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1.1 bash config

```
#copy first argument to clipborad ! ONLY WORK ON
     XORG !
alias clip="xclip —sel clip"
# compile the $1 parameter, if a $2 is provided
# the name will be the the binary output, if
# none is provided the binary name will be
  'a.out
comp() {
  echo ">> COMPILING $1 <<" 1>&2
  if [ $# -gt 1 ]; then
     outfile="${2}
   else
     outfile="a.out"
  time g++ -std=c++20 \
     -02 \
-g3 \
     -Wall
     —fsanitize=address,undefined \
     -fno-sanitize-recover \
-D LOCAL \
-0 "${outfile}" \
"$1"
  if [ $? -ne 0 ]; then
  echo ">> FAILED <<" 1>&2
return 1
  echo ">> DONE << " 1>&2
# run the binary given in $1, if none is
# given it will try to run the 'a.out'
# binary
run() {
  to_run=./a.out
if [ -n "$1" ]; then
  to_run="$1"
   time $to_run
  just comp and run your cpp file
# accpets <in1 >out and everything else
comprun() {
   comp "$1" "a" && run ./a ${@:2}
testall() {
  comp "$1" generator
  comp "$2" brute
  comp "$3" main
  input_counter=1
  while true; do
     echo "$input_counter"
     run ./generator >input
     run ./main <input >main_output.txt
     run ./brute <input >brute_output.txt
     diff brute_output.txt main_output.txt
     if [ $? —ne 0 ]; then
  echo "Outputs differ at input $input_counter"
        echo "Brute file output:"
       cat brute_output.txt
echo "Main file output:"
        cat main_output.txt
        echo "input used:
        cat input
        break
     ((input_counter++))
  done
}
touch_macro() {
  cat "$1"/template.cpp >"$2"
cat "$1"/run.cpp >>"$2"
  cp "$1"/debug.cpp .
# Creates a contest with hame $2
# Copies the macro and debug file from $1
# Already creates files a...z .cpp and .py
prepare_contest() {
```

```
mkdir "$2"
  cd "$2"
  for i in {a..z}; do
    touch_macro $1 $i.cpp
  done
}
```

1.2 debug

```
template <typename T>
concept Printable = requires(T t) {
  {
     std::cout << t
  } -> std::same_as<std::ostream &>;
template <Printable T>
void __print(const T &x) {
  cerr << x;</pre>
template <size_t T>
void __print(const bitset<T> &x) {
  cerr << x;</pre>
template <typename A, typename B>
void __print(const pair<A, B> &p);
template <typename... A>
void __print(const tuple<A...> &t);
template <typename T>
void __print(stack<T> s);
template <typename T>
void __print(queue<T> q);
template <typename T, typename... U>
void __print(priority_queue<T, U...> q);
template <typename A>
void __print(const A &x) {
  bool first = true;
  cerr << '{';</pre>
  for (const auto &i : x) {
  cerr << (first ? "" : ","), __print(i);</pre>
     first = false;
  cerr << '}';
template <typename A, typename B>
void \_\_print(const pair<A, B> &p) {
  cerr << '(
   __print(p.first);
  cerr <<
  __print(p.second);
cerr << ')';</pre>
template <typename... A>
void __print(const tuple<A...> &t) {
  bool first = true;
cerr << '(';</pre>
  apply(
       [&first](const auto &...args) {
    ((cerr << (first ? "" : ","),
           __print(args), first = false),
       },
t);
  cerr << ')';
template <typename T>
void __print(stack<T> s) {
  vector<T> debugVector;
  while (!s.empt\bar{y}()) {
     T t = s.top();
     debugVector.push_back(t);
     s.pop();
  reverse(debugVector.begin(), debugVector.end());
  __print(debugVector);
template <typename T>
void __print(queue<T> q) {
  vector<T> debugVector;
  while (!q.empty()) {
     T t = q.front();
     debugVector.push_back(t);
     q.pop();
```

```
for (common_type_t<decltype(a), decltype(b)> \
   _print(debugVector);
template <typename T, typename... U>
void __print(priority_queue<T, U...> q) {
  vector<T> debugVector;
                                                                #define pb push_back
  while (!q.empty()) {
                                                                #define pf push_front
     T t = q.top();
     debugVector.push_back(t);
                                                                #define ppb pop_back
                                                                #define ppf pop_front
#define eb emplace_back
     q.pop();
                                                                #define lb lower_bound
  __print(debugVector);
                                                                #define ub upper_bound
                                                                #define fi first
void _print() { cerr << "]\n"; }
template <typename Head, typename...</pre>
                                                                #define se second
#define emp emplace
                                           Tail>
void _print(const Head &H, const Tail &...T) {
                                                                #define ins insert
    _{	extsf{print}}(\mathsf{H});
  if (sizeof...(T)) cerr << ", ";
                                                                using str = string;
  _{\mathsf{print}}(\mathsf{T}\dots);
                                                                using ll = long long;
                                                                using ull = unsigned long long;
#define dbg(x...) \
cerr << "[" << #x << "] = ["; \
                                                                using ld = long double;
                                                                using vll = vector<ll>
  _print(x)
                                                                using pll = pair<ll, ll>
                                                                using vll2d = vector<vll>;
                                                                using vi = vector<int>;
                                                                using vi2d = vector<vi>
1.3 run
                                                                using pii = pair<int, int>;
                                                                using vpii = vector<pii>;
void run();
                                                                using vc = vector<char>;
int32_t main() {
                                                                 using vs = vector<str>
#ifndef LOCAL
  fastio;
#endif
                                                                template <typename T>
  int T = 1;
                                                                using pqmn =
  cin >> T;
  rep(t, 0, T) {
                                                                template <typename T>
     dbg(t);
     run();
                                                                   return (a < b ? a = b, 1 : 0);
void run() {}
                                                                   return (a > b ? a = b, 1 : 0);
1.4 short-template
#include <bits/stdc++.h>
using namespace std;
#define fastio
                                                                1.6 vim config
  ios_base::sync_with_stdio(0); \
  cin.tie(0);
                                                                set sta nu rnu sc cindent
set ts=2 sw=2
void run() {}
int32_t main(void) {
  fastio;
                                                                     timeoutlen=100
  int t;
                                                                 colorscheme default
                                                                syntax on
  // cin >> t;
while (t—) run();
                                                                 " in the vim clipboard
                                                                     register
1.5 template
                                                                     normal! gvy
                                                                     with the yanked text
#include <bits/stdc++.h>
using namespace std;
#ifdef LOCAL
#include "debug.cpp"
                                                                        | sha256sum')
#else
                                                                     let @" = l:hash
#define dbg(...)
                                                                endfunction
#endif
#define endl '\n'
#define fastio
                                                                    Data Structures
  ios_base::sync_with_stdio(0); \
cin.tie(0);
#define int long long
#define all(j) j.begin(), j.end()
#define rall(j) j.rbegin(), j.rend()
                                                                2.1
                                                                      SQRT decomposition
                                                                2.1.1 two-sequence-queries
```

```
i = (a);
i > (b); i—)
#define trav(xi, xs) for (auto &xi : xs)
#define rtrav(xi, xs) \
for (auto &xi : ranges::views::reverse(xs))
#define divc(a, b) ((a) + (b)-1ll) / (b)
template <typename T, typename T2>
using umap = unordered_map<T, T2>;
    priority_queue<T, vector<T>, greater<T>>;
using pqmx = priority_queue<T, vector<T>>;
template <typename_T, typename U>
inline bool chmax(T &a, U const &b) {
template <typename T, typename U>
inline bool chmin(T &a, U const &b) {
/*===========*/
set bg=dark ruler clipboard=unnamed,unnamedplus,
" Takes the hash of the selected text and put
function! HashSelectedText()
    " Yank the selected text to the unnamed
     " Use the system() function to call sha256sum
     let l:hash = system('echo ' . shellescape(@@) .
    " \dot{Y} ank the hash into \dot{V} im's unnamed register
using ll = long long;
const ll MOD = 998244353;
inline ll sum(const ll a, const ll b) {
  return (a + b) % MOD;
```

}

#define len(j) (int)j.size()
#define rep(i, a, b)

i = (a); i < (b); i++

#define rrep(i, a, b)

for (common_type_t<decltype(a), decltype(b)>

```
ll sub(const ll a, const ll b) {
  return (a - b + MOD) % MOD;
inline ll mul(const ll a, const ll b) {
  return (a * b) % MOD;
struct SqrtDecomposition {
  struct t_sqrt {
     int l, r;
    ll x, y;
    ll prod;
    ll sum_as, sum_bs;
    t_sqrt() {
       l = numeric_limits<int>::max();
       r = numeric_limits<int>::min();
       x = y = prod = sum_as = sum_bs = 0;
  };
  int sqrtLen;
  vector<t_sqrt> blocks;
  vector<ll> as, bs
  SqrtDecomposition(const vector<ll> &as_,
                        const vector<ll> &bs_) {
    int n = as_.size();
    sqrtLen = (int) sqrt(n + .0) + 1;
    blocks.resize(sqrtLen + 6.66);
    as = as_-;
    bs = bs_{-};
    for (int i = 0; i < n; i++) {
  auto &bi = blocks[i / sqrtLen];</pre>
       bi.l = min(bi.l, i);
                           i);
       bi.r = max(bi.r,
       bi.sum_as = sum(bi.sum_as, as[i]);
       bi.sum_bs = sum(bi.sum_bs, bs[i]);
bi.prod = sum(bi.prod, mul(as[i], bs[i]));
  }
  // adds x to a[i], and y to b[i], in range [l,
  auto &block = blocks[idx / sqrtLen];
       block.prod =
            sub(block.prod, mul(as[idx], bs[idx]));
       block.sum_as = sub(block.sum_as, as[idx]);
block.sum_bs = sub(block.sum_bs, bs[idx]);
       as[idx] = sum(as[idx], x);
bs[idx] = sum(bs[idx], y);
       block.prod =
            sum(block.prod, as[idx] * bs[idx]);
       block.sum_as = sum(block.sum_as, as[idx]);
block.sum_bs = sum(block.sum_bs, bs[idx]);
    };
    auto apply2 = [\&](int idx, ll x,
                          ll y) \rightarrow void {
       blocks[idx].x = sum(blocks[idx].x, x);
blocks[idx].y = sum(blocks[idx].y, y);
    int cl = l / sqrtLen, cr = r / sqrtLen;
    if (cl == cr) {
  for (int i = l; i <= r; i++) {</pre>
         apply1(i, x, y);
    } else {
  for (int i = l; i <= (cl + 1) * sqrtLen - 1;</pre>
             i++) {
         apply1(i, x, y);
       for (int i = cl + 1; i \le cr - 1; i++) {
         apply2(i, x, y);
       for (int i = cr * sqrtLen; i <= r; i++) {
         apply1(i, x, y);
  // sum of a[i]*b[i] in range [l r]
```

```
ll query(int l, int r) {
  auto eval1 = [&](int idx) -> ll {
       auto &block = blocks[idx / sqrtLen];
       auto eval2 = [\&](int idx) \rightarrow ll {
       auto &block = blocks[idx];
       ll ret = 0;
       ret =
           sum(ret
                mul(mul(block.x, block.y),
    sum(sub(block.r, block.l), 1)));
       ret = sum(ret, block.prod);
       ret = sum(ret, block.y * block.sum_as);
ret = sum(ret, block.x * block.sum_bs);
       return ret;
    };
    ll ret = 0;
    int cl = l / sqrtLen, cr = r / sqrtLen;
       (cl == cr) {
for (int i = l; i <= r; i++) {
         ret = sum(ret, eval1(i));
    } else {
       for (int i = l; i \le (cl + 1) * sqrtLen - 1;
             i++) {
         ret = sum(eval1(i), ret);
       }
       for (int i = cl + 1; i \le cr - 1; i++) { ret = sum(ret, eval2(i));
       for (int i = cr * sqrtLen; i <= r; i++) {
         ret = sum(ret, eval1(i));
    }
    return ret;
  }
};
2.2
     Segment tree (dynamic)
     Range Max Query Point Max Assignment
```

Description: Answers range queries in ranges until 10^9 (maybe more) **Time**: Query and update $O(n \cdot \log n)$

```
struct node
node *newNode();
struct node {
  node *left, *right;
  int lv, rv;
  node() : left(NULL), right(NULL), val(—oo) {}
  inline void init(int l, int r) {
    lv = l;
rv = r;
  inline void extend() {
    if (!left) {
  int m = (lv + rv) / 2;
       left = newNode()
      right = newNode();
      left->init(lv, m);
      right—>init(m + 1, rv);
  ll query(int l, int r)
    if (r < lv \mid | rv < l) {
      return 0;
    if (l <= lv && rv <= r) {
      return val;
    extend():
    return max(left->query(l, r),
                right—>query(l, r));
  void update(int p, ll newVal) {
    if (lv == rv)
      val = max(val, newVal);
```

```
return;
    }
    extend()
    (p <= left->rv ? left : right)
        ->update(p, newVal);
    val = max(left->val, right->val);
};
const int BUFFSZ(1e7);
node *newNode()
  static int bufSize = BUFFSZ
  static node buf[(int)BUFFSZ];
assert(bufSize);
  return &buf[—bufSize];
struct SegTree {
  int n;
  node *root;
  SegTree(int _n) : n(_n) {
    root = newNode();
    root->init(0, n);
  fl query(int l, int r) {
    return root—>query(l, r);
  void update(int p, ll v) { root->update(p, v); }
2.2.2 Range Sum Query Point Sum Update
Description: Answers range queries in ranges until 10<sup>9</sup> (maybe more)
Time: Query and update in O(n \cdot \log n)
struct node;
node *newNode();
struct node {
  node *left, *right;
  int lv, rv;
  ll val;
  node() : left(NULL), right(NULL), val(0) {}
  inline void init(int l, int r) {
  inline void extend() {
  if (!left) {
      int m = (rv -
                    – lv) / 2 + lv;
      left = newNode();
      right = newNode();
      left->init(lv, m);
      right—>init(m + 1, rv);
    }
  }
  ll query(int l, int r) {
    if (r < lv || rv < l) {
      return 0;
    if (l <= lv && rv <= r) {
      return val;
    extend();
return left->query(l, r) + right->query(l, r);
  void update(int p, ll newVal) {
    if (lv == rv)
      val += newVal;
      return;
    extend()
    (p <= left->rv ? left : right)
        ->update(p, newVal);
    val = left->val + right->val;
const int BUFFSZ(1.3e7);
node *newNode()
  static int bufSize = BUFFSZ;
static node buf[(int)BUFFSZ];
  // assert(bufSize)
  return &buf[—bufSize];
```

```
struct SegTree {
  int n;
  node *root;
  SegTree(int _n) : n(_n) {
    root = newNode();
    root->init(0, n);
  ll query(int l, int r) {
    return root—>query(l, r);
  void update(int p, ll v) { root->update(p, v); }
};
2.3
     Segment tree point update range query
2.3.1 Query GCD (bottom up)
using ll = long long;
struct Node {
  ll value
  bool undef;
  Node()
         value(1), undef(1){}; // Neutral element
  Node(ll v) : value(v), undef(0){};
};
inline Node combine(const Node &nl,
                      const Node &nr) {
  if (nl.undef) return nr
  if (nr.undef) return nl;
  Node m;
  m.value = gcd(nl.value, nr.value);
  m.undef = false;
  return m;
template <typename T = Node, auto F = combine>
struct SegTree {
  int n;
  vector<T> st;
  SegTree(int _n) : n(_n), st(n \ll 1) {}
  void assign(int p, const T &k) {
  for (st[p += n] = k; p >>= 1;)
    st[p] = F(st[p << 1], st[p << 1 | 1]);</pre>
  T query(int l, int r) {
    if (l & 1) ansl = F(ansl, st[l++]);
if (r & 1) ansr = F(st[—r], ansr);
    return F(ansl, ansr);
};
2.3.2 Query hash (top down)
const ll MOD = 1'000'000'009;
const | | P = 31;
const int MAXN = 2'000'000;
ll\ pows[MAXN + 1];
void computepows() {
  pows[0] = 1;
  for (int i = 1; i <= MAXN; i++)
    pows[i] = (pows[i - 1] \star P) \% MOD;
}
struct Node {
  ll hash;
Node() : hash(-1){}; // Neutral element
  Node(ll\ v) : hash(v){};
inline Node combine(Node &vl, Node &vr, int nl,
  int nr, int ql, int qr) {
if (vl.hash == -1) return vr;
  \overline{\mathsf{if}} (vr.hash == -1) return vl;
  Node vm;
  int nm = midpoint(nl, nr);
  int lsize = min(nm, qr) - max(nl, ql) + 1;
  vm.hash = (vl.hash +
```

```
((vr.hash * pows[lsize]) % MOD)) %
                                                            Node m:
            MOD;
  return vm;
template <typename T = Node, auto F = combine>
struct SegTree {
  int n;
                                                             int n;
  vector<T> st;
  SegTree(int n) : n(n), st(n << 2) {}
  void assign(int p, const T &v) {
    assign(1, 0, n - 1, p, v);
  void assign(int node, int l, int r, int p,
               const T &v) {
    if (l == r) {
      st[node] = v;
      return;
    int m = midpoint(l, r);
    if (p \ll m)
      assign(node << 1, l, m, p, v);
                                                          };
      assign(node << 1 | 1, m + 1, r, p, v);
    st[node] = F(st[node << 1], st[node << 1 | 1],
                  l, r, l, r);
  }
  inline T query(int l, int r) {
    return query(1, 0, n - 1, l, r);
  inline T query(int node, int nl, int nr, int l,
       int r) const {
(r < nl or nr < l) return T();</pre>
    if (l <= nl and nr <= r) return st[node];</pre>
    int m = midpoint(nl, nr);
                                                             Node m:
    auto a = query(node \ll 1, nl, m, l, r);
        query(node << 1 | 1, m + 1, nr, l, r);
    return F(a, b, nl, nr, l, r);
                                                             int n;
2.3.3 Query max subarray sum (bottom up)
struct Node {
  ll tot, suf, pref, best;
  // Neutral element
  Node()
      : tot(-oo),
suf(-oo),
        pref(-oo),
        best(-oo) {} // Neutral element
  // for assign
  Node(ll x) {
    tot = x, suf = x, pref = x,
                                                          };
    best = max(Oll, x);
};
                                                          2.4
Node combine(Node &nl, Node &nr) {
  if (nl.tot == -oo) return nr;
if (nr.tot == -oo) return nl;
  Node m;
  m.tot = nl.tot + nr.tot;
  m.pref = max({nl.pref, nl.tot + nr.pref});
  m.suf = max({nr.suf, nr.tot + nl.suf});
  m.best =
      max({nl.best, nr.best, nl.suf + nr.pref});
  return m;
2.3.4 Query min (bottom up)
                                                               Data()
struct Node {
  ll value:
  Node()
      : value(numeric_limits<
```

ll>::max()){}; // Neutral element

Node(ll v) : value(v){};

Node combine(Node &l, Node &r) {

```
m.value = min(l.value, r.value);
  return m:
template <typename T = Node, auto F = combine>
struct SegTree {
  vector<T> st;
  SegTree(int \underline{n}) : n(\underline{n}), st(n << 1) {}
  void assign(int p, const T &k)
     for (st[p += n] = k; p >= 1;)
       st[p] = F(st[p << 1], st[p << 1 | 1]);
  for (l += n, r += n + 1; l < r;
l >>= 1, r >>= 1) {
if (l & 1) ansl = F(ansl, st[l++]);
       if (r \& 1) ansr = F(st[-r], ansr);
     return F(ansl, ansr);
2.3.5 Query sum (bottom up)
struct Node {
  ll value;
  Node() : value(0){}; // Neutral element
  Node(ll v) : value(v){};
inline Node combine(const Node &nl
                        const Node &nr) {
  m.value = nl.value + nr.value;
  return m;
struct SegTree {
  vector<Node> st;
SegTree(int _n) : n(_n), st(n << 1) {}</pre>
  void assign(int p, const Node &k) {
  for (st[p += n] = k; p >>= 1;)
       st[p] = combine(st[p << 1], st[p << 1 | 1]);
  Node query(int l, int r) {
    Node ansl = Node(), ansr = Node();
     for (l += n, r += n + 1; l < r;
           l >>= 1, r >>= 1)
       if (l & 1) ansl = combine(ansl, st[l++]);
       if (r & 1) ansr = combine(st[—r], ansr);
     return combine(ansl, ansr);
    Segment tree range update range query
2.4.1 Arithmetic progression sum update, query sum
Description: Makes arithmetic progression updates in range and sum queries.
Usage: Considering PA(A, R) = [A + R, A + 2R, A + 3R, ...]
• update_set(l, r, A, R): sets [l, r] to PA(A, R)
• update_add(l, r, A, R): sum PA(A, R) in [l, r]
   • query(l, r): sum in range [l, r]
Time: build O(N), updates and queries O(log N)
const ll oo = 1e18;
struct SegTree {
  struct Ďata {
     ll sum;
     ll set_a, set_r, add_a, add_r;
         : sum(0),
            set_a(oo),
            set_r(0),
            add_a(0),
            add_r(0) {}
  int n;
vector<Data> seg;
  SegTree(int n_)
```

```
: n(n_{-}), seg(vector<Data>(4 * n)) {}
                                                                  ĺl query(int a, int b, int p, int l, int r) {
void prop(int p, int l, int r) {
  int sz = r - l + 1;
  ll &sum = seg[p].sum, &set_a = seg[p].set_a,
                                                                    prop(p, l, r);
if (b < l or r < a) return 0;
if (a <= l and r <= b) return seg[p].sum;</pre>
      &set_r = seg[p].set_r,
                                                                    add_a = seg[p].add_a,
      add_r = seg[p].add_r;
  if (set_a != oo) {
  set_a += add_a, set_r += add_r;
                                                                  il query(int l, int r) {
  return query(l, r, 1, 0, n - 1);
     sum =
         set_a * sz + set_r * sz * (sz + 1) / 2;
    if (l != r) {
 int m = (l + r) / 2;
                                                               };
       seg[2 * p].set_a = set_a;
       seg[2 * p].set_r = set_r;
                                                               2.4.2 Increment update query min & max (bottom up)
       seg[2 * p].add_a = seg[2 * p].add_r = 0;
       seq[2 * p + 1].set_a =
                                                               using SeqT = ll;
       set_a + set_r * (m - l + 1);
seg[2 * p + 1].set_r = set_r;
                                                               struct QueryT {
                                                                  SegT mx, mn;
       seg[2 * p + 1].add_a =
                                                                  QueryT()
            seg[2 * p + 1].add_r = 0;
                                                                       : mx(numeric_limits<SegT>::min()),
  set_a = oo, set_r = 0;
add_a = add_r = 0;
else if (add_a or add_r) {
                                                                         mn(numeric_limits<SegT>::max()) {}
                                                                  QueryT(SegT _v) : mx(_v), mn(_v) {}
                                                               };
     sum +=
                                                               inline QueryT combine(QueryT ln, QueryT rn,
    add_a * sz + add_r * sz * (sz + 1) / 2;
if (l != r) {
int m = (l + r) / 2;
                                                                                          pii lr1, pii lr2) {
                                                                  chmax(ln.mx, rn.mx);
chmin(ln.mn, rn.mn);
       seg[2 * p].add_a += add_a;
                                                                  return ln;
       seg[2 * p].add_r += add_r;
                                                               }
       seq[2 * p + 1].add_a +=
                                                               using LazyT = SeqT;
       add_a + add_r * (m - l + 1);

seg[2 * p + 1].add_r += add_r;
                                                               inline QueryT applyLazyInQuery(QueryT q, LazyT l,
                                                                                                     pii ĺr)
                                                                  if (q.mx == QueryT().mx) q.mx = SegT();
     add_a = add_r = 0;
                                                                  if (q.mn == QueryT().mn) q.mn = SegT();
                                                                  q.mx += l, q.mn += l;
}
                                                                  return q;
int inter(pii a, pii b) {
  if (a.first > b.first) swap(a, b);
                                                               }
                                                               inline LazyT applyLazyInLazy(LazyT a, LazyT b) {
  return max(
                                                                  return a + b;
       0, min(a.second, b.second) - b.first + 1);
                                                               using UpdateT = SegT;
ll set(int a, int b, ll aa, ll rr, int p, int l,
  int r) {
prop(p, l, r);
if (b < l or r < a) return seg[p].sum;</pre>
                                                               inline QueryT applyUpdateInQuery(QueryT q,
                                                                                                        UpdateT u,
                                                                                                       pii lr) {
  if (a \le l \text{ and } r \le b) {
                                                                  if (q.mx == QueryT().mx) q.mx = SegT();
     seg[p].set_a = aa;
                                                                  if (q.mn == QueryT().mn) q.mn = SegT();
q.mx += u, q.mn += u;
     seg[p].set_r = rr;
     prop(p, l, r);
                                                                  return q;
     return seg[p].sum;
                                                               }
                                                               inline LazyT applyUpdateInLazy(LazyT l, UpdateT u,
  int m = (l + r) / 2;
int tam_l = inter({l, m}, {a, b});
                                                                  return l + u;
  return seg[p].sum =
                                                               }
                set(a, b, aa, rr, 2 * p, l, m) + set(a, b, aa + rr * tam_l, rr,
                                                               template <typename Qt = QueryT,
                                                                            typename Lt = LazyT
                     2 * p + 1, m + 1, r);
                                                                            typename Ut = UpdateT, auto C = combine,
                                                                           auto ALQ = applyLazyInQuery,
void update_set(int l, int r, ll aa, ll rr) {
                                                                            auto ALL = applyLazyInLazy,
  set(l, r, aa, rr, 1, 0, n - 1);
                                                                            auto AUQ = applyUpdateInQuery,
                                                                           auto AUL = applyUpdateInLazy>
ll add(int a, int b, ll aa, ll rr, int p, int l,
        int r) {
                                                               struct LazySegmentTree {
  prop(p, l, r);
if (b < l or r < a) return seg[p].sum;
                                                                  int n, h;
vector<Qt> ts;
                                                                  vector<Lt> ds;
  if (a \le l \text{ and } r \le b) {
     seq[p].add_a += aa;
                                                                  vector<pii> lrs;
     seg[p].add_r += rr;
                                                                  LazySegmentTree(int _n)
     prop(p, l, r);
                                                                       : n(_n),
     return seg[p].sum;
                                                                         h(sizeof(int) * 8 - __builtin_clz(n)),
                                                                         ts(n \ll 1),
  int m = (l + r) / 2;
int tam_l = inter({l, m}, {a, b});
                                                                         ds(n),
                                                                    lrs(n << 1) {
rep(i, 0, n) lrs[i + n] = {i, i};}
  return seg[p].sum =
                                                                    rrep(i, n - 1, 0) {
                add(a, b, aa, rr, 2 * p, l, m) +
               add(a, b, aa + rr * tam_l, rr,
2 * p + 1, m + 1, r);
                                                                       lrs[i] = {lrs[i << 1].first,</pre>
                                                                                   lrs[i << 1 | 1].second};</pre>
void update_add(int l, int r, ll aa, ll rr) { add(l, r, aa, rr, 1, 0, n-1);
                                                                  LazySegmentTree(const vector<Qt> &xs)
```

```
: LazySegmentTree(len(xs)) {
   copy(all(xs), ts.begin() + n);
rep(i, 0, n) lrs[i + n] = {i, i};
rrep(i, n - 1, 0) {
      ts[i] = C(ts[i << 1], ts[i << 1 | 1]
                     lrs[i << 1], lrs[i << 1 | 1]);
}
void set(int p, Qt v) {
   ts[p + n] = v;
   build(p + n);
void upd(int l, int r, Ut v) {
   l += n, r += n + 1;
int l0 = l, r0 = r;
for (; l < r; l >>= 1, r >>= 1) {
      if (l & 1) apply(l++, v);
if (r & 1) apply(—r, v);
   build(l0), build(r0 - 1);
Qt qry(int l, int r) {
    l += n, r += n + 1;
    push(l), push(r - 1);
   Qt resl = Qt(), resr = Qt();
pii lr1 = {l, l}, lr2 = {r, r};
for (; l < r; l >>= 1, r >>= 1) {
   if (l & 1)
         resl = C(resl, ts[l], lr1, lrs[l]), l++;
      if (r & 1)
         \dot{r}, resr = C(ts[r], resr, lrs[r], lr2);
   return C(resl, resr, lr1, lr2);
void build(int p) {
   while (p > 1) {
      p >>= 1;
      ts[p] = ALQ(C(ts[p << 1], ts[p << 1 | 1], lrs[p << 1 | 1]),
                         ds[p], lrs[p]);
}
void push(int p) {
   rrep(s, h, 0) {
      int i = p >> s;
if (ds[i] != Lt()) {
  apply(i << 1, ds[i]),</pre>
               apply(i << 1 | 1, ds[i]);
         ds[i] = Lt();
}
inline void apply(int p, Ut v) {
  ts[p] = AUQ(ts[p], v, lrs[p]);
  if (p < n) ds[p] = AUL(ds[p], v, lrs[p]);</pre>
```

2.4.3 Increment update sum query (top down)

```
struct Lnode {
  ll v;
bool assign;
  Lnode() : v(), assign() {} // Neutral element
Lnode(ll _v, bool a = 0) : v(_v), assign(a){};
using Qnode = ll;
using Unode = Lnode;
struct LSegTree {
  int n, ql, qr;
  vector<Qnode> st;
vector<Lnode> lz;
  Qnode merge(Qnode lv, Qnode rv, int nl,
                  int nr) {
     return lv + rv;
  }
  void prop(int i, int l, int r) {
     if (lz[i].assign) {
```

```
st[i] = lz[i].v * (r - l + 1);

if(l!= r) lz[tol(i)] = lz[tor(i)] = lz[i];
    } else {
       st[i] += lz[i].v * (r - l + 1);
if (l != r)
         lz[i] = Lnode();
  void applyV(int i, Unode v) {
    if (v.assign) {
       lz[i] = v;
    } else {
    lz[i].v += v.v;
  /*-
  LSegTree() {}
  LSegTree(int _n)
: n(_n), st(_n << 2), lz(_n << 2) {}
  bool disjoint(int l, int r) {
     return qr < l or r < ql;
  bool contains(int l, int r) {
    return gl <= l and r <= gr;
  int tol(int i) { return i << 1; }
int tor(int i) { return i << 1 | 1; }</pre>
  void build(vector<Qnode> &v) {
    build(v, 1, 0, n - 1);
  void build(vector<Qnode> &v, int i, int l,
               int r) {
    if (l == r)
       st[i] = v[\bar{l}];
       return;
    int m = midpoint(l, r);
    build(v, tol(i), l, m);
build(v, tor(i), m + 1, r);
st[i] = merge(st[tol(i)], st[tor(i)], l, r);
  void upd(int l, int r, Unode v) {
    ql = l, qr = r;
upd(1, 0, n — 1, v);
  void upd(int i, int l, int r, Unode v) {
    prop(i, l, r);
    if (disjoint(l, r)) return;
    if (contains(l, r)) {
       applyV(i, v);
prop(i, l, r);
       return;
     int m = midpoint(l, r);
    upd(tol(i), l, m, v);
upd(tor(i), m + 1, r, v);
    st[i] = merge(st[tol(i)], st[tor(i)], l, r);
  Qnode qry(int l, int r) {
    ql = l, qr = r
     return qry(1, 0, n - 1);
  Qnode gry(int i, int l, int r) {
    prop(i, l, r);
    if (disjoint(l, r)) return Qnode();
if (contains(l, r)) return st[i];
     int m = midpoint(l, r);
     return merge(qry(tol(i), l, m),
                    qry(tor(i), m + 1, r), l, r);
};
```

Bitree 2D

Description: Given a 2D array you can increment an arbitrary position, and also query the subsum of a subgrid **Time**: Update and query in $O(log N^2)$

```
struct Bit2d {
  int n;
vll2d bit;
  Bit2d(int ni) : n(ni), bit(n + 1, vll(n + 1)) \{\}
```

```
Bit2d(int ni, vll2d &xs)
: n(ni), bit(n + 1, vll(n + 1)) {
   for (int i = 1; i <= n; i++) {
  for (int j = 1; j <= n; j++) {
    update(i, j, xs[i][j]);
}</pre>
  }
bit[x][i] += val;
\{1\} sum(int x, int y) {
   ll ans = 0;
   for (int i = x; i; i -= (i & (-i))) { for (int j = y; j; j -= (j & (-j))) { ans += bit[i][j];
   return ans;
il query(int x1, int y1, int x2, int y2) {
  return sum(x2, y2) - sum(x2, y1 - 1) -
             sum(x1 - 1, y2) + sum(x1 - 1, y1 - 1);
```

Convex Hull Trick / Line Container

Description: Container where you can add lines of the form mx + b, and query the maximum value at point x.

Usage: $insert_line(m, b)$ inserts the line $m \cdot x + b$ in the container.

eval(x) find the highest value among all lines in the point x.

