

Pegpp, a C++ PEG library

Table of Contents

The PEG library.....	2
PEG grammars.....	2
A simple example.....	4
A calculator.....	5
The rules.....	5
Left recursion.....	6
Parsing the grammar.....	6
The value stack.....	6
Implicit primaries.....	6
Action and semantic predicate macros.....	7
Checking a grammar.....	8
Unicode support.....	9
The parser classes.....	10
A multi-typed value stack parser.....	11

The PEG library

PEG (Parsing Expression Grammar) is a formalism for specifying recursive descent parsers with unlimited lookahead and backtracking. The theory of PEGs can be found in Wikipedia or in Bryan Ford's original paper <http://bford.info/pub/lang/peg.pdf>.

Like YACC, which is a compiler of context-free grammars, there are compilers of PEG grammars. The LEG tool by Ian Piumarta is a good example. LEG takes a PEG grammar with embedded actions and generates a parser in C.

Pegpp is inspired by LEG and tries to mimic its syntax, but it is not a compiler. Using this library PEG grammars can be embedded in C++ code and parsed without intermediate code generation.

The way of handling the value stack was inspired by YACC.

The library is a single header file (peg.h) and requires C++17.

PEG grammars

PEG grammars are written using a syntax similar to Extended Backus-Naur Form (EBNF).

A PEG grammar is a set of rules (non-terminals). One of them is selected as the start rule. Each rule is assigned a parsing expression, or just "expression", for short:

```
rule 1 = expression 1;
rule 2 = expression 2;
...
rule n = expression n;
```

Once assigned, a rule allows referring to an expression by name.

These are the most basic expressions:

- Primaries
- Rules
- Embedded actions
- Semantic predicates

Primaries match some basic input patterns. If they parse successfully they consume the matched input. Otherwise, no input is consumed. Input read during an unsuccessful parse is returned to the input stream.

Any() matches any character. It only fails when no character can be read, i.e. at end of file.

Lit(c) matches a single character c.

Lit(s) matches the character sequence in string s.

Ccl(s) matches any character contained in the character class defined by string s. Character - is special. If it appears between two other characters in s it defines a character range, otherwise, if it is the first or last character in s, it just represents itself. Ccl(" \t\r\n\f") may be a definition of white space. Ccl("0-9") is the class of decimal digits. If s begins with ^, the class is complemented or negated. The ^ is removed from s and the class contains all characters not included in the rest of s. Thus Ccl("^") is a class that contains all characters and matches any. It is equivalent to Any().

A *rule* is a reference to the expression assigned to it.

Embedded actions do not consume input and always succeed. They schedule actions to be executed after a successful parse, if the branch that contains them is part of the match. Actions scheduled during an unsuccessful parse are cancelled.

Do(f) schedules an action f. It takes an argument of type std::function<void()>, i.e. a function, function object or lambda that takes no arguments and returns nothing.

Semantic predicates do not consume input. They are executed during parsing and may be used as a guard in order to make parsing fail or succeed depending on certain semantic conditions.

Pred(f) takes an argument of type std::function<void(bool &)>, i.e. a function, function object or lambda that receives a reference to a boolean variable and returns nothing. The function passed to Pred() is executed when the predicate is parsed, and parsing succeeds or fails, depending on the value of the boolean variable when the function returns. The variable is initialized to true (success) before calling the function, so that Pred(f) succeeds by default if f does not modify its argument.

If *e* is an expression, other expressions may be obtained by applying some prefix and postfix operators:

- A text-capturing expression: `--e e--`
- A repetitive expression: `e[{n1, n2}] e[{n, n}] e[{n}] e[n] ~e *e +e`
- A syntactic predicate: `&e !e`

If *e1* and *e2* are expressions, other expressions may be obtained by composing them with some binary operators:

- A sequence: `e1 >> e2`
- An attachment: `e1(e2)`
- An ordered choice: `e1 | e2`

Text capture

`--e` or `e--` succeed or fail as *e* does. If they succeed they capture the text consumed by *e*. The captured text is available in `Do()` or `Pred()` code following the capture. Captures may be nested (which is useful for debugging).

Repetitions

`e[{n1, n2}]` means parse *e* no less than *n1* and no more than *n2* times. It succeeds if *e* is parsed at least *n1* times. After the first *n1* times, the parser keeps trying to parse *e* until a maximum of *n2* times. If *n2* == 0 there is no upper limit and the parser keeps trying to parse *e* indefinitely.

`e[{n, n}]`, `e[{n}]` and `e[n]` are equivalent. For *n* != 0 they try to parse *e* exactly *n* times. For *n* == 0 the meaning is to parse *e* as many times as possible and succeed.

`~e` means parse *e* 0 or 1 times. It is equivalent to `e[{0, 1}]`. It tries to parse *e* once and succeeds.

`*e` means parse *e* 0 or more times. It is equivalent to `e[{0, 0}]`, `e[{0}]` or `e[0]`. It tries to parse *e* as many times as possible and succeeds.

`+e` means parse *e* 1 or more times. It is equivalent to `e[{1, 0}]`. It tries to parse *e* as many times as possible and succeeds if *e* is parsed at least once.

Syntactic predicates

These are lookahead mechanisms that check if some expression *e* can be parsed at the current point. If *e* parses successfully, input eventually consumed by *e* is returned to the input stream and actions eventually scheduled by *e* are canceled.

`&e` ("and-predicate") succeeds if *e* can be parsed at the current point.

`!e` ("not-predicate") succeeds if parsing *e* at the current point would fail. This kind of predicate is most frequently used as a stop condition. `!Any()` means end of file.

Sequence

Sequences, built by concatenation in the original LEG syntax, are implemented with the binary operator `>>` in C++. Parsing `e1 >> e2` starts by parsing *e1*. If *e1* fails, the expression fails without parsing *e2*. If *e1* succeeds, the result is obtained by parsing *e2*.

Attachment

An attachment `e1(e2)` parses like the sequence `e1 >> e2`, but *e2* does not occupy any slots in the value stack. This will be explained later.

Choice

The prioritized choice operator is `|`. Parsing `e1 | e2` starts by parsing *e1*. If *e1* succeeds, the expression succeeds without parsing *e2*. If *e1* fails, the result is obtained by parsing *e2*. This short-circuit evaluation makes PEG grammars deterministic, since alternatives are always tried in fixed left-to-right order.

These are the available operators grouped by decreasing order of priority:

- `e-- e1(e2) e[{n1, n2}] e[{n}] e[n]`
- `--e &e !e ~e *e +e`
- `e1 >> e2`
- `e1 | e2`

Parentheses may be used for grouping in order to override priorities.

Unary operator `&` has been overridden, so that `&e` means "e can be parsed at this point" and not "the address of *e*". In the unlikely case that the address of an expression is needed, use `std::addressof`.

A simple example

This little example is a direct translation of a LEG parser taken from the LEG distribution. It copies its input to its output, replacing occurrences of the string "username" by the current user's login name.

This is the original LEG source:

```
%{
    #include <unistd.h>
    #include <stdio.h>
}%

start  = "username"      { printf("%s", getlogin()); }
      | < . >           { putchar(yytext[0]); }

%%

int main()
{
    while ( yyparse() )
        ;
    return 0;
}
```

This is the same parser using pegpp:

```
#include <unistd.h>
#include <iostream>

#include "peg.h"

using namespace std;
using namespace peg;

class parser : public Parser<>
{
    Rule start;

public:
    parser(istream &in = cin) : Parser(start, in)
    {
        start  = "username"_lit      _( cout << getlogin(); )
          | Any()--                  _( cout << text(); )
          ;
    }
};

int main()
{
    parser p;
    while ( p.parse() )
        p.accept();
}
```

This parser has only one rule (start). The empty parameter list in Parser<> indicates that it does not use a value stack; if it did, the types of the values returned by the rules would be listed between the angle brackets.

The constructor receives an argument with the input stream to read characters from, by default std::cin (standard input). The base class (Parser) needs to be initialized with the grammar's starting rule (start) and the input stream, which also defaults to std::cin if not passed.

The "username"_lit construction is a user-defined literal returning Lit("username"). It matches the string "username" literally. Characters not matching "username" are matched individually by Any().

Embedded actions are placed in the grammar using the _(...) construct, which replaces the pair of braces { ... } used in the original LEG syntax. This construct is a macro, which expands to (peg::Do([&]{ ... })), i.e. it schedules a capture-all-by-reference lambda. Note that the Do() expression is surrounded by parentheses. This makes it attach to the preceding expression, making an intermediate sequence >> operator unnecessary.

Operator-- captures the text matched by Any(), which is accessed by calling the text() method inherited from the Parser class.

