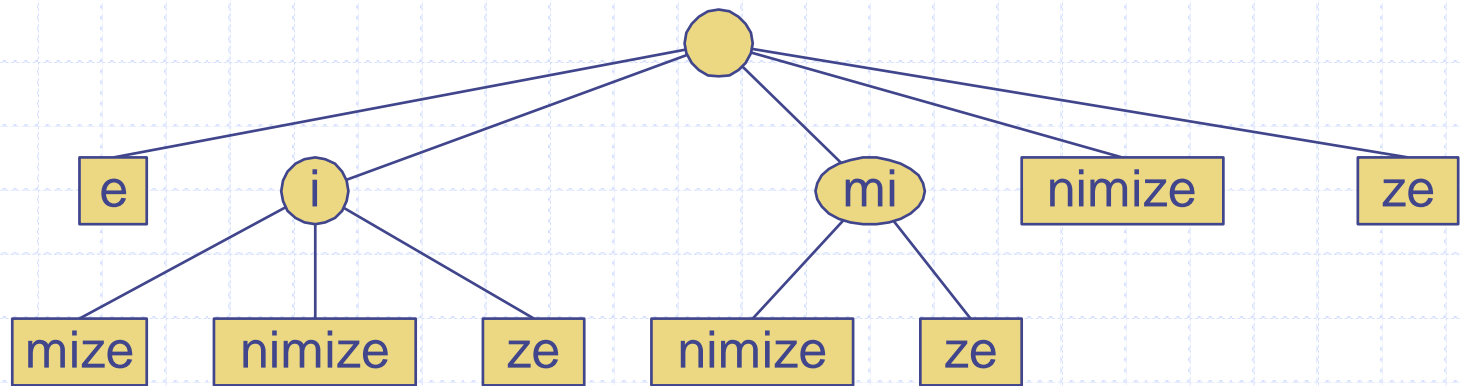


Tries

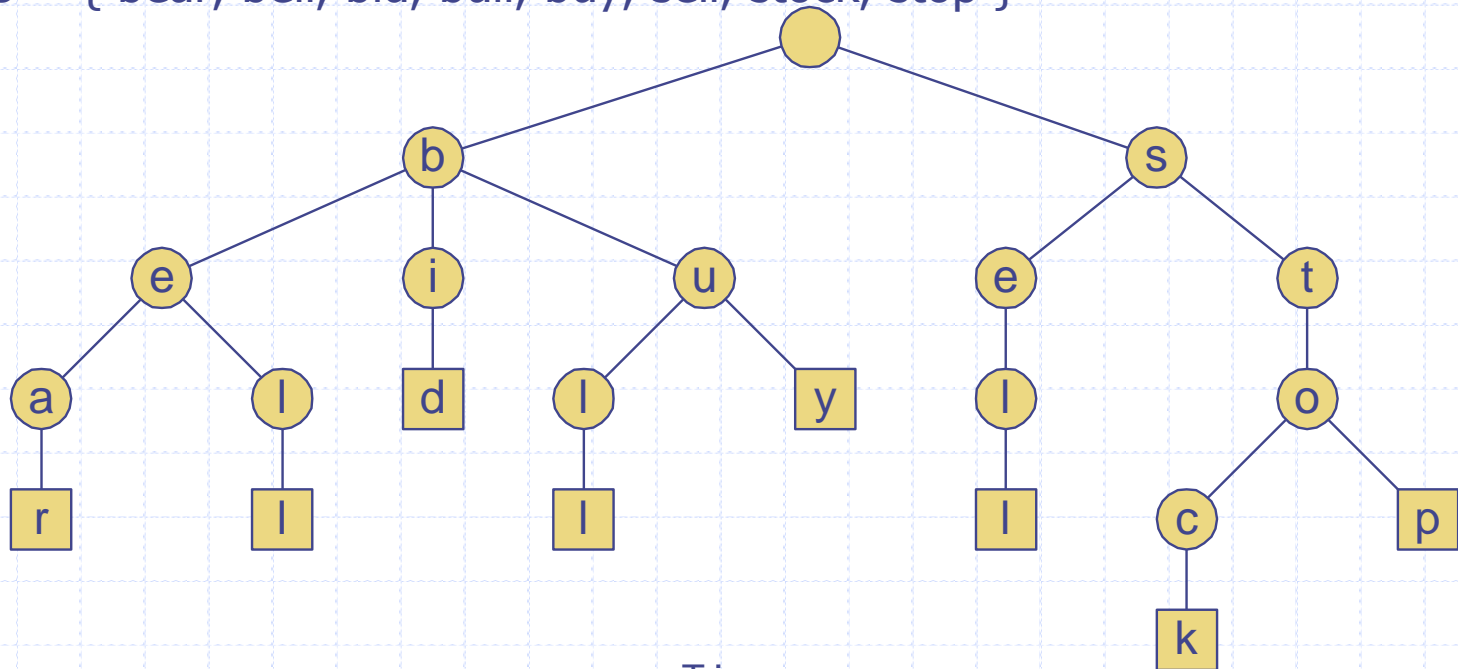


Preprocessing Strings

- ◆ Preprocessing the pattern speeds up pattern matching queries
 - After preprocessing the pattern, KMP's algorithm performs pattern matching in time proportional to the text size
- ◆ If the text is large, immutable and searched for often (e.g., works by Shakespeare), we may want to preprocess the text instead of the pattern
- ◆ A trie is a compact data structure for representing a set of strings, such as all the words in a text
 - A trie supports pattern matching queries in time proportional to the pattern size

