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BRIEFINGS

# URB Excalibur: The New VMware All-Platform VM Escapes

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# Who are we?

Security researchers at Ant Group Light-Year Security Lab

Escaped from virtual machine many times

Won the Pwnie Awards🦄 at 2023



蚂蚁安全实验室  
Ant Security Lab

蚂蚁光年  
Ant Light-Year

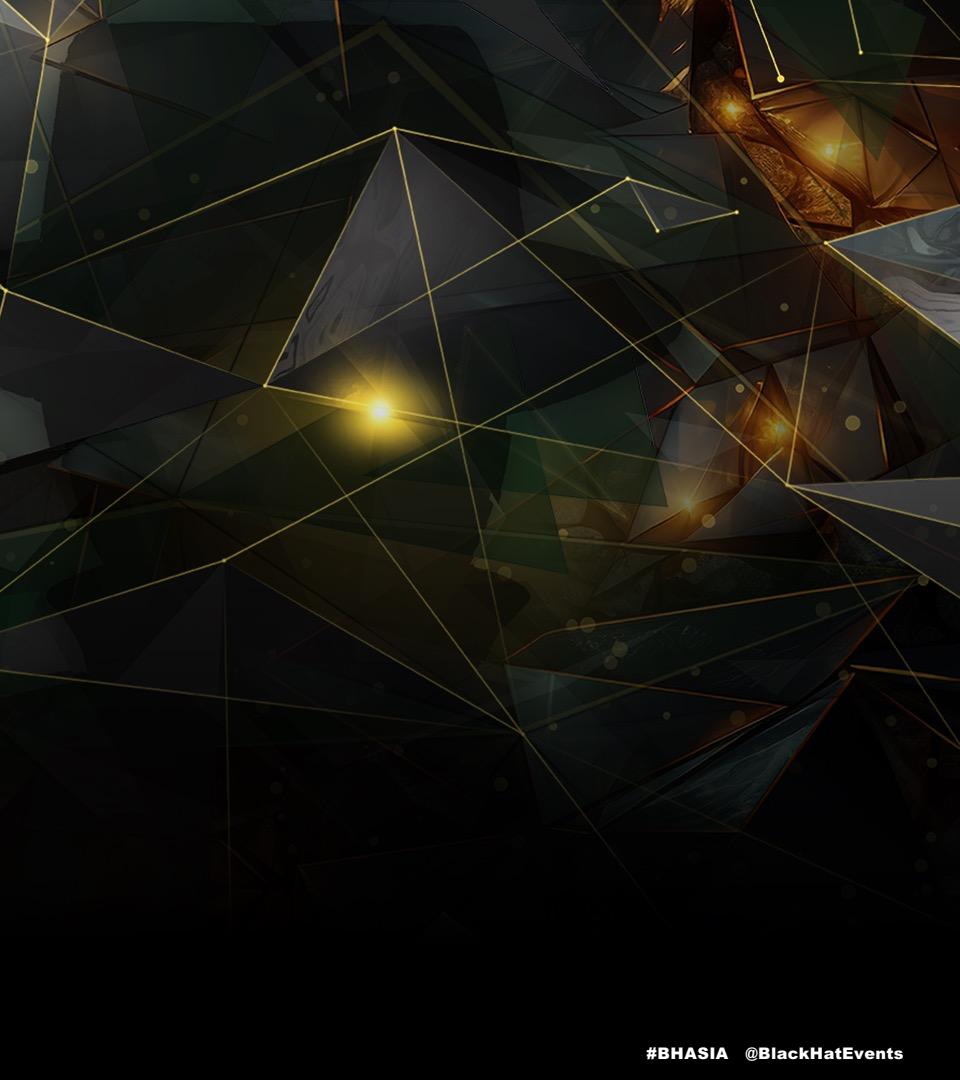
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# Talk Roadmap

- 1. Introduction**
- 2. A journey of finding vulnerabilities in VMware's hypervisor**
- 3. Exploit development of VMware VM escape**

# Introduction



# What is Virtual Machine escape and the danger of it

- Escape from the isolation sphere
- Take control over the whole hypervisor
- Network escape

**One of the most catastrophic threats to the Cloud**

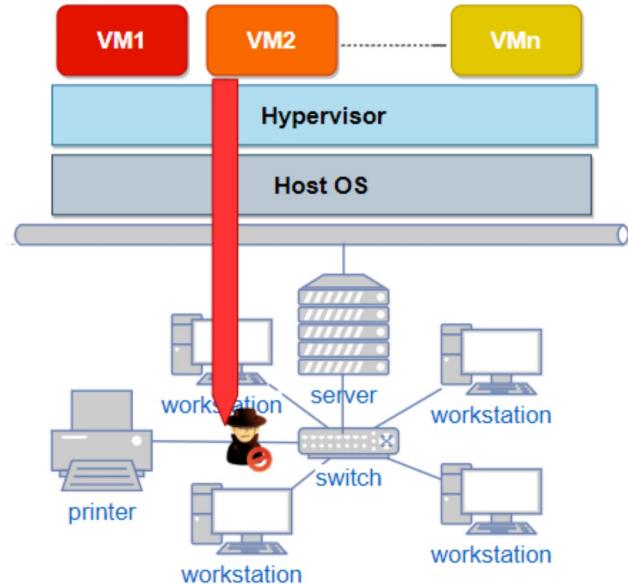


Fig. 6. Demonstration of VM Escape to Network.

# VMware's Architecture

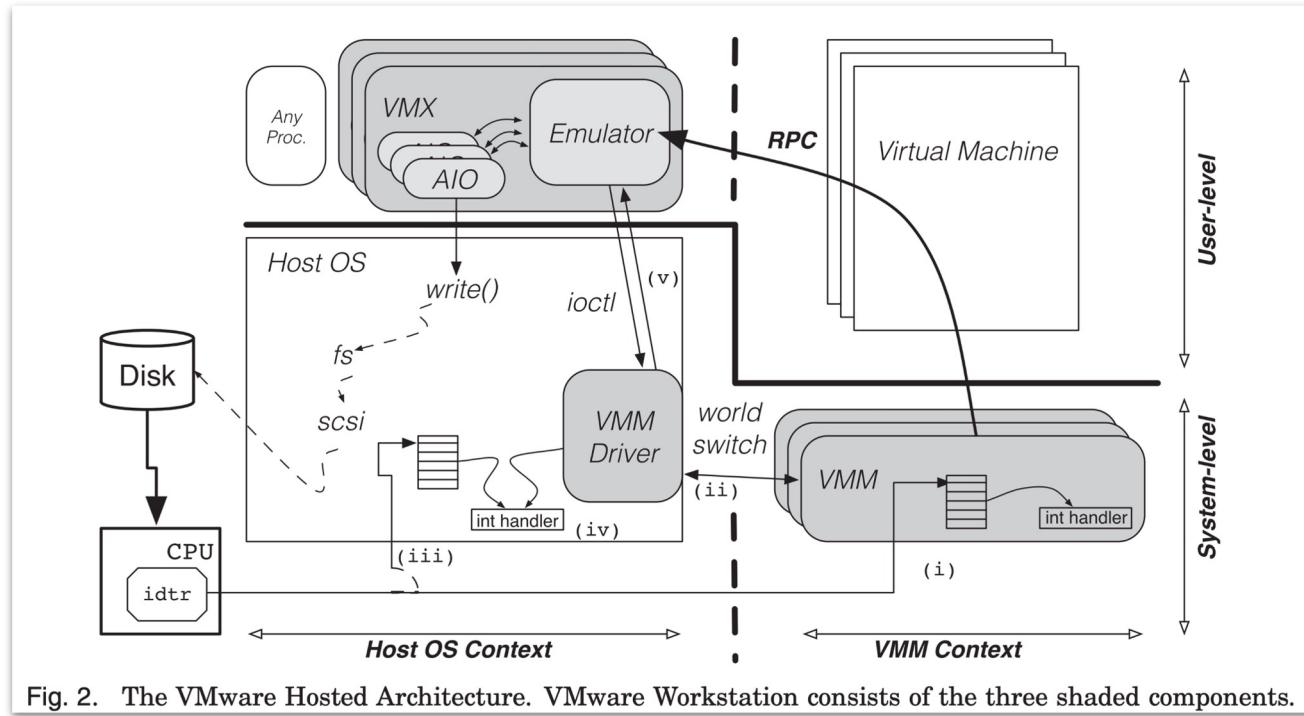


Fig. 2. The VMware Hosted Architecture. VMware Workstation consists of the three shaded components.

