$O(n^2)$ sorting algorithms

- Selection sort and insertion sort are both $O(n^2)$
- $O(n^2)$ sorting is infeasible for n over 5000

A different strategy?

- Divide array in two equal parts
- Separately sort left and right half
- Combine the two sorted halves to get the full array sorted

Combining sorted lists

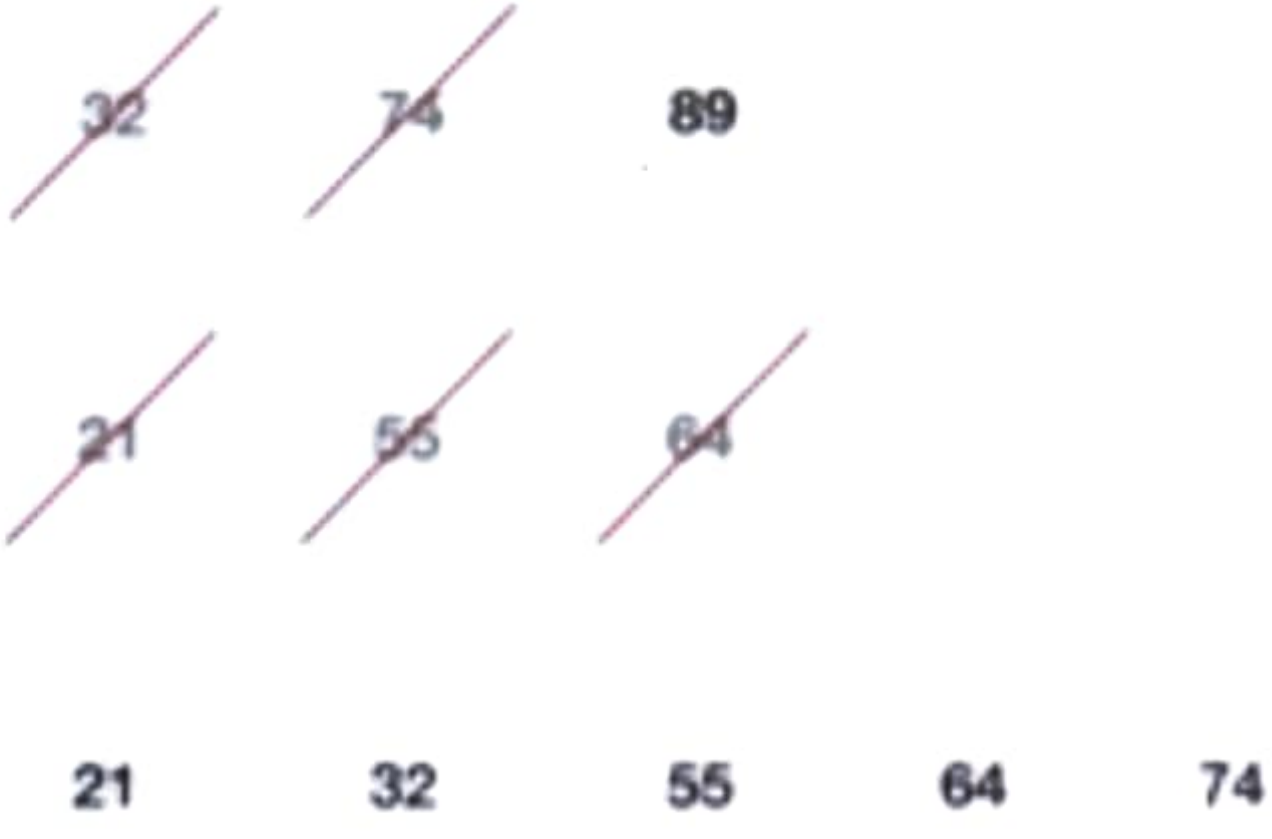
- Given two sorted lists A and B, combine into a sorted list C
  - Compare first element of A and B
  - Move it into C
  - Repeat until all elements in A and B are over
- Merging A and B

Merging two sorted lists

32 **74** **89**

21 **55** **64**

Merging two sorted lists



Merge Sort


43 32 22 78 63 57 91 13
43 32 22 78 63 57 91 13
43 32 22 78 63 57 91 13
43 32 22 78 63 57 91 13

Divide and conquer

- Break up problem into disjoint parts
- Solve each part separately
- Combine the solutions efficiently

Merging sorted lists

Combine two sorted lists A and B into C

- If A is empty, copy B into C
- If B is empty, copy A into C
- Otherwise, compare first element of A and B and move the smaller of the two into C
- Repeat until all elements in A and B have been moved

