Graph of Distributed Data Structures

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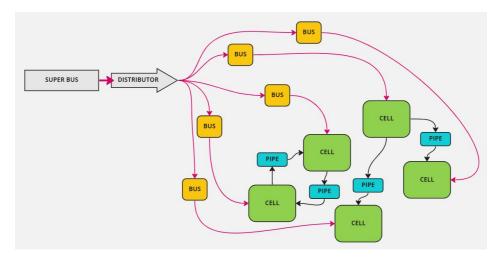
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1 How to get started

```
git clone git@github.com:proman3419/Graph-of-Distributed-Data-Structures.git cd Graph-of-Distributed-Data-Structures/goddslang mvn package alias godds="java -jar target/goddslang-1.0-SNAPSHOT-jar-with-dependencies.jar" godds "src\main\resources\code snippets\example01_hello_world.godds"
```

2 Model

A standard program written in GoDDS can be featured as:



2.1 Super bus

It's used to input data to the program, implemented as a queue; in code #INPUT.

2.2 Distributor

It's used to distirbute data from the super bus to separate buses for each cell; in code #INPUT_CELLS. How does it work? Consider the following header:

```
#CELLS_COUNT 4

#CELLS_GRAPH
0011
1011
0000
1010

#END

#INPUT 10 16 23 110 21 45

#INPUT_CELLS 0 2 3 1 0
```

The key notes:

- there are 4 cells (with indexes 0, 1, 2, 3)
- super bus' initial state is [10, 16, 23, 110, 21]
- the distributor's content is [0 2 3 1]

Steps of the algorithm:

- 1. 10 lands in cell 0. bus \rightarrow [10]
- 2. 16 lands in cell 2. bus \rightarrow [16]
- 3. 23 lands in cell 3. bus \rightarrow [23]
- 4. 110 lands in cell 1. bus \rightarrow [110]
- 5. 21 lands in cell 0. bus \rightarrow [10, 21]
- 6. 45 lands in cell 0. bus \rightarrow [10, 21, 45]

As we can see the value from the super bus at the ith index is added to the bus with id INPUT_CELLS[i%len(INPUT_CELLS)].

2.3 Bus

Input source for a cell, implemented as a queue.

After initialization no data can be written to it.

It can be accessed only by the cell which it's dedicated for.

When a cell reads from its bus it retrieves the first value that had been inserted.

The retrieved value is removed from the bus.

If there is no value in bus, reading is skipped.

2.4 Pipe

It's a kind of bus between two cells, it's directed, implemented as a queue.

The key difference is that it can be written to and read from.

2.5 Cell

Entity that has two registers (R0, R1) that can store integer values.

The R0 register is the active one - most of the functions work with its value and modify its value.

The R1 register is meant for storing data.

Each cell may have code which it executes.

2.5.1 States

- Running the cell is executing code
- Finished the cell has finished executing code

2.5.2 Options

- Daemon cell terminates if there are no more non-Daemon cells
- Inactive the cell doesn't run code (it's used only for storing it)

Options are prefixed with %.

2.6 Code execution

Program runs in steps, for each step a cell calls one function.

Id of a cell determines its order, the smaller it is the higher the priority.

In other words: cell with id X will perform a step before a cell with id (X + 1).

If all cells have changed their states to Finished the program ends.

2.7 Call stack

The cells obey the following rules regarding jumps:

- If the label of a jump refers to the same cell which owns this call then the point from the cell had jumped isn't kept on the stack
- If the label of a jump refers to a neighboring cell then the point from the cell had jumped is kept on the stack

Check this example to better understand the difference.

