

Chapter 2

Boolean Algebra and Logic Gates

Definitions

Theorems

Functions

Canonical and Standard Forms

Operations

Gates

Integrated Circuits

Canonical Forms

- A **canonical form** is a standard method for representing Boolean functions.
- The two canonical forms that are used are:
 - Sum of Minterms
 - Product of Maxterms
- These forms are sometimes considered the “brute force” method of representing functions as they seldom represent a function in a minimized form.

Minterms

- Any given binary variable can be represented in two forms:
 - x , its normal form, and
 - x' , its complement
- If we consider two binary variables and the AND operation, there are four combinations of the variables:
 - xy
 - xy'
 - $x'y$
 - $x'y'$
- Each of the above four AND terms is called a **minterm** or a **standard product**.
- n variables can be combined to form 2^n minterms.

Minterms Expressed

				Minterms	
x	y	z		Term	Designation
0	0	0		$x' y' z'$	m_0
0	0	1		$x' y' z$	m_1
0	1	0		$x' y z'$	m_2
0	1	1		$x' y z$	m_3
1	0	0		$x y' z'$	m_4
1	0	1		$x y' z$	m_5
1	1	0		$x y z'$	m_6
1	1	1		$x y z$	m_7

Maxterms

- Any given binary variable can be represented in two forms:
 - x , its normal form, and
 - x' , its complement
- If we consider two binary variables and the OR operation, there are four combinations of the variables:
 - $x + y$
 - $x + y'$
 - $x' + y$
 - $x' + y'$
- Each of the above four OR terms is called a **maxterm** or a **standard sum**.
- n variables can be combined to form 2^n maxterms.
- Each maxterm is the complement of its corresponding minterm and vice-versa.

Maxterms Expressed

				Maxterms	
x	y	z		Term	Designation
0	0	0		$x + y + z$	M_0
0	0	1		$x + y + z'$	M_1
0	1	0		$x + y' + z$	M_2
0	1	1		$x + y' + z'$	M_3
1	0	0		$x' + y + z$	M_4
1	0	1		$x' + y + z'$	M_5
1	1	0		$x' + y' + z$	M_6
1	1	1		$x' + y' + z'$	M_7

Truth Table to Expression (Sum of Minterms)

- Any Boolean function can be expressed as a **sum of minterms** or **sum of products** (i.e. the ORing of terms).
 - You can form the function algebraically by forming a **minterm** for each combination of the variables that produces a **1** in the function. (Each row with output of **1** becomes a **product term**)
 - Sum (OR)** product terms together.

x	y	z	G	
0	0	0	0	
0	0	1	0	
0	1	0	0	
0	1	1	1	└──────────┘
1	0	0	0	
1	0	1	0	
1	1	0	1	└────────┘
1	1	1	1	└──┘

$xyz + xyz' + x'yz$

Sum of Minterms Example

x	y	z		Function F ₁	Required Minterms
0	0	0		1	$x' y' z'$
0	0	1		0	
0	1	0		0	
0	1	1		1	$x' y z$
1	0	0		1	$x y' z'$
1	0	1		0	
1	1	0		0	
1	1	1		0	

$$F_1 = x' y' z' + x' y z + x y' z'$$

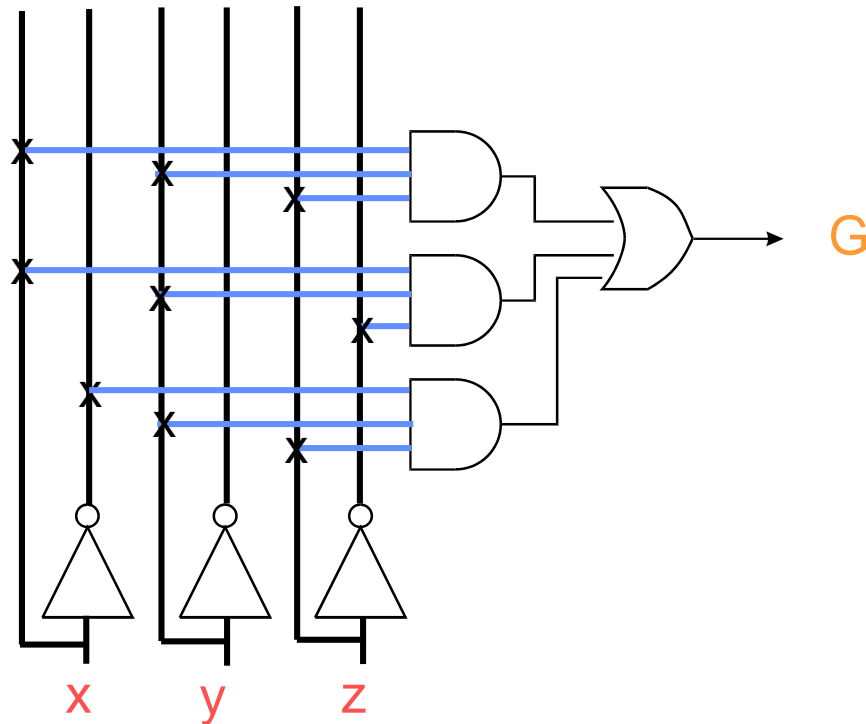
$$= m_0 + m_3 + m_4$$

$$= \sum(0, 3, 4)$$

Equivalent Representations of Circuits

- All three formats are equivalent
- Number of 1's in truth table output column equals AND terms for Sum-of-Products (SOP)

x	y	z	G
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	0
1	1	0	1
1	1	1	1



$$G = xyz + xyz' + x'yz$$

Truth Table to Expression (Product of Maxterms)