Merging

```
def merge(A,B): # Merge A[0:m],B[0:n]
  (C,m,n) = ([],len(A),len(B))
  (i,j) = (0,0) # Current positions in A,B
  while i+j < m+n: # i+j is number of elements merged so far
    if i == m:  # Case 1: A is empty
      C.append(B[j])
      j = j+1
    elif j == n: # Case 2: B is empty
      C.append(A[i])
      i = i+1
    elif A[i] <= B[j]: # Case 3: Head of A is smaller
      C.append(A[i])
      i = i+1
    elif A[i] > B[j]: # Case 4: Head of B is smaller
      C.append(B[j])
      j = j+1
  return(C)
```

Divide and conquer

- Break up problem into disjoint parts
- Solve each part separately
- Combine the solutions efficiently

Merging sorted lists

Combine two sorted lists A and B into C

- If A is empty, copy
- If B is empty, copy A
- Otherwise, compare first element of A and B and move the smaller of the two into C
- Repeat until all elements in A and B have been moved

Merging

```
def merge(A,B): # Merge A[0:m],B[0:n]
  (C,m,n) = ([],len(A),len(B))
  (i,j) = (0,0) # Current positions in A,B
  while i+j < m+n: # i+j is number of elements merged so far
    if i == m:  # Case 1: A is empty
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      C.append(B[j])
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  return(C)
```

```
adhavan@dolphinair:...ul/week4/python/mergesort$ more merge.py
ef merge(A,B): # Merge A[0:m],B[0:n]

    (C,m,n) = ([],len(A),len(B))
    (i,j) = (0,0) # Current positions in A,B

    while i+j < m+n: # i+j is number of elements merged so far

        if j == n:  # Case 1: A is empty
            C.append(A[i])
            i = i+1
        elif i == m:  # Case 2: B is empty
            C.append(B[j])
            j = j+1
        elif A[i] <= B[j]:  # Case 3: Head of A is smaller
            C.append(A[i])
            i = i+1
        elif A[i] > B[j]: # Case 4: Head of B is smaller
            C.append(B[j])
            j = j+1

    return(C)
```

```
>>> from merge import *
>>> a = list(range(0,100,2))
>>> b = list(range(1,75,2))
>>> len(a)
50
>>> len(b)
37
>>> a
[0, 2, 4, 6, 8, 10, 12, 14, 16, 18, 20, 22, 24, 26, 28, 30, 32, 34, 36, 38, 40, 42, 44
, 46, 48, 50, 52, 54, 56, 58, 60, 62, 64, 66, 68, 70, 72, 74, 76, 78, 80, 82, 84, 86,
88, 90, 92, 94, 96, 98]
>>> b
[1, 3, 5, 7, 9, 11, 13, 15, 17, 19, 21, 23, 25, 27, 29, 31, 33, 35, 37, 39, 41, 43, 45
, 47, 49, 51, 53, 55, 57, 59, 61, 63, 65, 67, 69, 71, 73]
>>> merge(a,b)
[0, 1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11, 12, 13, 14, 15, 16, 17, 18, 19, 20, 21, 22, 23,
 24, 25, 26, 27, 28, 29, 30, 31, 32, 33, 34, 35, 36, 37, 38, 39, 40, 41, 42, 43, 44, 4
5, 46, 47, 48, 49, 50, 51, 52, 53, 54, 55, 56, 57, 58, 59, 60, 61, 62, 63, 64, 65, 66,
 67, 68, 69, 70, 71, 72, 73, 74, 76, 78, 80, 82, 84, 86, 88, 90, 92, 94, 96, 98]
>>> len(merge(a,b))
87
>>>
```

Merging, wrong

```
def mergewrong(A,B): # Merge A[0:m],B[0:n]
  (C,m,n) = ([],len(A),len(B))
  (i,j) = (0,0) # Current positions in A,B

  while i+j < m+n:
  # i+j is number of elements merged so far

    # Combine Case 1, Case 4
    if i == m or A[i] > B[j]:
      C.append(B[j])
      j = j+1

    # Combine Case 2, Case 3:
    elif j == n or A[i] <= B[j]:
      C.append(A[i])
      i = i+1

  return(C)
```

```
madhavan@dolphinsir:...ul/week4/python/mergesort$ more mergewrong.py
def merge(A,B): # Merge A[0:m],B[0:n]

    (C,m,n) = ([],len(A),len(B))
    (i,j) = (0,0) # Current positions in A,B

    while i+j < mn: # i+j is number of elements merged so far

        if i == m or A[i] > B[j]:  # Combine Case 1 and 4
            C.append(B[j])
            j = j+1
        elif j == n or A[i] <= B[j]: # Combine Case 2 and 3
            C.append(A[i])
            i = i+1

    return(C)
madhavan@dolphinsir:...ul/week4/python/mergesort$
```

```
>>> a = [2,4,6]
>>> b = [1,3,5]
>>> merge(a,b)
3 3 0 0
3 3 0 1
3 3 1 1
3 3 1 2
3 3 2 2
3 3 3 3
Traceback (most recent call last):
  File "<stdin>", line 1, in <module>
  File "/Users/madhavan/mirror/projects/NPTEL/python-2016-jul/week4/python/mergesort/m
ergemerge.py", line 10, in merge
    if i == n or A[i] > B[j]:  # Combines Case 1 and 4
IndexError: list index out of range
>>>
```

