

## Reference

Friday, April 26, 2024 1:21 PM

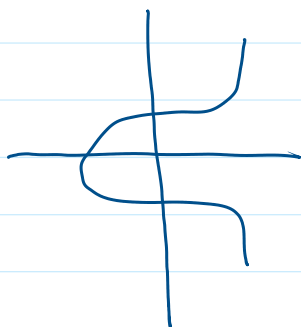
1. <https://andrea.corbellini.name/2015/05/17/elliptic-curve-cryptography-a-gentle-introduction/>
2. <https://andrea.corbellini.name/2015/05/23/elliptic-curve-cryptography-finite-fields-and-discrete-logarithms/>
3. <https://andrea.corbellini.name/2015/05/30/elliptic-curve-cryptography-ecdh-and-ecdsa/>
4. <https://www.rarekills.io/post/elliptic-curve-addition>
5. <https://www.rarekills.io/post/elliptic-curves-finite-fields>
6. <https://www.rarekills.io/post/bilinear-pairing>

Def: The set of points described by

$$y^2 = x^3 + ax + b, \text{ where } 4a^3 + 27b^2 \neq 0$$

$\Delta$ : discriminant

Weierstrass Curves



Symmetric about x-axis

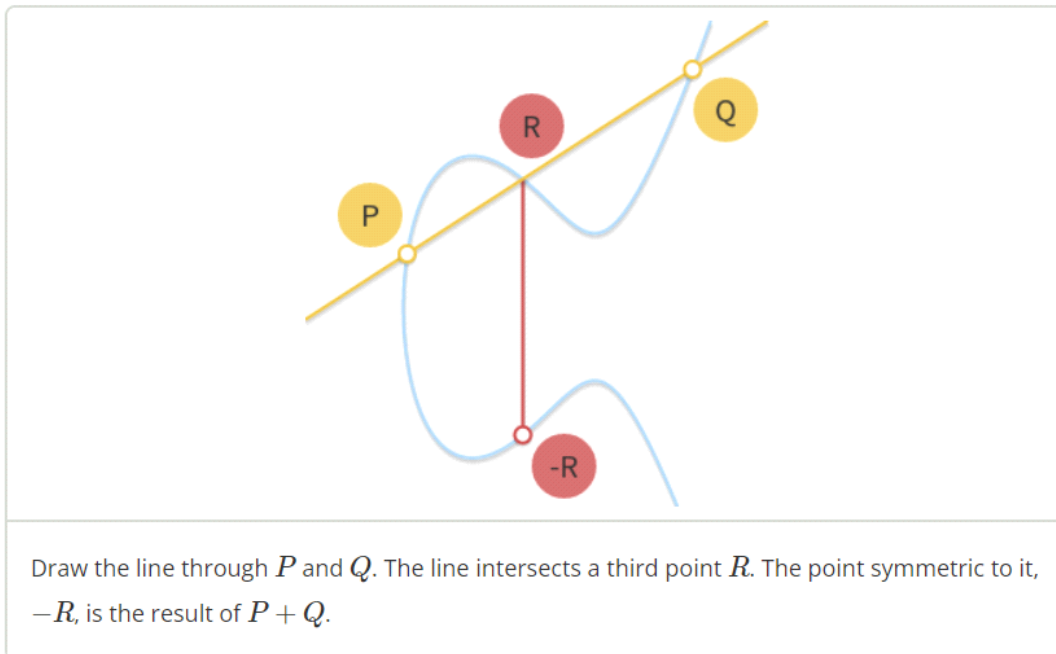
↑  
exclude singular curves  
(double or triple root  $x_0$   
and  $(x_0, 0)$  is a singular  
point)

↑  
hard to define addition

Def: Point at infinity  $O$  (acts as identity in group structure)

## Group Law for EC

- Elements in this group is just points on EC
- Identity is  $O$
- Inverse of point  $P$  is  $-P$  (symmetric about x-axis)
- Addition:



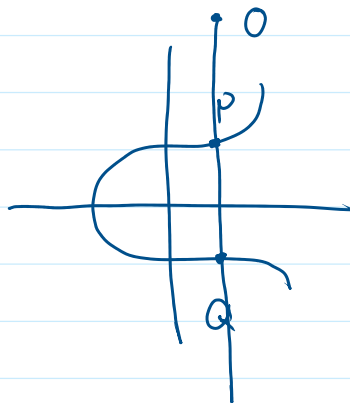
Claim: Group of EC is abelian group.

