

Ligetron and the Llama Inference

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Thanks Stealth Software Tech.

ZK evolution

 At the Beginning: Reduce to an NP-complete problem (eg, Graph Hamiltonicity, Graph 3-Coloring)

Then: Boolean and Arithmetic circuits

Since Blockchains: R1CS, Circom, Cairo, Gnark, etc...

Today: zkVMs

Instrumenting ZK today - zkVMs

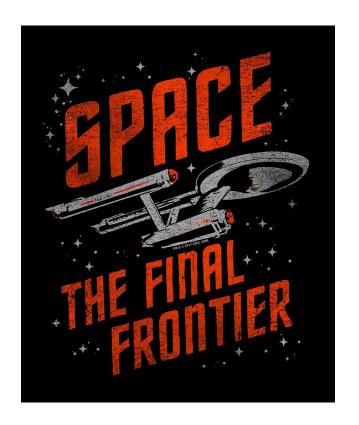
- 1. Code application in a high-level language (C,C++,Rust)
- 2. Compile down to popular VMs. Eg, RISC-V, WASM
- Prove correct execution of VM

Challenges in Scaling ZK

First, it was (circuit) representation

Then, optimizing prover running time

Today, optimize prover memory



Main Question

Do there exist time and space preserving ZK-SNARKs from minimal assumptions?

YES* - Based on hash functions [BBHV 2022] Prover time $\tilde{O}(T(n))$ and space $\tilde{O}(S(n))$. Proof Length $\tilde{O}(T(n)/S(n))$

Today: Concretely Efficient!

Introducing Ligetron

A Time and Space Efficient ZK-SNARK

Key Ingredients

- 1. WASM as an intermediate representation Ligetron is a zkWASM
- 2. A space-efficient variant of the Ligero ZKP [Ames et al, 2017]

Ligetron

Post quantum

Black-box

Non-interactive Yes

Succinct proof Sublinear

Succinct verification No*

Time efficient Yes (quasilinear)

Space efficient Yes (linear)

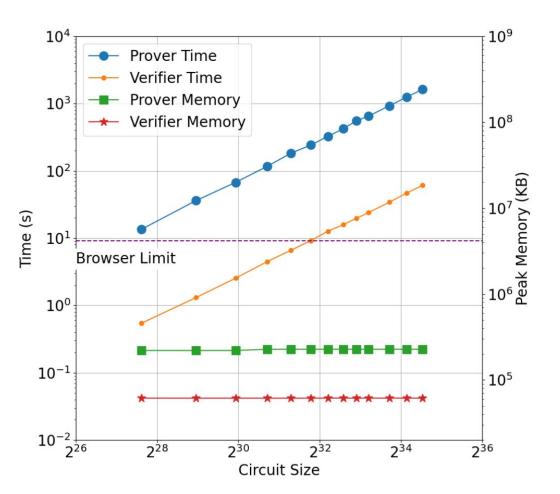
Implementable Yes (runs from a browser)

Why WASM?

- Space/time efficiency from operational semantics of WASM (stacks, blocks)
- Few set of instructions to compile
- Numerous compilers from high-level languages to WASM (C,C++,Rust,etc)
- Clean sandbox (no I/Os, no system calls, no malloc*)
- Extern function calls for gadgets / pre-compiles

Performance

- Prover time: 500 ns/g
- Verifier time: 250 ns/g
- ❖ Prover memory: < 100 MB</p>
- Verifier memory: <10 MB</p>



Performance

	Prover Time Batched	Verifier Time Batched	Prover Time Random	Verifier Time Random
Browser	1.2 us/g	20 ns/g	5 us/g	2.5 us/g
Macbook	65 ns/g	3 ns/g	500 ns/g	250 ns/g

Where we are?

Proof Length

- Unstructured circuit: Square root circuit size. But, parameterizable
- Structured circuit (M copies of size T circuit): O(M+k.T)

Verification

- Unstructured circuit: Quasilinear in circuit size
- Structured circuit: O(M+k.T)

Extent of WASM integration

- All 32-bit and 64-bit integer operations
- Oblivious control flow
- More recently, RAM, 32-bit floating operations

Are ZKVMs Scalable?

Jolt Risco SP1

VM	RISC-V	RISC-V	RISC-V
Prover Speed	150 kHz int32	25 kHz int32	40-150 kHz int32
Succinct Proof/Verification	○ Yes/Yes	Yes/Yes	○ Yes/Yes
Hardware	64 Core 512 GB RAM	64 Core 512 GB RAM	64 Core 512 GB RAM
Memory Efficient	O No	○ No	O No

Are ZKVMs Scalable?