Merge Sort

To sort `A[0:n]` into `B[0:n]`

- If n is 1, nothing to be done
- Otherwise
  - Sort `A[0:n//2]` into `L` (left)
  - Sort `A[n//2:n]` into `R` (right)
  - Merge `L` and `R` into `B`

```
        elif j == n:  # Case 2: B is empty
            C.append(A[i])
            i = i+1
        elif A[i] <= B[j]:  # Case 3: Head of A is smaller
            C.append(A[i])
            i = i+1
        elif A[i] > B[j]: # Case 4: Head of B is smaller
            C.append(B[j])
            j = j+1

    return(C)


def mergesort(A,left,right):    # Sort the slice A[left:right]

    if right - left <= 1: # Base case
        return(A[left:right])

    if right - left > 1:  # Recursive call
        mid = (left+right)//2
        L = mergesort(A,left,mid)
        R = mergesort(A,mid,right)
        return(merge(L,R))
mashaward@dolphinair:...ul/week4/python/mergesort$ python3.5
Python 3.5.2 (v3.5.2:4def2a2901a5, Jun 26 2016, 10:47:25)
[GCC 4.2.1 (Apple Inc. build 5666) (dot 3)] on darwin
Type "help", "copyright", "credits" or "license" for more information.
>>>
```

Analysis of Merge

How much time does Merge take?

- Merge A of size m, B of size n into C
- In each iteration, we add one element to C
  - Size of C is m+n
  - $m+n \leq 2 \max(m,n)$
- Hence $O(\max(m,n)) = O(n)$ if $m \approx n$

Merge Sort: Shortcomings

- Merging A and B cr[illegible]w array C
  - No obvious way t[illegible]ly merge in place
- Extra storage can be costly
- Inherently recursive
  - Recursive call and return are expensive

Alternative approach

- Extra space is required to merge
- Merging happens because elements in left half must move right and vice versa
- Can we divide so that everything to the left is smaller than everything to the right?
  - No need to merge!

Divide and conquer without merging

- Suppose the median value in A is m
- Move all values $\leq$ m to left half of A
  - Right half has values $>$ m
  - This shifting can be done in place, in time O(n)

Divide and conquer without merging

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- Recursively sort left and right halves
- A is now sorted! No need to merge
  - $T(n) = 2T(n/2) + n = O(n \log n)$

Quicksort in Python

```
def Quicksort(A,l,r): # Sort A[l:r]
  if r - l <= 1:  # Base case
     return ()

  # Partition with respect to pivot, a[l]
  yellow = l+1

  for green in range(l+1,r):
     if A[green] <= A[l]:
        (A[yellow],A[green]) = (A[green],A[yellow])
        yellow = yellow + 1

  # Move pivot into place
  (A[l],A[yellow-1]) = (A[yellow-1],A[l])

  Quicksort(A,l,yellow-1)  # Recursive calls
  Quicksort(A,yellow,r)
```

Merge Sort: Shortcomings

- Merging `A` and `B` creates a new array `C`
  - No obvious way to efficiently merge in place
- Extra storage can be costly
- Inherently recursive
  - Recursive call and return are expensive

Alternative approach

- Extra space is required to merge
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Divide and conquer without merging

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Quicksort

- Choose a pivot element
  - Typically the first value in the array
- Partition A into lower and upper parts with respect to pivot
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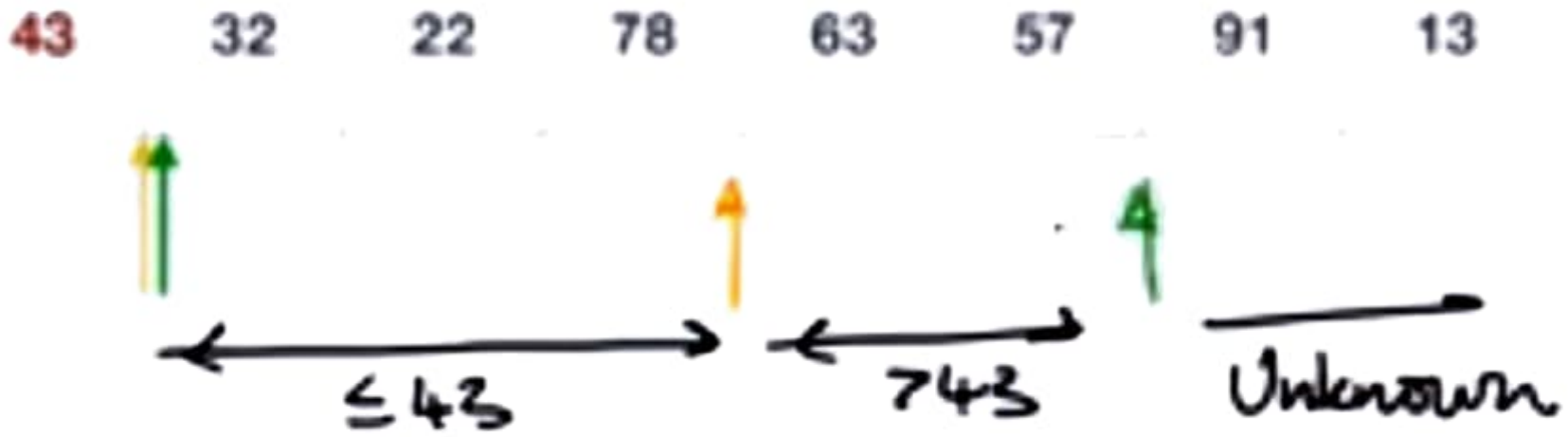
Quicksort

