Matchings II: Finding the maximum
COMS20017 (Algorithms and Data)

John Lapinskas, University of Bristol

An algorithm for maximum matchings

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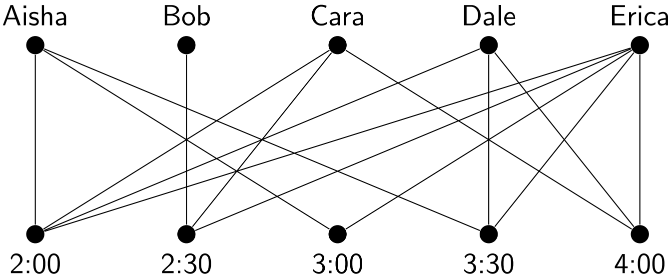
Can we form M by greedily adding edges? E.g.:



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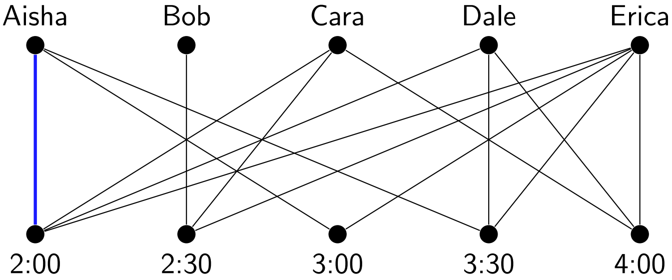

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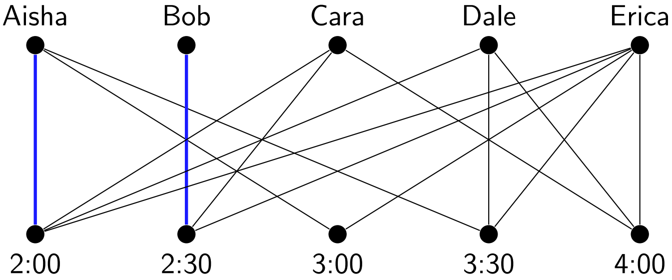

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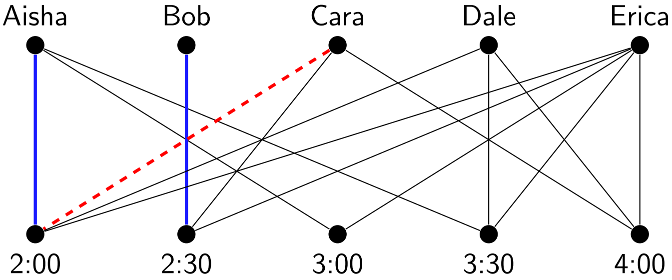

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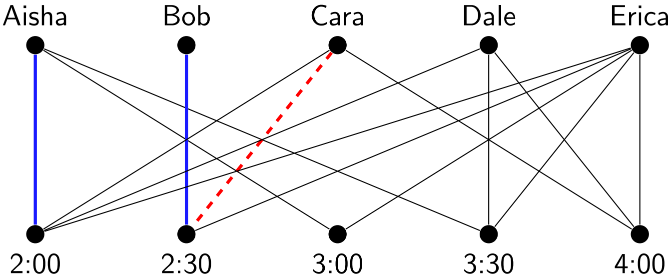

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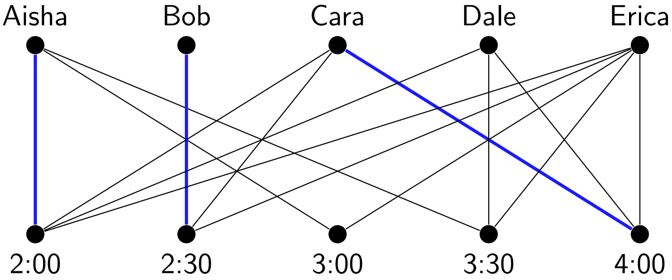

