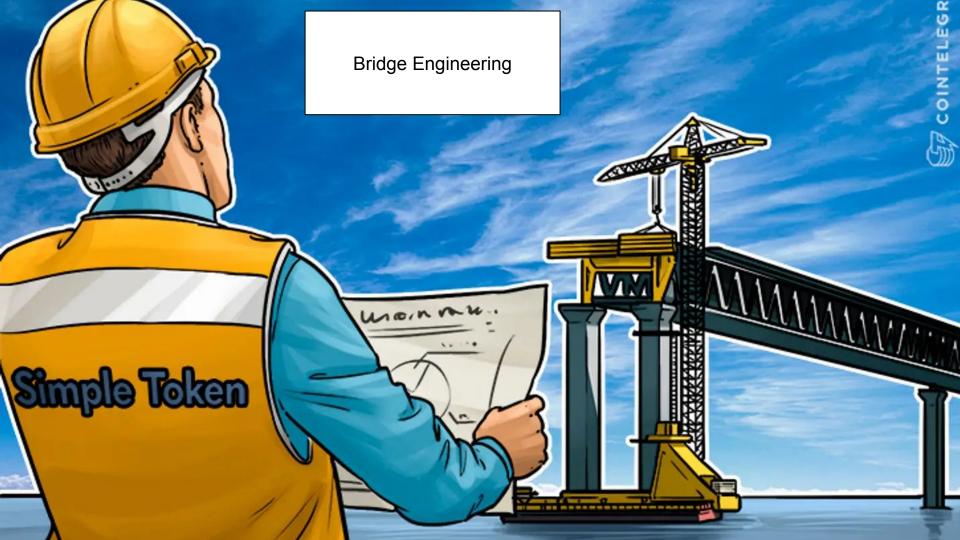


A Better Mental Model for Rollups, Plasma, and Validating Bridges

An Intro to Bridge Engineering

Patrick McCorry
Intern, Infura







coinbase

Ethereum & Users Blockchain network Bridge contract Holds user funds **Single authority**One ring to rule them all





coinbase

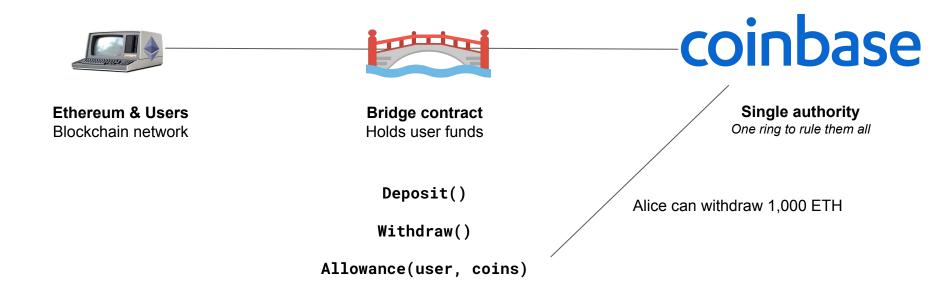
Ethereum & Users Blockchain network

Bridge contract Holds user funds **Single authority**One ring to rule them all

Deposit()

Withdraw()

Allowance(user, coins)





Ethereum & Users
Blockchain network



Bridge contract Holds user funds

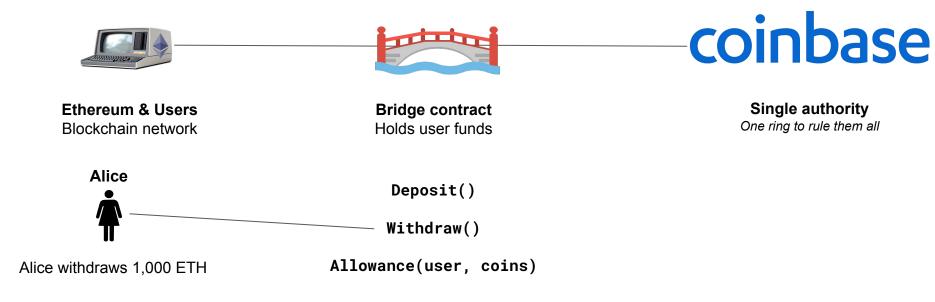
Deposit()

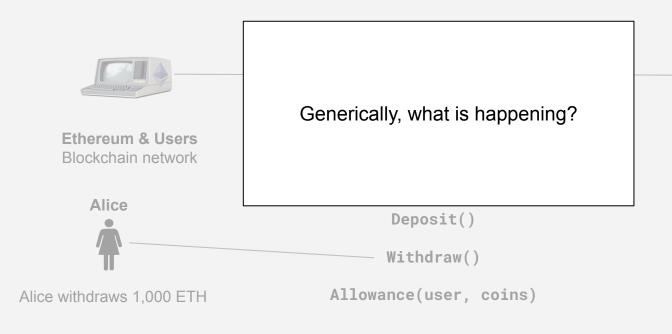
Withdraw()

Allowance(user, coins)

Coinbase has informed me that Alice can withdraw 1,000 ETH.

I trust Coinbase - the database must be OK

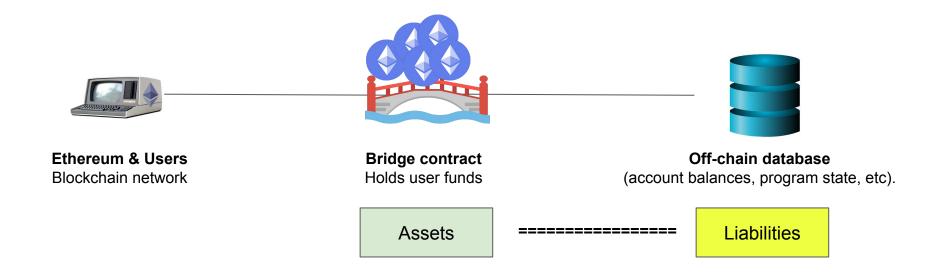




coinbase

Single authorityOne ring to rule them all

A bridge from Ethereum to an off-chain system

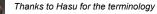


Trust assumption

Before processing a withdrawal, I need to check the database is OK

Trust assumption for bridges

have evolved over time



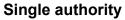












One ring to rule them all





coinbase Bitstamp

Multi-authority

K of N parties

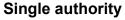






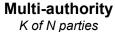






One ring to rule them all













Crypto-economic bridge

Staked investment in its success









Diberen

coinbase Bitstamp

Polygon's proof of stake bridge:

Binance 506,183,677

Stakin 322,033,445

All nodes 206,676,574

Web3Nodes 123,762,284

Anonymous 94 100,622,650

Decentral Games 70,018,093

Total 1,329,296,723

Attack Target: 1,283,657,130 matic

9/12/2021

Multi-authority

K of N parties

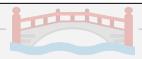




Crypto-economic bridgeStaked investment in its success







Single authority

One ring to rule them all







Trusting <10 parties to protect our funds

... sucks a bit right?















