# Teorema do chinês viajante

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Universidade de Brasília, faculdade do gama

# Coach

Edson Alves (edsomrjr)

# Team

André Macedo (andre\_rei)

Iago Rocha (iagorrr)

Thalisson Alves (Thalisson\_Alves)

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### 1.1 bash config

```
#copy first argument to clipborad ! ONLY WORK ON
     XORG !
alias clip="xclip —sel clip"
# compile the $1 parameter, if a $2 is provided
# the name will be the the binary output, if
# none is provided the binary name will be
  'a.out
comp() {
  echo ">> COMPILING $1 <<" 1>&2
  if [ $# -gt 1 ]; then
     outfile="${2}
   else
     outfile="a.out"
  time g++ -std=c++20 \
     -02 \
-g3 \
     -Wall
     —fsanitize=address,undefined \
     -fno-sanitize-recover \
-D LOCAL \
-0 "${outfile}" \
"$1"
  if [ $? -ne 0 ]; then
  echo ">> FAILED <<" 1>&2
return 1
  echo ">> DONE << " 1>&2
# run the binary given in $1, if none is
# given it will try to run the 'a.out'
# binary
run() {
  to_run=./a.out
if [ -n "$1" ]; then
  to_run="$1"
   time $to_run
  just comp and run your cpp file
# accpets <in1 >out and everything else
comprun() {
   comp "$1" "a" && run ./a ${@:2}
testall() {
  comp "$1" generator
  comp "$2" brute
  comp "$3" main
  input_counter=1
  while true; do
     echo "$input_counter"
     run ./generator >input
     run ./main <input >main_output.txt
     run ./brute <input >brute_output.txt
     diff brute_output.txt main_output.txt
     if [ $? —ne 0 ]; then
  echo "Outputs differ at input $input_counter"
        echo "Brute file output:"
       cat brute_output.txt
echo "Main file output:"
        cat main_output.txt
        echo "input used:
        cat input
        break
     ((input_counter++))
  done
}
touch_macro() {
  cat "$1"/template.cpp >"$2"
cat "$1"/run.cpp >>"$2"
  cp "$1"/debug.cpp .
# Creates a contest with hame $2
# Copies the macro and debug file from $1
# Already creates files a...z .cpp and .py
prepare_contest() {
```

```
mkdir "$2"
  cd "$2"
  for i in {a..z}; do
    touch_macro $1 $i.cpp
  done
}
```

### 1.2 debug

```
template <typename T>
concept Printable = requires(T t) {
  {
     std::cout << t
  } -> std::same_as<std::ostream &>;
template <Printable T>
void __print(const T &x) {
  cerr << x;</pre>
template <size_t T>
void __print(const bitset<T> &x) {
  cerr << x;</pre>
template <typename A, typename B>
void __print(const pair<A, B> &p);
template <typename... A>
void __print(const tuple<A...> &t);
template <typename T>
void __print(stack<T> s);
template <typename T>
void __print(queue<T> q);
template <typename T, typename... U>
void __print(priority_queue<T, U...> q);
template <typename A>
void __print(const A &x) {
  bool first = true;
  cerr << '{';</pre>
  for (const auto &i : x) {
  cerr << (first ? "" : ","), __print(i);</pre>
     first = false;
  cerr << '}';
template <typename A, typename B>
void \_\_print(const pair<A, B> &p) {
  cerr << '(
   __print(p.first);
  cerr <<
  __print(p.second);
cerr << ')';</pre>
template <typename... A>
void __print(const tuple<A...> &t) {
  bool first = true;
cerr << '(';</pre>
  apply(
       [&first](const auto &...args) {
    ((cerr << (first ? "" : ","),
           __print(args), first = false),
       },
t);
  cerr << ')';
template <typename T>
void __print(stack<T> s) {
  vector<T> debugVector;
  while (!s.empt\bar{y}()) {
     T t = s.top();
     debugVector.push_back(t);
     s.pop();
  reverse(debugVector.begin(), debugVector.end());
  __print(debugVector);
template <typename T>
void __print(queue<T> q) {
  vector<T> debugVector;
  while (!q.empty()) {
     T t = q.front();
     debugVector.push_back(t);
     q.pop();
```

```
_print(debugVector);
template <typename T, typename... U>
void __print(priority_queue<T, U...> q) {
  vector<T> debugVector;
  while (!q.empty()) {
     T t = q.top();
     debugVector.push_back(t);
     q.pop();
  __print(debugVector);
void _print() { cerr << "]\n"; }
template <typename Head, typename...</pre>
                                              Tail>
void _print(const Head &H, const Tail &...T) {
    _{	extsf{print}}(\mathsf{H});
  if (sizeof...(T)) cerr << ", ";
  _print(T...);
#define dbg(x...) \
cerr << "[" << #x << "] = ["; \
  _print(x)
1.3 run
void run();
int32_t main() {
#ifndef LOCAL
  fastio;
#endif
  int T = 1;
  cin >> T;
  rep(t, 0, T) {
     dbg(t);
     run();
void run() {}
1.4 short-template
                                                                    }
#include <bits/stdc++.h>
using namespace std;
#define fastio
  ios_base::sync_with_stdio(0); \
  cin.tie(0);
void run() {}
int32_t main(void) {
  fastio;
                                                                    syntax on
  int t;
  // cin >> t;
while (t—) run();
                                                                         register
1.5 template
#include <bits/stdc++.h>
using namespace std;
#ifdef LOCAL
#include "debug.cpp"
                                                                    endfunction
#else
#define dbg(...)
#endif
#define endl '\n
#define fastio
  ios_base::sync_with_stdio(0); \
                                                                    2.1
cin.tie(0);
#define all(j) j.begin(), j.end()
#define rall(j) j.rbegin(), j.rend()
#define len(j) (int)j.size()
                                                                    2.1.1
```

#define rep(i, a, b)
for (common\_type\_t<decltype(a), decltype(b)>

for (common\_type\_t<decltype(a), decltype(b)> \

i = (a);

i < (b); i++)
#define rrep(i, a, b)</pre>

```
i = (a);

i > (b); i - )
#define trav(xi, xs) for (auto &xi : xs)
#define rtrav(xi, xs) \
  for (auto &xi : ranges::views::reverse(xs))
using ll = long long;
#define inte l
#define pb push_back
#define pf push_front
#define ppb pop_back
#define ppf pop_front
#define eb emplace_back
#define lb lower_bound
#define ub upper_bound
#define fi first
#define se second
#define emp emplace
#define ins insert
#define divc(a, b) ((a) + (b)-1ll) / (b)
using str = string;
using ull = unsigned long long;
using ld = long double;
using vll = vector<ll>
using pll = pair<ll, ll>;
using vll2d = vector<vll>;
using vi = vector<int>
using vi2d = vector<vi>,
using pii = pair<int, int>;
using vpii = vector<pii>;
using vc = vector<char>;
using vs = vector<str>
template <typename T, typename T2>
using umap = unordered_map<T, T2>;
template <typename T>
using pgmn =
    priority_queue<T, vector<T>, greater<T>>;
template <typename T>
using pqmx = priority_queue<T, vector<T>>;
template <typename T, typename U>
inline bool chmax(T &a, U const &b) {
  return (a < b ? a = b, 1 : 0);
template <typename T, typename U>
inline bool chmin(T &a, U const &b) {
  return (a > b ? a = b, 1 : 0);
1.6 vim config
set sta nu rnu sc cindent
set ts=2 sw=2
set bg=dark ruler clipboard=unnamed,unnamedplus,
    timeoutlen=100
colorscheme default
" Takes the hash of the selected text and put
" in the vim clipboard
function! HashSelectedText()
    " Yank the selected text to the unnamed
    normal! gvy
    " Use the system() function to call sha256sum
    with the yanked text
    " Yank the hash into Vim's unnamed register
    let @" = l:hash
   Data Structures
     SQRT decomposition
      two-sequence-queries
using ll = long long;
const ll\ MOD = 998244353;
inline ll sum(const ll a, const ll b) {
  return (a + b) % MOD;
}
```

ll sub(const ll a, const ll b) {

```
return (a - b + MOD) % MOD;
inline ll mul(const ll a, const ll b) {
  return (a * b) % MOD;
struct SqrtDecomposition {
  struct t_sqrt {
     int l, r;
     ll x, y;
     ll prod;
     ll sum_as, sum_bs;
     t_sqrt() {
       l = numeric_limits<int>::max();
r = numeric_limits<int>::min();
       x = y = prod = sum_as = sum_bs = 0;
     };
  };
  int sqrtLen;
  vector<t_sqrt> blocks;
vector<ll> as, bs;
  SqrtDecomposition(const vector<ll> &as_,
                         const vector<ll> &bs_) {
     int n = as_.size();
sqrtLen = (int)sqrt(n + .0) + 1;
     blocks.resize(sqrtLen + 6.66);
     as = as_{-};
     bs = bs_{-};
     for (int i = 0; i < n; i++) {
  auto &bi = blocks[i / sqrtLen];</pre>
       bi.l = min(bi.l, i);
bi.r = max(bi.r, i);
bi.sum_as = sum(bi.sum_as, as[i]);
bi.sum_bs = sum(bi.sum_bs, bs[i]);
       bi.prod = sum(bi.prod, mul(as[i], bs[i]));
  }
  // adds x to a[i], and y to b[i], in range [l,
  // rl
  void update(int l, int r, ll x, ll y) {
    auto apply1 = [\&](int idx, ll x,
ll y) \rightarrow void {
       auto &block = blocks[idx / sqrtLen];
       block.prod =
            sub(block.prod, mul(as[idx], bs[idx]));
       block.sum_as = sub(block.sum_as, as[idx]);
block.sum_bs = sub(block.sum_bs, bs[idx]);
       as[idx] = sum(as[idx], x);
       bs[idx] = sum(bs[idx], y);
       block.prod =
            sum(block.prod, as[idx] * bs[idx]);
       block.sum_as = sum(block.sum_as, as[idx]);
block.sum_bs = sum(block.sum_bs, bs[idx]);
     auto apply2 = [\&](int idx, ll x,
                            ll y) -> void {
       blocks[idx].x = sum(blocks[idx].x, x)
       blocks[idx].y = sum(blocks[idx].y, y);
     };
     int cl = l / sqrtLen, cr = r / sqrtLen;
     if (cl == cr) {
       for (int i = l; i <= r; i++) {
   apply1(i, x, y);
     } else {
  for (int i = l; i <= (cl + 1) * sqrtLen - 1;</pre>
              i++) {
          apply1(i, x, y);
       for (int i = cl + 1; i \le cr - 1; i++) {
          apply2(i, x, y);
       for (int i = cr * sqrtLen; i \le r; i++) {
          apply1(i, x, y);
  }
  // sum of a[i]*b[i] in range [l r]
  il query(int l, int r) {
  auto eval1 = [&](int idx) -> ll {
```

```
auto eval2 = [\&](int idx) \rightarrow ll {
      auto &block = blocks[idx];
       ll ret = 0;
       ret =
           sum(ret
                mul(mul(block.x, block.y)
                    sum(sub(block.r, block.l), 1)));
      ret = sum(ret, block.prod);
ret = sum(ret, block.y * block.sum_as);
       ret = sum(ret, block.x * block.sum_bs);
       return ret;
    };
    ll ret = 0;
int cl = l / sqrtLen, cr = r / sqrtLen;
    if (cl == cr) {
       for (int i = l; i \le r; i++) {
         ret = sum(ret, eval1(i));
    } else {
  for (int i = l; i <= (cl + 1) * sqrtLen - 1;</pre>
         ret = sum(eval1(i), ret);
      }
       for (int i = cl + 1; i \le cr - 1; i++) {
         ret = sum(ret, eval2(i));
      for (int i = cr * sqrtLen; i \le r; i++) {
         ret = sum(ret, eval1(i));
    }
    return ret;
};
2.2
    Segment tree (dynamic)
2.2.1 Range Max Query Point Max Assignment
Description: Answers range queries in ranges until 10<sup>9</sup> (maybe more)
Time: Query and update O(n \cdot \log n)
struct node;
node *newNode();
struct node {
  node *left, *right;
  int lv, rv;
  ll val;
  node() : left(NULL), right(NULL), val(—oo) {}
  inline void init(int l, int r) {
  inline void extend() {
    if (!left) {
      int m = (lv + rv) / 2;
left = newNode();
       right = newNode();
       left->init(lv, m);
       right—>init(m' + 1, rv);
  ll query(int l, int r) {
  if (r < lv || rv < l) {</pre>
      return 0;
    if (l <= lv && rv <= r) {
       return val;
    extend();
    return max(left->query(l, r);
                 right—>query(l, r));
  void update(int p, ll newVal) {
    if (lv == rv)
       val = max(val, newVal);
```

auto &block = blocks[idx / sqrtLen];

return;

```
extend():
    (p <= left->rv ? left : right)
        ->update(p, newVal);
    val = max(left->val, right->val);
 }
};
const int BUFFSZ(1e7);
node *newNode()
  static int bufSize = BUFFSZ
  static node buf[(int)BUFFSZ];
assert(bufSize);
  return &buf[—bufSize];
struct SegTree {
  int n;
  node *root;
  SegTree(int _n) : n(_n) {
    root = newNode();
    root—>init(0, n);
  [[] fl query(int l, int r] {
    return root—>query(l, r);
  void update(int p, ll v) { root->update(p, v); }
2.2.2 Range Sum Query Point Sum Update
Description: Answers range queries in ranges until 10<sup>9</sup> (maybe more)
Time: Query and update in O(n \cdot \log n)
struct node:
node *newNode();
struct node {
  node *left, *right;
  int lv, rv;
  ll val;
  node() : left(NULL), right(NULL), val(0) {}
  inline void init(int l, int r) {
    rv = r:
  inline void extend() {
    if (!left) {
      int m = (rv - lv) / 2 + lv;
      left = newNode()
      right = newNode();
      left->init(lv, m);
      right—>init(m + 1, rv);
  }
  ll query(int l, int r) {
    if (r < lv || rv < l) {
      return 0;
    if (l <= lv && rv <= r) {
      return val;
    extend();
    return left->query(l, r) + right->query(l, r);
  void update(int p, ll newVal) {
    if (lv == rv)
      val += newVal;
      return;
    extend()
    (p <= left->rv ? left : right)
         ->update(p, newVal);
    val = left->val + right->val;
const int BUFFSZ(1.3e7);
node *newNode()
  static int bufSize = BUFFSZ
  static node buf[(int)BUFFSZ];
  // assert(bufSize)
  return &buf[—bufSize];
struct SegTree {
```

```
int n;
  node *root:
  SegTree(int _n) : n(_n) {
    root = newNode();
    root->init(0, n);
  ll query(int l, int r) {
    return root—>query(l, r);
  void update(int p, ll v) { root->update(p, v); }
};
2.3
     Segment tree point update range query
2.3.1 Query GCD (bottom up)
using ll = long long;
struct Node {
  ll value
  bool undéf;
  Node()
         value(1), undef(1){};
                                    // Neutral element
  Node(ll v) : value(v), undef(0){};
inline Node combine(const Node &nl, const Node &nr) {
  if (nl.undef) return nr;
if (nr.undef) return nl;
  Node m:
  m.value = gcd(nl.value, nr.value);
  m.undef = false;
  return m;
}
template <typename T = Node, auto F = combine>
struct SegTree {
  int n;
  vector<T> st;
  SegTree(int _n) : n(_n), st(n \ll 1) {}
  void assign(int p, const T &k) {
  for (st[p += n] = k; p >>= 1;)
       st[p] = F(st[p << 1], st[p << 1 | 1]);
  T query(int l, int r) {
    T ansl, ansr;
for (l += n, r += n + 1; l < r;
l >>= 1, r >>= 1) {
       if (l \& 1) ans l = F(ansl, st[l++]);
      if (r \& 1) ansr = F(st[-r], ansr);
    return F(ansl, ansr);
};
2.3.2 Query hash (top down)
const ll MOD = 1'000'000'009;
const ll P = 31;
const int MAXN = 2'000'000;
ll pows[MAXN + 1];
void computepows() {
  pows[0] = 1;
  for (int i = 1; i <= MAXN; i++) {
   pows[i] = (pows[i - 1] * P) % MOD;
}
struct Node {
  ll hash;
  Node()
          : hash(-1){};
                           // Neutral element
  Node(ll\ v) : hash(v){};
};
inline Node combine(Node &vl, Node &vr, int nl,
                       int nr, int ql, int qr) {
     (vl.hash == -1) return vr;

(vr.hash == -1) return vl;
  Node vm;
  int nm = midpoint(nl, nr);
  int lsize = min(nm, qr) - max(nl, ql) + 1;
```

vm.hash = (vl.hash +

((vr.hash \* pows[lsize]) % MOD)) %

```
MOD:
  return vm:
template <typename T = Node, auto F = combine>
struct SegTree {
  int n;
  vector<T> st;
  SegTree(int n) : n(n), st(n << 2) {}
  void assign(int p, const T &v) {
    assign(\bar{1}, 0, \bar{n} – 1, p, v);
  void assign(int node, int l, int r, int p,
               const T &v) {
    if (l == r) {
      st[node] = v;
      return;
    int m = midpoint(l, r);
    if (p \ll m)
      assign(node << 1, l, m, p, v);
      assign(node << 1 | 1, m + 1, r, p, v);
    st[node] = F(st[node << 1], st[node << 1 | 1],
                   l, r, l, r);
  inline T query(int l, int r) {
    return query(1, 0, n - 1, l, r);
  inline T query(int node, int nl, int nr, int l,
                  int r) const {
    if (r < nl or nr < l) return T();
if (l <= nl and nr <= r) return st[node];</pre>
    int m = midpoint(nl, nr);
    auto a = query(node \ll 1, nl, m, l, r);
         query(node << 1 | 1, m + 1, nr, l, r);
    return F(a, b, nl, nr, l, r);
};
2.3.3 Query max subarray sum (bottom up)
struct Node {
   ll tot, suf, pref, best;
  // Neutral element
  Node()
        tot(-oo),
         suf(-oo)
         pref(-oo)
         best(-oo) {} // Neutral element
  // for assign
  Node(ll x) {
  tot = x, suf = x, pref = x,
    best = max(0ll, x);
};
Node combine(Node &nl, Node &nr) {
  if (nl.tot == -oo) return nr;
     (nr.tot == -oo) return nl;
  Node m;
  m.tot = nl.tot + nr.tot;
m.pref = max({nl.pref, nl.tot + nr.pref});
  m.suf = max({nr.suf, nr.tot + nl.suf});
  m.best =
      max({nl.best, nr.best, nl.suf + nr.pref});
  return m;
2.3.4 Query min (bottom up)
```

struct Node {

ll value;

value(numeric\_limits<

Node(ll v) : value(v){};

Node combine(Node &l, Node &r) {

ll>::max()){};

Node()

Node m;

};

# m.value = min(l.value, r.value); return m: } template <typename T = Node, auto F = combine> struct SegTree { int n; vector<T> st; SegTree(int $_n$ ) : $n(_n)$ , st( $n \ll 1$ ) {} void assign(int p, const T &k) { for (st[p += n] = k; p >>= 1;) st[p] = F(st[p << 1], st[p << 1 | 1]); T query(int l, int r) { return F(ansl, ansr); } }; 2.3.5 Query sum (bottom up) struct Node { ll value; Node() : value(0){}; // Neutral element Node(ll v) : value(v){}; inline Node combine(const Node &nl const Node &nr) { m.value = nl.value + nr.value; return m: } struct SegTree { int n; vector<Node> st; SegTree(int $_n$ ) : $n(_n)$ , st( $n \ll 1$ ) {} void assign(int p, const Node &k) { for (st[p += n] = k; p >>= 1;) st[p] = combine(st[p << 1], st[p << 1 | 1]);Node query(int l, int r) { Node ansl = Node(), ansr = Node(); for (l += n, r += n + 1; l < r; l >>= 1, r >>= 1) { if (l & 1) ansl = combine(ansl, st[l++]); if (r & 1) ansr = combine(st[—r], ansr); return combine(ansl, ansr); }; Segment tree range update range query 2.4.1 Arithmetic progression sum update, query sum $\label{eq:Description: Makes arithmetic progression updates in range and sum queries.} \\ \textbf{Usage: Considering } PA(A,R) = [A+R,A+2R,A+3R,\ldots] \\ \bullet \ \ \textbf{update\_set(l, r, A, R): sets} \ [l, r] \ \ \textbf{to} \ PA(A,R) \\ \\$ • update\_add(l, r, A, R): sum PA(A, R) in [l, r] • query(l, r): sum in range [l, r] Time: build O(N), updates and queries $O(\log N)$ const ll oo = 1e18;struct SegTree { struct Data { ll sum; ll set\_a, set\_r, add\_a, add\_r; Data() : sum(0),

int n;

// Neutral element

set\_a(oo),

 $add_r(0)$  {}

: n(n\_), seg(vector<Data>(4 \* n)) {}

 $set_r(0)$ , add\_a(0),

vector<Data> seg;

SegTree(int  $n_{-}$ )

```
void prop(int p, int l, int r) {
  int sz = r - l + 1;
  ll &sum = seg[p].sum, &set_a = seg[p].set_a,
                                                                ĺl query(int a, int b, int p, int l, int r) {
                                                                 prop(p, l, r);
if (b < l or r < a) return 0;
if (a <= l and r <= b) return seg[p].sum;</pre>
      &set_r = seg[p].set_r,
                                                                 add_a = seg[p].add_a
      \&add_r = seg[p].add_r;
  if (set_a != oo) {
    set_a += add_a, set_r += add_r;
                                                               il query(int l, int r) {
  return query(l, r, 1, 0, n - 1);
     sum =
         set_a * sz + set_r * sz * (sz + 1) / 2;
    if (l != r) {
       int m = (l + r) / 2;
                                                             };
       seg[2 * p].set_a = set_a;
       seg[2 * p].set_r = set_r;
                                                             2.4.2 Increment update query min & max (bottom up)
       seg[2 * p].add_a = seg[2 * p].add_r = 0;
       seg[2 * p + 1].set_a =
                                                             using SeqT = ll;
       set_a + set_r * (m - l + 1);
seg[2 * p + 1].set_r = set_r;
seg[2 * p + 1].add_a =
                                                             struct QueryT {
                                                                SegT mx, mn;
           seq[2 * p + 1].add_r = 0;
                                                                QueryT()
                                                                    : mx(numeric_limits<SegT>::min()),
                                                                      mn(numeric_limits<SegT>::max()) {}
    set_a = oo, set_r = 0;
    add_a = add_r = 0;
                                                               QueryT(SegT _v) : mx(_v), mn(_v) {}
  } else if (add_a or add_r) {
  sum +=
                                                             };
                                                             inline QueryT combine(QueryT ln, QueryT rn,
         add_a * sz + add_r * sz * (sz + 1) / 2;
    if (l != r) {
 int m = (l + r) / 2;
                                                                                       pii lr1, pii lr2) {
                                                               chmax(ln.mx, rn.mx);
chmin(ln.mn, rn.mn);
       seg[2 * p].add_a += add_a;
                                                                return ln;
       seg[2 * p].add_r += add_r;
                                                             }
       seg[2 * p + 1].add_a +=
                                                             using LazyT = SegT;
       add_a + add_r * (m - l + 1);
seg[2 * p + 1].add_r += add_r;
                                                             inline QueryT applyLazyInQuery(QueryT q, LazyT l,
                                                                                                 pii ĺr)
                                                                if (q.mx == QueryT().mx) q.mx = SegT();
    add_a = add_r = 0;
                                                               if (q.mn == QueryT().mn) q.mn = SegT();
                                                               q.mx += l, q.mn += l;
}
                                                                return q;
int inter(pii a, pii b) {
                                                             }
  if (a.first > b.first) swap(a, b);
  return max(
                                                             inline LazyT applyLazyInLazy(LazyT a, LazyT b) {
       0, min(a.second, b.second) - b.first + 1);
                                                                return a + b;
ll set(int a, int b, ll aa, ll rr, int p, int l,
                                                             using UpdateT = SegT;
        int_r) {
  prop(p, l, r);
if (b < l or r < a) return seg[p].sum;</pre>
                                                             inline QueryT applyUpdateInQuery(QueryT q,
                                                                                                    UpdateT u,
                                                                                                    pii lr) {
  if (a <= l and r <= b) {
    seg[p].set_a = aa;</pre>
                                                                if (q.mx == QueryT().mx) q.mx = SegT();
                                                               if (q.mn == QueryT().mn) q.mn = SegT();
q.mx += u, q.mn += u;
    seg[p].set_r = rr;
    prop(p, l, r);
                                                                return q;
    return seg[p].sum;
                                                             }
  int m = (l + r) / 2;
int tam_l = inter({l, m}, {a, b});
                                                             inline LazyT applyUpdateInLazy(LazyT l, UpdateT u,
  return seg[p].sum =
                                                                return l + u;
               }
                                                             template <typename Qt = QueryT,
                                                                         typename Lt = LazyT
                                                                         typename Ut = UpdateT, auto C = combine,
void update_set(int l, int r, ll aa, ll rr) { set(l, r, aa, rr, 1, 0, n - 1);
                                                                         auto ALQ = applyLazyInQuery,
                                                                         auto ALL = applyLazyInLazy,
                                                                         auto AUQ = applyUpdateInQuery,
ĺl add(int a, int b, ll aa, ll rr, int p, int l,
                                                                         auto AUL = applyUpdateInLazy>
        int r)
                                                             struct LazySegmentTree {
  prop(p, l, r);
                                                               int n, h;
vector<Qt> ts;
  if (b < l or r < a) return seg[p].sum;</pre>
  if (a <= l and r <= b) {
                                                                vector<Lt> ds;
    seg[p].add_a += aa;
                                                                vector<pii> lrs;
    seg[p].add_r += rr;
                                                                LazySegmentTree(int _n)
    prop(p, l, r);
                                                                    : n(_n),
     return seg[p].sum;
                                                                      h(sizeof(int) * 8 - __builtin_clz(n)),
                                                                       ts(n \ll 1),
  int m = (l + r) / 2;
int tam_l = inter({l, m}, {a, b});
                                                                      ds(n),
                                                                  lrs(n << 1) {
rep(i, 0, n) lrs[i + n] = {i, i};}
  return seg[p].sum =
               add(a, b, aa, rr, 2 * p, l, m) +
                                                                  rrep(i, n - 1, 0) {
               add(a, b, aa + rr * tam_l, rr,
2 * p + 1, m + 1, r);
                                                                    lrs[i] = {lrs[i << 1].first,</pre>
                                                                                lrs[i << 1 | 1].second};</pre>
void update_add(int l, int r, ll aa, ll rr) {
   add(l, r, aa, rr, 1, 0, n - 1);
                                                                LazySegmentTree(const vector<Qt> &xs)
```

```
: LazySegmentTree(len(xs)) {
   copy(all(xs), ts.begin() + n);
rep(i, 0, n) lrs[i + n] = {i, i};
rrep(i, n - 1, 0) {
      ts[i] = C(ts[i << 1], ts[i << 1 | 1]
                     lrs[i << 1], lrs[i << 1 | 1]);
}
void set(int p, Qt v) {
   ts[p + n] = v;
   build(p + n);
void upd(int l, int r, Ut v) {
   l += n, r += n + 1;
int l0 = l, r0 = r;
for (; l < r; l >>= 1, r >>= 1) {
      if (l & 1) apply(l++, v);
if (r & 1) apply(—r, v);
   build(l0), build(r0 - 1);
Qt qry(int l, int r) {
    l += n, r += n + 1;
    push(l), push(r - 1);
   Qt resl = Qt(), resr = Qt();
pii lr1 = {l, l}, lr2 = {r, r};
for (; l < r; l >>= 1, r >>= 1) {
   if (l & 1)
         resl = C(resl, ts[l], lr1, lrs[l]), l++;
      if (r & 1)
         \dot{r}, resr = C(ts[r], resr, lrs[r], lr2);
   return C(resl, resr, lr1, lr2);
void build(int p) {
   while (p > 1) {
      p >>= 1;
      ts[p] = ALQ(C(ts[p << 1], ts[p << 1 | 1], lrs[p << 1 | 1]),
                         ds[p], lrs[p]);
}
void push(int p) {
   rrep(s, h, 0) {
      int i = p >> s;
if (ds[i] != Lt()) {
  apply(i << 1, ds[i]),</pre>
               apply(i << 1 | 1, ds[i]);
         ds[i] = Lt();
}
inline void apply(int p, Ut v) {
  ts[p] = AUQ(ts[p], v, lrs[p]);
  if (p < n) ds[p] = AUL(ds[p], v, lrs[p]);</pre>
```

### 2.4.3 Increment update sum query (top down)

```
struct Lnode {
  ll v;
bool assign;
  Lnode() : v(), assign() {} // Neutral element
Lnode(ll _v, bool a = 0) : v(_v), assign(a){};
using Qnode = ll;
using Unode = Lnode;
struct LSegTree {
  int n, ql, qr;
  vector<Qnode> st;
vector<Lnode> lz;
  Qnode merge(Qnode lv, Qnode rv, int nl,
                  int nr) {
     return lv + rv;
  }
  void prop(int i, int l, int r) {
     if (lz[i].assign) {
```

```
st[i] = lz[i].v * (r - l + 1);

if(l!= r) lz[tol(i)] = lz[tor(i)] = lz[i];
    } else {
       st[i] += lz[i].v * (r - l + 1);
if (l != r)
         lz[i] = Lnode();
  void applyV(int i, Unode v) {
    if (v.assign) {
       lz[i] = v;
    } else {
    lz[i].v += v.v;
  /*-
  LSegTree() {}
  LSegTree(int _n)
: n(_n), st(_n << 2), lz(_n << 2) {}
  bool disjoint(int l, int r) {
     return qr < l or r < ql;
  bool contains(int l, int r) {
    return gl <= l and r <= gr;
  int tol(int i) { return i << 1; }
int tor(int i) { return i << 1 | 1; }</pre>
  void build(vector<Qnode> &v) {
    build(v, 1, 0, n - 1);
  void build(vector<Qnode> &v, int i, int l,
               int r) {
    if (l == r)
       st[i] = v[\bar{l}];
       return;
    int m = midpoint(l, r);
    build(v, tol(i), l, m);
build(v, tor(i), m + 1, r);
st[i] = merge(st[tol(i)], st[tor(i)], l, r);
  void upd(int l, int r, Unode v) {
    ql = l, qr = r;
upd(1, 0, n — 1, v);
  void upd(int i, int l, int r, Unode v) {
    prop(i, l, r);
    if (disjoint(l, r)) return;
    if (contains(l, r)) {
       applyV(i, v);
prop(i, l, r);
       return;
     int m = midpoint(l, r);
    upd(tol(i), l, m, v);
upd(tor(i), m + 1, r, v);
    st[i] = merge(st[tol(i)], st[tor(i)], l, r);
  Qnode qry(int l, int r) {
    ql = l, qr = r
     return qry(1, 0, n - 1);
  Qnode gry(int i, int l, int r) {
    prop(i, l, r);
    if (disjoint(l, r)) return Qnode();
if (contains(l, r)) return st[i];
     int m = midpoint(l, r);
     return merge(qry(tol(i), l, m),
                    qry(tor(i), m + 1, r), l, r);
};
```

### Bitree 2D

Description: Given a 2D array you can increment an arbitrary position, and also query the subsum of a subgrid **Time**: Update and query in  $O(log N^2)$ 

```
struct Bit2d {
  int n;
vll2d bit;
  Bit2d(int ni) : n(ni), bit(n + 1, vll(n + 1)) \{\}
```

```
Bit2d(int ni, vll2d &xs)
    : n(ni), bit(n + 1, vll(n + 1)) {
    for (int i = 1; i <= n; i++) {
        for (int j = 1; j <= n; j++) {
            update(i, j, xs[i][j]);
        }
    }
}
void update(int x, int y, ll val) {
    for (; x <= n; x += (x & (-x))) {
        for (int i = y; i <= n; i += (i & (-i))) {
            bit[x][i] += val;
        }
    }
}
ll sum(int x, int y) {
    ll ans = 0;
    for (int i = x; i; i -= (i & (-i))) {
        for (int j = y; j; j -= (j & (-j))) {
            ans += bit[i][j];
        }
    return ans;
}
ll query(int x1, int y1, int x2, int y2) {
    return sum(x2, y2) - sum(x2, y1 - 1) -
            sum(x1 - 1, y2) + sum(x1 - 1, y1 - 1);
};</pre>
```

### 2.6 Convex Hull Trick / Line Container

**Description**: Container where you can add lines of the form mx + b, and query the maximum value at point x.

Usage:  $insert\_line(m, b)$  inserts the line  $m \cdot x + b$  in the container.

eval(x) find the highest value among all lines in the point x.

**Time**: Eval and insert in  $O(\log N)$ 

```
const ll LLINF = 1e18;
const ll is_query = -LLINF;
struct Line {
  ll m, b;
  mutable function<const Line *()> succ;
  bool operator<(const Line &rhs) const {
  if (rhs.b != is_query) return m < rhs.m;</pre>
    const Line *s = succ();
    if (!s) return 0;
ll x = rhs.m;
    return b - s \rightarrow b < (s \rightarrow m - m) * x;
  }
struct Cht : public multiset<Line> {
                                             // maintain
                                             // max m*x+b
  bool bad(iterator y) {
    auto z = next(y)
    if (y == begin()) {
          (z == end()) return 0;
       return y\rightarrow m == z\rightarrow m \&\& y\rightarrow b <= z\rightarrow b;
    auto x = prev(y);
    if (z == end())
       return y->m == x->m && y->b <= x->b;
     return (ld)(x\rightarrowb - y\rightarrowb) * (z\rightarrowm - y\rightarrowm) >=
             (ld)(y-b-z-b)*(y-m-x-m);
  void insert_line(
       ll m,
       il b) {
                 // min -> insert (-m,-b) -> -eval()
    auto y = insert({m, b});
    y -> succ = [=] {
       return next(y) == end() ? 0 : &*next(y);
    if (bad(y)) {
       erase(y);
       return;
    while (next(y) != end() && bad(next(y)))
       erase(next(y));
    while (y != begin() \&\& bad(prev(y)))
       erase(prev(y));
  ĺl eval(ll x) {
    auto l = *lower_bound((Line){x, is_query});
    return l.m * x + l.b;
```

```
};
```

## 2.7 DSU (with rollback)

**Description**: Performs every operation a regular DSU does, but you can roll back to a specific time.

Usage: int t = uf.time(); ...; uf.rollback(t); T

Time:  $O(\log(N))$ 

```
struct RollbackUF {
   vi e;
   vector<pii> st;
   RollbackUF(int n) : e(n, -1) {}
   int size(int x) { return -e[find(x)]; }
int find(int x) {
   int find(int x)
      return e[x] < 0 ? x : find(e[x]);
   int time() {    return len(st);    }
  void rollback(int t) {
  for (int i = time(); i— > t;)
      e[st[i].first] = st[i].second;
      st.resize(t);
   bool join(int a, int b)
     a = find(a), b = find(b);
if (a == b) return false;
     if (e[a] > e[b]) swap(a, b);
     st.push_back({a, e[a]});
st.push_back({b, e[b]});
     e[a] += e[b];
e[b] = a;
     return true;
};
```

### 2.8 DSU / UFDS

 $\textbf{Usage:} \ \ \text{You may discomment the commented parts to find online which nodes belong to each set, it makes the $union\_set$ method cost $O(log^2)$ instead $O(A)$ and $O(A)$ instead $O(A)$ and $O(A)$ instead $O(A)$ and $O(A)$ instead $O(A)$ and $O(A)$ instead $O(A)$ are also as a substitution of the s$ 

```
struct DSU {
  vi ps, sz;
  // vector<unordered_set<int>> sts;
  DSU(int N)
       : ps(N + 1),
         sz(N, 1) /*, sts(N) */
    iota(ps.begin(), ps.end(), 0);
// for (int i = 0; i < N; i++)
// sts[i].insert(i);</pre>
  int find_set(int x)
    return ps[x] == x ? x
                           : ps[x] = find_set(ps[x]);
  int size(int u) { return sz[find_set(u)]; }
bool same_set(int x, int y) {
     return find_set(x) == find_set(y);
  void union_set(int x, int y) {
    if (same_set(x, y)) return;
    int px = find_set(x);
     int py = find_set(y);
     if (sz[px] < sz[py]) swap(px, py);
     ps[py] = px;
    sz[px] += sz[py]
     // sts[px].merge(sts[py]);
};
```

# 2.9 Lichao Tree (dynamic)

**Description**: Lichao Tree that creates the nodes dynamically, allowing to query and update from range [MAXL, MAXR] Usage:

- query(x): find the highest point among all lines in the structure
- add(a,b): add a line of form y = ax + b in the structure
- addSegment(a,b,l,r): add a line segment of form y=ax+b which covers from range [l,r]

Time:  $O(\log N)$ 

```
template <typename T = ll, T MAXL = 0,
    T MAXR = 1 '000' 000'001>
struct LiChaoTree {
  bool first_best(T a, T b) { return a > b; }
T get_best(T a, T b) {
     return first_best(a, b) ? a : b;
  struct line {
     T m, b;
     T operator()(T x) { return m * x + b; }
  struct node {
     line li;
    left(núĺlptr)
            right(nullptr) {}
     ~node()
       delete left;
       delete right;
    }
  };
  node *root;
  LiChaoTree(line li = {0, inf})
: root(new node(li)) {}
~LiChaoTree() { delete root; }
  T query(T x, node *cur, T l, T r) \{
     if (cur == nullptr) return inf;
if (x < l or x > r) return inf;
     T mid = midpoint(l, r);
     T ans = cur \rightarrow li(x);
     ans = get_best(ans)
                       query(x, cur->left, l, mid));
     ans = get_best(
         ans, query(x, cur->right, mid + 1, r));
     return ans;
  Ť query(T x) {
     return query(x, root, MAXL, MAXR);
  void add(line li, node *&cur, T l, T r) {
  if (cur == nullptr) {
       cur = new node(li);
       return;
    if (first_best(li(mid), cur->li(mid)))
    swap(li, cur->li);
if (first_best(li(l), cur->li(l)))
  add(li, cur->left, l, mid);
if (first_best(li(r), cur->li(r)))
  add(li, cur->right, mid + 1, r);
  void add(T m, T b) {
  add({m, b}, root, MAXL, MAXR);
  if (r < lseg || l > rseg) return;
     if (cur == nullptr) cur = new node;
     if (lseg \leftarrow l && r \leftarrow rseg) {
       add(li, cur, l, r);
       return:
     \dot{T} mid = midpoint(l, r);
     if (l != r)
       addSegment(li, cur->left, l, mid, lseg,
                     rseg);
       addSegment(li, cur->right, mid + 1, r, lseg,
                     rseq);
  void addSegment(T a, T b, T l, T r)
    addSegment({a, b}, root, MAXL, MAXR, l, r);
```

### 2.10 Merge sort tree

**Description**: Like a segment tree but each node stores the ordered subsegment it represents. Usage:

```
• inrange(l,r,a,b) : counts the number of positions i,\ l \leq i \leq r such that
     a < x_i < b.
Time: Build O(N \log N^2), inrange O(\log N^2)
Memory: O(n \log N)
template <class T>
struct MergeSortTree {
  int n;
  vector<vector<T>> st;
  MergeSortTree(vector<T> &xs)
        : n(len(xs)), st(n << 1) {
     rep(i, 0, n) st[i + n] = vector<T>({xs[i]});
     rrep(i, n - 1, 0)
       st[i].resize(len(st[i << 1]) +
len(st[i << 1 | 1]));
       merge(all(st[i << 1]), all(st[i << 1 | 1]),
               st[i].begin());
     }
  int count(int i, T a, T b)
     return upper_bound(all(st[i]), b) -
              lower_bound(all(st[i]), a);
  int inrange(int l, int r, T a, T b) {
     int ans = 0;
     for (l += n, r += n + 1; l < r;
        l >>= 1, r >>= 1) {
if (l & 1) ans += count(l++, a, b);
       if (r \& 1) ans += count(—r, a, b);
     }
     return ans;
};
2.11 Mex with update
\bf Description: This DS allows you to mantain an array of elments, insert, and
remove, and query the MEX at any time.
Usage:
   • Mex(mxsz): Initialize the DS, mxsz must be the maximum number of
     elements that the structure may have.
   • add(x): just adds one copy of x.
   • rmv(x): just remove a copy of x.

    operator(): returns the MEX.

   • Mex(mxsz): O(\log mxsz)
• add(x): O(\log mxsz)
• rmv(x): O(\log mxsz)
   • operator(): O(1)
struct Mex {
  int mx_sz;
  vi hs;
  set<int> st;
  Mex(int_mx_sz): mx_sz(_mx_sz), hs(mx_sz + 1) {
     auto it = st.begin();
     rep(i, 0, mx_sz + 1) it = st.insert(it, i);
  void add(int x) {
     if (x > mx_sz) return;
     if (!hs[x]++) st.erase(x);
  void rmv(int x) {
  if (x > mx_sz) return;
  if (!—hs[x]) st.emplace(x);
  int operator()() const { return *st.begin(); }
     Optional, you can just create with size
     len(xs) add N elements :D
  Mex(const vi &xs, int _mx_sz = -1)
    : Mex(~_mx_sz ? _mx_sz : len(xs)) {
    for (auto xi : xs) add(xi);
};
```

## 2.12 Orderd Set (GNU PBDS)

**Usage**: If you need an ordered **multi** set you may add an id to each value. Using greater\_equal, or less\_equal is considered undefined behavior.

