

Groth16 is UC-secure

The **Brave New World**
of Global **Generic Groups**
and **UC-Secure Zero-Overhead SNARKs**

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INPUT | OUTPUT



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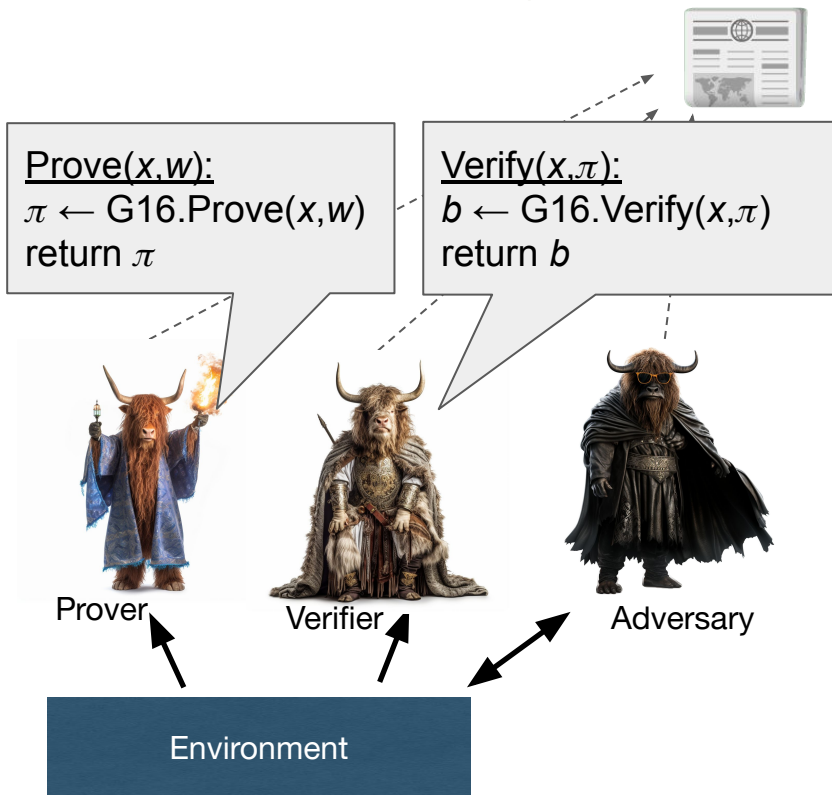
AlgoCRYPT CoE