	Name	Tokens	Comment
	MtGox (2014)	850k BTC	6% of all bitcoin
Single authority hacks	Bitcoinica (2011)	61k BTC	Linode hosting provider hacked
	Bitfloor (2012)	24k BTC	Wallets stored on server
Guarding custody of tokens is not trivial	Bitstamp (2015)	19k BTC	Hot wallet hacked
I ran out of space this is only a small sample of hacks. Taylor Monahan maintains a larger list https://docs.google.com/spreadsheets/d/1ZEEAmXjpN8 kL9BvITg9GKu-dbeUra6c14YLpLkCp5Zo/edit?usp=sha ring	BTER (2015)	7k BTC	Inside job
	Gatecoin (2015)	185k ETH	Hot wallet hacked
	Bitfinex (2016)	119k BTC	Compromised server
	Bithumb (2018)	2k BTC	Hot wallet hacked
	Zaif (2018)	6k BTC	Hot wallet hacked
	Coincheck (2018)	\$534m NEM tokens	Hot wallet hacked
	Coinbin (2019)	\$26m in tokens	Inside job
	CoinBene (2019)	\$45m in tokens	Hot wallet hacked
	Binance (2019)	7k BTC	Hot wallet hacked

Multi authority hacks

... trusting multiple folk to do the right thing ... is also not good enough

5 out of 9 validators compromised

(4 compromised validators controlled by 1 company)

RONIN BRIDGE HACKED



DEVS FIND OUT 6 DAYS LATE





Old school motto

Can we transact on a off-chain system, while still allowing users to maintain self-custody of their funds?

Enabling Blockchain Innovations with Pegged Sidechains

Adam Back, Matt Corallo, Luke Dashjr, Mark Friedenbach, Gregory Maxwell, Andrew Miller, Andrew Poelstra, Jorge Timón, and Pieter Wuille*† 2014-10-22 (commit 5620e43)

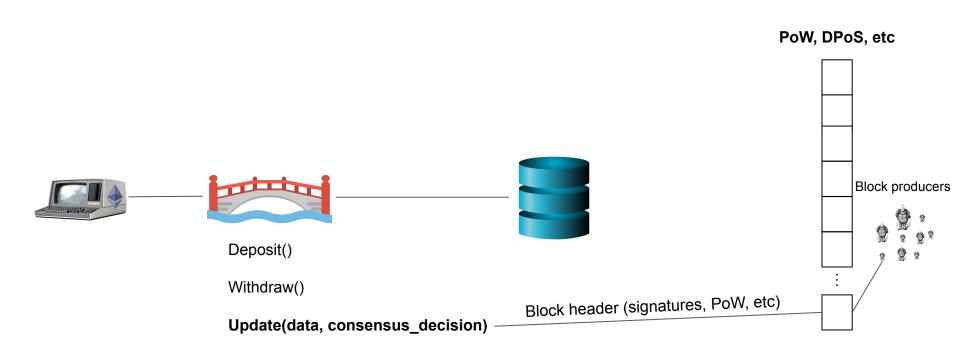
Abstract

Since the introduction of Bitcoin[Nak09] in 2009, and the multiple computer science and electronic cash innovations it brought, there has been great interest in the potential of decentralised cryptocurrencies. At the same time, implementation changes to the consensuscritical parts of Bitcoin must necessarily be handled very conservatively. As a result, Bitcoin has greater difficulty than other Internet protocols in adapting to new demands and accommodating new innovation.

We propose a new technology, pegged sidechains, which enables bitcoins and other ledger assets to be transferred between multiple blockchains. This gives users access to new and innovative cryptocurrency systems using the assets they already own. By reusing Bitcoin's currency, these systems can more easily interoperate with each other and with Bitcoin, avoiding the liquidity shortages and market fluctuations associated with new currencies. Since sidechains are separate systems, technical and economic innovation is not hindered. Despite bidirectional transferability between Bitcoin and pegged sidechains, they are isolated: in the case of a cryptographic break (or malicious design) in a sidechain, the damage is entirely confined to the sidechain itself.

This paper lays out pegged sidechains, their implementation requirements, and the work needed to fully benefit from the future of interconnected blockchains. At the heart of the original sidechain paper was a protocol to build *a trustless bridge*.

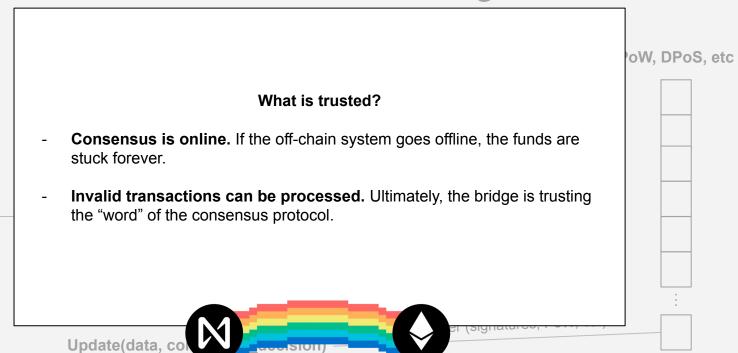
The "Consensus" Bridge



The "Consensus" Bridge

PoW, DPoS, etc What is a "consensus decision"? The judgement of a set of parties! For example, the PoW of a Bitcoin block header or a threshold of signatures from a set of validators. Update(data, consensus_decision) -

The "Consensus" Bridge



Can we really build a bridge that protects us from an all powerful authority?

Lightning strikes create plasma via a very strong jolt of electricity. Most of the Sun, and other stars, is in a plasma state. Certain regions of Earth's atmosphere contain some plasma created primarily by ultraviolet radiation from the Sun. Collectively, these regions are called the ionosphere.

https://scied.ucar.edu > learning-zone > sun-space-weather

Plasma - UCAR Center for Science Education



It all began with Plasma

Plasma: Scalable Autonomous Smart Contracts

Joseph Poon

Vitalik Buterin

joseph@lightning.network

vitalik@ethereum.org

August 11, 2017 WORKING DRAFT https://plasma.io/

Abstract

Plasma is a proposed framework for incentivized and enforced execution of smart contracts which is scalable to a significant amount of state updates per second (potentially billions) enabling the blockchain to be able to represent a significant amount of decentralized financial applications worldwide. These smart contracts are incentivized to continue operation autonomously via network transaction fees, which is ultimately reliant upon the underlying blockchain (e.g. Ethereum) to enforce transactional state transitions.

We propose a method for decentralized autonomous applications to scale to process not only financial activity, but also construct economic incentives for globally persistent data services, which may produce an alternative to centralized server farms.

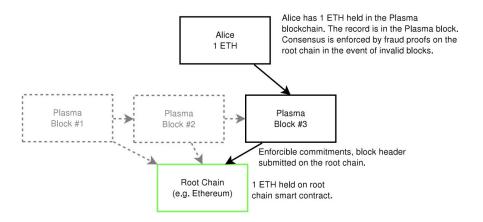
Plasma is composed of two key parts of the design: Reframing all blockchain computation into a set of MapReduce functions, and an optional method to do Proof-of-Stake token bonding on top of existing blockchains with the understanding that the Nakamoto Consensus incentives discourage block withholding.

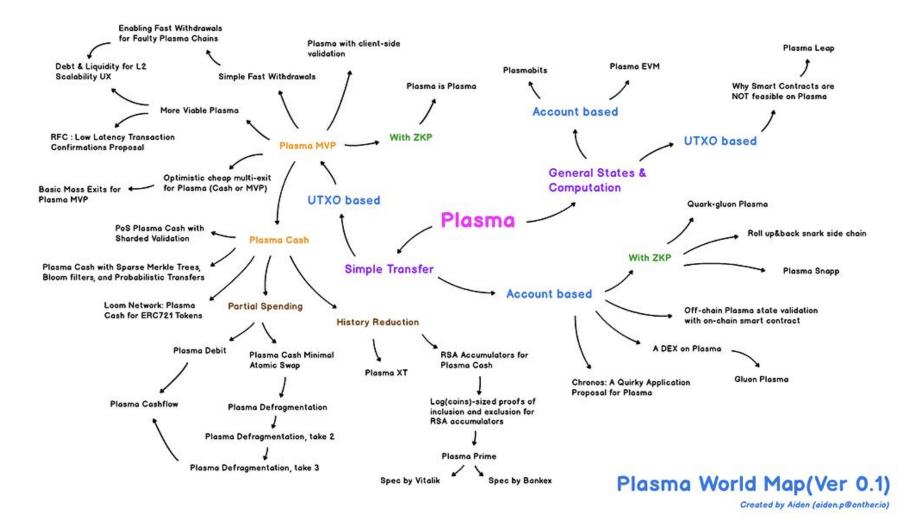
This construction is achieved by composing smart contracts on the main blockchain using fraud proofs whereby state transitions can be enforced on a parent blockchain. We compose blockchains into a tree hierarchy, and treat each as an individual branch blockchain with enforced blockchain history and MapReducible computation committed into merkle proofs. By framing one's ledger entry into a child blockchain which is enforced by the parent chain, one can enable incredible scale with minimized trust (presuming root blockchain availability and correctness).