Time: Eval and insert in $O(\log N)$

```
const ll LLINF = 1e18;
const ll is_query = -LLINF;
struct Line {
  ll m, b;
  mutable function<const Line *()> succ;
  bool operator<(const Line &rhs) const {
  if (rhs.b != is_query) return m < rhs.m;</pre>
    const Line *s = succ();
    if (!s) return 0;
ll x = rhs.m;
    return b - s \rightarrow b < (s \rightarrow m - m) * x;
  }
struct Cht : public multiset<Line> {
                                             // maintain
                                              // max m*x+b
  bool bad(iterator y) {
    auto z = next(y)
    if (y == begin()) {
          (z == end()) return 0;
       return y\rightarrow m == z\rightarrow m \&\& y\rightarrow b <= z\rightarrow b;
    auto x = prev(y);
    if (z == end())
       return y->m == x->m && y->b <= x->b;
     return (ld)(x\rightarrowb - y\rightarrowb) * (z\rightarrowm - y\rightarrowm) >=
             (ld)(y-b-z-b)*(y-m-x-m);
  void insert_line(
       ll m,
       ii b) {
                 // \min \rightarrow insert (-m,-b) \rightarrow -eval()
    auto y = insert(\{m, b\});
    y -> succ = [=] {
       return next(y) == end() ? 0 : &*next(y);
    if (bad(y)) {
       erase(y);
       return;
    while (next(y) != end() && bad(next(y)))
       erase(next(y));
    while (y != begin() && bad(prev(y)))
       erase(prev(y));
  ĺl eval(ll x) {
    auto l = *lower_bound((Line){x, is_query});
    return l.m * x + l.b;
```

```
2.7
    DSU / UFDS
```

};

Usage: You may discomment the commented parts to find online which nodes belong to each set, it makes the $union_set$ method cost $O(log^2)$ instead O(A)

```
struct DSU {
  vi ps, sz;
  // vector<unordered_set<int>> sts;
  DSU(int N)
       : ps(N + 1),
          sz(N, 1) /*, sts(N) */
    iota(ps.begin(), ps.end(), 0);
// for (int i = 0; i < N; i++)
// sts[i].insert(i);</pre>
  int find_set(int x) {
     return ps[x] == x ? x
                           : ps[x] = find_set(ps[x]);
  int size(int u) { return sz[find_set(u)]; }
bool same_set(int x, int y) {
     return find_set(x) == find_set(y);
  void union_set(int x, int y) {
     if (same_set(x, y)) return;
     int px = find_set(x);
     int py = find_set(y);
     if (sz[px] < sz[py]) swap(px, py);
    ps[py] = px;
sz[px] += sz[py];
     // sts[px].merge(sts[py]);
};
```

Lichao Tree (dynamic) 2.8

Description: Lichao Tree that creates the nodes dynamically, allowing to query and update from range [MAXL, MAXR]

- \bullet query(x): find the highest point among all lines in the structure
- add(a,b): add a line of form y=ax+b in the structure addSegment(a,b,l,r): add a line segment of form y=ax+b which covers from range [l, r]

Time: $O(\log N)$

```
template <typename T = ll, T MAXL = 0,
    T MAXR = 1 '000' 000'001>
struct LiChaoTree {
  bool first_best(T a, T b) { return a > b; }
T get_best(T a, T b) {
     return first_best(a, b) ? a : b;
   struct line {
     T operator()(T x) { return m * x + b; }
   struct node {
     line li;
     node *left, *right;
node(line _li = {0, inf})
    : li(_li),
    : li(_li),
              left(nullptr),
              right(nullptr) {}
     ~node() {
        delete left;
        delete right;
     }
  };
  node *root;
  LiChaoTree(line li = {0, inf})
: root(new node(li)) {}
~LiChaoTree() { delete root; }
  T query(T x, node *cur, T l,
     if (cur == nullptr) return inf;
     if (x < l or x > r) return inf;
T mid = midpoint(l, r);
     T ans = cur \rightarrow li(x);
```

```
ans = get_best(ans,
                        query(x, cur->left, l, mid));
     ans = get_best(
          ans, query(x, cur\rightarrowright, mid + 1, r));
     return ans;
  T query(T x) {
     return query(x, root, MAXL, MAXR);
  void add(line li, node *&cur, T l, T r) {
  if (cur == nullptr) {
       cur = new node(li);
       return;
     if (first_best(li(mid), cur->li(mid)))
       swap(li, cur->li);
     if (first_best(li(l), cur->li(l)))
add(li, cur->left, l, mid);
if (first_best(li(r), cur->li(r)))
add(li, cur->right, mid + 1, r);
  void add(T m, T b) {
  add({m, b}, root, MAXL, MAXR);
  if (cur == nullptr) cur = new node;
     if (lseg \leftarrow l && r \leftarrow rseg) {
       add(li, cur, l, r);
       return;
     }
T mid = midpoint(l, r);
     if (l != r)
       addSegment(li, cur—>left, l, mid, lseg,
                      rseg);
       addSegment(li, cur—>right, mid + 1, r, lseg,
                     rseg);
  void addSegment(T a, T b, T l, T r) {
  addSegment({a, b}, root, MAXL, MAXR, l, r);
     Merge sort tree
Description: Like a segment tree but each node stores the ordered
subsegment it represents.
Usage:
   • inrange(l, r, a, b): counts the number of positions i, l < i < r such that
     a \leq x_i \leq b.
Time: Build O(N \log N^2), inrange O(\log N^2)
Memory: O(n \log N)
template <class T>
struct MergeSortTree {
  int n;
  vector<vector<T>> st;
  MergeSortTree(vector<T> &xs)
        : n(len(xs)), st(n << 1) {
     rep(i, 0, n) st[i + n] = vector<T>(\{xs[i]\});
     rrep(i, n - 1, 0)  {
       st[i].resize(len(st[i << 1]) +
len(st[i << 1 | 1]));
       merge(all(st[i << 1]), all(st[i << 1 | 1]),
               st[i].begin());
    }
  int count(int i, T a, T b) {
   return upper_bound(all(st[i]), b) -
              lower_bound(all(st[i]), a);
  int inrange(int l, int r, T a, T b) {
     int ans = 0;
     for (l += n, r += n + 1; l < r;
    l >>= 1, r >>= 1) {
    if (l & 1) ans += count(l++, a, b);
}
       if (r \& 1) ans += count(--r, a, b);
```

return ans;

};

2.10 Mex with update

Description: This DS allows you to mantain an array of elments, insert, and remove, and query the MEX at any time.

Usage:

- Mex(mxsz): Initialize the DS, mxsz must be the maximum number of elements that the structure may have. \bullet add(x): just adds one copy of x.
- rmv(x): just remove a copy of x.
- operator(): returns the MEX.
- Time:
 Mex(mxsz): $O(\log mxsz)$
 - add(x): $O(\log mxsz)$ • rmv(x): $O(\log mxsz)$
 - operator(): O(1)

```
struct Mex {
   int mx_sz;
   vi hs;
   set<int> st;
   Mex(int _mx_sz) : mx_sz(_mx_sz), hs(mx_sz + 1) {
     auto it = st.begin();
rep(i, 0, mx_sz + 1) it = st.insert(it, i);
  void add(int x) {
  if (x > mx_sz) return;
  if (!hs[x]++) st.erase(x);
   void rmv(int x) {
  if (x > mx_sz) return;
         (!—hs[x]) st.emplace(x);
   int operator()() const { return *st.begin(); }
     Optional, you can just create with size
     len(xs) add N elements :D
  Mex(const vi &xs, int _mx_sz = -1)
    : Mex(~_mx_sz ? _mx_sz : len(xs)) {
    for (auto xi : xs) add(xi);
};
```

Orderd Set (GNU PBDS)

Usage: If you need an ordered **multi** set you may add an id to each value. Using greater_equal, or less_equal is considered undefined behavior.

- order of key (k): Number of items strictly smaller/greater than k .
- find by order(k): K-th element in a set (counting from zero).

Time: Both $O(\log N)$

Warning: Is 2 or 3 times slower then a regular set/map

```
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
template <typename T>
using ordered_set =
    tree<T, null_type, less<T>, rb_tree_tag,
         tree_order_statistics_node_update>;
```

2.12 Prefix Sum 2D

Description: Given an 2D array with N lines and M columns, find the sum of the subarray that have the left upper corner at (x1, y1) and right bottom corner at (x^2, y^2) .

Time: Build $O(N \cdot M)$, Query O(1).

```
template <typename T>
struct psum2d {
   vector<vector<T>> s;
   vector<vector<T>> psum;
   psum2d(vector<vector<T>>> &grid, int n, int m)
      : s(n + 1, vector<T>(m + 1)),
    psum(n + 1, vector<T>(m + 1)) {
for (int i = 1; i <= n; i++)
    for (int j = 1; j <= m; j++) {
        s[i][i] -</pre>
             s[i][j]
                     s[i][j-1] + grid[i-1][j-1]
              psum[i][j] = psum[i - 1][j] + s[i][j];
          }
   T query(int x1, int y1, int x2, int y2) { 
 T ans = psum[x2 + 1][y2 + 1] + psum[x1][y1]; 
 ans -= psum[x2 + 1][y1] + psum[x1][y2 + 1];
```

```
return ans;
```

2.13 Sparse table

```
template <typename T = ll,
              auto cmp =
                    [](T &src1, T &src2, T &dst) {
                      dst = min(src1, src2);
class SparseTable {
 private:
  int sz;
vi logs;
  vector<vector<T>> st;
 public:
  SparseTable(const vector<T> &v)
     : sz(len(v)), logs(sz + 1) { rep(i, 2, sz + 1) logs[i] = logs[i >> 1] + 1;
     st.resize(logs[sz] + 1, vector<T>(sz));
     rep(i, 0, sz) st[0][i] = v[i];

for (int k = 1; (1 << k) <= sz; k++) {

  for (int i = 0; i + (1 << k) <= sz; i++) {

    cmp(st[k - 1][i],

    st[k - 1][i] + (1 << (k - 1))],
                 st[k][i]);
        }
     }
  }
T query(int l, int r) {
     const int k = logs[r - l];
     T ret
     cmp(st[k][l], st[k][r - (1 << k)], ret);
     return ret;
```

2.14 Static range queries

```
template <typename T = ll,
          auto op =
               auto invop =
               [](const T &src1, const T &src2,
T &dst) { dst = src1 - src2; }> struct StaticRangeQueries {
  vector<T> acc;
  StaticRangeQueries(const vector<T> &XS)
       acc(len(XS)) {
    acc[0] = XS[0];
rep(i, 1, len(XS))
      op(acc[i-1], XS[i], acc[i]);
  T operator()(int l, int r) {
    T lv = (l ? acc[l - 1] : T());
    invop(acc[r], lv, ret);
    return ret;
};
```

Venice Set

Description: A container that you can insert q copies of element e, increment every element in the contaiver in x, query which is the best element and it's quantity and also remove k copies of the greatest element.

```
    add elment O(log N)

• remove O(\log N)
• update: \hat{O}(1)
• query O(1)
```

```
template <typename T = ll >
struct VeniceSet {
  using T2 = pair<T, ll>;
priority_queue<T2, vector<T2>, greater<T2>> pq;
  VeniceSet() : acc() {}
```

```
void add_element(const T &e, const ll q) {
    pq.emplace(e - acc, q);
  void update_all(const T &x) { acc += x; }
  T2 best() {
  auto ret = pq.top();
    ret.first += acc;
    return ret;
  void pop() { pq.pop(); }
  void pop_k(int k) {
    auto [e, q] = pq.top();
    pq.pop();
    q = k;
    if (q) pq.emplace(e, q);
};
```

2.16Venice Set (complete)

Description: A container which you can insert elements update all at once and also make a few queries Usage:

- $add_element(e, q)$: adds q copies of e, if no q is provided adds a single
- $update_all(x)$: increment every value by x
- erase(e): removes every copy of e, and returns how much was removed.
 count(e): returns the number of e in the container
- high()/low(): returns the hightest/lowest element, and it's quantity
- pop_low(q)/pop_high(q): removes q copies of the lowest/highest elements if no q is provided removes all copies of the lowest/highest element.

You may answer which is the K-th value and it's quantity using an $ordered_set$. Probably works with other operations

Time: Considering N the number of distinct numbers in the container

• $add_element(e,q)$: $O(\log(N))$

- $update_all(x):O(1)$ $erase(e): O(\log(N))$ $count(e): O(\log(N))$
- high()/low(): O(1)
- $pop_low(q)/pop_high(q)$: worst case is $O(N \cdot \log(N))$ if you remove all

elements and so on...
ing: There is no error handling if you try to pop more elements than exists or related stuff

```
struct VeniceSet {
  set<pll> st;
  ll acc;
  VeniceSet() : acc() {}
  ll add_element(ll e, ll q = 1) {
    q += erase(e);
    \vec{e} = acc
    st.emplace(e, q);
    return q;
  void update_all(ll x) { acc += x; }
  ll erase(ll e) {
    auto it = st.lb({e, LLONG_MIN});
if (it == end(st) || (*it).first != e)
      return 0;
    ll ret = (*it).second;
    st.erase(it);
    return rèt;
  ll count(ll x) {
    auto it = st.lb({x, LLONG_MIN});
if (it == end(st) || (*it).first != x)
      return 0;
    return (*it).second;
  pll high() { return *rbegin(st); }
  pll low() { return *begin(st); }
  void pop_high(ll q = -1) {
    if (q == -1) q = high().second;
    while (q) {
                eq] = high()
      auto [e,
      st.erase(prev(end(st)));
      if (eq > q) add_element(e, eq - q);
      q = max(0ll, q - eq);
  void pop_low(ll q = -1) {
```

```
if (q == -1) q = low().second;
    while (q) {
      auto [e, eq] = low();
      st.erase(st.begin());
      if (eq > q) add_element(e, eq - q);
      q = \max(011, q - eq);
  }
};
2.17
     Wavelet tree
```

```
using ll = long long;
template <typename T>
struct WaveletTree {
  struct Node {
    T lo, hi; int left_child, right_child;
    vector<int> pcnt;
    vector<ll> psum;
    Node(int lo_, int hi_)
: lo(lo_),
hi(hi_),
          left_child(0)
          right_child(0),
          pcnt(),
          psum() {}
  vector<Node> nodes;
  WaveletTree(vector<T> v) {
    nodes.reserve(2 * v.size());
    auto [mn, mx] =
    return;
      auto mid = midpoint(node.lo, node.hi);
      auto f = [\&mid](T x) \{ return x <= mid; \};
      node.pcnt.reserve(to
                              from + 1);
      node.pcnt.push_back(0);
                             - from + 1);
      node.psum.reserve(to -
      node.psum.push_back(0);
      T left_upper = node.lo
        right_lower = node.hi;
      for (auto it = from; it != to; it++) {
        auto value = f(*it)
        node.pcnt.push_back(node.pcnt.back() +
                             value);
        node.psum.push_back(node.psum.back() +
                             *it);
        if (value)
          left_upper = max(left_upper, *it);
        else
          right_lower = min(right_lower, *it);
      auto pivot = stable_partition(from, to, f);
      node.left_child =
          make_node(node.lo, left_upper);
                 nodes[node.left_child], from,
      self(self,
           pivot);
      node.right_child =
          make_node(right_lower, node.hi);
      self(self, nodes[node.right_child], pivot,
    build(build, nodes[make_node(*mn, *mx)],
          v.begin(), v.end());
 if (l > r) return 0;
if (node.lo == node.hi) return node.lo;
      int lb = node.pcnt[l],
    rb = node.pcnt[r + 1],
    left_size = rb - lb;
      return (left_size > k
? self(self,
                          nodes[node.left_child],
                   lb, r\tilde{b} - 1, k): self(self,
```

```
nodes[node.right_child],
                         l - lb, r - rb, k - left_size);
    return f(f, nodes[0], L, R, K);
  pair<int, ll> count_and_sum_in_range(
     int L, int R, T a, T b) const
    if (l > r or node.lo > b or node.hi < a)
  return {0, 0};
if (a <= node.lo and node.hi <= b)</pre>
        return \{r - l + 1,
                node.psum[l])};
      int lb = node.pcnt[l],
          rb = node.pcnt[r + 1];
     auto [right_cnt, right_sum] =
          self(self, nodes[node.right_child],
               l - lb, r - rb);
      return {left_cnt + right_cnt,
              left_sum + right_sum};
    return f(f, nodes[0], L, R);
  inline int count_in_range(int L, int R, T a,
                            T b) const {
    return count_and_sum_in_range(L, R, a, b)
        .first;
  inline ll sum_in_range(int L, int R, T a,
                         T b) const {
    return count_and_sum_in_range(L, R, a, b)
        .second;
 }
 private:
  int make_node(T lo, T hi) {
   int id = (int)nodes.size();
nodes.emplace_back(lo, hi);
    return id;
};
```

Dynamic Programming

3.1 Binary Knapsack (bottom up)

Description: Given the points each element have, and it repespective cost, computes the maximum points we can get if we can ignore/choose an element, in such way that the sum of costs don't exceed the maximum cost allowed. Time: O(N*W)

Warning: The vectors VS and WS starts at one, so it need an empty value at index 0

```
const int MAXN(1 '000), MAXCOST(1' 000 * 20); ll dp[MAXN + 1][MAXCOST + 1];
bool ps[MAXN + 1][MAXCOST + 1]
pair<ll, vi> knapsack(const vll &points,
                             const vi &costs,
                             int maxCost) {
  int n = len(points)
             1; // ELEMENTS START AT INDEX 1 !
  for (int m = 0; m \le maxCost; m++) {
     dp[0][m] = 0;
  for (int i = 1; i \le n; i++) { dp[i][0] = dp[i-1][0] +
                    (costs[i] == 0) * points[i];
     ps[i][0] = costs[i] == 0;
  for (int i = 1; i \le n; i++) {
     for (int m = 1; m <= maxCost; m++) {
  dp[i][m] = dp[i - 1][m], ps[i][m] = 0;</pre>
       int w = costs[i];
ll v = points[i];
        if (w \le m \text{ and }
```

```
dp[i - 1][m - w] + v > dp[i][m]) {
    dp[i][m] = dp[i - 1][m - w] + v,
    ps[i][m] = 1;
    }
}

vi is;
for (int i = n, m = maxCost; i >= 1; —i) {
    if (ps[i][m]) {
        is.emplace_back(i);
        m -= costs[i];
    }
}
return {dp[n][maxCost], is};
```

3.2 Edit Distance

Time: O(N*M)

```
int edit_distance(const string &a,
                    const string &b) {
  int n = a.size();
  int m = b.size();
  vector<vi> dp(n + 1, vi(m + 1, 0));
  int ADD = 1, DEL = 1, CHG = 1;
  for (int i = 0; i <= n; ++i) {
    dp[i][0] = i * DEL;
  for (int i = 1; i <= m; ++i) {
    dp[0][i] = ADD * i;
 for (int i = 1; i <= n; ++i) {
  for (int j = 1; j <= m; ++j) {
    int add = dp[i][j - 1] + ADD;
      int del = dp[i - 1][j] + DEL;
      int chg =
      }
  return dp[n][m];
```

3.3 Knapsack

Description: Finds the maximum score you can achieve, given that you have N items, each item has a cost, a point and a quantity, you can spent at most maxcost and buy each item the maximum quantity it has.

Time: $O(n \cdot maxcost \cdot \log maxqtd)$

 $\mathbf{Memory} \colon O(maxcost).$

```
ll knapsack(const vi &weight, const vll &value,
             const vi &qtd, int maxCost) {
  vi costs;
  vll values;
  for (int i = 0; i < len(weight); i++) {
    ll q = qtd[i];
    for (ll x = 1; x <= q; q -= x, x <<= 1) {
      costs.eb(x * weight[i]);
      values.eb(x * value[i]);
    if (q) {
      costs.eb(q * weight[i]);
      values.eb(q * value[i]);
  }
  vll dp(maxCost + 1);
  for (int i = 0; i < len(values); i++) {</pre>
    for (int j = maxCost; j > 0; j—) {
  if (j >= costs[i])
        dp[j] = max(dp[j],
                     values[i] + dp[j - costs[i]]);
    }
  return dp[maxCost];
```

3.4 Longest Increasing Subsequence

Description: Find the pair (sz, psx) where sz is the size of the longest subsequence and psx is a vector where psx_i tells the size of the longest increase subsequence that ends at position $i.\ get_i dx$ just tells which indices could be in the longest increasing subsequence.

Time: $O(n \log n)$

```
template <typename T>
pair<int, vi> lis(const vector<T> &xs, int n) {
  vector<T> dp(n + 1, numeric_limits<T>::max());
  dp[0] = numeric_limits<T>::min();
  int sz = 0;
  vi psx(n);
  rep(i, 0, n) {
    int pos =
        lower_bound(all(dp), xs[i]) - dp.begin();
    sz = max(sz, pos);
    dp[pos] = xs[i];
    psx[i] = pos;
  return {sz, psx};
}
template <typename T>
vi get_idx(vector<T> xs) {
  int n = xs.size();
  auto [sz1, psx1] = lis(xs, n);
  transform(rall(xs), xs.begin(),
             [](T x) \{ return -x; \});
  auto [sz2, psx2] = lis(xs, n);
 vi ans;

----(i, 0, n) {
  rep(i, 0, n) {
  int l = psx1[i];
    int r = psx2[n - i - 1];
    if (l + r - 1 == sz1) ans.eb(i);
  return ans;
}
```

3.5 Monery sum

Description: Find every possible sum using the given values only once.

```
set<int> money_sum(const vi &xs) {
  using vc = vector<char>;
  using vvc = vector<vc>
  int _m = accumulate(all(xs), 0);
  int _n = xs.size()
  vvc _dp(_n + 1, vc(_m + 1, 0));
  set<int> _ans;
   _dp[0][xs[0]] = 1;
  for (int i = 1; i < _n; ++i) {
  for (int j = 0; j <= _m; ++j) {
    if (j == 0 or _dp[i - 1][j]) {
          _{dp[i][j + xs[i]] = 1;}
          _dp[i][j] = 1;
       }
     }
  }
  for (int i = 0; i < _n; ++i)
for (int j = 0; j <= _m; ++j)
       if (_dp[i][j]) _ans.insert(j);
  return _ans;
}
```

3.6 Steiner tree

```
}
  }
rep(mask, 2, (1 << k)) {
  rep(i, 0, n) {
      if (inks[i]) continue;
      for (int mask2 = (mask - 1) & mask;
    mask2 >= 1;
    mask2 = (mask2 - 1) & mask) {
    int mask3 = mask ^ mask2;
         chmin(dp[i][mask]
                  dp[i][mask2] + dp[i][mask3]);
      rep(j, 0, n) {
         if (inks[j]) continue;
chmin(dp[j][mask],
                  dp[i][mask] + adj[i][j]);
  }
\bar{T} ans = inf; rep(i, 0, n) chmin(ans, dp[i][(1 << k) - 1]);
return ans;
```

3.7 Sum of Subsets

Description: Allows you to find if some mask X is a super mask of any of the given masks Usage: Call build with the masks then it returns a vector of bool V where

Coage. Can out with the masks then it returns a vector of bool V wh V_X says if X is a super mask of any of the initial maks. You can change it to count how many submasks of each mask exsists, by changing the bitwise or by a plus sign...

Time: $O(LOG \cdot 2^{LOG})$

Memory: $O(LOG^2 \cdot 2^{LOG})$ Warning: Remember to set LOG with the highest bit possible

```
const int LOG = 20;
vc build(const vi &masks) {
  vc ret(1 << LOG);
  trav(mi, masks) ret[mi] = 1;
  rep(b, 0, LOG) {
    rep(mask, 0, (1 << LOG)) {
  if (mask & (1 << b))
         ret[mask] |= ret[mask ^ (1 << b)];
  return ret;
```

Travelling Salesman Problem

```
Time: O(N^2 \cdot 2^N)
Memory: O(N^2 \cdot 2^N)
vll2d dist;
vll memo;
int tsp(int i, int mask, int N) {
  if (mask == (1 \le N) - 1) return dist[i][0]
  if (memo[i][mask] != -1) return memo[i][mask];
int ans = INT_MAX << 1;
for (int j = 0; j < N; ++j) {
    if (mask);</pre>
     if (mask & (1 << j)) continue;
      auto t =
            tsp(j, mask | (1 << j), N) + dist[i][j];
      ans = min(ans, t);
   return memo[i][mask] = ans;
```

Extras

4.1 Binary to gray

```
string binToGray(string bin)
 string gray(bin.size(),
                   ′0′);
 int n = bin.size() - 1;
 gray[0] = bin[0];
 return gray;
```

4.2 Get permutation cycles

Description: Receives a permutation [0, n-1] and return a vector 2D with

```
vll2d getPermutationCicles(const vll &ps) {
  ll n = len(ps);
  vector<char> visited(n);
  vector<vll> cicles;
  rep(i, 0, n)
    if (visited[i]) continue;
    vll cicle;
    ll pos = i;
    while (!visited[pos]) {
      cicle.pb(pos);
      visited[pos] = true;
      pos = ps[pos];
    cicles.push_back(vll(all(cicle)));
  return cicles;
```

4.3Max & Min Check

Description: Returns the min/max value in range [l, r] that satisfies the lambda function check, if there is no such value the 'nullopt' is returned. Usage: check must be a function that receives an integer and return a boolean. Time: $O(\log r - l + 1)$

```
template <typename T>
optional<T> maxCheck(T l, T r, auto check) {
  optional<T> ret;
  while (l <= r)
    T m = midpoint(l, r);
    if (check(m))
      ret ? chmax(ret, m) : ret = m, l = m + 1;
    else
      r = m - 1:
  return ret;
template <typename T>
optional<T> minCheck(T l, T r, auto check) {
  optional<T> ret;
  while (l \ll r)
    T m = midpoint(l, r);
    if (check(m))
      ret ? chmin(ret, m) : ret = m, r = m - 1;
    else
      l = m + 1:
  return ret;
```

4.4 Mo's algorithm

```
template <typename T, typename Tans>
struct Mo {
  struct Query
    int l, r, idx, block;
    Query(int _l, int _r, int _idx, int _block)
        : l(_l),
          r(_r)
          idx(_idx)
          block(_block) {}
    bool operator<(const Query &q) const {</pre>
      if (block != q.block)
        return block < q.block;</pre>
      return (block & 1 ? (r < q.r) : (r > q.r));
    }
  };
  vector<T> vs;
  vector<Query> qs;
  const int block_size;
  Mo(const vector<T> &a)
        vs(a),
        block_size((int)ceil(sqrt(a.size()))) {}
  void add_query(int l, int r) {
    qs.emplace_back(l, r, qs.size(),
```

```
l / block_size);
 }
 auto solve() {
    // get answer return type
    vector<Tans> answers(qs.size());
    sort(all(qs));
   int cur_l = 0, cur_r = -1;
for (auto q : qs) {
  while (cur_l > q.l) add(—cur_l);
      while (cur_r < q.r) add(++cur_r);</pre>
      while (cur_l < q.l) remove(cur_l++);
while (cur_r > q.r) remove(cur_r—);
      answers[q.idx] = get_answer();
   return answers;
 }
private:
 // add value at idx from data structure
inline void add(int idx) {}
    remove value at idx from data structure
 inline void remove(int idx) {}
 // extract current answer of the data structure inline Tans get_answer() \{\}
```

4.5 int 128t stream

```
void print(__int128 x) {
   if (x < 0) {
      cout << '-';
      x = -x;
   }
   if (x > 9) print(x / 10);
   cout << (char)((x % 10) + '0');
}
__int128 read() {
   string s;
   cin >> s;
   __int128 x = 0;
   for (auto c : s) {
      if (c != '-') x += c - '0';
      x *= 10;
   }
   x /= 10;
   if (s[0] == '-') x = -x;
   return x;
}
```

5 Geometry

5.1 Check if a point belong to line segment

```
// Verifica se o ponto P pertence ao segmento de
// reta AB
const ld EPS = 1e-9;
template <typename T>
struct Point {
  T x, y;
  Point(T_x, T_y) : x(_x), y(_y) \{ \}
template <typename T>
bool equals(const T a, const T b) {
  if (is_floating_point<T>) {
    return fabsl(a - b) <= EPS;
  return a == b;
}
   Verify if the segment AB contains point P
template <typename T>
bool contains(const Point<T> &A,
                const Point<T> &B
                const Point<T> &P) {
  auto xmin = min(A.x, B.x);
  auto xmax = max(A.x, B.x);
  auto ymin = min(A.y, B.y);
auto ymax = max(A.y, B.y);
```

```
if (P.x < xmin \mid\mid P.x > xmax \mid\mid P.y < ymin \mid\mid
       P.y > ymax)
     return false;
  return equals((P.y - A.y) * (B.x - A.x)
                   (P.x - A.x) * (B.y - A.y));
}
    Check if point is inside triangle
5.2
struct point {
  int x, y;
  int id;
  point operator—(const point &o) const {
     return \{x - o.x, y - o.y\};
  int operator^(const point &o) const {
     return x * o.y - y * o.x;
};
    Verify the direction that the point
  _e_ is in relation to the vector formed by the points a—>b
  -1 = right
  0 = collinear
1 = left
int ccw(point a, point b, point e) {
  int tmp = (b - a) ^ (e - a);
  return (tmp > 0) - (tmp < 0);
}
  Verify if the point e
  is inside the triangle formed by
  the points t1, t2, t3
bool inside_triangle(point t1, point t2, point t3,
                         point e) {
  int x = ccw(t1, t2, e);
  int y = ccw(t2, t3, e);

int z = ccw(t3, t1, e);

return !((x == 1 or y == 1 or z == 1) and
             (x == -1 \text{ or } y == -1 \text{ or } z == -1));
}
5.3 Convex hull
struct pt {
  double x, y;
  int id;
};
int orientation(pt a, pt b, pt c) {
  double v = a.x * (b.y - c.y) + b.x * (c.y - a.y) +
  c.x * (a.y - b.y);
if (v < 0) return -1; // clockwise
if (v > 0) return +1; // counter-clockwise
  return 0;
}
int o = orientation(a, b, c);
return o < 0 || (include_collinear && o == 0);</pre>
bool collinear(pt a, pt b, pt c)
  return orientation(a, b, c) == 0;
void convex_hull(vector<pt> &pts,
                    bool include_collinear = false) {
  pt p0 = *min_element(all(pts), [](pt a, pt b) {
    sort(all(pts), [&p0](const pt &a, const pt &b) {
```

int o = orientation(p0, a, b);

return (p0.x - a.x) * (p0.x - a.x) +

(p0.x - b.x) * (p0.x - b.x) +

(p0.y - a.y) * (p0.y - a.y) <

if (o == 0)

```
(p0.y - b.y) * (p0.y - b.y);
  return o < 0;
});
if (include_collinear) {
  int i = len(pts) - 1;
  while (i >= 0 \&\&
         collinear(p0, pts[i], pts.back()))
  reverse(pts.begin() + i + 1, pts.end());
vector<pt> st;
for (int i = 0; i < len(pts); i++) {
 while (st.size() > 1 \&\&
         !cw(st[len(st)
                          - 2], st.back(), pts[i],
             include_collinear))
    st.pop_back();
  st.push_back(pts[i]);
}
pts = st;
```

5.4 Polygon lattice points

```
ll cross(ll x1, ll y1, ll x2, ll y2) {
  return x1 * y2 - x2 * y1;
ll polygonArea(vector<pll> &pts) {
  ll ats = 0;
  for (int i = 2; i < len(pts); i++)
    ats +=
        cross(pts[i].first - pts[0].first,
    pts[i].second - pts[0].second,
    pts[i - 1].first - pts[0].first,
    pts[i - 1].second - pts[0].second);
  return abs(ats / 2ll);
ll boundary(vector<pll> &pts) {
  ll ats = pts.size();
  for (int i = 0; i < len(pts); i++) {
    ll deltay =
         (pts[i].second -
          pts[(i + 1) % pts.size()].second);
    ats += abs(\_gcd(deltax, deltay)) - 1;
  return ats;
pll latticePoints(vector<pll> &pts) {
  ll bounds = boundary(pts);
  ll area = polygonArea(pts);
  ll inside = area + 1ll - bounds / 2ll;
  return {inside, bounds};
```

5.5 Segment intersection

```
using ld = long double;
template <typename T = ld >
struct Point {
  T x, y;
  bool is_port;
template <typename T = ld>
bool operator == (const Point < T > &a,
                 const Point<T> &b) {
  return a.x == b.x and a.y == b.y;
}
template <typename T = ld >
struct Segment {
 Point<T> p1, p2;
template <typename T>
int orientation(Point<T> p, Point<T> q,
                Point<T> r) {
  int val = (q.y - p.y) * (r.x - q.x) -
            (q.x - p.x) * (r.y - q.y);
```

```
// TODO: if it's a float must use other way to
  // compare
  if (val == 0)
  return 0; '// colinear
else if (val > 0)
return ?
    return 1; // clockwise
    return 2; // counterclockwise
template <typename T>
bool do_segment_intersect(Segment<T> s1
                             Segment<T> s2)
  int o1 = orientation(s1.p1, s1.p2, s2.p1);
  int o2 = orientation(s1.p1, s1.p2, s2.p2);
  int o3 = orientation(s2.p1, s2.p2, s1.p1);
  int o4 = orientation(s2.p1, s2.p2, s1.p2);
  return (o1 != o2 and o3 != o4) or
           (01 == 0 \text{ and } 03 == 0) \text{ or }
          (o2 == 0 \text{ and } o4 == 0);
}
```

6 Graphs

6.1 Heavy-Light Decomposition (point update)

6.1.1 Maximum number on path

```
struct Node {
  ll value;
  Node()
       : value(numeric_limits<
                ll>::min()){}; // Neutral
                                 // element
  Node(ll v) : value(v){};
Node combine(Node l, Node r) {
  Node m;
m.value = max(l.value, r.value);
  return m;
template <typename T = Node, auto F = combine>
struct SegTree {
  int n;
  vector<T> st;
  SegTree(int _n) : n(_n), st(n \ll 1) {}
  void set(int p, const T &k) {
  for (st[p += n] = k; p >>= 1;)
      st[p] = F(st[p << 1], st[p << 1 | 1]);
  T query(int l, int r) {
    T ansl, ansr;
for (l += n, r += n + 1; l < r;
      l >>= 1, r >>= 1) {
if (l & 1) ansl = F(ansl, st[l++]);
      if (r \& 1) ansr = F(st[-r], ansr);
    return F(ansl, ansr);
  }
};
template <typename SegT = Node,
           auto SegOp = combine>
struct HeavyLightDecomposition {
  int n;
vi ps, ds, sz, heavy, head, pos;
  SegTree<SegT, Seg0p> seg;
  HeavyLightDecomposition(const vi2d &g,
                             const vector<SegT> &v,
                             int root = 0
       : n(len(g)), seg(n) {
    ps = ds = sz = heavy = head = pos = vi(n, -1);
    auto dfs = [&](auto &&self, int u) -> void {
      sz[u] = 1;
      int mx = 0;
for (auto x : g[u])
         if (x != ps[u]) {
           ps[x] = u;
           ds[x] = ds[u] + 1;
           self(self, x);
           sz[u] += sz[x]
           if (sz[x] > mx)
```

```
mx = sz[x], heavy[u] = x;
  };
  dfs(dfs, root);
  for (int i = 0, cur = 0; i < n;
    if(ps[i] == -1 \text{ or heavy}[ps[i]] != i)
      for (int j = i; j != -1; j = heavy[j]) {
        head[j] = i;
        pos[j] = cur++;
  rep(i, 0, n) seg.set(pos[i], v[i]);
vector<pii> disjoint_ranges(int u, int v) {
  vector<pii> ret;
  for (; head[u] != head[v]; v = ps[head[v]]) {
    if (ds[head[u]] > ds[head[v]]) swap(u, v);
    ret.eb(pos[head[v]], pos[v]);
  if (ds[u] > ds[v]) swap(u, v);
  ret.eb(pos[u], pos[v]);
  return ret;
SegT query_path(int u, int v) {
  SegT res;
  for (auto [l, r] : disjoint_ranges(u, v)) {
    res = SegOp(res, seg.query(l, r));
  return res;
}
SegT query_subtree(int u) const {
  return seg.query(pos[u], pos[u] + sz[u] - 1);
void set(int u, SegT x) { seg.set(pos[u], x); }
```

6.2 2-SAT

Description: Calculates a valid assignment to boolean variables $a,\,b,\,c,...$ to a 2-SAT problem, so that an expression of the type (a||b)&&(!a||c)&&(d||!b)&&... becomes true, or reports that it is unsatisfiable. **Usage**: Negated variables are represented by bit-inversions (x). Returns true iff it is solvable ts.values[0..N-1] holds the assigned values to the **Time**: O(N+E), where N is the number of boolean variables, and E is the number of clauses.

```
struct TwoSat {
  int N;
  vector<vi> gr;
vi values; // 0 = false, 1 = true
  TwoSat(int n = 0) : N(n), gr(2 * n) {}
  int addVar() { // (optional)
    gr.eb();
    ar.eb():
     return N++;
  void either(int f, int j) {
    f = max(2 * f, -1 - 2 * f);
    j = max(2 * j, -1 - gr[f].pb(j ^ 1); gr[j].pb(f ^ 1);
  void setValue(int x) { either(x, x); }
  void atMostOne(const vi &li) { // (optional)
    if (sz(li) <= 1) return;
int cur = ~li[0];
    rep(i, 2, sz(li))
       int next = addVar();
       either(cur, ~li[i]);
       either(cur, next);
either(~li[i], next);
       cur = ~next;
    either(cur, ~li[1]);
  vi val, comp, z;
  int time = 0;
  int dfs(int i)
    int low = val[i] = ++time, x;
     z.pb(i);
    for (int e : gr[i])
```

```
if (!comp[e])
        low = min(low, val[e] ?: dfs(e));
(low == val[i]) do {
           x = z.back();
           z.pop_back();
           comp[x] = low;
        if (values[x >> 1] == -1)
   values[x >> 1] = x & 1;
} while (x != i);
     return val[i] = low;
  bool solve() {
     values.assign(N, -1);
val.assign(2 * N, 0);
     comp = val;
rep(i, 0, 2 * N) if (!comp[i]) dfs(i);
     rep(i, 0, N) if (comp[2 * i] ==
                             comp[2 * i + 1]) return 0;
     return 1:
};
```

6.3BFS-01

Description: Similar to a Dijkstra given a weighted graph finds the distance from source s to every other node. Time: O(V + E)

Warning: Applicable only when the weight of the edges $\{0, x\}$

```
vector<pair<ll, int>> adj[maxn];
ll dists[maxn];
int s, n;
void bfs_01() +
  fill(dists, dists + n, oo);
  dist[s] = 0;
  deque<int> q;
  q.emplace_back(s);
  while (not q.empty()) {
    auto u = q.front();
     q.pop_front();
     for (auto [v, w] : adj[u]) {
  if (dist[v] <= dist[u] + w) continue;</pre>
       dist[v] = dist[u] + w;
       w ? q.emplace_back(v) : q.emplace_front(v);
}
```

Bellman ford

Description: Find shortest path from a single source to all other nodes. Can detect negative cycles

Time: $O(V \cdot E)$

```
bool bellman_ford(
    const vector<vector<pair<int, ll>>> &g, int s,
    vector<ll> &dist)
  int n = (int)g.size()
  dist.assign(n, LLONG_MAX);
  vector<int> count(n);
vector<char> in_queue(n);
  queue<int> q;
  dist[s] = 0
  q.push(s);
  in_queue[s] = true;
  while (not q.empty()) {
    int cur = q.front();
    q.pop();
    in_queue[cur] = false;
    for (auto [to, w] : g[cur]) {
       if (dist[cur] + w < dist[to]) {</pre>
         dist[to] = dist[cur] + w;
if (not in_queue[to]) {
           q.push(to);
           in_queue[to] = true;
           count[to]++
           if (count[to] > n) return false;
      }
    }
  }
```

```
return true;
}
```

6.5 Bellman-Ford (find negative cycle)

Description: Given a directed graph find a negative cycle by running n iterations, and if the last one produces a relaxation than there is a cycle. **Time**: $O(V \cdot E)$

```
const ll oo = 2500 * 1e9;
using graph = vector<vector<pair<int, ll>>>;
vi negative_cycle(graph &g, int n) {
  vll d(n, oo);
  vi p(n, -1);
  int x = -1;

d[0] = 0;
  for (int i = 0; i < n; i++) {
    for (int u = 0; u < n; u++)
       for (auto &[v, l] : g[u]) {
  if (d[u] + l < d[v]) {
           d[v] = d[u] + 1;

p[v] = u;
           X = V;
         }
      }
    }
  }
  if (x == -1)
    rèturn {};
  else {
   for (int i = 0; i < n; i++) x = p[x];
    vi cycle;
    for (int v = x;; v = p[v]) {
       cycle.eb(v);
       if (v == x \text{ and } len(cycle) > 1) break;
    reverse(all(cycle));
    return cycle;
}
```

6.6 Biconnected Components

Description: Build a vector of vectors, where the i-th vector correspond to the nodes of the i-th, biconnected component, a biconnected component is a subset of nodes and edges in which there is no cut point, also exist at least two distinct routes in vertex between any two vertex in the same biconnected component.

