

Indexed Search Tree (Trie)

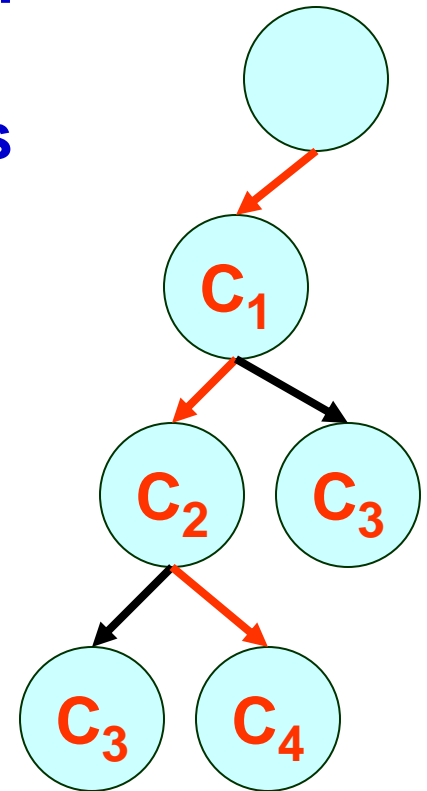


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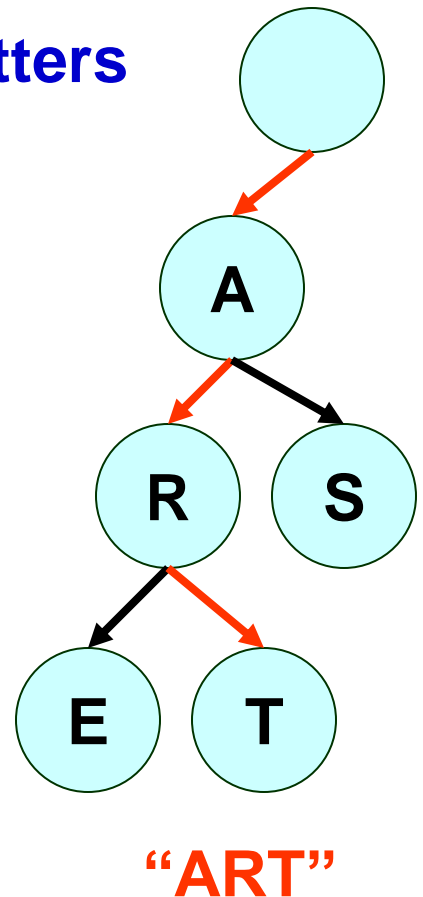
Indexed Search Tree (Trie)

- Special case of tree
- Applicable when
 - Key **C** can be decomposed into a sequence of subkeys **C₁, C₂, ... C_n**
 - Redundancy exists between subkeys
- Approach
 - Store subkey at each node
 - Path through trie yields full key
- Example
 - Huffman tree



Tries

- Useful for searching strings
 - String decomposes into sequence of letters
 - Example
 - “ART” \Rightarrow “A” “R” “T”
- Can be very fast
 - Less overhead than hashing
- May reduce memory
 - Exploiting redundancy
- May require more memory
 - Explicitly storing substrings



Types of Tries

- **Standard**

- Single character per node

- **Compressed**

- Eliminating chains of nodes

- **Compact**

- Stores indices into original string(s)

