details depending on specification)



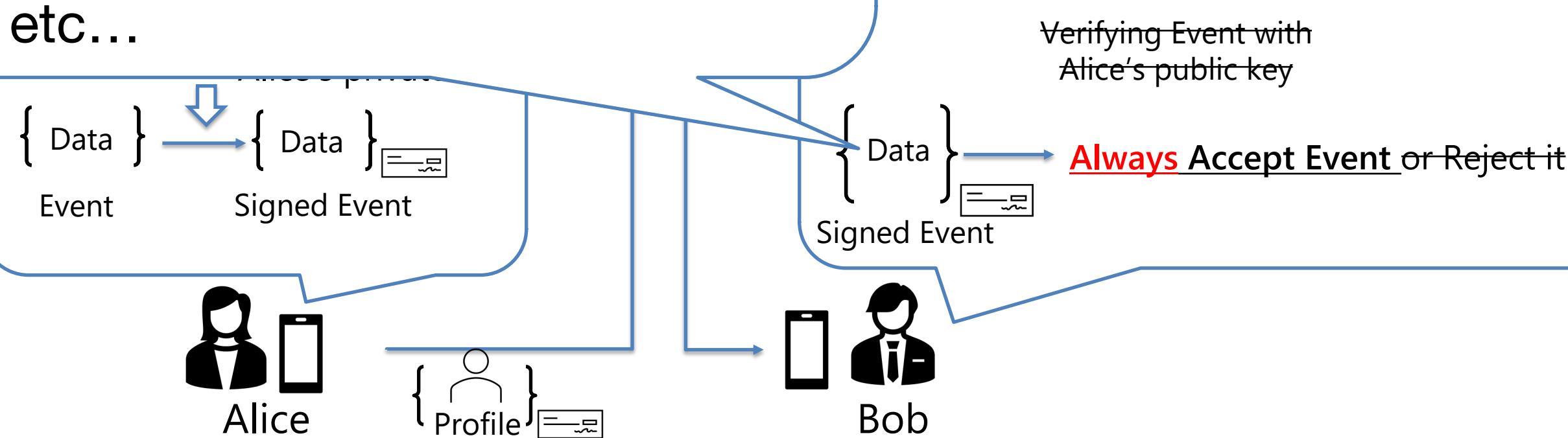
Breaking...
Plaintext integrity
(simple)

Details depending on many **actual** implementations



Why actual implementations

Event Type	Data
Profile	Name, Bio, BTC address
Encrypted DM	Encrypted Msg
Post	Plaintext Msg
etc...	



Breaking...
Plaintext integrity
(simple)

Details depending on many **actual** implementations

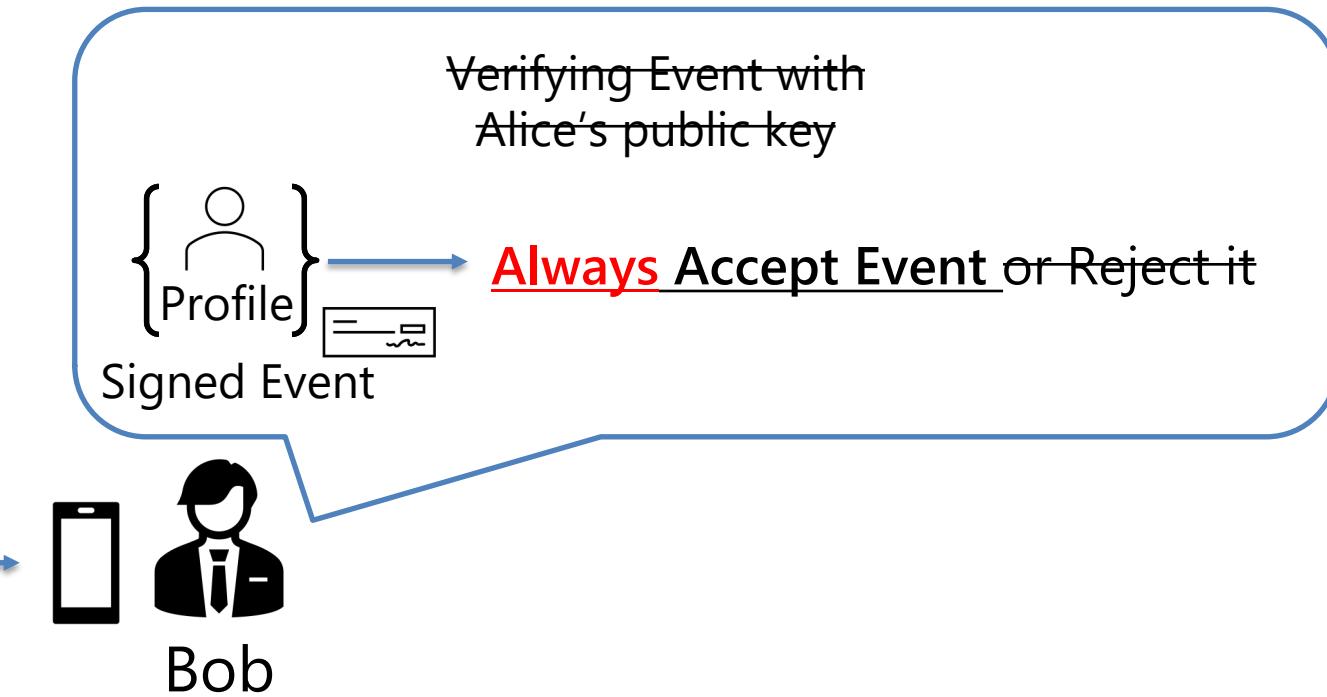
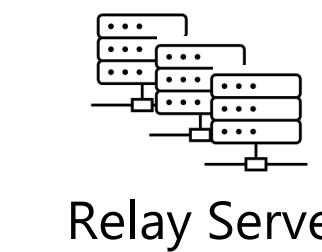
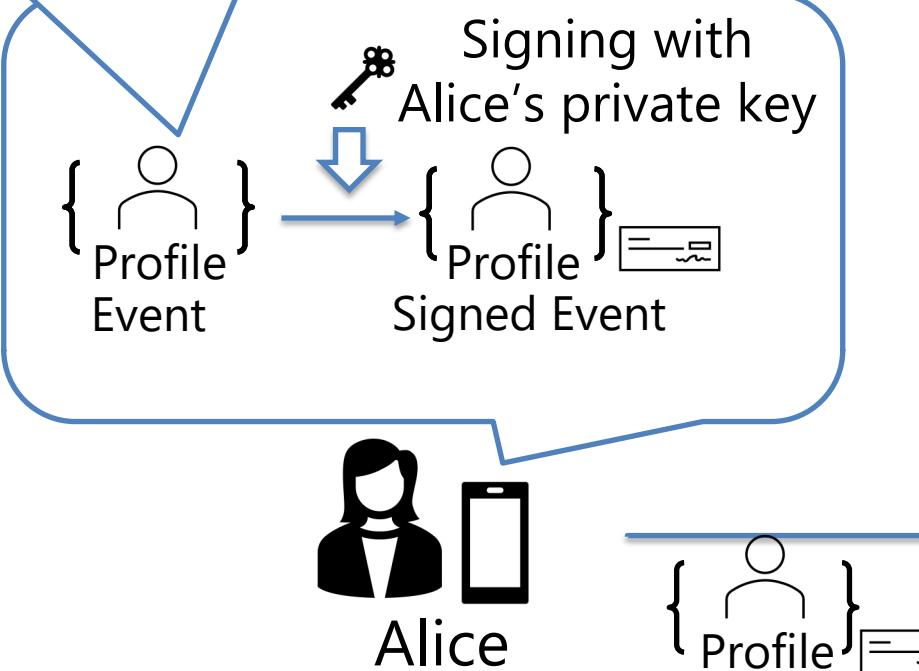
```
541
542     func handle_text_event(sub_id: String, _ ev: NostrEvent) {
543         guard should_show_event(contacts: damus_state.contacts, ev: ev) else {
544             return
545         }
546
547         process_image_metadata(cache: damus_state.events, ev: ev)
548         damus_state.replies.count_replies(ev)
549         damus_state.events.insert(ev)
550
551         if sub_id == home_subid {
552             insert_home_event(ev)
553         } else if sub_id == notifications_subid {
554             handle_notification(ev: ev)
555         }
556     }
557 }
```

There is no Verify(Sig) call
in the event handling!

Breaking...
Plaintext integrity
(simple)

Case : Alice publishes her Profile & Bob subscribes it

- Alice's display name
- Alice's bio
- Alice's Bitcoin(sat) address etc...



Verifying Event with
Alice's public key

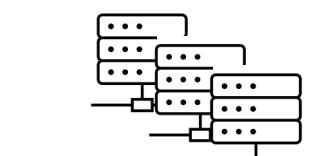
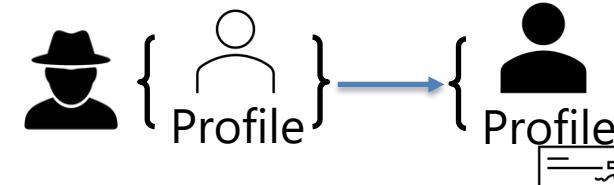
Always Accept Event or Reject it

Breaking...
Plaintext integrity
(simple)

Profile Forgery on Plebstr, FreeFrom Attacker also can publish Alice's Profile

- Alice's display name
- Alice's modified bio
- Attacker's Bitcoin(sat) address

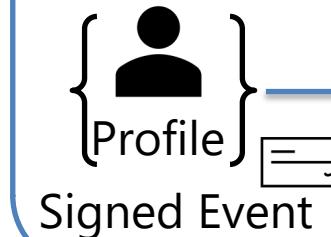
Copy Alice's Event and modified it



Relay Server



Verifying Event with
Alice's public key



Signed Event

Always Accept Event or Reject it

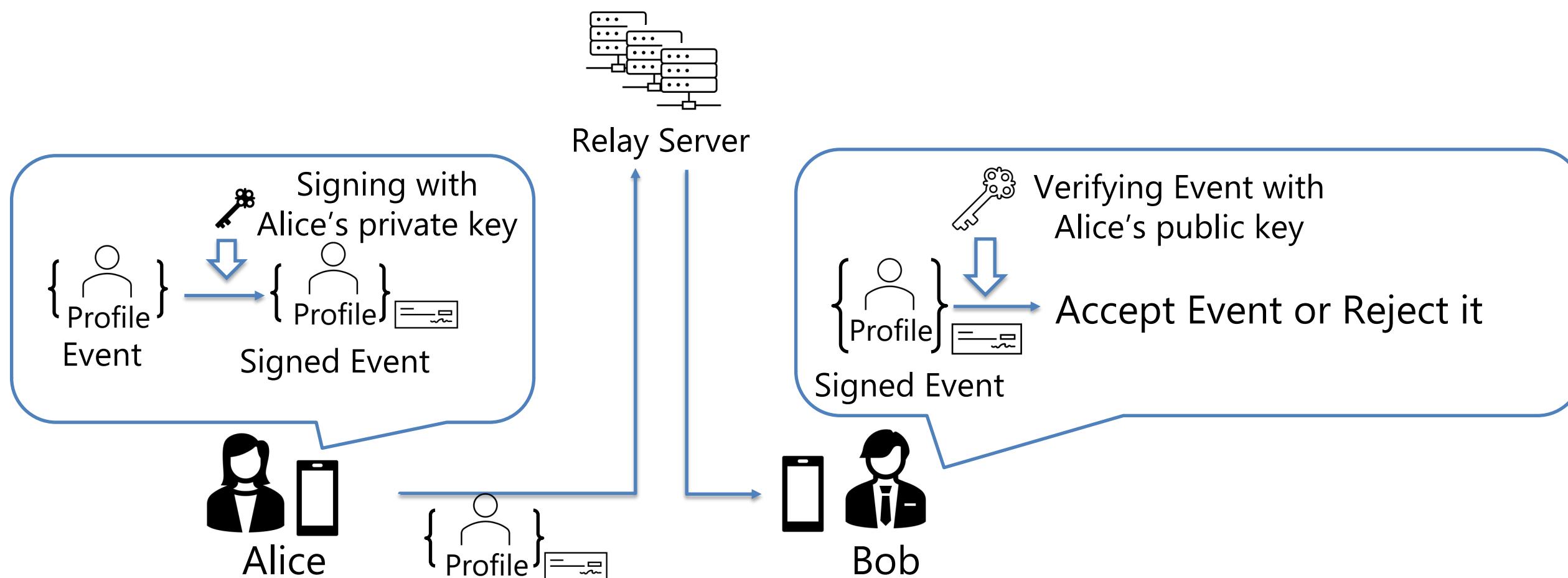


Bob

Breaking...
Plaintext integrity
(cache poisoning)

On the Profile validation of Damus (v1.5(8) & v1.6 (29))

Attack on a popular Nostr client



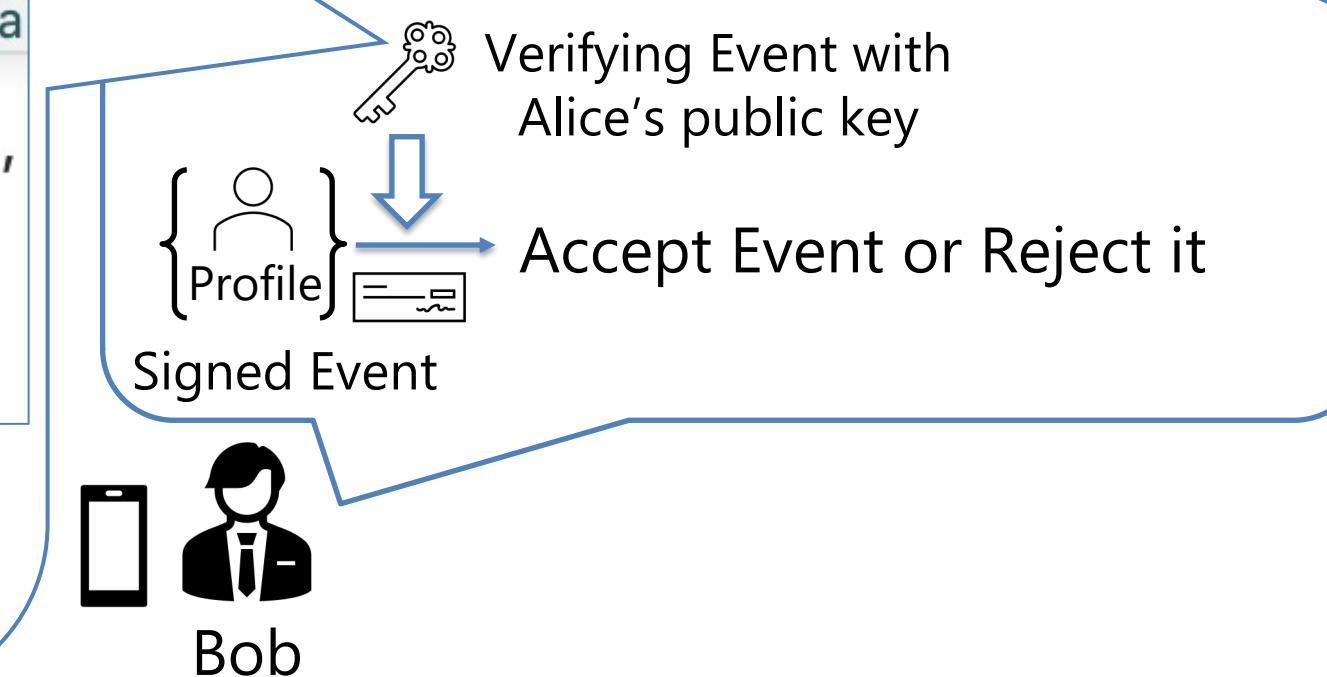
Breaking...
Plaintext integrity
(cache poisoning)

