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```
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```

10 lines

Contest (1)

```
template.cpp
#include <bits/stdc++.h>
#define all(x) begin(x), end(x)
using namespace std;
using 11 = long long;
int main() {
 cin.tie(0)->sync_with_stdio(0);
 cin.exceptions(cin.failbit);
```

Data structures (2)

BIT.h

Description: Query [l, r] sums, and point updates. kth() returns the smallest index i s.t. query(0, i) >= k**Time:** $\mathcal{O}(\log n)$ for all ops. 33f78c, 22 lines

template <typename T> struct BIT { vector<T> s; int n: BIT(int n) : s(n + 1), n(n) {} void update(int i, T v) { for (i++; i <= n; i += i & -i) s[i] += v; T query(int i) { T ans = 0; for (i++; i; i -= i & -i) ans += s[i]; return ans; T query(int 1, int r) { return query(r) - query(l - 1); } int kth(T k) { // returns n if k > sum of tree **if** (k <= 0) **return** -1; int i = 0;for (int pw = 1 << __lg(n); pw; pw >>= 1) **if** $(i + pw \le n \&\& s[i + pw] \le k) k -= s[i += pw];$ return i: };

KDBIT.h

```
Description: k-dimensional BIT. BIT<int, N, M> gives an N \times M BIT.
Query: bit.query(x1, x2, y1, y2) Update: bit.update(x, y, delta)
Time: O(\log^k n) Status: Tested
template <class T, int... Ns>
struct BIT {
 T val = 0;
  void update(T v) { val += v; }
 T query() { return val; }
template <class T, int N, int... Ns>
struct BIT<T, N, Ns...> {
 BIT<T, Ns...> bit[N + 1];
  // map < int, BIT < T, Ns... >> bit;
  // if the memory use is too high
  template <class... Args>
  void update(int i, Args... args) {
    for (i++; i <= N; i += i & -i) bit[i].update(args...);</pre>
  template <class... Args>
  T query(int i, Args... args) {
    for (i++; i; i -= i & -i) ans += bit[i].query(args...);
    return ans:
  template <class... Args,
            enable_if_t<(sizeof...(Args) ==
                          2 * sizeof...(Ns))>* = nullptr>
  T query (int 1, int r, Args... args) {
    return query(r, args...) - query(l - 1, args...);
};
DSU.h
Description: Maintains union of disjoint sets
Time: \mathcal{O}(\alpha(N))
                                                       c22586, 14 lines
struct DSU {
  vector<int> s:
  DSU(int n) : s(n, -1) {}
  int find(int i) { return s[i] < 0 ? i : s[i] = find(s[i]); }</pre>
  bool join(int a, int b) {
    a = find(a), b = find(b);
    if (a == b) return false;
    if (s[a] > s[b]) swap(a, b);
    s[a] += s[b], s[b] = a;
    return true;
  int size(int i) { return -s[find(i)]; }
  bool same(int a, int b) { return find(a) == find(b); }
RMQ.h
Description: Constant time subarray min/max queries for a fixed array
Time: O(nlogn) initialization and O(1) queries. Status: Tested 536eac, 15 lines
template <typename T, class Compare = less<T>>
struct RMQ {
  vector<vector<T>> t;
  Compare cmp;
  RMQ(vector<T>& a) : t(\underline{lg(a.size())} + 1, a) {
    int n = a.size(), lg = __lg(n);
    for (int k = 1, len = 1; k <= lq; k++, len <<= 1)</pre>
      for (int i = 0; i + 2 * len - 1 < n; i++)
        t[k][i] = min(t[k-1][i], t[k-1][i+len], cmp);
 T query(int a, int b) {
    int k = __lg(b - a + 1), len = 1 << k;</pre>
    return min(t[k][a], t[k][b - len + 1], cmp);
```

Splay.h

else

Description: An implicit balanced BST. You only need to change update () and prop().

