



Vampire, a Novel, Cheap to Verify, zkSNARK

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From **Count** to **Vampire**

Vampire

Novel zkSNARK for **R1CSLite**

Universal and **updatable**

Based on **Marlin**, but highly optimized

Only **4G + 2F** elements

Communication-wise: very close to **Groth16**

Only two sumchecks

Count

Univariate sumcheck

Best communication efficiency: **1G**

Very good computational complexity:

almost **linear**

zkSNARK **building blocks** and efficiency

Relation

Sumcheck

Lincheck

PCS

What relation we are showing?

Arithmetization matters

How many sumchecks do we make?

What sumcheck arguments do we use?

Do we even need them?

What do we commit to?

What commitments do we open?

zkSNARK **building blocks** and efficiency

Relation

Marlin

R1CS

Vampire

R1CSLite

Sumcheck

2x Aurora

1x Count + 1x Aurora

Lincheck

Yes

Not needed :)

**Polynomial commitment
scheme**

KZG, batched

KZG, highly batched



R1CS vs **R1CSLite**

R1CS

3 matrices

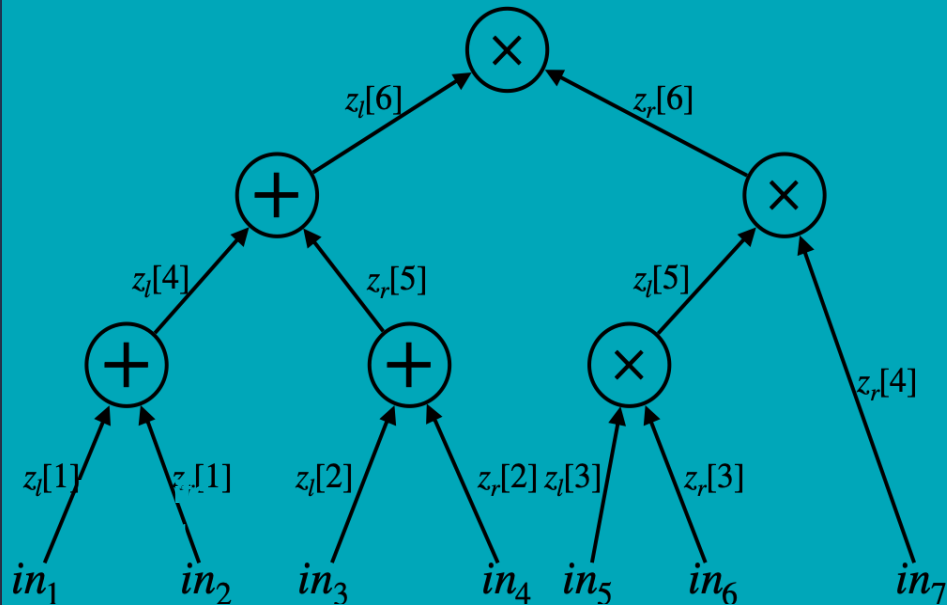
Need to show relation between the matrices **and** about the matrices

$$\underbrace{A \cdot z}_z \circ \underbrace{B \cdot z}_z = \underbrace{C \cdot z}_z$$

R1CSLite

1 matrix

$$Wz = 0$$



$$Wz = 0$$

$$\forall x \sum_y W[x, y] z[y] = 0$$

$$\forall x \sum_y W(x, y) z(y) = 0$$

...

$$\alpha \leftarrow_{\$} \mathbb{F}$$

$$\sum_y \underbrace{W(\alpha, y) z(y)}_{f(y)} = 0$$



Count - the new sumcheck

Count new sumcheck argument

$$\sum_{h \in H} f(h) = v$$

Aurora:

2G

Needs FFT

Count

1G

Doesn't need FFT

Easy to run in parallel

$$\sum_{h \in H} f(h) = v$$

$$\sum_{y \in H} f(y) = f(0) \cdot |H|$$

Assume that the prover cannot send polynomials with non-zero coefficient next to X^d

$$g(X) = g_0 + g_1 X + \dots g_{d-1} X^{d-1} + g_{d+1} X^{d+1} + \dots$$

\mathcal{P}

$$\xrightarrow{f'(X) = f(X)X^d - \frac{v}{|H|}X^d}$$

v

$$f(X) = f_0 + f_1 X + \dots f_k X^k$$

$$f'(X) = f_0 X^d - \frac{v}{|H|} X^d + f_1 X^{d+1} + \dots$$

Highly batched **KZG**

What we can **batch**?

Batch openings of **multiple** polynomials evaluated at a **single** point

$$f_1(z), f_2(z), \dots, f_k(z)$$

Batch openings of **multiple** polynomials at **multiple** points

$$f_1(z_1), f_2(z_2), \dots, f_k(z_k)$$

Counting Vampires: From Univariate Sumcheck to Updatable ZK-SNARK*

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Thank you!

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