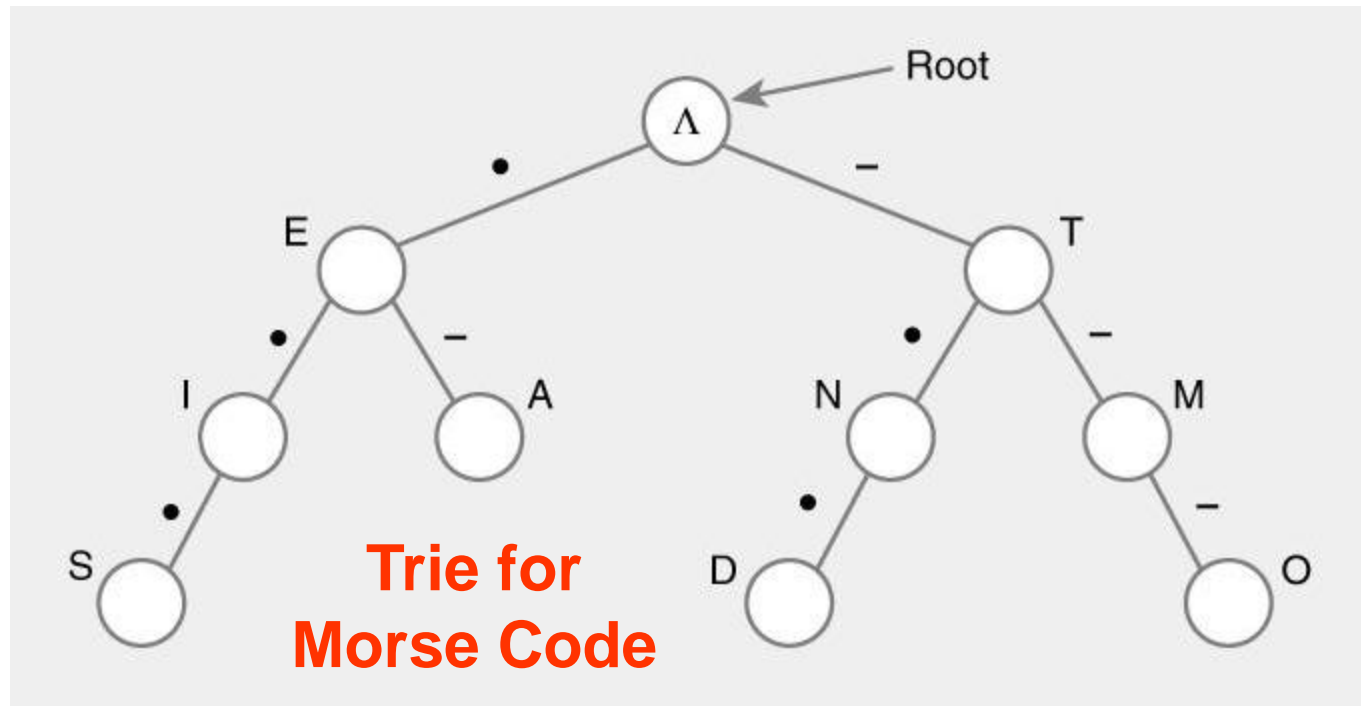
- **Suffix**

- Stores all suffixes of string

Standard Tries

■ Approach

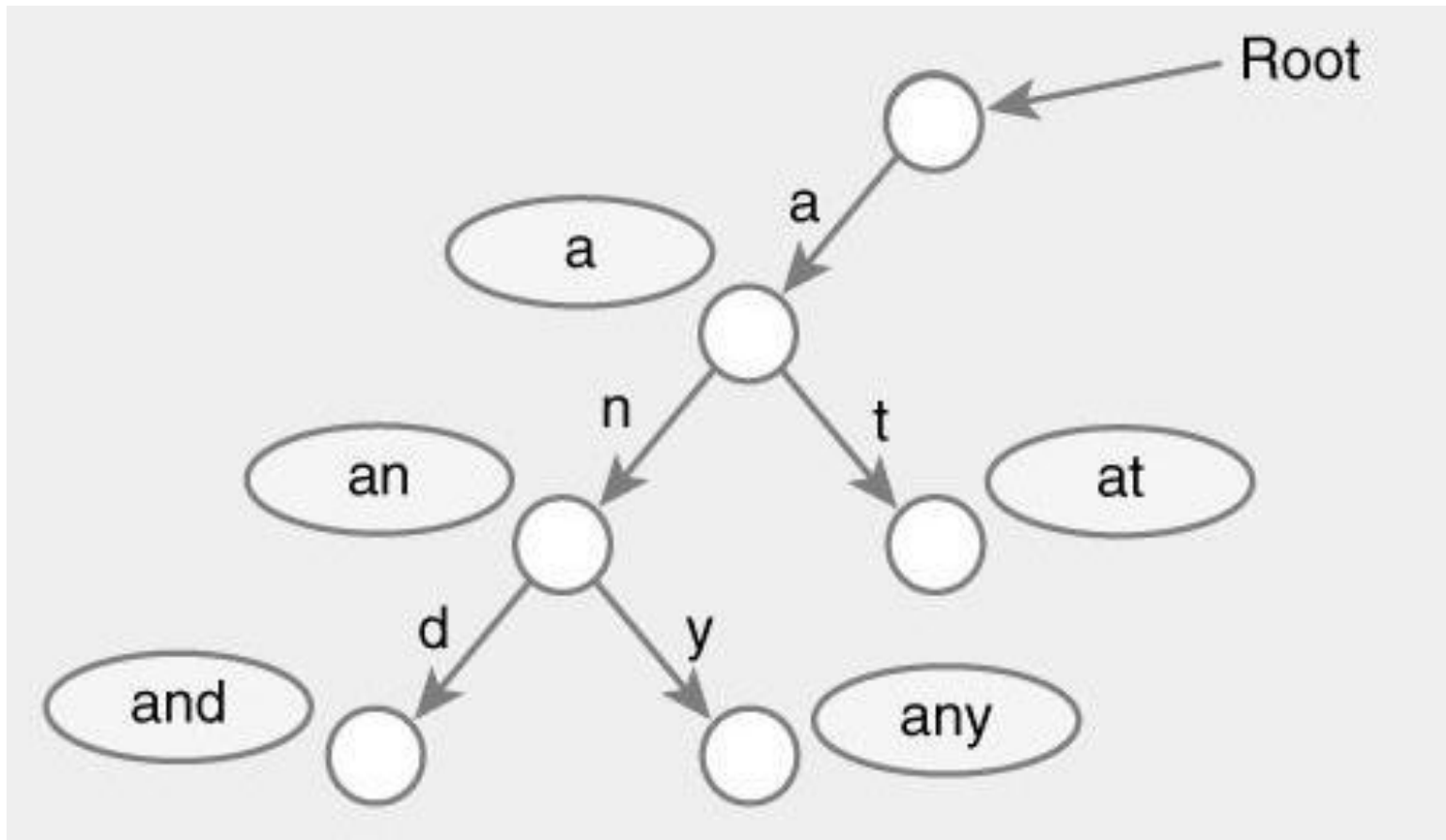
- Each node (except root) is labeled with a character
- Children of node are ordered (alphabetically)
- Paths from root to leaves yield all input strings