Intuition:  $P + Q = Q + P$  cause they share the same straight line.

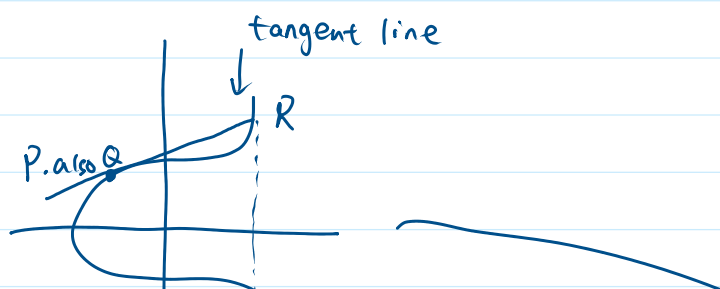
Special cases:

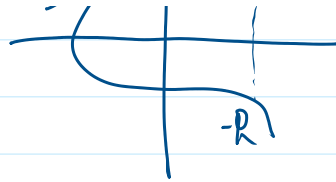
1.  $P = 0$  or  $Q = 0 \rightarrow$  identity

2.  $P = -Q$



3.  $P = Q$





- As long as we do not pick a perfectly vertical line, if we intersect two points in an elliptic curve, then we will also intersect a 3rd point on the elliptic curve.
- If a straight line crosses an elliptic curve at exactly two points, then it must be perfectly vertical.

$$P + P = \underbrace{2P}$$

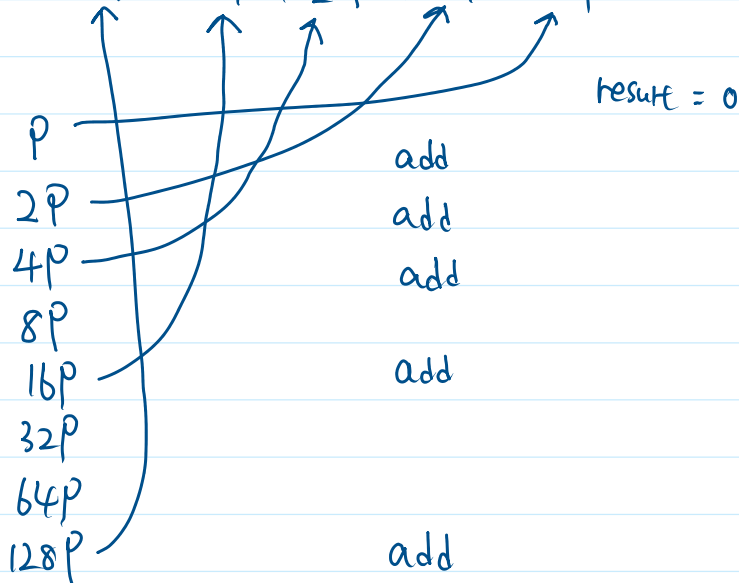
Scalar multiplication

$$nP = \underbrace{P + P + P + \dots + P}_{n \text{ times}}$$

$nP$  can be computed efficiently using "double and add" algorithm:

Example:  $n = 151$

$$151P = 2^7P + 2^4P + 2^2P + 2^1P + 2^0P$$

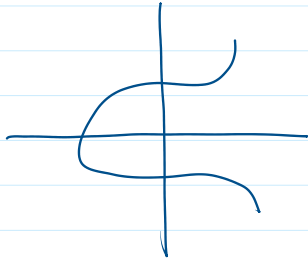


$8 + 5 = 13$  computations  $O(\log n)$   
↑    ↑  
double    add

Remember, EC over  $R$  is just an abelian group. Group theory simplifies EC.