- **High level view**

13	32	22	43	63	57	91	78

Quicksort: Partitioning



Quicksort in Python

```
def Quicksort(A,l,r): # Sort A[l:r]
  if r - l <= 1:   # Base case
     return ()

  # Partition with respect to pivot, a[l]
  yellow = l+1

  for green in range(l+1,r):
     if A[green] <= A[l]:
        (A[yellow],A[green]) = (A[green],A[yellow])
        yellow = yellow + 1

  # Move pivot into place
  (A[l],A[yellow-1]) = (A[yellow-1],A[l])

  Quicksort(A,l,yellow-1)  # Recursive calls
  Quicksort(A,yellow,r)
```

85, 86, 87, 88, 89, 90, 91, 92, 93, 94, 95, 96, 97, 98, 99, 100, 101, 10
104, 105, 106, 107, 108, 109, 110, 111, 112, 113, 114, 115, 116, 117, 11
120, 121, 122, 123, 124, 125, 126, 127, 128, 129, 130, 131, 132, 133, 13
136, 137, 138, 139, 140, 141, 142, 143, 144, 145, 146, 147, 148, 149, 15
152, 153, 154, 155, 156, 157, 158, 159, 160, 161, 162, 163, 164, 165, 16
168, 169, 170, 171, 172, 173, 174, 175, 176, 177, 178, 179, 180, 181, 18
184, 185, 186, 187, 188, 189, 190, 191, 192, 193, 194, 195, 196, 197, 19
200, 201, 202, 203, 204, 205, 206, 207, 208, 209, 210, 211, 212, 213, 21
216, 217, 218, 219, 220, 221, 222, 223, 224, 225, 226, 227, 228, 229, 23
232, 233, 234, 235, 236, 237, 238, 239, 240, 241, 242, 243, 244, 245, 24
248, 249, 250, 251, 252, 253, 254, 255, 256, 257, 258, 259, 260, 261, 26
264, 265, 266, 267, 268, 269, 270, 271, 272, 273, 274, 275, 276, 277, 27
280, 281, 282, 283, 284, 285, 286, 287, 288, 289, 290, 291, 292, 293, 29
296, 297, 298, 299, 300, 301, 302, 303, 304, 305, 306, 307, 308, 309, 31
312, 313, 314, 315, 316, 317, 318, 319, 320, 321, 322, 323, 324, 325, 32
328, 329, 330, 331, 332, 333, 334, 335, 336, 337, 338, 339, 340, 341, 34
344, 345, 346, 347, 348, 349, 350, 351, 352, 353, 354, 355, 356, 357, 35
360, 361, 362, 363, 364, 365, 366, 367, 368, 369, 370, 371, 372, 373, 37
376, 377, 378, 379, 380, 381, 382, 383, 384, 385, 386, 387, 388, 389, 39
392, 393, 394, 395, 396, 397, 398, 399, 400, 401, 402, 403, 404, 405, 40
408, 409, 410, 411, 412, 413, 414, 415, 416, 417, 418, 419, 420, 421, 42
424, 425, 426, 427, 428, 429, 430, 431, 432, 433, 434, 435, 436, 437, 43
440, 441, 442, 443, 444, 445, 446, 447, 448, 449, 450, 451, 452, 453, 45
456, 457, 458, 459, 460, 461, 462, 463, 464, 465, 466, 467, 468, 469, 47
472, 473, 474, 475, 476, 477, 478, 479, 480, 481, 482, 483, 484, 485, 48
488, 489, 490, 491, 492, 493, 494, 495, 496, 497, 498, 499, 500, 501, 50

```
    Quicksort(A,yellow,r)
  File "/Users/madhavan/mirror/projects/NPTEL/python-2016-jul/week4/python/quick
sort/quicksort.py", line 12, in Quicksort
    Quicksort(A,l,yellow-1)  # Recursive calls
  File "/Users/madhavan/mirror/projects/NPTEL/python-2016-jul/week4/python/quick
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  File "/Users/madhavan/mirror/projects/NPTEL/python-2016-jul/week4/python/quick
sort/quicksort.py", line 7, in Quicksort
    for green in range(l+1,r):
RecursionError: maximum recursion depth exceeded in comparison
>>>
```