The parser is executed by calling its parse() method. When parse() returns, accept() is called to execute the actions scheduled while parsing and discard the consumed input. It is not necessary to call accept() every time parse() returns, as is done here. Matched input and scheduled actions just queue up in the parser until accept() is called.

A calculator

This simple integer calculator supports the four basic operations and grouping with parentheses:

```
#include <iostream>
#include <string>

#include "peg.h"

using namespace std;
using namespace peg;

class calculator : public Parser<int>
{
    Rule WS, SIGN, DIGIT, NUMBER, LPAR, RPAR, ADD, SUB, MUL, DIV;
    Rule calc, expression, term, factor;

public:
    calculator(istream &in = cin) : Parser(calc, in)
    {
        // Lexical rules

        WS      = "*" \t\f\r\n" _ccl;
        SIGN    = "+-" _ccl;
        DIGIT    = "0-9" _ccl;
        NUMBER   = (~SIGN >> +DIGIT)-- >> WS  _( val (0) = stoi(text()); );
        LPAR     = '(' >> WS;
        RPAR     = ')' >> WS;
        ADD      = '+' >> WS;
        SUB      = '-' >> WS;
        MUL      = '*' >> WS;
        DIV      = '/' >> WS;

        // Calculator

        calc     = WS >> expression          _( cout << val(1) << endl; )
        ;
        expression = term >> *(
            ADD >> term          _( val(0) += val(2); )
            | SUB >> term        _( val(0) -= val(2); )
        )
        ;
        term      = factor >> *(
            MUL >> factor        _( val(0) *= val(2); )
            | DIV >> factor      _( val(0) /= val(2); )
        )
        ;
        factor    = NUMBER
        | LPAR >> expression >> RPAR  _( val(0) = val(1); )
        ;
    }
};

int main()
{
    calculator c;
    while ( c.parse() )
        c.accept();
}
```

The rules

Rules are instances of the Rule class. Assigning an expression to a rule builds a syntax tree, dynamically allocating memory for the structures in the nodes of the tree. Smart pointers of the type `std::shared_ptr` take care of releasing memory when the rules are no longer used.

When a grammar is built, rules may be assigned to in any arbitrary order. Any rule can refer to itself or to any other rule, whether it has already been initialized or not. This is necessary to build recursive grammars, which are very frequent.

When used as a parsing expression, a rule behaves as a reference to itself. For this reason no Rule temporaries can be created, since they would produce expressions with dangling references. The Rule class has no constructors except the default, and the copy constructor has been deleted.

Rules must be defined as uninitialized variables. The grammar is built by assigning parsing expressions to the rules.

Left recursion

The copy assignment operator of rules is non-standard. Normally, the statement `r = r` is supposed to do nothing. However, if `r` is a rule, it actually makes `r` left-recursive, meaning that the recursive-descent parser for `r` calls itself directly without reading input. Such a rule overflows the stack if parsed. This is an example of direct left recursion, where a rule is assigned an expression that starts by trying to parse the same rule without reading any input.

Left recursion can also be indirect, originating in cyclic references. The calculator has the following reference cycle:

`expression → term → factor → expression`

The first two references (`expression → term → factor`) are direct. The third one (`factor → expression`) closes the cycle but only after reading a LPAR (left parenthesis) token. Without this LPAR the grammar would be left-recursive and the parser would enter an infinite loop and blow the stack. This is an inherent limitation of PEG grammars: reference cycles that may close without consuming at least one character from the input are forbidden.

Parsing the grammar

Once built, a grammar is used by calling method `parse()` of the parser. It returns true if the input matches the grammar, and false otherwise. When `parse()` returns the parser holds in an internal buffer all input read during parsing, both the part that matches the grammar and the one that does not. For example, let's assume that the calculator is fed with this input:

`(1 + 2) * 3 Hello`

The first call to `parse()` returns true. Now the parser has advanced its input pointer to the H in Hello (the first character that does not match the grammar). This is the limit of the part of the input consumed by the parser. The parser also holds a vector of actions (scheduled by `Do()` during parsing). These actions must be executed in order to evaluate the calculation. Calling `accept()` executes the scheduled actions, resets the vector that holds them and discards the consumed part of the input. Input now starts at the H in Hello. The next call to `parse()` fails, since Hello does not match the grammar, and the input pointer does not advance.

The value stack

Expressions in a rule may return a value in a slot of the value stack, according to their position in the sequence. Rule actions (and semantic predicates) may return a value by assigning `val(0)`. If it does not assign `val(0)`, a rule returns the value of its first expression.

