# On Calculations Involving Perturbed Orthogonal Matrices

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#### Abstract

Many common algorithms employed in Data Science require the calculation of an Orthonormal Basis. Often, this is done using the Gram-Schmidt Algorithm, which has time complexity  $O(n^3)$  in the standard case. For low-dimensional spaces, this does not pose a significant computational burden. However, for high-dimensional vector spaces, calculating such a basis can prove unfeasible, so it could be desirable to find faster algorithms that approximate Orthonormal Bases up to some error tolerance  $\varepsilon$ . In this paper, we examine bounds on error propagation if such an " $\varepsilon$ -almost orthonormal basis" (defined below) is employed.

# 1. Introduction

#### 1.1. MOTIVATION

Orthogonal matrices (and their complex counterparts Unitary matrices) are of vital importance in linear algebra, both pure and applied. In a sense, they are the least "disruptive" class of deformations we can apply to a vector space without fixing all of its elements — every linear isometry operator on a finite-dimensional inner product space can be represented by an orthogonal matrix, and conversely, every orthogonal matrix represents an isometry encoded with respect to some particular basis. Thus, every orthogonal matrix Q is necessarily invertible, with  $Q^{-1} = Q^{\mathsf{T}}$ . This is very nice property, for both theoretical and computational purposes.

From the theoretical perspective, this trivially guarantees the existence of an inverse, and from  $1 = \det(I) = \det(QQ^{-1}) = \det(Q)\det(Q^{\mathsf{T}}) = \det(Q)^2$ , we get that  $\det(Q) = \pm 1.^1$  Additionally, because  $Q^{-1}Q = I = QQ^{-1}$ , orthogonal matrices encode normal operators, and the complex spectral theorem they are thus unitarily diagonalizable [Axl15]. In fact, we can guarantee an even stronger property — if Q is orthogonal, then the spectrum  $\sigma(Q) \subseteq \{z \in \mathbb{C} \mid |z| = 1\}$ . That is, every eigenvalue of Q lies on the unit circle in  $\mathbb{C}$ . As such, the group of  $n \times n$  orthogonal matrices (denoted O(n)) and its subgroups are of great interest from a theoretical perspective.

Orthogonal matrices also have many nice computational properties. Again, let Q be an orthogonal matrix. Then because all eigenvalues have unit modulus, orthogonal matrices cannot magnify errors in computa-

tions [TB97, p. 95]. Furthermore, because  $Q^{-1} = Q^{\mathsf{T}}$ , inverting Q is essentially trivial, as we can simply swap indexing schemes to yield a semantically faithful representation of  $Q^{-1}$ , without having to perform any major calculations. Finally, Orthogonal matrices are of great importance because of their occurrence in important matrix factorizations, such as QR decomposition.

At the time that this paper is being written, there is a large body of ongoing research in numerical linear algebra surrounding "sketching" methods — that is, fast, randomized algorithms for querying large matrices in a way that preserves relevant structure. One classic result in this vein is the *Johnson-Lindenstrauss Lemma* (hence referred to as the JL Lemma), which proves a bounds about the extent to which a family of randomized dimension-reduction algorithms can preserve metric structure. Such tools are necessitated by the ever-growing importance of Big Data and Big Data Analytics techniques in the tech industry.

For our project, we will examine how error resulting from a JL-like dimension reduction scheme can propagate through further computations. To do so, we will re-frame the problem in terms of examining how properties of orthogonal matrices are affected when we apply a small random perturbation matrix. We will examine both additive and multiplicative perturbations, and describe some possible connections to eigenvalue estimation problems.

But first, we will give an overview of any nonstandard definitions and notation that we will make use of in our paper. As our randomized algorithms will make some use of concepts from Random Matrix Theory, we will also include an overview of some of the

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<sup>&</sup>lt;sup>1</sup>note that the converse is not true — i.e., for some  $n \times n$  matrix M with  $\det(M) = \pm 1$ , it is not necessarily true that M is orthogonal

relevant concepts in the space below.

#### 1.2. Definitions and Notations

First, some general things: to make things easier for the reader to scan through our paper quickly, we will use  $\Delta$  to denote the end of a theorem, definition, remark, or similar, and  $\blacksquare$  to denote end of a proof. In general, bases for spaces will be given by  $\mathcal{B}$ , while the basis vectors themselves will be denoted  $\mathbf{e}_i$ .  $\delta_{ij}$  will, as usual, refer to the Kronecker delta. In general, we will use the variables i, j, k and  $\ell$  as indices variables, resorting to Greek symbols only if necessary. Unless otherwise stated,  $\mathcal{U}, \mathcal{V}$  and  $\mathcal{W}$  will refer to finite-dimensional inner product spaces. All other notation (except that described below) should be standard.

DEFINITION 1.1 (Random Matrix Family). We'll use some nonstandard notation for random matrices in our paper. Let  $\mathcal{A}$  be a probability distribution with parameters  $x_1, \ldots, x_n$ , and let  $M \in M_{m \times n}(\mathbb{R})$ . Then we say  $M \sim \mathcal{A}(X_1, \ldots, X_n)$  iff

$$m_{ij} \sim \mathcal{A}((\boldsymbol{X}_1)_{ij}, \dots, (\boldsymbol{X}_n)_{ij}).$$

for all i, j.  $\triangle$ 

EXAMPLE 1.1.1 (Gaussian random matrix). Let  $\mu \in M_{m \times n}(\mathbb{R})$ , and let  $\sigma \in M_{m \times n}(\mathbb{R}^{\geq 0})$ . Then we say  $\boldsymbol{\xi} \sim \mathcal{N}(\boldsymbol{\mu}, \boldsymbol{\sigma})$  iff

$$\xi_{ij} \sim \mathcal{N}(\mu_{ij}, \sigma_{ij})$$

for all i, j. We will sometimes denote this by

$$\begin{bmatrix} \xi_{1,1} & \cdots & \xi_{1,n} \\ \vdots & \ddots & \vdots \\ \xi_{m,1} & \cdots & \xi_{m,n} \end{bmatrix} \sim \begin{bmatrix} \mathcal{N}(\mu_{1,1}, \sigma_{1,1}) & \cdots & \mathcal{N}(\mu_{1,n}, \sigma_{1,n}) \\ \vdots & \ddots & \vdots \\ \mathcal{N}(\mu_{m,1}, \sigma_{m,1}) & \cdots & \mathcal{N}(\mu_{m,n}, \sigma_{m,n}) \end{bmatrix}$$

if it is more convenient to do so.

Oftentimes, we'll be interested in matrices where all of the entries are i.i.d. Hence, we'll introduce the following notation:

Definition 1.2. Let  $c \in \mathbb{R}$ . Then we define

$$c_{m \times n} = c \begin{bmatrix} 1 & \cdots & 1 \\ \vdots & \ddots & \vdots \\ 1 & \cdots & 1 \end{bmatrix} = \begin{bmatrix} c & \cdots & c \\ \vdots & \ddots & \vdots \\ c & \cdots & c \end{bmatrix}$$

Thus, we can define a  $m \times n$  matrix of i.i.d. Gaussian random variables by  $M \sim \mathcal{N}(\mu_{m \times n}, \sigma_{m \times n})$  for some  $\mu, \sigma \in \mathbb{R}$  (where  $\sigma \geq 0$ ).

DEFINITION 1.3. Let  $M \in M_{m \times n}(\mathbb{R})$ . If we have  $M_{ij} < \varepsilon$  for all i, j, then we write

$$M \prec \varepsilon$$

and say M is an  $\varepsilon$ -bounded matrix. We denote the set of all matrices M such that  $-\varepsilon \prec M \prec \varepsilon$  by  $B_{\varepsilon}(0_{m \times n})$ , because under the metric induced by the max norm on

 $M_{m \times n}(\mathbb{R})$ , the set of  $\varepsilon$ -bounded matrices is simply the open ball about the matrix  $0_{m \times n}$ .

In the case that we want  $M_{ij} \leq \varepsilon$ , we of course replace  $\prec$  with a  $\leq$  symbol, and denote the set of all such matrices by  $\overline{B_{\varepsilon}(0_{m \times n})}$ .

DEFINITION 1.4 (Orthogonal Group). The standard notation for the group of orthogonal matrices on  $\mathbb{R}^n$  is O(n). However, this looks confusingly similar to the big-O notation O(n). Hence, we will elect to denote the orthogonal group on  $\mathbb{R}^n$  by Orth(n).

DEFINITION 1.5. Let  $n \in \mathbb{N}$ , and let  $M \in M_{n \times n}(\mathbb{R})$ . Then if M can be expressed as the sum of an an orthogonal matrix Q and a matrix  $\xi$  with  $-\varepsilon \prec \xi \prec \varepsilon$ , then we call M a  $\varepsilon$ -almost orthogonal matrix.

DEFINITION 1.6. Let V be a finite-dimensional inner product space. Let  $\mathcal{B} = \{\mathbf{e}_1, \dots, \mathbf{e}_n\}$  be a basis for V, and suppose that  $\forall \mathbf{e}_i, \mathbf{e}_j \in \mathcal{B}$  with  $i \neq j$ , we have

$$|\langle \mathbf{e}_i, \mathbf{e}_j \rangle| < \varepsilon.$$

and  $\langle \mathbf{e}_i, \mathbf{e}_i \rangle < 1 + \varepsilon$  for every  $\mathbf{e}_i$ . Then we call call  $\mathcal{B}$  an  $\varepsilon$ -almost orthonormal basis of V. Note that it will often be useful to restate the conditions above as  $\forall \mathbf{e}_i, \mathbf{e}_i \in \mathcal{B}$ ,

$$|\langle \mathbf{e}_i, \mathbf{e}_j \rangle - \delta_{ij}| < \varepsilon$$

which will be particularly relevant when dealing with double summations.  $\triangle$ 

To contrast orthonormal bases with  $\varepsilon$ -almost orthonormal bases, we will often use  $\mathcal{B}_{\perp}$  to denote an orthonormal basis, and  $\mathcal{B}_{\varepsilon}$  to denote an  $\varepsilon$  almost-orthonormal basis.

