



INVERSED TECH



geometry
RESEARCH

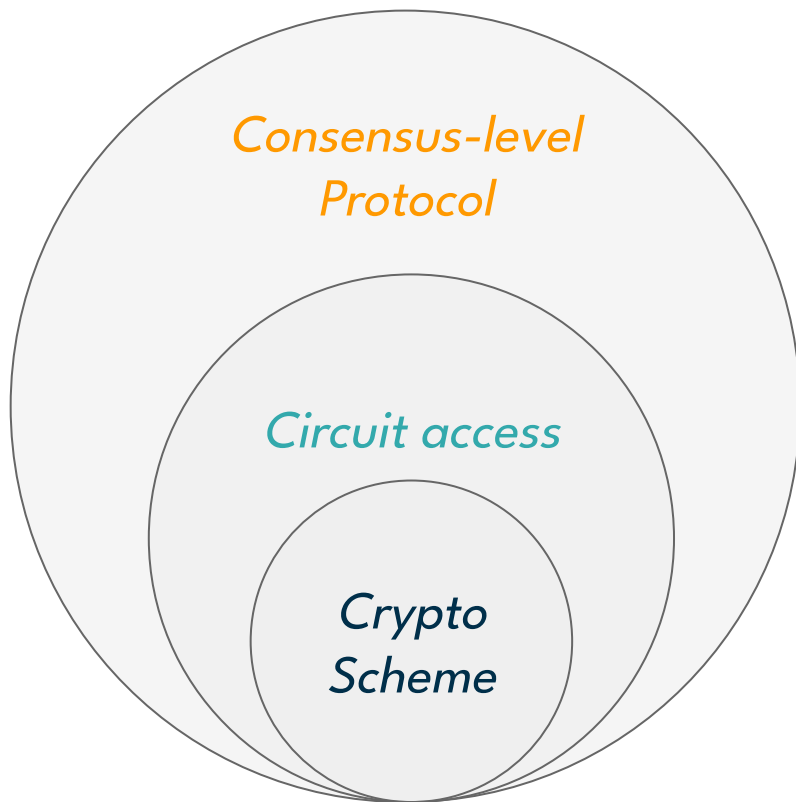
SoK: Programmable Privacy

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ZKProof6 Berlin | 2024

Additional collaborators:
Bryan Gillespie, Inversed Tech
Andrew Miller, UIUC

The programmability stack



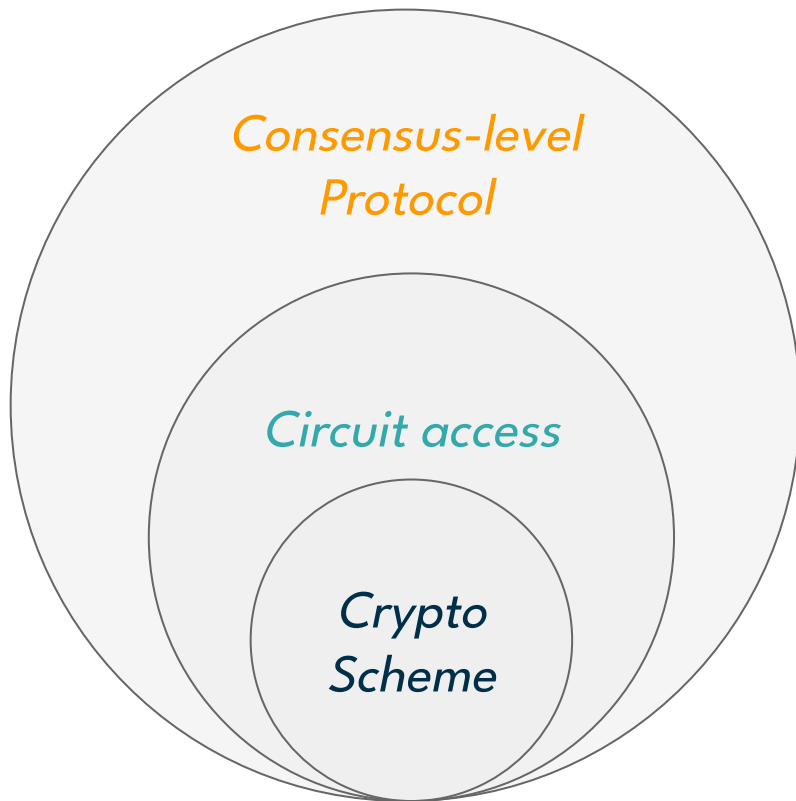
The programmability stack

e.g. *Penumbra*

aggregation
+ batch decryption

frequent batch auctions

homomorphic threshold
encryption



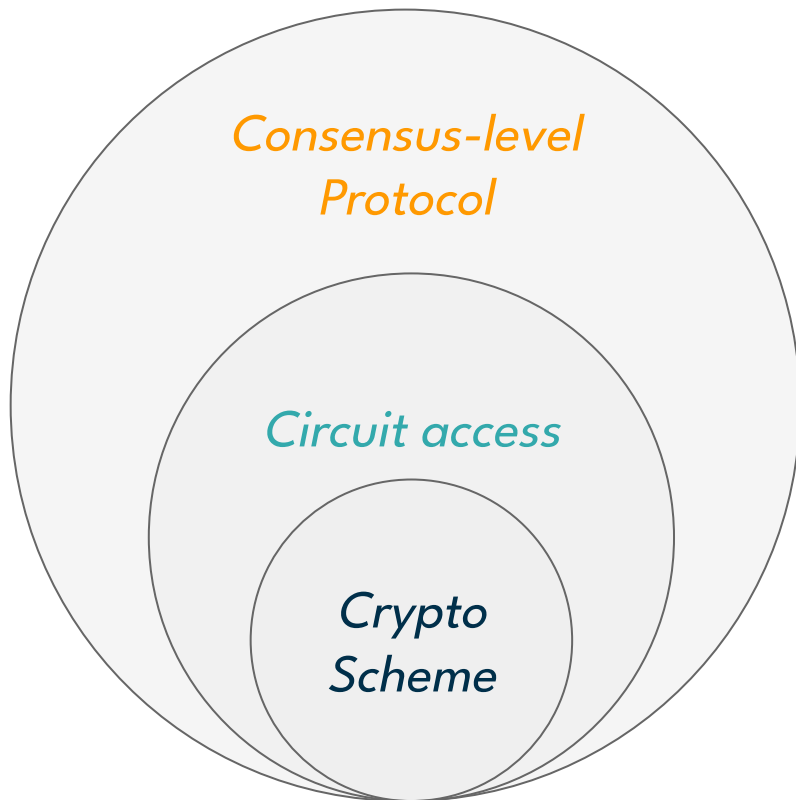
The programmability stack

e.g. *Penumbra*

aggregation
+ batch decryption

frequent batch auctions

homomorphic threshold
encryption



e.g. *renegade.fi*

proof verification
+ update balances

order-matching +
collaborative proving

two-party computation

Crypto Schemes



programmable disclosure / verification



programmable computation over private data

Consensus-level Protocol

Simple applications

- *individual private state*
- *sequential updates*
- *commutative updates*

(e.g. Zcash, voting, identity)

Interactive applications

- *shared private state*
- *atomic state updates*
- *non-commutative updates*

*(e.g. limit order auctions,
multilateral trade credit set-off)*

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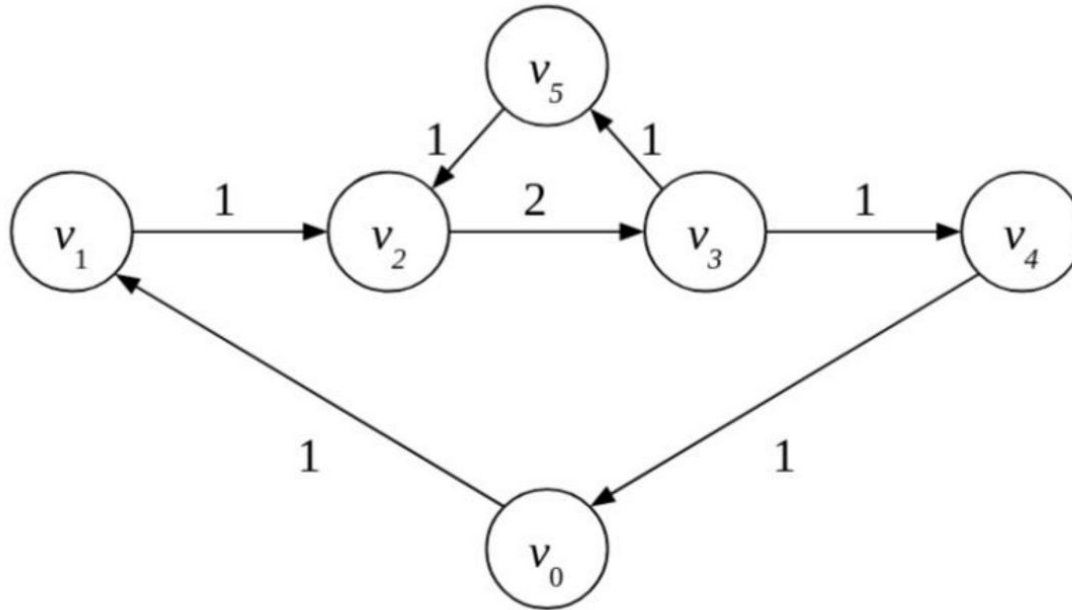
- *shared private state*
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- *non-commutative updates*