Quicksort

- Choose a pivot element
  - Typically the first value in the array
- Partition A into lower and upper parts with respect to pivot
- Move pivot between lower and upper partition
- Recursively sort the two partitions

Analysis of Quicksort

Worst case

- Pivot is either maximum or minimum
  - One partition is empty
  - Other has size n-1
  - $T(n) = T(n-1) + n = T(n-2) + (n-1) + n$
    $= \ldots = 1 + 2 + \ldots + n = O(n^2)$

```
madhavan@dolphinair:.../week4/python/quicksort$ more quicksort.py
def Quicksort(A,l,r): # Sort A[l:r]
  if r - l <= 1: # Base case
    return()

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  for green in range(l+1,r):
    if A[green] <= A[l]:
      (A[yellow],A[green]) = (A[green],A[yellow])
      yellow = yellow + 1
  (A[l],A[yellow-1]) = (A[yellow-1],A[l]) # Move pivot into place
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[GCC 4.2.1 (Apple Inc. build 5666) (dot 3)] on darwin
Type "help", "copyright", "credits" or "license" for more information.
>>> from quicksort import *
>>> l = 
```

Analysis of Quicksort

But …

- Average case is $O(n \log n)$
  - All permutations of n values, each equally likely
  - Average running time across all permutations
- Sorting is a rare example where average case can be computed

Quicksort

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Analysis of Quicksort

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  - Other has size $n-1$
  - $T(n) = T(n-1) + n = T(n-2) + (n-1) + n$
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Analysis of Quicksort

But …

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  - All permutations of n values, each equally likely
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- Sorting is a rare example where average case can be computed

Quicksort: randomization

- Worst case arises because of fixed choice of pivot
  - We chose the first element
  - For any fixed strategy (last element, midpoint), can work backwards to construct $O(n^2)$ worst case
- Instead, choose pivot randomly
  - Pick any index in `range(0,n)` with uniform probability
- Expected running time is again $O(n \log n)$

Quicksort in practice

- In practice, Quicksort is very fast
  - Typically the default algorithm for in-built sort functions
    - Spreadsheets
    - Built in sort function in programming languages

Quicksort in practice

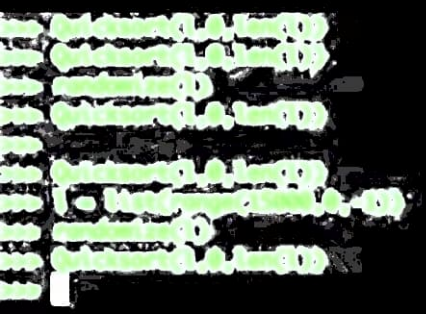
l.sort()

- In practice, Quicksort is very fast
  - Typically the default algorithm for in-built sort functions
    - Spreadsheets
    - Built in sort function in programming languages

```
mashaven@dolphinair:.../week4/python/quicksort$ more quicksort.py
def Quicksort(A,l,r): # Sort A[l:r]
  if r - l <= 1: # Base case
    return()

  # Partition with respect to pivot, a[l]
  yellow = l+1
  for green in range(l+1,r):
    if A[green] <= A[l]:
      (A[yellow],A[green]) = (A[green],A[yellow])
      yellow = yellow + 1
  (A[l],A[yellow-1]) = (A[yellow-1],A[l]) # Move pivot into place
  Quicksort(A,l,yellow-1)  # Recursive calls
  Quicksort(A,yellow,r)
mashaven@dolphinair:.../week4/python/quicksort$
```

```
[redacted]:~/Assets/python/quicksort$ nano randomize.py
import random
def randomize():
    for i in range(len(l)/2):
        j = random.randrange(0,len(l),1)
        k = random.randrange(0,len(l),1)
        l[j],l[k] = l[k],l[j]
[redacted]:~/Assets/python/quicksort$
```



Stable sorting

- Sorting on multiple criteria
- Assume students are listed in alphabetical order
- Now sort students by marks
  - After sorting, are students with equal marks still in alphabetical order?
- Stability is crucial in applications like spreadsheets
  - Sorting column B should not disturb previous sort on column A

Stable sorting ...

