

# Parameterized Quantum Circuits and Bayesian Optimization for violating the Bell Inequality

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## Introduction

### Aim

Win a formulation of the CHSH game

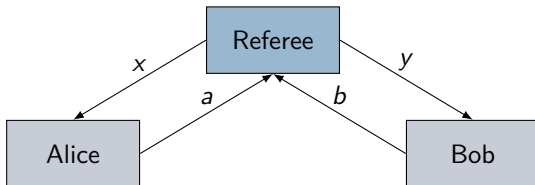
### Methodology

Training a simple parameterized quantum circuit

### Data

Outputs of quantum circuits simulations

## Problem Formulation: CHSH Game [1]



- 1 The referee sends a bit  $x$  to Alice and a bit  $y$  to Bob. These bits are drawn from two Bernoulli-distributed random variables:

$$X \sim \mathcal{B}(0.5), \quad Y \sim \mathcal{B}(0.5)$$

- 2 Alice and Bob send back, respectively, bits  $a$  and  $b$ .

Alice and Bob **win the game** if<sup>1</sup>:

$$x \wedge y = a \oplus b$$

(1)

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<sup>1</sup>Where  $\wedge$  denotes the logical *AND* and  $\oplus$  the logical *XOR* operations.

## Strategy [1]

Naming  $p_{AB|XY}(a, b|x, y)$  the conditional distribution which relates to the **strategy** operated by Alice and Bob, we define the winning function  $V : \{0, 1\}^4 \rightarrow \{0, 1\}$  as:

$$V(x, y, a, b) := \begin{cases} 1 & \text{if } x \wedge y = a \oplus b, \\ 0 & \text{otherwise.} \end{cases} \quad (2)$$

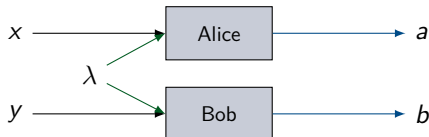
and, hence, express the probability of winning the CHSH game as

$$\begin{aligned} \Pr\{\text{win}\} &= \sum_{a,b,x,y} V(x, y, a, b) p_{AB|XY}(a, b|x, y) p_{XY}(x, y) = \\ &= \frac{1}{4} \sum_{a,b,x,y} V(x, y, a, b) p_{AB|XY}(a, b|x, y) \end{aligned} \quad (3)$$

**Assumption:** there is a random variable  $\Lambda$  that describes the strategy:

$$p_{AB|XY}(a, b|x, y) = \int d\lambda p_{AB|\Lambda XY}(a, b|\lambda, x, y) p_{\Lambda|XY}(\lambda|x, y) \quad (4)$$

## Classical Strategy [1]



We have conditional independence:

Also it is always possible to write:

$$\begin{cases} p_{\Lambda|XY} = p_{\Lambda}(\lambda) \\ p_{AB|\Lambda XY} = p_{A|\Lambda X}(a|\lambda, x) p_{B|\Lambda Y}(b|\lambda, y) \end{cases} \quad \begin{cases} p_{A|\Lambda X}(a|\lambda, x) = \int dn f(a|\lambda, x, n) p_N(n) \\ p_{B|\Lambda Y}(b|\lambda, y) = \int dm g(b|\lambda, y, m) p_M(m) \end{cases}$$

$$\begin{aligned} p_{AB|XY} &= \int \int \int d\lambda dn dm f(a|\lambda, x, n) g(b|\lambda, y, m) p_{\Lambda}(\lambda) p_N(n) p_M(m) \\ &= \int d\lambda f'(a|\lambda, x) g'(b|\lambda, y) p_{\Lambda}(\lambda) \end{aligned} \quad (5)$$

where  $f, g, f'$  and  $g'$  are deterministic (binary-valued) functions.

