

# Linear System

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# Outline

- Direct Method
  - Gaussian elimination and LU decomposition
  - Cholesky decomposition
- Iterative Method
  - Basic stationary methods
  - Smooth multigrid Solver
  - Krylov subspace methods

# Problem

- Solve Linear Equation

$$A\mathbf{x} = \mathbf{b},$$

# Gaussian elimination and LU decomposition

1. *LU decomposition:*  $A = LU$  ;
2. *Forward substitution:* solve  $L\mathbf{y} = \mathbf{b}$  ;
3. *Backward substitution:* solve  $U\mathbf{x} = \mathbf{y}$  .

Forming the LU decomposition takes  $\mathcal{O}(n^3)$  floating point operations for a general dense  $A \in \mathbb{R}^{n \times n}$ . It takes  $\mathcal{O}(n^2)$  floating point operations to perform backward or forward substitutions.

# Gaussian elimination and LU decomposition

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**Algorithm 1** The LU decomposition of a matrix  $A$ . Upon exit, the entries of  $A$  have been overwritten with the entries of  $L$  (below the main diagonal) and the entries of  $U$  (main diagonal and above). The diagonal entries of  $L$  are all equal to 1.

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```
for  $k = 1, \dots, n - 1$  do
    for  $i = k + 1, \dots, n$  do
         $a_{i,k} = \frac{a_{i,k}}{a_{k,k}}$ 
        for  $j = k + 1, \dots, n$  do
             $a_{i,j} = a_{i,j} - a_{i,k}a_{k,j}$ 
        end for
    end for
end for
```

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# Cholesky decomposition

- special method for symmetric positive definiteness matrix

$$A = FF^\top,$$

# Cholesky decomposition

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**Algorithm 2** The Cholesky decomposition of a symmetric positive definite matrix  $A$ . Upon exit, the entries of  $A$  on its diagonal and below it have been overwritten with the entries of the lower triangular Cholesky factor  $F$ .

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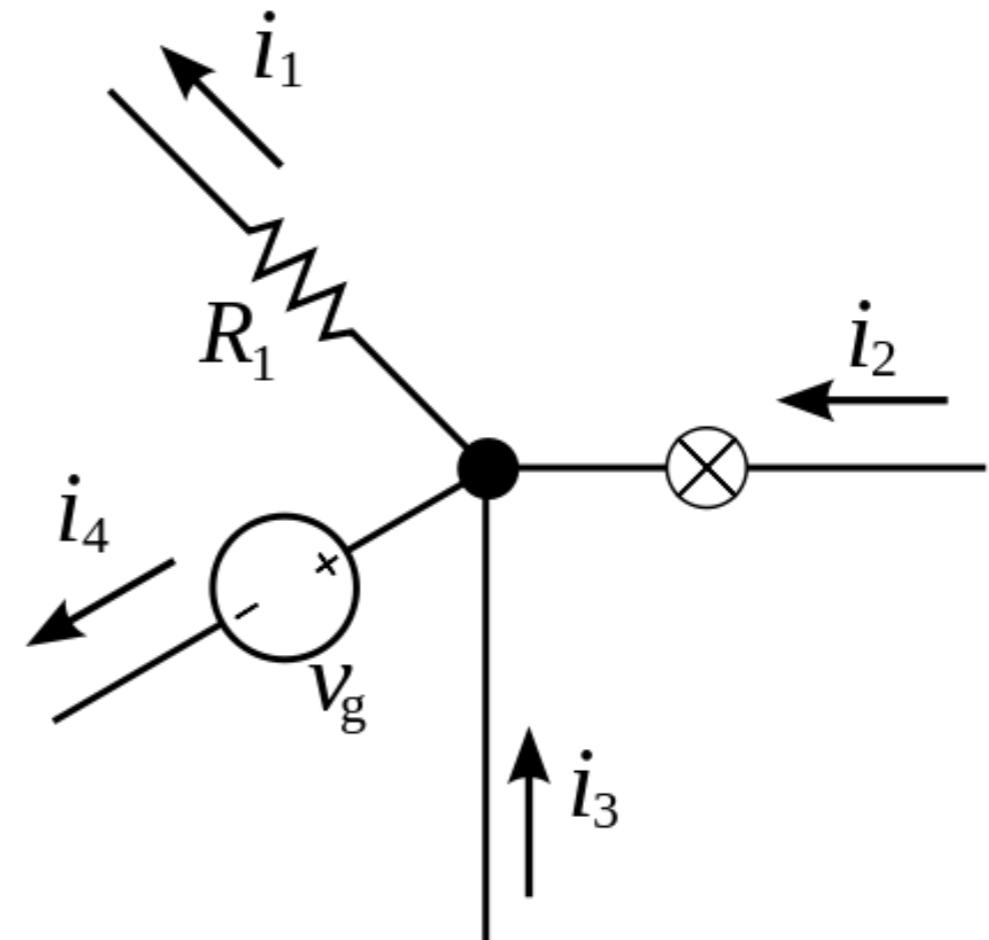
```
for  $k = 1, \dots, n$  do
     $a_{k,k} = \sqrt{a_{k,k}}$ 
    for  $i = k + 1, \dots, n$  do
         $a_{i,k} = \frac{a_{i,k}}{a_{k,k}}$ 
    end for
    for  $j = k + 1, \dots, n$  do
        for  $i = j, \dots, n$  do
             $a_{i,j} = a_{i,j} - a_{i,k}a_{k,j}$ 
        end for
    end for
end for
```

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# Analysis of an Electrical Network

Kirchhoff's circuit laws

$$\sum_{k=1}^n I_k = 0$$



# 1D Smooth filter

Let  $v$  be the sum of a smooth 1D signal  $u$  and IID Gaussian noise  $e$ :

$$v = u + e , \quad (1)$$

where  $u = (u_1, \dots, u_N)$ ,  $v = (v_1, \dots, v_N)$ , and  $e = (e_1, \dots, e_N)$ .

- minimize the following objective

$$E(u) = \sum_{n=1}^N (u_n - v_n)^2 + \lambda \sum_{n=1}^{N-1} (u_{n+1} - u_n)^2$$

# 1D Smooth filter

- minimize the following objective

$$E(u) = \sum_{n=1}^N (u_n - v_n)^2 + \lambda \sum_{n=1}^{N-1} (u_{n+1} - u_n)^2$$

- calculate the gradient

$$\frac{\partial E(u)}{\partial u_n} = 2(u_n - v_n) + 2\lambda(-u_{n-1} + 2u_n - u_{n+1})$$

$$u_n + \lambda(-u_{n-1} + 2u_n - u_{n+1}) = v_n$$

# 1D Smooth filter

$$\begin{pmatrix} 1 + \lambda & -\lambda & 0 & 0 & \dots & 0 \\ -\lambda & 1 + 2\lambda & -\lambda & 0 & \dots & 0 \\ 0 & -\lambda & 1 + 2\lambda & -\lambda & \dots & 0 \\ \ddots & \ddots & \ddots & & & \\ 0 & \dots & & 0 & -\lambda & 1 + \lambda \end{pmatrix} \begin{pmatrix} u_1 \\ u_2 \\ u_3 \\ \vdots \\ u_N \end{pmatrix} = \begin{pmatrix} v_1 \\ v_2 \\ v_3 \\ \vdots \\ v_N \end{pmatrix}$$

O(n) by using Gaussian Elimination

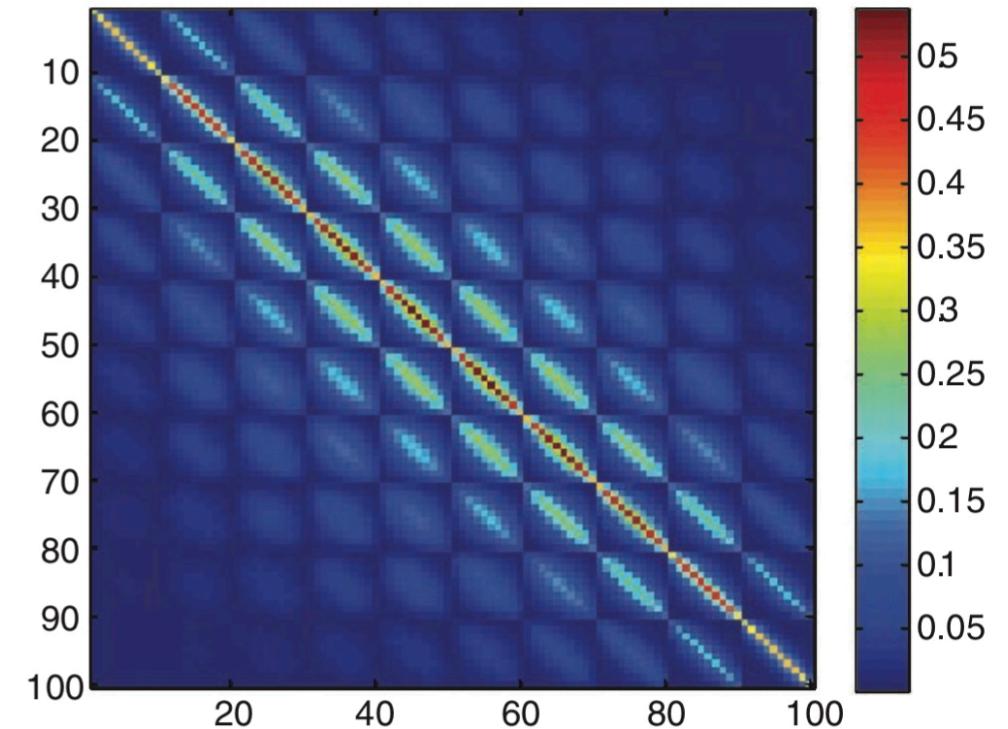
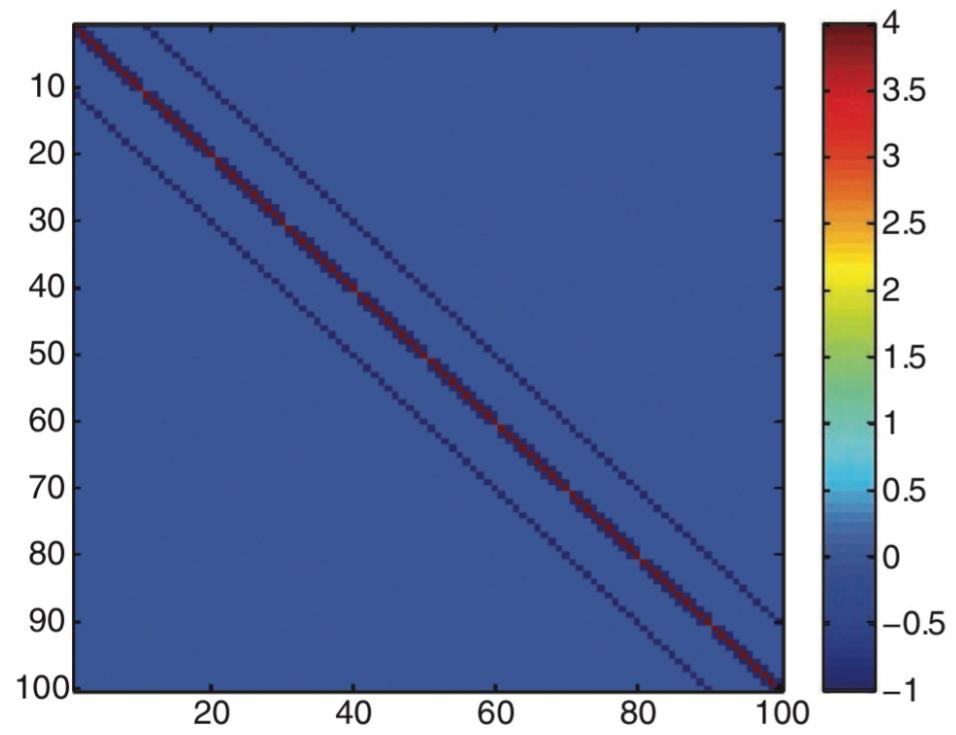
# Nonzero structure and sparsity pattern

- Narrow-banded matrices
  - tridiagonal matrices  $O(n)$  for Gaussian Elimination

$$T = \begin{pmatrix} a_1 & b_1 & & & \\ c_1 & a_2 & b_2 & & \\ & c_2 & \ddots & \ddots & \\ & & \ddots & \ddots & b_{n-1} \\ & & & c_{n-1} & a_n \end{pmatrix}$$

# nonzero structure and sparsity pattern

- the discrete two-dimensional Laplacian arising from discretization of the Poisson equation on a uniform mesh on a square domain
  - The matrix has approximately  $5n$  nonzero entries
  - but the Cholesky factor contains about  $O(n^*sqrt(n))$
  - the decomposition is  $O(n^2)$  floating-point operations



# Basic stationary methods

## Basic stationary methods

Given  $A\mathbf{x} = \mathbf{b}$ , suppose  $A = M - N$ . We have  $M\mathbf{x} = N\mathbf{x} + \mathbf{b}$ , which may lead to the iteration

$$M\mathbf{x}_{k+1} = N\mathbf{x}_k + \mathbf{b}.$$

This is a *stationary* or a *fixed-point* iteration.