*(e.g. limit order auctions,
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**For many interesting applications, ZKP is not enough to
provide privacy!**

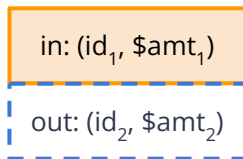
Motivating App: Multilateral trade credit set-off

Figure 21. Payment system with a chain and a cycle.



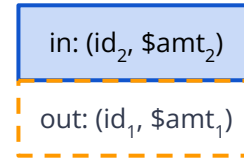
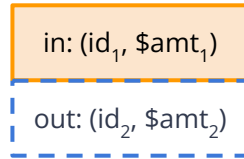
Motivating App: Limit Order Auctions

(order book)



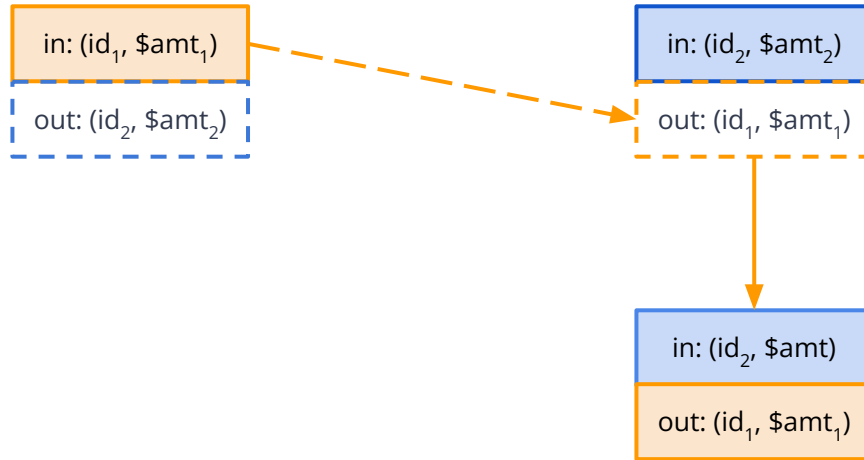
Motivating App: Limit Order Auctions

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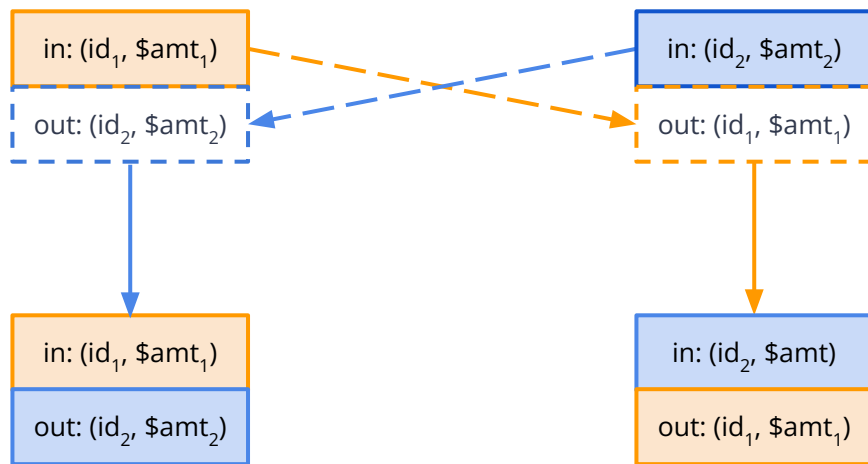
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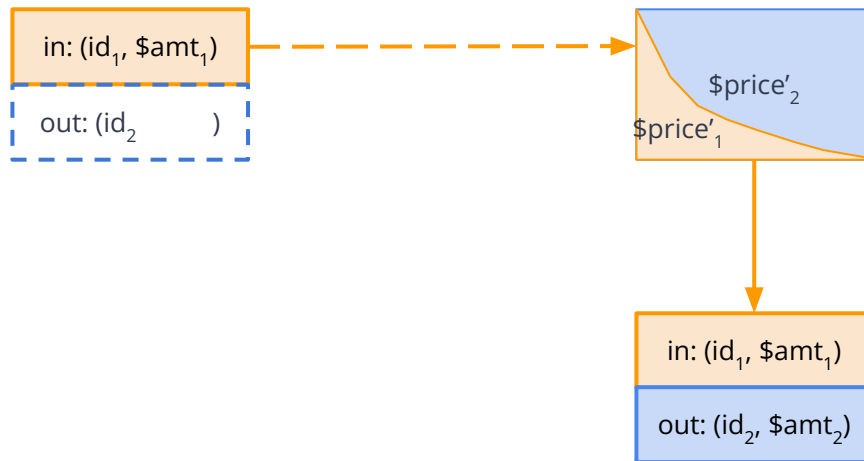
Motivating App: Limit Order Auctions

(CFMM)



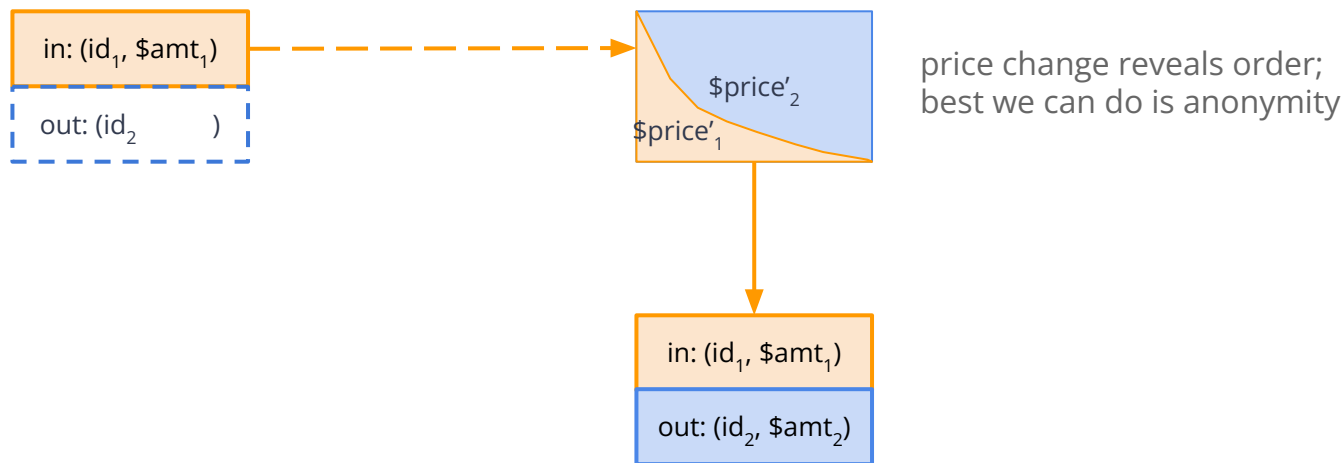
Motivating App: Limit Order Auctions

(CFMM)



Motivating App: Limit Order Auctions

(CFMM)



Three-Phase Computation Model



Phase 1: Independent computation

On input (Intent, Metadata) from party \mathcal{P} :

add (Intent, Metadata) to IntentSet

Leak $\mathcal{L}_{\text{Submit}}(\text{Intent}, \text{Metadata}, \text{CorruptionState})$ to \mathcal{A}

On input (Intent, Metadata) from \mathcal{A} :

add (Intent, Metadata) to IntentSet

Phase 2: Mediated computation

On input (Advice) from \mathcal{A} :

Result, NewTxs := MediatedComputation

assert ComputationAdviceIsValid(Result)

add elements of NewTxs to PendingTxs

If SatisfiesSecurityAssumptions(Corr

Leak $\mathcal{L}_{\text{Med}}^{\text{Expected}}(\text{Result}, \text{IntentSet}, \text{AppS}$