This should handle any outstanding notational quirks, so we will continue to the main body.

# 2. Results

Since we'll be building on results from the midterm project, we'll first reacquaint ourselves with the material therein. To avoid complete redundancy of content, we will take the section of this paper as opportunity to improve the quality and clarity of our exposition. In particular, we will add much greater justification and explanation for our proofs, as many were opaque and generally hard to read. Our analysis will be broken into two main parts: first, dealing with  $\varepsilon$ -almost orthonormal bases, and then later, dealing with  $\varepsilon$ -almost orthogonal matrices. These two topics are of course related, as we will detail later. The bulk of our interest will be on the latter, but we will include results about the former as well so as to show two different perspectives on the problem.

# 2.1. $\varepsilon$ -almost orthonormal bases

Let  $\mathcal{V}$  be an n-dimensional inner product space, and let  $\mathcal{B}_{\perp}$  be an orthonormal basis for  $\mathcal{V}$ . One nice property

of  $\mathcal{B}_{\perp}$  is that  $\forall x \in \mathcal{V}$ , we have the representation

$$x = \sum_{i=1}^{n} \langle x, \mathbf{e}_i \rangle \mathbf{e}_i$$

and hence

$$||x||_2^2 = \sum_{i=1}^n |\langle x, \mathbf{e}_i \rangle|^2.$$

To help us build intuition for how it "feels" to work with an  $\varepsilon$ -almost orthonormal basis instead of an orthonormal basis, we will first prove a result about the extent to which the results above fail.

THEOREM 2.1. Let V be finite-dimensional inner product space, and  $n = \dim(V)$ . Let  $\varepsilon > 0$  be given, with the constraint that  $\varepsilon < \frac{1}{n}$ , and let  $\mathcal{B}_{\varepsilon}$  be an  $\varepsilon$ -almost orthonormal basis of V. Now, let  $x \in V$ , and suppose that x has the representation

$$x = \sum_{i=1}^{n} c_i \langle x, \mathbf{e}_i \rangle$$

where  $c_i \in \mathbb{R}$ , and  $\mathbf{e}_i \in \mathcal{B}_{\varepsilon}$ . Then we have

$$\sum_{i=1}^{n} |c_i| \le \frac{\|x\|_2}{\sqrt{\frac{1}{n} - \varepsilon}}$$

*Proof.* First, we will prove a lower bound for  $||x||_2$  in terms of the  $c_i$ , and use this to obtain the desired result. Observe that by the triangle inequality, we have

$$\left| \sum_{i=1}^{n} c_i^2 \right| - \left| \langle x, x \rangle - \sum_{i=1}^{n} c_i^2 \right| \le \left| \langle x, x \rangle \right| = \left\| x \right\|_2^2.$$

Removing the superfluous absolute values, this is just

$$\sum_{i=1}^{n} c_i^2 - \left| \langle x, x \rangle - \sum_{i=1}^{n} c_i^2 \right| \le ||x||_2^2.$$

We seek a bound for the left-hand-side in terms of of the  $c_i$ . We'll work with the second term. Note that

$$\langle x, x \rangle = \sum_{i=1}^{n} \sum_{j=1}^{n} c_i c_j \langle \mathbf{e}_i, \mathbf{e}_j \rangle.$$

Hence we have

$$\left| \langle x, x \rangle - \sum_{i=1}^{n} c_i^2 \right| = \left| \sum_{i=1}^{n} \sum_{j=1}^{n} c_i c_j (\langle \mathbf{e}_i, \mathbf{e}_j \rangle - \delta_{ij}) \right|$$

$$\leq \sum_{i=1}^{n} \sum_{j=1}^{n} |c_i c_j (\langle \mathbf{e}_i, \mathbf{e}_j \rangle - \delta_{ij})|$$

$$= \sum_{i=1}^{n} \sum_{j=1}^{n} |c_i c_j| |(\langle \mathbf{e}_i, \mathbf{e}_j \rangle - \delta_{ij})|$$

Recall that  $\mathcal{B}_{\varepsilon}$  is an  $\varepsilon$ -almost orthonormal basis implies that  $\forall \mathbf{e}_i, \mathbf{e}_j \in \mathcal{B}_{\varepsilon}$ , we have  $|\langle \mathbf{e}_i, \mathbf{e}_j \rangle - \delta_{ij}| < \varepsilon$ , and hence we have

$$\leq \sum_{i=1}^{n} \sum_{j=1}^{n} |c_i c_j| \varepsilon$$

$$= \varepsilon \sum_{i=1}^{n} \sum_{j=1}^{n} |c_i c_j|$$

$$= \varepsilon \left( \sum_{i=1}^{n} |c_i| \right)^2.$$

This gives us the inequality

$$\sum_{i=1}^{n} c_i^2 - \varepsilon \left( \sum_{i=1}^{n} |c_i| \right)^2 \le \sum_{i=1}^{n} c_i^2 - \left| \langle x, x \rangle - \sum_{i=1}^{n} c_i^2 \right|, \quad (1)$$

and so we can chain together with our original bound to yield

$$\left(\sum_{i=1}^{n} c_i^2\right) - \varepsilon \left(\sum_{i=1}^{n} |c_i|\right)^2 \le \|x\|_2^2$$

By Hölder's inequality for sums, we have

$$\left(\sum_{i=1}^{n} |c_i|\right)^2 = \left(\sum_{i=1}^{n} c_i \operatorname{sgn}(c_i)\right)^2$$

$$\leq \left(\sum_{i=1}^{n} c_i^2\right) \left(\sum_{i=1}^{n} \operatorname{sgn}(c_i)^2\right)$$

$$= \left(\sum_{i=1}^{n} c_i^2\right) \cdot n$$

whence

 $\triangle$ 

$$\frac{1}{n} \left( \sum_{i=1}^{n} |c_i| \right)^2 \le \sum_{i=1}^{n} c_i^2$$

and so we can chain together with (1) to obtain

$$\frac{1}{n} \left( \sum_{i=1}^{n} |c_i| \right)^2 - \varepsilon \left( \sum_{i=1}^{n} |c_i| \right)^2 \le \|x\|_2^2$$

From which we obtain

$$\left(\sum_{i=1}^{n} |c_i|\right)^2 \le \frac{\|x\|_2^2}{\frac{1}{n} - \varepsilon}$$

$$\sum_{i=1}^{n} |c_i| \le \frac{\|x\|_2}{\sqrt{\frac{1}{n} - \varepsilon}}$$

as desired.

As a corollary, we have the following result:

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 $\begin{array}{llll} \text{Corollary} & 2.1.1. & \textit{With all variables quantified as} \\ \textit{above, we have} \end{array}$ 

$$|\langle x, \mathbf{e}_i \rangle - c_i| \le \frac{\varepsilon ||x||_2^2}{\sqrt{\frac{1}{n} - \varepsilon}}.$$

for every  $i = 1, \ldots, n$ .

Proof. Observe that

$$\langle x, \mathbf{e}_i \rangle = \left\langle \sum_{j=1}^n c_j \mathbf{e}_j, \ \mathbf{e}_i \right\rangle$$

Thus we have

$$|\langle x, \mathbf{e}_i \rangle - c_i| = \left| \left\langle \sum_{j=1}^n c_j \mathbf{e}_j, \ \mathbf{e}_i \right\rangle - c_i \right|$$

$$= \left| \left( \sum_{j=1}^n c_j \langle \mathbf{e}_j, \mathbf{e}_i \rangle \right) - c_i \right|$$

$$= \left| \sum_{j=1}^n c_j (\langle \mathbf{e}_j, \mathbf{e}_i \rangle - \delta_{ij}) \right|$$

$$\leq \varepsilon \left| \sum_{j=1}^n c_j \right|$$

$$\leq \varepsilon \sum_{j=1}^n |c_j|$$

$$\leq \frac{\varepsilon ||x||_2^2}{\sqrt{\frac{1}{n} - \varepsilon}}$$

as desired.

REMARK. Note that if  $\varepsilon \ll \frac{1}{n}$ , then the error is essentially  $O(\varepsilon)$ .

#### 2.2. $\varepsilon$ -Almost Orthogonal Matrices

We now turn our attention to  $\varepsilon$ -almost orthogonal matrices. We begin by showing the connection to  $\varepsilon$ -almost orthonormal bases, and then introducing the problem that will be our main focus for this paper.

#### 2.2.1. Introduction

THEOREM 2.2. Let V be a a finite-dimensional inner product space with  $\dim(V) = n$ . Let  $Q \in \operatorname{Orth}(n)$ , and  $\xi \in B_{\varepsilon}(0_{n \times n})$ . Then let  $M = Q + \xi$ , i.e.:

$$M = \begin{bmatrix} q_{11} + \xi_{11} & \dots & q_{1n} + \xi_{1n} \\ \vdots & \ddots & \vdots \\ q_{n1} + \xi_{n1} & \dots & q_{nn} + \xi_{nn} \end{bmatrix}$$

Let  $\mathbf{v}_i, \mathbf{v}_j$  be the  $i^{th}$  and  $j^{th}$  columns of M, respectively. Then

$$\langle \mathbf{v}_i \mathbf{v}_j \rangle \le 2\sqrt{n\varepsilon} + n\varepsilon^2.$$

Note that this bound is not tight. In particular, the coefficient of  $2\sqrt{n}$  on the  $\varepsilon$  term is somewhat crude and can be lowered.  $\triangle$ 

*Proof.* We have

$$\langle \mathbf{v}_{i}, \mathbf{v}_{j} \rangle = \sum_{k=1}^{n} (q_{ki} + \xi_{ki})(q_{kj} + \xi_{kj}) \langle \mathbf{e}_{k}, \mathbf{e}_{k} \rangle$$

$$= \sum_{k=1}^{n} q_{ki}q_{kj} + \xi_{ki}q_{kj} + q_{kj}\xi_{ki} + \xi_{ki}\xi_{kj}$$

$$= \langle \mathbf{q}_{i}, \mathbf{q}_{j} \rangle^{*} + \sum_{k=1}^{n} \xi_{ki}q_{kj} + q_{kj}\xi_{ki} + \xi_{ki}\xi_{kj}$$

$$\leq \sum_{k=1}^{n} \varepsilon(q_{kj} + q_{ki}) + \varepsilon^{2}$$

$$= \varepsilon \left(\sum_{k=1}^{n} (q_{kj} + q_{ki})\right) + n\varepsilon^{2}$$

$$\leq \varepsilon \left(2\sum_{k=1}^{n} q_{kj}\right) + n\varepsilon^{2}$$

$$\leq \varepsilon \left(2\sum_{k=1}^{n} \frac{1}{\sqrt{n}}\right) + n\varepsilon^{2}$$

$$= 2\sqrt{n}\varepsilon + n\varepsilon^{2}$$
(\*)

As desired. Note that to obtain (\*), we do not require that  $q_{kj} \leq \frac{1}{\sqrt{n}}$  for each k — we are bounding the sum, not the individual terms.