- order of key (k): Number of items strictly smaller/greater than k .
- $\operatorname{find}$ \_by\_order(k): K-th element in a set (counting from zero).

```
Time: Both O(\log N)
```

Warning: Is 2 or 3 times slower then a regular set/map.

```
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
template <typename T>
using ordered_set =
    tree<T, null_type, less<T>, rb_tree_tag,
         tree_order_statistics_node_update>;
```

### 2.13 Prefix Sum 2D

**Description**: Given an 2D array with N lines and M columns, find the sum of the subarray that have the left upper corner at (x1, y1) and right bottom corner at  $(x^2, y^2)$ 

**Time**: Build  $O(N \cdot M)$ , Query O(1).

```
template <typename T>
struct psum2d {
  vector<vector<T>> s;
  vector<vector<T>> psum;
  psum2d(vector<vector<T>> &grid, int n, int m)
    : s(n + 1, vector<T>(m + 1)),
psum(n + 1, vector<T>(m + 1)) {
for (int i = 1; i <= n; i++)
       for (int j = 1; j \le m; j++) {
          s[i][j]
               s[i][j-1] + grid[i-1][j-1];
          psum[i][j] = psum[i - 1][j] + s[i][j];
  }
  T query(int x1, int y1, int x2, int y2) {
  T ans = psum[x2 + 1][y2 + 1] + psum[x1][y1];
     ans -= psum[x2 + 1][y1] + psum[x1][y2 + 1];
     return ans;
```

### 2.14 Sparse table

```
template <typename T = ll,
           auto cmp =
                [](T &src1, T &src2, T &dst) {
                  dst = min(src1, src2);
class SparseTable {
 private:
  int sz;
  vi logs;
  vector<vector<T>> st;
 public:
  st.resize(logs[sz] + 1, vector<T>(sz));
    rep(i, 0, sz) st[0][i] = v[i];

for (int k = 1; (1 << k) <= sz; k++) {

  for (int i = 0; i + (1 << k) <= sz; i++) {

    cmp(st[k - 1][i],
                   -1][i + (1 << (k -1))],
             st[k][i]);
    }
  }
T query(int l, int r) {
    const int k = logs[r - l];
    T ret;
    cmp(st[k][l], st[k][r - (1 << k)], ret);
    return ret;
```

#### 2.15Static range queries

```
template <typename T = ll,
          auto op =
               [](const T &src1, const T &src2,
                  T \& dst) { dst = src1 + src2; },
          auto invop =
```

```
[](const T &src1, const T &src2,
    T &dst) { dst = src1 - src2; }>
struct StaticRangeQueries {
   vector<T> acc;
   StaticRangeQueries(const vector<T> &XS)
          acc(len(XS)) {
     acc[0] = XS[0];
rep(i, 1, len(XS))
        op(acc[i-1], XS[i], acc[i]);
  T operator()(int l, int r) {
        T lv = (l ? acc[l - 1] : T());
     T ret:
     invop(acc[r], lv, ret);
     return ret;
};
```

# 2.16 Venice Set

**Description**: A container that you can insert q copies of element e, increment every element in the contaiver in x, query which is the best element and it's quantity and also remove k copies of the greatest element.

Time:

- add elment O(log N) • remove  $O(\log N)$
- update: O(1)
- query *O*(1)

```
template <typename T = ll >
struct VeniceSet {
  using T2 = pair<T, ll>;
  priority_queue<T2, vector<T2>, greater<T2>> pq;
  T acc:
  VeniceSet() : acc() {}
  void add_element(const T &e, const ll q) {
    pq.emplace(e - acc, q);
  void update_all(const T &x) { acc += x; }
  T2 best() {
  auto ret = pq.top();
    ret.first += acc;
    return ret;
  void pop() { pq.pop(); }
  void pop_k(int k) {
    auto [e, q] = pq.top();
pq.pop();
    q = k;
    if (q) pq.emplace(e, q);
};
```

### Venice Set (complete)

Description: A container which you can insert elements update all at once and also make a few queries Usage:

- $add\_element(e, q)$ : adds q copies of e, if no q is provided adds a single
- $update \ all(x)$ : increment every value by x
- $erase(\overline{e})$ : removes every copy of e, and returns how much was removed.
- count(e): returns the number of e in the container
- high()/low(): returns the hightest/lowest element, and it's quantity
- $pop\_low(q)/pop\_high(q)$ : removes q copies of the lowest/highest elements if no q is provided removes all copies of the lowest/highest element.

element. You may answer which is the K-th value and it's quantity using an  $ordered\_set$ . Probably works with other operations

Time: Considering N the number of distinct numbers in the container

•  $add\_element(e,q)$ :  $O(\log(N))$ 

- $update\_all(x):O(1)$ •  $erase(\overline{e})$ :  $O(\log(N))$
- count(e):  $O(\log(N))$
- high()/low(): O(1)
- $pop\_low(q)/pop\_high(q)$ : worst case is  $O(N \cdot \log(N))$  if you remove all elements and so on...

  Warning: There is no error handling if you try to pop more elements than

exists or related stuff

```
struct VeniceSet {
  set<pll> st;
  ll acc;
  VeniceSet() : acc() {}
  ll add_element(ll e, ll q = 1) {
    q += erase(e);
```

```
e -= acc;
    st.emplace(e, q);
    return q;
  void update_all(ll x) { acc += x; }
  ll erase(ll e) {
       –= aċc;
    auto it = st.lb({e, LLONG_MIN});
if (it == end(st) || (*it).first != e)
       return 0;
    ll ret = (*it).second;
    st.erase(it);
    return ret;
  ll count(ll x) {
    auto it = st.lb({x, LLONG_MIN});
if (it == end(st) || (*it).first != x)
       return 0;
     return (*it).second;
  pll high() { return *rbegin(st); }
  pll low() { return *begin(st); }
  void pop_high(ll q = -1) {
    if (q == -1) q = high().second;
    while (q) {
                 eq] = high()
       auto [e,
       st.erase(prev(end(st)));
       if (eq > q) add_element(e, eq - q);
       q = max(0ll, q - eq);
  }
  void pop_low(ll q = -1) {
    if (q == -1) q = low().second;
while (q) {
    auto [e, eq] = low();
       st.erase(st.begin());
       if (eq > q) add_element(e, eq - q);
       q = max(Oll, q - eq);
  }
};
```

### 2.18 Wavelet tree

```
using ll = long long;
template <typename T>
struct WaveletTree {
  struct Node {
    T lo, hi;
int left_child, right_child;
    vector<int> pcnt;
    vector<ll> psum;
    Node(int lo_, int hi_)
: lo(lo_),
hi(hi_),
            left_child(0)
            right_child(0),
            pcnt(),
            psum() {}
  vector<Node> nodes;
  WaveletTree(vector<T> v) {
    nodes.reserve(2 * v.size());
    auto [mn, mx] =
    minmax_element(v.begin(), v.end());
auto build = [&](auto &&self, Node &node,
       auto from, auto to) {
if (node.lo == node.hi or from >= to)
         return:
       auto mid = midpoint(node.lo, node.hi);
       auto f = [&mid](T x) { return x <= mid; };
node.pcnt.reserve(to - from + 1);</pre>
       node.pcnt.push_back(0);
       node.psum.reserve(to
                                  from + 1);
       node.psum.push_back(0);
       T left_upper = node.lo,
         right_lower = node.hi;
       for (auto it = from; it != to; it++) {
         auto value = f(*it);
```

```
node.pcnt.push_back(node.pcnt.back() +
                             value);
        node.psum.push_back(node.psum.back() +
                             *it):
        if (value)
          left_upper = max(left_upper, *it);
        else
          right_lower = min(right_lower, *it);
      }
      auto pivot = stable_partition(from, to, f);
      node.left_child =
    make_node(node.lo, left_upper);
      self(self, nodes[node.left_child], from,
           pivot);
      node.right_child =
          make_node(right_lower, node.hi);
      self(self, nodes[node.right_child], pivot,
    build(build, nodes[make_node(*mn, *mx)],
          v.begin(), v.end());
  if (l > r) return 0;
if (node.lo == node.hi) return node.lo;
      int`lb = node.pcnt[l],
          rb = node.pcnt[r + 1],
left_size = rb - lb;
      return (left_size > k
? self(self,
                          nodes[node.left_child],
                  lb, rb — 1, k)
: self(self,
                          nodes[node.right_child],
                         l - lb, r - rb, k - left_size));
    return f(f, nodes[0], L, R, K);
  pair<int, ll> count_and_sum_in_range(
    if (l > r or node.lo > b or node.hi < a)</pre>
      return {0, 0}; if (a <= node.lo and node.hi <= b)
        node.psum[l])};
      int lb = node.pcnt[l],
      rb = node.pcnt[r + 1];
auto [left_cnt, left_sum] =
      self(self, nodes[node.right_child],
      l - lb, r - rb);
return {left_cnt + right_cnt,
              left_sum + right_sum};
    return f(f, nodes[0], L, R);
  inline int count_in_range(int L, int R, T a,
                             T b) const {
    return count_and_sum_in_range(L, R, a, b)
        .first;
  inline ll sum_in_range(int L, int R, T a,
                          T b) const {
    return count_and_sum_in_range(L, R, a, b)
        .second;
  }
 private:
  int make_node(T lo, T hi) {
    int id = (int)nodes.size();
nodes.emplace_back(lo, hi);
    return id;
};
```

### Dynamic Programming

### Binary Knapsack (bottom up)

Description: Given the points each element have, and it repespective cost, computes the maximum points we can get if we can ignore/choose an element, in such way that the sum of costs don't exceed the maximum cost allowed. Time: O(N\*W) Warning: The vectors VS and WS starts at one, so it need an empty value at

```
const int MAXN(1 '000), MAXCOST(1' 000 * 20); ll dp[MAXN + 1][MAXCOST + 1]; bool ps[MAXN + 1][MAXCOST + 1];
pair<ll, vi> knapsack(const vll &points,
                                 const vi &costs,
                                 int maxCost) {
   int n = len(points)
               1; // ELEMENTS START AT INDEX 1 !
   for (int m = 0; m <= maxCost; m++) {
     dp[0][m] = 0;
   for (int i = 1; i \le n; i++) { dp[i][0] = dp[i - 1][0] +
                      (costs[i] == 0) * points[i];
      ps[i][0] = costs[i] == 0;
  for (int i = 1; i <= n; i++) {
  for (int m = 1; m <= maxCost; m++) {
    dp[i][m] = dp[i - 1][m], ps[i][m] = 0;
    int w = costs[i];
    int m = costs[i];</pre>
         ll v = points[i];
         if (w \le m \text{ and } w \le m)
            dp[i-1][m-w] + v > dp[i][m])  { dp[i][m] = dp[i-1][m-w] + v,
            ps[i][m] = 1;
        }
     }
   }
   vi is;
   for (int i = n, m = maxCost; i >= 1; —i) {
     if (ps[i][m]) {
         is.emplace_back(i);
         m -= costs[i];
   }
   return {dp[n][maxCost], is};
```

### 3.2 Edit Distance

Time: O(N \* M)

```
int edit_distance(const string &a,
                          const string &b) {
   int n = a.size();
   int m = b.size();
  vector<vi> dp(n + 1, vi(m + 1, 0));
  int ADD = 1, DEL = 1, CHG = 1;
for (int i = 0; i <= n; ++i) {
  dp[i][0] = i * DEL;</pre>
   for
        (int i = 1; i \le m; ++i) {
      dp[0][i] = ADD * i;
  for (int i = 1; i <= n; ++i) {
  for (int j = 1; j <= m; ++j) {
    int add = dp[i][j - 1] + ADD;
    int del = dp[i - 1][j] + DEL;
}</pre>
        int cha =
              dp[i-1][j-1] +
               (a[i-1]) = b[j-1] ? 0 : 1) * CHG;
        dp[i][j] = min({add, del, chg});
  }
  return dp[n][m];
```

### 3.3 Knapsack

Description: Finds the maximum score you can achieve, given that you have N items, each item has a cost, a point and a quantity, you can spent at most maxcost and buy each item the maximum quantity it has.

```
Time: O(n \cdot maxcost \cdot \log maxqtd)
Memory: O(maxcost).
```

```
ll knapsack(const vi &weight, const vll &value,
               const vi &qtd, int maxCost) {
  vi costs;
  vll values;
for (int i = 0; i < len(weight); i++) {</pre>
     ll q = qtd[i];
     for (ll x = 1; x <= q; q -= x, x <<= 1) {
  costs.eb(x * weight[i]);</pre>
       values.eb(x * value[i]);
     if (q) {
        costs.eb(q * weight[i]);
        values.eb(q * value[i]);
  }
  vll dp(maxCost + 1);
for (int i = 0; i < len(values); i++) {
  for (int j = maxCost; j > 0; j—) {
       if (j \ge costs[i])
          dp[j] = max(dp[j],
                          values[i] + dp[j - costs[i]]);
  return dp[maxCost];
```

### 3.4 Longest Increasing Subsequence

**Description**: Find the pair (sz, psx) where sz is the size of the longest subsequence and psx is a vector where  $psx_i$  tells the size of the longest increase subsequence that ends at position  $i.\ get_i dx$  just tells which indices could be in the longest increasing subsequence. Time:  $O(n \log n)$ 

```
template <typename T>
pair<int, vi> lis(const vector<T> &xs, int n) {
  vector<T> dp(n + 1, numeric_limits<T>::max());
  dp[0] = numeric_limits<T>::min();
  int sz = 0;
  vi psx(n);
  rep(i, 0, n) {
    int pos =
        lower_bound(all(dp), xs[i]) - dp.begin();
    sz = max(sz, pos);
    dp[pos] = xs[i];
    psx[i] = pos;
  return {sz, psx};
}
template <typename T>
vi get_idx(vector<T> xs) {
  int n = xs.size()
  auto [sz1, psx1] = lis(xs, n);
 auto [sz2, psx2] = lis(xs, n);
  vi ans;
  rep(i, 0, n) {
  int l = psx1[i];
    int r = psx2[n - i - 1];
    if (l + r - 1 == sz1) ans.eb(i);
  return ans;
}
```

#### 3.5Monery sum

**Description**: Find every possible sum using the given values only once.

```
set<int> money_sum(const vi &xs) {
  using vc = vector<char>;
  using vvc = vector<vc>
  int _m = accumulate(all(xs), 0);
  int _n = xs.size();
vvc _dp(_n + 1, vc(_m + 1, 0));
  set<int> _ans;
```

```
_dp[0][xs[0]] = 1;
for (int i = 1; i < _n; ++i) {
  for (int j = 0; j <= _m; ++j) {
    if (j == 0 or _dp[i - 1][j]) }</pre>
            _dp[i][j + xs[i]] = 1;
_dp[i][j] = 1;
    }
}
for (int i = 0; i < _n; ++i)
for (int j = 0; j <= _m; ++j)
        if (_dp[i][j]) _ans.insert(j);
return _ans;
```

### 3.6 Steiner tree

```
template <typename T>
T steinerCost(const vector<vector<T>> &adj,
                   const vi ks,
T inf = numeric_limits<T>::max()) {
  int k = len(ks), n = len(adj);
  vector<vector<T>> dp(n, vector<T>(1 << k, inf));</pre>
  vi inks(n);
  trav(ki, ks) inks[ki] = 1;
  trav(ķi, ks)
     rep(j, 0, n) {
        if (count(all(ks), j) == 0)
           dp[j][1 << ki] = adj[ki][j];
     }
  rep(mask, 2, (1 << k)) {
  rep(i, 0, n) {
    if (inks[i]) continue;
}</pre>
        for (int mask2 = (mask - 1) \& mask;
               mask2 >= 1;
           \begin{array}{l} \text{mask2} = (\text{mask2} - 1) \& \text{mask}) \; \{\\ \text{int mask3} = \text{mask } ^* \text{mask2};\\ \text{chmin(dp[i][mask],} \end{array}
                   dp[i][mask2] + dp[i][mask3]);
        rep(j, 0, n) {
  if (inks[j]) continue;
           chmin(dp[j][mask],
                   dp[i][mask] + adj[i][j]);
     }
  }
T ans = inf;
  rep(i, 0, n) chmin(ans, dp[i][(1 << k) - 1]);
  return ans;
```

### 3.7 Sum of Subsets

**Description**: Allows you to find if some mask X is a super mask of any of the

Usage: Call build with the masks then it returns a vector of bool V where  $V_X$  says if X is a super mask of any of the initial maks

You can change it to count how many submasks of each mask exsists, by changing the bitwise or by a plus sign... Time:  $O(LOG \cdot 2^{LOG})$ 

Memory:  $O(LOG^2 \cdot 2^{LOG})$ 

Warning: Remember to set LOG with the highest bit possible

```
const int LOG = 20;
vc build(const vi &masks) {
  vc ret(1 << LOG);
  trav(mi, masks) ret[mi] = 1;
rep(b, 0, LOG) {
    rep(mask, 0, (1 << LOG)) {</pre>
      if (mask & (1 << b))
         ret[mask] |= ret[mask ^ (1 << b)];
    }
  return ret;
```

# 3.8 Travelling Salesman Problem

```
Time: O(N^2 \cdot 2^N)
Memory: O(N^2 \cdot 2^N)
```

```
vll2d dist;
vll memo:
int tsp(int i, int mask, int N) {
   if (mask == (1 << N) - 1) return dist[i][0];
   if (memo[i][mask] != -1) return memo[i][mask];
   int ans = INT_MAX << 1;
   for (int j = 0; j < N; ++j) {
      if (mask & (1 << j)) continue;
   }
}</pre>
         auto t =
                  tsp(j, mask | (1 << j), N) + dist[i][j];
         ans = min(ans, t);
     return memo[i][mask] = ans;
```

#### 4 Extras

### 4.1 Binary to gray

```
string binToGray(string bin)
   string gray(bin.size(),
   int n = bin.size() - 1;
   gray[0] = bin[0];
   gray[0] - bin[0],

for (int i = 1; i <= n; i++) {

   gray[i] = '0' + (bin[i - 1] == '1') ^

   (bin[i] == '1');
   return gray;
}
```

#### 4.2Get permutation cycles

Description: Receives a permutation [0, n-1] and return a vector 2D with each cycle.

```
vll2d getPermutationCicles(const vll &ps) {
  ll n = len(ps);
  vector<char> visited(n);
  vector<vll> cicles;
  rep(i, 0, n)
    if (visited[i]) continue;
    vll cicle;
    ll pos = i;
    while (!visited[pos]) {
      cicle.pb(pos);
      visited[pos] = true;
      pos = ps[pos];
    cicles.push_back(vll(all(cicle)));
  return cicles;
}
```

# Max & Min Check

**Description**: Returns the min/max value in range [l, r] that satisfies the lambda function check, if there is no such value the 'nullopt' is returned. Usage: check must be a function that receives an integer and return a boolean. Time:  $O(\log r - l + 1)$ 

```
template <typename T>
optional<T> maxCheck(T l, T r, auto check) {
  optional<T> ret;
  while (l \ll r)
    T m = midpoint(l, r);
    if (check(m))
      ret ? chmax(ret, m) : ret = m, l = m + 1;
    else
      r = m - 1:
  return ret;
}
template <typename T>
optional<T> minCheck(T l, T r, auto check) {
  optional<T> ret;
  while (l \ll r)
    T m = midpoint(l, r);
    if (check(m))
      ret ? chmin(ret, m) : ret = m, r = m - 1;
    else
      l = m + 1;
  return ret;
}
```

### 4.4 Mo's algorithm

```
template <typename T, typename Tans>
struct Mo
  struct Query {
    int l, r, idx, block;
    Query(int_l, int _r, int _idx, int _block)
          l(_l),
           r(_r)
           idx(_idx)
           block(_block) {}
    bool operator<(const Query &q) const {</pre>
      if (block != q.block)
       return block < q.block;
return (block & 1 ? (r < q.r) : (r > q.r));
  };
  vector<T> vs;
  vector<Query> qs;
  const int block_size;
  Mo(const vector<T> &a)
       : vs(a),
         block_size((int)ceil(sqrt(a.size()))) {}
  void add_query(int l, int r) {
    auto solve() {
    // get answer return type
    vector<Tans> answers(qs.size());
    sort(all(qs));
    int cur_l = 0, cur_r = -1;
for (auto q : qs) {
  while (cur_l > q.l) add(—cur_l);
      while (cur_r < q.r) add(++cur_r);
      while (cur_l < q.l) remove(cur_l++);
while (cur_r > q.r) remove(cur_r—);
      answers[q.idx] = get_answer();
    return answers;
  }
 private:
  // add value at idx from data structure
inline void add(int idx) {}
     remove value at idx from data structure
  inline void remove(int idx) {}
     extract current answer of the data structure
  inline Tans get_answer() {}
```

# 4.5 \_\_int128t stream

```
void print(__int128 x) {
  if (x < 0) {
    cout << '-';
    x = -x;
  }
  if (x > 9) print(x / 10);
  cout << (char)((x % 10) + '0');
}
__int128 read() {
  string s;
  cin >> s;
  __int128 x = 0;
  for (auto c : s) {
    if (c != '-') x += c - '0';
    x *= 10;
  }
  x /= 10;
  if (s[0] == '-') x = -x;
  return x;
}
```

### 5 Geometry

### 5.1 Check if a point belong to line segment

```
// Verifica se o ponto P pertence ao segmento de
// reta AB
const ld EPS = 1e-9;
template <typename T>
struct Point {
  Point(T _{-}x, T _{-}y) : x(_{-}x), y(_{-}y) {}
template <typename T>
bool equals(const T a, const T b) {
  if (is_floating_point<T>) {
    return fabsl(a — b) <= EPS;
  return a == b;
}
/*
   Verify if the segment AB contains point P
 */
template <typename T>
bool contains(const Point<T> &A, const Point<T> &B,
                const Point<T> &P) {
  auto xmin = min(A.x, B.x);
  auto xmax = max(A.x, B.x);
  auto ymin = min(A.y, B.y);
  auto ymax = max(A.y, B.y);
  if (P.x < xmin \mid\mid P.x > xmax \mid\mid P.y < ymin \mid\mid
       P.y > ymax)
     return false;
  return equals((P.y - A.y) * (B.x - A.x)
                  (P.x - A.x) * (B.y - A.y));
}
5.2 Check if point is inside triangle
struct point {
  int x, y;
int id;
  point operator—(const point &o) const {
    return \{x - o.x, y - o.y\};
  int operator^(const point &o) const {
    return x * o.y - y * o.x;
};
    Verify the direction that the point
```

# 5.3 Convex hull

```
struct pt {
  double x, y;
  int id;
};
int orientation(pt a, pt b, pt c) {
  double v = a.x * (b.y - c.y) +
```

```
b.x * (c.y - a.y) +
c.x * (a.y - b.y);
return -1; // clockwise
return +1; // counter-clockwise
  if (v < 0) return -1; if (v > 0) return +1;
  return 0;
}
bool cw(pt a, pt b, pt c,
         bool include_collinear) {
  int o = orientation(a, b, c)
  return o < 0 || (include_collinear && o == 0);</pre>
bool collinear(pt a, pt b, pt c)
  return orientation(a, b, c) == 0;
void convex_hull(vector<pt> &pts,
                   bool include_collinear = false) {
  pt p0 = *min_element(all(pts), [](pt a, pt b) {
    return make_pair(a.y, a.x) <
            make_pair(b.y, b.x);
  sort(all(pts), [&p0](const pt &a, const pt &b) {
    int o = orientation(p0, a, b);
    if (o == 0)
      return (p0.x - a.x) * (p0.x - a.x) +
               (p0.y - a.y) * (p0.y - a.y) < (p0.x - b.x) * (p0.x - b.x) +
                   (p0.y - b.y) * (p0.y - b.y);
    return o < 0;
  if (include_collinear) {
    int i = len(pts) - 1;
    while (i >= 0 \&\&
            collinear(p0, pts[i], pts.back()))
    reverse(pts.begin() + i + 1, pts.end());
  vector<pt> st;
  for (int i = 0; i < len(pts); i++) {
    while (st.size() > 1 \&\&
            !cw(st[len(st) - 2], st.back(), pts[i],
                 include_collinear))
       st.pop_back();
    st.push_back(pts[i]);
  }
  pts = st;
```

### 5.4 Polygon lattice points

```
ll cross(ll x1, ll y1, ll x2, ll y2) \{
   return x1 * y2 - x2 * y1;
ll polygonArea(vector<pll> &pts) {
   ll ats = 0;
  for (int i = 2; i < len(pts); i++)
     ats +=
           \begin{array}{c} {\sf cross(pts[i].first-pts[0].first,} \\ {\sf pts[i].second-pts[0].second,} \end{array}
                   \begin{array}{l} \mathsf{pts}[\mathsf{i}-\mathsf{1}].\mathsf{first}-\mathsf{pts}[\mathsf{0}].\mathsf{first},\\ \mathsf{pts}[\mathsf{i}-\mathsf{1}].\mathsf{second}-\mathsf{pts}[\mathsf{0}].\mathsf{second}); \end{array}
  return abs(ats / 2ll);
}
ll boundary(vector<pll> &pts) {
   ll ats = pts.size();
  for (int i = 0; i < len(pts); i++) {
     ll deltax = (pts[i].first
                       pts[(i + 1) % pts.size()].first);
     ll deltay =
           (pts[i].second -
            pts[(i + 1) % pts.size()].second);
     ats += abs(\_gcd(deltax, deltay)) - 1;
   return ats;
}
pll latticePoints(vector<pll> &pts) {
   ll bounds = boundary(pts);
   ll area = polygonArea(pts);
  ll inside = area + 1ll - bounds / 2ll;
   return {inside, bounds};
```

# 5.5 Segment intersection

```
using ld = long double;
template <typename T = ld >
struct Point {
  T x, y;
bool is_port;
template <typename T = ld >
bool operator==(const Point<T> &a,
                   const Point<T> &b) {
  return a.x == b.x and a.y == b.y;
template <typename T = ld >
struct Segment {
  Point<T> p1, p2;
template <typename T>
int orientation(Point<T> p, Point<T> q,
                   Point<T> r) {
  int val = (q.y - p.y) * (r.x - q.x) - (q.x - p.x) * (r.y - q.y);
  // TODO: if it's a float must use other way to
  // compare
  if (val == 0)
  return 0; // colinear else if (val > 0)
     return 1; // clockwise
  else
     return 2; // counterclockwise
}
template <typename T>
bool do_segment_intersect(Segment<T> s1,
                                Segment<T> s2)
  int o1 = orientation(s1.p1, s1.p2, s2.p1);
int o2 = orientation(s1.p1, s1.p2, s2.p2);
  int o3 = orientation(s2.p1, s2.p2, s1.p1);
int o4 = orientation(s2.p1, s2.p2, s1.p2);
  return (o1 != o2 and o3 != o4) or
           (o1 == 0 and o3 == 0) or
           (02 == 0 \text{ and } 04 == 0);
}
```

## 6 Graphs

### 6.1 Heavy-Light Decomposition (point update)

### 6.1.1 Maximum number on path

```
struct Node {
  ll value;
  Node()
      : value(numeric_limits<
               ll>::min()){}; // Neutral
                                 // element
  Node(ll v) : value(v){};
}:
Node combine(Node l, Node r) {
  Node m;
  m.value = max(l.value, r.value);
  return m;
template <typename T = Node, auto F = combine>
struct SegTree {
  int n;
  vector<T> st;
  SegTree(int _n) : n(_n), st(n \ll 1) {}
  void set(int p, const T &k) {
  for (st[p += n] = k; p >>= 1;)
      st[p] = F(st[p << 1], st[p << 1 | 1]);
  T query(int l, int r) {
    if (l & 1) ansl = F(ansl, st[l++]);
if (r & 1) ansr = F(st[—r], ansr);
```

```
return F(ansl, ansr);
};
template <typename SegT = Node,</pre>
            auto SegOp = combine>
struct HeavyLightDecomposition {
  int n;
  vi ps, ds, sz, heavy, head, pos;
SegTree<SegT, SegOp> seg;
  HeavyLightDecomposition(const vi2d &g,
                                const vector<SegT> &v,
                                int root = 0
       : n(len(g)), seg(n) {
    ps = ds = sz = heavy = head = pos = vi(n, -1);
     auto dfs = [&](auto &&self, int u) -> void {
       sz[u] = 1;
       int mx = 0;
for (auto x : g[u])
         if (x != ps[\bar{u}]) {
            ps[x] = u;
            ds[x] = ds[u] + 1;
self(self, x);
            sz[u] += sz[x];
            if(sz[x] > mx)
              mx = sz[x], heavy[u] = x;
    }:
    dfs(dfs, root);
    for (int i = 0, cur = 0; i < n; i++)
       if (ps[i] == -1 \text{ or heavy}[ps[i]] != i)
         for (int j = i; j != -1; j = heavy[j]) {
            head[j] = i;
            pos[j] = cur++;
    rep(i, 0, n) seg.set(pos[i], v[i]);
  vector<pii> disjoint_ranges(int u, int v) {
    vector<pii> ret;
    for (; head[u] != head[v]; v = ps[head[v]]) {
       if (ds[head[u]] > ds[head[v]]) swap(u, v);
       ret.eb(pos[head[v]], pos[v]);
    if (ds[u] > ds[v]) swap(u, v);
    ret.eb(pos[u], pos[v]);
    return ret;
  SegT query_path(int u, int v) {
    SegT res;
    for (auto [l, r] : disjoint_ranges(u, v)) {
       res = SegOp(res, seg.query(l, r));
    return res;
  SegT query_subtree(int u) const {
    return seg.query(pos[u], pos[u] + sz[u] - 1);
  void set(int u, SegT x) { seg.set(pos[u], x); }
6.2 2-SAT
Description: Calculates a valid assignment to boolean variables a, b, c,... to a
2-SAT problem, so that an expression of the type (a||b)\&\&(!a||c)\&\&(d||!b)\&\&... becomes true, or reports that it is unsatisfiable. Usage: Negated variables are represented by bit-inversions (x).
Time: O(N+E), where N is the number of boolean variables, and E is the
```

Returns true iff it is solvable ts.values[0..N-1] holds the assigned values to the

number of clauses.

```
struct TwoSat {
  int N;
  vector<vi> gr;
vi values; // 0 = false, 1 = true
  TwoSat(int n = 0) : N(n), gr(2 * n) {}
  int addVar() { // (optional)
    gr.eb();
    gr.eb();
    return N++;
```

```
void either(int f, int j) {
  f = max(2 * f, -1 - 2 * f);
  j = max(2 * j, -1 - 2 * j);
  gr[f].pb(j ^ 1);
  gr[j].pb(f ^ 1);
   void setValue(int x) { either(x, x); }
   void atMostOne(const vi &li) { // (optional)
     if (sz(li) <= 1) return;
int cur = ~li[0];
     rep(i, 2, sz(li))
        int next = addVar()
        either(cur, ~li[i]);
        either(cur, next);
either(~li[i], next);
        cur = ~next;
     either(cur, ~li[1]);
  }
  vi val, comp, z;
   int time = 0;
   int dfs(int i)
     int low = val[i] = ++time, x;
     z.pb(i);
     for (int e : gr[i])
        if (!comp[e])
     low = min(low, val[e] ?: dfs(e));
if (low == val[i]) do {
           x = z.back()
           z.pop_back();
           comp[x] = low;
           if (values[x >> 1] == -1)
        values[x >> 1] = x & 1;
} while (x != i);
     return val[i] = low;
  \begin{array}{c} \texttt{bool solve() \{} \\ \texttt{values.assign(N, } -1); \end{array}
     val.assign(2 * N, 0);
     comp = val;
rep(i, 0, 2 * N) if (!comp[i]) dfs(i);
     rep(i, 0, N) if (comp[2 * i] ==
                             comp[2 * i + 1]) return 0;
     return 1;
};
```

#### 6.3BFS-01

Description: Similar to a Dijkstra given a weighted graph finds the distance from source s to every other node.

Time: O(V + E)**Warning**: Applicable only when the weight of the edges  $\{0, x\}$ 

```
vector<pair<ll, int>> adj[maxn];
ll dists[maxn];
int s, n;
void bfs_01()
  fill(dists,
dist[s] = 0;
                 dists + n, oo);
  deque<int> q;
  q.emplace_back(s);
  while (not q.empty()) {
     auto u = q.front();
     q.pop_front();
     for (auto [v, w] : adj[u]) {
  if (dist[v] <= dist[u] + w) continue;
  dist[v] = dist[u] + w;</pre>
       w ? q.emplace_back(v) : q.emplace_front(v);
}
```

### 6.4 Bellman ford

Description: Find shortest path from a single source to all other nodes. Can detect negative cycles. **Time**:  $O(V \cdot E)$ 

```
bool bellman_ford(
    const vector<vector<pair<int, ll>>> &g, int s,
    vector<ll> &dist) {
```

```
int n = (int)g.size();
dist.assign(n, LLONG_MAX);
vector<int> count(n);
vector<char> in_queue(n);
queue<int> q;
dist[s] = 0;
q.push(s);
in_queue[s] = true;
while (not q.empty()) {
  int cur = q.front();
  q.pop();
  in_queue[cur] = false;
  for (auto [to, w] : g[cur]) {
   if (dist[cur] + w < dist[to]) {</pre>
       dist[to] = dist[cur] + w;
if (not in_queue[to]) {
         q.push(to);
         in_queue[to] = true;
         count[to]++
         if (count[to] > n) return false;
  }
}
return true;
```

### 6.5 Bellman-Ford (find negative cycle)

**Description**: Given a directed graph find a negative cycle by running n iterations, and if the last one produces a relaxation than there is a cycle. **Time**:  $O(V \cdot E)$ 

```
const ll oo = 2500 * 1e9;
using graph = vector<vector<pair<int, ll>>>;
vi negative_cycle(graph &g, int n) {
  vll d(n, oo);
vi p(n, -1);
  int x = -1;
  d[0] = 0;
for (int i = 0; i < n; i++) {
     x = -1
     for (int u = 0; u < n; u++)
       for (auto &[v, l] : g[u]) {
   if (d[u] + l < d[v]) {
      d[v] = d[u] + l;
      p[v] = u;
            \dot{x} = v;
         }
       }
    }
  }
  if (x == -1)
     return {};
  else {
     for (int i = 0; i < n; i++) x = p[x];
     vi cycle;
     for (int v = x; v = p[v]) {
       cycle.eb(v);
       if (v == x and len(cycle) > 1) break;
     reverse(all(cycle));
     return cycle;
  }
}
```

### 6.6 Biconnected Components

**Description:** Build a vector of vectors, where the i-th vector correspond to the nodes of the i-th, biconnected component, a biconnected component is a subset of nodes and edges in which there is no cut point, also exist at least two distinct routes in vertex between any two vertex in the same biconnected component. **Time:** O(N+M)

```
for (auto v : g[u]) {
   if (!tin[v]) {
      son_cnt++;
       int lowx = tarjan(v, u);
       lowu = min(lowu, lowx);
       if (lowx >= tin[u]) {
         while (top != -1 \&\& stck[top + 1] != v)
           nodes[bcc_cnt].emplace_back(
                stck[top—]);
         nodes[bcc_cnt++].emplace_back(u);
    } else {
       lowu = min(lowu, tin[v]);
    }
  if (p == -1 \&\& son_cnt == 0) {
    nodes[bcc_cnt++].emplace_back(u);
  return lowu;
}
void build_bccs() {
  timer = 1;
  top = -1;
  memset(tin, 0, sizeof(int) * n);
  for (int i = 0; i < n; i++) nodes[i] = {};
  bcc_cnt = 0;
  for (int u = 0; u < n; u++)
  if (!tin[u]) tarjan(u);</pre>
}
```

### 6.7 Binary Lifting/Jumping

**Description:** Given a function/successor grpah answers queries of the form which is the node after k moves starting from u. **Time:** Build  $O(N \cdot MAXLOG2)$ , Query O(MAXLOG2).

```
const int MAXN(2e5), MAXLOG2(30);
int bl[MAXN][MAXLOG2 + 1];
int N;
int jump(int u, ll k) {
   for (int i = 0; i <= MAXLOG2; i++) {
      if (k & (1ll << i)) u = bl[u][i];
   }
   return u;
}
void build() {
   for (int i = 1; i <= MAXLOG2; i++) {
      for (int j = 0; j < N; j++) {
        bl[j][i] = bl[bl[j][i - 1]][i - 1];
   }
}</pre>
```

# 6.8 Bipartite Graph

}

}

**Description**: Given a graph, find the 'left' and 'right' side if is a bipartite graph, if is not then a empty vi2d is returned **Time**: O(N+M)

```
vi2d bipartite_graph(vi2d &adj) {
  int n = len(adj);
   vi side(n,
                 -1);
  vi2d ret(2);
rep(u, 0, n) {
  if (side[u] == -1) {
        queue<int> q;
        q.emp(u);
        side[u] = 0
        ret[0].eb(u);
while (len(q)) {
           int u = q.front();
           q.pop();
           for (auto v : adj[u]) {
  if (side[v] == -1) {
    side[v] = side[u] ^ 1;
                ret[side[v]].eb(v);
                q.push(v);
              } else if (side[u] == side[v])
                return {};
           }
       }
     }
```

```
return ret;
```

### Block-Cut tree

```
struct block_cut_tree {
  int n;
  vector<int> id, is_cutpoint, tin, low, stk;
  vector<vector<int>> comps, tree;
  block_cut_tree(vector<vector<int>> &g)
       : n(g.size()),
         id(n),
         is_cutpoint(n),
         tin(n),
low(n) {
    // build comps
    for (int i = 0; i < n; i++) {
   if (!tin[i]) {
         int timer = 0;
         dfs(i, -1, timer, g);
    }
    int node_id = 0;
    for (int u = 0; u < n; u++) {
   if (is_cutnoin+[u]) {
         (is_cutpoint[u]) {
         id[u] = node_id++
         tree.push_back({});
      }
    for (auto &comp : comps) {
      int node = node_id++;
       tree.push_back({});
      for (int u : comp) {
         if (!is_cutpoint[u]) {
           id[u] = node;
         } else {
           tree[node].emplace_back(id[u]);
           tree[id[u]].emplace_back(node);
      }
    }
  void dfs(int u, int p, int &timer
            vector<vector<int>> &g) {
    tin[u] = low[u] = ++timer;
    stk.emplace_back(u)
    for (auto v : g[u]) {
  if (v == p) continue;
      if (!tin[v]) {
         dfs(v, u, timer, g);
low[u] = min(low[u], lo
if (low[v] >= tin[u]) {
                                 low[v]);
           is_cutpoint[u]
                (tin[u] > 1 \text{ or } tin[v] > 2);
           comps.push_back({u});
           while (comps.back().back() != v) {
             comps.back().emplace_back(stk.back());
             stk.pop_back();
        else
         low[u] = min(low[u], tin[v]);
    }
```

#### 6.10D'Escopo-Pape

**Description**: Is a single source shortest path that works faster than Dijkstra's algorithm and the Bellman-Ford algorithm in most cases, and will also work for negative edges. However not for negative cycles. There exists cases where it runs in exponential time.