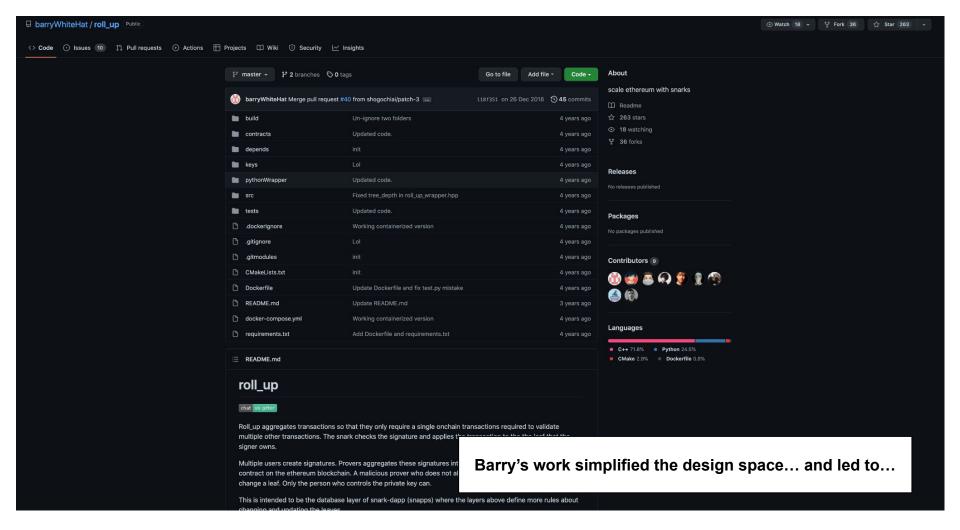
The greatest complexity around global enforcement of non-global data revolves around data availability and block withholding attacks, Plasma has mitigations for this issue by allowing for exiting faulty chains while also creating mechanisms to incentivize and enforce continued correct execution of data.

As only merkleized commitments are broadcast periodically to the root blockchain (i.e. Ethereum) during non-faulty states, this can allow for incredibly scalable, low cost transactions and computation. Plasma enables persistently operating decentralized applications at high scale.

Again, an impossible paper to read









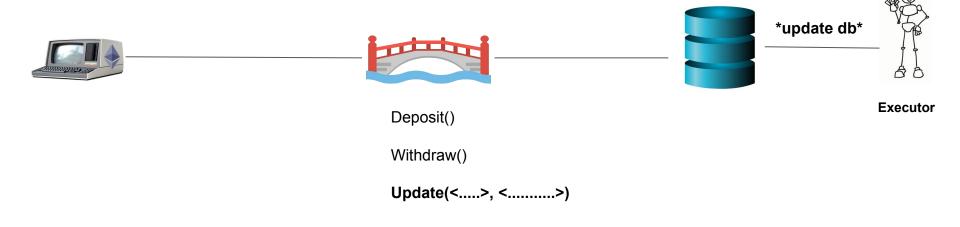


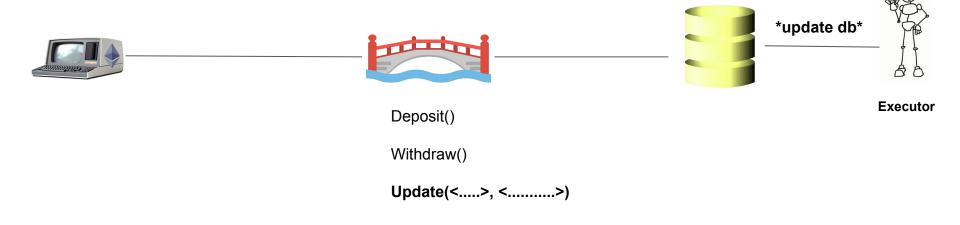


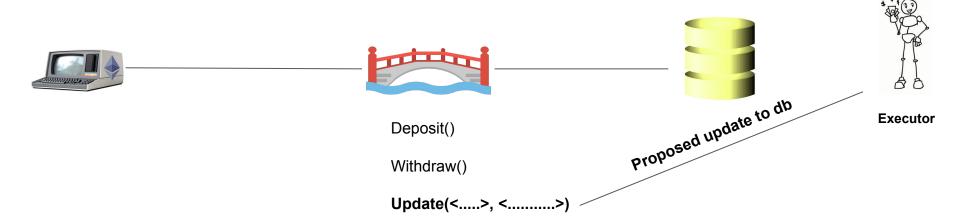
Deposit()

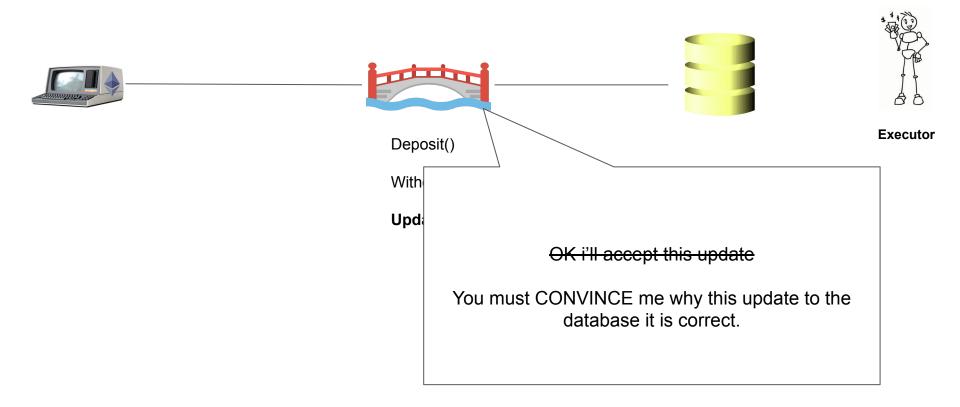
Withdraw()

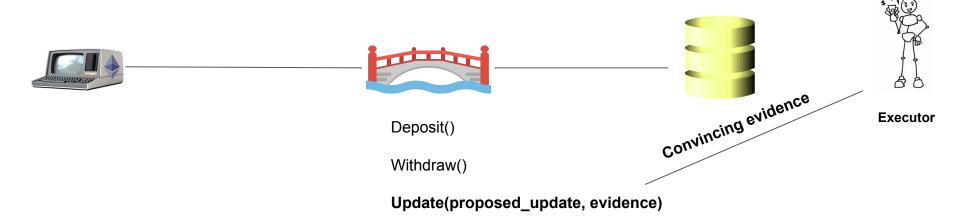
Update(<.....>)

