Contents

1 Data structures

1.1 Segtree Lazy (Atcoder)

```cpp
struct Node {
  // need an empty constructor with the neutral node
  Node() : {}
};

struct Lazy {
  // need an empty constructor with the neutral lazy
  Lazy() : {}
};

// how to merge two nodes
Node op(Node a, Node b) {}

// how to apply the lazy into a node
Node mapping(Lazy a, Node b, int, int) {}

// how to merge two lazy
Lazy comp(Lazy a, Lazy b) {}

template <typename T, auto op, typename L, auto mapping,
          auto composition>
struct SegTreeLazy {
  static_assert(
    is_convertible_v<decltype(op), function<T(T, T)>>,
    "op must be a function T(T, T)");
  static_assert(
    is_convertible_v<decltype(mapping),
                     function<T(L, T, int, int)>>,
    "mapping must be a function T(L, T, int, int)");
  static_assert(is_convertible_v<decltype(composition),
                                 function<L(L, L)>>,
                "composition must be a function L(L, L)");

  int N, size, height;
  const T eT;
  const L eL;
  vector<T> d;
  vector<L> lz;

  SegTreeLazy(const T &eT_ = T(), const L &eL_ = L())
    : SegTreeLazy(0, eT_, eL_) {}
  explicit SegTreeLazy(int n, const T &eT_ = T(),
                       const L &eL_ = L())
    : SegTreeLazy(vector<T>(n, eT_), eT_, eL_) {}
  explicit SegTreeLazy(const vector<T> &v,
                       const T &eT_ = T(),
                       const L &eL_ = L())
    : N(int(v.size())), eT(eT_), eL(eL_) {
    size = 1;
    height = 0;
    while (size < N) size <<= 1, height++;
    d = vector<T>(2 * size, eT);
    lz = vector<L>(size, eL);
    for (int i = 0; i < N; i++) d[size + i] = v[i];
    for (int i = size - 1; i >= 1; i--) {
      update(i);
    }
  }

  void set(int p, T x) {
    assert(0 <= p && p < N);
    p += size;
    for (int i = height; i >= 1; i--) push(p >> i);
    d[p] = x;
    for (int i = 1; i <= height; i++) update(p >> i);
  }

  T get(int p) {
    assert(0 <= p && p < N);
    p += size;
    for (int i = height; i >= 1; i--) push(p >> i);
    return d[p];
  }

  T query(int l, int r) {
    assert(0 <= l && l <= r && r < N);

    l += size;
    r += size;

    for (int i = height; i >= 1; i--) {
      if (((l >> i) << i) != l) push(l >> i);
      if ((((r + 1) >> i) << i) != (r + 1)) push(r >> i);
    }
```

```
  T sml = eT, smr = eT;
  while (l <= r) {
    if (l & 1) sml = op(sml, d[l++]);
    if (!(r & 1)) smr = op(d[r--], smr);
    l >>= 1;
    r >>= 1;
  }

  return op(sml, smr);
}

T query_all() { return d[1]; }

void update(int p, L f) {
  assert(0 <= p && p < N);
  p += size;
  for (int i = height; i >= 1; i--) push(p >> i);
  d[p] = mapping(f, d[p]);
  for (int i = 1; i <= height; i++) update(p >> i);
}
void update(int l, int r, L f) {
  assert(0 <= l && l <= r && r < N);

  l += size;
  r += size;

  for (int i = height; i >= 1; i--) {
    if (((l >> i) << i) != l) push(l >> i);
    if ((((r + 1) >> i) << i) != (r + 1)) push(r >> i);
  }

  {
    int l2 = l, r2 = r;
    while (l <= r) {
      if (l & 1) all_apply(l++, f);
      if (!(r & 1)) all_apply(r--, f);
      l >>= 1;
      r >>= 1;
    }
    l = l2;
    r = r2;
  }

  for (int i = 1; i <= height; i++) {
    if (((l >> i) << i) != l) update(l >> i);
      if ((((r + 1) >> i) << i) != (r + 1)) update(r >> i);
    }
  }
  pair<int, int> node_range(int k) const {
    int remain = height;
    for (int kk = k; kk >>= 1; --remain)
      ;
    int fst = k << remain;
    int lst = min(fst + (1 << remain) - 1, size + N - 1);
    return {fst - size, lst - size};
  }

 private:
  void update(int k) { d[k] = op(d[2 * k], d[2 * k + 1]); }
  void all_apply(int k, L f) {
    auto [fst, lst] = node_range(k);
    d[k] = mapping(f, d[k], fst, lst);
    if (k < size) lz[k] = composition(f, lz[k]);
  }
  void push(int k) {
    all_apply(2 * k, lz[k]);
    all_apply(2 * k + 1, lz[k]);
    lz[k] = eL;
  }
};
```

1.2 Hash module $2^{61} - 1$

```
const ll MOD = (1ll << 61) - 1;
ll mulmod(ll a, ll b) {
  const static ll LOWER = (1ll << 30) - 1,
                  GET31 = (1ll << 31) - 1;
  ll l1 = a & LOWER, h1 = a >> 30, l2 = b & LOWER,
     h2 = b >> 30;
  ll m = l1 * h2 + l2 * h1, h = h1 * h2;
  ll ans = l1 * l2 + (h >> 1) + ((h & 1) << 60) +
           (m >> 31) + ((m & GET31) << 30) + 1;
  ans = (ans & MOD) + (ans >> 61),
  ans = (ans & MOD) + (ans >> 61);
  return ans - 1;
}

mt19937_64 rng(
  chrono::steady_clock::now().time_since_epoch().count());
```

```cpp
ll uniform(ll l, ll r) {
  uniform_int_distribution<ll> uid(l, r);
  return uid(rng);
}

struct str_hash {
  static ll P;
  vector<ll> h, p;
  str_hash(string s) : h(s.size()), p(s.size()) {
    p[0] = 1, h[0] = s[0];
    for (int i = 1; i < s.size(); i++)
      p[i] = mulmod(p[i - 1], P),
      h[i] = (mulmod(h[i - 1], P) + s[i]) % MOD;
  }
  ll operator()(int l, int r) {  // retorna hash s[l...r]
    ll hash =
      h[r] - (l ? mulmod(h[l - 1], p[r - l + 1]) : 0);
    return hash < 0 ? hash + MOD : hash;
  }
};

ll str_hash::P = uniform(256, MOD - 1);  // l > |sigma|
```

1.3 Bitree 2D

Given a 2d array allow you to sum val to the position (x, y) and find the sum of the rectangle with left top corner $(x1, y1)$ and right bottom corner $(x2, y2)$
Update and query 1 indexed !
Time: update $O(logn^2)$, query $O(logn^2)$

```cpp
struct Bit2d {
  int n;
  vll2d bit;
  Bit2d(int ni) : n(ni), bit(n + 1, vll(n + 1)) {}
  Bit2d(int ni, vll2d &xs) : n(ni), bit(n + 1, vll(n + 1)) {
    for (int i = 1; i <= n; i++) {
      for (int j = 1; j <= n; j++) {
        update(i, j, xs[i][j]);
      }
    }
  }
  void update(int x, int y, ll val) {
    for (; x <= n; x += (x & (-x))) {
      for (int i = y; i <= n; i += (i & (-i))) {
        bit[x][i] += val;
      }
    }
  }
  ll sum(int x, int y) {
    ll ans = 0;

    for (int i = x; i; i -= (i & (-i))) {
      for (int j = y; j; j -= (j & (-j))) {
        ans += bit[i][j];
      }
    }
    return ans;
  }
  ll query(int x1, int y1, int x2, int y2) {
    return sum(x2, y2) - sum(x2, y1 - 1) - sum(x1 - 1, y2) +
           sum(x1 - 1, y1 - 1);
  }
};
```

1.4 Bitree

```cpp
template <typename T>
struct BITree {
  int N;
  vector<T> v;

  BITree(int n) : N(n), v(n + 1, 0) {}

  void update(int i, const T& x) {
    if (i == 0) return;
    for (; i <= N; i += i & -i) v[i] += x;
  }

  T range_sum(int i, int j) {
    return range_sum(j) - range_sum(i - 1);
  }

  T range_sum(int i) {
    T sum = 0;
    for (; i > 0; i -= i & -i) sum += v[i];
    return sum;
  }
};
```

1.5 Convex Hull Trick / Line Container

Container where you can add lines of the form $mx + b$, and query maximum value at point x.

insert_line(m, b) inserts the line $m \cdot x + b$ in the container.
eval(x) find the highest value among all lines in the point x.
both in $O(\log N)$

```
const ll LLINF = 1e18;
const ll is_query = -LLINF;
struct Line {
  ll m, b;
  mutable function<const Line*()> succ;
  bool operator<(const Line& rhs) const {
    if (rhs.b != is_query) return m < rhs.m;
    const Line* s = succ();
    if (!s) return 0;
    ll x = rhs.m;
    return b - s->b < (s->m - m) * x;
  }
};
struct Cht : public multiset<Line> {  // maintain max m*x+b
  bool bad(iterator y) {
    auto z = next(y);
    if (y == begin()) {
      if (z == end()) return 0;
      return y->m == z->m && y->b <= z->b;
    }
    auto x = prev(y);
    if (z == end()) return y->m == x->m && y->b <= x->b;
    return (ld)(x->b - y->b) * (z->m - y->m) >=
           (ld)(y->b - z->b) * (y->m - x->m);
  }
  void insert_line(
    ll m, ll b) {  // min -> insert (-m,-b) -> -eval()
    auto y = insert({m, b});
    y->succ = [=] {
      return next(y) == end() ? 0 : &*next(y);
    };
    if (bad(y)) {
      erase(y);
      return;
    }
    while (next(y) != end() && bad(next(y))) erase(next(y));
    while (y != begin() && bad(prev(y))) erase(prev(y));
  }
  ll eval(ll x) {
    auto l = *lower_bound((Line){x, is_query});
    return l.m * x + l.b;
  }
};
```

1.6 Disjoint Sparse Table

Answers queries of any monoid operation (i.e. has identity element and is associative)
Build: $O(N \log N)$, Query: $O(1)$

```
#define F(expr) [](auto a, auto b) { return expr; }
template <typename T>
struct DisjointSparseTable {
  using Operation = T (*)(T, T);

  vector<vector<T>> st;
  Operation f;
  T identity;

  static constexpr int log2_floor(
    unsigned long long i) noexcept {
    return i ? __builtin_clzll(1) - __builtin_clzll(i) : -1;
  }

  // Lazy loading constructor. Needs to call build!
  DisjointSparseTable(Operation op, const T neutral = T())
    : st(), f(op), identity(neutral) {}

  DisjointSparseTable(vector<T> v)
    : DisjointSparseTable(v, F(min(a, b))) {}

  DisjointSparseTable(vector<T> v, Operation op,
                      const T neutral = T())
    : st(), f(op), identity(neutral) {
    build(v);
  }

  void build(vector<T> v) {
    st.resize(log2_floor(v.size()) + 1,
              vector<T>(1ll << (log2_floor(v.size()) + 1)));
    v.resize(st[0].size(), identity);
    for (int level = 0; level < (int)st.size(); ++level) {
      for (int block = 0; block < (1 << level); ++block) {
        const auto l = block << (st.size() - level);
        const auto r = (block + 1) << (st.size() - level);
        const auto m = l + (r - l) / 2;

        st[level][m] = v[m];
        for (int i = m + 1; i < r; i++)
```

```
          st[level][i] = f(st[level][i - 1], v[i]);
        st[level][m - 1] = v[m - 1];
        for (int i = m - 2; i >= l; i--)
          st[level][i] = f(st[level][i + 1], v[i]);
      }
    }
  }

  T query(int l, int r) const {
    if (l > r) return identity;
    if (l == r) return st.back()[l];

    const auto k = log2_floor(l ^ r);
    const auto level = (int)st.size() - 1 - k;
    return f(st[level][l], st[level][r]);
  }
};
```

1.7 Dsu

```
struct DSU {
  vi ps, sz;

  // vector<unordered_set<int>> sts;

  DSU(int N) : ps(N + 1), sz(N, 1) /*, sts(N) */ {
    iota(all(ps), 0);
    // for (int i = 0; i < N; i++) sts[i].insert(i);
  }
  int find_set(int x) {
    return ps[x] == x ? x : ps[x] = find_set(ps[x]);
  }
  int size(int u) { return sz[find_set(u)]; }
  bool same_set(int x, int y) {
    return find_set(x) == find_set(y);
  }
  void union_set(int x, int y) {
    if (same_set(x, y)) return;

    int px = find_set(x);
    int py = find_set(y);

    if (sz[px] < sz[py]) swap(px, py);

    ps[py] = px;
    sz[px] += sz[py];
    // sts[px].merge(sts[py]);
  }
};
```

1.8 Lichao Tree (dynamic)

Lichao Tree that creates the nodes dynamically, allowing to query and update from range $[MAXL, MAXR]$
$query(x)$: find the highest point among all lines in the structure
$add(a, b)$: add a line of form $y = ax + b$ in the structure
$addSegment(a, b, l, r)$: add a line segment of form $y = ax + b$ which covers from range $[l, r]$
time: $O(\log N)$

```
template <typename T = ll, T MAXL = 0,
          T MAXR = 1'000'000'001>
struct LiChaoTree {
  static const T inf = -numeric_limits<T>::max() / 2;
  bool first_best(T a, T b) { return a > b; }
  T get_best(T a, T b) { return first_best(a, b) ? a : b; }
  struct line {
    T m, b;
    T operator()(T x) { return m * x + b; }
  };
  struct node {
    line li;
    node *left, *right;
    node(line _li = {0, inf})
      : li(_li), left(nullptr), right(nullptr) {}
    ~node() {
      delete left;
      delete right;
    }
  };
  node *root;
  LiChaoTree(line li = {0, inf}) : root(new node(li)) {}
  ~LiChaoTree() { delete root; }
  T query(T x, node *cur, T l, T r) {
    if (cur == nullptr) return inf;
    if (x < l or x > r) return inf;
    T mid = midpoint(l, r);
    T ans = cur->li(x);
    ans = get_best(ans, query(x, cur->left, l, mid));
    ans = get_best(ans, query(x, cur->right, mid + 1, r));
    return ans;
  }
  T query(T x) { return query(x, root, MAXL, MAXR); }
  void add(line li, node *&cur, T l, T r) {
```

```
    if (cur == nullptr) {
      cur = new node(li);
      return;
    }
    T mid = midpoint(l, r);
    if (first_best(li(mid), cur->li(mid)))
      swap(li, cur->li);
    if (first_best(li(l), cur->li(l)))
      add(li, cur->left, l, mid);
    if (first_best(li(r), cur->li(r)))
      add(li, cur->right, mid + 1, r);
  }
  void add(T m, T b) { add({m, b}, root, MAXL, MAXR); }
  void addSegment(line li, node *&cur, T l, T r, T lseg,
                  T rseg) {
    if (r < lseg || l > rseg) return;
    if (cur == nullptr) cur = new node;
    if (lseg <= l && r <= rseg) {
      add(li, cur, l, r);
      return;
    }
    T mid = midpoint(l, r);
    if (l != r) {
      addSegment(li, cur->left, l, mid, lseg, rseg);
      addSegment(li, cur->right, mid + 1, r, lseg, rseg);
    }
  }
  void addSegment(T a, T b, T l, T r) {
    addSegment({a, b}, root, MAXL, MAXR, l, r);
  }
};
```

1.9 Merge Sort Tree

Like a segment tree but each node st_i stores a sorted subarray

- $inrange(l, r, a, b)$: counts the number of elements $x \in [l, r]$ such that $a \leq x \leq b$.

Memory: $O(n \log N)$ Time: build $O(N \log N)$, inrange $O(\log N)$

```
template <class T>
struct MergeSortTree {
  int n;
  vector<vector<T>> st;
  MergeSortTree(vector<T> &xs) : n(len(xs)), st(n << 1) {
    for (int i = 0; i < n; i++)
      st[i + n] = vector<T>({xs[i]});
    for (int i = n - 1; i > 0; i--) {
      st[i].resize(len(st[i << 1]) + len(st[i << 1 | 1]));
      merge(all(st[i << 1]), all(st[i << 1 | 1]),
            st[i].begin());
    }
  }
  int count(int i, T a, T b) {
    return upper_bound(all(st[i]), b) -
           lower_bound(all(st[i]), a);
  }
  int inrange(int l, int r, T a, T b) {
    int ans = 0;

    for (l += n, r += n + 1; l < r; l >>= 1, r >>= 1) {
      if (l & 1) ans += count(l++, a, b);
      if (r & 1) ans += count(--r, a, b);
    }

    return ans;
  }
};
```

1.10 Ordered Set

If you need an ordered **multi**set you may add an id to each value. Using greater_equal, or less_equal is considered undefined behavior.

- **order_of_key (k)** : Number of items strictly smaller/greater than k .
- **find_by_order(k)** : K-th element in a set (counting from zero).

```
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;

template <typename T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
                         tree_order_statistics_node_update>;
```

1.11 Prefix Sum 2D

Given an 2d array with n lines and m columns, find the sum of the subarray that have the left upper corner at $(x1, y1)$ and right bottom corner at $(x2, y2)$.
Time: build $O(n \cdot m)$, query $O(1)$.

```
struct psum2d {
```

```
  vll2d s;
  vll2d psum;
  psum2d(vll2d &grid, int n, int m)
    : s(n + 1, vll(m + 1)), psum(n + 1, vll(m + 1)) {
    for (int i = 1; i <= n; i++)
      for (int j = 1; j <= m; j++)
        s[i][j] = s[i][j - 1] + grid[i][j];

    for (int i = 1; i <= n; i++)
      for (int j = 1; j <= m; j++)
        psum[i][j] = psum[i - 1][j] + s[i][j];
  }

  ll query(int x1, int y1, int x2, int y2) {
    ll ans = psum[x2][y2] + psum[x1 - 1][y1 - 1];
    ans -= psum[x2][y1 - 1] + psum[x1 - 1][y2];
    return ans;
  }
};
```

1.12 SegTree Range Sum Query Range PA sum/set Update

Makes arithmetic progression updates in range and sum queries.
Considering $PA(A, R) = [A + R, A + 2R, A + 3R, ...]$

- **update_set(l, r, A, R):** sets [l, r] to PA(A, R)
- **update_add(l, r, A, R):** sum PA(A, R) in [l, r]
- **query(l, r):** sum in range [l, r]

0 indexed !
Time: build $O(n)$, updates and queries $O(logn)$

```
const ll oo = 1e18;
struct SegTree {
  struct Data {
    ll sum;
    ll set_a, set_r, add_a, add_r;
    Data()
      : sum(0), set_a(oo), set_r(0), add_a(0), add_r(0) {}
  };
  int n;
  vector<Data> seg;
  SegTree(int n_) : n(n_), seg(vector<Data>(4 * n)) {}

  void prop(int p, int l, int r) {
    int sz = r - l + 1;
    ll &sum = seg[p].sum, &set_a = seg[p].set_a,
       &set_r = seg[p].set_r, &add_a = seg[p].add_a,
       &add_r = seg[p].add_r;

    if (set_a != oo) {
      set_a += add_a, set_r += add_r;
      sum = set_a * sz + set_r * sz * (sz + 1) / 2;
      if (l != r) {
        int m = (l + r) / 2;

        seg[2 * p].set_a = set_a;
        seg[2 * p].set_r = set_r;
        seg[2 * p].add_a = seg[2 * p].add_r = 0;

        seg[2 * p + 1].set_a = set_a + set_r * (m - l + 1);
        seg[2 * p + 1].set_r = set_r;
        seg[2 * p + 1].add_a = seg[2 * p + 1].add_r = 0;
      }
      set_a = oo, set_r = 0;
      add_a = add_r = 0;
    } else if (add_a or add_r) {
      sum += add_a * sz + add_r * sz * (sz + 1) / 2;
      if (l != r) {
        int m = (l + r) / 2;

        seg[2 * p].add_a += add_a;
        seg[2 * p].add_r += add_r;

        seg[2 * p + 1].add_a += add_a + add_r * (m - l + 1);
        seg[2 * p + 1].add_r += add_r;
      }
      add_a = add_r = 0;
    }
  }

  int inter(pii a, pii b) {
    if (a.first > b.first) swap(a, b);
    return max(0, min(a.second, b.second) - b.first + 1);
  }
  ll set(int a, int b, ll aa, ll rr, int p, int l, int r) {
    prop(p, l, r);
    if (b < l or r < a) return seg[p].sum;
    if (a <= l and r <= b) {
      seg[p].set_a = aa;
      seg[p].set_r = rr;
      prop(p, l, r);
      return seg[p].sum;
```

```
    }
    int m = (l + r) / 2;
    int tam_l = inter({l, m}, {a, b});
    return seg[p].sum = set(a, b, aa, rr, 2 * p, l, m) +
                        set(a, b, aa + rr * tam_l, rr,
                            2 * p + 1, m + 1, r);
  }
  void update_set(int l, int r, ll aa, ll rr) {
    set(l, r, aa, rr, 1, 0, n - 1);
  }
  ll add(int a, int b, ll aa, ll rr, int p, int l, int r) {
    prop(p, l, r);
    if (b < l or r < a) return seg[p].sum;
    if (a <= l and r <= b) {
      seg[p].add_a += aa;
      seg[p].add_r += rr;
      prop(p, l, r);
      return seg[p].sum;
    }
    int m = (l + r) / 2;
    int tam_l = inter({l, m}, {a, b});
    return seg[p].sum = add(a, b, aa, rr, 2 * p, l, m) +
                        add(a, b, aa + rr * tam_l, rr,
                            2 * p + 1, m + 1, r);
  }
  void update_add(int l, int r, ll aa, ll rr) {
    add(l, r, aa, rr, 1, 0, n - 1);
  }
  ll query(int a, int b, int p, int l, int r) {
    prop(p, l, r);
    if (b < l or r < a) return 0;
    if (a <= l and r <= b) return seg[p].sum;
    int m = (l + r) / 2;
    return query(a, b, 2 * p, l, m) +
           query(a, b, 2 * p + 1, m + 1, r);
  }
  ll query(int l, int r) {
    return query(l, r, 1, 0, n - 1);
  }
};
```

1.13 SegTree Point Update (dynamic function)

Answers queries of any monoid operation (i.e. has identity element and is associative)
Build: $O(N)$, Query: $O(\log N)$

```
#define F(expr) [](auto a, auto b) { return expr; }
template <typename T>
struct SegTree {
  using Operation = T (*)(T, T);

  int N;
  vector<T> ns;
  Operation operation;
  T identity;

  SegTree(int n, Operation op = F(a + b), T neutral = T())
    : N(n),
      ns(2 * N, neutral),
      operation(op),
      identity(neutral) {}

  SegTree(const vector<T> &v, Operation op = F(a + b),
          T neutral = T())
    : SegTree((int)v.size(), op, neutral) {
    copy(v.begin(), v.end(), ns.begin() + N);

    for (int i = N - 1; i > 0; --i)
      ns[i] = operation(ns[2 * i], ns[2 * i + 1]);
  }

  T query(size_t i) const { return ns[i + N]; }

  T query(size_t l, size_t r) const {
    auto a = l + N, b = r + N;
    auto ans = identity;
    // Non-associative operations needs to be processed
    // backwards
    stack<T> st;
    while (a <= b) {
      if (a & 1) ans = operation(ans, ns[a++]);
      if (not(b & 1)) st.push(ns[b--]);

      a >>= 1;
      b >>= 1;
    }

    for (; !st.empty(); st.pop())
      ans = operation(ans, st.top());

    return ans;
```

```
  }

  void update(size_t i, T value) {
    update_set(i, operation(ns[i + N], value));
  }

  void update_set(size_t i, T value) {
    auto a = i + N;

    ns[a] = value;
    while (a >>= 1)
      ns[a] = operation(ns[2 * a], ns[2 * a + 1]);
  }
};
```