Else :

Leak $\mathcal{L}_{\text{Med}}^{\text{Broken}}(\text{Result}, \text{IntentSet}, \text{AppSta}$

Phase 3: Global computation

On input (Advice) from \mathcal{A} :

Result, Txs := SequenceTransactions(PendingTxs, AppState, Advice)

assert SequencingAdviceIsValid(Result, CorruptionState)

For Tx in Txs:

AppState = UpdateState(AppState, Tx)

remove Tx from PendingTxs

append Tx to ExecutedTxs

Output GlobalView(AppState) to global communication channel

For each party \mathcal{P} :

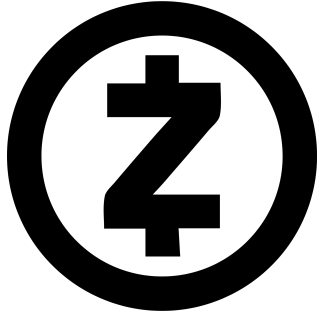
Output PrivatePartyView(AppState, \mathcal{P}) to \mathcal{P}

Leak PrivateAdversaryView(AppState) to \mathcal{A}

Mediated Computation in Privacy Protocols

<i>Programmability</i> <i>Privacy type</i>	General-purpose applications	Specialized functionality
k-of-N security	Sunscreen, Zama	Renegade, Penumbra
Hardware security	Secret Network, Obscuro/Ten, Oasis Sapphire, Phala Network, Automata	SUAVE
Extra-protocol	Aztec, Mina, Anoma, AlphaSwap (Aleo), Polygon Miden	ZSA Swap

Example Protocols



ZSA Swap



Renegade.fi



SecretSwap



Penumbra

Example 1: ZSA Swap

in: (id_1 , $\text{\$amt}_1$)

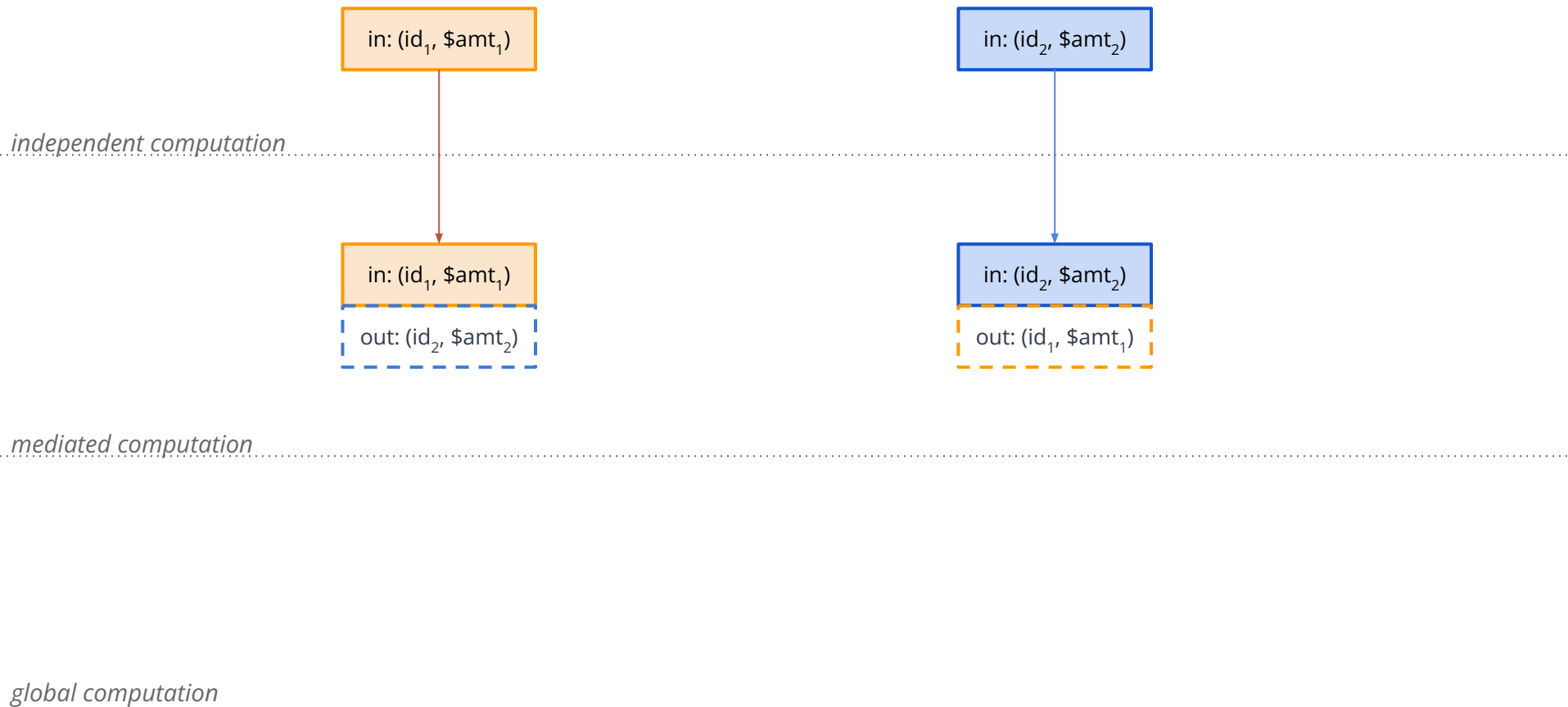
in: (id_2 , $\text{\$amt}_2$)