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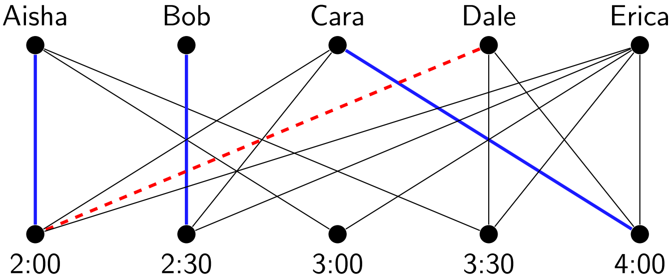

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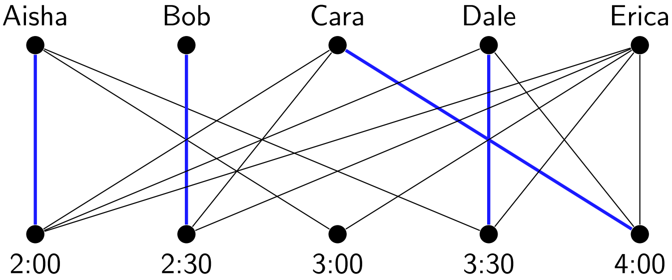

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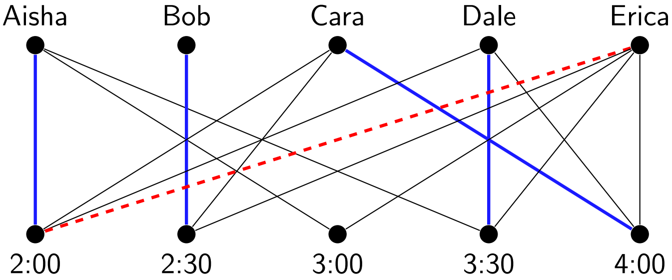

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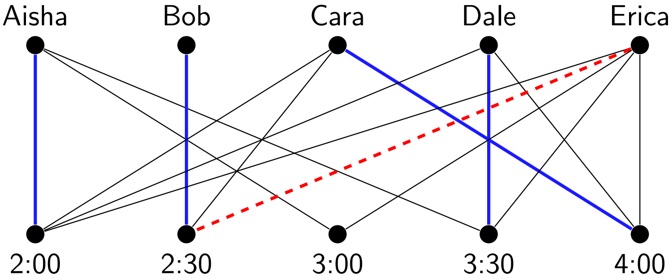

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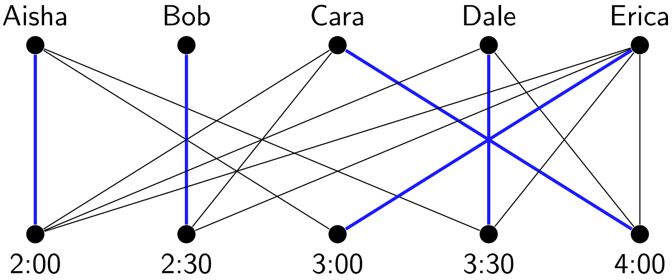

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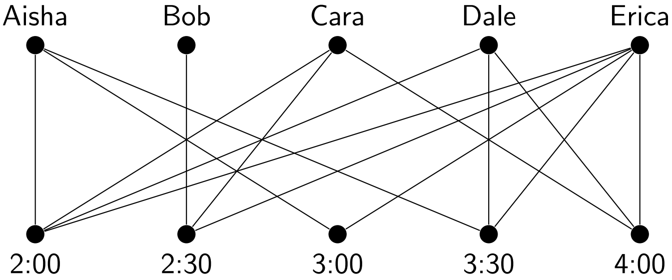

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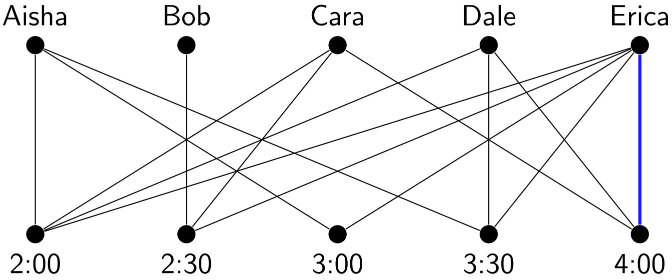

