

Lecture 7

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1.1 Acyclic subgraph

Give a factor 1/2 algorithm for the following. Given a directed graph $G = (V, E)$, pick a maximum cardinality set of edges from E so that the resulting subgraph is acyclic.

Solve:

Arbitrarily number the vertices and divide the edges into two sets, the forward-going edges F and the backward-going edges B . Pick the bigger set, denoted by S .

Obviously, edges in F or B will not form a cycle, because a cycle needs at least a forward-going edge and a backward-going edge.

$$\begin{aligned} |F| + |B| &= |E| \\ S = \max\{F, B\} &\geq \frac{1}{2}|E| \end{aligned}$$

And we have $OPT < |E|$, hence

$$|S| \geq \frac{1}{2}OPT$$

Vertex Cover and DFS tree

1.3 (R. Bar-Yehuda) Consider the following factor 2 approximation algorithm for the cardinality vertex cover problem. Find a depth first search tree in the given graph, G , and output the set, say S , of all the nonleaf vertices of this tree. Show that S is indeed a vertex cover for G and $|S| \leq 2 \cdot OPT$

Solve

First we prove that S is a vertex cover. Assume that there exists an edge $\langle u, v \rangle$ in G and not covered by S . Since all nonleaf vertices in DFS tree are in S , u and v must be leaf vertices. This is contradictory to the definition of DFS tree, hence S is indeed a vertex cover.

Then we show the ratio is $\frac{1}{2}$. Construct a matching set T in the following way.

```
1  initialize T = NULL
2  denote by D the nonleaf vertices in DFS tree
3  while D is not empty do
4      choose vertex v in D closest to root
5      if v has a nonleaf child u in DFS tree
6          add match <v,u> to T
7          remove u,v from D
8      else
9          (then v's children are leaf vertices)
10         add <v> to T as unmatched single vertex
11 endwhile
12 return T
```

Note that the match T is different from maximum match problem, here we allow unmatched single vertex because we have ignore leaf vertices from the origin graph G . To cover edges joint with leaf vertices, we have to include all unmatched single vertices in T . For all matches in T , at least one of its endpoint is included in vertex cover. Therefore the ratio is $\frac{1}{2}$.