Standard Trie Example

■ For strings

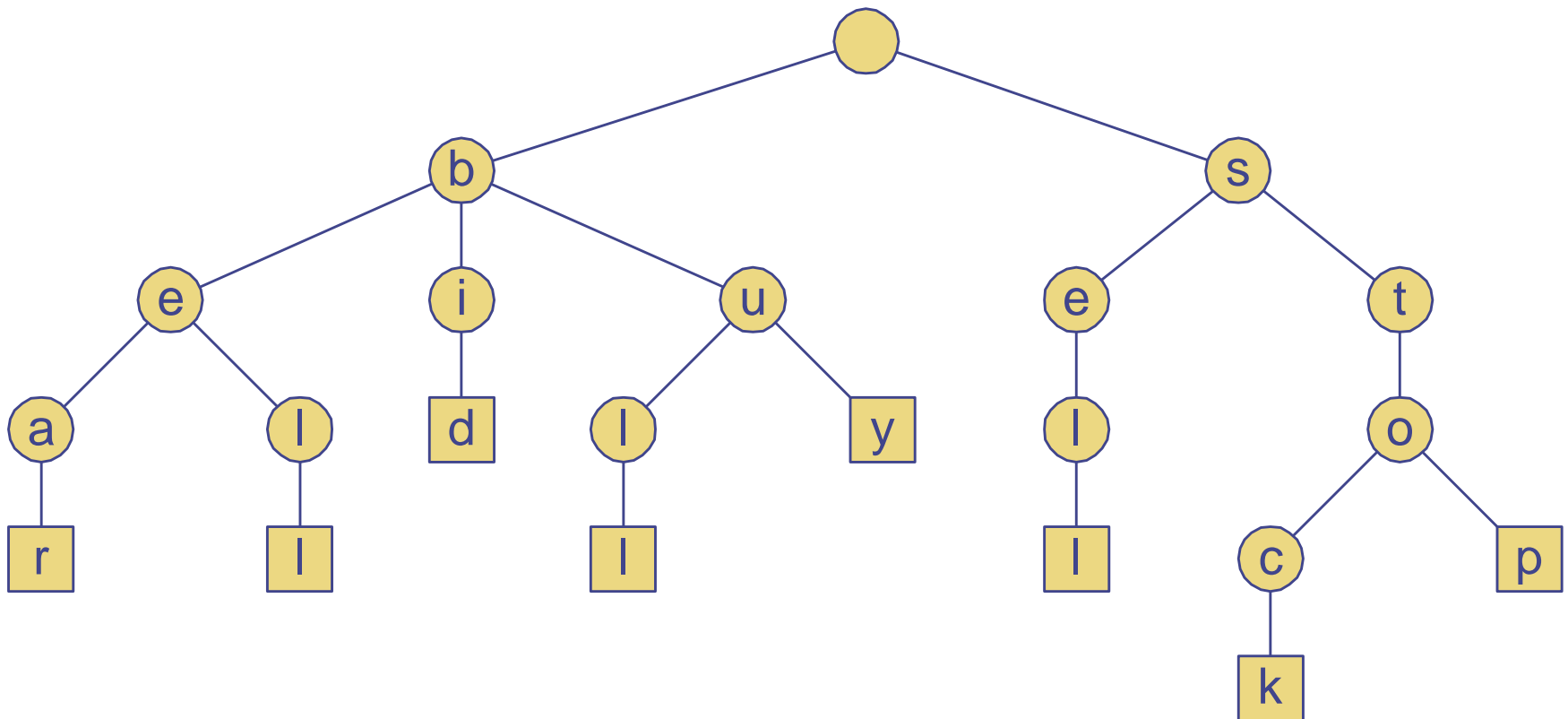
■ { a, an, and, any, at }



Standard Trie Example

■ For strings

■ { bear, bell, bid, bull, buy, sell, stock, stop }



Standard Tries

■ Node structure

- Value between $1 \dots m$
- Reference to m children
 - Array or linked list

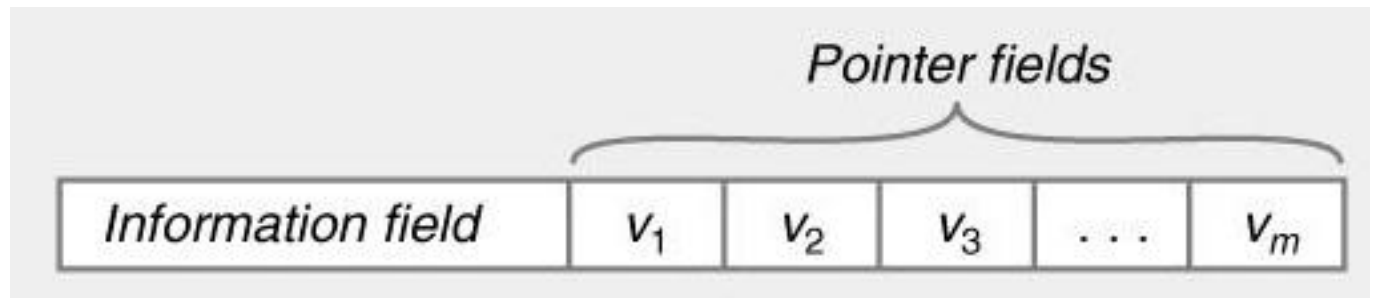
■ Example

Class Node {

Letter value; // Letter $V = \{ V_1, V_2, \dots V_m \}$

Node child[m];