1.14 Segtree Range Max Query Point Max Assign Update (dynamic)

Answers range queries in ranges until 10^9 (maybe more)
Time: query and update $O(n \cdot \log n)$

```
struct node;
node *newNode();

struct node {
  node *left, *right;
  int lv, rv;
  ll val;

  node() : left(NULL), right(NULL), val(-oo) {}

  inline void init(int l, int r) {
    lv = l;
    rv = r;
  }

  inline void extend() {
    if (!left) {
      int m = (lv + rv) / 2;
      left = newNode();
      right = newNode();
      left->init(lv, m);
      right->init(m + 1, rv);
    }
  }

  ll query(int l, int r) {
    if (r < lv || rv < l) {
      return 0;
    }

    if (l <= lv && rv <= r) {
      return val;
    }

    extend();
    return max(left->query(l, r), right->query(l, r));
  }

  void update(int p, ll newVal) {
    if (lv == rv) {
      val = max(val, newVal);
      return;
    }

    extend();
    (p <= left->rv ? left : right)->update(p, newVal);
    val = max(left->val, right->val);
  }
};

const int BUFFSZ(1e7);
node *newNode() {
  static int bufSize = BUFFSZ;
  static node buf[(int)BUFFSZ];
  assert(bufSize);
  return &buf[--bufSize];
}

struct SegTree {
  int n;
  node *root;
  SegTree(int _n) : n(_n) {
    root = newNode();
    root->init(0, n);
  }
  ll query(int l, int r) { return root->query(l, r); }
  void update(int p, ll v) { root->update(p, v); }
};
```

1.15 Segtree Range Max Query Range Max Update

```cpp
template <typename T = ll>
struct SegTree {
  int N;
  T nu, nq;
  vector<T> st, lazy;
  SegTree(const vector<T> &xs)
    : N(len(xs)),
      nu(numeric_limits<T>::min()),
      nq(numeric_limits<T>::min()),
      st(4 * N + 1, nu),
      lazy(4 * N + 1, nu) {
    for (int i = 0; i < len(xs); ++i) update(i, i, xs[i]);
  }

  void update(int l, int r, T value) {
    update(1, 0, N - 1, l, r, value);
  }

  T query(int l, int r) { return query(1, 0, N - 1, l, r); }

  void update(int node, int nl, int nr, int ql, int qr,
              T v) {
    propagation(node, nl, nr);

    if (ql > nr or qr < nl) return;

    st[node] = max(st[node], v);
    if (ql <= nl and nr <= qr) {
      if (nl < nr) {
        lazy[left(node)] = max(lazy[left(node)], v);
        lazy[right(node)] = max(lazy[right(node)], v);
      }
      return;
    }
    update(left(node), nl, mid(nl, nr), ql, qr, v);
    update(right(node), mid(nl, nr) + 1, nr, ql, qr, v);

    st[node] = max(st[left(node)], st[right(node)]);
  }

  T query(int node, int nl, int nr, int ql, int qr) {
    propagation(node, nl, nr);

    if (ql > nr or qr < nl) return nq;

    if (ql <= nl and nr <= qr) return st[node];

    T x = query(left(node), nl, mid(nl, nr), ql, qr);
    T y = query(right(node), mid(nl, nr) + 1, nr, ql, qr);

    return max(x, y);
  }

  void propagation(int node, int nl, int nr) {
    if (lazy[node] != nu) {
      st[node] = max(st[node], lazy[node]);

      if (nl < nr) {
        lazy[left(node)] =
          max(lazy[left(node)], lazy[node]);
        lazy[right(node)] =
          max(lazy[right(node)], lazy[node]);
      }

      lazy[node] = nu;
    }
  }

  int left(int p) { return p << 1; }
  int right(int p) { return (p << 1) + 1; }
  int mid(int l, int r) { return (r - l) / 2 + l; }
};
int main() {
  int n;
  cin >> n;
  vector<array<int, 3>> xs(n);
  for (int i = 0; i < n; ++i) {
    for (int j = 0; j < 3; ++j) {
      cin >> xs[i][j];
    }
  }
  vi aux(n, 0);
  SegTree<int> st(aux);
  for (int i = 0; i < n; ++i) {
    int a = min(i + xs[i][1], n);
    int b = min(i + xs[i][2], n);
    st.update(i, i, st.query(i, i) + xs[i][0]);
    int cur = st.query(i, i);
```

```
    st.update(a, b, cur);
  }

  cout << st.query(0, n) << '\n';
}
```

1.16 SegTree Range Min Query Point Assign Update

```
template <typename T = ll>
struct SegTree {
  int n;
  T nu, nq;
  vector<T> st;
  SegTree(const vector<T> &v)
    : n(len(v)),
      nu(0),
      nq(numeric_limits<T>::max()),
      st(n * 4 + 1, nu) {
    for (int i = 0; i < n; ++i) update(i, v[i]);
  }
  void update(int p, T v) { update(1, 0, n - 1, p, v); }
  T query(int l, int r) { return query(1, 0, n - 1, l, r); }

  void update(int node, int nl, int nr, int p, T v) {
    if (p < nl or p > nr) return;

    if (nl == nr) {
      st[node] = v;
      return;
    }

    update(left(node), nl, mid(nl, nr), p, v);
    update(right(node), mid(nl, nr) + 1, nr, p, v);

    st[node] = min(st[left(node)], st[right(node)]);
  }

  T query(int node, int nl, int nr, int ql, int qr) {
    if (ql <= nl and qr >= nr) return st[node];
    if (nl > qr or nr < ql) return nq;
    if (nl == nr) return st[node];

    return min(
      query(left(node), nl, mid(nl, nr), ql, qr),
      query(right(node), mid(nl, nr) + 1, nr, ql, qr));
  }

  int left(int p) { return p << 1; }
  int right(int p) { return (p << 1) + 1; }
  int mid(int l, int r) { return (r - l) / 2 + l; }
};
```

1.17 Segtree Range Sum Query Point Sum Update (dynamic)

Answers range queries in ranges until 10^9 (maybe more)
Time: query and update $O(n \cdot \log n)$

```
struct node;
node *newNode();

struct node {
  node *left, *right;
  int lv, rv;
  ll val;

  node() : left(NULL), right(NULL), val(0) {}

  inline void init(int l, int r) {
    lv = l;
    rv = r;
  }

  inline void extend() {
    if (!left) {
      int m = (rv - lv) / 2 + lv;
      left = newNode();
      right = newNode();
      left->init(lv, m);
      right->init(m + 1, rv);
    }
  }

  ll query(int l, int r) {
    if (r < lv || rv < l) {
      return 0;
    }

    if (l <= lv && rv <= r) {
      return val;
    }

    extend();
```

```
    return left->query(l, r) + right->query(l, r);
  }

  void update(int p, ll newVal) {
    if (lv == rv) {
      val += newVal;
      return;
    }

    extend();
    (p <= left->rv ? left : right)->update(p, newVal);
    val = left->val + right->val;
  }
};

const int BUFFSZ(1.3e7);
node *newNode() {
  static int bufSize = BUFFSZ;
  static node buf[(int)BUFFSZ];
  // assert(bufSize);
  return &buf[--bufSize];
}

struct SegTree {
  int n;
  node *root;
  SegTree(int _n) : n(_n) {
    root = newNode();
    root->init(0, n);
  }
  ll query(int l, int r) { return root->query(l, r); }
  void update(int p, ll v) { root->update(p, v); }
};
```

1.18 SegTree Range Xor query Point Assign Update

```
template <typename T = ll>
struct SegTree {
  int n;
  T nu, nq;
  vector<T> st;
  SegTree(const vector<T> &v)
    : n(len(v)), nu(0), nq(0), st(n * 4 + 1, nu) {
    for (int i = 0; i < n; ++i) update(i, v[i]);
  }
  void update(int p, T v) { update(1, 0, n - 1, p, v); }
  T query(int l, int r) { return query(1, 0, n - 1, l, r); }

  void update(int node, int nl, int nr, int p, T v) {
    if (p < nl or p > nr) return;

    if (nl == nr) {
      st[node] = v;
      return;
    }

    update(left(node), nl, mid(nl, nr), p, v);
    update(right(node), mid(nl, nr) + 1, nr, p, v);

    st[node] = st[left(node)] ^ st[right(node)];
  }

  T query(int node, int nl, int nr, int ql, int qr) {
    if (ql <= nl and qr >= nr) return st[node];
    if (nl > qr or nr < ql) return nq;
    if (nl == nr) return st[node];

    return query(left(node), nl, mid(nl, nr), ql, qr) ^
           query(right(node), mid(nl, nr) + 1, nr, ql, qr);
  }

  int left(int p) { return p << 1; }
  int right(int p) { return (p << 1) + 1; }
  int mid(int l, int r) { return (r - l) / 2 + l; }
};
```

1.19 SegTree Range Min Query Range Sum Update

```
template <typename t = ll>
struct SegTree {
  int n;
  t nu;
  t nq;
  vector<t> st, lazy;
  SegTree(const vector<t> &xs)
    : n(len(xs)),
      nu(0),
      nq(numeric_limits<t>::max()),
      st(4 * n, nu),
      lazy(4 * n, nu) {
```

```
  for (int i = 0; i < len(xs); ++i) update(i, i, xs[i]);
}

SegTree(int n) : n(n), st(4 * n, nu), lazy(4 * n, nu) {}

void update(int l, int r, ll value) {
  update(1, 0, n - 1, l, r, value);
}

t query(int l, int r) { return query(1, 0, n - 1, l, r); }

void update(int node, int nl, int nr, int ql, int qr,
            ll v) {
  propagation(node, nl, nr);

  if (ql > nr or qr < nl) return;

  if (ql <= nl and nr <= qr) {
    st[node] += (nr - nl + 1) * v;

    if (nl < nr) {
      lazy[left(node)] += v;
      lazy[right(node)] += v;
    }

    return;
  }

  update(left(node), nl, mid(nl, nr), ql, qr, v);
  update(right(node), mid(nl, nr) + 1, nr, ql, qr, v);

  st[node] = min(st[left(node)], st[right(node)]);
}

t query(int node, int nl, int nr, int ql, int qr) {
  propagation(node, nl, nr);

  if (ql > nr or qr < nl) return nq;

  if (ql <= nl and nr <= qr) return st[node];

  t x = query(left(node), nl, mid(nl, nr), ql, qr);
  t y = query(right(node), mid(nl, nr) + 1, nr, ql, qr);

  return min(x, y);
}

void propagation(int node, int nl, int nr) {
  if (lazy[node]) {
    st[node] += lazy[node];

    if (nl < nr) {
      lazy[left(node)] += lazy[node];
      lazy[right(node)] += lazy[node];
    }

    lazy[node] = nu;
  }
}

int left(int p) { return p << 1; }
int right(int p) { return (p << 1) + 1; }
int mid(int l, int r) { return (r - l) / 2 + l; }
};
```

1.20 SegTree Range Sum Query Range Sum Update

```
template <typename T = ll>
struct SegTree {
  int N;
  T nu;
  T nq;
  vector<T> st, lazy;
  SegTree(const vector<T> &xs)
    : N(len(xs)),
      nu(0),
      nq(0),
      st(4 * N, nu),
      lazy(4 * N, nu) {
    for (int i = 0; i < len(xs); ++i) update(i, i, xs[i]);
  }

  SegTree(int n)
    : N(n), nu(0), nq(0), st(4 * N, nu), lazy(4 * N, nu) {}

  void update(int l, int r, ll value) {
    update(1, 0, N - 1, l, r, value);
  }

  T query(int l, int r) { return query(1, 0, N - 1, l, r); }
```

```cpp
void update(int node, int nl, int nr, int ql, int qr,
            ll v) {
  propagation(node, nl, nr);

  if (ql > nr or qr < nl) return;

  if (ql <= nl and nr <= qr) {
    st[node] += (nr - nl + 1) * v;

    if (nl < nr) {
      lazy[left(node)] += v;
      lazy[right(node)] += v;
    }

    return;
  }

  update(left(node), nl, mid(nl, nr), ql, qr, v);
  update(right(node), mid(nl, nr) + 1, nr, ql, qr, v);

  st[node] = st[left(node)] + st[right(node)];
}

T query(int node, int nl, int nr, int ql, int qr) {
  propagation(node, nl, nr);

  if (ql > nr or qr < nl) return nq;

  if (ql <= nl and nr <= qr) return st[node];

  T x = query(left(node), nl, mid(nl, nr), ql, qr);
  T y = query(right(node), mid(nl, nr) + 1, nr, ql, qr);

  return x + y;
}

void propagation(int node, int nl, int nr) {
  if (lazy[node]) {
    st[node] += (nr - nl + 1) * lazy[node];

    if (nl < nr) {
      lazy[left(node)] += lazy[node];
      lazy[right(node)] += lazy[node];
    }

      lazy[node] = nu;
    }
  }

  int left(int p) { return p << 1; }
  int right(int p) { return (p << 1) + 1; }
  int mid(int l, int r) { return (r - l) / 2 + l; }
};
```

1.21 Sparse Table

Answer the range query defined at the function **op**.
Build: O(NlogN), Query: O(1)

```cpp
template <typename T>
struct SparseTable {
  vector<T> v;
  int n;
  static const int b = 30;
  vi mask, t;

  int op(int x, int y) { return v[x] < v[y] ? x : y; }
  int msb(int x) {
    return __builtin_clz(1) - __builtin_clz(x);
  }
  SparseTable() {}
  SparseTable(const vector<T>& v_)
    : v(v_), n(v.size()), mask(n), t(n) {
    for (int i = 0, at = 0; i < n; mask[i++] = at |= 1) {
      at = (at << 1) & ((1 << b) - 1);
      while (at and op(i, i - msb(at & -at)) == i)
        at ^= at & -at;
    }
    for (int i = 0; i < n / b; i++)
      t[i] = b * i + b - 1 - msb(mask[b * i + b - 1]);
    for (int j = 1; (1 << j) <= n / b; j++)
      for (int i = 0; i + (1 << j) <= n / b; i++)
        t[n / b * j + i] =
          op(t[n / b * (j - 1) + i],
             t[n / b * (j - 1) + i + (1 << (j - 1))]);
  }
  int small(int r, int sz = b) {
    return r - msb(mask[r] & ((1 << sz) - 1));
  }
  T query(int l, int r) {
    if (r - l + 1 <= b) return small(r, r - l + 1);
```

```
    int ans = op(small(l + b - 1), small(r));
    int x = l / b + 1, y = r / b - 1;
    if (x <= y) {
      int j = msb(y - x + 1);
      ans = op(ans, op(t[n / b * j + x],
                       t[n / b * j + y - (1 << j) + 1]));
    }
    return ans;
  }
};
```

2 Dynamic programming

2.1 Binary Knapsack (bottom up)

Given N items, each with its own value V_i and weight W_i and a maximum knapsack weight W, compute the maximum value of the items that we can carry, if we can either ignore or take a particular item.
Assume that $1 \leq n \leq 1000$, $1 \leq S \leq 10000$.
Time and space: $O(N * W)$
the vectors VS and WS starts at one, so it need an empty value at index 0.

```
const int MAXN(2010), MAXM(2010);
ll st[MAXN + 1][MAXM + 1];
char ps[MAXN + 1][MAXM + 1];
pair<ll, vi> knapsack(int M, const vll &VS, const vi &WS) {
  memset(st, 0, sizeof(st));
  memset(ps, 0, sizeof(ps));
  int N = len(VS) - 1;  // ELEMENTS START AT INDEX 1 !

  for (int i = 0; i <= N; ++i) st[i][0] = 0;

  for (int m = 0; m <= M; ++m) st[0][m] = 0;

  for (int i = 1; i <= N; ++i) {
    for (int m = 1; m <= M; ++m) {
      st[i][m] = st[i - 1][m];
      ps[i][m] = 0;
      int w = WS[i];
      ll v = VS[i];

      if (w <= m and st[i - 1][m - w] + v > st[i][m]) {
        st[i][m] = st[i - 1][m - w] + v;
        ps[i][m] = 1;
      }
    }
  }
  int m = M;
  vi is;
  for (int i = N; i >= 1; --i) {
    if (ps[i][m]) {
      is.emplace_back(i - 1);
      m -= WS[i];
    }
  }

  return {st[N][M], is};
}
```

2.2 Binary Knapsack (top down)

Given N items, each with its own value V_i and weight W_i and a maximum knapsack weight W, compute the maximum value of the items that we can carry, if we can either ignore or take a particular item.
Assume that $1 \leq n \leq 1000$, $1 \leq S \leq 10000$.
Time and space: $O(N * W)$
the bottom up version is 5 times faster !

```
const int MAXN(2000), MAXM(2000);
ll memo[MAXN][MAXM + 1];
char choosen[MAXN][MAXM + 1];
ll knapSack(int u, int w, vll &VS, vi &WS) {
  if (u < 0) return 0;
  if (memo[u][w] != -1) return memo[u][w];

  ll a = 0, b = 0;
  a = knapSack(u - 1, w, VS, WS);
  if (WS[u] <= w)
    b = knapSack(u - 1, w - WS[u], VS, WS) + VS[u];
  if (b > a) {
    choosen[u][w] = true;
  }
  return memo[u][w] = max(a, b);
}
pair<ll, vi> knapSack(int W, vll &VS, vi &WS) {
  memset(memo, -1, sizeof(memo));
  memset(choosen, 0, sizeof(choosen));
  int n = len(VS);
  ll v = knapSack(n - 1, W, VS, WS);
  ll cw = W;
  vi choosed;
  for (int i = n - 1; i >= 0; i--) {
    if (choosen[i][cw]) {
      cw -= WS[i];
```

```
      choosed.emplace_back(i);
    }
  }
  return {v, choosed};
}
```

2.3 Edit Distance

$O(N * M)$

```
int edit_distance(const string &a, const string &b) {
  int n = a.size();
  int m = b.size();
  vector<vi> dp(n + 1, vi(m + 1, 0));

  int ADD = 1, DEL = 1, CHG = 1;
  for (int i = 0; i <= n; ++i) {
    dp[i][0] = i * DEL;
  }
  for (int i = 1; i <= m; ++i) {
    dp[0][i] = ADD * i;
  }

  for (int i = 1; i <= n; ++i) {
    for (int j = 1; j <= m; ++j) {
      int add = dp[i][j - 1] + ADD;
      int del = dp[i - 1][j] + DEL;
      int chg = dp[i - 1][j - 1] +
                (a[i - 1] == b[j - 1] ? 0 : 1) * CHG;
      dp[i][j] = min({add, del, chg});
    }
  }

  return dp[n][m];
}
```

2.4 Kadane

Find the maximum subarray sum in a given a rray.

```
int kadane(const vi &as) {
  vi s(len(as));
  s[0] = as[0];

  for (int i = 1; i < len(as); ++i)
    s[i] = max(as[i], s[i - 1] + as[i]);
  return *max_element(all(s));
}
```

2.5 Longest Increasing Subsequence (LIS)

Finds the lenght of the longest subsequence in

$$O(n \log n)$$

.

```
int LIS(const vi& as) {
  const ll oo = 1e18;
  int n = len(as);
  vll lis(n + 1, oo);
  lis[0] = -oo;

  auto ans = 0;

  for (int i = 0; i < n; ++i) {
    auto it = lower_bound(all(lis), as[i]);
    auto pos = (int)(it - lis.begin());

    ans = max(ans, pos);
    lis[pos] = as[i];
  }

  return ans;
}
```

2.6 Money Sum (Bottom Up)

Find every possible sum using the given values only once.

```
set<int> money_sum(const vi &xs) {
  using vc = vector<char>;
  using vvc = vector<vc>;
  int _m = accumulate(all(xs), 0);
  int _n = xs.size();
  vvc _dp(_n + 1, vc(_m + 1, 0));
  set<int> _ans;
  _dp[0][xs[0]] = 1;
  for (int i = 1; i < _n; ++i) {
    for (int j = 0; j <= _m; ++j) {
      if (j == 0 or _dp[i - 1][j]) {
        _dp[i][j + xs[i]] = 1;
        _dp[i][j] = 1;
```

```
      }
    }
  }

  for (int i = 0; i < _n; ++i)
    for (int j = 0; j <= _m; ++j)
      if (_dp[i][j]) _ans.insert(j);
  return _ans;
}
```

2.7 Travelling Salesman Problem

```
using vi = vector<int>;
vector<vi> dist;
vector<vi> memo;
/* O ( N^2 * 2^N )*/
int tsp(int i, int mask, int N) {
  if (mask == (1 << N) - 1) return dist[i][0];
  if (memo[i][mask] != -1) return memo[i][mask];
  int ans = INT_MAX << 1;
  for (int j = 0; j < N; ++j) {
    if (mask & (1 << j)) continue;
    auto t = tsp(j, mask | (1 << j), N) + dist[i][j];
    ans = min(ans, t);
  }
  return memo[i][mask] = ans;
}
```

3 Geometry

3.1 Convex Hull

Given a set of points find the smallest convex polygon that contains all the given points.
Time: $O(N \log N)$
By default it removes the collinear points, set the boolean to true if you don't want that

```
struct pt {
  double x, y;
  int id;
};

int orientation(pt a, pt b, pt c) {
  double v = a.x * (b.y - c.y) + b.x * (c.y - a.y) +
             c.x * (a.y - b.y);
  if (v < 0) return -1;  // clockwise
  if (v > 0) return +1;  // counter-clockwise
  return 0;
}

bool cw(pt a, pt b, pt c, bool include_collinear) {
  int o = orientation(a, b, c);
  return o < 0 || (include_collinear && o == 0);
}
bool collinear(pt a, pt b, pt c) {
  return orientation(a, b, c) == 0;
}

void convex_hull(vector<pt>& pts,
                 bool include_collinear = false) {
  pt p0 = *min_element(all(pts), [](pt a, pt b) {
    return make_pair(a.y, a.x) < make_pair(b.y, b.x);
  });
  sort(all(pts), [&p0](const pt& a, const pt& b) {
    int o = orientation(p0, a, b);
    if (o == 0)
      return (p0.x - a.x) * (p0.x - a.x) +
               (p0.y - a.y) * (p0.y - a.y) <
             (p0.x - b.x) * (p0.x - b.x) +
               (p0.y - b.y) * (p0.y - b.y);
    return o < 0;
  });
  if (include_collinear) {
    int i = len(pts) - 1;
    while (i >= 0 && collinear(p0, pts[i], pts.back())) i--;
    reverse(pts.begin() + i + 1, pts.end());
  }

  vector<pt> st;
  for (int i = 0; i < len(pts); i++) {
    while (st.size() > 1 && !cw(st[len(st) - 2], st.back(),
                                pts[i], include_collinear))
      st.pop_back();
    st.push_back(pts[i]);
  }

  pts = st;
}
```

3.2 Determinant

```
#include "Point.cpp"
```

```cpp
template <typename T>
T D(const Point<T> &P, const Point<T> &Q,
    const Point<T> &R) {
  return (P.x * Q.y + P.y * R.x + Q.x * R.y) -
         (R.x * Q.y + R.y * P.x + Q.x * P.y);
}
```

3.3 Equals

```cpp
template <typename T>
bool equals(T a, T b) {
  const double EPS{1e-9};
  if (is_floating_point<T>::value)
    return fabs(a - b) < EPS;
  else
    return a == b;
}
```

3.4 Line

```cpp
#include <bits/stdc++.h>

#include "point-struct-and-utils.cpp"
using namespace std;

struct line {
  ld a, b, c;
};

// the answer is stored in the third parameter (pass by
// reference)
void pointsToLine(const point &p1, const point &p2,
                  line &l) {
  if (fabs(p1.x - p2.x) < EPS)
    // vertical line
    l = {1.0, 0.0, -p1.x};
  // default values
  else
    l = {-(ld)(p1.y - p2.y) / (p1.x - p2.x), 1.0,
         -(ld)(l.a * p1.x) - p1.y};
}
```

