## Yair Caplan

Intro To Algorithms

Drill #2

## Ordering Running Times

```
• Doubling size of input (n = old input size,
```

$$2n = doubled input size)$$

o 
$$n^2 \rightarrow slower by a factor of 4$$

• 
$$(2n)^2 = 4 \times n^2$$

$$\circ n^3 \rightarrow slower by a factor of 8$$

• 
$$(2n)^3 = 8 \times n^3$$

$$\circ$$
 100n<sup>2</sup>  $\rightarrow$  slower by a factor of 4

$$100(2n)^2 = 4 \times 100n^2$$

$$\circ$$
  $n \log(n) \rightarrow slower by a factor of 2$ 

$$(2n) \log(2n) = (2n)[\log(n) + \log(2)]$$

$$= 2 \times n \log(n) + 2n (assuming \log_2)$$

 $\circ$  2<sup>n</sup>  $\rightarrow$  slower by a factor of 2<sup>n</sup>

$$2^{2n} = 2^n \times 2^n$$

• Increase input by one (n = old input size,

$$n + 1 = new input size$$
)

$$\circ n^2 \rightarrow slower by a factor of 1$$

$$(n+1)^2 = n^2 + (2n+1)$$

$$o$$
  $n^3 \rightarrow slower by a factor of 1$ 

$$(n+1)^3 = n^3 + (3n^2 + 3n + 1)$$

$$\circ$$
 100n<sup>2</sup>  $\rightarrow$  slower by a factor of 1

$$100(n+1)^2 = 100n^2 + (200n+100)$$

○ 
$$n \log(n) \rightarrow slower \ by \ a \ factor \ of \ 1$$

•  $(n+1)\log(n+1) \cong n\log(n) + \log(n)$ 

○  $2^n \rightarrow slower \ by \ a \ factor \ of \ 2$ 

•  $2^{n+1} = 2 \times 2^n$ 

## Really Understanding Order-Of-Growth

- General Approach
  - Set running time equal to  $3600 \times 10^{10}$  (# of seconds in an hour times # of operations per second), solve for n. Solution is [n].
- $n^2 \rightarrow n = 6000000$ •  $n^2 = 3.6 \times 10^{13}$ ,  $n = \sqrt{3.6 \times 10^{13}} = 6000000$
- $n^3 \to n = 33019$ •  $n^3 = 3.6 \times 10^{13}$ , •  $n = \sqrt[3]{3.6 \times 10^{13}} = 33019.2724889$
- $100n^2 \rightarrow n = 600000$ •  $100n^2 = 3.6 \times 10^{13}$ ,  $n^2 = 3.6 \times 10^{11}$ ,  $n = \sqrt{3.6 \times 10^{11}} = 600000$
- $n \log(n) \rightarrow n = 2.8891 \times 10^{12}$  $\circ n \log(n) = 3.6 \times 10^{13}, n \cong 2.8891 \times 10^{12} (\log_{10})$
- $2^n \rightarrow n = 45$  $0.5 \times 10^{13}$ ,  $0.5 \times 10^{13}$ ,  $0.5 \times 10^{13} = 45.033$
- $2^{2^n} \to n = 5$  $\circ 2^{2^n} = 3.6 \times 10^{13}, \ n = \log_2(\log_2 3.6 \times 10^{13}) \cong 5.5$