



HUST

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HANOI UNIVERSITY OF SCIENCE AND TECHNOLOGY

ONE LOVE. ONE FUTURE.



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Applied Algorithm Lab

Nurse

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- Schedule working timetable for a nurse in N consecutive days $1, \dots, N$.

This is divided into series of consecutive working days (call **working period**)

- **Constraints:**

- Only 1 day off between 2 working period
- A working period has length in segment $[K1, K2]$

- **Input:** $N, K1, K2$

- **Output:** Number of valid ways to schedule

- Example: Input: 6 2 3

Output: 4

Explain:

110111

111011

110110

011011

- **Idea to solve #1:** Brute-force
 - A timetable can be considered as a binary sequence with length N
 - List all binary sequence and check if they are valid or not
 - Complexity: $O(n 2^n)$ (checking costs $O(n)$)
 - **Applying Branch and Bound technique:** after the first bit 1 must be K1 bit 1, if there are K2 bit 1 then the next bit must be 0 to separate 2 working periods

- **Idea to solve #2:** dynamic programming

- Consider the problem: Scheduling for day i depends on the earlier days
- Observe 2 cases: day i working and off
- Consider 2 arrays $S_0[N]$ and $S_1[N]$
 - $S_0[i]$: number of ways to schedule i days that day i off
 - $S_1[i]$: number of ways to schedule i days that day i work
- Recursive formula:

$$S_0[i] = S_1[i-1];$$

// day i off = day $i-1$ working

$$S_1[i] = \sum_{j=K1}^{K2} S_0[i-j]$$

// day i working = end of 1 working period

- Base cases: $S_0[1] = S_1[K1] = S_0[0] = 1$;
- return: $S_0[n] + S_1[n]$
- Complexity: $O(n(K2-K1))$

A large graphic on the left side of the slide. It features a dark blue background with a circular pattern of red dots of varying sizes, creating a sense of depth and movement. The word "HUST" is centered within this graphic in a bold, white, sans-serif font.

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THANK YOU !