

Linked List: Fast and Slow Pointer — Coding Interview Notes (Light Theme)

General Pattern Template

```
def fn(head):
    slow = head
    fast = head
    ans = 0

    while fast and fast.next:
        # do logic
        slow = slow.next
        fast = fast.next.next

    return ans
```

Concept:

The **Fast and Slow Pointer** (also called *Tortoise and Hare*) pattern uses two pointers that traverse a linked list at different speeds — typically one moves one node at a time (slow), and the other two nodes at a time (fast).

Purpose: Detect cycles, find middle nodes, measure lengths, or identify intersections efficiently.

Time Complexity: $O(n)$, **Space Complexity:** $O(1)$.

Key Ideas

- 1 Use two pointers moving at different speeds through the list.
- 2 When the fast pointer reaches the end, the slow pointer is at the midpoint.
- 3 If a cycle exists, fast and slow will eventually meet inside it.
- 4 Efficient for $O(1)$ space solutions to linked list problems.

Example 1: Find Middle of Linked List

Goal: Return the middle node of a singly linked list.

Approach: Move slow by one and fast by two steps. When fast reaches the end, slow is at the middle.

```
def find_middle(head):
    slow = fast = head
    while fast and fast.next:
        slow = slow.next
        fast = fast.next.next
    return slow
```

Example 2: Detect Cycle in Linked List (Floyd's Algorithm)

Goal: Detect whether a cycle exists in a linked list.

Approach: If fast and slow pointers ever meet, a cycle exists.

```
def has_cycle(head):
    slow = fast = head
    while fast and fast.next:
        slow = slow.next
        fast = fast.next.next
        if slow == fast:
            return True
    return False
```

Example 3: Find Start of Cycle

Goal: If a cycle exists, find the node where the cycle begins.

Approach: After detecting the meeting point, move one pointer to head. Move both at 1x speed — they meet at cycle start.

```
def detect_cycle_start(head):
    slow = fast = head

    # Step 1: Detect cycle
    while fast and fast.next:
        slow = slow.next
        fast = fast.next.next
        if slow == fast:
            break
    else:
        return None # No cycle

    # Step 2: Find cycle start
    slow = head
    while slow != fast:
        slow = slow.next
        fast = fast.next

    return slow
```

Summary Table

Problem	Approach	Complexity
Find middle of list	Fast 2x, slow 1x	$O(n)$
Detect cycle	Floyd's algorithm (fast meets slow)	$O(n)$
Find cycle start	Reset one pointer, move both 1x	$O(n)$