An idealistic treatment of quantifiers in categorical compositional distributional semantics
Extended Abstract

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Abstract. We show how one can formalise quantifiers in the categorical compositional distributional model of meaning. Our model is based on the generalised quantifier theory of Barwise and Cooper. We develop an abstract compact closed semantics and instantiate it in the category of relations.

1 Quantification in Sets

The category Rel of sets and relations is a basic example of a symmetric monoidal compact closed category. The tensor is Cartesian product, $\times$, and monoidal the unit is the one element set, 1, $A^* = A$. Monoidal closure is therefore just the product. As an example consider the correspondence between relations and subsets of the Cartesian product expressed formally as

$$\mathrm{Rel}(1, X \times Y) \;\cong\; \mathrm{Rel}(X, Y) \;\cong\; \mathrm{Rel}(X \times Y, 1)$$

where we are repeatedly using compact closure and the fact that relation beween 1 and Z are demonstrably subsets of Z.

O.R. This is probably best put where compact closed cats are introduced.

The frobenius maps Δ is just the diagonal relation $x \sim (x, x)$ and ∇ its transpose; $\perp : 1 \longrightarrow X$ is the so-called *fan* relation: $* \sim x,\ \forall x$, $\top$ its converse. The required axioms are easilly checked.

O.R. This is probably best put where Frobenius structure is introduced.

In order to introduce quantification in Rel, we first discuss the well known case in Set where the existential and universal quantifiers are left and right adjoints, respectively, to the inverse image functor. More formally this can be described as follows. For a set X, the set of functions from X to the two-element set, 2^X, is the powerset of X. For a function $f : X \longrightarrow Y$, precomposition with f is internally the map $2^f : 2^Y \longrightarrow 2^X$ that takes a h to hf. It has both left and right adjoints with respect to the pointwise subset ordering. In more detail, the powerset of X ordered by subset inclusion is isomorphic to the set of characteristic functions on X, 2^X, ordered pointwise by the ordering $0 < 1$ of 2. Any pre-ordered set can be seen as a category[1]. The function 2^f becomes a functor[2] from the category 2^Y to 2^X. This functor has both adjoints. It is well known that when f is the first projection $\pi_1 : X \times Y \longrightarrow X$, 2^{π_1} corresponds to *weakening* in the sense that it takes predicates over context X to predicates over the larger context, $X \times Y$, where the second component in Y doesn't occur. The left adjoint to 2^{π_1} interprets

[1] internal category?
[2] internal functor?

existential quantification, and the right adjoint interprets universal quantification. In more detail, the left adjoint, $\exists_f : 2^X \longrightarrow 2^Y$, is defined by

$$\exists_f(X' \subseteq X) = \{y \in Y \mid \exists x \in X. fx = y \wedge x \in X'\} \tag{1}$$

i.e. the image of X' under f. It is also known as the *direct image* of f. The right adjoint $\forall_f : 2^X \longrightarrow 2^Y$ is defined by

$$\forall_f(X' \subseteq X) = \{y \in Y \mid \forall x \in X. fx = y \implies x \in X'\}$$

There are other functions from $2^{X\times Y}$ to 2^X, such as the one sending an f to the g for which $g(x) = 1$ if and only if $f(x, y) = 1$ for exactly two distinct values of y. This quantifier could be called "two". Another example is "none" for which $g(x) = 1$ iff $f(x, y) = 0$ always. This exhibits the existential and universal quantifiers as two extreme cases of a spectrum of all possible quantifiers[3] .

In the following we assume for simplicity that the ambient context, X, is empty, i.e. $X = 1$, and then quantifiers become certain functions $2^Y \longrightarrow 2$. Everything we say generalises straightforwardly for arbitrary X. In this case, the existential quantifier becomes the characteristic function of nonempty subsets of Y; the universal quantifier is the characteristic function of the singleton $\{Y\}$, and "two" is the characteristic function of two-element sets.