From the theorem, we see that the idea of an  $\varepsilon$ -almost orthogonal matrix is related to that of an  $\varepsilon$ -almost orthonormal basis, in that the columns of the  $\varepsilon$ -almost orthogonal matrix are an  $O(\varepsilon)$  almost-orthonormal basis.

Theorem 2.3. Let  $\varepsilon > 0$  be given, and let  $\xi \in B_{\varepsilon}(0_{n \times n})$ . Then  $\|\xi\|_F \leq \varepsilon n$ .

*Proof.* Observe that

$$\|\xi\|_{F} = \sqrt{\sum_{i,j \le n} |\xi_{ij}|^{2}}$$

$$\leq \sqrt{\sum_{i,j \le n} \varepsilon^{2}}$$

$$= \sqrt{\varepsilon^{2} n^{2}}$$

$$= \varepsilon n$$

as desired.

Now, we'll introduce the problem that will be the main focus of our paper today. An important property of orthogonal matrices is that they can be used to perform fast matrix exponentiation for symmetric matrices. That is, let M be a real symmetric matrix. Then by the spectral theorem, M can be diagonalized by  $A\Lambda A^{\mathsf{T}}$ , where A is an orthogonal matrix. Using this, we can perform rapid exponentiation of M by utilizing the fact  $M^n = A\Lambda^n A^\mathsf{T}$ , which can be computed quite rapidly. However, we can imagine situations in which we have access to an almost orthogonal matrix in some matrix exponentiation problem, but not an orthogonal matrix. For instance, imagine we have an IMU measuring the body frame of some machine, say, a UAV, encoded as a matrix B. Say there is measurement error  $\xi \sim N(\mathbf{0}, \varepsilon_{m \times n})$  in the hardware, such that B is now an  $\varepsilon$ -Almost Orthogonal Matrix. In this situation, what kinds of errors can we expect in the exponentiation process if we treat B like it's orthogonal? That is, if we assume

$$(BDB^{\mathsf{T}})^n = BD^nB^{\mathsf{T}}$$

What kinds of errors will this introduce into our calculations? We examine this question below, and prove asymptotic bounds for the errors we might see.

# 2.2.2. Deterministic Matrices

We will begin by examining the problem in the restricted context of non-random matrices, as the proofs are much simpler. Here, we will be primarily interested in two error metrics, which we call  $\delta$  and  $\Delta$ , as defined below.

DEFINITION 2.1 (Error Metrics for Diagonalization). Let  $\mathcal V$  be a finite-dimensional inner product space of dimension  $n\in\mathbb N$ , and let  $\varepsilon>0$  be given. Let  $Q\in\operatorname{Orth}(n)$  and  $\xi\in B_\varepsilon(0_{n\times n})$ , and let  $M=A+\xi$ . Then we define our Diagonalization error metrics by

$$\delta(k) = \|QD^k Q^\mathsf{T} - (MDM^\mathsf{T})^k\|_F^2$$

and

$$\Delta(k) = \|MD^k M^{\mathsf{T}} - (MDM^{\mathsf{T}})^k\|_F^2.$$

We will refer to the stuff inside the norm in the expression for  $\delta$  as  $\mathbf{v}_{\delta}$ , and the stuff inside the norm in the expression for  $\Delta$  as  $\mathbf{v}_{\Delta}$ .

Before we begin working with these error metrics, it will be useful to prove some technical lemmas to aid in proving our first theorem.

Lemma 2.4. Let  $\Lambda, D$  be diagonal  $n \times n$  matrices, and let H be an orthogonal matrix. Then

$$\operatorname{tr}(HD^kH^\mathsf{T}\Lambda^k) = \sum_{i,j}^n h_{ij}^2 \lambda_i^k(D) \lambda_j^k(\Lambda)$$

*Proof.* We'll start from the left-hand-side, first calculating  $HD^kH^{\mathsf{T}}$ . Note that indexing variables might get a little wacky here, because we have a lot of matrices in the expression. Observe that, by the standard matrix multiplication formula  $c_{ij} = \sum_{k=1}^{n} a_{ik}b_{kj}$ , we have

$$(D^k H^\mathsf{T})_{ij} = \sum_{\alpha=1}^n d_{i\alpha} h_{j\alpha}$$
$$= \sum_{\alpha=1}^n \lambda_i^k(D) h_{j\alpha} \delta_{i\alpha}$$

Where  $\delta_{i\alpha}$  is the kronecker delta. Thus, applying this rule a second time, we have

$$(H(D^k H^\mathsf{T}))_{ij} = \sum_{\beta=1}^n (H)_{i\beta} (D^k H^\mathsf{T})_{\beta j}$$
$$= \sum_{\beta=1}^n h_{i\beta} \sum_{\alpha=1}^n \lambda_\beta^k(D) h_{j\alpha} \delta_{\beta \alpha}$$

And so finally,

$$(HD^{k}H^{\mathsf{T}}\Lambda^{k})_{ij} = \sum_{\gamma=1}^{n} (HD^{k}H^{\mathsf{T}})_{i\gamma}(\Lambda^{k})_{\gamma j}$$

$$= \sum_{\gamma=1}^{n} \left( \sum_{\beta=1}^{n} h_{i\beta} \sum_{\alpha=1}^{n} \lambda_{\beta}^{k}(D) h_{\gamma \alpha} \delta_{\beta \alpha} \right) \lambda_{\gamma}^{k}(\Lambda) \delta_{\gamma j}$$

$$= \sum_{\gamma=1}^{n} \sum_{\beta=1}^{n} \sum_{\alpha=1}^{n} h_{i\beta} h_{\gamma \alpha} \lambda_{\beta}^{k}(D) \lambda_{\gamma}^{k}(\Lambda) \delta_{\beta \alpha} \delta_{\gamma j}$$

Thus, the trace is given by

$$\operatorname{tr}(HD^{k}H^{\mathsf{T}}\Lambda^{k}) = \sum_{\ell=1}^{n} (HD^{k}H^{\mathsf{T}}\Lambda^{k})_{\ell\ell}$$

$$= \sum_{\ell=1}^{n} \sum_{\gamma=1}^{n} \sum_{\beta=1}^{n} \sum_{\alpha=1}^{n} h_{\ell\beta} h_{\gamma\alpha} \lambda_{\beta}^{k}(D) \lambda_{\gamma}^{k}(\Lambda) \delta_{\beta\alpha} \delta_{\gamma\ell}$$

$$= \sum_{\ell=1}^{n} \sum_{m=1}^{n} h_{\ell m} h_{\ell m} \lambda_{m}^{k}(D) \lambda_{\ell}^{k}(\Lambda)$$

$$= \sum_{\ell=1}^{n} \sum_{m=1}^{n} h_{\ell m}^{2} \lambda_{m}^{k}(D) \lambda_{\ell}^{k}(\Lambda)$$

as desired.

COROLLARY 2.4.1. Let  $\Lambda, D, H$  be as defined above. Then  $\operatorname{tr}(\Lambda^k H D^k H^\mathsf{T}) = \operatorname{tr}(H D^k H^\mathsf{T} \Lambda^k)$ .

*Proof.* This follows simply from the fact that trace is invariant under cyclic permutations.

Now, we proceed to the first of our main claims.

THEOREM 2.5. Let V be a finite-dimensional inner product space with  $\dim V = n \in \mathbb{N}$ , and let  $\varepsilon > 0$  be given. Let  $Q \in \operatorname{Orth}(n)$  and  $\xi \in B_{\varepsilon}(0_{n \times n})$ , and let  $M = Q + \xi$ . Now, let D be an  $n \times n$  diagonal matrix, and let  $\sigma_0 = \max\{|\sigma_i(D)|\}$ , and  $\tau_0 = \max\{|\sigma_j(MDM^{\mathsf{T}})|\}$  be the largest singular values of D and  $MDM^{\mathsf{T}}$ , respectively. Then if  $\nu = \max\{\sigma_0, \tau_0\}$ , then

$$\delta(k) \sim \alpha \nu^{2k}$$

for almost all cases, where  $\alpha$  is a constant that depends on D and  $MDM^T$ . We list edge cases as follows:

- (i) If the eigenvalues of D and  $MDM^{\mathsf{T}}$  are equal, then  $\delta(k) = 0$ .
- (ii) If the singular values of D and  $MDM^{\mathsf{T}}$  are equal but some of the eigenvalues are not (i.e., the eigenvalues have the same magnitudes but some have different signs), then  $\delta(k) = 0$  for even k and  $\delta(k) = \alpha c^{2k}$  for odd k.
- (iii) If the multiplicity of  $\nu$  in  $\sigma(D)$ ,  $\sigma(\Lambda)$  is the same, then whenever k is even,  $\delta(k) = o(\nu^{2k})$  (i.e.,  $\delta$  is subexponential in k).

