Modified Sic One Pass Assembler

1. Team members:

- Nada Hamada 20101043 Class 5
- Ahmed Elshennawy 20100483 Class 7
- Ahmed Amr 20100159 Class 5

2. <u>Detailed Description of the Code and Functions Used:</u>

This code is a one pass assembler that works on modified sic assembly language. Modified sic includes format 3, format 3 modified - which deals with immediate values, and format 1 instructions.

A one pass assembler generates the location counter and object code of an instruction while generating a corresponding text record and symbol table in one process.

Location counters are incremented according to the instruction format. Unless the instruction is a byte/word/resw/resb, these have special calculations which are covered in the code.

The object code is made up of 2 parts, opcode – result of mapping the instruction to the dictionary, and a target label – address/immediate value. In case of an immediate value, it is written as 4 hex characters. I case of an address; it could be referencing a line whose address is known (backward referencing) and therefore found in the symbol table so the target label is 4 hex characters of that address; or it could be referencing a line whose location is yet to be calculated (forward referencing). If an instruction is referencing an upcoming line, the target label is 4 0's, and the address of the bytes to be modified are added to a linked list so that when the line is reached, all the bytes where that address was supposed to be is modified. If this label was a combination of an address and an indexing register, a flag is raised so that during modification the target label is updated to be address + 0x8000.

To achieve this, a series of steps is required.

4 dictionaries are defined, 3 of them are for the instructions which map each key (instruction) to its corresponding opcode. The 4th dictionary contains the hexadecimal ascii code of uppercase and lowercase letters, numbers, and special characters.

Input file's (assembly) path is defined, and any output files are created during runtime.

Class Node is the linked list element definition. It is unlike the usual node definition as it contains "value" and "index". Value will contain the address to be modified. Index is a flag, which is set when the address has an indexing register.

Print_linked_list is a function that loops through the linked list and prints all values in it.

Remove_index is a function that extracts the target address from a label containing an indexing register if present. For example, BUFFER,X returns BUFFER.

Remove_immediate is function that removes an immediate value from a label. BUFFER,#4 returns BUFFER. We are not entirely sure if sic assembly allows such kind of a target address, however we've handled it either way.

Remove_hash is a function that removes the hashtag from an immediate value so that we can process the value as an integer. #4 returns 4.

Check_target_label is function that takes 4 parameters - symbol table, label, location counter, index. As we've mentioned before index is a flag that indicates whether the label originally had an indexing register. This function works by mapping the target label to the symbol table dictionary. The result could be one of two things, the target label has been previously added to the dictionary (could be defined by a backward reference in this case it has an integer value, or forward reference defined by a linked list),

or this is the first time it is being referenced (in this case added to the symbol table dictionary as a linked list).

Check_line_label is function that maps a line label to the symbol table, it is either found to be a linked list - in this case a node is added to the beginning of the list containing the location counter of this line; or it is not found - in this case it is defined as integer.

Processing the assembly code itself works by reading the input file line by line. The first line helps us generate part of the header, as it gives us the programme name and starting address. However, the header record is written once we reach the end so that we could calculate the size of the programme.

The rest of the lines are read through a for loop. Each line is split into elements separated by white spaces. Depending on the number of elements in the line, it is processed accordingly. The number of elements could either be 3, 2, or 1.

3 Elements: the line has a line label, instruction, and target label. In such a case, we need to map the line label to the symbol table by using the chack_line_label function. The instruction is initially mapped to the format 3 dictionary since format 1 instructions don't take target addresses. The target label is then processed using the functions we've created to check whether it is a simple label, contains an index register, or an immediate value. In the case of an immediate value the instruction is mapped to the modified dictionary. If the mapping to format 3 is unsuccessful, the instruction is tested to be word/byte/resw/resb.

2 Elements: the line is either format 1 and a line label, or format 3 (may be modified, too) and target label. The 2nd element (instruction) is first mapped to format 1 dictionary. If the mapping is unsuccessful it is then mapped to format 3 instructions and processed the same way as if they were 3 elements in the line (without the possibility of word/byte/resw/resb as these instructions require line and target labels).

1 Element: this is only an instruction, which could only be format 1, RSUB format 3 instruction, or END. Format 1 and 3 instructions have already been discussed. If the instruction was END, the header record is written, the size is calculated, the text records are copied from t.txt file to hte.txt file hen the end record is written which contains the address of the 1^{st} executable instruction. To find this address, a counter is incremented every time an instruction is mapped, then verified if the counter is ==1, if true the location address of this instruction is stored in a variable.

3. The Code:

Please note that is not a screenshot, this is the copied code from PyCharm.

```
creating the instruction set dictionary for format 3 without modification
ascii dictionary = {
```

```
assembly code file = open(assembly code, 'w')
symbol table file = open('symtable steps.txt', 'w')
symbol table final = open('symbol table.txt', 'w')
hte record = open('hte.txt', 'w')
```

```
def check target label(symbol table, label, counter, index):
       symbol table[label] = Node(counter + 1, index)
   if label in symbol table:
```

```
# checks if the line label was previously targeted but yet to be
   new node = Node(counter, 0)
```

```
search line label(symbol table, line label, loc)
```

```
assembly code file.write(f"{loc:04X}\t{line}\t\t{opcode:02X}{symbol table[tar
isinstance(symbol table[remove index(target label)], Node):
```

```
t size = 3
```

```
search line label(symbol table, line label, loc)
assembly code file.write(f"\{loc:04X\}\t\{line\}\t\}int(target label,
                         t size = 3
```

```
t size = 3
reloc bits = (int(reloc bits, 2))
```

```
flag = 1
```

```
first instruc loc = loc
c}\n")
```

```
remove index(target label), loc, 0)
```

```
c}\n")
                              t size = 3
```

```
flaq = 0
```

```
t record.write(f"{t size -
input file.close()
for key, value in symbol table.items():
```

```
print(f"Error: The file '{assembly}' was not found.")
except Exception as e:
   print(f"An error occurred: {e}")
```

4. Output Screenshot:

