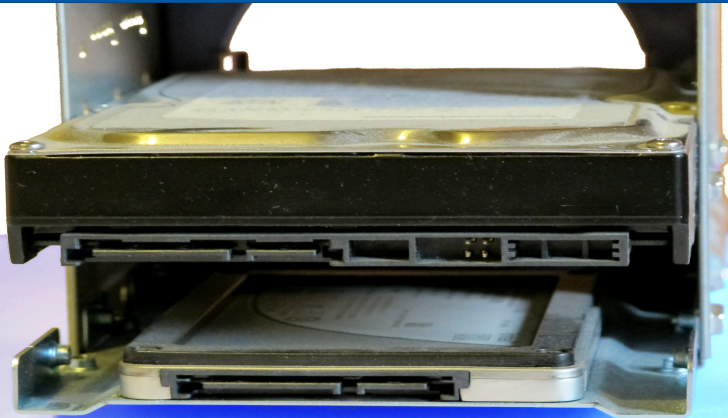


Making Cost-Based Query Optimization Asymmetry-Aware

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Asymmetry in new storage devices



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- Writing to Flash memory is slower than reading from it
- This also applies to emerging non-volatile memories (PCM, etc.)
- Small writes to random locations on Flash are even more slow
- Random reads from Flash are only $\frac{1}{3}$ slower than sequential reads¹

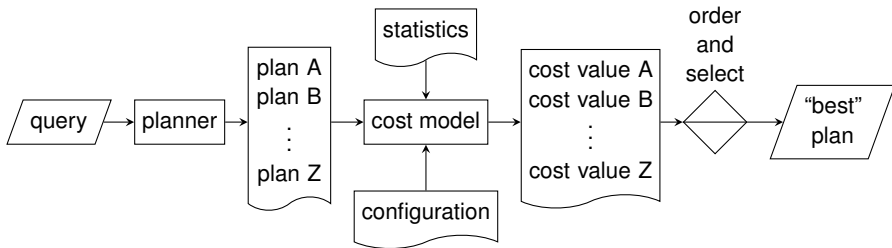
¹on Intel X25-E using full command queue

Cost-Based Query Optimization

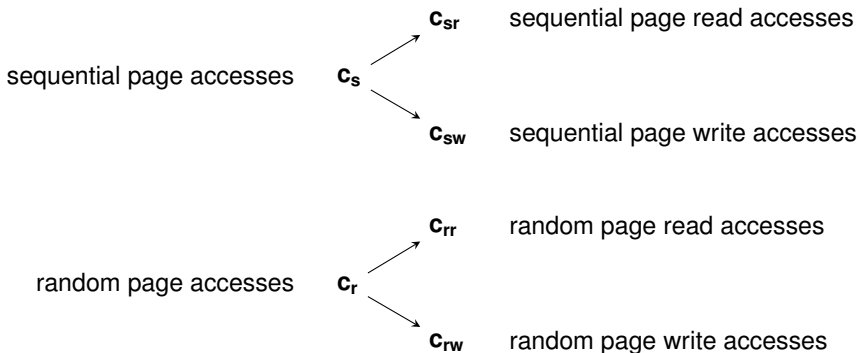


- estimation of run-time before real execution (e.g. in PostgreSQL)
- model comprised of functions like

$$c(\text{seqscan}) = \underbrace{c_s \|R\|_p}_{\text{I/O cost}} + \underbrace{q_{R,0} + (\dot{c}_{\text{cpu}} + \dot{q}_R) \|R\|_t}_{\text{CPU cost}}$$



Splitting Parameters



Cost functions for “pure load” algorithms



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cost function	kind	original	replacement
sequential scan	read only	C_s	C_{sr}
index scan	read only	C_s	C_{sr}
		C_r	C_{rr}
bitmap scan	read only	C_s	C_{sr}
		C_r	C_{rr}
TID scan	read only	C_r	$C_{rr} \cdot e$
materialization	write only	C_s	C_{sw}
re-scan	read only	C_s	C_{sr}