Usage: Returns a pair containing two vectors, the first one with the distance from s to every other node, and another one with the ancestor of each node, note that the ancestor of s is -1

```
using Edge = pair<ll, int>;
using Adj = vector<vector<Edge>>;
pair<vll, vi> desopo_pape(int s, int n,
                           const Adj &adj) {
  vll ds(n, LLONG_MAX), ps(n, -1);
  ds[s] = 0;
```

```
vi ms(n, 2);
deque<int> q;
q.eb(s);
while (len(q)) {
   int u = q.front();
   q.pop_front();
   ms[u] = 0;
      or (auto [w, v] : adj[u]) {
  if (chmin(ds[v], w + ds[u])) {
         ps[v] = u;
        if (ms[v] == 2)
  ms[v] = 1, q.pb(v);
else if (ms[v] == 0)
            ms[v] = 1, q.pf(v);
  }
return {ds, ps};
```

### 6.11 Dijkstra

```
const int MAXN = 1'00'000;
const ll MAXW = 1'000'000ll;
constexpr ll 00 = MAXW * MAXN + 1;
using Edge = pair<ll, int>; // { weigth, node}
using Adj = vector<vector<Edge>>;
template <typename T>
using min_heap =
    priority_queue<T, vector<T>, greater<T>>;
pair<vll, vi> dijkstra(const Adj \delta g, int s) { int n = len(g);
  min_heap<Edge> pq;
  vļl dṣ(n, 00);
  vi ps(n, -1)
  pq.emp(0, s);
  ds[s] = 0;
while (len(pq)) {
    auto [du, u] = pq.top();
    pq.pop();
     if (ds[u] < du) continue;
    for (auto [w, v] : g[u]) {
    ll ndv = du + w;
       if (chmin(ds[v], ndv)) {
         ps[v] = u;
         pq.emp(ndv, v);
    }
     return {ds, ps};
  // optional !
  vi recover_path(int source, int ending,
                     const vi &ps) {
     if (ps[ending] == -1) return \{\};
     int cur = ending;
    vi ans;
    while (cur !=-1) {
       ans.eb(cur);
       cur = ps[cur];
     reverse(all(ans));
     return ans;
```

# 6.12 Dijkstra (K-shortest pahts)

```
const ll oo = 1e9 * 1e5 + 1;
using adj = vector<vector<pll>>>;
vector<priority_queue<ll>> dijkstra(
     const vector<vector<pll>>> &g, int n, int s,
     int k)
  priority_queue<pll, vector<pll>, greater<pll>>
  vector<priority_queue<ll>> dist(n);
dist[0].emplace(0);
  pq.emplace(0, s);
  while (!pq.empty()) {
    auto [d1, v] = pq.top();
    pq.pop();
```

```
if (not dist[v].empty() and
       dist[v].top() < d1)
    continue:
  for (auto [d2, u] : g[v]) {
    if (len(dist[u]) < k) {
   pq.emplace(d2 + d1, u);</pre>
       dist[u].emplace(d2 + d1);
    } else {
       if (dist[u].top() > d1 + d2) {
         dist[u].pop();
         dist[u].emplace(d1 + d2);
         pq.emplace(d2 + d1, u);
    }
 }
}
return dist;
```

#### Extra Edges to Make Digraph Fully Strongly 6.13Connected

**Description**: Given a directed graph G find the necessary edges to add to make the graph a single strongly connected component. Time: O(N+M)

```
\mathbf{Memory} \colon O(N)
```

```
struct SCC {
  int n, num_sccs;
  vi2d adj;
  vi scc_id;
  SCC(int _n)
      : n(_n),
        num_sccs(0),
        adj(n)
         scc_id(n, -1) {}
  SCC(const vi2d &_adj) : SCC(len(_adj)) {
    adj = \_adj;
    find_sccs();
  void add_edge(int u, int v) { adj[u].eb(v); }
  void find_sccs() {
    int timer = 1
    vi tin(n), st;
    st.reserve(n)
    function<int(int)> dfs = [&](int u) -> int {
      int low = tin[u] = timer++, siz = len(st);
      st.eb(u);
      for (int'v : adj[u])
        if (scc_id[v] < 0)
           low =
               min(low, tin[v] ? tin[v] : dfs(v));
      st.resize(siz);
num_sccs++;
      return low;
    for (int i = 0; i < n; i++)
  if (!tin[i]) dfs(i);</pre>
 }
};
vector<array<int, 2>> extra_edges(
    const vi2d &adj) {
  SCC scc(adj);
  auto scc_id = scc.scc_id;
  auto num_sccs = scc.num_sccs;
  if (num_sccs == 1) return {};
  int n = len(adj);
  vi2d scc_adj(num_sccs);
  vi zero_in(num_sccs, 1);
  rep(u, 0, n) {
  for (int v : adj[u]) {
      if (scc_id[u] == scc_id[v]) continue;
scc_adj[scc_id[u]].eb(scc_id[v]);
      zero_in[scc_id[v]] = 0;
  }
  int random source :
      max_element(all(zero_in)) - zero_in.begin();
```

```
function<int(int)> dfs = [&](int u) {
     if (empty(scc_adj[u])) return u;
    for (int v : scc_adj[u])
  if (!vis[v]) {
         vis[v] = 1;
         int zero_out = dfs(v);
if (zero_out != -1) return zero_out;
     return (int)—1;
  };
  vector<array<int, 2>> edges;
  vi in_unused;
  rep(i, 0, num_sccs) {
  if (zero_in[i]) {
       vis[i] = 1
       int zero_out = dfs(i);
       if (zero_out != -1)
         edges.push_back({zero_out, i});
       else
         in_unused.push_back(i);
    }
  rep(i, 1, len(edges)) {
     swap(edges[i][0], edges[i - 1][0]);
  rep(i, 0, num_sccs) {
    if (scc_adj[i].empty() && !vis[i]) {
       if (!in_unused.empty()) {
         edges.push_back({i, in_unused.back()});
         in_unused.pop_back();
       } else {
         edges.push_back({i, random_source});
    }
  }
  for (int u : in_unused) edges.push_back({0, u});
  vi to_node(num_sccs)
  rep(i, 0, n) to_node[scc_id[i]] = i;
  for (auto &[u, v] : edges)
    u = to_node[u], v = to_node[v];
  return edges;
}
6.14 Find Articulation/Cut Points
Description: Given an undirected graph find it's articulation points.
Time: O(N+M)
Warning: A vertex u can be an articulation point if and only if has at least 2 adjascent vertex
const int MAXN(100);
int N;
vi2d G;
int timer
int tin[MAXN], low[MAXN];
set<int> cpoints;
int dfs(int u, int p = -1) {
  int cnt = 0;
  low[u] = tin[u] = timer++;
  for (auto v : G[u]) {
    if (not tin[v]) {
       dfs(v, u);
       if (low[v] >= tin[u]) cpoints.insert(u);
      low[u] = min(low[u], low[v]);
else if (v != p)
       low[u] = min(low[u], tin[v]);
  return cnt;
void getCutPoints() {
  memset(low, 0, sizeof(low));
memset(tin, 0, sizeof(tin));
cpoints.clear();
  timer = 1;
  for (int i = 0; i < N; i++) {
  if (tin[i]) continue:</pre>
    if (tin[i]) continue;
int cnt = dfs(i);
    if (cnt == 1) cpoints.erase(i);
```

vi vis(num\_sccs)

}

### 6.15 Find Bridge-Tree components

**Usage:** label2CC(u, p) finds the 2-edge connected component of every node. **Time:** O(n+m)

```
const int maxn(3 '00' 000);
int tin[maxn], compId[maxn], qtdComps;
vi g[maxn], stck;
int n
int dfs(int u, int p = -1) {
  int low = tin[u] = len(stck);
  stck.emplace_back(u);
  bool multEdge = false;
  for (auto v : g[u]) {
     if (v == p and !multEdge) {
       multEdge = 1;
        continue:
     low = min(low)
                  tin[v] == -1 ? dfs(v, u) : tin[v]);
  if (low == tin[u]) {
  for (int i = tin[u]; i < len(stck); i++)
    compId[stck[i]] = qtdComps;</pre>
     stck.resize(tin[u]);
     qtdComps++;
  return low;
}
void label2CC() {
  memset(compId, -1, sizeof(int) * n);
memset(tin, -1, sizeof(int) * n);
  stck.reserve(n);
for (int i = 0; i < n; i++) {
  if (tin[i] == -1) dfs(i);</pre>
```

### 6.16 Find Bridges

**Description**: Find every bridge in a **undirected** connected graph. **Warning**: Remember to read the graph as pair where the second is the id of the edge!

### 6.17 Find Centroid

**Description**: Given a tree (don't forget to make it 'undirected'), find it's centroid O(V)

```
bool iscentroid = true;
for (auto v : g[u])
   if (v != p) {
      dfs(v, u, n, g, sz, centroid);
      if (sz[v] > n / 2) iscentroid = false;
      sz[u] += sz[v];
   }
   if (n - sz[u] > n / 2) iscentroid = false;
   if (iscentroid) centroid.eb(u);
}
vi getCentroid(vi2d &g, int n) {
   vi centroid;
   vi sz(n);
   dfs(0, -1, n, g, sz, centroid);
   return centroid;
}
```

# 6.18 Find bridges (online)

```
// O((n+m)*log(n))
struct BridgeFinder {
  // 2ecc = 2 edge conected component
  // cc = conected component
  vector<int> parent, dsu_2ecc, dsu_cc,
  dsu_cc_size;
int bridges, lca_iteration;
vector<int> last_visit;
  BridgeFinder(int n)
       : parent(n, -1),
          dsu_2ecc(n),
          dsu_cc(n),
dsu_cc_size(n, 1),
          bridges(0),
          lca_iteration(0)
    last_visit(n) {
for (int i = 0; i < n; i++) {
  dsu_2ecc[i] = i;
  dsu_cc[i] = i;
  int find_2ecc(int v) {
     if (v == -1) return -1;
     return dsu_2ecc[v] == v
? v
                   : dsu_2ecc[v] =
                           find_2ecc(dsu_2ecc[v]);
  int find_cc(int v) {
    v = find_2ecc(v);
return dsu_cc[v] == v
                   : dsu_cc[v] = find_cc(dsu_cc[v]);
  void make_root(int v) {
     v = find_2ecc(v);
int root = v;
    int child = -1;
while (v != -1) {
  int p = find_2ecc(parent[v]);
       parent[v] = child;
dsu_cc[v] = root;
       child = v;
     dsu_cc_size[root] = dsu_cc_size[child];
  void merge_path(int a, int b) {
     ++lca_iteration;
     vector<int> path_a, path_b;
     int lca = -1;
    while (lca == -1) { if (a != -1) {
          a = find_2 ecc(a)
          path_a.push_back(a);
          if (last_visit[a] == lca_iteration) {
            lca = a;
            break;
          last_visit[a] = lca_iteration;
          a = parent[a];
       if (b != -1) {
```

```
b = find_2ecc(b);
path_b.push_back(b);
        if (last_visit[b] == lca_iteration) {
          l\dot{c}a = b;
          break:
        last_visit[b] = lca_iteration;
        b = parent[b];
  }
  for (auto v : path_a) {
  dsu_2ecc[v] = lca;
  if (v == lca) break;
      —bridges;
  for (auto v : path_b) {
  dsu_2ecc[v] = lca;
  if (v == lca) break;
       ⊸bridges;
}
void add_edge(int a, int b) {
  a = find_2ecc(a);
  b = find_2ecc(b);
  if (a == b) return;
  int ca = find_cc(a);
  int cb = find_cc(b);
  if (ca != cb) {
     ++bridges;
     if (dsu_cc_size[ca] > dsu_cc_size[cb]) {
        swap(a, b);
        swap(ca, cb);
     make_root(a);
parent[a] = dsu_cc[a] = b;
     dsu_cc_size[cb] += dsu_cc_size[a];
  } else {
     merge_path(a, b);
  }
```

### 6.19 Floyd Warshall

**Description**: Simply finds the minimal distance for each node to every other node.  $O(V^3)$ 

# 6.20 Functional/Successor Graph

**Description**: Given a functional graph find the vertice after k moves starting at u and also the distance between u and v, if it's impossible to reach v starting at u returns -1.

**Time**: build  $O(N \cdot MAXLOG2)$ , kth O(MAXLOG2), dist O(MAXLOG2)

```
const int MAXN(2 '000' 000), MAXLOG2(24);
int N;
vi2d succ(MAXN, vi(MAXLOG2 + 1));
vi dst(MAXN, 0);
int vis[MAXN];
void dfsbuild(int u) {
  if (vis[u]) return;
  vis[u] = 1;
  int v = succ[u][0];
  dfsbuild(v);
  dst[u] = dst[v] + 1;
}
void build() {
  for (int i = 0; i < N; i++) {</pre>
```

```
if (not vis[i]) dfsbuild(i);
   for (int k = 1; k \le MAXLOG2; k++) {
     for (int i = 0; i < N; i++)
       succ[i][k] = succ[succ[i][k-1]][k-1];
}
int kth(int u, ll k) {
  if (k <= 0) return u;
  for (int i = 0; i <= MAXLOG2; i++)
    if ((1ll << i) & k) u = succ[u][i];</pre>
   return`u;
}
int dist(int u, int v) {
  int cu = kth(u, dst[u]);
  if (kth(u, dst[u] - dst[v]) == v)
    return dst[u] - dst[v];
  else if (kth(cu, dst[cu] - dst[v]) ==
    return dst[u] + (dst[cu] - dst[v]);
}
                                    dst[v]) == v)
     return -1;
6.21 Heavy light decomposition (supreme)
struct HLD {
   int V;
   int id;
   int nb_heavy_path;
   std::vector<std::vector<int>> g;
   std::vector<pair<int, int>>
                               // edges of the tree
// par[i] = parent of
// vertex i (Default: -1)
        edges;
   std::vector<int> par;
   std::vector<int>
        depth; // depth[i] = distance between root
                  // and vertex i
   std::vector<int>
        subtree_sz;
                         // subtree_sz[i] = size of
                        // subtree whose root is i
   std::vector<int>
        heavy_child;
                          // heavy_child[i] = child of
                          // vertex i on heavy path
                          // (Default: -1)
   std::vector<int>
        tree_id;
                    // tree_id[i] = id of tree vertex
                     // i belongs to
   std::vector<int> aligned_id,
        aligned_id_inv;
                             // aligned_id[i] = aligned
                             // id for vertex i
// (consecutive on heavy
                              // edges)
   std::vector<int>
        head; // head[i] = id of vertex on heavy
                // path of vertex i, nearest to root
   std::vector<int>
        head_ids;
                     // consist of head vertex id's
   std::vector<int>
                            // heavy_path_id[i] =
        heavy_path_id;
                            // heavy_path_id for vertex
                            // [i]
  HLD(const std::vector<std::vector<int>> &e,
    vector<int> roots = {0})
        : HLD((int)e.size()) {
     build(roots);
   HLD(int sz = 0)
        : V(sz),
          id(0)
          nb_heavy_path(0),
          g(sz),
          par(sz)
          depth(sz),
          subtree_sz(sz)
          heavy_child(sz),
          tree_id(sz,
          aligned_id(sz),
          aligned_id_inv(sz),
          head(sz)
          heavy_path_id(sz, -1) {}
```

void add\_edge(int u, int v) {

```
edges.emplace_back(u, v);
  g[u].emplace_back(v);
  g[v].emplace_back(u);
void _build_dfs(int root) {
  std::stack<std::pair<int, int>> st;
par[root] = -1;
  depth[root] = 0;
  st.emplace(root, 0);
  while (!st.empty()) {
    int now = st.top().first;
    int &i = st.top().second;
    if (i < (int)g[now].size()) {
  int nxt = g[now][i++];</pre>
       if (nxt == par[now]) continue;
       par[nxt] = now;
       depth[nxt] = depth[now] + 1;
       st.emplace(nxt, 0);
    } else {
       st.pop();
       int max_sub_sz = 0;
subtree_sz[now] = 1;
heavy_child[now] = -1;
       for (auto nxt : g[now]) {
         if (nxt == par[now]) continue;
         subtree_sz[now] += subtree_sz[nxt];
if (max_sub_sz < subtree_sz[nxt])</pre>
           max_sub_sz = subtree_sz[nxt],
           heavy_child[now] = nxt;
    }
  }
void _build_bfs(int root, int tree_id_now) {
  std::queue<int> q({root});
  while (!q.empty())
    int h = q.front();
    q.pop();
    head_ids.emplace_back(h);
    for (int now = h; now !=-1;
          now = heavy_child[now]) {
       tree_id[now] = tree_id_now;
aligned_id[now] = id++;
       aligned_id_inv[aligned_id[now]] = now;
       heavy_path_id[now] = nb_heavy_path;
       head[now] = h;
       for (int nxt : g[now])
         if (nxt != par[now] and
    nxt != heavy_child[now])
            q.push(nxt);
    nb_heavy_path++;
}
void build(std::vector<int> roots = {0}) {
  int tree_id_now = 0;
  for (auto r : roots)
    _build_dfs(r), _build_bfs(r, tree_id_now++);
// data[i] = value of vertex i
template <class T>
std::vector<T> segtree_rearrange(
    const std::vector<T> &data) const {
  assert(int(data.size()) == V);
  std::vector<T> ret;
  ret.reserve(V);
  for (int i = 0; i < V; i++)
    ret.emplace_back(data[aligned_id_inv[i]]);
  return ret;
// data[i] = weight of edge[i]
template <class T>
std::vector<T> segtree_rearrange_weighted(
    const std::vector<T> &data) const {
  assert(data.size() == edges.size());
  vector<T> ret(V);
for (int i = 0; i < (int)edges.size(); i++) {</pre>
    auto [u, v] = edges[i];
    if (depth[u] > depth[v]) swap(u, v);
    ret[aligned_id[v]] = data[i];
  return ret;
```

```
}
int segtree_edge_index(int i) const {
 auto [u, v] = edges[i];
if (depth[u] > depth[v]) swap(u, v);
  return aligned_id[v];
// query for vertices on path [u, v] (INCLUSIVE)
static_assert(
      std::is_invocable_r_v<void, decltype(f),</pre>
                            int, int>);
 while (true) {
    if (aligned_id[u] > aligned_id[v])
  std::swap(u, v);
    f(std::max(aligned_id[head[v]],
               aligned_id[u]),
      aligned_id[v]);
    if (head[u] == head[v]) break;
    v = par[head[v]];
}
void for_each_vertex_noncommutative(
    int from, int to, const auto &fup,
    const auto &fdown) const {
  static_assert(
      static_assert(
      std::is_invocable_r_v<
  void, decltype(fdown), int, int>);
 int u = from, v = to;
 while (u \ge 0 \text{ and depth}[u] > dlca) {
    const int p =
        (depth[head[u]] > dlca ? head[u] : lca);
    fup(aligned_id[p] + (p == lca),
        aligned_id[u]),
        u = par[p];
  static std::vector<std::pair<int, int>> lrs;
  int sz = 0;
  while (v \ge 0 \text{ and depth}[v] \ge dlca) {
    const int p =
        (depth[head[v]] >= dlca ? head[v]
                                : lca);
    if (int(lrs.size()) == sz)
  lrs.emplace_back(0, 0);
    lrs.at(sz++) = \{p, v\}, v = par.at(p);
  while (sz-
    fdown(aligned_id[lrs.at(sz).first]
          aligned_id[lrs.at(sz).second]);
// query for edges on path [u, v]
void for_each_edge(int u, int v,
                   const auto &f) const {
  static_assert(
      std::is_invocable_r_v<void, decltype(f),</pre>
 while (true) {
  if (aligned_id[u] > aligned_id[v])
    std::swap(u, v);
if (head[u] != head[v]) {
      f(aligned_id[head[v]], aligned_id[v]);
      v = par[head[v]];
    } else {
  if (u != v)
    f(aligned_id[u] + 1, aligned_id[v]);
 }
// lowest_common_ancestor: 0(log V)
int lowest_common_ancestor(int u, int v) const {
```

```
while (true) {
      if (aligned_id[u] > aligned_id[v])
        std::swap(u, v);
      if (head[u] == head[v]) return u;
      v = par[head[v]];
   }
  }
 2 * depth[lowest_common_ancestor(u,
                                             v)1;
  // Level ancestor, O(log V)
  // if k-th parent is out of range, return -1
  int kth_parent(int v, int k) const {
    if (k < 0) return -1;
   while (v >= 0)
      int h = head.at(v)
          len = depth.at(v) - depth.at(h);
      if (k \le len)
        return aligned_id_inv.at(
            aligned_id.at(v) - k);
      k = len + 1, v = par.at(h);
    return -1;
  // Jump on tree, O(log V)
  if (k < 0) return -1;
if (k == 0) return s;
int lca = lowest_common_ancestor(s, t);</pre>
    if (k <= depth.at(s) - depth.at(lca))</pre>
      return kth_parent(s, k);
    return kth_parent(t,
                      depth.at(s) + depth.at(t) -
                           depth.at(lca) * 2 - k);
};
6.22 Kruskal
Description: Find the minimum spanning tree of a graph.
Time: O(E \log E)
struct UFDS {
  vector<int> ps, sz;
  int components;
  UFDS(int n)
    : ps(n + 1), sz(n + 1, 1), components(n) { iota(all(ps), 0);
  }
  int find_set(int x) {
    return (x == ps[x]
                : (ps[x] = find_set(ps[x])));
  bool same_set(int x, int y) {
  return find_set(x) == find_set(y);
  void union_set(int x, int y) {
   x = find_set(x);
    y = find_set(y);
    if (x == y) return;
    if (sz[x] < sz[y]) swap(x, y);
    ps[y] = x;
    sz[x] += sz[y];
    components—;
 }
};
vector<tuple<ll, int, int>> kruskal(
    int n, vector<tuple<ll, int, int>> &edges) {
```

UFDS ufds(n);

sort(all(edges));

vector<tuple<ll, int, int>> ans;

for (auto [a, b, c] : edges) {

```
if (ufds.same_set(b, c)) continue;
  ans.emplace_back(a, b, c);
  ufds.union_set(b, c);
}
  return ans;
}
```

### 6.23 Lowest Common Ancestor

 $\bf Description:$  Given two nodes of a tree find their lowest common ancestor, or their distance

```
template <typename T>
struct SparseTable {
   vector<T> v;
   int n;
   static const int b = 30;
   vi mask, t;
   int op(int x, int y) {
     return v[x] < v[y] ? x : y;
   int msb(int x)
     return __builtin_clz(1) — __builtin_clz(x);
   ŚparseTable() {}
   SparseTable(const vector<T> \&v_{-})
     : v(v_), n(v.size()), mask(n), t(n) {
for (int i = 0, at = 0; i < n;
    mask[i++] = at |= 1) {</pre>
        at = (at << 1) & ((1 << b) - 1);
while (at and op(i, i - msb(at & -at)) == i)
           at ^= at & —at;
     for (int i = 0; i < n / b; i++)

t[i] = b * i + b - 1 -

msb(mask[b * i + b - 1]);
     for (int j = 1; (1 << j) <= n / b; j++) for (int i = 0; i + (1 << j) <= <math>n / b; i++)
           t[n / b * j + i] =
                int small(int r, int sz = b) {
  return r - msb(mask[r] & ((1 << sz) - 1));</pre>
     query(int l, int r) if (r - l + 1 \le b)
     return small(r, r - l + 1);
int ans = op(small(l + b - 1), small(r));
int x = l / b + 1, y = r / b - 1;
     if (x \le y) \{
        int j = msb(y - x + 1); ans =
              op(ans,
                  op(t[n / b * j + x],
                      t[n / b * j + y - (1 << j) + 1]));
     return ans:
  }
};
struct LCA {
   SparseTable<int> st;
   int n;
  vi v, pos, dep;
   LCA(const vi2d &g, int root)
     : n(len(g)), pos(n) {
dfs(root, 0, -1, g);
     st = SparseTable<int>(vector<int>(all(dep)));
  void dfs(int i, int d, int p, const vi2d &g) {
     v.eb(len(dep)) = i, pos[i] = len(dep),
     dep.eb(d);
     for (auto j : g[i])
  if (j != p) {
    dfs(j, d + 1, i)
           dfs(j, d + 1, i, g);
v.eb(len(dep)) = i, dep.eb(d);
  int lca(int a, int b) {
  int l = min(pos[a], pos[b]);
     int r = max(pos[a], pos[b]);
```

```
return v[st.query(l, r)];
  int dist(int a, int b) {
  return dep[pos[a]] + dep[pos[b]] -
              2 * dep[pos[lca(a, b)]];
};
```

### 6.24 Lowest Common Ancestor (Binary Lifting)

Description: Given a directed tree, finds the LCA between two nodes using binary lifting, and answer a few queries with it. Usage:

• lca: returns the LCA between the two given nodes

ullet on path: fids if c is in the path from a to b

Time: build  $O(N \cdot MAXLOG2)$ , all queries O(MAXLOG2)

```
struct LCA {
  int n;
const int maxlog;
  vector<vector<int>> up;
  vector<int> depth;
  LCA(const vector<vector<int>> &tree)
       : n(tree.size()),
  maxlog(ceil(log2(n))),
         up(n, vector<int>(maxlog + 1)),
         depth(n, -1) {
    for (int i = 0; i < n; i++) {
       if (depth[i] == -1) {
         depth[i] = 0;
         dfs(i, -1, tree);
    }
  void dfs(int u, int p,
             const vector<vector<int>> &tree) {
    if (p != -1) {
       depth[u] = depth[p] + 1;
       up[u][0] = p;
for (int i = 1; i <= maxlog; i++) {
  up[u][i] = up[up[u][i - 1]][i - 1];
    for (int v : tree[u]) {
  if (v == p) continue;
       dfs(v, u, tree);
  int kth_jump(int u, int k) {
    for (int i = \max \log; i \ge 0; i—) {
       if^{((1 << i) \& k)} {
         u = up[u][i];
       }
    return u;
  }
  int lca(int u, int v) {
  if (depth[u] < depth[v]) swap(u, v);</pre>
    int diff = depth[u] - depth[v];
    u = kth_jump(u, diff);
    if (u == v) return u;
for (int i = maxlog; i >= 0; i—) {
       if (up[u][i] != up[v][i]) {
         u = up[u][i];
         v = up[v][i];
       }
    return up[u][0];
  bool on_path(int u, int v, int s) {
    int uv = lca(u, v), us = lca(u, s),
    vs = lca(v, s);
return (uv == s or (us == uv and vs == s) or
              (vs == uv and us == s));
  int dist(int u, int v)
    return depth[u] + depth[v] -
             2 * depth[lca(u, v)];
};
```

### 6.25 Maximum flow (Dinic)

**Description:** Finds the **maximum flow** in a graph network, given the **source** s and the **sink** t. Add edge from a to b with capcity c. **Time:** In general  $O(E \cdot V^2)$ , if every capacity is 1, and every vertice has in degree equal 1 or out degree equal 1 then  $O(E \cdot \sqrt{V})$ , Warning: Suffle the edges list for every vertice may take you out of the worst

```
struct Dinic {
  struct Edge {
    int to, rev;
    ll c, oc;
ll flow() {
      return max(oc - c, 0LL);
// if you need flows
  };
  vi lvl, ptṛ, q<u>;</u>
  vector<vector<Edge>> adj;
  Dinic(int n) : lvl(n), ptr(n), q(n), adj(n) {}
  void addEdge(int a, int b, ll c, ll rcap = 0) {
    adj[a].pb({b, len(adj[b]), c, c});
    adj[b].pb({a, len(adj[a]) - 1, rcap, rcap});
  ll dfs(int v, int t, ll f) {
  if (v == t || !f) return f;
  for (int &i = ptr[v]; i < len(adj[v]); i++) {</pre>
       Edge &e = adj[v][i];
          (lvl[e.to] == lvl[v] + 1)
         if (ll p = dfs(e.to, t, min(f, e.c))) {
           e.c —= p, adj[e.to][e.rev].c += p;
           return p;
    return 0;
  ll maxFlow(int s, int t) {
    ll flow = 0;
    q[0] = s;
    rep(L, 0, 31) {
       do { // 'int L=30' maybe faster for random
              // data
         lvl = ptr = vi(len(q));
         int qi = 0, qe = lvl[s] = 1;
         while (qi < qe \&\& !lvl[t]) {
           int v = q[qi++]
            for (Edge e : adj[v])
              if (!lvl[e.to] && e.c >> (30 - L))
q[qe++] = e.to,
                lvl[e.to] = lvl[v] + 1;
         while (ll p = dfs(s, t, LLONG_MAX))
            flow += p;
      } while (lvl[t]);
    return flow;
  bool leftOfMinCut(int a) { return lvl[a] != 0; }
};
```

### 6.26 Minimum Cost Flow

Description: Given a network find the minimum cost to achieve a flow of at most f. Works with **directed** and **undirected** graphs **Usage**:

• add(u, v, c, w): adds an edge from u to v with capacity c and cost w. • flow(f): return a pair (flow, cost) with the maximum flow until f with

source at s and sink at t, with the minimum cost possible. Time:  $O(N\cdot M + f\cdot m\log n)$ 

```
template <typename T>
struct MinCostFlow {
  struct Edge {
    int to;
    ll c, rc;
                 // capcity, residual capacity
    Tw;
                 // cost
  int n, s, t;
  vector<Edge> edges;
  vi2d g;
  vector<T> dist;
  vi pre;
  MinCostFlow() {}
  MinCostFlow(int n_, int _s, int _t
: n(n_), s(_s), t(_t), g(n) {}
```

```
void addEdge(int u, int v, ll c, T w) {
   g[u].pb(len(edges));
   edges.eb(v, c, 0, w);
   g[v].pb(len(edges));
   edges.eb(u, 0, 0, -w);
 // {flow, cost}
pair<ll, T> flow(ll flow_limit = LLONG_MAX) {
   ll flow = 0;
   T cost = 0;
   while (flow < flow_limit and dijkstra(s, t)) {</pre>
      ll aug = LLONG_MAX;
      for (int i = t; i != s;
    i = edges[pre[i] ^ 1].to) {
         aug = min({flow_limit — flow, aug,
                      edges[pre[i]].c});
      for (int i = t; i != s;
    i = edges[pre[i] ^ 1].to) {
        edges[pre[i]].c —= aug;
edges[pre[i] ^ 1].c += aug;
        edges[pre[i]].rc += aug;
edges[pre[i] ^ 1].rc -= aug;
      flow += aug;
      cost += (T)aug * dist[t];
   return {flow, cost};
 // Needs to be called after flow method
 vi2d paths() {
   vi2d p;
   for (;;) {
      int cur = s;
      auto &res = p.eb();
      res.pb(cur);
      while (cur != t) {
        bool found = false;
for (auto i : g[cur]) {
           auto &[v, _, c, cost] = edges[i];
if (c > 0) {
    __c;
             res.pb(cur = v);
             found = true;
        }
        if (!found) break;
      }
      if (cur != t) {
        p.ppb();
        break;
      }
   }
   return p;
private:
 bool bellman_ford(int s, int t) {
   dist.assign(n, numeric_limits<T>::max());
   pre.assign(n, -1);
   vc inq(n, false);
   queue<int> q;
   dist[s] = 0;
   q.push(s);
   while (len(q)) {
      int u = q.front();
      q.pop();
      inq[u] = false;
      for (int i : g[u]) {
  auto [v, c, w, _] = edges[i];
  auto new_dist = dist[u] + w;
  if (c > 0 and dist[v] > new_dist) {
           dist[v] = new_dist;
           pre[v] = i;
           if (not inq[v]) {
             inq[v] = true;
             q.push(v);
        }
```

```
}
     }
     return dist[t] != numeric_limits<T>::max();
  bool dijkstra(int s, int t) {
     dist.assign(n, numeric_limits<T>::max());
     pre.assign(n, -1);
     dist[s] = 0;
     using PQ = pair<T, int>;
     pqmn<PQ> pq;
     pq.emp(0, s);
     while (len(pq)) {
       auto [cost, u] = pq.top();
       pq.pop();
       if (cost != dist[u]) continue;
       for (int i : g[u]) {
  auto [v, c, _, w] = edges[i];
  auto new_dist = dist[u] + w;
  if (c > 0 and dist[v] > new_dist) {
            dist[v] = new_dist;
            pre[v] = i;
            pq.emp(new_dist, v);
         }
       }
     return dist[t] != numeric_limits<T>::max();
  }
};
```

### 6.27 Minimum Vertex Cover (already divided)

**Description:** Given a bipartite graph g with n vertices at left and m vertices at right, where g[i] are the possible right side matches of vertex i from left side, find a minimum vertex cover. The size is the same as the size of the maximum matching, and the complement is a maximum independent set.

```
vector<int> min_vertex_cover(
     vector<vector<int>> &g, int n, int m) {
  vector<int> match(m, -1), vis;
  auto find = [\&](auto \&\&self, int j) \rightarrow bool {
     if (match[j] == -1) return 1;
     vis[j] = 1;
     int di = match[j];
     for (int e : g[di])
  if (!vis[e] and self(self, e)) {
          match[e] = di;
          return 1;
     return 0;
  for (int i = 0; i < (int)g.size(); i++) {
     vis.assign(match.size(), 0);
     for (int j : g[i]) {
       if (find(find, j)) {
          match[j] = i;
          break;
       }
    }
  }
  int res =
        (int)match.size()
       (int)count(match.begin(), match.end(), -1);
  vector<char> lfound(n, true), seen(m);
  for (int it : match)
  if (it != -1) lfound[it] = false;
  vector<int> q, cover;
for (int i = 0; i < n; i++)
  if (lfound[i]) q.push_back(i);
  while (!q.empty()) {
  int i = q.back();
     q.pop_back();
     lfound[i] = 1;
for (int e : g[i])
       if (!seen[e] and match[e] !=-1) {
          seen[e] = true;
          q.push_back(match[e]);
  for (int i = 0; i < n; i++)
  if (!lfound[i]) cover.push_back(i);</pre>
```

```
for (int i = 0; i < m; i++)
  if (seen[i]) cover.push_back(n + i);</pre>
assert((int)size(cover) == res);
return cover;
```

### 6.28 Prim (MST)

**Description**: Given a graph with N vertex finds the minimum spanning tree. if there is no such three returns inf, it starts using the edges that connect with each  $s_i \in s$ , if none is provided than it starts with the edges of node 0.

```
const int MAXN(1 '00' 000);
int N:
vector<pair<ll, int>> G[MAXN];
ll prim(vi s = vi(1, 0)) {
  priority_queue<pair<ll, int>
                 vector<pair<ll, int>>,
                 greater<pair<ll, int>>>
      pq;
  vector<char> ingraph(MAXN);
  int ingraphcnt(0);
  for (auto si : s) {
    ingraphcnt++;
    ingraph[si] = true;
    for (auto &[w, v] : G[si]) pq.emplace(w, v);
  ll\ mstcost = 0;
  while (ingraphent < N and !pq.empty()) {</pre>
    ll w;
    int v;
    do {
      tie(w, v) = pq.top();
      pq.pop();
    } while (not pq.empty() and ingraph[v]);
    mstcost += w, ingraph[v] = true, ingraphcnt++;
    for (auto &[w2, v2] : G[v]) {
      pq.emplace(w2, v2);
  }
  return ingraphcnt == N ? mstcost : oo;
```

# Shortest Path With K-edges

**Description**: Given an adjacency matrix of a graph, and a number Kcomputes the shortest path between all nodes that uses exactly K edges, so for  $0 \le i, j \le N - 1$  ans[i][j] = "the shortest path between i and j that uses exactly K edges, remember to initialize the adjacency matrix with  $\infty$ . Time:  $O(N^3 \cdot \log K)$ 

```
template <typename T>
vector<vector<T>> prod(vector<vector<T>> &a,
                          vector<vector<T>> &b) {
  const T _oo = numeric_limits<T>::max();
  int n = a.size();
  vector<vector<T>> c(n, vector<T>(n, _oo));
  for (int i = 0; i < n; i++)
for (int j = 0; j < n; j++)
for (int k = 0; k < n; k++)
            (a[i][k] != _oo and b[k][j] != _oo)
         if
           c[i][j] =
                min(c[i][j], a[i][k] + b[k][j]);
  return c;
}
template <typename T>
vector<vector<T>>> shortest_with_k_moves(
    vector<vector<T>> adj, long long k)
  if (k == 1) return adj;
  auto ans = adj;
  while (k)
    ile (k) {
if (k & 1) ans = prod(ans, adj);
    k >>= 1;
    adj = prod(adj, adj);
  }
  return ans;
```

### 6.30 Strongly Connected Components (struct)

Description: Find the connected component for each edge (already in a topological order), some additional functions are also provided.

**Time**: Build: O(V + E)

```
struct SCC {
  int n, num_sccs;
  vi2d adj;
  vi scc_id;
  SCC(int _n)
: n(_n)
         num_sccs(0),
         adj(n)
         scc_id(n, -1) {}
  void add_edge(int u, int v) { adj[u].eb(v); }
  void find_sccs() {
    int timer = 1;
    vi tin(n), st;
     st.reserve(n)
    function<int(int)> dfs = [&](int u) -> int {
       int low = tin[u] = timer++, siz = len(st);
       st.eb(u)
       for (int v : adj[u])
    if (scc_id[v] < 0)
           low =
                min(low, tin[v] ? tin[v] : dfs(v));
       if (tin[u] == low)
         rep(i, siz, l
num_sccs;
                      len(st)) scc_id[st[i]] =
         st.resize(siz);
num_sccs++;
       return low;
    };
    for (int i = 0; i < n; i++)
       if (!tin[i]) dfs(i);
  vector<set<int>> build_gscc() {
    vector<set<int>> gscc;
    for (int i = 0; i < len(adj); ++i)
for (auto i : adilil)</pre>
       for (auto j : adj[i])
  if (scc_id[i] != scc_id[j])
           gscc[scc_id[i]].emplace(scc_id[j]);
    return gscc;
  vi2d per_comp() {
    vi2d ret(num_sccs);
    rep(i, 0, n) ret[scc_id[i]].eb(i);
     reverse(all(
                  // already in topological order ;)
         ret));
    return ret:
  }
};
```

#### Topological Sorting (Kahn) 6.31

Description: Finds the topological sorting in a DAG, if the given graph is not a DAG than an empty vector is returned, need to 'initialize' the INCNT as you build the graph.