AI Research

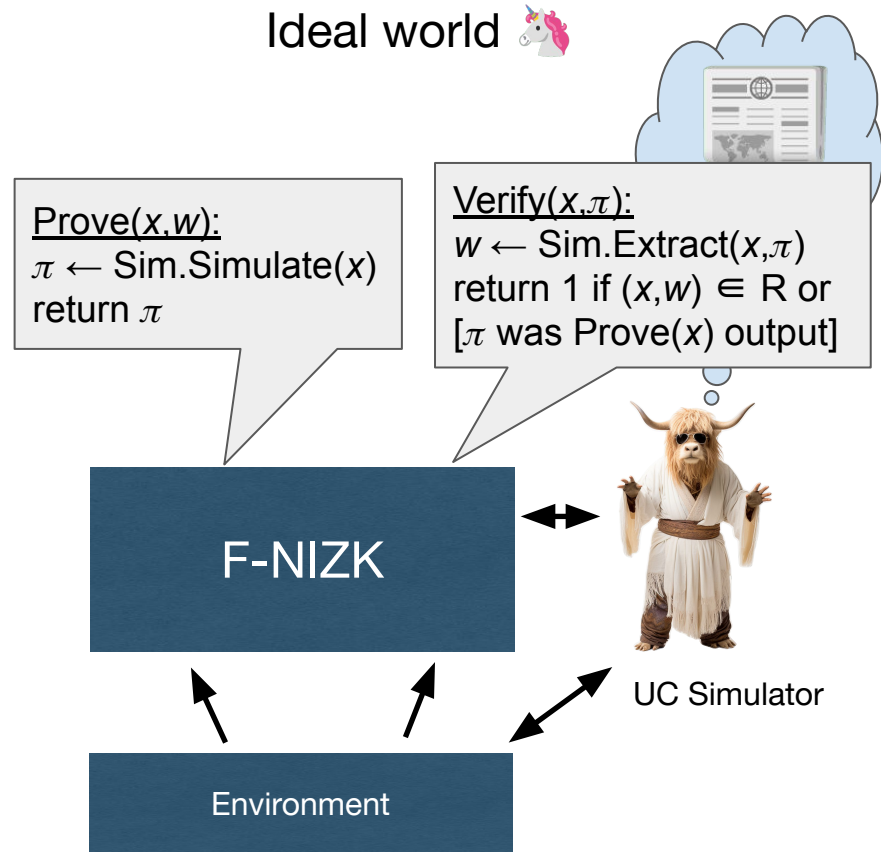


UC modeling of NIZKs

Real world 🐎

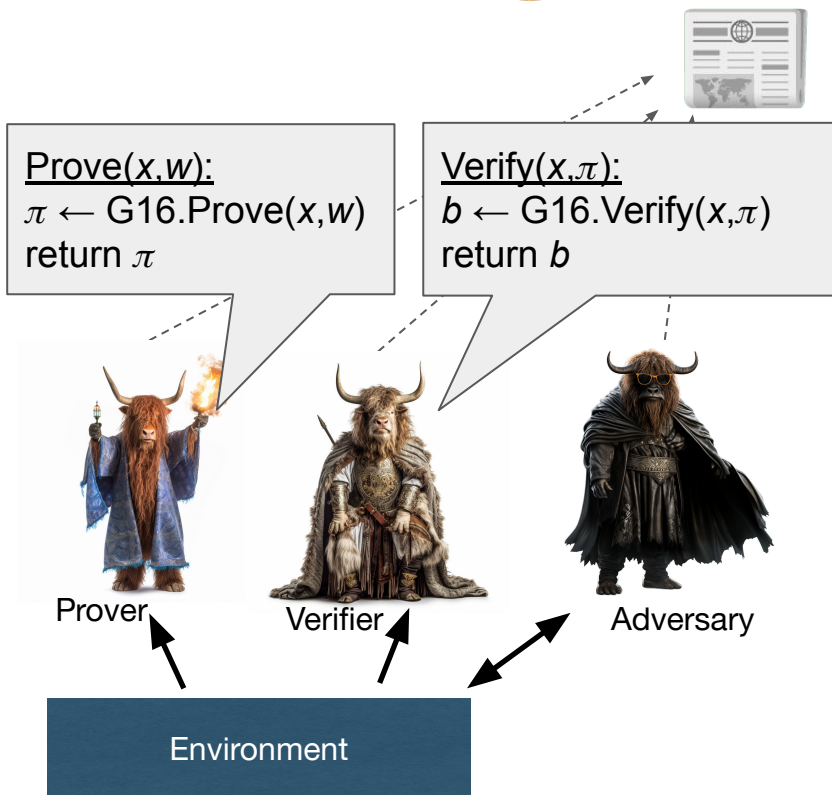


Ideal world 🦄

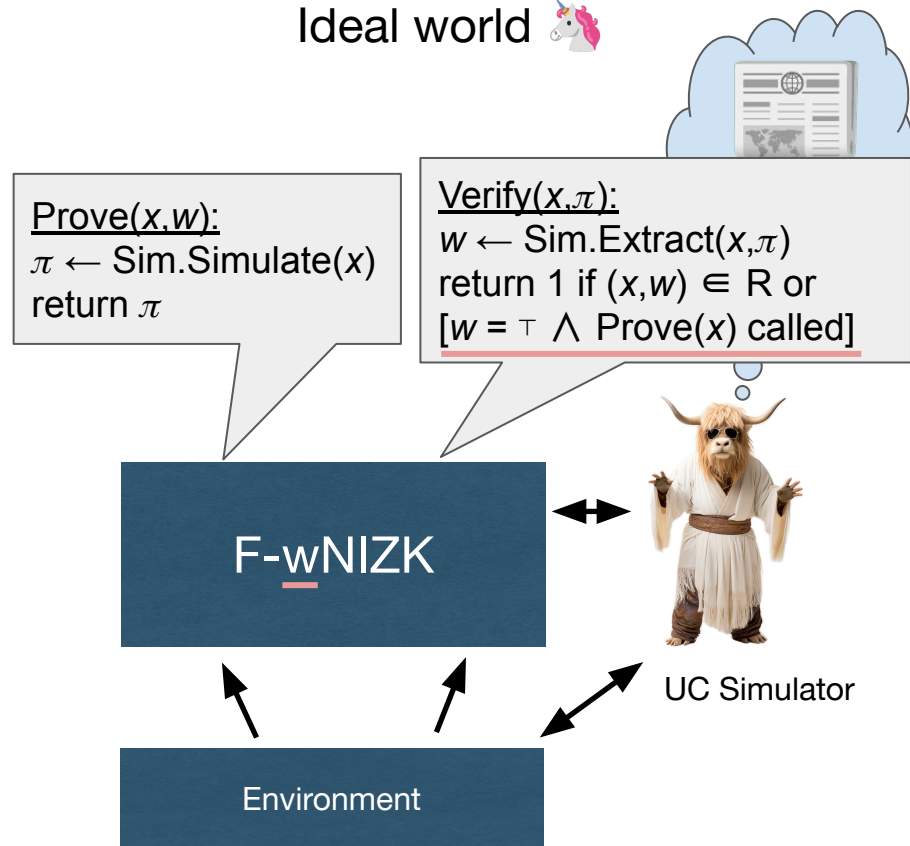


UC modeling of NIZKs

Real world 🐎

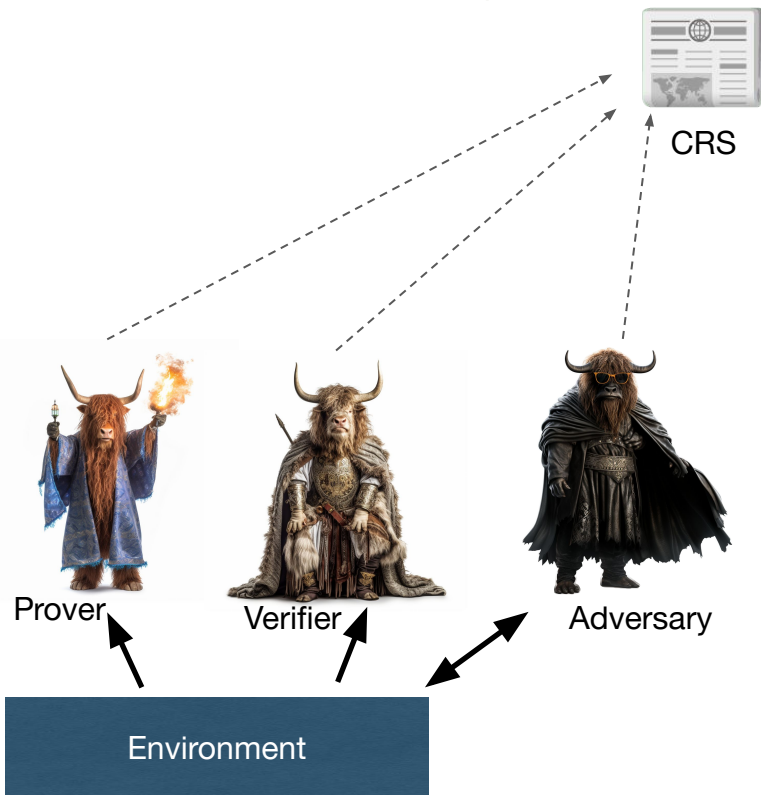