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*Proof.* Note that  $MDM^{\mathsf{T}}$  is symmetric:

$$(MDM^{\mathsf{T}}) = (M^{\mathsf{T}})^{\mathsf{T}} D^{\mathsf{T}} M^{\mathsf{T}}$$
$$= MDM^{\mathsf{T}}$$

Thus, by the real spectral theorem,  $MDM^{\mathsf{T}}$  is orthogonally diagonalizable, and so there exists an orthogonal matrix P and a diagonal matrix  $\Lambda$  such that  $MDM^{\mathsf{T}} = P\Lambda P^{\mathsf{T}}$ . Hence,

$$\delta(k) = \|QD^kQ^\mathsf{T} - (MDM^\mathsf{T})^k\|_F^2$$
$$= \|QD^kQ^\mathsf{T} - (P\Lambda P^\mathsf{T})^k\|_F^2$$
$$= \|QD^kQ^\mathsf{T} - P\Lambda^k P^\mathsf{T}\|_F^2$$

Since the Frobenius Norm is unitary invariant, it follows that

$$= \|P^{T}(QD^{k}Q^{\mathsf{T}} - P\Lambda^{k}P^{\mathsf{T}})P\|_{F}^{2}$$
$$= \|P^{\mathsf{T}}QD^{k}Q^{\mathsf{T}}P - \Lambda^{k}\|_{F}^{2}$$
$$= \operatorname{tr}\left(\left(P^{\mathsf{T}}QD^{k}Q^{\mathsf{T}}P - \Lambda^{k}\right)^{2}\right)$$

For the sake of concision, let  $H = P^{\mathsf{T}}Q$ . Then we can write this as

$$\delta(k) = \operatorname{tr}\left(\left(HD^{k}H^{\mathsf{T}} - \Lambda^{k}\right)^{2}\right)$$

$$= \operatorname{tr}\left(HD^{2k}H^{\mathsf{T}} - HD^{k}H^{\mathsf{T}}\Lambda^{k} - \Lambda^{k}HD^{k}H^{\mathsf{T}} + \Lambda^{2k}\right)$$
by the Lemma and its Corollary, we have
$$= \operatorname{tr}\left(HD^{2k}H^{\mathsf{T}}\right) + \operatorname{tr}\left(\Lambda^{2k}\right) - 2\operatorname{tr}\left(HD^{k}H^{\mathsf{T}}\Lambda^{k}\right)$$

$$= \left(\sum_{i=1}^{n} \lambda_{i}^{2k}(D) + \lambda_{i}^{2k}(\Lambda)\right) - 2\sum_{i,j \leq n} h_{ij}^{2} \lambda_{i}^{k}(D) \lambda_{i}^{k}(\Lambda).$$

Let  $\nu = \max \{\sigma_0, \tau_0\}$ . Suppose that for all  $i = 1, \ldots, i_0 \le n, j = 1, \ldots, j_0 \le n$ , we have

$$\nu = |\lambda_i(D)| = |\lambda_j(\Lambda)|,$$

and without loss of generality, suppose that  $i_0 \geq j_0$ , and none of the other eigenvalues satisfy this equality. Then we have two cases:  $j_0 < n$ , and  $j_0 = n$  (and hence  $i_0 = n$  as well).

(i) Suppose  $j_0 < n$ . We evaluate the left-hand term in our expression for  $\delta(k)$  first. Because  $\nu$  was chosen to be the max of the singular values, we have

$$\lim_{k \to \infty} \frac{\sum_{i=1}^{n} \lambda_i^{2k}(D) + \lambda_i^{2k}(\Lambda)}{\nu^{2k}} = i_0 + j_0.$$

Now, we examine the right-hand term. By similar reasoning,

RHS = 
$$\lim_{k \to \infty} -2 \frac{\sum_{i,j \le n} h_{ij}^2 \lambda_i^k(D) \lambda_i^k(\Lambda)}{\nu^{2k}}$$
= 
$$-2 \frac{\sum_{i,j \le j_0} h_{ij}^2 \lambda_i^k(D) \lambda_i^k(\Lambda)}{\nu^{2k}}$$
= 
$$-2 \sum_{i,j \le j_0} (-1)^{\beta_{ij}}$$

where  $\beta_{ij}$  is defined by

$$(-1)^{\beta_{ij}} = \frac{\lambda_i^k(D)\lambda_i^k(\Lambda)}{\nu^{2k}}.$$

(This is because for  $i, j \leq j_0$ , we have  $\left|\lambda_i^k(D)\lambda_i^k(\Lambda)\right| = \nu^{2k}$ , by definition of  $j_0$ ). Now, because H is orthogonal, the columns of H are orthonormal, hence we obtain the inequality

$$\sum_{i,j \le j_0} (-1)^{\beta_{ij}} h_{ij}^2 < \sum_{i=1}^n \sum_{j=1}^{j_0} h_{ij}^2$$

$$= \sum_{j=1}^{j_0} \|\mathbf{h}_j\|_2^2$$

$$= j_0$$

with the inequality being strict since  $j_0 < n$ . This yields the lower bound  $i_0 + j_0 - 2j_0 = i_0 - j_0$ :

$$0 \le i_0 - j_0 < \lim_{k \to \infty} \frac{\delta(k)}{\nu^{2k}}$$

we obtain an upper bound when we always have  $\beta_{ij} = 1$ , yielding

$$\lim_{k \to \infty} \frac{\delta(k)}{\nu^{2k}} < i_0 + 3j_0$$

thus exists a constant  $\alpha > 0$  such that

$$\lim_{k \to \infty} \frac{\delta(k)}{\nu^{2k}} = \alpha$$

hence

$$\delta(k) \sim \alpha \nu^{2k}$$

as desired.

(ii) Now, suppose  $i_0 = n = j_0$ . Then all of the singular values of D and  $\Lambda$  are the same. Suppose that k of the eigenvalues are of the same sign. If k is even, then  $\lambda_i^{2k}(D) = \lambda_i^{2k}(\Lambda)$  for all i, and hence  $(-1)^{\beta_{ij}} = 1$ . It follows that

$$\sum_{i,j \le n} (-1)^{\beta_{ij}} h_{ij}^2 = \sum_{i,j \le n} h_{ij}^2$$
$$= n$$

hence we have

$$\delta(k) = n + n - 2n$$
$$= 0$$

In fact, in general,  $\delta(k)$  is subexponential in k whenever the multiplicity of  $\nu$  is the same in both  $\sigma(D)$  and  $\sigma(\Lambda)$ .

If k is odd, then we have

$$\sum_{i,j \le n} (-1)^{\beta_{ij}} < h_{ij}^2 = n,$$

and so

$$\delta(k) = \left(2n - \sum_{i,j \le n} (-1)^{\beta_{ij}} h_{ij}^2\right) \nu^{2k} = \alpha \nu^{2k},$$

where

$$\alpha = \left(2n - \sum_{i,j \le n} (-1)^{\beta_{ij}} h_{ij}^2\right).$$

Corollary 2.5.1. Quantify all variables as above. Then  $\alpha \leq 4n$ .

*Proof.* Note that  $\forall i, j, |2h_{ij}\lambda_i^k(D)\lambda_i^k(\Lambda)| \leq 2h_{ij}^2\nu^{2k}$ . Then because  $\sum_{i,j\leq n}h_{ij}^2 = n$ , we have

$$\begin{split} \delta(k) &= \sum_{i=1}^n \lambda_i^{2k}(D) + \lambda_i^{2k}(\Lambda) - 2\sum_{i,j \leq n} h_{ij}^2 \lambda_i^k(D) \lambda_i^k(\Lambda) \\ &\leq 2n\nu^{2k} + 2n\nu^{2k} \\ &= 4n\nu^{2k} \end{split}$$

as desired.

Before deriving some results for  $\Delta$ , we should observe that nowhere in the proof above did we utilize the condition that  $M=Q+\xi$ . Surely, we conjecture, we should be able to use this fact to tighten our bounds. Indeed, this is the case. We outline the proof in the theorem below. But first, we have a small lemma.

LEMMA 2.6. Let  $A, B \in M_{n \times n}(\mathbb{R})$ . Then AB and BA share the same eigenvalues.  $\triangle$ 

*Proof.* Let v be an eigenvector of AB, with associated eigenvalue  $\lambda$ . Then we have

$$ABv = \lambda v$$
.

Left multiplying each side by B, we have

$$B(ABv) = B\lambda v$$
$$= \lambda Bv$$
$$= BA(Bv)$$

hence Bv is an eigenvector of BA with associated eigenvalue  $\lambda$ . Thus, every eigenvalue of AB is an eigenvalue of BA. Conversely, let u be an eigenvector of BA, with associated eigenvalue  $\tau$ . Then

$$A(BAu) = A\tau u$$
$$= \tau Au$$
$$= AB(Au)$$

hence Au is an eigenvector of AB with associated eigenvalue  $\tau$ . Thus, AB and BA have the same spectrum.

COROLLARY 2.6.1. Let  $Q \in \text{Orth}(n)$ , and let  $M \in M_{n \times n}(\mathbb{R})$ . Then M and  $QMQ^{\mathsf{T}}$  have the same spectrum.  $\triangle$ 

*Proof.* Take A = QM and  $B = Q^{\mathsf{T}}$ , then apply the lemma above.

In the theorem below, we will apply Weyl's inequality for perturbations, hence we give the statement below for completeness:

THEOREM 2.7 (Weyl). Let  $H, P \in M_{n \times n}(\mathbb{R})$ , and let M = H + P. Suppose M has eigenvalues

$$\mu_1 \ge \mu_2 \ge \cdots \ge \mu_n$$

H has eigenvalues

$$\eta_1 \geq \eta_2 \geq \cdots \geq \eta_n,$$

and P has eigenvalues

$$\rho_1 \ge \rho_2 \ge \cdots \ge \rho_n.$$

Then if any two of M, H, P are Hermitian,  $\forall i = 1, ..., n$ , we have

$$\eta_i + \rho_n \le \mu_i \le \eta_i + \rho_1.$$

Now the theorem.