Cost function of sort algorithm



$$c_{io}(sort) = \overbrace{2 \|S\|_p \lceil \log_m n \rceil \left(\frac{3}{4} \mathbf{c}_s + \frac{1}{4} \mathbf{c}_r \right)}^{\text{startup}}$$

blktrace stats	write		read	
	s	r	s	r
external sort of unordered data				
external sort of ordered data				
sort-merge join				

$$c_{io}(sort)_{rw} = \underbrace{\|S\|_p \left(\mathbf{c}_{sw} + (\lceil \log_m n \rceil - 1) \frac{\mathbf{c}_{sw} + \mathbf{c}_{rw}}{2} + \lceil \log_m n \rceil \frac{\mathbf{c}_{sr} + \mathbf{c}_{rr}}{2} \right)}_{\text{startup}}$$

Cost function of hash join

$$c_{io}(\text{hashjoin}) = \overbrace{\|P_i\|_p \mathbf{c}_s}^{\text{startup}} + \left(\|P_i\|_p + 2 \|P_o\|_p \right) \mathbf{c}_s$$

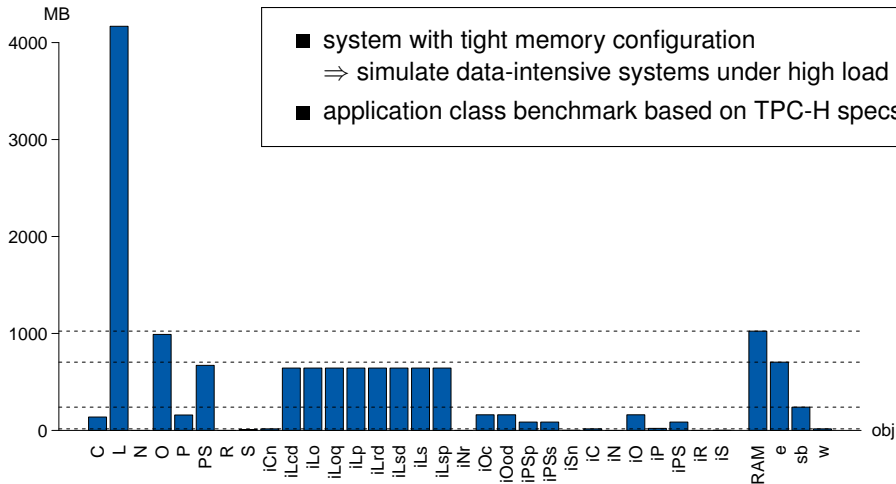
blktrace stats	write		read	
	s	r	s	r
hash join				

$$c_{io}(\text{hashjoin})_{rw} = \underbrace{\|P_i\|_p \mathbf{c}_{rw}}_{\text{startup}} + \|P_i\|_p \mathbf{c}_{sr} + \|P_o\|_p (\mathbf{c}_{rw} + \mathbf{c}_{sr})$$

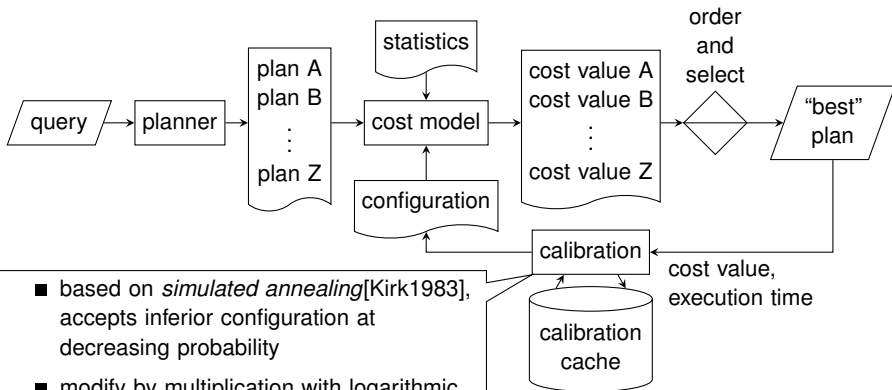
System and Load



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Calibration



- based on *simulated annealing*[Kirk1983], accepts inferior configuration at decreasing probability
- modify by multiplication with logarithmic normally distributed random variable
- cooling cycle of 100 iterations, restarted 100 times