The calculator inherits from `Parser<int>`, indicating that the value stack will hold only elements of type `int`. In this case, the `val(n)` method is all that is needed to access the stack. Multi-typed stacks require a special syntax, which will be explained later.

The index of each expression in a rule is easily calculated by counting how many `>>`'s separate it from the start of the rule. For example, in rule `term` of the calculator, the first factor is in position 0 and the second one in position 2. If an expression of the rule has alternatives, it uses as many stack slots as its longest alternative. For example:

```
r = e0 >> (
    a1 >> a2
    | b1 >> b2 >> b3 >> b4
) >> e5;
```

The expression between parentheses uses four stack slots. The result of `e5` always appears in `val(5)`, no matter which alternative matches when parsing (`a1 ... a2` or `b1 ... b4`).

This way of handling the value stack is similar to YACC's. Our `val(0)` is equivalent to YACC's `$$` and `$1`, `val(1)` is YACC's `$2`, etc.

Implicit primaries

The operators used to build the rules are overloaded so that they can accept values of the types taken by the primaries, and automatically convert them to the respective primary. The assignment operator of class `Rule` is similarly overloaded.

<code>std::string s</code>	converts to <code>Lit(s)</code>
<code>const char *s</code>	converts to <code>Lit(std::string(s))</code>
<code>char c</code>	converts to <code>Lit(c)</code>
<code>std::function<void()> f</code>	converts to <code>Do(f)</code>
<code>std::function<void(bool &)> f</code>	converts to <code>Pred(f)</code>

In many cases these automatic conversions allow replacing explicit calls to the primaries by their arguments. In the calculator we have eliminated explicit calls to `Lit()`, for example. It is not possible to eliminate calls to `Ccl(s)`, since a standalone string `s` converts to `Lit(s)`, not to `Ccl(s)`.

In the case of the binary operators, one of the arguments has to be an explicit expression; the other may convert automatically.

Calls to Lit(s), Lit(c) and Ccl(s) with literal arguments may also be replaced by user-defined literals:

"ab"_lit	is Lit(std::string("ab", 2))
'c'_lit	is Lit('c');
"a-z"_ccl	is Ccl(std::string("a-z", 3))

User-defined string literals allow embedded null characters, since the compiler passes the length of the literal to the operator. The operators are defined in inline namespace `peg::literals`.

Action and semantic predicate macros

Peg.h defines some macros for using capture-all-by-reference lambdas in embedded actions and semantic predicates with a very compact notation:

```
#define _(...) (peg::Do([&]{ __VA_ARGS__ })))
#define pa_(...) (peg::Pred([&](bool &){ __VA_ARGS__ })))
#define pr_(...) (peg::Pred([&](bool &r){ __r = [&]()->bool{ __VA_ARGS__ }(); })))
#define if_(...) (peg::Pred([&](bool &r){ __r = (__VA_ARGS__); })))
```

The `_(...)` macro defines a scheduled action. The argument is the body of a `void()` function.

The other macros define semantic predicates, which execute at parsing time.

The argument of `pa_(...)` is the body of a `void()` function. This predicate always succeeds.

The argument of `pr_(...)` is the body of a `bool()` function. This predicate succeeds if the function returns true, and fails otherwise.

The argument of `if_(...)` is a boolean expression. This predicate succeeds if the value of the expression is true, and fails otherwise.

For example, the following rules are equivalent and always fail:

```
Rule fail1, fail2;
fail1 = if_( false );
fail2 = pr_( return false; );
```

These macros expand to parenthesized expressions. If they are placed immediately following another expression, they attach to it, because the parentheses are interpreted as `operator()`. Otherwise, the parentheses are redundant and have no effect.

Attaching actions and semantic predicates is very useful because it allows adding or removing them anywhere in the rules without shifting the positions of the rule's expressions in the value stack.

Consider this simple palindrome recognizer as a use case for semantic predicates and parsing time actions:

```
#include <string>
#include <iostream>

#include "peg.h"

using namespace std;
using namespace peg;

class parser : public Parser<string>
{
    Rule start, pal, chr;

public:
    parser(istream &in = cin) : Parser(start, in)
    {
        start    = pal--                _( cout << text() << endl; )
        ;
        pal      = chr >> pal >> chr    if_( val(0) == val(2) )
        | chr >> chr                    if_( val(0) == val(1) )
        | chr
        ;
        chr      = Any()--              pa_( val(0) = text(); )
        ;
    }
};

int main()
{
    parser p;
    while ( p.parse() )
        p.accept();
}
```

The pal rule defines a palindrome as a symmetric string, including single characters. For lengths > 1 the symmetry is enforced by two if_(...) macros. Characters read from the input are placed in the value stack by the pa_(...) macro. In this parser the value stack is only used during parsing. At execution time, the _(...) macro in rule start outputs a direct text capture of each recognized palindrome, followed by a newline. Recognized palindromes are not necessarily the longest possible, especially when the input contains sequences of repeated characters.

