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Intro to Algorithms

Drill 1

# Pick Your Constants: Big-O

1)

Definition: 
$$T(n) = O(f(n))$$
 iff  $T(n) \le c \times f(n)$  for  $n \ge n_0$ 

Givens: 
$$T(n) = 32n^2 + 17n + 1$$
,  $f(n) = n^2$ 

Prove: 
$$T(n) = O(f(n))$$
,

or: 
$$32n^2 + 17n + 1 \le c \times n^2$$
 for  $n \ge n_0$ 

Constants: 
$$n_0 = 1$$
,  $c = 51$ 

Substitute: 
$$32n^2 + 17n + 1 \le 51n^2$$
 for  $n \ge 1$ 

Prove By Induction:

Proof for 
$$n = 1$$
:

$$(32 + 17 + 1 = 50) < 51$$

*Proof for* n + 1:

*Assume*: 
$$32n^2 + 17n + 1 < 51n^2$$

Show that: 
$$32(n+1)^2 + 17(n+1) + 1 < 51(n+1)^2$$

After expansion we have:

$$32n^2 + 64n + 32 + 17n + 17 + 1 < 51n^2 + 102n + 51$$

After reordering we have:

$$(32n^2 + 17n + 1) + (64n + 33) < (51n^2) + (102n + 51)$$

*Using assumption, and* 64n + 33 < 102n + 51,

if 
$$a < c$$
 and  $b < d$ , then  $a + b < c + d$ ,

therefore 
$$32(n+1)^2 + 17(n+1) + 1 < 51(n+1)^2$$

#### --For n

Definition: 
$$T(n) = O(f(n))$$
 iff  $T(n) \le c \times f(n)$  for  $n \ge n_0$ 

Givens: 
$$T(n) = 32n^2 + 17n + 1$$
,  $f(n) = n$ 

Prove: 
$$T(n) \neq O(f(n))$$
,

or,  $32n^2 + 17n + 1 \le c \times n$  for  $n \ge n_0$  cannot hold for any c and  $n_0$ .

No matter what c is, for all  $n \ge c$  the inequality does not hold, because  $32n^2 \ge 32nc > nc$  for all c > 0.

# --For nlogn

Definition: 
$$T(n) = O(f(n))$$
 iff  $T(n) \le c \times f(n)$  for  $n \ge n_0$ 

Givens: 
$$T(n) = 32n^2 + 17n + 1$$
,  $f(n) = n\log(n)$ 

Prove: 
$$T(n) \neq O(f(n))$$
,

or,  $32n^2 + 17n + 1 \le c \times n\log(n)$  for  $n \ge n_0$  cannot hold for any c and  $n_0$ .

No matter what c is, for  $n \ge 2^c$  the inequality does not hold, because  $32(2^c)^2 + 17(2^c) + 1 > c \times 2^c \log(2^c)$  for all  $c \ge 0$ , as follows from simplifying:

$$32(2^c)^2 + 17(2^c) + 1 > c^2 \times 2^c \rightarrow \text{pull c out of log, log(2)=1}$$

$$32(2^c) + 17 + \frac{1}{2^c} > c^2 \rightarrow$$
 divide both sides by  $2^c$ 

We know that this new inequality will hold if we prove that  $32(2^c) > c^2$ . Since, by the popular lemma,  $2^c > c^2$  for  $c \ge 4$ , and for  $0 \le c < 4$ ,  $32(2^c) \ge 32 > 16 > c^2$ ,

it follows that 
$$32(2^c) > c^2$$
 for all  $c \ge 0$ .

**Alternatively**, we can show it's not  $O(n\log(n))$  using limits, by showing that no matter the c, when n gets large enough the definition for being  $O(n\log(n))$  will fail:

$$32n^2 + 17n + 1 > n^2 > c \times n \log(n)$$
  $\rightarrow$  contradiction to T(n)=O(f(n))

$$\frac{n}{\log(n)} > c \rightarrow \text{divide both sides by } n \log(n)$$

$$\lim_{n\to\infty} \frac{n}{\log(n)} = \lim_{n\to\infty} n \ln(b) = \infty > c \quad \Rightarrow \text{L'Hopital (b is the base of the log)}$$

# Pick Your Constants: Big-Omega

1)

### --For n<sup>2</sup>

Definition: 
$$T(n) = \Omega(f(n))$$
 iff  $T(n) \ge c \times f(n)$  for  $n \ge n_0$ 

Givens: 
$$T(n) = 32n^2 + 17n + 1$$
,  $f(n) = n^2$ 

Prove: 
$$T(n) = \Omega(f(n))$$
,

or: 
$$32n^2 + 17n + 1 \ge c \times n^2$$
 for  $n \ge n_0$ 

Constants: 
$$n_0 = 1$$
,  $c = 1$ 

Substitute: 
$$32n^2 + 17n + 1 \ge n^2$$
 for  $n \ge 1$ 

$$31n^2 + 17n + 1 > 0$$
 for  $n \ge 1$  subtract  $n^2$  from both sides

#### --For n

Definition: 
$$T(n) = \Omega(f(n))$$
 if  $f(n) \ge c \times f(n)$  for  $n \ge n_0$ 