independent computation

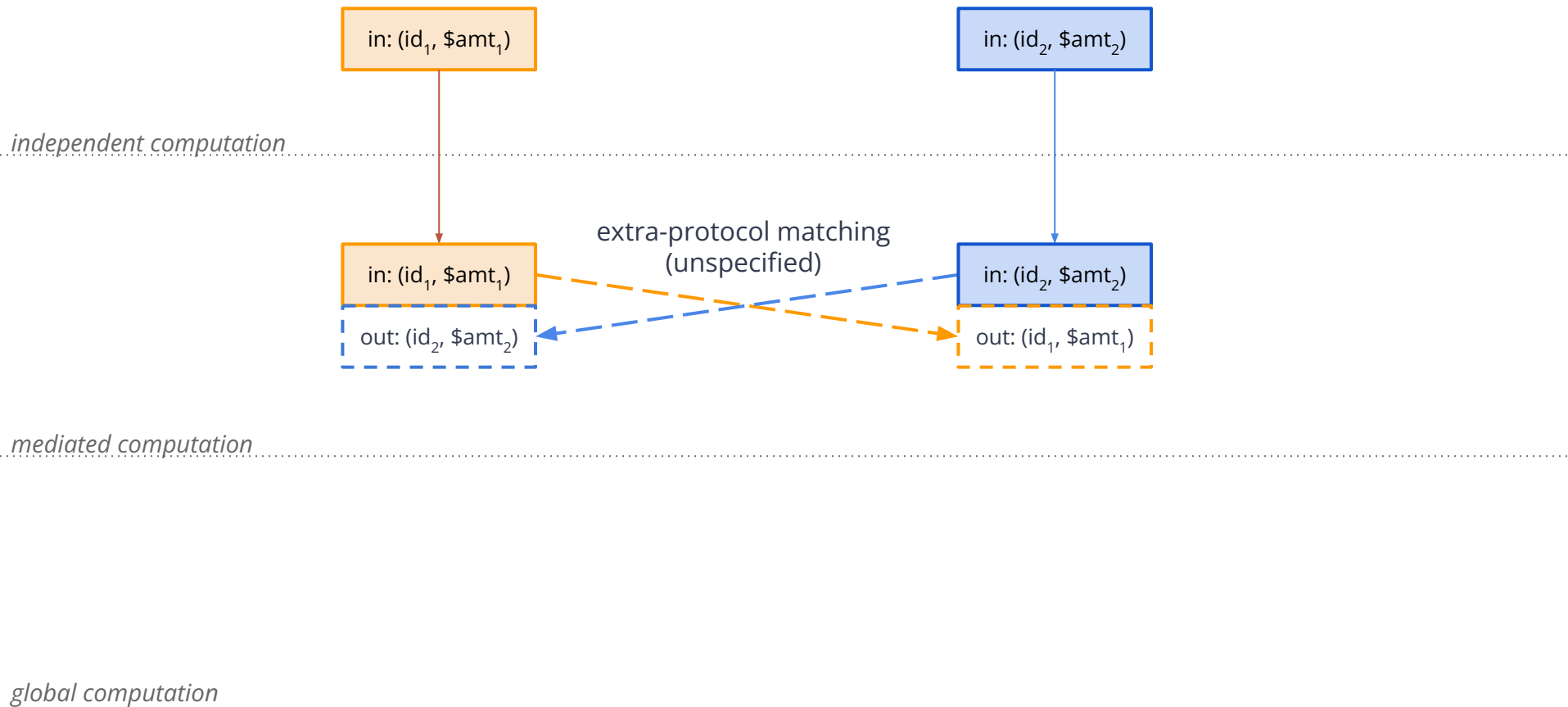
mediated computation

global computation

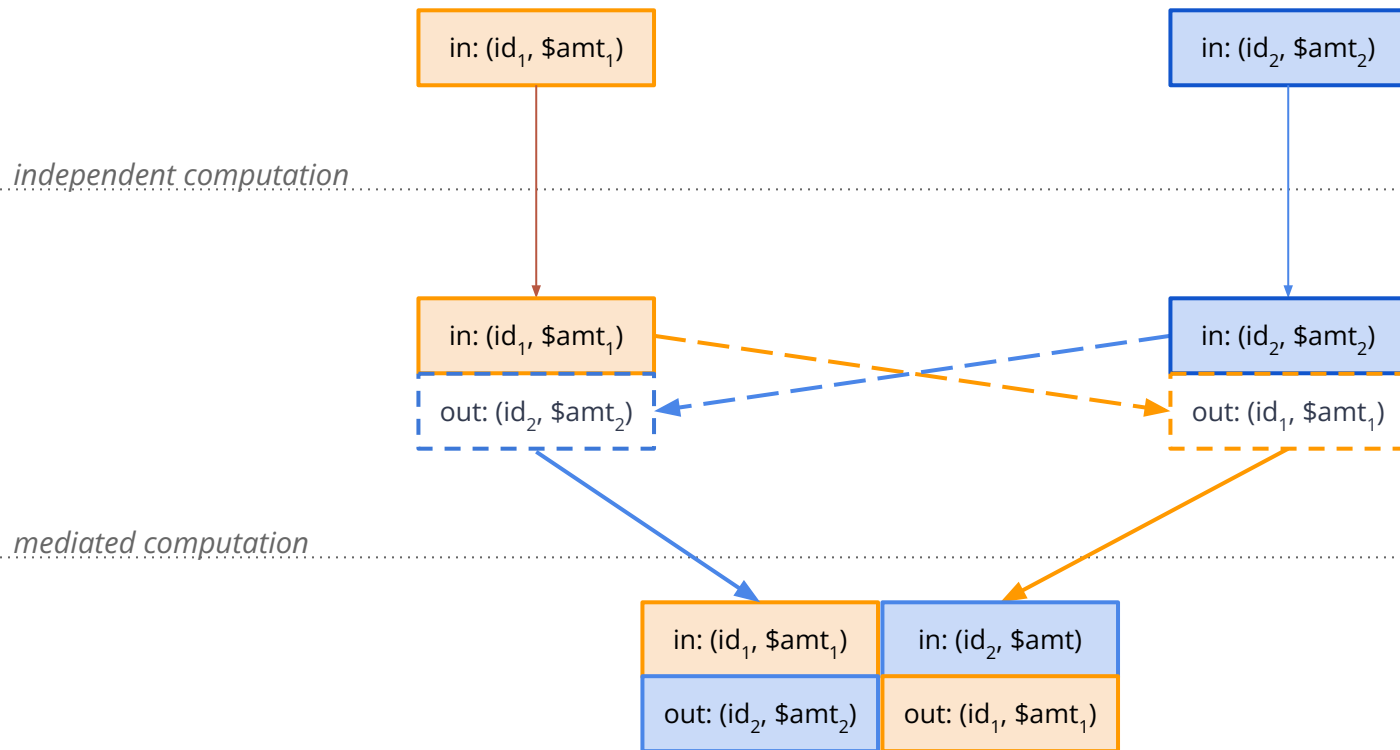
Example 1: ZSA Swap



Example 1: ZSA Swap



Example 1: ZSA Swap



global computation

Example 2: Renegade.fi

in: (id₁, \$amt₁)

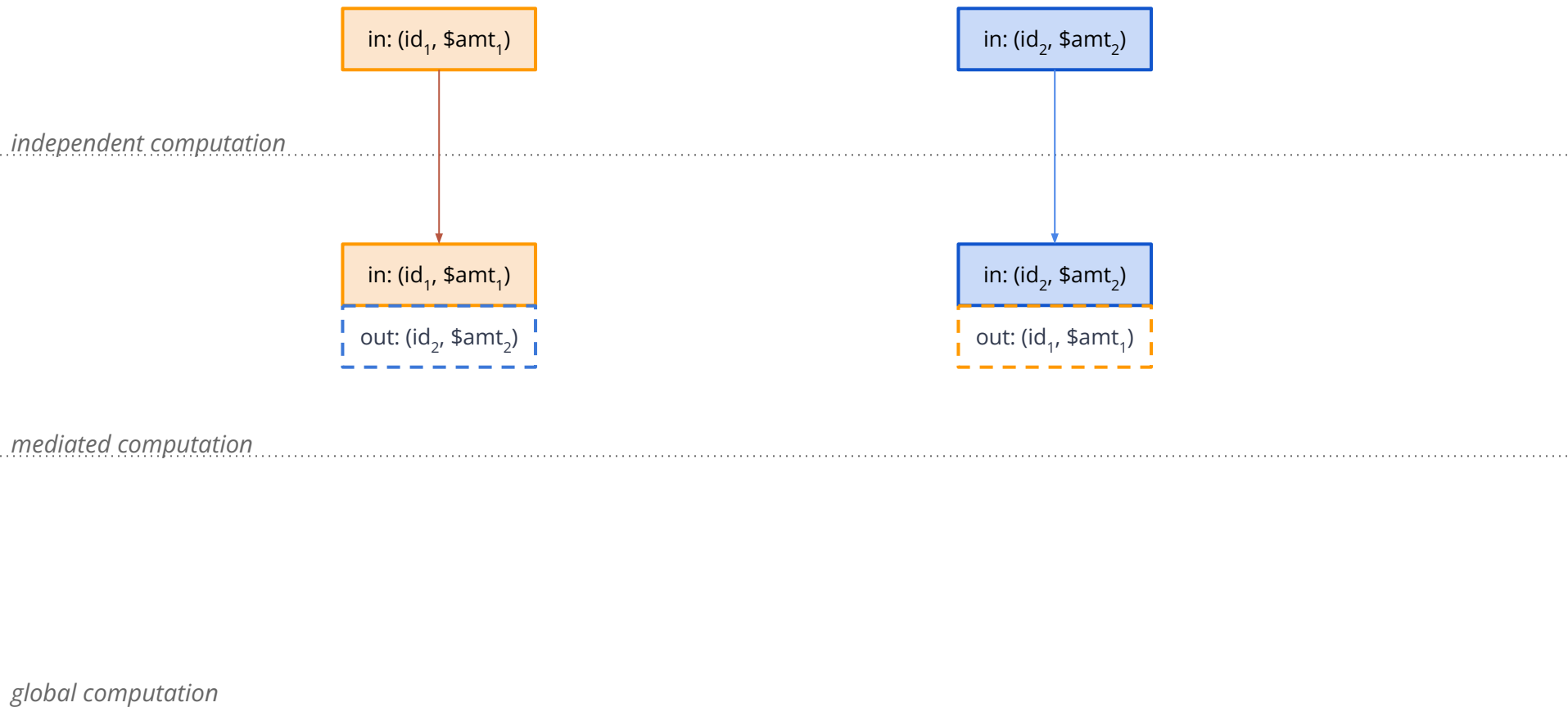
in: (id₂, \$amt₂)