THEOREM 2.8. Let  $Q \in \text{Orth}(n)$  and  $D \in M_{n \times n}(\mathbb{R})$  be diagonal, and let  $\xi \in B_{\varepsilon}(0_{n \times n})$ . Now, let  $M = Q + \xi$ , and let  $\lambda_1, \lambda_2, \ldots, \lambda_n$  be the eigenvalues of  $QDQ^{\mathsf{T}}$ , and let  $\tau_1, \tau_2, \ldots, \tau_n$  be the eigenvalues of  $MDM^{\mathsf{T}}$ . If  $\lambda = \max_i(|\lambda_i|)$ , then

$$|\lambda_i - \tau_i| \le \lambda (2\varepsilon n + \varepsilon^2 n^2)$$

Δ

 $\triangle$ 

*Proof.* For this proof, we will make use of the spectral norm, hence we first note some properties of the spectra of  $QDQ^\mathsf{T}$  and  $MDM^\mathsf{T}$ . By the corollary above, note that the eigenvalues of  $QDQ^\mathsf{T}$  are simply the entries of D. Similarly, note that the spectrum of  $MDM^\mathsf{T}$  is equivalent to the spectrum of  $M^\mathsf{T}MD$ . That is,

$$\sigma(MDM^{\mathsf{T}}) = \sigma(M^{\mathsf{T}}MD)$$

$$= \sigma((Q+\xi)^{\mathsf{T}}(Q+\xi)D)$$

$$= \sigma(D+\xi^{\mathsf{T}}QD+Q^{\mathsf{T}}\xi D+\xi^{\mathsf{T}}\xi D)$$

and thus

$$\begin{split} \left\| QDQ^\mathsf{T} - MDM^\mathsf{T} \right\|_2 &= \left\| D - M^\mathsf{T} M d \right\|_2 \\ &= \left\| \xi^\mathsf{T} QD + Q^\mathsf{T} \xi D + \xi^\mathsf{T} \xi D \right\|_2 \end{split}$$

Since  $QDQ^{\mathsf{T}}$  and  $MDM^{\mathsf{T}}$  are symmetric (and thus hermetian), we can apply Weyl's inequality by taking  $M = QDQ^{\mathsf{T}}$ ,  $H = MDM^{\mathsf{T}}$ , and  $P = QDQ^{\mathsf{T}} - MDM^{\mathsf{T}}$  to obtain the following bound:

$$|\lambda_i - \tau_i| \le ||D - M^\mathsf{T} M D||_2$$
$$= ||\xi^\mathsf{T} Q D + Q^\mathsf{T} \xi D + \xi^\mathsf{T} \xi D||_2$$

by submultiplicativity of the spectral norm, we have

$$\leq \|D\|_2 \|\xi^\mathsf{T} Q + Q^\mathsf{T} \xi + \xi^\mathsf{T} \xi\|_2$$
$$= \lambda \|\xi^\mathsf{T} Q + Q^\mathsf{T} \xi + \xi^\mathsf{T} \xi\|_2$$

and so by the triangle inequality,

$$\leq \lambda (\|\xi^{\mathsf{T}}Q\|_{2} + \|Q^{\mathsf{T}}\xi\|_{2} + \|\xi^{\mathsf{T}}\xi\|_{2})$$

$$\leq \lambda (\|\xi^{\mathsf{T}}\|_{2}\|Q\|_{2} + \|Q^{\mathsf{T}}\|_{2}\|\xi\|_{2} + \|\xi^{\mathsf{T}}\|_{2}\|\xi\|_{2})$$

$$= \lambda (\|\xi^{\mathsf{T}}\|_{2} + \|\xi\|_{2} + \|\xi^{\mathsf{T}}\|\|\xi\|)$$

$$= \lambda (2\|\xi\|_{2} + \|\xi\|_{2}^{2})2$$

$$\leq \lambda (2\|\xi\|_{F} + \|\xi\|_{F}^{2})2$$

$$\leq \lambda (2\varepsilon n + \varepsilon^{2}n^{2})$$

as desired.

We have the following corollary:

COROLLARY 2.8.1. Let  $Q \in \text{Orth}(n)$ , and let  $D \in M_{n \times n}(\mathbb{R})$  be diagonal. Suppose  $M = Q + \xi$  where  $\xi \in B_{\varepsilon}(0_{n \times n})$ . Let  $\lambda_1^k, \lambda_2^k, \ldots, \lambda_n^k$  be the eigenvalues of  $D^k$ , and let  $\tau_1(k), \tau_2(k), \ldots, \tau_n(k)$  be the eigenvalues of  $MD^kM^{\mathsf{T}}$  for  $k \in \mathbb{N}$ . If  $\lambda = \max_i(|\lambda_i|)$ , then

$$\left|\lambda_i^k - \tau_i(k)\right| \le \lambda^k \left(2\varepsilon n + \varepsilon^2 n^2\right)$$

 $\triangle$ 

*Proof.* This follows immediately by replacing D in the previous theorem by  $D^k$ .

Using this result, we can tighten some of the bounds we presented earlier.

COROLLARY 2.8.2. Let  $Q \in \text{Orth}(n)$  and let  $D \in M_{n \times n}(\mathbb{R})$  be diagonal. Suppose  $M = Q + \xi$  where  $\xi \in B_{\varepsilon}(0_{n \times n})$ . Then for sufficiently small  $\varepsilon$ ,

$$\delta(k) \le 4n(d(1+2\varepsilon n+\varepsilon^2 n^2))^{2k}$$
 where  $d = \max_i(|d_i|)$   $\triangle$ 

*Proof.* In light of Theorem 2.8 we can assume for small enough  $\varepsilon$  that the largest singular value  $\lambda_l$  in  $\sigma(QDQ^{\mathsf{T}})$  corresponds to the largest singular value  $\tau_l$  in  $\sigma(MDM^{\mathsf{T}})$ . Therefore, we get that  $\nu$  (the same  $\nu$  in Theorem 2.5) is bounded by

$$d + d(2\varepsilon n + \varepsilon^2 n^2)$$

Applying Corollary 2.5.1 gives us the desired bound. ■

Note that, for all of the arguments above, by the equivalence of matrix norms, the results also hold (up to a constant multiple) under the Frobenius norm. Thus, we can apply them to our original bounds.

COROLLARY 2.8.3. Let  $Q \in \text{Orth}(n)$ , and let  $\xi \in B_{\varepsilon}(0_{n \times n})$ . As usual, let  $M = Q + \xi$ , and let  $D \in M_{n \times n}(\mathbb{R})$  be diagonal, and let  $\lambda = \max_i (|\lambda_i|)$ . Then  $\left|\sqrt{\Delta(k)} - \sqrt{\delta(k)}\right| \leq \lambda^k (2 + \varepsilon n + \varepsilon^2 n^2)$ .

*Proof.* Proving this result is now easy. Let  $\lambda = \max_{i} \lambda_{i}(D)$ . By the reverse triangle inequality, we have

$$\left| \sqrt{\Delta(k)} - \sqrt{\delta(k)} \right| = \left| \left\| \mathbf{v}_{\Delta} \right\|_{F} - \left\| \mathbf{v}_{\delta} \right\| \right|_{F}$$

$$\leq \left\| \mathbf{v}_{\Delta} - \mathbf{v}_{\delta} \right\|$$

$$= \left\| M D^{k} M^{\mathsf{T}} - Q D^{k} Q^{\mathsf{T}} \right\|$$

$$\leq \lambda^{k} (2\varepsilon n + \varepsilon^{2} n^{2})$$

as desired.

We're interested in the bounds on  $\Delta(k)$ . We'll use a crude approach here to approximate it.

LEMMA 2.9. Let  $M = Q + \xi$ , where  $Q \in \text{Orth}(n)$ , and  $\xi \in B_{\varepsilon}(0_{n \times n})$ . Let  $D \in M_{n \times n}(\mathbb{R})$  be diagonal. Then for  $\varepsilon < 1$ ,  $\|(MDM^{\mathsf{T}})^k\|_2 \le \|D\|_2^k (1 + O(\varepsilon))$   $\triangle$ 

Proof.

$$\begin{split} \left\| \prod_{i=1}^k MDM^\mathsf{T} \right\|_2 &= \left\| \prod_{i=1}^k (Q+\xi)D(Q+\xi)^\mathsf{T} \right\|_2 \\ &\leq \prod_{i=1}^k \left\| (Q+\xi)D(Q+\xi)^\mathsf{T} \right\|_2 \\ &\leq \prod_{i=1}^k \left\| QDQ^\mathsf{T} \right\|_2 + \left\| \xi DQ^\mathsf{T} \right\|_2 \\ &+ \left\| QD\xi^\mathsf{T} \right\|_2 + \left\| \xi D\xi^\mathsf{T} \right\|_2 \\ &= \prod_{i=1}^k \left\| D \right\|_2 \left( 1 + 2\varepsilon + \varepsilon^2 \right) \end{split}$$

$$= \|D\|_2^k \sum_{i=1}^{2k} {2k \choose i} \varepsilon^i$$
$$= \|D\|_2^k (1 + O(\varepsilon)),$$

As desired.

REMARK. This bound is exceptionally crude. See computational results section for performance analysis.

# 3. Computational Results & Randomness

We made considerable efforts to obtain theoretical answers to the following questions:

- If we were to draw  $\xi$  from  $\mathcal{N}(0_{n \times n}, \varepsilon_{n \times n})$ , how much tighter would our bounds on  $\mathbb{E}[\delta(k)], \mathbb{E}[\Delta(k)]$  be than our current bounds on  $\delta(k), \Delta(k)$ ?
- On a related note, if we sample the entries of our random matrix  $\xi$  from another distribution, what
- How can we view the Johnson-Lindenstrauss Lemma from the perspective of almost orthogonal matrices?
- How can  $n \times k$  (where  $k \ll n$ ) almost-orthogonal matrices be used in
- Do we always have  $\delta(k) > \Delta(k)$ ?
- Is  $\delta(k) \sim \Delta(k)$ ?

However, after spending a lot of time trying to find literature on these topics and/or trying to prove the bounds ourselves from scratch, we realized we simply didn't have adequate machinery yet to try and approach these problems from a theoretical standpoint. As such, we decided to take a more computational approach, as is detailed in the following section.