Current Techniques for Memory Efficiency

- Recursive Composition
- Incrementally Verifiable Computation (IVC)
- 1. Break computation into small units
- 2. Prove units in distributed manner
- 3. Compose

Bottlenecks

- Need a lot of hardware (each unit needs to be run in parallel)
- Gadgets/pre-compiles are harder in RISC-V
- zkVMs today are really succinctVMs

Benchmark 1 (Structured Circuits): Rollups

Simple payment circuit

- EdDSA signatures
- Poseidon hash function for Merkle Trees





Starknet: 200 nodes with 400MB RAM

Benchmark: 2000 tps*

Ligetron: g5.xlarge (single A10 GPU)

Benchmark: 500tps

WASM extern for 256-bit

```
163
         fp256_class(int i) {_fp256_init(data_); _fp256_set_ui(data_, i); }
164
         fp256_class(uint32_t i) { _fp256_init(data_); _fp256_set_ui(data_, i); }
         fp256_class(const char *str, int base = 10)
165
166
             { _fp256_init(data_); _fp256_set_str(data_, str, base); }
167
         fp256_class(const fp256_t o)
168
169
             { _fp256_init(data_); _fp256_set_fp256(data_, o); }
170
         fp256_class(const fp256_class& o)
171
             { _fp256_init(data_); _fp256_set_fp256(data_, o.data_); }
172
173
         // fp256 class(fp256 class&& o)
                { _fp256_init(data_); std::swap(*data_, *o.data_); }
174
175
         fp256_class& operator=(const fp256_t o)
176
177
             { _fp256_set_fp256(data_, o); return *this; }
178
         fp256_class& operator=(const fp256_class& o)
             { _fp256_set_fp256(data_, o.data_); return *this; }
179
         // fp256_class& operator=(fp256_class&& o)
180
                { std::swap(*data_, *o.data_); return *this; }
181
182
183
         ~fp256_class() { _fp256_clear(data_); }
184
185
         __fp256_backend* data() { return data_; }
186
         const __fp256_backend* data() const { return data_; }
187
188
     protected:
189
         fp256_t data_;
190
     };
191
```

struct fp256_class {

fp256_class() { _fp256_init(data_); }

160 161

162

WASM extern for 256-bit

Implementations:

- EdDSA signatures
- Poseidon hash function

```
// Vector Initialization
     WASM EXTERN(ligetron-batch, fp256vec get size)
199
     uint64 t fp256vec get size();
200
201
     WASM_EXTERN(ligetron-batch, fp256vec_init)
202
     void _fp256vec_init(fp256vec_t v);
203
     WASM_EXTERN(ligetron-batch, fp256vec_clear)
204
205
     void _fp256vec_clear(fp256vec_t v);
206
     WASM_EXTERN(ligetron-batch, fp256vec_set_ui)
207
208
     void _fp256vec_set_ui(fp256vec_t v, uint32_t* num, uint64_t len);
209
     WASM_EXTERN(ligetron-batch, fp256vec_set_ui_scalar)
210
     void _fp256vec_set_ui(fp256vec_t v, uint32_t num);
211
212
213
     WASM_EXTERN(ligetron-batch, fp256vec_set_str)
     int _fp256vec_set_str(fp256vec_t v, const char *str[], uint64_t len, int base = 0);
214
215
216
     WASM_EXTERN(ligetron-batch, fp256vec_set_str_scalar)
     int _fp256vec_set_str(fp256vec_t v, const char *str, int base = 0);
217
218
     WASM_EXTERN(ligetron-batch, fp256vec_copy)
219
     void _fp256vec_copy(fp256vec_t out, const fp256vec_t in);
220
221
222
     WASM_EXTERN(ligetron-batch, fp256vec_print)
     void _fp256vec_print(const fp256vec_t v);
```

Benchmark 2 ("Unstructured"): LLM Inference



Modulus Labs: 200+ hours for verifying GPT2-XL 1.5B using 128-core 1TB CPU with 10TB disk space



Ligetron: 14 hours for verifying the Llama 7B using 8-core CPU with peak memory 10GB*