- Any Boolean function can be expressed as a **product of maxterms** or **product of sums** (i.e. the ANDing of terms).
 - You can form the function algebraically by forming a **maxterm** for each combination of the variables that produces a **0** in the function. (Each row with output of **0** becomes a **standard sums**)
 - **AND** these maxterms together.

Product of Maxterms Example

x	y	z		Function F₁	Required Maxterms
0	0	0		1	
0	0	1		0	$x + y + z'$
0	1	0		0	$x + y' + z$
0	1	1		1	
1	0	0		1	
1	0	1		0	$x' + y + z'$
1	1	0		0	$x' + y' + z$
1	1	1		0	$x' + y' + z'$

$$\begin{aligned}F_1 &= (x + y + z')(x + y' + z)(x' + y + z')(x' + y' + z)(x' + y' + z') \\&= M_1 M_2 M_5 M_6 M_7 \\&= \prod(1, 2, 5, 6, 7)\end{aligned}$$

Minterms and Maxterms

- Each variable in a Boolean expression is a **literal**
- Boolean variables can appear in normal (**x**) or complement form (**x'**)
- Each AND combination of terms is a **minterm**
- Each OR combination of terms is a **maxterm**
- Example:

Minterms

x	y	z	Minterm	
0	0	0	$x'y'z'$	m_0
0	0	1	$x'y'z$	m_1
...				
1	0	0	$xy'z'$	m_4
...				
1	1	1	xyz	m_7

Maxterms

x	y	z	Maxterm	
0	0	0	$x+y+z$	M_0
0	0	1	$x+y+z'$	M_1
...				
1	0	0	$x'+y+z$	M_4
...				
1	1	1	$x'+y'+z'$	M_7

Obtaining Sum of Minterms Form

A	B	C	$F = A'B + B' + C$	Required Minterms	Required Designations
0	0	0	1	$A'B'C'$	m_0
0	0	1	1	$A'B'C$	m_1
0	1	0	1	$A'BC'$	m_2
0	1	1	1	$A'BC$	m_3
1	0	0	1	$AB'C'$	m_4
1	0	1	1	$AB'C$	m_5
1	1	0	0		
1	1	1	1	ABC	m_7

$$F = A'B'C' + A'B'C + A'BC' + A'BC + AB'C' + AB'C + ABC$$

$$= m_0 + m_1 + m_2 + m_3 + m_4 + m_5 + m_7$$

$$F(A, B, C) = \sum(0, 1, 2, 3, 4, 5, 7)$$

Obtaining Product of Maxterms

A	B	C		$F = A'B + B'C$	Required Maxterms	Required Designations
0	0	0		0	$A + B + C$	M_0
0	0	1		1		
0	1	0		1		
0	1	1		1		
1	0	0		0	$A' + B + C$	M_4
1	0	1		1		
1	1	0		0	$A' + B' + C$	M_6
1	1	1		0	$A' + B' + C'$	M_7

$$F = (A + B + C)(A' + B + C)(A' + B' + C)(A' + B' + C')$$

$$= M_0 \cdot M_4 \cdot M_6 \cdot M_7$$

$$F(A, B, C) = \prod(0, 4, 6, 7)$$

Canonical Form Conversion

- A function represented as Sum of **minterms** can be represented as the Product of **maxterms** of the remaining terms.
- The complement of a function expressed in sum of minterms equals the sum of minterms missing from the original function
 - $F(A, B, C) = \sum(0, 3, 4) = m_0 + m_3 + m_4$
 - $F'(A, B, C) = \sum(1, 2, 5, 6, 7) = m_1 + m_2 + m_5 + m_6 + m_7$
- Now if we take the complement of F' using DeMorgan's theorem, we obtain F in the product of maxterms form:
 - $(F')' = (m_1 + m_2 + m_5 + m_6 + m_7)'$
 - $F = m_1' \cdot m_2' \cdot m_5' \cdot m_6' \cdot m_7'$ [Complement of minterms]
 - $= M_1 M_2 M_5 M_6 M_7$ [maxterms]
 - $= \prod(1, 2, 5, 6, 7)$
- This implies the following relation:
 - **$m'j = Mj$**
- So sum of minterms: $\sum(0, 3, 4) =$ product of maxterms: $\prod(1, 2, 5, 6, 7)$