Givens: 
$$T(n) = 32n^2 + 17n + 1$$
,  $f(n) = n$ 

Prove: 
$$T(n) = \Omega(f(n))$$
,

or: 
$$32n^2 + 17n + 1 \ge c \times n$$
 for  $n \ge n_0$ 

Constants:  $n_0 = 1$ , c = 1

Substitute: 
$$32n^2 + 17n + 1 \ge n$$
 for  $n \ge 1$ 

$$32n^2 + 16n + 1 > 0$$
 for  $n \ge 1$ 

2)

Definition: 
$$T(n) = \Omega(f(n))$$
 iff  $T(n) \ge c \times f(n)$  for  $n \ge n_0$ 

Givens: 
$$T(n) = 32n^2 + 17n + 1$$
,  $f(n) = n^3$ 

Prove:  $T(n) \neq \Omega(f(n))$ ,

or:  $32n^2 + 17n + 1 \ge c \times n^3$  for  $n \ge n_0$  cannot hold for any for any c and  $n_0$ 

*First we simplify*:

$$32 + \frac{17}{n} + \frac{1}{n^2} \ge c \times n$$
  $\rightarrow$  divide both sides by  $n^2$ 

Using limits, we will show that as n approaches  $\infty$ , the left will be smaller than the right side, proving  $T(n) \neq \Omega(f(n))$ :

$$\lim_{n \to \infty} 32 + \frac{17}{n} + \frac{1}{n^2} = 32 < \lim_{n \to \infty} c \times n = \infty \text{ (for } c > 0)$$

### Pick Your Constants: Big-Theta

1)

Definition: 
$$T(n) = \theta(f(n))$$
 iff  $c_1 \times f(n) \le T(n) \le c_2 \times f(n)$  for  $n \ge n_0$ 

Givens: 
$$T(n) = 32n^2 + 17n + 1$$
,  $f(n) = n^2$ 

Prove: 
$$T(n) = \theta(f(n))$$
,

or: 
$$c_1 n^2 \le 32n^2 + 17n + 1 \le c_2 n^2$$
 for  $n \ge n_0$ 

Constants: 
$$c_1 = 1, c_2 = 51, n_0 = 1$$

Substitute: 
$$n^2 \le 32n^2 + 17n + 1 \le 51n^2$$
 for  $n \ge 1$ 

Split it up:

$$32n^2 + 17n + 1 \le 51n^2$$
 for  $n \ge 1$ 

Using induction:

*Proof for* 
$$n = 1$$
:

$$(32 + 17 + 1 = 50) < 51$$

*Proof for* n + 1:

*Assume*:  $32n^2 + 17n + 1 < 51n^2$ 

Show that:  $32(n+1)^2 + 17(n+1) + 1 < 51(n+1)^2$ 

After expansion we have:

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After reordering we have:

$$(32n^2 + 17n + 1) + (64n + 33) < (51n^2) + (102n + 51)$$

*Using assumption, and* 64n + 33 < 102n + 51,

if a < c and b < d, then a + b < c + d,

therefore 
$$32(n+1)^2 + 17(n+1) + 1 < 51(n+1)^2$$

$$32n^2 + 17n + 1 \ge n^2$$
 for  $n \ge 1$ 

$$31n^2 + 17n + 1 > 0$$
 for  $n \ge 1$  subtract  $n^2$  from both sides

2)

#### --For n

Definition: 
$$T(n) = \theta(f(n))$$
 iff  $c_1 \times f(n) \le T(n) \le c_2 \times f(n)$  for  $n \ge n_0$ 

Givens: 
$$T(n) = 32n^2 + 17n + 1$$
,  $f(n) = n$ 

Prove:  $T(n) \neq \theta(f(n))$ ,

or:  $c_1 n \le 32n^2 + 17n + 1 \le c_2 n$  for  $n \ge n_0$  cannot hold for any c and  $n_0$ 

Simplify: just need to show that  $32n^2 + 17n + 1 \le c_2 n$  for  $n \ge n_0$  cannot hold for any  $c_2$  and  $n_0$ 

No matter what  $c_2$  is, for  $n \ge c$  the inequality does not hold, because  $32n^2 + 17n + 1 \ge 32nc > nc$  for all c > 0.

### --For n<sup>3</sup>

Definition:  $T(n) = \theta(f(n))$  iff  $c_1 \times f(n) \le T(n) \le c_2 \times f(n)$  for  $n \ge n_0$ 

Givens:  $T(n) = 32n^2 + 17n + 1$ ,  $f(n) = n^3$ 

Prove:  $T(n) \neq \theta(f(n))$ ,

or:  $c_1 n^3 \le 32n^2 + 17n + 1 \le c_2 n^3$  for  $n \ge n_0$  cannot hold for any c and  $n_0$ 

Simplify: just need to show that  $c_1 n^3 \le 32n^2 + 17n + 1$  for  $n \ge n_0$  cannot hold for any  $c_1$  and  $n_0$ 

*First we simplify*:

 $32 + \frac{17}{n} + \frac{1}{n^2} \ge c_1 n$   $\rightarrow$  divide both sides by  $n^2$ 

Using limits, we will show that as n approaches  $\infty$ , the left will be smaller than the right side, proving  $T(n) \neq \Omega(f(n))$ :

$$\lim_{n \to \infty} 32 + \frac{17}{n} + \frac{1}{n^2} = 32 < \lim_{n \to \infty} c_1 n = \infty \text{ (for } c_1 > 0)$$