In single sorted first order predicate logic, quantification ranges over the whole domain Y, which is simply assumed to be fixed, and often implicit, throught the formula. For example, in $\exists x.x > 0$, the bound variable, x, ranges over some ordered set. In linguistics, however, the range of quantification is explicitly stated, as in "all men sleep", where the word "men" restricts the range of the quantifier, "all", to just men. In other words, the quantifier is "all men", and "all" is a particle which expects a linguistic category to become a quantifier. Formally, when we fix a domain of all possible subjects of quantification, Y, such as all nouns, a quantifier Q must have type

$$Q \ : \ 2^Y \longrightarrow 2^{2^Y} \tag{2}$$

Here, the first argument to Q is the range of the quantification. The result is a quantifier which possibly ignores everything that is outside the range.

For instance, "some" ($\exists$) takes a subset $\mathrm{men} \subseteq X$ to all nonempty subsets of men; all ($\forall$) maps $\mathrm{men} \subseteq X$ to $\{\mathrm{men}\}$; "two" takes $\mathrm{men} \subseteq X$ to two-element subsets of men.

2 Relations

We show two alternative categorical views on relations. The first one understands a relation R between sets X and Y as a function into a powerset of Y, namely the function r that assigns to each x the set of all $y \in Y$ related to x. Formally, $r(x) = \{ \ y \in Y \mid x \sim_R y \ \}$. On the other hand the same relation R is a subset of the Cartesian product $X \times Y$ of those pairs (x, y) for which $x \sim_R y$. We now give the standard Categorical presentation of these two points of view on relations.

Relations as Kleisli arrows

The assignment $X \mapsto 2^X$ extends to a functor $2^{(-)} : \mathrm{Set} \longrightarrow \mathrm{Set}$ where the action on arrows is direct image (1). The functor is a monad where the unit, η, maps x to $\{x\}$ and multiplication, μ, is set union. Recall [2] that for any monad $(\mathbf{T}, \mu, \eta)$ on $\mathbb{C}$, the Kleisli category of $\mathbf{T}$, denoted $\mathbb{C}_\mathbf{T}$, has the same objects as $\mathbb{C}$. Arrows $X \longrightarrow Y$ in $\mathbb{C}_\mathbf{T}$ are arrows $X \longrightarrow \mathbf{T}Y$ in $\mathbb{C}$. Moreover, there is an adjunction

[3] Uhm, yes. That is precisely what it means to be a left and right adjoint.

$\mathbf{F} \dashv \mathbf{U} : \mathbb{C} \longrightarrow \mathbb{C}_{\mathbf{T}}$ such that $\mathbf{F}$ is identity on objects, and $f \mapsto \eta \cdot f$ on arrows and $\mathbf{U}$ is $\mathbf{T}$ on objects and $\mu \cdot \mathbf{T}f$ on arrows. Moreover, $\mathbf{UF} = \mathbf{T}$.

Thus by Rel one can understand the Kleisli category of $2^{(-)}$, which is understood as the category of sets and relations where a function $r : X \longrightarrow 2^Y$ defines a relation $\sim_R$ by

$$x \sim_R y \quad \equiv \quad y \in fx$$

Relations as spans in Set

On the other hand the same relation R defines a subset of $X \times Y$ by $R = \{ (x, y) \mid x \sim_R y \}$. Such a set is called the *tabulation of* R. In Set, R with the two projections is a *span* over X and Y:

$$\begin{tikzcd}
& R \arrow[dl, "p"'] \arrow[d, dashed, "\langle p,q\rangle"] \arrow[dr, "q"] & \\
X & X \times Y \arrow[l, "\pi_1"'] \arrow[r, "\pi_2"] & Y
\end{tikzcd}$$

Spans form a (bi)category where composition is defined using the pullback as in the following diagram:

$$\begin{tikzcd}
& & R \times_Y S \arrow[dl] \arrow[dr] & & \\
& R \arrow[dl, "p"'] \arrow[dr, "q"] & & S \arrow[dl, "r"'] \arrow[dr, "s"] & \\
X & & Y & & Z
\end{tikzcd}$$

where $R \times_Y S$ is the pullback of q and r. The unit is then just the span of identities:

$$\begin{tikzcd}
& X \arrow[dl, "\mathsf{id}"'] \arrow[dr, "\mathsf{id}"] & \\
X & & X
\end{tikzcd}$$