Time: O(N+M)

```
const int maxn(5 '00' 000);
int tin[maxn], stck[maxn], bcc_cnt, n, top = 0,
                                              timer = 1;
vector<int> g[maxn], nodes[maxn];
int tarjan(int u, int p = -1) {
  int lowu = tin[u] = timer++;
int son_cnt = 0;
  stck[++top] = u;
  for (auto v : g[u]) {
   if (!tin[v]) {
       son_cnt++;
       int lowx = tarjan(v, u);
       lowu = min(lowu, lowx);
       if (lowx >= tin[u])
         while (top !=-1 \&\& stck[top + 1] != v)
           nodes[bcc_cnt].emplace_back(
                stck[top—]);
         nodes[bcc_cnt++].emplace_back(u);
    } else {
       lowu = min(lowu, tin[v]);
  }
  if (p == -1 \&\& son\_cnt == 0) {
    nodes[bcc_cnt++].emplace_back(u);
  return lowu;
void build_bccs() {
  timer = 1;
top = -1;
  memset(tin, \theta, sizeof(int) * n);
for (int i = \theta; i < n; i++) nodes[i] = {};
```

```
bcc_cnt = 0;
for (int u = 0; u < n; u++)
   if (!tin[u]) tarjan(u);</pre>
```

6.7 Binary Lifting/Jumping

Description: Given a function/successor grpah answers queries of the form which is the node after k moves starting from u. **Time:** Build $O(N \cdot MAXLOG2)$, Query O(MAXLOG2).

```
const int MAXN(2e5), MAXLOG2(30);
int bl[MAXN][MAXLOG2 + 1];
int N;
int jump(int u, ll k) {
  for (int i = 0; i <= MAXLOG2; i++) {
    if (k & (1ll << i)) u = bl[u][i];
  }
  return u;
}
void build() {
  for (int i = 1; i <= MAXLOG2; i++) {
    for (int j = 0; j < N; j++) {
      bl[j][i] = bl[bl[j][i - 1]][i - 1];
    }
}</pre>
```

6.8 Bipartite Graph

Description: Given a graph, find the 'left' and 'right' side if is a bipartite graph, if is not then a empty vi2d is returned **Time**: O(N+M)

```
vi2d bipartite_graph(vi2d &adj) {
  int n = len(adj);
  vi side(n, -1); vi2d ret(2);
  rep(u, 0, n) {
    if (side[u] = -1) {
       queue<int> q;
       q.emp(u);
       side[u] = 0;
       ret[0].eb(u)
      while (len(q)) {
         int u = q.front();
         q.pop();
         for (auto v : adj[u])
           if (side[v] == -1) {
    side[v] = side[u] ^ 1;
              ret[side[v]].eb(v);
              q.push(v);
           } else if (side[u] == side[v])
              return {};
         }
      }
    }
  return ret;
```

6.9 Block-Cut tree

```
struct block_cut_tree {
  int n;
  vector<int> id, is_cutpoint, tin, low, stk;
  vector<vector<int>> comps, tree;
  block_cut_tree(vector<vector<int>> &g)
       : n(g.size()),
         id(n),
         is_cutpoint(n),
         tin(n),
    low(n) {
// build comps
    for (int i = 0; i < n; i++) {
  if (!tin[i]) {</pre>
         (!tin[i]) {
int timer = 0;
         dfs(i, -1, timer, g);
      }
    }
    int node_id = 0;
    for (int u = 0; u < n; u++) {
```

```
if (is_cutpoint[u]) {
         id[u] = node_id++
         tree.push_back({});
    }
    for (auto &comp : comps) {
       int node = node_id++;
       tree.push_back({});
       for (int u : comp)
         if (!is_cutpoint[u]) {
           id[u] = node;
         } else {
  tree[node].emplace_back(id[u]);
           tree[id[u]].emplace_back(node);
      }
    }
  tin[u] = low[u] = ++timer;
    stk.emplace_back(u);
    for (auto v : g[u]) {
  if (v == p) continue;
  if (!tin[v]) {
         dfs(v, u, timer, g);
low[u] = min(low[u], lo
if (low[v] >= tin[u]) {
                                 low[v]);
           is_cutpoint[u] =
           (tin[u] > 1 or tin[v] > 2);
comps.push_back({u});
           while (comps.back().back() != v) {
              comps.back().emplace_back(stk.back());
              stk.pop_back();
           }
       } else
         low[u] = min(low[u], tin[v]);
    }
};
```

6.10 D'Escopo-Pape

Description: Is a single source shortest path that works faster than Dijkstra's algorithm and the Bellman-Ford algorithm in most cases, and will also work for negative edges. However not for negative cycles. There exists cases where it runs in exponential time.

Usage: Returns a pair containing two vectors, the first one with the distance from s to every other node, and another one with the ancestor of each node, note that the ancestor of s is -1

```
using Edge = pair<ll, int>;
using Adj = vector<vector<Edge>>;
pair<vll, vi> desopo_pape(int s, int n,
                            const Adj &adj) {
  vll ds(n, LLONG_MAX), ps(n, -1);
  ds[s] = 0;
  vi ms(n, 2);
  deque<int> q;
  q.eb(s);
  while (len(q)) {
    int u = q.front();
    q.pop_front();
    ms[u] = 0;
    for (auto [w, v] : adj[u])
      if (chmin(ds[v], w + ds[u])) {
        ps[v] = u;
        if (ms[v] == 2)
ms[v] = 1, q.pb(v);
        else if (ms[v] == 0)
          ms[v] = 1, q.pf(v);
   }
  return {ds, ps};
```

6.11 Dijkstra

```
const int MAXN = 1'00'000;
const ll MAXW = 1'000'000ll;
```

```
constexpr ll 00 = MAXW * MAXN + 1;
using Edge = pair<ll, int>; // { weigth, node}
using Adj = vector<vector<Edge>>;
template <typename T>
using min_heap =
    priority_queue<T, vector<T>, greater<T>>;
pair<vll, vi> dijkstra(const Adj &g, int s) {
  int n = len(g);
  min_heap<Edge> pq;
  vll ds(n, 00);
  vi ps(n, -1);
  pq.emp(0, s);
  ds[s] = 0;
while (len(pq)) {
    auto [du, u] = pq.top();
    pq.pop();
    if (ds[u] < du) continue;
    for (auto [w, v] : g[u]) {
      ll ndv = du + w;
       if (chmin(ds[v], ndv)) {
         ps[v] = u;
         pq.emp(ndv, v);
      }
    }
    return {ds, ps};
  }
  // optional !
  vi recover_path(int source, int ending,
                    const vi &ps) {
    if (ps[ending] == -1) return \{\};
    int cur = ending;
    vi ans;
    while (cur !=-1) {
      ans.eb(cur);
       cur = ps[cur];
    reverse(all(ans));
    return ans;
```

6.12 Dijkstra (K-shortest pahts)

```
const ll oo = 1e9 * 1e5 + 1
using adj = vector<vector<pil>>>;
vector<priority_queue<ll>> dijkstra(
    const vector<vector<pll>>> &g, int n, int s,
    int k) {
  priority_queue<pll, vector<pll>, greater<pll>>
      pq;
  vector<priority_queue<ll>> dist(n);
  dist[0].emplace(0);
  pq.emplace(0, s)
  while (!pq.empty()) {
    auto [d1, v] = pq.top();
    pq.pop();
    if (not dist[v].empty() and
        dist[v].top() < d1)
      continue;
    for (auto [d2, u] : g[v]) {
      if (len(dist[u]) < k) {
   pq.emplace(d2 + d1, u);</pre>
        dist[u].emplace(d2 + d1);
      } else {
        if (dist[u].top() > d1 + d2) {
           dist[u].pop();
           dist[u].emplace(d1 + d2);
           pq.emplace(d2 + d1, u);
        }
      }
    }
  return dist;
```

$\begin{array}{ccc} \textbf{6.13} & \textbf{Extra Edges to Make Digraph Fully Strongly} \\ \textbf{Connected} & \end{array}$

Description: Given a directed graph G find the necessary edges to add to make the graph a single strongly connected component. **Time**: O(N+M)

```
struct SCC {
  int n, num_sccs;
vi2d adj;
  vi scc_id;
  SCC(int _n)
       : n(_n),
         num_sccs(0),
          adj(n)
          scc_id(n, -1) {}
  SCC(const vi2d &_adj) : SCC(len(_adj)) {
     adj = \_adj;
     find_sccs();
  void add_edge(int u, int v) { adj[u].eb(v); }
  void find_sccs() {
     int timer = 1;
     vi tin(n), st;
     st.reserve(n)
     function<int(int)> dfs = [\&](int u) \rightarrow int {
       int low = tin[u] = timer++, siz = len(st);
       st.eb(u);
       for (int v : adj[u])
   if (scc_id[v] < 0)</pre>
            low =
                 min(low, tin[v] ? tin[v] : dfs(v));
       if (tin[u] == low) {
  rep(i, siz, len(st)) scc_id[st[i]] =
              núm_sccs;
         st.resize(siz);
num_sccs++;
       return low;
    };
     for (int i = 0; i < n; i++)
  if (!tin[i]) dfs(i);</pre>
};
vector<array<int, 2>> extra_edges(
     const vi2d &adj) {
  SCC scc(adj);
  auto scc_id = scc.scc_id;
  auto num_sccs = scc.num_sccs;
  if (num_sccs == 1) return {};
  int n = len(adj);
  vi2d scc_adj(num_sccs);
  vi zero_in(num_sccs, 1);
  rep(u, 0, n) {
    for (int v : adj[u]) {
  if (scc_id[u] == scc_id[v]) continue;
  scc_adj[scc_id[u]].eb(scc_id[v]);
       zero_in[scc_id[v]] = 0;
    }
  }
  int random_source_
       max_element(all(zero_in)) - zero_in.begin();
  vi vis(num_sccs);
  function<int(int)> dfs = [&](int u) {
  if (empty(scc_adj[u])) return u;
     for (int v : scc_adj[u])
  if (!vis[v]) {
          vis[v] = 1;
         int zero_out = dfs(v);
if (zero_out != -1) return zero_out;
     return (int)—1;
  vector<array<int, 2>> edges;
  vi in_unused;
  rep(i, 0, num_sccs) {
  if (zero_in[i]) {
       vis[i] = 1;
       int zero_out = dfs(i);
          (zero\_out != -1)
         edges.push_back({zero_out, i});
       else
          in_unused.push_back(i);
  }
  rep(i, 1, len(edges)) {
     swap(edges[i][0], edges[i-1][0]);
```

```
}
  rep(i, 0, num_sccs) {
    if (scc_adj[i].empty() && !vis[i]) {
      if (!in_unused.empty()) {
        edges.push_back({i, in_unused.back()});
        in_unused.pop_back();
      } else {
        edges.push_back({i, random_source});
    }
  for (int u : in_unused) edges.push_back({0, u});
  vi to_node(num_sccs)
  rep(i, 0, n) to_node[scc_id[i]] = i;
  for (auto &[u, v] : edges)
    u = to_node[u], v = to_node[v];
  return edges:
}
6.14 Find Articulation/Cut Points
Description: Given an undirected graph find it's articulation points.
Time: O(N+M)
```

Warning: A vertex u can be an articulation point if and only if has at least 2 adjascent vertex

```
const int MAXN(100);
int N;
vi2d G;
int timer;
int tin[MAXN],
                   low[MAXN];
set<int> cpoints;
int dfs(int u, int p = -1) {
  int cnt = 0;
  low[u] = tin[u] = timer++;
for (auto v : G[u]) {
   if (not tin[v]) {
        dfs(v, u);
        if (low[v] >= tin[u]) cpoints.insert(u);
        low[u] = min(low[u], low[v]);
     } else if (v != p)
        low[u] = min(low[u], tin[v]);
  }
  return cnt;
void getCutPoints() {
  memset(low, 0, sizeof(low));
memset(tin, 0, sizeof(tin));
cpoints.clear();
   timer = 1;
   for (int i = 0; i < N; i++) {
     if (tin[i]) continue;
int cnt = dfs(i);
     if (cnt == 1) cpoints.erase(i);
}
```

Find Bridge-Tree components

Usage: label2CC(u, p) finds the 2-edge connected component of every node. $\bar{\mathbf{Time}}:\; O(n+m)$

```
const int maxn(3 '00' 000);
int tin[maxn], compId[maxn], qtdComps;
vi g[maxn], stck;
int n;
int dfs(int u, int p = -1) {
  int low = tin[u] = len(stck);
  stck.emplace_back(u);
  bool multEdge = false;
  for (auto v : g[u]) {
  if (v == p and !multEdge) {
      multEdge = 1;
      continue;
    low = min(low,
               tin[v] == -1 ? dfs(v, u) : tin[v]);
  if (low == tin[u])
    for (int i = tin[u]; i < len(stck); i++)
```

```
compId[stck[i]] = qtdComps;
    stck.resize(tin[u]);
    qtdComps++;
}
return low;
}
void label2CC() {
    memset(compId, -1, sizeof(int) * n);
    memset(tin, -1, sizeof(int) * n);
    stck.reserve(n);
    for (int i = 0; i < n; i++) {
        if (tin[i] == -1) dfs(i);
    }
}</pre>
```

6.16 Find Bridges

Description: Find every bridge in a **undirected** connected graph. **Warning**: Remember to read the graph as pair where the second is the id of the edge!

```
Time: \$0(N + M) $ const int MAXN(10000),
        MAXM(100000)
        M, clk, tin[MAXN], low[MAXN],
     isBridge[MAXM];
vector<pii> G[MAXN];
void dfs(int u, int p = -1) {
  tin[u] = low[u] = clk++;
  for (auto [v, i] : G[u]) {
  if (v == p) continue;
     if (tin[v])
       low[u] = min(low[u], tin[v]);
     } else {
       dfs(v, u);
low[u] = min(low[u], low[v]);
if (low[v] > tin[u]) {
          isBridge[i] = 1;
     }
  }
void findBridges() {
  fill(tin, tin + N, 0);
fill(low, low + N, 0);
fill(isBridge, isBridge + M, 0);
  clk = 1;
for (int i = 0; i < N; i++) {
    if (!tin[i]) dfs(i);
```

6.17 Find Centroid

Description: Given a tree (don't forget to make it 'undirected'), find it's centroids. @Time : O(V)

6.18 Find bridges (online)

```
// O((n+m)*log(n))
struct BridgeFinder {
```

```
// 2ecc = 2 edge conected component
// cc = conected component
vector<int> parent, dsu_2ecc, dsu_cc,
dsu_cc_size;
int bridges, lca_iteration;
vector<int> last_visit;
BridgeFinder(int n)
     \vdots parent(n, -1),
       dsu_2ecc(n),
       dsu_cc(n),
       dsu_cc_size(n, 1),
       bridges(0),
       lca_iteration(0),
       last_visit(n) -
  for (int i = 0; i < n; i++) {
    dsu_2ecc[i] = i;
    dsu_cc[i] = i;
  }
int find_2ecc(int v) {
  if (v == -1) return -1;
  return dsu_2ecc[v] == v
               : dsu_2ecc[v] =
                      find_2ecc(dsu_2ecc[v]);
int find_cc(int v) {
  v = find_2ecc(v);
  return dsu_cc[v] == v
               : dsu_cc[v] = find_cc(dsu_cc[v]);
void make_root(int v) {
  v = find_2ecc(v);
  int root = v;
  int child = -
  while (v != -1) {
    int p = find_2ecc(parent[v]);
    parent[v] = child;
    dsu_cc[v] = root;
    child = v;
  dsu_cc_size[root] = dsu_cc_size[child];
void merge_path(int a, int b) {
  ++lca_iteration;
  vector<int> path_a, path_b;
  int lca = -1;
  while (lca == -1) { if (a != -1) {
      a = find_2ecc(a)
       path_a.push_back(a);
       if (last_visit[a] == lca_iteration) {
         lca = a;
         break;
       last_visit[a] = lca_iteration;
       a = parent[a];
    if (b != -1) {
    b = find_2ecc(b);
       path_b.push_back(b);
       if (last_visit[b] == lca_iteration) {
         lca = b;
         break;
       last_visit[b] = lca_iteration;
       b = parent[b];
  for (auto v : path_a) {
    dsu_2ecc[v] = lca;
    if (v == lca) break;
     __bridges;
  for (auto v : path_b) {
  dsu_2ecc[v] = lca;
  if (v == lca) break;
      -bridges;
void add_edge(int a, int b) {
```

```
a = find_2ecc(a);
b = find_2ecc(b);
if (a == b) return;
int ca = find_cc(a);
int cb = find_cc(b);
if (ca != cb) {
    ++bridges;
    if (dsu_cc_size[ca] > dsu_cc_size[cb]) {
        swap(a, b);
        swap(ca, cb);
    }
    make_root(a);
    parent[a] = dsu_cc[a] = b;
    dsu_cc_size[cb] += dsu_cc_size[a];
} else {
    merge_path(a, b);
}
};
```

6.19 Floyd Warshall

Description: Simply finds the minimal distance for each node to every other node. $O(V^3)$

6.20 Functional/Successor Graph

Description: Given a functional graph find the vertice after k moves starting at u and also the distance between u and v, if it's impossible to reach v starting at u returns -1.

Time: build $O(N \cdot MAXLOG2)$, kth O(MAXLOG2), dist O(MAXLOG2)

```
const int MAXN(2 '000' 000), MAXLOG2(24);
int N:
vi2d succ(MAXN, vi(MAXLOG2 + 1));
vi dst(MAXN, 0);
int vis[MAXN]
void dfsbuild(int u) {
  if (vis[u]) return;
vis[u] = 1;
  int v = succ[u][0];
  dfsbuild(v);
  dst[u] = dst[v] + 1;
void build() {
  for (int i = 0; i < N; i++) {
    if (not vis[i]) dfsbuild(i);</pre>
  }
  for (int k = 1; k <= MAXLOG2; k++) {
  for (int i = 0; i < N; i++) {</pre>
       succ[i][k] = succ[succ[i][k-1]][k-1];
  }
int kth(int u, ll k) {
  return u;
int dist(int u, int v)
  int cu = kth(u, dst[u]);
if (kth(u, dst[u] - dst[v]) == v)
  return dst[u] - dst[v];
else if (kth(cu, dst[cu] - dst[v]) == v)
return dst[u] + (dst[cu] - dst[v]);
  else
     return -1;
```

6.21 Kruskal

Description: Find the minimum spanning tree of a graph. $\textbf{Time} \colon O(E \log E)$

```
struct UFDS {
  vector<int> ps, sz;
  int components;
  UFDS(int n)
        ps(n + 1), sz(n + 1, 1), components(n) {
    iota(all(ps), 0);
  int find_set(int x) {
    return (x == ps[x]
                 : (ps[x] = find_set(ps[x]));
  bool same_set(int x, int y) {
    return find_set(x) == find_set(y);
  void union_set(int x, int y) {
    x = find_set(x);
y = find_set(y);
    if (x == y) return;
    if (sz[x] < sz[y]) swap(x, y);
    ps[y] = x;
sz[x] += sz[y];
    components—;
};
vector<tuple<ll, int, int>> kruskal(
    int n, vector<tuple<ll, int, int>> &edges) {
  UFDS ufds(n);
  vector<tuple<ll, int, int>> ans;
  sort(all(edges));
  for (auto [a, b, c] : edges) {
    if (ufds.same_set(b, c)) continue;
    ans.emplace_back(a, b, c);
    ufds.union_set(b, c);
  return ans;
}
```

6.22 Lowest Common Ancestor

Description: Given two nodes of a tree find their lowest common ancestor, or their distance

```
template <typename T>
struct SparseTable {
   vector<T> v;
   int n;
   static const int b = 30;
   vi mask, t;
   int op(int x, int y) {
       return v[x] < v[y] ? x : y;
   int msb(int x)
       return __builtin_clz(1) - __builtin_clz(x);
   SparseTable() {}
   SparseTable(const vector<T> &v_)
       i v(v_), n(v.size()), mask(n), t(n) {
for (int i = 0, at = 0; i < n;
    mask[i++] = at |= 1) {
    at = (at << 1) & ((1 << b) - 1);
    while (at and op(i, i - msb(at & -at)) == i)</pre>
             at ^= at & —at;
       for (int i = 0; i < n / b; i++)
t[i] = b * i + b - 1 -
      \begin{array}{c} \text{msb}(\text{mask}[b*i+b-1]);\\ \text{for (int } j=1;\; (1<< j)<= n \ / \ b;\; j++)\\ \text{for (int } i=0;\; i+(1<< j)<= n \ / \ b;\; i++) \end{array}
             t[n / b * j + i] =
                    op(t[n / b * (j - 1) + i],
t[n / b * (j - 1) + i + i
                             (1 << (j - 1))]);
```

```
int small(int r, int sz = b) {
  return r - msb(mask[r] & ((1 << sz) - 1));</pre>
  T query(int l, int r) {
  if (r - l + 1 <= b)</pre>
     return small(r, r - l + 1);
int ans = op(small(l + b - 1), small(r));
int x = l / b + 1, y = r / b - 1;
if (x < - y) f
     if (x \le y) \{
        int j = msb(y - x + 1);
ans =
              op(ans
                  op(t[n / b * j + x],
                      t[n / b * j + y - (1 << j) + 1]));
     return ans;
  }
};
struct LCA {
  SparseTable<int> st;
   int n;
  vi v, pos, dep;
  LCA(const vi2d &g, int root)
     : n(len(g)), pos(n) { dfs(root, 0, -1, g);
     st = SparseTable<int>(vector<int>(all(dep)));
  void dfs(int i, int d, int p, const vi2d &g) {
  v.eb(len(dep)) = i, pos[i] = len(dep),
     dep.eb(d);
     for (auto j : g[i])
        if (j != p) {
    dfs(j, d + 1,
           dfs(j, d + 1, i, g);
v.eb(len(dep)) = i, dep.eb(d);
  }
  int lca(int a, int b) {
     int l = min(pos[a], pos[b]);
     int r = max(pos[a], pos[b]);
     return v[st.query(l, r)];
   int dist(int a, int b) {
  return dep[pos[a]] + dep[pos[b]] -
               2 * dep[pos[lca(a, b)]];
```

6.23 Lowest Common Ancestor (Binary Lifting)

Description: Given a directed tree, finds the LCA between two nodes using binary lifting, and answer a few queries with it Usage:

• lca: returns the LCA between the two given nodes

• on_path: fids if c is in the path from a to bTime: build $O(N \cdot MAXLOG2)$, all queries O(MAXLOG2)

```
struct LCA {
  int n;
  const int maxlog;
  vector<vector<int>> up;
  vector<int> depth;
  LCA(const vector<vector<int>> &tree)
      : n(tree.size()),
  maxlog(ceil(log2(n)))
         up(n, vector<int>(maxlog + 1)),
    depth(n, -1) { for (int i = 0; i < n; i++) {
         (depth[i] == -1) {
         depth[i] = 0;
         dfs(i, -1, tree);
    }
  void dfs(int u, int p,
    const vector<vector<int>> &tree) { if (p != -1) {
      depth[u] = depth[p] + 1;
      up[u][0] = p;
for (int i = 1; i <= maxlog; i++) {
         up[u][i] = up[up[u][i - 1]][i - 1];
    for (int v : tree[u]) {
```

```
if (v == p) continue;
       dfs(v, u, tree);
  int kth_jump(int u, int k) {
    for (int i = maxlog; i >= 0; i—) {
  if ((1 << i) & k) {
         u = up[u][i];
    return u;
  int lca(int u, int v) {
  if (depth[u] < depth[v]) swap(u, v);</pre>
    int diff = depth[u] - depth[v];
    u = kth_jump(u, diff);
    if (u == v) return u;
    for (int i = maxlog; i \ge 0; i—) {
       if (up[u][i] != up[v][i]) {
  u = up[u][i];
         v = up[v][i];
      }
    }
    return up[u][0];
  bool on_path(int u, int v, int s) {
    int uv = lca(u, v), us = lca(u, s),
vs = lca(v, s);
     return (uv == s or (us == uv and vs == s) or
              (vs == uv and us == s));
  int dist(int u, int v) {
    return depth[u] + depth[v]
             2 * depth[lca(u, v)];
};
```

6.24Maximum flow (Dinic)

Description: Finds the **maximum flow** in a graph network, given the **source** s and the **sink** t. Add edge from a to b with capcity c.

Time: In general $O(E\cdot V^2)$, if every capacity is 1, and every vertice has in degree equal 1 or out degree equal 1 then $O(E \cdot \sqrt{V})$,

Warning: Suffle the edges list for every vertice may take you out of the worst case

```
struct Dinic {
   struct Edge {
      int to, rev;
     ll c, oc;
ll flow() {
        return max(oc - c, OL)
// if you need flows
                                    0LL);
  };
  vi lvl, ptr, q;
  vector<vector<Edge>> adj;
  Dinic(int n) : lvl(n), ptr(n), q(n), adj(n) {}
  void addEdge(int a, int b, ll c, ll rcap = 0) {
   adj[a].pb({b, len(adj[b]), c, c});
   adj[b].pb({a, len(adj[a]) - 1, rcap, rcap});
  ll dfs(int v, int t, ll f) {
  if (v == t || !f) return f;
  for (int &i = ptr[v]; i < len(adj[v]); i++) {</pre>
        Edge &e = adj[v][i];
if (lvl[e.to] == lvl[v] + 1)
  if (ll p = dfs(e.to, t, min(f, e.c))) {
               e.c —= p, adj[e.to][e.rev].c += p;
               return p;
            }
      return 0;
   ll maxFlow(int s, int t) {
      Il flow = 0;
     q[0] = s;
      rep(L, 0, 31) { do { // 'int L=30' maybe faster for random
                  // data
            lvl = ptr = vi(len(q));
            int qi = 0, qe = lvl[s] = 1;
```

```
while (qi < qe \&\& !lvl[t]) {
        int v = q[qi++];
        for (Edge e : adj[v])
           if (!lvl[e.to] && e.c >> (30 - L))
q[qe++] = e.to,
             lvl[e.to] = lvl[v] + 1;
      while (ll p = dfs(s, t, LLONG_MAX))
        flow += p;
    } while (lvl[t]);
  return flow;
}
bool leftOfMinCut(int a) { return lvl[a] != 0; }
```

Minimum Cost Flow

Description: Given a network find the minimum cost to achieve a flow of at most f. Works with **directed** and **undirected** graphs Usage:

• add(u, v, c, w): adds an edge from u to v with capacity c and cost w. • flow(f): return a pair (flow, cost) with the maximum flow until f with source at s and sink at t, with the minimum cost possible. Time: $O(N \cdot M + f \cdot m \log n)$

```
template <typename T>
struct MinCostFlow {
  struct Edge {
    int to;
    ll c, rc;
                // capcity, residual capacity
                 // cost
    Tw;
  int n, s, t;
vector<Edge> edges;
  vi2d g;
  vector<T> dist;
  vi pre;
  MinCostFlow() {}
  MinCostFlow(int n_, int _s, int _
       : n(n_{-}), s(_{-}s), t(_{-}t), g(n) {}
  void addEdge(int u, int v, ll c, T w) {
    g[u].pb(len(edges));
    edges.eb(v, c, 0, w);
    g[v].pb(len(edges));
    edges.eb(u, 0, 0, -w);
  // {flow, cost}
pair<ll, T> flow(ll flow_limit = LLONG_MAX) {
    ll flow = 0;
    T cost = 0;
    while (flow < flow_limit and dijkstra(s, t)) {</pre>
      ll aug = LLONG_MAX;
      aug = min({flow_limit - flow, aug,
                     edges[pre[i]].c});
      for (int i = t; i != s;
    i = edges[pre[i] ^ 1].to) {
         edges[pre[i]].c = aug;
edges[pre[i] ^ 1].c += aug;
         edges[pre[i]].rc += aug;
edges[pre[i] ^ 1].rc -= aug;
       flow += aug;
      cost += (T)aug * dist[t];
    return {flow, cost};
   // Needs to be called after flow method
  vi2d paths() {
    vi2d p;
    for (;;) {
      int cur = s;
      auto \&res = p.eb();
      res.pb(cur);
      while (cur != t) {
         bool found = false;
for (auto i : g[cur]) {
           auto \&[v, \_, c, cost] = edges[i];
           if (c > 0) {
```

```
res.pb(cur = v);
               found = true;
               break;
          }
         if (!found) break;
       if (cur != t) {
          p.ppb();
          break;
     }
     return p;
 private:
  bool bellman_ford(int s, int t) {
  dist.assign(n, numeric_limits<T>::max());
     pre.assign(n, -1);
     vc inq(n, false);
     queue<int> q;
     dist[s] = 0;
     q.push(s);
     while (len(q)) {
       int u = q.front();
       q.pop();
       inq[u] = false;
       for (int i : g[u]) {
  auto [v, c, w, _] = edges[i];
  auto new_dist = dist[u] + w;
  if (c > 0 and dist[v] > new_dist) {
            dist[v] = new_dist;
            pre[v] = i;
            if (not inq[v]) {
               inq[v] = true;
               q.push(v);
         }
       }
     }
     return dist[t] != numeric_limits<T>::max();
  bool dijkstra(int s, int t) {
     dist.assign(n, numeric_limits<T>::max());
     pre.assign(n, -1);
     dist[s] = 0;
     using PQ = pair<T, int>;
     pqmn<PQ> pq;
     pq.emp(0, s);
     while (len(pq)) {
       auto [cost, u] = pq.top();
       pq.pop();
       if (cost != dist[u]) continue;
       for (int i : g[u]) {
  auto [v, c, _, w] = edges[i];
  auto new_dist = dist[u] + w;
          if (c > 0 \text{ and } dist[v] > new_dist) {
            dist[v] = new_dist;
            pre[v] = i;
            pq.emp(new_dist, v);
       }
    }
     return dist[t] != numeric_limits<T>::max();
  }
};
       Minimum Vertex Cover (already divided)
6.26
```

Description: Given a bipartite graph g with n vertices at left and m vertices at right, where g[i] are the possible right side matches of vertex i from left side, find a minimum vertex cover. The size is the same as the size of the maximum matching, and the complement is a maximum independent set.

```
vector<int> min_vertex_cover(
     vector<vector<int>> &g, int n, int m) {
  vector<int> match(m, -1), vis;
  auto find = [\&](auto &&self, int j) \rightarrow bool { if (match[j] == -1) return 1;
     vis[j] = 1;
```

```
int di = match[j]
  for (int e : g[di])
    if (!vis[e] and self(self, e)) {
  match[e] = di;
       return 1;
  return 0;
for (int i = 0; i < (int)g.size(); i++) {
  vis.assign(match.size(), 0);
  match[j] = i;
      break;
  }
}
int
     int)match.size() -
     (int)count(match.begin(), match.end(), -1);
vector<char> lfound(n, true), seen(m);
for (int it : match)
  if (it != -1) lfound[it] = false;
vector<int> q, cover;
for (int i = 0; i < n; i++)
  if (lfound[i]) q.push_back(i);
while (!q.empty()) {
  int i = q.back();
  q.pop_back();
  lfound[i] = 1;
  for (int e : g[i])
    if (!seen[e] and match[e] !=-1) { seen[e] = true;
       q.push_back(match[e]);
for (int i = 0; i < n; i++)
  if (!lfound[i]) cover.push_back(i);
for (int i = 0; i < m; i++)
  if (seen[i]) cover.push_back(n + i);
assert((int)size(cover) == res);
return cover;
```

6.27 Prim (MST)

Description: Given a graph with N vertex finds the minimum spanning tree, if there is no such three returns inf, it starts using the edges that connect with each $s_i \in s$, if none is provided than it starts with the edges of node 0. **Time**: $O(V \log E)$

```
const int MAXN(1 '00' 000);
int N:
vector<pair<ll, int>> G[MAXN];
ll prim(vi s = vi(1, 0)) {
  priority_queue<pair<ll, int>
                   vector<pair<ll, int>>,
                   greater<pair<ll, int>>>
  vector<char> ingraph(MAXN);
  int ingraphcnt(\tilde{0});
  for (auto si : s) {
    ingraphcnt++;
    ingraph[si] = true;
for (auto &[w, v] : G[si]) pq.emplace(w, v);
  ll\ mstcost = 0:
  while (ingraphent < N and !pq.empty()) {</pre>
    ll w;
    int v;
    do {
      tie(w, v) = pq.top();
      pq.pop();
    } while (not pq.empty() and ingraph[v]);
    mstcost += w, ingraph[v] = true, ingraphcnt++;
    for (auto &[w2, v2] : G[v]) {
      pq.emplace(w2, v2);
  return ingraphcnt == N ? mstcost : oo;
```

6.28 Shortest Path With K-edges

}

Description: Given an adjacency matrix of a graph, and a number K computes the shortest path between all nodes that uses exactly K edges, so for $0 \le i, j \le N-1$ ans[i][j] = "the shortest path between i and j that uses exactly K edges, remember to initialize the adjacency matrix with ∞ . **Time**: $O(N^3 \cdot \log K)$

```
template <typename T>
vector<vector<T>> prod(vector<vector<T>> &a,
                            vector<vector<T>> &b) {
  const T _oo = numeric_limits<T>::max();
  int n = a.size();
  vector<vector<T>> c(n, vector<T>(n, _oo));
  for (int i = 0; i < n; i++)
for (int j = 0; j < n; j++)
for (int k = 0; k < n; k++)
         if (a[i][k] != _oo and b[k][j] != _oo)
            c[i][j] =
                 min(c[i][j], a[i][k] + b[k][j]);
  return c;
}
template <typename T>
vector<vector<T>> shortest_with_k_moves(
   vector<vector<T>> adj, long long k) {
  if (k == 1) return adj;
  auto ans = adi;
  while (k)
    if (k \& 1) ans = prod(ans, adj);
    k >>= 1;
    adj = prod(adj, adj);
  return ans;
}
```

6.29 Strongly Connected Components (struct)

Description: Find the connected component for each edge (already in a topological order), some additional functions are also provided. **Time:** Build: O(V+E)

```
struct SCC {
  int n, num_sccs;
  vi2d adj;
  vi scc_id;
  SCC(int _n)
: n(_n)
         num_sccs(0),
          adj(n)
          scc_id(n, -1) {}
  void add_edge(int u, int v) { adj[u].eb(v); }
  void find_sccs() {
    int timer = 1;
    vị tin(n),
     st.reserve(n)
     function<int(int)> dfs = [\&](int u) \rightarrow int {
       int low = tin[u] = timer++, siz = len(st);
       st.eb(u);
       for (int v : adj[u])
          if
             (scc_id[v] < 0)
                 min(low, tin[v] ? tin[v] : dfs(v));
       if (tin[u] == low)
          rep(i, siz, len(st)) scc_id[st[i]] =
   num_sccs;
         st.resize(siz);
num_sccs++;
       return low;
    };
    for (int i = 0; i < n; i++)
  if (!tin[i]) dfs(i);</pre>
  vector<set<int>> build_gscc() {
    vector<set<int>> gscc;
    for (int i = 0; i < len(adj); ++i)
  for (auto j : adj[i])
    if (scc_id[i] != scc_id[j])</pre>
            gscc[scc_id[i]].emplace(scc_id[j]);
```

```
return gscc;
  }
  vi2d per_comp() {
    vi2d_ret(num_sccs);
    rep(i, 0, n) ret[scc_id[i]].eb(i);
    reverse(all(
        ret));
               // already in topological order ;)
    return ret;
};
```

6.30Topological Sorting (Kahn)

Description: Finds the topological sorting in a DAG, if the given graph is not a DAG than an empty vector is returned, need to 'initialize' the INCNT as you build the graph.