- Quicksort, as described, is not stable
  - Swap operation during partitioning disturbs original order
- Merge sort is stable if we merge carefully
  - Do not allow elements from right to overtake elements from left
  - Favour left list when breaking ties

Tuples

- Simultaneous assignments

```
(age,name,primes) = (23,"Kamal",[2,3,5])
```

- Can assign a "tuple" of values to a name

```
point = (3.5,4.8)
date = (16,7,2013)
```

- Extract positions, slices

```
xcoordinate = point[0]
monthyear = date[1:]
```

- Tuples are immutable

`date[1] = 8` is an error

Generalizing lists

- `l = [13, 46, 0, 25, 72]`
- View `l` as a function, associating values to positions
  - `l : {0,1,..,4} → integers`
  - `l(0) = 13, l(4) = 72`
- `0,1,..,4` are keys
- `l[0],l[1],..,l[4]` are corresponding values

Dictionaries

- Allow keys other than `range(0,n)`
- Key could be a string

```
test1["Dhawan"] = 84
test1["Pujara"] = 16
test1["Kohli"] = 200
```

- Python dictionary
  - Any immutable value can be a key

Dictionaries

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- Key could be a string

```
test1["Dhawan"] = 84
test1["Pujara"] = 16
test1["Kohli"] = 200
```

- Python dictionary
  - Any immutable value can be a key
  - Can update dictionaries in place —mutable, like lists

Dictionaries

- Can nest dictionaries

```
score["Test1"]["Dhawan"] = 84
score["Test2"]["Kohli"] = 200
score["Test2"]["Dhawan"] = 27
```

- Directly assign values to a dictionary

```
score = {"Dhawan":84, "Kohli":200}
score = {"Test1":{"Dhawan":84,
  "Kohli":200}, "Test2":{"Dhawan":50}}
```

```
>>> score = {}
>>> score
```

Operating on dictionaries

- `d.keys()` returns sequence of keys of dictionary `d`

```
for k in d.keys():
    # Process d[k]
```

- `d.keys()` is not in any predictable order

```
for k in sorted(d.keys()):
    # Process d[k]
```

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- `sorted(l)` returns sorted copy of `l`, `l.sort()` sorts `l` in place

Dictionaries

- Can nest dictionaries

```
score["Test1"]["Dhawan"] = 84
score["Test2"]["Kohli"] = 200
score["Test2"]["Dhawan"] = 27
```

Operating on dictionaries

- `d.keys()` returns sequence of keys of dictionary `d`

```
for k in d.keys():
    # Process d[k]
```

- `d.keys()` is not in any predictable order

```
for k in sorted(d.keys()):
    # Process d[k]
```

- `sorted(l)` returns sorted copy of `l`, `l.sort()` sorts `l` in place
- `d.keys()` is not a list —use `list(d.keys())`

```
>>> d = {}
>>> for l in "abcdefghi":
...     d[l] = l
...
>>> d["a"]
'a'
>>> d["i"]
'i'
>>> d.keys()
dict_keys(['e', 'i', 'g', 'b', 'f', 'd', 'a', 'c', 'h'])
>>>
```

Operating on dictionaries

- Similarly, `d.values()` is sequence of values in `d`

```
total = 0
for s in test1.values():
  total = total + test1
```

- Test for key using `in`, like list membership

```
for n in ["Dhawan","Kohli"]:
  total[n] = 0
  for match in score.keys():
    if n in score[match].keys():
      total[n] = total[n] + score[match][n]
```

Operating on dictionaries

- Similarly, `d.values()` is sequence of values in `d`

```
total = 0
for s in test1.values():
  total = total + test1
```

- Test for key using `in`, like list membership

```
for n in ["Dhawan","Kohli"]:
  total[n] = 0
  for match in score.keys():
    if n in score[match].keys():
      total[n] = total[n] + score[match][n]
```

Dictionaries vs lists

- Assigning to an unknown key inserts an entry

```
d = {}
d[0] = 7  # No problem, d == {0:7}
```

Dictionaries vs lists

- Assigning to an unknown key inserts an entry

```
d = {}
d[0] = 7  # No problem, d == {0:7}
```

- ... unlike a list

```
l = []
l[0] = 7  # IndexError!
```

Summary

- Dictionaries allow a flexible association of values to keys
  - Keys must be immutable values
- Structure of dictionary is internally optimized for key-based lookup
  - Use `sorted(d.keys())` to retrieve keys in predictable order
- Extremely useful for manipulating information from text files, tables … — use column headings as keys

Passing values to functions

- Argument value is substituted for name

```
def power(x,n):
    ans = 1
    for i in range(0,n):
        ans = ans*x
    return(ans)
```

```
power(3,5)
x = 3
n = 5
ans = 1
for i in range..
```

- Like an implicit assignment statement

Pass arguments by name

```
def power(x,n):
    ans = 1
    for i in range(0,n):
        ans = ans*x
    return(ans)
```

- Call power(n=5,x=4)

Default arguments

- Recall `int(s)` that converts string to integer
  - `int("76")` is 76
  - `int("A5")` generates an error
- Actually `int(s,b)` takes two arguments, string s and base b
  - b has default value `10`

int("76", 10)