## Classical Strategy - II [1]

We can then write

$$\begin{aligned}
 \Pr\{\text{win}\} &= \frac{1}{4} \sum_{a,b,x,y} V(x,y,a,b) p_{AB|XY}(a,b|x,y) \\
 &= \frac{1}{4} \sum_{a,b,x,y} V(x,y,a,b) \int d\lambda f'(a|\lambda,x) g'(b|\lambda,y) p_{\Lambda}(\lambda) \\
 &= \int d\lambda p_{\Lambda}(\lambda) \left[ \frac{1}{4} \sum_{a,b,x,y} V(x,y,a,b) f'(a|\lambda,x) g'(b|\lambda,y) \right] \\
 &\leq \frac{1}{4} \sum_{a,b,x,y} V(x,y,a,b) f'(a|\lambda^*,x) g'(b|\lambda^*,y) \tag{6}
 \end{aligned}$$

and conclude that deterministic strategies are optimal among all classical strategies.

## Deterministic Classical Strategy [1]

Naming  $a_x$  ( $b_y$ ) the bit that Alice (Bob) sends once received the bit  $x$  ( $y$ ), we can list all possible outcome of the CHSH game:

$x$	$y$	$x \wedge y$	$= a_x \oplus b_y$
0	0	0	$= a_0 \oplus b_0$
0	1	0	$= a_0 \oplus b_1$
1	0	0	$= a_1 \oplus b_1$
1	1	1	$= a_1 \oplus b_1$

We can see that, for a *deterministic classical strategy*, we have

$$\mathbf{Pr}\{\text{win}\} = \frac{1}{4} \sum_{a,b,x,y} V(x,y,a,b) p_{AB|XY}(a,b|x,y) \leq \frac{3}{4} \quad (7)$$

And that

$$\alpha := \mathbf{Pr}\{\text{win}\} - \mathbf{Pr}\{\text{lose}\} \implies \|4\alpha\| \leq 2 \quad (8)$$

## Quantum Strategy [1]

Parameter  $\lambda$  can correspond to a shared quantum state  $|\phi\rangle_{AB}$ .

Alice's  $x$ -dependent local projective measurements:  $\{\Pi_a^{(x)}\} : \sum_a \Pi_a^{(x)} = I$

Bob's  $y$ -dependent local projective measurements:  $\{\Pi_b^{(y)}\} : \sum_b \Pi_b^{(y)} = I$

We then have

$$p_{AB|XY}(a, b|x, y) = \langle \phi|_{AB} \Pi_a^{(x)} \otimes \Pi_b^{(y)} |\phi\rangle_{AB} \quad (9)$$

and we can re-write Eq. 4 as

$$\mathbf{Pr}\{\text{win}\} = \frac{1}{4} \sum_{a,b,x,y} V(x, y, a, b) \langle \phi|_{AB} \Pi_a^{(x)} \otimes \Pi_b^{(y)} |\phi\rangle_{AB} \quad (10)$$

We define  $A^{(x)} := \Pi_0^{(x)} - \Pi_1^{(x)}$  and  $B^{(y)} := \Pi_0^{(y)} - \Pi_1^{(y)}$  and:

$$C_{AB} := A^{(0)} \otimes B^{(0)} + A^{(0)} \otimes B^{(1)} + A^{(1)} \otimes B^{(0)} - A^{(1)} \otimes B^{(1)} \quad (11)$$



## Quantum Strategy - II [1], [2]

It is possible to show that

$$\alpha = \mathbf{Pr}\{\text{win}\} - \mathbf{Pr}\{\text{lose}\} = \frac{1}{4} \langle \phi |_{AB} C_{AB} | \phi \rangle_{AB} \quad (12)$$

and also that

$$\|C_{AB}\|_{\infty} \leq 2\sqrt{2} \quad (13)$$

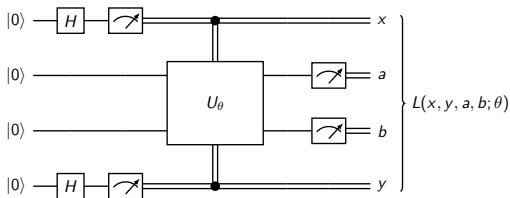
This implies that

$$\|4\alpha\|_{\infty} \leq 2\sqrt{2} \quad (14)$$

which is known as the *Tsirelson's bound*. This result translates to a larger maximum winning probability:

$$\mathbf{Pr}\{\text{win}\} \leq \frac{1}{2} + \frac{\sqrt{2}}{4} \approx 0.835 \quad (15)$$