Found “Optimal” Settings



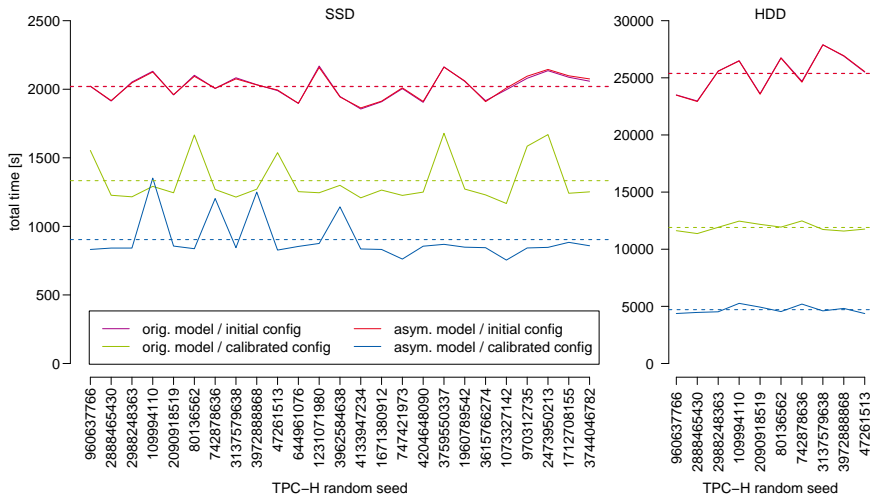
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SSD		HDD	
original model	asymmetric model	original model	asymmetric model
$\mathbf{c_s} = 1.00000$	$\mathbf{c_{sr}} = 1.00000$	$\mathbf{c_s} = 1.00000$	$\mathbf{c_{sr}} = 1.00000$
	$\mathbf{c_{sw}} = 49.91840$		$\mathbf{c_{sw}} = 110.21139$
	$\mathbf{c_{rr}} = 5.62724$		$\mathbf{c_{rr}} = 19.25494$
$\mathbf{c_r} = 6.77405$	$\mathbf{c_{rw}} = 19.08421$	$\mathbf{c_r} = 29.04790$	$\mathbf{c_{rw}} = 20.18467$
$\dot{\mathbf{c}}_{\text{cpu}} = 0.00121$	$\dot{\mathbf{c}}_{\text{cpu}} = 0.00003$	$\dot{\mathbf{c}}_{\text{cpu}} = 0.00280$	$\dot{\mathbf{c}}_{\text{cpu}} = 0.00082$
$\hat{\mathbf{c}}_{\text{cpu}} = 0.03658$	$\hat{\mathbf{c}}_{\text{cpu}} = 0.01608$	$\hat{\mathbf{c}}_{\text{cpu}} = 0.03718$	$\hat{\mathbf{c}}_{\text{cpu}} = 0.00045$
$\mathbf{c_{op}} = 0.00016$	$\mathbf{c_{op}} = 0.00008$	$\mathbf{c_{op}} = 0.00004$	$\mathbf{c_{op}} = 0.00119$