Standard Tries

- ◆ The standard trie for a set of strings S is an ordered tree such that:
 - Each node but the root is labeled with a character
 - The children of a node are alphabetically ordered
 - The paths from the external nodes to the root yield the strings of S
- ◆ Example: standard trie for the set of strings
 $S = \{ \text{bear, bell, bid, bull, buy, sell, stock, stop} \}$



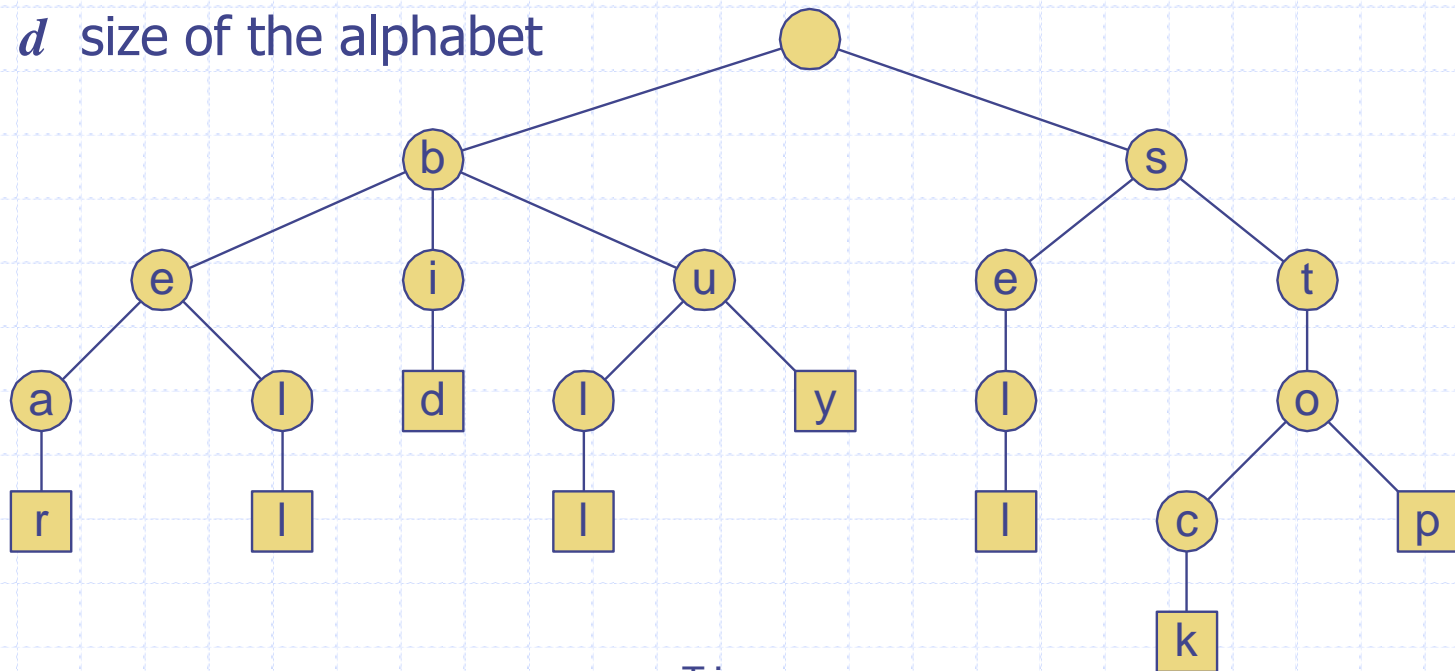
Analysis of Standard Tries

- ◆ A standard trie uses $O(n)$ space and supports searches, insertions and deletions in time $O(dm)$, where:

n total size of the strings in S

m size of the string parameter of the operation

d size of the alphabet

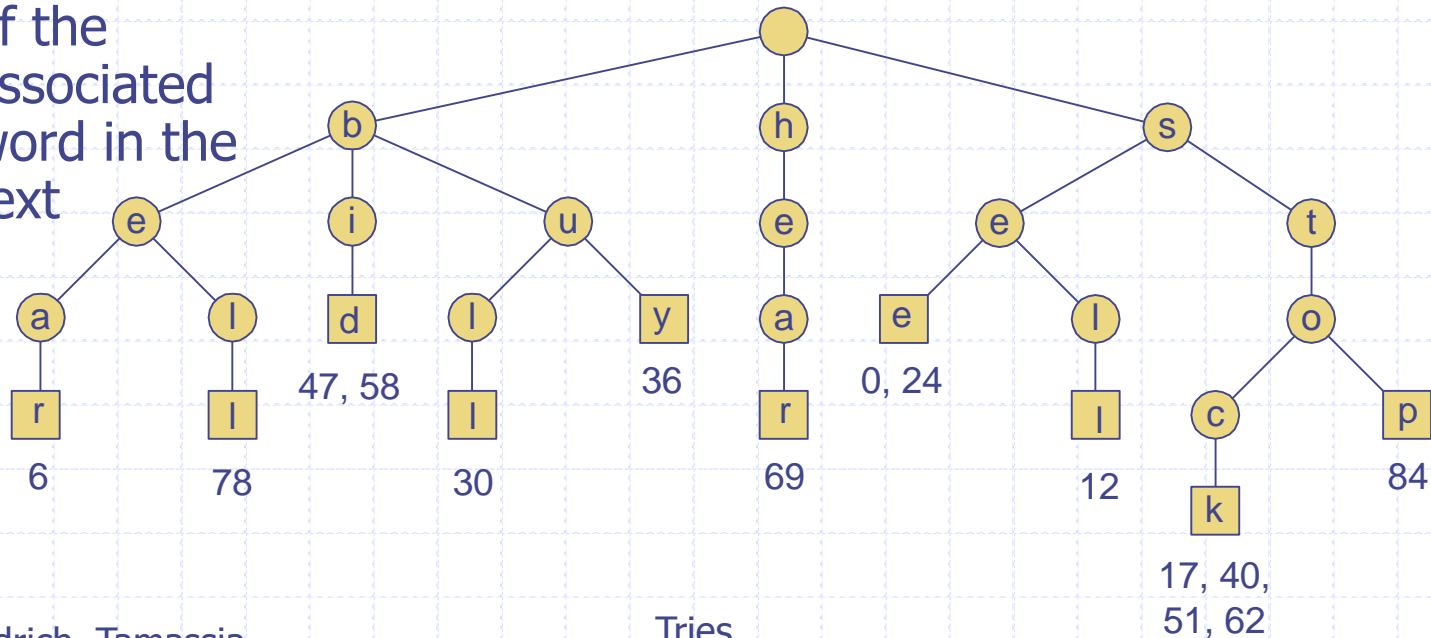


Word Matching with a Trie

- ◆ We insert the words of the text into a trie

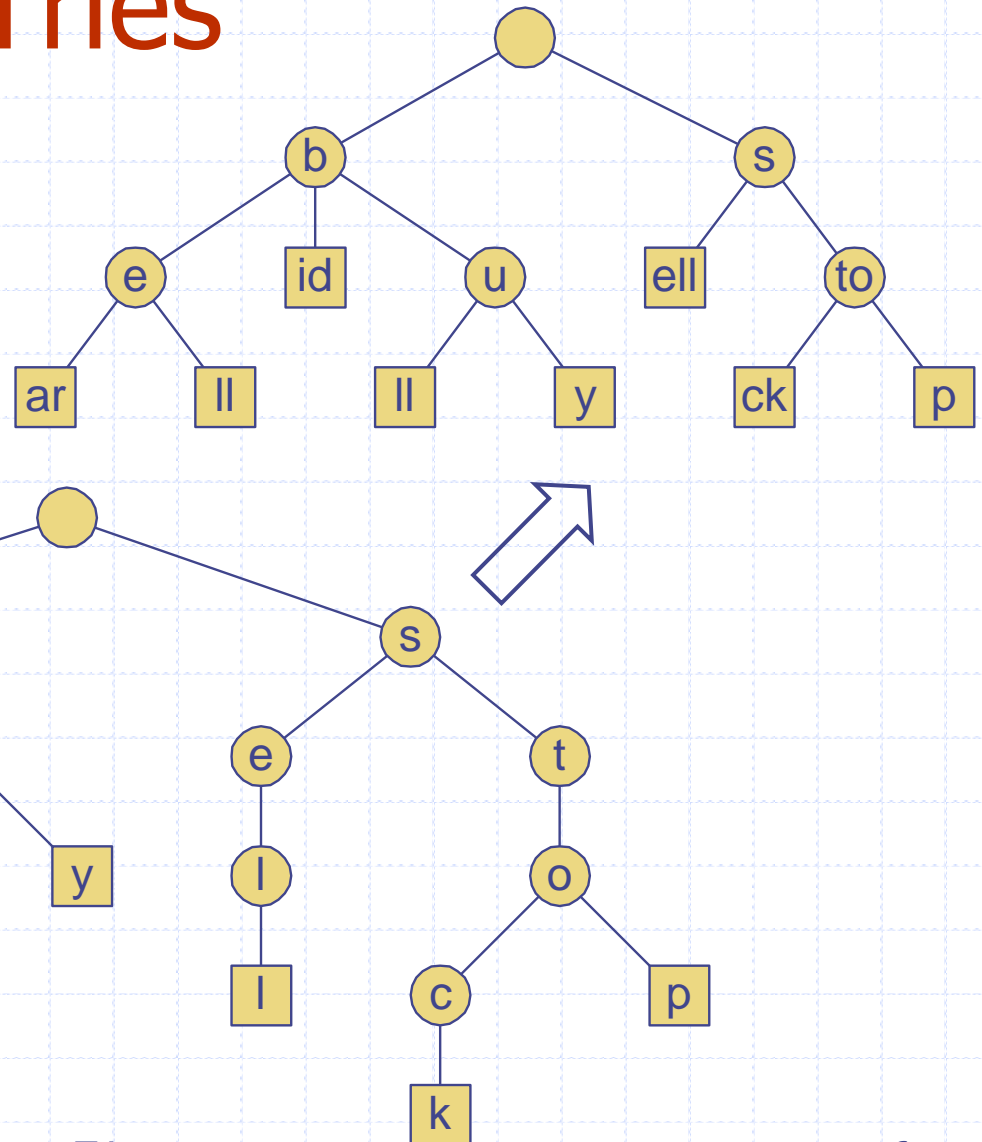
- ◆ Each leaf stores the occurrences of the associated word in the text

s	e	e		a		b	e	a	r	?		s	e	l	l		s	t	o	c	k	!		
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	
s	e	e		a		b	u	l	l	?		b	u	y			s	t	o	c	k	!		
24	25	26	27	28	29	30	31	32	33	34	35	36	37	38	39	40	41	42	43	44	45	46		
b	i	d			s	t	o	c	k	!		b	i	d			s	t	o	c	k	!		
47	48	49	50	51	52	53	54	55	56	57	58	59	60	61	62	63	64	65	66	67	68			
h	e	a	r			t	h	e		b	e	l	l	?			s	t	o	p	!			
69	70	71	72	73	74	75	76	77	78	79	80	81	82	83	84	85	86	87	88					