On the Profile validation of Damus (v1.5(8) & v1.6 (29))

Attack on a popular Nostr client

Sig verification in place 😊

```
905 func validate_event(ev: NostrEvent) -> Va  
929  
930     ok = secp256k1_schnorrsig_verify(ctx,  
         &xonly_pubkey) > 0  
931     return ok ? .ok : .bad_sig  
932 }
```

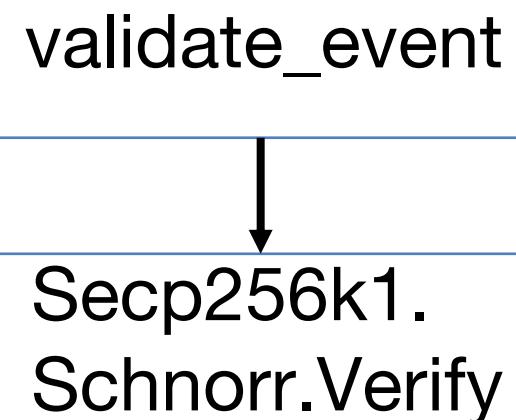


Breaking...
Plaintext integrity
(cache poisoning)

On the Profile validation of Damus (v1.5(8) & v1.6 (29)) Let's see stack trace

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933 }
```



Breaking...
Plaintext integrity
(cache poisoning)

On the Profile validation of Damus (v1.5(8) & v1.6 (29))

Let's see stack trace

EventCache & reference ...? 😐

```
func guard_valid_event(events: EventCache, ev: NostrEvent, callback: @escaping  
    let validated = events.is_event_valid(ev.id)
```

guard_valid_event



validate_event

Secp256k1.
Schnorr.Verify

Breaking...
Plaintext integrity
(cache poisoning)

On the Profile validation of Damus (v1.5(8) & v1.6 (29))

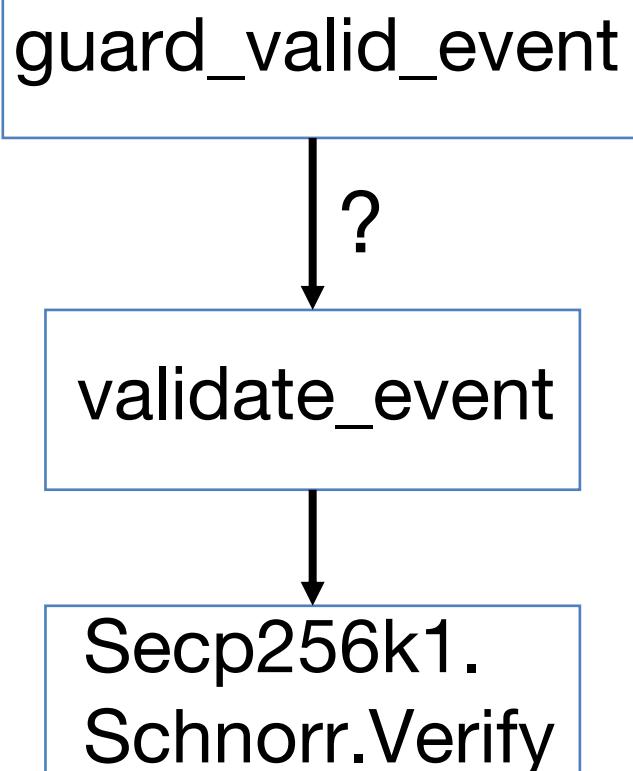
Let's see stack trace

EventCache.is_event_valid(ev.id)

Check past signature verification result

- Return **true** if the event is found and the past verification **succeeded**
- Return **false** otherwise.

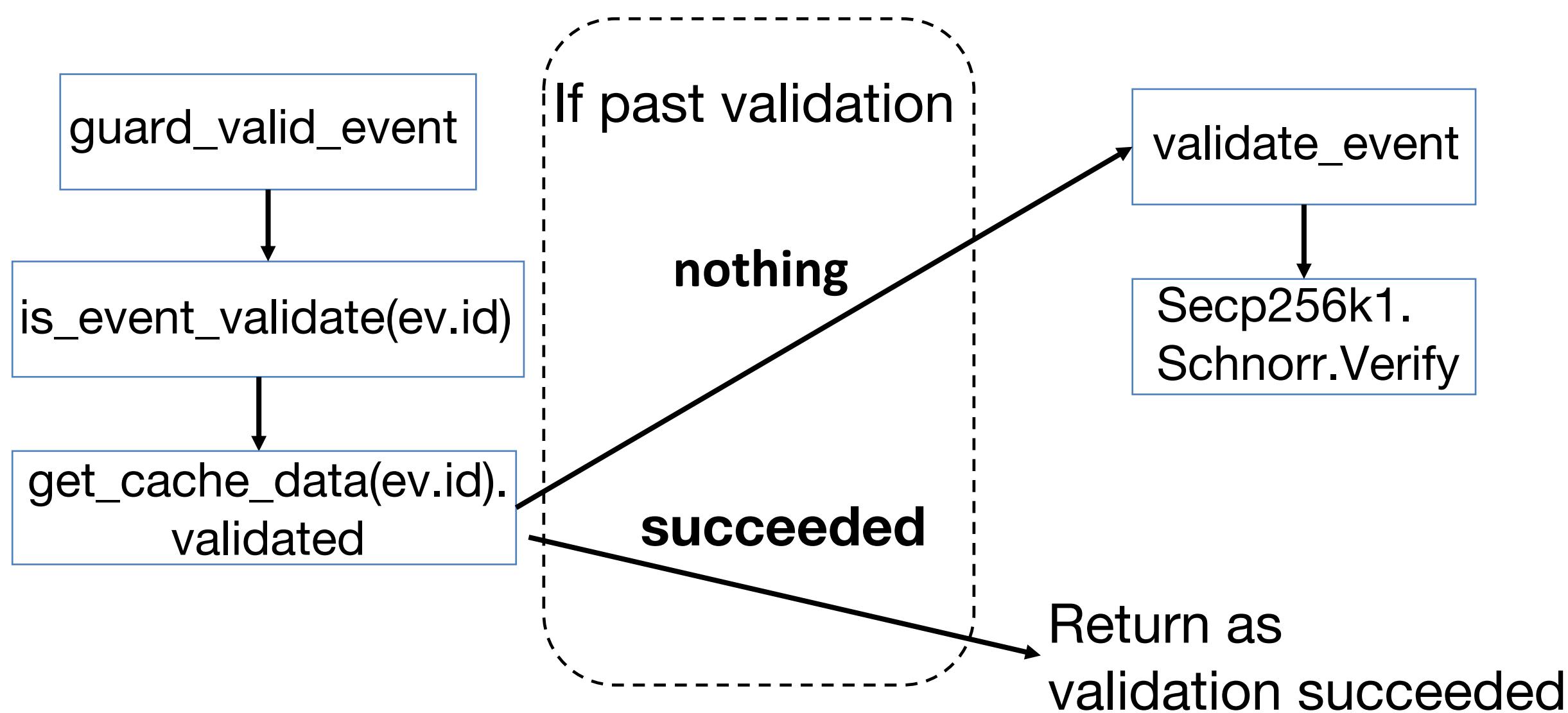
```
157     func get_cache_data(_ evid: String) -> EventData {  
158         guard let data = event_data[evid] else {  
159             let data = EventData()  
160             event_data[evid] = data  
161             return data  
162         }  
163         return data  
164     }  
165  
166     func is_event_valid(_ evid: String) -> ValidationResult {  
167         return get_cache_data(evid).validated  