Comparison



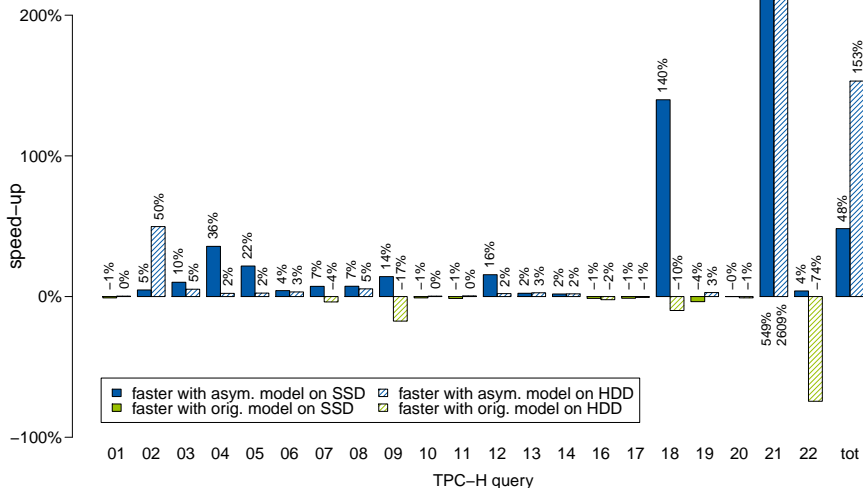
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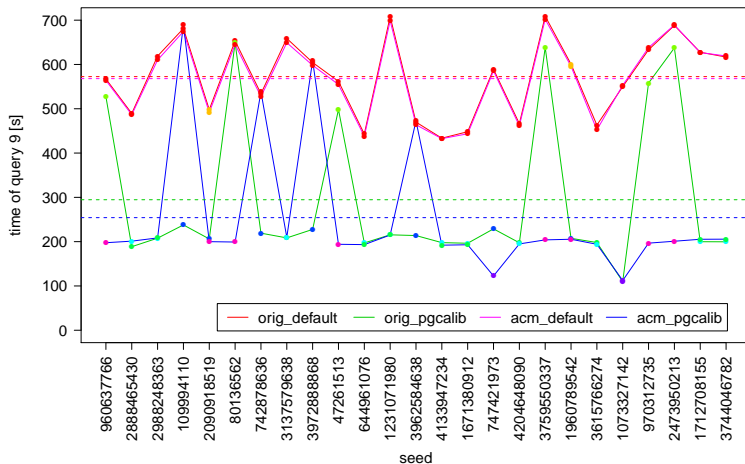
Individual Query Speed-Up



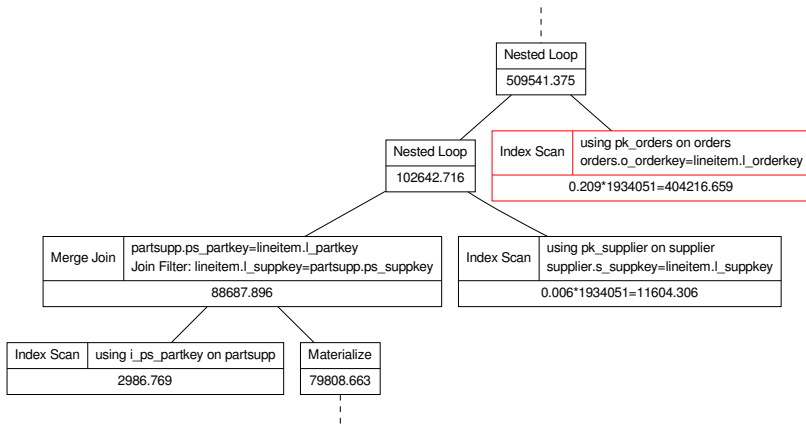
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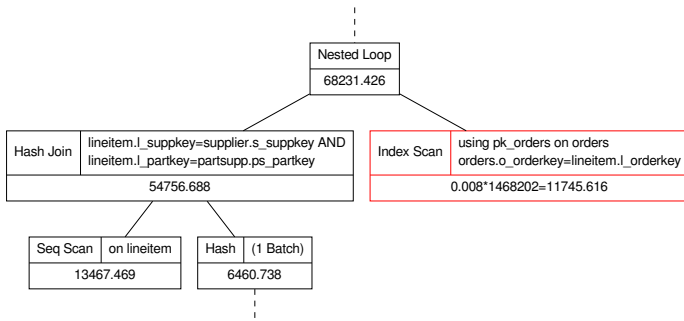
Discussion – Cause for Peaks \Rightarrow Query 9