# Jacobi Method

$$A = D + R \quad \text{where} \quad D = \begin{bmatrix} a_{11} & 0 & \cdots & 0 \\ 0 & a_{22} & \cdots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \cdots & a_{nn} \end{bmatrix} \quad \text{and} \quad R = \begin{bmatrix} 0 & a_{12} & \cdots & a_{1n} \\ a_{21} & 0 & \cdots & a_{2n} \\ \vdots & \vdots & \ddots & \vdots \\ a_{n1} & a_{n2} & \cdots & 0 \end{bmatrix}.$$

$$\mathbf{x}^{(k+1)} = D^{-1}(\mathbf{b} - R\mathbf{x}^{(k)}),$$

$$x_i^{(k+1)} = \frac{1}{a_{ii}} \left( b_i - \sum_{j \neq i} a_{ij} x_j^{(k)} \right), \quad i = 1, 2, \dots, n.$$

# Gauss–Seidel Method

$$A = L_* + U \quad \text{where} \quad L_* = \begin{bmatrix} a_{11} & 0 & \cdots & 0 \\ a_{21} & a_{22} & \cdots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ a_{n1} & a_{n2} & \cdots & a_{nn} \end{bmatrix}, \quad U = \begin{bmatrix} 0 & a_{12} & \cdots & a_{1n} \\ 0 & 0 & \cdots & a_{2n} \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \cdots & 0 \end{bmatrix}$$

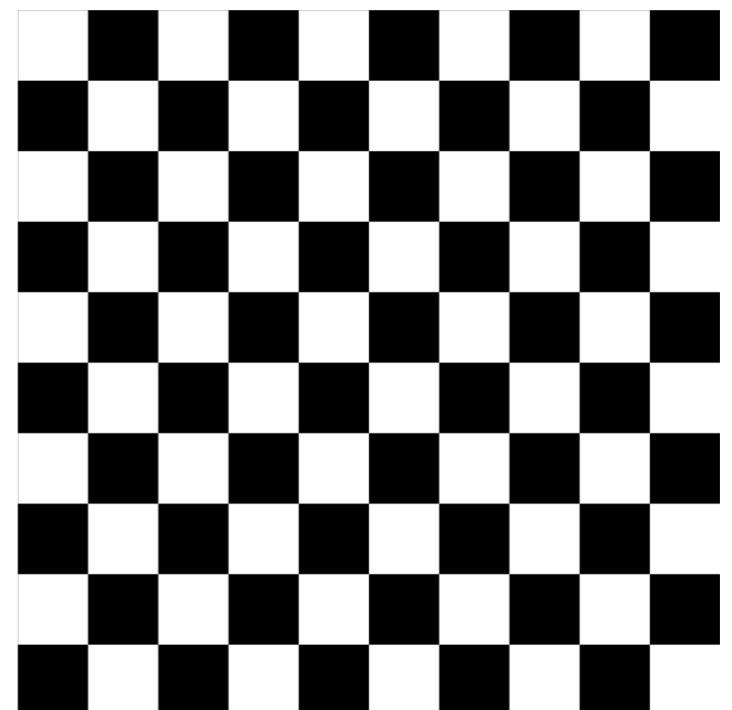
$$\mathbf{x}^{(k+1)} = L_*^{-1}(\mathbf{b} - U\mathbf{x}^{(k)}).$$

$$x_i^{(k+1)} = \frac{1}{a_{ii}} \left( b_i - \sum_{j=1}^{i-1} a_{ij} x_j^{(k+1)} - \sum_{j=i+1}^n a_{ij} x_j^{(k)} \right), \quad i = 1, 2, \dots, n.$$

# Basic stationary methods

- Jacobi Method: Highly Parallelizable
- Gauss–Seidel Method: Faster Convergence Speed
  - In some cases(eg. processing Images), Gauss–Seidel Method is also highly parallelizable
    - Black and white coloring

$$x_i^{(k+1)} = \frac{1}{a_{ii}} \left( b_i - \sum_{j=1}^{i-1} a_{ij} x_j^{(k+1)} - \sum_{j=i+1}^n a_{ij} x_j^{(k)} \right), \quad i = 1, 2, \dots, n.$$



# Convergence of iterative methods

- convergence is governed by the eigenvalues of the iteration matrix

$$T = M^{-1}N = I - M^{-1}A$$

- A necessary and sufficient condition for convergence is that the eigenvalues of  $T$  are all smaller than 1
- The smaller the maximal magnitude of the eigenvalues, the faster the convergence