Default arguments

- Recall `int(s)` that converts string to integer
  - `int("76")` is 76
  - `int("A5")` generates an error
- Actually `int(s,b)` takes two arguments, string s and base b
  - b has default value 10
  - `int("A5",16)` is 165 (10 x 16 + 5)

Default arguments

```
def int(s,b=10):
    . . .
```

- Default value is provided in function definition
- If parameter is omitted, default value is used
  - Default value must be available at definition time
  - `def Quicksort(A,l=0,r=len(A)):` does not work

Default arguments

```
def f(a,b,c=14,d=22):
    . . .
```

- `f(13,12)` is interpreted as `f(13,12,14,22)`
- `f(13,12,16)` is interpreted as `f(13,12,16,22)`
- Default values are identified by position, must come at the end

Default arguments

```
def f(a,b,c=14,d=22):
    . . .
```

- `f(13,12)` is interpreted as `f(13,12,14,22)`
- `f(13,12,16)` is interpreted as `f(13,12,16,22)`
- Default values are identified by position, must come at the end
  - Order is important

Function definitions

- `def` associates a function body with a name
- Flexible, like other value assignments to name
- Definition can be conditional

```
if condition:
    def f(a,b,c):
        . . .
else:
    def f(a,b,c):
        . . .
```

Function definitions

- Can assign a function to a new name

```
def f(a,b,c):
    . . .

g = f
```

- Now `g` is another name for `f`

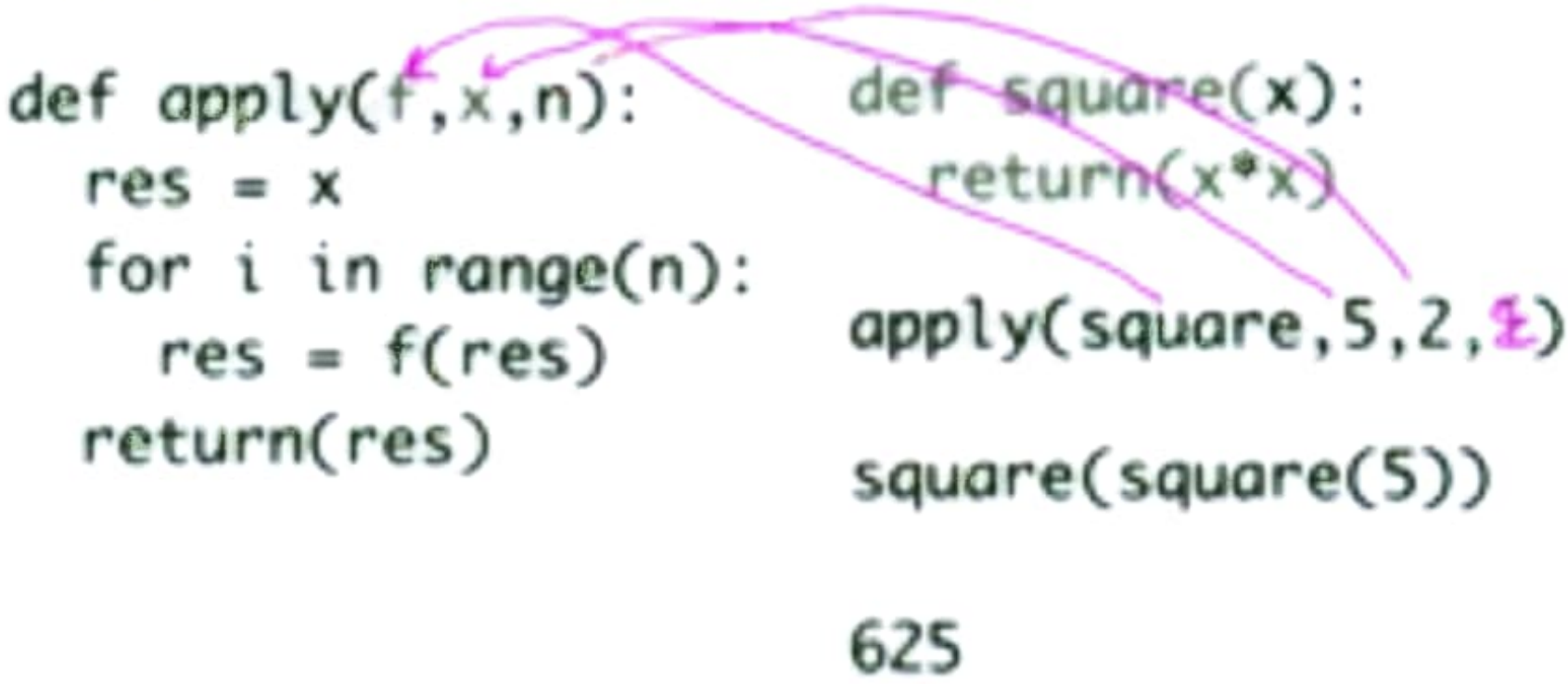
Can pass functions

- Apply `f` to `x` `n` times

```
def apply(f,x,n):
    res = x
    for i in range(n):
        res = f(res)
    return(res)
```