EC over  $\mathbb{R}$ :

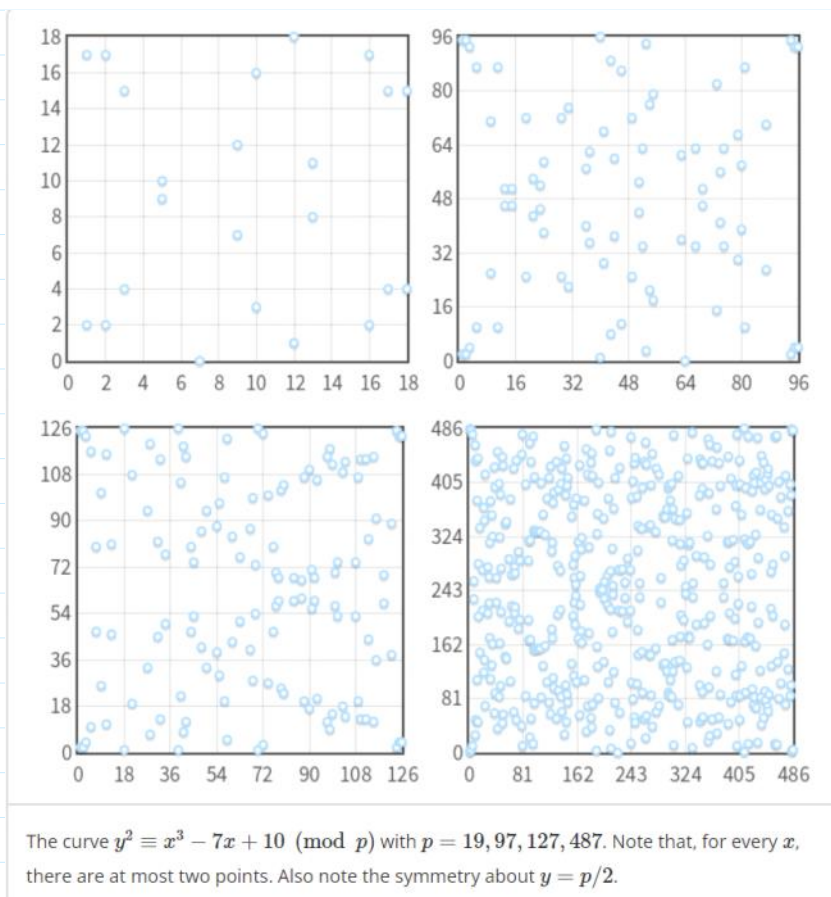
$$\{(x, y) \in \mathbb{R}^2 \mid y^2 = x^3 + ax + b, 4a^3 + 27b^2 \neq 0\} \cup \{0\}$$



EC over  $\mathbb{F}_p$ :

$$\{(x, y) \in \mathbb{F}_p^2 \mid y^2 = x^3 + ax + b \pmod{p},$$

$$4a^3 + 27b^2 \not\equiv 0 \pmod{p}\} \cup \{0\}$$



Claim: EC over  $\mathbb{F}_p$  is abelian group

Even better, EC over  $\mathbb{F}_p$  is cyclic group !!!

← recall that any  $G$  with  $|G| = p$  is cyclic

Even better, EC over  $\mathbb{F}_p$  is cyclic group !!!

← recall that any  $G$  with  $|G|=p$  is cyclic

→ generator  $G$ , operator is +

( $G \neq 0$ )

→ every non-identity element is generator

→ points can be written as  $nG$

→  $\{nG\}$  covers every points on EC

**Corollary 6.12.** Let  $|G| = p$  with  $p$  a prime number. Then  $G$  is cyclic and any  $g \in G$  such that  $g \neq e$  is a generator.