 $O(\log n)$ for all operations

```
If used for link-cut tree, code everything up to splay(). Time: amortized
struct node {
 node *ch[2] = \{0\}, *p = 0;
 int cnt = 1, val;
 node (int val, node * 1 = 0, node * r = 0)
   : ch{1, r}, val(val) {}
int cnt(node* x) { return x ? x->cnt : 0; }
int dir(node* p, node* x) { return p && p->ch[0] != x; }
void setLink(node* p, node* x, int d) {
 if (p) p \rightarrow ch[d] = x;
 if (x) x->p = p;
node* update(node* x) {
 if (!x) return 0;
 x - cnt = 1 + cnt(x - ch[0]) + cnt(x - ch[1]);
  setLink(x, x->ch[0], 0);
 setLink(x, x->ch[1], 1);
 return x;
void prop(node* x) {
 if (!x) return;
 // update(x); // needed if prop() can change subtree sizes
void rotate(node* x, int d) {
 if (!x || !x->ch[d]) return;
 node *y = x - > ch[d], *z = x - > p;
  setLink(x, y->ch[d ^ 1], d);
  setLink(y, x, d^{1});
 setLink(z, y, dir(z, x));
 update(x);
 update(y);
node* splay(node* x) {
  while (x && x->p) {
    node *y = x->p, *z = y->p;
    // prop(z), prop(y), prop(x); // needed for LCT
    int dy = dir(y, x), dz = dir(z, y);
    if (!z)
      rotate(y, dy);
    else if (dy == dz)
      rotate(z, dz), rotate(y, dy);
      rotate(y, dy), rotate(z, dz);
  return x;
// the returned node becomes the new root, update the root
// pointer!
node* nodeAt(node* x, int pos) {
 if (!x) return 0;
 while (prop(x), cnt(x->ch[0]) != pos)
    if (pos < cnt(x->ch[0]))
     x = x -> ch[0];
```

```
pos -= cnt(x->ch[0]) + 1, x = x->ch[1];
  return splay(x);
node* merge(node* 1, node* r) {
 if (!1 || !r) return 1 ?: r;
 1 = nodeAt(1, cnt(1) - 1);
  setLink(l, r, 1);
  return update(1);
// first is everything < pos, second is >= pos
pair<node*, node*> split(node* t, int pos) {
  if (pos <= 0 || !t) return {0, t};</pre>
  if (pos > cnt(t)) return {t, 0};
  node *1 = nodeAt(t, pos - 1), *r = 1->ch[1];
  if (r) 1 \rightarrow ch[1] = r \rightarrow p = 0;
  return {update(1), update(r)};
// insert a new node between pos-1 and pos
node* insert(node* t, int pos, int val) {
  auto [1, r] = split(t, pos);
  return update(new node(val, 1, r));
// apply lambda to all nodes in an inorder traversal
template <class F>
void each(node* x, F f) {
 if (x) \text{ prop}(x), \text{ each}(x->\text{ch}[0], f), f(x), \text{ each}(x->\text{ch}[1], f);
```

Geometry (3)

Graphs (4)

SCCTarjan.h

Description: Finds strongly connected components of a directed graph. Visits/indexes SCCs in reverse topological order.

Usage: scc(graph) returns an array that has the ID of each node's SCC. scc(graph, [&](vector<int>& v) { ... }) calls the lambda on each SCC, and returns the same array.

```
Time: \mathcal{O}(|V| + |E|)
                                                       358d18, 37 lines
namespace SCCTarjan {
 vector<int> val, comp, z, cont;
  int Time, ncomps;
  template <class G, class F>
  int dfs(int j, G& g, F& f) {
    int low = val[j] = ++Time, x;
    z.push_back(j);
    for (auto e : q[j])
     if (comp[e] < 0) low = min(low, val[e] ?: dfs(e, g, f));</pre>
    if (low == val[j]) {
     do {
       x = z.back();
        z.pop back();
        comp[x] = ncomps;
        cont.push_back(x);
      } while (x != j);
      f(cont);
      cont.clear();
     ncomps++;
    return val[j] = low;
  template <class G, class F>
```

```
vector<int> scc(G& g, F f) {
    int n = q.size();
   val.assign(n, 0);
    comp.assign(n, -1);
   Time = ncomps = 0;
   for (int i = 0; i < n; i++)</pre>
     if (comp[i] < 0) dfs(i, g, f);</pre>
   return comp;
 template <class G> // convenience function w/o lambda
 vector<int> scc(G& g) {
   return scc(g, [](auto& v) {});
} // namespace SCCTarjan
SCCKosaraju.h
Description: Finds strongly connected components of a directed graph.
Visits/indexes SCCs in topological order.
Usage: scc(graph) returns an array that has the ID
of each node's SCC.
Time: \mathcal{O}\left(|V| + |E|\right)
                                                      9b78e7, 29 lines
namespace SCCKosaraju {
 vector<vector<int>> adj, radj;
 vector<int> todo, comp;
 vector<bool> vis;
 void dfs1(int x) {
   vis[x] = 1;
    for (int y : adj[x])
      if (!vis[y]) dfs1(y);
    todo.push back(x);
 void dfs2(int x, int i) {
    comp[x] = i;
    for (int y : radj[x])
      if (comp[v] == -1) dfs2(v, i);
 vector<int> scc(vector<vector<int>>& _adj) {
   adi = adi;
    int time = 0, n = adj.size();
    comp.resize(n, -1), radj.resize(n), vis.resize(n);
    for (int x = 0; x < n; x++)
     for (int y : adj[x]) radj[y].push_back(x);
    for (int x = 0; x < n; x++)
      if (!vis[x]) dfs1(x);
    reverse(todo.begin(), todo.end());
    for (int x : todo)
     if (comp[x] == -1) dfs2(x, time++);
   return comp;
}; // namespace SCCKosaraju
Dinic.h
```