168     }  
169 }
```



Breaking...
Plaintext integrity
(cache poisoning)

On the Profile validation of Damus (v1.5(8) & v1.6 (29))

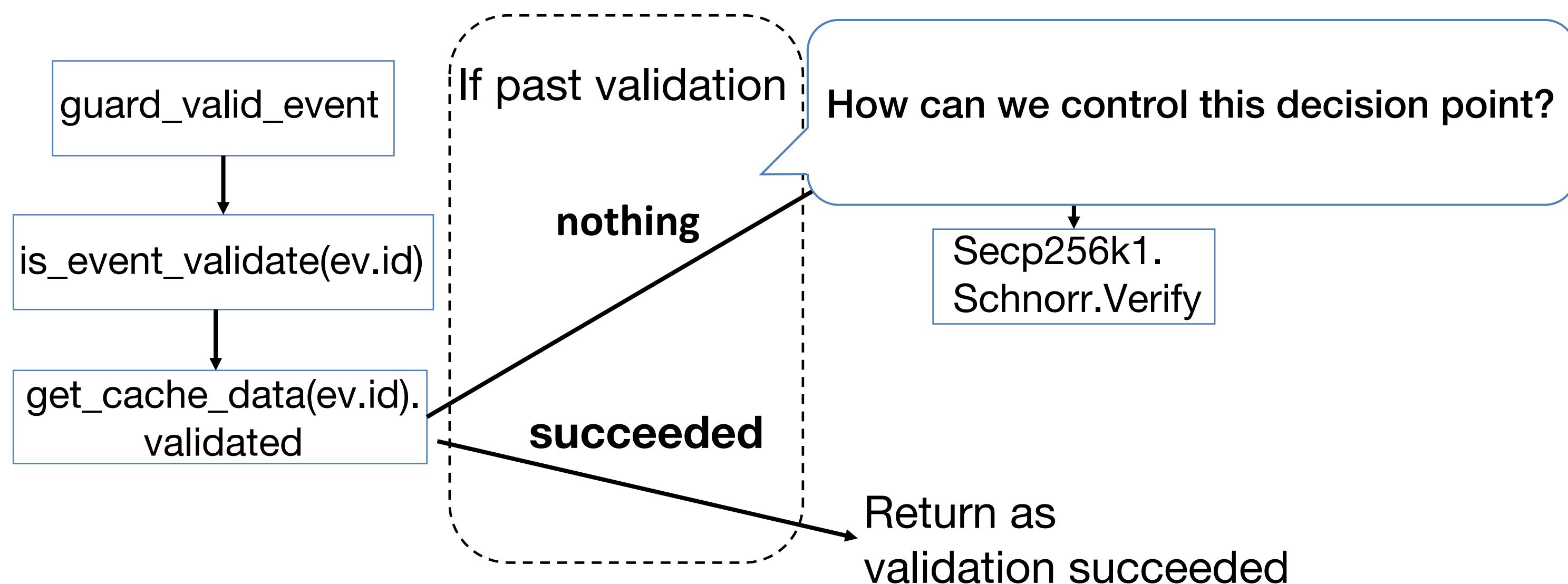
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Breaking...
Plaintext integrity
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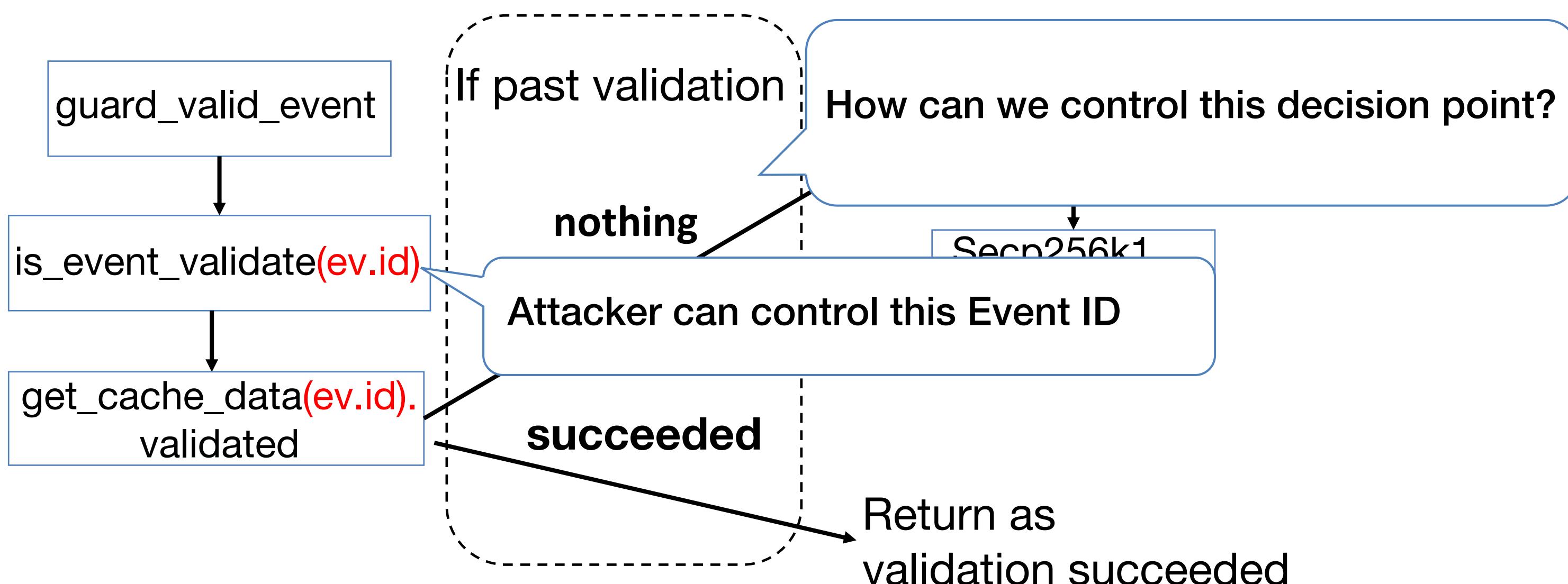
Let's see stack trace



Breaking...
Plaintext integrity
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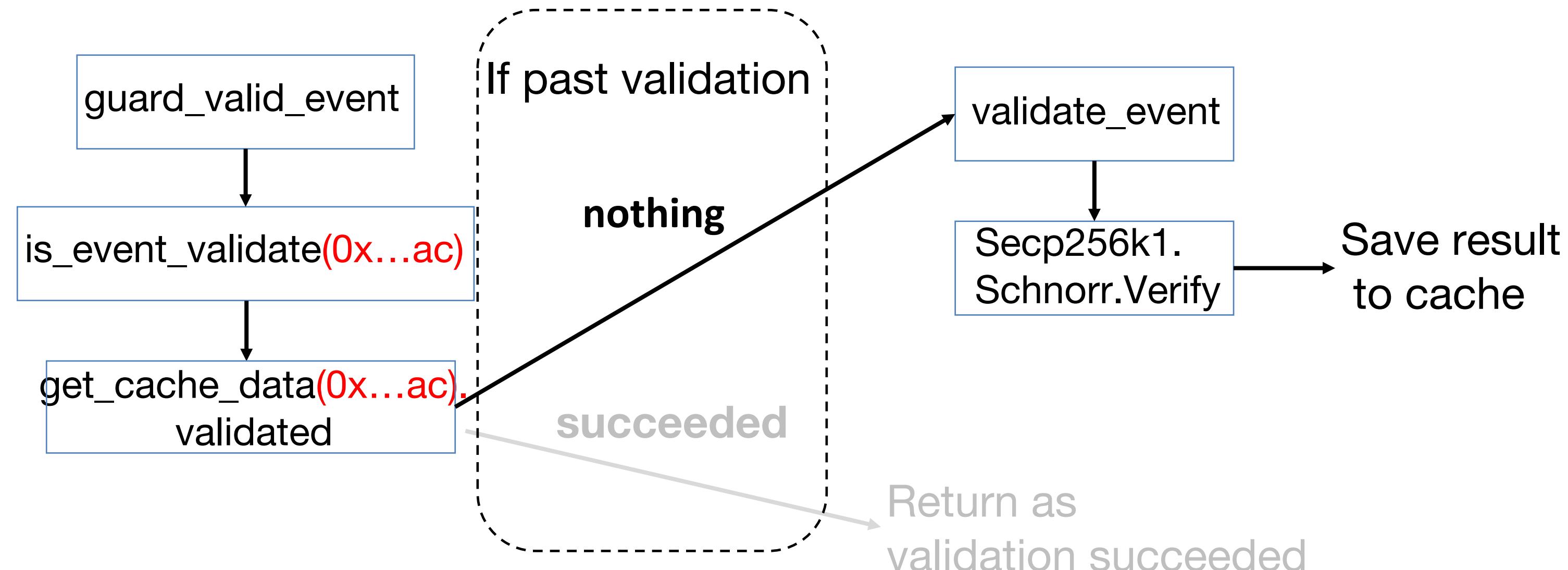
Let's see stack trace



Breaking...
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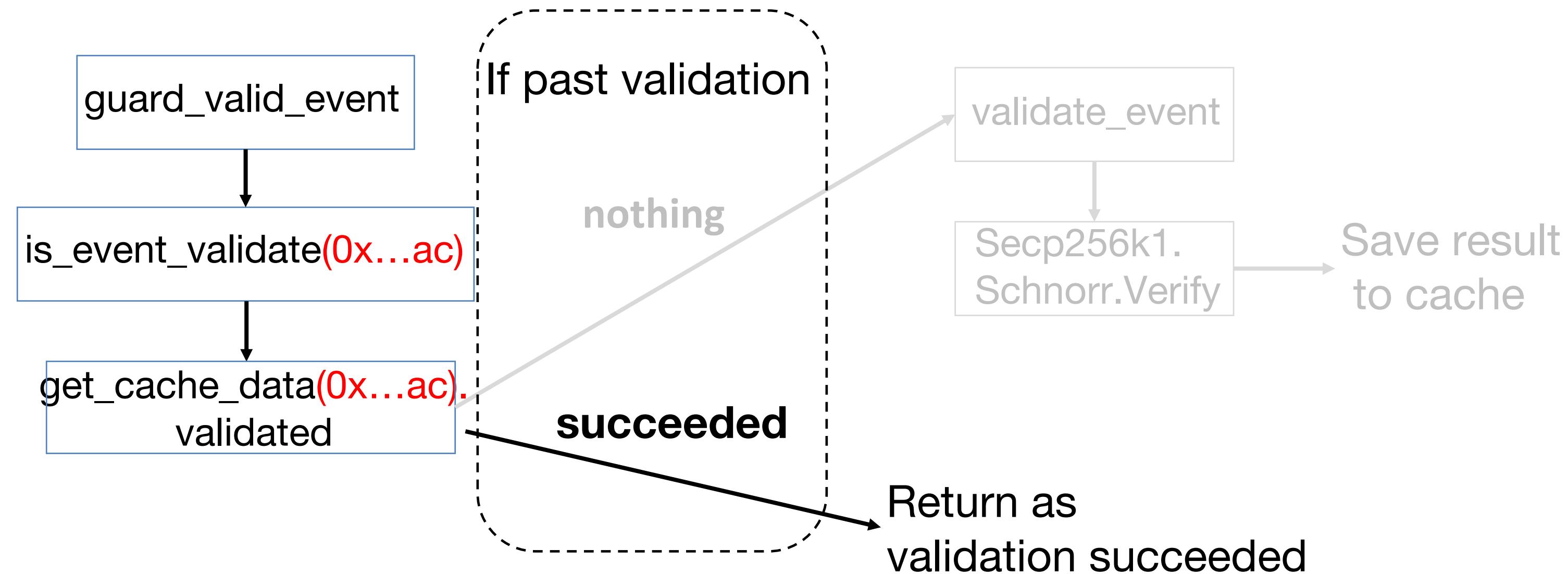
When Bob received an Alice's Event (id== 0x...ac)



Breaking...
Plaintext integrity
(cache poisoning)

On the Profile validation of Damus (v1.5(8) & v1.6 (29))

Attacker sends a fake event with an ID (0x...ac) to Bob

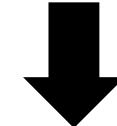


Breaking...
Plaintext integrity
(cache poisoning)

On the Profile validation of Damus (v1.5(8) & v1.6 (29))

How to derivate Event ID on Nostr

Event ID : $ev.id = \text{SHA-256}("0" || \{ev.data\})$



The event ID is a deterministic value derived from `ev.data`

Root cause: Refer to the cache using the ID without recalculating it



Mitigation

The **ID should be recalculated if {ev.data} is modified.**

Breaking...
Plaintext integrity
(cache poisoning)

On the Profile validation of Damus (v1.5(8) & v1.6 (29))

Mitigation: Event ID validation

Original: No ID validation

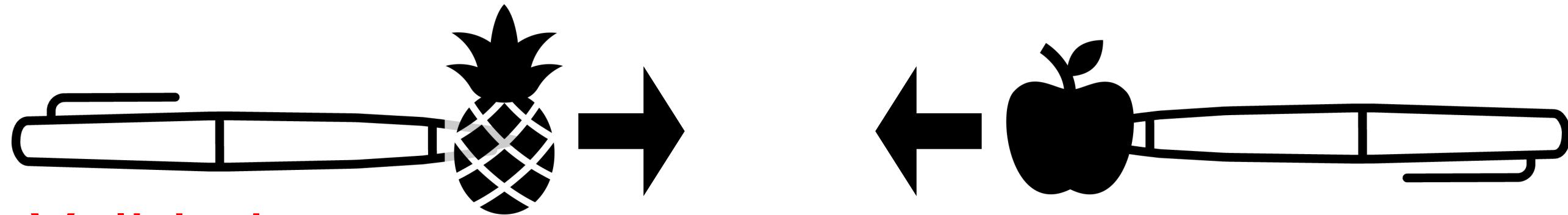
```
750 func guard_valid_event(events: EventCache, ev: NostrEvent,  
751     let validated = events.is_event_valid(ev.id)  
___
```

Patched: Ensure ID validation