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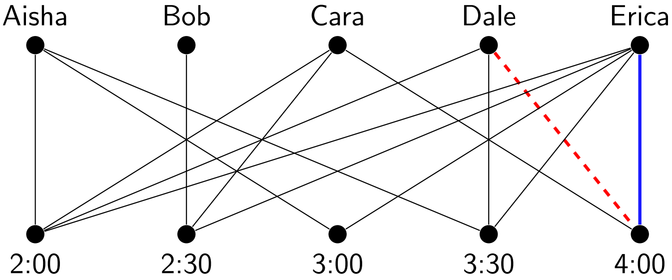

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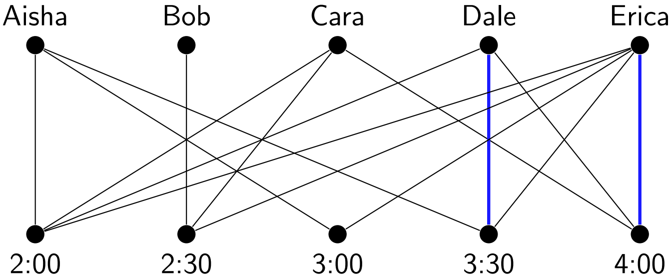

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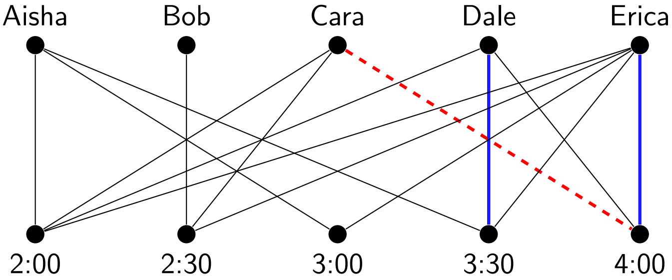

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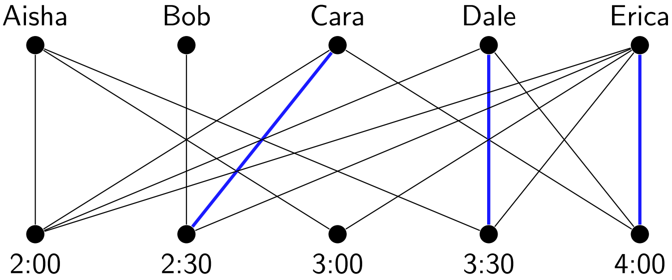

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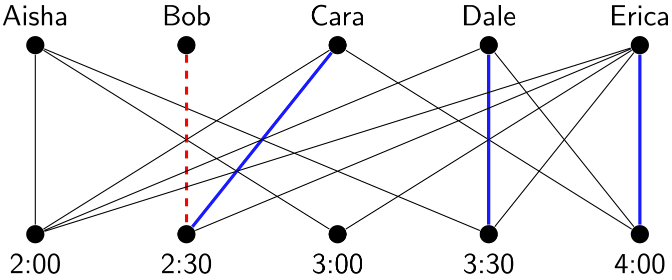

But if we had considered edges a different order... we wouldn't have been able to match Bob! So this algorithm **fails**.

Repairing poor decisions: an example

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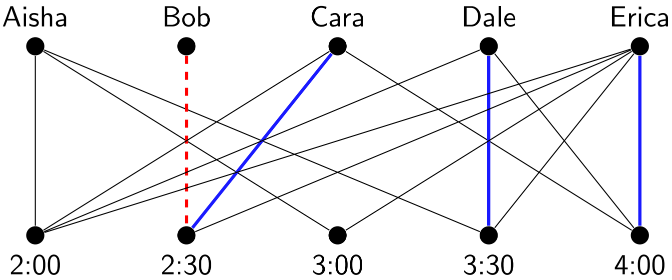