*Proof.*

Corollary 6.12 suggests that groups of prime order  $p$  must somehow look like  $\mathbb{Z}_p$ .

source:

<https://abstract.ups.edu/abstract.ups.edu/aata/cosets-section-lagranges-theorem.html>

bits of security

Ethereum precompiles use bn128 (bn254)

<https://hackmd.io/@jpw/bn254>

↖ #bits of prime  $p$

field\_modulus =

21888242871839275222246405745257275088696311157297823662689037894645226208583

$y^2 = x^3 + 3 \pmod{\text{field\_modulus}}$  ← prime  $p$

( $y^2 = x^3 + b$  is BN-curve equation)

(curve\_order can be computed by Schoof's algorithm: [https://en.wikipedia.org/wiki/Schoof%27s\\_algorithm](https://en.wikipedia.org/wiki/Schoof%27s_algorithm))

```
>>> from py_ecc.bn128 import G1, multiply, add, eq, neg
>>> G1
(1, 2)
>>> add(G1, G1)
(136801517948995470139040035907857969304351944733113978918064868415326638035, 9918110051302171585080402603319702774565515993150576347155970296011118125764)
>>> add(G1, G1) == multiply(G1, 2)
True
```

field\_modulus != curve\_order

```
>>> from py_ecc.bn128 import curve_order, field_modulus, G1, multiply, eq
>>> x = 5
>>> multiply(G1, x) == multiply(G1, x + curve_order)
True
>>> multiply(G1, x) == multiply(G1, x + field_modulus)
False
>>>
```

Can encode rational numbers :  $\frac{2}{3} = 2 \cdot 3^{-1}$ ,  $\frac{4}{3} = 4 \cdot 3^{-1}$ ,  $\frac{2}{3} + \frac{4}{3} = 2$

↑  
always exists in a field

```
>>> two_over_three = (2 * pow(3, -1, curve_order)) % curve_order
>>> four_over_three = (4 * pow(3, -1, curve_order)) % curve_order
>>> add(multiply(G1, two_over_three), multiply(G1, four_over_three)) == multiply(G1, 2)
True
>>>
```

ECDLP: Given  $Q = nP$ , can't solve for  $n$

Recall that DLP says given  $h = g^x \bmod p$  it is hard to solve for  $x$ . This is the foundation for Diffie-Hellman. Similarly, ECDLP is the foundation for ECDH:

<u>Alice</u>	Generator $G$ of EC	<u>Bob</u>
Generate random $sk_A$		--- $sk_B$
Compute $PK_A = \underbrace{sk_A \cdot G}$		--- $PK_B = \underbrace{sk_B \cdot G}$
$\xrightarrow{\text{exchange } PK}$		

Compute  $sk_A \cdot PK_B = sk_A \cdot sk_B \cdot G$       ---  $sk_B \cdot PK_A = sk_B \cdot sk_A \cdot G$

Correctness:  $(sk_A \cdot sk_B) \cdot G = (sk_B \cdot sk_A) \cdot G$

Attack: <https://www.ret2basic.me/2024/04/12/elliptic-curve-attacks-small-subgroup.html>



Example 1: Prover: I know  $x, y$  s.t.  $x+y=15$

Verifier: prove it, but don't send me  $x$  and  $y$

Prover computes  $xG$  and  $yG$ , send to verifier

Verifier checks if  $xG + yG = 15G$

Example: Prover: I know  $x$  s.t.  $23x = 161$

Verifier: prove it, but don't send me  $x$

Prover computes  $161 \cdot 23^{-1} G$ , send to verifier

verifier checks if  $23 \cdot (161 \cdot 23^{-1} G) = 23 \cdot 23^{-1} \cdot 161 G = 161 G$

In reality ZK property is achieved using bilinear pairing ("multiply 2 EC points from different dimensions") We will cover that next week.

## Further study

Thursday, May 16, 2024 6:26 PM

Read py\_ecc source code and play with it [https://github.com/ethereum/py\\_ecc](https://github.com/ethereum/py_ecc)