Time: O(V+E)

```
const int MAXN(2 '00' 000);
int INCNT[MAXN];
vi2d GOUT(MAXN);
int N;
vi toposort() {
  vi order
  queue<int> q;
  for (int i = 0; i < N; i++)
    if (!INCNT[i]) q.emplace(i);
  while (!q.empty())
    auto u = q.front();
    q.pop();
    order.emplace_back(u)
    for (auto v : GOUT[u]) {
       INCNT[v]
       if (INCNT[v] == 0) q.emplace(v);
    }
  return len(order) == N ? order : vi();
}
```

### 6.32 Topological Sorting (Tarjan)

**Description:** Finds a the topological order for the graph, if there is no such order it means the graph is cyclic, then it returns an empty vector **Time:** O(V + E)

```
const int maxn(1 '00' 000);
int n, m;
vi g[maxn];
int not_found = 0, found = 1, processed = 2;
int state[maxn];
bool dfs(int u, vi &order)
  if (state[u] == processed) return true;
if (state[u] == found) return false;
  state[u] = found;
  for (auto v : g[u]) {
    if (not dfs(v, order)) return false;
  state[u] = processed;
  order.emplace_back(u);
  return true;
vi topo_sort() {
  vi order
  memset(state, 0, sizeof state);
  for (int u = 0; u < n; u++) {
   if (state[u] == not_found and</pre>
         not dfs(u, order))
       return {};
  reverse(all(order));
  return order;
```

### 6.33 Tree Isomorphism (not rooted)

**Description**: Two trees are considered **isomorphic** if the hash given by thash() is the same. **Time**:  $O(V \cdot \log V)$ 

```
map<vi, int> mphash;
struct Tree {
  int n;
  vi2d g;
  vi sz, cs;
  Tree(int n_{-}) : n(n_{-}), g(n), sz(n) {}
  void add_edge(int u, int v) {
  g[u].emplace_back(v);
    g[v].emplace_back(u);
  void dfs_centroid(int v, int p) {
    sz[v] = 1;
bool cent = true
    for (int u : g[v])
      if (u != p)
         dfs_centroid(u, v);
         sz[v] += sz[u];
         cent &= not(sz[u] > n / 2);
    if (cent and n - sz[v] \ll n / 2)
      cs.push_back(v);
  int fhash(int v, int p) {
    vi h;
    for (int u : g[v])
  if (u != p) h.push_back(fhash(u, v));
    sort(all(h));
    if (!mphash.count(h))
  mphash[h] = mphash.size();
    return mphash[h];
  }
  ll thash()
    cs.clear();
    h2 = fhash(cs[1], cs[0]);
return (min(h1, h2) << 30ll) + max(h1, h2);
};
```

### 6.34 Tree Isomorphism (rooted)

**Description**: Given a rooted tree find the hash of each subtree, if two roots of two distinct trees have the same hash they are considered isomorphic **Time**: hash first time in  $O(\log N_v \cdot N_v)$  where  $(N_v)$  is the of the subtree of v

```
map<vi, int> hasher;
int hs = 0;
struct RootedTreeIso {
  int n:
  vi2d adj;
  vi hashes;
  RootedTreeIso(int _n)
       : n(_n), adj(_n), hashes(_n, -1){};
  void add_edge(int u, int v) {
    adj[u].emplace_back(v);
    adj[v].emplace_back(u);
  int hash(int u, int p = -1) {
    if (hashes[u] != -1) return hashes[u];
    vi children;
    for (auto v : adj[u])
  if (v != p)
        children.emplace_back(hash(v, u));
    sort(all(children));
    if (!hasher.count(children))
      hasher[children] = hs++;
    return hashes[u] = hasher[children];
};
```

### 6.35 Tree diameter (DP)

```
const int MAXN(1 '000' 000);
int N;
vi G[MAXN];
int diameter, toLeaf[MAXN];
void calcDiameter(int u = 0, int p = -1) {
  int d1, d2;
  d1 = d2 = -1;
  for (auto v : G[u]) {
    if (v != p) {
      calcDiameter(v, u);
      d1 = max(d1, toLeaf[v]);
      tie(d1, d2) = minmax({d1, d2});
    }
  }
  toLeaf[u] = d2 + 1;
  diameter = max(diameter, d1 + d2 + 2);
}
```

### 7 Linear Algebra

# 7.1 Matrix (primitive)

```
#include "../Contest/template.cpp"
template <typename T>
struct Matrix {
  int n, m;
  valarray<valarray<T>> v;
  Matrix(int _n, int _m, int id = 0)
        n(_n), m(_m), v(valarray<T>(m), n) {
    if (id)
      for (int i = 0; i < n; i++) v[i][i] = 1;
  valarray<T>& operator[](int x) { return v[x]; }
  Matrix transpose() {
    Matrix newv(m, n);
    for (int i = 0; i < n; i++)
for (int j = 0; j < m; j++)
        newv[j][i] = (*this)[i][j];
    return newv;
  Matrix operator+(Matrix& b) {
    Matrix ret(*this);
    return ret.v += b.v;
```

```
Matrix& operator+=(Matrix& b) {
  return v += b.v;
Matrix operator*(Matrix b) {
  Matrix ret(n, b.m);
  for (int i = 0;
                    _i < n; i++)
    for (int j = 0; j < m; j++)

for (int k = 0; k < b.m; k++)

ret[i][k] += v[i][j] * b.v[j][k];
  return ret;
}
Matrix& operator*=(Matrix b) {
  return *this = *this * b;
Matrix power(ll exp) {
  Matrix in = *this;
  Matrix ret(n, n, 1);
  while (exp) {
    if (exp & 1) ret *= in;
    in *= in;
    exp >>= 1;
  return ret;
Alters current matrix.
Does gaussian elimination and puts matrix in
upper echelon form (possibly reduced).
Returns the determinant of the square matrix
with side equal to the number of rows of the
original matrix.
*/
T gaussjordanize(int reduced = 0) {
  T \det = T(1);
  int line = 0;
  for (int col = 0; col < m; col++) {
    int pivot = line;
    while (pivot < n && v[pivot][col] == T(0))
       pivot++;
    if (pivot >= n) continue;
    swap(v[line], v[pivot]);
    if (line != pivot) det *= T(-1);
    det *= v[line][line];
    v[line] /= T(v[line][col]);
    if (reduced)
       for (int i = 0; i < line; i++)
         v[i] = T(v[i][col]) * v[line];
    for (int i = line + 1; i < n; i++) {
  v[i] -= T(v[i][col]) * v[line];</pre>
    line++;
  }
  return det * (line == n);
}
Needs to be called in a square matrix that
represents a system of linear equations. Returns
{possible solution, number of solutions (2 if
infinite solutions)}
pair<vector<T>, int> solve_system(
    vector<T> results) {
  Matrix aux(n, m + 1);
  for (int i = 0; i < n; i++) {
  for (int j = 0; j < m; j++)
    aux[i][j] = v[i][j];
</pre>
    aux[i][m] = results[i];
  }
  T det = aux.gaussjordanize(1);
  int ret = 1 + (det == T(0));
  int n = results.size();
  for (int i = n - 1; i \ge 0; i—) {
    int pivot = 0;
    while (pivot < n && aux[i][pivot] == T(0))
```

```
if (pivot == n) {
          if (aux[i][m] != T(0)) ret = 0;
        } else
          swap(aux[i], aux[pivot]);
     for (int i = n - 1; i >= 0; i—) {
  for (int j = i + 1; j < n; j++)
          aux[i][m] -= aux[i][j] * aux[j][m];
     for (int i = 0; i < n; i++)
       results[i] = aux[i][m];
     for (int i = 0; i < n; i++) {
  for (int j = 0; j < m; j++)
   v[i][j] = aux[i][j];</pre>
     return {results, ret};
 Does not alter current matrix.
Returns {inverse matrix, is curent matrix
 invertable}
  pair<Matrix<T>, bool> find_inverse() {
     int n = v.size();
Matrix<T> aug(n, 2 * n);
     for (int i = 0; i < n; i++)
for (int j = 0; j < n; j++)
          aug[i][j] = v[i][j];
     for (int i = 0; i < n; i++) aug[i][n + i] = 1;
     T det = aug.gaussjordanize(1);
     Matrix<T> ret(n, n);
     for (int i = 0; i < n; i++) {
       ret[i]
             valarray<T>(aug[i][slice(n, n, 1)]);
     return {ret, det != T(0)};
  // Returns rank of matrix. Does not alter it.
  int get_rank() const
     if (m == 0) return 0;
     Matrix<T> aux(*this);
     aux.gaussjordanize();
     int resp = 0;
     for (int i = 0; i < n; i++)
       resp += (aux[i] != valarray<T>(m)).sum();
     return resp;
};
    Math
      Arithmetic Progression Sum
8.1
Usage:
   • s: first term
• d: common difference
• n: number of terms
ll arithmeticProgressionSum(ll s, ll d, ll n) {
  return (s + (s + d * (n - 1))) * n / 2ll;
}
8.2 Binomial
Time: O(N \cdot K)
Memory: O(K)
ll binom(ll n, ll k) \{
  if (k > n) return 0;
vll dp(k + 1, 0);
  dp[0] = 1;
  for (ll i = 1; i <= n; i++)
  for (ll j = k; j > 0; j—)
    dp[j] = dp[j] + dp[j - 1];
  return dp[k];
}
```

pivot++;

### 8.3 Binomial MOD

**Description**: find  $\binom{n}{k}$  (mod MOD)

### 8.4 Chinese Remainder Theorem

```
Description: Find the solution X to the N modular equations. x\equiv a_1(modm_1)\\ \cdots\\ x\equiv a_n(modm_n) (1)
```

The  $m_i$  don't need to be coprime, if there is no solution then it returns -1.

```
tuple<ll, ll, ll> ext_gcd(ll a, ll b) {
  if (!a) return {b, 0, 1};
  auto [g, x, y] = ext_gcd(b % a, a);
  return \{g, y - b / a * x, x\};
template <typename T = ll >
struct crt {
  T a, m;
  crt() : a(0), m(1) {}
crt(T a_, T m_) : a(a_), m(m_) {}
crt operator*(crt C) {
     auto [g, x, y] = ext\_gcd(m, C.m);
if ((a - C.a) % g != 0) a = -1;
if (a == -1 \text{ or } C.a == -1) return crt(-1, 0);
     T lcm = m / g * C.m;
     T ans =
     a + (x * (C.a - a) / g % (C.m / g)) * m;
return crt((ans % lcm + lcm) % lcm, lcm);
};
template <typename T = ll >
struct Congruence {
  T a, m;
template <typename T = ll>
T chinese_remainder_theorem(
     const vector<Congruence<T>> &equations) {
  crt<T> ans;
  for (auto \&[a_-, m_-] : equations) {
     ans = ans * crt<T>(a_, m_);
  return ans.a;
```

### 8.5 Derangement / Matching Problem

```
Description: Computes the derangement of N, which is given by the formula: D_N = N! \left(1 - \frac{1}{1!} + \frac{1}{2!} - \frac{1}{3!} + \ldots + (-1)^N \frac{1}{N!}\right)
Time: O(N)
```

```
#warning Remember to call precompute !
const ll MOD = 1e9 + 7;
const int MAXN(1 '000' 000);
ll fats[MAXN + 1];
void precompute() {
```

```
fats[0] = 1;
for (ll i = 1; i <= MAXN; i++) {</pre>
     fats[i] = (fats[i - 1] * i) % MOD;
}
ll fastpow(ll a, ll p, ll m) {
  ll ret = 1;
  while (p)
     if (p \& 1) ret = (ret * a) % MOD;
     p >>= 1;
     a = (a * a) % MOD;
  return ret;
}
ll divmod(ll a, ll b) { return (a * fastpow(b, MOD - 2, MOD)) % MOD;
ll derangement(const ll n) {
  ll ans = fats[n];
for (ll i = 1; i <= n; i++) {
    ll k = divmod(fats[n], fats[i]);</pre>
     if (i & 1) {
       ans = (ans - k + MOD) % MOD;
     } else {
       ans = (ans + k) % MOD;
  return ans;
```

### 8.6 Euler Phi $\varphi(N)$

**Description**: Computes the number of positive integers less than N that are coprimes with N, in  $O(\sqrt{N})$ .

### 8.7 Euler phi $\varphi(N)$ (in range)

**Description:** Computes the number of positive integers less than n that are coprimes with N, in the range [1, N], in  $O(N \log N)$ .

```
const int MAX = 1e6;
vi range_phi(int n) {
  bitset<MAX> sieve;
  vi phi(n + 1);
  iota(phi.begin(), phi.end(), 0);
  sieve.set();
  for (int p = 2; p <= n; p += 2) phi[p] /= 2;
  for (int p = 3; p <= n; p += 2) {
    if (sieve[p]) {
      for (int j = p; j <= n; j += p) {
         sieve[j] = false;
         phi[j] /= p;
         phi[j] *= (p - 1);
      }
    }
  return phi;
}</pre>
```

### 8.8 FFT convolution and exponentiation

```
const ld PI = acos(-1);
/* change the ld to doulbe may increase
 * performance =D */
struct num {
  ld a{0.0}, b{0.0};
  num() {}
  num(ld na) : a{na} {}
```

```
num(ld na, ld nb) : a{na}, b{nb} {}
  const num operator+(const num &c) const {
    return num(a + c.a, b + c.b);
  const num operator—(const num &c) const {
    return num(a - c.a, b - c.b);
  const num operator*(const num &c) const {
     return num(a * c.a - b * c.b,
a * c.b + b * c.a);
  const num operator/(const ll &c) const {
     return num(a / c, b / c);
void fft(vector<num> &a, bool invert) {
  int n = len(a);
  for (int i = 1, j = 0; i < n; i++) {
  int bit = n >> 1;
     for (; j & bit; bit >>= 1) j ^= bit;
     j ^= bit;
     if (i < j) swap(a[i], a[j]);
  for (int sz = 2; sz <= n; sz <<= 1) {
    ld ang = 2 * PI / sz * (invert ? -1 : 1);
    num wsz(cos(ang), sin(ang));
    for (int in a continue);</pre>
     for (int i = 0; i < n; i += sz) {
       num w(1);
       rep(j, 0, sz / 2)  {
          num u = a[i + j],

v = a[i + j + sz / 2] * w;
          a[i + j] = u + v;

a[i + j + sz / 2] = u - v;

w = w * wsz;
       }
  if (invert)
     for (num \&x : a) x = x / n;
vi conv(vi const a, vi const b) {
  vector<num> fa(all(a));
  vector<num> fb(all(b));
  int n = 1; while (n < len(a) + len(b)) n <<= 1;
  fa.resize(n);
  fb.resize(n);
  fft(fa, false);
fft(fb, false);
rep(i, 0, n) fa[i] = fa[i] * fb[i];
  fft(fa, true);
  vi result(n);
  rep(i, 0, n) result[i] = round(fa[i].a);
  while (len(result) and result.back() == 0)
     result.pop_back();
  /* Unconment this line if you want a boolean
     convolution*/
  // for (auto &xi : result) xi = min(xi, 1ll);
  return result;
vll poly_exp(vll &ps, int k) {
  vll ret(len(ps));
  auto base = ps;
  ret[0] = 1;
  while (k) {
  if (k & 1) ret = conv(ret, base);
     k >>= 1;
     base = conv(base, base);
  }
  return ret;
8.9 Factorial Factorization
Description: Computes the factorization of N! in \varphi(N) * \log N
```

**Description**: Computes the factorization of N! in  $\varphi(N) * \log N$ **Time**:  $O(\varphi(N) \cdot \log N)$ 

```
ll E(ll n, ll p) {
    ll k = 0, b = p;
    while (b <= n) {
        k += n / b;
        b *= p;
}</pre>
```

```
return k;
}
map<ll, ll> factorial_factorization(
     ll n, const vll &primes) {
   map<ll, ll> fs;
   for (const auto &p : primes) {
  if (p > n) break;
     fs[p] = E(n, p);
   return fs;
8.10 Factorization
Description: Computes the factorization of N.
map<ll, ll> factorization(ll n) {
  map<ll, ll> ans;
  for (ll i = 2; i * i <= n; i++) {</pre>
     ll count = 0;
for (; n % i == 0; count++, n /= i)
     if (count) ans[i] = count;
   íf (n > 1) ans[n]++;
   return ans;
8.11 Factorization (Pollard's Rho)
Description: Factorizes a number into its prime factors.
Time: O(N^{(\frac{1}{4})} * \log(N)).
ll mul(ll a, ll b, ll m) {
   ll ret =
        a * b - (ll)((ld)1 / m * a * b + 0.5) * m;
   return ret < 0 ? ret + m : ret;
ll pow(ll a, ll b, ll m) {
  ll ans = 1;
for (; b > 0; b /= 2ll, a = mul(a, a, m)) {
  if (b % 2ll == 1) ans = mul(ans, a, m);
   return ans;
}
bool prime(ll n) {
   if (n < 2) return 0; if (n <= 3) return 1;
   if (n % 2 == 0) return 0;
  ll r = __builtin_ctzll(n - 1), d = n >> r;
for (int a : {2, 325, 9375, 28178, 450775,
9780504, 795265022}) {
     ll x = pow(a, d, n);
if (x == 1 \text{ or } x == n - 1 \text{ or a } \% n == 0)
        continue;
     for (int_j = 0; j < r - 1; j++) {
        x = mul(x, x, n);
if (x == n - 1) break;
     if (x != n - 1) return 0;
   return 1;
}
ll rho(ll n) {
  if (n == 1 or prime(n)) return n;
   auto f = [n](ll x) \{ return mul(x, x, n) + 1; \};
   ll x = 0, y = 0, t = 30, prd = 2, x0 = 1, q;
   while (t % 40 != 0 or gcd(prd, n) == 1) {
     if (x == y) x = ++x0, y = f(x);
     q = mul(prd, abs(x - y), n);
     if (q != 0) prd = q;
     x = f(x), y = f(f(y)), t++;
   return gcd(prd, n);
}
vector<ll> fact(ll n)
   if (n == 1) return {};
if (prime(n)) return {n};
  ll d = rho(n);
vector<ll> l = fact(d), r = fact(n / d);
   l.insert(l.end(), r.begin(), r.end());
```

```
return l;
}

8.12 Fast Pow
Description: Computes a<sup>b</sup> (mod m)
Time: O(log B).

Il fpow(ll a, ll b, ll m) {
    ll ret = 1;
    while (b) {
        if (b & 1) ret = (ret * a) % m;
        b >>= 1;
        a = (a * a) % m;
    }
    return ret;
```

### 8.13 Find diophantine equation solution

**Description:** Given a b, c finds the solution to the equation ax + by = c, the result will be stored in the reference variables x0 and y0. **Time:**  $O(\log min(a,b))$ 

```
template <typename T>
tuple<T, T, T> ext_gcd(T a, T b) {
  if (b == 0) return {a, 1, 0};
  auto [d, x1, y1] = ext_gcd(b, a % b);
  return {d, y1, x1 - y1 * (a / b)};
template <typename T>
tuple<bool, T, T> find_any_solution(T a, T b,
                                          T c) {
  assert(a != 0 or b != 0);
#warning Be careful with overflow, use __int128 if
   needed!
  auto [d, x0, y0] =
       ext_{gcd}(a < 0 ? -a : a, b < 0 ? -b : b);
  if (c % d) return {false, 0, 0};
  x0 *= c / d;

y0 *= c / d;
  if (a < 0) x0 = -x0;
if (b < 0) y0 = -y0;
  return {true, x0, y0};
```

### 8.14 Find linear recurrence (Berlekamp-Massey)

**Description**: Given the first N terms of a linear recurrence finds the smallest recurrence that matches the sequence.

Time:  $O(N^2)$ 

Warning: Works faster if the mod is const but can be also be a parameter. Absolute magic!

```
const ll \ mod = 998244353;
ll modpow(ll b, ll e) {
  ll ans = 1;
for (; e; b = b * b % mod, e /= 2)
  if (e & 1) ans = ans * b % mod;
  return ans;
vl berlekampMassey(vll s) {
  int n = len(s), L = 0, m = 0;
if (!n) return {};
vll C(n), B(n), T;
C[0] = B[0] = 1;
  ll b = 1;
rep(i, 0, n) {
     ++m;
ll d = s[i] % mod;
     rep(j, 1, \hat{L} + 1) \hat{d} = (d + C[j] * s[i - j]) % mod;
     if (!d) continue;
     T = C;
     ll coef = d * modpow(b, mod - 2) % mod;
     rep(j, m, n) C[j] =
           (C[j] - coef * B[j - m]) % mod;
     if (2 * L > i) continue;
     L = i + 1 - L;
     B = T;
     b = d;
     m = 0;
```

```
}
C.resize(L + 1);
C.erase(C.begin());
for (ll &x : C) x = (mod - x) % mod;
return C;
}
```

### 8.15 Find multiplicatinve inverse

### 8.16 GCD

```
ll gcd(ll a, ll b) {
  return b ? gcd(b, a % b) : a;
}
```

### 8.17 Gauss XOR elimination / XOR-SAT

**Description:** Execute gaussian elimination with xor over the system Ax=b in. The add method must receive a bitset indicating which variables are present in the equation, and the solution of the equation.

Time:  $O(\frac{nm^2}{64})$ 

```
const int MAXXI = 2009;
using Equation = bitset<MAXXI>;
struct GaussXor {
   vector<char> B
   vector<Equation> A;
   void add(const Equation &ai, bool bi) {
     A.push_back(ai);
     B.push_back(bi);
  pair<bool, Equation> solution() {
  int cnt = 0, n = A.size();
     Equation vis;
     vis.set();
     Equation x;
     for (int j = MAXXI - 1, i; j \ge 0; j—) {
        for (i = cnt; i < n; i++) {
  if (A[i][j]) break;</pre>
        if (i == n) continue;
        swap(A[i], A[cnt]), swap(B[i], B[cnt]);
        i = cnt++
        vis[j] = 0;
        for (int k = 0; k < n; k++) {
  if (i == k || !A[k][j]) continue;
  A[k] ^= A[i];
  B[k] ^= B[i];</pre>
       }
     }
     x = vis;
     for (int i = 0; i < n; i++) {
        int acum = 0;
for (int j = 0; j < MAXXI; j++) {
          if (!A[i][j]) continue; if (!vis[j]) {
             vis[j] = 1;
             x[j] = acum ^ B[i];
          acum ^= x[j];
        if (acum != B[i])
          return {false, Equation()};
     return {true, x};
};
```

# 8.18 Integer partition

```
Description: Find the total of ways to partition a given number N in such
way that none of the parts is greater than K.

Time: O(N \cdot min(N, K))
Memory: O(N)
```

Warning: Remember to memset everything to -1 before using it

```
const ll MOD = 10000000007;
const int MAXN(100);
               1];
ll memo[MAXN +
ll dp(ll n, ll k = oo) {
  if (n == 0) return 1;
  ll \&ans = memo[n];
  if (ans !=-1) return ans;
  for (int i = 1; i <= min(n, k); i++) {
   ans = (ans + dp(n - i, k)) % MOD;
 return ans;
```

### 8.19 LCM

```
ll gcd(ll a, ll b) {
  return b ? gcd(b, a % b) : a;
ll lcm(ll a, ll b) \{ return a / gcd(a, b) * b; \}
```

### 8.20 Linear Recurrence

Description: Find the n-th term of a linear recurrence, given the recurrence rec and the first K values of the recurrence, remember that first k[i] is the value of f(i), considering 0-indexing. **Time**:  $O(K^3 \log N)$ 

```
template <typename T>
vector<vector<T>> prod(vector<vector<T>> &a,
                          vector<vector<T>> &b,
                          const ll mod) {
  int n = a.size();
  vector<vector<T>> c(n, vector<T>(n));
  for (int i = 0; i < n; i++) {
  for (int j = 0; j < n; j++)</pre>
      for (int k = 0; k < n; k++) {
         c[i][j] = (c[i][j] +
                     ((a[i][k] * b[k][j]) % mod)) %
                    mod;
    }
  }
  return c;
template <typename T>
vector<vector<T>> fpow(vector<vector<T>> &xs,
                          ll p, ll mod) {
  vector<vector<T>> ans(xs.size(),
                           vector<T>(xs.size()));
  for (int i = 0; i < (int)xs.size(); i++)
    ans[i][i] = 1;
  for (auto b = xs; p;
    p >>= 1, b = prod(b, b, mod))
if (p & 1) ans = prod(ans, b, mod);
  return ans;
ll linear_req(vector<vector<ll>> rec,
                vector<ll> first_k, ll n, ll mod) {
  int k = first_k.size();
  if (n < k) {
    return first_k[n];
  ll n2 = n - k + 1;
rec = fpow(rec, n2, mod);
  for (int i = 0; i < k; i++) {
    ret = (ret +
             (rec.back()[i] * first_k[i]) % mod) %
  return ret;
```

### 8.21 List N elements choose K

**Description**: Process every possible combination of K elements from Nelements, thoose index marked as 1 in the index vector says which elments are choosed at that moment.

Time:  $O({N \choose K} \cdot O(process))$ 

```
void process(vi &index) {
  for (int i = 0; i < len(index); i++) {</pre>
    if (index[i])
       cout << i << " \n"[i == len(index) - 1];
}
void n_choose_k(int n, in k) {
  vi index(n):
  fill(index.end() - k, index.end(), 1);
  do {
    process(index);
  } while (next_permutation(all(index)));
```

### 8.22 List primes (Sieve of Eratosthenes)

```
const ll\ MAXN = 2e5;
vll list_primes(ll n = MAXN) {
  vll ps;
  bitset<MAXN + 1> sieve;
  sieve.set()
  sieve.reset(1);
for (ll i = 2; i <= n; ++i)
    if (sieve[i]) ps.push_back(i);
    for (ll j = i * 2; j <= n; j += i) {
      sieve.reset(j);
    }
  return ps;
```

#### 8.23Matrix exponentiation

```
const ll MOD = 1 '000' 000'007;
template <typename T>
vector<vector<T>> prod(vector<vector<T>> &a
                            vector<vector<T>> &b) {
  int n = len(a);
  vector<vector<T>> c(n, vector<T>(n));
  for (int i = 0; i < n; i++) {
  for (int j = 0; j < n; j++) {
    for (int k = 0; k < n; k++) {
        c[i][j] = (c[i][j] +</pre>
                       ((a[i][k] * b[k][j]) % MOD)) %
                      MOD:
    }
  }
  return c;
}
template <typename T>
vector<vector<T>> fpow(vector<vector<T>> &xs,
                            ll p) {
  vector<vector<T>> ans(len(xs)
                             vector<T>(len(xs)))
  for (int i = 0; i < len(xs); i++) ans[i][i] = 1;
  auto b = xs;
  while (p)
    if (p \& 1) ans = prod(ans, b);
     p >>= 1;
     b = prod(b, b);
  return ans;
}
```

```
Time:
```

```
• Convolution O(N \cdot \log N),
• Exponentiation: O(\log K \cdot N \cdot \log N)
```

```
template <int _mod>
struct mint {
  ll expo(ll b, ll e) {
     ll ret = 1;
while (e) {
  if (e % 2) ret = ret * b % _mod;
        e /= 2, b = b * b % _mod;
     return ret;
   ĺl inv(ll b) { return \mathsf{expo}(\mathsf{b},\ \_\mathsf{mod}\ -\ \mathsf{2}); }
  using m = mint;
  ll v;
mint()
           : v(0) {}
  mint(). V(0) {}
mint(|| V_) {
    if (v_ >= _mod or v_ <=
    if (v_ < 0) v_ += _mod;
    v = v_;
                   _mod or v_ <= -_mod) v_ %= _mod;
  m &operator+=(const m &a) {
     v += a.v;
if (v >= _mod) v -= _mod;
     return *this;
  m &operator—=(const m &a) {
         -= a.ν;
     if (v < 0) v += \_mod;
     return *this;
  m &operator*=(const m &a) {
     v = v * ll(a.v) % _mod;
     return *this;
  m &operator/=(const m &a) {
     v = v * inv(a.v) % _mod;
     return *this:
  m operator—() { return m(-v); }
  m & operator = (ll e) {
     if (e < 0)
          = inv(v);
        e = -e;
     v = expo(v, e);
     // possivel otimizacao:
     // cuidado com 0^0
// v = expo(v, e%(p-1));
     return *this:
  bool operator==(const m &a) {    return v == a.v;    }
  bool operator!=(const m &a) { return v != a.v; }
  friend istream &operator>>(istream &in, m &a) {
     ll val;
     in >> val
     a = m(val);
     return in;
  friend ostream &operator<<(ostream &out, m a) {
     return out << a.v;
  friend m operator+(m a, m b) { return a += b;
friend m operator-(m a, m b) { return a -= b;
  friend m operator*(m a, m b) { return a *= b; }
friend m operator/(m a, m b) { return a /= b; }
friend m operator^(m a, ll e) { return a ^= e; }
const ll MOD1 = 998244353;
const ll MOD2 = 754974721;
const ll MOD3 = 167772161;
template <int _mod>
void ntt(vector<mint<_mod>> &a, bool rev) {
  int n = len(a);
  auto b = a;
  assert(!(n \& (n-1)));
  mint<_mod> g = 1;
while ((g ^ (_mod / 2)) == 1) g += 1;
  if (rev)^{\overline{g}} = 1 / g;
  for (int step = n / 2; step; step /= 2) {
  mint<_mod> w = g ^ (_mod / (n / step)),
```

```
wn = 1;

for (int i = 0; i < n / 2; i += step) {

  for (int j = 0; j < step; j++) {

    auto u = a[2 * i + j],
                 v = wn * a[2 * i + j + step];
           b[i + j] = u + v;
          b[i + n' / 2 + j] = u - v;
        \dot{w}n = wn * w;
     swap(a, b);
  if (rev) {
  auto n1 = mint<_mod>(1) / n;
     for (auto \&x : a) x *= n1;
}
template <ll _mod>
vector<mint<_mod>> convolution(
     const vector<mint<_mod>> &a
     const vector<mint<_mod>> &b) {
  vector<mint<_mod>> l(all(a)), r(all(b));
int N = len(l) + len(r) - 1, n = 1;
while (n <= N) n *= 2;</pre>
   l.resize(n), r.resize(n);
  ntt(l, false), ntt(r, false);
for (int i = 0; i < n; i++) l[i] *= r[i];
ntt(l, true);
l.resize(N);</pre>
   // Uncommnent for a boolean convolution :)
  for (auto& li : l) {
    li.v = min(li.v, 1ll);
   */
  return l;
template <ll _mod>
vector<mint<_mod>> poly_exp(
     vector<mint<_mod>> &ps, int k) {
   vector<mint<_mod>> ret(len(ps));
  auto base = ps;
  ret[0] = 1;
  while (k) {
     if (k & 1) ret = convolution(ret, base);
     k > > = 1;
     base = convolution(base, base);
  return ret;
       NTT integer convolution and exponentiation (2)
8.25
        mods) modules)
\textbf{Description} \hbox{: } \textbf{Computes the convolution between the two polynomials and}.
Time: O(N \log N)
Warning: This is pure magic!
template <int _mod>
struct mint
   ruct mint {
ll expo(ll b, ll e) {
     ll ret = 1;
while (e) {
        if (e % 2) ret = ret * b % _mod;
        e /= 2, b = b * b % _mod;
     return ret;
   ll inv(ll b) \{ return expo(b, \_mod - 2); \}
   using m = mint;
  ll v;
mint() : v(0) {}
  mint(ll v_) {
    if (v_ >= _mod or v_ <= __mod) v_ %= _mod;
    if (v_ < 0) v_ += _mod;
    v = v_;
  m &operator+=(const m &a) {
     v += a.v;
if (v >= _mod) v -= _mod;
return *this;
  m &operator—=(const m &a) {
```

```
if (v < 0) v += \_mod;
     return *this;
  m &operator*=(const m &a) {
     v = v * ll(a.v) % _mod;
     return *this;
  m &operator/=(const m &a) {
     v = v * inv(a.v) % _mod;
     return *this;
  }
  m operator—() { return m(-v); }
  m &operator^=(ll e) {
     if (e < 0)
       v = inv(v);
       e = -e;
     v = expo(v, e);
     // possivel otimizacao:
     // cuidado com 0^0
// v = expo(v, e%(p-1));
     return *this;
  bool operator==(const m &a) {    return v == a.v;    }
  bool operator!=(const m &a) { return v != a.v; }
  friend istream &operator>>(istream &in, m &a) {
     ļl val;
     in >> val
     a = m(val);
     return in;
  friend ostream &operator<<(ostream &out, m a) {
     return out << a.v;
  friend m operator+(m a, m b) { return a += b; friend m operator-(m a, m b) { return a -= b;
  friend m operator*(m a, m b) { return a *= b; }
friend m operator/(m a, m b) { return a /= b; }
friend m operator^(m a, ll e) { return a ^= e; }
const ll MOD1 = 998244353;
const ll\ MOD2 = 754974721;
const ll\ MOD3 = 167772161;
template <int _mod>
void ntt(vector<mint<_mod>> &a, bool rev) {
  int n = len(a);
  auto b = a;
  assert(!(n \& (n-1)));
  mint<_mod> g = 1;
while ((g ^ (_mod / 2)) == 1) g += 1;
  if (rev)^{\circ}g = 1 / g;
  for (int step = n / 2; step; step /= 2) {
  mint<_mod> w = g_^ (_mod / (n / step)),
     wn = 1;
for (int i = 0; i < n / 2; i += step) {
       for (int j = 0; j < step; j++) {
          auto u = a[2 * i + j],
                v = wn * a[2 * i + j + step];
          b[i + j] = u + v;
          b[i + n / 2 + j] = u - v;
       \dot{w}n = wn * w;
     swap(a, b);
  if (rev) {
   auto n1 = mint<_mod>(1) / n;
     for (auto \&x : a) x *= n1;
tuple<ll, ll, ll> ext_gcd(ll a, ll b) {
  if (!a) return {b, 0, 1};
auto [g, x, y] = ext_gcd(b % a, a);
  return \{g, y - b / a * x, x\};
template <typename T = ll >
struct crt {
  T a, m;
  crt() : a(0), m(1) {}
crt(T a_, T m_) : a(a_), m(m_) {}
crt operator*(crt C) {
     auto [g, x, y] = ext_gcd(m, C.m);
```

```
if ((a - C.a) % g != 0) a = -1;
if (a == -1 \text{ or } C.a == -1) return crt(-1, 0);
     T \stackrel{\cdot}{lcm} = m / g * C.m;
     T ans =
     a + (x * (C.a - a) / g % (C.m / g)) * m;
return crt((ans % lcm + lcm) % lcm, lcm);
};
template <typename T = ll >
struct Congruence {
  T a, m;
template <typename T = ll >
T chinese_remainder_theorem(
     const vector<Congruence<T>> &equations) {
  crt<T> ans;
  for (auto \&[a_-, m_-] : equations) {
    ans = ans * crt<T>(a_{-}, m_{-});
  return ans.a;
}
#define int long long
template <ll m1, ll m2>
vll merge_two_mods(const vector<mint<m1>> &a,
                       const vector<mint<m2>> &b) {
  int n = len(a);
  vll ans(n);
  for (int i = 0; i < n; i++) {
     auto cur = crt<ll>();
     auto ai = a[i].v;
     auto bi = b[i].v
     cur = cur * crt<ll>(ai, m1);
     cur = cur * crt<ll>(bi, m2);
     ans[i] = cur.a;
  return ans;
}
vector<mint<MOD1>> l(all(a)), r(all(b));
  int N = len(l) + len(r) -1, n = 1; while (n <= N) n *= 2;
  l.resize(n), r.resize(n);
  ntt(l, false), ntt(r, false);
for (int i = 0; i < n; i++) l[i] *= r[i];</pre>
  ntt(l, true)
l.resize(N);
          true);
  vector<mint<MOD2>> l2(all(a)), r2(all(b));
  12.resize(n), r2.resize(n)
  ntt(l2, false), ntt(r2, fals
rep(i, 0, n) l2[i] *= r2[i];
  ntt(l2, true);
  l2.resize(N);
  return merge_two_mods(l, l2);
vll poly_exp(const vll &xs, ll k) {
  vll_ret(len(xs));
  ret[0] = 1;
  auto base = xs;
  while (k)
     if (k \& 1) ret = convolution_2mods(ret, base);
     k >>= 1:
     base = convolution_2mods(base, base);
  return ret;
```

# 8.26 Polyominoes

**Usage**: buildPolyominoes(x) creates every polyomino until size x, and put it in polyominoes[x], access polyomino.v to find the vector of pairs representing the coordinates of each piece, considering that the polyomino was 'rooted' in coordinate (0,0).

Warning: note that when accessing polyominoes[x] only the first x coordinates are valid.

```
struct Polyominoes {
  pii v[MAXP];
  ll id;
  int n:
  Polyominoes() {
    n = 1;
v[0] = {0, 0};
normalize();
  pii &operator[](int i) { return v[i]; }
  bool add(int a, int b) {
  for (int i = 0; i < n; i++)
    if (v[i].first == a and v[i].second == b)</pre>
          return false;
     v[n++] = pii(a, b);
     normalize();
     return true;
  id = 0;
     for (int i = 0; i < n; i++) {
  v[i].first -= mnx, v[i].second -= mny;
  id |= (1LL << (pos[v[i].first] +</pre>
                           v[i].second));
  }
vector<Polyominoes> polyominoes[MAXP + 1];
void buildPolyominoes(int mxN = 10) {
  vector<pair<int, int>> dt(
  \{\{1, 0\}, \{-1, 0\}, \{0, -1\}, \{0, 1\}\}\}; for (int i = 0; i <= mxN; i++) polyominoes[i].clear();
  Polyominoes init;
  queue<Polyominoes> q;
  unordered_set<int64_t> used;
  q.push(init);
  used.insert(init.id);
  while (!q.empty()) {
     Polyominoes u = q.front();
     q.pop();
     polyominoes[u.n].push_back(u);
     if (u.n == mxN) continue;
for (int i = 0; i < u.n; i++) {
       for (auto [dx, dy] : dt) {
  Polyominoes to = u;
          bool ok = to.add(to[i].first + dx,
                                 to[i].second + dy);
          if (ok and !used.count(to.id)) {
             q.push(to);
             used.insert(to.id);
       }
    }
  }
```

# 9 Outside

# 9.1 alien trick

```
int n, k, l;
string s;
pi solve(vector<int> &v, int lambda) {
    // associar um custo lambda para ser subtraido
    // quando realizamos uma çãoperao dp[i] — melhor
    // profit que tivemos considerando as i
    // primeiras çõposies cnt[i] — quantas çõoperaes
    // utilizamos para chegarno valor de dp[i]
    vector<int> dp(n + 1);
    vector<int> cnt(n + 1);
    dp[0] = 0;
    cnt[0] = 0;
    for (int i = 1; i <= n; i++) {
        dp[i] = dp[i - 1];
        cnt[i] = cnt[i - 1];
        int id = i - 1;
        dp[i] += v[id];</pre>
```

```
int lo = max(0ll, id - l + 1);
int s = dp[lo] + (id - lo + 1) - lambda;
     if (s > dp[i]) {
        dp[i] = s;
        cnt[i] = cnt[lo] + 1;
   return {dp[n], cnt[n]};
int aliens_trick(vector<int> &v) {
  int l = 0, r = n;
while (l < r) {
     int mid = (l + r) \gg 1;
     pi ans = solve(v, mid);
(ans.sec > k) ? l = mid + 1 : r = mid;
  pi ans = solve(v, l);
   return ans.fir + (l * k);
signed main() {
   ios_base::sync_with_stdio(false);
   cin.tie(NULĹ);
  cin >> n >> k >> l >> s;
vector<int> a(n);
   vector<int> b(n);
  for (int i = 0; i < n; i++) {
    a[i] = 1, b[i] = 0;
    if (s[i] >= 'A' && s[i] <= 'Z') {
        a[i] ^= 1;
        b[i] ^= 1;
     }
   cout << n - max(aliens_trick(a)</pre>
                         aliens_trick(b))
         << endl;
   return 0;
// https://codeforces.com/contest/1279/problem/F
```

#### 9.2 catalan

```
Description: Recursive formula: C_0 = C_1 = 1 C_n = \sum_{k=0}^{n-1} C_k C_{n-1-k}, n \geq 2 Analytical formula: C_n = \binom{2n}{n} - \binom{2n}{n-1} = \frac{1}{n+1} \binom{2n}{n}, n \geq 0 The first few numbers Catalan numbers, C_n (starting from zero): 1, 1, 2, 5, 14, 42, 132, 429, 1430, \ldots The Catalan number C_n is the solution for:
```

- $\bullet\,$  Number of correct bracket sequence consisting of n opening and n closing brackets.
- The number of rooted full binary trees with n+1 leaves (vertices are not numbered). A rooted binary tree is full if every vertex has either two children or no children.
- The number of ways to completely parenthesize n+1 factors.
- The number of triangulations of a convex polygon with n+2 sides (i.e. the number of partitions of polygon into disjoint triangles by using the diagonals)
- $\bullet\,$  The number of ways to connect the 2n points on a circle to form n disjoint chords.
- ullet The number of non-isomorphic full binary trees with n internal nodes (i.e. nodes having at least one son).
- The number of monotonic lattice paths from point (0,0) to point (n,n) in a square lattice of size  $n \times n$ , which do not pass above the main diagonal (i.e. connecting (0,0) to (n,n)).
- Number of permutations of length n that can be stack sorted (i.e. it can be shown that the rearrangement is stack sorted if and only if there is no such index i < j < k, such that  $a_k < a_i < a_j$ ).
- The number of non-crossing partitions of a set of n elements.
- The number of ways to cover the ladder  $1 \dots n$  using n rectangles (The ladder consists of n columns, where  $i^{th}$  column has a height i).

```
#include <bits/stdc++.h>
using namespace std;
const int MOD = 10000000007;
typedef long long ll;
ll extGcd(ll a, ll b, ll &x, ll &y) {
  if (b == 0) {
    x = 1, y = 0;
    return a;
} else {
    ll g = extGcd(b, a % b, y, x);
    y -= (a / b) * x;
```

```
return g;
   }
ĺl inv(ll a) {
   ll inv_x, y;
extGcd(a, MOD, inv_x, y);
return (inv_x % MOD + MOD) % MOD;
const int MAXN = 4000010;
ll fat[MAXN], ifat[MAXN];
void init() {
   fat[0] = 1;
   fact[0] = 1;
for (int i = 1; i < MAXN; i++)
  fat[i] = (fat[i - 1] * i) % MOD;
ifat[MAXN - 1] = inv(fat[MAXN - 1]);
for (int i = MAXN - 2; i >= 0; i—)
  ifat[i] = (ifat[i + 1] * (i + 1)) % MOD;
assert(ifat[0] == 1);
Il C(int n, int k) {
   if (k > n) return 0;
   return (fat[n] *
               ((ifat[k] * ifat[n - k]) % MOD)) %
ĺl catalan(int n) {
   return (\dot{C}(2 * \dot{n}, \dot{n}) - C(2 * \dot{n}, \dot{n} - 1) + MOD) %
             MOD:
Il f(int x1, int y1, int x2, int y2) {
  int y = y2 - y1, x = x2 - x1;
  if (y < 0 or x < 0) return 0;</pre>
   return C(x + y, x);
int x = 0 + k - c;
if (x < 0) return 0;
   MOD) %
             MOD;
}
9.3 centroid
#include <bits/stdc++.h>
using namespace std;
const int MAXN = 500010;
typedef pair<int, int> pii;
namespace Centroid
vector<int> adj[MAXN];
int sub[MAXN];
int n:
void init(int n1) {
   n = n1;
   for (int i = 0; i < n; i++) adj[i].clear();
void addEdge(int a, int b) {
  adj[a].push_back(b);
  adj[b].push_back(a);
int dfsS(int u, int p) {
   sub[u] = 1;
   for (int to : adj[u]) {
```

if (to != p) sub[u] += dfsS(to, u);

if (to != p and sub[to] > n / 2)

for (int to : adj[u]) {
 if (to != p and (sub[to] \* 2) == n)

return sub[u];

pii dfsC(int u, int p) {
 for (int to : adj[u]) {

return pii(u, u);

return dfsC(to, u);

return pii(u, to);

```
pii findCentroid() {
  dfsS(0,
  return dfsC(0, -1);
   // namespace Centroid
9.4 centroid decomposition
#include <bits/stdc++.h>
using namespace std;
typedef long long ll;
// O(N*log(N))
// Centroid Decomposition
const int MAXN = 200010;
namespace CD
vector<int> adj[MAXN];
int dad[MAXN], sub[MAXN];
bool rem[MAXN];
int centroidRoot, n;
void init(int n1) {
  n = n1;
  for (int i = 0; i < n; i++) {
    adj[i].clear();
    rem[i] = false;
int dfs(int u, int p) {
  sub[u] = 1;
  for (int to : adj[u])
    if (!rem[to] and to != p)
       sub[u] += dfs(to, u);
  return sub[u];
int centroid(int u, int p, int sz) {
  for (auto to : adj[u])
    if (!rem[to] \text{ and to } != p \text{ and sub}[to] > sz / 2)
       return centroid(to, u, sz);
  return u;
void getChildren(int u, int p, int d,
                   vector<int> &v) {
  v.push_back(d);
  for (int to : adj[u]) {
  if (rem[to] or to == p) continue;
    getChildren(to, u, d + 1, v);
il ans = 0;
int k;
int decomp(int u, int p) {
  int sz = dfs(u, p);
  int c = centroid(u, p, sz);
  if (p == -1) p = c;
  dad[c] = p;
  rem[c] = true;
  // Begin
  vector<int> f(sz + 1, 0);
  f[0] = 1;
  for (auto to : adj[c])
    if (!rem[to]) {
       vector<int> v
       getChildren(to, c, 1, v);
for (int d : v) { // Query
  if (d <= k and k - d <= sz)</pre>
           ans += f[k - d];
       for (int d : v) // Update
         f[d]++;
  //Énd
  for (auto to : adj[c]) {
  if (!rem[to]) decomp(to, c);
  return c;
void addEdge(int a, int b) {
  adj[a].push_back(b);
  adj[b].push_back(a);
// Number of k—size paths: O(N * log(N))
```

ll solve(int k1) {

assert(n > 0);

```
ans = 0, k = k1;
centroidRoot = decomp(0, -1);
  return ans;
   // namespace CD
    checking bipartiteness online
const int N = 500010;
pii parent[N];
int rk[N];
int bipartite[N];
void make_set(int v)
  parent[v] = pii(v, 0);
  rk[v] = 0;
  bipartite[v] = true;
pii find_set(int v) {
  if (v != parent[v].first) {
    int parity = parent[v].second;
    parent[v] = find_set(parent[v].first);
    parent[v].second ^= parity;
  return parent[v];
void add_edge(int a, int b) {
  int x, y;
tie(a, x) = find_set(a);
tie(b, y) = find_set(b);
                                                             9.8
  if (a == b) {
    if (x == y) bipartite[a] = false;
  } else {
    if (rk[a] < rk[b]) swap(a, b);
parent[b] = pii(a, x ^ y ^ 1);
bipartite[a] &= bipartite[b];</pre>
    if(rk[a] == rk[b]) ++rk[a];
bool is_bipartite(int v) {
  return bipartite[find_set(v).first];
9.6 chinese remainder theorem
                                                               }
#include <bits/stdc++.h>
#include "extended_euclidean.h"
using namespace std;
typedef long long ll;
namespace CRT {
inline ll normalize(ll x, ll mod) {
  x \% = mod;
  if (x < 0) \times += mod;
  return x;
ĺl solve(vector<ll> a, vector<ll> m) {
  int n = a.size();
for (int i = 0; i < n;</pre>
                           i++)
    normalize(a[i], m[i]);
  ll ans = a[0]
  ll lcm1 = m[0];
for (int i = 1; i < n; i++) {
    ans + ((((a[i] - ans) / g) * x) %
                 (m[i] / g)) *
                    lcm1
                                                                 }
         (lcm1 / g) * m[i]);
    lcm1 = (lcm1 / g) * m[i]; // lcm(lcm1, m[i]);
  return ans;
   // namespace CRT
9.7 counting inversions
```

#include <bits/stdc++.h>

using namespace std;

```
typedef long long ll;
const int INF = 0x3f3f3f3f3f;
 / Counting Inversions: O(N*log(N))
ll ci(vector<int> &v) {
  int n = v.size();
  ll inv = 0LL;
if (n == 1) return 0;
  vector<int> u1, u2;
for (int i = 0; i < n / 2; i++)
    u1.push_back(v[i]);</pre>
  for (int i = n / 2; i < n; i++)
  u2.push_back(v[i]);</pre>
  inv += ci(u1)
  inv += ci(u2);
u1.push_back(INF);
  u2.push_back(INF);
  int ini1 = 0, ini2 = 0;
for (int i = 0; i < n; i++) {
  if (u1[ini1] <= u2[ini2]) {</pre>
     v[i] = u1[ini1++];
} else {
       v[i] = u2[ini2++];
       inv += u1.size() - ini1 - 1;
  return inv;
     custom hash
#include <bits/stdc++.h>
using namespace std;
struct custom_hash {
  static uint64_t splitmix64(uint64_t x) {
     x += 0x9e3779b97f4a7c15;
    x = (x ^ (x >> 30)) * 0xbf58476d1ce4e5b9;
x = (x ^ (x >> 27)) * 0x94d049bb133111eb;
return x ^ (x >> 31);
  size_t operator()(uint64_t x) const {
     .time_since_epoch()
                .count()
     return splitmix64(x + FIXED_RANDOM);
typedef unordered_map<int, int, custom_hash> umap;
9.9 dinic
#include <bits/stdc++.h>
using namespace std; // O((V^2)*E): for generic graph.
// O(sqrt(V)*E): on unit networks. A unit network
// is a network in which all the edges have unit
// capacity, and for any vertex except s and t
// either incoming or outgoing edge is unique.
// That's exactly the case with the network we
// build to solve the maximum matching problem
 // with flows
template <typename flow_t>
struct Dinic
  struct FlowEdge {
     int from, to, id;
flow_t cap, flow = 0;
     FlowEdge(int f, int t, flow_t c, int id1)
    : from(f), to(t), cap(c) {
       id = id1;
  const flow_t flow_inf =
    numeric_limits<flow_t>::max();
  vector<FlowEdge> edges;
  vector<vector<int>> adj;
  int n, m = 0;
  int s, t;
vector<int> level, ptr;
  queue<int> q;
  bool bfs()
     while (!q.empty()) {
        int u = q.front();
```

```
q.pop();
for (int id : adj[u]) {
        if (edges[id].cap - edges[id].flow < 1)
          continue
        if (level[edges[id].to] !=-1) continue;
        level[edges[id].to] = level[u] + 1;
        q.push(edges[id].to);
    }
    return level[t] !=-1;
  flow_t dfs(int u, flow_t pushed) {
    if (pushed == 0) return 0;
if (u == t) return pushed;
    for (int &cid = ptr[u];
         cid < (int)adj[u].size(); cid++) {</pre>
      int id = adj[u][cid];
      int to = edges[id].to;
      if (level[u] + 1 != level[to]
          edges[id].cap - edges[id].flow < 1)
        continue;
      flow_t tr = dfs(
          to, min(pushed, edges[id].cap
                                edges[id].flow));
      if (tr == 0) continue;
      edges[id].flow += tr;
      edges[id ^ 1].flow —= tr;
      return tr;
    return 0;
  // Public:
 Dinic() {}
  void init(int _n) {
    adj.resize(n);
    level.resize(n);
    ptr.resize(n);
 assert(n > 0);
    edges.emplace_back(from, to, cap, id);
    edges.emplace_back(to, from, 0, -id);
    adj[from].push_back(m)
    adj[to].push_back(m + 1);
    m += 2;
  void resetFlow() {
  for (int i = 0; i < m; i++) edges[i].flow = 0;</pre>
  flow_t maxFlow(int s1, int t1) {
    s = s1, t = t1;
flow_t f = 0;
    while (true) {
      level.assign(n, -1);
      level[s] = 0;
      q.push(s);
      if (!bfs()) break;
      ptr.assign(n, 0);
while (flow_t pushed = dfs(s, flow_inf))
        f += pushed;
    return f;
// Returns the minimum cut edge IDs
vector<int> recoverCut(Dinic<int> &d) {
 vector<bool> seen(d.n, false);
  queue<int> q;
  q.push(d.s);
 seen[d.s] = true;
 while (!q.empty())
    int u = q.front();
    q.pop();
    for (int idx : d.adj[u]) {
      auto e = d.edges[idx];
      if (e.cap == e.flow) continue;
      if (!seen[e.to]) {
        q.push(e.to);
        seen[e.to] = true;
 }
```

```
vector<int> ans;
for (auto e : d.edges) {
   if (e.id >= 0) ans.push_back(e.id);
    }
  return ans;
typedef long long ll;
typedef tuple<int, int, ll> tp;
                                // (u, to, cap)
#define all(x) x.begin(), x.end()
// 0(V*E*log(MAXC))
ll maxFlowWithScaling(int n, vector<tp> edges,
                      int s, int t) {
 Dinic<ll> graph;
graph.init(n);
  sort(all(edges), [&](tp a, tp b) {
    return get<2>(a) < get<2>(b);
 auto [u, to, cap] = edges.back();
if (cap >= l) {
        graph.addEdge(u, to, cap);
        edges.pop_back();
      } else {
        break;
     }
    ans += graph.maxFlow(s, t);
  return ans;
}
9.10 dynamic median
#include <bits/stdc++.h>
using namespace std; class DinamicMedian {
  typedef int t_median;
  priority_queue<t_median> mn;
  priority_queue<t_median, vector<t_median>,
                 greater<t_median>>
      mx:
```

```
public:
  double median() {
  if (mn.size() > mx.size())
       return mn.top();
     else
       return (mn.top() + mx.top()) / 2.0;
  void push(t_{median} x) {
     if (mn.size() <= mx.size())</pre>
       mn.push(x);
     else
       mx.push(x);
     if ((!mx.empty()) and (!mn.empty())) {
  while (mn.top() > mx.top()) {
          t_{median} = mx.top();
          mx.pop();
          t_median b = mn.top();
         mn.pop();
         mx.push(b);
         mn.push(a);
    }
};
```

#### 9.11 dynamic wavelet tree

```
#include <bits/stdc++.h>
using namespace std;
struct SplayTree {
   struct Node {
    int x, y, s;
   Node *p = 0;
   Node *l = 0;
```

```
Node *r = 0;
                                                                                       p = p \rightarrow l;
  Node(int v) {
    x = v;
    y = v;
                                                                                    else
                                                                                       p = p - r, idx - lf + 1;
     s = 1;
                                                                                 splay(p);
                                                                                  return p;
   void upd() {
     s = 1;

y = x;
                                                                               int insert(int idx, int x) {
                                                                                 Node *l, *r;
tie(l, r) = split(root, idx - 1);
        f (l) {
y += l—>y;
      if (l)
                                                                                 int v = l ? l \rightarrow y : 0;
        s += l \rightarrow s;
                                                                                  root = join(l, join(new Node(x), r));
                                                                                  return v;
          (r) {
'+= r->y;
     if
                                                                               void erase(int idx) {
                                                                                 Node *l, *r;
tie(l, r) = split(root, idx);
                                                                                 root = join(l->l, r);
delete l;
   int left_size() { return l ? l->s : 0; }
Node *root = 0;
void rot(Node *c) {
                                                                               int rank(int idx) {
                                                                                 Node *l, *r;
tie(l, r) = split(root, idx);
  auto p = c \rightarrow p;
   auto g = p \rightarrow p;
                                                                                 int x = (l && l \rightarrow l ? l \rightarrow b \rightarrow y : 0);
root = join(l, r);
   if (g) (g \rightarrow l == p ? g \rightarrow l : g \rightarrow r) = c;
   if (p\rightarrow l) == c {
     p\rightarrow l = c\rightarrow r;
                                                                                  return x;
      c—>r = p;
     if (p\rightarrow l) p\rightarrow l\rightarrow p=p;
                                                                               int operator[](int idx) { return rank(idx); }
   } else {
   p->r = c->l;
                                                                               ~SplayTree() {
                                                                                 if (!root) return;
     \dot{c} = p;
                                                                                  vector<Node *> nodes{root};
                                                                                 while (nodes.size())
     if (p\rightarrow r) p\rightarrow r\rightarrow p=p;
                                                                                    auto u = nodes.back();
                                                                                    nodes.pop_back();
   c \rightarrow p = g;
                                                                                    if (u\rightarrow 1) nodes.emplace_back(u\rightarrow 1); if (u\rightarrow r) nodes.emplace_back(u\rightarrow r);
   p \rightarrow upd();
   c->upd();
                                                                                    delete u;
                                                                                 }
void splay(Node *c) {
                                                                              }
  while (c\rightarrow p) {
     auto p = c \rightarrow p;
auto g = p \rightarrow p;
                                                                            class WaveletTree {
                                                                             private:
                                                                              int lo, hi;
WaveletTree *l = 0;
        rot((g->l == p) == (p->l == c) ? p : c);
                                                                               WaveletTree *r = 0;
     rot(c);
                                                                               SplayTree b;
   c->upd();
                                                                             public:
   root = c;
                                                                               WaveletTree(int min_value, int max_value) {
                                                                                 lo = min_value;
Node *join(Node *l, Node *r) {
                                                                                 hi = max_value;
  if (not l) return r;
if (not r) return l;
while (l->r) l = l->r;
                                                                                 b.insert(0, 0);
                                                                               ~WaveletTree() {
   splay(l);
                                                                                 delete l;
   r->p = 1;
l->r = r;
                                                                                 delete r;
   l\rightarrow upd();
                                                                               // 0—indexed
                                                                              void insert(int idx, int x) {
   if (lo >= hi) return;
   int mid = (lo + hi - 1) / 2;
}
   return l;
pair<Node *, Node *> split(Node *p, int idx) {
  if (not p) return make_pair(nullptr, nullptr);
                                                                                  if (x \le mid)
                                                                                    l = l ?: new WaveletTree(lo, mid);
   if (idx < 0) return make_pair(nullptr, p);</pre>
                                                                                    l->insert(b.insert(idx, 1), x);
   if (idx >= p->s) return make_pair(p, nullptr);
for (int lf = p->left_size(); idx != lf;
                                                                                 } else {
  r = r ?: new WaveletTree(mid + 1, hi);
          lf = p \rightarrow left_size())  {
                                                                                     r->insert(idx - b.insert(idx, 0), x);
     if (idx < lf)
                                                                                 }
        p = p \rightarrow l;
     else
                                                                               // 0—indexed
        p = p \rightarrow r, idx = lf + 1;
                                                                               void erase(int idx) {
  if (lo == hi) return;
   splay(p);
                                                                                 auto p = b.get(idx);
int lf = p->l ? p->l->y : 0;
   Node *l = p;
   Node *r = p \rightarrow r;
                                                                                  int x = p \rightarrow x;
   if (r) { l->r = r->p = 0;
                                                                                 b.erase(idx);
                                                                                 if (x == 1)
    l->erase(lf);
     l \rightarrow upd();
                                                                                  else
   return make_pair(l, r);
                                                                                    r\rightarrow erase(idx - lf);
Node *get(int idx) {
                                                                               // kth smallest element in range [i, j[
                                                                               // 0—indexed
   auto p = root;
   for (int lf = p->left_size(); idx != lf;
    lf = p->left_size()) {
                                                                               int kth(int i, int j, int k) {
                                                                                 if (i >= j) return 0;
if (lo == hi) return lo;
     if (idx < lf)
```

```
int x = b.rank(i);
int y = b.rank(j);
     if (k \le y - x)
return l\rightarrow kth(x, y, k);
     else
       return r\rightarrow kth(i-x, j-y, k-(y-x));
  // Amount of numbers in the range [i, j[ Less
  // than or equal to k 0—indexed
  int lte(int i, int j, int k) {
  if (i >= j or k < lo) return 0;</pre>
     if (hi \le k) return j - i;
     int x = b.rank(i);
     int y = b.rank(j);
    return l->lte(x, y, k) +
r->lte(i - x, j - y, k);
  // Amount of numbers in the range [i, j[ equal
   // to k 0—indexed
  int count(int i, int j, int k) {
  if (i >= j or k < lo or k > hi) return 0;
     if (lo == hi) return j
     int mid = (lo + hi - 1) / 2;
     int x = b.rank(i);
     int y = b.rank(j);
     if (k <= mid) return l—>count(x, y, k);
     return r\rightarrowcount(i - x, j - y, k);
   // 0—indexed
  int get(int idx) {
  return kth(idx, idx + 1, 1);
9.12 edmond blossoms
    master/Codigo/Grafos/blossom.cpp
  Edmond's Blossoms algorithm give a maximum
```

# #include <bits/stdc++.h> using namespace std; const int MAXN = 510; // Adaptado de: // https://github.com/brunomaletta/Biblioteca/blob/ // matching in general graphs (non-bipartite) // 0(N^3) namespace EdmondBlossoms { vector<int> adj[MAXN]; int match[MAXN] int n, pai[MAXN], base[MAXN], vis[MAXN]; queue<int> q; void init(int n1) { for (int i = 0; i < n; i++) adj[i].clear(); void addEdge(int a, int b) { adj[a].push\_back(b); adj[b].push\_back(a); void contract(int u, int v, bool first = 1) { static vector<bool> bloss; static int l; if (first) { bloss = vector<bool>(n, 0); vector<bool> teve(n, 0); int k = u; $\bar{1} = v;$ while (1) { teve[k = base[k]] = 1; if (match[k] == -1) break; k = pai[match[k]]; while (!teve[l = base[l]]) l = pai[match[l]]; while (base[u] != l) { bloss[base[u]] = bloss[base[match[u]]] = 1; pai[u] = v;v = match[u]u = pai[match[u]]; if (!first) return; contract(v, u, 0); for (int i = 0; i < n; i if (bloss[base[i]]) {</pre>

```
base[i] = l;
if (!vis[i]) q.push(i);
        vis[i] = 1;
int getpath(int s) {
  for (int i = 0; i < n; i++)
    base[i] = i, pai[i] = -1, vis[i] = 0;</pre>
   vis[s] = 1;
   q = queue<int>();
   q.push(s);
   while (q.size()) {
     int u = q.front();
      q.pop();
      for (int i : adj[u]) {
        if (base[i] == base[u] or match[u] == i)
           continue;
        if (i == s \text{ or } (match[i] != -1 \text{ and }
                            pai[match[i]] != -1))
           contract(u
        else if (pai[i]'==-1) {
           pai[i] = u;
           if (match[i] == -1) return i;
i = match[i];
           vis[i] = 1;
           q.push(i);
        }
     }
   return -1;
typedef pair<int, int> pii;
vector<pii> maximumMatching() {
   vector<pii> ans;
   memset(match, -1, sizeof(match));
for (int i = 0; i < n; i++)
  if (match[i] == -1)
    for (int i = adj[i])</pre>
           if (match[j] == -1) {
              match[i] = j;
              match[j] = i;
              break;
   for (int i = 0; i < n; i++)
  if (match[i] == -1) {</pre>
        int j = getpath(i);
        if (j == -1) continue; while (j != -1) {
           int p = pai[j], pp = match[p];
           match[p] = j;
           match[j] = p;
           j = pp;
        }
     }
   for (int i = 0; i < n; i++)
  if (i < match[i])</pre>
        ans.emplace_back(i, match[i]);
   return ans;
};
     // namespace EdmondBlossoms
9.13 extended euclidean
```

```
#include <bits/stdc++.h>
using namespace std;
typedef long long ll;
ll extGcd(ll a, ll b, ll &x, ll &y) {
  if (b == 0) {
    x = 1, y = 0;
return a;
  } else {
     ll g = extGcd(b, a % b, y, x);
        = (a / b) * x;
     return g;
  }
// a*x + b*y = g
// a*(x-(b/g)*k) + b*(y+(a/g)*k) = g
bool dioEq([l a, ll b, ll c, ll &x0, ll &y0,
             ll &g)
  g = extGcd(abs(a), abs(b), x0, y0);
if (c % g) return false;
  x0 *= c / g;
```

```
y0 *= c / g;
if (a < 0) x0 = -x0;
if (b < 0) y0 = -y0;
  return true;
inline void shift(ll &x, ll &y, ll a, ll b,
                         ll cnt) {
  x += cnt * b:
  y = cnt * a;
// a1 + m1*x = a2 + m2*y
// Find the first moment that both are equal
ll findMinimum(ll a1, ll m1, ll a2, ll m2) {
    ll a = m1, b = -m2, c = a2 - a1;
  ll x, y, \dot{g}; if (!dioEq(a, b, c, x, y, g)) return -1;
  a /= g;
  b /= g;
  int sa = a > 0 ? +1 : -1;
  int sa = a > 0; +1; -1; shift(x, y, a, b, -x / b); if (x < 0) shift(x, y, a, b, sb);
  if (y < 0) {
     shift(x, y, a, b, y / a);
if (y < 0) shift(x, y, a, b, —sa);
if (x < 0) return —1;
   return a * x * g;
return (c == 0) * (maxx - minx + 1) *
                  (maxy - miny + 1);
     if (a == 0)
        return (c % b == 0) * (maxx - minx + 1) * (miny <= c / b and c / b <= maxy);
     return (c % a == 0) *
               (minx <= c / a and c / a <= maxx) * (maxy - miny + 1);
  if (!dioEq(a, b, c, x, y, g)) return 0;
  a /= g;
  b /= \bar{g};
  int sign_a = a > 0 ? +1 : -1;
int sign_b = b > 0 ? +1 : -1;
  shift(x, y, a, b, (minx - x) / b);
if (x < minx) shift(x, y, a, b, sign_b);
  if (x > maxx) return 0;
ll lx1 = x;
shift(x, y, a, b, (maxx - x) / b);
if (x > maxx) shift(x, y, a, b, -sign_b);
  ll rx1 = x;
  shift(x, y, a, b, -(miny - y) / a);
if (y < miny) shift(x, y, a, b, -sign_a);
if (y > maxy) return 0;
  \bar{l}l \hat{l}x2 = x;
  shift(x, y, a, b, -(maxy - y) / a);
if (y > maxy) shift(x, y, a, b, sign_a);
  ll rx2 = x;
       (lx2 > rx2) swap(lx2, rx2);
  ll lx = max(lx1, lx2);
  ll rx = min(rx1, rx2);
  if (lx > rx) return 0;
  return (rx - lx) / abs(b) + 1;
9.14 flow with demand
#include "dinic.h"
using namespace std;
template <typename flow_t>
struct MaxFlowEdgeDemands {
```

```
#include "dinic.h"
using namespace std;
template <typename flow_t>
struct MaxFlowEdgeDemands {
   Dinic<flow_t> mf;
   vector<flow_t> ind, outd;
   flow_t D;
   int n;
   MaxFlowEdgeDemands(int n) : n(n) {
    D = 0;
    mf.init(n + 2);
   ind.assign(n, 0);
```

```
#include <bits/stdc++.h>
using namespace std;
typedef long long f_type;
// Representation of the a/b
struct Fraction {
  f_type a, b;
Fraction(f_type _a = 0) : a(_a), b(1) {}
Fraction(f_type _a, f_type _b) {
    f_type g = __gcd(_a, _b);
a = _a / g;
b = _b / g;
    if (b < 0) {
 a = -a;
    }
  Fraction operator+(Fraction oth) {
    return Fraction(a * oth.b + oth.a * b,
                        b * oth.b);
  Fraction operator—(Fraction oth) {
     return Fraction(a * oth.b - oth.a * b,
                        b * oth.b);
  Fraction operator*(Fraction oth) {
     return Fraction(a * oth.a, b * oth.b);
  Fraction operator/(Fraction oth) {
     return Fraction(a * oth.b, b * oth.a);
  bool operator>=(Fraction oth)
     return ((*this) - oth).a >= 0;
  bool operator==(Fraction oth) {
     return a == oth.a and b == oth.b;
  operator f_type() { return a / b; }
operator double() { return double(a) / b; }
```

## 9.16 function root using newton

```
#include <bits/stdc++.h>
using namespace std;
typedef long double ld;
struct Poly {
  vector<ld> v;
  Poly(vector<ld> &v1) : v(v1) {}
  // return f(x)
  ld f(ld x) {
    ld ans = 0;
    ld e = 1;
    int n = v.size();
    for (int i = 0; i < n; i++) {
        ans += v[i] * e;
        e *= x;
    }
  return ans;
}
// return f'(x)</pre>
```

```
ld df(ld x) {
     ld ans = 0;
     ld e = 1;
     int n = v.size();
for (int i = 1; i < n; i++) {
   ans += i * v[i] * e;
   e *= x;</pre>
     return ans;
   // takes some root of the polynomial
  ld root(ld x0 = 1) {
  const ld eps = 1E-10;
     ld x = x0;
     for (;;) {
        d \ln x = x - (f(x) / df(x));
if (abs(x - nx) < eps) break;
        x = nx:
     return x;
  // div f(x) by (x-a)
  void div(ld a) {
  int g = (int)v.size() - 1;
     vector<ld> aux(g);
     for (int i = g; i >= 1; i—) {
        aux[i-1] = v[i];
        v[i-1] += a * aux[i-1];
      \dot{v} = aux;
  }
};
9.17
       gauss
#include <bits/stdc++.h>
using namespace std;
const int INF = 0x3f3f3f3f3f;
typedef long double ld;
const ld EPS = 1e-9;
int gauss(vector<vector<ld>> a, vector<ld> &ans) {
   int n = (int)a.size();
int m = (int)a[0].size() - 1;
   vector<int> where(m, -1);
   for (int col = 0, row = 0; col < m && row < n;
          col++) {
     int sel = row;
     for (int i = row; i < n; i++)
  if (abs(a[i][col]) > abs(a[sel][col]))
     if (abs(a[sel][col]) < EPS) continue;</pre>
     for (int i = col; i <= m; i++)
  swap(a[sel][i], a[row][i]);</pre>
     where[col] = row;
for (int i = 0; i < n; i++) {
        if (i != row)
           ld c = a[i][col] / a[row][col];
for (int j = col; j <= m; j++)
a[i][j] -= a[row][j] * c;
        }
     row++;
  ans.assign(m, 0);
for (int i = 0; i < m; i++)
  if (where[i] != -1)</pre>
  ans[i] = a[where[i]][m] / a[where[i]][i];
for (int i = 0; i < n; i++) {
     ld sum = 0;
     for (int j = 0; j < m; j++)
sum += ans[j] * a[i][j];
     if (abs(sum - a[i][m]) > EPS) return 0;
   for (int i = 0; i < m; i++)
     if (where[i] == -1) return INF;
   return 1;
9.18 gauss_xor
#include <bits/stdc++.h>
using namespace std;
const int MAXB = 30;
struct GaussXOR {
```

```
int table[MAXB];
  GaussXOR() {
   for (int i = 0; i < MAXB; i++) {
       table[i] = 0;
   int size() {
     int ans = 0;
for (int i = 0; i < MAXB; i++) {
  if (table[i]) ans++;
     return ans;
  bool isComb(int x) {
  for (int i = MAXB - 1; i >= 0; i—) {
    x = std::min(x, x ^ table[i]);
     return x == 0;
  void add(int x) {
  for (int i = MAXB - 1; i >= 0; i---) {
        if ((table[i] == 0)) and ((x >> i) & 1)) {
          table[i] = x;
          x = 0;
        } else {
          x = std::min(x, x ^ table[i]);
        }
     }
   int max() {
     int ans = 0;
for (int i = MAXB - 1; i >= 0; i-
        ans = std::max(ans, ans ^ table[i]);
     return ans;
};
```

#### 9.19 graph theorem

```
#include <bits/stdc++.h>
#define all(x) x.begin(), x.end()
using namespace std;
using ll = long long;
using pii = pair<int, int>;
namespace GraphTheorem {
// return if a sequence of integers d can be
// represented as the degree sequence of a finite
// simple graph on n vertices
bool ErdosGallai(vector<int> d) {
   int n = d.size();
  sort(all(d), greater<int>());
ll sum1 = 0, sum2 = 0;
  int mn = n - 1;
for (int k = 1; k \le n; k++) {
     sum1 += d[k - 1];
     while (k \le mn \text{ and } k > d[mn]) \text{ sum2} += d[mn]; if (mn + 1 < k) \text{ sum2} == d[mn++]; ll a = \text{sum1}, b = k * (ll)mn + \text{sum2};
     if (a > b) return false;
  return sum1 % 2 == 0;
vector<pii> recoverErdosGallai(vector<int> d) {
  int n = d.size();
  priority_queue<pii> pq;
for (int i = 0; i < n; i++) pq.emplace(d[i], i);</pre>
   vector<pii> edges
  while (!pq.empty()) {
     auto [g, u] = pq.top();
     pq.pop();
     vector<pii> aux(g);
for (int i = 0; i < g; i++) {
   if (pq.empty()) return {};
}</pre>
        auto [g2, u2] = pq.top();
        pq.pop();
        if (g2 == 0) return {};
        edges.emplace_back(u, u2);
aux[i] = pii(g2 - 1, u2);
     for (auto [g2, u2] : aux) pq.emplace(g2, u2);
   return edges;
```

```
do {
                                                                              used[k0] = 1;
   // namespace GraphTheorem
9.20 gray code
int grayCode(int nth) { return nth ^ (nth >> 1); }
int revGrayCode(int g) {
  int nth = 0;
                                                                                      way[j] = k0;
  for (; g > 0; g >>= 1) nth ^= g;
  return nth;
                                                                                      k1 = j;
9.21 histogram
#include <bits/stdc++.h>
using namespace std;
typedef long long ll;
  Largest Kectangular Area in a Histogram
                                                                                   v[j] = delta;
                                                                                 } else {
ll histogram(vector<int> v) {
  int n = v.size();
v.push_back(0);
  ll ans = 0;
  stack<int> st;
for (int i = 0; i <= n; i++) {
  while (st.size() && v[st.top()] >= v[i]) {
                                                                              k0 = k1:
                                                                           } while (p[k0]);
                                                                           do {
                                                                              int k1 = way[k0];
       int idx = st.top();
       st.pop();
                                                                              p[k0] = p[k1];
       int L = st.size() ? st.top() : -1;
                                                                              k0 = k1;
       ans = \max(ans, (i - L - 1) * (ll)v[idx]);
                                                                           } while (k0);
     st.push(i);
  return ans;
                                                                           ans[p[j] - 1] = j - 1;
// Largest Rectangular Area formed only by 1
                                                                         return \{ans, -v[0]\};
int maxArea1(vector<vector<bool>> mat) {
  int n = mat.size();
if (n == 0) return 0;
int m = mat[0].size();
                                                                      9.23 implicit treap
  vector<int> v(m, 0);
  int ans = 0;

for (int i = 0; i < n; i++) {

  for (int j = 0; j < m; j++) {
                                                                      #include <bits/stdc++.h>
                                                                      using namespace std;
       if (mat[i][j])
                                                                      namespace ITreap
          v[j]++;
                                                                      const int N = 500010;
       else
                                                                      typedef long long treap_t;
treap_t X[N];
          v[j] = 0;
     ans = max(ans, (int)histogram(v));
                                                                      treap_t op_val[N];
  return ans;
                                                                      bool rev[N];
                                                                         return a + b + c;
9.22 hungarian
                                                                      void calc(
#include <bits/stdc++.h>
using namespace std;
// input: matrix n x m, n <= m
// return vector p of size n, where p[i] is the // match for i
    and minimum cost
                                                                         op_val[u] =
// time complexity: 0(n^2 * m)
const int ms = 310, INF = 0x3f3f3f3f;
int u[ms], v[ms], p[ms], way[ms], minv[ms];
                                                                      void unlaze(int u) {
bool used[ms];
pair<vector<int>, int> solve(
     const vector<vector<int>> &matrix) {
                                                                         if (rev[u]) {
  int n = matrix.size();
if (n == 0) return {vector<int>(), 0};
  int m = matrix[0].size();
  assert(n <= m);</pre>
                                                                           rev[u] = false;
  memset(u, 0, (n + 1) * sizeof(int));
memset(v, 0, (m + 1) * sizeof(int));
  memset(p, 0, (m + 1) * sizeof(int));
                                                                      void split(
  memset(p, 0, (m + 1) * Sizeof(int));
for (int i = 1; i <= n; i++) {
   memset(minv, 0x3f, (m + 1) * sizeof(int));
   memset(way, 0, (m + 1) * sizeof(int));
   for (int j = 0; j <= m; j++) used[j] = 0;
   p[0] = i;
   int k0 = 0;</pre>
                                                                         unlaze(u);
```

```
int i0 = p[k0], delta = INF, k1 = 0;
for (int j = 1; j <= m; j++) {
  if (!used[j]) {</pre>
                int cur = matrix[i0 - 1][j - 1] -
                                u[i0] - v[j];
                if (cur < minv[j]) {
  minv[j] = cur;</pre>
                if (minv[j] < delta) {</pre>
                   delta = minv[j];
         for (int j = 0; j <= m; j++) {
  if (used[j]) {
                u[p[j]] += delta;
                minv[j] -= delta;
   vector<int> ans(n, -1);
for (int j = 1; j <= m; j++) {
  if (!p[j]) continue;</pre>
int en = 1, Y[N], sz[N], L[N], R[N], P[N], root;
const treap_t neutral = 0;
int u) { // update node given children info if (L[u]) P[L[u]] = u; if (R[u]) P[R[u]] = u;
   sz[u] = sz[L[u]] + 1 + sz[R[u]];
// code here, no recursion
         join(op_val[L[u]], X[u], op_val[R[u]]);
   if (!u) return;
// code here, no recursion
      if (L[u]) rev[L[u]] ^= rev[u];
if (R[u]) rev[R[u]] ^= rev[u];
swap(L[u], R[u]);
      int u, int s, int &l,
int &r) { // l gets first s, r gets remaining
   if (!u) return (void)(l = r = 0);
if (sz[L[u]] < s) {
   split(R[u], s - sz[L[u]] - 1, l, r);</pre>
```

int k0 = 0;

```
R[u] = l;
    l_= u;
  } else {
    split(L[u], s, l, r);
    L[u] = r;
r = u;
  P[u] = 0;
  calc(u);
int merge(int l, int r) { // els on l <= els on r
  unlaze(l);
  unlaze(r);
  if (!l || !r) return l + r;
  int u;
  if (Y[l] > Y[r]) {
   R[l] = merge(R[l], r);
    u = l;
  } else {
   L[r] = merge(l, L[r]);
   u = r;
  P[u] = 0; calc(u);
  return u;
int new_node(treap_t x) {
  P[en] = 0;
  X[en] = x
  op_val[en] = x;
  rev[en] = false;
  return en++;
int nth(int u, int idx) {
  if (!u) return 0;
  unlaze(u);
  if (idx \le sz[L[u]])
    return nth(L[u], idx);
  else if (idx = sz[L[u]] + 1)
    return u;
  else
    return nth(R[u], idx - sz[L[u]] - 1);
// Public
void init(
    int n = N
             1) {
                    // call before using other funcs
  // init position 0
  sz[0] = 0;
  op_val[0] = neutral;
  // init Treap
  root = 0;
  std::mt19937 rng(
       (int)std::chrono::steady_clock::now()
            .time_since_epoch()
            .count());
  for (int i = en = 1; i <= n; i++) {
    Y[i] = i;
    sz[i] = 1
    L[i] = R[i] = 0;
  shuffle(Y + 1, Y + n + 1, rng);
// 0—indexed
int insert(int idx, int val) {
  int a, b;
  split(root, idx, a, b);
int node = new_node(val);
  root = merge(merge(a, node), b);
  return node;
// 0—indexed
void erase(int idx) {
  int a, b, c, d;
  split(root, idx, a, b);
split(b, 1, c, d);
  root = merge(a, d);
// 0—indexed
treap_t nth(int idx) {
  int u = nth(root, idx + 1);
  return X[u];
^{\prime}// 0-indexed [l, r]
treap_t query(int 1, int r) {
  if (l > r) swap(l, r);
```

```
int a, b, c, d;
    split(root, l, a, d);
    split(d, r - l + 1, b, c);
    treap_t ans = op_val[b];
    root = merge(a, merge(b, c));
    return ans;
}
// 0-indexed [l, r]
void reverse(int l, int r) {
    if (l > r) swap(l, r);
    int a, b, c, d;
    split(root, l, a, d);
    split(d, r - l + 1, b, c);
    if (b) rev[b] ^= 1;
    root = merge(a, merge(b, c));
}
int getRoot(int x) {
    while (P[x]) x = P[x];
    return x;
}
int getPos(int node) {
    int ans = sz[L[node]];
    while (P[node]) {
        if (L[P[node]] == node) {
            node = P[node];
            ans += sz[L[node]] + 1;
        }
    }
    return ans;
}
; // namespace ITreap
```