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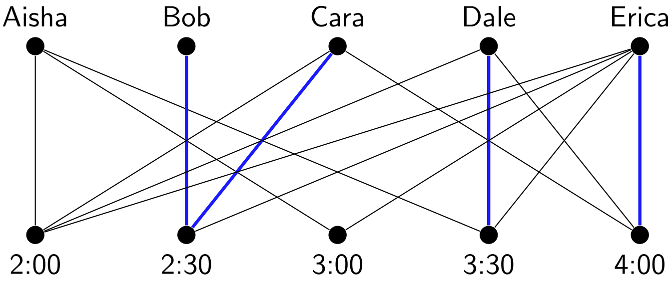

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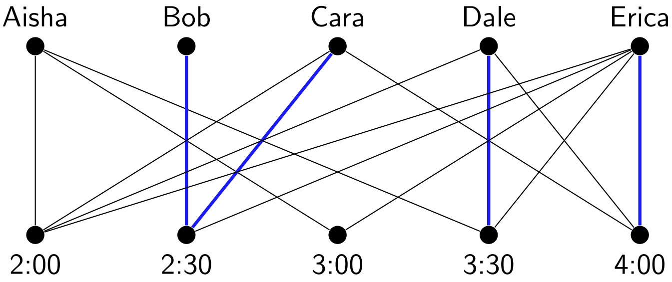

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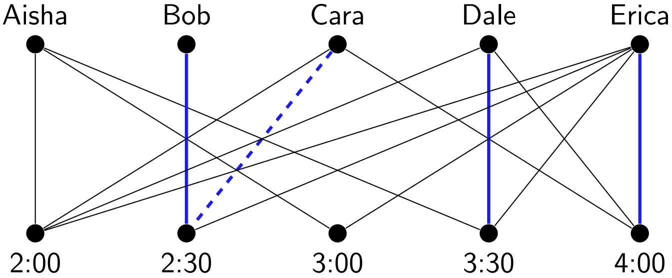

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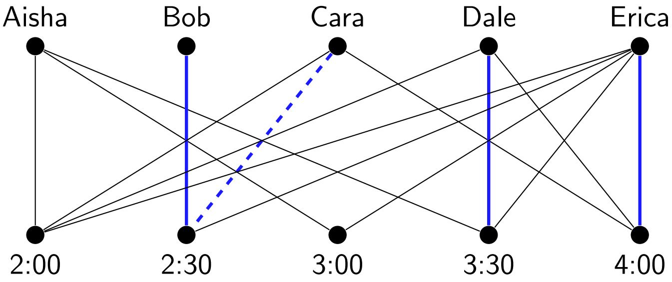

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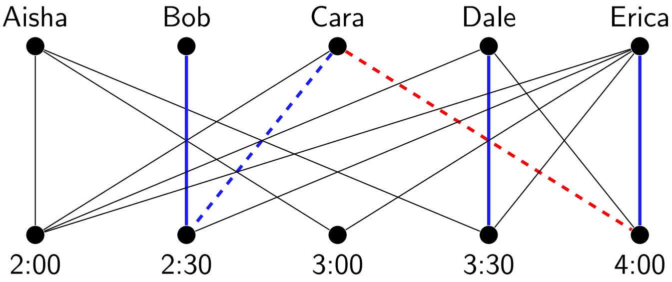

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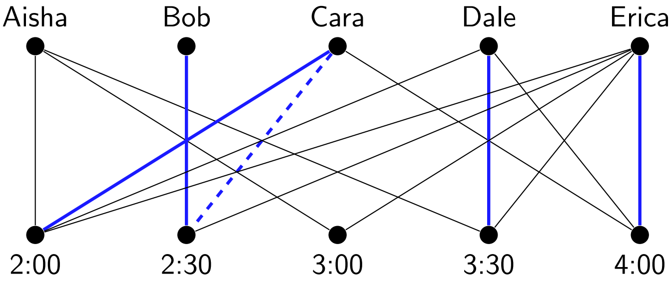