Ideal world 🦄

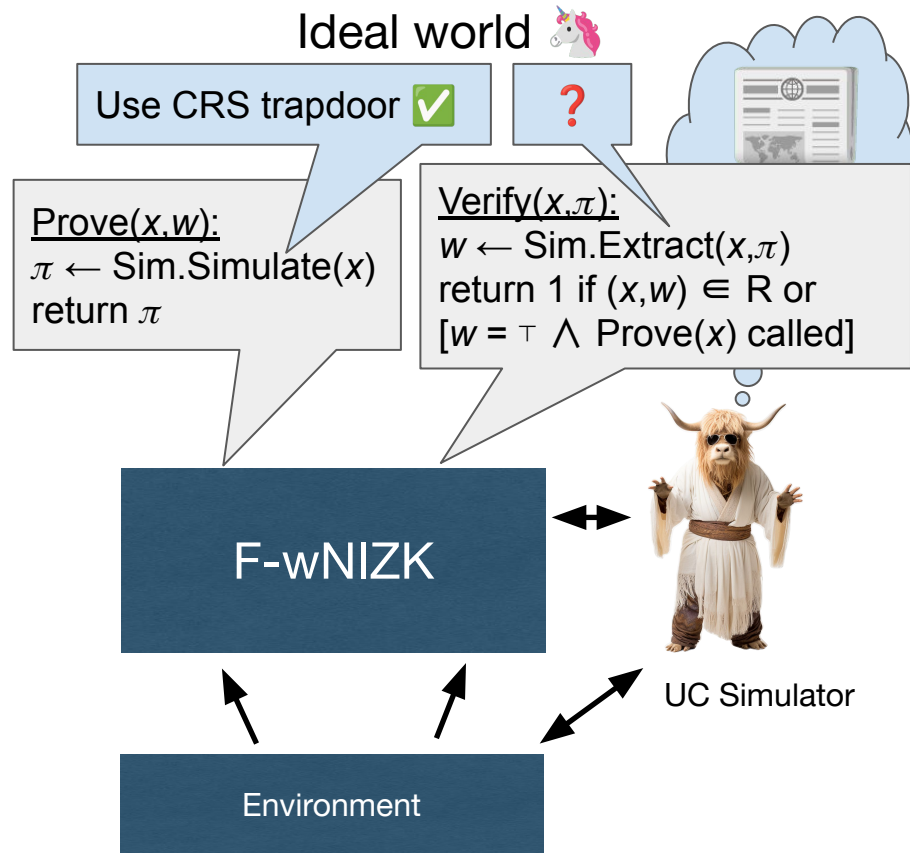


UC modeling of NIZKs

Real world 🐎









Ideal world 🦄





Extraction

Extractor advantage	Groth16 extraction	UC modeling
 GGM Generic group model	✓ [Gro16]	✗ (only folklore <i>local</i> F-GGM)