Discussion – Query 9 – the slow plan



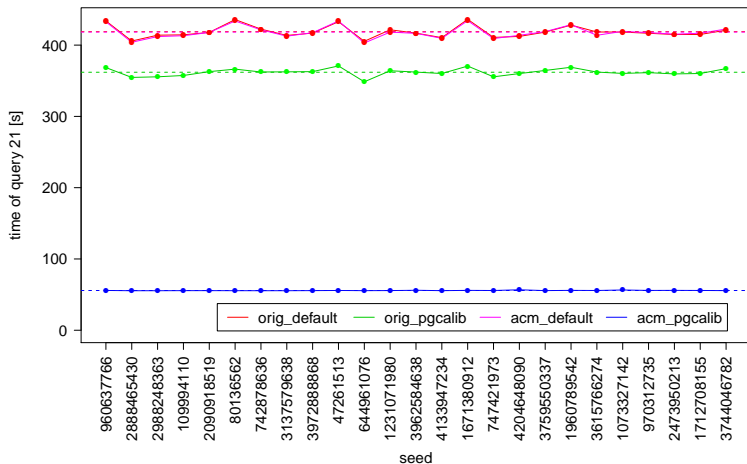
Discussion – Query 9 – the fast plan



Discussion – Biggest Improvement \Rightarrow Query 21



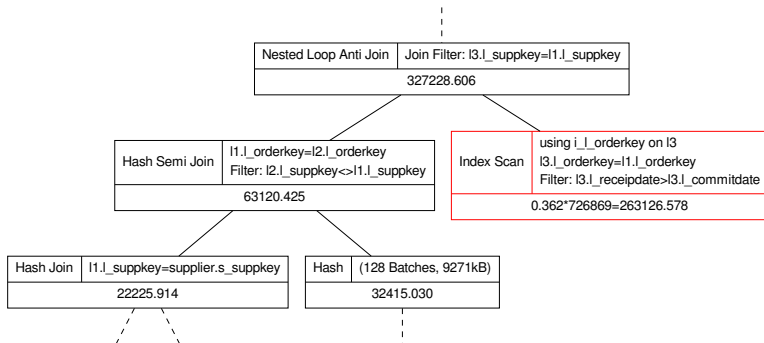
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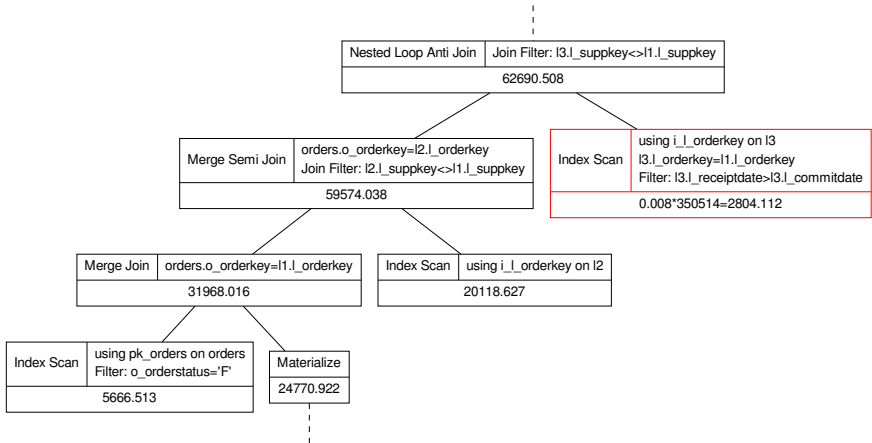
Discussion – Query 21 – the slow plan



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Discussion – Query 21 – the fast plan



- Derived an asymmetry-aware model from facts and observable behavior
 - Storage properties
 - Algorithm access behavior
- Comparison shows improved performance on application-class benchmark
 - Average speed-up at 48%
 - Individual query speed-up by up to 549%
- Calibration is attracted by strong effects
 - TPC-H features properties not respected in PostgreSQL's optimizer
 - Available degrees of freedom may get abused to compensate deficiencies
- Additional experiments are required to decide upon original hypothesis
 - Using a special load (i.e. a custom micro-benchmark)
 - Focussing on optimization problems in which asymmetry matters explicitly
- Calibration may be a useful tool find optimal configurations in general
 - Typical queries need to be combined in a repeatable benchmark
 - Can cope with any number of variables

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