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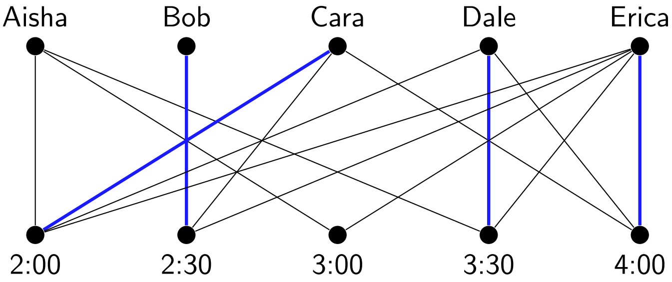

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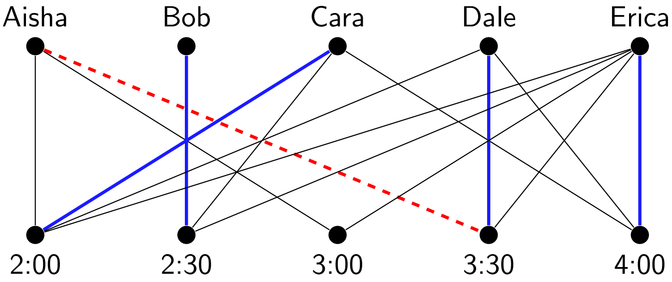

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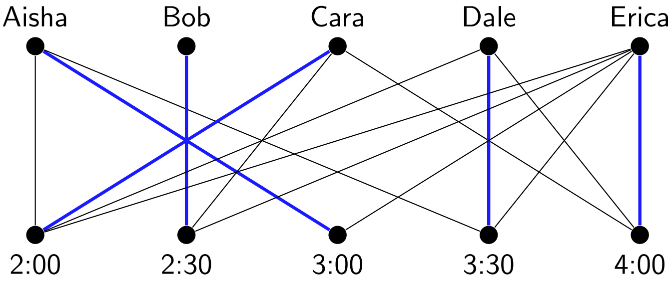

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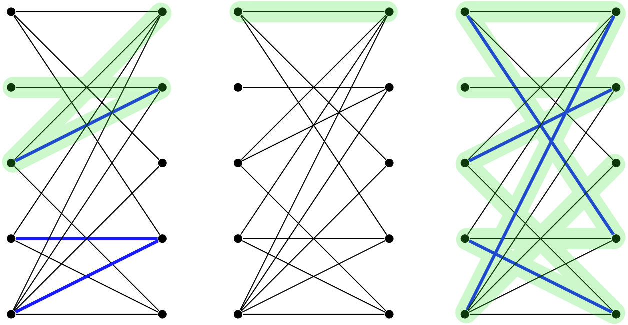
Formally, writing $P = v_0 \dots v_k$, we require $\{v_i, v_{i+1}\} \in M$ for all odd i, $\{v_i, v_{i+1}\} \notin M$ for all even i, and $v_0, v_k \notin \bigcup_{e \in M} e$. For example:



Given a matching M in a bipartite graph G, an **augmenting path** for M is a path $P = v_0 \dots v_k$ such that:

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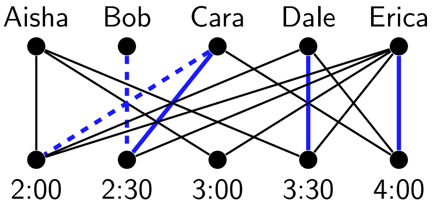
If P is an augmenting path for M, we define

$$\text{Switch}(M, P) = M - \{\{v_i, v_{i+1}\} : i \text{ is odd}\} \cup \{\{v_i, v_{i+1}\} : i \text{ is even}\}.$$

Then Switch(M, P) is a matching containing one more edge than M.



