

# Optional Task Scheduling

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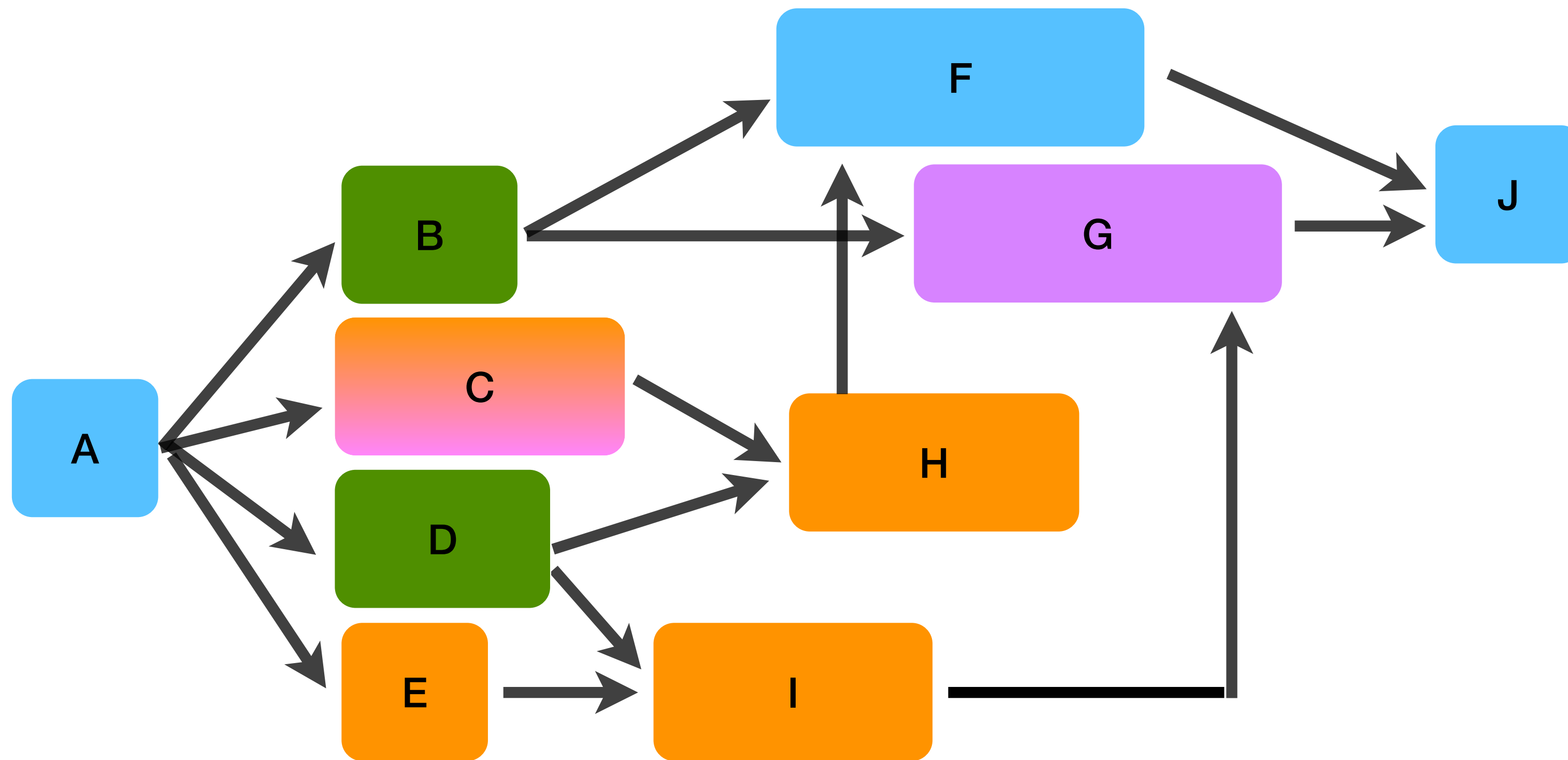
# Optional Tasks