  - If some of the eigenvalues of  $T$  are very close to 1, then we may experience trouble. Unfortunately, this is often the case in many applications, particularly in the solution of discretized PDE.

# Krylov subspace methods

- Find solution in Krylov sub space

$$\mathbf{x}_k \in \mathbf{x}_0 + \mathcal{K}^k(A; \mathbf{r}_0) \equiv \mathbf{x}_0 + \text{span}\{\mathbf{r}_0, A\mathbf{r}_0, A^2\mathbf{r}_0, \dots, A^{k-1}\mathbf{r}_0\}$$

- Conjugate Gradient

- Minimize the following quadratic function

$$f(\mathbf{x}) = \frac{1}{2} \mathbf{x}^\top \mathbf{A} \mathbf{x} - \mathbf{x}^\top \mathbf{b}, \quad \mathbf{x} \in \mathbf{R}^n.$$

$$\nabla^2 f(\mathbf{x}) = \mathbf{A}, \quad \nabla f(\mathbf{x}) = \mathbf{A}\mathbf{x} - \mathbf{b}.$$

# Direct Conjugate Gradient Method

$$\mathbf{x}_* = \sum_{i=1}^n \alpha_i \mathbf{p}_i. \quad \forall i \neq k : \langle \mathbf{p}_k, \mathbf{p}_i \rangle_{\mathbf{A}} = 0$$

Based on this expansion we calculate:

$$\mathbf{Ax}_* = \sum_{i=1}^n \alpha_i \mathbf{Ap}_i. \quad \langle \mathbf{p}_k, \mathbf{b} \rangle = \alpha_k \langle \mathbf{p}_k, \mathbf{p}_k \rangle_{\mathbf{A}}$$

Left-multiplying by  $\mathbf{p}_k^T$ :

$$\mathbf{p}_k^T \mathbf{Ax}_* = \sum_{i=1}^n \alpha_i \mathbf{p}_k^T \mathbf{Ap}_i, \quad \alpha_k = \frac{\langle \mathbf{p}_k, \mathbf{b} \rangle}{\langle \mathbf{p}_k, \mathbf{p}_k \rangle_{\mathbf{A}}}$$

substituting  $\mathbf{Ax}_* = \mathbf{b}$  and  $\mathbf{u}^T \mathbf{Av} = \langle \mathbf{u}, \mathbf{v} \rangle_{\mathbf{A}}$ :

$$\mathbf{p}_k^T \mathbf{b} = \sum_{i=1}^n \alpha_i \langle \mathbf{p}_k, \mathbf{p}_i \rangle_{\mathbf{A}},$$

# Conjugate Gradient Method

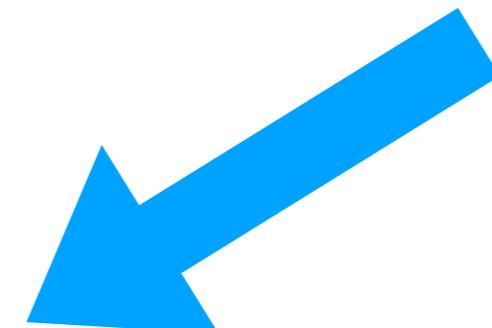
$$\mathbf{r}_k = \mathbf{b} - \mathbf{A}\mathbf{x}_k.$$

$$\mathbf{p}_k = \mathbf{r}_k - \sum_{i < k} \frac{\mathbf{p}_i^\top \mathbf{A} \mathbf{r}_k}{\mathbf{p}_i^\top \mathbf{A} \mathbf{p}_i} \mathbf{p}_i$$

$$\mathbf{x}_{k+1} = \mathbf{x}_k + \alpha_k \mathbf{p}_k$$