Can pass functions

- Apply f to x n times



```
def apply(f,x,n):
    res = x
    for i in range(n):
        res = f(res)
    return(res)
```

```
def square(x):
    return(x*x)

apply(square,5,2,1)

square(square(5))

625
```

Passing functions

- Useful for customizing functions such as sort
- Define `cmp(x,y)` that returns `-1` if `x < y`, `0` if `x == y` and `1` if `x > y`
  - `cmp("aab","ab")` is `-1` in dictionary order
  - `cmp("aab","ab")` is `1` if we compare by length
- `def sortfunction(l,cmpfn=defaultcmpfn):`

Summary

- Function definitions behave like other assignments of values to names
- Can reassign a new definition, define conditionally ...
- Can pass function names to other functions

Built in function map()

- `map(f,l)` applies `f` to each element of `l`
- Output of `map(f,l)` is not a list!
  - Use `list(map(f,l))` to get a list
  - Can be used directly in a `for` loop

```
for i in list(map(f,l)):
```

  - Like `range(i,j)`, `d.keys()`

Selecting a sublist

- Extract list of primes from list `numberlist`

```
primelist = []
for i in numberlist:
  if isprime(i):
    primelist.append(i)
return(primelist)
```

Selecting a sublist

- In general

```
def select(property,l):
    sublist = []
    for x in l:
        if property(x):
            sublist.append(x)
    return(sublist)
```

Combining map and filter

- Sum of squares of even numbers from 0 to 99

```
list(map(square,filter(iseven,range(100))

def square(x):
  return(x*x)

def iseven(x):
  return(x%2 == 0)
```

Combining map and filter

- Sum of **squares of even numbers from 0 to 99**

```
list(map(square,filter(iseven,range(100))))

def square(x):
  return(x*x)

def iseven(x):
  return(x%2 == 0)
```

List comprehension

- Pythagorean triple: $x^2 + y^2 = z^2$
- All Pythagorean triples (x,y,z) with values below n

$$\{ (x,y,z) \mid 1 \leq x,y,z \leq n,\ x^2 + y^2 = z^2 \}$$

List comprehension

- Pythagorean triple: $x^2 + y^2 = z^2$
- All Pythagorean triples (x,y,z) with values below n

$$\{ (x,y,z) \mid 1 \leq x,y,z \leq n,\ x^2 + y^2 = z^2 \}$$

- In set theory, this is called set comprehension
  - Building a new set from existing sets

List comprehension

- Squares of even numbers below 100

```
[square(x) for i in range(100) if iseven(x)]
```

map generator filter

Multiple generators

- Pythagorean triples with x,y,z below 100

```
[(x,y,z) for x in range(100)
         for y in range(100)
          for z in range(100)
           if x*x + y*y == z*z]
```

- Order of x,y,z is like nested for loop

```
for x in range(100):
  for y in range(100):
    for z in range(100):
```

Multiple generators

- Later generators can depend on earlier ones
- Pythagorean triples with x,y,z below 100, no duplicates

```
[(x,y,z) for x in range(100)
         for y in range(x,100)
          for z in range(y,100)
           if x*x + y*y == z*z]
```

Useful for initialising lists

- Initialise a 4 x 3 matrix
  - 4 rows, 3 columns
  - Stored row-wise

```
l = [ [ 0 for i in range(3) ]
      for j in range(4)]
```

↳ for each row

```
>>> zerolist = [ 0 for i in range(3) ]
>>> l = [ zerolist for j in range(4) ]
>>> l
[[0, 0, 0], [0, 0, 0], [0, 0, 0], [0, 0, 0]]
>>> l[1][1] = 7
>>>
```

Warning

- **What's happening here?**

```
>>> zerolist = [ 0 for i in range(3) ]
>>> l = [ zerolist for j in range(4) ]

>>> l[1][1] = 7
```

Warning

- **What's happening here?**

```
>>> zerolist = [ 0 for i in range(3) ]
>>> l = [ zerolist for j in range(4) ]

>>> l[1][1] = 7

>>> l
[[0,7,0],[0,7,0],[0,7,0],[0,7,0]]
```

- **Each row in `l` points to same list `zerolist`**

Warning

- **What's happening here?**

```
>>> zerolist = [ 0 for i in range(3) ]
>>> l = [ zerolist for j in range(4) ]

>>> l[1][1] = 7

>>> l
[[0,7,0],[0,7,0],[0,7,0],[0,7,0]]
```

- **Each row in `l` points to same list `zerolist`**

Summary

- `map` and `filter` are useful functions to manipulate lists
- List comprehension provides a useful notation for combining `map` and `filter`