# 3.1. Testing $\delta$ , $\Delta$ Bounds

To investigate the first question, we created some python code to test the  $\mathbb{E}[\delta(k)]$  and  $\mathbb{E}[\Delta(k)]$  when  $\xi \sim \mathcal{N}(0_{n \times n}, \varepsilon)$  as a function of n, k, and  $\varepsilon$ . The  $\varepsilon$  dependence was not terribly interesting, hence we'll focus mainly on the k, n dependence here. Implementation details can be found in the appendix.

*Proof.* (For result 3.1; plots on next page)

The first big hurdle we have to work past is how we can test the average-case performance of the algorithm. Generating acceptably-distributed  $\xi$  is trivial, since we're given a particular distribution we want them to follow. However, we also want to make sure that our orthogonal matrices are also generated in an evenly-distributed manner, to ensure that we're not skewing the results by our particular choice of Q. We guarantee this as follows:

**Data:** Number of dimensions, n**Result:**  $Q \in \text{Orth}(n)$ 

 $M \leftarrow \text{construct a random } n \times n \text{ matrix with uniformly distributed entries;}$ 

 $Q, R \leftarrow \text{apply } QR \text{ decomposition to } M;$ return Q

**Algorithm 1:** Algorithm for generating our orthogonal matrix Q

Then, we perform the following computational scheme:

**Data:** Number of dimensions n,  $\sigma$ , k list, tests per k

**Result:** Plot of  $\log(\delta(k)/\Delta(k))$  initialization of  $\delta$  list,  $\Delta$  list;

 $Q, R \leftarrow \text{apply } QR \text{ decomposition to } M;$ 

return Q for  $k \leftarrow \in k$  list do

initialize list of  $\delta(k)$ ,  $\Delta(k)$  results;

for  $0 \le i \le tests \ per \ k$  do

 $Q \leftarrow$  generate random orthogonal matrix

 $\xi \leftarrow \text{sample from } \mathcal{N}(0_{n \times n}, \sigma_{n \times n});$ 

 $D \leftarrow$  generate a random diagonal matrix D:

calculate  $\delta(k)$ ,  $\Delta(k)$  in the usual way; Push to list of  $\delta(k)$ ,  $\Delta(k)$  results;

end

take the average of the  $\delta(k)$ ,  $\Delta(k)$  lists and push to the  $\delta$ ,  $\Delta$  lists;

 $\mathbf{end}$ 

perform linear regression on the  $\log(\delta)$ ,  $\log(\Delta)$  lists to verify the bound is tight;

**Algorithm 2:** Algorithm for testing the  $\delta$ ,  $\Delta$  bounds

Interestingly, upon making plots of  $\delta(k) - \Delta(k)$  and  $\log\left(\frac{\delta(k)}{\Delta(k)}\right)$ , we obtained the following empirical result:

Result 3.1. 
$$\delta(k) - \Delta(k) = o(\delta(k)/\Delta(k))$$

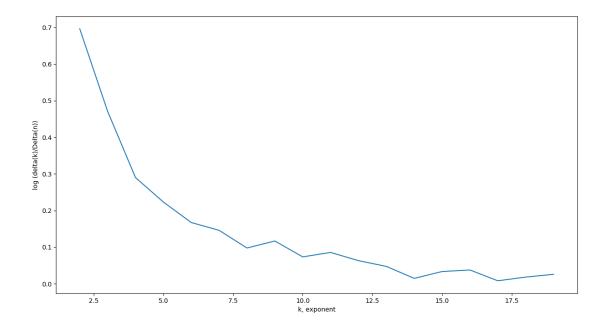


Figure 1:  $\log \left( \frac{\delta(k)}{\Delta(k)} \right)$ 

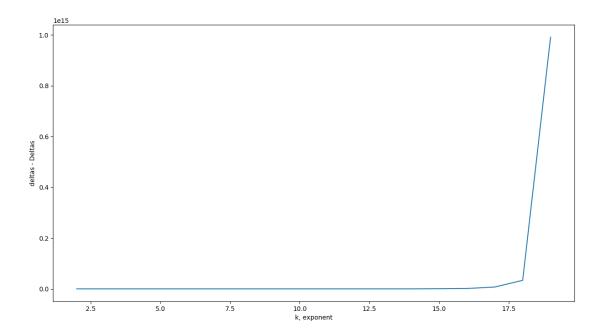


Figure 2:  $\log\left(\frac{\delta(k)}{\Delta(k)}\right)$ 

hence we see  $\log\left(\frac{\delta(k)}{\Delta(k)}\right) \to 1$ , but  $\delta(k) - \Delta(k)$  is unbounded, hence we have the desired result.

Similarly, observe the following empirical verification of the exponential nature of the bounds:

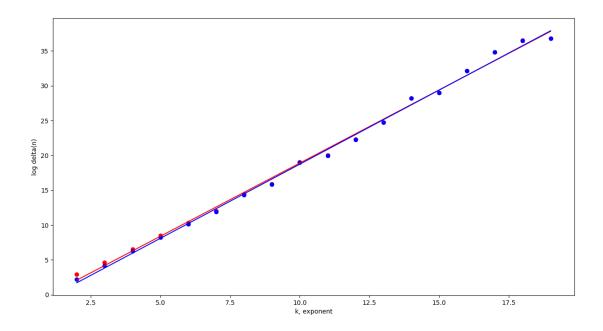


Figure 3:  $\log(\delta(k))$ 

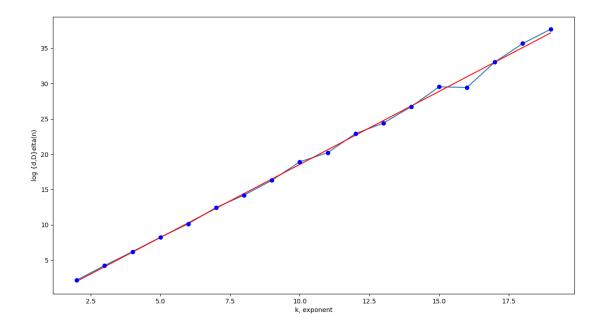


Figure 4:  $\log(\Delta(k))$ 

hence we see  $\log\left(\frac{\delta(k)}{\Delta(k)}\right) \to 1$ , but  $\delta(k) - \Delta(k)$  is unbounded, hence we have the desired result.