## General Idea



- ▶ The idea is to learn an optimal unitary such that some loss function is maximized.
- ▶ Relate the loss function to the possibility of winning the CHSH game.

An obvious choice is just the winning frequency on  $N$  simulations:

$$L(x, y, a, b; \theta) = \frac{1}{N} \sum_{i=1}^N V_\theta(x, y, a, b) \quad (16)$$

## Parameterized Quantum Circuits [3]

$$|\psi_0\rangle \equiv \boxed{U_\theta} \equiv |\psi_\theta\rangle$$

We define a parameterized quantum circuit (PQC) as a parameterized unitary operation on  $n$  qubits,  $U_\theta$ :

$$|\psi_\theta\rangle = U_\theta |\psi_0\rangle \quad (17)$$

where, usually,  $|\psi_0\rangle = |0\rangle^{\otimes n}$ . Sometimes compared to neural networks:

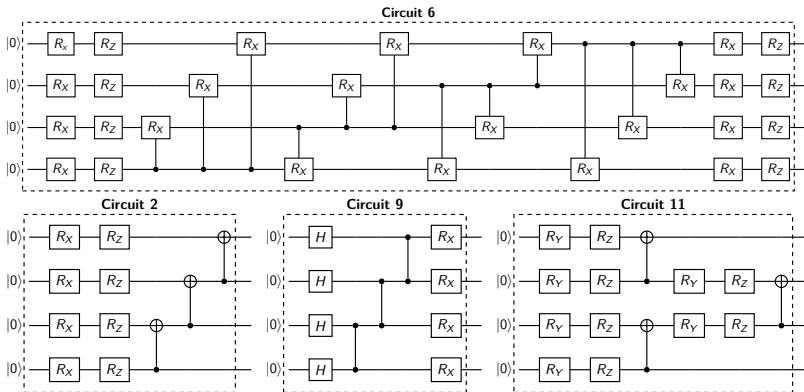
### Similarities

- ▶ Universal approximator.
- ▶ Number of parameters scales polynomially in  $n$ .

### Differences

- ▶ Unitary.
- ▶ Internal state not accessible.

## Parameterized Quantum Circuits: Some Examples [4]



**Expressibility** and **Entangling Capability** studied on 19 (4-qubits) circuits wrt the increasing number of consecutive layers (from 1 to 4).



## Circuit Descriptors: Entangling Capability [4], [5]

The entangling capability of a PQC is estimated through the Meyer-Wallach measure, here presented in the form given by Brennen. For a state  $|\psi\rangle \in \mathbb{C}^{2^n}$  we define

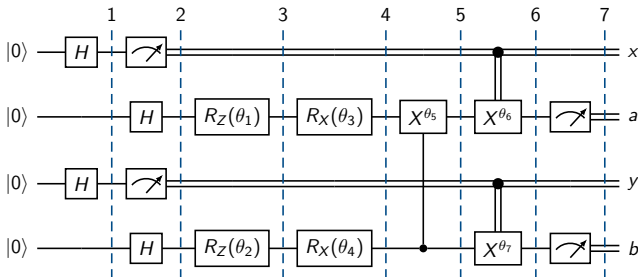
$$Q(|\psi\rangle) := \frac{1}{n} \sum_{k=1}^n 2(1 - \text{Tr}[\rho_k^2]) \quad (19)$$

where  $\rho_k$  is the one-qubit reduced density operator of the  $k$ -th qubit after tracing out the rest. Being  $|\psi_\theta\rangle = U_\theta |\psi_0\rangle$  we define the entangling capability as

$$\text{Ent} := \frac{1}{|S|} \sum_{\theta \in S} Q(|\psi_\theta\rangle) \quad (20)$$

where  $S = \{\theta_i\}_{i=0}^N$  is the set of sampled circuit parameters.