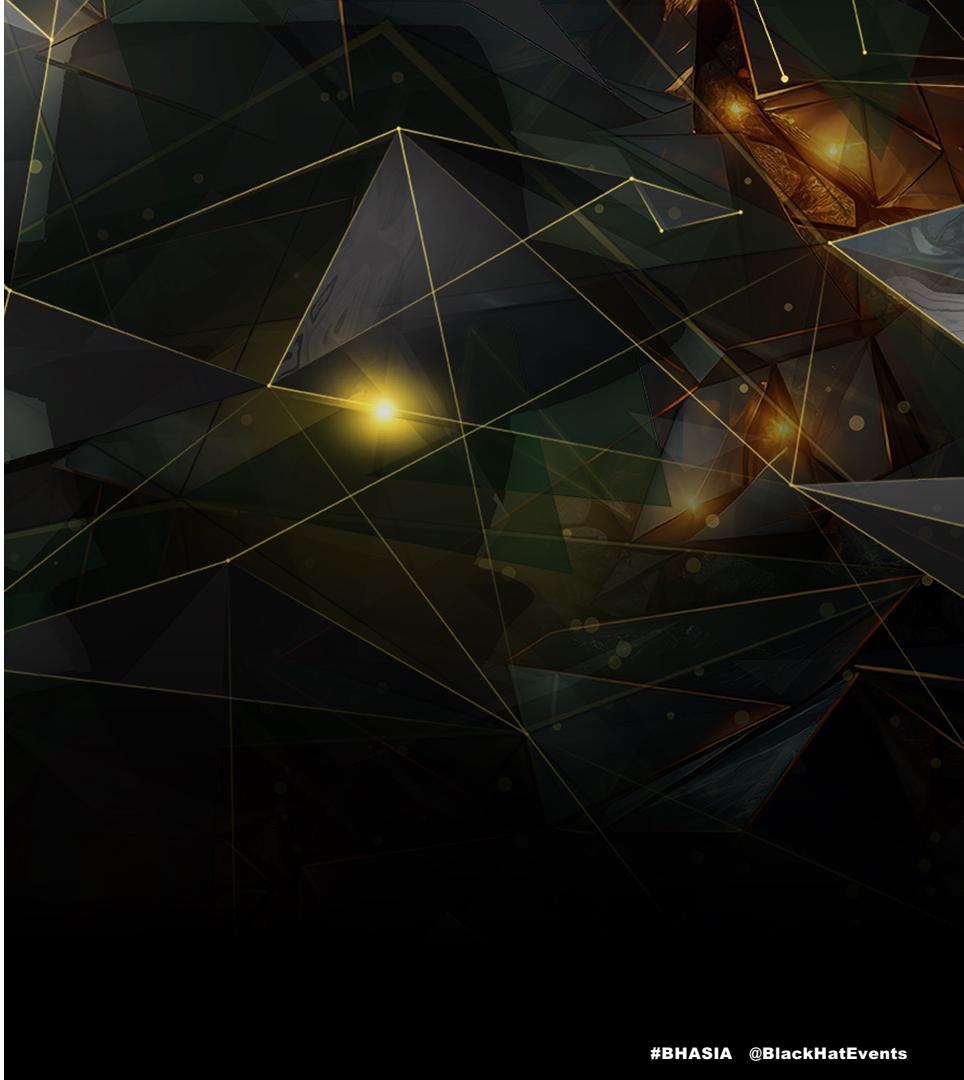
# VMware hypervisors' attack surface



Hard Disk	LSI Logic NVMe	
Network Adapter	E1000/E1000e VMXNET3	
Virtual Device	USB Controller	UHCI Tianfu Cup 2021 Workstation (CVE-2021-22041), Tianfu Cup 2023 Workstation (CVE-2024-22253, CVE-22255)
		EHCI GeekPwn 2022 Fusion (CVE-2022-31705)
		XHCI Tianfu Cup 2021 ESXi (CVE-2021-22040), Tianfu Cup 2023 ESXi (CVE-2024-22252)
	USB Device	HID (mouse) Bluetooth Pwn2Own 2023 Workstation (CVE-2023-20869, CVE-2023-20870)
		...
	GPU	SVGA 2D SVGA 3D
	Sound Card	ES1371
	TPM	vTPM
		...
	GuestRPC	Backdoor
	VMM	

# Vulnerability Discovery

A journey of finding  
vulnerabilities in VMware's  
hypervisor



# Start vulnerability discovery in VMware

First encounter with VMware, closed-source hypervisor

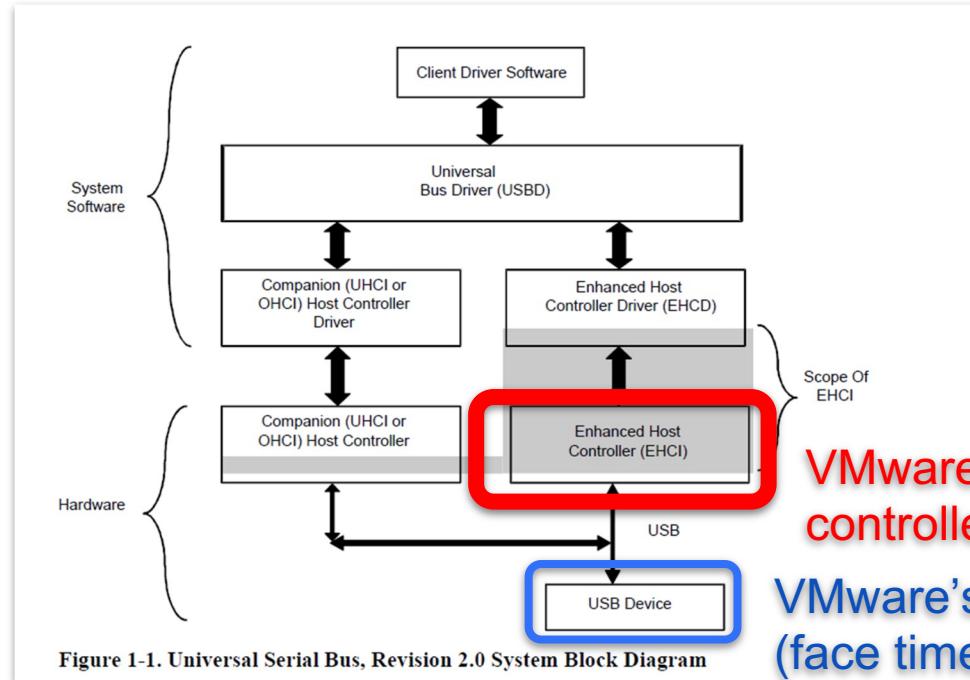
## 1. Focusing on an interesting and potentially risky attack surface

- Having studied QEMU EHCI vulnerabilities
- Interested in VMware's EHCI implementation

## 2. Reverse engineering

- Using string search as an entry point
- Understanding EHCI specification and QEMU code while reverse engineering VMware

# EHCI / USB 2.0 Controller



VMware's virtual EHCI  
controller

VMware's virtual video device  
(face time)