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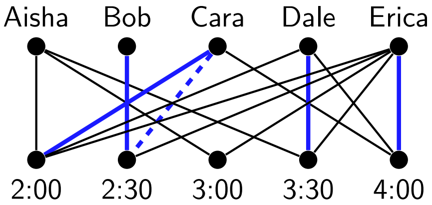

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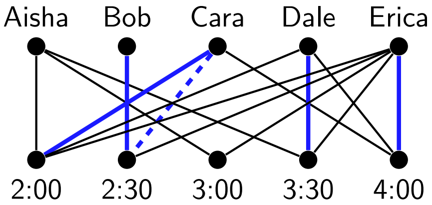

This suggests a new greedy algorithm!

A **correct** algorithm for maximum matchings

Algorithm: MAXMATCHING (SKETCH)

Input : A bipartite graph $G = (V, E)$.

Output : A list of edges forming a matching in G of maximum size.

begin
- Initialise $M \leftarrow []$, the empty matching.
- **while** *G contains an augmenting path for M* **do**
  - Find an augmenting path P for M.
  - Update $M \leftarrow \text{Switch}(M, P)$.
- Return M.

To make this work, we need to do two things:

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To make this work, we need to do two things:

- Find an efficient way to find an augmenting path whenever one exists.
- Prove that if M has no augmenting paths, then M is maximum.

Finding augmenting paths efficiently

If we search by brute force, this could take $\Theta(|V|!)$ time! Let's not.

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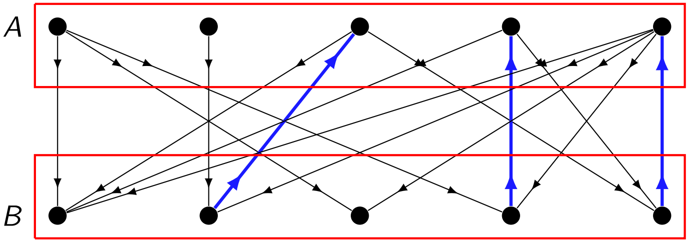
Here, we reduce the problem of finding an augmenting path to a problem we can already solve: finding a path from one set to another in a **directed** graph, via breadth-first search.

(For how to apply breadth-first search to sets instead of vertices, see last week's problem sheet — this is itself a reduction!)

Suppose $G = (V, E)$ has a matching M and a bipartition (A, B).
Turn G into an auxiliary digraph $D_{G,M}$ by directing non-matching edges from A to B and matching edges from B to A. Formally:

$$\begin{aligned} V(D_{G,M}) &:= V, \\ E(D_{G,M}) &:= \{(a,b) \colon a \in A,\ b \in B,\ \{a,b\} \in E \setminus M\} \cup \\ &\qquad \{(b,a) \colon a \in A,\ b \in B,\ \{a,b\} \in M\}. \end{aligned}$$



$D_{G,M}$ is defined by directing edges outside M from A to B,
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Lemma: A path in G is augmenting for M if and only if it's also a path from $U \cap A$ to $U \cap B$ in $D_{G,M}$.

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Proof: First note any augmenting path in G has endpoints in $U \cap A$ and $U \cap B$, since it has an odd number of edges and G is bipartite.

So let $P = v_0 \dots v_k$ be any path in G with $v_0 \in U \cap A$, $v_k \in U \cap B$.
We show P is augmenting for M iff it is also a path in $D_{G,M}$.

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Lemma: A path in G is augmenting for M if and only if it's also a path from $U \cap A$ to $U \cap B$ in $D_{G,M}$.

Proof: First note any augmenting path in G has endpoints in $U \cap A$ and $U \cap B$, since it has an odd number of edges and G is bipartite.

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Algorithm: MaxMatching

Input : A bipartite graph $G = (V, E)$.

Output : A list of edges forming a matching in G of maximum size.

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- Find a bipartition (A, B) of G. Initialise $M \leftarrow [\,]$.
- **repeat**
  - Form the graph $D_{G,M}$.
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Invariant: At the start of the ith loop iteration, M is a matching with $i - 1$ edges. M can have at most $|V|/2$ edges in total, so $\textsc{MaxMatching}$ outputs a matching with no augmenting paths.

