Notes for ECE 27000 - Introduction to Digital System Design

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Contents

Course Description

An introduction to digital system design, with an emphasis on practical design techniques and circuit implementation.

Verilog

SystemVerilog is a hardware description and verification language (HDVL) that extends Verilog by adding high-level programming constructs. It is widely used for modeling, simulating, and verifying digital systems.

The listing below is an example of a NAND gate expressed in SystemVerilog.

Listing 1: NAND

```
module nand_gate (
    input   logic a,
    input   logic b,
    output  logic y
);
    assign y = ~(a & b);
endmodule
```

In SystemVerilog, numbers are written in format `[size]'[base][number]`, for example:

- 4'b1001 (binary, 9 in decimal, bit width 4 bits)

- 8'hf1 (hex, equals 421, bit width 8 bits)

- 3'o3 (octal, 3, bit width 3 bits)

- 32'b1001_1101_0101_1111 (binary, 40255, bit width 32 bits)

Number Systems

In daily life, we primarily interact with the familiar base-10 numbers. However, when interaction with digital systems, we must also concern ourselves with base-2, base-8, base-16, and other bases which are friendly to binary states. Unless completely unambiguous, the base of a number is written as a right subscript such as 144_{10} for base-10 or 1001_2 for base-2.

For binary numbers, each digit represents a power of two. To convert a binary number to decimal, you sum the products of each binary digit with its corresponding power of two. For example, the binary number 1001 is calculated as

$$2^3 \times 1 + 2^2 \times 0 + 2^1 \times 0 + 2^0 \times 1 = 8 + 0 + 0 + 1 \qquad (1)$$
$$= 9_{10} \qquad (2)$$

To convert from hexadecimal to decimal, each digit represents a power of sixteen. For instance, the hexadecimal number f1 is calculated as

$$15 \times 161 + 1 \times 160 = 240 + 1 \qquad (3)$$
$$= 241_{10} \qquad (4)$$

where f represents the decimal value 15. When converting to another base, reverse the process by dividing the decimal number by the target base, recording the remainder, and repeating with the quotient until it reaches zero. The remainders give you the digits of the number in the new base, read in reverse order.

To convert a decimal number into binary, for example, you repeatedly divide the number by 2 and record the remainders. For the decimal number 9, dividing by 2 gives a quotient of 4 and a remainder of 1. Dividing 4 by 2 gives a quotient of 2 and a remainder of 0. Dividing 2 by 2 gives a quotient of 1 and a remainder of 0, and finally, dividing 1 by 2 gives a quotient of 0 and a remainder of 1. Reading the remainders from bottom to top, the binary representation of 9 is 1001.

Boolean Algebra

Computers operate in binary. To represent the state of a computer we require a suitable mathematical framework, provided by boolean algebra. In boolean algebra, variables can only take on two values: 0 and 1.

Rule	Expression
Commutativity	$X + Y = Y + X$
	$X \cdot Y = Y \cdot X$
Associativity	$(X + Y) + Z = X + (Y + Z)$
	$(X \cdot Y) \cdot Z = X \cdot (Y \cdot Z)$
Distributivity	$X \cdot Y + X \cdot Z = X \cdot (Y + Z)$
	$(X + Y) \cdot (X + Z) = X + Y \cdot Z$
Covering	$X + X \cdot Y = X$
	$X \cdot (X + Y) = X$
Combining	$X \cdot Y + X \cdot Y = X$
	$(X + Y) \cdot (X + Y) = X$
Consensus	$X \cdot Y + X \cdot Z + Y \cdot Z = X \cdot Y + X' \cdot Z$
	$(X + Y) \cdot (X + Z) \cdot (Y + Z) = (X + Y) \cdot (X + Z)$
Generalized Idempotency	$X + X + \cdots + X = X$
	$X \cdot X \cdot \cdots \cdot X = X$
DeMorgan's Theorems	$(X_1 \cdot X_2 \cdot \cdots \cdot X_n)' = X_1' + X_2' + \cdots + X_n'$
	$(X_1 + X_2 + \cdots + X_n)' = X_1' \cdot X_2' \cdot \cdots \cdot X_n'$
Generalized DeMorgan's	$F(X_1, X_2, \ldots, X_n, +, \cdot) = F(X_1, X_2, \ldots, X_n, \cdot, +)'$
Shannon's Expansion	$F(X_1, X_2, \ldots, X_n) = X_1 \cdot F(1, X_2, \ldots, X_n) + X_1' \cdot F(0, X_2, \ldots, X_n)$
	$F(X_1, X_2, \ldots, X_n) = [X_1 + F(0, X_2, \ldots, X_n)] \cdot [X_1' + F(1, X_2, \ldots, X_n)]$

Table 1: Boolean Algebra

An interesting and useful property in boolean algebra is "duality", where replacing all ANDs with ORs and all 1s with 0s gives a valid and equivalent theorem. For instance,

X AND 0 = 0	X OR 1 = 1
X OR 0 = X	X AND 1 = X

Table 2: Boolean Duality

Any logic can be implemented using just the following:

- AND, OR, and NOT gates

- NAND gates

- NOR gates

Logic Gates

We represent boolean operations in circuit diagrams with gates. Below are their symbols and corresponding truth tables.