Checking a grammar

If the macro PEG_DEBUG is defined before including peg.h, the rules are compiled with a method that allows checking the grammar. Checking is obtained by calling method check() at the end of the constructor, after the grammar is built.

Check() visits all grammar paths originating in the starting rule and throws an exception of type Rule::bad_rule if it detects uninitialized rules or potential left recursion, i.e. closed paths that may be traversed without consuming input.

It cannot be called more than once on the same grammar, which would be pointless anyway. Once the grammar is checked, both the macro and the call to check() can be removed.

Additionally, in debug mode each rule may be assigned a name (a const char *). For example,

```
calc.name = "calc";
```

In debug mode peg.h defines the macro peg_debug() as follows:

```
#define peg_debug(rule) rule.name = #rule
```

So in order to assign a debugging name to calc, this is the easiest way:

```
peg_debug(calc);
```

Rules that have names assigned will print debugging messages to standard error when visited by check().

For example, we define PEG_DEBUG in the calculator. After the grammar is built, before the end of the constructor, we write:

```
peg_debug(NUMBER);
peg_debug(LPAR);
peg_debug(RPAR);
peg_debug(ADD);
peg_debug(SUB);
peg_debug(MUL);
peg_debug(DIV);
peg_debug(calc);
peg_debug(expression);
peg_debug(term);
peg_debug(factor);

check();
```

When the parser is constructed, we get the following output on standard error:

```
calc
| expression
| | term
| | | factor
| | | | NUMBER
| | | | LPAR
| | | | expression (r)
| | | | RPAR
| | | | MUL
| | | | factor (v)
| | | | DIV
| | | | factor (v)
| | | | ADD
| | | | term (v)
| | | | SUB
| | | | term (v)
calc: check OK
```

The name of each rule is printed when visited, with an indentation that reflects the level of nesting. The syntax tree of each rule is visited only once. Rules that have already been visited are marked (v). Recursive calls are marked (r).

Now we introduce a small change in the grammar: in rule factor we make LPAR optional (~LPAR). This change makes the grammar left-recursive, because now the cycle expression → term → factor → expression can close without consuming input. Check() detects the problem:

```
calc
| expression
| | term
| | | factor
| | | | NUMBER
| | | | LPAR
| | | | expression (r)
terminate called after throwing an instance of 'peg::Rule::bad_rule'
what(): Left-recursive rule
```

Some low-level lexical rules like WS (white space) have not been debugged, because they would clutter the output of the debugger. Note that check() always checks the whole grammar, independently of which rules are debugged, if any. Normally check() is first called without debugging. If the grammar checks OK, nothing else need be done. If check() throws, some rules may be gradually added to the debugger, until the cause of the problem is clear.

Unicode support

Pegpp supports Unicode.

Strings and input streams should be encoded in UTF8. This is the default in Linux. In other environments, it may be necessary to prefix string literals containing non-ascii characters with the u8 prefix to force UTF8 encoding. The u (16 bit), U (32 bit) and L (wide char) prefixes cannot be used on string literals, since wide strings are not supported. This will produce a compilation error. Examples:

"aa"	(ok, ascii string)
u8"áá"	(ok, explicit UTF8: 0xc3, 0xa1, 0xc3, 0xa1, 0)
u8"\U000000e1\U000000e1"	(same, using code points)
"áá"	(ok in Linux)
U"áá"	(compilation error, wide strings not supported)

Valid string literals may be suffixed with _lit or _ccl as necessary.

Character values are 32 bits wide (`char32_t`) and can hold any Unicode code point. Ascii (7-bit) literals like `'a'` are automatically promoted to 32 bits.

Non-ascii character literals must be explicitly defined as `char32_t` by prepending them with the `U` prefix. Omitting the prefix is always incorrect. Depending on the context, it may generate compilation errors or warnings about using multichar constants, or just go undetected.