Our Approach for Llama inference



Pure C implementation of Llama 7B

https://github.com/karpathy/llama2.c

llama2.c

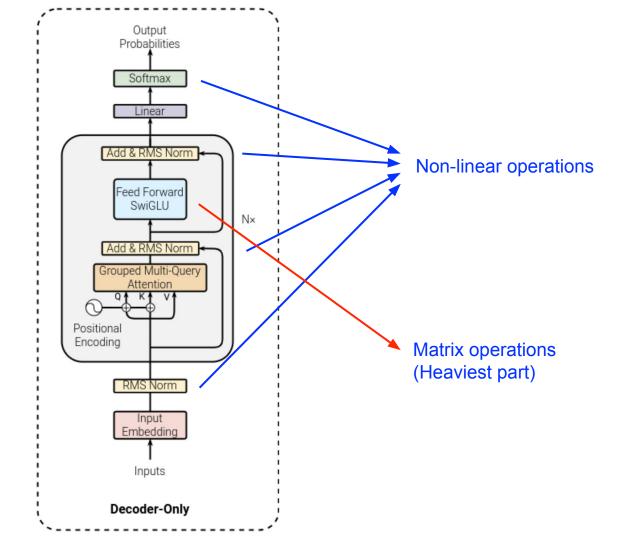
Main Challenge floating point ops

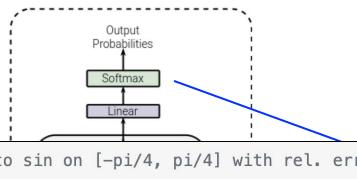


Have you ever wanted to inference a baby Llama 2 model in pure C? No? Well, now you can!

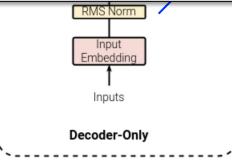
Train the Llama 2 LLM architecture in PyTorch then inference it with one simple 700-line C file (run.c). You might think that you need many billion parameter LLMs to do anything useful, but in fact very small LLMs can have surprisingly strong performance if you make the domain narrow enough (ref: TinyStories paper). This repo is a "fullstack" train + inference solution for Llama 2 LLM, with focus on minimalism and simplicity.

Llama 7B





```
lama 7B
/* minimax approximation to sin on [-pi/4, pi/4] with rel. err. ~= 5.5e−12 */
                                                                                   bns
double sin core (double x)
  double x4, x2, t;
  x2 = x * x;
  x4 = x2 * x2;
  /* evaluate polynomial using a mix of Estrin's and Horner's scheme */
  return ((2.7181216275479732e-6 * x2 - 1.9839312269456257e-4) * x4 +
           (8.3333293048425631e-3 * x2 - 1.6666666640797048e-1)) * x2 * x + x;
(eg, sin, exp, etc)
                                    RMS Norm
                                      Input
                                    Embedding
```



Llama 7B

Operation distribution

fp32 add (100 million)

~1000 constraints

fp32 mul (200 million)

~200 constraints

Total constraints

~130 billion add

~78 billion mul

Num Linear constraints: 125591511068

Num quadratic constraints: 78174961265

Num quadratic constraints (padded): 78174968624

f32.fnn_le: 32001

f32.convert_i32_u: 1

f32.fnn_ne: 1

f32.convert_i32_s: 103235588

f32.const: 6875505

i32.reinterpret_f32: 749634

i32.trunc_f32_u: 385280

f32.fnn_abs: 1499136

i32.trunc_f32_s: 749568

f32.reinterpret_i32: 385346

f32.fnn_div: 1178626

f32.fnn_mul: 209896969

f32.fnn add: 104927810

f32.fnn_gt: 401281

f32.fnn_lt: 2285694

f32.fnn_sub: 164097

f32.fnn_ge: 385281

Benchmark 2 ("Unstructured"): LLM Inf.



Modulus Labs: 200+ hours for verifying GPT2-XL 1.5B using 128-core 1TB CPU with 10TB disk space



Ligetron: 14 hours for verifying the Llama 7B using 8-core CPU with peak mem. 10GB

