

Find Top K Elements with Heap — Coding Interview Notes (Light Theme)

General Pattern Template

```
import heapq

def fn(arr, k):
    heap = []
    for num in arr:
        # do some logic to push onto heap according to problem's criteria
        heapq.heappush(heap, (CRITERIA, num))
        if len(heap) > k:
            heapq.heappop(heap)

    return [num for num in heap]
```

Concept:

To extract the **top K** items by a criterion (largest values, most frequent, closest to target, etc.), maintain a **size-k heap** while streaming through the data. In Python, `heapq` is a *min-heap* by default.

Recipe: Push tuples of the form (key, item), where key encodes the ranking criterion. Keep heap size $\leq k$; pop when size exceeds k so the heap stores the current top K by your criterion.

Complexity: $O(n \log k)$ time, $O(k)$ space.

Key Ideas

- 1 Python's `heapq` is a min-heap; store keys so the smallest key corresponds to the worst among current top K .
- 2 For 'largest K ', use `key = value` (min-heap keeps the smallest of the top K at the root).
- 3 For 'smallest K ', you can store negative values or use a max-heap emulation with `(-key, item)`.
- 4 For frequency problems, precompute counts and heapify pairs (count, item).
- 5 At the end, the heap contains K items; sort if you need ordered output.

Example 1: Top K Largest Elements

Goal: Return the K largest numbers from the array.

Approach: Maintain a size- k *min-heap* of values. Pop when size exceeds k .

```
import heapq

def top_k_largest(nums, k):
    heap = []
```

```

    for x in nums:
        heapq.heappush(heap, x)      # key is the value itself
        if len(heap) > k:
            heapq.heappop(heap)      # remove current smallest among top-K
    # heap now holds k largest in arbitrary heap order
    return sorted(heap, reverse=True)

# Example
print(top_k_largest([3,1,5,12,2,11], 3)) # Output: [12, 11, 5]

```

Example 2: Top K Frequent Elements

Goal: Return the K elements with the highest frequency.

Approach: Count with a dictionary, then keep a size-k min-heap of (count, value).

```

import heapq
from collections import Counter

def top_k_frequent(nums, k):
    freq = Counter(nums)
    heap = []
    for val, cnt in freq.items():
        heapq.heappush(heap, (cnt, val)) # min-heap by count
        if len(heap) > k:
            heapq.heappop(heap)
    # Extract values and sort by count descending (optional)
    return [val for cnt, val in sorted(heap, key=lambda x: -x[0])]

# Example
print(top_k_frequent([1,1,1,2,2,3], 2)) # Output: [1, 2]

```

Example 3: K Closest Numbers to Target

Goal: Return K numbers closest to a target t by absolute difference.

Approach: Keep a size-k *max-heap* of (distance, value) by pushing negative distances.

```

import heapq

def k_closest(nums, k, t):
    heap = [] # will store (-distance, value) so root is the farthest among kept
    for x in nums:
        key = abs(x - t)
        heapq.heappush(heap, (-key, x))
        if len(heap) > k:
            heapq.heappop(heap) # remove farthest among the kept
    # Extract values; result order not guaranteed unless sorted
    return [x for _, x in sorted(heap, key=lambda p: abs(p[1] - t))]

# Example
print(k_closest([5,6,7,8,9], 3, 7)) # Output: [7, 6, 8] (order may vary)

```

Summary Table

Problem	Heap	Content	Key/Trick	Complexity
K largest	Min-heap of values	Pop when size > k		$O(n \log k)$
most frequent	Min-heap of (count, val)	Use Counter, key by count		$O(u \log k)$ (u=#unique)
K closest to t	Max-heap via negatives	Push $(- x-t , x)$		$O(n \log k)$