 AGM Algebraic group model	✓ ~[FLK18, BKS21]	✗ UC-AGM [ABKLX21] Unsuitable for NIZKs
 KA Knowledge assumptions	✓ ~[GM17, BFHK23]	✗ In CC framework [KKK21] Unsuitable, problems like AGM
 CRS Common reference string	📦 via compiler (like C \emptyset C \emptyset) [KZMQCPRsS15], [...]	✓ F-CRS hybrid model
 ROM Random oracle model	📦 via compiler [GKOPTT23]	✓ Global observable RO [CDGLN18] (e.g., [CF24])
 Rewinding	✗ (?)	✗ (does not compose)

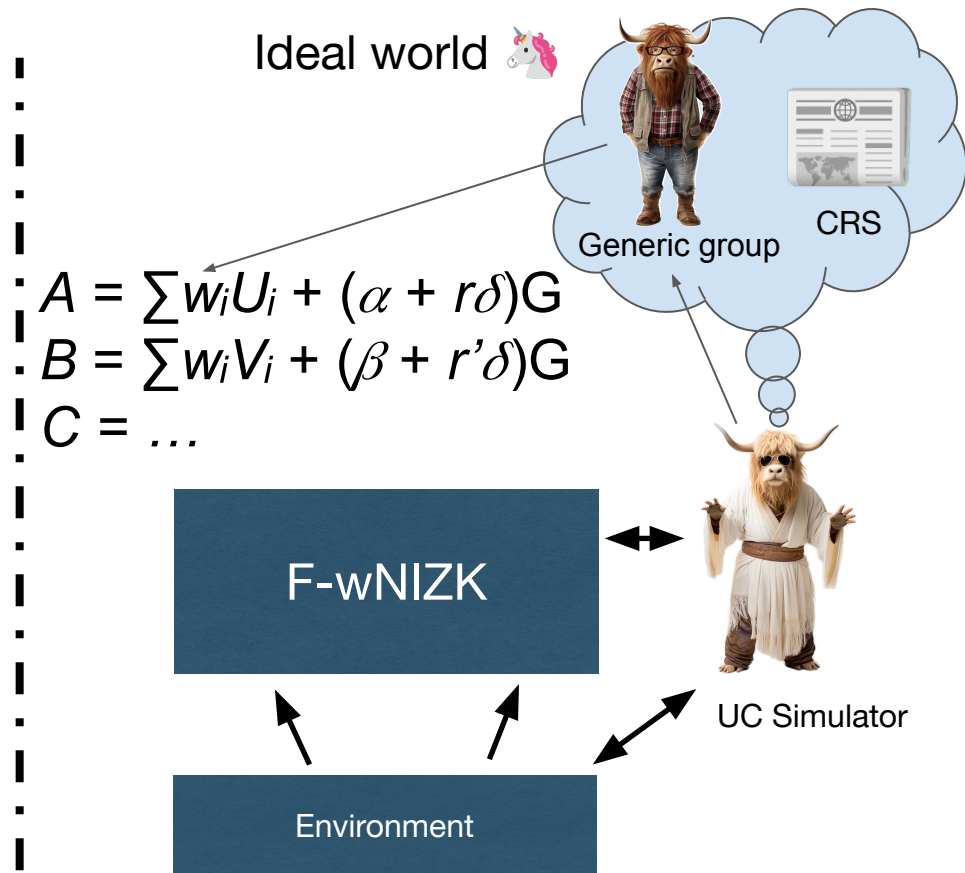
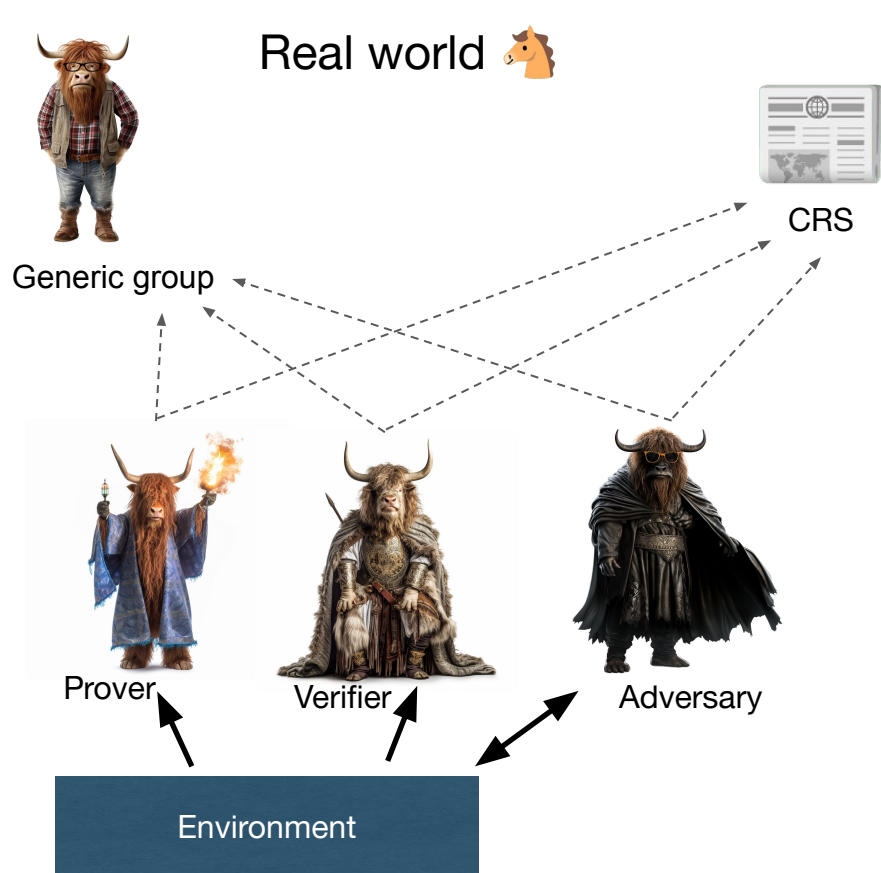
The generic group model