- Sometimes we have two or more ways to complete the same task
  - e.g. the same task could be performed on two different machines (but with different duration)
- The usual way of modelling this in CP is using **optional tasks**
  - e.g. create two tasks (one for each machine) but only require one to be executed.
- We can model optional execution by
  - **Optional start time**:  $\langle \rangle$  if the task is not executed
  - **Variable duration**: 0 if not executed

# Optional Tasks

- Optional tasks if not executed
  - **Act as if they dont exist**
- Care must be taken to model this correctly
  - e.g. precedences among optional tasks
    - Should always hold if one of the tasks is not executed.
  - e.g. if t1 before t2 and t2 before t3 and t2 doesnt run?
    - One interpretation: no constraints
    - Another interpretation: t1 before t3 (transitive closure)

# Scheduling Instance Redux



- Suppose only one of C or G is required
- And only one of B and D

# Duration modelling of Optionality

- Optional tasks that don't execute have 0 duration
  - Implicitly implements the transitive closure

```

array[int] of set of TASK: options; % data

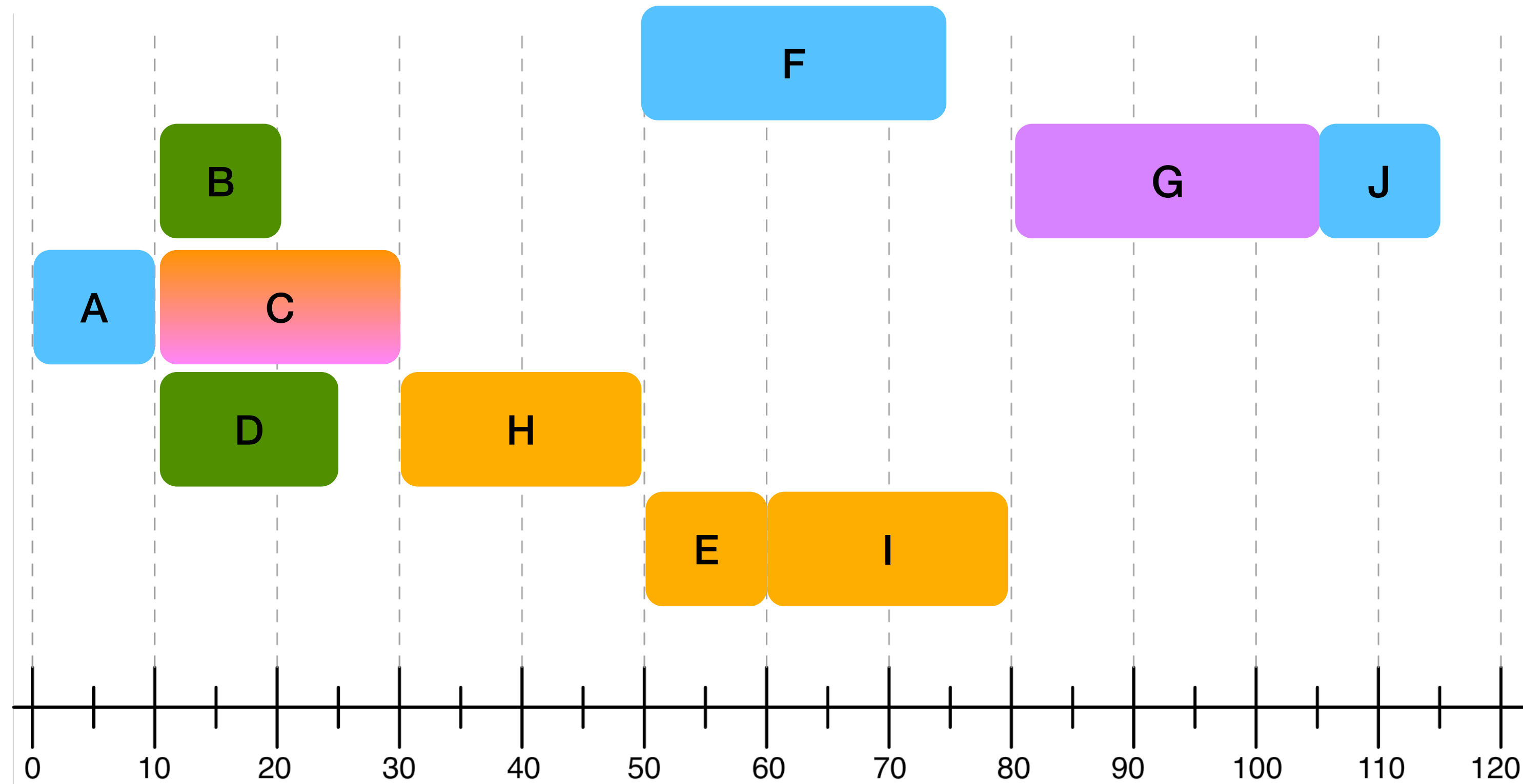
array[TASK] of var 0..e: start;
array[TASK] of var bool: runs; % which tasks run
array[TASK] of var 0..max(duration): aduration =
  [ runs[t]*duration[t] | t in TASK ];
  % replace duration by aduration in precs/disj

% every non optional task always runs
constraint forall(t in TASK diff array_union(options))
  (runs[t]);

% (at least) one task of each option set runs
constraint forall(i in index_set(options))
  (exists(t in options[i])(runs[t]));