$$\alpha_k = \frac{\mathbf{p}_k^\top (\mathbf{b} - \mathbf{A}\mathbf{x}_k)}{\mathbf{p}_k^\top \mathbf{A} \mathbf{p}_k} = \frac{\mathbf{p}_k^\top \mathbf{r}_k}{\mathbf{p}_k^\top \mathbf{A} \mathbf{p}_k}$$

$$f(\mathbf{x}_{k+1}) = f(\mathbf{x}_k + \alpha_k \mathbf{p}_k) =: g(\alpha_k)$$
$$g'(\alpha_k) \stackrel{!}{=} 0 \quad \Rightarrow \quad \alpha_k = \frac{\mathbf{p}_k^\top (\mathbf{b} - \mathbf{A}\mathbf{x}_k)}{\mathbf{p}_k^\top \mathbf{A} \mathbf{p}_k}$$



# Conjugate Gradient Method

$$\mathbf{r}_0 := \mathbf{b} - \mathbf{Ax}_0$$

$$\mathbf{p}_0 := \mathbf{r}_0$$

$$k := 0$$

repeat

$$\alpha_k := \frac{\mathbf{r}_k^\top \mathbf{r}_k}{\mathbf{p}_k^\top \mathbf{A} \mathbf{p}_k}$$

$$\mathbf{x}_{k+1} := \mathbf{x}_k + \alpha_k \mathbf{p}_k$$

$$\mathbf{r}_{k+1} := \mathbf{r}_k - \alpha_k \mathbf{A} \mathbf{p}_k$$

if  $r_{k+1}$  is sufficiently small, then exit loop

$$\beta_k := \frac{\mathbf{r}_{k+1}^\top \mathbf{r}_{k+1}}{\mathbf{r}_k^\top \mathbf{r}_k}$$

$$\mathbf{p}_{k+1} := \mathbf{r}_{k+1} + \beta_k \mathbf{p}_k$$

$$k := k + 1$$

end repeat

The result is  $\mathbf{x}_{k+1}$

## The Algorithm

# Conjugate Gradient Method

## Convergence theorem



Define a subset of polynomials as

$$\Pi_k^* := \{ p \in \Pi_k : p(0) = 1 \} ,$$

where  $\Pi_k$  is the set of [polynomials](#) of maximal degree  $k$ .

Let  $(\mathbf{x}_k)_k$  be the iterative approximations of the exact solution  $\mathbf{x}_*$ , and define the errors as  $\mathbf{e}_k := \mathbf{x}_k - \mathbf{x}_*$ . Now, the rate of convergence can be approximated as [\[5\]](#)

$$\begin{aligned} \|\mathbf{e}_k\|_{\mathbf{A}} &= \min_{p \in \Pi_k^*} \|p(\mathbf{A})\mathbf{e}_0\|_{\mathbf{A}} \\ &\leq \min_{p \in \Pi_k^*} \max_{\lambda \in \sigma(\mathbf{A})} |p(\lambda)| \|\mathbf{e}_0\|_{\mathbf{A}} \\ &\leq 2 \left( \frac{\sqrt{\kappa(\mathbf{A})} - 1}{\sqrt{\kappa(\mathbf{A})} + 1} \right)^k \|\mathbf{e}_0\|_{\mathbf{A}}, \end{aligned}$$

$$\kappa(A) = \|A\| \|A^{-1}\|$$

where  $\sigma(\mathbf{A})$  denotes the [spectrum](#), and  $\kappa(\mathbf{A})$  denotes the [condition number](#).