## The PQC in detail



Parameters  $\theta_1, \theta_2, \theta_3$  and  $\theta_4$  have been constrained inside  $[0, 2\pi]$  while  $\theta_5, \theta_6$  and  $\theta_7$  inside  $[-2, 2]$ .

**Cost:** There are 11 gates: 8 single-qubit gates and 3 two-qubit gates. The parametric gates are 4 rotations and 3 CNOTs.

**Expressibility:**  $\text{Exp} \approx 0.05$ .

**Entangling capability:**  $\text{Ent} \approx 0.2$ .

# Bayesian Optimization [6]

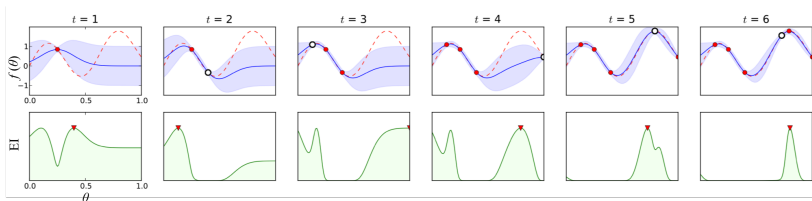
**Aim:** Find

$$\theta^* = \arg \max_{\theta \in \Theta} L(x, y, a, b; \theta) \quad (21)$$

**Idea:** Attack the problem with Bayesian Optimization:

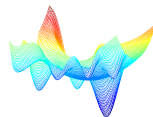
- ▶ Objective function is costly to evaluate.
- ▶ We don't have access to derivatives.
- ▶ The parameter space is "small" (usually  $\dim(\Theta) < 20$ ).

Actors involved: a prior distribution, a posterior distribution and an *acquisition function*.





## Simulation Setup

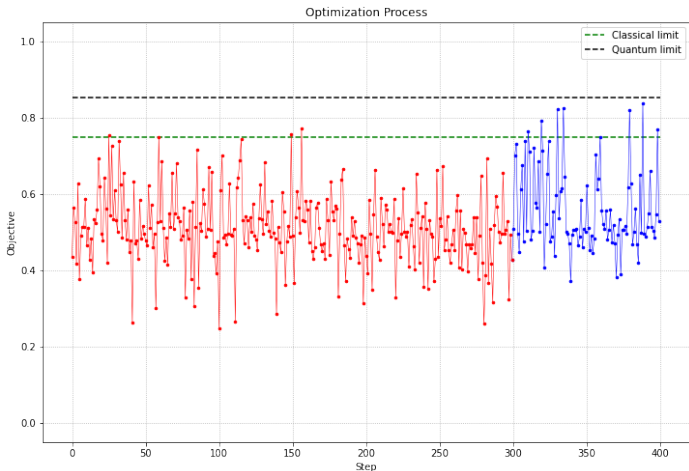


### Pseudo-code

- 1: Init PQC
- 2: Init Simulator
- 3: **for**  $i = 1, \dots, N$  **do**
- 4:     Sample  $\theta_i \in \Theta$  according to the acquisition function.
- 5:     **for**  $j = 1, \dots, M$  **do**
- 6:         Run the simulation and collect data.
- 7:     **end for**
- 8:     Update prior and posterior based on data.
- 9: **end for**
- 10: **return** Best set of parameters  $\theta^*$ .