Time: O(V + E)

```
const int MAXN(2 '00' 000);
int INCNT[MAXN];
vi2d GOUT(MAXN);
int N;
vi toposort() {
  vi order
  queue<int> q;
  for (int i = 0; i < N; i++)
    if (!INCNT[i]) q.emplace(i);
  while (!q.empty())
    auto u = q.front();
    q.pop();
    order.emplace_back(u);
    for (auto v : GOUT[u]) {
      if (INCNT[v] == 0) q.emplace(v);
    }
  return len(order) == N ? order : vi();
```

6.31 Topological Sorting (Tarjan)

Description: Finds a the topological order for the graph, if there is no such order it means the graph is cyclic, then it returns an empty vector Time: O(V+E)

```
const int maxn(1 '00' 000);
int n. m:
vi g[maxn];
int not_found = 0, found = 1, processed = 2;
int state[maxn];
bool dfs(int u, vi &order) {
  if (state[u] == processed) return true;
  if (state[u] == found) return false;
  state[u] = found;
  for (auto v : g[u]) {
     if (not dfs(v, order)) return false;
  state[u] = processed;
  order.emplace_back(u);
  return true;
vi topo_sort() {
  vi order
  memset(state, 0, sizeof state);
  for (int u = 0; u < n; u++) {
     if (state[u] == not_found and
    not dfs(u, order))
        return {};
  reverse(all(order));
  return order;
```

Tree Isomorphism (not rooted)

Description: Two trees are considered isomorphic if the hash given by thash() is the same.

```
Time: O(V \cdot \log V)
```

```
map<vi, int> mphash;
struct Tree {
  int n;
  vi2d g;
  vi sz, cs;
  Tree(int n_-) : n(n_-), g(n), sz(n) {}
  void add_edge(int u, int v) {
    g[u].emplace_back(v);
     g[v].emplace_back(u);
  void_dfs_centroid(int v, int p) {
     sz[v] = 1;
     bool cent = true;
     for (int u : g[v])
if (u != p) {
         dfs_centroid(u, v);
         sz[v] += sz[u]
         cent &= not(sz[u] > n / 2);
    if
        (cent and n - sz[v] \ll n / 2)
       cs.push_back(v);
  int fhash(int v, int p) {
    vi h;
for (int u : g[v])
       if (u != p) h.push_back(fhash(u, v));
     sort(all(h));
     if (!mphash.count(h))
       mphash[h] = mphash.size();
     return mphash[h];
  ll thash()
     cs.clear()
    dfs_centroid(0, -1);
if (cs.size() == 1) return fhash(cs[0], -1);
    ll h1 = fhash(cs[0], cs[1]),
h2 = fhash(cs[1], cs[0]);
return (min(h1, h2) << 30ll) + max(h1, h2);
};
```

Tree Isomorphism (rooted)

Description: Given a rooted tree find the hash of each subtree, if two roots of two distinct trees have the same hash they are considered isomorphic **Time**: hash first time in $O(\log N_v \cdot N_v)$ where (N_v) is the of the subtree of v

```
map<vi, int> hasher;
int hs = 0;
struct RootedTreeIso {
  int n;
  vi2d adj;
  vi hashes;
  RootedTreeIso(int _n)
        : n(_n), adj(_n), hashes(_n, -1){};
  void add_edge(int u, int v) {
     adj[u].emplace_back(v)
     adj[v].emplace_back(u);
  int hash(int u, int p = -1) {
  if (hashes[u] != -1) return hashes[u];
     vi children;
for (auto v : adj[u])
  if (v != p)
          children.emplace_back(hash(v, u));
     sort(all(children));
     if (!hasher.count(children))
  hasher[children] = hs++;
     return hashes[u] = hasher[children];
};
```

Tree diameter (DP) 6.34

```
const int MAXN(1 '000' 000);
int N;
vi G[MAXN];
int diameter, toLeaf[MAXN];
void calcDiameter(int u = 0, int p = -1) {
  int d1, d2;
```

```
d1 = d2 = -1;
   for (auto v : G[u]) {
      if (v != p)
         calcDiameter(v, u);
         d1 = max(d1, toLeaf[v]);
tie(d1, d2) = minmax({d1, d2});
   toLeaf[u] = d2 + 1;
   diameter = max(diameter, d1 + d2 + 2);
     Math
7.1 Arithmetic Progression Sum
Usage:
    • s: first term
• d: common difference
• n: number of terms
ll arithmeticProgressionSum(ll s, ll d, ll n) {
   return (s + (s + d * (n - 1))) * n / 2ll;
7.2 Binomial
Time: O(N \cdot K)
Memory: O(K)
ll binom(ll n, ll k) {
   if (k > n) return 0;
vll dp(k + 1, 0);
   dp[0] = 1;
  for (ll i = 1; i <= n; i++)
  for (ll j = k; j > 0; j—)
    dp[j] = dp[j] + dp[j - 1];
   return dp[k];
7.3 Binomial MOD
Description: find \binom{n}{k} (mod MOD)
Time:
• precompute: on first call it takes O(MAXNBIN) to precompute the factorials
• query: O(1).

Memory: O(MAXNBIN)

Warning: Remember to set MAXNBIN properly!
const ll\ MOD = 998244353;
inline ll binom(ll n, ll k) {
  static const int BINMAX = 2'000'000;
   static vll FAC(BINMAX + 1), FINV(BINMAX + 1);
static bool done = false;
   if (!done)
      vli INV(BINMAX + 1);
FAC[0] = FAC[1] = INV[1] = FINV[0] = FINV[1] =
      for (int i = 2; i <= BINMAX; i++) {
  FAC[i] = FAC[i - 1] * i % MOD;
  INV[i] = MOD - MOD / i * INV[MOD % i] % MOD;
  FINV[i] = FINV[i - 1] * INV[i] % MOD;
}</pre>
      done = true;
   if (n < k || n < 0 || k < 0) return 0;
   return FAC[n] * FINV[k] % MOD * FINV[n - k] %
              MOD;
```

7.4 Chinese Remainder Theorem

```
Description: Find the solution X to the N modular equations. x\equiv a_1(modm_1)\\ \cdots\\ x\equiv a_n(modm_n) (1)
```

The m_i don't need to be coprime, if there is no solution then it returns -1.

```
tuple<ll, ll, ll> ext_gcd(ll a, ll b) {
   if (!a) return {b, 0, 1};
   auto [g, x, y] = ext_gcd(b % a, a);
   return {g, y - b / a * x, x};
}
template <typename T = ll>
struct crt {
```

```
T a, m;
  crt() : a(0), m(1) \{ \} 

crt(T a_, T m_) : a(a_), m(m_) \{ \} 
  crt operator*(crt C) {
     auto [g, x, y] = ext_gcd(m, C.m);
     if ((a - C.a) % g != 0) a = -1;
if (a == -1 \text{ or } C.a == -1) return crt(-1, 0);
     T \stackrel{\cdot}{lcm} = m / g * C.m;
     T ans =
     a + (x * (C.a - a) / g % (C.m / g)) * m;
return crt((ans % lcm + lcm) % lcm, lcm);
};
template <typename T = ll >
struct Congruence {
  T a, m;
template <typename T = ll >
T chinese_remainder_theorem(
     const vector<Congruence<T>> &equations) {
  crt<T> ans;
  for (auto &[a_, m_] : equations) {
     ans = ans * crt<T>(a_, m_);
  return ans.a;
}
```

7.5 Derangement / Matching Problem

```
Description: Computes the derangement of N, which is given by the formula: D_N = N! \left(1 - \frac{1}{1!} + \frac{1}{2!} - \frac{1}{3!} + \ldots + (-1)^N \frac{1}{N!}\right)
Time: O(N)
```

```
#warning Remember to call precompute !
const ll MOD = 1e9 + 7;
const int MAXN(1 '000' 000);
ll fats[MAXN + 1];
void precompute()
   fats[0] = 1;
for (ll i = 1; i <= MAXN; i++) {
  fats[i] = (fats[i - 1] * i) % MOD;</pre>
}
ll fastpow(ll a, ll p, ll m) {
   ll ret = 1;
   while (p) {
     if (p \& 1) ret = (ret * a) % MOD;
     p >>= 1;
     a = (a * a) % MOD;
   return ret;
}
ll divmod(ll a, ll b) {
   return (a * fastpow(b, MOD - 2, MOD)) % MOD;
ll derangement(const ll n) {
   ll ans = fats[n];
for (ll i = 1; i <= n; i++) {</pre>
     ll k = divmod(fats[n], fats[i]);
     if (i & 1) {
     ans = (ans - k + MOD) % MOD;
} else {
        ans = (ans + k) % MOD;
   return ans;
}
```

7.6 Euler Phi $\varphi(N)$

Description: Computes the number of positive integers less than N that are coprimes with N, in $O(\sqrt{N})$.

```
int phi(int n) {
   if (n == 1) return 1;
   auto fs = factorization(
        n);   // a vctor of pair or a map
   auto res = n;
   for (auto [p, k] : fs) {
      res /= p;
      res *= (p - 1);
```

```
}
return res;
}
```

7.7 Euler phi $\varphi(N)$ (in range)

Description: Computes the number of positive integers less than n that are coprimes with N, in the range [1, N], in $O(N \log N)$.

```
const int MAX = 1e6;
vi range_phi(int n) {
  bitset<MAX> sieve;
  vi phi(n + 1);
  iota(phi.begin(), phi.end(), 0);
  sieve.set();
  for (int p = 2; p <= n; p += 2) phi[p] /= 2;
  for (int p = 3; p <= n; p += 2) {
    if (sieve[p]) {
      for (int j = p; j <= n; j += p) {
         sieve[j] = false;
         phi[j] /= p;
         phi[j] *= (p - 1);
      }
    }
  }
  return phi;
}</pre>
```

7.8 FFT convolution and exponentiation

```
const ld PI = acos(-1);
/* change the ld to doulbe may increase
 * performance =D */
struct num {
   ld a{0.0}, b{0.0};
  num() {}
num(ld na) : a{na} {}
num(ld na, ld nb) : a{na}, b{nb} {}
  const num operator+(const num &c) const {
     return num(a + c.a, b + c.b);
  const num operator—(const num &c) const {
     return num(a - c.a, b - c.b);
  const num operator*(const num &c) const {
     return num(a * c.a - b * c.b,
                   a * c.b + b * c.a);
  const num operator/(const ll &c) const {
     return num(a / c, b / c);
void fft(vector<num> &a, bool invert) {
  int n = len(a);
for (int i = 1, j = 0; i < n; i++) {
     int bit = n \gg 1;
     for (; j & bit; bit >>= 1) j ^= bit;
       ^= bit;
     if (i < j) swap(a[i], a[j]);
  for (int sz = 2; sz <= n; sz <<= 1) {
  ld ang = 2 * PI / sz * (invert ? -1 : 1);
  num wsz(cos(ang), sin(ang));
  fan (int in ang);</pre>
     for (int i = 0; i < n; i += sz) {
       num w(1);
        rep(j, 0, sz / 2)  {
          num u = a[i + j],
v = a[i + j + sz / 2] * w;
          a[i + j] = u + v;
a[i + j + sz / 2] = u - v;
w = w * wsz;
     }
  if (invert)
     for (num \& x : a) x = x / n;
vi conv(vi const a, vi const b) {
  vector<num> fa(all(a));
  vector<num> fb(all(b));
```

```
int n = 1; while (n < len(a) + len(b)) n <<= 1;
  fa.resize(n)
  fb.resize(n)
  fft(fa, false);
fft(fb, false);
rep(i, 0, n) fa[i] = fa[i] * fb[i];
  fft(fa, true);
  vi result(n);
rep(i, 0, n) result[i] = round(fa[i].a);
  while (len(result) and result.back() == 0)
    result.pop_back();
  /* Unconment this line if you want a boolean
   * convolution*/
  // for (auto &xi : result) xi = min(xi, 1ll);
  return result;
vll poly_exp(vll &ps, int k) {
  vll ret(len(ps));
  auto base = ps;
  ret[0] = 1;
  while (k) {
  if (k & 1) ret = conv(ret, base);
    k >>= 1;
    base = conv(base, base);
  return ret;
```

7.9 Factorial Factorization

Description: Computes the factorization of N! in $\varphi(N) * \log N$ Time: $O(\varphi(N) \cdot \log N)$

```
ll E(ll n, ll p) {
    ll k = 0, b = p;
    while (b <= n) {
        k += n / b;
        b *= p;
    }
    return k;
}
map<ll, ll> factorial_factorization(
        ll n, const vll &primes) {
        map<ll, ll> fs;
        for (const auto &p : primes) {
            if (p > n) break;
            fs[p] = E(n, p);
        }
        return fs;
}
```

7.10 Factorization

Description: Computes the factorization of N. **Time**: $O(\sqrt{n})$.

```
map<ll, ll> factorization(ll n) {
   map<ll, ll> ans;
   for (ll i = 2; i * i <= n; i++) {
        ll count = 0;
        for (; n % i == 0; count++, n /= i)
        if (count) ans[i] = count;
   }
   if (n > 1) ans[n]++;
   return ans;
}
```

7.11 Factorization (Pollard's Rho)

Description: Factorizes a number into its prime factors. **Time:** $O(N^{(\frac{1}{4})} * \log(N))$.

```
ll mul(ll a, ll b, ll m) {
    ll ret =
        a * b - (ll)((ld)1 / m * a * b + 0.5) * m;
    return ret < 0 ? ret + m : ret;
}
ll pow(ll a, ll b, ll m) {
    ll ans = 1;
    for (; b > 0; b /= 2ll, a = mul(a, a, m)) {
        if (b % 2ll == 1) ans = mul(ans, a, m);
}
```

```
return ans;
bool prime(ll n) {
  if (n < 2) return 0;
if (n <= 3) return 1;
if (n % 2 == 0) return 0;
  ll x = pow(a, d, n);
     if (x == 1 \text{ or } x == n - 1 \text{ or a } \% n == 0)
       continue:
    for (int j = 0; j < r - 1; j++) { x = mul(x, x, n); if (x == n - 1) break;
    if (x != n - 1) return 0;
  return 1:
}
ll rho(ll n) {
  if (n == 1 or prime(n)) return n;
  auto f = [n](ll x) \{ return mul(x, x, n) + 1; \};
  ll x = 0, y = 0, t = 30, prd = 2, x0 = 1, q;
  while (t \% 40 != 0 \text{ or gcd(prd, n)} == 1) {
     if (x == y) x = ++x0, y = f(x);
     q = mul(prd, abs(x - y), n);
    if (q != 0) prd = q;
x = f(x), y = f(f(y)), t++;
  return gcd(prd, n);
}
vector<ll> fact(ll n) {
  if (n == 1) return {};
  if (prime(n)) return {n};
  il d = rho(n);
  vector<ll> l =
  vector<ll> l = fact(d), r = fact(n / d);
l.insert(l.end(), r.begin(), r.end());
  return l;
7.12 Fast Pow
Description: Computes a^b \pmod{m}
Time: O(\log B).
ll fpow(ll a, ll b, ll m) {
  ll ret = 1;
  while (b) {
   if (b & 1) ret = (ret * a) % m;
     b > = 1;
    a = (a * a) % m;
```

```
return ret;
```

7.13 Find diophantine equation solution

Description: Given a b, c finds the solution to the equation ax + by = c, the result will be stored in the reference variables x0 and y0. Time: $O(\log min(a, b))$

```
template <typename T>
tuple<T, T, T> ext_gcd(T a, T b) {
  if (b == 0) return {a, 1, 0};
  auto [d, x1, y1] = ext_gcd(b, a % b);
  return {d, y1, x1 - y1 * (a / b)};
template <typename T>
tuple<bool, T, T> find_any_solution(T a, T b,
                                           T c) {
  assert(a != 0 or b != 0);
#warning Be careful with overflow, use __int128 if
    needed!
  auto [d, x0, y0] =
       ext_{gcd}(a < 0 ? -a : a, b < 0 ? -b : b);
  if (c % d) return {false, 0, 0};
  x0 *= c / d;
y0 *= c / d;
  if (a < 0) x0 = -x0;
```

```
if (b < 0) y0 = -y0;
return {true, x0, y0};
```

7.14 Find linear recurrence (Berlekamp-Massey)

 ${\bf Description} .$ Given the first N terms of a linear recurrence finds the smallest recurrence that matches the sequence.

Time: $O(N^2)$

Warning: Works faster if the mod is const but can be also be a parameter. Absolute magic!

```
const ll \ mod = 998244353;
ll modpow(ll b, ll e) {
  ll ans = 1;
for (; e; b = b * b % mod, e /= 2)
  if (e & 1) ans = ans * b % mod;
   return ans;
}
vl berlekampMassey(vll s) {
  int n = len(s), L = 0, m = 0;
if (!n) return {};
vll C(n), B(n), T;
   C[0] = B[0] = 1;
   ll b = 1;
   rep(i, 0, n) {
     ++m;
ll d = s[i] % mod;
     rep(j, 1, L + 1) d = (d + C[j] * s[i - j]) % mod;
     if (!d) continue;
     ll coef = d * modpow(b, mod - 2) % mod;
     rep(j, m, n) C[j] =
           (C[j] - coef * B[j - m]) % mod;
     if (2 * L > i) continue;
L = i + 1 - L;
     B = T;
     b = d;
     m = 0:
  C.resize(L + 1);
C.erase(C.begin());
  for (ll &x : C) x = (mod - x) \% mod;
   return C;
```

7.15 Find multiplicatinve inverse

```
ll inv(ll a, ll m) {
  return a > 1ll ? m - inv(m % a, a) * m / a
                        : 1ll;
}
```

7.16 GCD

```
ll gcd(ll a, ll b) {
  return b ? gcd(b, a % b) : a;
}
```

7.17 Gauss XOR elimination / XOR-SAT

Description: Execute gaussian elimination with xor over the system Ax = bin. The add method must receive a bitset indicating which variables are present in the equation, and the solution of the equation.

 $\underline{\mathbf{Time}} : O(\tfrac{nm^2}{64})$

```
const int MAXXI = 2009;
using Equation = bitset<MAXXI>;
struct GaussXor {
  vector<char> B;
  vector<Equation> A;
  void add(const Equation &ai, bool bi) {
    A.push_back(ai);
    B.push_back(bi);
  pair<bool, Equation> solution() {
    int cnt = 0, n = A.size();
    Equation vis;
    vis.set();
```

```
Equation x;
for (int j = MAXXI - 1, i; j >= 0; j—) {
  for (i = cnt; i < n; i++) {
    if (A[i][j]) break;
}</pre>
  if (i == n) continue;
  swap(A[i], A[cnt]), swap(B[i], B[cnt]);
  i = cnt++
  vis[j] = 0;
  for (int k = 0; k < n; k++) {
  if (i == k || !A[k][j]) continue;</pre>
     A[k] ^= A[i];
     B[k] ^= B[i];
  = vis;
for (int i = 0; i < n; i++) {
  int acum = 0;
for (int j = 0; j < MAXXI; j++) {
     if (!A[i][j]) continue;
     if (!vis[j]) {
       vis[j] = 1;
       x[j] = acum ^ B[i];
     acum = x[j];
  if (acum != B[i])
     return {false, Equation()};
return {true, x};
```

7.18 Integer partition

Description: Find the total of ways to partition a given number N in such way that none of the parts is greater than K. Time: $O(N \cdot min(N, K))$

Memory: O(N)

Warning: Remember to memset everything to -1 before using it

```
const ll MOD = 1000000007;
const int MAXN(100);
ll memo[MAXN + 1];
ll dp(ll n, ll k = oo)
  if (n == 0) return 1;
  ll &ans = memo[n]
  if (ans !=-1) return ans;
  for (int i = 1; i <= min(n, k); i++) {
    ans = (ans + dp(n - i, k)) % MOD;
  return ans;
```

7.19 LCM

```
ll gcd(ll a, ll b) {
  return b ? gcd(b, a % b) : a;
1 lcm(ll a, ll b) { return a / gcd(a, b) * b; }
```

7.20 Linear Recurrence

Description: Find the n-th term of a linear recurrence, given the recurrence rec and the first K values of the recurrence, remember that first_k[i] is the value of f(i), considering 0-indexing.

Time: $O(K^3 \log N)$

```
template <typename T>
vector<vector<T>> prod(vector<vector<T>> &a,
                                  vector<vector<T>> &b,
                                 const ll mod) {
   int n = a.size();
   vector<vector<T>>> c(n, vector<T>(n));
  for (int i = 0; i < n; i++) {
  for (int j = 0; j < n; j++) {
    for (int k = 0; k < n; k++) {
        c[i][j] = (c[i][j] +</pre>
                            ((a[i][k] * b[k][j]) % mod)) %
                          mod:
     }
```

```
}
  return c;
}
template <typename T>
vector<vector<T>> fpow(vector<vector<T>> &xs,
                        ll p, ll mod) {
  vector<vector<T>> ans(xs.size()
                         vector<T>(xs.size()));
  for (int i = 0; i < (int)xs.size(); i++)
    ans[i][i] = 1;
  for (auto b = xs; p;
       p \gg 1, b = prod(b, b, mod)
    if (p \& 1) ans = prod(ans, b, mod);
  return ans;
}
ll linear_req(vector<vector<ll>>> rec,
              vector<ll> first_k, ll n, ll mod) {
  int k = first_k.size();
  if (n < k) {
    return first_k[n];
  ll n2 = n - k + 1;
  rec = fpow(rec, n2, mod);
  ll ret = 0;
  for (int i = 0; i < k; i++) {
    ret = (ret +
           (rec.back()[i] * first_k[i]) % mod) %
          mod:
  return ret;
```

7.21 List N elements choose K

Description: Process every possible combination of K elements from Nelements, thoose index marked as 1 in the index vector says which elments are choosed at that moment.

Time: $O({N \choose K} \cdot O(process))$

```
void process(vi &index) {
  for (int i = 0; i < len(index); i++) {
    if (index[i])
      cout \ll i \ll " \n"[i == len(index) - 1];
}
void n_choose_k(int n, in k) {
  vi index(n);
  fill(index.end() - k, index.end(), 1);
  do {
    process(index);
   while (next_permutation(all(index)));
```

7.22 List primes (Sieve of Eratosthenes)

```
const ll\ MAXN = 2e5;
vll list_primes(ll n = MAXN) {
  vll ps;
  bitset<MAXN + 1> sieve;
  sieve.set();
  sieve.reset(1);
for (ll i = 2; i <= n; ++i) {
   if (sieve[i]) ps.push_back(i);</pre>
     for (ll j = i * 2; j <= n; j += i) {
        sieve.reset(j);
  return ps;
}
```

7.23Matrix exponentiation

```
const ll MOD = 1 '000' 000'007;
template <typename T>
vector<vector<T>> prod(vector<vector<T>> &a
                       vector<vector<T>> &b) {
  int n = len(a)
  vector<vector<T>> c(n, vector<T>(n));
```

```
for (int i = 0; i < n; i++) {
  for (int j = 0; j < n; j++) {
    for (int k = 0; k < n; k++) {
      c[i][j] = (c[i][j] +</pre>
                         ((a[i][k] * b[k][j]) % MOD)) %
     }
  }
  return c;
template <typename T>
vector<vector<T>> fpow(vector<vector<T>> &xs,
                              ll p) {
  vector<vector<T>> ans(len(xs)
                                vector<f>(len(xs)))
  for (int i = 0; i < len(xs); i++) ans[i][i] = 1;
  auto b = xs;
  while (p) { if (p & 1) ans = prod(ans, b);
     p >>= 1;
     b = prod(b, b);
  return ans;
7.24 NTT integer convolution and exponentiation
    • Convolution O(N \cdot \log N),
   • Exponentiation: O(\log K \cdot N \cdot \log N)
template <int _mod>
struct mint
  ll expo(ll b, ll e) {
```

```
ll ret = 1;
  while (e) {
   if (e % 2) ret = ret * b % _mod;
     e /= 2, b = b * b % _mod;
  return ret;
ll inv(ll b) { return expo(b, \_mod - 2); }
using m = mint;
ll v;
mint() : v(0) {}
mint(ll v_) {
   if (v_ >= _mod or v_ <= __mod) v_ %= _mod;
   if (v_ < 0) v_ += _mod;
   v = v_;</pre>
m &operator+=(const m &a) {
  v += a.v;
if (v >= _mod) v -= _mod;
  return *this;
m &operator—=(const m &a) {
  v = a.v;
if (v < 0) v += _mod;
  return *this;
m &operator*=(const m &a) {
  v = v * ll(a.v) % _mod;
  return *this;
m &operator/=(const m &a) {
  v = v * inv(a.v) % _mod;
  return *this;
m operator_() { return m(-v); }
m & operator^=(ll e) {
  if (e < 0) {
     v = inv(v);
  \dot{v} = \exp(v, e);
  // possivel otimizacao:
   // cuidado com 0^0
  // v = \exp(v, e^{(p-1)});
  return *this;
bool operator==(const m &a) { return v == a.v; }
bool operator!=(const m &a) { return v != a.v; }
friend istream &operator>>(istream &in, m &a) {
```

```
ll val;
     in >> val
     a = m(val);
     return in;
  friend ostream &operator<<(ostream &out, m a) {
     return out << a.v;
  friend m operator+(m a, m b) { return a += b; friend m operator-(m a, m b) { return a -= b;
  friend m operator—(m a, m b)
  friend m operator*(m a, m b) { return a *= b; }
  friend m operator/(m a, m b) { return a /= b; }
friend m operator^(m a, ll e) { return a ^= e; }
};
const ll MOD1 = 998244353;
const ll MOD2 = 754974721;
const ll\ MOD3 = 167772161;
template <int _mod>
void ntt(vector<mint<_mod>> &a, bool rev) {
  int n = len(a);
  auto b = a;
  assert(!(n & (n-1)));
  mint<_mod> g = 1;
while ((g ^ (_mod / 2)) == 1) g += 1;
  if (rev)^{\overline{g}} = 1 / g;
  for (int step = n / 2; step; step /= 2) {
  mint<_mod> w = g_^ (_mod / (n / step)),
                  wn = 1;
     for (int i = 0; i' < n / 2; i += step) {
       for (int j = 0; j < step; j++) {
          auto u = a[2 * i + j],
                v = wn * a[2 * i + j + step];
          b[i + j] = u + v;
          b[i + n / 2 + j] = u - v;
       \dot{w}n = wn * w;
     swap(a, b);
  if (rev) {
  auto n1 = mint<_mod>(1) / n;
     for (auto \&x : a) x *= n1;
template <ll _mod>
vector<mint<_mod>> convolution(
     const vector<mint<_mod>> &a
     const vector<mint<_mod>> &b) {
  vector<mint<_mod>> l(all(a)), r(all(b));
int N = len(l) + len(r) - 1, n = 1;
  while (n \le N) n *= 2;
  l.resize(n), r.resize(n);
  ntt(l, false), ntt(r, false);
for (int i = 0; i < n; i++) l[i] *= r[i];
ntt(l, true);</pre>
  l.resize(N);
  // Uncommnent for a boolean convolution :)
  for (auto& li : l)
    li.v = min(li.v, 1ll);
  */
  return l;
template <ll _mod>
vector<mint<_mod>> poly_exp(
     vector<mint<_mod>> &ps, int k) {
  vector<mint<_mod>> ret(len(ps));
  auto base = ps;
  ret[0] = 1;
  while (k) {
     if (k & 1) ret = convolution(ret, base);
     k >>= 1;
     base = convolution(base, base);
  return ret;
```

7.25 NTT integer convolution and exponentiation (2) mods) modules)

Description: Computes the convolution between the two polynomials and.

```
Time: O(N \log N)
Warning: This is pure magic!
template <int _mod>
struct mint {
  ll expo(ll b, ll e) {
     ll ret = 1;
     while (e) {
   if (e % 2) ret = ret * b % _mod;
        e /= 2, b' = b * b % _mod;
     return ret;
   ll inv(ll b) { return expo(b, \_mod - 2); }
  using m = mint;
  mint() : v(0) {}
  mint(ll v_) {
     if (v_ >= _mod or v_ <=
if (v_ < 0) v_ += _mod;
v = v_;</pre>
                   m &operator+=(const m &a) {
     v += a.v;
if (v >= _mod) v -= _mod;
     return *this;
  m &operator—=(const m &a) {
     v = a.v;
if (v < 0) v += _mod;
     return *this;
  m &operator*=(const m &a) {
     v = v * ll(a.v) % _mod;
     return *this;
  m &operator/=(const m &a) {
     v = v * inv(a.v) % _mod;
     return *this;
  m operator—() { return m(-v); }
  m &operator^=(ll e) {
     if(e < 0)
        v = inv(v);
e = -e;
     v = expo(v, e);
// possivel otimizacao:
     // cuidado com 0^0
// v = expo(v, e%(p-1));
     return *this;
  bool operator==(const m &a) { return v == a.v;
bool operator!=(const m &a) { return v != a.v;
  friend istream &operator>>(istream &in, m &a) {
     ll val;
     in >> val
     a = m(val);
     return in;
  friend ostream &operator<<(ostream &out, m a) {</pre>
     return out << a.v;
   friend m operator+(m a, m b) { return a += b; }
  friend m operator-(m a, m b) { return a -= b; }
friend m operator*(m a, m b) { return a *= b; }
friend m operator/(m a, m b) { return a /= b; }
friend m operator^(m a, ll e) { return a ^= e; }
const ll MOD1 = 998244353;
const ll MOD2 = 754974721;
const ll MOD3 = 167772161;
template <int _mod>
void ntt(vector<mint<_mod>> &a, bool rev) {
  int n = len(a);
  auto b = a;

assert(!(n & (n - 1)));

mint<_mod> g = 1;

while ((g ^ (_mod / 2)) == 1) g += 1;

if (row) a = 1 / a;
  if (rev)^{\circ}g = 1 / g;
```

for (int step = n / 2; step; step /= 2) {

```
mint < mod > w = g ^ (mod / (n / step)),
     for (int i = 0; i < n / 2; i += step) {
       for (int j = 0; j < step; j++) {
  auto u = a[2 * i + j],
                v = wn * a[2 * i + j + step];
          b[i + j] = u + v;
          b[i + n / 2 + j] = u - v;
       \dot{w}n = wn * w;
     swap(a, b);
  if (rev) {
  auto n1 = mint<_mod>(1) / n;
     for (auto \&x : a) x *= n1;
}
tuple<ll, ll, ll> ext_gcd(ll a, ll b) {
  if (!a) return {b, 0, 1};
  auto [g, x, y] = ext_gcd(b % a, a);
   return \{g, y - b / a * x, x\};
template <typename T = ll >
struct crt {
  Ta, m;
  crt(\underline{)} : a(\underline{0}), m(1) {}
  crt(T a_, T m_) : a(a_), m(m_) {}
crt operator*(crt C) {
     auto [g, x, y] = ext\_gcd(m, C.m);
if ((a - C.a) % g != 0) a = -1;
     if (a == -1 \text{ or } C.a == -1) \text{ return } crt(-1, 0);

\underline{T} lcm = m / g * C.m;
          a + (x * (C.a - a) / g % (C.m / g)) * m;
     return crt((ans % lcm + lcm) % lcm, lcm);
};
template <typename T = ll>
struct Congruence {
  T a, m;
template <typename T = ll >
T chinese_remainder_theorem(
     const vector<Congruence<T>> &equations) {
   crt<T> ans;
   for (auto &[a_, m_] : equations) {
     ans = ans * crt<T>(a_, m_);
  return ans.a;
}
#define int long long
template <ll m1, ll m2>
vll merge_two_mods(const vector<mint<m1>> &a,
                       const vector<mint<m2>> &b) {
   int n = len(a);
   vll ans(n);
   for (int i = 0; i < n; i++) {
     auto cur = crt<ll>();
     auto ai = a[i].v;
     auto bi = b[i].v
     cur = cur * crt<ll>(ai, m1)
     cur = cur * crt<ll>(bi, m2);
     ans[i] = cur.a;
  return ans;
}
l.resize(n), r.resize(n);
  ntt(l, false), ntt(r, false);
for (int i = 0; i < n; i++) l[i] *= r[i];
ntt(l, true);</pre>
   l.resize(N)
   vector<mint<MOD2>> l2(all(a)), r2(all(b));
  l2.resize(n), r2.resize(n);
ntt(l2, false), ntt(r2, false);
rep(i, 0, n) l2[i] *= r2[i];
```

```
ntt(l2, true);
l2.resize(N);
return merge_two_mods(l, l2);
}
vll poly_exp(const vll &xs, ll k) {
  vll ret(len(xs));
  ret[0] = 1;
  auto base = xs;
  while (k) {
    if (k & 1) ret = convolution_2mods(ret, base);
    k >>= 1;
    base = convolution_2mods(base, base);
}
return ret;
}
```

7.26 Polyominoes

Usage: buildPolyominoes(x) creates every polyomino until size x, and put it in polyominoes[x], access polyomino.v to find the vector of pairs representing the coordinates of each piece, considering that the polyomino was 'rooted' in coordinate (0,0).