# EHCI / USB 2.0 Controller

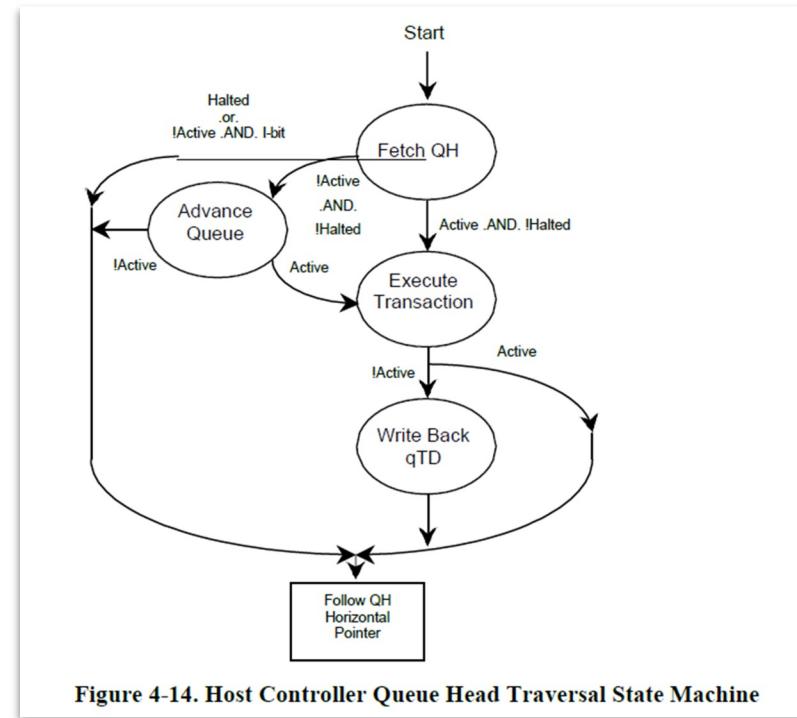
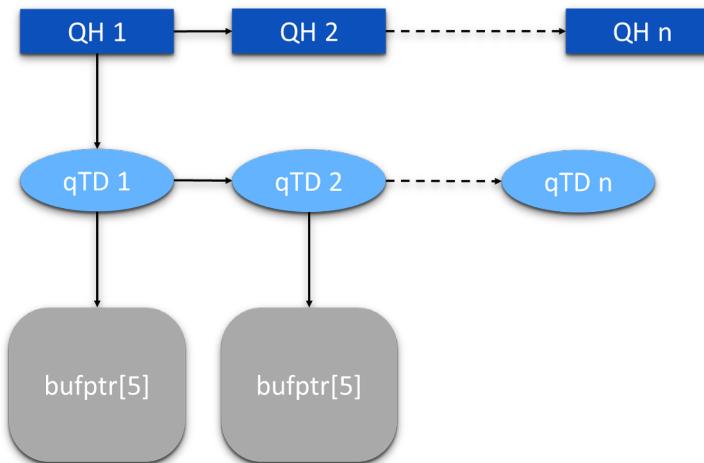


Figure 4-14. Host Controller Queue Head Traversal State Machine

# EHCI / USB 2.0 Controller

## Endpoint/Pipe:

- Control
- Bulk
- Interrupt
- Isochronous

## Token:

- Setup
- In: Device -> Software
- Out: Software -> Device

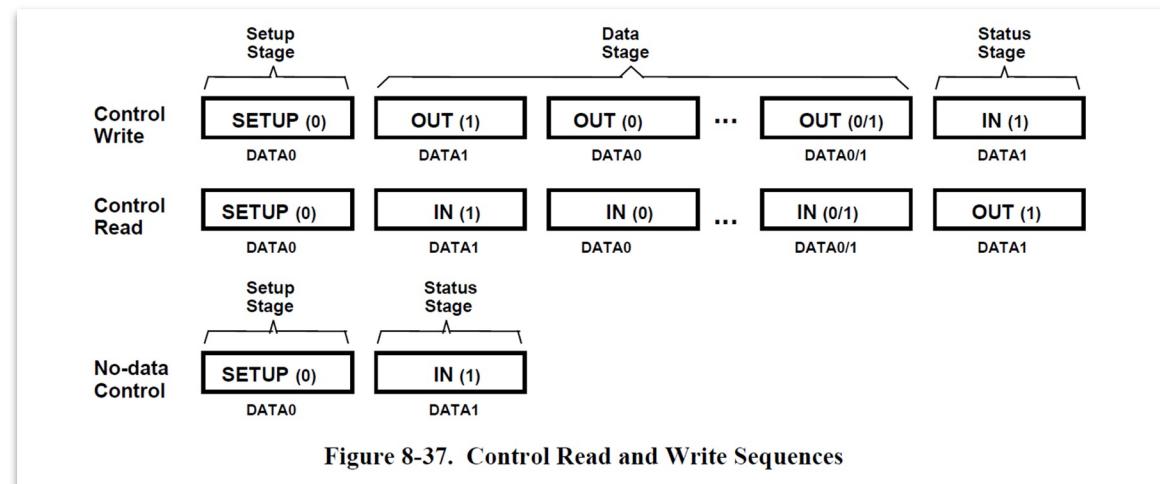


Figure 8-37. Control Read and Write Sequences

# How the data flow

QEMU

ehci\_state\_fetchqh

ehci\_state\_fetchqtd

usb\_handle\_packet

CVE-2020-14364

do\_token\_setup

do\_token\_in

do\_token\_out

Could there be  
bugs here?

YES!

VMware

ehci\_control\_transfer

while (1)  
setup  
in  
out

VUsb\_NewUrb

urb\_submit

# CVE-2022-31705

## 3. Heap out-of-bounds write vulnerability in EHCI controller (CVE-2022-31705)

### Description

VMware ESXi, Workstation, and Fusion contain a heap out-of-bounds write vulnerability in the USB 2.0 controller (EHCI). VMware has evaluated the severity of this issue to be in the [Critical severity range](#) with a maximum CVSSv3 base score of [9.3](#).

### Known Attack Vectors

A malicious actor with local administrative privileges on a virtual machine may exploit this issue to execute code as the virtual machine's VMX process running on the host. On ESXi, the exploitation is contained within the VMX sandbox whereas, on Workstation and Fusion, this may lead to code execution on the machine where Workstation or Fusion is installed.

### Resolution

To remediate CVE-2022-31705 apply the patches listed in the 'Fixed Version' column of the 'Response Matrix' found below.

### Workarounds

Workarounds for CVE-2022-31705 have been listed in the 'Workarounds' column of the 'Response Matrix' below.

### Additional Documentation

None.

### Acknowledgements

VMware would like to thank the organizers of GeekPwn 2022 and Yuhao Jiang for reporting this issue to us.

### Notes

None.