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# 4. Code

```
1 | import scipy
  import scipy.linalg
   # import numpy as np
3
   import numpy as np
4
   import matplotlib.pyplot as plt
5
6
   from tqdm import tqdm
   from mpl_toolkits import mplot3d
8
9
10
   class Distribution:
        Just a simple object wrapper so that we can partially instantiate
       a distribution for ease of use in defining perturbations below.
14
15
       def __init__(self, distribution, *args, **kwargs):
16
17
18
            distribution should be a numpy.random.<distribution_name>
19
20
            *args should be whatever parameters the distribution needs to
21
            instantiate
22
23
            self.distribution = distribution
24
25
            self.args = args
            self.kwargs = kwargs
26
            # self.kwargs = kwargs
27
28
       def __call__(self, **extra_kwargs):
29
30
            samples the attached distribution given the args,
31
32
            return self.distribution(*self.args, **self.kwargs, **extra_kwargs)
33
34
35
   # Functions for applying distributions to matrices and stuff
36
37
   def perturb(vec, distribution):
38
        Create a correctly sized perturbation vector and add it to the
39
       input vector
40
41
42
        return vec + distribution(size=vec.shape)
43
   def perturb_mat(M, distribution):
44
45
46
        Create a correctly sized perturbation matrix and add it to the
47
        input matrix
48
       size = M.shape
49
        return M + distribution(size = size)
50
51
52
   def construct_almost_basis(n, distribution):
53
54
55
        Use a given distribution to perturb basis vectors yielding an
        (almost) orthonormal basis
56
57
        I = np.eye(n)
58
        for i, row in enumerate(I):
    I[i] = perturb(row, distribution)
59
60
       return I
61
62
   def orthog_error(A):
63
64
        return error from the square root process as well as the
65
        calculation from the "closest" orthogonal matrix
66
67
68
        sqrt = scipy.linalg.sqrtm(A.T @ A)
        invsqrt = np.linalg.inv(sqrt)
69
       R = A @ invsqrt
70
       return np.linalg.norm(A - R)
71
72
73
   # Helper functions for plot generation
   def f(x,y):
74
75
```

```
76
         Example z_func.
77
         the transformation we'll apply to x and y in our surface plot below
78
79
80
             # return 3*(x ** 2) + (y ** 2)
81
         return x * y / (x**2 + y**2)
except ZeroDivisionError as e:
82
83
84
             return 0
85
86
    def surface_plot(x, y, z_func=f, x_label="x", y_label="y", z_label="z"):
87
         z_func should be a map from R^2 to R
88
89
         # Clear any previous drawings
90
91
         plt.clf()
92
        \# np.meshgrid basically gives an object such that we can apply f over the \# object as seen below instead of looping and stuff
93
94
        X, Y = np.meshgrid(x, y)
95
96
        # Apply f over all of (X,Y)
97
98
        Z = z_func(X, Y)
99
100
         # Initialize figure object
        fig = plt.figure()
101
         # Inform matplotlib we'll have a 3d surface plot
103
        ax = plt.axes(projection='3d')
         # Plot the surface
106
         surface = ax.plot_surface(X, Y, Z, cmap='viridis')
107
108
         # Make a scale thing
        fig.colorbar(surface)
        # axis labels
112
        ax.set_xlabel(x_label)
         ax.set_ylabel(y_label)
114
         ax.set_zlabel(z_label)
117
        # Return object so that it can be used later. Also allows us to
        # decide between "show" and "save"
118
         return (fig, ax)
119
120
121
    def plot_normal_3delbow(min_sigma=0.001, max_sigma=10,
122
                               num_tests=500, samps_per=20):
124
         # This is going to be the first axis thing to iterate over (in my
125
         # case, this is the stuff I'm gonna put on the x axis later)
126
         dims = np.array(range(2,20))
127
128
129
130
         sample_sigma = np.linspace(min_sigma, max_sigma, num=num_tests)
131
         def z_func(dim, sigma):
132
133
             Function to be applied to each entry
134
135
             normal = Distribution(np.random.normal, loc=0.0, scale=sigma)
136
137
             sub_test = []
             for i in range(samps_per):
138
                 almost = construct_almost_basis(dim, normal)
139
                 sub_test += [orthog_error(almost)]
140
                 # cond_num = np.linalg.cond(almost)
141
                 # sub_test += [cond_num]
142
             sub_test = np.array(sub_test)
143
             # print(sub_test)
144
             return np.mean(sub_test)
145
146
        X, Y = np.meshgrid(dims, sample_sigma)
147
148
        z_func = np.vectorize(z_func)
149
150
151
         fig, ax = surface_plot(dims, sample_sigma, z_func=z_func,
                               x_label="dimension of matrix", y_label="std." +
                               "dev. of normal perturbation"
153
```

```
154
                               z_label="distance to closest orthogonal matrix")
         plt.show()
156
    def exp_error(n, A, diag_matrix):
158
159
        n: exponent to raise to
160
161
         Return the error as propogated through when taking the matrix
162
163
         exponential and treating the input
164
         true_matrix = np.linalg.matrix_power(A @ diag_matrix @ A.T, n)
165
         est_matrix = A @ (diag_matrix ** n) @ A.T
166
167
         error = np.linalg.norm(true_matrix - est_matrix)
         return error/(np.linalg.norm(A)**n)
168
169
    def plot_normal_exp_elbow(min_val=0.0001, max_val=0.5,
171
                                 num_tests=500, samps_per=50):
         dim = 15
174
         sigmas = np.linspace(min_val, max_val, num=num_tests)
176
         Z = []
177
         exps = range(2,70)
178
180
         diag_matrix = np.diag(np.random.normal(loc=1.0, scale=.01, size=(dim,)))
181
182
         for exp in tqdm(exps):
             # Reverse the order to start with big errors then go to small
183
184
             y_vals = []
185
186
             # y_errs = []
             for sigma in sigmas:
187
                 normal = Distribution(np.random.normal, loc=0.0, scale=sigma)
188
189
                 sub_test = []
                  for i in range(samps_per):
190
                      # print(normal, dims)
191
                      almost = construct_almost_basis(dim, normal)
192
193
                      sub_test += [exp_error(exp, almost, diag_matrix)]
194
                 # Convert to numpy array so we can use nice easy stats without # writing helper functions
195
196
                 sub_test = np.array(sub_test)
197
198
                 # Use the average value
199
                 y_vals += [np.mean(sub_test)]
200
201
                 # std_dev = np.std(sub_test)
202
                 # y_errs += [std_dev/samps_per]
203
204
             Z += [y_vals]
205
206
207
         Z = np.array(Z).T
208
        X, Y = np.meshgrid(exps, sigmas)
         # print(X.shape, Y.shape, Z.shape)
209
211
        fig = plt.figure()
         ax = plt.axes(projection='3d')
212
        # ax.contour3D(X, Y, Z, 50, cmap='binary')
ax.plot_surface(X, Y, np.log(Z), cmap='binary')
213
214
215
         ax.set_xlabel('exponent')
         ax.set_ylabel('stard deviation of perturbation')
216
        ax.set_zlabel('log normed exponent error')
217
        plt.show()
218
220
221
222
    def plot_uniform_exp_elbow(min_val=0.0001, max_val=20,
                                 num_tests=100, samps_per=100):
224
         dim = 15
225
226
         sample_range = np.linspace(min_val, max_val, num=num_tests)
227
228
229
         exps = range(2,20)
230
         diag_matrix = np.diag(np.random.normal(loc=1.0, scale=0.5, size=(dim,)))
231
```

```
232
         for exp in tqdm(exps):
233
             # Reverse the order to start with big errors then go to small
234
235
             y_vals = []
236
             # y_errs = []
237
             for val in sample_range:
238
                  uniform = Distribution(np.random.uniform, low=-1*val, high=val)
239
                  sub_test = []
240
241
                  for i in range(samps_per):
242
                      # print(normal, dims)
                      almost = construct_almost_basis(dim, uniform)
243
                      sub_test += [exp_error(exp, almost, diag_matrix)]
244
245
                  # Convert to numpy array so we can use nice easy stats without
246
                  # writing helper functions
247
                  sub_test = np.array(sub_test)
248
                  # Use the average value
250
                 y_vals += [np.mean(sub_test)]
251
252
                 # std_dev = np.std(sub_test)
253
                  # y_errs += [std_dev/samps_per]
254
255
             Z += [y_vals]
256
257
        Z = np.array(Z).T
258
         X, Y = np.meshgrid(exps, sample_range)
259
260
         # print(X.shape, Y.shape, Z.shape)
261
262
         fig = plt.figure()
         ax = plt.axes(projection='3d')
263
        # ax.contour3D(X, Y, Z, 50, cmap='binary')
ax.plot_surface(X, Y, np.log(Z+1), cmap='binary')
ax.set_xlabel('exponent')
264
265
266
         ax.set_ylabel('range of uniform perturbation (-y to y)')
267
         ax.set_zlabel('log normed (exponent error +1)')
268
         plt.show()
269
270
271
272
    def plot_uniform_3delbow(min_val=0.001, max_val=10,
273
                               num_tests=100, samps_per=5):
274
276
         dims = np.array(range(2,30))
         sample_range = np.linspace(min_val, max_val, num=num_tests)
277
278
         7. = 11
         for dim in tqdm(dims):
             # Reverse the order to start with big errors then go to small
280
281
282
             y_vals = []
             # y_errs = []
283
             for val in sample_range:
284
285
                  uniform = Distribution(np.random.uniform, low=-1*val, high=val)
286
                  sub_test = []
287
                  for i in range(samps_per):
288
                      # print(normal, dims)
289
                      almost = construct_almost_basis(dim, uniform)
                      sub_test += [orthog_error(almost)]
290
291
                  # Convert to numpy array so we can use nice easy stats without
292
                  # writing helper functions
293
                  sub_test = np.array(sub_test)
294
295
                  # Use the average value
296
                 y_vals += [np.mean(sub_test)]
297
298
                 # std_dev = np.std(sub_test)
299
                  # y_errs += [std_dev/samps_per]
300
301
             Z += [y_vals]
302
303
304
         Z = np.array(Z).T
         X, Y = np.meshgrid(dims, sample_range)
305
         # print(X.shape, Y.shape, Z.shape)
306
307
         fig = plt.figure()
308
        ax = plt.axes(projection='3d')
309
```

```
310
         ax.contour3D(X, Y, Z, 50, cmap='binary')
         ax.set_xlabel('dimension of matrix')
311
         ax.set_ylabel('range of normal perturbation (-y to y)')
312
         ax.set_zlabel('distance to closest orthogonal matrix')
313
314
         plt.show()
315
    def plot_beta_3delbow(dim=10, min_val=0.001, max_val=10,
316
                               num_tests=100, samps_per=10):
317
318
319
         a_range = np.linspace(min_val, max_val, num=num_tests)
320
         b_range = a_range.copy()
321
         7. = [7]
322
323
         for a in tqdm(a_range):
             # Reverse the order to start with big errors then go to small
324
325
             y_vals = []
326
             # y_errs = []
327
             for b in b_range:
328
                 beta = Distribution(np.random.beta, a, b)
329
                  sub_test = []
330
                  for i in range(samps_per):
331
                      # print(normal, dims)
332
                      almost = construct_almost_basis(dim, beta)
333
                      sub_test += [orthog_error(almost)]
334
335
336