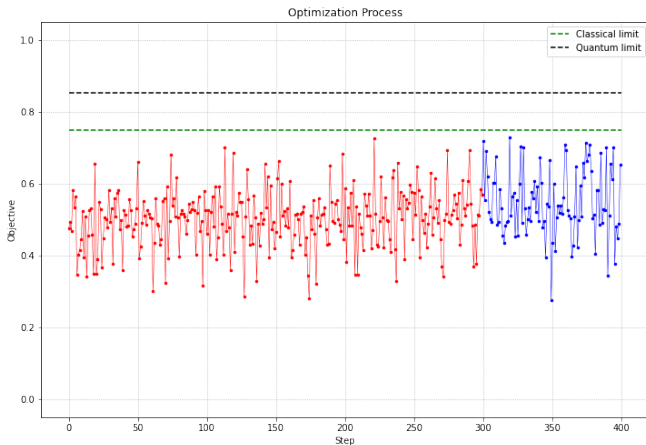
## Results

Using  $N = 400$  optimization steps and  $M = 1000$  circuit simulations per step:



## Results - II

Run the previous algorithm with a circuit with no entangling capability ( $\text{Ent} = 0.$ ) and comparable expressibility ( $\text{Exp} \approx 0.04$ ). This PCQ has 22 gates: 20 one-qubit gates and 2 two-qubit gates.

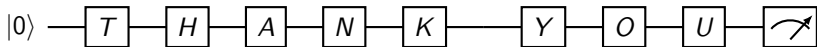


## Conclusions

- ▶ Joint work of PQC and classical optimization → win the CHSH game
- ▶ Goal: show that entangle can emerge naturally from the setup
- ▶ More of a PoC than a real experiment
- ▶ Provide a simple setup for future and more complex ideas

Look at the code

[https://github.com/w00zie/pqc\\_chsh](https://github.com/w00zie/pqc_chsh)



## References I

- [1] M. M. Wilde, “Preface to the second edition,” *Quantum Information Theory*, pp. xi–xii, DOI: 10.1017/9781316809976.001. [Online]. Available: <http://dx.doi.org/10.1017/9781316809976.001>.
- [2] B. S. Cirel'Son, “Quantum generalizations of Bell's inequality,” *Letters in Mathematical Physics*, vol. 4, no. 2, pp. 93–100, Mar. 1980. DOI: 10.1007/BF00417500.
- [3] M. Benedetti, E. Lloyd, S. Sack, and M. Fiorentini, “Parameterized quantum circuits as machine learning models,” *Quantum Science and Technology*, vol. 4, no. 4, p. 043 001, Nov. 2019, ISSN: 2058-9565. DOI: 10.1088/2058-9565/ab4eb5. [Online]. Available: <http://dx.doi.org/10.1088/2058-9565/ab4eb5>.

## References II

- [4] S. Sim, P. D. Johnson, and A. Aspuru-Guzik, “Expressibility and entangling capability of parameterized quantum circuits for hybrid quantum-classical algorithms,” *Advanced Quantum Technologies*, vol. 2, no. 12, p. 1900070, Oct. 2019, ISSN: 2511-9044. DOI: 10.1002/qute.201900070. [Online]. Available: <http://dx.doi.org/10.1002/qute.201900070>.
- [5] G. K. Brennen, “An observable measure of entanglement for pure states of multi-qubit systems,” *Quantum Info. Comput.*, vol. 3, no. 6, pp. 619–626, Nov. 2003, ISSN: 1533-7146.
- [6] P. I. Frazier, *A tutorial on bayesian optimization*, 2018. arXiv: 1807.02811 [stat.ML].