Warning: note that when accessing polyominoes[x] only the first x coordinates are valid.

```
const int MAXP = 10;
using pii = pair<int, int>;
// This implementation considers the rotations as
// distinct
0, 10, 10+9, 10+9+8...
int pos[11] = {0, 10, 19, 27, 34, 40,
45, 49, 52, 54, 55};
struct Polyominoes {
   pii v[MAXP];
ll id;
   int n;
   Polyominoes() {
     n = 1;
v[0] = {0, 0};
normalize();
  pii &operator[](int i) { return v[i]; }
bool add(int a, int b) {
  for (int i = 0; i < n; i++)
    if (v[i].first == a and v[i].second == b)</pre>
            return false
      v[n++] = pii(a, b);
      normalize();
      return true;
   void normalize() {
      int mnx = 100, mny = 100;
      for (int i = 0; i < n; i++)
  mnx = min(mnx, v[i].first),
  mny = min(mny, v[i].second);</pre>
      id = 0;
      for (int i = 0; i < n; i++) {
  v[i].first = mnx, v[i].second = mny;
  id |= (1LL << (pos[v[i].first] +</pre>
                               v[i].second));
  }
vector<Polyominoes> polyominoes[MAXP + 1];
void buildPolyominoes(int mxN = 10) {
  vector<pair<int, int>> dt(
    {{1, 0}, {-1, 0}, {0, -1}, {0, 1}});
for (int i = 0; i <= mxN; i++)</pre>
      polyominoes[i].clear();
   Polyominoes init;
   queue<Polyominoes> q;
   unordered_set<int64_t> used;
   q.push(init)
   used.insert(init.id);
   while (!q.empty()) {
      Polyominoes u = q.front();
      q.pop();
      polyominoes[u.n].push_back(u);
      if (u.n == mxN) continue;
for (int i = 0; i < u.n; i++) {
         for (auto [dx, dy] : dt) {
  Polyominoes to = u;
            bool ok = to.add(to[i].first + dx
                                     to[i].second + dy);
```

8 Outside

8.1 alien_trick

```
int n, k, l;
string s;
pi solve(vector<int> &v, int lambda) {
   // associar um custo lambda para ser subtraido
   // quando realizamos uma çãoperao dp[i] — melhor
   // profit que tivemos considerando as i
   // primeiras çõposies cnt[i] — quantas çõoperaes
  // utilizamos para chegarno valor de dp[i]
vector<int> dp(n + 1);
   vector<int> cnt(n + 1);
  dp[0] = 0;
cnt[0] = 0;
  for (int i = 1; i <= n; i++) {
  dp[i] = dp[i - 1];
  cnt[i] = cnt[i - 1];
  int id = i - 1;
  drid = i - 1;</pre>
     dp[i] += v[id]
     int lo = max(0ll, id - l + 1);
int s = dp[lo] + (id - lo + 1) - lambda;
     if (s > dp[i]) {
        dp[i] = s;
        cnt[i] = cnt[lo] + 1;
   return {dp[n], cnt[n]};
int aliens_trick(vector<int> &v) {
  int l = 0, r = n;
while (l < r) {
     int mid = (l + r) \gg 1;
     pi ans = solve(v, mid);
(ans.sec > k) ? l = mid + 1 : r = mid;
  pi ans = solve(v, l);
return ans.fir + (l * k);
signed main() {
   ios_base::sync_with_stdio(false);
   cin.tie(NULL);
   cin >> n >> k >> l >> s;
   vector<int> a(n);
   vector<int> b(n);
  for (int i = 0; i < n; i++) {
    a[i] = 1, b[i] = 0;
    if (s[i] >= 'A' && s[i] <= 'Z') {
        a[i] ^= 1;
        b[i] ^= 1;
     }
   cout << n - max(aliens_trick(a)</pre>
                        aliens_trick(b))
         << endl;
   return 0;
// https://codeforces.com/contest/1279/problem/F
```

8.2 catalan

```
Description: Recursive formula: C_0 = C_1 = 1 C_n = \sum_{k=0}^{n-1} C_k C_{n-1-k}, n \geq 2 Analytical formula: C_n = \binom{2n}{n} - \binom{2n}{n-1} = \frac{1}{n+1} \binom{2n}{n}, n \geq 0 The first few numbers Catalan numbers, C_n (starting from zero): 1, 1, 2, 5, 14, 42, 132, 429, 1430, \dots The Catalan number C_n is the solution for:
```

- \bullet Number of correct bracket sequence consisting of n opening and n closing brackets.
- The number of rooted full binary trees with n+1 leaves (vertices are not numbered). A rooted binary tree is full if every vertex has either two children or no children.

- $\bullet\,$ The number of ways to completely parenthesize n+1 factors.
- The number of triangulations of a convex polygon with n+2 sides (i.e. the number of partitions of polygon into disjoint triangles by using the diagonals).
- The number of ways to connect the 2n points on a circle to form n disjoint chords.
- The number of non-isomorphic full binary trees with n internal nodes (i.e. nodes having at least one son).
- The number of monotonic lattice paths from point (0,0) to point (n,n) in a square lattice of size $n \times n$, which do not pass above the main diagonal (i.e. connecting (0,0) to (n,n)).
- Number of permutations of length n that can be stack sorted (i.e. it can be shown that the rearrangement is stack sorted if and only if there is no such index i < j < k, such that $a_k < a_i < a_j$).
- \bullet The number of non-crossing partitions of a set of n elements.
- The number of ways to cover the ladder $1 \dots n$ using n rectangles (The ladder consists of n columns, where i^{th} column has a height i).

```
#include <bits/stdc++.h>
using namespace std;
const int MOD = 1000000007;
typedef long long ll;
ll extGcd(ll a, ll b, ll &x, ll &y) {
   if (b == 0) {
 x = 1, y = 0;
      return a;
   } else {
      ll g = extGcd(b, a % b, y, x);
y -= (a / b) * x;
       return g;
ĺl inv(ll a) {
   ll inv_x, y;
extGcd(a, MOD, inv_x, y);
return (inv_x % MOD + MOD) % MOD;
const int MAXN = 4000010;
ll fat[MAXN], ifat[MAXN];
void init() {
  fat[0] = 1;
   for (int i = 1; i < MAXN; i++)
  fat[i] = (fat[i - 1] * i) % MOD;
ifat[MAXN - 1] = inv(fat[MAXN - 1]);
fan (int i - MAXN)</pre>
   for (int i = MAXN - 2; i >= 0; i__)
  ifat[i] = (ifat[i + 1] * (i + 1)) % MOD;
assert(ifat[0] == 1);
Il C(int n, int k) {
  if (k > n) return 0;
  return (fat[n] *
                 ((ifat[k] * ifat[n - k]) % MOD)) %
fl catalan(int n) {
   return (C(2 * n, n) - C(2 * n, n - 1) + MOD) %
fl f(int x1, int y1, int x2, int y2) {
   int y = y^2 - y^1, x = x^2 - x^1; if (y < 0 \text{ or } x < 0) return 0;
   return C(x + y, x);
// o = number of '(', c = number of ')', k = fixed
// prefix of '(' extra Catalan Generalization,
// open[i] >= close[i] for each 0 <= i < o + c + k</pre>
// where open[i] is number of '(' in prefix until
// i and close[i] is number of ')'
ll catalan2(int o, int c, int k) {
   MOD;
}
```

8.3 centroid

```
#include <bits/stdc++.h>
using namespace std;
const int MAXN = 500010;
typedef pair<int, int> pii;
namespace Centroid
vector<int> adj[MAXN];
int sub[MAXN];
int n;
void init(int n1) {
  n = n1;
  for (int i = 0; i < n; i++) adj[i].clear();
void addEdge(int a, int b) {
  adj[a].push_back(b);
  adj[b].push_back(a);
int dfsS(int u, int p) {
  sub[u] = 1;
  for (int to : adj[u]) {
    if (to != p) sub[u] += dfsS(to, u);
  return sub[u];
pii dfsC(int u, int p) {
  for (int to : adj[u])
    if (to != p and sub[to] > n / 2)
      return dfsC(to, u);
  for (int to : adj[u]) {
   if (to != p and (sub[to] * 2) == n)
      return pii(u, to);
  return pii(u, u);
pii findCentroid() {
  dfsS(0,
  return dfsC(0, -1);
   // namespace Centroid
```

8.4 centroid decomposition

```
#include <bits/stdc++.h>
using namespace std;
typedef long long li;
// O(N*log(N))
// Centroid Decomposition
const int MAXN = 200010;
namespace CD {
vector<int> adj[MAXN];
int dad[MAXN], sub[MAXN];
bool rem[MAXN];
int centroidRoot, n;
void init(int n1) {
  n = n1;
  for (int i = 0; i < n; i++) {
  adj[i].clear();</pre>
     rem[i] = false;
int dfs(int u, int p) {
  sub[u] = 1;
  for (int to : adj[u])
     if (!rem[to] and to != p)
       sub[u] += dfs(to, u);
  return sub[u];
int centroid(int u, int p, int sz) {
  for (auto to : adj[u])
  if (!rem[to] and to != p and sub[to] > sz / 2)
       return centroid(to, u, sz);
  return u:
void getChildren(int u, int p, int d,
                     vector<int> &v) {
  v.push_back(d);
  for (int to : adj[u]) {
    if (rem[to] or to == p) continue;
getChildren(to, u, d + 1, v);
```

```
ĺl ans = 0;
int k;
int decomp(int u, int p) {
  int sz = dfs(u, p);
  int c = centroid(u, p, sz);
  if (p == -1) p = c;
  dad[c] = p;
  rem[c] = true;
  // Begin
  vector<int> f(sz + 1, 0);
   f[0] = 1;
       (auto to : adj[c])
     if (!rem[to]) {
       vector<int> v;
       getChildren(to, c, 1, v);
for (int d : v) { // Query
  if (d <= k and k - d <= sz)</pre>
            ans += f[k - d];
       for (int d : v) // Update
          f[d]++;
  }
// End
  for (auto to : adj[c]) {
    if (!rem[to]) decomp(to, c);
  return c;
void addEdge(int a, int b) {
  adj[a].push_back(b);
  adj[b].push_back(a);
^{\prime}// Number of k-size paths: O(N * log(N))
ll solve(int k1) {
  assert(n > 0);
  ans = 0, k = k1;
  centroidRoot = decomp(0, -1);
  return ans:
    // namespace CD
8.5 checking_bipartiteness_online
const int N = 500010;
pii parent[N];
int rk[N];
int bipartite[N];
void make_set(int v)
  parent[v] = pii(v, 0);
  rk[v] = 0:
  bipartite(v) = true;
pii find_set(int v) {
  if (v != parent[v].first) {
    int parity = parent[v].second;
parent[v] = find_set(parent[v].first);
parent[v].second ^= parity;
  return parent[v];
void add_edge(int a, int b) {
  int x, y;
tie(a, x) = find_set(a);
tie(b, y) = find_set(b);
  if (a = b) {
     if (x == y) bipartite[a] = false;
  } else {
    if (rk[a] < rk[b]) swap(a, b);
parent[b] = pii(a, x ^ y ^ 1);
bipartite[a] &= bipartite[b];</pre>
     if (rk[a] == rk[b]) ++rk[a];
bool is_bipartite(int v) {
  return bipartite[find_set(v).first];
     chinese remainder theorem
```

#include <bits/stdc++.h>

```
#include "extended_euclidean.h"
using namespace std;
typedef long long ll;
namespace CRT {
inline ll normalize(ll x, ll mod) {
  x \% = mod;
if (x < 0) x += mod;
  retùrn x;
il solve(vector<!!> a, vector<!!> m) {
  for (int i = 0; i < n; i
    normalize(a[i], m[i]);</pre>
  ll ans = a[0];
ll lcm1 = m[0];
  for (int i = 1; i < n; i++) {
    ll x, y;
     ll g = extGcd(lcm1, m[i], x, y);
     if ((a[i] - ans) % g != 0) return -1;
     ans = normalize(
          ans + ((((a[i] - ans) / g) * x) %
                   (m[i] / g)) *
                       lcm1
          (lcm1 / g) * m[i]);
     lcm1 = (lcm1 / g) * m[i]; // lcm(lcm1, m[i]);
  return ans;
}
   // namespace CRT
8.7 counting inversions
#include <bits/stdc++.h>
using namespace std;
typedef long long ll;
const int INF = 0x3f3f3f3f3f;
 / Counting Inversions: O(N*log(N))
ll ci(vector<int> &v) {
  int n = v.size();
ll inv = OLL;
if (n == 1) return 0;
  vector<int> u1, u2;
for (int i = 0; i < n / 2; i++)
    u1.push_back(v[i]);</pre>
  for (int i = n / 2; i < n; i++)
    u2.push_back(v[i]);
  inv += ci(u1)
  inv += ci(u2)
  u1.push_back(INF);
  u2.push_back(INF);
  int ini1 = 0, ini2 = 0;
for (int i = 0; i < n; i++) {
  if (u1[ini1] <= u2[ini2]) {</pre>
     v[i] = u1[ini1++];
} else {
       v[i] = u2[ini2++];
       inv += u1.size() - ini1 - 1;
  return inv;
}
8.8 custom hash
```

```
#include <bits/stdc++.h>
using namespace std;
// O((V^2)*\dot{E}): for generic graph.
// O(sqrt(V)*E): on unit networks. A unit network
// is a network in which all the edges have unit
// capacity, and for any vertex except s and t
// either incoming or outgoing edge is unique.
// That's exactly the case with the network we
// build to solve the maximum matching problem
// with flows
template <typename flow_t>
struct Dinic
  struct FlowEdge {
    id = id1;
    }
  const flow_t flow_inf =
    numeric_limits<flow_t>::max();
vector<FlowEdge> edges;
  vector<vector<int>> adj;
  int n, m = 0;
int s, t;
vector<int> level, ptr;
  queue<int> q;
  bool bfs() {
    while (!q.empty()) {
       int u = q.front();
       q.pop();
       for (int id : adj[u]) {
         if (edges[id].cap - edges[id].flow < 1)
         if (level[edges[id].to] !=-1) continue;
         level[edges[id].to] = level[u] + 1;
         q.push(edges[id].to);
    return level[t] !=-1;
  flow_t dfs(int u, flow_t pushed) {
  if (pushed == 0) return 0;
    if (u == t) return pushed;
    for (int &cid = ptr[u];
          cid < (int)adj[u].size(); cid++) {</pre>
       int id = adj[u][cid];
       int to = edges[id].to;
       if (level[u] + 1 != level[to]
            edges[id].cap — edges[id].flow < 1)
         continue;
       flow_t tr = dfs(
    to, min(pushed, edges[id].cap -
                                    edges[id].flow));
      if (tr == 0) continue;
edges[id].flow += tr;
edges[id ^ 1].flow -= tr;
       return tr;
    return 0;
   // Public:
  Dinic() {}
void init(int _n) {
       = _n;
    adj.resize(n);
    level.resize(n);
    ptr.resize(n);
  void addEdge(int from, int to, flow_t cap,
                  int id = 0) {
    assert(n > 0);
    edges.emplace_back(from, to, cap, id);
edges.emplace_back(to, from, 0, -id);
    adj[from].push_back(m);
    adj[to].push_back(m + 1);
    m += 2;
  void resetFlow() {
  for (int i = 0; i < m; i++) edges[i].flow = 0;</pre>
```

```
flow_t maxFlow(int s1, int t1) {
    s = s1, t = t1;
flow_t f = 0;
while (true) {
      level.assign(n, -1);
      level[s] = 0;
      q.push(s);
      if (!bfs()) break;
      ptr.assign(n, 0);
while (flow_t pushed = dfs(s, flow_inf))
        f += pushed;
    return f;
  }
vector<int> recoverCut(Dinic<int> &d) {
  vector<bool> seen(d.n, false);
  queue<int> q;
  q.push(d.s);
  seen[d.s] = true;
while (!q.empty()) {
    int u = q.front();
    q.pop();
    for (int idx : d.adj[u]) {
      auto e = d.edges[idx];
      if (e.cap == e.flow) continue;
if (!seen[e.to]) {
        q.push(e.to);
         seen[e.to] = true;
    }
  vector<int> ans;
  for (auto e : d.edges) {
    if (e.id >= 0) ans.push_back(e.id);
    }
  return ans;
typedef long long ll;
typedef tuple<int, int, ll> tp;
                                   // (u, to, cap)
#define all(x) x.begin(), x.end()
// O(V*E*log(MAXC))
ll maxFlowWithScaling(int n, vector<tp> edges,
                        int s, int t) {
  Dinic<ll> graph;
  graph.init(n);
  sort(all(edges), [&](tp a, tp b) {
    return get<2>(a) < get<2>(b);
  ll ans = 0;
  for (int l = (1 << 30); l > 0; l >>= 1) {
    while (!edges.empty()) {
      auto [u, to, cap] = edges.back();
if (cap >= l) {
        graph.addEdge(u, to, cap);
        edges.pop_back();
      } else {
        break;
      }
    ans += graph.maxFlow(s, t);
  return ans;
8.10 dynamic median
```

```
#include <bits/stdc++.h>
using namespace std;
class DinamicMedian {
  typedef int t_median;
 private:
  priority_queue<t_median> mn;
  priority_queue<t_median, vector<t_median>,
                   greater<t_median>>
 public:
  double median() {
```

```
if (mn.size() > mx.size())
      return mn.top();
    else
      return (mn.top() + mx.top()) / 2.0;
  void push(t_median x) {
    if (mn.size() <= mx.size())</pre>
      mn.push(x);
    else
      mx.push(x);
    if ((!mx.empty()) and (!mn.empty())) {
      while (mn.top() > mx.top()) {
         t_median a = mx.top();
        mx.pop();
         t_{median} b = mn.top();
        mn.pop();
        mx.push(b);
        mn.push(a);
   }
};
```

8.11 dynamic wavelet tree

```
#include <bits/stdc++.h>
using namespace std;
struct SplayTree {
   struct Node {
     int x, y, s;
Node *p = 0;
Node *l = 0;
      Node *r = 0;
      Node(int v) {
        x = v;
y = v;
        s = 1;
      void upd() {
        s = 1;
y = x;
         if (l)
           y += l->y;
            s += l->s;
            (r) {
y += r->y;
        if
            s += r \rightarrow s;
      int left_size() { return l ? l->s : 0; }
  Node *root = 0;
   void rot(Node *c) {
     auto p = c \rightarrow p;
      auto g = p \rightarrow p;
      if (g) (g->l == p ? g->l : g->r) = c;
if (p->l == c) {
        p->l = c->r;
c->r = p;
        if (p\rightarrow l) p\rightarrow l\rightarrow p = p;
      } else {
        p\rightarrow r = c\rightarrow l;
        c\rightarrow l = p;
        if (p\rightarrow r) p\rightarrow r\rightarrow p=p;
      \dot{p} \rightarrow p = c;
      c \rightarrow p = g;
     p\rightarrow upd();
      c->upd();
   void splay(Node *c) {
     while (c\rightarrow p) {
        auto p = c \rightarrow p;
        auto g = p \rightarrow p;
        if (g)
            rot((g \rightarrow l == p) == (p \rightarrow l == c) ? p : c);
        rot(c);
      c->upd();
      root = c;
  Node *join(Node *l, Node *r) {
      if (not l) return r;
```

```
if (not r) return l;
while (l->r) l = l->r;
      splay(1);
      r\rightarrow p = l;
      l \rightarrow r = r;
      l—>upd();
      return l;
   pair<Node *, Node *> split(Node *p, int idx) {
  if (not p) return make_pair(nullptr, nullptr);
  if (idx < 0) return make_pair(nullptr, p);</pre>
      if (idx >= p->s) return make_pair(p, nullptr);
for (int lf = p->left_size(); idx != lf;
             lf = p->left_size()) {
         if (idx < lf)
           p = p \rightarrow l;
         else
           p = p \rightarrow r, idx = lf + 1;
      splay(p);
      Node *l = p;
      Node *r = p \rightarrow r;
      if (r) { l \rightarrow r = r \rightarrow p = 0;
        l->upd();
      return make_pair(l, r);
   Node *get(int idx) {
      auto p = root;
      for (int lf = p->left_size(); idx != lf;
    lf = p->left_size()) {
         if (idx < lf)
           p = p \rightarrow l;
         else
           p = p - r, idx - lf + 1;
      splay(p);
      return p;
   int insert(int idx, int x) {
     Node *l, *r;
tie(l, r) = split(root, idx - 1);
int v = l ? l->y : 0;
      root = join(l, join(new Node(x), r));
      return v;
   void erase(int idx) {
     Node *l, *r;
tie(l, r) = split(root, idx);
     root = join(l->l, r);
delete l;
   int rank(int idx) {
     Node *l, *r;
tie(l, r) = split(root, idx);
      int x = (l \&\& l \rightarrow l ? l \rightarrow l \rightarrow y : 0);
      root = join(l, r);
      return x;
   int operator[](int idx) { return rank(idx); }
   ~SplayTree() {
  if (!root) return;
      vector<Node *> nodes{root};
      while (nodes.size())
         auto u = nodes.back();
         nodes.pop_back();
        if (u->l) nodes.emplace_back(u->l);
if (u->r) nodes.emplace_back(u->r);
         delete u;
  }
};
class WaveletTree {
 private:
   int lo, hi;
WaveletTree *l = 0;
   WaveletTree *r = 0;
   SplayTree b;
 public:
   WaveletTree(int min_value, int max_value) {
      lo = min_value;
      hi = max_value;
      b.insert(0, 0);
```

```
queue<int> q;
   ~WaveletTree() {
                                                                   void init(int n1) {
     delete l;
                                                                      n = n1;
                                                                      for (int i = 0; i < n; i++) adj[i].clear();
     delete r;
  // 0—indexed
void insert(int idx, int x) {
                                                                   void addEdge(int a, int b) {
                                                                     adj[a].push_back(b);
    if (lo >= hi) return;
int mid = (lo + hi - 1) / 2;
                                                                      adj[b].push_back(a);
     if (x \le mid) {
                                                                   void contract(int u, int v, bool first = 1) {
  static vector<bool> bloss;
       l = l ?: new WaveletTree(lo, mid);
       l->insert(b.insert(idx, 1), x);
                                                                      static int l;
    (first) {
                                                                        bloss = vector<bool>(n, 0);
                                                                        vector<bool> teve(n, 0);
    }
                                                                        int k = u;
                                                                        l = v;
  // 0—indexed
                                                                        while (1) {
  teve[k = base[k]] = 1;
  if (match[k] == -1) break;
  void erase(int idx) {
  if (lo == hi) return;
    auto p = b.get(idx);
int lf = p->l ? p->l->y : 0;
                                                                           k = pai[match[k]];
     int x = p \rightarrow x;
     b.erase(idx);
    if (x == 1)
l->erase(lf);
                                                                      while (base[u] != l) {
     else
       r—>erase(idx - lf);
                                                                        v = match[u];
                                                                        u = pai[match[u]];
  // kth smallest element in range [i, j[
  // 0—indexed
                                                                      if (!first) return;
  int kth(int i, int j, int k) {
  if (i >= j) return 0;
                                                                     contract(v, u, 0);
for (int i = 0; i < n; i++)
  if (bloss[base[i]]) {</pre>
     if (lo == hi) return lo;
    int x = b.rank(i);
int y = b.rank(j);
if (k <= y - x)
  return l->kth(x, y, k);
                                                                           base[i] = l
                                                                           if (!vis[i]) q.push(i);
                                                                           vis[i] = 1;
     else
       return r\rightarrowkth(i - x, j - y, k - (y - x));
                                                                   int getpath(int s) {
  // Amount of numbers in the range [i, j[ Less
  // than or equal to k 0—indexed
                                                                      vis[s] = 1;
  int lte(int i, int j, int k) {
  if (i >= j or k < lo) return 0;</pre>
                                                                      q = queue<int>();
                                                                      q.push(s);
     if (hi \le k) return j - i;
                                                                      while (q.size()) {
     int x = b.rank(i);
                                                                        int u = q.front();
     int y = b.rank(j);
                                                                        q.pop()
     return l\rightarrow lte(x, y, k) + r\rightarrow lte(i - x, j - y, k);
                                                                        for (int i : adj[u]) {
                                                                             continue;
  // Amount of numbers in the range [i, j[ equal
  // to k 0—indexed
  int count(int i, int j, int k) {
  if (i >= j or k < lo or k > hi) return 0;
    if (lo == hi) return j - i;
int mid = (lo + hi - 1) / 2;
                                                                             pai[i] = u;
     int x = b.rank(i);
     int y = b.rank(j)
                                                                             vis[i] = 1;
     if (k <= mid) return l->count(x, y, k);
                                                                             q.push(i);
     return r\rightarrowcount(i - x, j - y, k);
                                                                          }
                                                                        }
  // 0—indexed
  int get(int idx) {
                                                                      return -1;
    return kth(idx, idx + 1, 1);
                                                                   typedef pair<int, int> pii;
};
                                                                      vector<pii> ans;
8.12 edmond blossoms
                                                                             if (match[j] ==
#include <bits/stdc++.h>
using namespace std;
const int MAXN = 510;
                                                                               match[i] = j;
                                                                                match[j] = i;
// Adaptado de:
                                                                                break;
// https://github.com/brunomaletta/Biblioteca/blob/
    master/Codigo/Grafos/blossom.cpp
// Edmond's Blossoms algorithm give a maximum
```

```
while (!teve[l = base[l]]) l = pai[match[l]];
     bloss[base[u]] = bloss[base[match[u]]] = 1;
pai[u] = v;
   for (int i = 0; i < n; i++)

base[i] = i, pai[i] = -1, vis[i] = 0;
        if (base[i] == base[u] or match[u] == i)
         if (i == s or (match[i] != -1 and
                              pai[match[i]] != -1))
        contract(u, i);
else if (pai[i] == -1) {
           if (match[i] == -1) return i;
i = match[i];
vector<pii> maximumMatching() {
  memset(match, -1, sizeof(match));
for (int i = 0; i < n; i++)
  if (match[i] == -1)
    for (int j : adj[i])</pre>
   for (int i = 0; i < n; i++) if (match[i] == -1) {
        int j = getpath(i);
if (j == -1) continue;
        while (j != -1) {
           int p = pai[j], pp = match[p];
           match[p] = j;
```

// matching in general graphs (non-bipartite)

int n, pai[MAXN], base[MAXN], vis[MAXN];

// 0(N^3)

int match[MAXN]

namespace EdmondBlossoms { vector<int> adj[MAXN];

```
match[j] = p;
          j = pp;
  for (int i = 0; i < n; i++)
  if (i < match[i])</pre>
       ans.emplace_back(i, match[i]);
  return ans;
    // namespace EdmondBlossoms
8.13 extended euclidean
#include <bits/stdc++.h>
using namespace std;
typedef long long li;
ll extGcd(ll a, ll b, ll &x, ll &y) {
  if (b == 0) {
     x = 1, y = 0;
     return a;
  } else {
     ll g = extGcd(b, a % b, y, x);
y _= (a / b) * x;
     return g;
  }
// a*x + b*y = g
// a*(x-(b/g)*k) + b*(y+(a/g)*k) = g
bool dioEq(ll a, ll b, ll c, ll &x0, ll &y0,
              ll &g) {
  g = extGcd(abs(a), abs(b), x0, y0);
if (c % g) return false;
  x0 *= c / g;
  y0 *= c / g;

if (a < 0) x0 = -x0;

if (b < 0) y0 = -y0;
  return true;
inline void shift(ll &x, ll &y, ll a, ll b,
                        ll cnt) {
  x += cnt * b;
  y = cnt * a;
// a1 + m1*x = a2 + m2*y
// Find the first moment that both are equal
ll findMinimum(ll a1, ll m1, ll a2, ll m2) {
    ll a = m1, b = -m2, c = a2 - a1;
  if (!dioEq(a, b, c, x, y, g)) return -1;
  a /= g;
  b /= \bar{g};
  int sa = a > 0 ? +1 : -1;

int sb = b > 0 ? +1 : -1;

shift(x, y, a, b, -x / b);

if (x < 0) shift(x, y, a, b, sb);

if (y < 0) {
     shift(x, y, a, b, y / a);
if (y < 0) shift(x, y, a, b, —sa);
if (x < 0) return —1;
  return a * x * g;
ll x, y, g;
if (a == 0 or b == 0)
     if (a == 0 \text{ and } b == 0)
        return (c == 0) * (maxx - minx + 1) *
                  (maxy - miny + 1);
     if (a == 0)
        return (c % b == 0) * (maxx - minx + 1) * (miny <= c / b and c / b <= maxy);
     return (c % a == 0) *
               (minx \le c / a and c / a \le maxx) * (maxy - miny + 1);
  if (!dioEq(a, b, c, x, y, g)) return 0;
  a /= g;
  b /= \bar{g};
  int sign_a = a > 0 ? +1 : -1;
  int sign_b = b > 0 ? +1 : -1;
  shift(x, y, a, b, (minx - x) / b);
  if (x < minx) shift(x, y, a, b, sign_b);</pre>
```

```
if (x > maxx) return 0;
ll lx1 = x;
shift(x, y, a, b, (maxx - x) / b);
if (x > maxx) shift(x, y, a, b, -sign_b);
ll rx1 = x;
shift(x, y, a, b, -(miny - y) / a);
if (y < miny) shift(x, y, a, b, -sign_a);
if (y > maxy) return 0;
ll lx2 = x;
shift(x, y, a, b, -(maxy - y) / a);
if (y > maxy) shift(x, y, a, b, sign_a);
ll rx2 = x;
if (lx2 > rx2) swap(lx2, rx2);
ll rx = max(lx1, lx2);
ll rx = min(rx1, rx2);
if (lx > rx) return 0;
return (rx - lx) / abs(b) + 1;
}
```

8.14 flow_with_demand

```
#include "dinic.h"
using namespace std;
template <typename flow_t>
struct MaxFlowEdgeDemands {
  Dinic<flow_t> mf;
  vector<flow_t> ind, outd;
  flow_t D;
  int n;
MaxFlowEdgeDemands(int n) : n(n) {
    D = 0;
    mf.init(n + 2)
    ind.assign(n, 0);
    outd.assign(n, 0);
  mf.addEdge(a, b, cap — demands);
    D += demands;
    ind[b] += demands;
    outd[a] += demands;
  bool solve(int s, int t) {
    mf.addEdge(t, s,
               numeric_limits<flow_t>::max());
    for (int i = 0; i < n; i++) {
  if (ind[i]) mf.addEdge(n, i, ind[i]);</pre>
      if (outd[i]) mf.addEdge(i, n + 1, outd[i]);
    return mf.maxFlow(n, n + 1) == D;
};
```

8.15 fraction

```
#include <bits/stdc++.h>
using namespace std;
typedef long long f_type;
// Representation of the a/b
struct Fraction {
  f_type a, b;
Fraction(f_type _a = 0) : a(_a), b(1) {}
Fraction(f_type _a, f_type _b) {
    f_type g = __gcd(_a, _b);
a = _a / g;
b = _b / g;
     if (b < 0) { a = -a;
       b = -b;
  Fraction operator+(Fraction oth) {
     return Fraction(a * oth.b + oth.a * b,
                         b * oth.b);
  Fraction operator—(Fraction oth) {
     return Fraction(a * oth.b - oth.a * b,
                         b * oth.b);
  Fraction operator*(Fraction oth) {
     return Fraction(a * oth.a, b * oth.b);
```

```
Fraction operator/(Fraction oth) {
     return Fraction(a * oth.b, b * oth.a);
  bool operator>=(Fraction oth) {
     return ((*this) - oth).a >= 0;
                                                                             if (i != row)
  bool operator==(Fraction oth) {
     return a == oth.a and b == oth.b;
  operator f_type() { return a / b; }
operator double() { return double(a) / b; }
                                                                             }
                                                                           row++;
      function_root_using_newton
8.16
#include <bits/stdc++.h>
using namespace std;
typedef long double ld;
struct Poly {
   vector<ld> v
  Poly(vector<ld> &v1) : v(v1) {}
   // return f(x)
   ld f(ld x)
     ld ans = 0;
ld e = 1;
                                                                        return 1;
                                                                     }
     int n = v.size();
for (int i = 0; i < n; i++) {
   ans += v[i] * e;
                                                                     8.18 \quad gauss\_xor
     return ans;
                                                                     using namespace std;
                                                                     const int MAXB = 30;
  // return f'(x)
ld_df(ld x) {
                                                                     struct GaussXOR {
  int table[MAXB];
     ld ans = 0;
     ld e = 1;
int n = v.size();
                                                                             table[i] = 0;
     for (int i = 1; i < n; i++) {
                                                                           }
       ans += i * v[i] * e;
e *= x;
                                                                        int size() {
     return ans;
  // takes some root of the polynomial
  ld root(ld x0 = 1) {
                                                                           return ans;
     const ld eps = 1E-10;
     ld x = x0;
     for (;;) {
       ld nx = x - (f(x) / df(x));
       if (abs(x - nx) < eps) break;
                                                                           return x == 0;
     return x;
                                                                        void add(int x) {
  // div f(x) by (x-a)
void div(ld a) {
     int g = (int)v.size() - 1;
                                                                                table[i] = x;
     vector<ld> aux(g);
for (int i = g; i >= 1; i—) {
   aux[i - 1] = v[i];
}
                                                                                x = 0
                                                                             } else {
                                                                             }
       v[i - 1] += a * aux[i - 1];
                                                                          }
     \dot{v} = aux:
                                                                        int max() {
8.17 gauss
                                                                           return ans;
                                                                        }
                                                                     };
#include <bits/stdc++.h>
using namespace std;
const int INF = 0x3f3f3f3f;
typedef long double ld;
const ld EPS = 1e-9;
                                                                     8.19 graph theorem
int gauss(vector<vector<ld>> a, vector<ld> &ans) {
  int n = (int)a.size();
int m = (int)a[0].size() - 1;
                                                                     using namespace std;
using ll = long long;
  vector<int> where(m, -1);
  for (int col = 0, row = 0; col < m && row < n;
     col++) {
int sel = row;
                                                                     namespace GraphTheorem {
     for (int i = row; i < n; i++)
  if (abs(a[i][coll) > abs(a[
                                                                     // return if a sequence of integers d can be
// represented as the degree sequence of a finite
           (abs(a[i][col]) > abs(a[sel][col]))
```

```
if (abs(a[sel][col]) < EPS) continue;
for (int i = col; i <= m; i++)
          swap(a[sel][i], a[row][i]);
      where[col] = row;
for (int i = 0; i < n; i++) {
             ld c = a[ij[col] / a[row][col];
for (int j = col; j <= m; j++)</pre>
                a[i][j] = a[row][j] * c;
   ans.assign(m, 0);
for (int i = 0; i < m; i++)
if (where[i] != -1)
   ans[i] = a[where[i]][m] / a[where[i]][i];
for (int i = 0; i < n; i++) {
  ld sum = 0;
  for (int j = 0; j < m; j++)
    sum += ans[j] * a[i][j];
  if (abs(sum = a[i][m]) > EPS) return 0;
      if (abs(sum - a[i][m]) > EPS) return 0;
   for (int i = 0; i < m; i++)
if (where[i] == -1) return INF;
#include <bits/stdc++.h>
   GaussXOR() {
   for (int i = 0; i < MAXB; i++) {
      int ans = 0;
for (int i = 0; i < MAXB; i++) {
         if (table[i]) ans++;
   bool isComb(int x) {
  for (int i = MAXB - 1; i >= 0; i—) {
    x = std::min(x, x ^ table[i]);
       for (int i = MAXB - 1; i >= 0; i—) { if ((table[i] == 0) and ((x >> i) & 1)) {
             x = std::min(x, x ^ table[i]);
      int ans = 0;
for (int i = MAXB - 1; i \ge 0; i \ge 0
          ans = std::max(ans, ans ^ table[i]);
#include <bits/stdc++.h>
#define all(x) x.begin(), x.end()
using pii = pair<int, int>;
```