# CVE-2022-31705

urb's size =  
0x98 + 8 + setup\_len

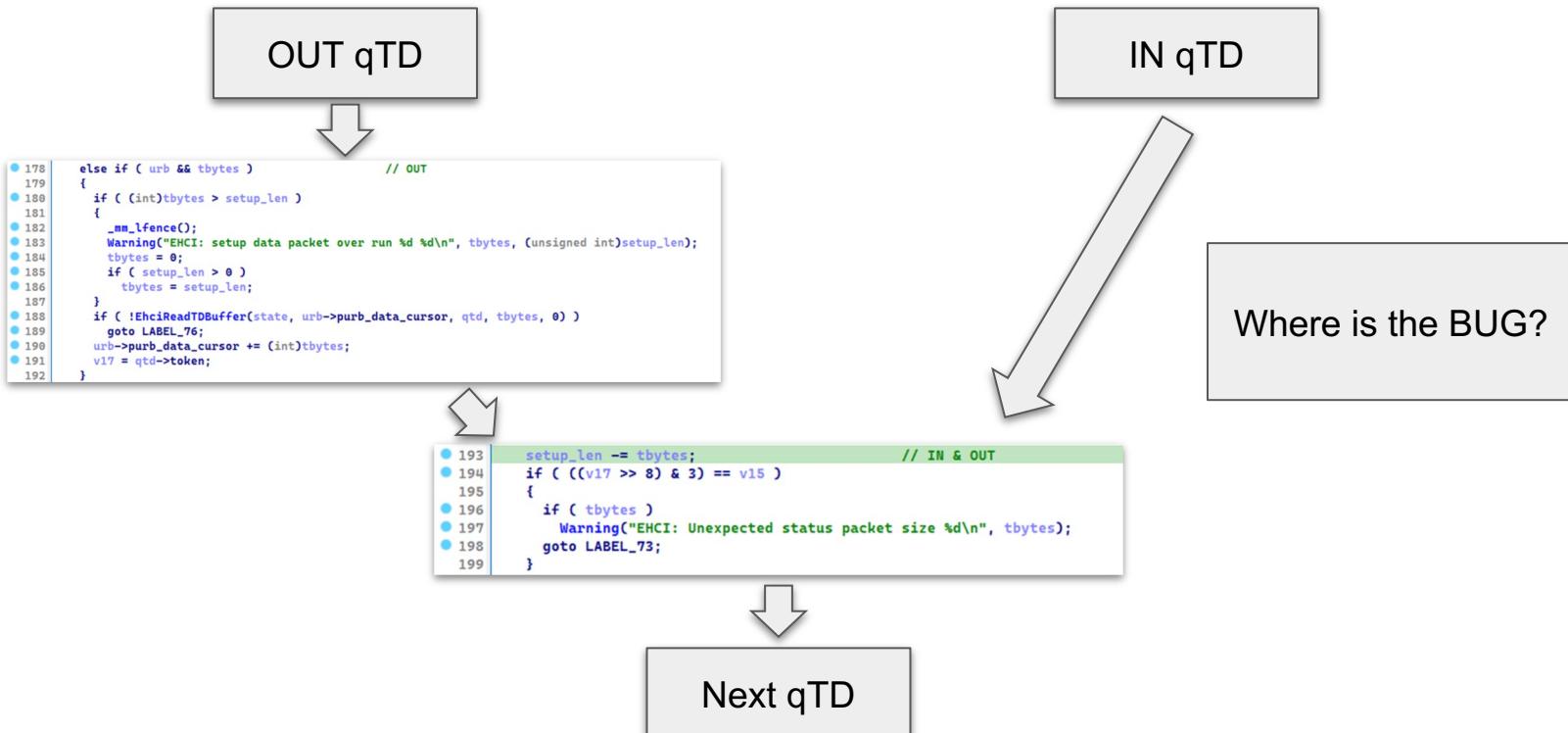
```
1 void __fastcall ehci_control_transfer(__int64 state, pipe *pipe, EHCIqh *qh)
2 {
3     // [COLLAPSED LOCAL DECLARATIONS. PRESS KEYPAD CTRL- "+" TO EXPAND]
4
5     urb_link_first = pipe->urb_link_first;
6     p_next_qtd = &qh->next_qtd;
7     urb = @i64;
8     v6 = 0;
```

SETUP qTD

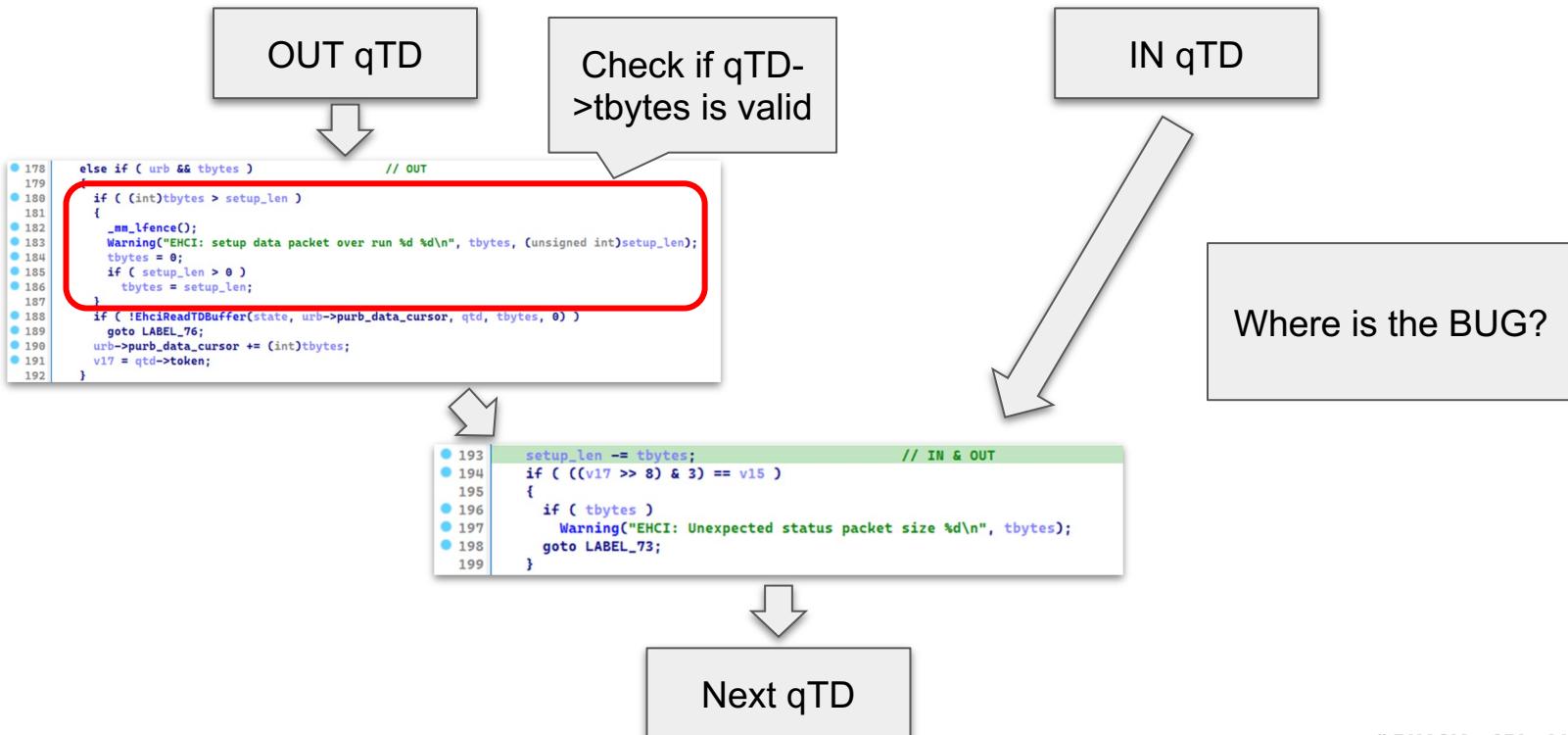
```
141     v14 = 0;
142     setup_len = *(unsigned __int16 *)&setup_buf[6];
143     data_len = *(unsigned __int16 *)&setup_buf[6] + tbytes;
144     v32 = v14;
```

```
157     urb = VUsb_NewUrb(pipe, 0, data_len);
158     purb_data_cursor = urb->purb_data_cursor;
159     urb->interrupt_pid = v14;
160     v14 = 1;
161     urb->num_packets = 1;
162     urb->datalen = data_len;
```