#### 9.24 kadane

```
#include <bits/stdc++.h>
using namespace std;
typedef long long ll;
   Largest Sum Contiguous Subarray: O(N)
ll kadane(vector<ll> &v) {
  ll ans = 0, bigger = 0;
for (int i = 0; i < (int)v.size(); i++) {
  bigger = max(OLL, bigger + v[i]);</pre>
    ans = max(ans, bigger);
  return ans;
// Largest Sum Submatrix: 0(N^3)
il kadane2d(vector<vector<int>> &mat) {
  if (mat.size() == 0) return 0;
  int n = mat.size(), m = mat[0].size();
  ll ans = 0;
  v[k] += mat[b][k]
      ans = max(ans, kadane(v));
  return ans;
Il circularKadane(vector<ll> v) {
  ll ans1 = kadane(v);
  ll sum = 0;
for (int i = 0; i < (int)v.size(); i++) {
    sum += v[i]
    v[i] = -\bar{v}[i];
  return max(ans1, sum + kadane(v));
```

#### 9.25 karatsuba

```
#include <bits/stdc++.h>
using namespace std;
// Source:
// #pragma GCC optimize("Ofast")
// #pragma GCC target ("avx,avx2")
template <typename T>
void kar(T* a, T* b, int n, T* r, T* tmp) {
```

```
if (n <= 64) {
  for (int i = 0; i < n; i++)
    for (int j = 0; j < n; j++)
        r[i + j] += a[i] * b[j];</pre>
      return:
  int mid = n / 2;

T *atmp = tmp, *btmp = tmp + mid, *E = tmp + n;
  memset(E, 0, sizeof(E[0]) * n);
for (int i = 0; i < mid; i++) {
  atmp[i] = a[i] + a[i + mid];</pre>
      btmp[i] = b[i] + b[i + mid];
  kar(atmp, btmp, mid, E, tmp + 2 * n);

kar(a, b, mid, r, tmp + 2 * n);

kar(a + mid, b + mid, mid, r + n, tmp + 2 * n);
  for (int i = 0; i < mid; i++) {
     T \text{ temp} = r[i + mid];
      r[i + mid] += E[i] - r[i] - r[i + 2 * mid];
     r[i + 2 * mid] +=
E[i + mid] - temp - r[i + 3 * mid];
  }
^{\prime}// 0(n^{1.58}), Advantages: you can add any module
template <typename T>
vector<T> karatsuba(vector<T> a, vector<T> b) {
  int n = max(a.size(), b.size());
while (n & (n - 1)) n++;
  a.resize(n), b.resize(n);
vector<T> ret(2 * n), tmp(4 * n);
kar(&a[0], &b[0], n, &ret[0], &tmp[0]);
   return ret;
9.26 kmp
#include <bits/stdc++.h>
using namespace std;
// "abcabcd" is [0,0,0,1,2,3,0]
// "aabaaab" is [0,1,0,1,2,2,3]
vector<int> kmp(string s) {
   int n = (int)s.length();
   // pi[i] is the length of the longest proper
// prefix of the substring s[0..i] which is also
   // a suffix of this substring.
  vector<int> pi(n);
  for (int i = 1; i < n; i++) {
  int j = pi[i - 1];</pre>
      while (j > 0 \text{ and } s[i] != s[j]) j = pi[j-1];
      if(s[i] == s[j]) j++;
     pi[i] = j;
   return pi;
// The ans[i] count the amount of occurrence of
// the prefix s[0..i] in s
vector<int> prefix0ccurrences(string &s) {
   auto pi = kmp(s);
   int n = pi.size();
  vector<int> ans(n + 1);
for (int i = 0; i < n; i++) ans[pi[i]]++;
for (int i = n - 1; i > 0; i—)
    ans[pi[i - 1]] += ans[i];
   for (int i = 1; i \le n; i++)
ans[i - 1] = ans[i] + 1;
   ans.pop_back();
   return ans;
int K = 26;
inline int getID(char c) { return c - 'a'; }
vector<vector<int>> computeAutomaton(string s) {
  int n = s.size();
vector<int> pi = kmp(s);
   vector<vector<int>> aut(n, vector<int>(26));
  for (int i = 0; i < n; i++) {
  for (int c = 0; c < K; c++)
        if (i > 0 \text{ and } c != getID(s[i]))
            aut[i][c] = aut[pi[i-1]][c];
        else
            aut[i][c] = i + (c == getID(s[i]));
  }
```

```
return aut;
9.27 lagrange
#include <bits/stdc++.h>
using namespace std;
typedef long double ld;
struct PointValue {
   ld x, y
   PointValue(ld x0 = 0, ld y0 = 0)
         : x(x0), y(y0) {}
void mul(vector<ld> &A, int x\theta) { // multiply A(x) by (x - x\theta)
   int n = A.size();
   A.push_back(0);
  auto B = A;
for (int i = n; i >= 1; i—) {
  A[i] = A[i - 1];
   \dot{A}[0] = 0;
  for (int i = 0; i < n + 1; i++)
A[i] -= B[i] * x0;
void div(vector<ld> &A, int x0) \{ // multiply A(x) by (x - x0)
   int g = (int)A.size() - 1;
  vector<ld> aux(g);
for (int i = g; i >= 1; i—) {
   aux[i - 1] = A[i];
   A[i - 1] += x0 * aux[i - 1];
   A = aux;
// Change Polynomial Representation from
// Point-Value to Coefficient 0(n^2)
vector<ld> LagrangeInterpolation(
      vector<PointValue> vp) {
   vector<ld> A(1, 1);
   int n = vp.size();
   for (int i = 0; i < n; i++) mul(A, vp[i].x);
   vector<ld> ans(n, 0);
for (int i = 0; i < n; i++) {</pre>
     ld x = vp[i].x, y = vp[i].y;
      div(A, x);
      for (int j = 0; j < n; j++) {
  if (j != i) d *= (x - vp[j].x);
      for (int j = 0; j < n; j++)
ans[j] += A[j] * (y / d);
     mul(A, vp[i].x);
   return ans;
9.28 lagrange_poly
#include "modular_int.h"
namespace LagrangePoly {
const int MAXN = 100010;
modInt den[MAXN], fat[MAXN], ifat[MAXN], l[MAXN],
      r[MAXN]
void build(int n) {
   fat[0] = 1;
for (int i = 1; i <= n; i++)
  fat[i] = fat[i - 1] * i;
   ifat[n] = fat[n].inv();
for (int i = n - 1; i >
     or (int i = n - 1; i >= 0; i—)
ifat[i] = ifat[i + 1] * (i + 1);
// f(i) = y[i]
// return f(x0)
modInt getVal(vector<modInt> &y, ll x0) {
   int n = y.size();
   assert(fat[n - 1] != 0);
   modInt x = x0;
  for (int i = 0; i < n; i++) {
    den[i] = ifat[n - i - 1] * ifat[i];
    if ((n - i - 1) % 2 == 1) {
```

den[i] = -den[i];

```
}
                                                                                access(v);
                                                                                return access(w);
  for (int i = 1; i < n; i++) {
    l[i] = l[i - 1] * (x - (i - 1));</pre>
                                                                                 // namespace LCT
  r[n-1] = 1;
for (int i = n - 2; i >= 0; i—) {
r[i] = r[i + 1] * (x - (i + 1));
                                                                             9.30 lct edge
                                                                             #include <bits/stdc++.h>
  modInt ans = 0;
for (int i = 0; i < n; i++) {
  modInt li = l[i] * r[i] * den[i];
  ans = (ans + (y[i] * li));
                                                                             using namespace std;
                                                                             // Link—Cut Tree — Edge, undirected version.
// All operations are O(log(n)) amortized.
typedef long long ll;
                                                                             typedef pair<int, int> pii;
const int MAXN = 100010, MAXQ = 100010;
   return ans;
                                                                             namespace LCT {
                                                                             struct node {
  int p, ch[2];
  ll val, sub;
    // namespace LagrangePoly
                                                                                bool rév;
9.29 lct
                                                                                int sz, ár;
                                                                                ll lazy;
#include <bits/stdc++.h>
                                                                                node() {}
using namespace std;
                                                                                node(int v, int ar_)
: p(-1),
// Link-Cut Tree, directed version.
// All operations are O(log(n)) amortized. const int MAXN = 200010;
                                                                                         val(v),
                                                                                         sub(v)
namespace LCT {
                                                                                         rev(0),
struct node
                                                                                         sz(ar_{-}),
  int p, ch[2];
                                                                                         ar(ar
  node() { p = ch[0] = ch[1] = -1; }
                                                                                         ar(ar_),
lazy(0) {
                                                                                   ch[0] = ch[1] = -1;
node t[MAXN];
bool isRoot(int x) {
                                                                             };
  return t[x].p == -1 or (t[t[x].p].ch[0] != x and t[t[x].p].ch[1] != x);
                                                                             node t[MAXN + MAXQ]; // MAXN + MAXQ
                                                                             map<pii, int> edges;
                                                                             int sz;
void rotate(int x) {
                                                                             void prop(int x)
  int p = t[x].p, pp = t[p].p;
if (!isRoot(p)) t[pp].ch[t[pp].ch[1] == p] = x;
                                                                                if (t[x].lazy) {
                                                                                   if (t[x].ar) t[x].val += t[x].lazy;
  bool d = t[p].ch[0] == x;
t[p].ch[!d] = t[x].ch[d], t[x].ch[d] = p;
if (t[p].ch[!d] + 1) t[t[p].ch[!d]].p = p;
                                                                                   t[x].sub += t[x].lazy * t[x].sz;
                                                                                   if (t[x].ch[0] + 1)
  t[t[x].ch[0]].lazy += t[x].lazy;
  t[x].p = pp, t[p].p = x;
                                                                                      f (t[x].ch[1] + 1)
t[t[x].ch[1]].lazy += t[x].lazy;
void splay(int x)
  while (!isRoot(x)) {
                                                                                if (t[x].rev) {
   swap(t[x].ch[0], t[x].ch[1]);
   if (t[x].ch[0] + 1) t[t[x].ch[0]].rev ^= 1;
   if (t[x].ch[1] + 1) t[t[x].ch[1]].rev ^= 1;
     int p = t[x].p, pp = t[p].p;
if (!isRoot(p))
        rotate((t[pp].ch[0] == p) ^
                             (t[p].ch[0] == x)
                        ? x
                                                                                t[x].lazy = 0, t[x].rev = 0;
                        : p);
     rotate(x);
                                                                             void update(int x) {
  }
                                                                                t[x].sz = t[x].ar, t[x].sub = t[x].val;
for (int i = 0; i < 2; i++)
if (t[x].ch[i] + 1) {
int access(int v) {
  int last = -1;
                                                                                      prop(t[x].ch[i])
   for (int w = v; w + 1;
last = w, splay(v), w = t[v].p)
                                                                                      t[x].sz += t[t[x].ch[i]].sz;
t[x].sub += t[t[x].ch[i]].sub;
     splay(w), t[w].ch[1] = (last == -1 ? -1 : v);
  return last;
                                                                             bool is_root(int x) { return t[x].p == -1 or (t[t[x].p].ch[0] != x and
// Public:
void init(int n) {
  for (int i = 0; i <= n; i++) t[i] = node();</pre>
                                                                                                                  t[t[x].p].ch[1] != x);
                                                                             void rotate(int x) {
                                                                                int p = t[x].p, pp = t[p].p;
if (!is_root(p)) t[pp].ch[t[pp].ch[1] == p] = x;
bool d = t[p].ch[0] == x;
int findRoot(int v) {
  access(v);
  while (t[v].ch[0] + 1) v = t[v].ch[0];
   return splay(v), v;
                                                                                t[p].ch[!d] = t[x].ch[d], t[x].ch[d] = p;
                                                                                if (t[p].ch[!d] + 1) t[t[p].ch[!d]].p = p;
t[x].p = pp, t[p].p = x;
^{\prime}/ V must be root. W will be the dad of V.
void link(int v, int w) {
                                                                                update(p), update(x);
  access(v);
   t[v].p`='w;
                                                                             int splay(int x) {
  while (!is_root(x)) {
    int p = t[x].p, pp = t[p].p;
    if (!is_root(p)) prop(pp);
// Removes edge (v, dad[v])
void cut(int v) {
  access(v);
                                                                                   prop(p), prop(x);
   if (t[v].ch[0] == -1) return;
   t[v].ch[0] = t[t[v].ch[0]].p = -1;
                                                                                   if (!is_root(p))
                                                                                      rotate((t[pp].ch[0] == p) ^
int lca(int v, int w) {
  if (findRoot(v) != findRoot(w)) return -1;
                                                                                                           (t[p].ch[0] == x)
```

```
struct node {
  int p, ch[2];
                     : p);
    rotate(x);
                                                                      ll val, sub;
  return prop(x), x;
                                                                      bool rev;
                                                                      int sz;
                                                                      ll lazy;
int access(int v) {
                                                                      node() {}
  int last = -1;
  for (int w = v; w + 1;
    update(last = w), splay(v), w = t[v].p)
    splay(w), t[w].ch[1] = (last == -1 ? -1 : v);
                                                                      node(int v)
                                                                            : p(-1)
                                                                              val(v),
                                                                              sub(v),
  return last;
                                                                              rev(0),
                                                                         sz(1),
lazy(0) {
ch[0] = ch[1] = -1;
void rootify(int v);
void link_(int v, int w) {
  rootify(w);
  t[w].p = v;
                                                                    node t[MAXN];
void cut_(int v, int w) {
                                                                    void prop(int x)
  rootify(w), access(v);
                                                                      if (t[x].lazy)
  t[v].ch[0] = t[t[v].ch[0]].p = -1;
                                                                         t[x].val += t[x].lazy,
t[x].sub += t[x].lazy * t[x].sz;
void makeTree(int v, int w = 0, int ar = 0) {
                                                                         t[v] = node(w, ar);
// Public:
                                                                           t[t[x].ch[1]].lazy += t[x].lazy;
void init(int n) {
  edges.clear();
                                                                      sz = 0;
  for (int i = 0; i <= n; i++) makeTree(i);</pre>
                                                                         if (t[x].ch[1] + 1) t[t[x].ch[1]].rev ^= 1;
int findRoot(int v) {
  access(v), prop(v);
while (t[v].ch[0] + 1) v = t[v].ch[0], prop(v);
                                                                      t[x].lazy = 0, t[x].rev = 0;
  return splay(v);
                                                                    void update(int x) {
                                                                      t[x].sz = 1, t[x].sub = t[x].val;
for (int i = 0; i < 2; i++)
if (t[x].ch[i] + 1) {
// Checks if v and w are connected
bool connected(int v, int w) {
  access(v), access(w);
return v == w ? true : t[v].p != -1;
                                                                           prop(t[x].ch[i]);
                                                                           t[x].sz += t[t[x].ch[i]].sz;
t[x].sub += t[t[x].ch[i]].sub;
// Change v to be root
void rootify(int v) {
  access(v);
t[v].rev ^= 1;
                                                                    bool is_root(int x) {
                                                                      return t[x].p == -1 or (t[t[x].p].ch[0] != x and t[t[x].p].ch[1] != x);
^{\prime\prime} Sum of the edges in path from v to w
ll query(int v, int w)
                                                                    void rotate(int x) {
  rootify(w), access(v);
                                                                      int p = t[x].p, pp = t[p].p;
if (!is\_root(p)) t[pp].ch[t[pp].ch[1] == p] = x;
  return t[v].sub;
                                                                      bool d = t[p].ch[0] == x;
// Sum +x in path from v to w
                                                                      t[p].ch[!d] = t[x].ch[d], t[x].ch[d] = p;
if (t[p].ch[!d] + 1) t[t[p].ch[!d]].p = p;
void update(int v, int w, int x) {
  rootify(w), access(v);
  t[v].lazy += x;
                                                                      t[x].p = pp, t[p].p = x;
                                                                      update(p), update(x);
^{\prime}// Add edge (v, w) with weight x
void link(int v, int w, int x) {
  int id = MAXN + sz++;
                                                                    int splay(int x) {
                                                                      while (!is_root(x)) {
                                                                         int p = t[x].p, pp = t[p].p;
if (!is_root(p)) prop(pp);
  edges[pii(v, w)] = id;
makeTree(id, x, 1);
link_(v, id), link_(id, w);
                                                                         prop(p), prop(x);
if (!is_root(p))
// Remove edge (v, w)
                                                                           rotate((t[pp].ch[0] == p) ^
void cut(int v, int w) {
  int id = edges[pii(v, w)];
                                                                                              (t[p].ch[0] == x)
  cut_(v, id), cut_(id, w);
                                                                                         : p);
                                                                         rotate(x);
int lca(int v, int w) {
  access(v);
                                                                      return prop(x), x;
  return access(w);
                                                                    int access(int v) {
   // namespace LCT
                                                                      int last = -1;
                                                                      for (int w = v; w + 1;
     update(last = w), splay(v), w = t[v].p)
     splay(w), t[w].ch[1] = (last == -1 ? -1 : v);
9.31 lct_vertex
                                                                      return last;
#include <bits/stdc++.h>
                                                                    // Public:
using namespace std;
                                                                    void makeTree(int v, int w) { t[v] = node(w); }
// Link-Cut Tree - Vertex, undirected version.
                                                                    int findRoot(int v) {
// All operations are O(log(n)) amortized.
                                                                      access(v), prop(v);
typedef long long ll;
                                                                      while (t[v].ch[0] + 1) v = t[v].ch[0], prop(v);
typedef pair<int, int> pii;
const int MAXN = 200010;
                                                                      return splay(v);
                                                                    }
namespace lct {
```

```
// Checks if v and w are connected
bool connected(int v, int w) {
                                                                             using namespace std;
                                                                             typedef long long ll;
const int MOD = 104857601;
   access(v), access(w);
   return v == w? true : t[v].p != -1;
                                                                             // Work only to prime MOD
// Change v to be root
                                                                             namespace LinearSeq {
void rootify(int v) {
                                                                             const int MAXN = 233333;
ll fastPow(ll a, ll b) {
  access(v);
t[v].rev ^= 1;
                                                                                ll x =
                                                                                a %= MOD
  / \mathsf{Sum} of the weight in \mathsf{path} from \mathsf{v} to \mathsf{w}
                                                                               while (b)
                                                                                  if (b \& 1) x = (x * a) % MOD;
ll query(int v, int w) {
   rootify(w), access(v);
return t[v].sub;
                                                                                   a = (a * a) % MOD;
^{\prime\prime} Sum +x in path from v to w
                                                                                return x;
void update(int v, int w, int x) {
   rootify(w), access(v);
   t[v].lazy += x;
// Add edge (v, w)
                                                                                  ll t = 0;
void link(int v, int w) {
  rootify(w);
   t[w].p = v;
// Remove edge (v, w)
                                                                                     cur.resize(i + 1);
void cut(int v, int w) {
                                                                                      lf = i
  rootify(w), access(v);
t[v].ch[0] = t[t[v].ch[0]].p = -1;
                                                                                     id = (t - x[i]) % MOD;
                                                                                     continue;
int lca(int v, int w) {
  access(v);
   return access(w);
                                                                                  c.push_back(k);
    // namespace lct
                                                                                   if (c.size() < cur.size())</pre>
9.32 line container
                                                                                     c.resize(cur.size());
#include <bits/stdc++.h>
                                                                                  if (i - lf + (int)ls.size() >=
     (int)cur.size())
#pragma once
using ll = long long;
using namespace std;
                                                                                   cur = c;
struct Line {
  mutable ll k, m, p;
   bool operator<(const Line& o) const {</pre>
     return k < o.k;
                                                                             int m;
   bool operator<(ll x) const { return p < x; }</pre>
struct LineContainer : multiset<Line, less<>> {
  // (for doubles, use inf = 1/.0, div(a,b) = a/b
  static const ll inf = LLONG_MAX;
  ll div(ll a, ll b) { // floored division
    return a / b - ((a ^ b) < 0 && a % b);
}</pre>
                                                                                       (p[i])
                                                                                     for (int j = 0; j < m; ++j)
  t2[i + j] =</pre>
   bool isect(iterator x, iterator y) {
     if (y == end()) return x \rightarrow p = inf, 0;
                                                                                for (int
     if (x\rightarrow k == y\rightarrow k)
                                                                                   if (t2[i])
        x \rightarrow p = x \rightarrow m > y \rightarrow m ? inf : -inf;
                                                                                     for (int j = m - 1; \sim j; \longrightarrow j)
                                                                                                  j — 1] =
                                                                                        t2[i -
        x \rightarrow p = div(y \rightarrow m - x \rightarrow m, x \rightarrow k - y \rightarrow k);
     return x\rightarrow p >= y\rightarrow p;
  void add(ll k, ll m) {
     auto z = insert(\{k, m, 0\}), y = z++, x = y;
while (isect(y, z)) z = erase(z);
                                                                                s[0] = 1;
if (m != 1)
     if (x != begin() \&\& isect(\_x, y))
        isect(x, y = erase(y))
     while ((y = x) != begin() \&\& (-x)->p >= y->p)
                                                                                  t[1] = 1;
                                                                                else
t[0] = h[0];
        isect(x, erase(y));
                                                                               while (K) {
  if (K, & 1) mull(s, t);
   ll getMax(ll x) -
     assert(!empty());
                                                                                  mull(t, t);
     auto l = *lower_bound(x);
return l.k * x + l.m;
                                                                                  K >>= 1;
                                                                                ĺl su = 0;
};
                                                                               for (int i = 0; i < m; ++i)

su = (su + s[i] * a[i]) % MOD;

return (su % MOD + MOD) % MOD;
9.33 linear sequence with berlekamp massey
```

```
// Source: https://codeforces.com/blog/entry/61306
inline vector<int> BM(vector<int> x) {
   vector<int> ls, cur;
int lf = 0, ld = 0;
for (int i = 0; i < int(x.size()); i++) {</pre>
      for (int j = 0; j < int(cur.size()); ++j)
        t = (t + x[i - j - 1] * (ll)cur[j]) % MOD;
      if ((t - x[i]) % MOD == 0) continue; if (!cur.size()) {
      -(x[i]-t)*fastPow(ld, MOD-2)% MOD; vector<int> c(i-lf-1);
      for (int j = 0; j < int(ls.size()); ++j)
  c.push_back(-ls[j] * k % MOD);</pre>
      for (int j = 0; j < int(cur.size()); ++j)
  c[j] = (c[j] + cur[j]) % MOD;</pre>
         ls = cur, lf = i, ld = (t - x[i]) % MOD;
   for (int i = 0; i < int(cur.size()); ++i)
  cur[i] = (cur[i] % MOD + MOD) % MOD;
return cur;</pre>
ll a[MAXN], h[MAXN], t2[MAXN], s[MAXN], t[MAXN];
inline void mull(ll* p, ll* q) {
  for (int i = 0; i < m + m; ++i) t2[i] = 0;
  for (int i = 0; i < m; ++i)</pre>
                (t2[i + j] + p[i] * q[j]) % MOD;

i = m + m - 1; i >= m; --i)
                   (t2[i-j-1] + t2[i] * h[j]) % MOD;
   for (int i = 0; i < m; ++i) p[i] = t2[i];
inline ll calc(ll K) {
  for (int i = m; ~i; —i) s[i] = t[i] = 0;
  // init
// O(MAXN + |x|^2 * log(N))
```

#include <bits/stdc++.h>

```
inline int findElementInPositionN(vector<int> x,
                                                                               return (ans + M) % M;
                                              ll n) {
  if (n < int(x.size())) return x[n];</pre>
  vector<int> v = BM(x);
m = v.size();
                                                                             static vec ReedsSloane(const vec &s,
                                                                                                            int64 mod) {
  if (!m) return 0;
for (int i = 0; i < m; ++i)
  h[i] = v[i], a[i] = x[i];
return calc(n);</pre>
                                                                               auto inverse = [](int64 a, int64 m) {
                                                                                  int64 d, x, y;
                                                                                  exgcd(a, m, d, x, y);
return d == 1 ? (x % m + m) % m : -1;
                                                                               auto L = [](const vec &a, const vec &b) {
  int da = (a.size() > 1 ||
   // namespace LinearSeq
                                                                                                (a.size() == 1 \&\& a[0]))
                                                                                                    ? a.size() -
9.34 linear_sequence_with_reeds_sloane
                                                                                                        1000;
                                                                                  int db = (b.size() > 1 | |
#include <bits/stdc++.h>
                                                                                                (b.size() == 1 \& b[0])
using namespace std;
                                                                                                    ? b.size() - 1
// Source:
// https://github.com/zimpha/algorithmic—library/
                                                                                                        -1000;
                                                                                  return std::max(da, db + 1);
    blob/master/cpp/mathematics/linear—recurrence.
                                                                               auto prime_power = [&](const vec &s,
struct LinearRecurrence {
                                                                                                               int64 mod, int64 p,
   using int64 = long long;
                                                                                                               int64 e)
  using vec = std::vector<int64>;
                                                                                  // linear feedback shift register mod p^e, p
  static void extend(vec &a, size_t d,
                                                                                  // is prime
                            int64 value = 0) {
                                                                                  std::vector<vec> a(e), b(e), an(e), bn(e),
     if (d <= a.size()) return;</pre>
                                                                                  ao(e), bo(e);
vec t(e), u(e), r(e), to(e, 1), uo(e),
pw(e + 1, 1);
     a.resize(d, value);
  static vec BerlekampMassey(const vec &s,
                                       int64 mod) {
                                                                                  for (int i = 1; i <= e; ++i) {
  pw[i] = pw[i - 1] * p;
  assert(pw[i] <= mod);</pre>
     std::function<int64(int64)> inverse =
           [\&](int64 a) {
             return a ==
                            ? 1 : (int64)(mod - mod / a) *
                                                                                  for (int64 i = 0; i < e; ++i) {
  a[i] = {pw[i]}, an[i] = {pw[i]};
  b[i] = {0}, bn[i] = {s[0] * pw[i] % mod};
  t[i] = s[0] * pw[i] % mod;
}</pre>
                                     inverse(mod % a) % mod;
     };
vec A = {1}, B = {1};
int64 b = s[0];
assert(b != 0);
                                                                                         (t[i] == 0)
                                                                                       t[i] = 1, u[i] = e;
                                                                                     } else {
     for (size_t i = 1, m = 1; i < s.size();
        ++i, m++) {
int64 d = 0;
                                                                                       for (u[i] = 0; t[i] % p == 0;
                                                                                              t[i] /= p, ++u[i]
        for (size_t j = 0; j < A.size(); ++j) {
                                                                                    }
          d += A[j] * s[i - j] % mod;
                                                                                  for (size_t k = 1; k < s.size(); ++k) {
        if (!(d %= mod)) continue;
if (2 * (A.size() - 1) <= i) {
                                                                                    for (int g = 0; g < e; ++g) {
   if (L(an[g], bn[g]) > L(a[g], b[g])) {
      ao[g] = a[e - 1 - u[g]];
          auto temp = A;
           extend(A, B.size() + m);
          int64 coef = d * inverse(b) % mod;
for (size_t j = 0; j < B.size(); ++j) {
    A[j + m] = coef * B[j] % mod;
    if (A[i + m] - coef * A[i + m] * mod;
                                                                                          bo[g] = b[e - 1 - u[g]];
                                                                                          to[g] = t[e - 1 - u[g]];

uo[g] = u[e - 1 - u[g]];
                                                                                          r[g] = k - 1;
             if (A[j + m] < 0) A[j + m] += mod;
                                                                                       }
           \dot{B} = temp, b = d, m = 0;
                                                                                    }
                                                                                     a = an, b = bn;
        } else
                                                                                     for (int o = 0; o < e; ++o) {
           extend(A, B.size() + m);
          int64 coef = d * inverse(b) % mod;
for (size_t j = 0; j < B.size(); ++j) {
    A[j + m] == coef * B[j] % mod;
                                                                                       int64 d = 0;
for (size_t i = 0;
                                                                                              i < a[o].size() && i <= k; ++i) {
                                                                                          d = (d + a[o][i] * s[k - i]) % mod;
             if (A[j + m] < 0) A[j + m] += mod;
                                                                                       if (d == 0) {
       }
                                                                                          t[o] = 1, u[o] = e;
                                                                                       } else {
     return A:
                                                                                          for (u[o] = 0, t[o] = d;

t[o] % p == 0; t[o] /= p, ++u[o])
  static void exgcd(int64 a, int64 b, int64 &g,
                           int64 &x, int64 &y) {
                                                                                          in't g = e - 1 - u[o];
if (L(a[g], b[g]) == 0) {
  extend(bn[o], k + 1);
  bn[o][k] = (bn[o][k] + d) % mod;
} older
     if (!b)
        x = 1, y = 0, g = a;
     else {
        exgcd(b, a % b, g, y, x);
                                                                                          } else {
  int64 coef =
       y = x * (a / b);
                                                                                                  t[o] * inverse(to[g], mod) %
                                                                                                  mod * pw[u[o] - uo[g]] % mod;
  static int64 crt(const vec &c, const vec &m) {
                                                                                             int m = k - r[g];
     int n = c.size();
                                                                                             assert(m >= 0);
     int64 M = 1, ans = 0;
for (int i = 0; i < n; ++i) M *= m[i];
for (int i = 0; i < n; ++i) {
                                                                                             extend(an[o], ao[g].size() + m);
                                                                                             extend(bn[o], bo[g].size() + m);
for (size_t i = 0; i < ao[g].size();
        int64 x, y, g, tm = M / m[i];
exgcd(tm, m[i], g, x, y);
                                                                                                an[o][i + m] =
        ans = (ans + tm * x * c[i] % M) % M;
```

```
coef * ao[g][i] % mod;
                    if (an[o][i + m] < 0)
                      an[o][i + m] += mod;
                 while (an[o].size() \&\& an[o].back() == 0)
                    an[o].pop_back();
                 for (size_t i = 0; i < bo[g].size();
                   ++i) {
bn[o][i + m] -=
                         coef * bo[g][i] % mod;
                    if (bn[o][i + m] < 0)
  bn[o][i + m] -= mod;</pre>
                 while (bn[o].size() && bn[o].back() == 0)
                    bn[o].pop_back();
           }
        }
     return std::make_pair(an[0], bn[0]);
  };
   std::vector<std::tuple<int64, int64, int>>
         fac
   for (int64 i = 2; i * i <= mod; ++i)
if (mod % i == 0) {
        int64 cnt = 0, pw = 1;
while (mod % i == 0)
mod /= i, ++cnt, pw *= i;
        fac.emplace_back(pw, i, cnt);
   if (mod > 1) fac.emplace_back(mod, mod, 1);
   std::vector<vec> as;
   size_t n = 0;
for (auto &&x : fac) {
     int64 mod, p, e;
     vec a, b;
     std::tie(mod, p, e) = x;
     auto ss = s
     for (auto &\&x : ss) x %= mod;
std::tie(a, b) = prime_power(ss, mod, p, e);
     as.emplace_back(a);
     n = std::max(n, a.size());
  vec a(n), c(as.size()), m(as.size());
for (size_t i = 0; i < n; ++i) {
  for (size_t j = 0; j < as.size(); ++j) {</pre>
        m[j] = std::get<0>(fac[j]);
c[j] = i < as[j].size() ? as[j][i] : 0;
     a[i] = crt(c, m);
   return a;
LinearRecurrence(const vec &s, const vec &c,
                        int64 mod)
      : init(s)
        trans(c),
        mod(mod)
m(s.size()) {}
LinearRecurrence(const_vec &s, int64 mod,
                        bool is_prime = true)
      : mod(mod)
   assert(s.size() % 2 == 0);
   vec A;
   if (is_prime)
     A = BerlekampMassey(s, mod);
     A = ReedsSloane(s, mod);
  m = s.size() / 2;
A.resize(m + 1, 0);
trans.resize(m);
   for (int i = 0; i < m; ++i) {
trans[i] = (mod - A[i + 1]) % mod;
  if (m == 0) m = 1, trans = {1};
std::reverse(trans.begin(), trans.end());
   init = {s.begin(), s.begin() + m};
int64 calc(int64 n) {
  if (mod == 1) return 0;
if (n < m) return init[n];
vec v(m), u(m << 1);
int64 msk = !!n;</pre>
```

```
for (int64 m = n; m > 1; m >>= 1) msk <<= 1; v[0] = 1 \% mod;
  for (int64 x = 0; msk; msk >>= 1, x <<= 1) {
    std::fill_n(u.begin(), m * 2, 0);
    x = !!(n \& msk);
    if(x < m)
    u[t] = (u[t] + v[i] * v[j]) % mod;
        }
      for (int i = m * 2 - 1; i >= m; —i) {
        for (int j = 0, t = i - m; j < m; ++j, ++t) {
          u[t] = (u[t] + trans[j] * u[i]) % mod;
        }
      }
    v = \{u.begin(), u.begin() + m\};
  int64 ret = 0;
  for (int i = 0; i < m; ++i)
   or (int 1 = 0; 1 < m; ++1) {
    ret = (ret + v[i] * init[i]) % mod;
  return ret;
vec init, trans;
int64 mod;
int m:
```

#### 9.35 min cyclic string

```
#include <bits/stdc++.h>
using namespace std;
string min_cyclic_string(string s) {
    s += s;
  int n = s.size();
  int i = 0, ans = 0; while (i < n / 2) {
    ans = i;
int j = i + 1, k = i;
    while (j < n \&\& s[k] <= s[j]) {
       if (s[k] < s[j])
         k = i;
       else
         k++;
       j++;
    while (i \le k) i += j - k;
  return s.substr(ans, n / 2);
```