// simple graph on n vertices

```
bool ErdosGallai(vector<int> d) {
  int n = d.size();
  sort(all(d), greater<int>());
ll sum1 = 0, sum2 = 0;
int mn = n - 1;
  for (int k = 1; k \le n; k++) {
     sum1 += d[k - 1];
    while (k \le mn \text{ and } k > d[mn]) \text{ sum2} += d[mn-];
    if (mn + 1 < k) sum2 -= d[mn++];
ll a = sum1, b = k * (ll)mn + sum2;
if (a > b) return false;
  return sum1 % 2 == 0;
vector<pii> recoverErdosGallai(vector<int> d) {
  int n = d.size();
  priority_queue<pii> pq;
  for (int i = 0; i < n; i++) pq.emplace(d[i], i);
  vector<pii> edges;
  while (!pq.empty()) {
    auto [g, u] = pq.top();
pq.pop();
     vector<pii> aux(g);
     for (int i = 0; i < g; i++) {
  if (pq.empty()) return {};</pre>
       auto [g2, u2] = pq.top();
       pq.pop();
       if (g2 == 0) return {};
       edges.emplace_back(u, u2);
       aux[i] = pii(g2 - 1, u2);
     for (auto [g2, u2] : aux) pq.emplace(g2, u2);
  return edges;
    // namespace GraphTheorem
8.20 gray code
int grayCode(int nth) { return nth ^ (nth >> 1); }
int revGrayCode(int g) {
  int nth = 0;
for (; g > 0; g >>= 1) nth ^= g;
  return nth;
8.21 histogram
#include <bits/stdc++.h>
using namespace std;
typedef long long li;
  Largest Rectangular Area in a Histogram
ll histogram(vector<int> v) {
  int n = v.size();
  v.push_back(0);
  ll ans = 0;
  stack<int>'st
  for (int i = 0; i <= n; i++) {
  while (st.size() && v[st.top()] >= v[i]) {
       int idx = st.top();
       st.pop();
       int L = st.size() ? st.top() : -1;
ans = max(ans, (i - L - 1) * (ll)v[idx]);
    st.push(i);
  return ans;
// Largest Rectangular Area formed only by 1
int maxArea1(vector<vector<bool>> mat) {
  int n = mat.size();
if (n == 0) return 0;
  int m = mat[0].size();
  vector<int> v(m, 0);
int ans = 0;
for (int i = 0; i < n; i++) {
  for (int j = 0; j < m; j++) {</pre>
       if (mat[i][j])
          v[j]++;
       else
          v[j] = 0;
```

```
ans = max(ans, (int)histogram(v));
   return ans;
8.22 hungarian
#include <bits/stdc++.h>
using namespace std;
// input: matrix n x m, n <= m
// return vector p of size n, where p[i] is the
// match for i
    and minimum cost
int u[ms], v[ms], p[ms], way[ms], minv[ms];
bool used[ms];
pair<vector<int>, int> solve(
     const vector<vector<int>> &matrix) {
   int n = matrix.size();
   if (n == 0) return {vector<int>(), 0};
   int m = matrix[0].size();
   assert(n <= m);
  memset(u, 0, (n + 1) * sizeof(int));
memset(v, 0, (m + 1) * sizeof(int));
  memset(v, v, (iii + 1) * S12e01(III));
memset(p, 0, (m + 1) * sizeof(int));
for (int i = 1; i <= n; i++) {
    memset(minv, 0x3f, (m + 1) * sizeof(int));
    memset(way, 0, (m + 1) * sizeof(int));
    for (int j = 0; j <= m; j++) used[j] = 0;
    n[0] = i;</pre>
     p[0] = i;
      int k0 = 0;
     do {
        used[k0] = 1;
int i0 = p[k0], delta = INF, k1 = 0;
for (int j = 1; j <= m; j++) {</pre>
           if (!used[j]) {
               int cur = matrix[i0 - 1][j - 1] -
                             u[i0] - v[j];
               if (cur < minv[j]) {</pre>
                 minv[j] = cur;
                 way[j] = k0;
               if (minv[j] < delta) {
                 delta = minv[j];
                 k1 = j;
```

for (int j = 0; j <= m; j++) {

u[p[j]] += delta;

v[j] = delta;

} else {
 minv[j] = delta;

if (used[j]) {

8.23 implicit treap

}

}

do {

k0 = k1;

k0 = k1;} while (k0);

} while (p[k0]);

int k1 = way[k0];

if (!p[j]) continue; ans[p[j] - 1] = j - 1;

return $\{ans, -v[0]\};$

vector<int> ans(n, -1); for (int j = 1; j <= m; j++) {

p[k0] = p[k1];

#include <bits/stdc++.h>
using namespace std;
namespace ITreap {
const int N = 500010;
typedef long long treap_t;

}

```
treap_t X[N];
int en = 1, Y[N], SZ[N], L[N], R[N], P[N], root;
const treap_t_neutral = 0;
treap_t op_val[N];
bool rev[N];
inline treap_t join(treap_t a, treap_t b,
                         treap_t c) {
  return a + b + c;
void calc(
  // update node given children info
  // code here, no recursion
  op_val[u] =
       join(op_val[L[u]], X[u], op_val[R[u]]);
void unlaze(int u) {
  if (!u) return;
// code here, no recursion
if (rev[u]) {
    if (L[u]) rev[L[u]] ^= rev[u];
if (R[u]) rev[R[u]] ^= rev[u];
swap(L[u], R[u]);
    rev[u] = false;
  }
void split(
    int u, int s, int &l,
int &r) { // l gets first s, r gets remaining
  unlaze(u);
  if (!u) return (void)(l = r = 0);
if (sz[L[u]] < s) {
   split(R[u], s - sz[L[u]] - 1, l, r);</pre>
     R[u] = l;
     l = u;
  } else
     split(L[u], s, l, r);
    L[u] = r;
r = u;
  P[u] = 0;
  calc(u);
int merge(int l, int r) { // els on l <= els on r
  unlaze(l);
  unlaze(r)
  if (!l || !r) return l + r;
  int u;
if (Y[l] > Y[r])
    R[l] = merge(R[l], r);
    u = l;
  } else {
  L[r] = merge(l, L[r]);
  u = r;
  \dot{P}[u] = 0;
  calc(u);
  return u;
int new_node(treap_t x) {
  P[en] = 0;
  X[en] = x;
op_val[en] = x;
  rev[en] = false;
  return en++;
int nth(int u, int idx) {
  if (!u) return 0;
  unlaze(u);
if (idx <= sz[L[u]])
  return nth(L[u], idx);
else if (idx == sz[L[u]] + 1)
     return u;
     return nth(R[u], idx - sz[L[u]] - 1);
// Public
void init(
     int n = N
               1) { // call before using other funcs
  // init position 0
  sz[0] = 0;
op_val[0] = neutral;
  // init Treap
```

```
root = 0;
   std::mt19937 rng(
        (int)std::chrono::steady_clock::now()
              .time_since_epoch()
              .count());
   for (int i = en = 1; i \le n; i++) {
     Y[i] = i;
sz[i] = 1;
     L[i] = R[i] = 0;
   shuffle(Y + 1, Y + n + 1, rng);
// 0—indexed
int insert(int idx, int val) {
   int a, b;
split(root, idx, a, b);
   int node = new_node(val)
   root = merge(merge(a, node), b);
   return node;
// 0—indexed
void erase(int idx) {
   int a, b, c, d;
split(root, idx, a, b);
   split(b, 1, c, d);
   root = merge(a, d);
// 0—indexed
treap_t nth(int idx) {
   int u = nth(root, idx + 1);
   return X[u];
// 0—indexed [l, r]
treap_t query(int l, int r) {
   if (l > r) swap(l, r);
   int a, b, c, d;
split(root, l, a, d);
split(d, r - l + 1, b, c);
   treap_t ans = op_val[b];
   root = merge(a, merge(b, c));
   return ans;
// 0—indexed [l, r]
void reverse(int l, int r) {
  if (l > r) swap(l, r);
   int a, b, c, d;
split(root, l, a, d);
split(d, r - l + 1, b, c);
if (b) rev[b] ^= 1;
   root = merge(a, merge(b, c));
int getRoot(int x) {
   while (P[x]) x = P[x];
   return x;
int getPos(int node) {
  int ans = sz[L[node]];
  while (P[node]) {
    if (L[P[node]] == node) {
      node = P[node];
    }
}
      } else {
        node = P[node];
        ans += sz[L[node]] + 1;
   return ans;
};
    // namespace ITreap
8.24 kadane
```

```
#include <bits/stdc++.h>
using namespace std;
typedef long long ll;
// Largest Sum Contiguous Subarray: 0(N)
ll kadane(vector<ll> &v) {
    ll ans = 0, bigger = 0;
    for (int i = 0; i < (int)v.size(); i++) {
        bigger = max(0LL, bigger + v[i]);
        ans = max(ans, bigger);
    }
    return ans;
}
// Largest Sum Submatrix: 0(N^3)</pre>
```

8.25 karatsuba

```
#include <bits/stdc++.h>
using namespace std;
// Source:
// #pragma GCC optimize("0fast")
// #pragma GCC target ("avx,avx2")
template <typename T>
void kar(T* a, T* b, int n, T* r, T* tmp) {
  if (n <= 64) {</pre>
      for (int i = 0; i < n; i++)
for (int j = 0; j < n; j++)
r[i + j] += a[i] * b[j];
      return;
  int mid = n / 2;
T *atmp = tmp, *btmp = tmp + mid, *E = tmp + n;
  memset(E, 0, sizeof(E[0]) * n);
for (int i = 0; i < mid; i++) {
  atmp[i] = a[i] + a[i + mid];
  btmp[i] = b[i] + b[i + mid];</pre>
   kar(atmp, btmp, mid, E, tmp + 2 * n);
   kar(a, b, mid, r, tmp + 2 * n);

kar(a + mid, b + mid, mid, r + n, tmp + 2 * n);
   for (int i = 0; i < mid; i++) {
  T temp = r[i + mid];</pre>
      r[i + mid] += E[i] - r[i] - r[i + 2 * mid];

r[i + 2 * mid] +=

E[i + mid] - temp - r[i + 3 * mid];
  }
// O(n^1.58), Advantages: you can add any module
template <typename T>
vector<T> karatsuba(vector<T> a, vector<T> b) {
   int n = max(a.size(), b.size());
   while (n \& (n-1)) n++;
  a.resize(n), b.resize(n);
vector<T> ret(2 * n), tmp(4 * n);
kar(&a[0], &b[0], n, &ret[0], &tmp[0]);
   return ret;
```

8.26 kmp

```
#include <bits/stdc++.h>
using namespace std;
// "abcabcd" is [0,0,0,1,2,3,0]
// "aabaaab" is [0,1,0,1,2,2,3]
vector<int> kmp(string s) {
  int n = (int)s.length();
  // pi[i] is the length of the longest proper
  // prefix of the substring s[0..i] which is also
  // a suffix of this substring.
  vector<int> pi(n);
  for (int i = 1; i < n; i++) {
    int j = pi[i - 1];</pre>
```

```
while (j > 0 \text{ and } s[i] != s[j]) j = pi[j - 1]; if (s[i] == s[j]) j++;
     pi[i] = j;
   return pi;
// The ans[i] count the amount of occurrence of
// the prefix s[0..i] in s
vector<int> prefix0ccurrences(string &s) {
   auto pi = kmp(s)
   int n = pi.size();
  vector<int> ans(n + 1);
for (int i = 0; i < n; i++) ans[pi[i]]++;
for (int i = n - 1; i > 0; i—)
     ans[pi[i-1]] += ans[i];
  for (int i = 1; i \le n; i++)
ans[i - 1] = ans[i] + 1;
   ans.pop_back();
   return ans;
int K = 26;
inline int getID(char c) { return c - 'a'; }
vector<vector<int>> computeAutomaton(string s) {
   s += '#';
  int n = s.size();
vector<int> pi = kmp(s);
  vector<vector<int>> aut(n, vector<int>(26));
for (int i = 0; i < n; i++) {
  for (int c = 0; c < K; c++) {</pre>
        if(i > 0 \text{ and } c != getID(s[i]))
          aut[i][c] = aut[pi[i-1]][c];
          aut[i][c] = i + (c == getID(s[i]));
     }
   return aut;
8.27 lagrange
#include <bits/stdc++.h>
using namespace std;
typedef long double ld;
struct PointValue {
   ld x,
   PointValue(ld x0 = 0, ld y0 = 0)
        : x(x0), y(y0) {}
void mul(vector<ld> &A,
           int x0) { // multiply A(x) by (x - x0)
   int n = A.size();
   A.push_back(0);
  auto B = A;
for (int i = n; i >= 1; i—) {
     A[i] = A[i - 1];
  A[0] = 0;

for (int i = 0; i < n + 1; i++)

A[i] = B[i] * x0;
void div(vector<ld> &A,
	int x0) { // multiply A(x) by (x - x0)
	int g = (int)A.size() - 1;
   vector<ld> aux(g);
  for (int i = g; i >= 1; i—) {
    aux[i - 1] = A[i];
     A[i-1] += x0 * aux[i-1];
  A = aux;
// Change Polynomial Representation from
// Point—Value to Coefficient O(n^2)
vector<ld> LagrangeInterpolation(
     vector<PointValue> vp) {
   vector<ld> A(1, 1);
   int n = vp.size();
   for (int i = 0; i < n; i++) mul(A, vp[i].x);
  vector<ld> ans(n, 0);
for (int i = 0; i < n; i++) {
     ld x = vp[i].x, y = vp[i].y;
     div(A, x);
ld d = 1;
for (int j = 0; j < n; j++)
        if (j != i) d *= (x - vp[j].x);
```

```
(t[p].ch[0] == x)
     for (int j = 0; j < n; j++)
ans[j] += A[j] * (y / d);
                                                                                                     ? x
                                                                                                     : p);
                                                                                  rotate(x);
     mul(A, vp[i].x);
                                                                               }
   return ans;
                                                                             int access(int v) {
                                                                                int last = -1;
                                                                                for (int w = v; w + 1;
last = w, splay(v), w = t[v].p)
8.28 lagrange poly
                                                                                   splay(w), t[w].ch[1] = (last == -1 ? -1 : v);
                                                                                return last;
#include "modular_int.h"
namespace LagrangePoly {
const int MAXN = 100010;
modInt den[MAXN], fat[MAXN], ifat[MAXN], l[MAXN],
                                                                            // Public:
void init(int n) {
  for (int i = 0; i <= n; i++) t[i] = node();</pre>
     r[MAXN]
void build(int n) {
                                                                             int findRoot(int v) {
                                                                               access(v); while (t[v].ch[0] + 1) v = t[v].ch[0];
   fat[0] = 1;
for (int i = 1; i <= n; i++)
  fat[i] = fat[i - 1] * i;</pre>
                                                                                return splay(v), v;
   ifat[n] = fat[n].inv();
for (int i = n - 1; i >= 0; i—
                                                                             // V must be root. W will be the dad of V.
     ifat[i] = ifat[i + 1] * (i + 1);
                                                                             void link(int v, int w) {
                                                                                access(v);
// f(i) = y[i]
// return f(x0)
                                                                                t[v].p = w;
modInt getVal(vector<modInt> &y, ll x0) {
                                                                             // Removes edge (v, dad[v])
   int n = y.size();
                                                                             void cut(int v) {
   assert(fat[n-1] != 0);
                                                                                access(v);
  modInt x = x0;
for (int i = 0; i < n; i++) {
                                                                                if (t[v].ch[0] == -1) return;
                                                                                t[v].ch[0] = t[t[v].ch[0]].p' = -1;
     den[i] = ifat[n - i - 1] * if ((n - i - 1) % 2 == 1) {
                                            ifat[i];
                                                                             int lca(int v, int w) {
  if (findRoot(v) != findRoot(w)) return -1;
        den[i] = -den[i];
                                                                                access(v);
                                                                                return access(w);
   \begin{cases} [0] = 1; \\ \text{for (int } i = 1; i < n; i++) \\ [i] = l[i-1] * (x - (i-1)); \end{cases} 
                                                                                 // namespace LCT
  r[n-1] = 1;
for (int i = n - 2; i >= 0; i—) {
r[i] = r[i+1] * (x - (i+1));
                                                                             8.30 lct edge
                                                                             #include <bits/stdc++.h>
  modInt ans = 0;
for (int i = 0; i < n; i++) {
  modInt li = l[i] * r[i] * den[i];
  ans = (ans + (y[i] * li));
                                                                             using namespace std;
                                                                             // Link—Cut Tree — Edge, undirected version.
// All operations are O(log(n)) amortized.
                                                                             typedef long long ll;
                                                                             typedef pair<int, int> pii;
const int MAXN = 100010, MAXQ = 100010;
   return ans;
                                                                             namespace LCT {
     // namespace LagrangePoly
                                                                             struct node
                                                                                int p, ch[2];
ll val, sub;
8.29 lct
                                                                                bool rev;
                                                                                int sz, ar;
                                                                                ll lazy;
#include <bits/stdc++.h>
                                                                               node() {}
node(int v, int ar_)
: p(-1),
using namespace std;
// Link—Cut Tree, directed version.
// All operations are O(log(n)) amortized.
                                                                                        val(v),
const int MAXN = 200010;
                                                                                        sub(v)
namespace LCT {
                                                                                        rev(0)
struct node {
  int p, ch[2];
                                                                                        sz(ar_{-}),
                                                                                        ar(ar_),
lazy(0)
  node() { p = ch[0] = ch[1] = -1; }
                                                                                  lazy(0) { ch[0] = ch[1] = -1;
node t[MAXN]
                                                                               }
bool isRoot(int x) {
  return t[x].p == -1 or (t[t[x].p].ch[0] != x and
                                                                             node t[MAXN + MAXQ]; // MAXN + MAXQ
                                     t[t[x].p].ch[1] != x);
                                                                             map<pii, int> edges;
                                                                             int sz;
void rotate(int x) {
                                                                             void prop(int x)
  int p = t[x].p, pp = t[p].p;
if (!isRoot(p)) t[pp].ch[t[pp].ch[1] == p] = x;
bool d = t[p].ch[0] == x;
t[p].ch[!d] = t[x].ch[d], t[x].ch[d] = p;
if (t[p].ch[!d] + 1) t[t[p].ch[!d]].p = p;
                                                                                if (t[x].lazy)
                                                                                  if (t[x].ar) t[x].val += t[x].lazy;
                                                                                  t[x].sub += t[x].lazy * t[x].sz;
if (t[x].ch[0] + 1)
t[t[x].ch[0]].lazy += t[x].lazy;
                                                                                   if (t[x].ch[1] + 1)
   t[x].p = pp, t[p].p = x;
                                                                                     t[t[x].ch[1]].lazy += t[x].lazy;
void splay(int x) {
                                                                                if (t[x].rev) {
  swap(t[x].ch[0], t[x].ch[1]);
  while (!isRoot(x)) {
     int p = t[x].p, pp = t[p].p;
if (!isRoot(p))
```

 $rotate((t[pp].ch[0] == p) ^$

if (t[x].ch[0] + 1) t[t[x].ch[0]].rev ^= 1;
if (t[x].ch[1] + 1) t[t[x].ch[1]].rev ^= 1;

```
t[x].lazy = 0, t[x].rev = 0;
void update(int x) {
  t(x).sz = t[x].ar, t[x].sub = t[x].val;
for (int i = 0; i < 2; i++)
  if (t[x].ch[i] + 1) {</pre>
        prop(t[x].ch[i]);
        t[x].sz += t[t[x].ch[i]].sz;
        t[x].sub += t[t[x].ch[i]].sub;
bool is_root(int x) {
  return t[x].p == -1 or (t[t[x].p].ch[0] != x and
                                  t[t[x].p].ch[1] != x);
void rotate(int x) {
  int p = t[x].p, pp = t[p].p;
if (!is_root(p)) t[pp].ch[t[pp].ch[1] == p] = x;
  bool d = t[p].ch[0] == x;
  t[p].ch[!d] = t[x].ch[d], t[x].ch[d] = p;
if (t[p].ch[!d] + 1) t[t[p].ch[!d]].p = p;
  t[x].p = pp, t[p].p = x;
  update(p), update(x);
int splay(int x) {
  while (!is_root(x)) {
     int p = t[x].p, pp = t[p].p;
if (!is_root(p)) prop(pp);
     prop(p), prop(x);
if (!is_root(p))
        rotate((t[pp].ch[0] == p) ^
                           (t[p].ch[0] == x)
                       : p);
     rotate(x);
  return prop(x), x;
int access(int v) {
  int last = -1;
for (int w = v; w + 1;
     update(last = w), splay(v), w = t[v].p)
    splay(w), t[w].ch[1] = (last == -1 ? -1 : v);
  return last;
void rootify(int v);
void link_{-}(int v, int w) {
  rootify(w);
  t[w].p = v;
void cut_(int v, int w) {
  rootify(w), access(v);
t[v].ch[0] = t[t[v].ch[0]].p = -1;
void makeTree(int v, int w = 0, int ar = 0) {
  t[v] = node(w, ar);
// Public:
void init(int n) {
  edges.clear();
  for (int i = 0; i \le n; i++) makeTree(i);
int findRoot(int v)
  access(v), prop(v);
while (t[v].ch[0] + 1) v = t[v].ch[0], prop(v);
  return splay(v);
// Checks if v and w are connected
bool connected(int v, int w) {
  access(v), access(w);
  return v == w ? true : t[v].p != -1;
// Change v to be root
void rootify(int v) {
  access(v);
t[v].rev ^= 1;
^{\prime}// Sum of the edges in path from v to w
ll query(int v, int w) {
  rootify(w), access(v);
return t[v].sub;
```

```
// Sum +x in path from v to w
void update(int v, int w, int x) {
   rootify(w), access(v);
   t[v].lazy += x;
^{\prime}// Add edge (v, w) with weight x
void link(int v, int w, int x) {
  int id = MAXN + sz++;
   edges[pii(v, w)] = id;
  makeTree(id, x, 1);
link_(v, id), link_(id, w);
// Remove edge (v, w)
void cut(int v, int w) {
   int id = edges[pii(v, w)];
   cut_(v, id), cut_(id, w);
int lca(int v, int w) {
  access(v);
  return access(w);
   // namespace LCT
8.31 lct vertex
#include <bits/stdc++.h>
using namespace std;
// Link—Cut Tree — Vertex, undirected version.
// All operations are O(log(n)) amortized.
typedef long long ll;
typedef pair<int, int> pii;
const int MAXN = 200010;
namespace lct {
struct node
  int p, ch[2];
ll val, sub;
   bool rev;
   int sz;
  ll lazy;
node() {}
   node(int v)
        : p(-1)
          val(v),
          sub(v),
          rev(0),
          sz(1)
          lazy(0) {
     ch[0] = ch[1] = -1;
};
node t[MAXN];
void prop(int x)
  if (t[x].lazy)
     t[x].val += t[x].lazy,
          t[x].sub += t[x].lazy * t[x].sz;
       f (t[x].ch[0] + 1)
t[t[x].ch[0]].lazy += t[x].lazy;
     if (t[x].ch[1] + 1)
  t[t[x].ch[1]].lazy += t[x].lazy;
   if (t[x].rev)
     swap(t[x].ch[0], t[x].ch[1]);
     if (t[x].ch[0] + 1) t[t[x].ch[0]].rev ^= 1;
     if (t[x].ch[1] + 1) t[t[x].ch[1]].rev ^= 1;
   t[x].lazy = 0, t[x].rev = 0;
void update(int x) {
  t[x].sz = 1, t[x].sub = t[x].val;
for (int i = 0; i < 2; i++)
if (t[x].ch[i] + 1) {
       prop(t[x].ch[i]);
        t[x].sz += t[t[x].ch[i]].sz;
       t[x].sub += t[t[x].ch[i]].sub;
bool is_root(int x) {
   return t[x].p == -1 or (t[t[x].p].ch[0] != x and
                                 t[t[x].p].ch[1] != x);
void rotate(int x) {
  int p = t[x].p, pp = t[p].p;
if (!is\_root(p)) t[pp].ch[t[pp].ch[1] == p] = x;
```

```
bool d = t[p].ch[0] == x;

t[p].ch[!d] = t[x].ch[d], t[x].ch[d] = p;
  if (t[p].ch[!d] + 1) t[t[p].ch[!d]].p = p;
  t[x].p = pp, t[p].p = x;
  update(p), update(x);
int splay(int x)
  while (!is_root(x)) {
    int p = t[x].p, pp = t[p].p;
    if (!is_root(p)) prop(pp);
    prop(p), prop(x);
    if (!is_root(p))
       rotate((t[pp].ch[0] == p) ^
                        (t[p].ch[0] == x)
                    : p);
    rotate(x);
  return prop(x), x;
int access(int v) {
  int last = -1;
  for (int w = v; w + 1;
	update(last = w), splay(v), w = t[v].p)
	splay(w), t[w].ch[1] = (last == -1 ? -1 : v);
                                                                  ll getMax(ll x)
  return last;
// Public:
                                                               };
void makeTree(int v, int w) { t[v] = node(w); }
int findRoot(int v) {
  access(v), prop(v);
  while (t[v].ch[0] + 1) v = t[v].ch[0], prop(v);
  return splay(v);
// Checks if v and w are connected
bool connected(int v, int w) {
  access(v), access(w);
  return v == w? true : t[v].p != -1;
// Change v to be root
void rootify(int v) {
  access(v);
t[v].rev ^= 1;
                                                                  ll x =
                                                                  a %= MOD
                                                                 while (b)
^{\prime}// Sum of the weight in path from {	t v} to {	t w}
11 query(int v, int w) {
  rootify(w), access(v);
                                                                    b >>= 1;
  return t[v].sub;
                                                                  return x;
^{\prime}// Sum +x in path from v to w
void update(int v, int w, int x) {
  rootify(w), access(v);
  t[v].lazy += x;
// Add edge (v, w)
                                                                    ll t = 0;
void link(int v, int w) {
  rootify(w);
  t[w].p = v;
// Remove edge (v, w)
void cut(int v, int w) {
                                                                       lf = i;
  rootify(w), access(v);
                                                                      continue;
  t[v].ch[0] = t[t[v].ch[0]].p = -1;
int lca(int v, int w) {
  access(v):
  return access(w);
   // namespace lct
8.32 line container
#include <bits/stdc++.h>
#pragma once
using ll = long long;
using namespace std;
                                                                    cur = c;
struct Line {
  mutable ll k, m, p;
  bool operator<(const Line& o) const {</pre>
```

return k < o.k;

bool operator<(ll x) const { return p < x; }</pre>

```
struct LineContainer : multiset<Line, less<>> {
  // (for doubles, use inf = 1/.0, div(a,b) = a/b
  static const ll inf = LLONG_MAX;
  ll div(ll a, ll b) { // floored division
    return a / b - ((a ^ b) < 0 && a % b);
}</pre>
   bool isect(iterator x, iterator y)
      if (y == end()) return x \rightarrow p = inf, 0;
      if (x->k == y->k)
         x - p = x - m > y - m ? inf : -inf;
         x \rightarrow p = div(y \rightarrow m - x \rightarrow m, x \rightarrow k - y \rightarrow k);
      return x\rightarrow p >= y\rightarrow p;
   void add(ll k, ll m) {
  auto z = insert({k, m, 0}), y = z++, x = y;
  while (isect(y, z)) z = erase(z);
      if (x != begin() \&\& isect(-x, y))
         isect(x, y = erase(y))
      while ((y = x) != begin() \&\& (--x)->p >= y->p)
         isect(x, erase(y));
      assert(!empty());
      auto l = *lower_bound(x);
      return l.k * x + l.m;
8.33 linear sequence with berlekamp massey
#include <bits/stdc++.h>
using namespace std;
// Source: https://codeforces.com/blog/entry/61306
typedef long long ll;
const int MOD = 104857601;
// Work only to prime MOD
namespace LinearSeq
const int MAXN = 233333;
ll fastPow(ll a, ll b) {
      if (b \& 1) x = (x * a) % MOD;

a = (a * a) % MOD;
inline vector<int> BM(vector<int> x) {
   vector<int> ls, cur;
int lf = 0, ld = 0;
for (int i = 0; i < int(x.size()); i++) {</pre>
      for (int j = 0; j < int(cur.size()); ++j)
  t = (t + x[i - j - 1] * (ll)cur[j]) % MOD;</pre>
      if ((t - x[i]) % MOD == 0) continue; if (!cur.size()) {
         cur.resize(i + 1);
         ld = (t - x[i]) % MOD;
      -(x[i] - t) * fastPow(ld, MOD - 2) % MOD;
vector<int> c(i - lf - 1);
      c.push_back(k);
      for (int j = 0; j < int(ls.size()); ++j)
  c.push_back(-ls[j] * k % MOD);</pre>
      if (c.size() < cur.size())</pre>
         c.resize(cur.size())
      for (int j = 0; j < int(cur.size()); ++j)
  c[j] = (c[j] + cur[j]) % MOD;</pre>
      if (i - lf + (int)ls.size() >=
            (int)cur.size())
= cur, lf = i, ld = (t - x[i]) % MOD;
   for (int i = 0; i < int(cur.size()); ++i)
      cur[i] = (cur[i] % MOD + MOD) % MOD;
   return cur;
```