```
750 func guard_valid_event(events: EventCache, ev: NostrEvent,  
751     guard ev.id==calculate_event_id(ev: ev) else {  
752         return  
753     }  
754     let validated = events.is_event_valid(ev.id)
```

Breaking...
Plaintext integrity

Takeaway : Plaintext Integrity

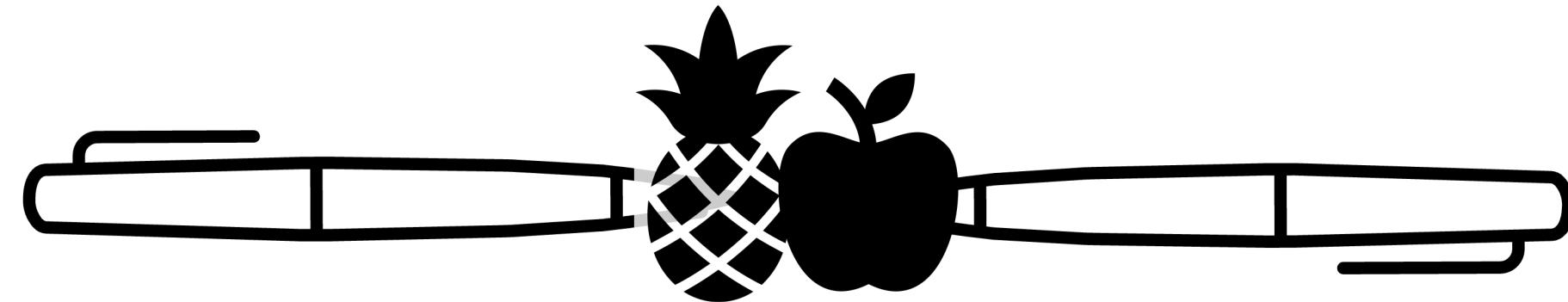


Developer should do integrated security test !

Breaking...

Plaintext integrity

Takeaway : Plaintext Integrity



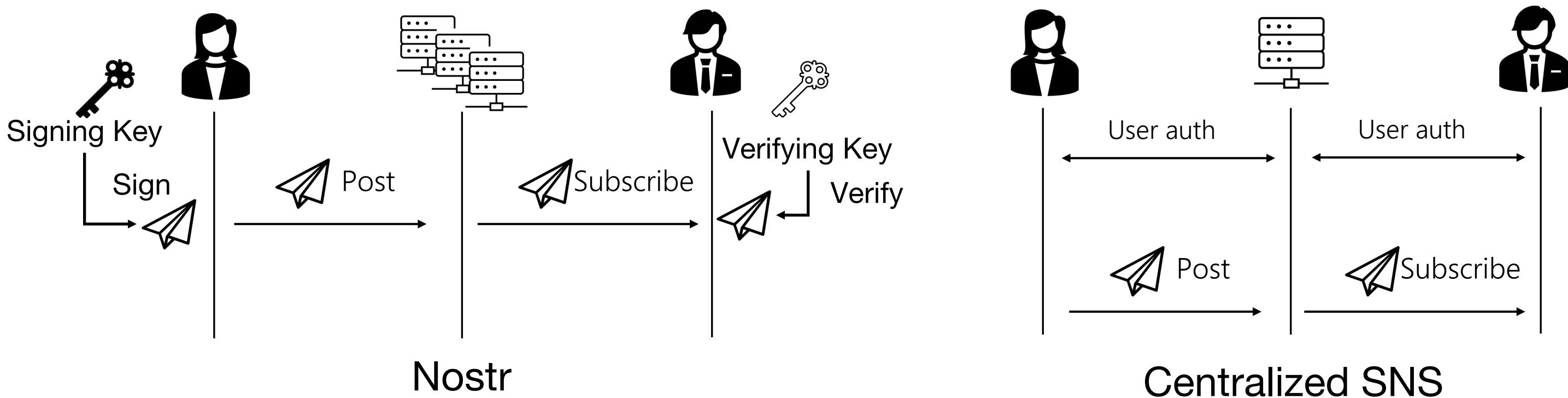
Authentication
Bypass

Developer should do integrated security test !

Breaking...
Plaintext integrity

Takeaway : Plaintext Integrity (2)

- In centralized settings, cryptographic flaws often remain “potential risks”
- In self-sovereign decentralized systems like Nostr, they become immediately exploitable
 - Nostr does not have centralized authority
 - Nostr does not provide user authentication by default



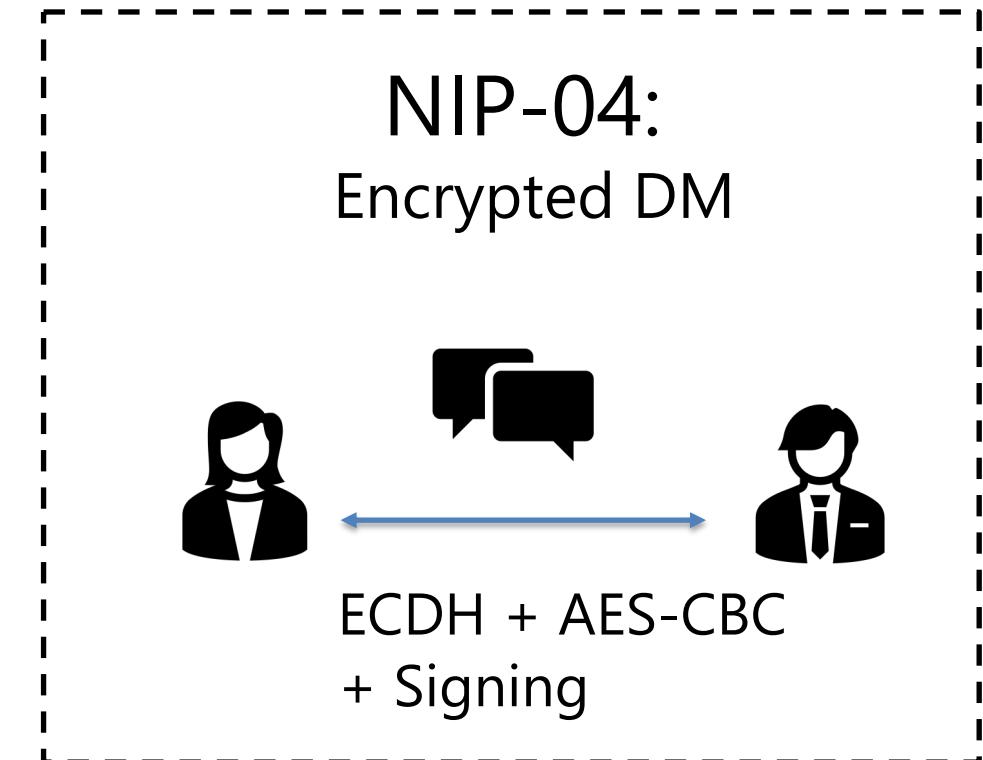
Step by step attack tracing Breaking...

Plaintext integrity
(simple / cache poisoning)

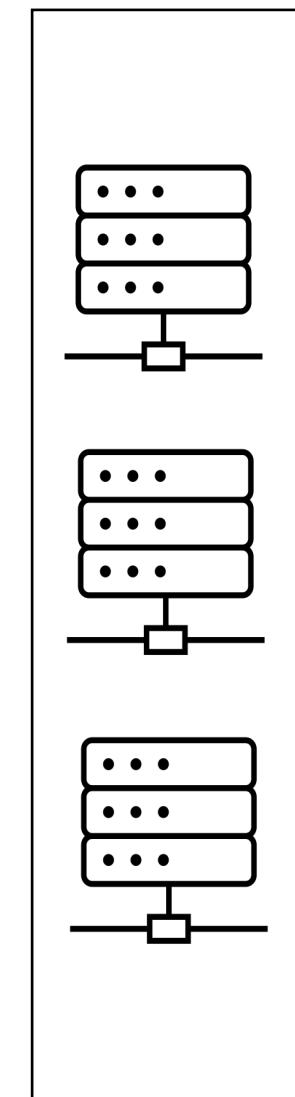
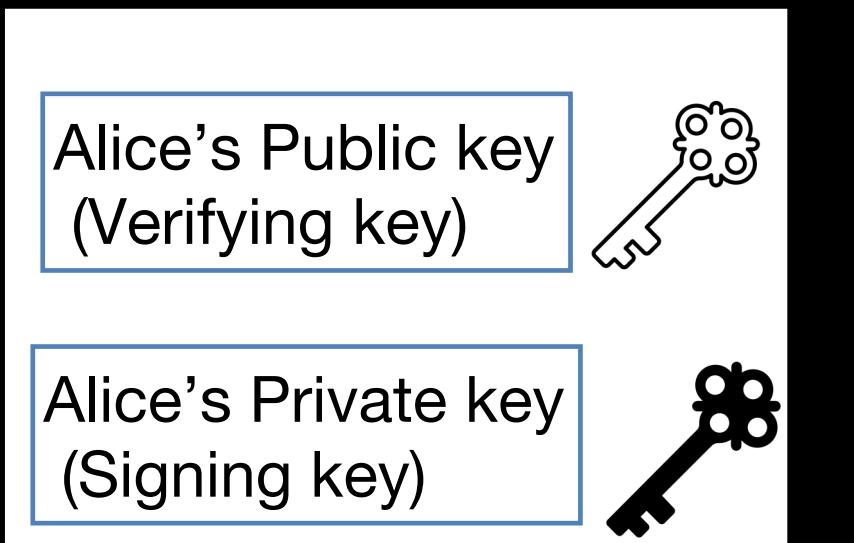
Ciphertext integrity

Ciphertext confidentiality

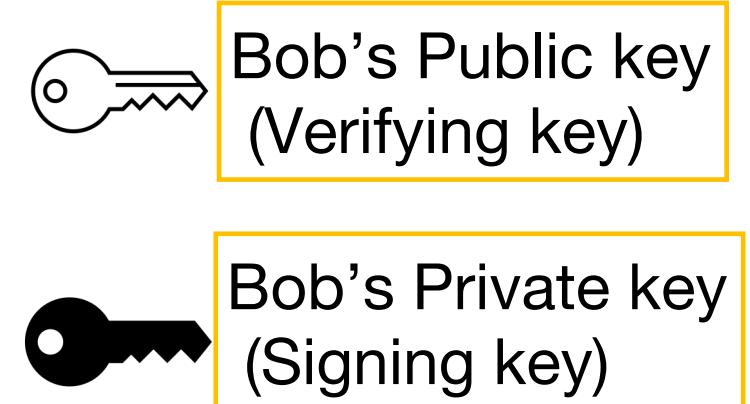
Remark: Encrypted Direct Messages
specification (simplified)



Encrypted DM Spec

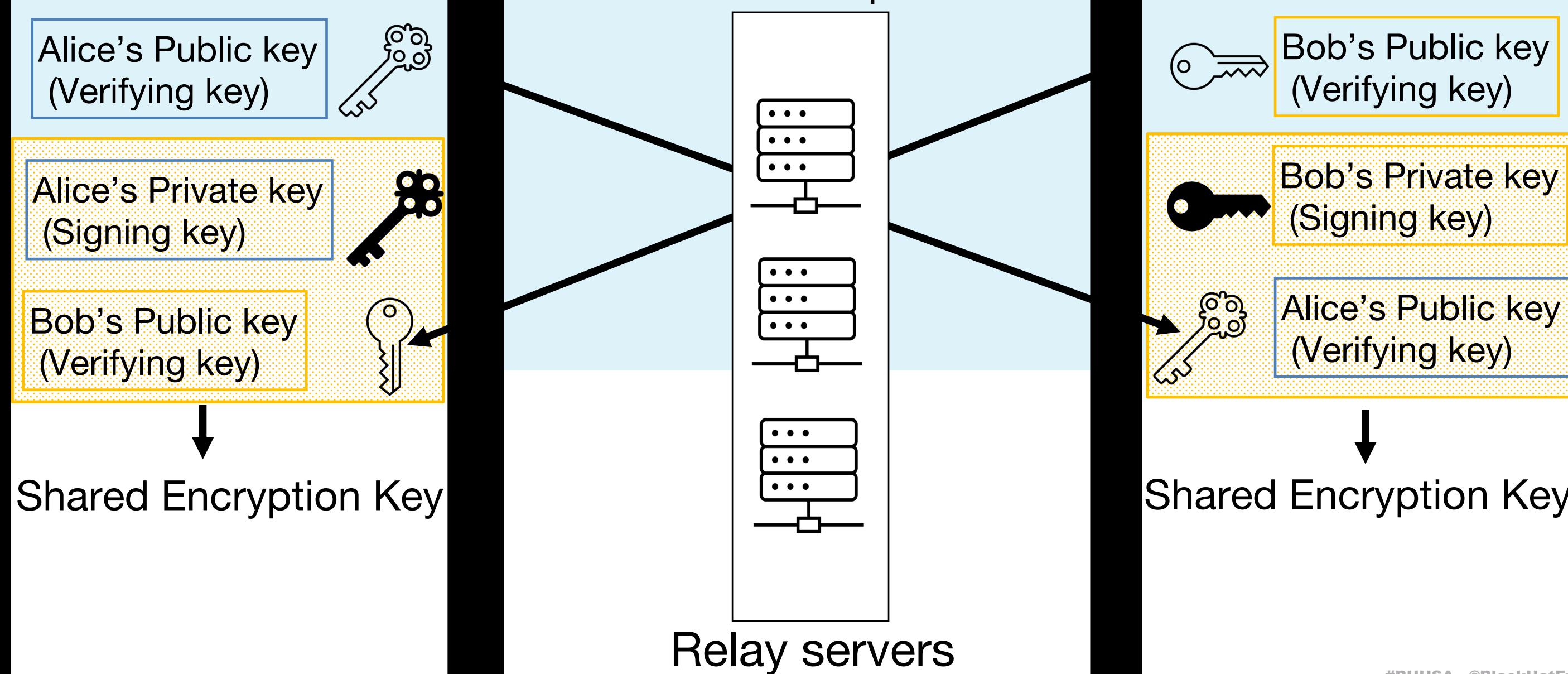


Relay servers



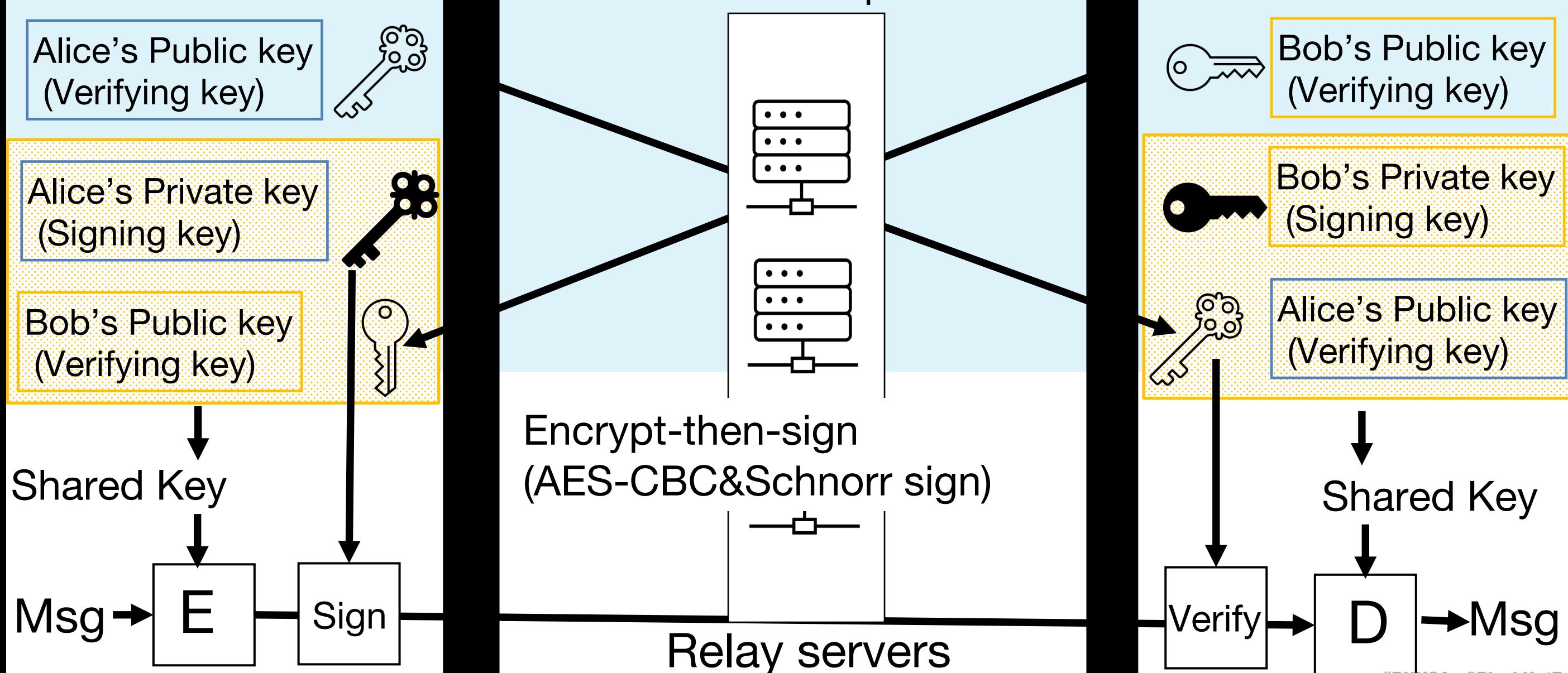
Encrypted DM Spec

ECDH over secp256k1



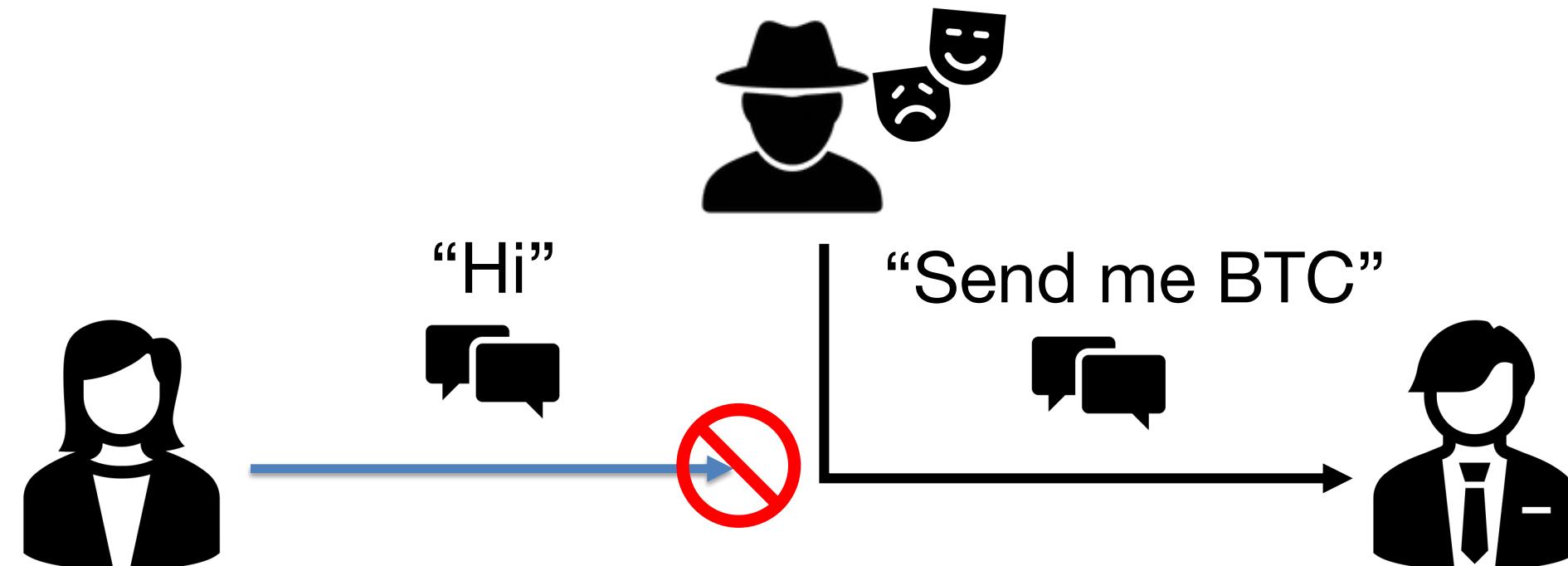
Encrypted DM Spec

ECDH over secp256k1



Encrypted DM Forgery

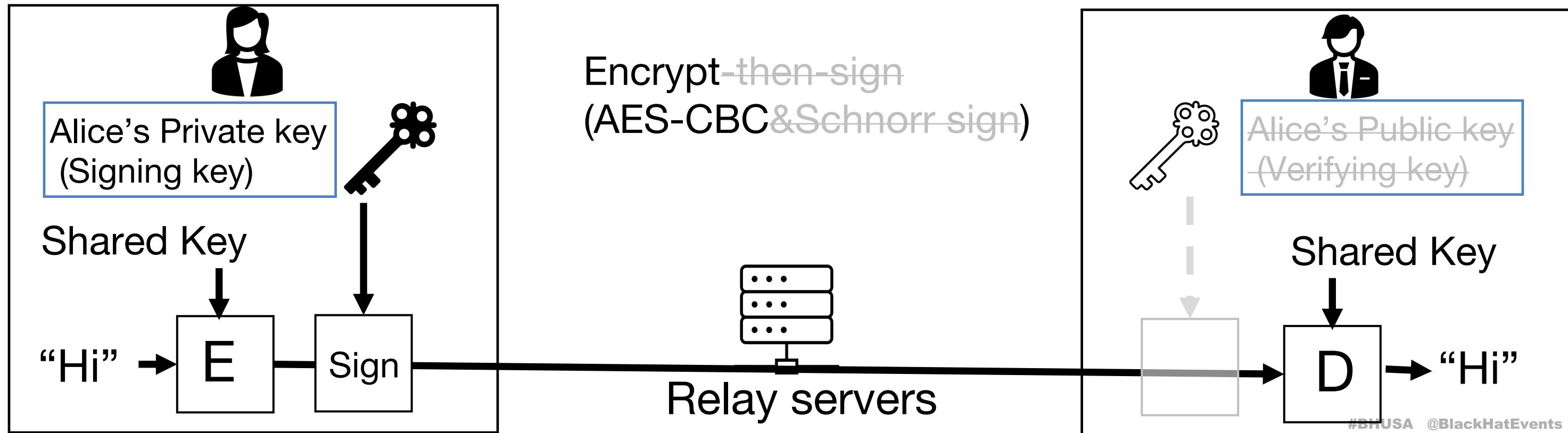
Attacker's Goal : Change decrypted Msg to attacker's Msg_{adv}
e.g., "Send me BTC"



ECDH + AES-CBC
+ Signing

Encrypted DM Forgery

Assumption1: Signature verification is skipped on the implementation (explained earlier)



Encrypted DM Forgery

Assumption2: Threat model

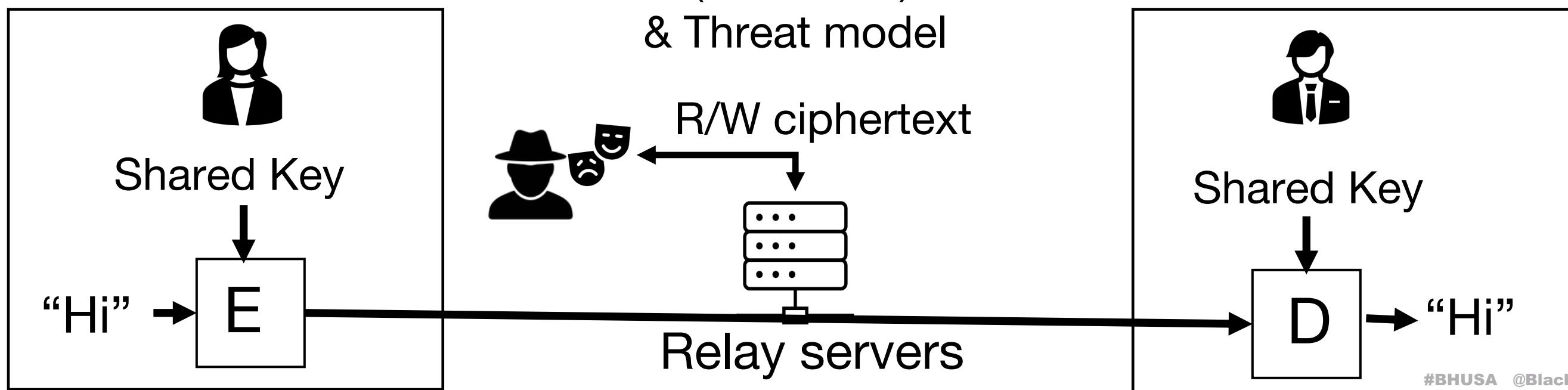
- Attacker is a user of Nostr
- Attacker cannot read/write to “Shared Key”
- Attacker can **freely fetch ciphertext** from relay servers

Nostr does not include user authentication on servers by default

Simplified encryption specs

(AES-CBC)

& Threat model



Encrypted DM Forgery

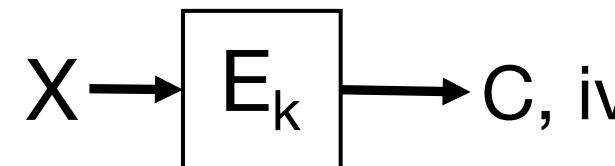
Problem : Verification bypass is not enough to achieve practical forgery on DMs

Reason : CBC Allows Bit Flipping – But decryption result blinds for the attacker

(simplified) Bit flipping on Message Encryption

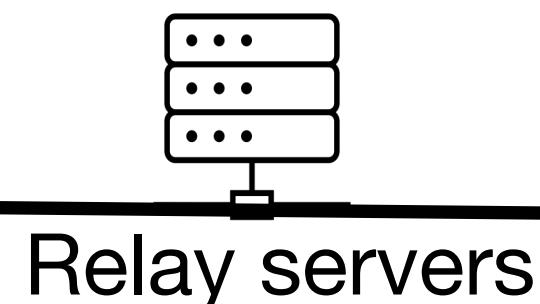
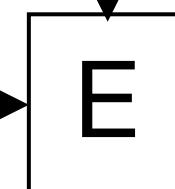
1 block CBC-mode encryption

$$X \leftarrow iv \oplus "Hi" \parallel pad$$