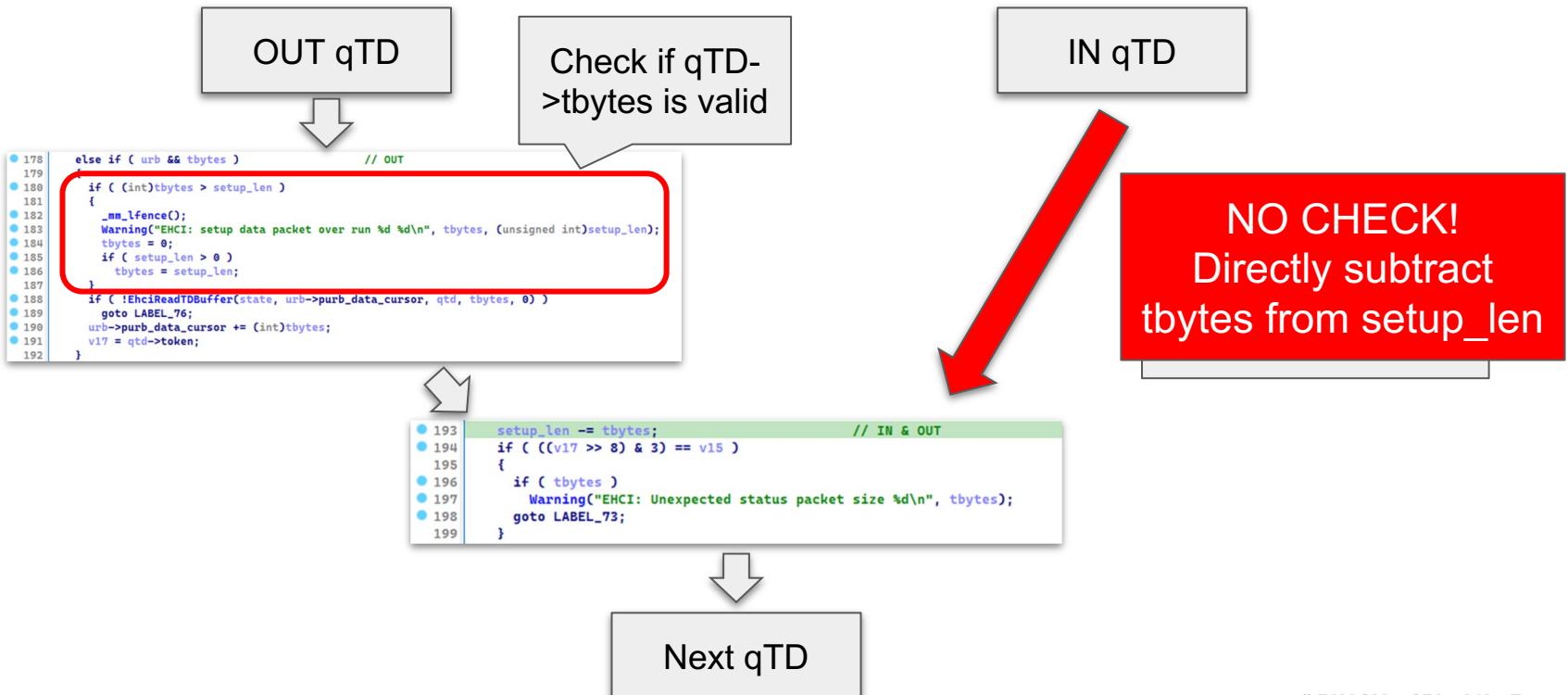
# CVE-2022-31705



# CVE-2022-31705

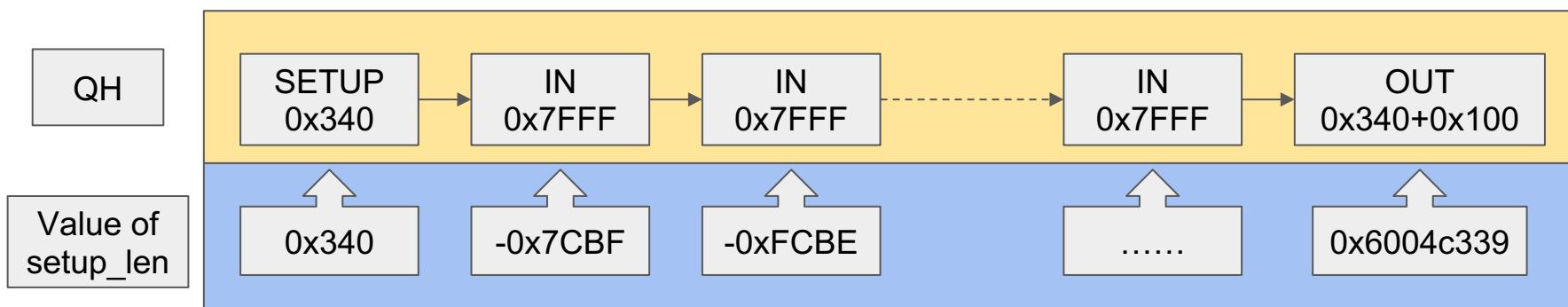


# CVE-2022-31705



# CVE-2022-31705

- Missing tbytes check when handling IN qTD
- setup\_len downward integer overflow



- setup\_len is much larger than the size of urb
- Use OUT qTD to obtain heap out-of-bounds write

# What else did we find? BUG 1: Out-of-bounds read vulnerability

- Pipe type confusion (Control  $\leftrightarrow$  ISOC)
- Handle urb incorrectly

```
34     pipe = *(v10 + 8 * (v9 & 0xF | (16 * ((v9 >> 7) & 1))) + 0x10);
35     if ( !pipe )
36         return 0;
37     }
38     if ( !sub_1F44B0(state, pipe, index, &purb_data_cursor, &cur_packet_len) )
39         return 1;
```

```
34     if ( (urb->cur_packet - urb - 0x98) / 12 < urb->num_packets )
35     {
36         cur_packet = urb->cur_packet;
37         _cur_packet_len = cur_packet_len;
38         *cur_packet_len = cur_packet->len;
39         *purb_data_cursor = urb->purb_data_cursor;
40         v15 = 8 * (1 << (10 - ((*(state + 1880) + 16
41         if ( (v15 & (index - cur_packet->maybe_id)) >
42             goto LABEL_21;
43         urb->purb_data_cursor += cur_packet->total_transfer_length;
44         pipe->all_size -= urb->datalen;
45         --pipe->num_packets;
46         ++urb->cur_packet;
47         if ( cur_packet->maybe_id == index )
48             return 1;
```

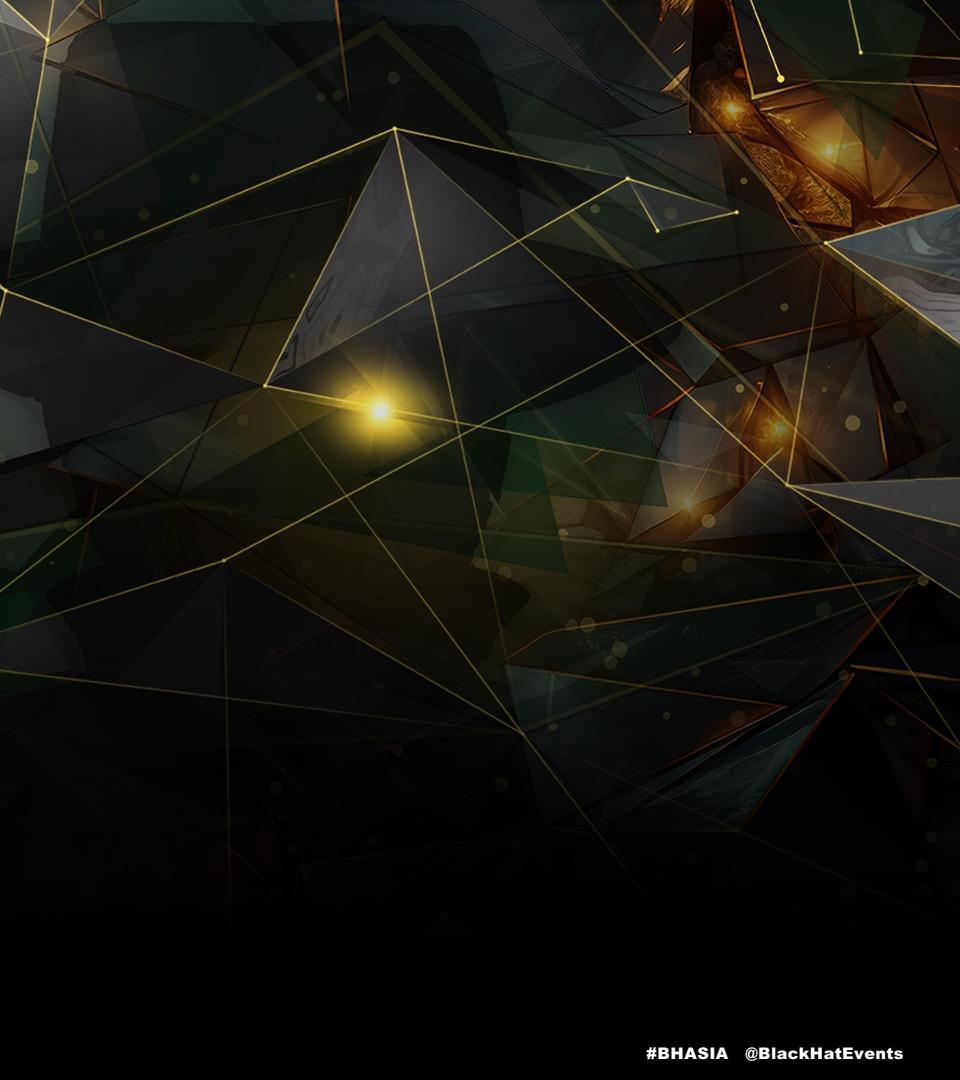
Always be 1 in  
control pipe's urb

# What else did we find?

## BUG 2: Information disclosure vulnerability

- In many virtual USB devices (USB Audio, USB Video, USB RNG...)
- No memset, writeback\_len is set to the data size of urb.