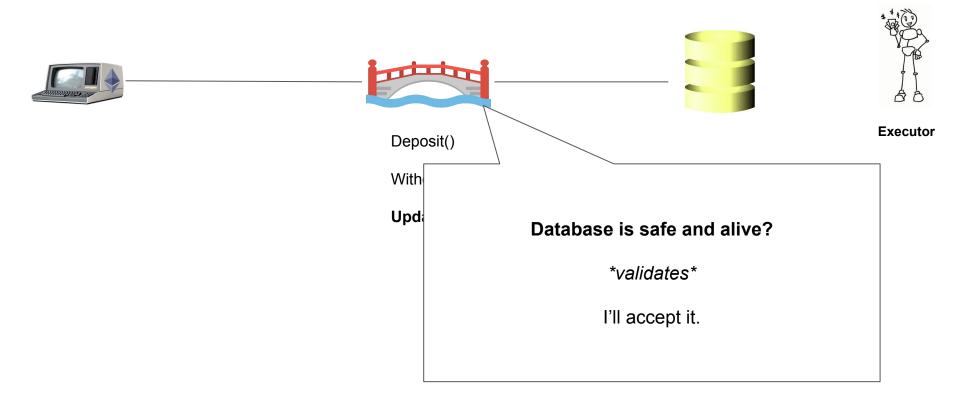








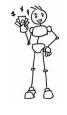












Executor

Deposit()

Withdraw()

Update(proposed_update, evidence)





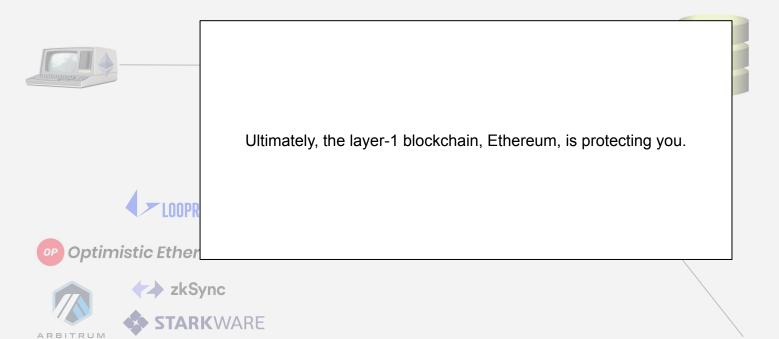






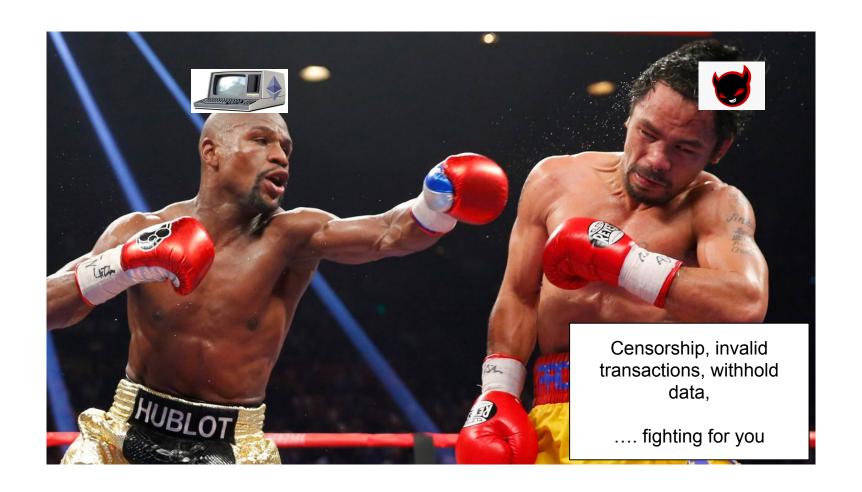
Continuously check the integrity of all proposed updates

The Validating Bridge (rollups)





Executor



Sounds so cool....

... but how do validating bridges work?



Let's try to define the environment

- Agents
 - Who are the players?
- Overview of a validating bridge
 - How does it work at a high level?
- Threat Model and Security properties
 - Who is our adversary? And what special powers do they have?
 - What are we trying to secure?

Agents



Honest user

Likes mooncats



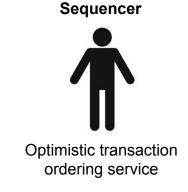
Sequencer

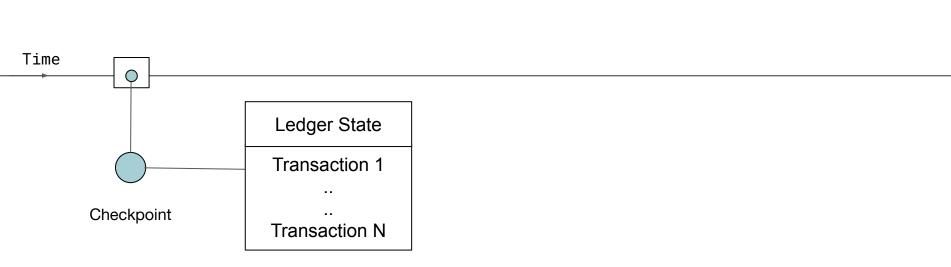
Orders transactions off-chain

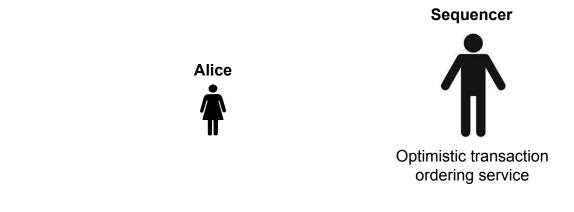


Executor

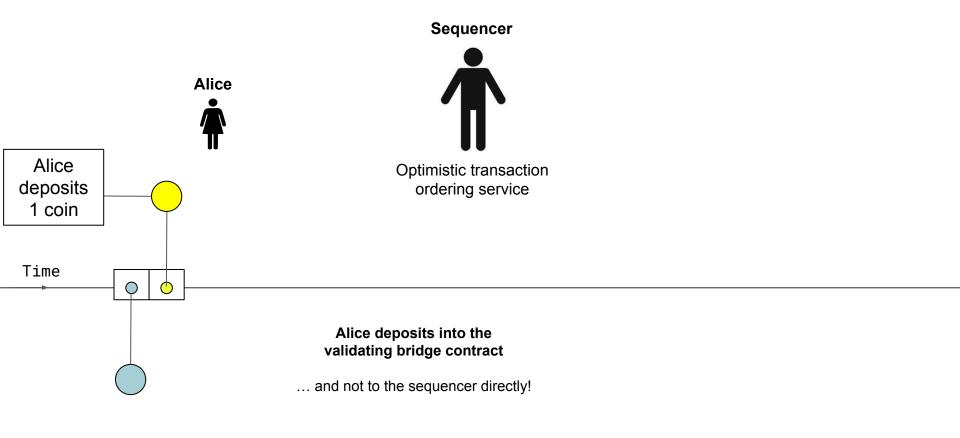
Forces bridge contract to execute transactions