```

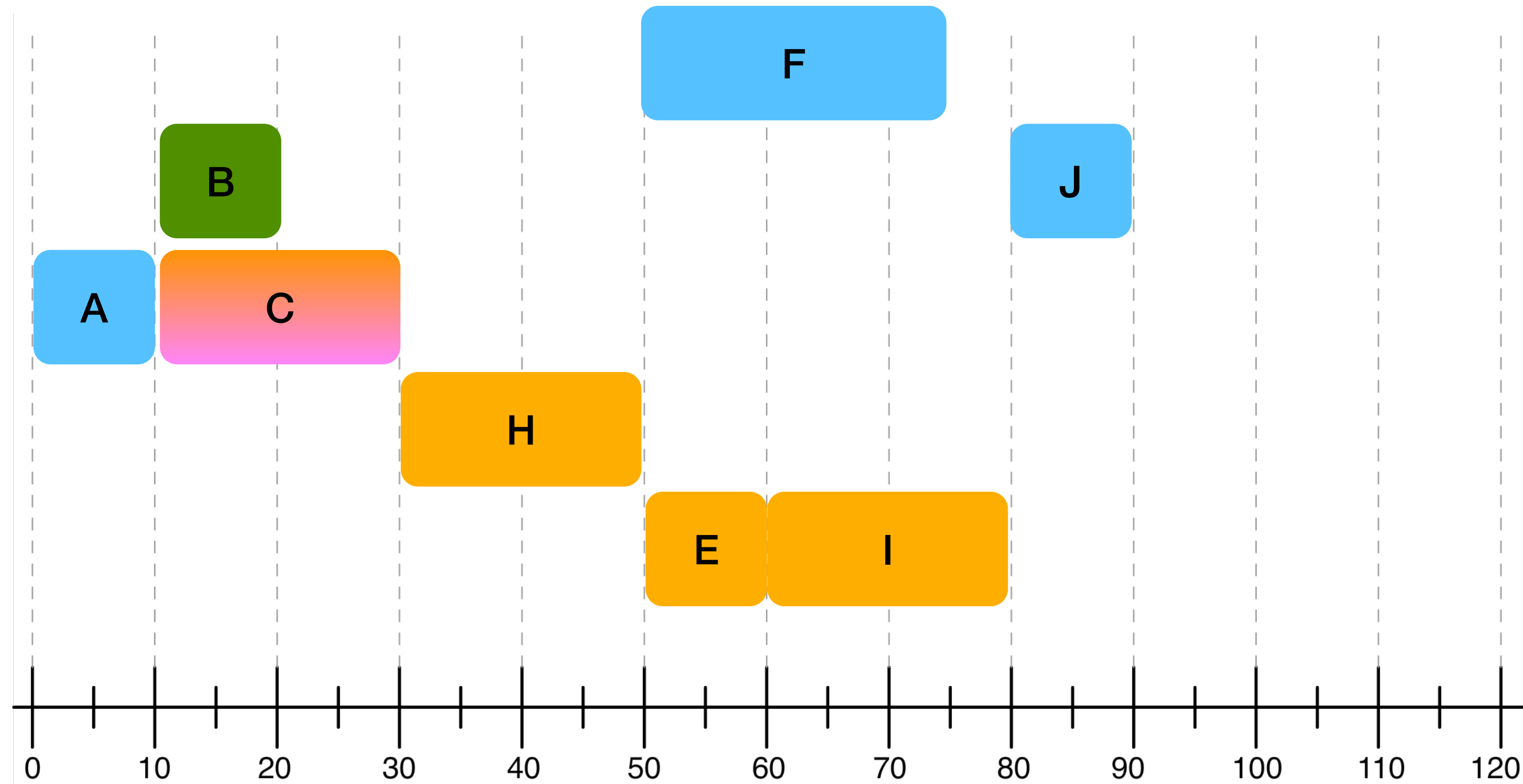
# Schedule without Optionality



makespan =  $[0, 10, 10, 10, 50, 50, 80, 30, 60, 105]$

**115**

# Scheduling with Optionality (0 duration)



makespan = ~~90~~ = [0, 10, 10, 10, 50, 50, 80, 30, 60, 80]

A    B    C    D    E    F    G    H    I    J

# Optional Start Time modelling of Optionality

- MiniZinc supports optional variables
  - Extra value `<>` means “absent”
  - Variables taking value `<>` act as if “not there”
  - `occurs(x)` forces `x` to not take value `<>`
- The optional variable modelling simply changes

```
array[int] of set of TASK: options; % data
```

```
array[TASK] of var opt 0..e: start;
```

```
% every non optional task always occurs
```

```
constraint forall(t in TASK diff array_union(options))
    (occurs(start[t]));
```

```
% (at least) one task of each option set occurs
```