private random injective $\tau: \mathbb{G} \rightarrow S$
public generator $g \in S$

$\text{op}(g_1, g_2)$:
return $\tau(\tau^{-1}(g_1) + \tau^{-1}(g_2))$



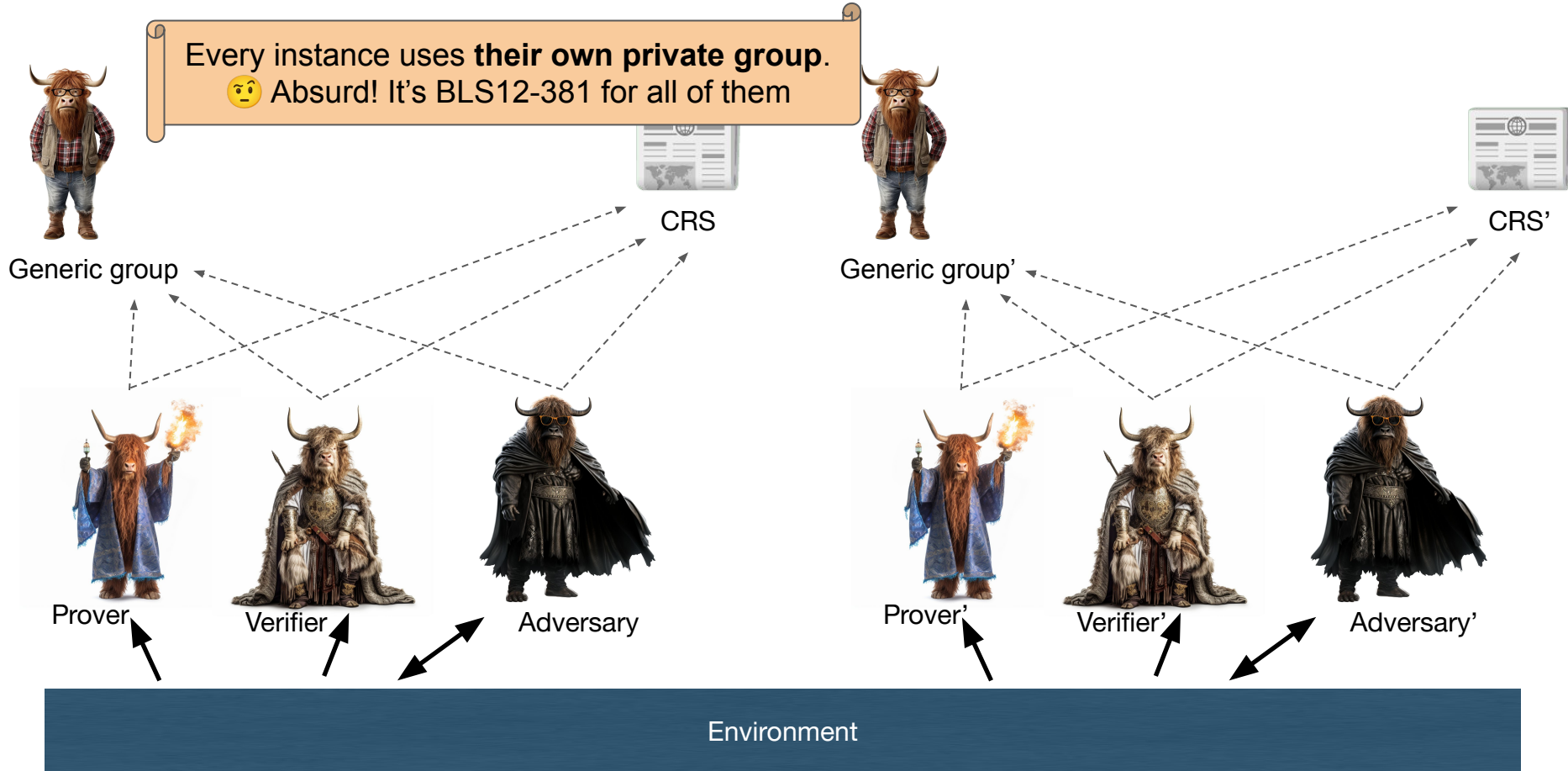
(Local) generic groups in UC



(Local) generic groups in UC







Every instance uses **their own private group**.