Shared Key (k)

"Hi"



Shared Key

D

Msg_{dec}

Encrypted DM Forgery

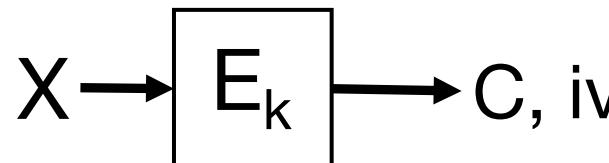
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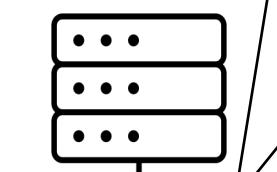


Shared Key (k)

"Hi"



$$iv' \leftarrow iv \oplus \text{Flip}$$



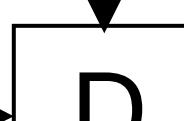
Relay servers

1 block CBC-mode decryption

$$C \rightarrow D_k \rightarrow X \oplus iv'$$

$Msg_{dec} = ??$ (unknown)

Shared Key



What does the attacker need to control the decryption result?

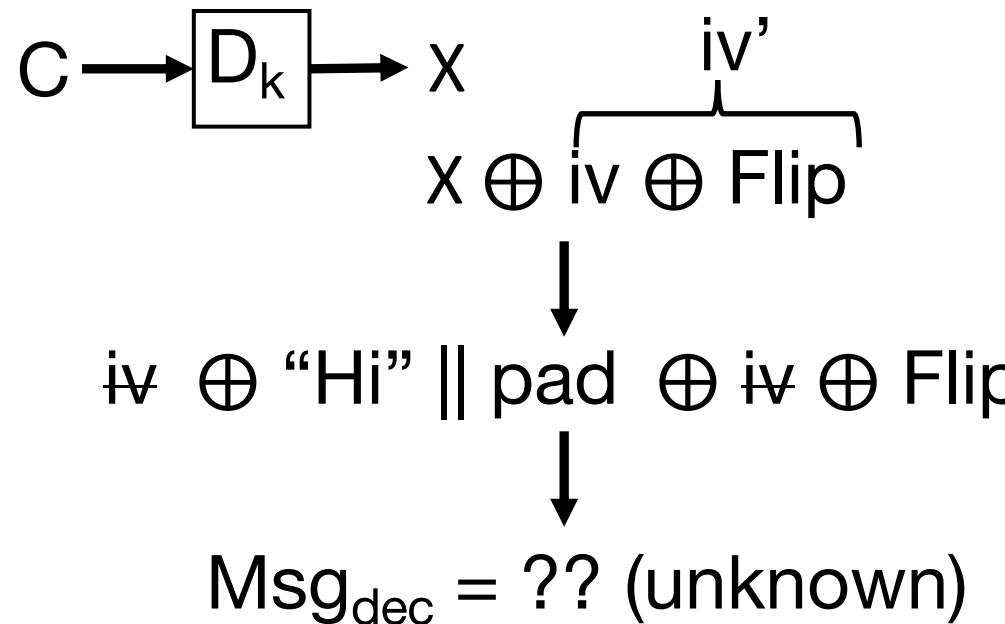
To craft a forged ciphertext, the attacker needs a reference point:

→ a known plaintext/ciphertext (C_{ref} , Msg_{ref}) pair with the same shared key (k)

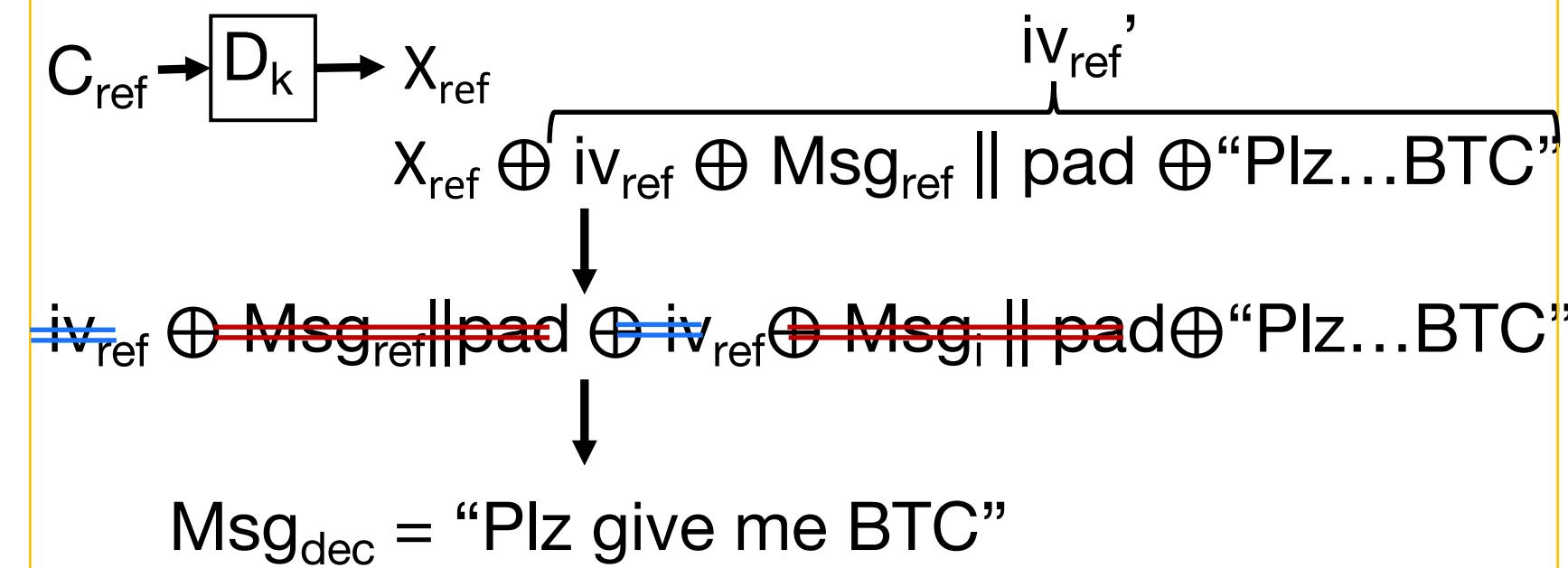
Cf. Encryption: $X \leftarrow iv \oplus Msg \parallel pad$ $C \leftarrow E_k(X)$, send iv & C



Random bit-flipping forgery



Practical forgery using a known (C_{ref}, Msg_{ref}) pair



Move from Bit Flipping Forgery to Controlled Practical Forgery



Random bit-flipping forgery

- No decryption knowledge
- Can't control decrypted message
- Just makes noise



Practical forgery using a known $(C_{ref}, \text{Msg}_{ref})$ pair

- Known plaintext/ciphertext block
- XOR trick enables precision
- Delivers chosen message to victim

Move from Bit Flipping Forgery to Controlled Practical Forgery



Random bit-flipping forgery

- No decryption knowledge
- Can't control decrypted m
- Just makes noise

Problem:
How can we get it ?



Practical forgery using a known $(C_{ref}, \text{Msg}_{ref})$ pair

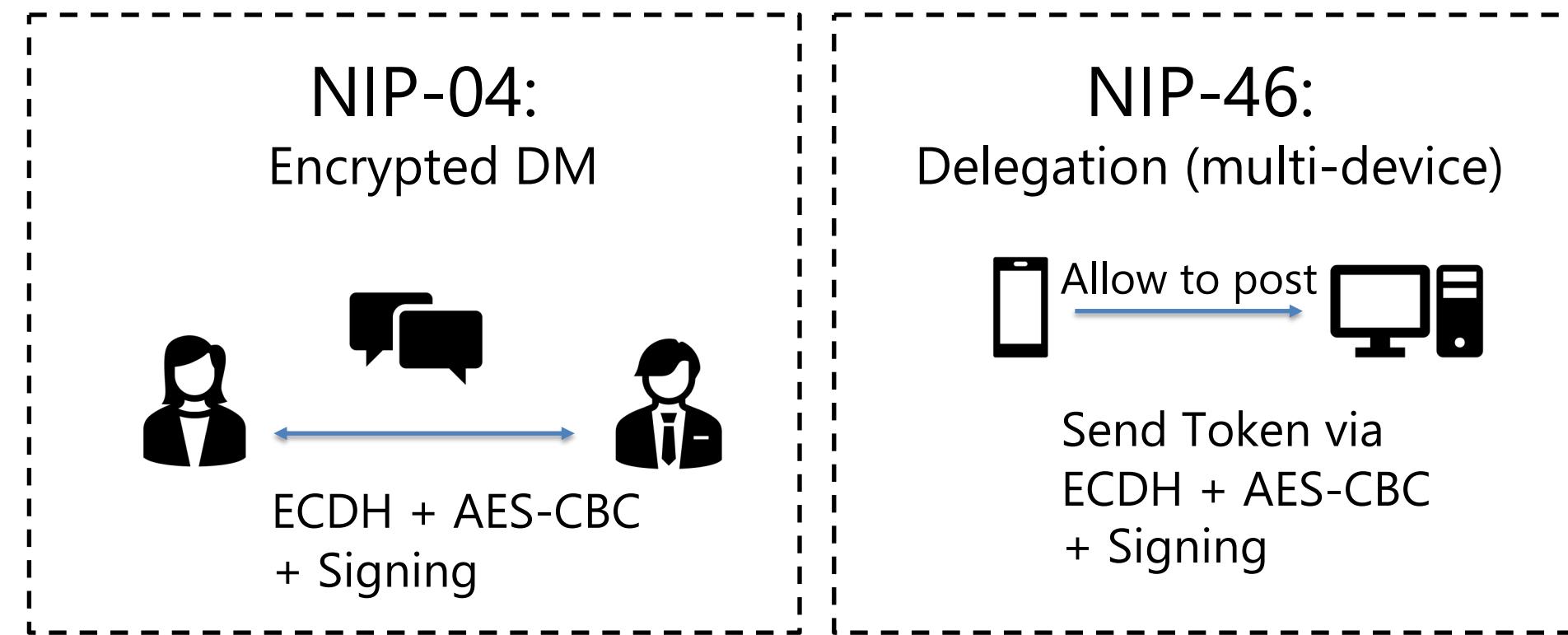
- Known plaintext/ciphertext block
- XOR trick enables precision
- Delivers chosen message to victim

Encrypted DM Forgery via Cross Protocol Attack

Solution : Breaking the Barrier via “Cross Protocol” Attack

Observation:

Delegation (NIP-46) uses same keying & encryption algorithms as DMs (NIP-04)
NIP-46 encrypts known metadata using the same shared key as DMs (NIP-04)



Encrypted DM Forgery via Cross Protocol Attack

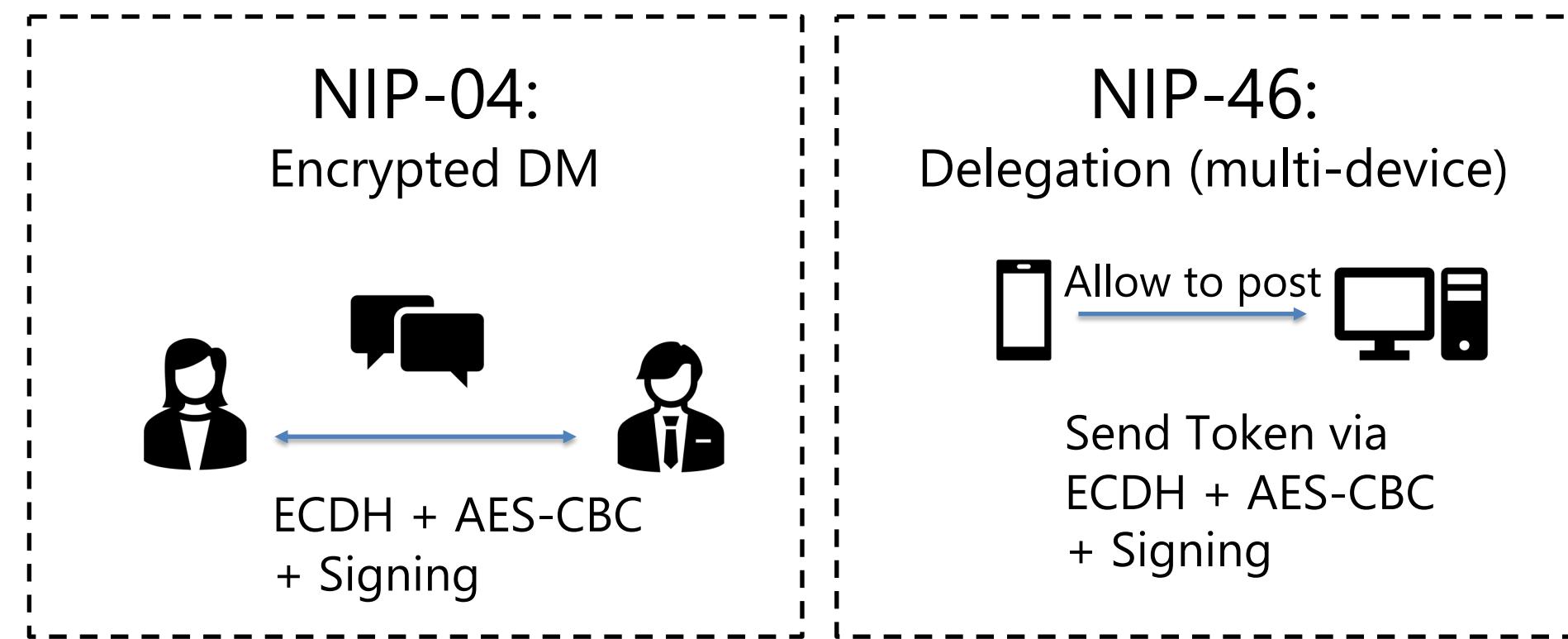
Solution : Breaking the Barrier via “Cross Protocol” Attack

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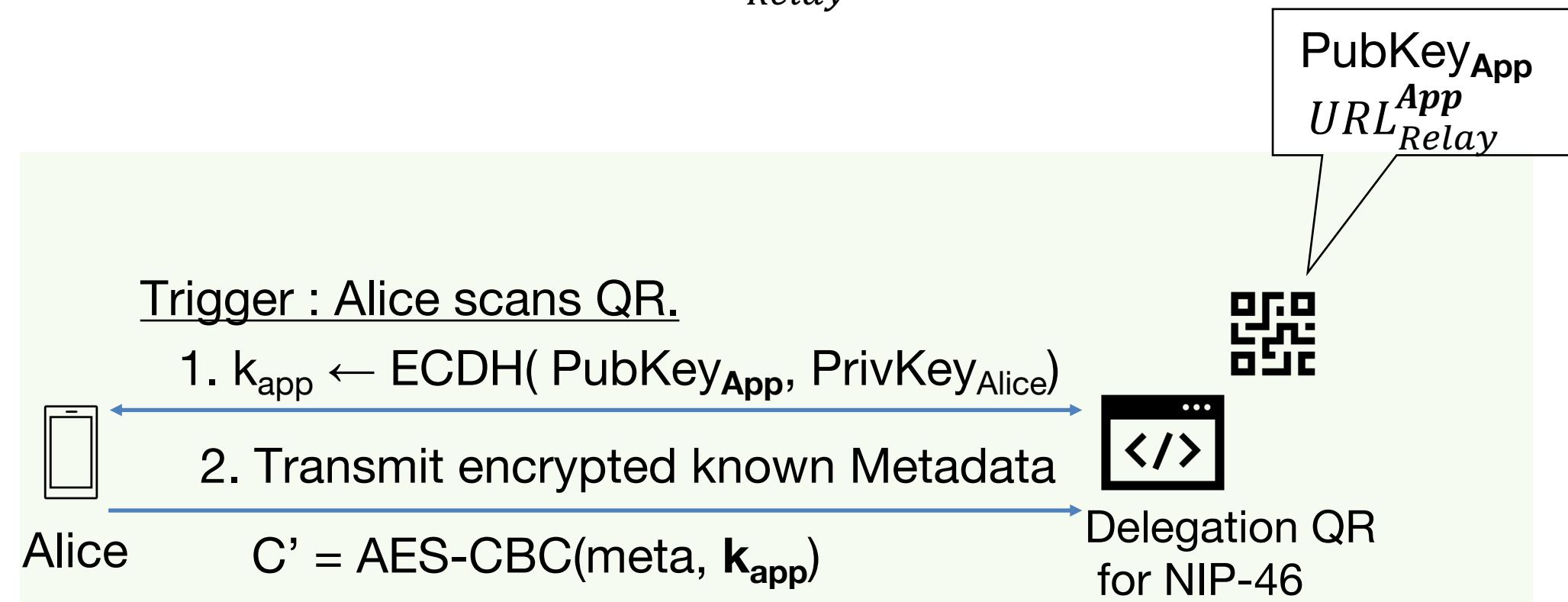
→ makes known plaintext → makes known ciphertext



Encrypted DM Forgery via Cross Protocol Attack

Normal Delegation initial sequence

- ECDH with the public key obtained from the QR
- Sends encrypted known metadata to URL_{Relay}^{App} from the QR



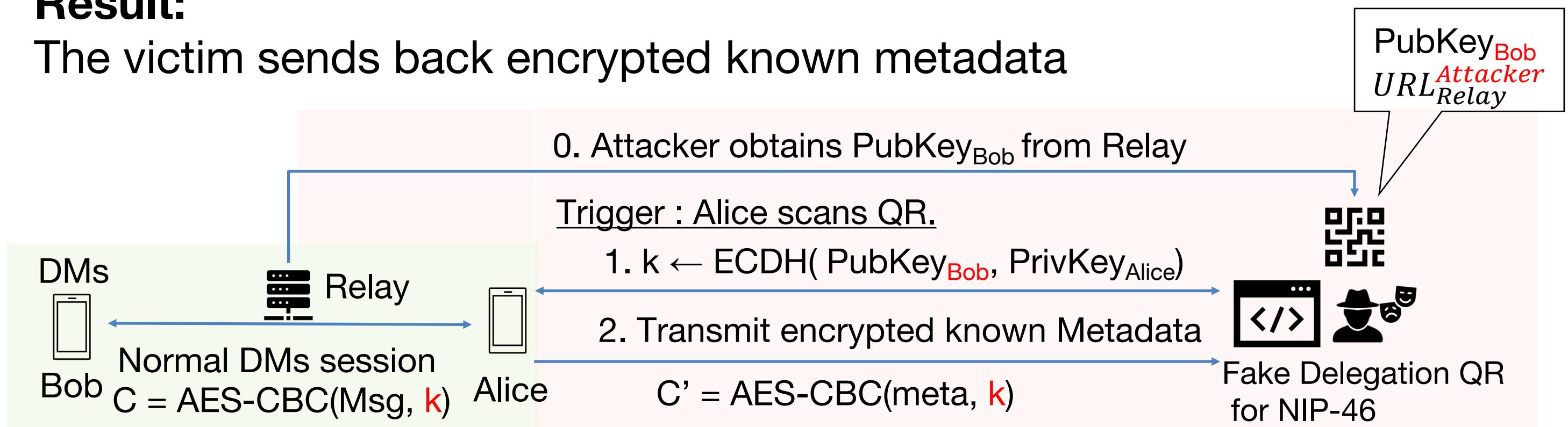
Encrypted DM Forgery via Cross Protocol Attack

Strategy:

The attacker starts a NIP-46 session with the victim (as a fake delegation app)
The attack puts PubKey_{Bob} to the QR

Result:

The victim sends back encrypted known metadata



Takeaway : Ciphertext Integrity

- **Should use Authenticated Encryption (AE)**
 - E.g., AES-GCM, ChaCha20-Poly1305
 - Don't use malleable encryption without MAC
- **Should separate key between sub-protocols**
 - Similar issues also occurred in Threema[PST23], Matrix[ACDJ23]

[PST23] Paterson, Scarlata and Truong, “Three Lessons From Threema: Analysis of a Secure Messenger”, USENIX Security’23

[ACDJ23] Albrecht, Celi, Dowling and Jones, “Practically-exploitable Cryptographic Vulnerabilities in Matrix”, IEEE S&P’23

(Also, Black Hat Europe’22)

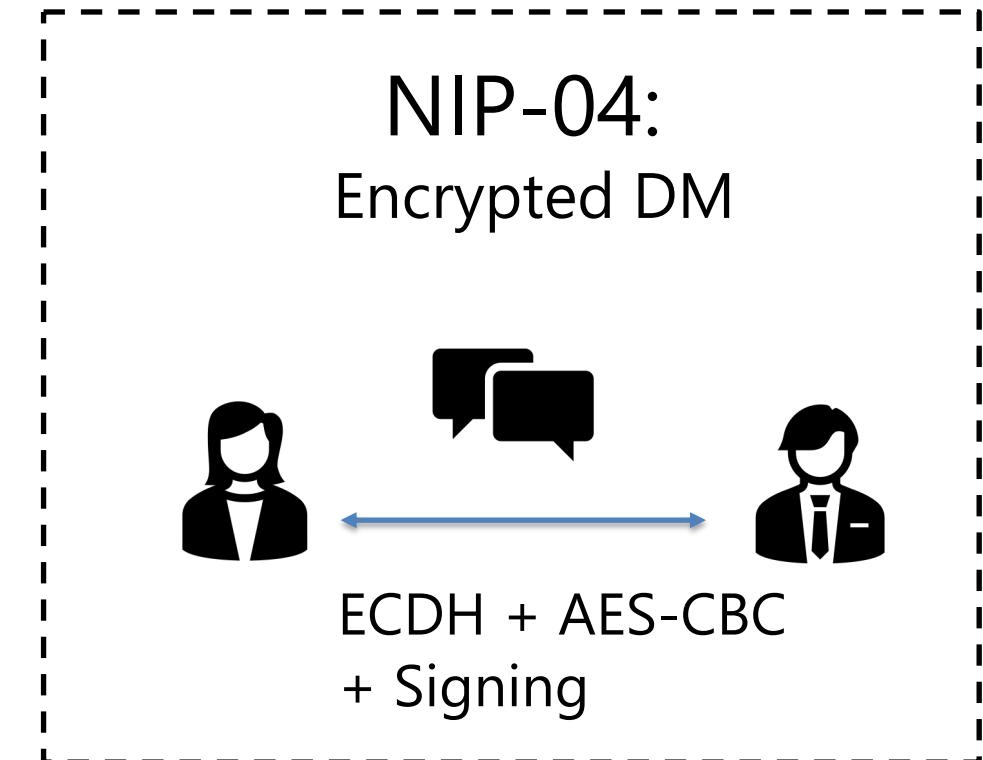
Step by step attack tracing Breaking...

Plaintext integrity
(simple / cache poisoning)

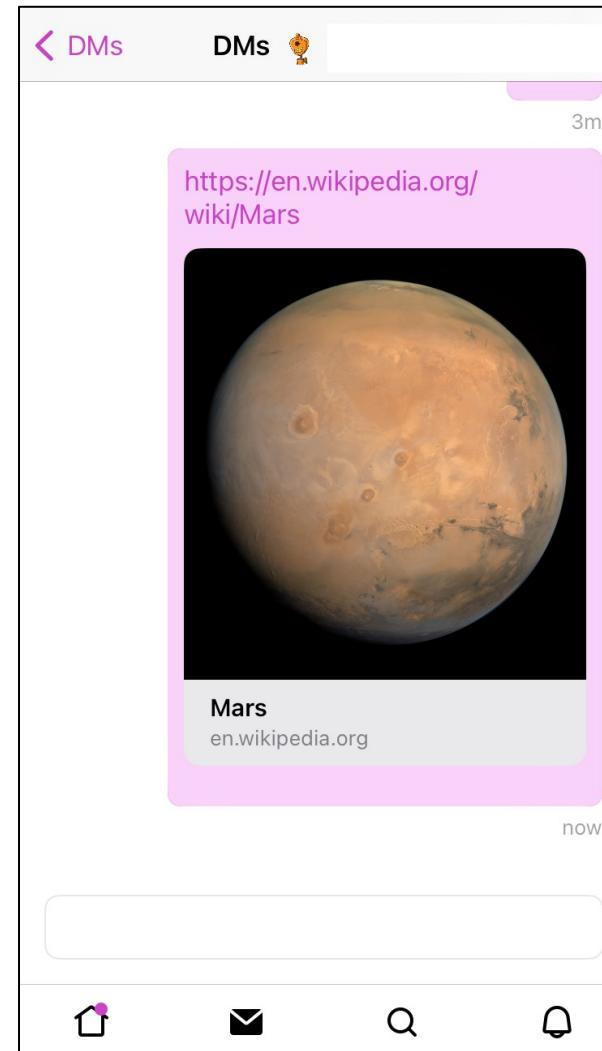
Ciphertext integrity

Ciphertext confidentiality

Remark: Encrypted Direct Messages
specification (simplified)

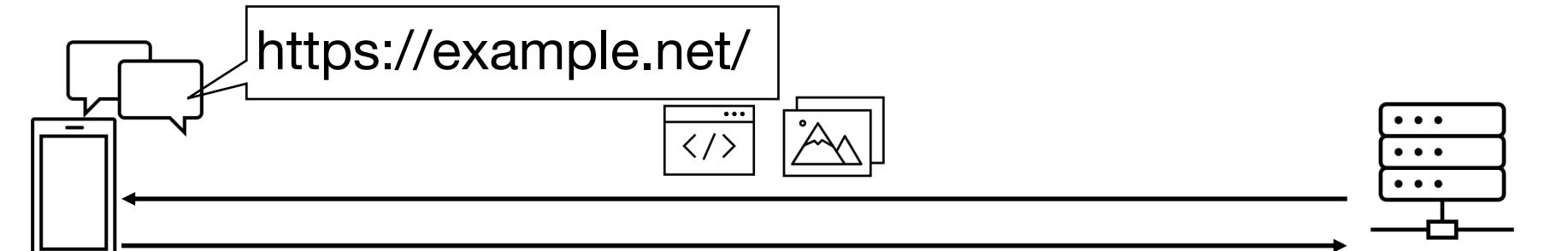


Link Preview in Messaging

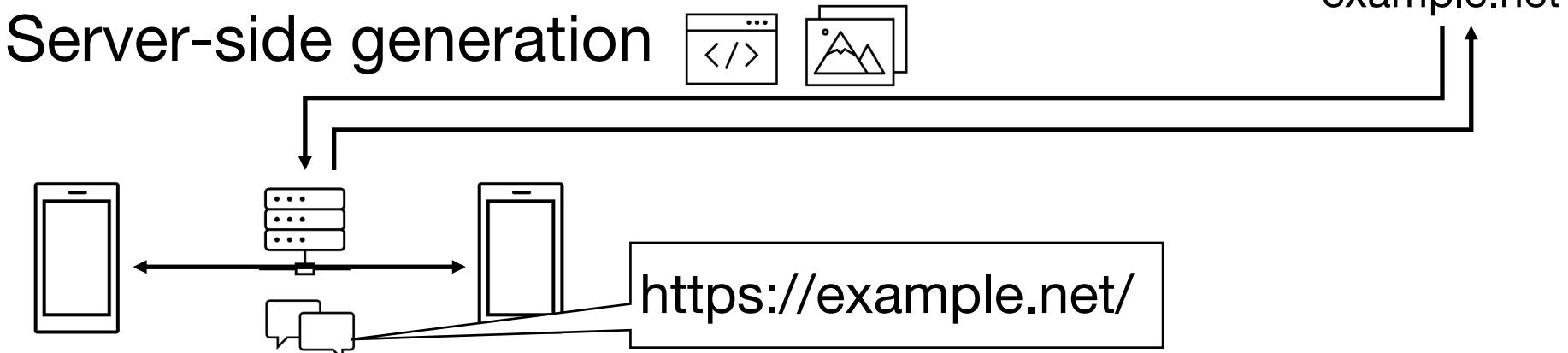


- Automatically retrieves and displays elements from the webpage
E.g., The webpage's title, part of its content, and images
- Someone must retrieve the page content (a sender, a receiver or a server)

Client-side generation

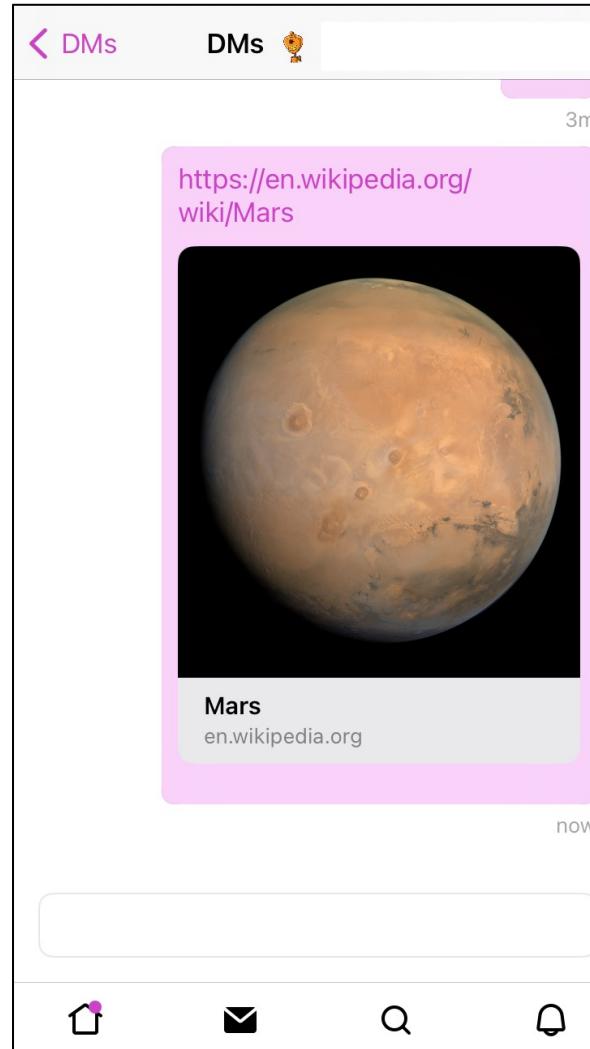


Server-side generation



*Non E2EE msg

Link Preview generation in **Encrypted Messaging**



Best practice

- Generate preview **ONLY** on the sender-side

Bad practice

- Generate preview on the receiver-side
- Known privacy issues (IP leakage): <https://mysk.blog/2020/10/25/link-previews/>

Is there any chance we can use it?

Many Nostr Clients

- Generate preview on the both sender-side and receiver-side

Thinking about plaintext recovery in the real-world **Encrypted Messaging**

- Hard to break cryptographic primitive standard
- But what if the recipient *helps* the attacker reveal an encrypted msg?
- How to win ? → Distinguishes & leaks decryption errors
 - Padding Oracle Attacks often appear in toy environments like CTFs

Q. Can we reproduce such an oracle in real-world systems?

Q. Can we reproduce such an oracle in real-world systems?

- Yes, we can! Receiver-side Link Preview generation helps us
- We finally find 3 attacks to break encrypted message confidentiality

URL recovery attack

Attacker's goal: disclose the authentication token in the URL
E.g., shared URL of cloud storage, web conference tools

`https://us04web.zoom.us/j/ [REDACTED] ?pwd=[REDACTED]`