Time to benchmark: 15 days



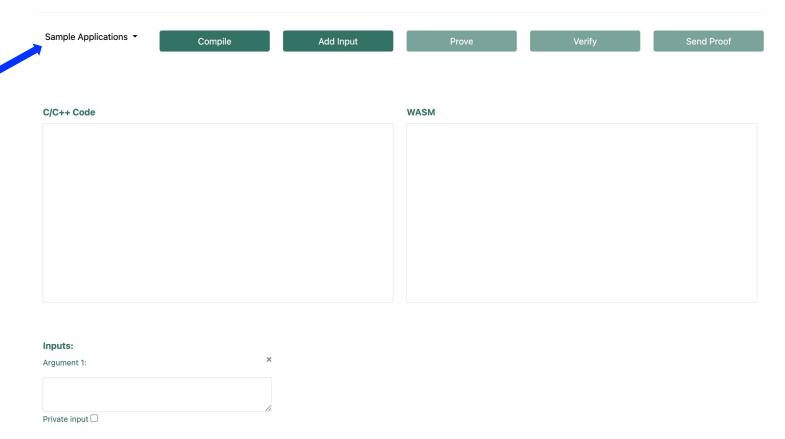
Ligetron 1.0 Compared to Other zkVMs



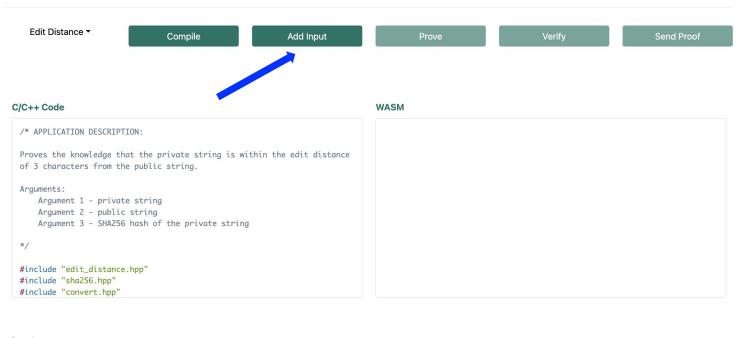
	Ligertron 1.0	O Jolt	Risc0	○ SP1
VM	● WASM	RISC-V	○ RISC-V	○ RISC-V
Prover Speed	● 50 kHz int32, int64, fp32	150 kHz int32	25 kHz int32	40-150 kHz int32
Succinct Proof/Verification	Yes/No	Yes/Yes	O Yes/Yes	O Yes/Yes
Hardware	8 Core 16GM RAM	64 Core 512 GB RAM	64 Core 512 GB RAM	64 Core 512 GB RAM
Memory Efficient	Yes	○ No	○ No	○ No

Demo

Build a ZK Application



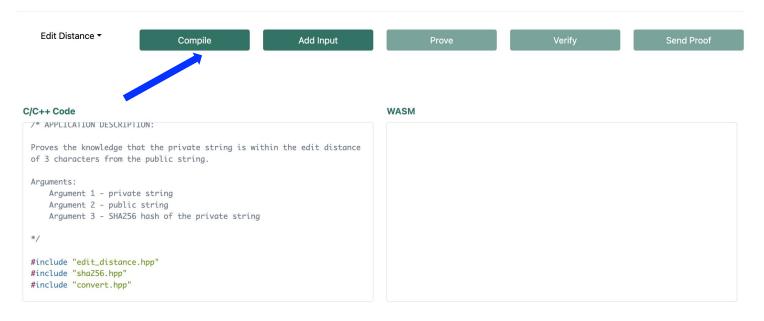
Build a ZK Application



Inputs:

Argument 1:

Build a ZK Application



Inputs:

Argument 1: X

abcd

Private input

Argument 2: X

abcdef

Build a ZK Application



Inputs:

Argument 1:

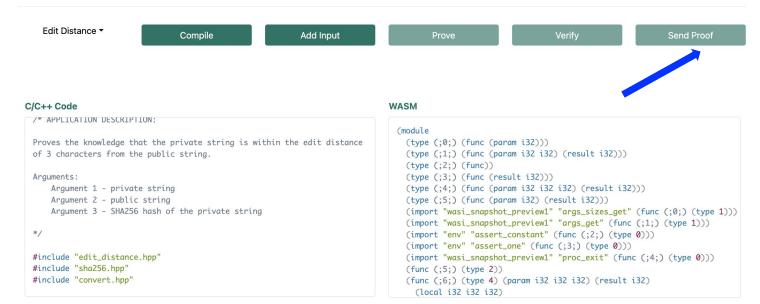
abcd

Private input

Argument 2:

abcdef

Build a ZK Application



abcdef

Results:

Prove result: true

Prove execution time: 5.61 s

Circuit size: 425,015

Verify result: true

Verify execution time: 2.71 s

Our Roadmap

- 1. You don't have succinct verification, duh!

 Pre-processing Ligetron for succinct verification
- 2. But your proofs are not short, duh!
 Compose with SNARKs to verify on-chain (Groth16/Halo2)
- 3. But you can't handle non-oblivious code, duh!
- https://eprint.iacr.org/2023/1257 (CCS 2023 Distinguished Paper Award)
- https://eprint.iacr.org/2024/456 (Precompiles)
- 4. What about lookup arguments? We can do that as well



Develop on Ligetron:

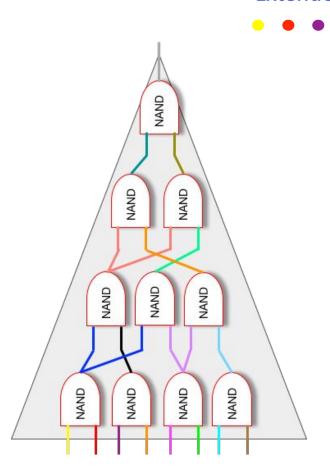
Build your ZK app at ligetron.com

Develop with Ligetron:

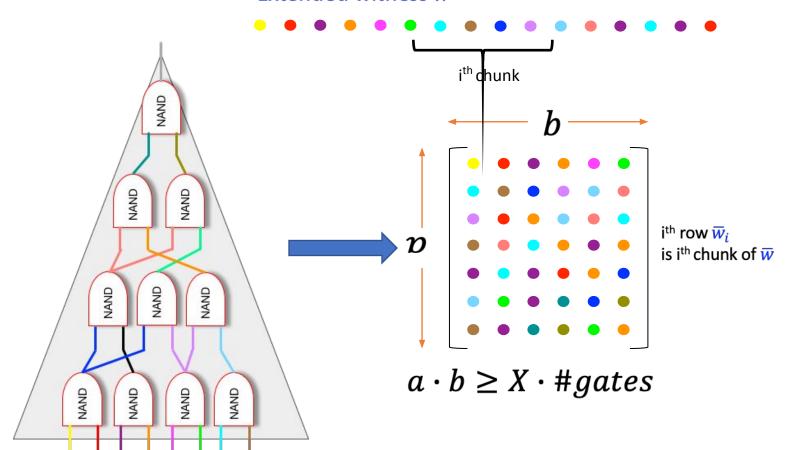
We are looking for developers Please email muthu@ligero-inc.com