                  # Convert to numpy array so we can use nice easy stats without
                  # writing helper functions
337
338
                  sub_test = np.array(sub_test)
339
                 # Use the average value
y_vals += [np.mean(sub_test)]
340
341
342
                 # std_dev = np.std(sub_test)
# y_errs += [std_dev/samps_per]
343
344
345
             Z += [y_vals]
346
347
         Z = np.array(Z).T
348
349
         X, Y = np.meshgrid(a_range, b_range)
         # print(X.shape, Y.shape, Z.shape)
350
351
         fig = plt.figure()
352
         ax = plt.axes(projection='3d')
353
         ax.plot_surface(X, Y, Z, cmap='viridis')
354
         ax.set_xlabel('alpha (beta distribution parameter)')
355
         ax.set_ylabel('beta (beta distribution parameter)')
356
         ax.set_zlabel('distance to closest orthogonal matrix')
357
358
         plt.show()
359
    def plot_gamma_3delbow(dim=10, min_val=0.001, max_val=10,
360
                               num_tests=100, samps_per=10):
361
362
363
         k_range = np.linspace(min_val, max_val, num=num_tests)
364
         theta_range = k_range.copy()
365
         \mathbf{Z} = []
366
         for k in tqdm(k_range):
367
             # Reverse the order to start with big errors then go to small
368
369
             y_vals = []
370
             # y_errs = []
371
             for t in theta_range:
372
                 gamma = Distribution(np.random.gamma, k, scale=t)
373
                  sub test = []
374
                  for i in range(samps_per):
375
376
                      # print(normal, dims)
                      almost = construct_almost_basis(dim, gamma)
377
                      sub_test += [orthog_error(almost)]
378
379
                 # Convert to numpy array so we can use nice easy stats without
380
                  # writing helper functions
381
382
                  sub_test = np.array(sub_test)
383
                 # Use the average value
384
385
                 y_vals += [np.mean(sub_test)]
386
387
                  # std_dev = np.std(sub_test)
```

```
# y_errs += [std_dev/samps_per]
388
389
             Z += [v vals]
390
391
        Z = np.array(Z).T
392
        X, Y = np.meshgrid(k_range, theta_range)
393
         # print(X.shape, Y.shape, Z.shape)
394
395
396
        fig = plt.figure()
        ax = plt.axes(projection='3d')
397
398
         ax.plot_surface(X, Y, Z, cmap='viridis')
        ax.set_xlabel('k (gamma distribution parameter)')
399
         ax.set_ylabel('theta (gamma distribution parameter)')
400
401
         ax.set_zlabel('distance to closest orthogonal matrix')
        plt.show()
402
403
404
405
    def frobenius_tester():
         dimensions = range(2,100)
406
        sigma = .1
407
        normal = Distribution(np.random.normal, loc=0.0, scale=sigma)
408
409
         tests_per_dim = 10000
410
411
        norms = []
412
413
414
        for dim in tqdm(dimensions):
             this_dim_norms = []
415
416
             for test in range(tests_per_dim):
                 this_dim_norms += [np.linalg.norm(normal(size=(dim,dim)))]
417
418
419
             norms += [np.mean(np.array(this_dim_norms)) - (dim*sigma)]
420
        print(np.mean(norms))
plt.plot(dimensions, norms)
421
422
423
         plt.xlabel("Dimension of xi matrix")
        plt.ylabel("Error ||xi||_F relative to (dim * sigma)")
424
         plt.title("Error (||xi||_F - dim * sigma); 10000 tests per dim, dim in" +
425
          [2..200]; sigma = 0.1")
426
427
         plt.show()
428
429
    def frobenius_tester():
430
         dimensions = range(2,100)
431
432
         sigma = .1
        normal = Distribution(np.random.normal, loc=0.0, scale=sigma)
433
434
        tests_per_dim = 10000
435
436
        norms = []
437
438
        for dim in tqdm(dimensions):
439
             this_dim_norms = []
440
441
             for test in range(tests_per_dim):
                 this_dim_norms += [np.linalg.norm(normal(size=(dim,dim)))]
442
443
             norms += [np.mean(np.array(this_dim_norms)) - (dim*sigma)]
444
445
446
        print(np.mean(norms))
         plt.plot(dimensions, norms)
447
         plt xlabel("Dimension of xi matrix")
448
        plt.ylabel("Error ||xi||_F relative to (dim * sigma)")
449
         plt.title("Error (||xi||_F - dim * sigma); 10000 tests per dim, dim in" +
450
          [2..200]; sigma = 0.1")
451
         plt.show()
452
453
    def delta(n, A, xi, diag_matrix):
454
455
456
        n: exponent to raise to
457
        Return the error as propogated through when taking the matrix
458
         exponential and treating the input
459
460
        M = A + xi
461
        true_matrix = np.linalg.matrix_power(M @ diag_matrix @ M.T, n)
462
463
         est_matrix = A @ (diag_matrix ** n) @ A.T
         error = np.linalg.norm(true_matrix - est_matrix)
464
        return error
465
```

```
466
    def Delta(n, A, xi, diag_matrix):
467
468
         n: exponent to raise to
469
470
         Return the error as propogated through when taking the matrix
471
         exponential and treating the input
472
473
         M = A + xi
474
475
         true_matrix = np.linalg.matrix_power(M @ diag_matrix @ M.T, n)
         est_matrix = M @ (diag_matrix ** n) @ M.T
error = np.linalg.norm(true_matrix - est_matrix)
476
477
         return error
478
479
480
481
    def get_rms(W, X, y):
             gets the root mean square error """
482
483
         return np.linalg.norm(y - D@W, ord=2)/np.sqrt(D.shape[0])
484
    def lin_reg(X, y):
485
         """ Takes two numpy arrays as inputs, and calculates the coefficients minimizing the least-square error between them """
486
487
         # Solve for optimal weight parameters minimizing regression
488
         X = np.column_stack((np.ones_like(X), X))
489
         W_opt = np.linalg.solve(X.T @ X, X.T @ y)
490
         return W_opt
491
492
    def gen_rand_orthog(dim):
493
         rand_mat = np.random.rand(dim,dim)
q, r = np.linalg.qr(rand_mat)
494
495
         return q
496
497
498
    def test_theorem_4(sigma=.2, plot=True):
499
         dim = 20
500
501
         # dimensions = np.array(range(2,30))
         exps = range(2,50)
503
504
505
         normal = Distribution(np.random.normal, loc=0.0, scale=sigma)
506
507
         tests_per_exp = 100
508
         deltas = []
509
         Deltas = []
510
511
512
         A = gen_rand_orthog(dim)
         xi = normal(size=(dim,dim))
514
         for exp in exps:
515
         # for exp in tqdm(exps):
              this_dim_deltas = []
this_dim_Deltas = []
518
              for test in range(tests_per_exp):
520
                   test_mat = np.diag(np.random.normal(loc=1.0, scale=2,
                                                            size=(dim,)))
521
                  this_dim_deltas += [delta(exp, A, xi, test_mat)]
this_dim_Deltas += [Delta(exp, A, xi, test_mat)]
523
524
525
              deltas += [np.mean(np.array(this_dim_deltas))]
              Deltas += [np.mean(np.array(this_dim_Deltas))]
526
527
         deltas = np.array(deltas)
528
         Deltas = np.array(Deltas)
530
         logdeltas = np.log(deltas)
logDeltas = np.log(Deltas)
531
532
534
         exps = np.array(exps)
535
         bd, md = lin_reg(exps, logdeltas)
536
         bD, mD = lin_reg(exps, logDeltas)
537
538
             540
541
542
              plt.ylabel("log delta(n)")
543
```

```
plt.show()
544
545
             plt.plot(exps, np.log(np.array(Deltas)))
plt.plot(exps, logDeltas, "bo")
546
547
             plt.xlabel("n, exponent")
548
             plt.ylabel("log {d,D}elta(n)")
549
             plt.show()
551
             plt.plot(exps, np.log(np.log(deltas/Deltas)))
             plt.xlabel("n, exponent")
plt.ylabel("diff of log deltas (n)")
553
554
             plt.show()
556
557
         return (bd, md, bD, mD)
558
    def test_ms():
         plt.clf()
560
561
         num\_tests = 50
562
         min_sigma = 0.001
         max_sigma = 3
563
         sigmas = np.linspace(min_sigma, max_sigma, num=num_tests)
564
565
566
         bds = []
         mds = []
567
         bDs = []
568
         mDs = []
569
         for sigma in tqdm(sigmas):
571
572
             bd, md, bD, mD = test_theorem_4(sigma=sigma, plot=False)
             bds += [bd]
574
575
             mds += [md]
576
             bDs += [bD]
             mDs += [md]
577
578
         bds = np.array(bds)
         mds = np.array(mds)
580
         bDs = np.array(bDs)
581
         mDs = np.array(mDs)
582
583
         print(sigmas.shape, bds.shape, mds.shape, bDs.shape, mDs.shape)
584
585
        586
587
588
         plt.ylabel("slope of best linear fit to log deltas")
589
         plt.xlabel("sigma")
590
         # plt.savefig("test.png")
591
         plt.show()
592
593
594
    def closest_orthog(A):
         sqrt = scipy.linalg.sqrtm(A.T @ A)
invsqrt = np.linalg.inv(sqrt)
595
596
         R = \bar{A} \otimes invsqrt
597
598
         return A - R
599
    def test_orthog_trace():
600
601
         numtests = range(100)
602
603
         dims = range(2,50)
604
605
         traces = []
         # trace_errs = []
606
607
         traces2 = []
608
         # trace_errs2 = []
609
610
         for dim in tqdm(dims):
611
             this_dim_traces = []
612
             this_dim_traces2 = []
613
614
             for test in numtests:
615
                  rand_mat = 2*np.random.randn(dim,dim) + 1
616
                  this_dim_traces += [np.trace(closest_orthog(rand_mat))]
617
618
619
                  Q, _ = np.linalg.qr(rand_mat)
                  this_dim_traces2 += [np.trace(Q)]
620
621
```

```
this_dim_traces = np.array(this_dim_traces)
this_dim_traces2 = np.array(this_dim_traces2)
622
623
624
              # sigma = np.std(this_dim_traces)
625
              mu = np.mean(this_dim_traces)
626
              traces += [mu]
627
              # trace_errs += [sigma]
628
629
              # sigma2 = np.std(this_dim_traces2)
630
              mu2 = np.mean(this_dim_traces2)
traces2 += [mu2]
631
632
              # trace_errs2 += [sigma2]
633
634
635
         plt.plot(dims, traces, "b", dims, traces2, "r")
636
         plt.show()
637
638
639
     def plot_beta_3delbow(dim=10, min_val=0.001, max_val=10,
640
                                 num_tests=100, samps_per=10):
641
642
         a_range = np.linspace(min_val, max_val, num=num_tests)
643
644
         b_range = a_range.copy()
645
646
         for a in tqdm(a_range):
647
              # Reverse the order to start with big errors then go to small
648
649
650
              y_vals = []
              # y_errs = []
651
              for b in b_range:
652
653
                   beta = Distribution(np.random.beta, a, b)
654
                   sub_test = []
                   for i in range(samps_per):
655
                       # print(normal, dims)
656
657
                       almost = construct_almost_basis(dim, beta)
                       sub_test += [orthog_error(almost)]
658
659
                  # Convert to numpy array so we can use nice easy stats without
660
                  # writing helper functions
661
                   sub_test = np.array(sub_test)
662
663
                   # Use the average value
664
                  y_vals += [np.mean(sub_test)]
665
666
                   # std_dev = np.std(sub_test)
667
                   # y_errs += [std_dev/samps_per]
668
669
              Z += [y_vals]
670
671
672
         Z = np.array(Z).T
         X, Y = np.meshgrid(a_range, b_range)
# print(X.shape, Y.shape, Z.shape)
673
674
675
676
         fig = plt.figure()
         ax = plt.axes(projection='3d')
677
         ax.plot_surface(X, Y, Z, cmap='viridis')
ax.set_xlabel('alpha (beta distribution parameter)')
678
679
         ax.set_ylabel('beta (beta distribution parameter)')
680
         ax.set_zlabel('distance to closest orthogonal matrix')
681
         plt.show()
682
```

almost.py