3.5 Point Struct And Utils (2d)

```cpp
#include <bits/stdc++.h>
using namespace std;
using ld = long double;

struct point {
  ld x, y;
  int id;
  point(ld x = 0.0, ld y = 0.0, int id = -1)
    : x(x), y(y), id(id) {}

  point& operator+=(const point& t) {
    x += t.x;
    y += t.y;
    return *this;
  }
  point& operator-=(const point& t) {
    x -= t.x;
    y -= t.y;
    return *this;
  }
  point& operator*=(ld t) {
    x *= t;
    y *= t;
    return *this;
  }
  point& operator/=(ld t) {
    x /= t;
    y /= t;
    return *this;
  }
  point operator+(const point& t) const {
    return point(*this) += t;
  }
  point operator-(const point& t) const {
    return point(*this) -= t;
  }
  point operator*(ld t) const { return point(*this) *= t; }
  point operator/(ld t) const { return point(*this) /= t; }
};

ld dot(point& a, point& b) { return a.x * b.x + a.y * b.y; }

ld norm(point& a) { return dot(a, a); }

ld abs(point a) { return sqrt(norm(a)); }
```

```cpp
ld proj(point a, point b) { return dot(a, b) / abs(b); }

ld angle(point a, point b) {
  return acos(dot(a, b) / abs(a) / abs(b));
}

ld cross(point a, point b) { return a.x * b.y - a.y * b.x; }
```

3.6 Polygon Lattice Points (Pick's Theorem)

Given a polygon with N points finds the number of lattice points inside and on boundaries. Time : $O(N)$

```cpp
ll cross(ll x1, ll y1, ll x2, ll y2) {
  return x1 * y2 - x2 * y1;
}

ll polygonArea(vector<pll>& pts) {
  ll ats = 0;
  for (int i = 2; i < len(pts); i++)
    ats += cross(pts[i].first - pts[0].first,
                 pts[i].second - pts[0].second,
                 pts[i - 1].first - pts[0].first,
                 pts[i - 1].second - pts[0].second);
  return abs(ats / 2ll);
}

ll boundary(vector<pll>& pts) {
  ll ats = pts.size();
  for (int i = 0; i < len(pts); i++) {
    ll deltax =
      (pts[i].first - pts[(i + 1) % pts.size()].first);
    ll deltay =
      (pts[i].second - pts[(i + 1) % pts.size()].second);
    ats += abs(__gcd(deltax, deltay)) - 1;
  }
  return ats;
}

pll latticePoints(vector<pll>& pts) {
  ll bounds = boundary(pts);
  ll area = polygonArea(pts);
  ll inside = area + 1ll - bounds / 2ll;

  return {inside, bounds};
}
```

3.7 Segment

```cpp
#include "Line.cpp"
#include "Point.cpp"
#include "equals.cpp"

template <typename T>
struct segment {
  Point<T> A, B;

  bool contains(const Point<T> &P) const;

  Point<T> closest(const Point<T> &p) const;
};

template <typename T>
bool segment<T>::contains(const Point<T> &P) const {
  // verifica se P áest contido na reta
  double dAB = Point<T>::dist(A, B),
         dAP = Point<T>::dist(A, P),
         dPB = Point<T>::dist(P, B);

  return equals(dAP + dPB, dAB);
}

template <typename T>
Point<T> segment<T>::closest(const Point<T> &P) const {
  Line<T> R(A, B);
  auto Q = R.closest(P);

  if (this->contains(Q)) return Q;

  auto distA = Point<T>::dist(P, A);
  auto distB = Point<T>::dist(P, B);

  if (distA <= distB)
    return A;
  else
    return B;
}
```

3.8 Template Line

```cpp
#include "template-point.cpp"

template <typename T>
```

```
struct Line {
  T a, b, c;

  Line(T av, T bv, T cv) : a(av), b(bv), c(cv) {}

  Line(const Point<T> &P, const Point<T> &Q)
    : a(P.y - Q.y),
      b(Q.x - P.x),
      c(P.x * Q.y - Q.x * P.y) {}

  // verify if a point belongs to the line
  bool contains(const Point<T> &P) {
    return equals(a * P.x + b * P.y + c, 0);
  }

  // shortest distance between P and a point Q that belongs
  // to this line
  double distance(const Point<T> &P) const {
    return fabs(a * P.x + b * P.y + c) / hypot(a, b);
  }

  // the closest point in this line to the given point
  Point<T> closest(const Point<T> &P) const {
    auto den = (a * a) + (b * b);

    auto x = (b * (b * P.x - a * P.y) - a * c) / den;
    auto y = (a * (-b * P.x + a * P.y) - b * c) / den;

    return Point<T>{x, y};
  }
};
```

3.9 Template Point

```
template <typename T>
struct Point {
  T x, y;

  Point(T xv = 0, T yv = 0) : x(xv), y(yv) {}

  double distance(const Point<T> &P) const {
    return hypot(static_cast<double>(P.x - this->x),
                 static_cast<double>(P.y - this->y));
  }
};
```

3.10 Template Segment

```
#include "equals.cpp"
#include "template-line.cpp"
#include "template-point.cpp"

template <typename T>
struct Segment {
  Point<T> A, B;

  Segment(const Point<T> &a, const Point<T> &b)
    : A(a), B(b) {}

  /*
   * Verify if a given point P belongs to the segment,
   * considering that P belongs to the line defined with A
   * and B
   */
  bool contains(const Point<T> &P) const {
    return equals(A.x, B.x)
             ? min(A.y, B.y) <= P.y and P.y <= max(A.y, B.y)
             : min(A.x, B.x) <= P.x and
                 P.x <= max(A.x, B.x);
  }

  /*
   * Verify if P belongs to the segment AB,
   * even if P don't belong to the line defined with A and B
   * */
  bool contains2(const Point<T> &P) const {
    double dAB = dist(A, B), dAP = dist(A, P),
           dPB = dist(P, B);
    return equals(dAP + dPB, dAB);
  }

  /*
   * Find the closest point in P that belongs to the segment
   * */
  Point<T> closest(const Point<T> &P) {
    Line<T> r(A, B);
    auto Q = r.closest(P);

    if (this->contains(Q)) return Q;

    auto distA = P.distance(A);
```

```
    auto distB = P.distance(B);

    return distA <= distB ? A : B;
  }

  double distToClosest(const Point<T> &P) {
    return closest(P).distance(P);
  }
};
```

4 Graphs

4.1 2 SAT

```
struct SAT {
  int n;
  vi2d g, tg;
  vi vis;
  vi order, comp;
  vc assignment;
  bool solvable;
  int qtdcomp;

  SAT(int _n)
    : n(2 * _n),
      g(n),
      tg(n),
      vis(n),
      comp(n, -1),
      assignment(n / 2) {}

  void dfs1(int u) {
    vis[u] = 1;
    for (auto v : g[u]) {
      if (!vis[v]) {
        dfs1(v);
      }
    }
    order.emplace_back(u);
  }

  void dfs2(int u) {
    comp[u] = qtdcomp;
    for (auto v : tg[u]) {
      if (comp[v] == -1) {
        dfs2(v);
      }
    }
  }

  bool solve2sat() {
    for (int i = 0; i < n; i++) {
      if (!vis[i]) dfs1(i);
    }

    reverse(all(order));
    qtdcomp = 0;
    for (auto u : order) {
      if (comp[u] == -1) {
        dfs2(u);
        qtdcomp++;
      }
    }

    assignment.assign(n / 2, false);
    for (int i = 0; i < n; i += 2) {
      if (comp[i] == comp[i + 1]) {
        solvable = false;
        return false;
      }

      assignment[i / 2] = comp[i] < comp[i + 1];
    }

    solvable = 1;
    return solvable;
  }

  void add_dis(int a, bool va, int b, bool vb) {  // a V b
    va = !va, vb = !vb;
    a = (2 * a) ^ va, b = (2 * b) ^ vb;
    int nota = a ^ 1, notb = b ^ 1;
    g[nota].emplace_back(b), g[notb].emplace_back(a),
      tg[b].emplace_back(nota), tg[a].emplace_back(notb);
  }

  void add_impl(int a, bool va, int b, int vb) {  // a -> b
    add_dis(a, !va, b, vb);
  }
```

```
  void add_equiv(int a, bool va, int b,
                 bool vb) {  // a <-> b
    add_impl(a, 1, b, 1);
    add_impl(b, 1, a, 1);
    add_impl(a, 0, b, 0);
    add_impl(b, 0, a, 0);
  }

  void add_xor(int a, bool va, int b, bool vb) {  // a xor b
    add_impl(a, 1, b, 0);
    add_impl(a, 0, b, 1);
    add_impl(b, 1, a, 0);
    add_impl(b, 0, a, 1);
  }
};
```

4.2 Cycle Distances

Given a vertex s finds the longest cycle that end's in s, note that the vector **dist** will contain the distance that each vertex u needs to reach s.
Time: $O(N)$

```
using adj = vector<vector<pair<int, ll>>>;
ll cycleDistances(int u, int n, int s, vc &vis, adj &g,
                  vll &dist) {
  vis[u] = 1;

  for (auto [v, d] : g[u]) {
    if (v == s) {
      dist[u] = max(dist[u], d);
      continue;
    }

    if (vis[v] == 1) {
      continue;
    }

    if (vis[v] == 2) {
      dist[u] = max(dist[u], dist[v] + d);
    } else {
      ll d2 = cycleDistances(v, n, s, vis, g, dist);
      if (d2 != -oo) {
        dist[u] = max(dist[u], d2 + d);
      }
    }
  }
  vis[u] = 2;
  return dist[u];
}
```

4.3 SCC (struct)

Build the condensation graph based in the strongly connected components.
tiem: $O(V + E)$

```
struct SCC {
  int N, totscc;
  vi2d g, tg;
  vi todo, comp;
  vector<set<ll>> gscc;
  vc vis;
  SCC(int _N)
    : N(_N),
      totscc(0),
      g(_N),
      tg(_N),
      comp(_N, -1),
      gscc(_N),
      vis(_N) {}

  void add_edge(int x, int y) { g[x].eb(y), tg[y].eb(x); }

  void dfs(int x) {
    vis[x] = 1;
    for (auto y : g[x])
      if (!vis[y]) dfs(y);
    todo.pb(x);
  }

  void dfs2(ll x) {
    comp[x] = totscc;
    for (auto y : tg[x])
      if (comp[y] == -1) dfs2(y);
  }

  void build() {
    for (int i = 0; i < N; ++i)
      if (!vis[i]) dfs(i);

    reverse(all(todo));
    for (auto &x : todo)
      if (comp[x] == -1) {
        dfs2(x);
```

```
        totscc++;
      }

    for (int i = 0; i < N; ++i)
      for (auto j : g[i])
        if (comp[i] != comp[j])
          gscc[comp[i]].insert(comp[j]);
  }
};
```

4.4 Bellman-Ford (find negative cycle)

Given a directed graph find a negative cycle by running n iterations, and if the last one produces a relaxation than there is a cycle.
Time: $O(V \cdot E)$

```
const ll oo = 2500 * 1e9;

using graph = vector<vector<pair<int, ll>>>;
vi negative_cycle(graph &g, int n) {
  vll d(n, oo);
  vi p(n, -1);
  int x = -1;
  d[0] = 0;
  for (int i = 0; i < n; i++) {
    x = -1;
    for (int u = 0; u < n; u++) {
      for (auto &[v, l] : g[u]) {
        if (d[u] + l < d[v]) {
          d[v] = d[u] + l;
          p[v] = u;
          x = v;
        }
      }
    }
  }

  if (x == -1)
    return {};
  else {
    for (int i = 0; i < n; i++) x = p[x];
    vi cycle;
    for (int v = x;; v = p[v]) {
      cycle.eb(v);
      if (v == x and len(cycle) > 1) break;
    }
    reverse(all(cycle));
    return cycle;
  }
}
```

4.5 Bellman Ford

Find shortest path from a single source to all other nodes. Can detect negative cycles.
Time: $O(V * E)$

```
bool bellman_ford(const vector<vector<pair<int, ll>>> &g,
                  int s, vector<ll> &dist) {
  int n = (int)g.size();
  dist.assign(n, LLONG_MAX);

  vector<int> count(n);
  vector<char> in_queue(n);
  queue<int> q;

  dist[s] = 0;
  q.push(s);
  in_queue[s] = true;

  while (not q.empty()) {
    int cur = q.front();
    q.pop();
    in_queue[cur] = false;

    for (auto [to, w] : g[cur]) {
      if (dist[cur] + w < dist[to]) {
        dist[to] = dist[cur] + w;
        if (not in_queue[to]) {
          q.push(to);
          in_queue[to] = true;
          count[to]++;
          if (count[to] > n) return false;
        }
      }
    }
  }

  return true;
}
```

4.6 BFS 01

Similar to a Dijkstra given a weighted graph finds the distance from source s to every other node (SSSP).
Applicable only when the weight of the edges $\in \{0, x\}$
Time: $O(V + E)$

```
vector<pair<ll, int>> adj[maxn];
ll dists[maxn];
int s, n;
void bfs_01() {
  fill(dists, dists + n, oo);
  dist[s] = 0;

  deque<int> q;
  q.emplace_back(s);

  while (not q.empty()) {
    auto u = q.front();
    q.pop_front();

    for (auto [v, w] : adj[u]) {
      if (dist[v] <= dist[u] + w) continue;
      dist[v] = dist[u] + w;
      w ? q.emplace_back(v) : q.emplace_front(v);
    }
  }
}
```

4.7 Binary Lifting/Jumping

Given a function/successor grpah answers queries of the form which is the node after k moves starting from u.
Time: build $O(N \cdot MAXLOG2)$, query $O(MAXLOG2)$.

```
const int MAXN(2e5), MAXLOG2(30);
int bl[MAXN][MAXLOG2 + 1];
int N;

int jump(int u, ll k) {
  for (int i = 0; i <= MAXLOG2; i++) {
    if (k & (1ll << i)) u = bl[u][i];
  }
  return u;
}
void build() {
  for (int i = 1; i <= MAXLOG2; i++) {
    for (int j = 0; j < N; j++) {
      bl[j][i] = bl[bl[j][i - 1]][i - 1];
    }
  }
}
```

4.8 Block Cut Tree

```
// O(n + m)
struct BlockCutTree {
  vector<vector<int>> blocks, tree;
  vector<vector<pair<int, int>>> block_edges;
  vector<int> articulation, pos;

  BlockCutTree(const vector<vector<int>> &g)
    : articulation(g.size()), pos(g.size()) {
    int t = 0;
    vector<int> id(g.size(), -1);
    stack<int> s1;
    stack<pair<int, int>> s2;
    for (int i = 0; i < (int)g.size(); i++)
      if (id[i] == -1) dfs(g, i, -1, t, id, s1, s2);

    tree.resize(blocks.size());
    for (int i = 0; i < (int)g.size(); i++)
      if (articulation[i])
        pos[i] = (int)tree.size(), tree.emplace_back();

    for (int i = 0; i < (int)blocks.size(); i++) {
      for (auto j : blocks[i]) {
        if (not articulation[j])
          pos[j] = i;
        else
          tree[i].push_back(pos[j]),
            tree[pos[j]].push_back(i);
      }
    }
  }

 private:
  int dfs(const vector<vector<int>> &g, int i, int p,
          int &t, vector<int> &id, stack<int> &s1,
          stack<pair<int, int>> &s2) {
    int lo = id[i] = t++;
    s1.push(i);

    if (p != -1) s2.emplace(i, p);
```

```
    for (auto j : g[i])
      if (j != p and id[j] != -1) s2.emplace(i, j);

    for (auto j : g[i])
      if (j != p) {
        if (id[j] == -1) {
          int val = dfs(g, j, i, t, id, s1, s2);
          lo = min(lo, val);

          if (val >= id[i]) {
            articulation[i]++;
            blocks.emplace_back(1, i);
            for (; blocks.back().back() != j; s1.pop())
              blocks.back().push_back(s1.top());

            block_edges.emplace_back(1, s2.top());
            s2.pop();
            for (; block_edges.back().back() !=
                   make_pair(j, i);
                 s2.pop())
              block_edges.back().push_back(s2.top());
          }
        } else {
          lo = min(lo, id[j]);
        }
      }

    if (p == -1 and articulation[i]) --articulation[i];
    return lo;
  }
};
```

4.9 Check Bipartitie

$O(V)$

```
vi2d G;
int N, M;

bool check() {
  vi side(N, -1);
  queue<int> q;
  for (int st = 0; st < N; st++) {
    if (side[st] == -1) {
      q.emplace(st);
      side[st] = 0;
      while (not q.empty()) {
        int u = q.front();
        q.pop();
        for (auto v : G[u]) {
          if (side[v] == -1) {
            side[v] = side[u] ^ 1;
            q.push(v);
          } else if (side[u] == side[v])
            return false;
        }
      }
    }
  }
  return true;
}
```

4.10 Dijkstra (k Shortest Paths)

```
const ll oo = 1e9 * 1e5 + 1;
using adj = vector<vector<pll>>;
vector<priority_queue<ll>> dijkstra(
  const vector<vector<pll>> &g, int n, int s, int k) {
  priority_queue<pll, vector<pll>, greater<pll>> pq;

  vector<priority_queue<ll>> dist(n);
  dist[0].emplace(0);
  pq.emplace(0, s);
  while (!pq.empty()) {
    auto [d1, v] = pq.top();
    pq.pop();

    if (not dist[v].empty() and dist[v].top() < d1)
      continue;

    for (auto [d2, u] : g[v]) {
      if (len(dist[u]) < k) {
        pq.emplace(d2 + d1, u);
        dist[u].emplace(d2 + d1);
      } else {
        if (dist[u].top() > d1 + d2) {
          dist[u].pop();
          dist[u].emplace(d1 + d2);
          pq.emplace(d2 + d1, u);
        }
      }
    }
```

```
  }
  return dist;
}
```

4.11 Dijkstra

Finds the shortest path from s to every other node, and keep the 'parent' tracking.
Time: $O(E \cdot \log V)$

```
pair<vll, vi> dijkstra(const vector<vector<pll>> &g, int n,
                       int s) {
  priority_queue<pll, vector<pll>, greater<pll>> pq;
  vll dist(n, oo);
  vi p(n, -1);
  pq.emplace(0, s);
  dist[s] = 0;
  while (!pq.empty()) {
    auto [d1, v] = pq.top();
    pq.pop();
    if (dist[v] < d1) continue;

    for (auto [d2, u] : g[v]) {
      if (dist[u] > d1 + d2) {
        dist[u] = d1 + d2;
        p[u] = v;
        pq.emplace(dist[u], u);
      }
    }
  }
  return {dist, p};
}
```

4.12 Disjoint Edges Path (Maxflow)

Given a directed graph find's every path with disjoint edges that starts at s and ends at t
Time : $O(E \cdot V^2)$

```
struct DisjointPaths {
  int n;
  vi2d g, capacity;
  vector<vc> isedge;

  DisjointPaths(int _n)
    : n(_n), g(n), capacity(n, vi(n)), isedge(n, vc(n)) {}

  void add(int u, int v, int w = 1) {
    g[u].emplace_back(v);
    g[v].emplace_back(u);
    capacity[u][v] += w;
    isedge[u][v] = true;
  }

  // finds the new flow to insert
  int bfs(int s, int t, vi &parent) {
    fill(all(parent), -1);
    parent[s] = -2;
    queue<pair<int, int>> q;
    q.push({oo, s});

    while (!q.empty()) {
      auto [flow, cur] = q.front();
      q.pop();

      for (auto next : g[cur]) {
        if (parent[next] == -1 and capacity[cur][next]) {
          parent[next] = cur;
          ll new_flow = min(flow, capacity[cur][next]);
          if (next == t) return new_flow;
          q.push({new_flow, next});
        }
      }
    }

    return 0;
  }

  int maxflow(int s, int t) {
    int flow = 0;
    vi parent(n);
    int new_flow;

    while ((new_flow = bfs(s, t, parent))) {
      flow += new_flow;
      int cur = t;
      while (cur != s) {
        int prev = parent[cur];
        capacity[prev][cur] -= new_flow;
        capacity[cur][prev] += new_flow;
        cur = prev;
      }
    }

    return flow;
```

```
  }

  // build the distinct routes based in the capacity set by
  // maxflow
  void dfs(int u, int t, vc2d &vis, vi &route,
           vi2d &routes) {
    route.eb(u);
    if (u == t) {
      routes.emplace_back(route);
      route.pop_back();
      return;
    }

    for (auto &v : g[u]) {
      if (capacity[u][v] == 0 and isedge[u][v] and
          not vis[u][v]) {
        vis[u][v] = true;
        dfs(v, t, vis, route, routes);
        route.pop_back();
        return;
      }
    }
  }

  vi2d disjoint_paths(int s, int t) {
    int mf = maxflow(s, t);
    vi2d routes;
    vi route;
    vc2d vis(n, vc(n));
    for (int i = 0; i < mf; i++)
      dfs(s, t, vis, route, routes);
    return routes;
  }
};
```

4.13 Euler Path (directed)

Given a **directed** graph finds a path that visits every edge exactly once.
Time: $O(E)$

```
vector<int> euler_cycle(vector<vector<int>> &g, int u) {
  vector<int> res;

  stack<int> st;
  st.push(u);
  while (!st.empty()) {
    auto cur = st.top();
    if (g[cur].empty()) {
      res.push_back(cur);
      st.pop();
    } else {
      auto next = g[cur].back();
      st.push(next);

      g[cur].pop_back();
    }
  }

  for (auto &x : g)
    if (!x.empty()) return {};

  return res;
}

vector<int> euler_path(vector<vector<int>> &g, int first) {
  {
    int n = (int)g.size();
    vector<int> in(n), out(n);
    for (int i = 0; i < n; i++)
      for (auto x : g[i]) in[x]++, out[i]++;

    int a = 0, b = 0, c = 0;
    for (int i = 0; i < n; i++)
      if (in[i] == out[i])
        c++;
      else if (in[i] - out[i] == 1)
        b++;
      else if (in[i] - out[i] == -1)
        a++;

    if (c != n - 2 or a != 1 or b != 1) return {};
  }

  auto res = euler_cycle(g, first);
  if (res.empty()) return res;

  reverse(all(res));
  return res;
}
```

4.14 Euler Path (undirected)

Given a **undirected** graph finds a path that visits every edge exactly once.
Time: $O(E)$

```
vector<int> euler_cycle(vector<vector<int>> &g, int u) {
  vector<int> res;
  multiset<pair<int, int>> vis;

  stack<int> st;
  st.push(u);
  while (!st.empty()) {
    auto cur = st.top();

    while (!g[cur].empty()) {
      auto it = vis.find(make_pair(cur, g[cur].back()));
      if (it == vis.end()) break;
      g[cur].pop_back();
      vis.erase(it);
    }

    if (g[cur].empty()) {
      res.push_back(cur);
      st.pop();
    } else {
      auto next = g[cur].back();
      st.push(next);

      vis.emplace(next, cur);
      g[cur].pop_back();
    }
  }

  for (auto &x : g)
    if (!x.empty()) return {};

  return res;
}

vector<int> euler_path(vector<vector<int>> &g, int first) {
  int n = (int)g.size();
  int v1 = -1, v2 = -1;
  {
    bool bad = false;
    for (int i = 0; i < n; i++)
      if (g[i].size() & 1) {
        if (v1 == -1)
          v1 = i;
        else if (v2 == -1)
          v2 = i;
        else
          bad = true;
      }

    if (bad or (v1 != -1 and v2 == -1)) return {};
  }

  if (v2 != -1) {
    // insert cycle
    g[v1].push_back(v2);
    g[v2].push_back(v1);
  }

  auto res = euler_cycle(g, first);
  if (res.empty()) return res;

  if (v1 != -1) {
    for (int i = 0; i + 1 < (int)res.size(); i++) {
      if ((res[i] == v1 and res[i + 1] == v2) ||
          (res[i] == v2 and res[i + 1] == v1)) {
        vector<int> res2;
        for (int j = i + 1; j < (int)res.size(); j++)
          res2.push_back(res[j]);
        for (int j = 1; j <= i; j++) res2.push_back(res[j]);
        res = res2;
        break;
      }
    }
  }

  reverse(all(res));
  return res;
}
```