# Exploit Development



# The problem

- [Again] Closed-source
- No public exploit code and rarely disclosed exploit flow
- Most of past exploit primitives have been patched
- Few code paths that can be controlled with in the guest OS.

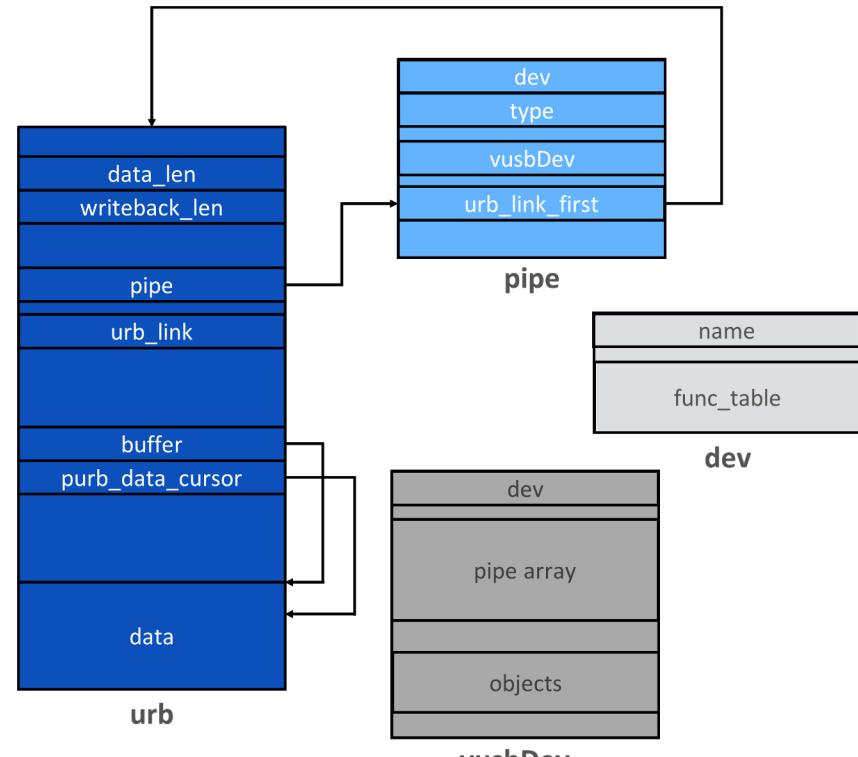
# Some patches for old primitives

- DnD/CP objects in **backdoor module** (2017)
  - a. VMware remove dynamic allocation and release of DnD/CP objects
- ResourceContainer in **SVGA backend module** (2018)
  - a. VMware first removed the function pointer table in the ResourceContainer in 15.5.7
  - b. VMware moves SVGA backend module into sandbox (mksSandbox) in 16
- GMR in **SVGA front-end module** (2021)
  - a. VMware adds a check at the head of GMR chunk (MKSMemMgrSafeMalloc)
- .....

# URB: Powerful Excalibur

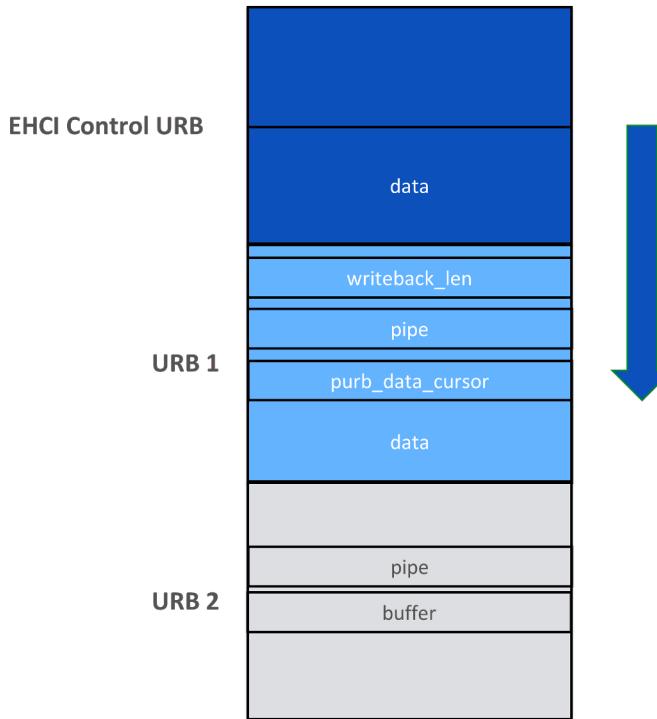
- USB Request Block
- Used by all virtual USB controllers

- Dynamic allocate and free
- Has:
  - A variable length data array
  - A member to control length to read
  - A data pointer
  - A pipe pointer
- ...



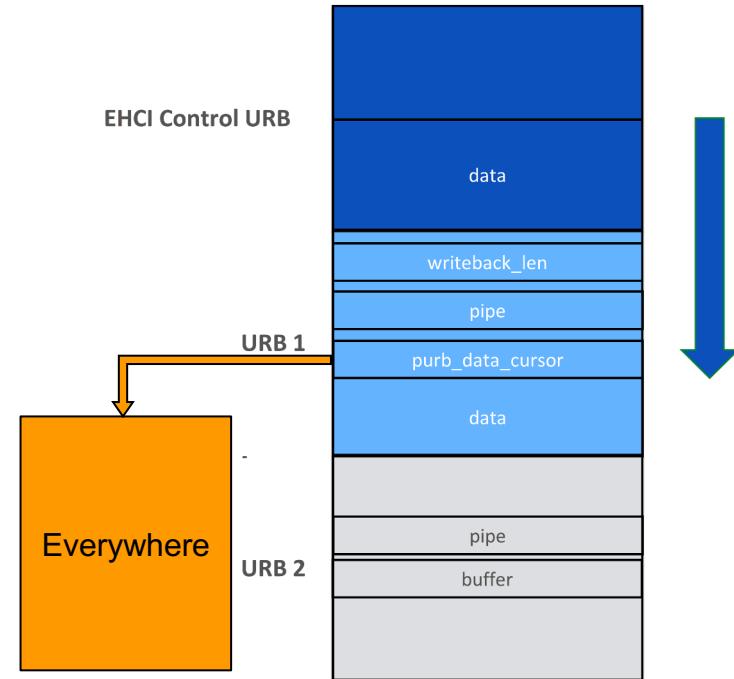
# Out-of-bounds write -> Out-of-bounds Read

1. Allocate URB1 and URB2, leaving space for EHCI Control URB
2. Allocate EHCI Control URB, then overwrites writeback\_len of URB1
3. Read back URB1, we can read the buffer address and pipe address