independent computation

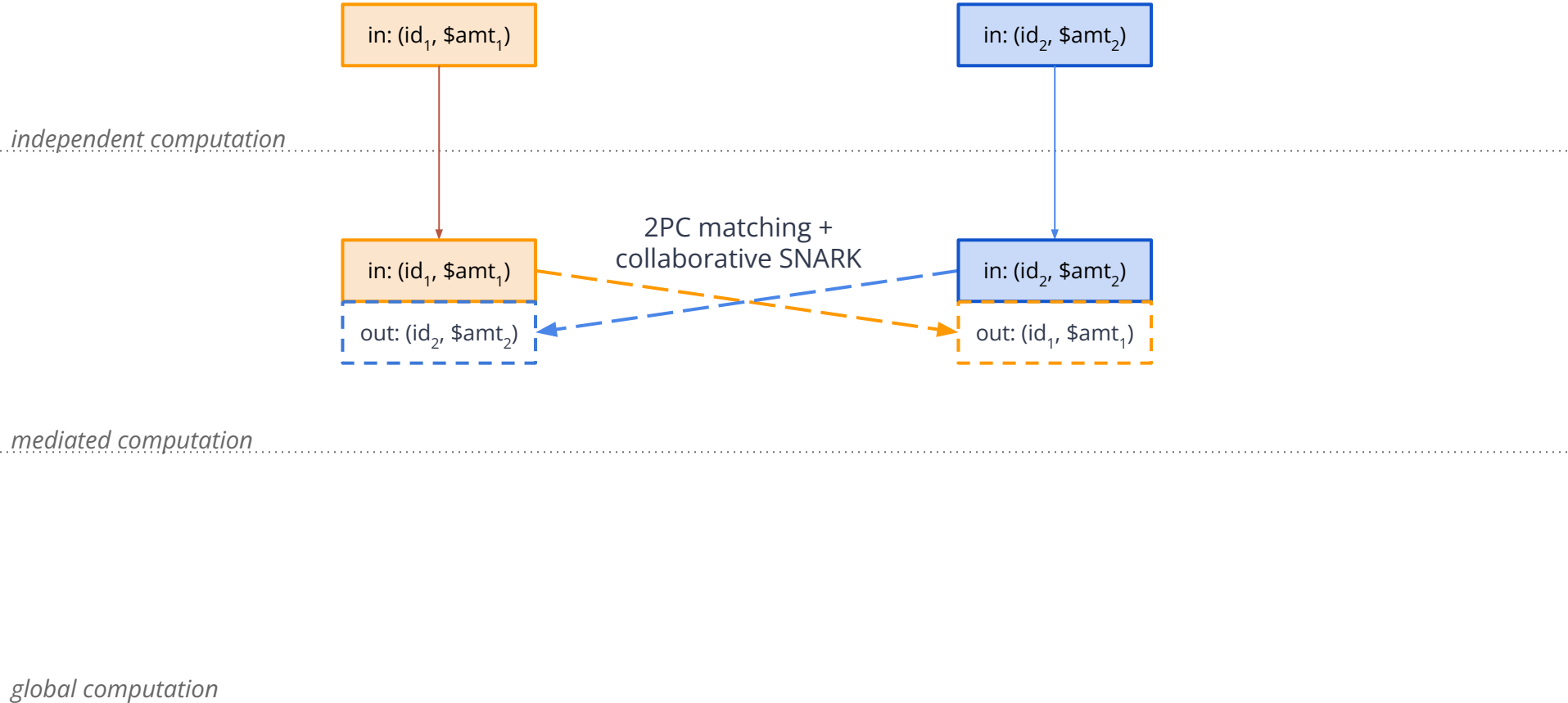
mediated computation

global computation

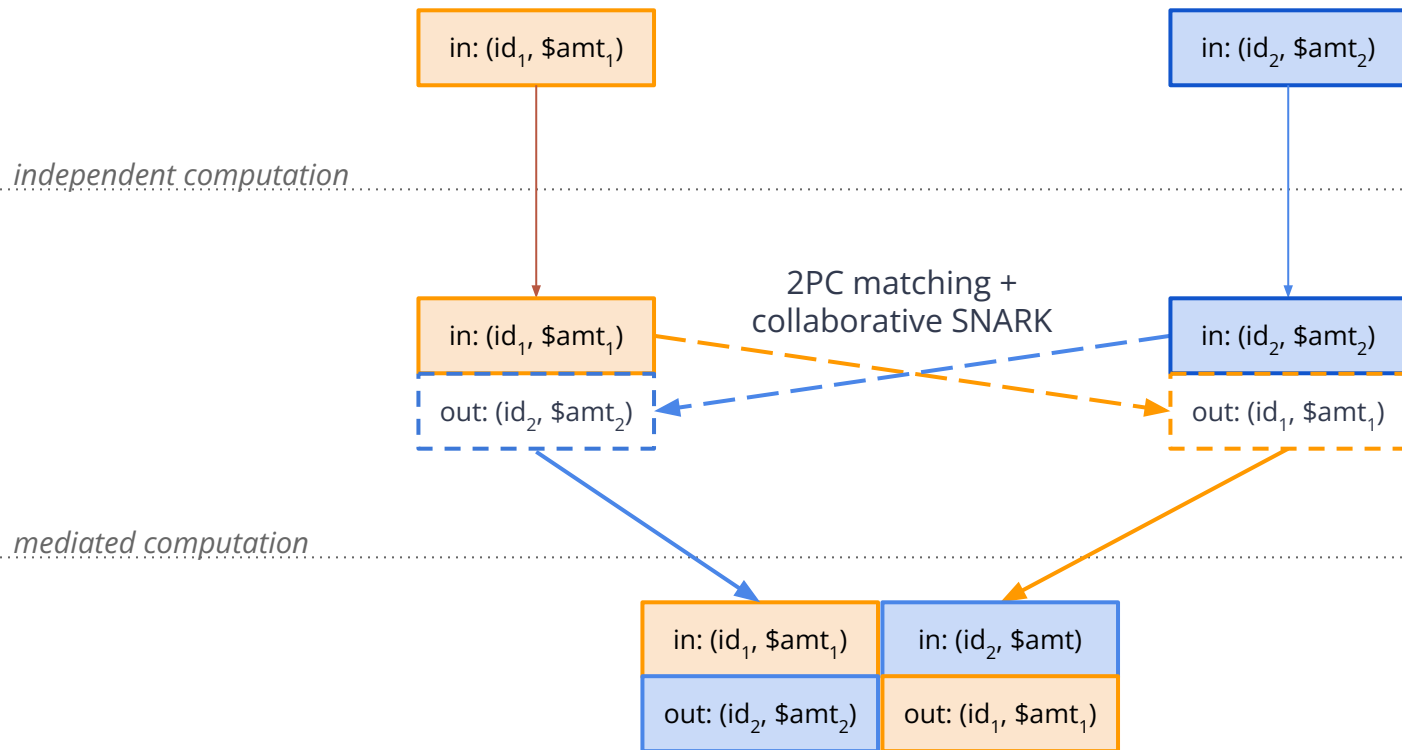
Example 2: Renegade.fi



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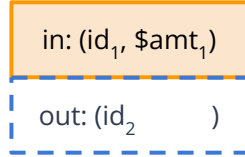