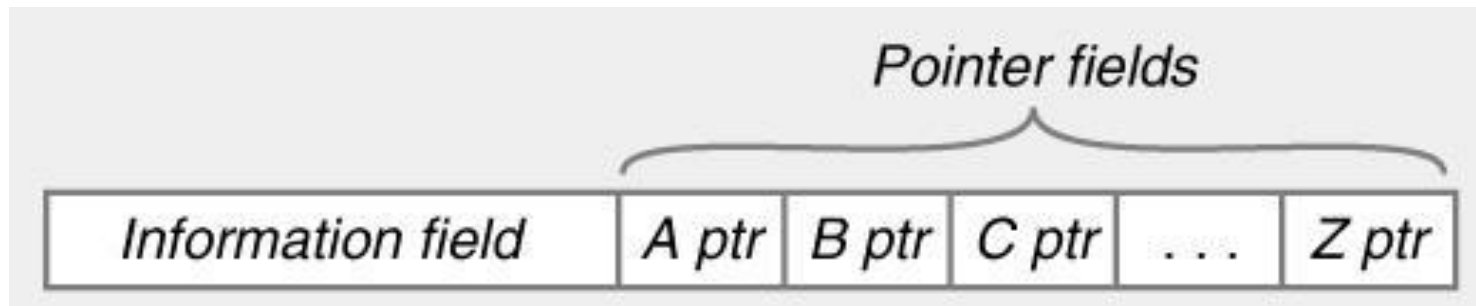
}



Standard Tries

■ Efficiency

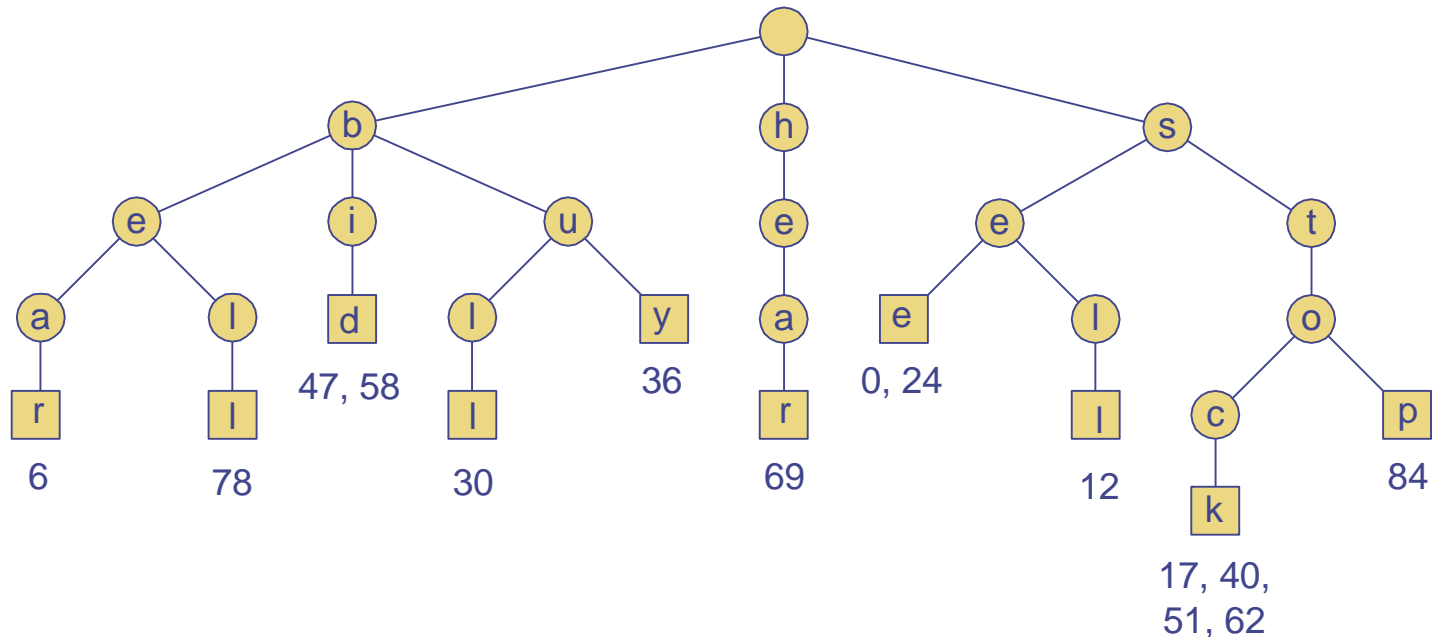
- Uses $O(n)$ space
- Supports search / insert / delete in $O(d \times m)$ time
- For
 - n total size of strings indexed by trie
 - d length of the parameter string
 - m size of the alphabet



Word Matching Trie

- Insert words into trie
- Each leaf stores occurrences of word in the text

s	e	e		a		b	e	a	r	?		s	e	l	l		s	t	o	c	k	!		
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	
s	e	e		a		b	u	l	l	?		b	u	y			s	t	o	c	k	!		
24	25	26	27	28	29	30	31	32	33	34	35	36	37	38	39	40	41	42	43	44	45	46		
b	i	d		s	t	o	c	k	!		b	i	d		s	t	o	c	k	!				
47	48	49	50	51	52	53	54	55	56	57	58	59	60	61	62	63	64	65	66	67	68			
h	e	a	r		t	h	e		b	e	l	l	?		s	t	o	p	!					
69	70	71	72	73	74	75	76	77	78	79	80	81	82	83	84	85	86	87	88					