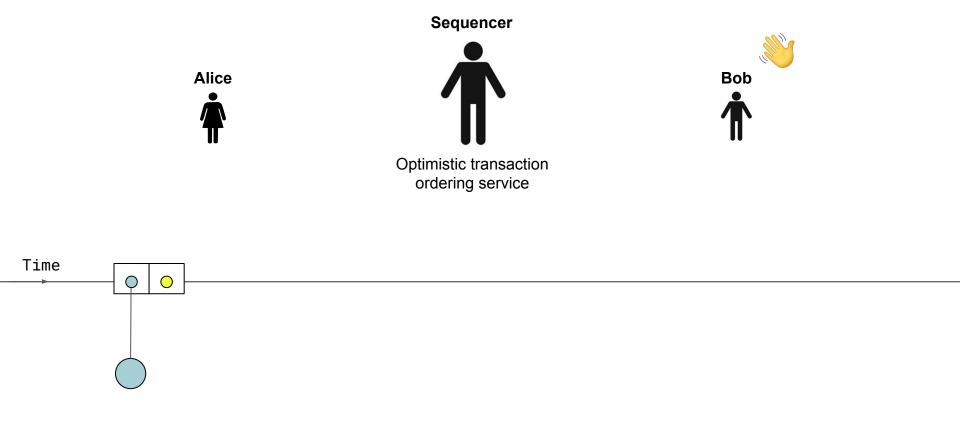


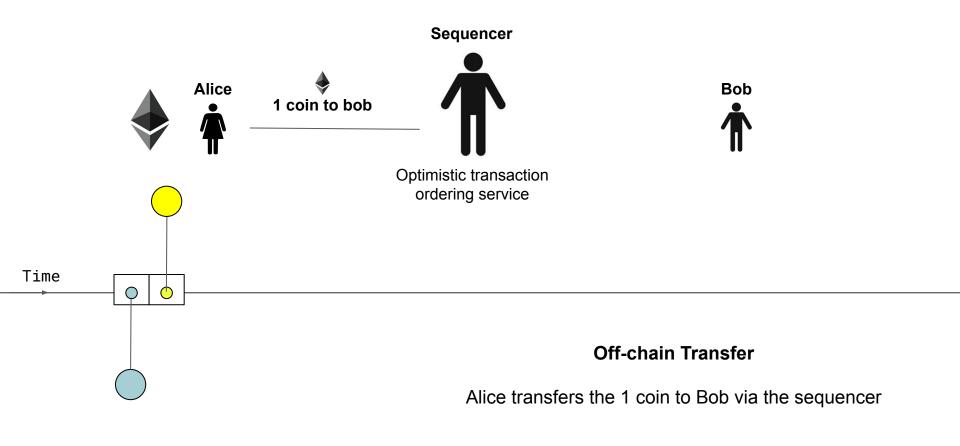


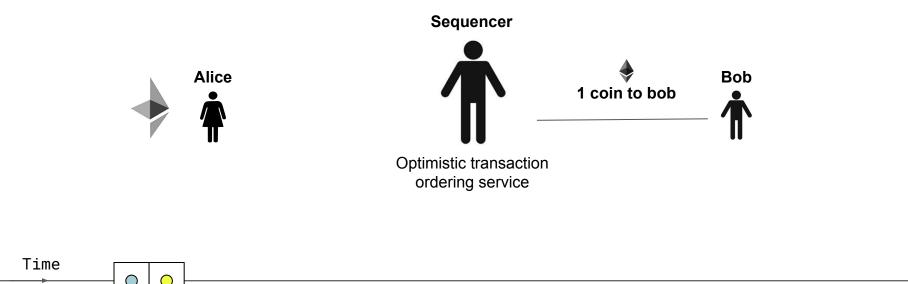














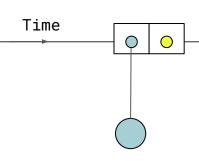
Alice transfers the 1 coin to Bob via the sequencer





Sequencer

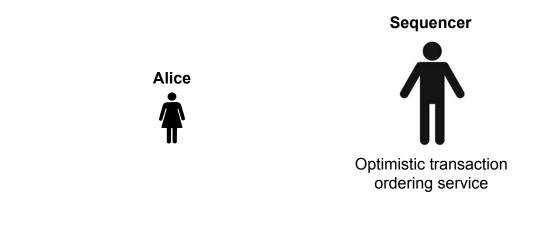




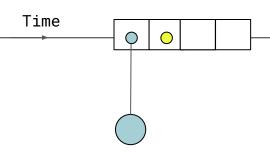
Sequencer waits around....
For more off-chain transfers...

Alice to Bob transfer is "pending" and not yet confirmed.







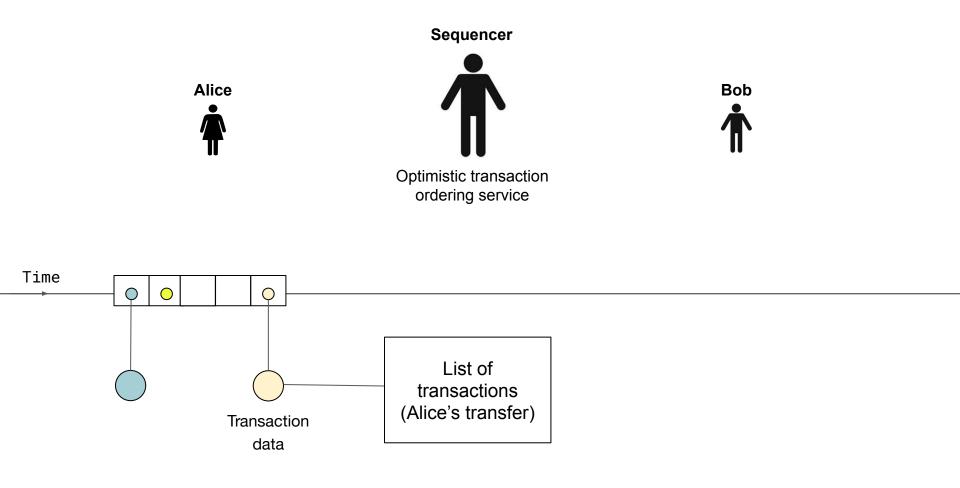


Sequencer waits around....
For more off-chain transfers...

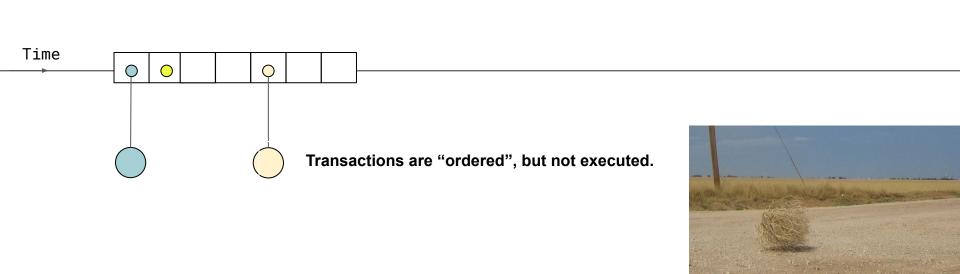
Alice to Bob transfer is pending and not yet confirmed.



