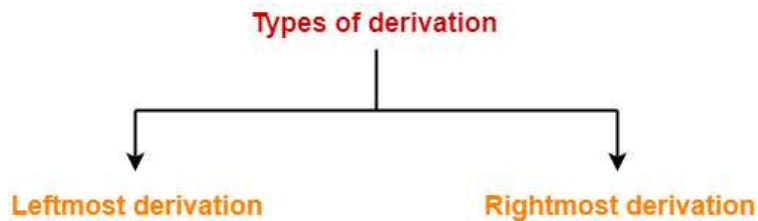


## Parse Tree-

- The process of deriving a string is called as **derivation**.
- The geometrical representation of a derivation is called as a **parse tree** or **derivation tree**.



### 1. Leftmost Derivation-

- The process of deriving a string by expanding the leftmost non-terminal at each step is called as **leftmost derivation**.
- The geometrical representation of leftmost derivation is called as a **leftmost derivation tree**.

### Example-

Consider the following grammar-

$$S \rightarrow aB / bA$$

$$S \rightarrow aS / bAA / a$$

$$B \rightarrow bS / aBB / b$$

(Unambiguous Grammar)

Let us consider a string  $w = aaabbabbba$

Now, let us derive the string  $w$  using leftmost derivation.

### Leftmost Derivation-

$$S \rightarrow aB$$

→ aa**B**B (Using  $B \rightarrow aBB$ )

→ aaa**B**BB (Using  $B \rightarrow aBB$ )

→ aaab**BB** (Using  $B \rightarrow b$ )

→ aaabb**B** (Using  $B \rightarrow b$ )

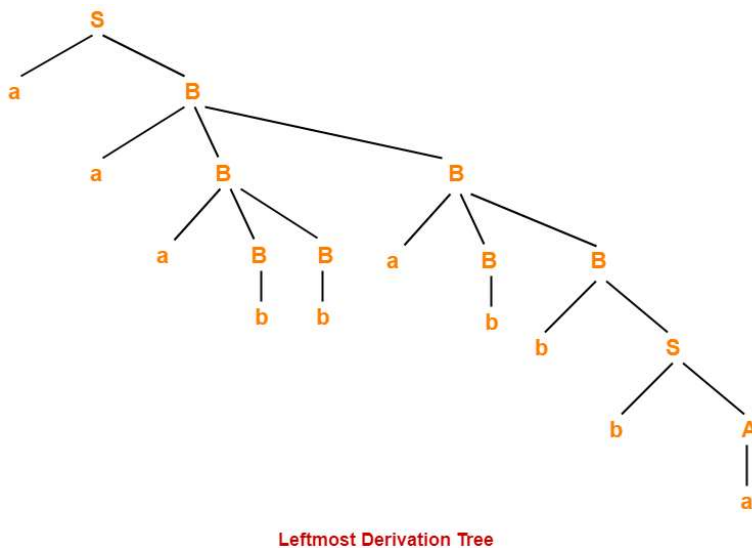
→ aaabba**BB** (Using  $B \rightarrow aBB$ )

→ aaabbab**B** (Using  $B \rightarrow b$ )

→ aaabbabb**S** (Using  $B \rightarrow bS$ )

→ aaabbabbb**A** (Using  $S \rightarrow bA$ )

→ aaabbabbba (Using  $A \rightarrow a$ )



## 2. Rightmost Derivation-

- The process of deriving a string by expanding the rightmost non-terminal at each step is called as **rightmost derivation**.
- The geometrical representation of rightmost derivation is called as a **rightmost derivation tree**.

### Example-

Consider the following grammar-

$$S \rightarrow aB / bA$$

$$S \rightarrow aS / bAA / a$$

$$B \rightarrow bS / aBB / b$$

**(Unambiguous Grammar)**

Let us consider a string  $w = aaabbabbba$

Now, let us derive the string  $w$  using rightmost derivation.

### **Rightmost Derivation-**

$$S \rightarrow aB$$

$$\rightarrow aaBB \text{ (Using } B \rightarrow aBB \text{)}$$

$$\rightarrow aaBaBB \text{ (Using } B \rightarrow aBB \text{)}$$

$$\rightarrow aaBaBbS \text{ (Using } B \rightarrow bS \text{)}$$

$$\rightarrow aaBaBbbA \text{ (Using } S \rightarrow bA \text{)}$$

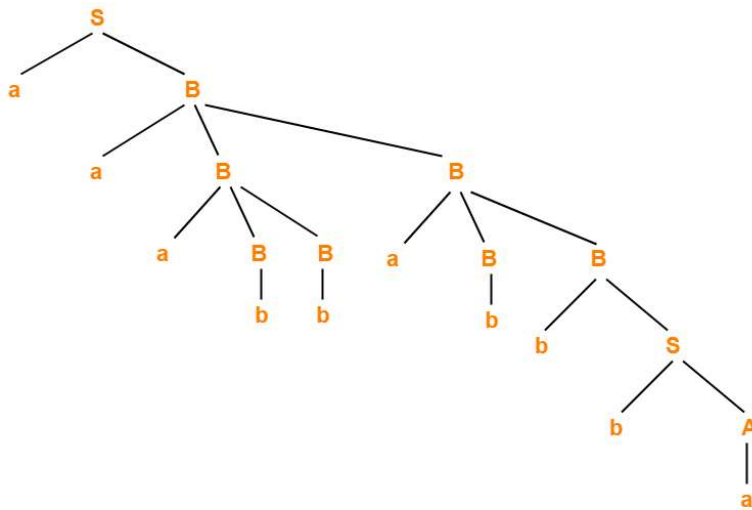
$$\rightarrow aaBaBbbba \text{ (Using } A \rightarrow a \text{)}$$

$$\rightarrow aaBabbba \text{ (Using } B \rightarrow b \text{)}$$

$$\rightarrow aaaBabbba \text{ (Using } B \rightarrow aBB \text{)}$$

$$\rightarrow aaaBbabba \text{ (Using } B \rightarrow b \text{)}$$

$$\rightarrow aaabbabbba \text{ (Using } B \rightarrow b \text{)}$$



### Rightmost Derivation Tree

## NOTES

- For unambiguous grammars, Leftmost derivation and Rightmost derivation represents the same parse tree.
- For ambiguous grammars, Leftmost derivation and Rightmost derivation represents different parse trees.

Here,

- The given grammar was unambiguous.
- That is why, leftmost derivation and rightmost derivation represents the same parse tree.

**Leftmost Derivation Tree = Rightmost Derivation Tree**

**Also Read- [Ambiguous Grammar](#)**

## Properties Of Parse Tree-