ll a[MAXN], h[MAXN], t2[MAXN], s[MAXN], t[MAXN];

```
inline void mull(ll* p, ll* q) {
  for (int i = 0; i < m + m; ++i) t2[i] = 0;
  for (int i = 0; i < m; ++i)</pre>
                                                                                      for (size_t j = 0; j < B.size(); ++j) { A[j + m] = coef * B[j] % mod;
                                                                                         if(A[j + m] < 0) A[j + m] += mod;
     if (p[i])
for (int j = 0; j < m; ++j)
                                                                                      B = temp, b = d, m = 0;
           t2[i + j] =
                                                                                   } else {
                                                                                      extend(A, B.size() + m);
int64 coef = d * inverse(b) % mod;
for (size_t j = 0; j < B.size(); ++j) {
    A[j + m] = coef * B[j] % mod;
    if (A[i + m] < 0) A[i + m] += mod;
                (t2[i + j] + p[i] * q[j]) % MOD;
  for (int i =
   if (t2[i])
               i = m + m - 1; i >= m;
        for (int j = m - 1; ~j; —j)
t2[i - j - 1] =
                                                                                        if (A[j + m] < 0) A[j + m] += mod;
  }
                                                                                }
inline ll calc(ll K) {
  for (int i = m; ~i; —i) s[i] = t[i] = 0;
  // ainit
                                                                                return A;
                                                                              static void exgcd(int64 a, int64 b, int64 &g,
  s[0] = 1;
if (m != 1)
                                                                                                      int64 &x, int64 &y) {
                                                                                if (!b)
     t[1] = 1;
                                                                                  x = 1, y = 0, g = a;
  else
t[0] = h[0];
                                                                                else {
                                                                                   exgcd(b, a % b, g, y, x);
  t[0] = n[0],
while (K) {
  if (K & 1) mull(s, t);
  mull(t, t);
                                                                                   y = x * (a / b);
     K >>= 1;
                                                                             static int64 crt(const vec &c, const vec &m) {
  fl su = 0;
for (int i = 0; i < m; ++i)
  su = (su + s[i] * a[i]) % MOD;
return (su % MOD + MOD) % MOD;</pre>
                                                                                int n = c.size();
                                                                                int64 M = 1, ans = 0;

for (int i = 0; i < n; ++i) M *= m[i];

for (int i = 0; i < n; ++i) {

  int64 x, y, g, tm = M / m[i];

  exgcd(tm, m[i], g, x, y);
// Public:
                                                                                   ans = (ans + tm * x * c[i] % M) % M;
// 0(MAXN + |x|^2 * log(N))
inline int findElementInPositionN(vector<int> x,
                                                                                return (ans + M) % M;
  if (n < int(x.size())) return x[n];
vector<int> v = BM(x);
                                                                              static vec ReedsSloane(const vec &s,
                                                                                                             int64 mod)
  m = v.size();
                                                                                auto inverse = [](int64 a, int64 m) {
  if (!m) return 0;
for (int i = 0; i < m; ++i)
  h[i] = v[i], a[i] = x[i];
return calc(n);
                                                                                   int64 d, x, y;
exgcd(a, m, d, x, y);
return d == 1 ? (x % m + m) % m : -1;
                                                                                // namespace LinearSeq
                                                                                                     ? a.size() —
8.34 linear_sequence_with_reeds_sloane
                                                                                                        -1000;
                                                                                   int db = (b.size() > 1 ||
#include <bits/stdc++.h>
                                                                                                 (b.size() == 1 \&\& b[0]))
using namespace std;
                                                                                                     ? b.size() - 1
// Source:
                                                                                                       -1000;
// https://github.com/zimpha/algorithmic—library/
                                                                                   return std::max(da, db + 1);
     blob/master/cpp/mathematics/linear—recurrence.
                                                                                auto prime_power = [&](const vec &s,
struct LinearRecurrence {
                                                                                                                int64 mod, int64 p,
  using int64 = long long;
                                                                                                                int64 e) {
  using vec = std::vector<int64>;
                                                                                   // linear feedback shift register mod p^e, p
// is prime
  static void extend(vec &a, size_t d,
                             int64 value = 0) {
                                                                                   std::vector<vec> a(e), b(e), an(e), bn(e),
     if (d <= a.size()) return;</pre>
                                                                                   ao(e), bo(e);
vec t(e), u(e), r(e), to(e, 1), uo(e),
pw(e + 1, 1);
     a.resize(d, value);
  static vec BerlekampMassey(const vec &s,
                                        int64 mod) {
                                                                                   for (int i = 1; i <= e; ++i) {
    pw[i] = pw[i - 1] * p;
     std::function<int64(int64)> inverse =
  [&](int64 a) {
                                                                                      assert(pw[i] <= mod);</pre>
             return a == 1
                             ? 1
: (int64)(mod - mod / a) *
                                                                                   for (int64 i = 0; i < e; ++i)
                                                                                      a[i] = {pw[i]}, an[i] = {pw[i]};
b[i] = {0}, bn[i] = {s[0] * pw[i] % mod};
t[i] = s[0] * pw[i] % mod;
                                     inverse(mod % a) % mod;
     };
vec A = {1}, B = {1};
int64 b = s[0];
assert(b != 0);
                                                                                      if (t[i] == 0) {
  t[i] = 1, u[i] = e;
} else {
     for (size_t i = 1, m = 1; i < s.size();
                                                                                         for (u[i] = 0; t[i] % p == 0;
            ++i, m++) {
        int64 d = 0;
                                                                                               t[i] /= p, ++u[i])
        for (size_t j = 0; j < A.size(); ++j) {
   d += A[j] * s[i - j] % mod;
                                                                                      }
                                                                                   for (size_t k = 1; k < s.size(); ++k) {
        if (!(d %= mod)) continue;
                                                                                      for (int g = 0; g < e; ++g) {
  if (L(an[g], bn[g]) > L(a[g], b[g])) {
        if (2 * (A.size() - 1) <= i) {
           auto temp = A;
extend(A, B.size() + m);
int64 coef = d * inverse(b) % mod;
                                                                                           ao[g] = a[e - 1 - u[g]];

bo[g] = b[e - 1 - u[g]];
```

```
to[g] = t[e - 1 - u[g]];

uo[g] = u[e - 1 - u[g]];
           r[g] = k - 1;
     a = an, b = bn;
     for (int o = 0; o < e; ++o) {
        int64 d = 0;
for (size_t i = 0;
           i < a[o].size() \&\& i <= k; ++i) { d = (d + a[o][i] * s[k - i]) % mod; }
        if (d == 0) {
           t[o] = 1, u[o] = e;
        } else {
           for (u[o] = 0, t[o] = d;
t[o] % p == 0; t[o] /= p, ++u[o])
           int g = e - 1 - u[o];
if (L(a[g], b[g]) == 0) {
  extend(bn[o], k + 1);
  bn[o][k] = (bn[o][k] + d) % mod;
} older f
           } else {
              int64 coef =
                    t[o] * inverse(to[g], mod) %
                    mod * pw[u[o] - uo[g]] % mod;
              int m = k - r[g];
              assert(m >= 0);
              extend(an[o], ao[g].size() + m);
extend(bn[o], bo[g].size() + m);
              for (size_t i = 0; i < ao[g].size();
                     ++i) {
                 an[o][i + m] -=

coef * ao[g][i] % mod;

if (an[o][i + m] < 0)
                    an[o][i+m] += mod;
              while (an[o].size() \&\& an[o].back() == 0)
                 an[o].pop_back();
              for (size_t i = 0; i < bo[g].size();
                      ++i) {
                 bn[o][i + m] =
                       coef * bo[g][i] % mod;
                 if (bn[o][i + m] < 0)
  bn[o][i + m] = mod;</pre>
              while (bn[o].size() &&
                        bn[o].back() == 0)
                 bn[o].pop_back();
        }
  return std::make_pair(an[0], bn[0]);
std::vector<std::tuple<int64, int64, int>>
     fac
for (int64 i = 2; i * i <= mod; ++i)
  if (mod % i == 0) {</pre>
     int64 cnt = 0, pw = 1;
while (mod % i == 0)
mod /= i, ++cnt, pw *= i;
     fac.emplace_back(pw, i, cnt);
if (mod > 1) fac.emplace_back(mod, mod, 1);
std::vector<vec> as;
size_t n = 0;
for (auto &&x : fac) {
  int64 mod, p, e;
  vec a, b;
std::tie(mod, p, e) = x;
  auto ss = s;
for (auto &&x : ss) x %= mod;
std::tie(a, b) = prime_power(ss, mod, p, e);
  n = std::max(n, a.size());
vec a(n), c(as.size()), m(as.size());
for (size_t i = 0; i < n; ++i) {
  for (size_t j = 0; j < as.size(); ++j) {</pre>
     m[j] = std::get<0>(fac[j]);
     c[j] = i < as[j].size() ? as[j][i] : 0;
  a[i] = crt(c, m);
```

```
return a;
  LinearRecurrence(const vec &s, const vec &c,
                          int64 mod)
        : init(s)
           trans(c),
           mod(mod)
           m(s.size()) {}
   LinearRecurrence(const vec &s, int64 mod,
                          bool is_prime = true)
        : mod(mod) -
     assert(s.size() % 2 == 0);
     vec A;
     if (is_prime)
        A = BerlekampMassey(s, mod);
     else
        A = ReedsSloane(s, mod);
     m = s.size() / 2;
A.resize(m + 1, 0);
trans.resize(m);
for (int i = 0; i < m; ++i) {
   trans[i] = (mod - A[i + 1]) % mod;</pre>
     if (m == 0) m = 1, trans = {1};
     std::reverse(trans.begin(), trans.end());
init = {s.begin(), s.begin() + m};
   int64 calc(int64 n)
     if (mod == 1) return 0;
     if (n < m) return init[n];</pre>
     vec`v(m), u(m << 1);
int64 msk = !!n;</pre>
          (int64 m = n; m > 1; m >>= 1) msk <<= 1;
     v[0] = 1 % mod;
for (int64 x = 0; msk; msk >>= 1, x <<= 1) {
    std::fill_n(u.begin(), m * 2, 0);</pre>
        x = !!(n \& msk);
        if (x < m)

u[x] = 1 % mod;

else { // can be optimized by fft/ntt
           for (int i = 0; i < m; ++i) {
  for (int j = 0, t = i + (x \& 1); j < m;
                ++j, ++t) {
u[t] = (u[t] + v[i] * v[j]) % mod;
             }
           for (int i = m * 2 - 1; i >= m; —i) { for (int j = 0, t = i - m; j < m;
                    ++j, ++t) {
                u[t] = (u[t] + trans[j] * u[i]) % mod;
          }
        v = \{u.begin(), u.begin() + m\};
     int64 ret = 0;
     for (int i = 0; i < m; ++i)
        ret = (ret + v[i] * init[i]) % mod;
     return ret;
  vec init, trans;
   int64 mod;
   int m;
};
8.35 min cyclic string
   int n = s.size();
```

```
#include <bits/stdc++.h>
using namespace std;
string min_cyclic_string(string s) {
   s += s;
  int i = 0, ans = 0; while (i < n / 2) {
     ans = i;
int j = i + 1, k = i;
while (j < n && s[k] <= s[j]) {
        if (s[k] < s[j])
           k = i;
        else
           k++;
        j++;
```

```
return l < oth.l;
if (r / BLOCK_SIZE != oth.r / BLOCK_SIZE)</pre>
    while (i <= k) i += j - k;
                                                                         return r < oth.r;
  return s.substr(ans, n / 2);
                                                                       return t < oth.t;
                                                                     }
                                                                  struct Update {
                                                                    int pos, newV, oldV, t;
8.36 mincut
                                                                  // 0(Q * N^(2/3)): N=10^5 -> 1.5s
#include <bits/stdc++.h>
                                                                  vector<int> mo_s_algorithm(vector<Query> vq,
using namespace std;
                                                                                                    vector<Update> vu) {
typedef long long ll;
                                                                     vector<pii> answers;
// This algorithm finds the Global Min—Cut in
                                                                     sort(all(vq));
// 0(|V|^3)
                                                                     clearAnswer()
namespace MinCut {
const int MAXN = 510;
bool exist[MAXN], in_a[MAXN];
ll g[MAXN][MAXN], w[MAXN];
                                                                     int L = 0, R = 0, T = 0, szT = vu.size();
                                                                     add(v[0]);
                                                                     for (Query q : vq) {
  while (q.l < L) add(v[—L]);</pre>
vector<int> v[MAXN];
                                                                       while (R < q.r) add(v[++R]);
int n;
                                                                       while (L < q.l) remove(v[L++]); while (q.r < R) remove(v[R-]);
void init(int n1) {
  n = n1;
                                                                       while (T < szT \text{ and } vu[T].t <= q.t) {
  memset(g, 0, sizeof(g));
                                                                         Update \&u = vu[T++];
                                                                         if (L <= u.pos and u.pos <= R) {
  remove(u.oldV);</pre>
void addEdge(int a, int b, int w1) {
  if (a == b) return;
g[a][b] += w1;
                                                                            add(u.newV);
  g[b][a] += w1;
                                                                          v[u.pos] = u.newV;
pair<ll, vector<int>> mincut() {
                                                                       while (T > 0 \text{ and } vu[T - 1].t > q.t) {
  ll best_cost = 0x3f3f3f3f3f3f3f3fLL;
  vector<int> best_cut;
for (int i = 0; i < n; ++i) v[i].assign(1, i);</pre>
                                                                         Update \&u = vu[-T];
                                                                          if (L \le u.pos and u.pos \le R) {
  remove(u.newV);
                                                                            add(u.oldV);
                                                                         v[u.pos] = u.oldV;
                                                                       answers.emplace_back(q.t, getAnswer());
           ++it) {
       int sel = -1;

for (int i = 0; i < n; ++i)

if (exist[i] && !in_a[i] &&
                                                                     sort(all(answers));
                                                                     vector<int> ret;
for (auto [t, x] : answers) ret.push_back(x);
               (sel == -1 \mid | w[i] > w[sel]))
            sel̀ = i;
                                                                     return ret;
       if (it == n - ph - 1) {
          if (w[sel] < best_cost)</pre>
         best_cost = w[sel], best_cut = v[sel];
v[prev].insert(v[prev].end(),
                                                                  8.38 nearest pair of points
                            v[sel].begin(),
         v[sel].end());
for (int i = 0; i < n; ++i)
  g[prev][i] = g[i][prev] += g[sel][i];</pre>
                                                                  #include <bits/stdc++.h>
                                                                  using namespace std;
                                                                  struct pt {
          exist[sel] = false;
                                                                     long long x, y, id;
       } else {
  in_a[sel] = true;
  for (int i = 0; i < n; ++i)
    w[i] += g[sel][i];</pre>
                                                                     pt() {}
                                                                     prev = sel;
                                                                  namespace NearestPairOfPoints {
                                                                  struct cmp_x {
    }
                                                                     bool operator()(const pt& a,
                                                                                        const pt& b) const {
  return {best_cost, best_cut};
                                                                       return a.x < b.x \mid | (a.x == b.x && a.y < b.y);
                                                                     }
    // namespace MinCut
                                                                  struct cmp_y {
                                                                     bool operator()(const pt& a,
8.37 mo with update
                                                                                         const pt& b) const {
                                                                       return a.y < b.y;
                                                                     }
#include <bits/stdc++.h>
#define all(x) x.begin(), x.end()
                                                                  int n;
using namespace std;
                                                                  vector<pt> v;
using pii = pair<int, int>;

const int INF = 0x3f3f3f3f;

const int BLOCK_SIZE = 2800; // (2*N^2)^(1/3)

const int MAXN = 100010;
                                                                  vector<pt> t;
                                                                  double mindist;
pair<int, int> best_pair;
                                                                  void upd_ans(const pt& a, const pt& b) {
int v[MAXN];
void remove(int x);
void add(int x);
                                                                     double dist = sqrt((a.x - b.x) * (a.x - b.x) +
                                                                                            (a.y - b.y) * (a.y - b.y);
void clearAnswer();
                                                                     if (dist < mindist)</pre>
                                                                       mindist = dist;
int getAnswer();
struct Query {
                                                                       best_pair = {a.id, b.id};
  int_l, r, t;
  bool operator<(const Query &oth) const</pre>
```

void rec(int l, int r) {

if (l / BLOCK_SIZE != oth.l / BLOCK_SIZE)

```
if (r - l <= 3) {
  for (int i = l; i < r; ++i) {
    for (int j = i + 1; j < r; ++j) {</pre>
                          upd_ans(v[i], v[j]);
              sort(v.begin() + l, v.begin() + r, cmp_y());
              return;
       int m = (l + r) >> 1;
int midx = v[m].x;
       rec(l, m);
rec(m, r);
       merge(v.begin() + l, v.begin() + m,
                             v.begin() + m, v.begin() + r, t.begin(),
                            cmp_y());
       copy(t.begin(), t.begin() + r - l,
                         v.begin() + l);
       int tsz = 0;
       for (int i = l; i < r; ++i) {
   if (abs(v[i].x - midx) < midulum interpretation of the context of the con
                     (abs(v[i].x - midx) < mindist) {
for (int j = tsz - 1;</pre>
                                       j >= 0 \& v[i].y - t[j].y < mindist;
                             upd_ans(v[i], t[j]);
                     t[tsz++] = v[i];
      }
pair<int, int> solve(vector<pt> _v) {
       n = v.size();
       t.resize(n)
       sort(v.begin(), v.end(), cmp_x());
       mindist = 1E20;
       rec(0, n);
       return best_pair;
            // namespace NearestPairOfPoints
8.39 parallel binary search
#include <bits/stdc++.h>
using namespace std;
const int MAXN = 100010;
int ans[MAXN];
       if (v.empty()) return;
       if (i == j) {
              for (int x : v) ans[x] = i;
              return;
      int mid = (i + j) / 2;
for (int k = i; k <= mid; k++) add(k);
vector<int> left, right;
       for (int x : v) {
   if (test(x))
                     left.push_back(x);
```

bool test(int x); void add(int k); void remove(int k); void solve(int i, int j, vector<int> &v) { right.push_back(x); solve(mid + 1, j, right); for (int k = mid; k >= i; k—) remove(k); // Or roolback(); solve(i, mid, left);

8.40 permutation

```
#include <bits/stdc++.h>
using namespace std;
using ll = long long;
mt19937_64 rng(
     (int)std::chrono::steady_clock::now()
          .time_since_epoch()
          .count())
namespace Permutation {
const int MAXN = 500010;
ll mp[MAXN], sumXor[MAXN], p[MAXN + 1], inv[MAXN];
```

```
void init(vector<int> v)
   sumXor[0] = inv[0] = p[0] = 0;
   for (int i = 0; i < MAXN; i++) {
  mp[i] = rng() + 1;
  p[i + 1] = p[i] ^ mp[i];</pre>
  for (int i = 0; i < v.size(); i++
  if (v[i] < 0 or v[i] >= MAXN) {
    inv[i + 1] = 1 + inv[i];
}
                                            i++) {
        sumXor[i + 1] = sumXor[i];
     } else {
        inv[i + 1] = inv[i]
        sumXor[i + 1] = sumXor[i] ^ mp[v[i]];
  }
// Verify if {v[l], v[l+1], ..., v[r]} is {0, 1, // ..., r-l+1} 0-indexed;
// ... , r_{-l+1} , u_{-1} indexed, bool isPermutation(int l, int r) {
  }; // namespace Permutation
8.41 range color
#include <bits/stdc++.h>
using namespace std;
class RangeColor {
 private:
   typedef long long ll;
   struct Node {
     ll l, r;
int color;
     Node() {}
Node(ll l1, ll r1, int color1)
: l(l1), r(r1), color(color1) {}
bool operator<(const Node &oth) const {
        return r < oth.r;
   std::set<Node> st;
   vector<ll> ans;
 public:
   RangeColor(ll first, ll last, int maxColor) {
     ans.resize(maxColor + 1);
ans[0] = last - first + 1LL;
      st.insert(Node(first, last, 0));
   // get color in position x
   int get(ll x) {
     auto p = st.upper_bound(Node(0, x - 1LL, -1));
      return p—>color;
  // set newColor in [a, b]
void set(ll a, ll b, int newColor) {
  auto p = st.upper_bound(Node(0, a - 1LL, -1));
      assert(p != st.end());
      ll l = p \rightarrow l;
     ll r = p->r;
int oldColor = p->color;
ans[oldColor] == (r - l + 1LL);
     p = st.erase(p);
if (l < a) {</pre>
        ans[oldColor] += (a - l);
st.insert(Node(l, a - 1LL, oldColor));
     if (b < r)
        ans[oldColor] += (r - b);
st.insert(Node(b + 1LL, r, oldColor));
     while ((p != st.end()) and (p->l <= b)) {
        l = p \rightarrow l;

r = p \rightarrow r;
        oldColor = p->color;
         ans[oldColor] -= (r - l + 1LL);
           ans[oldColor] += (r - b);
           st.erase(p);
           st.insert(Node(b + 1LL, r, oldColor));
```

```
if (l & 1) ans.push_back(l++);
if (r & 1) ans.push_back(—r);
            } else {
  p = st.erase(p);
                                                                                                                       return ans:
        ans[newColor] += (b - a + 1LL);
        st.insert(Node(a, b, newColor));
                                                                                                                  Node query(int x1, int y1, int x2, int y2) {
                                                                                                                      auto c1 = getCover(x1, x2 + 1, n);
auto c2 = getCover(y1, y2 + 1, m);
    ll countColor(int x) { return ans[x]; }
                                                                                                                       Node ans = neutral
                                                                                                                      for (auto i : c1) {
  for (auto j : c2) {
8.42 rank matrix
                                                                                                                              ans = join(ans, seg[i][j]);
#include <bits/stdc++.h>
                                                                                                                       return ans;
using namespace std;
typedef long double ld;
                                                                                                              };
const ld EPS = 1e-9;
int compute_rank(vector<vector<ld>>> A) {
    int n = A.size();
                                                                                                              8.44 simpson integration
    int m = A[0].size();
    int rank = max(n, m);
vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector<br/>vector
                                                                                                              #include <bits/stdc++.h>
                                                              false):
                                                                                                              using namespace std;
    for (int i = 0; i < m; ++i) {
                                                                                                               double f(double x)
        int j;
for (j = 0; j < n; ++j) {
                                                                                                               const int N = 1000000;
                                                                                                              double simpson_integration(double a, double b) {
            if (!row_selected[j] && abs(A[j][i]) > EPS)
                                                                                                                  double h = (b - a) / N;

double s = f(a) + f(b); // a = x_0 and b = x_2n

for (int i = 1; i <= N - 1;

++i) { // Refer to final Simpson's formula

double x = a + h * i;

s += f(x) * ((i \& 1) ? 4 : 2);
                break;
        if (j == n) {
        rank—;
} else {
            row_selected[j] = true;
            for (int p = i + 1; p < m; p++)
                                                                                                                   s *= h / 3;
                A[j][p] /= A[j][i];
                                                                                                                   return s;
            for (int k = 0; k < n; k++) {
  if (k != j && abs(A[k][i]) > EPS) {
                                                                                                              }
                    for (int p = i + 1; p < m; p++)
                        A[k][p] = A[j][p] * A[k][i];
                                                                                                              8.45 sqrt_decomposition
           }
                                                                                                              #include <bits/stdc++.h>
       }
                                                                                                              using namespace std;
                                                                                                              struct SqrtDecomposition {
    return rank;
                                                                                                                   typedef long long t_sqrt;
                                                                                                                   int sqrtLen;
                                                                                                                   vector<t_sqrt> block;
                                                                                                                   vector<t_sqrt> v;
8.43 segment tree2d
                                                                                                                   template <class MyIterator>
                                                                                                                   SqrtDecomposition(MyIterator begin,
#include <bits/stdc++.h>
                                                                                                                                                       MyIterator end) {
using namespace std;
                                                                                                                      int n = end - begin;
struct SegTree2D {
                                                                                                                       sqrtLen = (int)sqrt(n + .0) + 1;
  private:
                                                                                                                       v.resize(n);
    int n, m;
typedef int Node;
                                                                                                                       block.resize(sqrtLen + 5);
                                                                                                                              (int i = 0; i < n; i++, begin++) {
    Node neutral = -0x3f3f3f3f;
                                                                                                                          v[i] = (*begin);
    vector<vector<Node>> seg;
                                                                                                                          block[i / sqrtLen] += v[i];
    Node join(Node a, Node \vec{b}) { return max(a, b); }
                                                                                                                   // 0—indexed
    SegTree2D(int n1, int m1) {
                                                                                                                   void update(int idx, t_sqrt new_value) {
        n = n1, m = m1;
                                                                                                                      t_sqrt d = new_value - v[idx];
v[idx] += d;
block[idx / sqrtLen] += d;
        seg.assign(2 * n, vector<Node>(2 * m, 0));
    void update(int x, int y, int val)
        assert(0 \le x \&\& x < n \&\& 0 \le y \&\& y < m);
                                                                                                                  // 0-indexed [l, r]
t_sqrt query(int l, int r) {
        x += n, y += m;

seg[x][y] = val;
                                                                                                                       t_sqrt sum = 0;
        for (int j = y / 2; j > 0; j /= 2)
        join(seg[x][2 * j], seg[x][2 * j + 1]);
for (x /= 2; x > 0; x /= 2) {
  seg[x][y] =
                                                                                                                       int c_l = l / sqrtLen, c_r = r / sqrtLen;
                                                                                                                          ' (c_l == c_r) {
for (int i = l; i <= r; i++) sum += v[i];</pre>
                                                                                                                       } else
                                                                                                                          for (int i = l,
end = (c_l + 1) * sqrtLen - 1;
            join(seg[2 * x][y], seg[2 * x + 1][y]);
for (int j = y / 2; j > 0; j /= 2) {
  seg[x][j] = join(seg[x][2 * j],
                                                                                                                                     i <= end; i++)
                                                                                                                          sum += v[i];
for (int i = c
                                                  seg[x][2 * j + 1]);
                                                                                                                                                         _l + 1; i <= c_r - 1; i++)
            }
                                                                                                                               sum += block[i];
       }
                                                                                                                           for (int i = c_r * sqrtLen; i <= r; i++)
                                                                                                                               sum += v[i];
    vector<int> getCover(int l, int r, int N) {
        l = std::max(0, l);
                                                                                                                       return sum;
        r = std::min(N, r);
```

};

vector<int> ans

for (l += N, r += N; l < r; l /= 2, r /= 2) {

```
#include <algorithm>
#include <iostream>
constexpr int N = 510, p = 998244353;
int fp(int a, int b) {
  int ans = 1, off = a;
  while (b)
    if (b \& 1) ans = 1ll * ans * off % p;
    off = 1ll * off * off % p;
    b >>= 1;
  return ans;
}
int gauss(int (&dat)[N][N], int (&ans)[N],
            int (&basis)[N][N], int n, int m) {
  int k = 1;
  static int col[N];
  for (int i = 1; i \le m \&\& k \le n; ++i) {
    int pos = 0;
    for (int j = k; j <= n; ++j)
       if (dat[j][i]) {
         pos = j;
         break;
    if (!pos) continue;
    col[k] = i;
    if (pos != k) {
       for (int j = i; j \le m + 1; ++j)
         std::swap(dat[pos][j], dat[k][j]);
    int rv = fp(dat[k][i], p - 2);
    for (int j = i; j <= m + 1; ++j)
dat[k][j] = 1ll * dat[k][j] * rv % p;
    for (int j = k + 1; j \le n; ++j)
       if (dat[j][i]) {
   int num = p — dat[j][i];
          for (int t = i; t \le m + 1; ++t)
            dat[j][t] = (dat[j][t] +
                            1ll * num * dat[k][t]) %
       }
  for (int i = k; i <= n; ++i)
  if (dat[i][m + 1]) return -1;</pre>
   —k;
  int R = m - k;
  for (int i = 1; i \le m; ++i) ans[i] = 0;
  for (int i = k; i; —i) {
    ans[col[i]] = dat[i][m + 1];
    for (int j = i + 1; j <= k; ++j)
  ans[col[i]] = (ans[col[i]] +</pre>
                         1ll * (p - dat[i][col[j]]) *
                              ans[col[j]]) %
  for (int i = 1; i \le R; ++i)
for (int j = 1; j \le m; ++j) basis[i][j] = 0;
  col[k + 1] = m + 1;
  col[0] = 0;
  for (int i = 0, t = 0; i <= k; ++i) {
  for (int j = col[i] + 1; j < col[i + 1];
       ++t;
       for (int l = i; l; —l) {
  int c = dat[l][j];
            or (int r = l + 1; r <= i; ++r
c = (c + 1ll * dat[l][col[r]]
         for (int r = l + 1;
                            basis[t][col[r]]) %
         basis[t][col[l]] = (p - c) \% p;
       basis[t][j] = 1;
    }
```

```
}
  return R;
}
int main() {
  int n, m;
  static int dat[N][N], ans[N], basis[N][N];
  std::ios::sync_with_stdio(false);
  std::cin.tie(0);
  std::cin >> n >> m;
  for (int i = 1; i <= n; ++i)
for (int j = 1; j <= m; ++j)
      std::cin >> dat[i][j];
  for (int j = 1; j <= n; ++j)
    std::cin >> dat[j][m + 1];
  int R = gauss(dat, ans, basis, n, m);
  if (!(~R)) {
   std::cout << "-1\n";</pre>
    return 0;
  std::cout << R << '\n';
  for (int j = 1; j \le m; ++j)
    std::cout << ans[j] << '</pre>
  std::cout << '\n';
  for (int i = 1; i \le R; ++i) {
    for (int j = 1; j \le m; ++j)
      std::cout << basis[i][j] << ' ';
    std::cout << '\n';</pre>
  return 0;
}
8.47 treap
#include <bits/stdc++.h>
using namespace std;
namespace Treap {
const int N = 500010;
typedef long long treap_t;
treap_t X[N];
int en = 1, Y[N], sz[N], L[N], R[N], root;
const treap_t neutral = 0;
treap_t op_val[N];
inline treap_t join(treap_t a, treap_t b,
                      treap_t c) {
  return a + b + c;
void calc(
    int u) {
               // update node given children info
  sz[u] = sz[L[u]] + 1 + sz[R[u]];
  // code here, no recursion
  op_val[u] =
      join(op_val[L[u]], X[u], op_val[R[u]]);
void unlaze(int u) {
  if (!u) return;
  // code here, no recursion
void split(int u, treap_t x, int &l,
            int &r) { // l gets <= x, r gets > x
  unlaze(u);
  if (!u) return (void)(l = r = 0);
if (X[u] <= x) {
    split(R[u], x, l, r);
    R[u] = l;
      = u;
  } else -
    split(L[u], x, l, r);
    L[u] = r;
r = u:
  calc(u);
void split_sz(
    int u, int s, int &l,
int &r) { // l gets first s, r gets remaining
  unlaze(u);
  if (!u) return (void)(l = r = 0);
if (sz[L[u]] < s) {</pre>
```

 $split_sz(R[u], s - sz[L[u]] - 1, l, r);$

```
R[u] = l;
    l_= u;
  } else {
    split_sz(L[u], s, l, r);
    L[u] = r;
r = u;
  calc(u);
int merge(int l, int r) { // els on l <= els on r
  unlaze(l);
  unlaze(r):
  if (!l || !r) return l + r;
  if (Y[l] > Y[r])
    R[l] = merge(R[l], r);
  } else {
   L[r] = merge(l, L[r]);
   u = r;
  calc(u);
  return u;
int new_node(treap_t x) {
  X[en] = x;
op_val[en] = x;
  return en++;
int nth(int u, int idx) {
  if (!u) return 0;
  unlaze(u);
  if (idx <= sz[L[u]])
  return nth(L[u], idx);
else if (idx == sz[L[u]] + 1)
    return`u;
    return nth(R[u], idx - sz[L[u]] - 1);
// Public
void init(
    int n = N
             1) { // call before using other funcs
  // init position 0
  sz[0] = 0;
  op_val[0] = neutral;
  // init Treap
  root = 0;
  std::mt19937 rng(
       (int)std::chrono::steady_clock::now()
           .time_since_epoch()
  .count());
for (int i = en = 1; i <= n; i++) {
    Y[i] = i;
sz[i] = 1;
    L[i] = R[i] = 0;
  shuffle(Y + 1, Y + n + 1, rng);
void insert(treap_t x) {
  split(root, x, a, b);
  root = merge(merge(a, new_node(x)), b);
void erase(treap_t x) {
  int a, b, c, d; split(root, x-1, a, b);
  split(b, x, c, d);
split_sz(c, 1, b, c);
  root = merge(a, merge(c, d));
int count(treap_t x) {
  int a, b, c, d;
split(root, x - 1, a, b);
  split(b, x, c, d);
  int ans = sz[c];
  root = merge(a, merge(c, d));
  return ans;
int size() { return sz[root]; }
// 0—indexed
treap_t nth(int idx) {
  int u = nth(root, idx + 1);
  return X[u];
// Query in k smallest elements
```

```
treap_t query(int k) {
   int a, b;
split_sz(root, k, a, b);
treap_t ans = op_val[a];
   root = merge(a, b);
   return ans;
};
    // namespace Treap
8.48 union_find_persistent
#include <bits/stdc++.h>
using namespace std;
namespace UnionFind
const int MAXN = 200010;
int n, p[MAXN], sz[MAXN], ti[MAXN], T;
void build(int n0) {
   T = -1, n = n0;
for (int i = 0; i < n; i++) {
     p[i] = i;
      sz[\bar{i}] = \bar{1};

ti[i] = -1;
   }
int find(int k, int t) {
  if (p[k] == k or ti[k] > t) return k;
   return find(p[k], t);
bool join(int a, int b, int t) {
   assert(T <= t);
   a = find(a, t);
  b = find(b, t);
if (a == b) return false;
if (sz[a] > sz[b]) swap(a, b);
   sz[b] += sz[a];
   p[a] = b;
   \underline{t}i[a] = t;
   T = t;
   return true;
bool isSame(int a, int b, int t) {
   return find(a, t) == find(b, t);
    // namespace UnionFind
8.49 union find with rollback
#include <bits/stdc++.h>
using namespace std;
struct RollbackUF {
   vector<int> e;
   vector<tuple<int, int, int, int>> st;
RollbackUF(int n) : e(n, -1) {}
int size(int x) { return -e[find(x)]; }
   int size(int x) {
   int find(int x)
      return e[x] < 0 ? x : find(e[x]);
   int time() { return st.size(); }
   void rollback(int t)
      while (st.size() > t)
        auto [a1, v1, a2, v2] = st.back();
e[a1] = v1;
e[a2] = v2;
        st.pop_back();
      }
   bool unite(int a, int b) {
    a = find(a), b = find(b);
    if (a == b) return false;
    if (e[a] > e[b]) swap(a, b);
    ct push back(b) scale b eff
      st.push_back({a, e[a], b, e[b]});
     e[a] += e[b];
e[b] = a;
      return true;
};
8.50 vertex cover in tree
```

```
#include <bits/stdc++.h>
using namespace std;
const int MAXN = 200010;
                                                                                     // Amount of numbers in the range [i, j] Less than
                                                                                     // or equal to k 1—indexed
                                                                                    int dp[MAXN][2];
vector<int> adj[MAXN];
// vertexCover(node current, free to choose, dad) int vertexCover(int u, bool color = true, int p = -1) {
   if (dp[u][color] != -1) return dp[u][color];
   int case1 = 1, case2 = 0;
for (int to : adj[u]) {
                                                                                                   lte(i2, j2, k, r[u]);
      if (to == p) continue;
                                                                                     // Amount of numbers in the range [i, j] equal to
      case1 += vertexCover(to, true, u);
case2 += vertexCover(to, false, u);
                                                                                     // k 1—indexed
                                                                                     int count(int i, int j, int k, int u = 1) {
  if (i > j or k < lo[u] or k > hi[u]) return 0;
  if (lo[u] == hi[u]) return j - i + 1;
  t_wavelet mid = (lo[u] + hi[u] - 1) / 2;
  int i1 = a[u][i - 1] + 1, j1 = a[u][j];
  int i2 = b(u, i - 1) + 1, j2 = b(u, j);
  if (k = mid) return count(i1 i1, k = 1[u]);
}
   if (color)
      return dp[u][color] = min(case1, case2);
   else
      return dp[u][color] = case1;
                                                                                        if (k <= mid) return count(i1, j1, k, l[u]);</pre>
                                                                                        return count(i2, j2, k, r[u]);
8.51 wavelet tree
                                                                                     // swap v[i] with v[i+1]
                                                                                     // 1—indexed
                                                                                     // l-indexed
void swp(int i, int u = 1) {
   if (lo[u] == hi[u] or a[u].size() <= 2) return;
   if (a[u][i - 1] + 1 == a[u][i] and
       a[u][i] + 1 == a[u][i + 1])
       swp(a[u][i], l[u]);
   else if (b(u, i - 1) + 1 == b(u, i) and
            b(u, i) + 1 == b(u, i + 1))
       swp(b(u, i), r[u]);
   else if (a[u][i - 1] + 1 == a[u][i])
   a[u][i]...</pre>
#include <bits/stdc++.h>
using namespace std;
namespace WaveletTree {
const int MAXN = 100010,
               MAXW =
                     MAXN * 30; // MAXN * LOG(maxX—MinX)
typedef int t_wavelet;
int last;
int v[MAXN], aux[MAXN];
int lo[MAXW], hi[MAXW], l[MAXW], r[MAXW];
vector<t_wavelet> a[MAXW];
                                                                                           a[u][i]-
                                                                                        else
                                                                                           a[u][i]++;
int stable_partition(int i, int j
                                 t_wavelet mid) {
                                                                                          // namespace WaveletTree
   int pivot = 0;
   for (int k = i; k < j; k++)
aux[k] = v[k], pivot += (v[k] <= mid);
                                                                                     8.52 xor_and_or_convolution
   int i1 = i, i2 = i + pivot;
   for (int k = i; k < j; k++) {
      if (aux[k] \le mid)
                                                                                     #include <bits/stdc++.h>
         v[i1++] = aux[k];
                                                                                     using namespace std;
      else
                                                                                     typedef long long li;
void xorFWHT(vector<ll> &P,
         v[i2++] = aux[k];
                                                                                                         bool inverse = false) {
                                                                                        int n = P.size();
for (int len = 1; 2 * len <= n; len <<= 1) {
   for (int i = 0; i < n; i += 2 * len) {
      for (int j = 0; j < len; j++) {
        ll u = P[i + j];
        ll v = P[i + len + j];
        pri + il - u + v.</pre>
   return i1;
lo[u] = minX, hi[u] = maxX;
if (lo[u] == hi[u] or i >= j) return;
t_wavelet mid = (minX + maxX - 1) / 2;
a[u].resize(j - i + 1);
                                                                                                  P[i + j] = u + v;
                                                                                                 P[i + len + j] = u - v;
   a[u][0] = 0;
   for (int k = i; k < j; k++)
a[u][k - i + 1] = a[u][k - i] + (v[k] \leftarrow mid);
                                                                                           }
  int pivot = stable_partition(i, j, mid);
l[u] = last++, r[u] = last++;
build(l[u], i, pivot, minX, mid);
build(r[u], pivot, j, mid + 1, maxX);
                                                                                        if (inverse) {
                                                                                           for (int i = 0; i < n; i++) {
                                                                                              P[i] /= n;
inline int b(int u, int i) { return i - a[u][i]; }
                                                                                     void orFWHT(vector<ll> &P, bool inverse = false) {
template <class MyIterator>
                                                                                        void init(MyIterator begin, MyIterator end,
               t_wavelet minX, t_wavelet maxX) {
   last = 1;
   int n = end - begin;
                                                                                                 if (inverse)
   for (int i = 0; i < n; i++, begin++)
                                                                                                    P[i + len + j] = P[i + j];
     v[i] = *begin;
   build(last++, 0, n, minX, maxX);
                                                                                                     P[i + len + j] += P[i + j];
                                                                                              }
// kth smallest element in range [i, j]
                                                                                           }
// 1—indexed
                                                                                        }
int kth(int i, int j, int k, int u = 1) {
  if (i > j) return 0;
if (lo[u] == hi[u]) return lo[u];
int inLeft = a[u][j] - a[u][i - 1];
                                                                                     void andFWHT(yector<ll> &P,
                                                                                                         bool inverse = false) {
                                                                                        int n = P.size();
for (int len = 1; 2 * len <= n; len <<= 1) {
  for (int i = 0; i < n; i += 2 * len) {
    for (int j = 0; j < len; j++) {
        ll u = P[i + j];
    }
}</pre>
  int i1 = a[u][i - 1] + 1, j1 = a[u][j];
int i2 = b(u, i - 1) + 1, j2 = b(u, j);
if (k <= interf) return kth(i1, j1, k, l[u]);
   return kth(i2, j2, k - inLeft, r[u]);
```

```
ll v = P[i + len + j];
if (inverse) {
            P[i + j] = v - u;
P[i + len + j] = u;
          } else {
            P[i + j] = v;
P[i + len + j] = u + v;
      }
    }
  }
int mx = max(a.size(), b.size());
  int n = 1;
  while (n < mx) n <<= 1;
  a.resize(n, 0);
b.resize(n, 0);
  xorFWHT(a);
xorFWHT(b);
  for (int i = 0; i < n; i++) a[i] *= b[i];
  xorFWHT(a, true);
  return a:
8.53 xor_trie
#include <bits/stdc++.h>
using namespace std;
using ll = long long;
struct Vertex {
  int next[2];
  int leaf;
int count;
  Vertex() {
  next[0] = next[1] = -1;
  leaf = count = 0;
const int MAXB = 20;
struct Trie {
  vector<Vertex> trie;
  ll lazy;
  Trie()
     trie.emplace_back();
     lazy = 0;
  void add(ll x) {
     int v = 0;
     trie[v].count++;
     for (int i = MAXB; i >= 0; i—) {
  int c = (x >> i) & 1;
  if (trie[v].next[c] == -1) {
    trie[v].next[c] = trie.size();
}
          trie.emplace_back();
       v = trie[v].next[c];
       trie[v].count++;
     trie[v].leaf++;
  void apply(ll x) { lazy ^= x; }
  ll min() {
     int \dot{v} = 0
     ll ans = 0;
for (int i = MAXB; i >= 0; i—) {
  int b = (lazy >> i) & 1;
       int to1 = trie[v].next[b];
int to2 = trie[v].next[b ^ 1];
       if (to1 != -1)^{-}{
       v = to1;
} else if (to2 != -1) {
          ans |= (1LL << i);
          v = to2;
       } else {
          return -1;
     return ans;
  ĺl max() {
     int v = 0
     ll ans = 0;
```

```
for (int i = MAXB; i >= 0; i—) { int b = (lazy >> i) & 1;
         int to1 = trie[v].next[b];
int to2 = trie[v].next[b ^ 1];
if (to2 != -1) {
            ans |= (1LL << i);
         v = to2;
} else if (to1 != -1) {
            v = to1;
         } else {
            return -1;
         }
      return ans;
   int countLE(ll x) {
  int v = 0, ans = 0;
  for (int i = MAXB; i >= 0; i—) {
    int c = (x >> i) & 1;
    int b = (lazy >> i) & 1;
}
         if (c == 0)
            if (trie[v].next[c ^ b] == -1) return ans;
v = trie[v].next[c ^ b];
             int to = trie[v].next[c ^ b ^ 1];
             if (to !=-1) ans += trie[to].count;
            if (trie[v].next[c ^ b] == -1) return ans; v = trie[v].next[c ^ b];
      }
      ans += trie[v].leaf;
      return ans;
};
8.54 z_function
```

```
#include <bits/stdc++.h>
using namespace std;
// z[i] is the length of the longest common prefix
// between s[0..(n-1)] and the suffix of
// s[i..(n-1)]. z[0] is generally not well
// defined. "aaabaab" - [0,2,1,0,2,1,0] "abacaba"
// - [0,0,1,0,3,0,1]
vector<int> z_function(string s) {
  int n = (int)s.length();
  vector<int> z(n);
  for (int i = 1, l = 0, r = 0; i < n; i++) {
    if (i <= r) z[i] = min(r - i + 1, z[i - l]);
    while (i + z[i] < n && s[z[i]] == s[i + z[i]])
    z[i]++;
  if (i + z[i] - 1 > r) l = i, r = i + z[i] - 1;
  }
  return z;
}
```

