

# Monotonic Queue — Coding Interview Notes (Light Theme)

## General Pattern Template

```
from collections import deque

def fn(arr, k):
    dq = deque() # will store indices; maintain monotonicity by values
    ans = []

    for right, x in enumerate(arr):
        # Pop from back while it violates monotonic property (for max: pop smaller)
        while dq and arr[dq[-1]] <= x: # for min queue, flip <= to >=
            dq.pop()

        dq.append(right)

        # Remove out-of-window indices from front
        if dq[0] <= right - k:
            dq.popleft()

        # Record answer when window of size k has formed
        if right + 1 >= k:
            ans.append(arr[dq[0]]) # max in window (for min, use min-queue rule)

    return ans
```

### Concept:

A **monotonic queue** (deque) maintains window candidates in sorted order by value while sliding a window. It supports  $O(1)$  amortized push/pop and  $O(1)$  query for the window's min/max.

**Time Complexity:**  $O(n)$  for all windows — each index enters and leaves the deque once.

## Key Ideas

- 1 Store indices, compare by values; pop from back while invariant breaks.
- 2 Drop indices that fall out of the window from the front.
- 3 Use  $\leq/\geq$  carefully to avoid duplicate outdated candidates.
- 4 Works for both max-queue and min-queue by flipping inequalities.

## Example 1: Sliding Window Maximum

**Goal:** For each window of size k, output the maximum.

**Approach:** Maintain a decreasing deque (front is max).

```
from collections import deque

def sliding_window_max(nums, k):
    dq = deque()
    ans = []
    for i, x in enumerate(nums):
        while dq and nums[dq[-1]] <= x:
            dq.pop()
        dq.append(i)
        if dq[0] <= i - k:
            dq.popleft()
        if i + 1 >= k:
            ans.append(nums[dq[0]])
    return ans
```

## Example 2: Shortest Subarray with Sum $\geq K$

**Goal:** Find the shortest subarray with sum  $\geq K$  (handles negatives).

**Approach:** Use prefix sums and maintain an *increasing* deque of prefix indices.

```
from collections import deque

def shortest_subarray_at_least_k(nums, K):
    n = len(nums)
    prefix = [0]
    for x in nums:
        prefix.append(prefix[-1] + x)

    ans = n + 1
    dq = deque()

    for i, s in enumerate(prefix):
        # Try to satisfy sum >= K
        while dq and s - prefix[dq[0]] >= K:
            ans = min(ans, i - dq.popleft())

        # Maintain increasing deque of prefix sums
        while dq and prefix[dq[-1]] >= s:
            dq.pop()

        dq.append(i)

    return ans if ans <= n else -1
```

## Example 3: First Negative Number in Every Window (Size k)

**Goal:** Report the first negative number for each window.

**Approach:** Maintain a deque of indices of negative numbers; pop from front when out of window.

```

from collections import deque

def first_negative_in_windows(nums, k):
    dq = deque()
    ans = []

    for i, x in enumerate(nums):
        if x < 0:
            dq.append(i)
        # remove out-of-window from front
        while dq and dq[0] <= i - k:
            dq.popleft()
        if i + 1 >= k:
            ans.append(nums[dq[0]] if dq else 0) # 0 if none present
    return ans

```

## Summary Table

Problem	Deque Invariant	Answer per Window	Complexity	Sliding window maximum	Values
decreasing	Front = max	$O(n)$	Shortest subarray $\geq K$	Prefix sums increasing	Shrink from front
$O(n)$	First negative per window	Indices of negatives	Front = first negative	$O(n)$	