

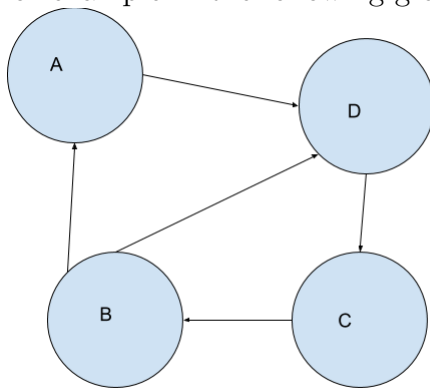
Graph Theory

What is a graph

A graph G is defined by a combination of a vertex set V and an edge set $E \subseteq V \times V$. It is common to see the notation

$$G = (V, E)$$

For example in the following graph



We will mathematically write $G = (V, E)$ where
 $V = \{A, B, C, D\}$ and
 $E = \{(A, D), (D, C), (C, B), (B, A), (B, D)\}$

Terminology

The two vertices that are part of an edge are its endpoints.

When a vertex is an endpoint of an edge we say the edge is incident on the vertex.

Note that two vertices connected by an edge is describing a relation (recall what that means from a previous part of the course).

A graph can either be directed meaning that there can be an edge from vertex u to vertex v without having an edge from vertex v to vertex u or undirected meaning that $(u, v) \in E \leftrightarrow (v, u) \in E$.

If we view this in terms of relations, an undirected graph is describing a symmetric relationship.

Loop(or sometimes self-loop) - an edge that joins a vertex to itself.

Complete graph - a graph where every vertex is connected to every other vertex.

A complete graph on n vertices is denoted as K_n .

Question

How many edges does a complete graph have?

That is the same as $\binom{n}{2}$ since for any pair of distinct vertices they have to be connected by an edge.

Paths and walks

For the remainder of the course we will be focusing on the following type of graph

- undirected
- no self loops
- no parallel edges
- finite number of edges

whenever we have to deviate from this type of graph it will be made clear.

Path - A path is an alternating sequence of vertices and edges that satisfy the following

1. it starts and ends with a vertex.
2. each edge joins the vertex before it in the sequence to the vertex after it in the sequence.
3. no vertex appears more than once.

A **walk** satisfies the first two conditions but not the third

Result

If there is a walk between two vertices u and v , there has to be a path between vertices u and v .

Proof

Consider any walk between vertices u and v . There are two cases

Case 1 - There is no vertex that is repeated in the walk. In this case then by definition this walk is a path and we are done.

Case 2 - There exists a vertex x in this walk such that it is repeated. For this vertex, we can make a shorter walk by removing the part of the walk that is between the first and last occurrences of the the vertex. That is, include the last occurrence of x but not the first. Repeat the process of finding these repeated vertices and making shorter walks. Note that in each case we are eliminating one of the vertices from the list of repeated vertices. Therefore eventually we will have no repeated vertices in our walk and our walk will reduce to a path. (draw a picture to complete this proof).

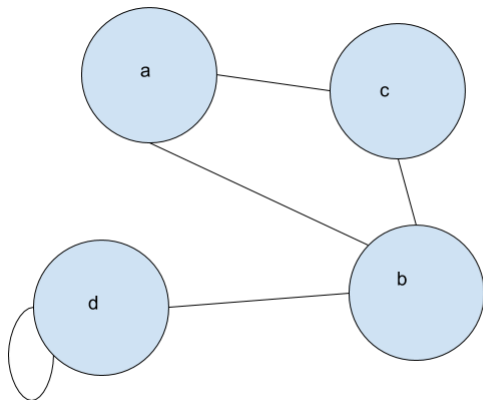
The length of a path is the number of edges in the path.

There can be multiple paths between two vertices. An algorithms course is likely to spend a fair number of lectures talking about different methods to find shortest paths in graphs.

Degrees

The degree of a vertex is the number of edges that are incident on it.

Note that if you have a self-loop, it counts twice to the degree of the vertex. We are allowing self loops in the discussion of degrees.



In the above example the degree of vertex d, denoted by $\deg(d) = 3$. $\deg(a) = 2$ etc.

Theorem 1. *Sum of the degrees of the vertices is twice the number of edges.*

Proof

Proof by induction on the number of edges. Consider that we are given a graph with e edges and we want to prove the result holds for any e .

Base case: If the number of edges is 0 then the sum of the degrees is 0. Twice the number of edges is also 0. Done.

Induction hypothesis: Suppose that for $0 < k \leq e$ the result holds.

Now consider a graph with $e > 0$ edges. Let e_1 be some edge in this graph. Consider what happens when we delete this edge (note that deleting an edge means just deleting the edge and not deleting the associated vertices). We get a graph G' which has 1 less edge. So we can apply the induction hypothesis to it.

Applying the induction hypothesis says that the sum of degrees of the vertices in G' is $2(e - 1)$.

Now let us consider 2 cases for the edge that we removed

Case 1: the removed edge was a self loop. When we add this edge back in the degree of one single vertex will go up by 2. So the sum of the degrees will be $2(e - 1) + 2 = 2e$.

Case 2: the removed edge joined two distinct vertices. When we add this edge back the degree of two vertices goes up by 1 each. Again the sum of the degrees will be $2e$.

Combining this with the base case of 0 edges, we have a complete proof by induction. QED!