Example 2: Renegade.fi



global computation

Example 3: SecretSwap

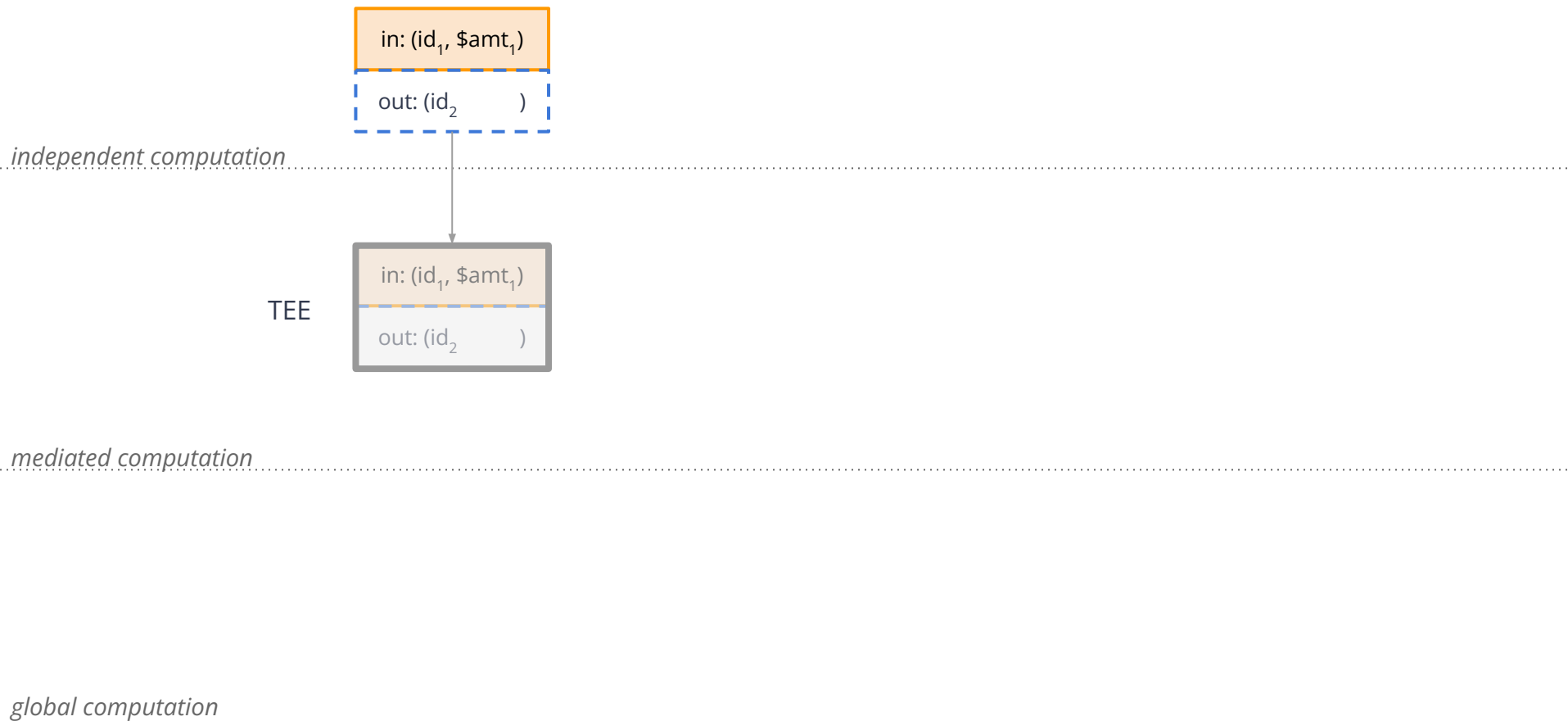


independent computation

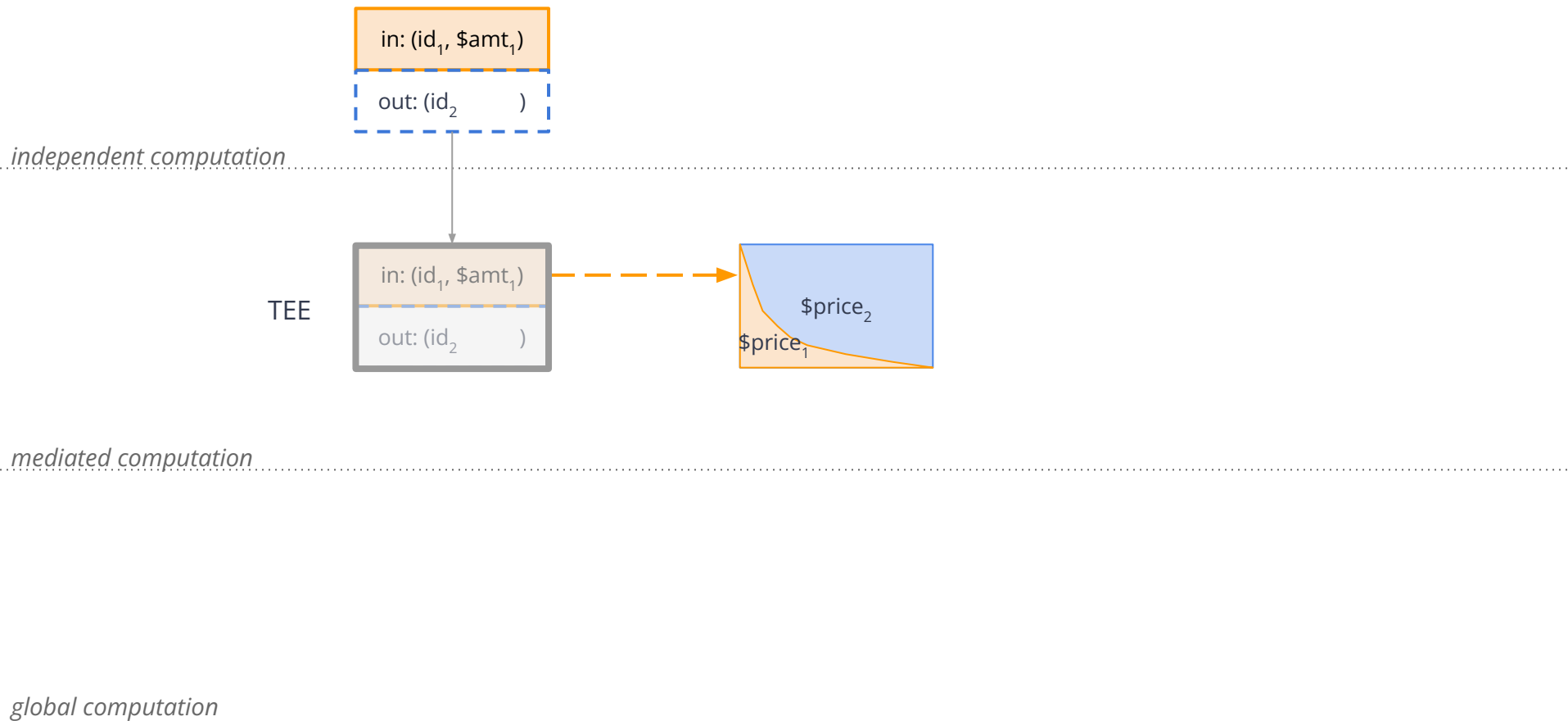
mediated computation

global computation

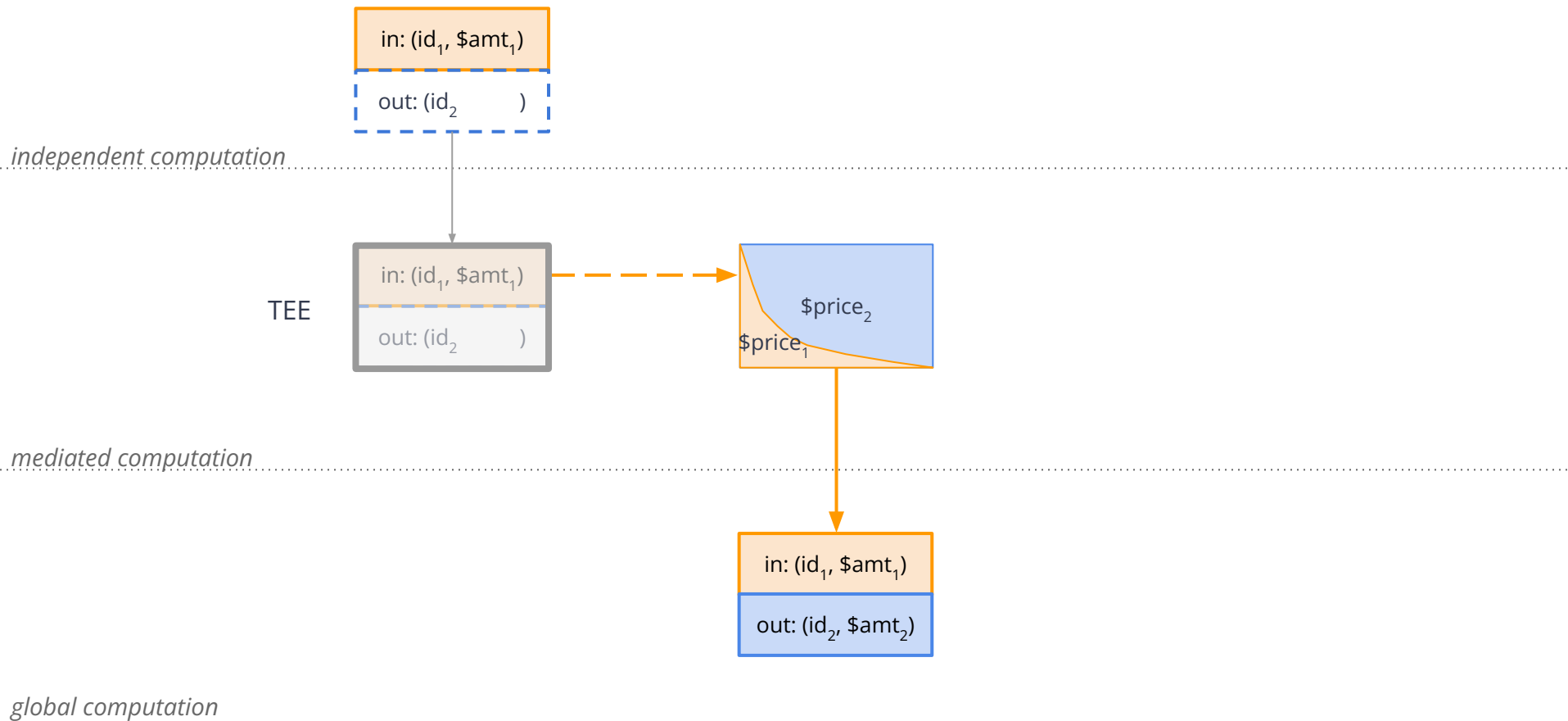
Example 3: SecretSwap



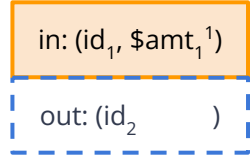
Example 3: SecretSwap



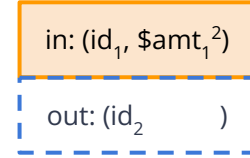
Example 3: SecretSwap



Example 4: Penumbra



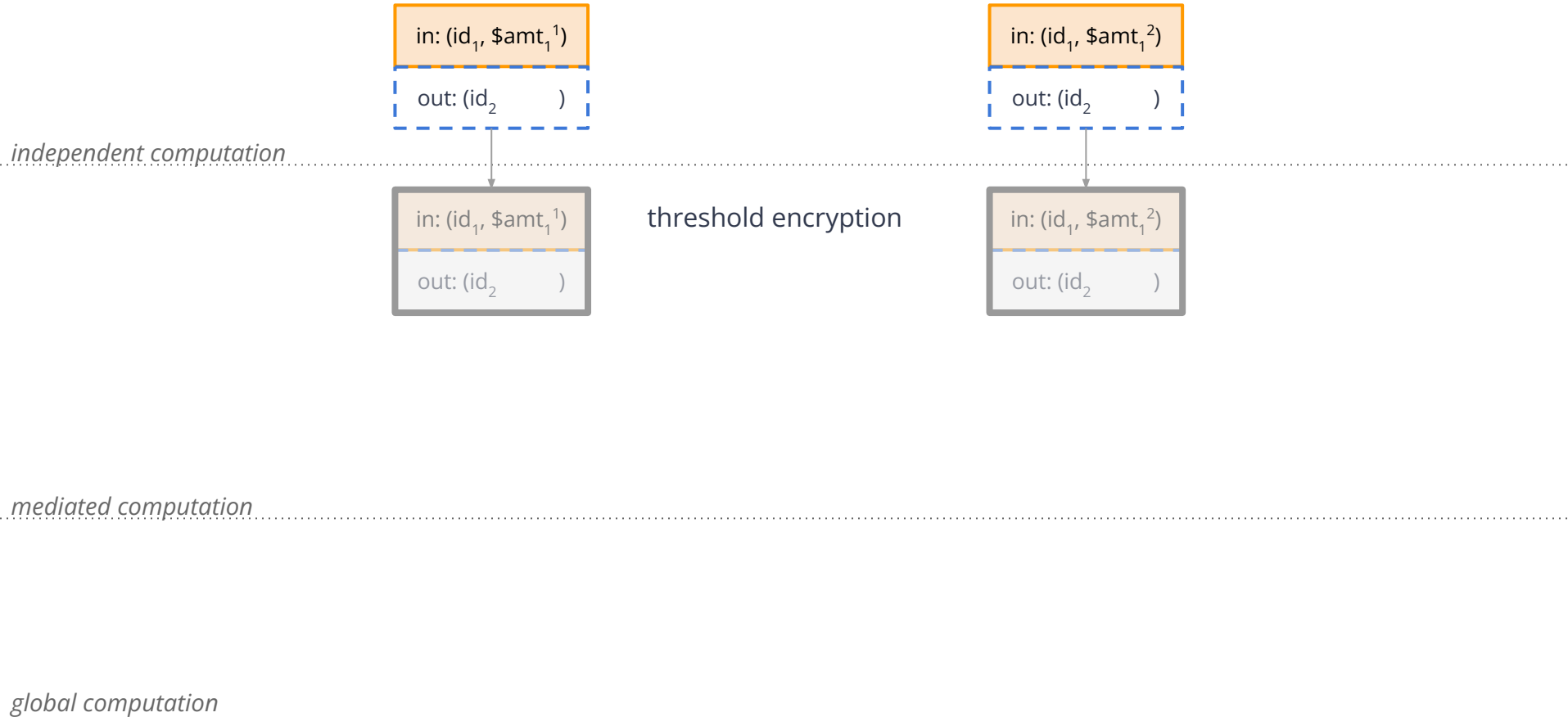
independent computation



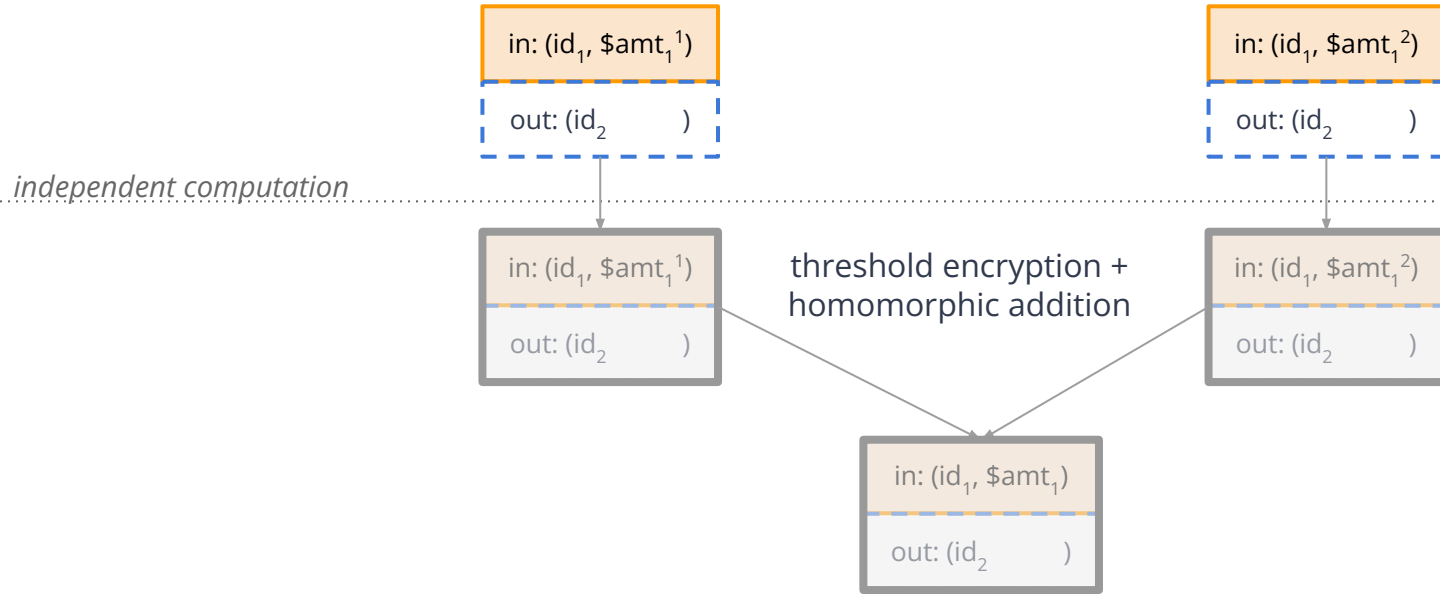
mediated computation

global computation

Example 4: Penumbra



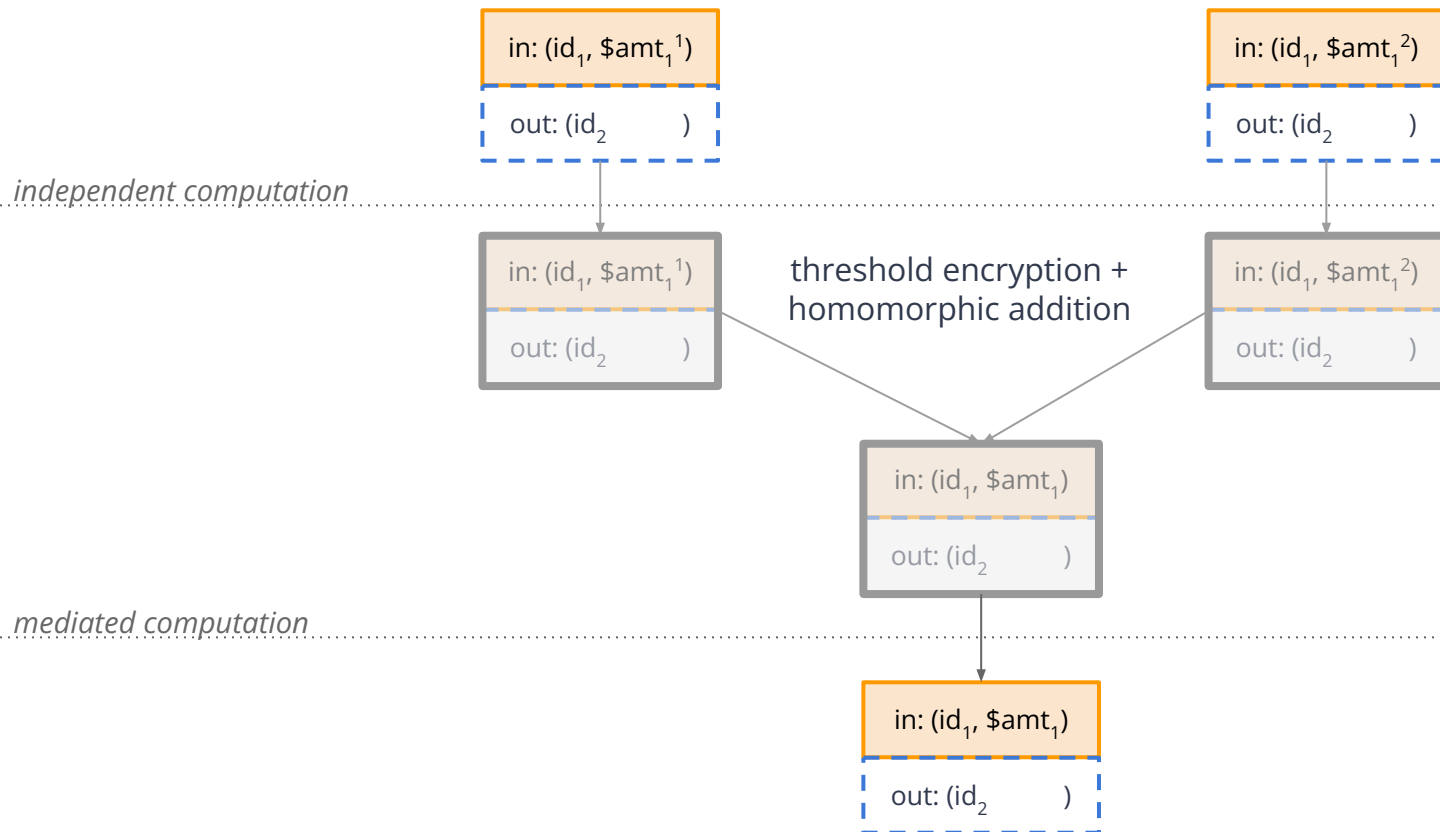
Example 4: Penumbra



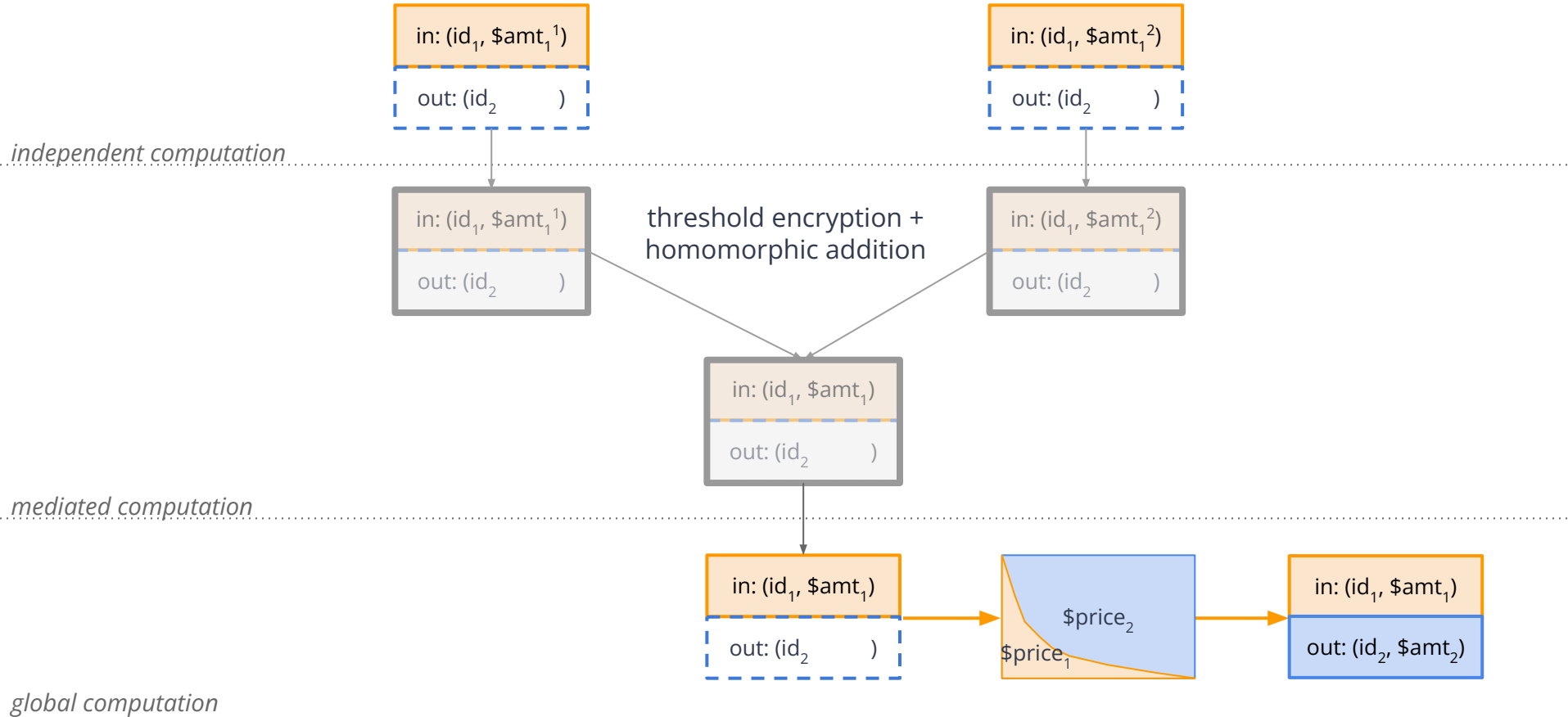
mediated computation

global computation

Example 4: Penumbra



Example 4: Penumbra



Tradeoff: Security vs. Expressivity

- **In-protocol** mediated computation \Rightarrow **resource-constrained**
 - e.g. Penumbra, threshold homomorphic encryption
