

Basic Scheduling

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Basic Scheduling

- Scheduling is an important class of discrete optimisation problems
- Basic scheduling involves:
 - tasks with a **duration**
 - **precedences** between tasks: one task must complete before another can start
- The aim is to schedule the tasks and usually to minimize the latest end time

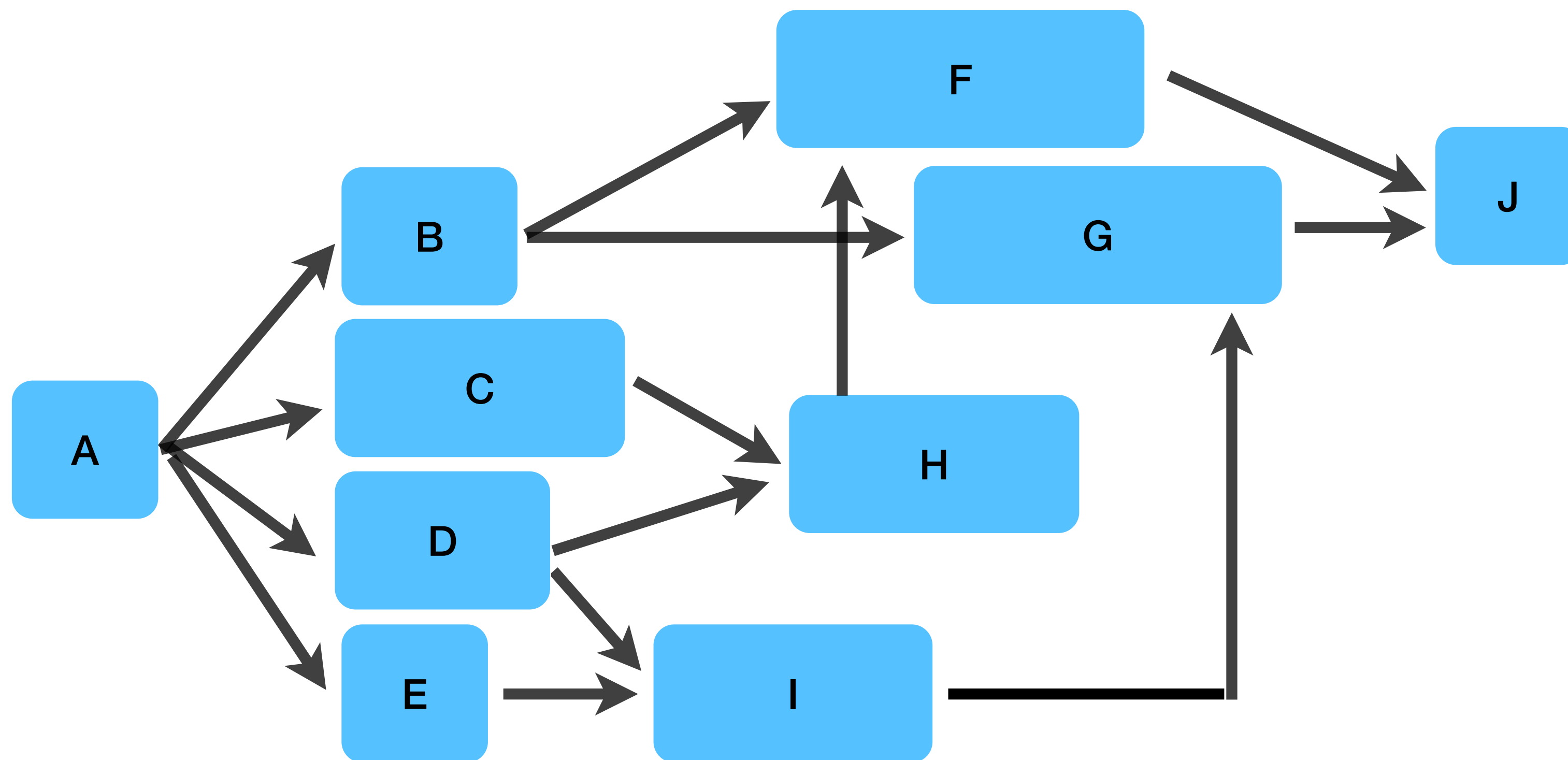
Modeling Time

- In discrete optimization time is modeled by integers (not continuous)
- Time variables tend to have VERY large ranges
 - e.g. start times on the minute for a 7 day schedule
- We typically only care about
 - earliest time, or
 - latest time
- when reasoning (not about all possible times)

Basic Scheduling Data & Decisions

```
enum TASK;  
  
array[TASK] of int: duration;  
array[int,1..2] of TASK: pre;  
set of int: PREC = index_set_1of2(pre);  
  
int: e = sum(duration);  
array[TASK] of var 0..e: start;
```

Basic Scheduling Instance



- Length indicates durations
- Arcs indicate precedences

Basic Scheduling Data & Decisions

TASK = {A, B, C, D, E, F, G, H, I, J};

duration = [10,10,20,15,10,25,20,20,20,10];

pre =

	A, B		A, C	
	A, D		A, E	
	B, F		B, G	
	F, J		G, J	
	C, H		D, H	
	D, I		E, I	
	H, F		I, G];

Constraints & Objective

- We can force the precedences of the tasks using

```
constraint forall(i in PREC)  
    (start[pre[i,1]] + duration[pre[i,1]]  
     <= start[pre[i,2]]);
```