3 Sample program

3.1 Code

```
#CELLS_COUNT 3
#CELLS_GRAPH
    010
    101
    010
#END
#INPUT 5 72 101 108 108 111
                             // The first value n = messages' length
                              // followed by n values - the message
#INPUT_CELLS 0
                              // Write to cell 0
    #CELL 0 Client1
                              // Read n
        READ_BUS
        SWAP
                              // Store it
        LABEL send_loop
            READ_BUS
                              // Read a character of the message
            WRITE_CELL Server // Send it to the server
            SWAP
            SUB 1
                              // Mark the character as sent
            IFEZ $
                              // Check if message has been sent
            SWAP
            JUMP send_loop
        TERMINATE
    #END
    #CELL 1 Server %Daemon
        READ_CELL Client1
                              // Get a character
        WRITE_CELL Client2
                              // Forward the character
        JUMP ^
    #END
    #CELL 2 Client2 %Daemon
        READ_CELL Server
                              // Get a character
        PRINT "Received " 1
        PRINTNL_CHAR RO O
                              // Print the character
        JUMP ^
    #END
#END_ALL
```

3.2 Preamble

```
#CELLS_COUNT VAL

#CELLS_GRAPH

G[0,0] G[0,1] ... G[0,CELLS_COUNT-1]

...

G[CELLS_COUNT,0] G[CELLS_COUNT,1] ... G[CELLS_COUNT,CELLS_COUNT-1]

#END
```

Defines cell count and edges for the graph of cells that will be created for usage in code. Each row has to have as many values as cells count VAL and there have to be as many rows as cells count VAL.

```
#INPUT VAL(s)
#INPUT_CELLS VAL(s)
```

Optional element, defines bus of values and order in which they will be read, VAL in input cells defines id of cell which will receive VAL from on corresponding index in input. Bus loops around if there are more cell ids given in input cells than values in input.

3.3 Cell code

```
#CELL ID LABEL
     (Any combination of keywords creating cell code)
#END
```

Cell starts with #CELL keyword requiring cell id and cell label, and ends with #END. Cell code can be written between those two keywords. Cell ids have to be a permutation of numbers from the set: $\{0,1,..., < CELLS_COUNT>-1\}$. Cell label has to be unique when compared to other already existing cells.

3.4 Code end

#END_ALL

After inputting all cell codes, #END_ALL has to be used, signifying the end of executable code.

4 Used nomenclature

Refers to a register's name.

```
<_REG_VALUE_>
```

Refers to a value stored in a register.

Refers to a math function argument where VAL can be any integer.

Refers to an argument for jump function, either local label for intracell jumps or label of different cell and label local to it's content for intercell jumps.

```
<_PRINT_ARG_> can be one of:
    <_REG_NAME_>
    "Any string"
```

Refers to an argument for print functions.

5 Functions

5.1 Supers

```
#CELLS_COUNT VAL
```

Defines cell count for the program.

Defines adjecency graph representing neighbourhood of all cells.

```
#INPUT VAL(s)
#INPUT_CELLS VAL(s)
```

Optional element, defines bus of values and order in which they will be read, VAL in input cells defines id of cell which will receive VAL from on corresponding index in input. Bus loops around if there are more cell ids given in input cells than values in input.

```
#CELL ID LABEL
```

Defines start of executable code for cell with ID and LABEL.

#END

Defines end of executable code for cell.

#END_ALL

Defines end for executable code for entire program.

5.2 Modifying registry values

```
ADD <_ARG_>
where <_ARG_> can be one of:
     <_REG_NAME_>
     <VAL>
```

Adds the value <VAL> or <_REG_VAL_> to the value stored in R0 then saves the result in R0.

In short: R0 = R0 + < ARG >

SUB <_ARG_>

Subtracts the value <VAL> or <_REG_VAL $_>$ from the value stored in R0 then saves the result in R0

In short: $R0 = R0 - \langle ARG \rangle$

MUL <_ARG_>

Multiplies the value stored in R0 by <VAL> or <REG_VAL> then saves the result in R0.

In short: R0 = R0 * < ARG >

DIV <_ARG_>

Divides the value stored in R0 by <VAL> or <REG_VAL> then saves the result (as integer) in R0. In short: R0 = |R0 / <ARG $_>|$

MOD <_ARG_>

Executes modulo operation using the value stored in R0 and <VAL> or <_REG_VAL_> then saves the result in R0.