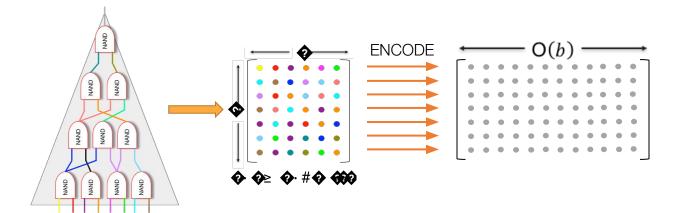
Our Techniques

Extended witness \overline{w}



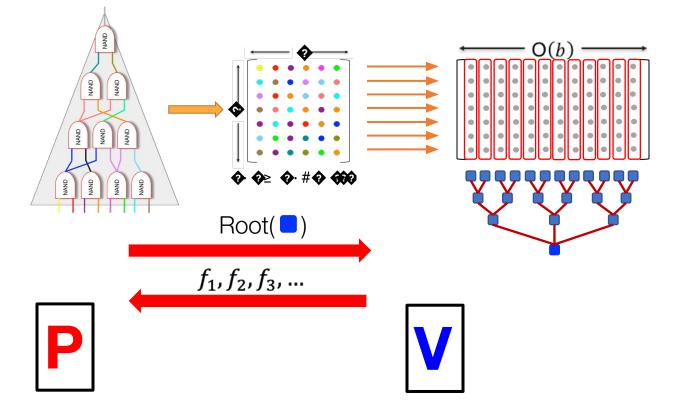
Extended witness \overline{w}

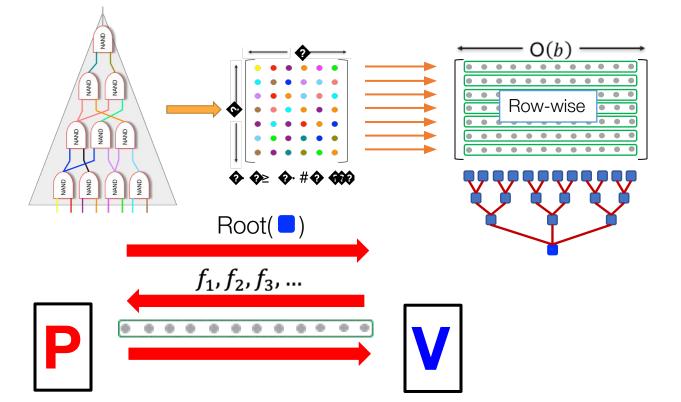


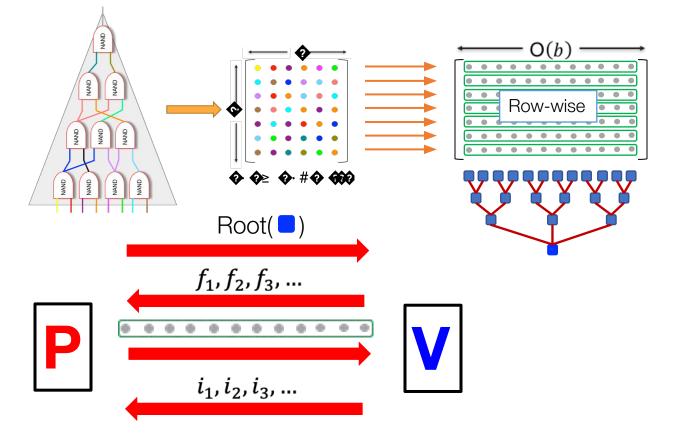


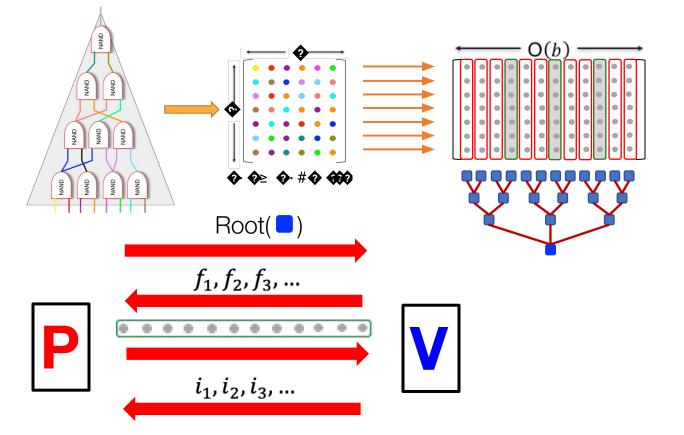


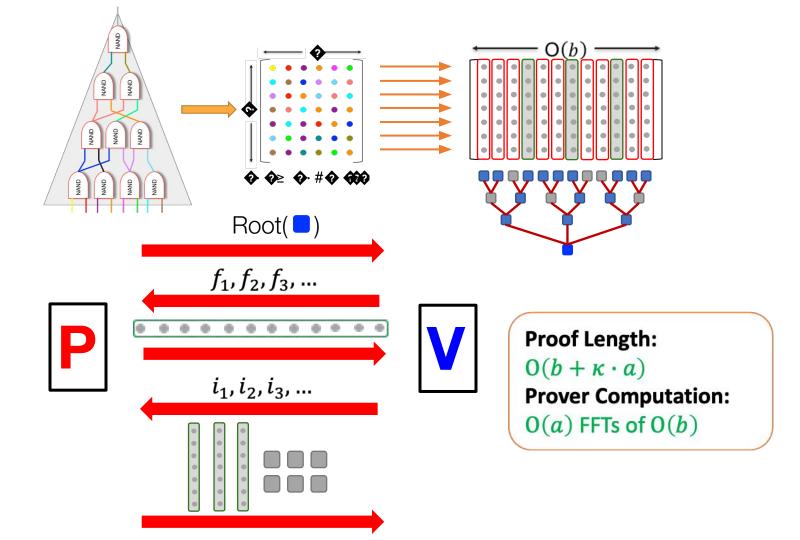






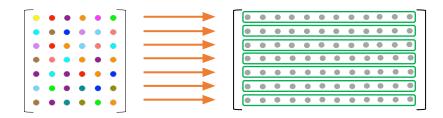




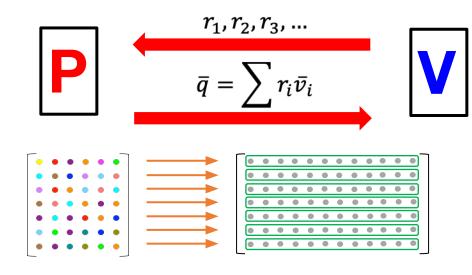


What about row aggregates?

- 1. Code test the prover encoded each row correctly
- 2. Quadratic test the prover computed multiplication gates correctly
- 3. Linear test the prover computed "linear" gates correctly



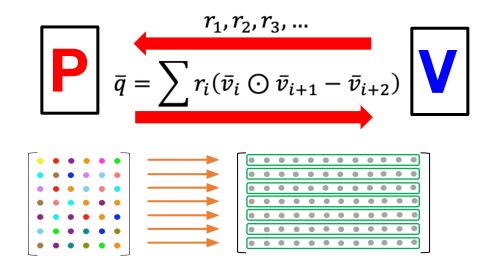
1. Code test – prover encoded correctly



Check if \bar{q} is a valid codeword and agrees on random columns

2. Multiplication gates were correct

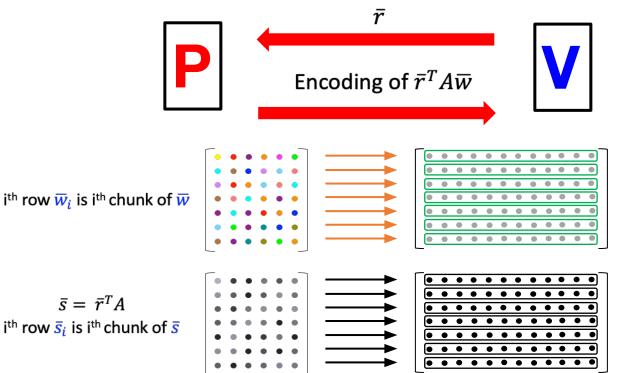
Arrange values so that all constraints for multiplication gates are so that $\overline{w}_i \odot \overline{w}_{i+1} = \overline{w}_{i+2}$ for some set of values i