<code>'a'</code>	(ok, ascii value implicitly promoted to <code>char32_t</code>)
<code>U'a'</code>	(ok, explicit <code>char32_t</code>)
<code>U'\xe1'</code>	(same, using code point)
<code>U'a'_lit</code>	(ok, explicit <code>char32_t</code>)
<code>'á'</code>	(incorrect, possible compiler warning)
<code>'á'_lit</code>	(compilation error, operator <code>""_lit(int)</code> not defined)

Valid character literals accept the `_lit` suffix.

The parser classes

The parser classes are:

<code>BasicParser</code>	No value stack.
<code>Parser<T...></code>	Extends <code>BasicParser</code> adding a value stack with variant type elements.
<code>Parser<T></code>	Extends <code>BasicParser</code> adding a value stack with elements of type <code>T</code> .
<code>Parser<></code>	Same as <code>BasicParser</code> , no value stack. Provided for uniformity.

Value stacks are implemented by default with an auto-resizing vector. If the value stack is very sparse and the default constructor or copy constructor of a stack element is very costly, it might be more efficient to use a map (associative array). To try this option, `#define PEG_USE_MAP` before including `peg.h`.

This is the interface of the `BasicParser` class:

```
// Constructor
BasicParser(Rule &r, std::istream &in = std::cin);

// Parsing methods
bool parse();
void accept();
void clear();
std::string text() const;

// Grammar check
#ifdef PEG_DEBUG
void check() const;
#endif
```

The constructor takes one or two arguments. The first argument is a reference to the grammar's starting rule. The second (optional) argument is a reference to the input stream, by default standard input.

- Method `parse()` parses the starting rule.
- Method `accept()` executes the scheduled actions and discards matched input.
- Method `clear()` discards all input and all scheduled actions.
- Method `text()` returns a string with the contents of the most recently closed text capture.
- If `PEG_DEBUG` is defined, `check()` checks the grammar from the starting rule.

The `Parser<T...>` class adds the following interface for accessing the value stack:

```
element_type &val(std::size_t idx);
template <typename U> U &val(std::size_t idx);
```

where `element_type` is the type of the elements of the value stack, defined as `std::variant<T...>`.

Method `val(n)` returns a reference to the variant in slot `n` of the stack. Rules should return their values by assigning `val(0)`. For example, in a `Parser<int, std::string>` a rule could execute any of the following:

```
val(0) = val(2);    // assign type and value of slot 2 to slot 0
val(0) = 33;        // assign slot 0 type int and value 33
val(0) = "hello";    // assign slot 0 type string and value "hello"
```

When accessing the value contained in a slot, it is necessary to use the second form of `val()`, with an explicit type qualification: `val<T>(n)` returns a reference to the value of type `T` contained in the variant in slot `n` of the value stack, or throws `std::bad_variant_access` if the slot does not currently hold a value of type `T`.

For example, knowing that slot 3 contains a string and slot 2 contains an int, these are legal:

```
std::cout << val<std::string>(3) << std::endl;
val<int>(2) += 100;
```

The `Parser<T>` class has a value stack with elements of a single type `T`. It has the same interface as `Parser<T...>`, but now `element_type` is `T`, not a variant.

The `val<T>(n)` method is redundant and returns the same reference as `val(n)`. However, its use is encouraged for compatibility with the variant case. Incorrect invocations, like calling `val<std::string>(n)` on a `Parser<int>` do not throw, since no variants are used, but generate compilation errors.

The `Parser<>` class has no value stack and is just a `BasicParser`. When writing a parser it is handy to make it inherit from `Parser<>` initially and populate the empty types list as necessary when the grammar is developed. The first type in the type list has to be default constructible. Otherwise, start the type list with `std::monostate`.

A multi-typed value stack parser

This parser copies its input to its output, except when it finds embedded integers or sums of integers, which are replaced by their values. It is an artificial example, just for illustration.

```
#include <iostream>
#include <string>
#include "peg.h"

using namespace std;
using namespace peg;

class numsum : public Parser<int, string>
{
    Rule start, sum, other, number;

public:
    numsum(istream &in = cin) : Parser(start, in)
    {
        start = sum
            | other
            ;
        sum = number >> *(
            '+' >> number
            )
            ;
        number = ("0-9"_ccl)--
            | Any()--
            ;
    }
};

int main()
{
    numsum ns;
    while ( ns.parse() )
        ns.accept();
}
```

Input: aaa123+001+02bbb00044cc

Output: aaa126bbb44cc