# Arbitrary Address Read

1. Allocate EHCI Control URB again
2. This time overwrite  
`purb_data_cursor` to any location
3. Read back URB1



# Arbitrary Address Write

```

1 char __fastcall uhci_check_and_writeback(__int64 a1, pipe *pipe)
2 {
3     // [COLLAPSED LOCAL DECLARATIONS. PRESS KEYPAD CTRL- "+" TO EXPAND]
4
5     p_urb_link_first = &pipe->urb_link_first;
6     v48 = p_urb_link_first;
7     while ( 2 )
8     {
9         urb = 0i64;
10        if ( *p_urb_link_first != p_urb_link_first )
11            urb = &(*p_urb_link_first)[-3].prev;
12        if ( !urb )
13            return 0;
14        _pipe = urb->pipe;
15        if ( urb->writeback_flag != 2 )
16            return 0;
17        idx = _pipe->idx;
18        v47 = idx;
19        if ( idx >= _pipe->num_frames )
20            return 0;
21        for ( i = idx; ; ++i )
22        {
23            frame = &_pipe->frames[i];
24            qh = frame->mem_and_qh.qh;
25            if ( !qh || (*qh + 1) & 0xFFFFFFFF0 != *frame->mem_and_qh.mem )
26                goto LABEL_37;

```



```

142     ++v20;
143     while ( v19 < _pipe->num_frames );
144 }
145
146
147 *((&frame->mem_and_qh.qh + 4i64) = *((frame->mem_and_qh.mem + 1);
148 frame->td.ctrl &= ~0x800000u;
149

```

- Write from a pointer in frame to another pointer in frame
- frame is a member in pipe
- We can fake the pipe in urb using out-of-bounds write



001EFB98 088	mov rax, [rbx]
001EFB9B 088	mov rdx, [rbx+8]
001EFB9F 088	mov rcx, [rax+8]
001EFBA3 088	mov eax, [rcx]
001EFBA5 088	mov [rdx+4], eax
001EFBA8	

# Control the RIP

1. A dynamically allocated object that holds function pointers
2. We can trigger a call to the function pointer



pipe

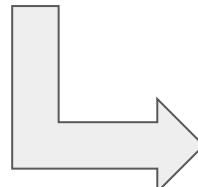
```
1 __int64 __fastcall cancel_pipe(pipe *pipe)
2 {
3     // [COLLAPSED LOCAL DECLARATIONS. PRESS KEYPAD CTRL- "+" TO EXPAND]
4
5     error(6, "UsbDev: DevID(%I64x): Cancel pipe(%p).\n", *(pipe->vusbDev + 51), pipe);
6     (*(*(*pipe->vusbDev + 40) + 16i64) + 16i64))(pipe->vusbDev, pipe->endpt);
7     urb_link_next = pipe->urb_link_first;
8     result = 0i64;
9     if ( urb_link_next != &pipe->urb_link_first )
10    {
```

## Control the RIP: Path 1

- The pipe when calling `cancel_pipe` in `ehci_check_and_writeback` comes from the pointer of urb
  - We can use out-of-bounds write to forge the `urb->pipe` to implement arbitrary address calls.

```
1 char __fastcall ehci_check_and_writeback(__int64 state, EHCIqh *qh, urb *_urb)
2 {
3 // [COLLAPSED] LOCAL DECLARATIONS. PRESS KEYPAD CTRL- "+" TO EXPAND]
4
5 pipe = _urb->pipe;

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94
95 v13 = qtd->token & 0x300;
96 v43 = _writeback_len;
97 if ( v13 != 0x100 )
98 {
99     if ( !*&pipe->type && v13 == 0x200 && !ehci_checkModified_setup(state, urb, qtd) )
100    {
101        cancel_pipe(pipe);
102        return 0;
103    }
104    goto LABEL_35;
105 }
```



# Control the RIP

## Path 2

- Fake a new pipe directly in vusbDev by arbitrary address write



**Use Path 2 when we can't reserve EHCI urb, although it needs more actions**

# What's more? We need heap grooming

Heap spraying and grooming primitive: **SVGA\_3D\_CMD\_SET\_SHADER**

Allocate and free in large quantities, the heap size is `sizelnBytes+8`

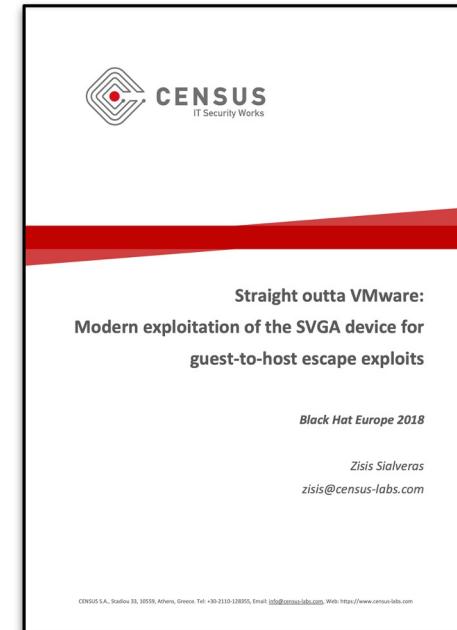
```
svga_3d_cmd_define_gb_shader(shid, SVGA3D_SHADERTYPE_MIN, sizelnBytes);
```

```
svga_3d_cmd_bind_gb_shader(shid, mobid, 0);
```

```
svga_3d_cmd_set_shader(cid, SVGA3D_SHADERTYPE_MIN, shid);
```

```
svga_3d_cmd_destroy_gb_shader(shid);
```

<https://census-labs.com/media/straightouttavmware-wp.pdf>



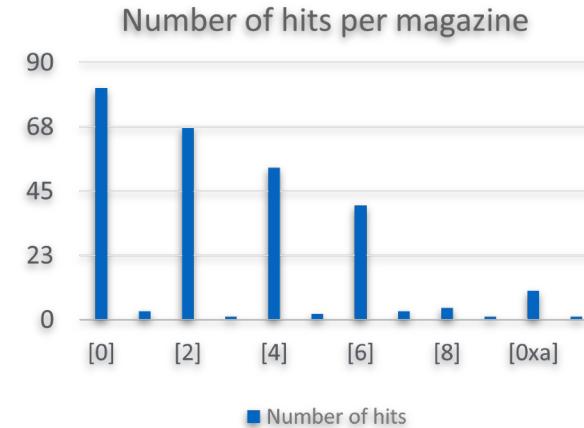
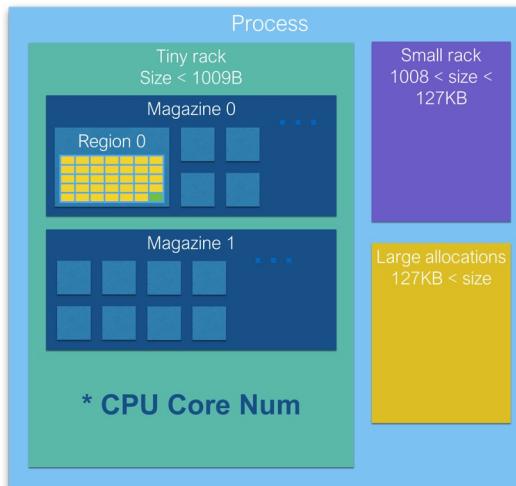
# Try on the VMware Fusion!

1 out-of-bounds read, 3 arbitrary address reads, and 2 arbitrary address writes

1. Heap grooming
2. Leak pipe address and heap address
3. Leak the program base address (pipe->dev)
4. Leak ehci state address (in .data)
5. Leak vusbdev address (in ehci state)
6. Write the upper 4 bytes of the fake pipe to vusbdev
7. Write the lower 4 bytes of the fake pipe to vusbdev
8. Trigger cancel pipe
9. Escape

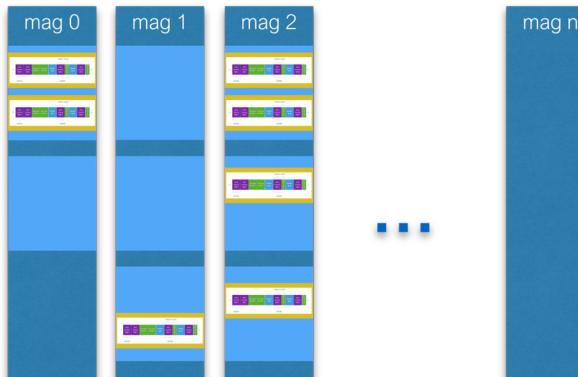
# Big problem. Magazine

- MacOS's libmalloc uses magazines to manage heap blocks
- Each CPU core will have a unique corresponding magazine