Check if \overline{q} is a valid codeword encoding 0s and agrees on revealed columns

2. Linear gates were correct

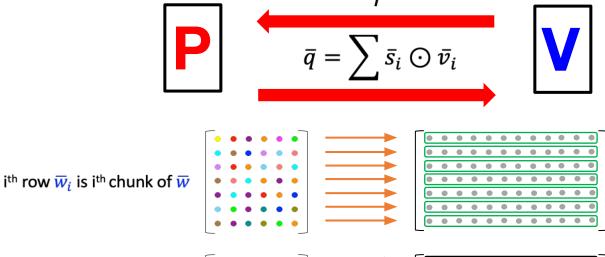
Linear constraints can be expressed as $A\overline{w} = \overline{b}$



Check if response encodes $\bar{r}^T \bar{b}$

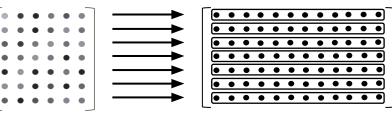
3. Linear gates were correct

Linear constraints can be expressed as $A\overline{w} = \overline{b}$



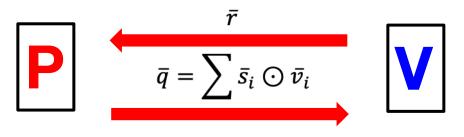
Check if response encodes $\bar{r}^T \bar{b}$

 $ar{s} = ar{r}^T A$ ith row $ar{s}_i$ is ith chunk of $ar{s}$



3. Linear gates were correct

Linear constraints can be expressed as $A\overline{w} = \overline{b}$

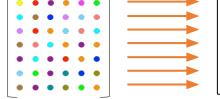


Check if \bar{q} is a valid codeword encoding \bar{z} s.t.

$$\mathbf{1}^T \bar{z} = \bar{r}^T \bar{b}$$

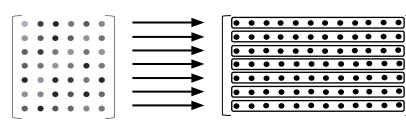
and agrees on revealed columns

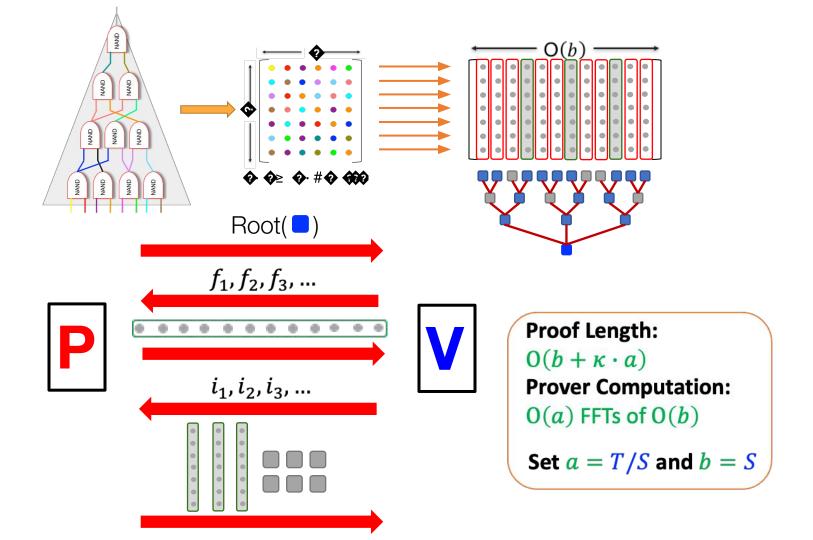
 i^{th} row \overline{w}_i is i^{th} chunk of \overline{w}