In short: R0 = R0 % < ARG >

ABS

Overwrites R0's value with its absolute value. In short: R0 = |R0|

SET <_REG_NAME_> <VAL>

```
Sets the value of register <_REG_NAME_> to <VAL>.
```

```
In short: <_REG_NAME_> = <VAL>
```

```
COMP <VAL>
```

Compares the value stored in R0 to <VAL>, then saves the result (as integer) in R0. The value saved is either -1, 0 or 1 depending on the values of R0 and <VAL>.

In short

```
<VAL> < R0 \rightarrow R0 = -1
<VAL> \equiv R0 \rightarrow R0 = 0
<VAL> > R0 \rightarrow R0 = 1
```

SWAP

Swaps the values stored in R0 and R1.

In short: $R0 \leftrightarrow R1$

COPY

Copies the value stored in R0 to R1.

In short: $R0 \rightarrow R1$

5.3 Modifying order of calls

JUMP <_EXTENDED_LABEL_>

```
LABEL <NAME>
```

Defines a label that can be referenced by <NAME> when using the JUMP function.

<CELL_LABEL>@^ // Start of a specific cell <CELL_LABEL>@\$ // End of a specific cell <CELL_LABEL>@<LABEL> // Label in a specific cell

Jump to the label <_EXTENDED_LABEL_> then continue execution from the point after its name. Jump can be either an intracell one, if no <CELL_LABEL> is given, or an intercell one if full <_EXTENDED_LABEL_> is used.

```
IFEZ <_EXTENDED_LABEL_>
```

If the value stored in R0 is equal to 0 then JUMP <_EXTENDED_LABEL_> will be performed. In short: if R0 \equiv 0 then JUMP <_EXTENDED_LABEL_>

```
IFNZ <_EXTENDED_LABEL_>
```

If the value stored in R0 is not equal to 0 then JUMP <_EXTENDED_LABEL_> will be performed. In short: if R0 \neq 0 then JUMP <_EXTENDED_LABEL_>

```
IFLZ <_EXTENDED_LABEL_>
```

If the value stored in R0 is less than 0 then JUMP <_EXTENDED_LABEL_> will be performed. In short: if R0 < 0 then JUMP <_EXTENDED_LABEL_>

```
IFGZ <_EXTENDED_LABEL_>
```

If the value stored in R0 is greater than 0 then JUMP <_EXTENDED_LABEL_> will be performed. In short: if R0 > 0 then JUMP <_EXTENDED_LABEL_>

EXIT

Stops execution of code from this cell. If the cell executed code from a different cell through an <_EXTENDED_LABEL_> this execution will halt and resume in earlier cell after the JUMP call.

TERMINATE

Stops execution of all code by origin cell and changes its state to Finished.

```
PASS <VAL>
```

After encountering this keyword cell will stop it's execution for $\langle VAL \rangle$ steps. It can be used without an argument in which case it defaults to 0.

5.4 Output

```
PRINT <_PRINT_ARG_> <FLAG>
       where <_PRINT_ARG_> can be one of:
             <_REG_NAME_>
             "Any string"
In short:
if \langle FLAG \rangle \equiv 0 then print(\langle PRINT\_ARG\_ \rangle)
if \langle FLAG \rangle \equiv 1 then print(\langle CELL\_LABEL \rangle: \langle PRINT\_ARG\_ \rangle)
if \langle FLAG \rangle \equiv 2 then print(\langle CELL\_LABEL \rangle \rightarrow \langle \_REG\_NAME\_ \rangle: \langle \_REG\_VALUE\_ \rangle)
      PRINTNL <_PRINT_ARG_> <FLAG>
In short:
if \langle FLAG \rangle \equiv 0 then println(\langle PRINT\_ARG\_ \rangle)
if \langle FLAG \rangle \equiv 1 then println(\langle CELL\_LABEL \rangle: \langle PRINT\_ARG\_\rangle)
if \langle FLAG \rangle \equiv 2 then println(\langle CELL\_LABEL \rangle \rightarrow \langle \_REG\_NAME\_ \rangle: \langle \_REG\_VALUE\_ \rangle)
      PRINT_CHAR <_REG_NAME_> <FLAG>
In short:
if \langle FLAG \rangle \equiv 0 then print(chr(\langle REG\_VALUE\_ \rangle))
if \langle FLAG \rangle \equiv 1 then print(\langle CELL\_LABEL \rangle: chr(\langle \_REG\_VALUE\_ \rangle))
if \langle FLAG \rangle \equiv 2 then print(\langle CELL\_LABEL \rangle \rightarrow \langle \_REG\_NAME\_ \rangle: chr(\langle \_REG\_VALUE\_ \rangle))
```

```
PRINTNL_CHAR <_REG_NAME_> <FLAG>
```