For an arbitrary category $\mathcal{C}$, the category of spans in $\mathcal{C}$, formally $\mathrm{Span}(\mathcal{C})$, is a bicategory where a 2-cell $R \Rightarrow S : X \leftrightarrow Y$ is the vertical arrow in the following commuting diagram:

$$\begin{tikzcd}
& R \arrow[dl, "p"'] \arrow[dr, "q"] \arrow[dd, dashed] & \\
X & & Y \\
& S \arrow[ul, "r"] \arrow[ur, "s"'] &
\end{tikzcd}$$

We leave out the details that this indeed forms a bicategory. We leave the compact closed monoidal structure of $\mathrm{Span}(\mathrm{Set})$ as an easy exercise.

Finally, note that the correspondence between spans and relations is an equivalence rather than an isomorphism, as there may be many spans representing (tabulating) the same relation. For instance, whenever $X \xleftarrow{p} R \xrightarrow{q} Y$ tabulates a relation, $X \xleftarrow{[p,p]} R + R \xrightarrow{[q,q]} Y$ tabulates the same relation. However, all such spans are equivalent via suitable 2-cells[4]. Thus, strictly speaking Rel is the quotient posetal[5] bicategory $\mathrm{Span}(\mathrm{Set})/_{\simeq}$ where $\simeq$ is the above-described equivalence.

2.1 Relations in Compact Closed Categories

The construction of the category of relations $\mathrm{Span}(\mathcal{C})$ can be generalised to an arbitrary category with pullbacks. The resulting category is compact closed if $\mathcal{C}$ is. Here are some of the details.

3 Powerset Relations

In this section we study relations over powersets of sets, i.e. relations between sets 2^X and 2^Y for some sets X, Y.

3.1 Definition

The following is the key observation of this section:

Proposition 1. *A relation between 2^X and 2^Y is equivalently a set T together with a pair of* relations $p : T \longrightarrow X$ *and* $q : T \longrightarrow Y$*, i.e. a span of relations, i.e. a relation in* Rel.

Proof. Given a relation R define $T = \{\, x \in 2^X, y \in 2^Y \mid x \sim_R y\}$. Such T comes with two projections $\pi_X : T \longrightarrow 2^X$ and $\pi_y : T \longrightarrow 2^Y$, which are relations in the first sense. In a diagram, in Rel (!):

$$\begin{tikzcd} & \{\, x \in 2^X, y \in 2^Y \mid x \sim_R y\} \arrow[dl, "\pi_X"'] \arrow[dr, "\pi_2"] & \\ X & & Y \end{tikzcd}$$

On the other hand given a span of relations:

$$\begin{tikzcd} & R \arrow[dl, "\pi_X"'] \arrow[dr, "\pi_2"] & \\ X & & Y \end{tikzcd}$$

define a relation between 2^X and 2^Y as in the general case by

$$x \subseteq X \ \sim_R \ y \subseteq Y \quad \equiv \quad \exists r \in R.\, \pi_X(r) = x \wedge \pi_Y(r) = y$$

In summary, whereas the category Rel is $\mathrm{Span}(\mathrm{Set})$, powerset relations are $\mathrm{Span}(\mathrm{Rel}) = \mathrm{Span}(\mathrm{Span}(\mathrm{Set}))$. We denote this bicategory $\wp\mathrm{Rel}$.

[4] Not true! $R + R$ is not $\cong R$

[5] where between each pair $f, g : X \longrightarrow Y$ of 1-cells is at most one 2-cell.

3.2 Compact-closed structure of $\wp\mathrm{Rel}$

4 Quantifiers as Relations

It follows from what was said thus far that equation (2) exhibits a quantifier as a relation $2^Y \longrightarrow 2^Y$.

So in order to interpret generailed quantifiers in Rel, we look into relations over powersets.

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