Note, the important limit when  $\kappa(\mathbf{A})$  tends to  $\infty$

$$\frac{\sqrt{\kappa(\mathbf{A})} - 1}{\sqrt{\kappa(\mathbf{A})} + 1} \approx 1 - \frac{2}{\sqrt{\kappa(\mathbf{A})}} \quad \text{for } \kappa(\mathbf{A}) \gg 1.$$

This limit shows a faster convergence rate compared to the iterative methods of [Jacobi](#) or [Gauss-Seidel](#) which scale as  $\approx 1 - \frac{2}{\kappa(\mathbf{A})}$ .

# Krylov subspace methods

- CG proceeds by defining special search directions and minimizing

$$\|\mathbf{e}_k\|_A = \sqrt{\mathbf{e}_k^T A \mathbf{e}_k}, \text{ where } \mathbf{e}_k = \mathbf{x} - \mathbf{x}_k$$

$$f(\mathbf{x}) = \frac{1}{2} \mathbf{x}^T \mathbf{A} \mathbf{x} - \mathbf{x}^T \mathbf{b}, \quad \mathbf{x} \in \mathbf{R}^n.$$

- We can also minimize the norm of the residual

$$\|\mathbf{b} - \mathbf{A} \mathbf{x}_k\|_2$$

- MINRES or GMRES

# Image Operator with Laplacian Equation

$$\begin{aligned}\Delta f &= f_{xx} + f_{yy} \\&= f(x+1, y) - 2f(x, y) + f(x-1, y) + f(x, y+1) \\&\quad - 2f(x, y) + f(x, y-1) \\&= f(x+1, y) + f(x-1, y) + f(x, y+1) + f(x, y-1) - 4f(x, y).\end{aligned}$$

$$\left[ \begin{array}{ccccccccc} 1 & 1 & -4 & 1 & 1 & \dots & \dots \\ \dots & 1 & 1 & -4 & 1 & 1 & \dots \\ & & & \vdots & & & \\ \dots & \dots & 1 & 1 & -4 & 1 & 1 \end{array} \right] \left[ \begin{array}{c} f_1 \\ \vdots \\ f_n \end{array} \right] = \left[ \begin{array}{c} 0 \\ \vdots \\ c_1 \\ \vdots \\ c_k \end{array} \right]$$

$I_{K \times K}$