🤔 Absurd! It's BLS12-381 for all of them

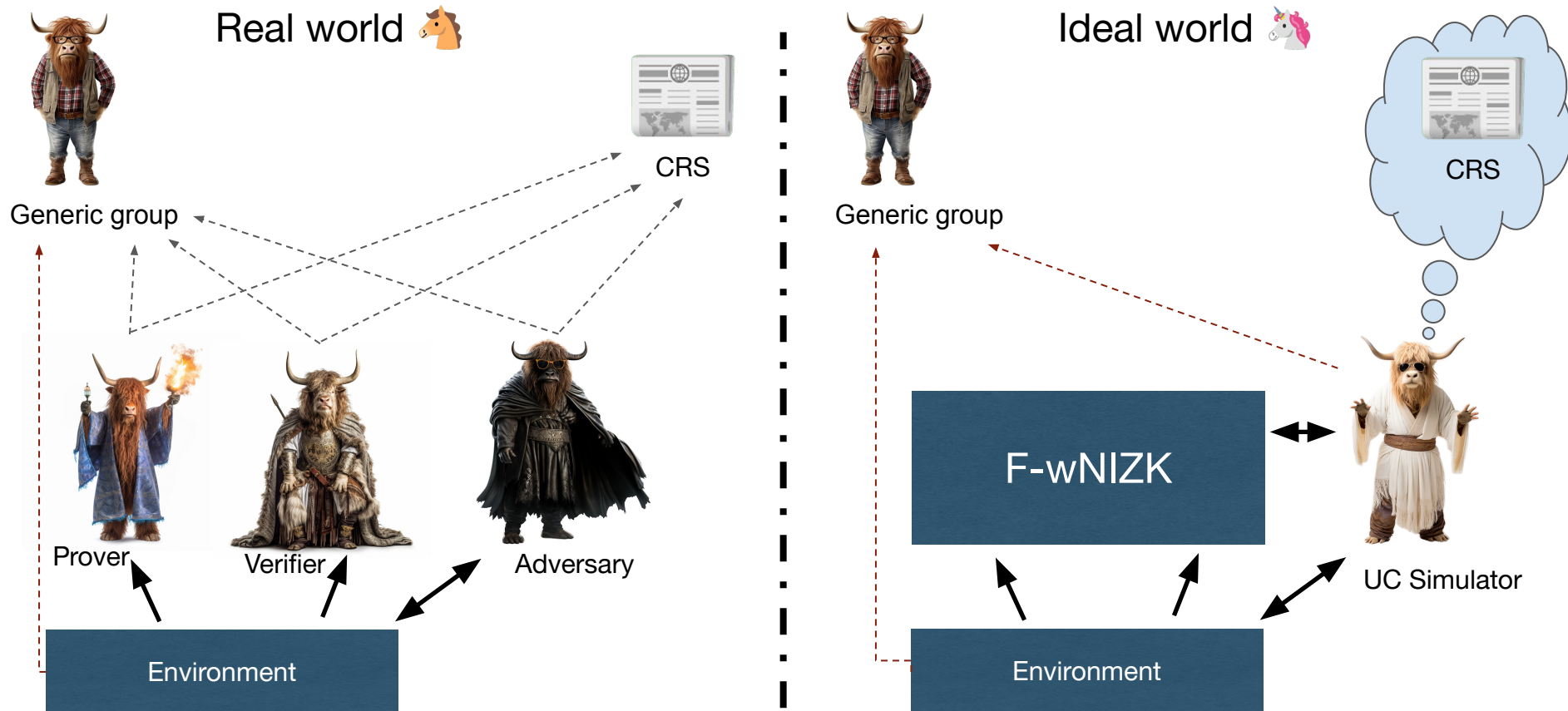




Extraction

Extractor advantage	Groth16 extraction	UC modeling
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Global generic groups in UC



Global generic groups in UC

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public generator $g \in \mathbb{S}$

$\text{op}(g_1, g_2)$:
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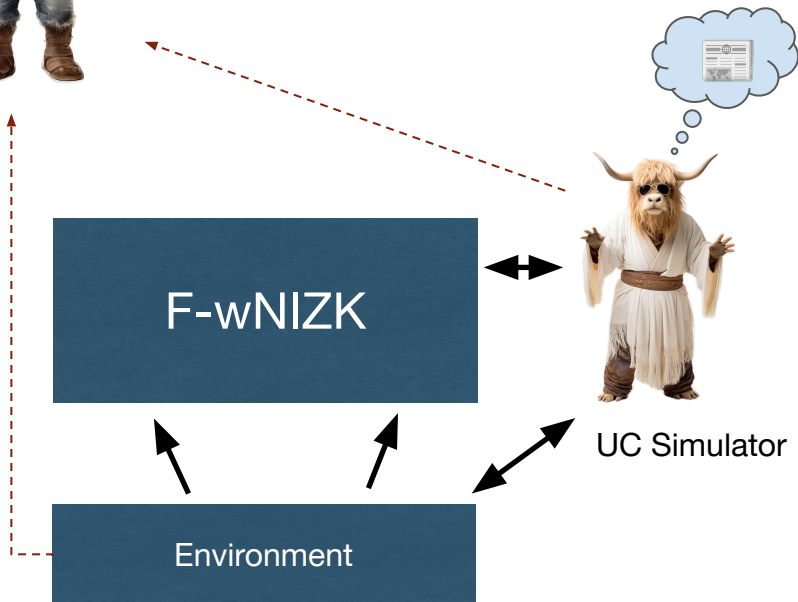
Issue: UC-simulator does not see the environment's group operations.



Attack: Env honestly computes proof π , then calls Verify. Extraction fails.



Ideal world 🦄



Fully observable global generic groups in UC

private random injective $\tau: \mathbb{G} \rightarrow \mathbb{S}$
public generator $g \in \mathbb{S}$
public observation list Obs ←

$op(g_1, g_2)$:

$Obs \leftarrow Obs \parallel (g_1, g_2)$ ←

return $\tau(\tau^{-1}(g_1) + \tau^{-1}(g_2))$



Good: UC-simulator can extract π using Obs !



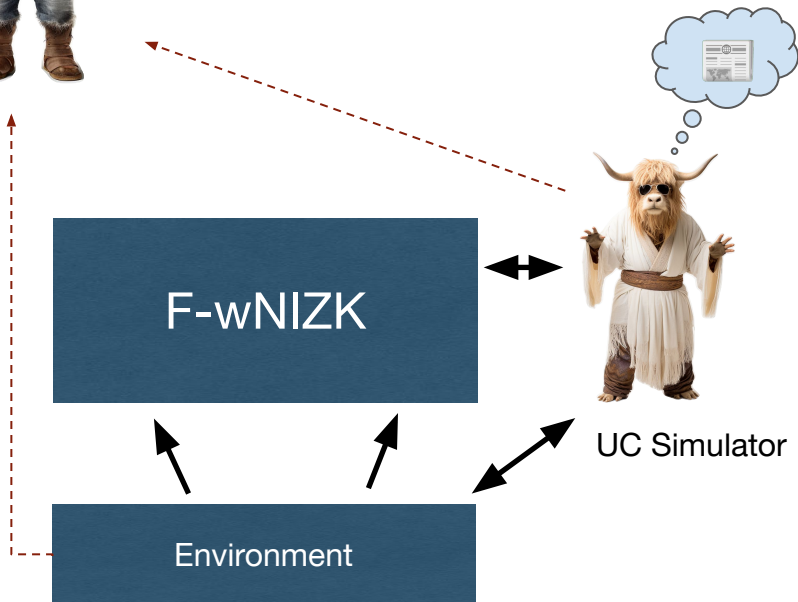
Issue: No ZK



Attack: Check Obs to see whether π was honestly computed or via simulation trapdoor



Ideal world 🦄



Design challenges



Requirements

- 👁 Simulator **must see** group operations made by environment.
 - Required for extraction.
- 🤖 Environment **must not see** group operations made by the simulator.
 - Required for ZK.
- 💡 Solution: Partial observability via **domain separation**.