 - e.g. Renegade.fi, 2PC's between relayers
- **Extra-protocol** mediated computation \Rightarrow **cheaper**, but
 - extra-protocol privacy guarantees fall on application designers
 - potentially fragments state / liquidity across different applications
- Decisions on trade-offs must be **informed by application**
 - e.g. how much order information to reveal is a tradeoff between price efficiency vs. frontrunning prevention
 - tension between enabling **general** applications vs. **optimising** for a single one

Future work / research questions

- **TODO: UC proofs for privacy properties**
 - [done] ideal functionality for three-phase epoch model
 - [done] ideal functionality and protocol descriptions for ZSA Swap, renegade.fi, SecretSwap, Penumbra
- **TODO: Fair metrics**
 - **efficiency**: for a fixed program (e.g. order matching), how much work does a consensus node need to do?
 - **security**: economic cost of attack?
- **TODO: Best practices**
 - defense-in-depth: collaborative proofs; proof of encryption; “remote attestation” zk proof, using TEEs to store key shares in MPC / FHE
 - DSL/compiler safeguards: e.g. prevent writing disclosive MPC circuits

Future work / research questions

TODO: design optimal programmable privacy protocol that fits

1. Fits the three-phase computation model
2. Enables expressiveness for user defined programs while ensuring security of application
3. Provides cryptographic-level privacy

POTENTIAL DIRECTION

- Include different cryptographic schemes as the “engines” for each phase (e.g.: ZK for phase 1 + FHE or MPC for phase 2)
- Build a combined DSL that knows to speak to both engines with a fixed separation as soon as inputs need to be aggregated
- Define privacy **invariants** at the language level to prevent unintentional leakage
- Provide high-level primitive implementations for users to use off-the-shelf, increasing security

THANK YOU QUESTIONS?



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