# Image Operator with Possion Equation

$$\left[ \begin{array}{ccccccccc} 1 & 1 & -4 & 1 & 1 & \dots & \dots \\ \dots & 1 & 1 & -4 & 1 & 1 & \dots \\ & & & \vdots & & & \\ & & & 1 & 1 & -4 & 1 & 1 \\ \hline 0 & & & & & I_{K \times K} & & \end{array} \right] \begin{bmatrix} f_1 \\ \vdots \\ f_n \end{bmatrix} = \begin{bmatrix} \text{div } \mathbf{g} \\ c_1 \\ \vdots \\ c_k \end{bmatrix}.$$

# Poisson Image Matting



$$I = \alpha F + (1 - \alpha)B$$

# Poisson Image Matting

$$\nabla I = (F - B) \nabla \alpha + \alpha \nabla F + (1 - \alpha) \nabla B$$

- Assume that  $F$  and  $B$  is very smooth

$$\nabla \alpha \approx \frac{1}{F - B} \nabla I$$

# Poisson Image Matting

- We minimize the following variational problem with dirichlet boundary  $\alpha|_{\partial\Omega} = \hat{\alpha}|_{\partial\Omega}$

$$\alpha^* = \arg \min_{\alpha} \iint_{p \in \Omega} \left\| \nabla \alpha_p - \frac{1}{F_p - B_p} \nabla I_p \right\|^2 dp$$

- Use Euler-Lagrange Equation

$$J = \int_a^b F(x, f(x), f'(x)) dx \quad \xrightarrow{\hspace{1cm}} \quad \frac{\partial F}{\partial f} - \frac{d}{dx} \frac{\partial F}{\partial f'} = 0$$

# Poisson Image Matting

- Use Euler-Lagrange Equation

$$J = \int_a^b F(x, f(x), f'(x)) dx \quad \xrightarrow{\text{blue arrow}} \quad \frac{\partial F}{\partial f} - \frac{d}{dx} \frac{\partial F}{\partial f'} = 0$$

- Poisson Equation

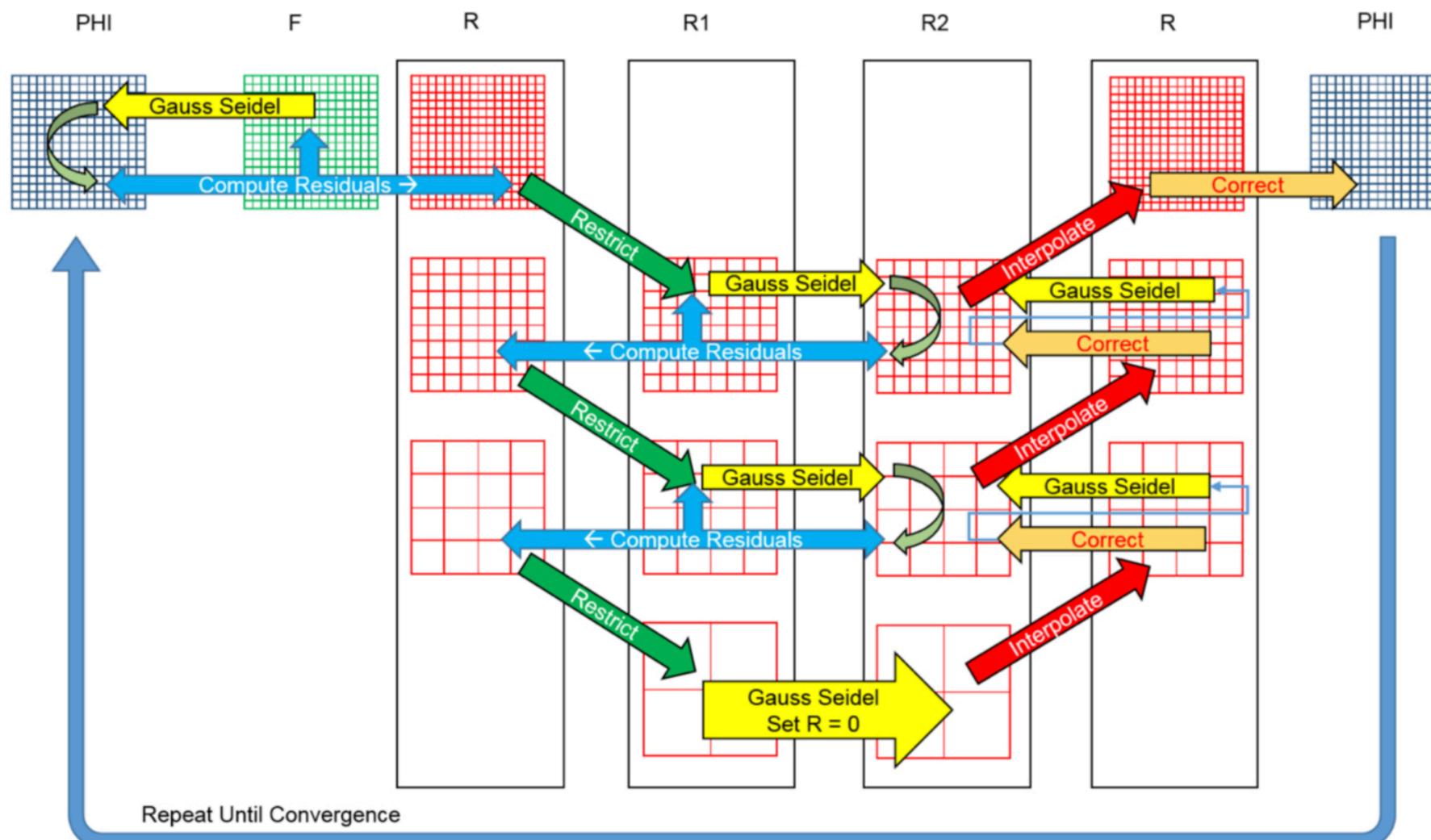
$$\Delta \alpha = \operatorname{div}\left(\frac{\nabla I}{F - B}\right)$$

Linear system !!!