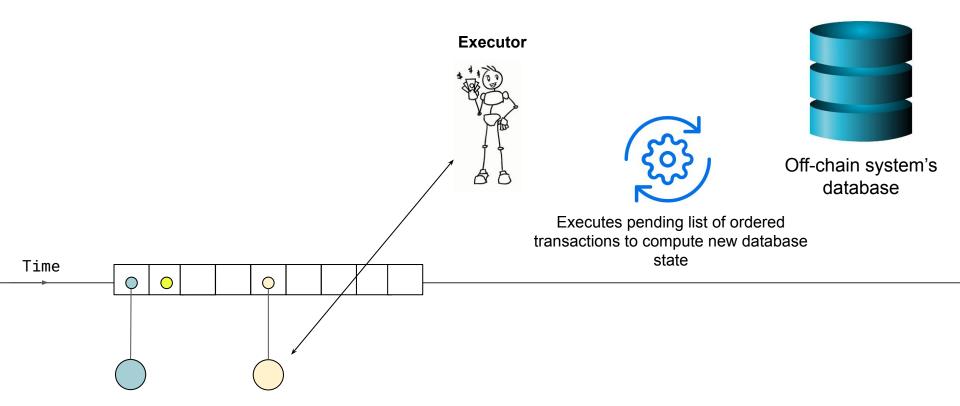
Bridge contract orders the pending transactions



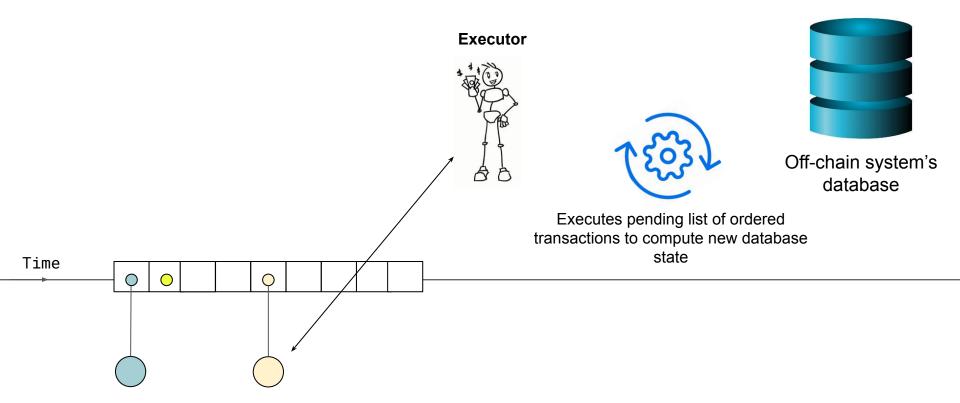
Bridge contract orders the pending transactions



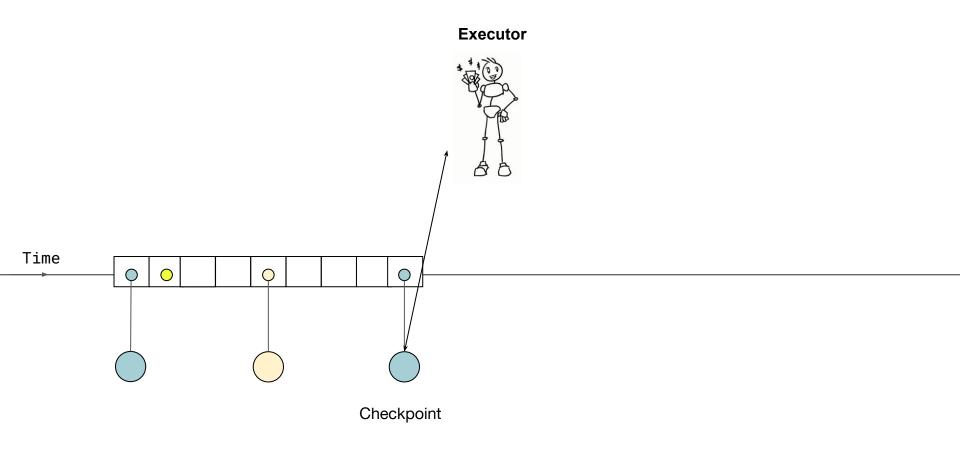
Convince a validating bridge of final execution



Convince a validating bridge of final execution

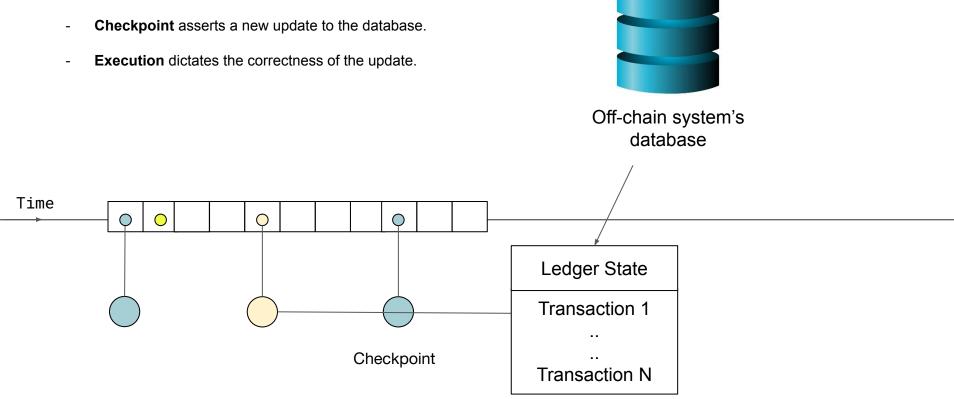


Convince a validating bridge of final execution

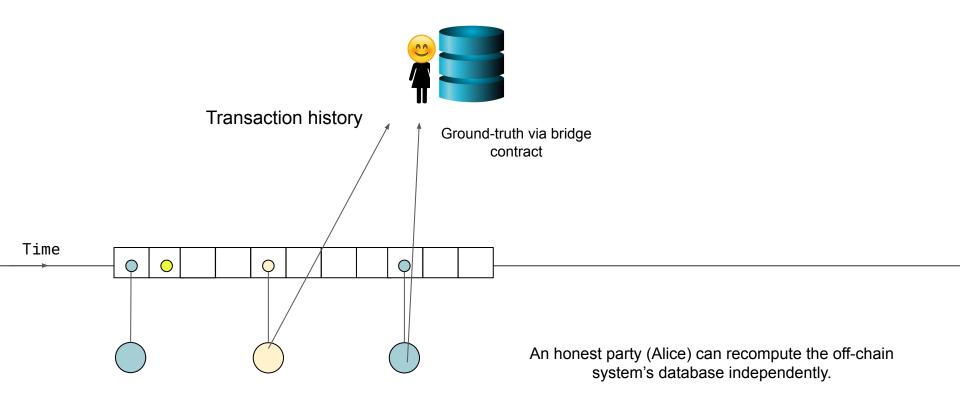


Continuously convince a validating bridge

It is a continuous process that never really ends



Proof of reserves and fully auditable by default











Adversarial threat model

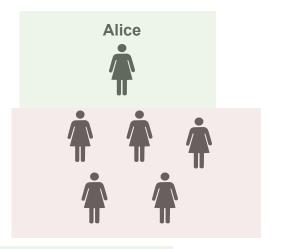


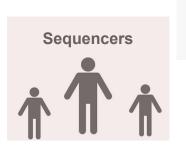


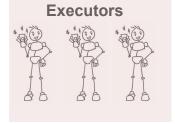


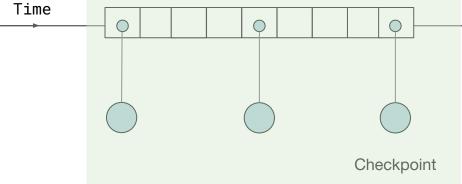
Adversarial Model Alice SAFE: Transaction to layer-1 CENSORABLE: Layer-2 tx Sequencer Time Message flow control. Adversary can view, order and drop all messages except for transactions sent to the Checkpoint layer-1 blockchain.

Adversarial Model









Corrupt nearly all parties

An honest user, optionally a challenger, and the blockchain (smart contract) vs everyone else.