## Maximum Quantum Winning Probability - 1

When the inputs  $x$  and  $y$  are 00, 01, 10 Alice and Bob win if they report back the same result. The probability is

$$\langle \phi |_{AB} \Pi_0^{(x)} \otimes \Pi_0^{(y)} | \phi \rangle_{AB} + \langle \phi |_{AB} \Pi_1^{(x)} \otimes \Pi_1^{(y)} | \phi \rangle_{AB}$$

The probability for it not to happen is

$$\langle \phi |_{AB} \Pi_0^{(x)} \otimes \Pi_1^{(y)} | \phi \rangle_{AB} + \langle \phi |_{AB} \Pi_1^{(x)} \otimes \Pi_0^{(y)} | \phi \rangle_{AB}$$

So in this case  $\Pr\{\text{win}\} - \Pr\{\text{lose}\} = \langle \phi |_{AB} A^{(x)} \otimes B^{(y)} | \phi \rangle_{AB}$

When  $x = 1$  and  $y = 1$  we have that

$$\Pr\{\text{win}\} - \Pr\{\text{lose}\} = - \langle \phi |_{AB} A^{(1)} \otimes B^{(1)} | \phi \rangle_{AB}$$

So, averaging over all values of the input bits

$$\alpha := \Pr\{\text{win}\} - \Pr\{\text{lose}\} = \frac{1}{4} \langle \phi |_{AB} C_{AB} | \phi \rangle_{AB}$$



## Maximum Quantum Winning Probability - 2

We have

$$C_{AB}^2 = 4I_{AB} - [A^{(0)}, A^{(1)}] \otimes [B^{(0)}, B^{(1)}]$$

and

$$\begin{aligned} \|C_{AB}^2\|_{\infty} &= \|4I_{AB} - [A^{(0)}, A^{(1)}] \otimes [B^{(0)}, B^{(1)}]\|_{\infty} \\ &\leq 4\|I_{AB}\|_{\infty} + \|[A^{(0)}, A^{(1)}] \otimes [B^{(0)}, B^{(1)}]\|_{\infty} \\ &= 4 + \|[A^{(0)}, A^{(1)}]\|_{\infty} \|[B^{(0)}, B^{(1)}]\|_{\infty} \\ &\leq 4 + 2 \cdot 2 = 8 \end{aligned}$$

So  $\|C_{AB}\|_{\infty} \leq 2\sqrt{2}$  and

$$\|4\alpha\|_{\infty} = \|\langle \phi|_{AB} C_{AB} |\phi\rangle_{AB}\|_{\infty} \leq 2\sqrt{2}$$

## Expressibility Calculation

Expressibility is

$$\text{Exp} := D_{\text{KL}}(P_{\text{PQC}}(F; \theta) \| P_{\text{Haar}}(F))$$

where, being  $d = \dim(\mathcal{H})$

$$P_{\text{Haar}}(F) = (d - 1)(1 - F)^{d-2} \quad (22)$$

### Expressibility

- 1: Init PQC
- 2: Init Simulator
- 3: **for**  $i = 1, \dots, M$  **do**
- 4:     Sample two sets of parameters  $\theta_i, \theta_j$  uniformly in  $\Theta$ .
- 5:     Calculate fidelity between the two output states.
- 6: **end for**
- 7: Calculate KL between fidelities and (22)

## Expressibility Calculation - 2

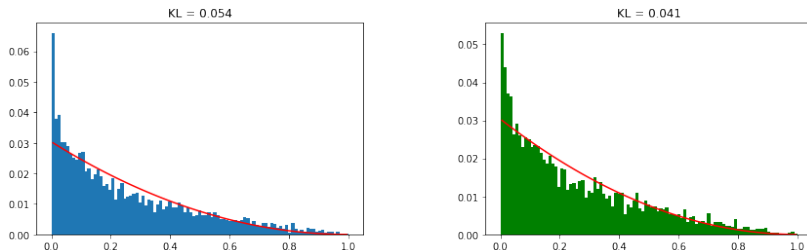


Figure 1: Entangling circuit (left) and non-entangling circuit (right)

Estimated with  $M = 5000$  samples per circuit, compared to the theoretical density function (22) in red.