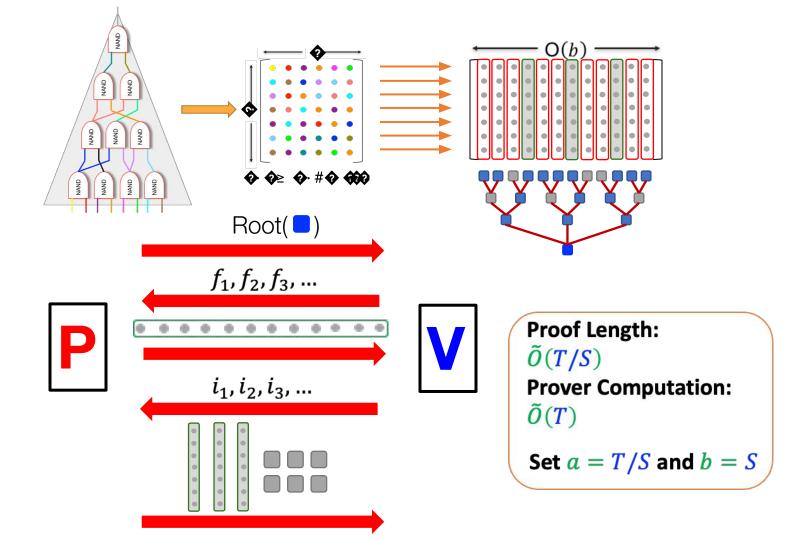


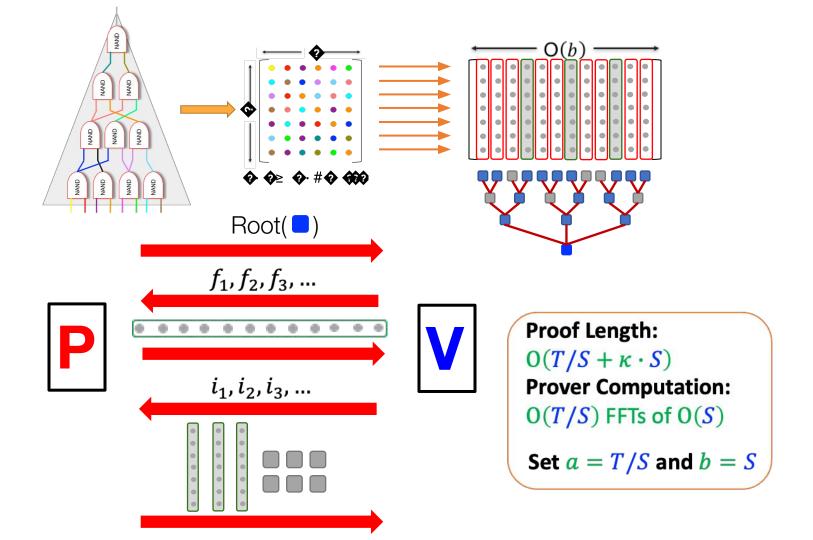
•	•	•	•	•	•	•	•	•	•	•
•	•	•		•	•		•	•	•	•
•	•			•	•		•	•	•	•
•	•	•	•	•	•	•	•	•	•	•
•	•	•	•	•	•	•	•	•	•	•
•	•	•		•	•		•	•	•	•
•	•	•	•	•	•	•	•	•		•

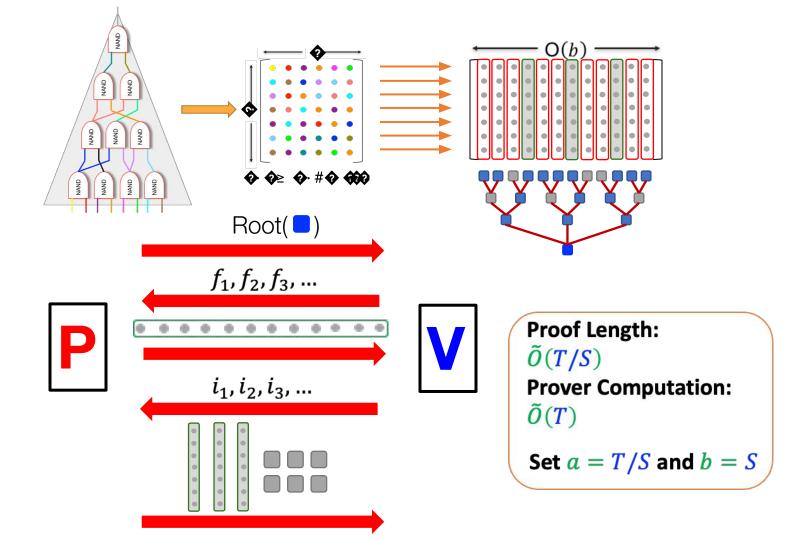
$$ar{s} = ar{r}^T A$$
 \mathbf{i}^{th} row $ar{s}_i$ is \mathbf{i}^{th} chunk of $ar{s}$











Build your ZK app at ligetron.com



Scan this to verify proof on your device