Threat model (power of adversary)

- **Message flow control.** Control the order (and drop) all messages at will except for messages destined for the parent blockchain (eventual delivery protocol assumption).
- **Corrupt nearly all parties.** Adversary can corrupt all sequencers and N-1 users. They cannot corrupt one honest user and the parent blockchain.
- **Financially motivated (optional):** Adversary may require to place a security bond in the parent blockchain that is slashed if fraudulent behaviour is detected.
- Cannot break cryptography: Weak against hashes, signatures, SNARKS

We have described the most POWERFUL adversary and some rollups lack the tools to fully constrain or out-right defeat it.



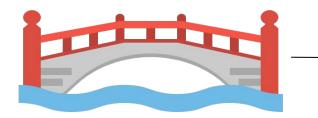
Security properties



sca0225 www.fotosearch.com



Goal: Protecting the safety & liveness of the off-chain database.



What the validating bridge checks

Data availability

Are all state updates to the database publicly available?

State transition integrity

 Are all state updates to the database valid and well-formed?

Censorship-resistance

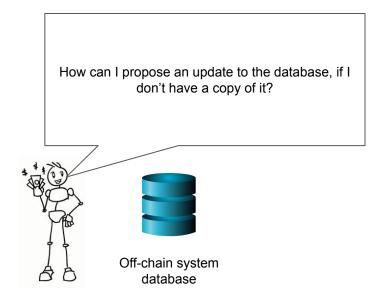
 Can the user self-enforce that a transaction will eventually execute?

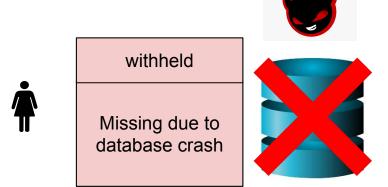
The Data Availability Problem

Data Availability Problem

- Why does the data need to be publicly available?
- What data needs to be publicly available?
- How do we guarantee it is publicly available?

Why does the data need to be publicly available?





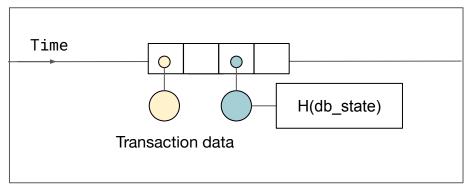
1 honest party (assistant) assumption

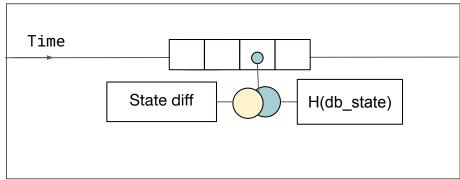
We need to assume there is one party, somewhere on the web, who will have a copy of the database and propose an update.

Adversary winning: Safety & Liveness issues

Adversary can freeze the system, potentially steal funds and lie about entries in the database.

What data needs to be publicly available?





Transaction history

Enforces the ordering of all transactions and its execution

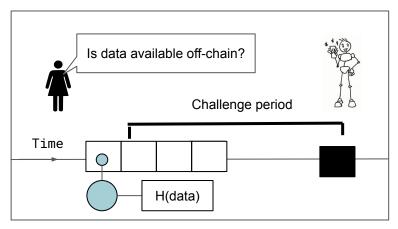
Honest party: Computes all transactions to get a copy of the database

State diffs

Bridge is not aware of individual transactions, just their aggregation

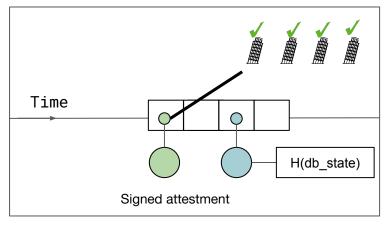
Honest party: Computes all state diffs to get a copy of the database (updates storage slots)

How do we guarantee the data is publicly available?



On-chain data availability challenge

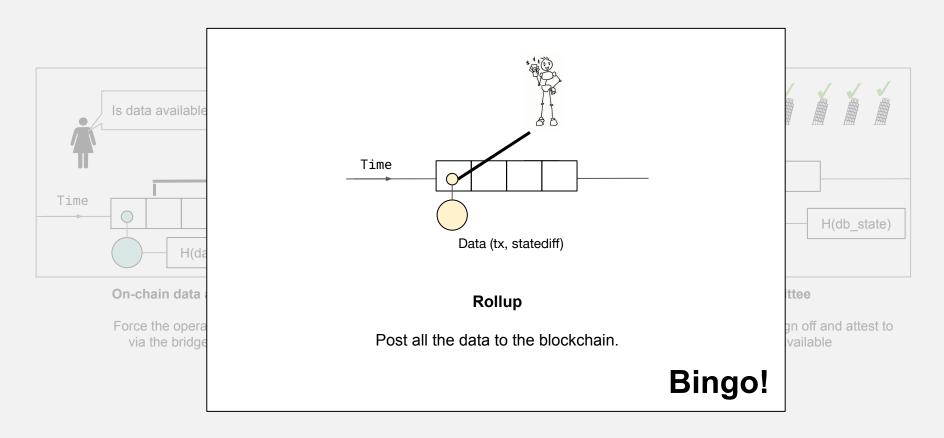
Force the operators to reveal the data via the bridge in a timely manner



Data availability committee

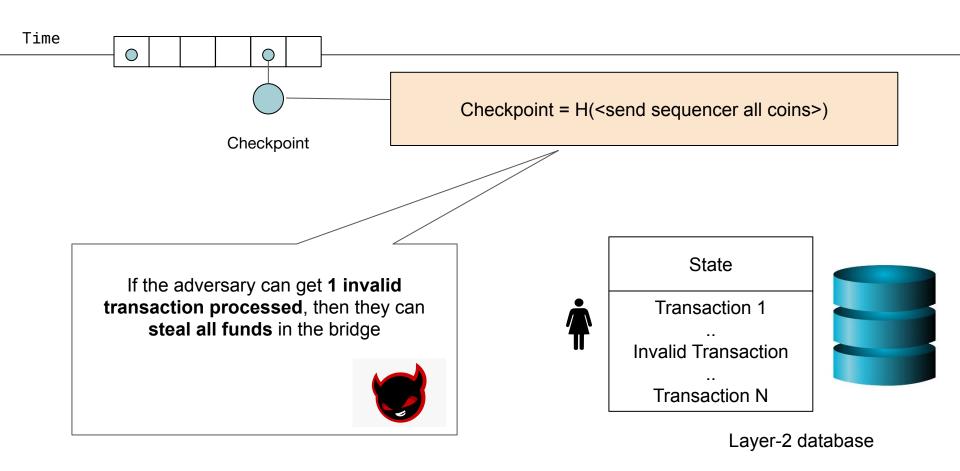
K of N data availability providers will sign off and attest to the fact the data is publicly available

How do we guarantee the data is publicly available?