Buffer

A	Output
0	0
1	1

AND

A	B	Output
0	0	0
0	1	0
1	0	0
1	1	1

OR

A	B	Output
0	0	0
0	1	1
1	0	1
1	1	1

NOT

A	Output
0	1
1	0

NAND

A	B	Output
0	0	1
0	1	1
1	0	1
1	1	0

NOR

A	B	Output
0	0	1
0	1	0
1	0	0
1	1	0

XOR

A	B	Output
0	0	0
0	1	1
1	0	1
1	1	0

XNOR

A	B	Output
0	0	1
0	1	0
1	0	0
1	1	1

CMOS

Complementary Metal-Oxide-Semiconductor (CMOS) technology is the dominant semiconductor technology for modern integrated circuits. CMOS combines both n-type (NMOS) and p-type (PMOS) metal-oxide-semiconductor field-effect transistors (MOSFETs) to create the familiar logic gates such as AND, NOT, XOR, etc.

PMOS
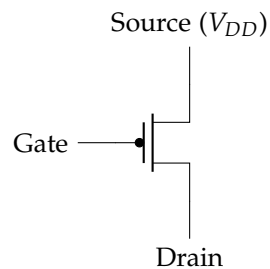
Source (V_{DD})

Gate ———

Drain

Figure 1:   PMOS transistor circuit symbol

PMOS transistors consist of:

- p+ source and drain regions

- n-type substrate (body)

- SiO₂ gate dielectric

- Polysilicon gate electrode

PMOS operates with negative gate-to-source voltage (V_{GS}):

- **Cut-off Region** ($V_{GS} > V_{th,p}$):

$$I_D = 0$$

- **Linear Region** ($V_{GS} \leq V_{th,p}$ and $V_{DS} \geq V_{GS} - V_{th,p}$):

$$I_D = -\mu_p C_{ox} \frac{W}{L} \left[(V_{GS} - V_{th,p}) V_{DS} - \frac{V_{DS}^2}{2} \right]$$

- **Saturation Region** ($V_{GS} \leq V_{th,p}$ and $V_{DS} < V_{GS} - V_{th,p}$):

$$I_D = -\frac{1}{2} \mu_p C_{ox} \frac{W}{L} (V_{GS} - V_{th,p})^2$$

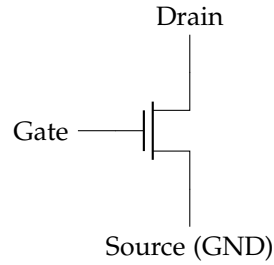### NMOS

NMOS transistors consist of:

Drain

Gate

Source (GND)

- n+ source and drain regions

- p-type substrate (body)

- SiO₂ gate dielectric

- Polysilicon gate electrode

  NMOS operates with positive gate-to-source voltage (V_{GS}):

- **Cut-off Region** ($V_{GS} < V_{th,n}$):

$$I_D = 0$$

- **Linear Region** ($V_{GS} \geq V_{th,n}$ and $V_{DS} \leq V_{GS} - V_{th,n}$):

$$I_D = \mu_n C_{ox} \frac{W}{L} \left[(V_{GS} - V_{th,n}) V_{DS} - \frac{V_{DS}^2}{2} \right]$$

- **Saturation Region** ($V_{GS} \geq V_{th,n}$ and $V_{DS} > V_{GS} - V_{th,n}$):

$$I_D = \frac{1}{2} \mu_n C_{ox} \frac{W}{L} (V_{GS} - V_{th,n})^2$$

Parameter	PMOS	NMOS
Majority Carrier	Holes	Electrons
Substrate Type	n-type	p-type
Threshold Voltage	Negative	Positive
Mobility (μ)	Lower ($\approx 150 \frac{cm^2}{V_s}$)	Higher ($\approx 400 \frac{cm^2}{V_s}$)
Speed	Slower	Faster

Table 3: PMOS vs NMOS characteristics