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So overall the running time is $O(|E||V|)$.

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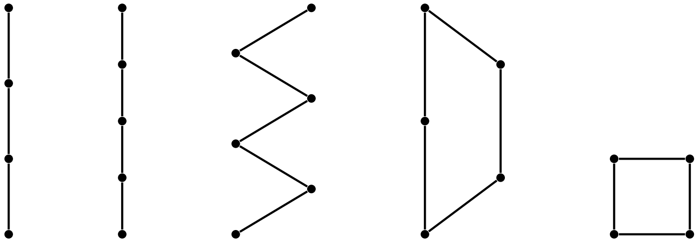
Let M' be another matching which **is** maximum, so $|M'| > |M|$.
Consider the symmetric difference $S = M \triangle M'$, i.e. the graph formed of edges contained in either M or M' but not both.



Since each vertex is in at most one M edge and at most one M' edge, S has maximum degree at most 2.

So S is a disjoint union of path and cycle components.

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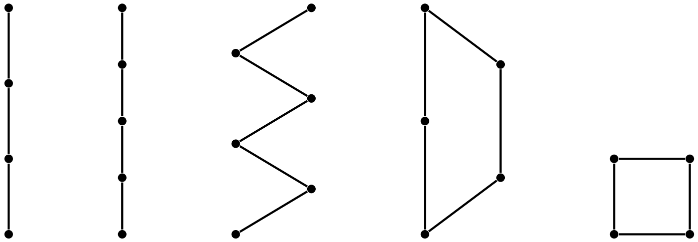

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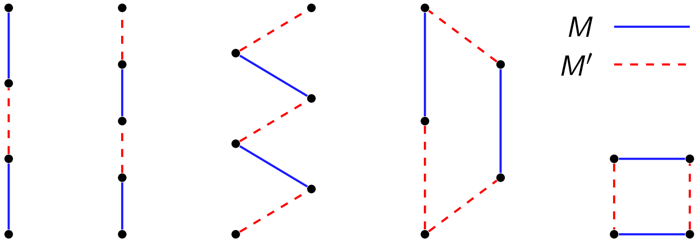

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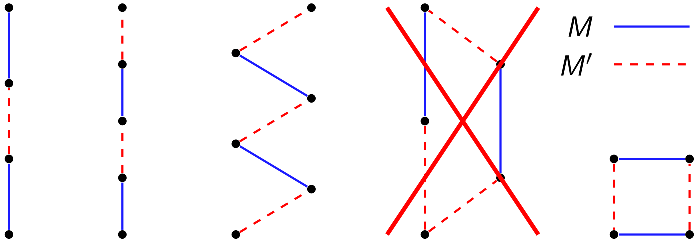

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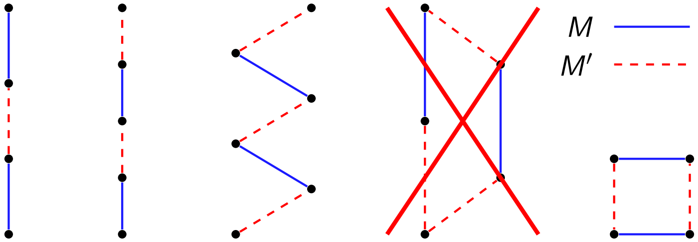

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Since M'-edges and M-edges alternate, it has exactly **one** more M'-edge.

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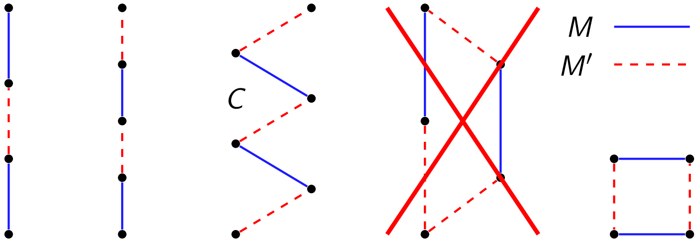

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Let us celebrate with a matching pair of kittens.

D'awww.