4.15 Find Articulation/Cut Points

Given an **undirected** graph find it's articulation points.
articulation point (or cut vertex): is defined as a **vertex** which, when removed along with associated edges, increases thee number of connected components in the graph.
A vertex u can be an articulation point if and only if has at least 2 adjacent vertex
Time: $O(N + M)$

```
const int MAXN(100);
int N;
vi2d G;
int timer;
int tin[MAXN], low[MAXN];
set<int> cpoints;

int dfs(int u, int p = -1) {
  int cnt = 0;
  low[u] = tin[u] = timer++;
  for (auto v : G[u]) {
    if (not tin[v]) {
      cnt++;
      dfs(v, u);

      if (low[v] >= tin[u]) cpoints.insert(u);
      low[u] = min(low[u], low[v]);
    } else if (v != p)
      low[u] = min(low[u], tin[v]);
  }

  return cnt;
}
void getCutPoints() {
  memset(low, 0, sizeof(low));
  memset(tin, 0, sizeof(tin));
  cpoints.clear();

  timer = 1;
  for (int i = 0; i < N; i++) {
    if (tin[i]) continue;
    int cnt = dfs(i);
    if (cnt == 1) cpoints.erase(i);
  }
}
```

4.16 Find Bridge Tree Components

$dfsb(u, p)$ finds the component of the connected coponent of u.
time: $O(n + m)$

```
int n;
const int MAXN(3'00'000);
vi g[MAXN], vi stck;
int tin[MAXN], low[MAXN], comp[MAXN], qtdcomps, clk;
```

```
void dfsb(int u, int p) {
  low[u] = tin[u] = ++clk;
  stck.emplace_back(u);

  for (auto v : g[u]) {
    if (!tin[v]) {
      dfsb(v, u);
      low[u] = min(low[u], low[v]);
    } else if (v != p) {
      low[u] = min(low[u], tin[v]);
    }
  }

  if (low[u] == tin[u]) {
    qtdcomps++;
    int v2;
    do {
      v2 = stck.back();
      comp[v2] = qtdcomps;
      stck.pop_back();
    } while (v2 != u);
  }
}
```

4.17 Find Bridges (online)

```
// O((n+m)*log(n))
struct BridgeFinder {
  // 2ecc = 2 edge conected component
  // cc = conected component
  vector<int> parent, dsu_2ecc, dsu_cc, dsu_cc_size;
  int bridges, lca_iteration;
  vector<int> last_visit;

  BridgeFinder(int n)
    : parent(n, -1),
      dsu_2ecc(n),
      dsu_cc(n),
      dsu_cc_size(n, 1),
      bridges(0),
      lca_iteration(0),
      last_visit(n) {
    for (int i = 0; i < n; i++) {
      dsu_2ecc[i] = i;
      dsu_cc[i] = i;
```

```cpp
    }
}

int find_2ecc(int v) {
  if (v == -1) return -1;
  return dsu_2ecc[v] == v
           ? v
           : dsu_2ecc[v] = find_2ecc(dsu_2ecc[v]);
}

int find_cc(int v) {
  v = find_2ecc(v);
  return dsu_cc[v] == v ? v
                        : dsu_cc[v] = find_cc(dsu_cc[v]);
}

void make_root(int v) {
  v = find_2ecc(v);
  int root = v;
  int child = -1;
  while (v != -1) {
    int p = find_2ecc(parent[v]);
    parent[v] = child;
    dsu_cc[v] = root;
    child = v;
    v = p;
  }
  dsu_cc_size[root] = dsu_cc_size[child];
}

void merge_path(int a, int b) {
  ++lca_iteration;
  vector<int> path_a, path_b;
  int lca = -1;
  while (lca == -1) {
    if (a != -1) {
      a = find_2ecc(a);
      path_a.push_back(a);
      if (last_visit[a] == lca_iteration) {
        lca = a;
        break;
      }
      last_visit[a] = lca_iteration;
      a = parent[a];
    }
    if (b != -1) {
      b = find_2ecc(b);
      path_b.push_back(b);
      if (last_visit[b] == lca_iteration) {
        lca = b;
        break;
      }
      last_visit[b] = lca_iteration;
      b = parent[b];
    }
  }

  for (auto v : path_a) {
    dsu_2ecc[v] = lca;
    if (v == lca) break;
    --bridges;
  }
  for (auto v : path_b) {
    dsu_2ecc[v] = lca;
    if (v == lca) break;
    --bridges;
  }
}

void add_edge(int a, int b) {
  a = find_2ecc(a);
  b = find_2ecc(b);

  if (a == b) return;

  int ca = find_cc(a);
  int cb = find_cc(b);

  if (ca != cb) {
    ++bridges;
    if (dsu_cc_size[ca] > dsu_cc_size[cb]) {
      swap(a, b);
      swap(ca, cb);
    }
    make_root(a);
    parent[a] = dsu_cc[a] = b;
    dsu_cc_size[cb] += dsu_cc_size[a];
  } else {
    merge_path(a, b);
  }
```

```
  }
};
```

4.18 Find Bridges

Find every bridge in a **undirected** connected graph.
bridge: A bridge is defined as an **edge** which, when removed, increases the number of connected components in the graph.
Remember to read the graph as pair where the second is the id of the edge !
Time: $O(N + M)$

```
const int MAXN(10000), MAXM(100000);
int N, M, clk, tin[MAXN], low[MAXN], isBridge[MAXM];
vector<pii> G[MAXN];

void dfs(int u, int p = -1) {
  tin[u] = low[u] = clk++;

  for (auto [v, i] : G[u]) {
    if (v == p) continue;
    if (tin[v]) {
      low[u] = min(low[u], tin[v]);
    } else {
      dfs(v, u);
      low[u] = min(low[u], low[v]);
      if (low[v] > tin[u]) {
        isBridge[i] = 1;
      }
    }
  }
}

void findBridges() {
  fill(tin, tin + N, 0);
  fill(low, low + N, 0);
  fill(isBridge, isBridge + M, 0);
  clk = 1;
  for (int i = 0; i < N; i++) {
    if (!tin[i]) dfs(i);
  }
}
```

4.19 Find Centroid

Given a tree (don't forget to make it 'undirected'), find it's centroids.
Time: $O(V)$

```
void dfs(int u, int p, int n, vi2d &g, vi &sz,
         vi &centroid) {
  sz[u] = 1;

  bool iscentroid = true;
  for (auto v : g[u])
    if (v != p) {
      dfs(v, u, n, g, sz, centroid);
      if (sz[v] > n / 2) iscentroid = false;
      sz[u] += sz[v];
    }

  if (n - sz[u] > n / 2) iscentroid = false;
  if (iscentroid) centroid.eb(u);
}
vi getCentroid(vi2d &g, int n) {
  vi centroid;
  vi sz(n);
  dfs(0, -1, n, g, sz, centroid);
  return centroid;
}
```

4.20 Floyd Warshall

Simply finds the minimal distance for each node to every other node. $O(V^3)$

```
vector<vll> floyd_warshall(const vector<vll> &adj, ll n) {
  auto dist = adj;

  for (int i = 0; i < n; ++i) {
    for (int j = 0; j < n; ++j) {
      for (int k = 0; k < n; ++k) {
        dist[j][k] =
          min(dist[j][k], dist[j][i] + dist[i][k]);
      }
    }
  }
  return dist;
}
```

4.21 Functional/Successor Graph

Given a functional graph find the vertice after k moves starting at u and also the distance between u and v, if it's impossible to reach v starting at u returns -1.
Time: build $O(N \cdot MAXLOG2)$, kth $O(MAXLOG2)$, dist $O(MAXLOG2)$

```
const int MAXN(2'000'000), MAXLOG2(24);
int N;
```

```
vi2d succ(MAXN, vi(MAXLOG2 + 1));
vi dst(MAXN, 0);

int vis[MAXN];
void dfsbuild(int u) {
  if (vis[u]) return;
  vis[u] = 1;
  int v = succ[u][0];
  dfsbuild(v);
  dst[u] = dst[v] + 1;
}

void build() {
  for (int i = 0; i < N; i++) {
    if (not vis[i]) dfsbuild(i);
  }

  for (int k = 1; k <= MAXLOG2; k++) {
    for (int i = 0; i < N; i++) {
      succ[i][k] = succ[succ[i][k - 1]][k - 1];
    }
  }
}

int kth(int u, ll k) {
  if (k <= 0) return u;
  for (int i = 0; i <= MAXLOG2; i++)
    if ((1ll << i) & k) u = succ[u][i];
  return u;
}

int dist(int u, int v) {
  int cu = kth(u, dst[u]);
  if (kth(u, dst[u] - dst[v]) == v)
    return dst[u] - dst[v];
  else if (kth(cu, dst[cu] - dst[v]) == v)
    return dst[u] + (dst[cu] - dst[v]);
  else
    return -1;
}
```

4.22 Graph Cycle (directed)

Given a directed graph finds a cycle (or not).
Time : $O(E)$

```
bool dfs(int v, vi2d &adj, vc &visited, vi &parent,
         vc &color, int &cycle_start, int &cycle_end) {
  color[v] = 1;
  for (int u : adj[v]) {
    if (color[u] == 0) {
      parent[u] = v;
      if (dfs(u, adj, visited, parent, color, cycle_start,
              cycle_end))
        return true;
    } else if (color[u] == 1) {
      cycle_end = v;
      cycle_start = u;
      return true;
    }
  }
  color[v] = 2;
  return false;
}

vi find_cycle(vi2d &g, int n) {
  vc visited(n);
  vi parent(n);
  vc color(n);
  int cycle_start, cycle_end;
  color.assign(n, 0);
  parent.assign(n, -1);
  cycle_start = -1;

  for (int v = 0; v < n; v++) {
    if (color[v] == 0 && dfs(v, g, visited, parent, color,
                             cycle_start, cycle_end))
      break;
  }

  if (cycle_start == -1) {
    return {};
  } else {
    vector<int> cycle;
    cycle.push_back(cycle_start);
    for (int v = cycle_end; v != cycle_start; v = parent[v])
      cycle.push_back(v);
    cycle.push_back(cycle_start);
    reverse(cycle.begin(), cycle.end());
    return cycle;
  }
```

```
}
```

4.23 Graph Cycle (undirected)

Detects if a graph contains a cycle. If path parameter is not null, it will contain the cycle if one exists.
Time: $O(V+E)$

```
void graph_cycles(const vector<vector<int>> &g, int u,
                  int p, vector<int> &ps,
                  vector<int> &color, int &cn,
                  vector<vector<int>> &cycles) {
  if (color[u] == 2) {
    return;
  }

  if (color[u] == 1) {
    cn++;
    int cur = p;
    cycles.emplace_back();
    auto &v = cycles.back();
    v.push_back(cur);
    while (cur != u) {
      cur = ps[cur];
      v.push_back(cur);
    }
    reverse(all(v));
    return;
  }

  ps[u] = p;
  color[u] = 1;
  for (auto v : g[u]) {
    if (v != p)
      graph_cycles(g, v, u, ps, color, cn, cycles);
  }

  color[u] = 2;
}

vector<vector<int>> graph_cycles(
  const vector<vector<int>> &g) {
  vector<int> ps(g.size(), -1), color(g.size());
  int cn = 0;
  vector<vector<int>> cycles;
  for (int i = 0; i < (int)g.size(); i++)
    graph_cycles(g, i, -1, ps, color, cn, cycles);
  return cycles;
}
```

4.24 Heavy Light Decomposition

```
struct HeavyLightDecomposition {
  vector<int> parent, depth, size, heavy, head, pos;

  using SegT = int;
  static SegT op(SegT a, SegT b) { return max(a, b); }
  SegTree<SegT, op> seg;

  HeavyLightDecomposition(const vector<vector<int>> &g,
                          const vector<int> &v,
                          int root = 0)
    : parent(g.size()),
      depth(g.size()),
      size(g.size()),
      heavy(g.size(), -1),
      head(g.size()),
      pos(g.size()),
      seg((int)g.size()) {
    dfs(g, root);
    int cur_pos = 0;
    decompose(g, root, root, cur_pos);

    for (int i = 0; i < (int)g.size(); i++) {
      seg.set(pos[i], v[i]);
    }
  }

  SegT query_path(int a, int b) const {
    int res = 0;
    for (; head[a] != head[b]; b = parent[head[b]]) {
      if (depth[head[a]] > depth[head[b]]) swap(a, b);
      res = op(res, seg.query(pos[head[b]], pos[b]));
    }
    if (depth[a] > depth[b]) swap(a, b);
    return op(res, seg.query(pos[a], pos[b]));
  }

  SegT query_subtree(int a) const {
    return seg.query(pos[a], pos[a] + size[a] - 1);
  }
```

```cpp
  void set(int a, int x) { seg.set(pos[a], x); }

 private:
  void dfs(const vector<vector<int>> &g, int u) {
    size[u] = 1;
    int mx_child_size = 0;
    for (auto x : g[u])
      if (x != parent[u]) {
        parent[x] = u;
        depth[x] = depth[u] + 1;
        dfs(g, x);
        size[u] += size[x];
        if (size[x] > mx_child_size)
          mx_child_size = size[x], heavy[u] = x;
      }
  }

  void decompose(const vector<vector<int>> &g, int u, int h,
                 int &cur_pos) {
    head[u] = h;
    pos[u] = cur_pos++;
    if (heavy[u] != -1) decompose(g, heavy[u], h, cur_pos);

    for (auto x : g[u])
      if (x != parent[u] and x != heavy[u]) {
        decompose(g, x, x, cur_pos);
      }
  }
};
```

4.25 Kruskal

Find the minimum spanning tree of a graph.
Time: $O(E \log E)$
can be used to find the maximum spanning tree by changing the comparison operator in the sort

```cpp
struct UFDS {
  vector<int> ps, sz;
  int components;

  UFDS(int n) : ps(n + 1), sz(n + 1, 1), components(n) {
    iota(all(ps), 0);
  }

  int find_set(int x) {
    return (x == ps[x] ? x : (ps[x] = find_set(ps[x])));
  }

  bool same_set(int x, int y) {
    return find_set(x) == find_set(y);
  }

  void union_set(int x, int y) {
    x = find_set(x);
    y = find_set(y);

    if (x == y) return;

    if (sz[x] < sz[y]) swap(x, y);

    ps[y] = x;
    sz[x] += sz[y];

    components--;
  }
};

vector<tuple<ll, int, int>> kruskal(
  int n, vector<tuple<ll, int, int>> &edges) {
  UFDS ufds(n);
  vector<tuple<ll, int, int>> ans;

  sort(all(edges));
  for (auto [a, b, c] : edges) {
    if (ufds.same_set(b, c)) continue;

    ans.emplace_back(a, b, c);
    ufds.union_set(b, c);
  }

  return ans;
}
```

4.26 Lowest Common Ancestor (Binary Lifting)

Finds the LCA between two nodes using binary lifting
Time: build $O(N \cdot MAXLOG2)$ query $O(MAXLOG2)$

```cpp
const int MAXLOG2 = 20, MAXN(2'00'000);
int N;
int G[MAXN];
int depth[MAXN];
int up[MAXN][MAXLOG2 + 1];
```

```cpp
vi GT[MAXN];

void build(int u = 0) {
  for (int i = 1; i <= MAXLOG2; i++)
    up[u][i] = up[up[u][i - 1]][i - 1];

  for (int v : GT[u])
    if (v != up[u][0]) {
      depth[v] = depth[up[v][0] = u] + 1;
      build(v);
    }
}

int jump(int u, ll k) {
  for (ll i = 0; i <= MAXLOG2; i++)
    if (k & (1ll << i)) u = up[u][i];

  return u;
}

int lca(int a, int b) {
  if (depth[a] < depth[b]) swap(a, b);

  a = jump(a, depth[a] - depth[b]);

  if (b == a) return a;

  for (int i = MAXLOG2; i >= 0; i--) {
    int at = up[a][i], bt = up[b][i];
    if (at != bt) a = at, b = bt;
  }

  return up[a][0];
}
```

4.27 Lowest Common Ancestor

Given two nodes of a tree find their lowest common ancestor, or their distance
Build : $O(V)$, Queries: $O(1)$

```cpp
template <typename T>
struct SparseTable {
  vector<T> v;
  int n;
  static const int b = 30;
  vi mask, t;

  int op(int x, int y) { return v[x] < v[y] ? x : y; }
  int msb(int x) {
    return __builtin_clz(1) - __builtin_clz(x);
  }
  SparseTable() {}
  SparseTable(const vector<T>& v_)
    : v(v_), n(v.size()), mask(n), t(n) {
    for (int i = 0, at = 0; i < n; mask[i++] = at |= 1) {
      at = (at << 1) & ((1 << b) - 1);
      while (at and op(i, i - msb(at & -at)) == i)
        at ^= at & -at;
    }
    for (int i = 0; i < n / b; i++)
      t[i] = b * i + b - 1 - msb(mask[b * i + b - 1]);
    for (int j = 1; (1 << j) <= n / b; j++)
      for (int i = 0; i + (1 << j) <= n / b; i++)
        t[n / b * j + i] =
          op(t[n / b * (j - 1) + i],
             t[n / b * (j - 1) + i + (1 << (j - 1))]);
  }
  int small(int r, int sz = b) {
    return r - msb(mask[r] & ((1 << sz) - 1));
  }
  T query(int l, int r) {
    if (r - l + 1 <= b) return small(r, r - l + 1);
    int ans = op(small(l + b - 1), small(r));
    int x = l / b + 1, y = r / b - 1;
    if (x <= y) {
      int j = msb(y - x + 1);
      ans = op(ans, op(t[n / b * j + x],
                       t[n / b * j + y - (1 << j) + 1]));
    }
    return ans;
  }
};

struct LCA {
  SparseTable<int> st;
  int n;
  vi v, pos, dep;

  LCA(const vi2d& g, int root) : n(len(g)), pos(n) {
    dfs(root, 0, -1, g);
    st = SparseTable<int>(vector<int>(all(dep)));
```

```
  }

  void dfs(int i, int d, int p, const vi2d& g) {
    v.eb(len(dep)) = i, pos[i] = len(dep), dep.eb(d);
    for (auto j : g[i])
      if (j != p) {
        dfs(j, d + 1, i, g);
        v.eb(len(dep)) = i, dep.eb(d);
      }
  }

  int lca(int a, int b) {
    int l = min(pos[a], pos[b]);
    int r = max(pos[a], pos[b]);
    return v[st.query(l, r)];
  }
  int dist(int a, int b) {
    return dep[pos[a]] + dep[pos[b]] -
           2 * dep[pos[lca(a, b)]];
  }
};
```

4.28 Maximum Flow (Edmonds-Karp)

Finds the **maximum flow** in a graph network, given the **source** s and the **sink** t.
Time: $O(V \cdot E^2)$

```
struct maxflow {
  int n;
  vi2d g;
  vll2d cps;
  vi ps;
  vector<vector<char>> isedge;

  maxflow(int _n)
    : n(_n),
      g(n),
      cps(n, vll(n)),
      ps(n),
      isedge(n, vc(n)) {}

  void add(int u, int v, ll c, bool set = true) {
    g[u].emplace_back(v);
    g[v].emplace_back(u);
    cps[u][v] = cps[u][v] * (!set) + c;
    isedge[u][v] = true;
  }

  ll bfs(int s, int t) {
    fill(all(ps), -1);
    ps[s] = -2;
    queue<pair<ll, int>> q;
    q.emplace(oo, s);

    while (!q.empty()) {
      auto [flow, cur] = q.front();
      q.pop();

      for (auto next : g[cur]) {
        if (ps[next] == -1 and cps[cur][next]) {
          ps[next] = cur;
          ll new_flow = min(flow, cps[cur][next]);
          if (next == t) return new_flow;
          q.emplace(new_flow, next);
        }
      }
    }

    return 0ll;
  }

  ll flow(int s, int t) {
    ll flow = 0;
    ll new_flow;

    while ((new_flow = bfs(s, t))) {
      flow += new_flow;
      int cur = t;
      while (cur != s) {
        int prev = ps[cur];
        cps[prev][cur] -= new_flow;
        cps[cur][prev] += new_flow;
        cur = prev;
      }
    }

    return flow;
  }

  vector<pii> get_used() {
    vector<pii> used;
```

```
    for (int i = 0; i < n; i++) {
      for (int j = 0; j < n; j++) {
        if (isedge[i][j] and cps[i][j] == 0)
          used.emplace_back(i, j);
      }
    }
    return used;
  }
};
```

4.29 Minimum Cost Flow

Given a network find the minimum cost to achieve a flow of at most f. Works with **directed** and **undirected** graphs

- **add(u, v, w, c):** adds an edge from u to v with capacity w and cost c.
- **flow(s, t, f):** return a pair $(flow, cost)$ with the maximum flow until f with source at s and sink at t, with the minimum cost possible.

