

Funcons-beta: Binding *

The P_LanCompS Project

Binding.cbs | PLAIN | PRETTY

OUTLINE

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Binding

- [*Type* environments
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- Datatype* identifiers
- Alias* ids
- Funcon* identifier-tagged
- Alias* id-tagged
- Funcon* fresh-identifier
- Entity* environment
- Alias* env
- Funcon* initialise-binding
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- Alias* bind
- Funcon* unbind
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- Funcon* bound-value
- Alias* bound
- Funcon* closed
- Funcon* scope
- Funcon* accumulate
- Funcon* collateral
- Funcon* bind-recursively
- Funcon* recursive]

Meta-variables $T <:$ values

*Suggestions for improvement: plancomps@gmail.com.
Reports of issues: <https://github.com/plancomps/CBS-beta/issues>.

Environments

Type `environments` \rightsquigarrow `maps(identifiers, values?)`

Alias `envs` = `environments`

An environment represents bindings of identifiers to values. Mapping an identifier to `()` represents that its binding is hidden.

Circularity in environments (due to recursive bindings) is represented using bindings to cut-points called `links`. Funcons are provided for making declarations recursive and for referring to bound values without explicit mention of links, so their existence can generally be ignored.

Datatype `identifiers` ::= `{_ : strings}` | `identifier-tagged(_ : identifiers, _ : values)`

Alias `ids` = `identifiers`

Alias `id-tagged` = `identifier-tagged`

An identifier is either a string of characters, or an identifier tagged with some value (e.g., with the identifier of a namespace).

Funcon `fresh-identifier` : \Rightarrow `identifiers`

`fresh-identifier` computes an identifier distinct from all previously computed identifiers.

Rule `fresh-identifier` \rightsquigarrow `identifier-tagged("generated", fresh-atom)`

Current bindings

Entity `environment(_ : environments)` \vdash `_` \longrightarrow `_`

Alias `env` = `environment`

The environment entity allows a computation to refer to the current bindings of identifiers to values.

Funcon `initialise-binding`(`X` : \Rightarrow `T`) : \Rightarrow `T`
 \rightsquigarrow `initialise-linking(initialise-generating(closed(X)))`

`initialise-binding`(`X`) ensures that `X` does not depend on non-local bindings. It also ensures that the linking entity (used to represent potentially cyclic bindings) and the generating entity (for creating fresh identifiers) are initialised.

Funcon `bind-value`(`I` : `identifiers`, `V` : `values`) : \Rightarrow `environments`
 \rightsquigarrow `{I \mapsto V}`

Alias `bind` = `bind-value`

`bind-value`(`I`, `X`) computes the environment that binds only `I` to the value computed by `X`.

Funcon `unbind`(`I` : `identifiers`) : \Rightarrow `environments`
 \rightsquigarrow `{I \mapsto ()}`

`unbind`(`I`) computes the environment that hides the binding of `I`.

Funcon `bound-directly`($_ : \text{identifiers}$) : $\Rightarrow \text{values}$

`bound-directly`(I) returns the value to which I is currently bound, if any, and otherwise fails.

`bound-directly`(I) does *not* follow links. It is used only in connection with recursively-bound values when references are not encapsulated in abstractions.

Rule
$$\frac{\text{lookup}(\rho, I) \rightsquigarrow (V : \text{values})}{\text{environment}(\rho) \vdash \text{bound-directly}(I : \text{identifiers}) \longrightarrow V}$$

Rule
$$\frac{\text{lookup}(\rho, I) \rightsquigarrow ()}{\text{environment}(\rho) \vdash \text{bound-directly}(I : \text{identifiers}) \longrightarrow \text{fail}}$$

Funcon `bound-value`($I : \text{identifiers}$) : $\Rightarrow \text{values}$
 $\rightsquigarrow \text{follow-if-link}(\text{bound-directly}(I))$

Alias `bound` = `bound-value`

`bound-value`(I) inspects the value to which I is currently bound, if any, and otherwise fails. If the value is a link, `bound-value`(I) returns the value obtained by following the link, if any, and otherwise fails. If the inspected value is not a link, `bound-value`(I) returns it.

`bound-value`(I) is used for references to non-recursive bindings and to recursively-bound values when references are encapsulated in abstractions.

Scope

Funcon `closed`($X : \Rightarrow T$) : $\Rightarrow T$

`closed`(X) ensures that X does not depend on non-local bindings.

Rule
$$\frac{\text{environment}(\text{map}(\)) \vdash X \longrightarrow X'}{\text{environment}(_) \vdash \text{closed}(X) \longrightarrow \text{closed}(X')}$$

Rule
$$\text{closed}(V : T) \rightsquigarrow V$$

Funcon `scope`($_ : \text{environments}, _ : \Rightarrow T$) : $\Rightarrow T$

`scope`(D, X) executes D with the current bindings, to compute an environment ρ representing local bindings. It then executes X to compute the result, with the current bindings extended by ρ , which may shadow or hide previous bindings.