Attacker wants to know

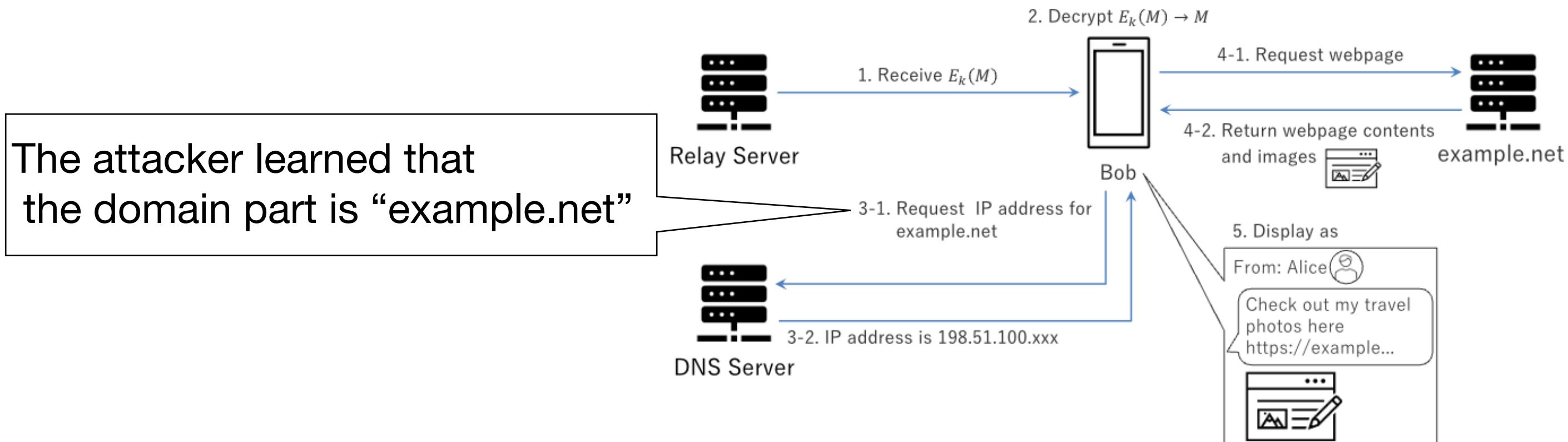
$$E_k(M) : M = \text{https://\{unknown domain\}/\{unknown part\}}$$

Authentication token

URL recovery attack

Disclose domain part

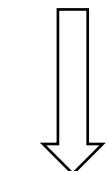
- Attacker can obtain domain part of URL via DNS or TLS SNI field
- Just by opening the message, DNS queries and TLS ClientHello packets are sent due to the automatic execution of link previews.



URL recovery attack

Disclose authentication token

- Force the authentication token to be sent to the attacker's server
- Generate a modified ciphertext $E_k(M')$ where the domain is changed to a malicious one
- When the victim receives $E_k(M')$, the token is sent to the malicious URL via Link Preview

$$E_k(M) : M = \text{https://example.net/}\{\text{unknown part}\}$$


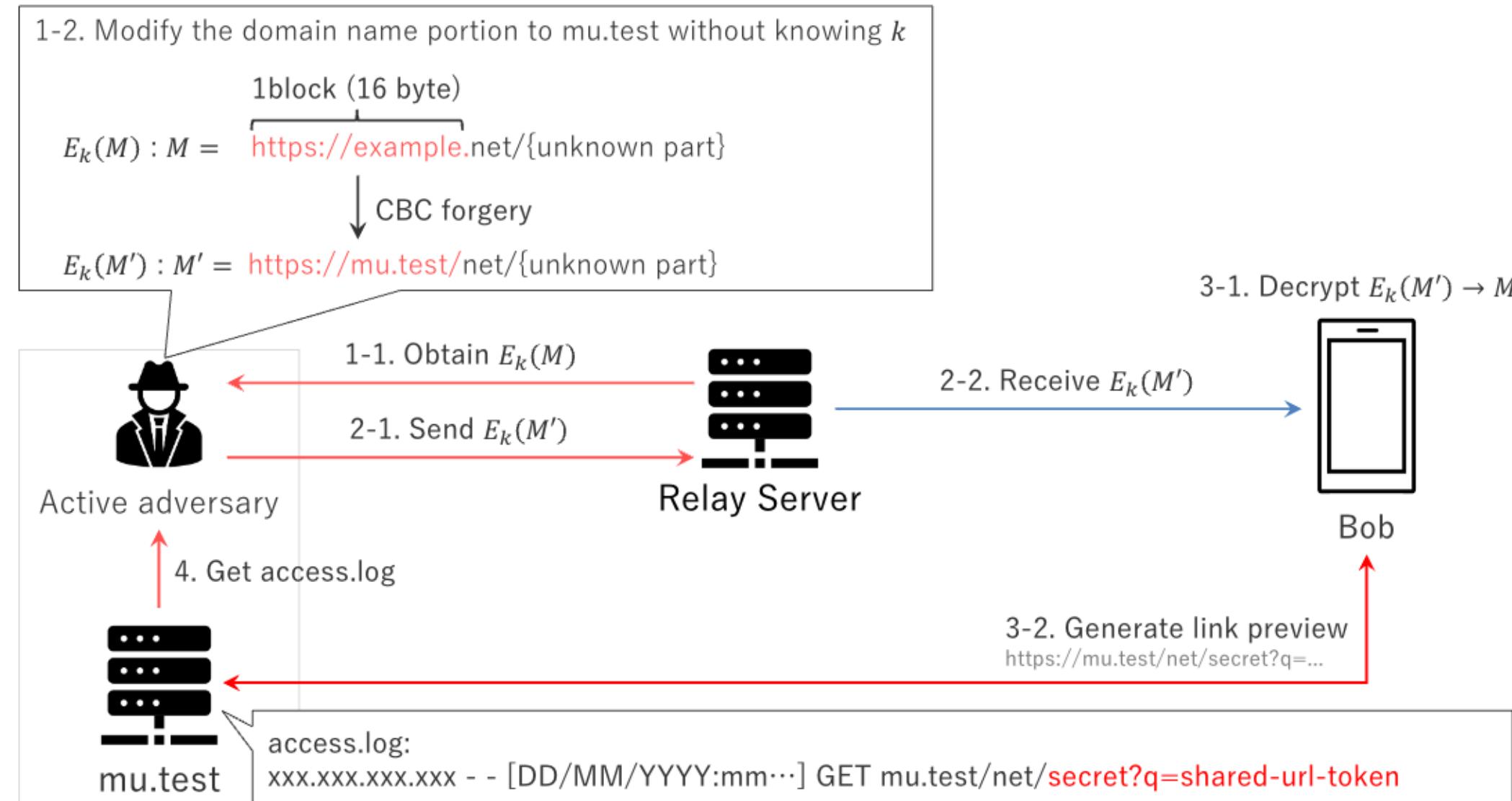
Encrypted DM forgery

$$E_k(M') : M' = \text{https://mu.test/}\{\text{unknown part}\}$$


1Block (16Byte)

URL recovery attack

Disclose authentication token



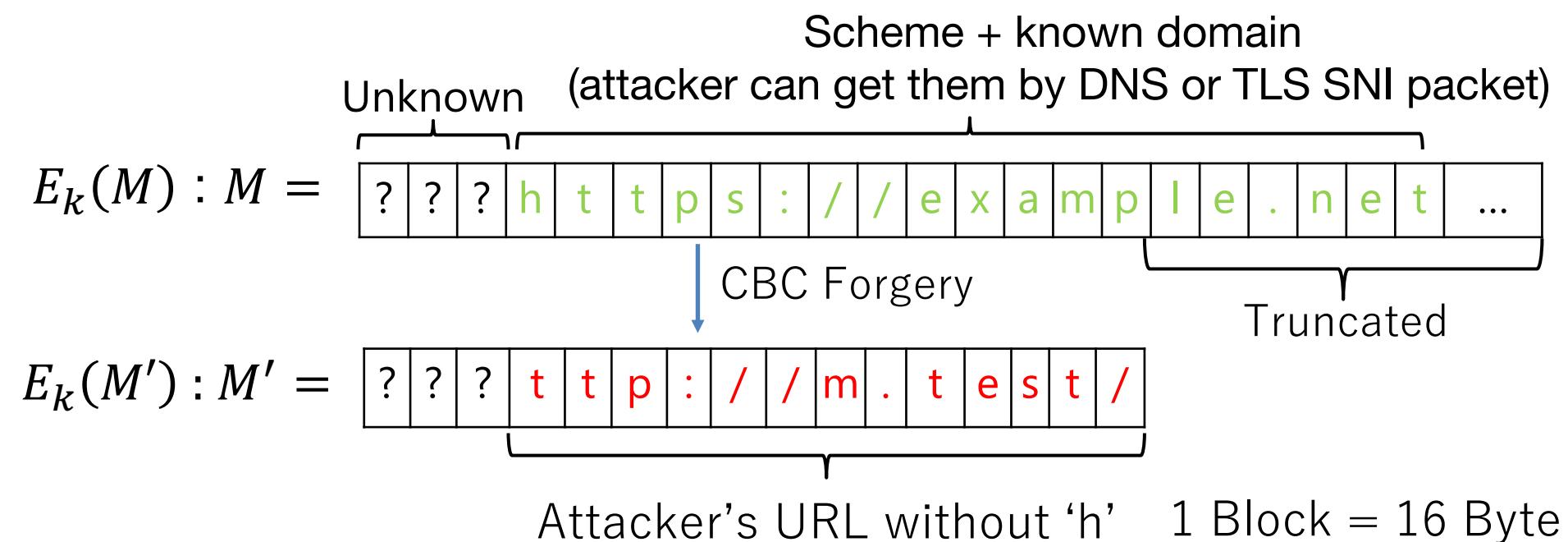
Link Preview Oracle Attack

Attack overview

Attacker recover an encrypted message **before** the encrypted URL.

It works like a padding oracle attack.

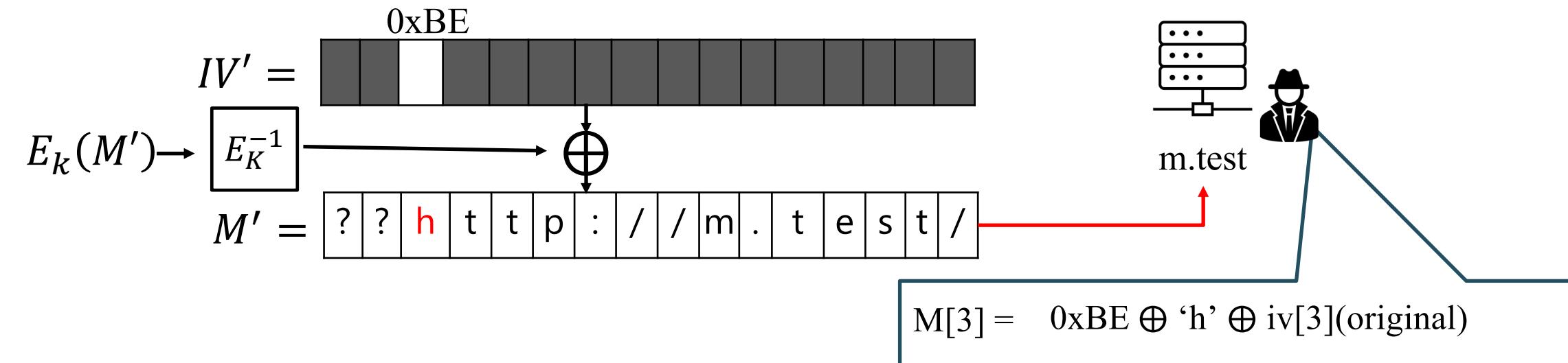
Step1. Modify the encrypted via a CBC malleability,
producing a **partially attacker-controlled URL**



Link Preview Oracle Attack

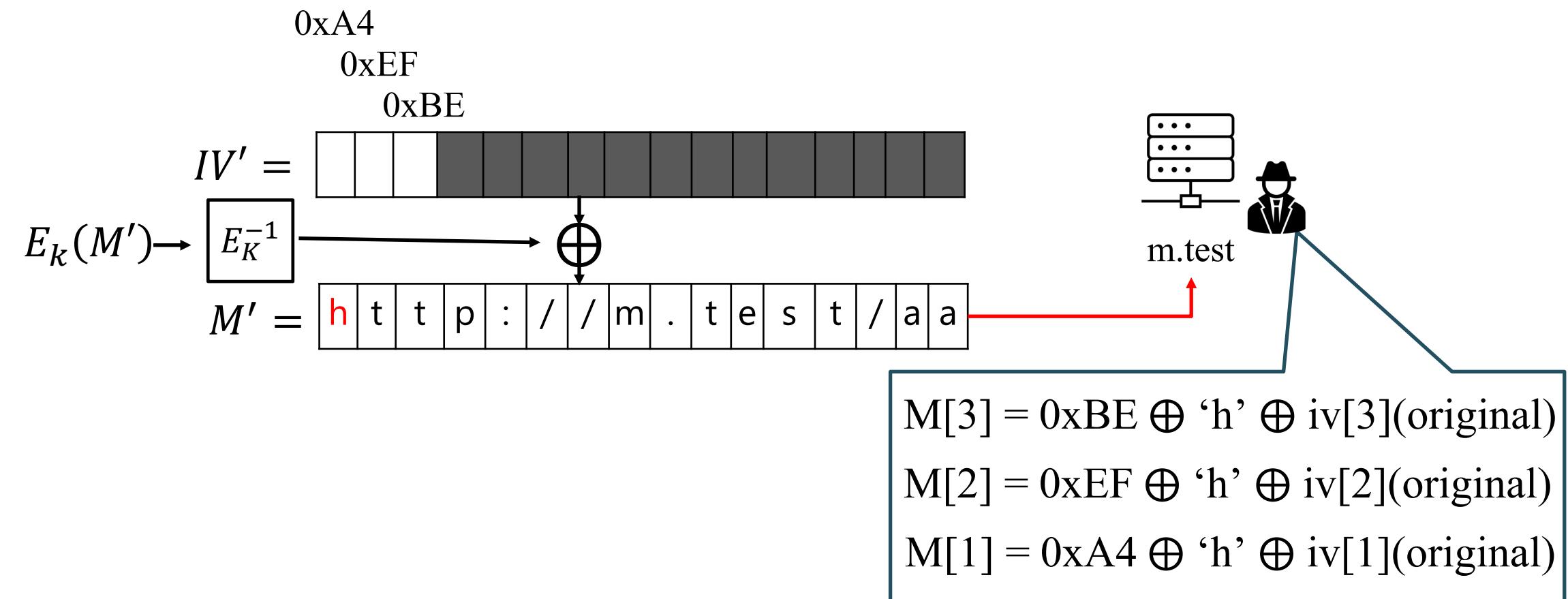
Step2. Seek an IV' such that the 3rd byte of M' becomes “h”.

When “h” appears, the client fires a link preview, allowing the attacker to detect $'h' \leftarrow IV'[3](0xBE) \oplus E_K^{-1}(M')[3]$



Link Preview Oracle Attack

Step3. Repeat **Step 2** for the second and first bytes



*Index starts with 1

Takeaway : Ciphertext Confidentiality

- **Remark: SHOULD use Authenticated Encryption (AE)**
 - E.g., AES-GCM, ChaCha20-Poly1305
 - Don't use malleable encryption without MAC
- **SHOULD generate preview ONLY on the sender-side**

3 Takeaways : Whole of this presentation

1. Decentralized Architecture's Untapped Risks and Rewards

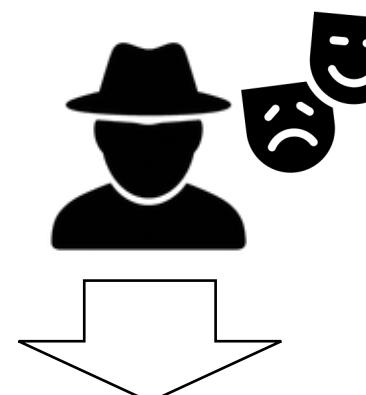
- Removing a central authentication server in Nostr brings new freedoms but also introduces subtle security pitfalls
- Multi-layered security are lost, and cryptographic weaknesses are immediately upgraded to practical attacks.

3 Takeaways : Whole of this presentation

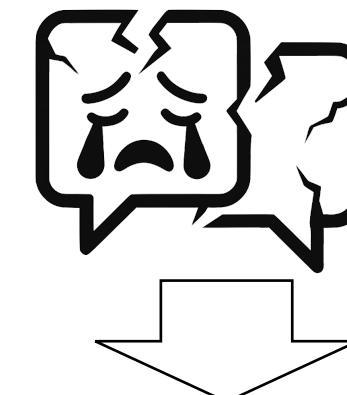
2. Hands-On Attacks & Immediate Mitigation

We guided our footsteps, and you learn how to destroy integrity & confidentiality

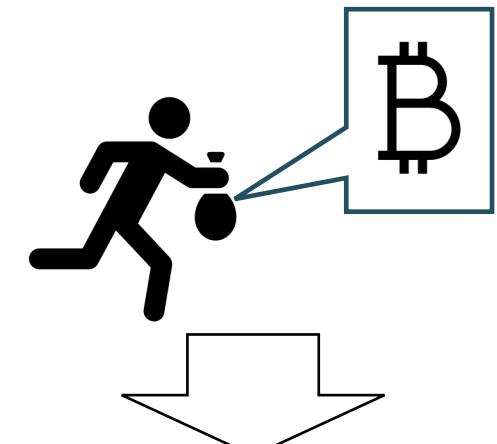