Time : $O(N \cdot M + f \cdot m \log n)$

```
template <typename T>
struct mcmf {
  struct edge {
    int to, rev, flow, cap;
    bool res;  // if it's a reverse edge
    T cost;    // cost per unity of flow
    edge()
      : to(0),
        rev(0),
        flow(0),
        cap(0),
        cost(0),
        res(false) {}
    edge(int to_, int rev_, int flow_, int cap_, T cost_,
         bool res_)
      : to(to_),
        rev(rev_),
        flow(flow_),
        cap(cap_),
        res(res_),
        cost(cost_) {}
  };

  vector<vector<edge>> g;
  vector<int> par_idx, par;
  T inf;
  vector<T> dist;

  mcmf(int n)
    : g(n),
      par_idx(n),
      par(n),
      inf(numeric_limits<T>::max() / 3) {}

  void add(int u, int v, int w, T cost) {
    edge a = edge(v, g[v].size(), 0, w, cost, false);
    edge b = edge(u, g[u].size(), 0, 0, -cost, true);

    g[u].push_back(a);
    g[v].push_back(b);
  }

  vector<T> spfa(int s) {  // don't code it if there isn't
                           // negative cycles
    deque<int> q;
    vector<bool> is_inside(g.size(), 0);
    dist = vector<T>(g.size(), inf);

    dist[s] = 0;
    q.push_back(s);
    is_inside[s] = true;

    while (!q.empty()) {
      int v = q.front();
      q.pop_front();
      is_inside[v] = false;

      for (int i = 0; i < g[v].size(); i++) {
        auto [to, rev, flow, cap, res, cost] = g[v][i];
        if (flow < cap and dist[v] + cost < dist[to]) {
          dist[to] = dist[v] + cost;

          if (is_inside[to]) continue;
          if (!q.empty() and dist[to] > dist[q.front()])
            q.push_back(to);
          else
            q.push_front(to);
          is_inside[to] = true;
        }
      }
    }
    return dist;
```

```cpp
}
bool dijkstra(int s, int t, vector<T>& pot) {
  priority_queue<pair<T, int>, vector<pair<T, int>>,
                 greater<>>
    q;
  dist = vector<T>(g.size(), inf);
  dist[s] = 0;
  q.emplace(0, s);
  while (q.size()) {
    auto [d, v] = q.top();
    q.pop();
    if (dist[v] < d) continue;
    for (int i = 0; i < g[v].size(); i++) {
      auto [to, rev, flow, cap, res, cost] = g[v][i];
      cost += pot[v] - pot[to];
      if (flow < cap and dist[v] + cost < dist[to]) {
        dist[to] = dist[v] + cost;
        q.emplace(dist[to], to);
        par_idx[to] = i, par[to] = v;
      }
    }
  }
  return dist[t] < inf;
}

pair<int, T> min_cost_flow(int s, int t, int flow = inf) {
  vector<T> pot(g.size(), 0);
  pot = spfa(s);  // comment this line if there isn't
                  // negative cycles

  int f = 0;
  T ret = 0;
  while (f < flow and dijkstra(s, t, pot)) {
    for (int i = 0; i < g.size(); i++)
      if (dist[i] < inf) pot[i] += dist[i];

    int mn_flow = flow - f, u = t;
    while (u != s) {
      mn_flow =
        min(mn_flow, g[par[u]][par_idx[u]].cap -
                       g[par[u]][par_idx[u]].flow);
      u = par[u];
    }

    ret += pot[t] * mn_flow;
```

```cpp
      u = t;
      while (u != s) {
        g[par[u]][par_idx[u]].flow += mn_flow;
        g[u][g[par[u]][par_idx[u]].rev].flow -= mn_flow;
        u = par[u];
      }

      f += mn_flow;
    }

    return make_pair(f, ret);
  }
};
```

4.30 Minimum Cut (unweighted)

After build the **direct/undirected** graph find the minimum of edges needed to be removed to make the sink t unreachable from the source s.
Time: $O(V \cdot E^2)$

```cpp
struct Mincut {
  int n;
  vi2d g;
  vii edges;
  vll2d capacity;
  vi ps, vis;

  Mincut(int _n)
    : n(_n), g(n), capacity(n, vll(n)), ps(n), vis(n) {}

  void add(int u, int v, ll c = 1, bool directed = false,
           bool set = false) {
    edges.emplace_back(u, v);
    g[u].emplace_back(v);

    if (not set)
      capacity[u][v] += c;
    else
      capacity[u][v] = c;

    if (not directed) {
      g[v].emplace_back(u);

      if (not set)
        capacity[v][u] += c;
```

```
    else
      capacity[v][u] = c;
  }
}

ll bfs(int s, int t) {
  fill(all(ps), -1);
  ps[s] = -2;
  queue<pair<ll, int>> q;
  q.push({oo, s});

  while (!q.empty()) {
    auto [flow, cur] = q.front();
    q.pop();

    for (auto next : g[cur]) {
      if (ps[next] == -1 and capacity[cur][next]) {
        ps[next] = cur;
        ll new_flow = min(flow, capacity[cur][next]);
        if (next == t) return new_flow;
        q.push({new_flow, next});
      }
    }
  }

  return 0ll;
}

ll maxflow(int s, int t) {
  ll flow = 0;
  ll new_flow;

  while ((new_flow = bfs(s, t))) {
    flow += new_flow;
    int cur = t;
    while (cur != s) {
      int prev = ps[cur];
      capacity[prev][cur] -= new_flow;
      capacity[cur][prev] += new_flow;
      cur = prev;
    }
  }

  return flow;
}
```

```
  void dfs(int u) {
    vis[u] = true;

    for (auto v : g[u]) {
      if (capacity[u][v] > 0 and not vis[v]) {
        dfs(v);
      }
    }
  }

  vii mincut(int s, int t) {
    maxflow(s, t);
    fill(all(vis), 0);
    dfs(s);

    vii removed;
    for (auto &[u, v] : edges) {
      if ((vis[u] and not vis[v]) or
          (vis[v] and not vis[u]))
        removed.emplace_back(u, v);
    }

    return removed;
  }
};
```

4.31 Prim (MST)

Given a graph with N vertex finds the minimum spanning tree, if there is no such three returns inf, it starts using the edges that connect with each $s_i \in s$, if none is provided than it starts with the edges of node 0. Time: $O(V \log E)$

```
const int MAXN(1'00'000);
int N;
vector<pair<ll, int>> G[MAXN];
ll prim(vi s = vi(1, 0)) {
  priority_queue<pair<ll, int>, vector<pair<ll, int>>,
                 greater<pair<ll, int>>>
    pq;
  vector<char> ingraph(MAXN);
  int ingraphcnt(0);
  for (auto si : s) {
    ingraphcnt++;
    ingraph[si] = true;
    for (auto &[w, v] : G[si]) pq.emplace(w, v);
  }
```

```cpp
  ll mstcost = 0;
  while (ingraphcnt < N and !pq.empty()) {
    ll w;
    int v;

    do {
      tie(w, v) = pq.top();
      pq.pop();
    } while (not pq.empty() and ingraph[v]);

    mstcost += w, ingraph[v] = true, ingraphcnt++;

    for (auto &[w2, v2] : G[v]) {
      pq.emplace(w2, v2);
    }
  }

  return ingraphcnt == N ? mstcost : oo;
}
```

4.32 Small to Large

Answer queries of the form "How many vertices in the subtree of vertex v have property P?"
* this implementation answers how many distinct $values[i]$ are in the subtree starting at u.
Build: $O(N)$, Query: $O(N \log N)$

```cpp
struct SmallToLarge {
  int n;
  vi2d tree, vis_childs;
  vi sizes, values, ans;
  set<int> cnt;

  SmallToLarge(vi2d &g, vi &v)
    : tree(g),
      vis_childs(len(g)),
      sizes(len(g)),
      values(v),
      ans(len(g)) {
    get_size(0);
    dfs(0);
  }

  inline void add_value(int u) { cnt.insert(values[u]); }

  inline void remove_value(int u) { cnt.erase(values[u]); }

  inline void update_ans(int u) { ans[u] = len(cnt); }

  void dfs(int u, int p = -1, bool keep = true) {
    int mx = -1;
    for (auto x : tree[u]) {
      if (x == p) continue;

      if (mx == -1 or sizes[mx] < sizes[x]) mx = x;
    }

    for (auto x : tree[u]) {
      if (x != p and x != mx) dfs(x, u, false);
    }

    if (mx != -1) {
      dfs(mx, u, true);
      swap(vis_childs[u], vis_childs[mx]);
    }

    vis_childs[u].push_back(u);
    add_value(u);

    for (auto x : tree[u]) {
      if (x != p and x != mx) {
        for (auto y : vis_childs[x]) {
          add_value(y);
          vis_childs[u].push_back(y);
        }
      }
    }

    update_ans(u);

    if (!keep) {
      for (auto x : vis_childs[u]) remove_value(x);
    }
  }

  void get_size(int u, int p = -1) {
    sizes[u] = 1;
    for (auto x : tree[u])
      if (x != p) {
        get_size(x, u);
        sizes[u] += sizes[x];
```

```
    }
  }
};
```

4.33 Successor Graph-(struct)

```
struct SuccessorGraph {
  vector<vector<int>> paths;
  vector<int> path_num, pos;
  vector<char> is_cycle;

  SuccessorGraph(const vector<int> &v)
    : path_num(v.size()), pos(v.size()) {
    paths.reserve(v.size());
    is_cycle.reserve(v.size());

    vector<char> vis(v.size());
    for (auto i : topological_order(v)) {
      if (vis[i]) continue;

      vector<int> path;
      int cur;
      for (cur = i; not vis[cur]; cur = v[cur]) {
        path.push_back(cur);
        vis[cur] = 1;
      }

      int cycle_start = 0;
      for (; cycle_start < (int)path.size() and
             path[cycle_start] != cur;
           cycle_start++)
        ;

      if (cycle_start > 0) {
        paths.emplace_back();
        for (int j = 0; j < cycle_start; j++) {
          paths.back().push_back(path[j]);
          pos[path[j]] = j;
          path_num[path[j]] = (int)paths.size() - 1;
        }
        paths.back().push_back(cur);
        is_cycle.push_back(false);
      }

      if (cycle_start < (int)path.size()) {
        paths.emplace_back();
        for (int j = cycle_start; j < (int)path.size();
             j++) {
          paths.back().push_back(path[j]);
          pos[path[j]] = j - cycle_start;
          path_num[path[j]] = (int)paths.size() - 1;
        }
        is_cycle.push_back(true);
      }
    }
  }

  const vector<int> &path(int cur) const {
    return paths[path_num[cur]];
  }

  int kth_pos(int cur, ll k) const {
    while (not is_cycle[path_num[cur]]) {
      auto &p = path(cur);
      int remain = (int)p.size() - pos[cur] - 1;
      if (k <= remain) return p[pos[cur] + k];
      cur = p.back();
      k -= remain;
    }

    auto &p = path(cur);
    return p[(pos[cur] + k) % p.size()];
  }

  // {element, number_of_moves}
  pair<int, int> go_to_cycle(int cur) const {
    int moves = 0;
    while (not is_cycle[path_num[cur]]) {
      auto &p = path(cur);
      moves += (int)p.size() - pos[cur] - 1;
      cur = p.back();
    }
    return {cur, moves};
  }

  // min cost to reach dest from cur
  int reach(int cur, int dest) const {
    int moves = 0;
    while (not is_cycle[path_num[cur]] and
           path_num[cur] != path_num[dest]) {
```

```cpp
      auto &p = path(cur);
      moves += (int)p.size() - pos[cur] - 1;
      cur = p.back();
    }

    if (path_num[cur] != path_num[dest]) return -1;

    if (pos[cur] <= pos[dest])
      return moves + pos[dest] - pos[cur];

    if (not is_cycle[path_num[cur]]) return -1;

    return moves + pos[dest] + (int)path(cur).size() -
           pos[cur];
  }

 private:
  void topological_order(const vector<int> &g,
                         vector<char> &vis,
                         vector<int> &order, int u) {
    vis[u] = true;
    if (not vis[g[u]])
      topological_order(g, vis, order, g[u]);
    order.push_back(u);
  }

  vector<int> topological_order(const vector<int> &g) {
    vector<char> vis(g.size(), false);
    vector<int> order;
    for (auto i = 0; i < (int)g.size(); i++)
      if (not vis[i]) topological_order(g, vis, order, i);
    reverse(order.begin(), order.end());
    return order;
  }
};
```

4.34 Sum every node distance

Given a **tree**, for each node i find the sum of distance from i to every other node.
don't forget to set the tree as undirected, that's needed to choose an arbitrary root
Time: $O(N)$

```cpp
void getRoot(int u, int p, vi2d &g, vll &d, vll &cnt) {
  for (int i = 0; i < len(g[u]); i++) {
    int v = g[u][i];
    if (v == p) continue;
    getRoot(v, u, g, d, cnt);
    d[u] += d[v] + cnt[v];
    cnt[u] += cnt[v];
  }
}

void dfs(int u, int p, vi2d &g, vll &cnt, vll &ansd,
         int n) {
  for (int i = 0; i < len(g[u]); i++) {
    int v = g[u][i];
    if (v == p) continue;

    ansd[v] = ansd[u] - cnt[v] + (n - cnt[v]);
    dfs(v, u, g, cnt, ansd, n);
  }
}
vll fromToAll(vi2d &g, int n) {
  vll d(n);
  vll cnt(n, 1);
  getRoot(0, -1, g, d, cnt);

  vll ansdist(n);
  ansdist[0] = d[0];

  dfs(0, -1, g, cnt, ansdist, n);
  return ansdist;
}
```

4.35 Topological Labelling (Kahn)

The same thing as topological sorting but over every possible order gives lexicographically smaller
Time: $O(E + V \cdot \log V)$

```cpp
const int MAXN(1'00'000);
int OUTCNT[MAXN];
vi2d GIN(MAXN);
int N;

vi toposort() {
  vi order;
  priority_queue<int> q;

  for (int i = 0; i < N; i++)
    if (!OUTCNT[i]) q.emplace(i);

  while (!q.empty()) {
```

```
    auto u = q.top();
    q.pop();
    order.emplace_back(u);
    for (auto v : GIN[u]) {
      OUTCNT[v]--;
      if (OUTCNT[v] == 0) q.emplace(v);
    }
  }

  reverse(all(order));
  return len(order) == N ? order : vi();
}
```

4.36 Topological Sorting (Kahn)

Finds the topological sorting in a **DAG**, if the given graph is not a **DAG** than an empty vector is returned, need to 'initialize' the **INCNT** as you build the graph.

Time: $O(V+E)$

```
const int MAXN(2'00'000);
int INCNT[MAXN];
vi2d GOUT(MAXN);
int N;

vi toposort() {
  vi order;
  queue<int> q;

  for (int i = 0; i < N; i++)
    if (!INCNT[i]) q.emplace(i);

  while (!q.empty()) {
    auto u = q.front();
    q.pop();
    order.emplace_back(u);
    for (auto v : GOUT[u]) {
      INCNT[v]--;
      if (INCNT[v] == 0) q.emplace(v);
    }
  }

  return len(order) == N ? order : vi();
}
```

4.37 Topological Sorting (Tarjan)

Finds a the topological order for the graph, if there is no such order it means the graph is cyclic, then it returns an empty vector

$O(V+E)$

```
const int maxn(1'00'000);
int n, m;
vi g[maxn];

int not_found = 0, found = 1, processed = 2;
int state[maxn];

bool dfs(int u, vi &order) {
  if (state[u] == processed) return true;
  if (state[u] == found) return false;

  state[u] = found;

  for (auto v : g[u]) {
    if (not dfs(v, order)) return false;
  }

  state[u] = processed;
  order.emplace_back(u);

  return true;
}

vi topo_sort() {
  vi order;
  memset(state, 0, sizeof state);

  for (int u = 0; u < n; u++) {
    if (state[u] == not_found and not dfs(u, order))
      return {};
  }

  reverse(all(order));
  return order;
}
```

4.38 Tree Diameter (DP)

```
const int MAXN(1'000'000);
int N;
```

```
vi G[MAXN];

int diameter, toLeaf[MAXN];
void calcDiameter(int u = 0, int p = -1) {
  int d1, d2;
  d1 = d2 = -1;

  for (auto v : G[u]) {
    if (v != p) {
      calcDiameter(v, u);
      d1 = max(d1, toLeaf[v]);
      tie(d1, d2) = minmax({d1, d2});
    }
  }
  toLeaf[u] = d2 + 1;
  diameter = max(diameter, d1 + d2 + 2);
}
```

4.39 Tree Isomorphism (not rooted)

Two trees are considered **isomorphic** if the hash given by *thash*() is the same.
Time: $O(V \cdot \log V)$

```
map<vi, int> mphash;

struct Tree {
  int n;
  vi2d g;
  vi sz, cs;

  Tree(int n_) : n(n_), g(n), sz(n) {}

  void add_edge(int u, int v) {
    g[u].emplace_back(v);
    g[v].emplace_back(u);
  }

  void dfs_centroid(int v, int p) {
    sz[v] = 1;
    bool cent = true;
    for (int u : g[v])
      if (u != p) {
        dfs_centroid(u, v);
        sz[v] += sz[u];
        cent &= not(sz[u] > n / 2);
      }
    if (cent and n - sz[v] <= n / 2) cs.push_back(v);
  }

  int fhash(int v, int p) {
    vi h;
    for (int u : g[v])
      if (u != p) h.push_back(fhash(u, v));
    sort(all(h));
    if (!mphash.count(h)) mphash[h] = mphash.size();
    return mphash[h];
  }

  ll thash() {
    cs.clear();
    dfs_centroid(0, -1);
    if (cs.size() == 1) return fhash(cs[0], -1);
    ll h1 = fhash(cs[0], cs[1]), h2 = fhash(cs[1], cs[0]);
    return (min(h1, h2) << 30ll) + max(h1, h2);
  }
};
```

4.40 Tree Isomorphism (rooted)

Given a rooted tree find the hash of each subtree, if two roots of two distinct trees have the same hash they are considered isomorphic
hash first time in $O(\log N_v \cdot N_v)$ where (N_v) is the of the subtree of v

```
map<vi, int> hasher;
int hs = 0;
struct RootedTreeIso {
  int n;
  vi2d adj;
  vi hashes;
  RootedTreeIso(int _n) : n(_n), adj(_n), hashes(_n, -1){};

  void add_edge(int u, int v) {
    adj[u].emplace_back(v);
    adj[v].emplace_back(u);
  }

  int hash(int u, int p = -1) {
    if (hashes[u] != -1) return hashes[u];

    vi children;
    for (auto v : adj[u])
      if (v != p) children.emplace_back(hash(v, u));
```

```
    sort(all(children));
    if (!hasher.count(children)) hasher[children] = hs++;

    return hashes[u] = hasher[children];
  }
};
```

4.41 Tree Maximum Distance

Returns the maximum distance from every node to any other node in the tree.
$O(6V) = O(V)$

```
pll mostDistantFrom(const vector<vll> &adj, ll n, ll root) {
  // O(V)
  // 0 indexed
  ll mostDistantNode = root;
  ll nodeDistance = 0;
  queue<pll> q;
  vector<char> vis(n);
  q.emplace(root, 0);
  vis[root] = true;
  while (!q.empty()) {
    auto [node, dist] = q.front();
    q.pop();
    if (dist > nodeDistance) {
      nodeDistance = dist;
      mostDistantNode = node;
    }
    for (auto u : adj[node]) {
      if (!vis[u]) {
        vis[u] = true;
        q.emplace(u, dist + 1);
      }
    }
  }
  return {mostDistantNode, nodeDistance};
}

ll twoNodesDist(const vector<vll> &adj, ll n, ll a, ll b) {
  queue<pll> q;
  vector<char> vis(n);
  q.emplace(a, 0);
  while (!q.empty()) {
    auto [node, dist] = q.front();
    q.pop();
    if (node == b) return dist;
    for (auto u : adj[node]) {
      if (!vis[u]) {
        vis[u] = true;
        q.emplace(u, dist + 1);
      }
    }
  }
  return -1;
}

tuple<ll, ll, ll> tree_diameter(const vector<vll> &adj,
                                ll n) {
  // returns two points of the diameter and the diameter
  // itself
  auto [node1, dist1] = mostDistantFrom(adj, n, 0);  // O(V)
  auto [node2, dist2] =
    mostDistantFrom(adj, n, node1);  // O(V)
  auto diameter =
    twoNodesDist(adj, n, node1, node2);  // O(V)
  return make_tuple(node1, node2, diameter);
}

vll everyDistanceFromNode(const vector<vll> &adj, ll n,
                          ll root) {
  // Single Source Shortest Path, from a given root
  queue<pair<ll, ll>> q;
  vll ans(n, -1);
  ans[root] = 0;
  q.emplace(root, 0);
  while (!q.empty()) {
    auto [u, d] = q.front();
    q.pop();

    for (auto w : adj[u]) {
      if (ans[w] != -1) continue;
      ans[w] = d + 1;
      q.emplace(w, d + 1);
    }
  }
  return ans;
}

vll maxDistances(const vector<vll> &adj, ll n) {
  auto [node1, node2, diameter] =
    tree_diameter(adj, n);  // O(3V)
```

```
  auto distances1 =
    everyDistanceFromNode(adj, n, node1);  // O(V)
  auto distances2 =
    everyDistanceFromNode(adj, n, node2);  // O(V)
  vll ans(n);
  for (int i = 0; i < n; ++i)
    ans[i] = max(distances1[i], distances2[i]);  // O(V)
  return ans;
}
```

4.42 Tree Flatten

```
void tree_flatten(const vector<vector<int>> &g, int u,
                  int p, vector<int> &pre, vector<int> &pos,
                  int &idx) {
  ++idx;
  pre.push_back(u);
  for (auto x : g[u])
    if (x != p) tree_flatten(g, x, u, pre, pos, idx);
  pos[u] = idx;
}

pair<vector<int>, vector<int>> tree_flatten(
  const vector<vector<int>> &g, int root = 0) {
  vector<int> first(g.size()), last(g.size()), pre;
  int timer = -1;
  tree_flatten(g, root, -1, pre, last, timer);
  for (int i = 0; i < (int)g.size(); i++) first[pre[i]] = i;
  return {first, last};
}
```

5 Math

5.1 GCD (with factorization)

$O(\sqrt{n})$ due to factorization.

```
ll gcd_with_factorization(ll a, ll b) {
  map<ll, ll> fa = factorization(a);
  map<ll, ll> fb = factorization(b);
  ll ans = 1;
  for (auto fai : fa) {
    ll k = min(fai.second, fb[fai.first]);
    while (k--) ans *= fai.first;
  }
  return ans;
}
```

5.2 GCD

```
ll gcd(ll a, ll b) { return b ? gcd(b, a % b) : a; }
```

5.3 LCM (with factorization)

$O(\sqrt{n})$ due to factorization.

```
ll lcm_with_factorization(ll a, ll b) {
  map<ll, ll> fa = factorization(a);
  map<ll, ll> fb = factorization(b);
  ll ans = 1;
  for (auto fai : fa) {
    ll k = max(fai.second, fb[fai.first]);
    while (k--) ans *= fai.first;
  }
  return ans;
}
```

5.4 LCM

```
ll gcd(ll a, ll b) { return b ? gcd(b, a % b) : a; }
ll lcm(ll a, ll b) { return a / gcd(a, b) * b; }
```

5.5 Arithmetic Progression Sum

- s : first term
- d : common difference
- n : number of terms

```
ll arithmeticProgressionSum(ll s, ll d, ll n) {
  return (s + (s + d * (n - 1))) * n / 2ll;
}
```

5.6 Binomial MOD

Precompute every factorial until $maxn$ ($O(maxn)$) allowing to answer the $\binom{n}{k}$ in $O(\log mod)$ time, due to the fastpow. Note that it needs $O(maxn)$ in memory.

```
const ll MOD = 1e9 + 7;
const ll maxn = 2 * 1e6;
vll fats(maxn + 1, -1);
void precompute() {
  fats[0] = 1;
  for (ll i = 1; i <= maxn; i++) {
    fats[i] = (fats[i - 1] * i) % MOD;
```

```
  }
}

ll fpow(ll a, ll n, ll mod = LLONG_MAX) {
  if (n == 0ll) return 1ll;
  if (n == 1ll) return a;
  ll x = fpow(a, n / 2ll, mod) % mod;
  return ((x * x) % mod * (n & 1ll ? a : 1ll)) % mod;
}

ll binommod(ll n, ll k) {
  ll upper = fats[n];
  ll lower = (fats[k] * fats[n - k]) % MOD;
  return (upper * fpow(lower, MOD - 2ll, MOD)) % MOD;
}
```

5.7 Binomial

O(nm) time, O(m) space
Equal to n choose k

```
ll binom(ll n, ll k) {
  if (k > n) return 0;
  vll dp(k + 1, 0);
  dp[0] = 1;
  for (ll i = 1; i <= n; i++)
    for (ll j = k; j > 0; j--) dp[j] = dp[j] + dp[j - 1];
  return dp[k];
}
```

5.8 Chinese Remainder Theorem

Finds the solution x to the n modular equations.

$$\begin{gathered} x \equiv a_1 (mod m_1) \\ \ldots \\ x \equiv a_n (mod m_n) \end{gathered} \tag{1}$$

The m_i don't need to be coprime, if there is no solution then it returns -1.

```
tuple<ll, ll, ll> ext_gcd(ll a, ll b) {
  if (!a) return {b, 0, 1};
  auto [g, x, y] = ext_gcd(b % a, a);
  return {g, y - b / a * x, x};
}

template <typename T = ll>
struct crt {
  T a, m;

  crt() : a(0), m(1) {}
  crt(T a_, T m_) : a(a_), m(m_) {}
  crt operator*(crt C) {
    auto [g, x, y] = ext_gcd(m, C.m);
    if ((a - C.a) % g != 0) a = -1;
    if (a == -1 or C.a == -1) return crt(-1, 0);
    T lcm = m / g * C.m;
    T ans = a + (x * (C.a - a) / g % (C.m / g)) * m;
    return crt((ans % lcm + lcm) % lcm, lcm);
  }
};

template <typename T = ll>
struct Congruence {
  T a, m;
};

template <typename T = ll>
T chinese_remainder_theorem(
  const vector<Congruence<T>> &equations) {
  crt<T> ans;

  for (auto &[a_, m_] : equations) {
    ans = ans * crt<T>(a_, m_);
  }

  return ans.a;
}
```