`closed`(`scope`(ρ, X)) ensures that X can reference only the bindings provided by ρ .

Rule
$$\frac{\text{environment}(\text{map-override}(\rho_1, \rho_0)) \vdash X \longrightarrow X'}{\text{environment}(\rho_0) \vdash \text{scope}(\rho_1 : \text{environments}, X) \longrightarrow \text{scope}(\rho_1, X')}$$

Rule
$$\text{scope}(_ : \text{environments}, V : T) \rightsquigarrow V$$

Funcon `accumulate`($_ : (\Rightarrow \text{environments})^*$) : $\Rightarrow \text{environments}$

`accumulate`(D_1, D_2) executes D_1 with the current bindings, to compute an environment ρ_1 representing some local bindings. It then executes D_2 to compute an environment ρ_2 representing further local bindings, with the current bindings extended by ρ_1 , which may shadow or hide previous current bindings. The result is ρ_1 extended by ρ_2 , which may shadow or hide the bindings of ρ_1 .

`accumulate`($_, _$) is associative, with `map`($_$) as unit, and extends to any number of arguments.

Rule $\frac{D_1 \rightarrow D'_1}{\text{accumulate}(D_1, D_2) \rightarrow \text{accumulate}(D'_1, D_2)}$
Rule $\text{accumulate}(\rho_1 : \text{environments}, D_2) \rightsquigarrow \text{scope}(\rho_1, \text{map-override}(D_2, \rho_1))$
Rule $\text{accumulate}() \rightsquigarrow \text{map}()$
Rule $\text{accumulate}(D_1) \rightsquigarrow D_1$
Rule $\text{accumulate}(D_1, D_2, D^+) \rightsquigarrow \text{accumulate}(D_1, \text{accumulate}(D_2, D^+))$

Funcon $\text{collateral}(\rho^* : \text{environments}^*) : \Rightarrow \text{environments}$
 $\rightsquigarrow \text{checked map-unite}(\rho^*)$

$\text{collateral}(D_1, \dots)$ pre-evaluates its arguments with the current bindings, and unites the resulting maps, which fails if the domains are not pairwise disjoint.

$\text{collateral}(D_1, D_2)$ is associative and commutative with $\text{map}()$ as unit, and extends to any number of arguments.

Recurse

Funcon $\text{bind-recursively}(I : \text{identifiers}, E : \Rightarrow \text{values}) : \Rightarrow \text{environments}$
 $\rightsquigarrow \text{recursive}(\{I\}, \text{bind-value}(I, E))$

$\text{bind-recursively}(I, E)$ binds I to a link that refers to the value of E , representing a recursive binding of I to the value of E . Since $\text{bound-value}(I)$ follows links, it should not be executed during the evaluation of E .

Funcon $\text{recursive}(SI : \text{sets}(\text{identifiers}), D : \Rightarrow \text{environments}) : \Rightarrow \text{environments}$
 $\rightsquigarrow \text{re-close}(\text{bind-to-forward-links}(SI), D)$

$\text{recursive}(SI, D)$ executes D with potential recursion on the bindings of the identifiers in the set SI (which need not be the same as the set of identifiers bound by D).

Auxiliary Funcon $\text{re-close}(M : \text{maps}(\text{identifiers}, \text{links}), D : \Rightarrow \text{environments}) : \Rightarrow \text{environments}$
 $\rightsquigarrow \text{accumulate}(\text{scope}(M, D), \text{sequential}(\text{set-forward-links}(M), \text{map}()))$

$\text{re-close}(M, D)$ first executes D in the scope M , which maps identifiers to freshly allocated links. This computes an environment ρ where the bound values may contain links, or implicit references to links in abstraction values. It then sets the link for each identifier in the domain of M to refer to its bound value in ρ , and returns ρ as the result.

Auxiliary Funcon $\text{bind-to-forward-links}(SI : \text{sets}(\text{identifiers})) : \Rightarrow \text{maps}(\text{identifiers}, \text{links})$
 $\rightsquigarrow \text{map-unite}(\text{interleave-map}(\text{bind-value}(\text{given}, \text{fresh-link}(\text{values})), \text{set-elements}(SI)))$

$\text{bind-to-forward-links}(SI)$ binds each identifier in the set SI to a freshly allocated link.

Auxiliary Funcon $\text{set-forward-links}(M : \text{maps}(\text{identifiers}, \text{links})) : \Rightarrow \text{null-type}$
 $\rightsquigarrow \text{effect}(\text{interleave-map}(\text{set-link}(\text{map-lookup}(M, \text{given}), \text{bound-value}(\text{given})), \text{set-elements}(\text{map-domain}(M))))$

For each identifier I in the domain of M , $\text{set-forward-links}(M)$ sets the link to which I is mapped by M to the current bound value of I .