Description: Max flow algorithm. Can find a valid circulation given vertex and/or edge demands. Time: $O(VE \log U)$ 9ad873, 84 lines

```
struct Dinic {
 // disable scaling when max flow/capacity is small, or
 // sometimes on random data
 const static bool SCALING = true;
 struct Edge {
   int v, dual;
   ll cap, res;
   constexpr 11 flow() { return max(cap - res, 011); }
 int n, s, t;
 vector<int> lvl, q, blk;
 vector<vector<Edge>> adj;
 vector<pair<int, int>> edges;
```

```
Dinic(int n): n(n + 2), s(n++), t(n++), q(n), adj(n) {}
int add(int u, int v, 11 cap, 11 flow = 0) {
  adj[u].push_back({v, int(adj[v].size()), cap, cap - flow});
  adj[v].push_back({u, int(adj[u].size()) - 1, 0, 0});
  edges.emplace_back(u, adj[u].size() - 1);
  return edges.size() - 1; // this Edge's ID
11 dfs(int u, 11 in) {
  if (u == t || !in) return in;
  11 \text{ flow} = 0:
  for (auto& e : adj[u]) {
    if (e.res && !blk[e.v] && lvl[e.v] == lvl[u] - 1)
      if (ll out = dfs(e.v, min(in, e.res))) {
        flow += out, in -= out, e.res -= out;
        adj[e.v][e.dual].res += out;
        if (!in) return flow;
  blk[u] = 1;
  return flow;
ll flow() {
  11 \text{ flow} = 0;
  q[0] = t;
  for (int B = SCALING * 30; B >= 0; B--) do {
      lvl = blk = vector<int>(n);
      int qi = 0, qe = lvl[t] = 1;
      while (qi < qe && !lvl[s]) {
        int u = q[qi++];
        for (auto& e : adj[u])
          if (!lvl[e.v] && adj[e.v][e.dual].res >> B)
            q[qe++] = e.v, lvl[e.v] = lvl[u] + 1;
      if (lvl[s]) flow += dfs(s, LLONG_MAX);
    } while (lvl[s]);
  return flow;
Edge& get(int id) { // get Edge object from its ID
  return adj[edges[id].first][edges[id].second];
void clear() {
  for (auto& it : adj)
    for (auto& e : it) e.res = e.cap;
bool leftOfMinCut(int u) { return lvl[u] == 0; }
//d is a list of vertex demands, d[u] = flow in -flow out
// negative if u is a source, positive if u is a sink
bool circulation(vector<ll> d = {}) {
  d.resize(n);
  vector<int> circEdges;
  Dinic q(n);
  for (int u = 0; u < n; u++)
    for (auto& e : adj[u]) {
      d[u] += e.flow(), d[e.v] -= e.flow();
      if (e.res) circEdges.push_back(g.add(u, e.v, e.res));
  int tylerEdge = g.add(t, s, LLONG_MAX, 0);
  11 \text{ flow} = 0;
  for (int u = 0; u < n; u++)
    if (d[u] < 0)
      g.add(g.s, u, -d[u]);
    else if (d[u] > 0)
      g.add(u, g.t, d[u]), flow += d[u];
  if (flow != q.flow()) return false;
  int i = 0; // reconstruct the flow into this graph
  for (int u = 0; u < n; u++)
    for (auto& e : adj[u])
      if (e.res) e.res -= g.get(circEdges[i++]).flow();
  return true;
```

```
}
;
```

Mathematics (5)

Fraction.h

Description: Struct for representing fractions/rationals. All ops are $O(\log N)$ due to GCD in constructor. Uses cross multiplication alde34, 27 lines

```
template <typename T>
struct 0 {
  T a, b;
  Q(T p, T q = 1) {
   T g = gcd(p, q);
    a = p / q;
   b = q / g;
   if (b < 0) a = -a, b = -b;
  T gcd(T x, T y) const { return __gcd(x, y); }
  Q operator+(const Q& o) const {
    return {a * o.b + o.a * b, b * o.b};
  O operator-(const O& o) const {
    return *this + Q(-o.a, o.b);
  Q operator*(const Q& o) const { return {a * o.a, b * o.b}; }
  O operator/(const O& O) const { return *this * O(o.b, o.a); }
  Q recip() const { return {b, a}; }
  int signum() const { return (a > 0) - (a < 0); }</pre>
  bool operator<(const Q& o) const {
    return a * o.b < o.a * b;
  friend ostream& operator<<(ostream& cout, const Q& o) {</pre>
   return cout << o.a << "/" << o.b;
};
```