# Big problem. Magazine. How we deal with it

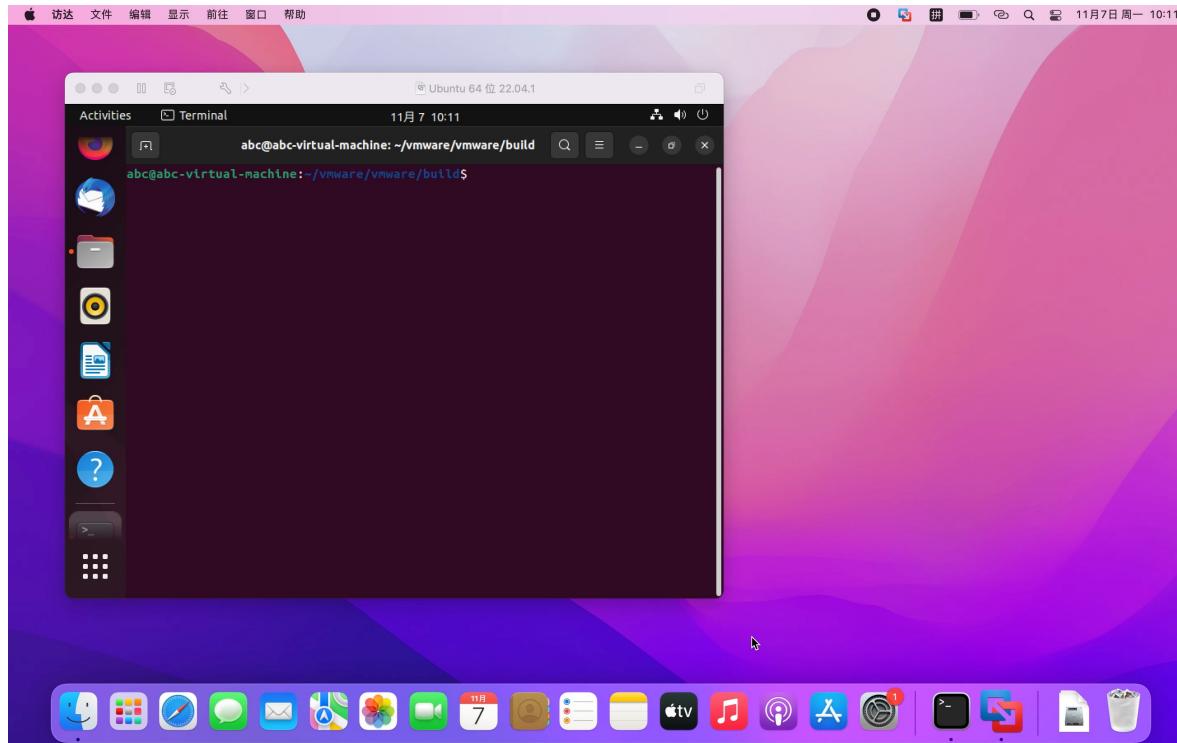
- Repeat the basic heap layout, and try to have at least one layout on each magazine
- Try a large number of times for every step (place objects, do oob read...)
- How to ensure that all magazines are occupied?
  - Add sleep between each allocation X
  - Increase cpu's occupancy and try to increase cpu core switching X



- Remove sleep, speed up exploit
- Use a huge number of spray rounds (0x1000)

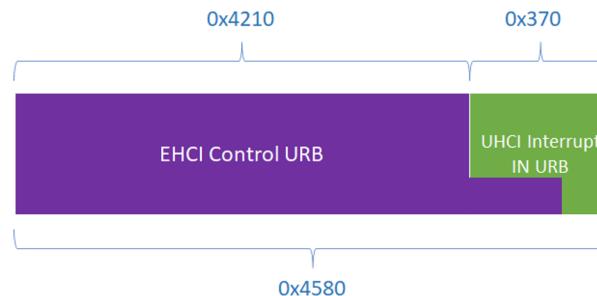
**Success rate > 80%**

# Demo



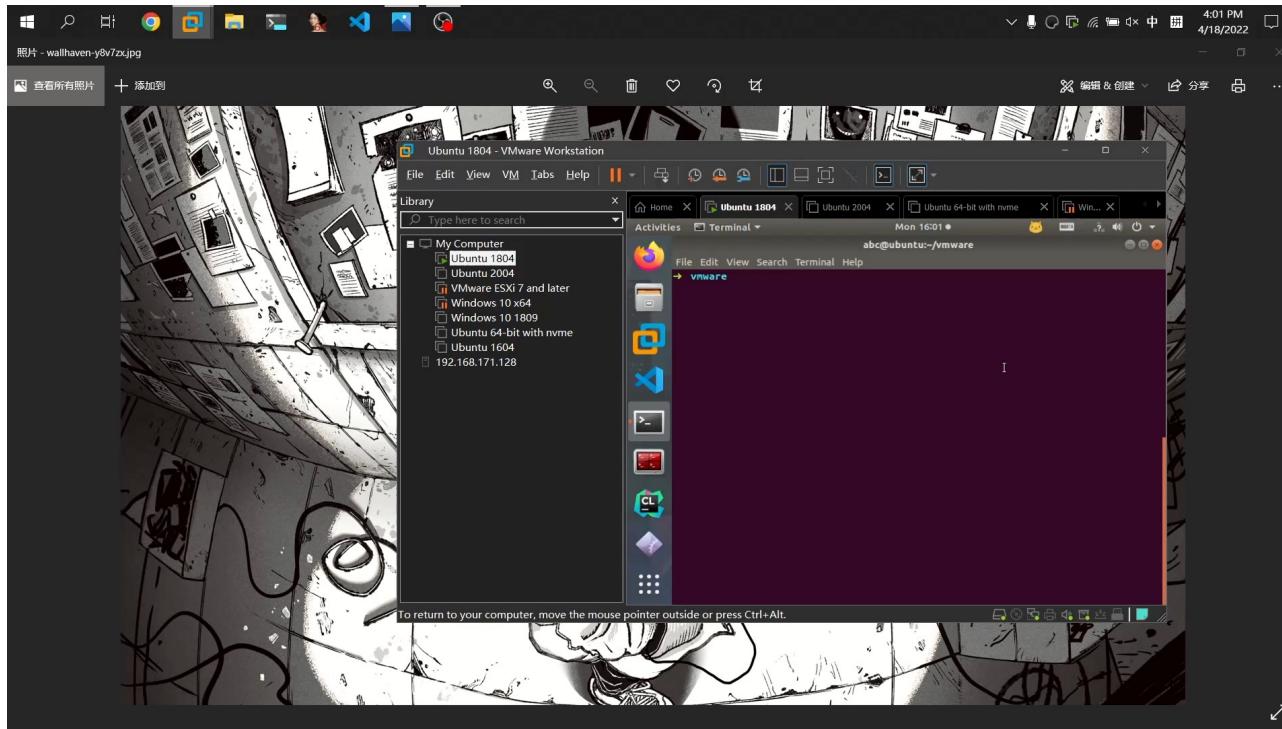
# On VMware Workstation

- In the default configuration, there will be no device on the EHCI
  - Plug in a usb device to connect to ehci
- To avoid the randomization of LFH:
  - Use chunks larger than 0x4000
  - Select a size that has not been used by LFH when we can't allocate larger than 0x4000



1. Leak heap address
2. Leak process base address
3. Leak the address of createProcessW (KERNEL32.dll)
4. Call WinExec

# Demo



## On ESXi

- Same as Workstation, no default device on EHCI
- Similar to CentOS 7, use very old glibc-2.17 (2.28 after ESXi 8.0.2)
- Basically the same as on Fusion (No need to face magazines)
- Use GMR instead of Shader

# Takeaways

- Where bugs have arisen with similar software, there may be new bugs
- When looking for exploit primitives, try to look for objects related to the vulnerability
- Virtual devices, especially USB-related devices, are now a popular attack surface

# Questions?