In short:

```
 \begin{split} & \text{if } <\text{FLAG}>\equiv 0 \text{ then } \\ & \text{println}(\text{chr}(<&\text{REG\_VALUE\_>})) \\ & \text{if } <\text{FLAG}>\equiv 1 \text{ then } \\ & \text{println}(<&\text{CELL\_LABEL}>: \\ & \text{chr}(<&\text{REG\_VALUE\_>})) \\ & \text{if } <\text{FLAG}>\equiv 2 \text{ then } \\ & \text{println}(<&\text{CELL\_LABEL}>\rightarrow <&\text{REG\_NAME\_>}: \\ & \text{chr}(<&\text{REG\_VALUE\_>})) \\ \end{aligned}
```

5.5 Intercell operations

```
WRITE_CELL <CELL_LABEL>
```

Writes value stored in R0 to the pipe between this cell and the one specified with <CELL_LABEL>.

```
READ_BUS
```

Reads a value from this cell's bus.

```
READ_CELL <CELL_LABEL>
```

Reads a value from the pipe between this cell and the one specified with <CELL_LABEL>. saving it to R0, if no value is found in the pipe then the cell awaits on this call.

```
COPY_CELL <CELL_LABEL>
```

Copy the value stored in R0 to R0 in cell specified with <CELL_LABEL>.

```
PRINT_LABEL_NAME <CELL_LABEL>
```

Prints label of cell specified with <CELL_LABEL>.

Pretty much useless, left over from changing the argument from <CELL_ID>.

6 Grammar

```
grammar Grammar;
start
   : preamble cells SUPER_END_ALL
preamble
   : cellsCount cellsGraph input
cellsCount
   : SUPER_CELLS_COUNT numberArgument
cellsGraph
   : SUPER_CELLS_GRAPH numberArgument+ SUPER_END
input
   : (inputVals inputCells)?
inputVals
   : SUPER_INPUT numberArgument+
inputCells
   : SUPER_INPUT_CELLS numberArgument+
cells
  : cell+
// Cell -----
cell
   : cellHeader cellCode? SUPER_END
cellHeader
   : SUPER_CELL numberArgument idArgument cellOptions
cellOptions
   : cellOption*
cellOption
   : CELL_OPTION_DAEMON
```

```
| CELL_OPTION_INACTIVE
cellCode
   : cellCodePart+
cellCodePart
   : functionCall cellCodePart*
CELL_OPTION_DAEMON
   : '%Daemon'
   ;
CELL_OPTION_INACTIVE
   : '%Inactive'
functionCall
   : functionAdd
   | functionSub
   | functionMul
   | functionDiv
   | functionMod
   | functionAbs
   | functionSet
   | functionComp
   | functionSwap
   | functionCopy
   | functionDefineLabel
   | functionJump
   | functionCheckIFEZ
   | functionCheckIFNZ
   | functionCheckIFLZ
   | functionCheckIFGZ
   | functionExit
   | functionPrint
   | functionPrintNL
   | functionPrintChar
   | functionPrintNLChar
   | functionWriteCell
   | functionReadCell
   | functionReadBus
   | functionCopyCell
   | functionPrintLabelName
   | functionPass
   | functionTerminate
```