FractionOverflow.h

Description: Safer struct for representing fractions/rationals. Comparison is 100% overflow safe; other ops are safer but can still overflow. All ops are $O(\log N)$.

```
template <typename T>
struct 00
 T a, b;
  QO(T p, T q = 1) {
   T g = gcd(p, q);
   a = p / g;
   b = q / g;
   if (b < 0) a = -a, b = -b;
  T gcd(T x, T y) const { return __gcd(x, y); }
  QO operator+(const QO& o) const {
   T g = gcd(b, o.b), bb = b / g, obb = o.b / g;
    return {a * obb + o.a * bb, o.b * obb};
  QO operator-(const QO& o) const {
    return *this + 00(-o.a, o.b);
  QO operator*(const QO& o) const {
   T g1 = gcd(a, o.b), g2 = gcd(o.a, b);
   return { (a / q1) * (o.a / q2), (b / q2) * (o.b / q1) };
  QO operator/(const QO& o) const {
    return *this * QO(o.b, o.a);
  QO recip() const { return {b, a}; }
  int signum() const { return (a > 0) - (a < 0); }</pre>
```

```
static bool lessThan(T a, T b, T x, T y) {
   if (a / b != x / y) return a / b < x / y;
   if (x % y == 0) return false;
   if (a % b == 0) return true;
   return lessThan(y, x % y, b, a % b);
}
bool operator<(const QO& o) const {
   if (this->signum() != o.signum() || a == 0) return a < o.a;
   if (a < 0)
     return lessThan(abs(o.a), o.b, abs(a), b);
   else
     return lessThan(a, b, o.a, o.b);
}
friend ostream& operator<<(ostream& cout, const QO& o) {
    return cout << o.a << "/" << o.b;
}
};</pre>
```

PrimeSieve.h

Description: Prime sieve for generating all primes up to a certain limit. is prime [i] is true iff i is a prime.

Time: $\lim_{m \to 100'000'000} \approx 0.8$ s. Runs 30% faster if only odd indices are stored.

```
const int MAX_PR = 5'000'000;
bitset<MAX_PR> isprime;
vector<int> primeSieve(int lim) {
  isprime.set();
  isprime[0] = isprime[1] = 0;
  for (int i = 4; i < lim; i += 2) isprime[i] = 0;
  for (int i = 3; i * i < lim; i += 2)
    if (isprime[i])
      for (int j = i * i; j < lim; j += i * 2) isprime[j] = 0;</pre>
```

PrimeSieveFast.h

return pr;

vector<int> pr;

for (**int** i = 2; i < lim; i++)

if (isprime[i]) pr.push_back(i);

Description: Prime sieve for generating all primes smaller than LIM. Time: LIM= $1e9 \approx 1.5s$

```
const int LIM = 1e8;
bitset<LIM> isPrime;
vector<int> primeSieve() {
  const int S = round(sqrt(LIM)), R = LIM / 2;
  vector < int > pr = \{2\}, sieve(S + 1);
  pr.reserve(int(LIM / log(LIM) * 1.1));
  vector<pair<int, int>> cp;
  for (int i = 3; i <= S; i += 2)
    if (!sieve[i]) {
      cp.push back(\{i, i * i / 2\});
      for (int j = i * i; j <= S; j += 2 * i) sieve[j] = 1;</pre>
  for (int L = 1; L <= R; L += S) {</pre>
    array<bool, S> block{};
    for (auto& [p, idx] : cp)
      for (int i = idx; i < S + L; idx = (i += p))</pre>
        block[i - L] = 1;
    for (int i = 0; i < min(S, R - L); i++)</pre>
      if (!block[i]) pr.push back((L + i) \star 2 + 1);
  for (int i : pr) isPrime[i] = 1;
  return pr:
```

Miscellaneous (6)

NDimensionalVector.h

3c0f61, 12 lines

Submasks.h

```
for (int mask = 0; mask < (1 << n); mask++)
for (int sub = mask; sub; sub = (sub - 1) & mask)
// do thing</pre>
```

Strings (7)

ZValues.h

151ee3, 10 lines

35424b, 3 lines

```
vector<int> zValues(string& s) {
  int n = ( int )s.length();
  vector<int> z(n);
  for (int i = 1, l = 0, r = 0; i < n; ++i) {
    if (i <= r) z[i] = min(r - i + 1, z[i - 1]);
    while (i + z[i] < n && s[z[i]] == s[i + z[i]]) ++z[i];
    if (i + z[i] - 1 > r) l = i, r = i + z[i] - 1;
  }
  return z;
}
```