Compressed Trie

■ Observation

- Internal node v of T is redundant if v has one child and is not the root

■ Approach

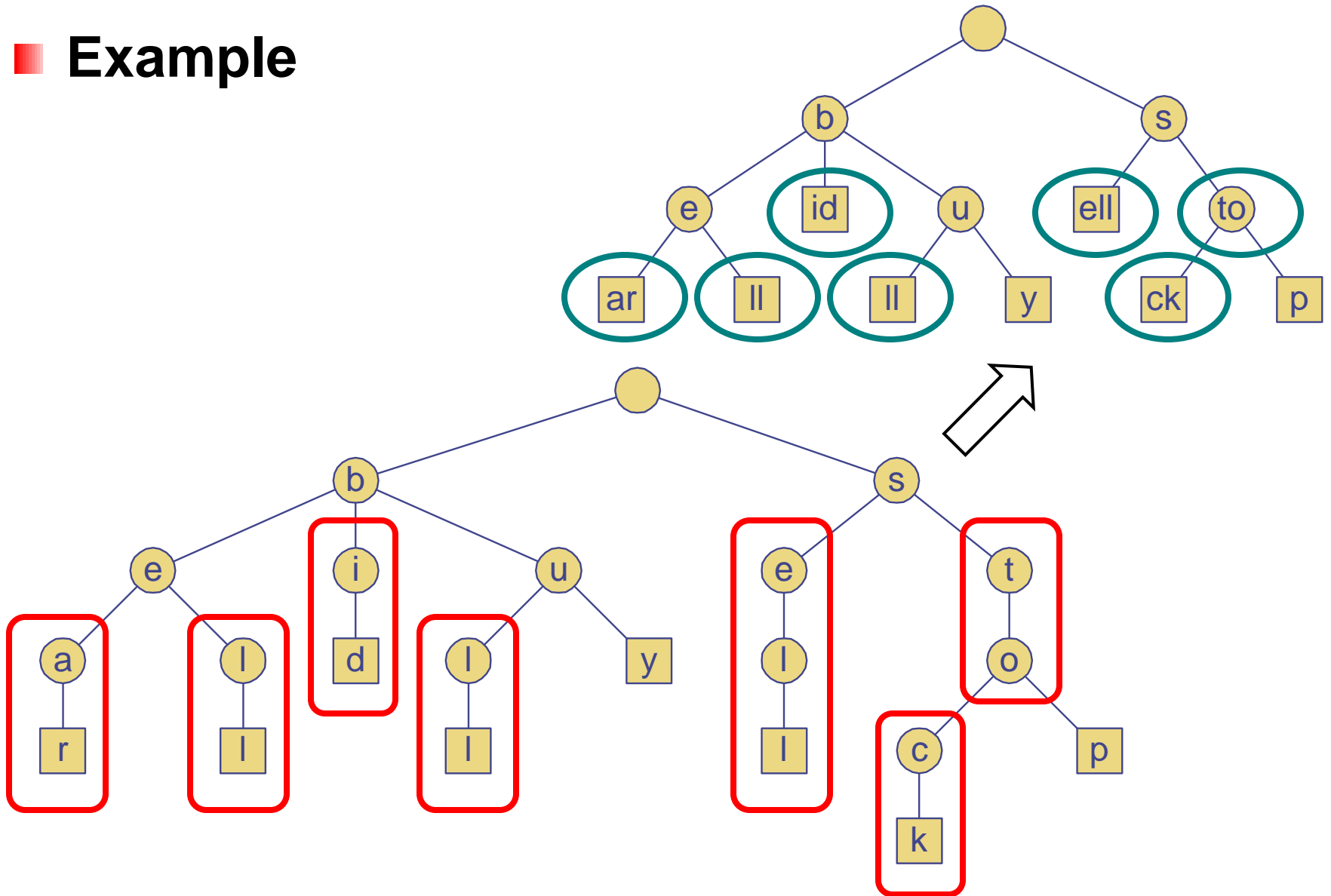
- A chain of redundant nodes can be compressed
 - Replace chain with single node
 - Include concatenation of labels from chain