5.9 Derangement / Matching Problem

Computes the derangement of N, which is given by the formula :
$D_N = N! \left(1 - \frac{1}{1!} + \frac{1}{2!} - \frac{1}{3!} + \ldots + (-1)^N \frac{1}{N!}\right)$
time: $O(N)$

```
#warning Remember to call precompute !
const ll MOD = 1e9 + 7;
const int MAXN(1'000'000);
ll fats[MAXN + 1];
void precompute() {
  fats[0] = 1;
  for (ll i = 1; i <= MAXN; i++) {
    fats[i] = (fats[i - 1] * i) % MOD;
  }
```

```
}

ll fastpow(ll a, ll p, ll m) {
  ll ret = 1;
  while (p) {
    if (p & 1) ret = (ret * a) % MOD;
    p >>= 1;
    a = (a * a) % MOD;
  }
  return ret;
}

ll divmod(ll a, ll b) {
  return (a * fastpow(b, MOD - 2, MOD)) % MOD;
}

ll derangement(const ll n) {
  ll ans = fats[n];
  for (ll i = 1; i <= n; i++) {
    ll k = divmod(fats[n], fats[i]);
    if (i & 1) {
      ans = (ans - k + MOD) % MOD;
    } else {
      ans = (ans + k) % MOD;
    }
  }
  return ans;
}
```

5.10 Euler phi $\varphi(n)$ (in range)

Computes the number of positive integers less than n that are coprimes with n, in the range $[1, n]$, in $O(N \log N)$.

```
const int MAX = 1e6;
vi range_phi(int n) {
  bitset<MAX> sieve;
  vi phi(n + 1);

  iota(phi.begin(), phi.end(), 0);
  sieve.set();

  for (int p = 2; p <= n; p += 2) phi[p] /= 2;
  for (int p = 3; p <= n; p += 2) {
    if (sieve[p]) {
      for (int j = p; j <= n; j += p) {
        sieve[j] = false;
        phi[j] /= p;
        phi[j] *= (p - 1);
      }
    }
  }

  return phi;
}
```

5.11 Euler phi $\varphi(n)$

Computes the number of positive integers less than n that are coprimes with n, in $O(\sqrt{N})$.

```
int phi(int n) {
  if (n == 1) return 1;

  auto fs = factorization(n);  // a vctor of pair or a map
  auto res = n;

  for (auto [p, k] : fs) {
    res /= p;
    res *= (p - 1);
  }

  return res;
}
```

5.12 Factorial Factorization

Computes the factorization of $n!$ in $\pi(N) * \log n$

```
// O(logN)
ll E(ll n, ll p) {
  ll k = 0, b = p;
  while (b <= n) {
    k += n / b;
    b *= p;
  }
  return k;
}

// O(pi(N)*logN)
map<ll, ll> factorial_factorization(ll n,
                                    const vll &primes) {
  map<ll, ll> fs;
```

```
  for (const auto &p : primes) {
    if (p > n) break;
    fs[p] = E(n, p);
  }
  return fs;
}
```

5.13 Factorial

```
const ll MAX = 18;
vll fv(MAX, -1);
ll factorial(ll n) {
  if (fv[n] != -1) return fv[n];
  if (n == 0) return 1;
  return n * factorial(n - 1);
}
```

5.14 Factorization (Pollard Rho)

Factorizes a number into its prime factors in $O(n^{(\frac{1}{4})} * \log(n))$.

```
ll mul(ll a, ll b, ll m) {
  ll ret = a * b - (ll)((ld)1 / m * a * b + 0.5) * m;
  return ret < 0 ? ret + m : ret;
}

ll pow(ll a, ll b, ll m) {
  ll ans = 1;
  for (; b > 0; b /= 2ll, a = mul(a, a, m)) {
    if (b % 2ll == 1) ans = mul(ans, a, m);
  }
  return ans;
}

bool prime(ll n) {
  if (n < 2) return 0;
  if (n <= 3) return 1;
  if (n % 2 == 0) return 0;

  ll r = __builtin_ctzll(n - 1), d = n >> r;
  for (int a :
       {2, 325, 9375, 28178, 450775, 9780504, 795265022}) {
    ll x = pow(a, d, n);
    if (x == 1 or x == n - 1 or a % n == 0) continue;

    for (int j = 0; j < r - 1; j++) {
      x = mul(x, x, n);
      if (x == n - 1) break;
    }
    if (x != n - 1) return 0;
  }
  return 1;
}

ll rho(ll n) {
  if (n == 1 or prime(n)) return n;
  auto f = [n](ll x) { return mul(x, x, n) + 1; };

  ll x = 0, y = 0, t = 30, prd = 2, x0 = 1, q;
  while (t % 40 != 0 or gcd(prd, n) == 1) {
    if (x == y) x = ++x0, y = f(x);
    q = mul(prd, abs(x - y), n);
    if (q != 0) prd = q;
    x = f(x), y = f(f(y)), t++;
  }
  return gcd(prd, n);
}

vll fact(ll n) {
  if (n == 1) return {};
  if (prime(n)) return {n};
  ll d = rho(n);
  vll l = fact(d), r = fact(n / d);
  l.insert(l.end(), r.begin(), r.end());
  return l;
}
```

5.15 Factorization

Computes the factorization of n in $O(\sqrt{n})$.

```
map<ll, ll> factorization(ll n) {
  map<ll, ll> ans;
  for (ll i = 2; i * i <= n; i++) {
    ll count = 0;
    for (; n % i == 0; count++, n /= i)
      ;
    if (count) ans[i] = count;
  }
  if (n > 1) ans[n]++;
  return ans;
}
```

5.16 Fast pow

Computes $a^b \pmod m$ in $O(\log N)$.

```
ll fpow(ll a, ll b, ll m) {
  ll ret = 1;
  while (b) {
    if (b & 1) ret = (ret * a) % m;
    b >>= 1;
    a = (a * a) % m;
  }
  return ret;
}

ll fpow(ll a, ll b, ll m) {
  if (!b) return 1;
  ll ans = fpow2((a * a) % m, b / 2ll, m);
  return b & 1 ? (a * ans) % m : ans;
}
```

5.17 FFT Convolution

Performs convolution in a vector duh !

```
const ld PI = acos(-1);

/* change the ld to doulbe may increase performance =D */
struct num {
  ld a{0.0}, b{0.0};
  num() {}
  num(ld na) : a{na} {}
  num(ld na, ld nb) : a{na}, b{nb} {}

  const num operator+(const num& c) const {
    return num(a + c.a, b + c.b);
  }
  const num operator-(const num& c) const {
    return num(a - c.a, b - c.b);
  }
  const num operator*(const num& c) const {
    return num(a * c.a - b * c.b, a * c.b + b * c.a);
  }
  const num operator/(const ll& c) const {
    return num(a / c, b / c);
  }
};
```

```
void fft(vector<num>& a, bool invert) {
  int n = len(a);
  for (int i = 1, j = 0; i < n; i++) {
    int bit = n >> 1;
    for (; j & bit; bit >>= 1) j ^= bit;
    j ^= bit;
    if (i < j) swap(a[i], a[j]);
  }
  for (int sz = 2; sz <= n; sz <<= 1) {
    ld ang = 2 * PI / sz * (invert ? -1 : 1);
    num wsz(cos(ang), sin(ang));
    for (int i = 0; i < n; i += sz) {
      num w(1);
      rep(j, 0, sz / 2) {
        num u = a[i + j], v = a[i + j + sz / 2] * w;
        a[i + j] = u + v;
        a[i + j + sz / 2] = u - v;
        w = w * wsz;
      }
    }
  }
  if (invert)
    for (num& x : a) x = x / n;
}

vi conv(vi const a, vi const b) {
  vector<num> fa(all(a));
  vector<num> fb(all(b));
  int n = 1;
  while (n < len(a) + len(b)) n <<= 1;
  fa.resize(n);
  fb.resize(n);
  fft(fa, false);
  fft(fb, false);
  rep(i, 0, n) fa[i] = fa[i] * fb[i];
  fft(fa, true);
  vi result(n);
  rep(i, 0, n) result[i] = round(fa[i].a);
  while (len(result) and result.back() == 0)
    result.pop_back();

  /* Unconment this line if you want a boolean convolution*/
  for (auto& xi : result) xi = min(xi, 1ll);

  return result;
```

```
}
vll poly_exp(vll& ps, int k) {
  vll ret(len(ps));
  auto base = ps;
  ret[0] = 1;

  while (k) {
    if (k & 1) ret = conv(ret, base);
    k >>= 1;
    base = conv(base, base);
  }

  return ret;
}
```

5.18 Find Multiplicative Inverse

```
ll inv(ll a, ll m) {
  return a > 1ll ? m - inv(m % a, a) * m / a : 1ll;
}
```

5.19 Linear Diophantine Equation: Find any solution

Given a b, c finds the solution to the equation $ax + by = c$, the result will be stored in the reference variables $x0$ and $y0$
time: $O(\log min(a, b))$

```
template <typename T>
tuple<T, T, T> ext_gcd(T a, T b) {
  if (b == 0) return {a, 1, 0};

  auto [d, x1, y1] = ext_gcd(b, a % b);

  return {d, y1, x1 - y1 * (a / b)};
}

template <typename T>
tuple<bool, T, T> find_any_solution(T a, T b, T c) {
  assert(a != 0 or b != 0);
#warning Be careful with overflow, use __int128 if needed !

  auto [d, x0, y0] =
    ext_gcd(a < 0 ? -a : a, b < 0 ? -b : b);

  if (c % d) return {false, 0, 0};

  x0 *= c / d;
  y0 *= c / d;
  if (a < 0) x0 = -x0;
  if (b < 0) y0 = -y0;

  return {true, x0, y0};
}

// optional if you want to use __int128
void print(__int128 x) {
  if (x < 0) {
    cout << '-';
    x = -x;
  }
  if (x > 9) print(x / 10);
  cout << (char)((x % 10) + '0');
}

__int128 read() {
  string s;
  cin >> s;
  __int128 x = 0;
  for (auto c : s) {
    if (c != '-') x += c - '0';
    x *= 10;
  }
  x /= 10;
  if (s[0] == '-') x = -x;
  return x;
}
```

5.20 Gauss Elimination

```
template <size_t Dim>
struct GaussianElimination {
  vector<ll> basis;
  size_t size;

  GaussianElimination() : basis(Dim + 1), size(0) {}

  void insert(ll x) {
    for (ll i = Dim; i >= 0; i--) {
      if ((x & 1ll << i) == 0) continue;

      if (!basis[i]) {
        basis[i] = x;
```

```
      size++;
      break;
    }

    x ^= basis[i];
  }
}

void normalize() {
  for (ll i = Dim; i >= 0; i--)
    for (ll j = i - 1; j >= 0; j--)
      if (basis[i] & 1ll << j) basis[i] ^= basis[j];
}

bool check(ll x) {
  for (ll i = Dim; i >= 0; i--) {
    if ((x & 1ll << i) == 0) continue;

    if (!basis[i]) return false;

    x ^= basis[i];
  }

  return true;
}

auto operator[](ll k) { return at(k); }

ll at(ll k) {
  ll ans = 0;
  ll total = 1ll << size;
  for (ll i = Dim; ~i; i--) {
    if (!basis[i]) continue;

    ll mid = total >> 1ll;
    if ((mid < k and (ans & 1ll << i) == 0) ||
        (k <= mid and (ans & 1ll << i)))
      ans ^= basis[i];

    if (mid < k) k -= mid;

    total >>= 1ll;
  }
  return ans;
}

ll at_normalized(ll k) {
  ll ans = 0;
  k--;
  for (size_t i = 0; i <= Dim; i++) {
    if (!basis[i]) continue;
    if (k & 1) ans ^= basis[i];
    k >>= 1;
  }
  return ans;
}
};
```

5.21 Integer Partition

Find the total of ways to partition a given number N in such way that none of the parts is greater than K. Remember to memset everything to -1 before using it
time: $O(N \cdot min(N, K))$
memory: $O(N)$

```
const ll MOD = 1000000007;
const int MAXN(100);
ll memo[MAXN + 1];
ll dp(ll n, ll k = oo) {
  if (n == 0) return 1;
  ll &ans = memo[n];
  if (ans != -1) return ans;

  ans = 0;
  for (int i = 1; i <= min(n, k); i++) {
    ans = (ans + dp(n - i, k)) % MOD;
  }

  return ans;
}
```

5.22 Integer Mod

```
const ll INF = 1e18;
const ll mod = 998244353;
template <ll MOD = mod>
struct Modular {
  ll value;
  static const ll MOD_value = MOD;

  Modular(ll v = 0) {
    value = v % MOD;
```

```cpp
    if (value < 0) value += MOD;
  }
  Modular(ll a, ll b) : value(0) {
    *this += a;
    *this /= b;
  }

  Modular& operator+=(Modular const& b) {
    value += b.value;
    if (value >= MOD) value -= MOD;
    return *this;
  }
  Modular& operator-=(Modular const& b) {
    value -= b.value;
    if (value < 0) value += MOD;
    return *this;
  }
  Modular& operator*=(Modular const& b) {
    value = (ll)value * b.value % MOD;
    return *this;
  }

  friend Modular mexp(Modular a, ll e) {
    Modular res = 1;
    while (e) {
      if (e & 1) res *= a;
      a *= a;
      e >>= 1;
    }
    return res;
  }
  friend Modular inverse(Modular a) {
    return mexp(a, MOD - 2);
  }

  Modular& operator/=(Modular const& b) {
    return *this *= inverse(b);
  }
  friend Modular operator+(Modular a, Modular const b) {
    return a += b;
  }
  Modular operator++(int) {
    return this->value = (this->value + 1) % MOD;
  }
  Modular operator++() {
    return this->value = (this->value + 1) % MOD;
  }
  friend Modular operator-(Modular a, Modular const b) {
    return a -= b;
  }
  friend Modular operator-(Modular const a) {
    return 0 - a;
  }
  Modular operator--(int) {
    return this->value = (this->value - 1 + MOD) % MOD;
  }

  Modular operator--() {
    return this->value = (this->value - 1 + MOD) % MOD;
  }
  friend Modular operator*(Modular a, Modular const b) {
    return a *= b;
  }
  friend Modular operator/(Modular a, Modular const b) {
    return a /= b;
  }
  friend std::ostream& operator<<(std::ostream& os,
                                  Modular const& a) {
    return os << a.value;
  }
  friend bool operator==(Modular const& a,
                         Modular const& b) {
    return a.value == b.value;
  }
  friend bool operator!=(Modular const& a,
                         Modular const& b) {
    return a.value != b.value;
  }
};
```

5.23 Matrix Exponentiation

```cpp
ll MOD = 1'000'000'007;

template <typename T>
vector<vector<T>> prod(vector<vector<T>> &a,
                       vector<vector<T>> &b) {
  int n = len(a);
  vector<vector<T>> c(n, vector<T>(n));
```

```
  for (int i = 0; i < n; i++) {
    for (int j = 0; j < n; j++) {
      for (int k = 0; k < n; k++) {
        c[i][j] =
          (c[i][j] + ((a[i][k] * b[k][j]) % MOD)) % MOD;
      }
    }
  }

  return c;
}

template <typename T>
vector<vector<T>> fpow(vector<vector<T>> &xs, ll p) {
  vector<vector<T>> ans(len(xs), vector<T>(len(xs)));
  for (int i = 0; i < len(xs); i++) ans[i][i] = 1;

  auto b = xs;
  while (p) {
    if (p & 1) ans = prod(ans, b);
    p >>= 1;
    b = prod(b, b);
  }
  return ans;
}
```

5.24 N Choose K (elements)

process every possible combination of K elements from N elements, thoose index marked as 1 in the index vector says which elments are choosed at that moment.
Time : O($\binom{N}{K} \cdot O(process)$)

```
void process(vi &index) {
  for (int i = 0; i < len(index); i++) {
    if (index[i]) cout << i << " \n"[i == len(index) - 1];
  }
}

void n_choose_k(int n, in k) {
  vi index(n);
  fill(index.end() - k, index.end(), 1);

  do {
    process(index);
  } while (next_permutation(all(index)));
}
```

5.25 NTT integer convolution and exponentiation

Convolution finds the product a and b, and exp finds a^k
time: convolution $O(N \cdot \log N)$, exponentiation: $O(\log K \cdot N \cdot \log N)$

```
/*
    ================================================================
    */
/*===================== MINT ===================== */
template <int _mod>
struct mint {
  ll expo(ll b, ll e) {
    ll ret = 1;
    while (e) {
      if (e % 2) ret = ret * b % _mod;
      e /= 2, b = b * b % _mod;
    }
    return ret;
  }
  ll inv(ll b) { return expo(b, _mod - 2); }

  using m = mint;
  ll v;
  mint() : v(0) {}
  mint(ll v_) {
    if (v_ >= _mod or v_ <= -_mod) v_ %= _mod;
    if (v_ < 0) v_ += _mod;
    v = v_;
  }
  m& operator+=(const m& a) {
    v += a.v;
    if (v >= _mod) v -= _mod;
    return *this;
  }
  m& operator-=(const m& a) {
    v -= a.v;
    if (v < 0) v += _mod;
    return *this;
  }
  m& operator*=(const m& a) {
    v = v * ll(a.v) % _mod;
    return *this;
  }
  m& operator/=(const m& a) {
    v = v * inv(a.v) % _mod;
    return *this;
```

```
  }
  m operator-() { return m(-v); }
  m& operator^=(ll e) {
    if (e < 0) {
      v = inv(v);
      e = -e;
    }
    v = expo(v, e);
    // possivel otimizacao:
    // cuidado com 0^0
    // v = expo(v, e%(p-1));
    return *this;
  }
  bool operator==(const m& a) { return v == a.v; }
  bool operator!=(const m& a) { return v != a.v; }

  friend istream& operator>>(istream& in, m& a) {
    ll val;
    in >> val;
    a = m(val);
    return in;
  }
  friend ostream& operator<<(ostream& out, m a) {
    return out << a.v;
  }
  friend m operator+(m a, m b) { return a += b; }
  friend m operator-(m a, m b) { return a -= b; }
  friend m operator*(m a, m b) { return a *= b; }
  friend m operator/(m a, m b) { return a /= b; }
  friend m operator^(m a, ll e) { return a ^= e; }
};
/*==================== END MINT ====================== */

/*================= ntt int convolution =============*/
const ll MOD1 = 998244353;
const ll MOD2 = 754974721;
const ll MOD3 = 167772161;

template <int _mod>
void ntt(vector<mint<_mod>>& a, bool rev) {
  int n = len(a);
  auto b = a;
  assert(!(n & (n - 1)));
  mint<_mod> g = 1;
  while ((g ^ (_mod / 2)) == 1) g += 1;
  if (rev) g = 1 / g;

  for (int step = n / 2; step; step /= 2) {
    mint<_mod> w = g ^ (_mod / (n / step)), wn = 1;
    for (int i = 0; i < n / 2; i += step) {
      for (int j = 0; j < step; j++) {
        auto u = a[2 * i + j], v = wn * a[2 * i + j + step];
        b[i + j] = u + v;
        b[i + n / 2 + j] = u - v;
      }
      wn = wn * w;
    }
    swap(a, b);
  }
  if (rev) {
    auto n1 = mint<_mod>(1) / n;
    for (auto& x : a) x *= n1;
  }
}

template <ll _mod>
vector<mint<_mod>> convolution(
  const vector<mint<_mod>>& a,
  const vector<mint<_mod>>& b) {
  vector<mint<_mod>> l(all(a)), r(all(b));
  int N = len(l) + len(r) - 1, n = 1;
  while (n <= N) n *= 2;
  l.resize(n), r.resize(n);
  ntt(l, false), ntt(r, false);
  for (int i = 0; i < n; i++) l[i] *= r[i];
  ntt(l, true);
  l.resize(N);

  // Uncommnent for a boolean convolution :)
  /*

  for (auto& li : l) {
    li.v = min(li.v, 1ll);
  }
  */

  return l;
}

template <ll _mod>
```

```
vector<mint<_mod>> poly_exp(vector<mint<_mod>>& ps, int k) {
  vector<mint<_mod>> ret(len(ps));
  auto base = ps;
  ret[0] = 1;

  while (k) {
    if (k & 1) ret = convolution(ret, base);
    k >>= 1;
    base = convolution(base, base);
  }

  return ret;
}
```

5.26 NTT Integer Convolution (combine 2 modules)

Computes the convolution between polynomials (vectors) a and b
This is pure magic !
time: $O(N \log N)$

```
/*
    ====================================================================================================
    */
/*===================== MINT ===================== */
template <int _mod>
struct mint {
  ll expo(ll b, ll e) {
    ll ret = 1;
    while (e) {
      if (e % 2) ret = ret * b % _mod;
      e /= 2, b = b * b % _mod;
    }
    return ret;
  }
  ll inv(ll b) { return expo(b, _mod - 2); }

  using m = mint;
  ll v;
  mint() : v(0) {}
  mint(ll v_) {
    if (v_ >= _mod or v_ <= -_mod) v_ %= _mod;
    if (v_ < 0) v_ += _mod;
    v = v_;
  }
  m& operator+=(const m& a) {
    v += a.v;
    if (v >= _mod) v -= _mod;
    return *this;
  }
  m& operator-=(const m& a) {
    v -= a.v;
    if (v < 0) v += _mod;
    return *this;
  }
  m& operator*=(const m& a) {
    v = v * ll(a.v) % _mod;
    return *this;
  }
  m& operator/=(const m& a) {
    v = v * inv(a.v) % _mod;
    return *this;
  }
  m operator-() { return m(-v); }
  m& operator^=(ll e) {
    if (e < 0) {
      v = inv(v);
      e = -e;
    }
    v = expo(v, e);
    // possivel otimizacao:
    // cuidado com 0^0
    // v = expo(v, e%(p-1));
    return *this;
  }
  bool operator==(const m& a) { return v == a.v; }
  bool operator!=(const m& a) { return v != a.v; }

  friend istream& operator>>(istream& in, m& a) {
    ll val;
    in >> val;
    a = m(val);
    return in;
  }
  friend ostream& operator<<(ostream& out, m a) {
    return out << a.v;
  }
  friend m operator+(m a, m b) { return a += b; }
  friend m operator-(m a, m b) { return a -= b; }
  friend m operator*(m a, m b) { return a *= b; }
  friend m operator/(m a, m b) { return a /= b; }
  friend m operator^(m a, ll e) { return a ^= e; }
```

```cpp
};
/*==================== END MINT ====================== */

/*================== ntt int convolution ==============*/
const ll MOD1 = 998244353;
const ll MOD2 = 754974721;
const ll MOD3 = 167772161;

template <int _mod>
void ntt(vector<mint<_mod>>& a, bool rev) {
  int n = len(a);
  auto b = a;
  assert(!(n & (n - 1)));
  mint<_mod> g = 1;
  while ((g ^ (_mod / 2)) == 1) g += 1;
  if (rev) g = 1 / g;

  for (int step = n / 2; step; step /= 2) {
    mint<_mod> w = g ^ (_mod / (n / step)), wn = 1;
    for (int i = 0; i < n / 2; i += step) {
      for (int j = 0; j < step; j++) {
        auto u = a[2 * i + j], v = wn * a[2 * i + j + step];
        b[i + j] = u + v;
        b[i + n / 2 + j] = u - v;
      }
      wn = wn * w;
    }
    swap(a, b);
  }
  if (rev) {
    auto n1 = mint<_mod>(1) / n;
    for (auto& x : a) x *= n1;
  }
}

tuple<ll, ll, ll> ext_gcd(ll a, ll b) {
  if (!a) return {b, 0, 1};
  auto [g, x, y] = ext_gcd(b % a, a);
  return {g, y - b / a * x, x};
}

template <typename T = ll>
struct crt {
  T a, m;
  crt() : a(0), m(1) {}
  crt(T a_, T m_) : a(a_), m(m_) {}
  crt operator*(crt C) {
    auto [g, x, y] = ext_gcd(m, C.m);
    if ((a - C.a) % g != 0) a = -1;
    if (a == -1 or C.a == -1) return crt(-1, 0);
    T lcm = m / g * C.m;
    T ans = a + (x * (C.a - a) / g % (C.m / g)) * m;
    return crt((ans % lcm + lcm) % lcm, lcm);
  }
};

template <typename T = ll>
struct Congruence {
  T a, m;
};

template <typename T = ll>
T chinese_remainder_theorem(
  const vector<Congruence<T>>& equations) {
  crt<T> ans;

  for (auto& [a_, m_] : equations) {
    ans = ans * crt<T>(a_, m_);
  }

  return ans.a;
}

#define int long long
template <ll m1, ll m2>
vll merge_two_mods(const vector<mint<m1>>& a,
                   const vector<mint<m2>>& b) {
  int n = len(a);
  vll ans(n);
  for (int i = 0; i < n; i++) {
    auto cur = crt<ll>();
    auto ai = a[i].v;
    auto bi = b[i].v;
    cur = cur * crt<ll>(ai, m1);
    cur = cur * crt<ll>(bi, m2);
    ans[i] = cur.a;
  }

  return ans;
```

```
}

vll convolution_2mods(const vll& a, const vll& b) {
  vector<mint<MOD1>> l(all(a)), r(all(b));
  int N = len(l) + len(r) - 1, n = 1;
  while (n <= N) n *= 2;
  l.resize(n), r.resize(n);
  ntt(l, false), ntt(r, false);
  for (int i = 0; i < n; i++) l[i] *= r[i];
  ntt(l, true);
  l.resize(N);

  vector<mint<MOD2>> l2(all(a)), r2(all(b));
  l2.resize(n), r2.resize(n);
  ntt(l2, false), ntt(r2, false);
  rep(i, 0, n) l2[i] *= r2[i];
  ntt(l2, true);
  l2.resize(N);

  return merge_two_mods(l, l2);
}

vll poly_exp(const vll& xs, ll k) {
  vll ret(len(xs));
  ret[0] = 1;
  auto base = xs;
  while (k) {
    if (k & 1) ret = convolution_2mods(ret, base);
    k >>= 1;
    base = convolution_2mods(base, base);
  }

  return ret;
}
```