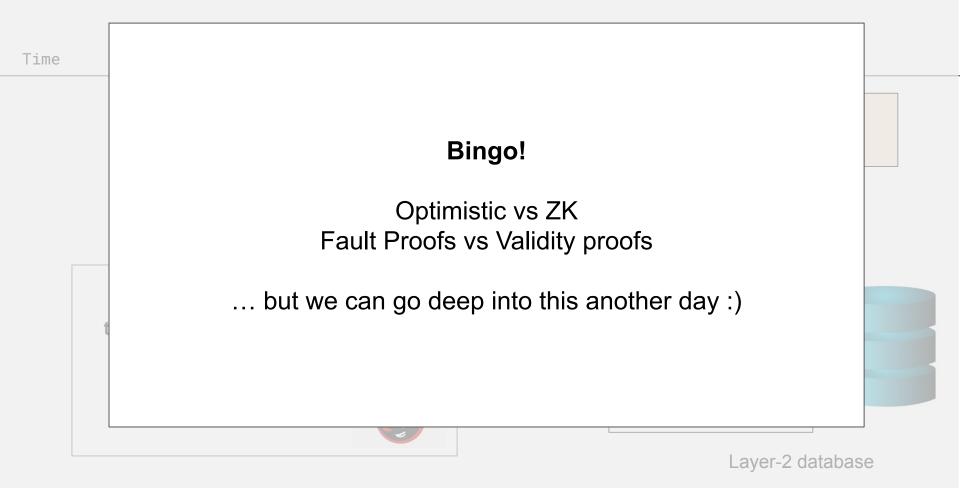


The State Transition Integrity Problem

State transition integrity (protecting the layer-2 database)



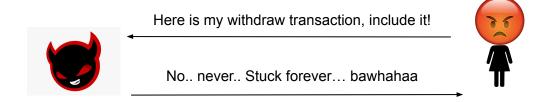
State transition integrity (protecting the layer-2 database)



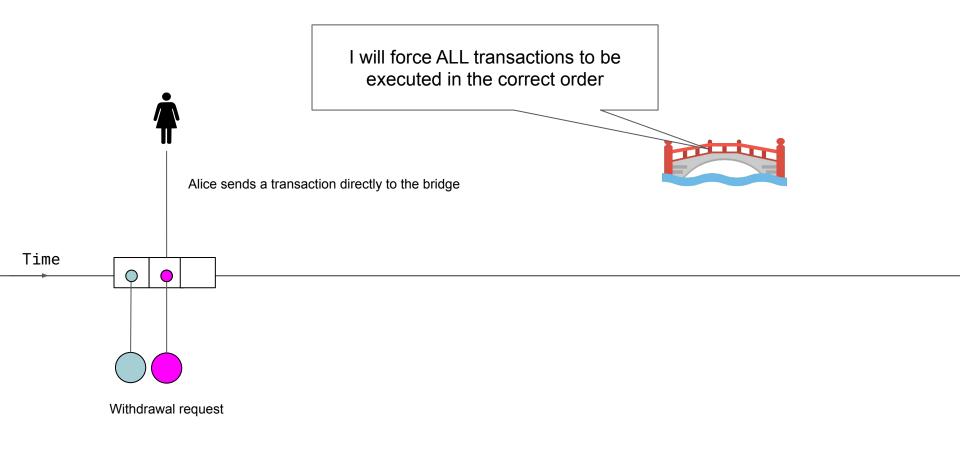
Enforcing censorship resistance

Censorship resistance

How can I withdraw my funds if the sequencer does not cooperate?



Forced inclusion: Bridge forces ordering of execution



Execution liveness (and the fast path)



Offers the "fast-path" and should have nothing to do with censorship-resistance

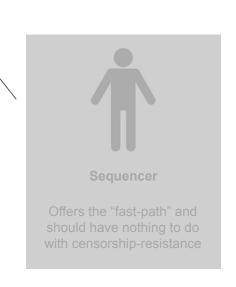


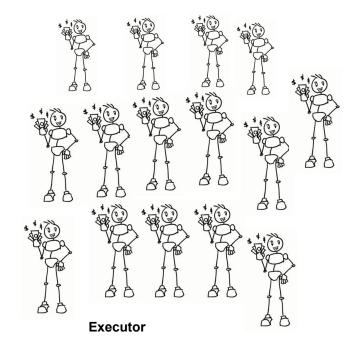
Executor

Trusted with liveness of execution (i.e., a transaction is eventually executed

Execution liveness (and the fast path)

Sequencer can be fully centralized and the off-chain system remains censorship resistant

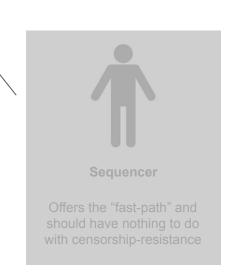


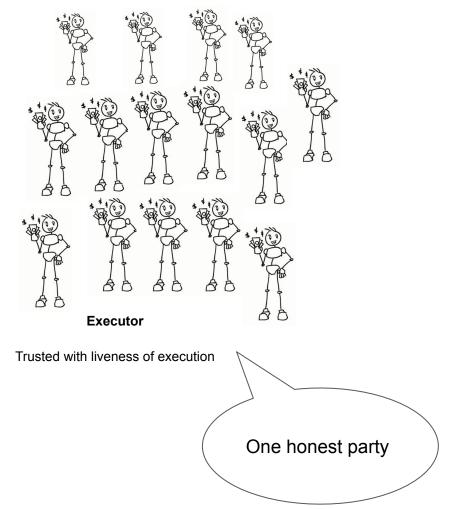


Trusted with liveness of execution

Execution liveness (and the fast path)

Sequencer can be fully centralized and the off-chain system remains censorship resistant





Security properties (summarised)

- Data availability. How does an honest user access the transaction history and recompute the same layer-2 ledger as everyone else?
- **State transition integrity.** How can we convince the layer-1 blockchain that all transactions in the layer-2 blockchain are valid?
- Censorship resistance. How can an honest user withdraw their funds from the layer-2 blockchain without the sequencer cooperation?

If we can satisfy the above properties....



Other problems emerge

Fragmentation of Assets & Interoperability

- Bypass bridge on L1 and send funds across rollups
- Gracefully handle failures while routing with smart contract execution
- Minimise trust for passive liquidity providers

Return of the data availability challenge?

- Posting data on-chain is still expensive
- EIP-4844 will help, but can optimistically avoid sending data on-chain?
- Only obstacle is the "Fisherman problem"

Experimental virtual machines on L2

- EVM-equivalence, compatibility or native?
- ZK-friendly virtual machines like Cairo?
- Compile to a simple virtual machine or build for every "opcode" of the machine?

Censorship-resistance is non-trivial

- Delay attacks by the executors to "hold out" execution of a tx
- Adversary may abuse race-condition to minimise computation
- Proving invalidity of a transaction for zkrollups (circuit overhead)

Sequencer's privilege and MEV

- Only sequencer has access to the "ordered mempool"
- Amble time to order transactions for maximum extraction
- Can we defeat MEV? Smooth MEV? Or Constrain MEV?

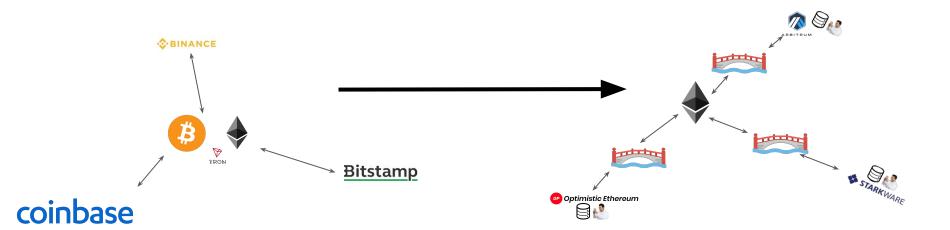
A formal model and evaluation of the "ideal bridge"

- Can we combine tx history and state diffs for data availability?
- How can we rate-limit who is an executor while upholding the 1 honest party assumption?
- Should a bridge enforce the transaction fees? Minimum quantity of execution?

Is it still worth it?

Welcome to Web3

Rise of public databases to replace custodial services (and exchanges)



Welcome to Web3

Rise of public databases to Custody is a liability for most off-chain systems Rollups, and validating bridges, will replicate the same user-experience but without the liability of custody TRON **Bitstamp** Optimistic Ethereum coinbase



It is VERY difficult to replicate human processes to secure billions of dollars



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Users do not care about the custody issues.

Operators do and they'll drive its adoption.

Custody is an unnecessary liability.