■ Result

- Internal nodes have at least 2 children
- Some nodes have multiple characters

Compressed Trie

■ Example



Compact Tries

■ Compact representation of a compressed trie

■ Approach

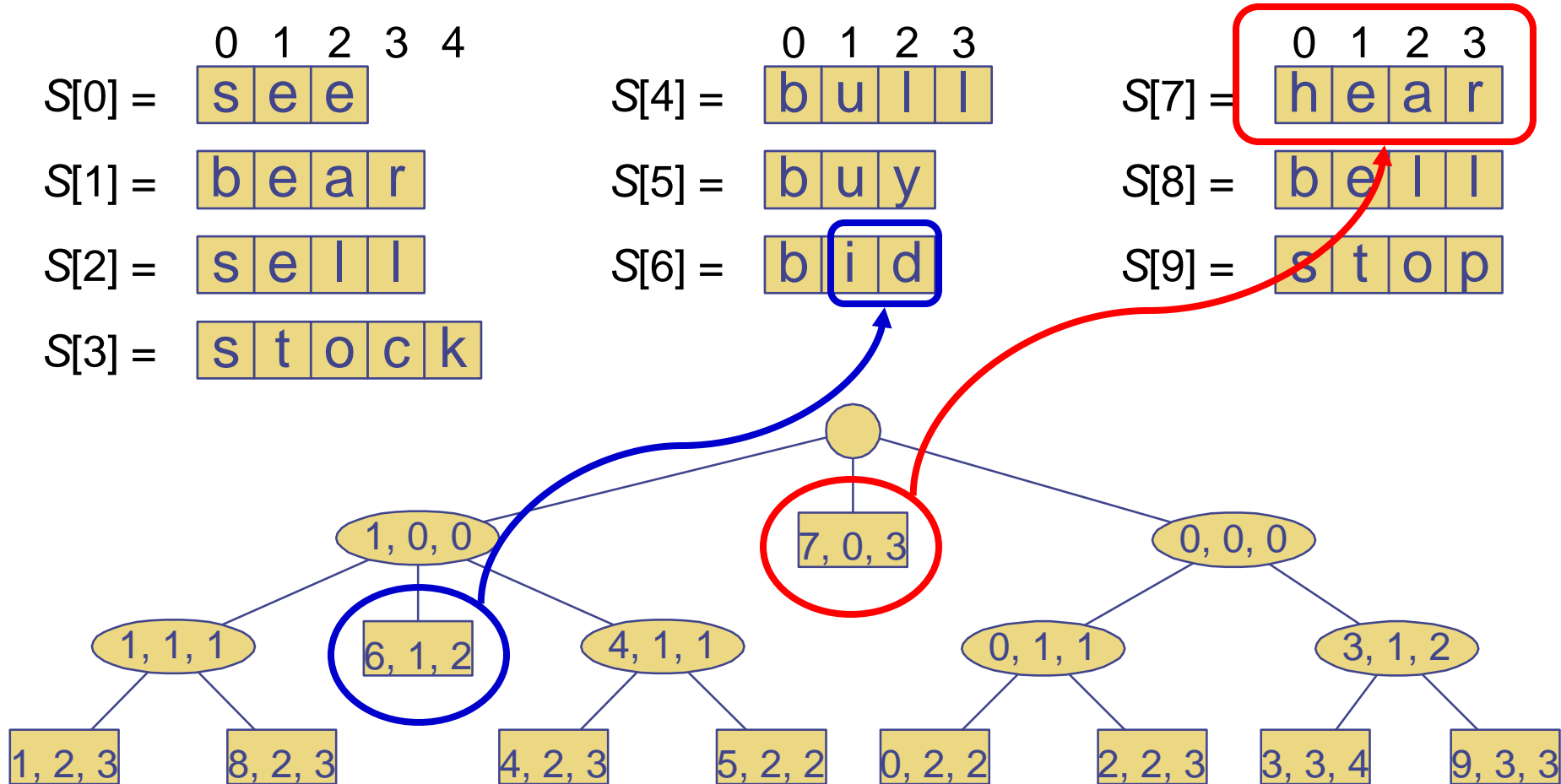
- For an array of strings $S = S[0], \dots S[s-1]$
- Store ranges of indices at each node
 - Instead of substring
- Represent as a triplet of integers (i, j, k)
 - Such that $X = s[i][j..k]$
- Example: $S[0] = \text{"abcd"}$, $(0,1,2) = \text{"bc"}$

■ Properties

- Uses $O(s)$ space, where $s = \#$ of strings in the array
- Serves as an auxiliary index structure

Compact Representation

■ Example



Suffix Trie

■ Compressed trie of all suffixes of text

■ Example: “IPDPS”