5.27 Number Of Divisors (sieve)

```
ll divisors(ll n) {
  ll ans = 1;
  for (auto p : primes) {
    if (p * p * p > n) break;

    int count = 1;
    while (n % p == 0) {
      n /= p;
      count++;
    }

    ans *= count;
  }

  if (is_prime[n])
    ans *= 2;
  else if (is_prime_square[n])
    ans *= 3;
  else if (n != 1)
    ans *= 4;

  return ans;
}
```

5.28 Number of Divisors $\tau(n)$

Find the total of divisors of N in $O(\sqrt{N})$

```
ll number_of_divisors(ll n) {
  ll res = 0;

  for (ll d = 1; d * d <= n; ++d) {
    if (n % d == 0) res += (d == n / d ? 1 : 2);
  }

  return res;
}
```

5.29 Power Sum

Calculates $K^0 + K^1 + ... + K^n$

```
ll powersum(ll n, ll k) {
  return (fastpow(n, k + 1) - 1) / (n - 1);
}
```

5.30 Sieve list primes

List every prime until MAXN, $O(N \log N)$ in time and $O(MAXN)$ in memory.

```
const ll MAXN = 2e5;
vll list_primes(ll n = MAXN) {
  vll ps;
  bitset<MAXN + 1> sieve;
  sieve.set();
  sieve.reset(1);
```

```
  for (ll i = 2; i <= n; ++i) {
    if (sieve[i]) ps.push_back(i);
    for (ll j = i * 2; j <= n; j += i) {
      sieve.reset(j);
    }
  }
  return ps;
}
```

5.31 Sum of Divisors $\sigma(n)$

Computes the sum of each divisor of n in $O(\sqrt{n})$.

```
ll sum_of_divisors(long long n) {
  ll res = 0;

  for (ll d = 1; d * d <= n; ++d) {
    if (n % d == 0) {
      ll k = n / d;

      res += (d == k ? d : d + k);
    }
  }

  return res;
}
```

6 Primitives

6.1 Bigint

```
const int maxn = 1e2 + 14, lg = 15;
const int base = 1000000000;
const int base_digits = 9;
struct bigint {
  vi a;
  int sign;

  int size() {
    if (a.empty()) return 0;
    int ans = (a.size() - 1) * base_digits;
    int ca = a.back();
    while (ca) ans++, ca /= 10;
    return ans;
  }
  bigint operator^(const bigint &v) {
    bigint ans = 1, a = *this, b = v;
    while (!b.isZero()) {
      if (b % 2) ans *= a;
      a *= a, b /= 2;
    }
    return ans;
  }
  string to_string() {
    stringstream ss;
    ss << *this;
    string s;
    ss >> s;
    return s;
  }
  int sumof() {
    string s = to_string();
    int ans = 0;
    for (auto c : s) ans += c - '0';
    return ans;
  }
  /*</arpa>*/
  bigint() : sign(1) {}

  bigint(long long v) { *this = v; }

  bigint(const string &s) { read(s); }

  void operator=(const bigint &v) {
    sign = v.sign;
    a = v.a;
  }

  void operator=(long long v) {
    sign = 1;
    a.clear();
    if (v < 0) sign = -1, v = -v;
    for (; v > 0; v = v / base) a.push_back(v % base);
  }

  bigint operator+(const bigint &v) const {
    if (sign == v.sign) {
      bigint res = v;

      for (int i = 0, carry = 0;
```

```
         i < (int)max(a.size(), v.a.size()) || carry;
         ++i) {
      if (i == (int)res.a.size()) res.a.push_back(0);
      res.a[i] += carry + (i < (int)a.size() ? a[i] : 0);
      carry = res.a[i] >= base;
      if (carry) res.a[i] -= base;
    }
    return res;
  }
  return *this - (-v);
}

bigint operator-(const bigint &v) const {
  if (sign == v.sign) {
    if (abs() >= v.abs()) {
      bigint res = *this;
      for (int i = 0, carry = 0;
           i < (int)v.a.size() || carry; ++i) {
        res.a[i] -=
          carry + (i < (int)v.a.size() ? v.a[i] : 0);
        carry = res.a[i] < 0;
        if (carry) res.a[i] += base;
      }
      res.trim();
      return res;
    }
    return -(v - *this);
  }
  return *this + (-v);
}

void operator*=(int v) {
  if (v < 0) sign = -sign, v = -v;
  for (int i = 0, carry = 0; i < (int)a.size() || carry;
       ++i) {
    if (i == (int)a.size()) a.push_back(0);
    long long cur = a[i] * (long long)v + carry;
    carry = (int)(cur / base);
    a[i] = (int)(cur % base);
    // asm("divl %%ecx" : "=a"(carry), "=d"(a[i]) :
    // "A"(cur), "c"(base));
  }
  trim();
}
```

```
bigint operator*(int v) const {
  bigint res = *this;
  res *= v;
  return res;
}

void operator*=(long long v) {
  if (v < 0) sign = -sign, v = -v;
  if (v > base) {
    *this =
      *this * (v / base) * base + *this * (v % base);
    return;
  }
  for (int i = 0, carry = 0; i < (int)a.size() || carry;
       ++i) {
    if (i == (int)a.size()) a.push_back(0);
    long long cur = a[i] * (long long)v + carry;
    carry = (int)(cur / base);
    a[i] = (int)(cur % base);
    // asm("divl %%ecx" : "=a"(carry), "=d"(a[i]) :
    // "A"(cur), "c"(base));
  }
  trim();
}

bigint operator*(long long v) const {
  bigint res = *this;
  res *= v;
  return res;
}

friend pair<bigint, bigint> divmod(const bigint &a1,
                                   const bigint &b1) {
  int norm = base / (b1.a.back() + 1);
  bigint a = a1.abs() * norm;
  bigint b = b1.abs() * norm;
  bigint q, r;
  q.a.resize(a.a.size());

  for (int i = a.a.size() - 1; i >= 0; i--) {
    r *= base;
    r += a.a[i];
    int s1 =
      r.a.size() <= b.a.size() ? 0 : r.a[b.a.size()];
    int s2 = r.a.size() <= b.a.size() - 1
```

```
              ? 0
              : r.a[b.a.size() - 1];
    int d = ((long long)base * s1 + s2) / b.a.back();
    r -= b * d;
    while (r < 0) r += b, --d;
    q.a[i] = d;
  }

  q.sign = a1.sign * b1.sign;
  r.sign = a1.sign;
  q.trim();
  r.trim();
  return make_pair(q, r / norm);
}

bigint operator/(const bigint &v) const {
  return divmod(*this, v).first;
}

bigint operator%(const bigint &v) const {
  return divmod(*this, v).second;
}

void operator/=(int v) {
  if (v < 0) sign = -sign, v = -v;
  for (int i = (int)a.size() - 1, rem = 0; i >= 0; --i) {
    long long cur = a[i] + rem * (long long)base;
    a[i] = (int)(cur / v);
    rem = (int)(cur % v);
  }
  trim();
}

bigint operator/(int v) const {
  bigint res = *this;
  res /= v;
  return res;
}

int operator%(int v) const {
  if (v < 0) v = -v;
  int m = 0;
  for (int i = a.size() - 1; i >= 0; --i)
    m = (a[i] + m * (long long)base) % v;
  return m * sign;
}

void operator+=(const bigint &v) { *this = *this + v; }
void operator-=(const bigint &v) { *this = *this - v; }
void operator*=(const bigint &v) { *this = *this * v; }
void operator/=(const bigint &v) { *this = *this / v; }

bool operator<(const bigint &v) const {
  if (sign != v.sign) return sign < v.sign;
  if (a.size() != v.a.size())
    return a.size() * sign < v.a.size() * v.sign;
  for (int i = a.size() - 1; i >= 0; i--)
    if (a[i] != v.a[i])
      return a[i] * sign < v.a[i] * sign;
  return false;
}

bool operator>(const bigint &v) const {
  return v < *this;
}
bool operator<=(const bigint &v) const {
  return !(v < *this);
}
bool operator>=(const bigint &v) const {
  return !(*this < v);
}
bool operator==(const bigint &v) const {
  return !(*this < v) && !(v < *this);
}
bool operator!=(const bigint &v) const {
  return *this < v || v < *this;
}

void trim() {
  while (!a.empty() && !a.back()) a.pop_back();
  if (a.empty()) sign = 1;
}

bool isZero() const {
  return a.empty() || (a.size() == 1 && !a[0]);
}

bigint operator-() const {
  bigint res = *this;
  res.sign = -sign;
```

```cpp
    return res;
}

bigint abs() const {
    bigint res = *this;
    res.sign *= res.sign;
    return res;
}

long long longValue() const {
    long long res = 0;
    for (int i = a.size() - 1; i >= 0; i--)
        res = res * base + a[i];
    return res * sign;
}

friend bigint gcd(const bigint &a, const bigint &b) {
    return b.isZero() ? a : gcd(b, a % b);
}
friend bigint lcm(const bigint &a, const bigint &b) {
    return a / gcd(a, b) * b;
}

void read(const string &s) {
    sign = 1;
    a.clear();
    int pos = 0;
    while (pos < (int)s.size() &&
           (s[pos] == '-' || s[pos] == '+')) {
        if (s[pos] == '-') sign = -sign;
        ++pos;
    }
    for (int i = s.size() - 1; i >= pos; i -= base_digits) {
        int x = 0;
        for (int j = max(pos, i - base_digits + 1); j <= i;
             j++)
            x = x * 10 + s[j] - '0';
        a.push_back(x);
    }
    trim();
}

friend istream &operator>>(istream &stream, bigint &v) {
    string s;
    stream >> s;
    v.read(s);
    return stream;
}

friend ostream &operator<<(ostream &stream,
                           const bigint &v) {
    if (v.sign == -1) stream << '-';
    stream << (v.a.empty() ? 0 : v.a.back());
    for (int i = (int)v.a.size() - 2; i >= 0; --i)
        stream << setw(base_digits) << setfill('0') << v.a[i];
    return stream;
}

static vector<int> convert_base(const vector<int> &a,
                                int old_digits,
                                int new_digits) {
    vector<long long> p(max(old_digits, new_digits) + 1);
    p[0] = 1;
    for (int i = 1; i < (int)p.size(); i++)
        p[i] = p[i - 1] * 10;
    vector<int> res;
    long long cur = 0;
    int cur_digits = 0;
    for (int i = 0; i < (int)a.size(); i++) {
        cur += a[i] * p[cur_digits];
        cur_digits += old_digits;
        while (cur_digits >= new_digits) {
            res.push_back(int(cur % p[new_digits]));
            cur /= p[new_digits];
            cur_digits -= new_digits;
        }
    }
    res.push_back((int)cur);
    while (!res.empty() && !res.back()) res.pop_back();
    return res;
}

typedef vector<long long> vll;

static vll karatsubaMultiply(const vll &a, const vll &b) {
    int n = a.size();
    vll res(n + n);
    if (n <= 32) {
        for (int i = 0; i < n; i++)
            for (int j = 0; j < n; j++)
```

```
        res[i + j] += a[i] * b[j];
    return res;
  }

  int k = n >> 1;
  vll a1(a.begin(), a.begin() + k);
  vll a2(a.begin() + k, a.end());
  vll b1(b.begin(), b.begin() + k);
  vll b2(b.begin() + k, b.end());

  vll a1b1 = karatsubaMultiply(a1, b1);
  vll a2b2 = karatsubaMultiply(a2, b2);

  for (int i = 0; i < k; i++) a2[i] += a1[i];
  for (int i = 0; i < k; i++) b2[i] += b1[i];

  vll r = karatsubaMultiply(a2, b2);
  for (int i = 0; i < (int)a1b1.size(); i++)
    r[i] -= a1b1[i];
  for (int i = 0; i < (int)a2b2.size(); i++)
    r[i] -= a2b2[i];

  for (int i = 0; i < (int)r.size(); i++)
    res[i + k] += r[i];
  for (int i = 0; i < (int)a1b1.size(); i++)
    res[i] += a1b1[i];
  for (int i = 0; i < (int)a2b2.size(); i++)
    res[i + n] += a2b2[i];
  return res;
}

bigint operator*(const bigint &v) const {
  vector<int> a6 = convert_base(this->a, base_digits, 6);
  vector<int> b6 = convert_base(v.a, base_digits, 6);
  vll a(a6.begin(), a6.end());
  vll b(b6.begin(), b6.end());
  while (a.size() < b.size()) a.push_back(0);
  while (b.size() < a.size()) b.push_back(0);
  while (a.size() & (a.size() - 1))
    a.push_back(0), b.push_back(0);
  vll c = karatsubaMultiply(a, b);
  bigint res;
  res.sign = sign * v.sign;
  for (int i = 0, carry = 0; i < (int)c.size(); i++) {
    long long cur = c[i] + carry;
      res.a.push_back((int)(cur % 1000000));
      carry = (int)(cur / 1000000);
    }
    res.a = convert_base(res.a, 6, base_digits);
    res.trim();
    return res;
  }
};
```

6.2 Integer Mod

```
const ll MOD = 1'000'000'000 + 7;
template <ll _mod = MOD>
struct mint {
  ll value;
  static const ll MOD_value = _mod;

  mint(ll v = 0) {
    value = v % _mod;
    if (value < 0) value += _mod;
  }
  mint(ll a, ll b) : value(0) {
    *this += a;
    *this /= b;
  }

  mint& operator+=(mint const& b) {
    value += b.value;
    if (value >= _mod) value -= _mod;
    return *this;
  }
  mint& operator-=(mint const& b) {
    value -= b.value;
    if (value < 0) value += _mod;
    return *this;
  }
  mint& operator*=(mint const& b) {
    value = (ll)value * b.value % _mod;
    return *this;
  }

  friend mint mexp(mint a, ll e) {
    mint res = 1;
    while (e) {
      if (e & 1) res *= a;
```

```
    a *= a;
    e >>= 1;
  }
  return res;
}
friend mint inverse(mint a) { return mexp(a, _mod - 2); }

mint& operator/=(mint const& b) {
  return *this *= inverse(b);
}
friend mint operator+(mint a, mint const b) {
  return a += b;
}
mint operator++(int) {
  return this->value = (this->value + 1) % _mod;
}
mint operator++() {
  return this->value = (this->value + 1) % _mod;
}
friend mint operator-(mint a, mint const b) {
  return a -= b;
}
friend mint operator-(mint const a) { return 0 - a; }
mint operator--(int) {
  return this->value = (this->value - 1 + _mod) % _mod;
}

mint operator--() {
  return this->value = (this->value - 1 + _mod) % _mod;
}
friend mint operator*(mint a, mint const b) {
  return a *= b;
}
friend mint operator/(mint a, mint const b) {
  return a /= b;
}
friend std::ostream& operator<<(std::ostream& os,
                                mint const& a) {
  return os << a.value;
}
friend bool operator==(mint const& a, mint const& b) {
  return a.value == b.value;
}
friend bool operator!=(mint const& a, mint const& b) {
  return a.value != b.value;
  }
};
```

6.3 Matrix

```
template <typename T>
struct Matrix {
  vector<vector<T>> d;

  Matrix() : Matrix(0) {}
  Matrix(int n) : Matrix(n, n) {}
  Matrix(int n, int m)
    : Matrix(vector<vector<T>>(n, vector<T>(m))) {}
  Matrix(const vector<vector<T>> &v) : d(v) {}

  constexpr int n() const { return (int)d.size(); }
  constexpr int m() const {
    return n() ? (int)d[0].size() : 0;
  }

  void rotate() { *this = rotated(); }

  Matrix<T> rotated() const {
    Matrix<T> res(m(), n());
    for (int i = 0; i < m(); i++) {
      for (int j = 0; j < n(); j++) {
        res[i][j] = d[n() - j - 1][i];
      }
    }
    return res;
  }

  Matrix<T> pow(int power) const {
    assert(n() == m());

    auto res = Matrix<T>::identity(n());
    auto b = *this;
    while (power) {
      if (power & 1) res *= b;
      b *= b;
      power >>= 1;
    }
    return res;
  }
```

```
Matrix<T> submatrix(int start_i, int start_j,
                    int rows = INT_MAX,
                    int cols = INT_MAX) const {
  rows = min(rows, n() - start_i);
  cols = min(cols, m() - start_j);
  if (rows <= 0 or cols <= 0) return {};

  Matrix<T> res(rows, cols);
  for (int i = 0; i < rows; i++)
    for (int j = 0; j < cols; j++)
      res[i][j] = d[i + start_i][j + start_j];
  return res;
}

Matrix<T> translated(int x, int y) const {
  Matrix<T> res(n(), m());
  for (int i = 0; i < n(); i++) {
    for (int j = 0; j < m(); j++) {
      if (i + x < 0 or i + x >= n() or j + y < 0 or
          j + y >= m())
        continue;
      res[i + x][j + y] = d[i][j];
    }
  }
  return res;
}

static Matrix<T> identity(int n) {
  Matrix<T> res(n);
  for (int i = 0; i < n; i++) res[i][i] = 1;
  return res;
}

vector<T> &operator[](int i) { return d[i]; }
const vector<T> &operator[](int i) const { return d[i]; }
Matrix<T> &operator+=(T value) {
  for (auto &row : d) {
    for (auto &x : row) x += value;
  }
  return *this;
}
Matrix<T> operator+(T value) const {
  auto res = *this;
  for (auto &row : res) {
    for (auto &x : row) x = x + value;
  }
  return res;
}
Matrix<T> &operator-=(T value) {
  for (auto &row : d) {
    for (auto &x : row) x -= value;
  }
  return *this;
}
Matrix<T> operator-(T value) const {
  auto res = *this;
  for (auto &row : res) {
    for (auto &x : row) x = x - value;
  }
  return res;
}
Matrix<T> &operator*=(T value) {
  for (auto &row : d) {
    for (auto &x : row) x *= value;
  }
  return *this;
}
Matrix<T> operator*(T value) const {
  auto res = *this;
  for (auto &row : res) {
    for (auto &x : row) x = x * value;
  }
  return res;
}
Matrix<T> &operator/=(T value) {
  for (auto &row : d) {
    for (auto &x : row) x /= value;
  }
  return *this;
}
Matrix<T> operator/(T value) const {
  auto res = *this;
  for (auto &row : res) {
    for (auto &x : row) x = x / value;
  }
  return res;
}
Matrix<T> &operator+=(const Matrix<T> &o) {
  assert(n() == o.n() and m() == o.m());
  for (int i = 0; i < n(); i++) {
```

```cpp
    for (int j = 0; j < m(); j++) {
      d[i][j] += o[i][j];
    }
  }
  return *this;
}
Matrix<T> operator+(const Matrix<T> &o) const {
  assert(n() == o.n() and m() == o.m());
  auto res = *this;
  for (int i = 0; i < n(); i++) {
    for (int j = 0; j < m(); j++) {
      res[i][j] = res[i][j] + o[i][j];
    }
  }
  return res;
}
Matrix<T> &operator-=(const Matrix<T> &o) {
  assert(n() == o.n() and m() == o.m());
  for (int i = 0; i < n(); i++) {
    for (int j = 0; j < m(); j++) {
      d[i][j] -= o[i][j];
    }
  }
  return *this;
}
Matrix<T> operator-(const Matrix<T> &o) const {
  assert(n() == o.n() and m() == o.m());
  auto res = *this;
  for (int i = 0; i < n(); i++) {
    for (int j = 0; j < m(); j++) {
      res[i][j] = res[i][j] - o[i][j];
    }
  }
  return res;
}
Matrix<T> &operator*=(const Matrix<T> &o) {
  *this = *this * o;
  return *this;
}
Matrix<T> operator*(const Matrix<T> &o) const {
  assert(m() == o.n());
  Matrix<T> res(n(), o.m());
  for (int i = 0; i < res.n(); i++) {
    for (int j = 0; j < res.m(); j++) {
      auto &x = res[i][j];
        for (int k = 0; k < m(); k++) {
          x += (d[i][k] * o[k][j]);
        }
      }
    }
    return res;
  }

  friend istream &operator>>(istream &is, Matrix<T> &mat) {
    for (auto &row : mat)
      for (auto &x : row) is >> x;
    return is;
  }
  friend ostream &operator<<(ostream &os,
                             const Matrix<T> &mat) {
    bool frow = 1;
    for (auto &row : mat) {
      if (not frow) os << '\n';
      bool first = 1;
      for (auto &x : row) {
        if (not first) os << ' ';
        os << x;
        first = 0;
      }

      frow = 0;
    }
    return os;
  }

  auto begin() { return d.begin(); }
  auto end() { return d.end(); }
  auto rbegin() { return d.rbegin(); }
  auto rend() { return d.rend(); }

  auto begin() const { return d.begin(); }
  auto end() const { return d.end(); }
  auto rbegin() const { return d.rbegin(); }
  auto rend() const { return d.rend(); }
};
```

7 Problems

7.1 Hanoi Tower

Let T_n be the total of moves to solve a hanoi tower, we know that $T_n >= 2 \cdot T_{n-1} + 1$, for $n > 0$, and $T_0 = 0$. By induction it's easy to see that $T_n = 2^n - 1$, for $n > 0$.
The following algorithm finds the necessary steps to solve the game for 3 stacks and n disks.

```
void move(int a, int b) { cout << a << ' ' << b << endl; }
void solve(int n, int s, int e) {
  if (n == 0) return;
  if (n == 1) {
    move(s, e);
    return;
  }
  solve(n - 1, s, 6 - s - e);
  move(s, e);
  solve(n - 1, 6 - s - e, e);
}
```