Design challenges



Requirements

- 👁 Simulator **must see** ^{relevant} group operations made by environment.
 - Required for extraction.
- 🧐 Environment **must not see** group operations made by the simulator.
 - Required for ZK.
- 💡 Solution: Partial observability via **domain separation**.

Observation rules (intuition)

- 📜 Every session s gets its own group generator h_s
- 👍 Legal/unobservable: Session s operates on h_s
- 👎 Illegal/observable: Session s' operates on h_s



Restricted observable global generic groups in UC (simplified)

```
private random injective  $\tau: \mathbb{G} \rightarrow \mathbb{S}$   
public generators  $h_s \in \mathbb{S}$  for sessions  $s$   
public observation list  $Obs$   
  
op( $g_1, g_2$ ):  
  if  $g_1$  or  $g_2$  contains some foreign  $h_s$  then  
     $Obs \leftarrow Obs \parallel (g_1, g_2)$   
  return  $\tau(\tau^{-1}(g_1) + \tau^{-1}(g_2))$ 
```



UC-simulator **can simulate** unobserved.



UC-simulator **can extract** π using Obs .



Ideal world 🦄

s^*

$s' \neq s^*$



UC Simulator



The actual G-oGG model in the paper

- 📖 Bookkeeping with polynomials.
 - (“does g contain some foreign h_s ?”)
- 👤 Multiple generators per session
- 🧐 Oblivious sampling
- ✨ Pairing groups



Our result

Theorem 1. *Π -G16 UC-realizes \mathcal{F} -wNIZK in the \mathcal{F} -CRS-hybrid model in the presence of \mathcal{G} -oGG. Concretely, for any PPT adversary \mathcal{A} , there exists a PPT simulator \mathcal{S}_{G16} such that for every \mathcal{Z} that makes at most $q_{\mathcal{Z}}$ queries to \mathcal{G} -oGG, $q_{\mathcal{P}}$ queries to the PROVE interface, and $q_{\mathcal{V}}$ queries to the VERIFY interface,*

$$\begin{aligned} & |\Pr[\text{EXEC}_{\mathcal{F}\text{-wNIZK}, \mathcal{Z}, \mathcal{S}_{\text{G16}}, \mathcal{G}\text{-oGG}}(\lambda, z) = 1] - \Pr[\text{EXEC}_{\Pi\text{-G16}, \mathcal{Z}, \mathcal{A}, \mathcal{G}\text{-oGG}}(\lambda, z) = 1]| \\ & \leq 72 \cdot d \cdot (m + d + q_{\mathcal{Z}} + (m + d)q_{\mathcal{P}} + \ell q_{\mathcal{V}} + 1)^2 / (p - 1) \end{aligned}$$

and \mathcal{S}_{G16} performs in total the following operations:

- at most $3q_{\mathcal{P}} + 9q_{\mathcal{V}} + 2q_{\mathcal{Z}} + 3d + m + 8$ queries to \mathcal{G} -oGG
- at most $(2\ell + 8)q_{\mathcal{P}} + (3q_{\mathcal{Z}} + 2\ell + 2)q_{\mathcal{V}} + (d + 1)(3m + 11)$ field operations where d, m, ℓ depend on the circuit size (see Functionality 6).

Groth16 UC proof challenges



Strategy:

Observe group operations to extract coefficients = w

Extraction

Simulation



Strategy:

Use CRS trapdoor to create proof without w



Challenge:

Can only observe *my* session's generators



Challenge:

Simulator's operations must not be observable



Solution:

Argue: Foreign generators don't mess up extraction.



Solution:

Easy: Simulator legally operates on CRS elements



Consequences

- 🎉 **Groth16 (as-is) is UC-secure**
 - No need for compilers.
- 🌟 Our model **also enables** KZG/IOPs, Schnorr, ... in UC
- 🚫 **Restrictions:**
 - GGM → **No circuit** for group operations
 - Some protocols **cannot deal** with (partial) observability. Workaround in paper.