■ Suffixes

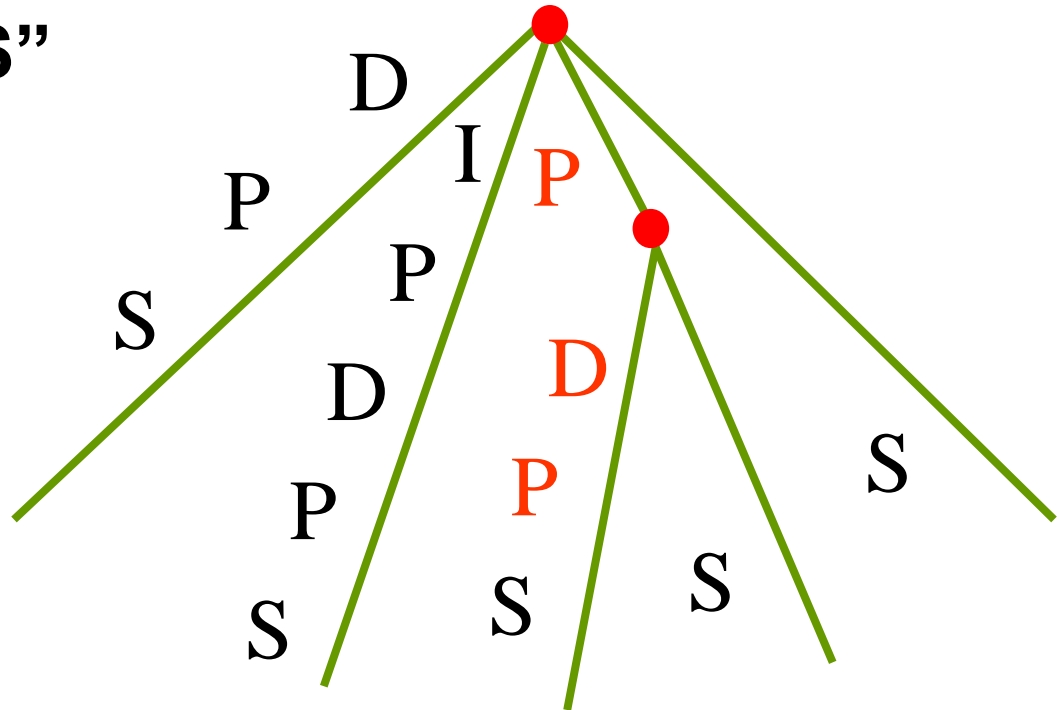
■ IPDPS

■ PDPS

■ DPS

■ PS

■ S



■ Useful for finding pattern in any part of text

■ Occurrence \Rightarrow prefix of some suffix

■ Example: find **PDP** in **IPDPS**

Suffix Trie

■ Properties

■ For

- String X with length n
- Alphabet of size m
- Pattern P with length d

■ Uses $O(n)$ space

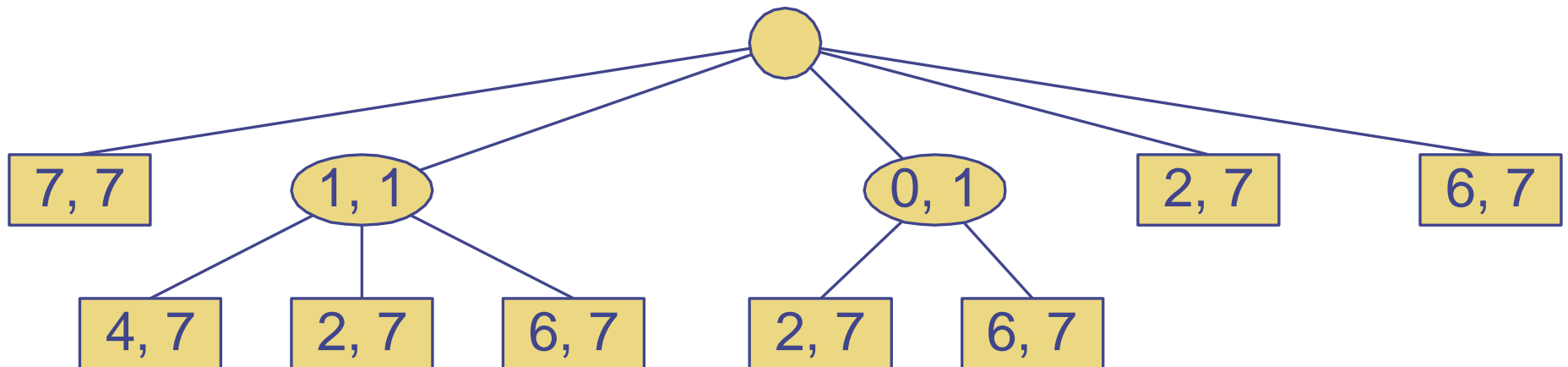
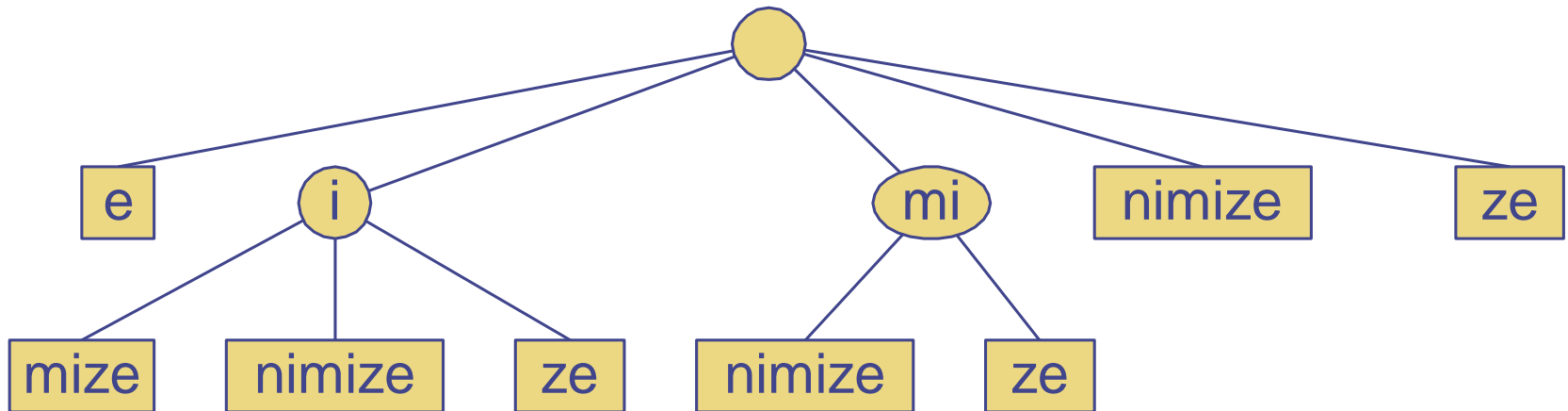
■ Can be constructed in $O(n)$ time