9 Primitives

9.1 Bigint

```
stringstream ss;
  ss << *this;
  string s;
ss >> s;
  return s;
int sumof() {
  string s = to_string();
  int ans = 0;
  for (auto c: s) ans += c - '0';
  return ans;
/*</arpa>*/
bigint() : sign(1) {}
bigint(long long v) { *this = v; }
bigint(const string &s) { read(s); }
void operator=(const bigint &v) {
  sign = v.sign;
      v.a;
}
void operator=(long long v) {
  sign = 1;
  a.clear()
  if (v < 0) sign = -1, v = -v;
for (; v > 0; v = v / base)
  a.push_back(v % base);
bigint operator+(const bigint &v) const {
  if (sign == v.sign) {
    bigint res = v;
    for (int i = 0, carry = 0;
          i < (int)max(a.size(), v.a.size()) ||
          carry;
       ++i) {
if (i == (int)res.a.size())
         res.a.push_back(0);
       res.a[i] += carry +
            (i < (int)a.size() ? a[i] : 0);
       carry = res.a[i] >= base;
       if (carry) res.a[i] -= base;
    return res;
  return *this - (-v);
}
bigint operator—(const bigint &v) const {
  if (sign == v.sign)
    if (abs() >= v.abs()) {
       bigint res = *this;
       for (int i = 0, carry = 0;
             i < (int)v.a.size() || carry; ++i) {
         res.a[i] -=
              carry +
         (i < (int)v.a.size() ? v.a[i] : 0);
carry = res.a[i] < 0;
         if (carry) res.a[i] += base;
       res.trim();
       return res;
     return -(v - *this);
  return *this + (-v);
}
void operator*=(int v) {
  if (v < 0) sign = -\text{sign},
                               v = -v;
  for (int i = 0, carry = 0;
        i < (int)a.size() || carry; ++i) {
    if (i == (int)a.size()) a.push_back(0);
    long long cur = a[i] * (long long)v + carry;
carry = (int)(cur / base);
    a[i] = (int)(cur % base);

// asm("divl %%ecx" : "=a"(carry),

// "=d"(a[i]) : "A"(cur), "c"(base));
  trim();
bigint operator*(int v) const {
  bigint res = *this;
res *= v;
  return rés;
```

```
void operator*=(long long v) {
  if (v < 0) sign = -\text{sign}, v = -v;
  if (v > base) {
    return:
  for (int i = 0, carry = 0;
    i < (int)a.size() || carry; ++i) {</pre>
    if (i == (int)a.size()) a.push_back(0);
    long long cur = a[i] * (long long)v + carry;
    carry = (int)(cur / base);
    a[i] = (int)(cur % base);

// asm("divl %%ecx" : "=a"(carry),

// "=d"(a[i]) : "A"(cur), "c"(base));
  trim();
bigint operator*(long long v) const {
  bigint res = *this;
res *= v;
  return res;
friend pair<bigint, bigint> divmod(
    const bigint &a1, const bigint &b1) {
  int norm = base / (bl.a.back() + 1);
bigint a = al.abs() * norm;
  bigint b = b1.abs() * norm;
  bigint q, r;
  q.a.resize(a.a.size());
  for (int i = a.a.size() - 1; i >= 0; i—) {
    r *= base
    r += a.a[i];
    int s1 = r.a.size() <= b.a.size()</pre>
                    ? 0
                    : r.a[b.a.size()];
    int s2 = r.a.size() \ll b.a.size() - 1
                   ? 0
                      r.a[b.a.size() - 1]
    int d = ((long long)base * s1 + s2) /
             b.a.back();
    r = b * d;
    while (r < 0) r += b, —d; q.a[i] = d;
  q.sign = a1.sign * b1.sign;
  r.sign = a1.sign;
  q.trim();
  r.trim();
  return make_pair(q, r / norm);
bigint operator/(const bigint &v) const {
  return divmod(*this, v).first;
bigint operator%(const bigint &v) const {
  return divmod(*this, v).second;
void operator/=(int v) {
  if (v < 0) sign = -\text{sign}, v = -v;
  for (int i = (int)a.size() - 1, rem = 0;
 i >= 0; -i) {
    long long cur =
         a[i] + rem * (long long)base;
    a[i] = (int)(cur / v);
    rem = (int)(cur % v);
  trim();
bigint operator/(int v) const {
  bigint res = *this;
  res /= v;
  return res;
int operator%(int v) const {
  if (v < 0) v = -v;
  int m = 0;
for (int i = a.size() -1; i >= 0; —i)
    m = (a[i] + m * (long long)base) % v;
  return m * sign;
```

```
void operator+=(const bigint &v) {
  *this = *this + v;
void operator—=(const bigint &v) {
  *this = *this - v;
void operator*=(const bigint &v) {
 *this = *this * v;
void operator/=(const bigint &v) {
  *this = *this / v;
bool operator<(const bigint &v) const {</pre>
  if (sign != v.sign) return sign < v.sign;
if (a.size() != v.a.size())</pre>
    return a.size() * sign <
            v.a.size() * v.sign;
  for (int i = a.size() - 1; i \ge 0; i—)
    if (a[i] != v.a[i])
      return a[i] * sign < v.a[i] * sign;</pre>
  return false;
}
bool operator>(const bigint &v) const {
  return v < *this;
bool operator<=(const bigint &v) const {
  return !(v < *this);
bool operator>=(const bigint &v) const {
  return !(*this < v);
bool operator==(const bigint &v) const {
 return !(*this < v) && !(v < *this);
bool operator!=(const bigint &v) const {
  return *this < v || v < *this;
void trim() {
 while (!a.empty() && !a.back()) a.pop_back();
  if (a.empty()) sign = 1;
bool isZero() const
  return a.empty() \mid \mid (a.size() == 1 && !a[0]);
bigint operator—() const {
  bigint res = *this;
  res.sign = -sign;
  return res;
}
bigint abs() const {
  bigint res = *this;
  res.sign *= res.sign;
  return res;
long longValue() const {
  long long res = 0;
  for (int i = a.size() -1; i >= 0; i—)
    res = res * base + a[i];
  return res * sign;
}
friend bigint gcd(const bigint &a,
                   const bigint &b) {
  return b.isZero() ? a : gcd(b, a % b);
friend bigint lcm(const bigint &a,
                   const bigint &b) {
  return a / gcd(a, b) * b;
}
void read(const string &s) {
  sign = 1;
  a.clear();
int pos = 0;
  while (pos < (int)s.size() &&
	(s[pos] == '-' || s[pos] == '+')) {
    if (s[pos] == '-') sign = -sign;
    ++pos;
  for (int i = s.size() - 1; i \ge pos;
       i —= base_digits) {
    int x = 0;
    for (int j = max(pos, i - base\_digits + 1);
```

```
j \le i; j++)

x = x * 10 + s[j] - '0';
    a.push_back(x);
  trim();
}
friend istream &operator>>(istream &stream,
                                bigint &v) {
  string s;
  stream >> s;
  v.read(s);
  return stream;
friend ostream & operator << (ostream & stream,
                                const bigint &v) {
  if (v.sign == -1) stream << '-';
  stream << (v.a.empty() ? 0 : v.a.back());</pre>
  for (int i = (int)v.a.size() - 2; i \ge 0;
    stream << setw(base_digits) << setfill('0')</pre>
             << v.a[i];
  return stream;
static vector<int> convert_base(
    const vector<int> &a, int old_digits,
     int new_digits) {
  vector<long long> p(
       max(old_digits, new_digits) + 1);
  p[0] = 1;
  for (int i = 1; i < (int)p.size(); i++)
p[i] = p[i - 1] * 10;</pre>
  vector<int> res;
  long long cur = 0;
  int cur_digits = 0;
  for (int i = 0; i < (int)a.size(); i++) {
    cur += a[i] * p[cur_digits];
    cur_digits += old_digits;
    while (cur_digits >= new_digits) {
       res.push_back(int(cur % p[new_digits]));
       cur /= p[new_digits];
       cur_digits —= new_digits;
  res.push_back((int)cur);
  while (!res.empty() && !res.back())
    res.pop_back();
  return res;
typedef vector<long long> vll;
static vll karatsubaMultiply(const vll &a,
                                   const vll &b) {
  int n = a.size();
  vll res(n + n);
  if (n <= 32) {
  for (int i = 0; i < n; i++)
    for (int j = 0; j < n; j++)
         res[i + j] += a[i] * b[j];
    return res;
  int k = n \gg 1;
  vll a1(a.begin(), a.begin() + k);
vll a2(a.begin() + k, a.end());
  vll b1(b.begin(), b.begin() + k);
  vll b2(b.begin() + k, b.end());
  vll a1b1 = karatsubaMultiply(a1, b1);
  vll a2b2 = karatsubaMultiply(a2, b2);
  for (int i = 0; i < k; i++) a2[i] += a1[i];
for (int i = 0; i < k; i++) b2[i] += b1[i];
  vll r = karatsubaMultiply(a2, b2);
  for (int i = 0; i < (int)alb1.size(); i++)
  r[i] -= alb1[i];</pre>
  for (int i = 0; i < (int)a2b2.size(); i++)
r[i] -= a2b2[i];</pre>
  for (int i = 0; i < (int)r.size(); i++)
  res[i + k] += r[i];
for (int i = 0; i < (int)alb1.size(); i</pre>
                        < (int)a1b1.size(); i++)
    res[i] += alb1[i];
  for (int i = 0; i < (int)a2b2.size(); i++)
     res[i + n] + = a2b2[i];
  return res;
bigint operator*(const bigint &v) const {
```

```
vector<int> a6 =
    convert_base(this->a, base_digits, 6);
     vector<int> b6 =
           convert_base(v.a, base_digits, 6);
     vll a(a6.begin(), a6.end());
     vll b(b6.begin(), b6.end());
while (a.size() < b.size()) a.push_back(0);</pre>
     while (b.size() < a.size()) b.push_back(0);</pre>
     while (a.size() & (a.size() - 1))
a.push_back(0), b.push_back(0);
     vll c = karatsubaMultiply(a, b);
     bigint res;
     res.sign = sign * v.sign;
for (int i = 0, carry = 0; i < (int)c.size();
            i++) {
        long long cur = c[i] + carry;
res.a.push_back((int)(cur % 1000000));
        carry = (int)(cur / 1000000);
     res.a = convert_base(res.a, 6, base_digits);
     res.trim();
     return res;
};
```

9.2 Integer Mod

```
const ll MOD = 1 '000' 000'000 + 7;
template <ll _mod = MOD>
struct mint {
  ll value;
 static const ll MOD_value = _mod;
 mint(ll v = 0) {
    value = v % _mod;
if (value < 0) value += _mod;</pre>
 mint(ll a, ll b) : value(0) {
    *this += a;
*this /= b;
 mint &operator+=(mint const &b) {
    value += b.value;
    if (value >= _mod) value -= _mod;
    return *this;
 mint &operator—=(mint const &b) {
    value == b.value;
    if (value < 0) value += _mod;
    return *this;
 mint &operator*=(mint const &b) {
    value = (ll)value * b.value % _mod;
    return *this;
  friend mint mexp(mint a, ll e) {
    mint res = 1;
    while (e) {
      if (e & 1) res *= a;
a *= a;
      e >>= 1;
    return res;
  friend mint inverse(mint a) {
    return mexp(a, \_mod - 2);
 mint &operator/=(mint const &b) {
    return *this *= inverse(b);
  friend mint operator+(mint a, mint const b) {
    return a += b;
 mint operator++(int) {
    return this->value = (this->value + 1) % _mod;
 mint operator++() {
    return this—>value = (this—>value + 1) % _mod;
  friend mint operator—(mint a, mint const b) {
    return a —= b;
  friend mint operator—(mint const a) {
    return 0 - a;
```

```
mint operator—(int) {
    return this->value =
                (this->value - 1 + mod) % mod;
  mint operator—() {
    return this->value =
               (this\rightarrowvalue -1 + \_mod) % \_mod;
  friend mint operator*(mint a, mint const b) {
    return a *= b;
  friend mint operator/(mint a, mint const b) {
    return a /= b;
  friend std::ostream &operator<<(
      std::ostream &os, mint const &a) {
    return os << a.value;
  friend bool operator==(mint const &a,
                          mint const &b) {
    return a.value == b.value;
  friend bool operator!=(mint const &a,
                          mint const &b) {
    return a.value != b.value;
};
```

9.3 Matrix

```
template <typename T>
struct Matrix {
  vector<vector<T>> d;
  Matrix() : Matrix(0) {}
  Matrix(int n) : Matrix(n, n) {}
Matrix(int n, int m)
        : Matrix(
                vector<vector<T>>(n, vector<T>(m))) {}
  Matrix(const vector<vector<T>> &v) : d(v) {}
  constexpr int n() const {
     return (int)d.size();
  constexpr int m() const {
     return n() ? (int)d[0].size() : 0;
  void rotate() { *this = rotated(); }
  Matrix<T> rotated() const {
     Matrix<T> res(m(), n());
for (int i = 0; i < m(); i++) {
   for (int j = 0; j < n(); j++) {
     res[i][j] = d[n() - j - 1][i];</pre>
     return res;
  Matrix<T> pow(int power) const {
     assert(n() == m());
     auto res = Matrix<T>::identity(n());
     auto b = *this;
     while (power)
        if (power \& 1) res *= b;
        b *= b;
       power >>= 1;
     return res;
  Matrix<T> submatrix(int start_i, int start_j,
                             int rows = INT_MA>
                             int cols = INT_MAX) const {
     rows = min(rows, n() - start_i);
cols = min(cols, m() - start_j);
     if (rows <= 0 or cols <= 0) return {};
     Matrix<T> res(rows, cols);
for (int i = 0; i < rows; i++)
  for (int j = 0; j < cols; j++)
    res[i][j] = d[i + start_i][j + start_j];</pre>
     return res;
  Matrix<T> translated(int x, int y) const {
```

```
Matrix<T> res(n(), m());
for (int i = 0; i < n(); i++) {
  for (int j = 0; j < m(); j++) {
    if (i + x < 0 or i + x >= n() or
        j + y < 0 or j + y >= m())
       continue;
res[i + x][j + y] = d[i][j];
    }
  return res;
static Matrix<T> identity(int n) {
  Matrix<T> res(n);
  for (int i = 0; i < n; i++) res[i][i] = 1;
  return res;
vector<T> &operator[](int i) { return d[i]; }
const vector<T> &operator[](int i) const {
  return d[i]:
Matrix<T> &operator+=(T value) {
  for (auto &row : d) {
     for (auto &x : row) x += value;
  return *this;
Matrix<T> operator+(T value) const {
  auto res = *this;
for (auto &row : res)
    for (auto \&x : row) x = x + value;
  return res:
Matrix<T> &operator—=(T value) {
  for (auto &row : d)
     for (auto \&x : row) x = value;
  return *this;
Matrix<T> operator—(T value) const {
  auto res = *this;
  for (auto &row : res) {
  for (auto &x : row) x = x - value;
  return res;
Matrix<T> &operator*=(T value) {
  for (auto &row : d)
    for (auto &x : row) x *= value;
  return *this;
Matrix<T> operator*(T value) const {
  auto res = *this;
for (auto &row : res) {
  for (auto &x : row) x = x * value;
  return res;
Matrix<T> & operator/=(T value) {
  for (auto &row : d)
    for (auto &x : row) x /= value;
  return *this;
Matrix<T> operator/(T value) const {
  auto res = *this;
  for (auto &row : res) {
  for (auto &x : row) x = x / value;
  return res;
Matrix<T> &operator+=(const Matrix<T> &o) {
  assert(n() == o.n() and m() == o.m());
  for (int i = 0; i < n(); i++) {
  for (int j = 0; j < m(); j++) {
      d[i][j] += o[i][j];
  return *this;
Matrix<T> operator+(const Matrix<T> &o) const {
  assert(n() == o.n() and m() == o.m());
  auto res = *this;
for (int i = 0; i < n(); i++) {
```

```
for (int j = 0; j < m(); j++) {
  res[i][j] = res[i][j] + o[i][j];</pre>
      return res;
   Matrix<T> &operator—=(const Matrix<T> &o) {
      assert(n() == 0.n() and m() == 0.m());
for (int i = 0; i < n(); i++) {
  for (int j = 0; j < m(); j++) {
    d[i][j] == 0[i][j];
      return *this;
   Matrix<T> operator—(const Matrix<T> &o) const {
      assert(n() == o.n() and m() == o.m());
      auto res = *this;
for (int i = 0; i < n(); i++) {
  for (int j = 0; j < m(); j++)
            res[i][j] = res[i][j] - o[i][j];
      return res;
   Matrix<T> &operator*=(const Matrix<T> &o) {
      *this = *this * o;
      return *this;
   Matrix<T> operator*(const Matrix<T> &o) const {
      assert(m() == o.n());
      Matrix<T> res(n(), o.m());
for (int i = 0; i < res.n(); i++) {
  for (int j = 0; j < res.m(); j++) {
    auto &x = res[i][j];
    for (int j = res[i][j];</pre>
            for (int k = 0; k < m(); k++) {
    x += (d[i][k] * o[k][j]);
        }
      return res;
   friend istream &operator>>(istream &is,
                                             Matrix<T> &mat) {
      for (auto &row: mat)
         for (auto &x : row) is >> x;
      return is;
   friend ostream &operator<<(
         ostream &os, const Matrix<T> &mat) {
      bool frow = 1;
      for (auto &row : mat) {
  if (not frow) os << '\n';</pre>
         bool first = 1;
         for (auto &x : row) {
   if (not first) os << ' ';
   os << x;</pre>
            first = 0;
         frow = 0;
      return os;
  auto begin() { return d.begin(); }
auto end() { return d.end(); }
auto rbegin() { return d.rbegin(); }
auto rend() { return d.rend(); }
   auto begin() const { return d.begin(); }
   auto end() const { return d.end(); }
auto rbegin() const { return d.rbegin(); }
   auto rend() const { return d.rend(); }
};
10 Strings
```

10.1 Count distinct anagrams

```
const ll MOD = 1e9 + 7;
const int maxn = 1e6;
vll fs(maxn + 1);
void precompute()
  fs[0] = 1;
for (ll i = 1; i <= maxn; i++) {
```

```
fs[i] = (fs[i-1] * i) % MOD;
}
ll fpow(ll a, int n, ll mod = LLONG_MAX) {
  if (n == 0) return 1;
if (n == 1) return a;
  ll x = fpow(a, n / 2, mod) % mod;
  return ((x * x) % mod * (n & 1 ? a : 1ll)) %
           mod:
ll distinctAnagrams(const string &s) {
  precompute();
vi hist('z' - 'a' + 1, 0);
for (auto &c : s) hist[c - 'a']++;
  ll ans = fs[len(s)];
for (auto &q : hist) {
    ans = (ans * fpow(fs[q], MOD - 2, MOD)) % MOD;
  return ans;
```

10.2 Double hash range query

```
using ll = long long;
using vll = vector<ll>
using pll = pair<ll, ll>;
const int MAXN(1 '000' 000);
const ll MOD = 1000027957
       11 \text{ MOD2} = 1000015187;
const ll P = 31;
ll p[MAXN + 1], p2[MAXN + 1];
void precompute() {
  p[0] = p2[0] = 1;
for (int i = 1; i <= MAXN; i++)
     p[i] = (P * p[i - 1]) % MOD,
     p2[i] = (P * p2[i - 1]) % MOD2;
struct Hash {
  int n;
vll h, h2, hi, hi2;
Hash() {}
  Hash(const string &s)
     in(const string ds)
: n(s.size()), h(n), h2(n), hi(n), hi2(n) {
h[0] = h2[0] = s[0];
for (int i = 1; i < n; i++)
   h[i] = (s[i] + h[i - 1] * P) % MOD,
   h2[i] = (s[i] + h2[i - 1] * P) % MOD2;</pre>
     \begin{array}{lll} \text{hi}[\mathsf{n}-1] = \text{hi}2[\mathsf{n}-1] = \text{s}[\mathsf{n}-1];\\ \text{for (int i = n-2; i >= 0; i--)}\\ \text{hi}[\underline{i}] = (\text{s}[\underline{i}] + \text{hi}[\underline{i}+1] * \text{P}) \% \text{ MOD} \end{array}
        hi2[i] = (s[i] + hi2[i + 1] * P) % MOD2;
  pll query(int l, int r) {
     ll\ hash =
             (\bar{l} ? h[l-1] * p[r-l+1] % MOD : 0));
     ll\ hash2 =
           (h2[r]
             (l ? h2[l - 1] * p2[r - l + 1] % MOD2
                   0));
     pll query_inv(int l, int r) {
     ll\ hash = (hi[l]
                      (r+1 < n ? hi[r+1] *
                                                      l + 1] % MOD
                                              p[r -
                                      : 0));
     ll hash2
           (hi2[l] -
             (r + 1 < n)
                      hi2[r + 1] * p2[r - l + 1] % MOD2
     : 0));
return {(hash < 0 ? hash + MOD : hash)
                 \{hash2 < 0 ? hash2 + MOD2 : hash2\};
};
```

10.3 Hash range query

```
const ll P = 31;
const ll MOD = 1e9 + 9;
const int MAXN(1e6);
ll ppow[MAXN + 1];
void pre_calc() {
  ppow[0] = 1;
  for (int i = 1; i <= MAXN; i++)
     ppow[i] = (ppow[i - 1] * P) % MOD;
struct Hash {
  int n;
vll h, hi;
  Hash(const string &s)
        : n(s.size()), h(n), hi(n) {
     h[0] = s[0];

hi[n - 1] = s[n - 1];

for (int i = 1; i < n; i++) {

h[i] = (s[i] + h[i - 1] * P) % MOD;

hi[n - i - 1] =
             (s[\bar{n} - \bar{i} - 1] + hi[n - i - 1] * P) %
            MOD:
     }
  }
  ll qry(int l, int r) {
     ll\ hash =
          (h[r] —
  (l ? h[l — 1] * ppow[r — l + 1] % MOD
               : 0));
     return hash < 0 ? hash + MOD : hash;
  ll gry_inv(int l, int r) {
     ll hash =
          (hi[l]
           (r +
                  ? hi[r + 1] * ppow[r - l + 1] % MOD
     : 0));
return hash < 0 ? hash + MOD : hash;
};
10.4 Hash unsigned long long 2^{64} - 1
```

Description: Arithmetic mod $2^{64} - 1$. 2x slower than mod 2^{64} and more code, but works on evil test data (e.g. Thue-Morse, where ABBA... and BAAB... of length 2^{10} hash the same mod 2^{64}).

```
"typedef ull H;" instead if you think test data is random
typedef uint64_t ull;
struct H {
  H(ull'x = 0) : x(x) {}
  H operator+(H o) { return x + o.x + (x + o.x < x);
  H operator-(H o) { return *this + ~o.x; }
H operator*(H o) {
     auto m = (\_uint128\_t)x * o.x;
     return H((ull)m) + (ull)(m >> 64);
  ull get() const { return x + !~x; }
bool operator==(H o) const {
     return get() == o.get();
  bool operator<(H o) const {
    return get() < o.get();
static const H C =
     (long long)le11 +
3; // (order ~ 3e9; random also ok)
struct Hash {
  int n:
  vector<H> ha, pw;
  Hash(string &str)
       : n(str.size())
         ha((int)str.size() + 1),
         pw(ha) {
     pw[0] = 1;
     for (int i = 0; i < (int)str.size(); i++)
  ha[i + 1] = ha[i] * C + str[i],</pre>
               pw[i + 1] = pw[i] * C;
```

10.5 K-th digit in digit string

Description: Find the k-th digit in a digit string, only works for $1 <= k <= 10^{18}$! **Time**: precompute O(1), query O(1)

```
using vull = vector<ull>;
vull pow10;
vector<array<ull, 4>> memo;
void precompute(int maxpow = 18) {
  ull qtd = 1;
ull start = 1;
  ull end = 9;
  ull curlenght = 9;
  ull startstr = 1;
  ull endstr = 9;
  for (ull i = 0, j = 1ll; (int)i < maxpow;
        i++, j *= 10ll)
    pow10.eb(j);
  for (ull i = 0; i < maxpow - 1ull; i++) {
    memo.push_back(
         {start, end, startstr, endstr});
    start = end + 1ll;
    end = end + (9ll * pow10[qtd]);
    curlenght = end - start + 1ull;
    qtd++;
    startstr = endstr + 1ull;
    endstr =
         (endstr + 1ull) + (curlenght)*gtd — 1ull;
  }
char kthDigit(ull k) {
  int qtd = 1;
  for (auto [s, e, ss, es] : memo) {
  if (k >= ss and k <= es) {
   ull pos = k - ss;
}</pre>
       ull index = pos / qtd;
      ull nmr = s + index;
int i = k - ss - qtd * index;
       return ((nmr / pow10[qtd - i - 1]) % 10) + '0';
    qtd++;
  }
  return 'X';
```

10.6 Longest Palindrome Substring (Manacher)

Description: Finds the longest palindrome substring, manacher returns a vector where the i-th position is how much is possible to grow the string to the left and the right of i and keep it a palindrome. Time: O(N)

```
vi manacher(const string &s) {
  int n = len(s) - 2;
  vi p(n + 2);
  int l = 1, r = 1;
  for (int i = 1; i <= n; i++) {
    p[i] = max(0, min(r - i, p[l + (r - i)]));
    while (s[i - p[i]] == s[i + p[i]]) p[i]++;
    if (i + p[i] > r) l = i - p[i], r = i + p[i];
```

```
p[i]—;
}
return p;
}
string longest_palindrome(const string &s) {
    string t("$#");
    for (auto c : s)
        t.push_back(c), t.push_back('#');
    t.push_back('^');
    vi xs = manacher(t);
    int mpos = max_element(all(xs)) - xs.begin();
    string p;
    for (int k = xs[mpos], i = mpos - k;
        i <= mpos + k; i++)
        if (t[i] != '#') p.push_back(t[i]);
    return p;
}</pre>
```

10.7 Longest palindrome

```
string longest_palindrome(const string &s) {
   int n = (int)s.size();
   vector<array<int, 2>> dp(n);
   pii odd(0, -1), even(0, -1);
   pii ans;
   for (int i = 0; i < n; i++) {
     int k = 0;
      if (i > odd.second)
        k = 1;
      else
        k = min(dp[odd.first + odd.second - i][0],
     \begin{array}{c} \text{odd.second} - \text{ i + 1);} \\ \text{while } (\text{i} - \text{k} >= 0 \text{ and i + k} < \text{n and} \\ \text{s[i - k]} == \text{s[i + k])} \end{array}
     dp[i][0] = k
     if (i + k > odd.second) odd = \{i - k, i + k\}; if (2 * dp[i][0] - 1 > ans.second) ans = \{i - k, 2 * dp[i][0] - 1\};
      k = 0;
      if (i <= even.second)</pre>
        k = min(
              dp[even.first + even.second - i + 1][1],
     even.second - i + 1); while (i - k - 1 >= 0 and i + k < n and
                s[i - k - 1] == s[i + k]
     dp[i][1] = k--;
     if (i + k > even.second)
even = \{i - k - 1, i + k\};
      if (2 * dp[i][1] > ans.second)
        ans = \{i - k - 1, 2 * dp[i][1]\};
   return s.substr(ans.first, ans.second);
```

10.8 Rabin-Karp

```
(hp.second + (p[i] - 'a' + 1)) % q2;
  }
  size_t occ = 0;
   for (size_t i = 0; i < n - m; i++) {
     occ += (hs == hp);
     int fi = s[i] - 'a' + 1;
int fm = s[i + m] - 'a' + 1;
     hs.first = (hs.first - fi + q1) % q1;
     hs.first = (hs.first * p1_1) % q1;
hs.first = (hs.first + fm * p1_2) % q1;
     hs.second = (hs.second - fi + q2) % q2;
     hs.second = (hs.second * p2_1) % q2;
hs.second = (hs.second + fm * p2_2) % q2;
  occ += hs == hp;
  return occ;
10.9 Suffix array
vector<int> sort_cyclic_shifts(string const &s) {
  int n = s.size();
  const int alphabet = 128;
  vector<int> p(n), c(n),
        cnt(max(alphabet, n), 0);
  for (int i = 0; i < n; i++) cnt[s[i]]++;
for (int i = 1; i < alphabet; i++)
  cnt[i] += cnt[i - 1];</pre>
  for (int i = 0; i < n; i++) p[-cnt[s[i]]] = i;
  c[p[0]] = 0;
   int classes = 1;
  for (int i = 1; i < n; i++) {
  if (s[p[i]] != s[p[i - 1]]) classes++;</pre>
     c[p[i]] = classes - 1;
  vector<int> pn(n), cn(n);
  for (int h = 0; (1 << h) < n; ++h) {
  for (int i = 0; i < n; i++) {
    pn[i] = p[i] - (1 << h);
}
        if (pn[i] < 0) pn[i] += n;
     fill(cnt.begin(), cnt.begin() + classes, 0);
     for (int i = 0; i < n; i++) cnt[c[pn[i]]]++;
for (int i = 1; i < classes; i++)
  cnt[i] += cnt[i - 1];
for (int i = n - 1; i >= 0; i—)
  p[—cnt[c[pn[i]]]] = pn[i];
     cn[p[0]] = 0;
     classes = 1;
for (int i = 1; i < n; i++) {</pre>
        pair<int, int> cur = {
             c[p[i]], c[(p[i] + (1 << h)) % n]};
        pair<int, int> prev = {
             c[p[i-1]],
             c[(p[i-1] + (1 << h)) % n]);
        if (cur != prev) ++classes;
        cn[p[i]] = classes - 1;
     c.swap(cn);
  }
  return p;
}
vector<int> suffix_array(string s) {
  s += "$";
  vector<int> p = sort_cyclic_shifts(s);
  p.erase(p.begin());
  return p;
10.10 Suffix automaton
struct state {
  int len, link, cnt, firstpos;
// this can be optimized using a vector with the
```

// alphabet size map<char, int> next;

vi inv_link;

```
struct SuffixAutomaton {
  vector<state> st;
  int sz = 0;
int last;
  vc cloned;
  SuffixAutomaton(const string &s, int maxlen)
        st(maxlen * 2), cloned(maxlen * 2) {
    st[0].len = 0;
    st[0].link = -1;
sz++;
    last = 0;
    for (auto &c : s) add_char(c);
     // precompute for count occurences
    for (int i = 1; i < sz; i++) {
  st[i].cnt = !cloned[i];</pre>
    vector<pair<state, int>> aux;
    for (int i = 0; i < sz; i++) {
      aux.push_back({st[i], i});
    return a.fst.len > b.fst.len;
    });
    for (auto &[stt, id] : aux) { if (stt.link !=-1) {
         st[stt.link].cnt += st[id].cnt;
    }
     // for find every occurende position
    for (int v = 1; v < sz; v++)
      st[st[v].link].inv_link.push_back(v);
  void add_char(char c) {
    int cur = sz++;
st[cur].len = st[last].len + 1;
    st[cur].firstpos = st[cur].len - 1;
    int p = last;
    // follow the suffix link until find a // transition to c while (p != -1 and !st[p].next.count(c)) {
      st[p].next[c] = cur;
      p = st[p].link;
    /// there was no transition to c so create and // leave if (p == -1) {
      st[cur].link = 0;
       last = cur;
       return;
    int q = st[p].next[c];
    if (st[p].len + 1 == st[q].len) {
      st[cur].link = q;
    } else {
      int clone = sz++;
cloned[clone] = true;
       st[clone].len = st[p].len + 1;
       st[clone].next = st[q].next;
       st[clone].link = st[q].link;
       st[clone].firstpos = st[q].firstpos;
      while (p != -1 \text{ and } st[p].next[c] == q)  {
         st[p].next[c] = clone;
         p = st[p].link;
      st[q].link = st[cur].link = clone;
     last = cur;
  bool checkOccurrence(
      const string &t) { // O(len(t))
    int cur = 0;
for (auto &c : t) {
   if (!st[cur].next.count(c)) return false;
      cur = st[cur].next[c];
    return true;
  Il totalSubstrings() { // distinct, O(len(s))
    ll tot = 0;
for (int i = 1; i < sz; i++)
       tot += st[i].len - st[st[i].link].len;
```

```
return tot;
  // count occurences of a given string t
  int countOccurences(const string &t) {
    int cur = 0;
    for (auto &c : t) {
       if (!st[cur].next.count(c)) return 0;
       cur = st[cur].next[c];
    return st[cur].cnt;
  // find the first index where t appears a
  // substring O(len(t))
  int firstOccurence(const string &t) {
    int cur = 0;
    for (auto c : t) {
      if (!st[cur].next.count(c)) return -1;
cur = st[cur].next[c];
    return st[cur].firstpos - len(t) + 1;
  vi every0ccurence(const string &t) {
    int cur = 0;
for (auto c : t) {
  if (!st[cur].next.count(c)) return {};
       cur = st[cur].next[c];
    vi ans;
    getEveryOccurence(cur, len(t), ans);
    return ans;
  void getEveryOccurence(int v, int P_length,
                             vi &ans) {
    if (!cloned[v])
      ans.pb(st[v].firstpos - P_length + 1);
    for (int u : st[v].inv_link)
       getEveryOccurence(u, P_length, ans);
};
10.11 Trie
Description:
   • build with the size of the alphabet (sigma) and the first char (norm)
   • insert(s) insert the string in the trie O(|s| * sigma)
   • erase(s) remove the string from the trie O(|s|)
   • find(s) return the last node from the string s, 0 if not found O(|s|)
struct trie {
  vi2d to;
  vi end, pref;
  int sigma;
  char norm;
  trie(int sigma_ = 26, char norm_ = 'a')
    : sigma(sigma_), norm(norm_) {
to = {vector<int>(sigma)};
    end = \{0\}, pref = \{0\};
  int next(int node, char key) {
```

return to[node][key - norm];

```
void insert(const string &s) {
    int x = 0;
    for (auto c : s)
      int &nxt = to[x][c - norm];
      if (!nxt) {
        nxt = len(to);
        to.push_back(vi(sigma));
        end.emplace_back(0), pref.emplace_back(0);
      x = nxt, pref[x]++;
    end[x]++, pref[0]++;
  void erase(const string &s) {
    int x = 0;
for (char c : s)
      int &nxt = to[x][c - norm];
      x = nxt, pref[x]
      if (!pref[x]) nxt = 0;
    end[x]—, pref[0]—;
  int find(const string &s) {
    int x = 0; for (auto c
                 : s) {
      x = to[x][c - norm];
      if (!x) return 0;
    return x;
};
```

10.12 Z-function get occurence positions

Time: O(len(s) + len(p))

```
vi getOccPos(string &s, string &p) {
  // Z—function
char delim = '#'
  string t{p + delim + s};
  vi zs(len(t));
  zs[i]++;
    if (r < i + zs[i] - 1)
      l = i, r = i + zs[i] - 1;
  // Iterate over the results of Z—function to get
  // ranges
  vi ans
  int start = len(p) + 1 + 1 - 1;
for (int i = start; i < len(zs); i++) {
  if (zs[i] == len(p)) {</pre>
      int l = i - start
      ans.emplace_back(l);
  return ans;
}
```