## **9.36** mincut

```
#include <bits/stdc++.h>
using namespace std;
typedef long long ll;
// This algorithm finds the Global Min—Cut in // O(|V|^3)
namespace MinCut {
const int MAXN = 510;
bool exist[MAXN], in_a[MAXN];
ll g[MAXN][MAXN], w[MAXN];
vector<int> v[MAXN];
int n:
void init(int n1) {
  n = n1;
  memset(g, 0, sizeof(g));
void addEdge(int a, int b, int w1) {
  if (a == b) return;
g[a][b] += w1;
g[b][a] += w1;
pair<ll, vector<int>> mincut() {
  ll best_cost = 0x3f3f3f3f3f3f3f3f3fLL;
  vector<int> best_cut;
```

```
for (int i = 0; i < n; ++i) v[i].assign(1, i);
                                                                         if (L \le u.pos and u.pos \le R) {
  memset(exist, true, sizeof(exist));
for (int ph = 0; ph < n - 1; ++ph) {
                                                                            remove(u.newV);
                                                                            add(u.oldV);
    memset(in_a, false, sizeof in_a);
memset(w, 0, sizeof w);
for (int it = 0, prev = 0; it < n - ph;</pre>
                                                                         v[u.pos] = u.oldV;
       answers.emplace_back(q.t, getAnswer());
                                                                    sort(all(answers));
                                                                    vector<int> ret
                                                                    for (auto [t, x] : answers) ret.push_back(x);
            sel = i;
                                                                    return ret;
       if (it == n - ph - 1) {
                                                                  }
         if (w[sel] < best_cost)
  best_cost = w[sel], best_cut = v[sel];
v[prev].insert(v[prev].end(),</pre>
                                                                  9.38 nearest pair of points
                            v[sel].begin(),
         v[sel].end());
for (int i = 0; i < n; ++i)
                                                                  #include <bits/stdc++.h>
                                                                  using namespace std;
            g[prev][i] = g[i][prev] += g[sel][i];
                                                                  struct pt {
          exist[sel] = false;
                                                                    long long x, y, id;
pt() {}
       } else {
         in_a[sel] = true;
for (int i = 0; i < n; ++i)
w[i] += g[sel][i];
                                                                    pt(int _x, int _y, int _id = -1)
                                                                         : x(_x), y(_y), id(_id) {}
          prev = sel;
                                                                  namespace NearestPairOfPoints {
    }
                                                                  struct cmp_x {
                                                                    bool operator()(const pt& a,
                                                                                        const pt& b) const {
  return {best_cost, best_cut};
                                                                       return a.x < b.x \mid | (a.x == b.x && a.y < b.y);
    // namespace MinCut
                                                                    }
                                                                  };
                                                                  struct cmp_y {
                                                                    bool operator()(const pt& a,
9.37 mo_with_update
                                                                                        const pt& b) const {
                                                                       return a.y < b.y;
#include <bits/stdc++.h>
#define all(x) x.begin(), x.end()
                                                                    }
using namespace std;
                                                                  int n;
using pii = pair<int, int>;

const int INF = 0x3f3f3f3f;

const int BLOCK_SIZE = 2800; // (2*N^2)^(1/3)

const int MAXN = 100010;
                                                                  vector<pt> v;
                                                                  vector<pt> t;
                                                                  double mindist;
pair<int, int> best_pair;
void upd_ans(const pt& a, const pt& b) {
int v[MAXN];
void remove(int x);
                                                                    double dist = sqrt((a.x - b.x) * (a.x - b.x) +
void add(int x);
                                                                                            (a.y - b.y) * (a.y - b.y);
void clearAnswer();
                                                                    if (dist < mindist)</pre>
int getAnswer();
                                                                       mindist = dist;
struct Query {
                                                                       best_pair = {a.id, b.id};
  int l, r, t;
                                                                    }
  bool operator<(const Query &oth) const {
    if (l / BLOCK_SIZE != oth.l / BLOCK_SIZE)
                                                                  void rec(int l, int r) {
    return l < oth.l;
if (r / BLOCK_SIZE != oth.r / BLOCK_SIZE)
                                                                    if (r - l <= 3) {
  for (int i = l; i < r; ++i) {
    for (int j = i + 1; j < r; ++j) {
      upd_ans(v[i], v[j]);
}</pre>
       return r < oth.r;
    return t < oth.t;
  }
                                                                         }
struct Update {
                                                                       sort(v.begin() + l, v.begin() + r, cmp_y());
  int pos, newV, oldV, t;
                                                                       return;
// 0(0 * N^{(2/3)}): N=10^5 \rightarrow 1.5s
                                                                    int m = (l + r) >> 1;
vector<int> mo_s_algorithm(vector<Query> vq,
                                                                    int midx = v[m].x;
                                 vector<Update> vu) {
                                                                    rec(l, m);
  vector<pii> answers;
                                                                    rec(m, r)
  sort(all(vq));
                                                                    merge(v.begin() + l, v.begin() + m,
  clearAnswer();
                                                                            v.begin() + m, v.begin() + r, t.begin(),
  int L = 0, R = 0, T = 0, szT = vu.size();
                                                                            cmp_y());
  add(\overline{v}[0]);
                                                                    copy(t.begin(), t.begin() + r - l,
  for (Query q : vq) {
  while (q.l < L) add(v[—L]);</pre>
                                                                          v.begin() + l);
                                                                    int tsz = 0;
    while (R < q.r) add(v[++R]);
                                                                    for (int i = l; i < r; ++i) {
    while (L < q.l) remove(v[L++]); while (q.r < R) remove(v[R-]);
                                                                       if (abs(v[i].x - midx) < mindist) {</pre>
                                                                         while (T < szT \text{ and } vu[T].t <= q.t) {
       Update \&u = vu[T++];
                                                                                  -i)
       if (L \leq u.pos and u.pos \leq R) {
                                                                            upd_ans(v[i], t[j]);
         remove(u.oldV);
                                                                         t[tsz++] = v[i];
         add(u.newV);
                                                                    }
       v[u.pos] = u.newV;
                                                                  }
                                                                  pair<int, int> solve(vector<pt> _v) {
    while (T > 0 \text{ and } vu[T - 1].t > q.t) {
       Update \&u = vu[-T];
                                                                    n = v.size();
```

```
t.resize(n);
sort(v.begin(), v.end(), cmp_x());
                                                                            class RangeColor {
                                                                              private:
  mindist = 1E20;
                                                                                typedef long long ll;
  rec(0, n);
return best_pair;
                                                                                struct Node {
                                                                                  ll l, r;
                                                                                  int color;
                                                                                  Node() {}
Node(ll l1, ll r1, int color1)
    : l(l1), r(r1), color(color1) {}
bool operator<(const Node &oth) const {</pre>
    // namespace NearestPairOfPoints
9.39 parallel binary search
                                                                                     return r < oth.r;
                                                                                  }
#include <bits/stdc++.h>
using namespace std;
const int MAXN = 100010;
                                                                                std::set<Node> st;
                                                                               vector<ll> ans;
int ans[MAXN];
                                                                              public:
bool test(int x);
void add(int k);
                                                                                RangeColor(ll first, ll last, int maxColor) {
                                                                                  ans.resize(maxColor + 1);
ans[0] = last - first + 1LL;
void remove(int k);
void solve(int i, int j, vector<int> &v) {
                                                                                  st.insert(Node(first, last, 0));
  if (v.empty()) return;
if (i == j) {
  for (int x : v) ans[x] = i;
                                                                                // get color in position x
                                                                               int get(ll x) {
     return;
                                                                                  auto p = st.upper_bound(Node(0, x - 1LL, -1));
  int mid = (i + j) / 2;
for (int k = i; k <= mid; k++) add(k);
vector<int> left, right;
                                                                                  return p—>color;
                                                                               // set newColor in [a, b]
void set(ll a, ll b, int newColor) {
  auto p = st.upper_bound(Node(0, a - 1LL, -1));
  for (int x : v) {
     if (test(x))
                                                                                  assert(p != st.end());
        left.push_back(x);
                                                                                  ll l = p->l;
ll r = p->r;
     else
        right.push_back(x);
                                                                                  int oldColor = p->color;
ans[oldColor] -= (r - l + 1LL);
  solve(mid + 1, j, right);
for (int k = mid; k >= i; k—)
  remove(k); // Or roolback();
solve(i, mid, left);
                                                                                  p = st.erase(p);
                                                                                  if (l < a) {
                                                                                     ans[oldColor] += (a - l);
st.insert(Node(l, a - 1LL, oldColor));
                                                                                  if (b < r) {
    ans[oldColor] += (r - b);
9.40 permutation
                                                                                     st.insert(Node(b + 1LL, r, oldColor));
#include <bits/stdc++.h>
                                                                                  while ((p != st.end()) and (p->l <= b)) {
using namespace std;
                                                                                     l = p \rightarrow l;

r = p \rightarrow r;
using ll = long long;
mt19937_64 rng(
                                                                                     oldColor = p->color;
     (int)std::chrono::steady_clock::now()
                                                                                     ans[oldColor] = (r - l + 1LL);
           .time_since_epoch()
                                                                                        ans[oldColor] += (r - b);
           .count())
                                                                                        st.erase(p);
namespace Permutation
const int MAXN = 500010;
                                                                                        st.insert(Node(b + 1LL, r, oldColor));
ll mp[MAXN], sumXor[MAXN], p[MAXN + 1], inv[MAXN];
                                                                                        break;
void init(vector<int> v) {
  sumXor[0] = inv[0] = p[0] = 0;
                                                                                     } else
                                                                                        p = st.erase(p);
  for (int i = 0; i < MAXN; i++) {
    mp[i] = rng() + 1;
    p[i + 1] = p[i] ^ mp[i];
                                                                                  ans[newColor] += (b - a + 1LL);
st.insert(Node(a, b, newColor));
     or (int i = 0; i < v.size(); i++)
if (v[i] < 0 or v[i] >= MAXN) {
  inv[i + 1] = 1 + inv[i];
  sumXor[i + 1] = sumXor[i];
                                             i++) {
                                                                                ll countColor(int x) { return ans[x]; }
                                                                            };
     } else {
        inv[i + 1] = inv[i];

sumXor[i + 1] = sumXor[i] ^ mp[v[i]];
                                                                            9.42 rank matrix
     }
                                                                            #include <bits/stdc++.h>
  }
                                                                            using namespace std;
typedef long double ld;
// Verify if {v[l], v[l+1], ..., v[r]} is {0, 1, // ... , r-l+1} 0-indexed;
                                                                             const ld EPS = 1e-9;
                                                                             int compute_rank(vector<vector<ld>>> A) {
bool isPermutation(int l, int r) {
                                                                               int n = A.size();
int m = A[0].size();
  l++, r++;
if (inv[r] - inv[l-1] > 0) return false;
return p[r-l+1] ==
                                                                               int rank = max(n, m);
vector<bool> row_selected(n, false);
            (sumXor[r] ^ sumXor[l - 1]);
                                                                                for (int i = 0; i < m; ++i) {
                                                                                  int j;
}; // namespace Permutation
                                                                                  for (j = 0; j < n; ++j) {
   if (!row\_selected[j] && abs(A[j][i]) > EPS)
                                                                                        break;
9.41 range_color
                                                                                  if (j == n) {
#include <bits/stdc++.h>
                                                                                     rank-
                                                                                  rank—;
} else {
```

using namespace std;

```
row_selected[j] = true;
      for (int p = i + 1; p < m; p++)
        A[j][p] /= A[j][i];
      for (int k = 0; k < n; k++)
           (k != j \&\& abs(A[k][i]) > EPS) {
for (int p = i + 1; p < m; p++)
             A[k][p] = A[j][p] * A[k][i];
      }
    }
  return rank;
9.43 segment_tree2d
#include <bits/stdc++.h>
using namespace std;
struct SegTree2D {
 private:
  int n, m;
  typedéf int Node;
  Node neutral = -0x3f3f3f3f;
  vector<vector<Node>> seg;
  Node join(Node a, Node b) { return max(a, b); }
 public:
  SegTree2D(int n1, int m1) {
    n = n1, m = m1;
    seq.assign(2 * n, vector<Node>(2 * m, 0));
  void update(int x, int y, int val) {
    assert(0 \le x \& x < n \& 0 \le y \& y < m);
    x += n, y += m;

seg[x][y] = val;
    for (int j = y / 2; j > 0; j /= 2)
      seg[x][j] =
           join(seg[x][2 * j], seg[x][2 * j + 1]);
    for (x /= 2; x > 0; x /= 2) {
    seg[x][y] =
      }
    }
  vector<int> getCover(int l, int r, int N) {
    l = std::max(0, l);
    r = std::min(N, r);
    vector<int> ans;
    for (l += N, r += N; l < r; l /= 2, r /= 2) {
  if (l & 1) ans.push_back(l++);
  if (r & 1) ans.push_back(—r);</pre>
                                                            };
    return ans;
  Node query(int x1, int y1, int x2, int y2) {
    auto c1 = getCover(x1, x2 + 1, n);
    auto c2 = getCover(y1, y2 + 1, m);
    Node ans = neutral;
    for (auto i : c1) {
   for (auto j : c2) {
        ans = join(ans, seg[i][j]);
    return ans;
9.44 simpson_integration
#include <bits/stdc++.h>
using namespace std;
double f(double x)
const int N = 1000000;
double_simpson_integration(double a, double b) {
```

# double h = (b - a) / N; double s = f(a) + f(b); $// a = x_0 and b = x_2 n$ for (int i = 1; i <= N - 1; ++i) { // Refer to final Simpson's formuladouble x = a + h \* i;

```
s += f(x) * ((i \& 1) ? 4 : 2);
  s *= h / 3;
  return s;
9.45 sqrt decomposition
#include <bits/stdc++.h>
using namespace std;
struct SqrtDecomposition {
  typedef long long t_sqrt;
  int sqrtLen;
  vector<t_sqrt> block;
  vector<t_sqrt> v;
  template <class MyIterator>
  SqrtDecomposition(MyIterator begin,
                     MyIterator end) {
    int n = end - begin;
    sqrtLen = (int)sqrt(n + .0) + 1;
    v.resize(n);
    block.resize(sqrtLen + 5);
    for (int i = 0; i < n; i++, begin++) {
      v[i] = (*begin);
      block[i / sqrtLen] += v[i];
  // 0—indexed
  void update(int idx, t_sqrt new_value) {
    t_sqrt d = new_value - v[idx];
    v[idx] += d;
    block[idx / sqrtLen] += d;
  // 0-indexed [l, r]
t_sqrt query(int l, int r) {
    t_sqrt sum = 0;
    int c_l = l / sqrtLen, c_r = r / sqrtLen;
    if (c_l == c_r) {
  for (int i = l; i <= r; i++) sum += v[i];</pre>
     else {
      for (int i = l,
end = (c_l + 1) * sqrtLen - 1;
            i <= end; i++)
      sum += v[i];
for (int i = c_1 + 1; i <= c_r - 1; i++)
        sum += block[i];
      for (int i = c_r * sqrtLen; i \le r; i++)
        sum += v[i];
    return sum;
9.46 system_of_linear_equations
#include <algorithm>
#include <iostream>
constexpr int N = 510, p = 998244353;
int fp(int a, int b) {
  int ans = 1, off = a;
  while (b) {
  if (b & 1) ans = 1ll * ans * off % p;
    off = 1ll * off * off % p;
    b >>= 1:
  return ans;
int gauss(int (&dat)[N][N], int (&ans)[N],
           int (&basis)[N][N], int n, int m) {
  int k = 1;
  static int col[N];
  for (int i = 1; i \le m \&\& k \le n; ++i) {
    int pos = 0;
    for (int j = k; j \le n; ++j)
      if (dat[j][i]) {
        pos = j;
        break;
```

if (!pos) continue;

```
col[k] = i;
    if (pos != k) {
      for (int j = i; j \le m + 1; ++j)
         std::swap(dat[pos][j], dat[k][j]);
    int rv = fp(dat[k][i], p - 2);
    for (int j = i; j \le m + 1; ++j)
      dat[k][j] = 1ll * dat[k][j] * rv % p;
    for (int j = k + 1; j \le n; ++j)
      if (dat[j][i]) {
         int num = p - dat[j][i];
         for (int t = i; t \le m + 1; ++t)
           dat[j][t] = (dat[j][t] +
                          1ll * num * dat[k][t]) %
      }
    ++k;
  }
  for (int i = k; i \le n; ++i)
    if (dat[i][m + 1]) return -1;
  ---k:
  int R = m - k;
  for (int i = 1; i \le m; ++i) ans[i] = 0;
  for (int i = k; i; —i) {
  ans[col[i]] = dat[i][m + 1];
    for (int j = i + 1; j \le k;
      ans[col[j]]) %
                       p;
  for (int i = 1; i \le R; ++i)
for (int j = 1; j \le m; ++j) basis[i][j] = 0;
  col[k + 1] = m + 1;
  col[0] = 0;
  for (int i = 0, t = 0; i \le k; ++i) {
    for (int j = col[i] + 1; j < col[i + 1];
          ++j) {
      ++t;
      for (int l = i; l; —l) {
  int c = dat[l][j];
         for (int r = l + 1; r <= i; ++r)
    c = (c + 1ll * dat[l][col[r]]</pre>
                          basis[t][col[r]]) %
         basis[t][col[l]] = (p - c) % p;
      basis[t][j] = 1;
  }
  return R;
int main() {
  int n, m;
  static int dat[N][N], ans[N], basis[N][N];
  std::ios::sync_with_stdio(false);
  std::cin.tie(0);
  std::cin >> n >> m;
  for (int i = 1; i <= n; ++i)
for (int j = 1; j <= m; ++j)
      std::cin >> dat[i][i];
  for (int j = 1; j <= n; ++j)
    std::cin >> dat[j][m + 1];
  int R = gauss(dat, ans, basis, n, m);
  if (!(~R)) {
    std::cout << "-1\n";
    return 0;
  std::cout << R << '\n';
  for (int j = 1; j <= m; ++j)
  std::cout << ans[j] << ' '</pre>
  std::cout << '\n';
  for (int i = 1; i \le R; ++i)
    for (int j = 1; j <= m; ++j)
```

```
std::cout << basis[i][j] << ' ';
std::cout << '\n';
  return 0;
}
9.47 treap
#include <bits/stdc++.h>
using namespace std;
namespace Treap -
const int N = 500010;
typedef long long treap_t;
treap_t X[N];
int en = 1, Y[N], sz[N], L[N], R[N], root;
const treap_t neutral = 0;
treap_t op_val[N];
inline treap_t join(treap_t a, treap_t b,
                       treap_t c) {
  return a + b + c;
void calc(
    int u) {
                // update node given children info
  sz[u] = sz[L[u]] + 1 + sz[R[u]];
  // code here, no recursion
  op_val[u] =
      join(op_val[L[u]], X[u], op_val[R[u]]);
void unlaze(int u) {
  if (!u) return;
  // code here, no recursion
void split(int u, treap_t x, int &l,
             int &r) \{ // l \text{ gets} \le x, r \text{ gets} > x \}
  unlaze(u):
  if (!u) return (void)(l = r = 0);
if (X[u] <= x) {</pre>
     split(R[u], x, l, r);
    R[u] = l;
  } else -
    split(L[u], x, l, r);
    L[u] = r;
r = u;
  calc(u);
void split_sz(
    int u, int s, int &l,
int &r) { // l gets first s, r gets remaining
  unlaze(u);
  if (!u) return (void)(l = r = 0);
if (sz[L[u]] < s) {</pre>
    split_sz(R[u], s - sz[L[u]] - 1, l, r);
    R[u] = l;
  } else {
    split_sz(L[u], s, l, r);
    L[u] = r;
r = u;
  calc(u);
int merge(int l, int r) { // els on l <= els on r</pre>
  unlaze(l);
  unlaze(r);
  if (!l || !r) return l + r;
  int`u;
if (Y[l] > Y[r])
    R[l] = merge(R[l], r);
    u = 1:
  } else {
   L[r] = merge(l, L[r]);
  calc(u);
```

return u;

X[en] = x;

op\_val[en] = x; return en++;

int new\_node(treap\_t x) {

int nth(int u, int idx) {

```
if (!u) return 0;
                                                             void unite(int a, int b) {
  unlaze(u);
  if (idx \le sz[L[u]])
                                                               a = get(a);
  return nth(L[u], idx);
else if (idx == sz[L[u]] + 1)
                                                               b = qet(b);
                                                               if (sz[a] < sz[b]) {
    return u;
                                                                 swap(a, b);
    return nth(R[u], idx - sz[L[u]] - 1);
                                                               // add to history
// Public
                                                               past_parent.push_back({b, p[b]});
void init(
    int n = N
                                                               past_size.push_back({a, sz[a]});
             1) { // call before using other funcs
                                                               if (a != b) {
  // init position 0
                                                                 p[b] = a;
  sz[0] = 0;
  op_val[0] = neutral;
                                                                 sz[a] += sz[b];
  // init Treap
                                                               }
                                                             }
  root = 0
  std::mt19937 rng(
                                                             bool sameset(int a, int b) {
      (int)std::chrono::steady_clock::now()
                                                               return get(a) == get(b);
          .time_since_epoch()
  .count()); for (int i = en = 1; i <= n; i++) {
                                                             // Reverts to previous DSU state.
    Y[i] = i;
                                                             void rollback() {
  p[past_parent.back().first] =
    sz[i] = i;
    L[i] = R[i] = 0;
                                                                   past_parent.back().second;
                                                               sz[past_size.back().first] =
  shuffle(Y + 1, Y + n + 1, rng);
                                                                   past_size.back().second;
                                                               past_parent.pop_back();
void insert(treap_t x) {
  int a, b;
                                                               past_size.pop_back();
  split(root, x, a, b);
  root = merge(merge(a, new_node(x)), b);
                                                          };
                                                          struct Query {
void erase(treap_t x) {
                                                             int t, k, u, v;
 int a, b, c, d;
split(root, x - 1, a, b);
  split(b, x, c, d);
split_sz(c, 1, b, c);
                                                          int main()
                                                             cin.tie(0)—>sync_with_stdio(false);
  root = merge(a, merge(c, d));
                                                             int n, q;
                                                             cin >> n >> q;
int count(treap_t x) {
 int a, b, c, d;
split(root, x - 1, a, b);
                                                             vector<Query> queries(q + 1);
                                                             vector<vector<int>> queries_using_k(q + 1);
  split(b, x, c,
                 d);
                                                             for (int i = 1; i \le q; i++) {
  int ans = sz[c];
                                                               Query tmp;
  root = merge(a, merge(c, d));
  return ans;
                                                               cin >> tmp.t >> tmp.k >> tmp.u >> tmp.v;
int size() { return sz[root]; }
                                                               // construct query graph — add an edge between
// 0—indexed
treap_t nth(int idx)
                                                               // k and itself
  int u = nth(root, idx + 1);
                                                               queries_using_k[tmp.k].push_back(i);
  return X[u];
                                                               queries[i] = tmp;
// Query in k smallest elements
treap_t query(int k) {
                                                             vector<bool> ans(q + 1);
 int a, b;
split_sz(root, k, a, b);
treap_t ans = op_val[a];
                                                             DSU dsu(n);
                                                             // process queries in a dfs manner
  root = merge(a, b);
                                                             auto dfs = [&](auto self, int idx) -> void {
  return ans;
                                                               if (queries[idx].t == 0) {
                                                                 dsu.unite(queries[idx].u, queries[idx].v);
   // namespace Treap
                                                               } else {
                                                                 if (dsu.sameset(queries[idx].u,
                                                                                   queries[idx].v)) {
9.48 union find persistent
                                                                   ans[idx] = true;
                                                                 }
#include <bits/stdc++.h>
using namespace std;
                                                               // answer all the queries that use this dsu
struct DSU {
                                                               for (int i : queries_using_k[idx]) {
  vector<int> p, sz;
                                                                 self(self, i);
  // stores info from the past
                                                               // roll back merges
  vector<pair<int, int>> past_parent, past_size;
                                                               if (queries[idx].t == 0) {
  DSU(int n)
                                                                 dsu.rollback();
    p.resize(n);
    sz.resize(n, 1);
    iota(p.begin(), p.end(), 0);
                                                             dfs(dfs, 0);
```

for (int i = 1; i <= q; i++) {

cout << ans[i] << "\n";</pre>

if (queries[i].t) {

}

int get(int x) {

return (p[x] == x) ? x : get(p[x]);

```
else
v[i2++] = aux[k];
    }
                                                                                                                                             return i1:
                                                                                                                                        9.49 union find with rollback
                                                                                                                                             lo[u] = minX, hi[u] = maxX;
if (lo[u] == hi[u] or i >= j) return;
t_wavelet mid = (minX + maxX - 1) / 2;
a[u].resize(j - i + 1);
#include <bits/stdc++.h>
using namespace std;
stručt RollbackUF {
     vector<int> e;
                                                                                                                                              a[u][0] = 0;
    vector<tuple<int, int, int, int>> st;
RollbackUF(int n) : e(n, -1) {}
int size(int x) { return -e[find(x)]; }
int find(int x) {
                                                                                                                                             for (int k = i; k < j; k++)
a[u][k-i+1] = a[u][k-i] + (v[k] <= mid);
                                                                                                                                              int pivot = stable_partition(i, j, mid);
                                                                                                                                             l[u] = last++, r[u] = last++;
build(l[u], i, pivot, minX, mid);
build(r[u], pivot, j, mid + 1, maxX);
          return e[x] < 0 ? x : find(e[x]);
     int time() { return st.size(); }
    void rollback(int t) {
  while (st.size() > t) {
    auto [a1, v1, a2, v2] = st.back();
    e[a1] = v1;
    e[a2] = v2;
    true the content of the con
                                                                                                                                        inline int b(int u, int i) { return i - a[u][i]; }
                                                                                                                                         // Public
                                                                                                                                         template <class MyIterator>
                                                                                                                                        void init(MyIterator begin, MyIterator end,
                                                                                                                                                                  t_wavelet minX, t_wavelet maxX) {
               st.pop_back();
                                                                                                                                              last = 1;
         }
                                                                                                                                              int n = end - begin;
                                                                                                                                              for (int i = 0; i < n; i++, begin++)
    bool unite(int a, int b) {
    a = find(a), b = find(b);
    if (a == b) return false;
    if (e[a] > e[b]) swap(a, b);
    return false;
                                                                                                                                                  v[i] = *begin;
                                                                                                                                             build(last++, 0, n, minX, maxX);
                                                                                                                                        // kth smallest element in range [i, j]
          st.push_back({a, e[a], b, e[b]});
          e[a] += e[b];
e[b] = a;
                                                                                                                                         // 1—indexed
                                                                                                                                        int kth(int i, int j, int k, int u = 1) {
                                                                                                                                             if (i > j) return 0;
if (lo[u] == hi[u]) return lo[u];
int inLeft = a[u][j] - a[u][i - 1];
          return true;
                                                                                                                                             int i1 = a[u][i - 1] + 1, j1 = a[u][j];
int i2 = b(u, i - 1) + 1, j2 = b(u, j);
if (k <= inLeft) return kth(i1, j1, k, l[u]);
return kth(i2, j2, k - inLeft, r[u]);</pre>
9.50 vertex cover in tree
#include <bits/stdc++.h>
                                                                                                                                         // Amount of numbers in the range [i, j] Less than
using namespace std;
                                                                                                                                        // or equal to k 1—indexed
const int MAXN = 200010;
                                                                                                                                        int lte(int i, int j, int k, int u = 1) {
  if (i > j or k < lo[u]) return 0;
  if (hi[u] <= k) return j - i + 1;
  int i1 = a[u][i - 1] + 1, j1 = a[u][j];
  int i2 = h(u i - 1) + 1, i2 = h(u, i);</pre>
int dp[MAXN][2]
vector<int> adj[MAXN];
// vertexCover(node current, free to choose, dad) int vertexCover(int u, bool color = true, int p = -1) {
                                                                                                                                             int i2 = b(u, i - 1) + 1, j2 = b(u, j);
return lte(i1, j1, k, l[u]) + lte(i2, j2, k, r[u]);
     if (dp[u][color] != -1) return dp[u][color];
     int case1 = 1, case2 = 0;
for (int to : adj[u]) {
                                                                                                                                        }
// Amount of numbers in the range [i, j] equal to
          if (to == p) continue;
                                                                                                                                         // k 1—indexed
          case1 += vertexCover(to, true, u);
case2 += vertexCover(to, false, u);
                                                                                                                                        // K 1—Indexed
int count(int i, int j, int k, int u = 1) {
   if (i > j or k < lo[u] or k > hi[u]) return 0;
   if (lo[u] == hi[u]) return j - i + 1;
   t_wavelet mid = (lo[u] + hi[u] - 1) / 2;
   int i1 = a[u][i - 1] + 1, j1 = a[u][j];
   int i2 = b(u, i - 1) + 1, j2 = b(u, j);
   if (k <= mid) return count(i1 i1 k = 1[u]);
}</pre>
     if (color)
          return dp[u][color] = min(case1, case2);
          return dp[u][color] = case1;
                                                                                                                                              if (k <= mid) return count(i1, j1, k, l[u]);</pre>
                                                                                                                                              return count(i2, j2, k, r[u]);
9.51 wavelet_tree
                                                                                                                                        // swap v[i] with v[i+1]
// 1—indexed
                                                                                                                                       // 1—indexed
void swp(int i, int u = 1) {
   if (lo[u] == hi[u] or a[u].size() <= 2) return;
   if (a[u][i - 1] + 1 == a[u][i] and
       a[u][i] + 1 == a[u][i + 1])
       swp(a[u][i], l[u]);
   else if (b(u, i - 1) + 1 == b(u, i) and
            b(u, i) + 1 == b(u, i + 1))
       swp(b(u, i), r[u]);
   else if (a[u][i - 1] + 1 == a[u][i])</pre>
#include <bits/stdc++.h>
using namespace std;
namespace WaveletTree
const int MAXN = 100010,
                         MAXW =
                                   MAXN * 30; // MAXN * LOG(maxX-MinX)
typedef int t_wavelet;
int last
int tast,
int v[MAXN], aux[MAXN];
int lo[MAXW], hi[MAXW], l[MAXW], r[MAXW];
vector<t_wavelet> a[MAXW];
                                                                                                                                              else if (a[u][i-1] + 1 == a[u][i])
                                                                                                                                                  a[u][i]-
                                                                                                                                             elsē
                                                                                                                                                  a[u][i]++;
int stable_partition(int i, int j
                                                     t_wavelet mid) {
                                                                                                                                                 // namespace WaveletTree
     int pivot = 0;
for (int k = i; k < j; k++)
   aux[k] = v[k], pivot += (v[k] <= mid);</pre>
     int i1 = i, i2 = i + pivot;
     for (int k = i; k < j; k++) {
                                                                                                                                        9.52 xor_and_or_convolution
          if (aux[k] \le mid)
```

v[i1++] = aux[k];

```
#include <bits/stdc++.h>
using namespace std;
typedef long long ll;
void xorFWHT(vector<ll> &P,
                     bool inverse = false) {
   int n = P.size();
for (int len = 1; 2 * len <= n; len <<= 1) {
   for (int i = 0; i < n; i += 2 * len) {
      for (int j = 0; j < len; j++) {
        ll u = P[i + j];
        ll v = P[i + len + j];
        ll v = P[i + len + j];</pre>
            P[i + j] = u + v;

P[i + len + j] = u - v;
      }
   if (inverse) {
  for (int i = 0; i < n; i++) {</pre>
         P[i] /= n;
void orFWHT(vector<ll> &P, bool inverse = false) {
   int n = P.size();
   for (int len = 1; 2 * len <= n; len <<= 1) {
    for (int i = 0; i < n; i += 2 * len) {
        for (int j = 0; j < len; j++) {
             if (inverse)
                P[i + len' + j] = P[i + j];
             else
                P[i + len + j] += P[i + j];
     }
  }
void andFWHT(vector<ll> &P,
   bool inverse = false) {
int n = P.size();
   for (int len = 1; 2 * len <= n; len <<= 1) {
  for (int i = 0; i < n; i += 2 * len) {
    for (int j = 0; j < len; j++) {
        ll u = P[i + j];
    }
}</pre>
             ll v = P[i + \overline{len} + j];
             if (inverse) {
               P[i + j] = v - u;

P[i + len + j] = u;
             } else {
               P[i + j] = v;

P[i + len + j] = u + v;
        }
     }
  }
vector<ll> convolution(vector<ll> a
                                     vector<ll> b) {
   int mx = max(a.size(), b.size());
   int n = 1;
   while (n < mx) n <<= 1;
   a.resize(n, 0);
b.resize(n, 0);
   xorFWHT(a)
   xorFWHT(b);
   for (int i = 0; i < n; i++) a[i] *= b[i];
xorFWHT(a, true);</pre>
   return a;
9.53 xor_trie
#include <bits/stdc++.h>
using namespace std;
using ll = long long;
struct Vertex {
   int next[2];
   int leaf
   int count;
   Vertex() { next[0] = next[1] = -1;
      leaf = count = 0;
  }
                                                                                        // between s[0..(n-1)] and the suffix of // s[i..(n-1)]. z[0] is generally not well
const int MAXB = 20;
struct Trie {
```

```
vector<Vertex> trie;
ll lazy;
  Trie() {
     trie.emplace_back();
     lazy = 0;
   void add(ll x) {
     int v = 0;
     trie[v].count++;
     for (int i = MAXB; i >= 0; i—) {
  int c = (x >> i) & 1;
        if (trie[v].next[c] == -1) {
   trie[v].next[c] = trie.size();
           trie.emplace_back();
        v = trie[v].next[c];
        trie[v].count++;
     trie[v].leaf++;
   void apply(ll x) { lazy ^= x; }
   ll min() {
     int v = 0
     ll ans = 0;
     for (int i = MAXB; i \ge 0; i—) {
        int b = (lazy >> i) & 1;
        int b = (td2y >> 1) d 1,
int to1 = trie[v].next[b];
int to2 = trie[v].next[b ^ 1];
if (to1 != -1) {
          v = to1;
        } else if (to2 != -1) {
  ans |= (1LL << i);</pre>
           v = to2;
        } else {
          return -1;
     return ans;
   ll max() {
     int v = 0;
     ll ans = 0;
     for (int i = MAXB; i \ge 0; i—) {
       int b = (lazy >> i) & 1;
int to1 = trie[v].next[b];
int to2 = trie[v].next[b ^ 1];
if (to2 != -1) {
          ans |= (1LL << i);
        v = to2;
} else if (to1 != -1) {
           v = to1;
        } else {
           return -1;
        }
     return ans;
   int countLE(ll x) {
     int v = 0, ans = 0;

for (int i = MAXB; i >= 0; i—) {

  int c = (x >> i) & 1;

  int b = (lazy >> i) & 1;

  if (c -- 0) {
        if (c == 0)
           if (trie[v].next[c ^ b] == -1) return ans;
           v = trie[v].next[c ^ b];
        } else {
           int to = trie[v].next[c ^ b ^ 1];
           if (to !=-1) ans += trie[to].count;
           if (trie[v].next[c ^ b] == -1) return ans;
v = trie[v].next[c ^ b];
     }
     ans += trie[v].leaf;
     return ans;
};
9.54 z function
#include <bits/stdc++.h>
using namespace std;
```

// z[i] is the length of the longest common prefix

```
// defined. "aaabaab" - [0,2,1,0,2,1,0] "abacaba"
// - [0,0,1,0,3,0,1]
vector<int> z_function(string s) {
   int n = (int)s.length();
   vector<int> z(n);
   for (int i = 1, l = 0, r = 0; i < n; i++) {
      if (i <= r) z[i] = min(r - i + 1, z[i - l]);
      while (i + z[i] < n && s[z[i]] == s[i + z[i]])
      z[i]++;
   if (i + z[i] - 1 > r) l = i, r = i + z[i] - 1;
   }
  return z;
}
```

#### 10 Primitives

#### 10.1 Bigint

```
const int maxn = 1e2 + 14, lg = 15;
const int base = 1000000000;
const int base_digits = 9;
struct bigint {
  vi a;
  int sign;
  int size() {
    if (a.empty()) return 0;
    int ans = (a.size() - 1) * base_digits;
    int ca = a.back();
    while (ca) ans++, ca \neq 10;
    return ans;
  bigint operator^(const bigint &v) {
    bigint ans = 1, a = *this, b = v;
    while (!b.isZero()) {
      if (b % 2) ans *= a;
a *= a, b /= 2;
    return ans;
  string to_string() {
    stringstream ss;
    ss << *this;
    string s;
ss >> s;
    return s;
  int sumof() {
    string s = to_string();
    int ans = 0;
    for (auto c: s) ans += c - '0';
    return ans;
  /*</arpa>*/
  bigint() : sign(1) {}
  bigint(long long v) { *this = v; }
  bigint(const string &s) { read(s); }
  void operator=(const bigint &v) {
    sign = v.sign;
        v.a;
  void operator=(long long v) {
    sign = 1;
    a.clear()
    if (v < 0) sign = -1, v = -v;
for (; v > 0; v = v / base)
      a.push_back(v % base);
  bigint operator+(const bigint &v) const {
    if (sign == v.sign) {
      bigint res = v;
      for (int i = 0, carry = 0;
            i < (int)max(a.size(), v.a.size()) ||
            carry;
        ++i) {
if (i == (int)res.a.size())
           res.a.push_back(0);
         res.a[i] += carry +
             (i < (int)a.size() ? a[i] : 0);
         carry = res.a[i] >= base;
         if (carry) res.a[i] = base;
```

```
return res;
  return *this - (-v);
bigint operator—(const bigint &v) const {
  if (sign == v.sign) {
    if (abs() >= v.abs()) {
       bigint res = *this;
       for (int i = 0, carry = 0;
             i < (int)v.a.size() || carry; ++i) {
         res.a[i] —=
carry +
              (i < (int)v.a.size() ? v.a[i] : 0);
         carry = res.a[i] < 0;
         if (carry) res.a[i] += base;
       res.trim();
       return res;
    return -(v - *this);
  return *this + (-v);
void operator*=(int v) {
  if (v < 0) sign = -\text{sign}, v = -v;
  for (int i = 0, carry = 0;
    i < (int)a.size() || carry; ++i) {
    if (i == (int)a.size()) a.push_back(0);</pre>
    long long cur = a[i] * (long long)v + carry;
    carry = (int)(cur / base);
    a[i] = (int)(cur % base);
// asm("divl %ecx" : "=a"(carry),
    // "=d"(a[i]) : "A"(cur), "c"(base));
  trim();
bigint operator*(int v) const {
  bigint res = *this;
res *= v:
  return res;
void operator*=(long long v) {
  if (v < 0) sign = -\text{sign}, v = -v; if (v > \text{base}) {
    return;
  for (int i = 0, carry = 0;
    i < (int)a.size() || carry; ++i) {
if (i == (int)a.size()) a.push_back(0);</pre>
    long long cur = a[i] * (long long)v + carry;
carry = (int)(cur / base);
    a[i] = (int)(cur % base);

// asm("divl %%ecx" : "=a"(carry),

// "=d"(a[i]) : "A"(cur), "c"(base));
  trim();
bigint operator*(long long v) const {
  bigint res = *this;
res *= v;
  return res;
friend pair<bigint, bigint> divmod(
    const bigint &a1, const bigint &b1) {
  int norm = base / (b1.a.back() + 1);
bigint a = a1.abs() * norm;
  bigint b = b1.abs() * norm;
  bigint q, r;
  q.a.resize(a.a.size());
  for (int i = a.a.size() -1; i >= 0; i—) {
    r *= base
    r += a.a[i];
    int s1 = r.a.size() \le b.a.size()
                     : r.a[b.a.size()];
    int s2 = r.a.size() \leq b.a.size() -1
                     ? 0
                       r.a[b.a.size() - 1]
    int d = ((long long)base * s1 + s2) /
              b.a.back();
     r = b * d;
```

```
while (r < 0) r += b, -d; q.a[i] = d;
  }
  q.sign = a1.sign * b1.sign;
  r.sign = al.sign;
  q.trim();
  r.trim();
  return make_pair(q, r / norm);
bigint operator/(const bigint &v) const {
  return divmod(*this, v).first;
bigint operator%(const bigint &v) const {
  return divmod(*this, v).second;
void operator/=(int v) {
  if (v < 0) sign = -\text{sign}, v = -v;
  for (int i = (int)a.size() - 1, rem = 0;
i >= 0; —i) {
    long long cur =
       a[i] + rem * (long long)base;
    a[i] = (int)(cur / v);
    rem = (int)(cur % v);
  trim();
}
bigint operator/(int v) const {
  bigint res = *this;
  res /= v
  return res;
int operator%(int v) const {
  if (v < 0) v = -v;
  int m = 0;
  for (int i = a.size() - 1; i >= 0; -
    m = (a[i] + m * (long long)base) % v;
  return m * sign;
void operator+=(const bigint &v) {
 *this = *this + v;
void operator—=(const bigint &v) {
  *this = *this - v;
void operator*=(const bigint &v) {
  *this = *this * v;
void operator/=(const bigint &v) {
  *this = *this / v;
bool operator<(const bigint &v) const {</pre>
  if (sign != v.sign) return sign < v.sign;
  if (a.size() != v.a.size())
    return a.size() * sign <
           v.a.size() * v.sign;
  for (int i = a.size() -1; i >= 0; i—)
    if (a[i] != v.a[i])
return a[i] * sign < v.a[i] * sign;
  return false;
bool operator>(const bigint &v) const {
  return v < *this;
bool operator<=(const bigint &v) const {
  return !(v < *this);
bool operator>=(const bigint &v) const {
  return !(*this < v);</pre>
bool operator==(const bigint &v) const {
  return !(*this < v) && !(v < *this);
bool operator!=(const bigint &v) const {
  return *this < v || v < *this;
void trim() {
 while (!a.empty() && !a.back()) a.pop_back();
  if (a.empty()) sign = 1;
bool isZero() const
```

```
bigint operator—() const {
  bigint res = *this;
  res.sign = -sign;
  return res;
bigint abs() const {
  bigint res = *this;
  res.sign *= res.sign;
  return res;
long longValue() const {
  long long res = 0;
  for (int i = a.size() - 1; i >= 0; i—)
    res = res * base + a[i];
  return res * sign;
friend bigint gcd(const bigint &a,
                   const bigint &b) {
  return b.isZero() ? a : gcd(b, a % b);
friend bigint lcm(const bigint &a,
                   const bigint &b) {
  return a / \gcd(a, b) * b;
void read(const string &s) {
  sign = 1
  a.clear();
  int pos = 0;
 while (pos < (int)s.size() &&
	(s[pos] == '-' || s[pos] == '+')) {
    if (s[pos] == '-') sign = -sign;
    ++pos;
  for (int i = s.size() - 1; i \ge pos;
       i —= base_digits) {
    int x = 0;
for (int j = max(pos, i - base_digits + 1);
      j <= i; j++)
x = x * 10 + s[j] - '0';
    a.push_back(x);
  trim();
}
friend istream &operator>>(istream &stream,
                             bigint &v) {
  string s;
  stream >> s;
  v.read(s);
  return stream;
friend ostream &operator<<(ostream &stream,
                             const bigint &v) {
  if (v.sign == -1) stream << '-';
  stream << (v.a.empty() ? 0 : v.a.back());
  for (int i = (int)v.a.size() - 2; i >= 0;
    stream << setw(base_digits) << setfill('0')</pre>
           << v.a[i];
  return stream;
static vector<int> convert_base(
    const vector<int> &a, int old_digits,
    int new_digits) {
  vector<long long> p(
      max(old_digits, new_digits) + 1);
  p[0] = 1;
  for (int i = 1; i < (int)p.size(); i++)
   p[i] = p[i - 1] * 10;
  vector<int> res;
  long long cur = 0;
  int cur_digits = 0;
for (int i = 0; i < (int)a.size(); i++) {</pre>
    cur += a[i] * p[cur_digits];
    cur_digits += old_digits;
    while (cur_digits >= new_digits) {
      res.push_back(int(cur % p[new_digits]));
      cur /= p[new_digits];
      cur_digits —= new_digits;
    }
  res.push_back((int)cur);
```

```
while (!res.empty() && !res.back())
        res.pop_back();
     return res;
  typedef vector<long long> vll;
  static vll karatsubaMultiply(const vll &a,
                                          const vll &b) {
     int n = a.size();
     vll res(n + n);
     if (n <= 32) {
  for (int i = 0; i < n; i++)
    for (int j = 0; j < n; j++)
             res[i + j] += a[i] * b[j];
        return res:
     }
     int k = n >> 1;
vll al(a.begin(), a.begin() + k);
vll a2(a.begin() + k, a.end());
     vll b1(b.begin(), b.begin() + k);
vll b2(b.begin() + k, b.end());
     vll a1b1 = karatsubaMultiply(a1, b1);
     vll a2b2 = karatsubaMultiply(a2, b2);
     for (int i = 0; i < k; i++) a2[i] += a1[i]; for (int i = 0; i < k; i++) b2[i] += b1[i];
     vll r = karatsubaMultiply(a2, b2);
       or (int i = 0; i < (int)alb1.size(); i++)
r[i] -= alb1[i];
     for
          (int i = 0;
                             < (int)a2b2.size(); i++)</pre>
        r[\dot{i}] = a2b2[\dot{i}];
     for (int i = 0; i < (int)r.size(); i++)
     res[i + k] += r[i];
for (int i = 0; i < (int)alb1.size(); i++)
res[i] += alb1[i];
for (int i = 0; i < (int)a2b2.size(); i++)
res[i + n] += a2b2[i];
     return res;
  bigint operator*(const bigint &v) const {
     vector<int> a6 =
   convert_base(this->a, base_digits, 6);
     vector<int> b6 =
          convert_base(v.a, base_digits, 6);
     vll a(a6.begin(), a6.end());
vll b(b6.begin(), b6.end());
     while (a.size() < b.size()) a.push_back(0);</pre>
     while (b.size() < a.size()) b.push_back(0);
while (a.size() & (a.size() - 1))
  a.push_back(0), b.push_back(0);</pre>
     vll c = karatsubaMultiply(a, b);
     bigint res;
     res.sign = sign * v.sign;
     for (int i = \bar{0}, carry = 0; i < (int)c.size();
            i++) {
        long long cur = c[i] + carry;
        res.a.push_back((int)(cur % 1000000));
        carry = (int)(cur / 1000000);
     res.a = convert_base(res.a, 6, base_digits);
     res.trim();
     return res;
10.2 Integer Mod
   ll value;
  static const ll MOD_value = _mod;
```

```
const ll MOD = 1 '000' 000'000 + 7;
template <ll _mod = MOD>
struct mint {
  mint(ll v = 0) {
    value = v % _mod;
if (value < 0) value += _mod;</pre>
  mint(ll a, ll b) : value(0) {
    *this += a;
    *this /= b;
  mint &operator+=(mint const &b) {
    value += b.value;
```

```
if (value >= _mod) value -= _mod;
    return *this;
  mint &operator—=(mint const &b) {
    value -= b.value;
    if (value < 0) value += _{mod};
    return *this;
  mint & operator*=(mint const & b) {
    value = (ll)value * b.value % _mod;
    return *this;
  friend mint mexp(mint a, ll e) {
    mint res = 1;
    while (e) {
      if (e & 1) res *= a;
a *= a;
      e >>= 1;
    return res;
  friend mint inverse(mint a) {
    return mexp(a, \_mod - 2);
  mint &operator/=(mint const &b) {
    return *this *= inverse(b);
  friend mint operator+(mint a, mint const b) {
    return a += b;
  mint operator++(int) {
    return this—>value = (this—>value + 1) % _mod;
  mint operator++() {
    return this—>value = (this—>value + 1) % _mod;
  friend mint operator—(mint a, mint const b) {
    return a -= b;
  friend mint operator—(mint const a) {
    return 0 - a;
  mint operator—(int) {
    return this->value =
                (this\rightarrowvalue -1 + \_mod) % \_mod;
  mint operator—() {
    return this—>value =
                (this\rightarrowvalue -1 + \_mod) % \_mod;
  friend mint operator*(mint a, mint const b) {
    return a *= b;
  friend mint operator/(mint a, mint const b) {
    return a /= b;
  friend std::ostream &operator<<(</pre>
      std::ostream &os, mint const &a) {
    return os << a.value;
  friend bool operator==(mint const &a,
                          mint const &b) {
    return a.value == b.value;
  friend bool operator!=(mint const &a,
                          mint const &b) {
    return a.value != b.value;
};
```