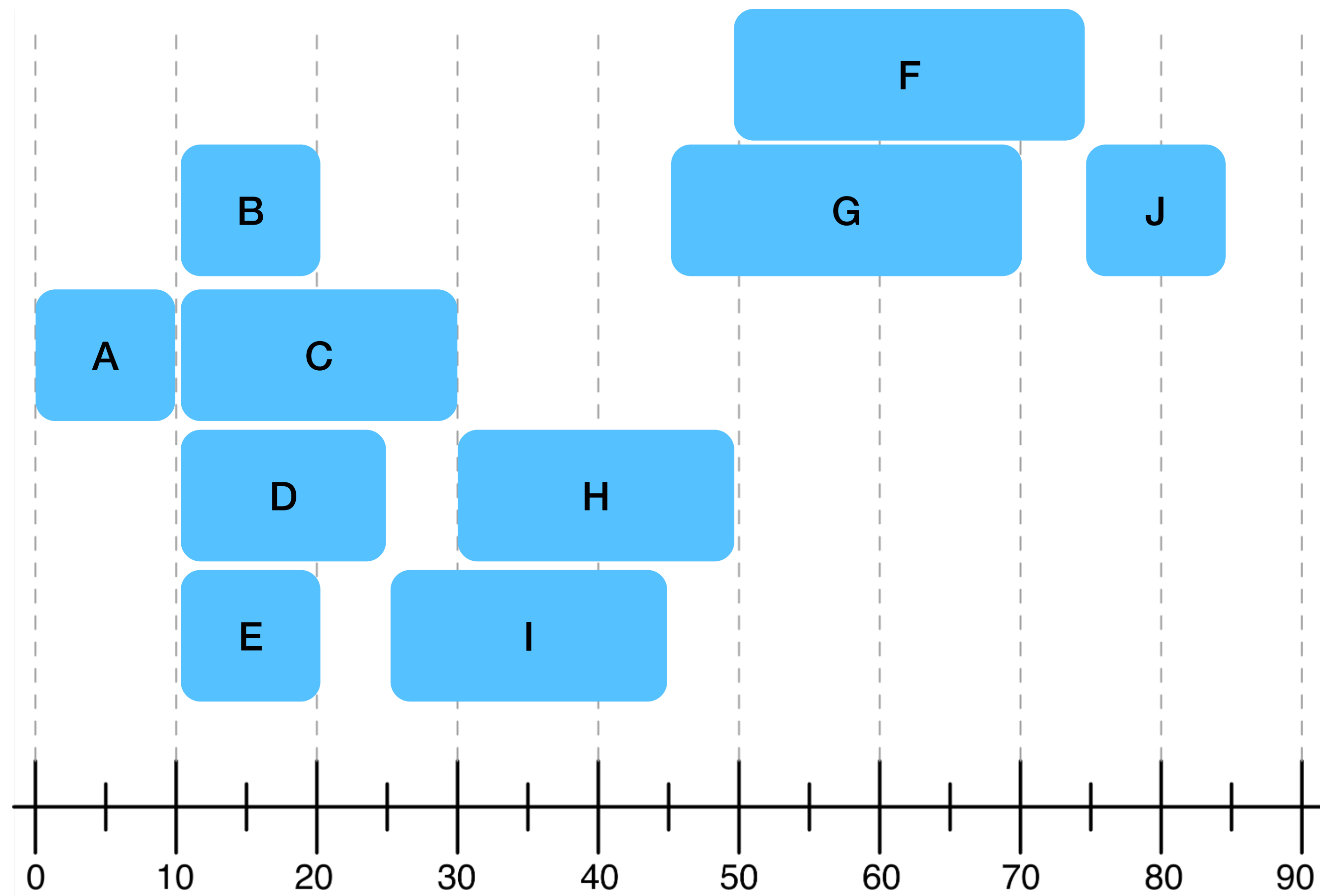
- To minimize the total time taken we use

```
var 0..e: makespan =  
    max(t in TASK)(start[t] + duration[t]);  
solve minimize makespan;
```

Constraints Generated

```
start[A] + 10 <= start[B]
start[A] + 10 <= start[C]
start[A] + 10 <= start[D]
start[A] + 10 <= start[E]
start[B] + 10 <= start[F]
start[B] + 10 <= start[G]
start[F] + 25 <= start[J]
start[G] + 20 <= start[J]
start[C] + 20 <= start[H]
start[D] + 15 <= start[H]
start[D] + 15 <= start[I]
start[E] + 10 <= start[I]
start[H] + 20 <= start[F]
start[I] + 20 <= start[J]
```


Basic Scheduling Solution



makespan = 85
 = [0, 10, 10, 10, 10, 50, 45, 30, 25, 75]

Difference Logic Constraints

- **Difference logic constraints** take the form $x + d \leq y$, where d is constant
- Note $x + d = y \leftrightarrow x + d \leq y \wedge y + (-d) \leq x$
- A problem consists solely of difference logic constraints can be rapidly solved as a **longest/shortest path problem**
- But adding extra constraints means this advantage disappears
 - e.g., at most two tasks can run simultaneously

Summary

- Basic scheduling problems are a common part of many complex discrete optimisation problems
 - tasks with precedences
- The constraints needed to model this are a **simple form** of linear constraints: **difference logic constraints**
- Problems involving only these constraints can be solved **very efficiently**

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