functionAdd

```
: ADD (numberArgument | idArgument)
functionSub
    : SUB (numberArgument | idArgument)
functionMul
    : MUL (numberArgument | idArgument)
{\tt functionDiv}
    : DIV (numberArgument | idArgument)
functionMod
    : MOD (numberArgument | idArgument)
functionAbs
    : ABS
functionSet
    : SET idArgument numberArgument
    ;
functionComp
    : COMP numberArgument
functionSwap
    : SWAP
functionCopy
    : COPY
{\tt functionDefineLabel}
    : DEFINE_LABEL idArgument
functionJump
    : JUMP extendedDefinedLabel
functionCheckIFEZ
    : CHECK_IFEZ extendedDefinedLabel
functionCheckIFNZ
```

```
: CHECK_IFNZ extendedDefinedLabel
functionCheckIFLZ
    : CHECK_IFLZ extendedDefinedLabel
{\tt functionCheckIFGZ}
    : CHECK_IFGZ extendedDefinedLabel
functionExit
   : EXIT
functionPrint
    : PRINT printArguments
{\tt functionPrintNL}
    : PRINTNL printArguments
{\tt functionPrintChar}
    : PRINT_CHAR printArguments
{\tt functionPrintNLChar}
    : PRINTNL_CHAR printArguments
functionWriteCell
    : WRITE_CELL idArgument
{\tt functionReadCell}
    : READ_CELL idArgument
{\tt functionReadBus}
    : READ_BUS
{\tt functionCopyCell}
    : COPY_CELL idArgument
functionPrintLabelName
    : PRINT_LABEL_NAME numberArgument
```

functionPass

```
: PASS numberArgument
functionTerminate
   : TERMINATE
SUPER_CELLS_COUNT
   : '#CELLS_COUNT'
SUPER_CELLS_GRAPH
  : '#CELLS_GRAPH'
SUPER_INPUT
  : '#INPUT'
SUPER_INPUT_CELLS
   : '#INPUT_CELLS'
SUPER_CELL
  : '#CELL'
  ;
SUPER_END
  : '#END'
SUPER_END_ALL
  : '#END_ALL'
ADD
  : 'ADD'
SUB
 : 'SUB'
MUL
 : 'MUL'
DIV
: 'DIV'
```

MOD

: 'MOD'

ABS

: 'ABS' ;

SET

: 'SET' ;

COMP

: 'COMP' ;

SWAP

: 'SWAP' ;

COPY

: 'COPY'

DEFINE_LABEL

: 'LABEL'

JUMP

: 'JUMP'

CHECK_IFEZ

: 'IFEZ'

CHECK_IFNZ

: 'IFNZ'

CHECK_IFLZ

: 'IFLZ'

CHECK_IFGZ

: 'IFGZ'

EXIT

: 'EXIT'

PRINT

```
: 'PRINT'
PRINTNL
   : 'PRINTNL'
PRINT_CHAR
   : 'PRINT_CHAR'
PRINTNL_CHAR
   : 'PRINTNL_CHAR'
WRITE_CELL
  : 'WRITE_CELL'
READ_CELL
  : 'READ_CELL'
READ_BUS
   : 'READ_BUS'
COPY_CELL
  : 'COPY_CELL'
PRINT_LABEL_NAME
   : 'PRINT_LABEL_NAME'
PASS
  : 'PASS'
TERMINATE
  : 'TERMINATE'
// Arguments ===========
arguments
   : (idArgument+ numberArgument*)*
   | (idArgument* numberArgument+)+
idArgument
  : ID
```

```
;
numberArgument
    : NUMBER
{\tt extendedDefinedLabel}
    : definedLabel
    | ID'@'definedLabel
{\tt definedLabel}
   : '^'
    | '$'
    | ID
printArguments
    : STRING* numberArgument
    | idArgument numberArgument
    | numberArgument
NUMBER
    : '-'?('0'..'9')+
ID
    : ('_'|'-'|[a-zA-Z])+('0'..'9')*
STRING
    : '"' (~["\r\n] | '""')* '"'
// Ignore white space characters
WS
    : [ \t \n] + -> skip
// Comment starts with '//'
  : '//' ~( '\r' | '\n' )* -> skip
```

7 Technology stack

- ANTLR4
- Java 17
- Maven