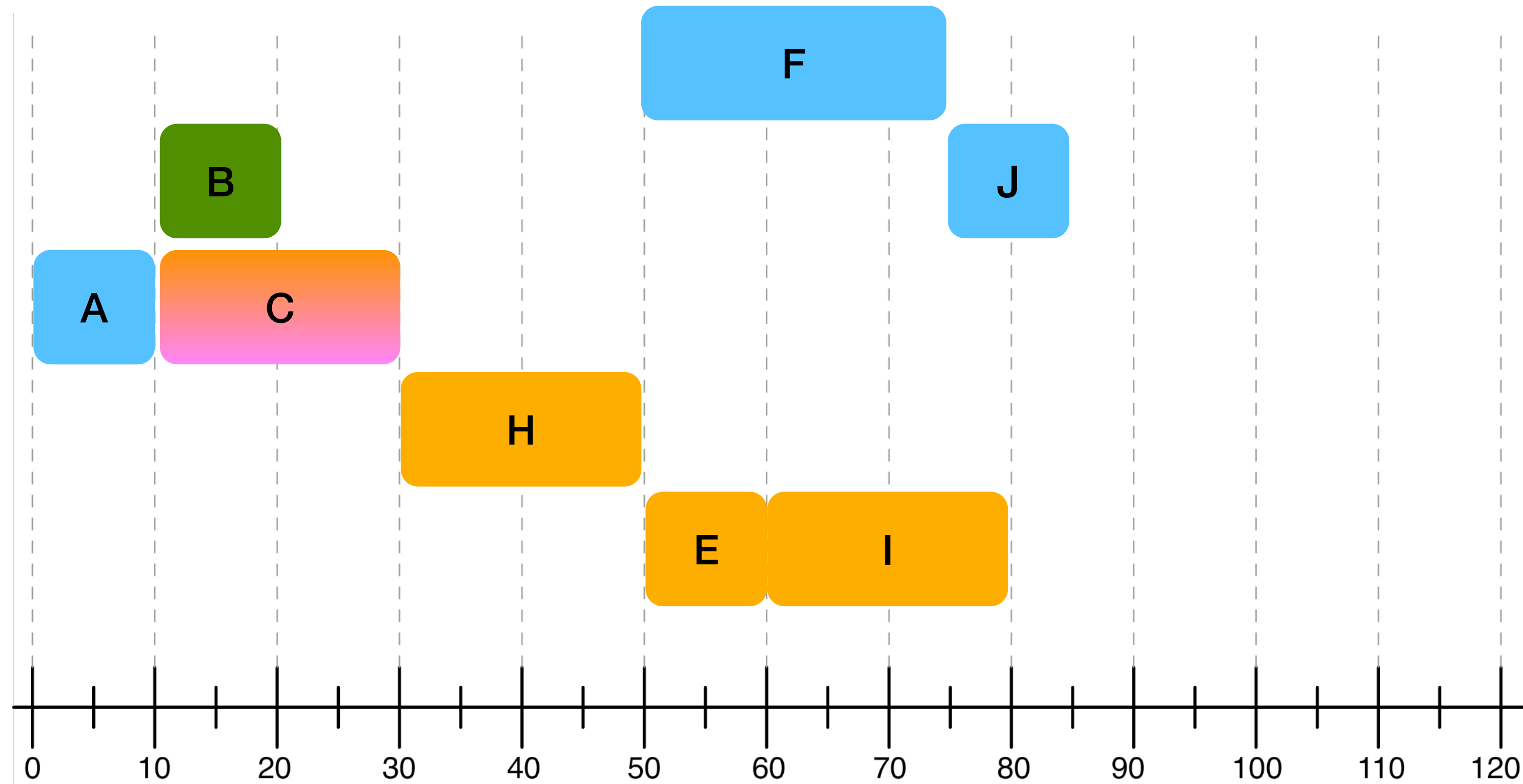
```
constraint forall(i in index_set(options))
    (exists(t in options[I])(occurs(start[t])));
```

Optional keyword

occurs



# Scheduling with Optionality (opt)



Lost the transitive constraint  
 $I \ll G \ll J$

makespan =  $\begin{matrix} A & B & C & D & E & F & G & H & I & J \\ 85 & = & [0, & 10, & 10, & <>, & 50, & 50, & <>, & 30, & 60, & 75] \end{matrix}$

# Task Variables

- Many CP solvers support **task variables**
- Record of: (occurs, start, duration, end)
  - Optionality
  - $\text{end} = \text{start} + \text{duration}$
- They also have complex propagators to reason about them
  - Particular interaction of optionality + resources

# Optional Task Scheduling vs Temporal Planning

- Suppose we generate  $k_a$  optional tasks for each possible action  $a$
- The optional task scheduling problem can be used to represent a **(bounded) temporal planning problem**.
- **Good idea:**
  - Small bounds, and shared resources for tasks
- **Bad idea:**
  - Complex fluents, large numbers  $k_a$

# Summary

- Optional Tasks / Task Variables
  - Are a key to modelling complex scheduling problems
  - Commercial CP solvers spend a lot of effort to deal with them well
- Optional Task Scheduling  $\approx$  Planning
- Task variables are coming to MiniZinc soonish!

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