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Point values are assigned for each question.

Points earned: _____ / 100, = _____ %

1. Find an upper bound for $f(n) = n^4 + 10n^2 + 5$. Write your answer here: $O(n^4)$ (4 points)

Prove your answer by giving values for the constants c and n_0 . Choose the smallest integral value possible for c. (4 points)

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Using
$$c = 2$$

 $0 \le n^4 + 10n^2 + 5 \le 2n^4$
Solving for n_0
 $n_0 = 4$

So this means that the upper bound is $O(n^4)$

2. Find an asymptotically tight bound for $f(n) = 3n^3 - 2n$. Write your answer here: $\Theta(n^3)$ (4 points)

Prove your answer by giving values for the constants c_1 , c_2 , and n_0 . Choose the tightest integral values possible for c_1 and c_2 . (6 points)

$$c_{1} = 2$$

$$c_{2} = 3$$

$$n_0 = 2$$

$$2n^{3} \le 3n^{3} - 2n \le 3n^{3}$$

when $n_{0} = 2$

$$16 \leq 20 \leq 24$$

3. Is $3n - 4 \in \Omega(n^2)$? Circle your answer: yes / **no**. (2 points)

If yes, prove your answer by giving values for the constants c and n_0 . Choose the smallest integral value possible for c. If not, derive a contradiction. (4 points)

3n-4 cannot be within $\Omega(n^2)$ because it's better than the best case which does not make sense.

$$cn^{2} \le 3n - 4$$

$$cn^{2} \le 3n - 4 \le 3n$$

$$cn^{2} \le 3n$$

$$cn \le 3$$

$$n \le \frac{3}{c} \, \forall n \ge n_0$$

This shows how 3n-4 is not within the best case of n^2 because $n \leq \frac{3}{c}$ is not possible. Proved by contradiction

4. Write the following asymptotic efficiency classes in **increasing** order of magnitude. (2 points each)

$$O(n^2), O(2^n), O(1), O(n\lg n), O(n), O(n!), O(n^3), O(\lg n), O(n^n), O(n^2\lg n)$$
 (2 points each)

$$O(1), O(lg(n)), O(n), O(nlg(n)), O(n^2), O(n^2lg(n)), O(n^3), O(2^n), O(n!), O(n^n)$$

- 5. Determine the largest size n of a problem that can be solved in time t, assuming that the algorithm takes f(n) milliseconds. Write your answer for n as an integer. (2 points each)
 - a. f(n) = n, t = 1 second 1,000
 - b. f(n) = nlg(n), t = 1 hour 204,094
 - c. $f(n) = n^2$, t = 1 hour 1,897
 - d. $f(n) = n^3$, t = 1 day 442
 - e. f(n) = n!, t = 1 minute 8
- 6. Suppose we are comparing two sorting algorithms and that for all inputs of size n the first algorithm runs in $4n^3$ seconds, while the second algorithm runs in 64n lg n seconds. For which integral values of n does the first algorithm beat the second algorithm?

$$2 \le n \le 6$$

(4 points)

Explain how you got your answer or paste code that solves the problem (2 point):

```
C: > Users > samps > Downloads > 💠 HA1Algo.py > ...
     import math
     for x in range(1, 11):
        print("First algorith: " + str(4 * x ** 3) + ", Second algorithm: " + str(64 * x * math.log(x, 2)))
PROBLEMS OUTPUT DEBUG CONSOLE TERMINAL
Windows PowerShell
Copyright (C) Microsoft Corporation. All rights reserved.
Try the new cross-platform PowerShell https://aka.ms/pscore6
PS C:\Users\samps\OneDrive\Desktop\CS306> & 'C:\Users\samps\AppData\Local\Microsoft\WindowsApps\python3.10.exe' 'c:\Users\samps
First algorith: 4, Second algorithm: 0.0
First algorith: 32, Second algorithm: 128.0
First algorith: 108, Second algorithm: 304.312800138462
First algorith: 256, Second algorithm: 512.0
First algorith: 500, Second algorithm: 743.0169903639559
First algorith: 864, Second algorithm: 992.6256002769239
First algorith: 1372, Second algorithm: 1257.6950050818066
First algorith: 2048, Second algorithm: 1536.0
First algorith: 2916, Second algorithm: 1825.876800830772
First algorith: 4000, Second algorithm: 2126.033980727912
PS C:\Users\samps\OneDrive\Desktop\CS306>
```

As we can see, the value for the first algorithm is smaller only on intervals from 2 to 6 (inclusive). Everything outside of these bounds, the second algorithm is faster.

7. Give the complexity of the following methods. Choose the most appropriate notation from among 0, 0, and Ω . (8 points each)

```
int function1(int n) {
    int count = 0;
    for (int i = n / 2; i <= n; i++) {
        for (int j = 1; j <= n; j *= 2) {
            count++;
        }
    }
    return count;
}
Answer: \theta(nlog(n))

int function2(int n) {
    int count = 0;
    for (int i = 1; i * i * i <= n; i++) {
        count++;
    }
    return count;
}
Answer: \theta(\sqrt[3]{n})

int function3(int n) {
```

```
int count = 0;
    for (int i = 1; i <= n; i++) {</pre>
         for (int j = 1; j <= n; j++) {</pre>
             for (int k = 1; k <= n; k++) {</pre>
                  count++;
             }
         }
    return count;
}
Answer: \theta(n^3)
int function4(int n) {
    int count = 0;
    for (int i = 1; i <= n; i++) {</pre>
         for (int j = 1; j <= n; j++) {</pre>
             count++;
             break;
         }
    return count;
}
Answer: \theta(n)
int function5(int n) {
    int count = 0;
    for (int i = 1; i <= n; i++) {</pre>
         count++;
    for (int j = 1; j <= n; j++) {</pre>
         count++;
    return count;
}
Answer: \theta(n)
```