# Multigrid Method(多重网格法)

- Use Gauss–Seidel as Smoother

Multigrid V-Cycle: Solving **PHI** in PDE  $f(\text{PHI}) = \mathbf{F}$

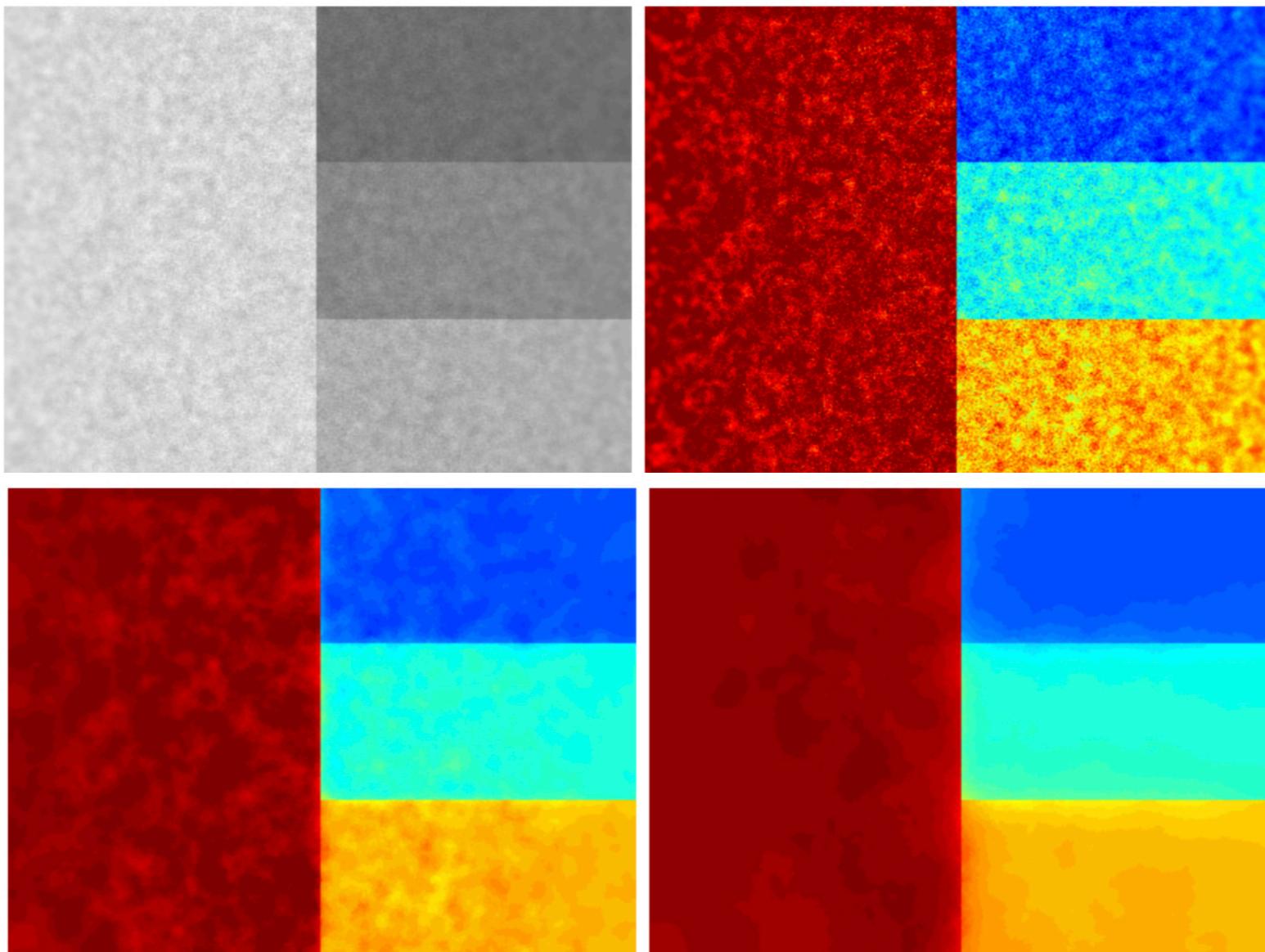


# Multigrid Method(多重网格法)

- Use Gauss–Seidel as Smoother
- The Complexity of this algorithm is  $O(n)$
- It is widely used in image processing and solving PDE

# Weighted Least Squares Filter

- Edge preserving Smoother by WLS



WLS:  $\alpha = 1.2, \lambda = 0.25$

WLS:  $\alpha = 1.8, \lambda = 0.35$

# Weighted Least Squares Filter

- Edge preserving Smoother by WLS

$$\sum_p \left( (u_p - g_p)^2 + \lambda \left( a_{x,p}(g) \left( \frac{\partial u}{\partial x} \right)_p^2 + a_{y,p}(g) \left( \frac{\partial u}{\partial y} \right)_p^2 \right) \right)$$

$$= (u - g)^T (u - g) + \lambda \left( u^T D_x^T A_x D_x u + u^T D_y^T A_y D_y u \right)$$

$$(I + \lambda L_g) u = g \quad L_g = D_x^T A_x D_x + D_y^T A_y D_y$$

$$a_{x,p}(g) = \left( \left| \frac{\partial \ell}{\partial x}(p) \right|^{\alpha} + \varepsilon \right)^{-1} \quad a_{y,p}(g) = \left( \left| \frac{\partial \ell}{\partial y}(p) \right|^{\alpha} + \varepsilon \right)^{-1}$$

# Weighted Least Squares Filter

- Edge preserving Smoother by WLS

$$(I + \lambda L_g)u = g \quad \text{Linear system !!!}$$

- Use Multigrid Method to solve in  $O(H \times W)$  time

# which method should we use?

|                             | Dense or Small | symmetric | positive definite |
|-----------------------------|----------------|-----------|-------------------|
| Gaussian Elimination        | +              |           |                   |
| Cholesky decomposition      | +              | +         | +                 |
| Iterative Method            | -              |           |                   |
| Conjugate gradient          | -              | +         | +                 |
| the Bunch–Kaufman algorithm | +              | +         | -                 |
| MINRES                      | -              | +         | -                 |
| GMRES                       | -              | -         |                   |

# Combination of these methods

- use Multigrid method to get a good initial solution
- then use conjugate gradient method to refine the solution