■ Find pattern P in X in $O(d \times m)$ time

- Proportional to length of pattern, not text

Suffix Trie Example

m	i	n	i	m	i	z	e
0	1	2	3	4	5	6	7



Tries and Web Search Engines

- **Search engine index**
 - Collection of all searchable words
 - Stored in compressed trie
- **Each leaf of trie**
 - Associated with a word
 - List of pages (URLs) containing that word
 - Called occurrence list
- **Trie is kept in memory (fast)**
- **Occurrence lists kept in external memory**
 - Ranked by relevance

Computational Biology

■ DNA

- Sequence of 4 different nucleotides (ATCG)
- Portions of DNA sequence produce proteins (genes)

■ Genome

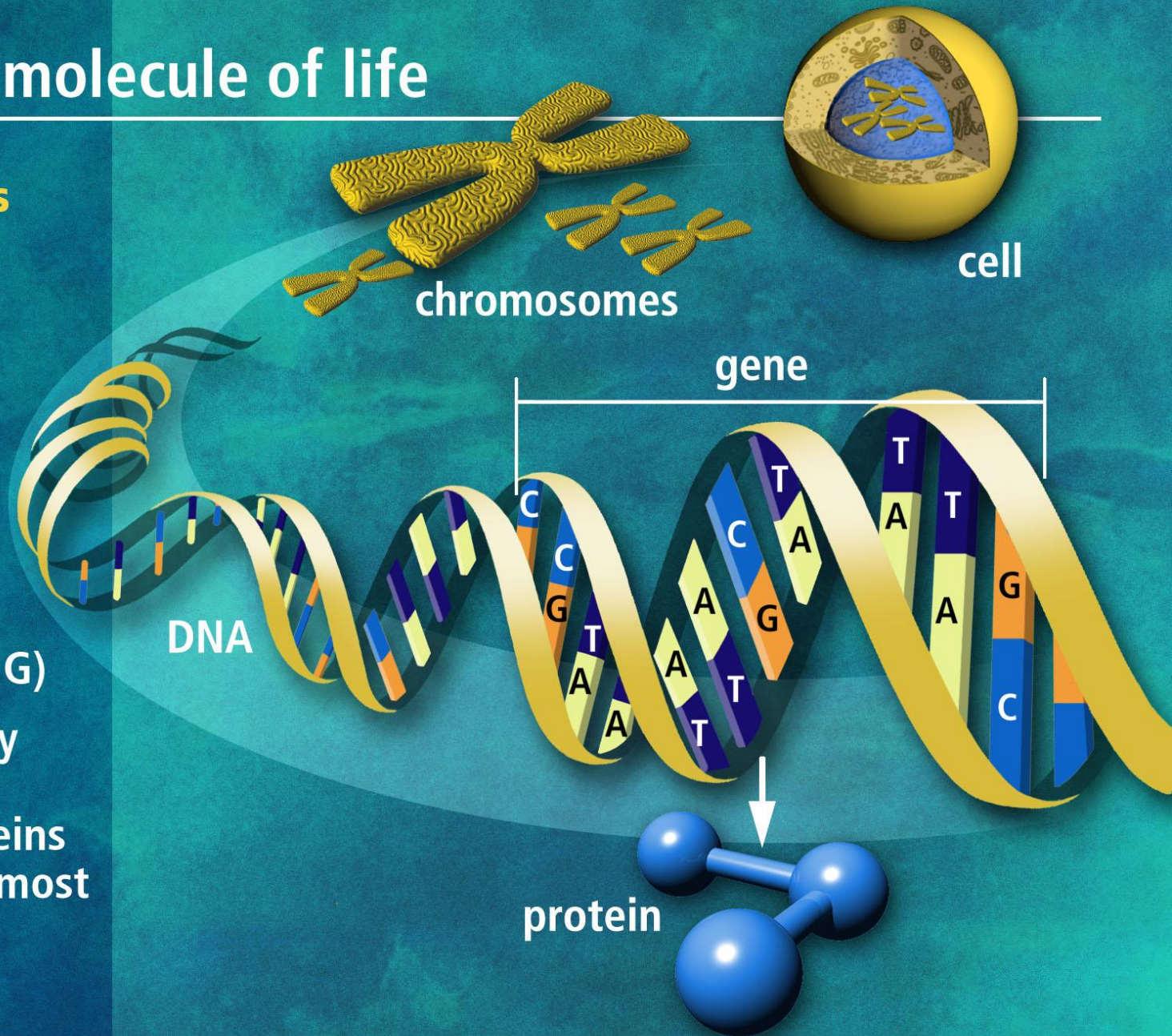
- Master DNA sequence for organism
- For Human
 - 46 chromosomes
 - 3 billion nucleotides

DNA the molecule of life

Trillions of cells

Each cell:

- 46 human chromosomes
- 2 meters of DNA
- 3 billion DNA subunits (the bases: A, T, C, G)
- Approximately 30,000 genes code for proteins that perform most life functions



Tries and Computational Biology

■ ESTs

- Fragments of expressed DNA
- Indicator for genes (& location)
- 5.5 million sequences at NIH

■ ESTmapper

- Build suffix trie of genome
 - 8 hours, 60 Gbytes
- Search for ESTs in suffix trie
 - 11 hours w/ 8 processor Sun

■ Search genome w/ BLAST

- 5⁺ years (predicted)