Identify the root cause and understand mitigation



- Signature verification
Bypass



- Lack of key separation
- Receiver-side preview generation



- Verification Bypass

3 Takeaways : Whole of this presentation

3. Blueprint for Future-Ready Decentralized Systems

Items	Nostr	Blueprint
Signature	<p>Signing is mandatory. But there is no concrete specs for verifying.</p>	<p>Signing & verifying are mandatory</p>
Link Preview	<p>No specs (Mostly receiver-side generation)</p>	<p>Sender-side generation</p>
Public Key Authenticity	<p>No specs (NIP-05 Badge is available, but an authenticity is out of scope)</p>	<ul style="list-style-type: none">• Out-of-band authentication• Key Transparency

Summary

- First cryptographic deep-dive into Nostr, a distributed SNS.
- Find practical attacks caused by cryptographic & implementation flaw.
- Client is the trust anchor.
- Mandatory signature checks, key-separation, and AEAD.
- Responsible disclosure, and patches

Our Paper

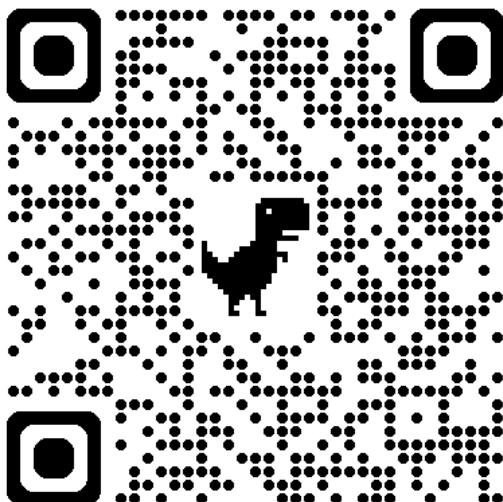
2025 10th IEEE European Symposium on Security and Privacy (EuroS&P)

Not in The Prophecies: Practical Attacks on Nostr

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Shogo Shiraki <i>University of Hyogo Hyogo, Japan 4w3tag185mpja@gmail.com</i>	Takanori Isobe <i>The University of Osaka Osaka, Japan takanori.isobe@ist.osaka-u.ac.jp</i>	

Abstract—Distributed social networking services (SNSs) recently received significant attention as an alternative to traditional, centralized SNSs, which have inherent limitations on user privacy and freedom. We provide the first in-depth security analysis of Nostr, an open-source, distributed SNS protocol developed in 2019 with more than 1.1 million registered users. We investigate the specification of Nostr and the client implementations and present a number of practical attacks allowing forgeries on various objects, such as encrypted direct messages (DMs), by a malicious user or a malicious server. Even more, we show a confidentiality attack against encrypted DMs by a malicious user exploiting a flaw in the link preview mechanism and the CBC malleability. Our attacks are due to cryptographic flaws in the protocol specification and client implementation, some of which in combination elevate the forgery attack to a violation of confidentiality. We verify the practicality of our attacks via Proof-of-Concept implementations and discuss how to mitigate them.

Index Terms—Nostr, plaintext recovery attack, forgery attack, key replace attack, Cache-based Forgery Attack, CBC-mode



<https://crypto-sec-n.github.io/>