## 10.3 Matrix

```
template <typename T>
struct Matrix {
  vector<vector<T>> d;
  Matrix() : Matrix(0) {}
  Matrix(int n) : Matrix(n, n) {}
Matrix(int n, int m)
      : Matrix(
             vector<vector<T>>(n, vector<T>(m))) {}
  Matrix(const vector<vector<T>>> &v) : d(v) {}
  constexpr int n() const {
    return (int)d.size();
```

```
constexpr int m() const {
  return n() ? (int)d[0].size() : 0;
void rotate() { *this = rotated(); }
Matrix<T> rotated() const {
  Matrix<T> res(m(), n());
for (int i = 0; i < m(); i++) {
   for (int j = 0; j < n(); j++) {
     res[i][j] = d[n() - j - 1][i];
}</pre>
  return res;
Matrix<T> pow(int power) const {
  assert(n() == m());
  auto res = Matrix<T>::identity(n());
  auto b = *this;
  while (power)
     if (power \& 1) res *= b;
     b *= b;
     power >>= 1;
  return res;
}
Matrix<T> submatrix(int start_i, int start_j,
                         int rows = INT_MAX,
int cols = INT_MAX) const {
  rows = min(rows, n() - start_i);
cols = min(cols, m() - start_j);
  if (rows <= 0 or cols <= 0) return {};
  Matrix<T> res(rows, cols);
for (int i = 0; i < rows; i++)
  for (int j = 0; j < cols; j++)</pre>
       res[i][j] = d[i + start_i][j + start_j];
  return res;
Matrix<T> translated(int x, int y) const {
  Matrix<T> res(n(), m());
for (int i = 0; i < n(); i++) {
  for (int j = 0; j < m(); j++) {
   if (i + x < 0 or i + x >= n() or
               + y < 0 \text{ or } j + y >= m())
          continue;
       res[i + x][j + y] = d[i][j];
    }
  return res;
}
static Matrix<T> identity(int n) {
  Matrix<T> res(n);
  for (int i = 0; i < n; i++) res[i][i] = 1;
  return res;
vector<T> &operator[](int i) { return d[i]; }
const vector<T> &operator[](int i) const {
  return d[i];
Matrix<T> &operator+=(T value) {
  for (auto &row : d)
    for (auto &x : row) x += value;
  return *this:
Matrix<T> operator+(T value) const {
  auto res = *this;
  for (auto &row : res) {
  for (auto &x : row) x = x + value;
  return res;
Matrix<T> &operator—=(T value) {
  for (auto &row : d) {
  for (auto &x : row) x -= value;
  return *this;
Matrix<T> operator—(T value) const {
  auto res = *this;
  for (auto &row : res)
     for (auto \&x : row) x = x - value;
```

```
return res;
Matrix<T> &operator*=(T value) {
      for (auto &row : d) {
  for (auto &x : row) x *= value;
      return *this;
Matrix<T> operator*(T value) const {
      auto res = *this;
      for (auto &row : res) {
           for (auto \&x : row) x = x * value;
      return res;
Matrix<T> &operator/=(T value) {
      for (auto &row : d)
           for (auto &x : row) x /= value;
      return *this:
Matrix<T> operator/(T value) const {
      auto res = *this;
      for (auto &row : res) {
           for (auto \&x : row) x = x / value;
       return res;
Matrix<T> &operator+=(const Matrix<T> &o) {
     assert(n() == o.n() and m() == o.m());
for (int i = 0; i < n(); i++) {
  for (int j = 0; j < m(); j++) {
    d[i][j] += o[i][j];
      return *this;
Matrix<T> operator+(const Matrix<T> &o) const {
      assert(n() == o.n() and m() == o.m());
     auto res = *this;
for (int i = 0; i < n(); i++) {
  for (int j = 0; j < m(); j++)
                res[i][j] = res[i][j] + o[i][j];
           }
       return res;
Matrix<T> &operator—=(const Matrix<T> &o) {
     assert(n() == o.n() and m() == o.m());
for (int i = 0; i < n(); i++) {
  for (int j = 0; j < m(); j++) {
    d[i][j] -= o[i][j];
      return *this;
Matrix<T> operator—(const Matrix<T> &o) const {
      assert(n() == o.n() and m() == o.m());
      auto rès = *this;
      for (int i = 0; i < n(); i++) {
  for (int j = 0; j < m(); j++) {
    res[i][j] = res[i][j] - o[i][j];</pre>
      return res;
Matrix<T> &operator*=(const Matrix<T> &o) {
      *this = *this * o;
      return *this;
Matrix<T> operator*(const Matrix<T> &o) const {
      assert(m() == o.n());
     Matrix<T> res(n(), o.m());
for (int i = 0; i < res.n(); i++) {
  for (int j = 0; j < res.m(); j++) {
    auto &x = res[i][j];
    continue of the continue of 
                 for (int k = 0; k < m(); k++) {
  x += (d[i][k] * o[k][j]);
                 }
           }
      return res;
friend istream &operator>>(istream &is,
                                                                                Matrix<T> &mat) {
```

```
for (auto &row : mat)
    for (auto &x : row) is >> x;
    return is;
}
friend ostream &operator<<(
      ostream &os, const Matrix<T> &mat) {
    bool frow = 1;
    for (auto &row : mat) {
      if (not frow) os << '\n';
      bool first = 1;
      for (auto &x : row) {
        if (not first) os << '';
        os << x;
        first = 0;
    }
    frow = 0;
}
    return os;
}
auto begin() { return d.begin(); }
auto rend() { return d.rend(); }
auto rend() { return d.rend(); }
auto begin() const { return d.begin(); }
auto begin() const { return d.end(); }
auto rend() const { return d.end(); }
auto rend() const { return d.rend(); }
</pre>
```

# 11 Strings

#### 11.1 Count distinct anagrams

```
const ll MOD = 1e9 + 7;
const int maxn = 1e6;
vll fs(maxn + 1);
void precompute() {
    fs[0] = 1;
    for (ll i = 1; i <= maxn; i++) {
        fs[i] = (fs[i - 1] * i) % MOD;
    }
}
ll fpow(ll a, int n, ll mod = LLONG_MAX) {
    if (n == 0) return 1;
    if (n == 1) return a;
    ll x = fpow(a, n / 2, mod) % mod;
    return ((x * x) % mod * (n & 1 ? a : 1ll)) %
        mod;
}
ll distinctAnagrams(const string &s) {
    precompute();
    vi hist('z' - 'a' + 1, 0);
    for (auto &c : s) hist[c - 'a']++;
    ll ans = fs[len(s)];
    for (auto &q : hist) {
        ans = (ans * fpow(fs[q], MOD - 2, MOD)) % MOD;
    }
    return ans;
}</pre>
```

#### 11.2 Double hash range query

```
using ll = long long;
using vll = vector<ll>;
using pll = pair<ll, ll>;
const int MAXN(1 '000' 000);
const ll MOD = 1000027957;
const ll MOD2 = 1000015187;
const ll P = 31;
ll p[MAXN + 1], p2[MAXN + 1];
void precompute() {
  p[0] = p2[0] = 1;
  for (int i = 1; i <= MAXN; i++)
    p[i] = (P * p[i - 1]) % MOD,
    p2[i] = (P * p2[i - 1]) % MOD2;
}
struct Hash {
  int n;
  vll h, h2, hi, hi2;
  Hash() {}</pre>
```

```
Hash(const string &s)
    : n(s.size()), h(n), h2(n), hi(n), hi2(n) {
   h[0] = h2[0] = s[0];

for (int i = 1; i < n; i++)

h[i] = (s[i] + h[i - 1] * P) % MOD,

h2[i] = (s[i] + h2[i - 1] * P) % MOD2;
   pll query(int l, int r) {
    ll hash =
   (h[r] - (l? h[l-1] * p[r-l+1] % MOD : 0));
lhash2 =
        (h2[r]
        (l^?h2[l-1] * p2[r-l+1] % MOD2
           : 0));
   pll query_inv(int l, int r) {
    ll\ hash = (hi[l] -
              (r+1 < n ? hi[r+1] *
                              p[r - l + 1] % MOD
                         : 0));
    ll\ hash2 =
        (hi2[l] -
        (r + 1)
             ? hi2[r + 1] * p2[r - l + 1] % MOD2
    };
```

#### 11.3 Hash range query

```
const ll P = 31;
const ll MOD = 1e9 + 9;
const int MAXN(1e6);
ll ppow[MAXN + 1];
void pre_calc() {
   ppow[0] = 1;
   for (int i = 1; i <= MAXN; i++)
   ppow[i] = (ppow[i - 1] * P) % MOD;</pre>
}
struct Hash {
  int n;
vll h, hi;
   Hash(const string &s)
         : n(s.size()), h(n), hi(n) {
     : n(s.size(),, n(n), n-.,
h[0] = s[0];
hi[n - 1] = s[n - 1];
for (int i = 1; i < n; i++) {
  h[i] = (s[i] + h[i - 1] * P) % MOD;
  hi[n - i - 1] =
        (s[n - i - 1] + hi[n - i - 1] * P) %</pre>
     }
   ll qry(int l, int r) {
      ll\ hash =
            (h[r]
             (\bar{1}, h[l-1] * ppow[r-l+1] % MOD
                   : 0));
      return hash < 0 ? hash + MOD : hash;
   ll qry_inv(int l, int r) {
      ll hash =
            (hi[l]
                    1 < n
? hi[r + 1] * ppow[r — l + 1] % MOD
      : 0); return hash < 0 ? hash + MOD : hash;
};
```

# 11.4 Hash unsigned long long $2^{64} - 1$

**Description:** Arithmetic mod  $2^{64}-1$ . 2x slower than mod  $2^{64}$  and more code, but works on evil test data (e.g. Thue-Morse, where ABBA... and BAAB... of length  $2^{10}$  hash the same mod  $2^{64}$ ). "typedef ull H;" instead if you think test data is random.

```
typedef uint64_t ull;
struct H {
  ull x;
  H(ull'x = 0) : x(x) {}
  H operator+(H o) {
    return x + o.x + (x + o.x < x);
  H operator—(H o) { return *this + ~o.x; }
  H operator*(H o) {
    auto m = (__uint128_t)x * o.x;
return H((ull)m) + (ull)(m >> 64);
  ull get() const { return x + !~x; }
  bool operator==(H o) const {
    return get() == o.get();
  bool operator<(H o) const {
    return get() < o.get();
static const H C =
    (long long)1e11 +
    3; // (order ~ 3e9; random also ok)
struct Hash {
  int n;
  vector<H> ha, pw;
  Hash(string &str)
       : n(str.size())
         ha((int)str.size() + 1),
         pw(ha) {
    pw[0] = 1;
    H query(int a, int b) { // hash [a, b]
    return ha[b] - ha[a] * pw[b - a];
vector<H> getHashes(string &str, int length) {
  if ((int)str.size() < length) return {};
  H h = 0, pw = 1;
  for (int i = 0; i < length; i++)
    h = h * C + str[i], pw = pw * C;
  vector<H> ret = {h};
for (int i = length; i < (int)str.size(); i++)
  ret.push_back(h = h * C + str[i] -</pre>
                        pw * str[i - length]);
  return ret;
H hashString(string &s) {
  H h{};
  for (char c : s) h = h * C + c;
  return h;
```

## 11.5 K-th digit in digit string

**Description**: Find the k-th digit in a *digit string*, only works for  $1 <= k <= 10^{18}$ ! **Time**: precompute O(1), query O(1)

```
memo.push_back(
         {start, end, startstr, endstr});
     start = end + 1ll;
    end = end + (9ll * pow10[qtd]);
     curlenght = end - start + 1ull;
    startstr = endstr + 1ull;
     endstr
          (endstr + 1ull) + (curlenght)*qtd - 1ull;
  }
char kthDigit(ull k) {
  int qtd = 1;
for (auto [s, e, ss, es] : memo) {
  if (k >= ss and k <= es) {</pre>
       ull pos = k - ss;
ull index = pos / qtd;
       ull nmr = s + index;
       int i = k - ss - qtd * index;
       return ((nmr / pow10[qtd -i - 1]) % 10) + '0';
    qtd++;
  }
  return 'X';
}
```

# 11.6 Longest Palindrome Substring (Manacher)

**Description**: Finds the longest palindrome substring, manacher returns a vector where the i-th position is how much is possible to grow the string to the left and the right of i and keep it a palindrome.  $\mathbf{Time}: O(N)$ 

```
vi manacher(const string &s) {
   int n = len(s) - 2;
   vi p(n + 2);
   int l = 1, r = 1;
   for (int i = 1; i <= n; i++) {
      p[i] = max(0, min(r - i, p[l + (r - i)]));
      while (s[i - p[i]] == s[i + p[i]]) p[i]++;
      if (i + p[i] > r) l = i - p[i], r = i + p[i];
      p[i]—;
   }
   return p;
}
string longest_palindrome(const string &s) {
    string t("$#");
   for (auto c : s)
      t.push_back(c), t.push_back('#');
      t.push_back('^');
   vi xs = manacher(t);
   int mpos = max_element(all(xs)) - xs.begin();
   string p;
   for (int k = xs[mpos], i = mpos - k;
      i <= mpos + k; i++)
      if (t[i] != '#') p.push_back(t[i]);
   return p;
}</pre>
```

#### 11.7 Longest palindrome

#### 11.8 Lyndon factorization

```
vi lyndon_factorization(string S) {
   auto sa = suffix_array(S);
   vi ans;
   vi mex(len(S) + 1, 0);
   int p = 0;
   rtrav(si, sa) {
      if (si == p) {
        ans.eb(si);
      }
      mex[si] = 1;
      while (mex[p]) p++;
   }
   ans.eb(len(S));
   return ans;
}
```

#### 11.9 Rabin-Karp

```
size_t rabin_karp(const string &s,
                         const string &p) {
  if (s.size() < p.size()) return 0;</pre>
  const ll p1_1 = fpow(p1, q1 - 2, q1),
               p1_2 = fpow(p1, m - 1, q1);
  const ll p2_1 = fpow(p2, q2 - 2, q2),

p2_2 = fpow(p2, m - 1, q2);
  pair<ll, ll> hs, hp;
  for (int i = (int)m - 1; ~i; —i) {
  hs.first = (hs.first * p1) % q1;
  hs.first = (hs.first + (s[i] - 'a' + 1)) % q1;
  hs.second = (hs.second * p2) % q2;
     hs.second =
           (hs.second + (s[i] - 'a' + 1)) % q2;
     hp.first = (hp.first * p1) % q1;
hp.first = (hp.first + (p[i] - 'a' + 1)) % q1;
     hp.second = (hp.second * p2) % q2;
     hp.second =
           (hp.second + (p[i] - 'a' + 1)) % q2;
  }
  size_t occ = 0;
  for (size_t i = 0; i < n - m; i++) {
     occ += (hs == hp);
     int fi = s[i] - 'a' + 1;
int fm = s[i + m] - 'a' + 1;
     hs.first = (hs.first - fi + q1) % q1;
     hs.first = (hs.first * pl_1) % q1;
hs.first = (hs.first + fm * pl_2) % q1;
     hs.second = (hs.second - fi + q2) % q2;
hs.second = (hs.second * p2_1) % q2;
hs.second = (hs.second + fm * p2_2) % q2;
  occ += hs == hp;
  return occ;
```

# 11.10 Suffix array

```
#include <bits/stdc++.h>
using namespace std;
#ifdef LOCAL
#include "debug.cpp"
#else
#define dbg(...)
#endif
#define endl '\n'
#define fastio
   ios_base::sync_with_stdio(0); \
   cin.tie(0);
#define int long long
#define all(j) j.begin(), j.end()
#define rall(j) j.rbegin(), j.rend()
#define rall(j) (int)j.size()
#define rep(i, a, b)
   for (common_type_t<decltype(a), decltype(b)>
               i = (a);
i < (b); i++)
#define rrep(i, a, b)
   for (common_type_t<decltype(a), decltype(b)>
i = (a);
i > (b); i—)
#define trav(xi, xs) for (auto &xi : xs)
#define rtrav(xi, xs) \
for (auto &xi : ranges::views::reverse(xs))
#define pb push_back
\#define\ pf\ push\_front
#define ppb pop_back
#define ppf pop_front
#define eb emplace_back
#define lb lower_bound
#define ub upper_bound
#define fi first
#define se second
#define emp emplace
#define ins insert #define divc(a, b) ((a) + (b)-111) / (b)
using str = string;
using ll = long long;
using ull = unsigned long long;
using ld = long double;
using vll = vector<ll>;
using pll = pair<ll, ll>;
using vll2d = vector<vll>;
using vi = vector<int>;
using vi2d = vector<vi>;
using pii = pair<int, int>;
using vpii = vector<pii>;
using vc = vector<char>;
using vs = vector<str>;
template <typename T, typename T2>
using umap = unordered_map<T, T2>;
template <typename T>
using pqmn =
     priority_queue<<u>T</u>, vector<T>, greater<T>>;
template <typename T>
using pqmx = priority_queue<T, vector<T>>;
template <typename T, typename U>
inline bool chmax(T &a, U const &b) {
   return (a < b ? a = b, 1 : 0);
template <typename T, typename U>
inline bool chmin(T &a, U const &b) {
   return (a > b ? a = b, 1 : 0);
vector<int> sort_cyclic_shifts(string const &s) {
   int n = s.size();
   const int alphabet = 128;
   vector<int> p(n), c(n),
  cnt(max(alphabet, n), 0);
for (int i = 0; i < n; i++) cnt[s[i]]++;
for (int i = 1; i < alphabet; i++)</pre>
   cnt[i] += cnt[i-1];
for (int i = 0; i < n; i++) p[—cnt[s[i]]] = i;
   c[p[0]] = 0;
  int classes = 1;
for (int i = 1; i < n; i++) {
  if (s[p[i]] != s[p[i - 1]]) classes++;
  c[p[i]] = classes - 1;</pre>
```

```
if (pn[i] < 0) pn[i] += n;
    fill(cnt.begin(), cnt.begin() + classes, \theta);
for (int i = \theta; i < n; i++) cnt[c[pn[i]]]++;
for (int i = 1; i < 1; classes; i++)
    cnt[i] += cnt[i-1];
for (int i = n - 1; i >= 0; i—)
       p[—cnt[c[pn[i]]]] = pn[i];
    cn[p[0]] = 0;
    classes = 1;
    for (int i = 1; i < n; i++) {
  pair<int, int> cur = {
           c[p[i]], c[(p[i] + (1 << h)) % n]};
       pair<int, int> prev = {
           c[p[i-1]],
           c[(p[i-1] + (1 << h)) % n]);
       if (cur != prev) ++classes;
       cn[p[i]] = classes - 1;
    c.swap(cn);
  }
  return p;
}
vector<int> suffix_array(string s) {
  vector<int> p = sort_cyclic_shifts(s);
  p.erase(p.begin());
  return p;
vector<int> longestCommonPrefix(
    const string &s, const vector<int> &suf) {
  int n = s.size()
  vector<int> isuf(n), res(n -1);
for (int i = 0; i < n; ++i) isuf[suf[i]] = i;
  int k = 0;
  for (; isuf[k] != n - 1; ++k) {
    int cmp_i = suf[isuf[k] + 1];
int r = k == 0 ? 0
                      : \max(\text{res}[\text{isuf}[k-1]] - 1,
    while (k + r < n \&\& cmp_i + r < n \&\&
            s[k + r] == s[cmp_i + r])
    res[isuf[k]] = r;
  ,
++k;
  for (int i = k; i < n; ++i) {
  int cmp_i = suf[isuf[i] + 1];</pre>
    int r = i == k ? 0
                      : max(res[isuf[i-1]] - 1,
                              (int)0);
    while (i + r < n \&\& cmp_i + r < n \&\&
            s[i + r] == s[cmp_i + r])
    res[isuf[i]] = r;
  return res;
ll distinct_substrings(const string &s,
                           const vi &sa) {
  int_n = len(s);
  vi lcp = longestCommonPrefix(s, sa);
  ll ans = n - sa[0];
  rep(i, 1, n) { ans += n - sa[i] - lcp[i - 1]; }
  return ans;
void run();
int32_t main() {
#ifndef LOCAL
  fastio;
#endif
  int T = 1;
  /*cin >> T;*/
  rep(t, 0, T) {
    dbg(t);
    run();
```

```
void run() {
  string S;
  cin >> S;
  auto sa = suffix_array(S);
  cout << distinct_substrings(S, sa) << endl;</pre>
11.11 Suffix array (supreme)
template <typename T = ll,
           auto cmp :
                [](T &src1, T &src2, T &dst) {
                  dst = min(src1, src2);
class SparseTable {
 private:
  int_sz;
  vi logs;
  vector<vector<T>> st;
 public:
  SparseTable() {}
  rep(i, 0, sz) st[0][i] = v[i];
    for (int k = 1; (1 << k) <= sz; k++) {
  for (int i = 0; i + (1 << k) <= sz; i++) {
    cmp(st[k - 1][i],
    st[k - 1][i] + (1 << (k - 1))],
              st[k][i]);
    }
  T query(int l, int r) {
    const int k = logs[r - l];
    T ret;
    cmp(st[k][l], st[k][r - (1 << k)], ret);
    return ret;
};
template <typename T>
using RMQ = SparseTable<T, [](T &a, T &b, T &c) {
  c = min(a, b);
// éCrditos: ShahjalalShohag
// O(N)
struct SA {
  string s;
  int n;
  vector<int> sa, lcp, pos;
  RMQ<int> rmq;
  void induced_sort(vector<int> &vec, int val,
                       vector<int> &sa
                       vector<bool> &sí
                       vector<int> &lms) {
    vector<int> l(val), r(val);
for (int c : vec) {
    for (int c : vec)
      if (c + 1 < val) l[c + 1] ++;
      r[c]++;
    partial_sum(l.begin(), l.end(), l.begin());
partial_sum(r.begin(), r.end(), r.begin());
    fill(sa.begin(), sa.end(), -1);
    for (int i = lms.size() - 1; i >
    sa[--r[vec[lms[i]]]] = lms[i];
                                      i >= 0; i--)
    for (int i : sa) {
  if (i >= 1 && sl[i - 1])
    sa[l[vec[i - 1]]++] = i - 1;
    fill(r.begin(), r.end(), 0);
    for (int c : vec) r[c]++;
    sa[-r[vec[i-1]]] = i-1;
  }
  vector<int> build_sa(vector<int> &vec,
```

```
int n = vec.size();
vector<int> sa(n), lms;
vector<bool> sl(n);
       sl[n-1] = false;
for (int i = n - 2; i >= 0; i—) {
             sl[i] =
                           (\text{vec}[i] > \text{vec}[i + 1] \mid \mid
             (\text{vec}[i] == \text{vec}[i + 1] \& \text{sl}[i + 1]));
if (\text{sl}[i] \& \& !\text{sl}[i + 1])
                     lms.push_back(i + 1);
      reverse(lms.begin(), lms.end());
induced_sort(vec, val, sa, sl, lms);
vector<int> new_lms(lms.size()),
       lms_vec(lms.size());
for (int i = 0, k = 0; i < n; i++) {
  if (!sl[sa[i]] && sa[i] >= 1 &&
     sl[sa[i] - 1])
     rectified the same of the sa
                    new_lims[k++] = sa[i];
       int cur = 0;
       sa[n - 1] = cur;
for (int k = 1; k < (int)new_lms.size();</pre>
                       k++) {
              int i = \text{new\_lms}[k - 1], j = \text{new\_lms}[k];
             if (vec[i] != vec[j]) {
                    sa[j] = ++cur;
                    continue;
             bool flag = false;
             for (int a = i + 1, b = j + 1;; ++a, ++b) {
  if (vec[a] != vec[b]) {
                          flag = true;
                          break;
                    if ((!sl[a] && sl[a - 1]) ||
    (!sl[b] && sl[b - 1])) {
    flag = !((!sl[a] && sl[a - 1]) &&
                                                         (!sl[b] \&\& sl[b-1]));
                          break:
                   }
             sa[j] = (flag ? ++cur : cur);
       for (int i = 0; i < (int)lms.size(); i++)
  lms_vec[i] = sa[lms[i]];</pre>
       if (cur + 1 < (int)lms.size()) {
             auto lms_sa = build_sa(lms_vec, cur + 1);
             for (int i = 0; i < (int)lms.size(); i++)
new_lms[i] = lms[lms_sa[i]];</pre>
       induced_sort(vec, val, sa, sl, new_lms);
       return sa;
vector<int> suffix_array() {
      vector<int> vec(n + 1);
copy(begin(s), end(s), begin(vec));
vec.back() = '$';
auto sa = build_sa(vec, 256);
       sa.erase(sa.begin());
       return sa;
vector<int> build_lcp() {
  int n = (int)s.size(), k = 0;
  vector<int> rank(n), lcp(n);
  for (int i = 0; i < n; i++) rank[sa[i]] = i;</pre>
       for (int i = 0; i < n; i++, k = !!k) {
             if (rank[i] = n - 1) {
                   k = 0;
                    continue;
            int j = sa[rank[i] + 1];

while (i + k < n \&\& j + k < n \&\&

s[i + k] == s[j + k])
            lcp[rank[i]] = k;
       return lcp;
}
SA() {}
SA(string_s) : s(s), n(len(s)), pos(n) {
       sa = suffix_array();
       lcp = build_lcp();
       rmq = RMQ<int>(lcp);
```

```
for (int i = 0; i < n; i++) pos[sa[i]] = i;
  int get_lcp(
      int i
                 // lcp na çãposio i, indica o lcp
      int j) {
                 // das çõposies i e i+1 do sa
    if (i == j) return n - i;
    int l = pos[i], r = pos[j];
    if (l > r) swap(l, r);
    return rmq.query(l, r);
  // string s = a + '+' + b;
tuple<int, int, int> lcs(
      int n) { // m é o tamanho da string a
    int m = len(s) - n - 1;
    int best_len = 0;
    int index_s = 0;
int index_t = 0;
for (int i = 0; i < n + m; ++i) {</pre>
      if ((sa[i] < n \&\& sa[i + 1] >= n + 1) ||
        (sa[i] >= n + 1 \&\& sa[i + 1] < n)) { if (lcp[i] > best_len) {}
           best_len = lcp[i];
           index_s = min(sa[i], sa[i + 1]);
index_t = max(sa[i], sa[i + 1]) - n - 1;
      }
    /*
             maior = lcp[i - 1], pos = sa[i];*/
    /*}*/
    /*return {maior, pos};*/
    return {best_len, index_s, index_t};
  // n*(n+1)/2 - sum(lcp[i])
    return resp — accumulate(lcp.begin()
                               lcp.end(), OLL);
};
```

#### 11.12 Suffix automaton

```
#include <bits/stdc++.h>
using namespace std;
#ifdef LOCAL
#include "debug.cpp"
#else
#define dbg(...)
#endif
#define endl '\n'
#define fastio
   ios_base::sync_with_stdio(0); \
cin.tie(0);
#define int long long
#define all(j) j.begin(), j.end()
#define rall(j) j.rbegin(), j.rend()
#define len(j) (int)j.size()
#define rep(i, a, b)
   for (common_type_t<decltype(a), decltype(b)>
           i = (a);
i < (b); i++)
#define rrep(i, a, b)
   for (common_type_t<decltype(a), decltype(b)>
          i = (a);
i > (b); i—
#define trav(xi, xs) for (auto &xi : xs)
#define rtrav(xi, xs) \
  for (auto &xi : ranges::views::reverse(xs))
#define pb push_back
#define pf push_front
#define ppb pop_back
#define ppf pop_front
#define eb emplace_back
#define lb lower_bound
#define ub upper_bound
#define fi first
```

```
#define se second
#define emp emplace
#define ins insert
#define divc(a, b) ((a) + (b)-111) / (b)
using str = string;
using ll = long long;
using ull = unsigned long long;
using ld = long double;
using vll = vector<ll>
using pll = pair<ll, li>;
using vll2d = vector<vll>;
using vi = vector<int>;
using vi2d = vector<vi>;
using pii = pair<int, int>;
using vpii = vector<pii>;
using vc = vector<char>;
using vs = vector<str>;
template <typename T, typename T2>
using umap = unordered_map<T, T2>;
template <typename T>
using pqmn =
    priority_queue<T, vector<T>, greater<T>>;
template <typename T>
using pqmx = priority_queue<T, vector<T>>;
template <typename T, typename U>
inline bool chmax(T &a, U const &b) {
  return (a < b ? a = b, 1 : 0);
template <typename T, typename U>
inline bool chmin(T &a, U const &b) {
  return (a > b ? a = b, 1 : 0);
struct SuffixAutomaton {
  struct state {
    int len, link, cnt, firstpos;
    // this can be optimized using a vector with
// the alphabet size
    map<char, int> next;
    vi inv_link;
  vector<state> st;
  int sz = 0;
int last;
  vc cloned;
  SuffixAutomaton(const string &s, int maxlen)
    . sc(maxten *
st[0].len = 0;
st[0].link = -1;
sz++;
         st(maxlen * 2), cloned(maxlen * 2) {
     last = 0;
     for (auto &c : s) add_char(c);
     // precompute for count occurences
     for (int i = 1; i < sz; i++) {
  st[i].cnt = !cloned[i];</pre>
     vector<pair<state, int>> aux;
     for (int i = 0; i < sz; i++) {
  aux.push_back({st[i], i});</pre>
     sort(all(aux), [](const pair<state, int> &a;
                          const pair<state, int> &b) {
       return a.fi.len > b.fi.len;
     });
     for (auto &[stt, id] : aux) { if (stt.link !=-1) {
         st[stt.link].cnt += st[id].cnt;
     }
     // for find every occurende position
     for (int v = 1; v < sz; v++)
       st[st[v].link].inv_link.push_back(v);
  void add_char(char c) {
     int cur = sz++;
st[cur].len = st[last].len + 1;
     st[cur].firstpos = st[cur].len - 1;
     int p = last;
     // follow the suffix link until find a
// transition to c
     while (p != -1 \text{ and } !st[p].next.count(c)) {
       st[p].next[c] = cur;
```

```
p = st[p].link;
    // there was no transition to c so create and // leave
    if (p == -1) {
       st[cur].link = 0;
       last = cur;
       return;
    int q = st[p].next[c];
    if (st[p].len + 1 == st[q].len) {
      st[cur].link = q;
    } else {
       int clone = sz++
       cloned[clone] = true;
       st[clone].len = st[p].len + 1;
      st[clone].next = st[q].next;
      st[clone].link = st[q].link;
       st[clone].firstpos = st[q].firstpos;
      while (p != -1 \text{ and } st[p].next[c] == q) {
         st[p].next[c] = clone;
         p = st[p].link;
      st[q].link = st[cur].link = clone;
    last = cur;
  }
  bool checkOccurrence(
      const string &t) { // O(len(t))
    int cur = 0;
for (auto &c : t) {
   if (!st[cur].next.count(c)) return false;
      cur = st[cur].next[c];
    return true;
  ll totalSubstrings() { // distinct, O(len(s))
    ll tot = 0;
for (int i = 1; i < sz; i++) {
  tot += st[i].len - st[st[i].link].len;
    return tot;
  // count occurences of a given string t
  int countOccurences(const string &t) {
    int cur = 0;
    for (auto &c : t) {
   if (!st[cur].next.count(c)) return 0;
      cur = st[cur].next[c];
    return st[cur].cnt;
  // find the first index where t appears a
  // substring O(len(t))
  int firstOccurence(const string &t) {
    int cur = 0;
    for (auto c : t) {
  if (!st[cur].next.count(c)) return -1;
      cur = st[cur].next[c];
    return st[cur].firstpos - len(t) + 1;
  vi everyOccurence(const string &t) {
    int cur = 0;
for (auto c : t) {
   if (!st[cur].next.count(c)) return {};
      cur = st[cur].next[c];
    vi ans;
    getEveryOccurence(cur, len(t), ans);
    return ans;
  void getEveryOccurence(int v, int P_length,
                            vi &ans) {
    if (!cloned[v])
      ans.pb(st[v].firstpos - P_length + 1);
    for (int u : st[v].inv_link)
      getEveryOccurence(u, P_length, ans);
  }
};
void run();
int32_t main() {
#ifndef LOCAL
```

```
fastio;
#endif
  int T = 1;
  /*cin >> T;*/
  rep(t, 0, T) {
    dbg(t);
     run();
void run() {
  string S;
cin >> S;
  SuffixAutomaton sa(S, len(S));
  cout << sa.totalSubstrings() << endl;</pre>
11.13 Trie
Description:
   • build with the size of the alphabet (sigma) and the first char (norm)
   • insert(s) insert the string in the trie O(|s|*sigma)
   • erase(s) remove the string from the trie O(|s|)
   • find(s) return the last node from the string s, 0 if not found O(|s|)
struct trie {
  vi2d to;
  vi end, pref;
  int sigma;
  char norm;
  trie(int sigma_ = 26, char norm_ = 'a')
     : sigma(sigma_), norm(norm_) {
to = {vector<int>(sigma)};
     end = \{0\}, pref = \{0\};
  int next(int node, char key) {
  return to[node][key - norm];
  void insert(const string &s) {
     int x = 0;
     for (auto c : s) {
       int &nxt = to[x][c - norm];
       if (!nxt) {
          nxt = len(to);
          to.push_back(vi(sigma));
          end.emplace_back(0), pref.emplace_back(0);
       x = nxt, pref[x]++;
     end[x]++, pref[0]++;
  void erase(const string &s) {
     int x = 0;
for (char c : s) {
       int &nxt = to[x][c - norm];
x = nxt, pref[x]—;
       if (!pref[x]) nxt = 0;
     end[x]—, pref[0]—;
  int find(const string &s) {
     int x = 0;
     for (auto c : s) {
```

x = to[x][c - norm];if (!x) return 0;

```
return x;
};
```

#### 11.14 Z-function get occurrence positions

```
Time: O(len(s) + len(p))
vi getOccPos(string& s, string& p) {
  // Z—function
char delim = '#'
  string t{p + delim + s};
  vi zs(len(t));
  t[zs[i]] == t[i + zs[i]])
      zs[i]++;
    if (r < i + zs[i] - 1)
      l = i, r = i + zs[i] - 1;
  // Iterate over the results of Z-function to get
  // ranges
  vi ans;
  int start = len(p) + 1 + 1 - 1;
  for (int i = start; i < len(zs); i++) {</pre>
    if (zs[i] == len(p)) {
  int l = i - start;
      ans.emplace_back(1);
  return ans;
template <class T>
std::vector<int> z_algorithm(
    const std::vector<T>& s) {
  int n = int(s.size());
if (n == 0) return {};
  std::vector<int> z(n);
  k = (j + z[j]) \le i)
             : std::min(j + z[j] - i, z[i - j]);
    while (i + k < n \&\& s[k] == s[i + k]) k++; if (j + z[j] < i + z[i]) j = i;
  z[0] = n;
  return z;
std::vector<int> z_algorithm(
    const std::string& s) {
  int n = int(s.size());
  std::vector<int> s2(n);
for (int i = 0; i < n; i++) {
    s2[i] = s[i];
  return z_algorithm(s2);
```

# 12 Trees

}