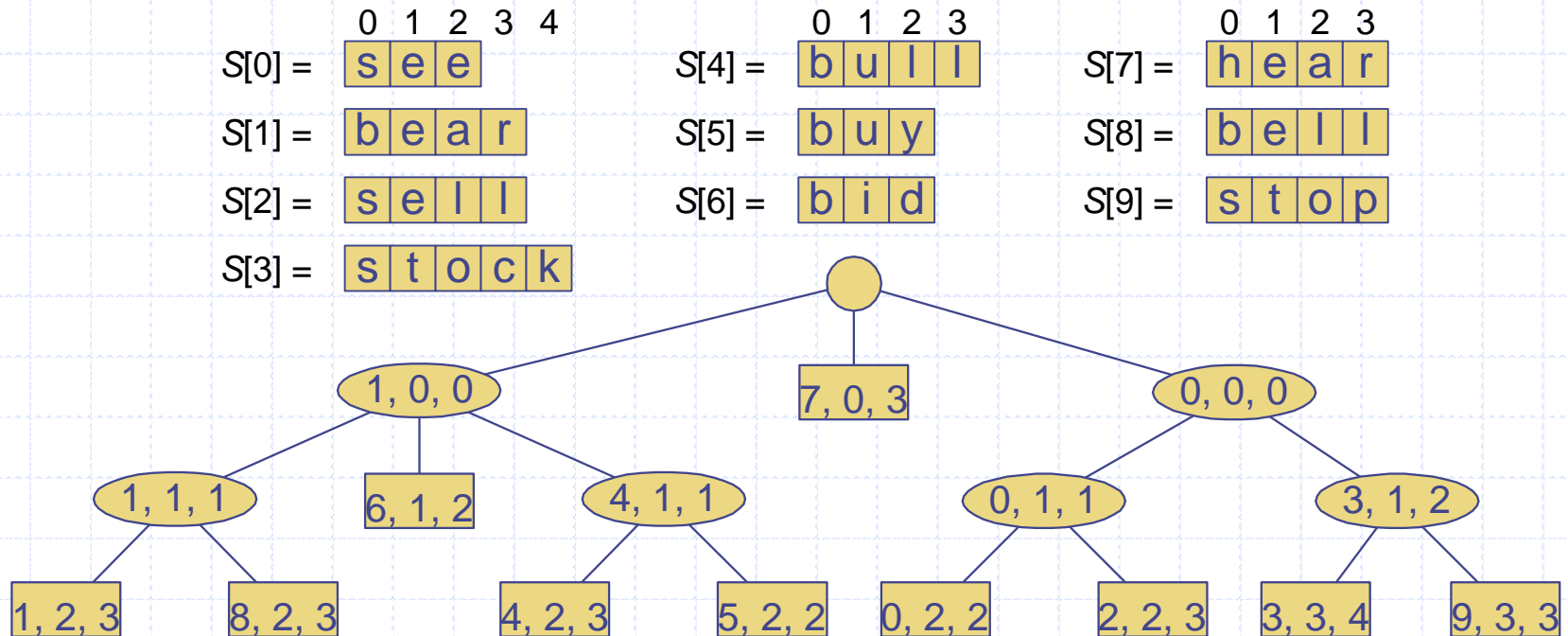
Compressed Tries

- ◆ A compressed trie has internal nodes of degree at least two
- ◆ It is obtained from standard trie by compressing chains of “redundant” nodes



Compact Representation

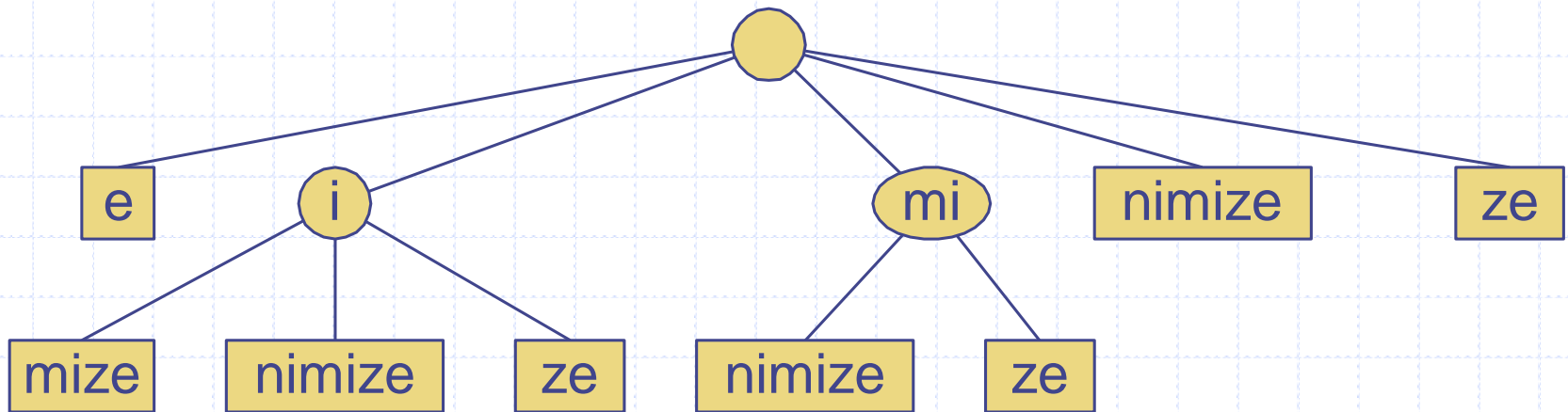
- ◆ Compact representation of a compressed trie for an array of strings:
 - Stores at the nodes ranges of indices instead of substrings
 - Uses $O(s)$ space, where s is the number of strings in the array
 - Serves as an auxiliary index structure



Suffix Trie

- ◆ The suffix trie of a string X is the compressed trie of all the suffixes of X

m	i	n	i	m	i	z	e
0	1	2	3	4	5	6	7



Analysis of Suffix Tries

- ◆ Compact representation of the suffix trie for a string X of size n from an alphabet of size d
 - Uses $O(n)$ space
 - Supports arbitrary pattern matching queries in X in $O(dm)$ time, where m is the size of the pattern
 - Can be constructed in $O(n)$ time