8 Searching

8.1 Meet in the middle

Answers the query how many subsets of the vector xs have sum equal x.
Time: $O(N \cdot 2^{\frac{N}{2}})$

```
vll get_subset_sums(int l, int r, vll &a) {
  int len = r - l + 1;
  vll res;

  for (int i = 0; i < (1 << len); i++) {
    ll sum = 0;
    for (int j = 0; j < len; j++) {
      if (i & (1 << j)) {
        sum += a[l + j];
      }
    }
    res.push_back(sum);
  }
  return res;
};

ll count(vll &xs, ll x) {
  int n = len(xs);
  vll left = get_subset_sums(0, n / 2 - 1, xs);
  vll right = get_subset_sums(n / 2, n - 1, xs);
  sort(all(left));
  sort(all(right));
  ll ans = 0;
  for (ll i : left) {
    auto start_index =
      lower_bound(right.begin(), right.end(), x - i) -
      right.begin();
    auto end_index =
      upper_bound(right.begin(), right.end(), x - i) -
      right.begin();
    ans += end_index - start_index;
  }
  return ans;
}
```

8.2 Ternary Search Recursive

```
const double eps = 1e-6;

// IT MUST BE AN UNIMODAL FUNCTION
double f(int x) { return x * x + 2 * x + 4; }

double ternary_search(double l, double r) {
  if (fabs(f(l) - f(r)) < eps)
    return f((l + (r - l) / 2.0));

  auto third = (r - l) / 3.0;
  auto m1 = l + third;
  auto m2 = r - third;

  // change the signal to find the maximum point.
  return m1 < m2 ? ternary_search(m1, r)
                 : ternary_search(l, m2);
}
```

9 Strings

9.1 Count Distinct Anagrams

```
const ll MOD = 1e9 + 7;
const int maxn = 1e6;
vll fs(maxn + 1);
void precompute() {
  fs[0] = 1;
  for (ll i = 1; i <= maxn; i++) {
```

```
    fs[i] = (fs[i - 1] * i) % MOD;
  }
}

ll fpow(ll a, int n, ll mod = LLONG_MAX) {
  if (n == 0) return 1;
  if (n == 1) return a;
  ll x = fpow(a, n / 2, mod) % mod;
  return ((x * x) % mod * (n & 1 ? a : 1ll)) % mod;
}

ll distinctAnagrams(const string &s) {
  precompute();
  vi hist('z' - 'a' + 1, 0);
  for (auto &c : s) hist[c - 'a']++;
  ll ans = fs[len(s)];
  for (auto &q : hist) {
    ans = (ans * fpow(fs[q], MOD - 2, MOD)) % MOD;
  }
  return ans;
}
```

9.2 Double Hash Range Query

```
const ll MOD = 1000027957;
const int MOD2 = 1000015187;

struct Hash {
  const ll P = 31;
  int n;
  string s;
  vll h, h2, hi, hi2, p, p2;
  Hash() {}
  Hash(string _s)
    : s(_s),
      n(len(_s)),
      h(n),
      h2(n),
      hi(n),
      hi2(n),
      p(n),
      p2(n) {
    for (int i = 0; i < n; i++)
      p[i] = (i ? P * p[i - 1] : 1) % MOD;
    for (int i = 0; i < n; i++)
      p2[i] = (i ? P * p2[i - 1] : 1) % MOD2;
    for (int i = 0; i < n; i++)
      h[i] = (s[i] + (i ? h[i - 1] : 0) * P) % MOD;
    for (int i = 0; i < n; i++)
      h2[i] = (s[i] + (i ? h2[i - 1] : 0) * P) % MOD2;
    for (int i = n - 1; i >= 0; i--)
      hi[i] =
        (s[i] + (i + 1 < n ? hi[i + 1] : 0) * P) % MOD;
    for (int i = n - 1; i >= 0; i--)
      hi2[i] =
        (s[i] + (i + 1 < n ? hi2[i + 1] : 0) * P) % MOD2;
  }
  pii query(int l, int r) {
    ll hash =
      (h[r] - (l ? h[l - 1] * p[r - l + 1] % MOD : 0));
    ll hash2 =
      (h2[r] - (l ? h2[l - 1] * p2[r - l + 1] % MOD2 : 0));

    return {(hash < 0 ? hash + MOD : hash),
            (hash2 < 0 ? hash2 + MOD2 : hash2)};
  }
  pii query_inv(int l, int r) {
    ll hash =
      (hi[l] -
       (r + 1 < n ? hi[r + 1] * p[r - l + 1] % MOD : 0));
    ll hash2 =
      (hi2[l] -
       (r + 1 < n ? hi2[r + 1] * p2[r - l + 1] % MOD2 : 0));
    return {(hash < 0 ? hash + MOD : hash),
            (hash2 < 0 ? hash2 + MOD2 : hash2)};
  }
};
```

9.3 Hash Range Query

```
const ll P = 31;
const ll MOD = 1e9 + 9;
const int MAXN(1e6);

ll ppow[MAXN + 1];
void pre_calc() {
  ppow[0] = 1;
  for (int i = 1; i <= MAXN; i++)
    ppow[i] = (ppow[i - 1] * P) % MOD;
}
```

```
struct Hash {
  int n;
  vll h, hi;
  Hash(const string &s) : n(s.size()), h(n), hi(n) {
    h[0] = s[0];
    hi[n - 1] = s[n - 1];
    for (int i = 1; i < n; i++) {
      h[i] = (s[i] + h[i - 1] * P) % MOD;
      hi[n - i - 1] =
        (s[n - i - 1] + hi[n - i - 1] * P) % MOD;
    }
  }

  ll qry(int l, int r) {
    ll hash =
      (h[r] - (l ? h[l - 1] * ppow[r - l + 1] % MOD : 0));
    return hash < 0 ? hash + MOD : hash;
  }

  ll qry_inv(int l, int r) {
    ll hash =
      (hi[l] -
       (r + 1 < n ? hi[r + 1] * ppow[r - l + 1] % MOD : 0));
    return hash < 0 ? hash + MOD : hash;
  }
};
```

9.4 K-th digit in digit string

Find the k-th digit in a *digit string*, only works for $1 <= k <= 10^{18}$!
Time: precompute $O(1)$, query $O(1)$

```
using vull = vector<ull>;
vull pow10;
vector<array<ull, 4>> memo;
void precompute(int maxpow = 18) {
  ull qtd = 1;
  ull start = 1;
  ull end = 9;
  ull curlenght = 9;
  ull startstr = 1;
  ull endstr = 9;

  for (ull i = 0, j = 1ll; (int)i < maxpow; i++, j *= 10ll)
    pow10.eb(j);
  for (ull i = 0; i < maxpow - 1ull; i++) {
    memo.push_back({start, end, startstr, endstr});

    start = end + 1ll;
    end = end + (9ll * pow10[qtd]);
    curlenght = end - start + 1ull;

    qtd++;
    startstr = endstr + 1ull;
    endstr = (endstr + 1ull) + (curlenght)*qtd - 1ull;
  }
}
char kthDigit(ull k) {
  int qtd = 1;
  for (auto [s, e, ss, es] : memo) {
    if (k >= ss and k <= es) {
      ull pos = k - ss;
      ull index = pos / qtd;
      ull nmr = s + index;
      int i = k - ss - qtd * index;

      return ((nmr / pow10[qtd - i - 1]) % 10) + '0';
    }
    qtd++;
  }

  return 'X';
}
```

9.5 Longest Palindrome Substring (Manacher)

Finds the longest palindrome substring, manacher returns a vector where the i-th position is how much is possible to grow the string to the left and the right of i and keep it a palindrome.
Time: $O(N)$

```
vi manacher(const string &s) {
  int n = len(s) - 2;
  vi p(n + 2);
  int l = 1, r = 1;
  for (int i = 1; i <= n; i++) {
    p[i] = max(0, min(r - i, p[l + (r - i)]));
    while (s[i - p[i]] == s[i + p[i]]) p[i]++;
    if (i + p[i] > r) l = i - p[i], r = i + p[i];
    p[i]--;
  }
  return p;
```

```
}
string longest_palindrome(const string &s) {
  string t("$#");
  for (auto c : s) t.push_back(c), t.push_back('#');
  t.push_back('^');
  vi xs = manacher(t);
  int mpos = max_element(all(xs)) - xs.begin();
  string p;
  for (int k = xs[mpos], i = mpos - k; i <= mpos + k; i++)
    if (t[i] != '#') p.push_back(t[i]);
  return p;
}
```

9.6 Longest Palindrome

```
string longest_palindrome(const string &s) {
  int n = (int)s.size();
  vector<array<int, 2>> dp(n);

  pii odd(0, -1), even(0, -1);
  pii ans;
  for (int i = 0; i < n; i++) {
    int k = 0;
    if (i > odd.second)
      k = 1;
    else
      k = min(dp[odd.first + odd.second - i][0],
              odd.second - i + 1);
    while (i - k >= 0 and i + k < n and
           s[i - k] == s[i + k])
      k++;
    dp[i][0] = k--;
    if (i + k > odd.second) odd = {i - k, i + k};
    if (2 * dp[i][0] - 1 > ans.second)
      ans = {i - k, 2 * dp[i][0] - 1};

    k = 0;
    if (i <= even.second)
      k = min(dp[even.first + even.second - i + 1][1],
              even.second - i + 1);
    while (i - k - 1 >= 0 and i + k < n and
           s[i - k - 1] == s[i + k])
      k++;
    dp[i][1] = k--;
    if (i + k > even.second) even = {i - k - 1, i + k};
    if (2 * dp[i][1] > ans.second)
      ans = {i - k - 1, 2 * dp[i][1]};
  }
  return s.substr(ans.first, ans.second);
}
```

9.7 Rabin Karp

```
size_t rabin_karp(const string &s, const string &p) {
  if (s.size() < p.size()) return 0;

  auto n = s.size(), m = p.size();
  const ll p1 = 31, p2 = 29, q1 = 1e9 + 7, q2 = 1e9 + 9;
  const ll p1_1 = fpow(p1, q1 - 2, q1),
           p1_2 = fpow(p1, m - 1, q1);
  const ll p2_1 = fpow(p2, q2 - 2, q2),
           p2_2 = fpow(p2, m - 1, q2);

  pair<ll, ll> hs, hp;
  for (int i = (int)m - 1; ~i; --i) {
    hs.first = (hs.first * p1) % q1;
    hs.first = (hs.first + (s[i] - 'a' + 1)) % q1;
    hs.second = (hs.second * p2) % q2;
    hs.second = (hs.second + (s[i] - 'a' + 1)) % q2;

    hp.first = (hp.first * p1) % q1;
    hp.first = (hp.first + (p[i] - 'a' + 1)) % q1;
    hp.second = (hp.second * p2) % q2;
    hp.second = (hp.second + (p[i] - 'a' + 1)) % q2;
  }

  size_t occ = 0;
  for (size_t i = 0; i < n - m; i++) {
    occ += (hs == hp);

    int fi = s[i] - 'a' + 1;
    int fm = s[i + m] - 'a' + 1;

    hs.first = (hs.first - fi + q1) % q1;
    hs.first = (hs.first * p1_1) % q1;
    hs.first = (hs.first + fm * p1_2) % q1;
    hs.second = (hs.second - fi + q2) % q2;
    hs.second = (hs.second * p2_1) % q2;
    hs.second = (hs.second + fm * p2_2) % q2;
  }
```

```
  occ += hs == hp;

  return occ;
}
```

9.8 String Psum

```
struct strPsum {
  ll n;
  ll k;
  vector<vll> psum;
  strPsum(const string &s)
    : n(s.size()), k(100), psum(k, vll(n + 1)) {
    for (ll i = 1; i <= n; ++i) {
      for (ll j = 0; j < k; ++j) {
        psum[j][i] = psum[j][i - 1];
      }
      psum[s[i - 1]][i]++;
    }
  }

  ll qtd(ll l, ll r, char c) {  // [0,n-1]
    return psum[c][r + 1] - psum[c][l];
  }
}
```

9.9 Suffix Automaton (complete)

```
struct state {
  int len, link, cnt, firstpos;
  // this can be optimized using a vector with the alphabet
  // size
  map<char, int> next;
  vi inv_link;
};
struct SuffixAutomaton {
  vector<state> st;
  int sz = 0;
  int last;
  vc cloned;

  SuffixAutomaton(const string &s, int maxlen)
    : st(maxlen * 2), cloned(maxlen * 2) {
    st[0].len = 0;
    st[0].link = -1;
    sz++;
    last = 0;
    for (auto &c : s) add_char(c);

    // precompute for count occurences
    for (int i = 1; i < sz; i++) {
      st[i].cnt = !cloned[i];
    }
    vector<pair<state, int>> aux;
    for (int i = 0; i < sz; i++) {
      aux.push_back({st[i], i});
    }

    sort(all(aux), [](const pair<state, int> &a,
                      const pair<state, int> &b) {
      return a.fst.len > b.fst.len;
    });

    for (auto &[stt, id] : aux) {
      if (stt.link != -1) {
        st[stt.link].cnt += st[id].cnt;
      }
    }

    // for find every occurende position
    for (int v = 1; v < sz; v++) {
      st[st[v].link].inv_link.push_back(v);
    }
  }

  void add_char(char c) {
    int cur = sz++;
    st[cur].len = st[last].len + 1;
    st[cur].firstpos = st[cur].len - 1;
    int p = last;
    // follow the suffix link until find a transition to c
    while (p != -1 and !st[p].next.count(c)) {
      st[p].next[c] = cur;
      p = st[p].link;
    }
    // there was no transition to c so create and leave
    if (p == -1) {
      st[cur].link = 0;
      last = cur;
      return;
    }
```

```cpp
  int q = st[p].next[c];
  if (st[p].len + 1 == st[q].len) {
    st[cur].link = q;
  } else {
    int clone = sz++;
    cloned[clone] = true;
    st[clone].len = st[p].len + 1;
    st[clone].next = st[q].next;
    st[clone].link = st[q].link;
    st[clone].firstpos = st[q].firstpos;
    while (p != -1 and st[p].next[c] == q) {
      st[p].next[c] = clone;
      p = st[p].link;
    }
    st[q].link = st[cur].link = clone;
  }
  last = cur;
}

bool checkOccurrence(const string &t) {  // O(len(t))
  int cur = 0;
  for (auto &c : t) {
    if (!st[cur].next.count(c)) return false;
    cur = st[cur].next[c];
  }
  return true;
}
ll totalSubstrings() {  // distinct, O(len(s))
  ll tot = 0;
  for (int i = 1; i < sz; i++) {
    tot += st[i].len - st[st[i].link].len;
  }
  return tot;
}

// count occurences of a given string t
int countOccurences(const string &t) {
  int cur = 0;
  for (auto &c : t) {
    if (!st[cur].next.count(c)) return 0;
    cur = st[cur].next[c];
  }
  return st[cur].cnt;
}

  // find the first index where t appears a substring
  // O(len(t))
  int firstOccurence(const string &t) {
    int cur = 0;
    for (auto c : t) {
      if (!st[cur].next.count(c)) return -1;
      cur = st[cur].next[c];
    }
    return st[cur].firstpos - len(t) + 1;
  }

  vi everyOccurence(const string &t) {
    int cur = 0;
    for (auto c : t) {
      if (!st[cur].next.count(c)) return {};
      cur = st[cur].next[c];
    }
    vi ans;
    getEveryOccurence(cur, len(t), ans);
    return ans;
  }

  void getEveryOccurence(int v, int P_length, vi &ans) {
    if (!cloned[v]) ans.pb(st[v].firstpos - P_length + 1);
    for (int u : st[v].inv_link)
      getEveryOccurence(u, P_length, ans);
  }
};
```

9.10 Trie

- build with the size of the alphabet (*sigma*) and the first char (*norm*)
- *insert*(s) insert the string in the trie $O(|s| * sigma)$
- *erase*(s) remove the string from the trie $O(|s|)$
- *find*(s) return the last node from the string s, 0 if not found $O(|s|)$

```cpp
struct trie {
  vi2d to;
  vi end, pref;
  int sigma;
  char norm;

  trie(int sigma_ = 26, char norm_ = 'a')
    : sigma(sigma_), norm(norm_) {
    to = {vector<int>(sigma)};
```

```
    end = {0}, pref = {0};
  }

  int next(int node, char key) {
    return to[node][key - norm];
  }

  void insert(const string &s) {
    int x = 0;
    for (auto c : s) {
      int &nxt = to[x][c - norm];
      if (!nxt) {
        nxt = len(to);
        to.push_back(vi(sigma));
        end.emplace_back(0), pref.emplace_back(0);
      }
      x = nxt, pref[x]++;
    }
    end[x]++, pref[0]++;
  }
  void erase(const string &s) {
    int x = 0;
    for (char c : s) {
      int &nxt = to[x][c - norm];
      x = nxt, pref[x]--;
      if (!pref[x]) nxt = 0;
    }
    end[x]--, pref[0]--;
  }
  int find(const string &s) {
    int x = 0;
    for (auto c : s) {
      x = to[x][c - norm];
      if (!x) return 0;
    }
    return x;
  }
};
```

9.11 Z-function get occurence positions

$O(len(s) + len(p))$

```
vi getOccPos(string &s, string &p) {
  // Z-function
  char delim = '#';
  string t{p + delim + s};
  vi zs(len(t));

  int l = 0, r = 0;
  for (int i = 1; i < len(t); i++) {
    if (i <= r) zs[i] = min(zs[i - 1], r - i + 1);
    while (zs[i] + i < len(t) and t[zs[i]] == t[i + zs[i]])
      zs[i]++;
    if (r < i + zs[i] - 1) l = i, r = i + zs[i] - 1;
  }

  // Iterate over the results of Z-function to get ranges
  vi ans;
  int start = len(p) + 1 + 1 - 1;
  for (int i = start; i < len(zs); i++) {
    if (zs[i] == len(p)) {
      int l = i - start;
      ans.emplace_back(l);
    }
  }
  return ans;
}
```

10 Settings and macros

10.1 .vimrc

```
set ts=4 sw=4 sta nu rnu sc cindent
set bg=dark ruler clipboard=unnamed,unnamedplus, timeoutlen=100
colorscheme default

nnoremap <C-j> :botright belowright term bash <CR>
syntax on
```

10.2 macro.cpp

```
#include <bits/stdc++.h>
using namespace std;
#ifdef LOCAL
#include "debug.cpp"
#else
#define dbg(...)
#endif
#define endl '\n'
#define fastio                  \
```

```
  ios_base::sync_with_stdio(0); \
  cin.tie(0);
// #define int long long
#define len(__x) (int)__x.size()
using ll = long long;
using ull = unsigned long long;
using ld = long double;
using vll = vector<ll>;
using pll = pair<ll, ll>;
using vll2d = vector<vll>;
using vi = vector<int>;
using vi2d = vector<vi>;
using pii = pair<int, int>;
using vii = vector<pii>;
using vc = vector<char>;
#define all(a) a.begin(), a.end()
#define rall(a) a.rbegin(), a.rend()
#define pb push_back
#define eb emplace_back
#define ff first
#define ss second
#define rep(i, begin, end)                          \
  for (__typeof(begin) i = (begin) - ((begin) > (end)); \
       i != (end) - ((begin) > (end));              \
       i += 1 - 2 * ((begin) > (end)))

int lg2(ll x) {
  return __builtin_clzll(1) - __builtin_clzll(x);
}

// vector<string> dir({"LU", "U", "RU", "R", "RD", "D",
// "LD", "L"}); int dx[] = {-1, -1, -1, 0, 1, 1, 1, 0}; int
// dy[] = {-1, 0, 1, 1, 1, 0, -1, -1};
vector<string> dir({"U", "R", "D", "L"});
int dx[] = {-1, 0, 1, 0};
int dy[] = {0, 1, 0, -1};

const ll oo = 1e18;
int T(1);
const int MAXN(1'000'000);

auto run() {}

int32_t main(void) {
#ifndef LOCAL
  fastio;
#endif

  // cin >> T;

  for (int i = 1; i <= T; i++) {
    run();
  }
}
```

10.3 debug.cpp

```
#include <bits/stdc++.h>
using namespace std;
/******** Debug Code *******/
template <typename T>
concept Printable = requires(T t) {
    { std::cout << t } -> std::same_as<std::ostream &>;
};
template <Printable T>
void __print(const T &x) {
    cerr << x;
}
template <size_t T>
void __print(const bitset<T> &x) {
    cerr << x;
}
template <typename A, typename B>
void __print(const pair<A, B> &p);
template <typename... A>
void __print(const tuple<A...> &t);
template <typename T>
void __print(stack<T> s);
template <typename T>
void __print(queue<T> q);
template <typename T, typename... U>
void __print(priority_queue<T, U...> q);
template <typename A>
void __print(const A &x) {
    bool first = true;
    cerr << '{';
    for (const auto &i : x) {
        cerr << (first ? "" : ","), __print(i);
        first = false;
    }
```

```cpp
    cerr << '}';
}
template <typename A, typename B>
void __print(const pair<A, B> &p) {
    cerr << '(';
    __print(p.first);
    cerr << ',';
    __print(p.second);
    cerr << ')';
}
template <typename... A>
void __print(const tuple<A...> &t) {
    bool first = true;
    cerr << '(';
    apply(
        [&first](const auto &...args) {
            ((cerr << (first ? "" : ","), __print(args), first
  = false), ...);
        },
        t);
    cerr << ')';
}
template <typename T>
void __print(stack<T> s) {
    vector<T> debugVector;
    while (!s.empty()) {
        T t = s.top();
        debugVector.push_back(t);
        s.pop();
    }
    reverse(debugVector.begin(), debugVector.end());
    __print(debugVector);
}
template <typename T>
void __print(queue<T> q) {
    vector<T> debugVector;
    while (!q.empty()) {
        T t = q.front();
        debugVector.push_back(t);
        q.pop();
    }
    __print(debugVector);
}
template <typename T, typename... U>
void __print(priority_queue<T, U...> q) {
    vector<T> debugVector;
    while (!q.empty()) {
        T t = q.top();
        debugVector.push_back(t);
        q.pop();
    }
    __print(debugVector);
}
void _print() { cerr << "]\n"; }
template <typename Head, typename... Tail>
void _print(const Head &H, const Tail &...T) {
    __print(H);
    if (sizeof...(T)) cerr << ", ";
    _print(T...);
}

#define dbg(x...)                       \
    cerr << "[" << #x << "] = ["; \
    _print(x)
```

10.4 .bashrc

```bash
cpp() {
  g++ -std=c++20 -fsanitize=address,undefined -Wall $1 && time
   ./a.out
}

cpp() {
  echo ">> COMPILING <<" 1>&2
  g++ -std=c++17 \
      -O2 \
      -g \
      -g3 \
      -Wextra \
      -Wshadow \
      -Wformat=2 \
      -Wconversion \
      -fsanitize=address,undefined \
      -fno-sanitize-recover \
      -Wfatal-errors  \
      $1

  if [ $? -ne 0 ]; then
      echo ">> FAILED <<" 1>&2
      return 1
```

```
  fi
  echo ">> DONE << " 1>&2
  time  ./a.out ${@:2}
}

prepare() {
    cp debug.cpp ./
    for i in {a..z}
    do
        cp macro.cpp $i.cpp
        touch $i.py
    done

    for i in {1..10}
    do
        touch in${i}
        touch out${i}
        touch ans${i}
    done
}
```

10.5 short-macro.cpp

```
#include <bits/stdc++.h>
using namespace std;
#define fastio                    \
  ios_base::sync_with_stdio(0); \
  cin.tie(0);

void run() {}

int32_t main(void) {
  fastio;
  int t;
  t = 1;
  // cin >> t;
  while (t--) run();
}
```