

QuantA&M Book Club Meeting 7

QuantA&M of Texas A&M University



Brief Recap Quantum Fourier Transform



Large Summations

Examples of sigma notation:

$$\sum_{n=1}^{3} n \qquad \qquad \sum_{n=1}^{3} x^{n}$$

$$= 1 + 2 + 3 \qquad = x^{1} + x^{2} + x^{3}$$



Large Products

Examples of pi notation:

$$\prod_{n=1}^{3} n \qquad \qquad \prod_{n=1}^{3} x^{n} \\
= 1 * 2 * 3 \qquad \qquad = x^{1} * x^{2} * x^{3}$$

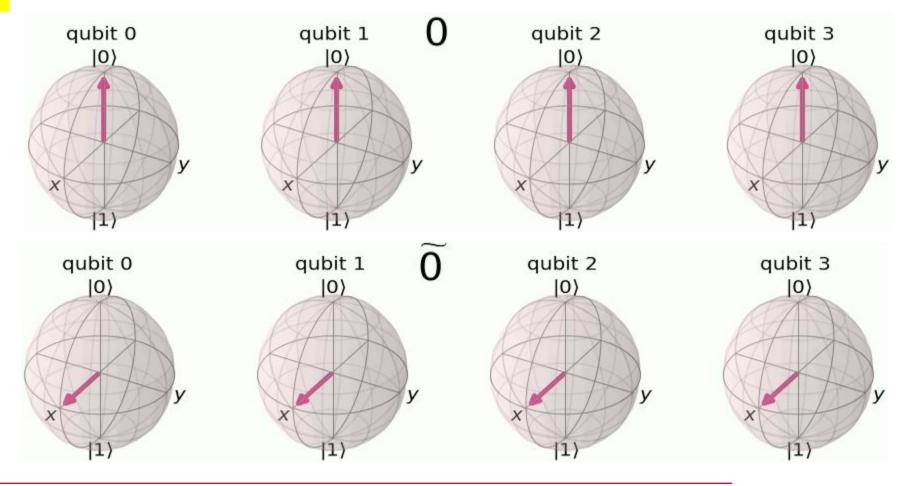
Large Tensor Products

New tensor notation:

$$\bigotimes_{n=1}^{3} e^{i\pi*k} |1\rangle$$

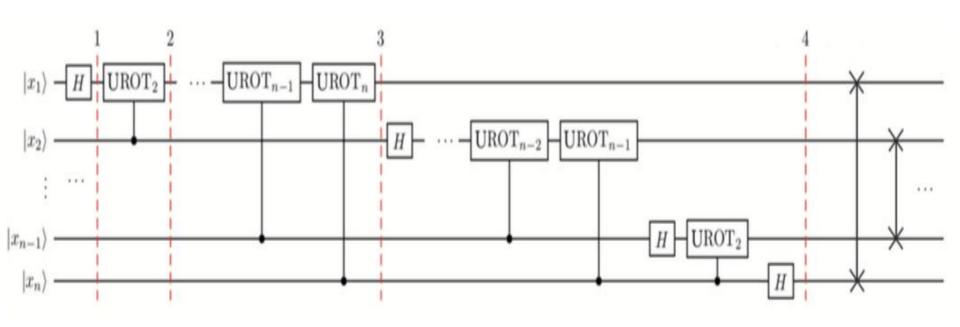
$$= e^{i\pi*1} |1\rangle \otimes e^{i\pi*2} |1\rangle \otimes e^{i\pi*3} |1\rangle$$

What does a QFT do?





The QFT Circuit





The 'Unitary Rotation' Gate

The URot-gate rotates an amount depending on the value of k.

'k' is the power of the control bit.

$$UROT_k = egin{bmatrix} 1 & 0 \ 0 & \exp\left(rac{2\pi i}{2^k}
ight) \end{bmatrix}$$

The QFT Formula

The circuit we built gives us the formula for the QFT:

$$\frac{1}{\sqrt{2}} \bigg[|0\rangle + \exp\bigg(\frac{2\pi i}{2^n} x\bigg) |1\rangle \bigg] \otimes \frac{1}{\sqrt{2}} \bigg[|0\rangle + \exp\bigg(\frac{2\pi i}{2^{n-1}} x\bigg) |1\rangle \bigg] \otimes \ldots \otimes \frac{1}{\sqrt{2}} \bigg[|0\rangle + \exp\bigg(\frac{2\pi i}{2^2} x\bigg) |1\rangle \bigg] \otimes \frac{1}{\sqrt{2}} \bigg[|0\rangle + \exp\bigg(\frac{2\pi i}{2^1} x\bigg) |1\rangle \bigg] \otimes \frac{1}{\sqrt{2}} \bigg[|0\rangle + \exp\bigg(\frac{2\pi i}{2^n} x\bigg) |1\rangle \bigg] \otimes \frac{1}{\sqrt{2}} \bigg[|0\rangle + \exp\bigg(\frac{2\pi i}{2^n} x\bigg) |1\rangle \bigg] \otimes \frac{1}{\sqrt{2}} \bigg[|0\rangle + \exp\bigg(\frac{2\pi i}{2^n} x\bigg) |1\rangle \bigg] \otimes \frac{1}{\sqrt{2}} \bigg[|0\rangle + \exp\bigg(\frac{2\pi i}{2^n} x\bigg) |1\rangle \bigg] \otimes \frac{1}{\sqrt{2}} \bigg[|0\rangle + \exp\bigg(\frac{2\pi i}{2^n} x\bigg) |1\rangle \bigg] \otimes \frac{1}{\sqrt{2}} \bigg[|0\rangle + \exp\bigg(\frac{2\pi i}{2^n} x\bigg) |1\rangle \bigg] \otimes \frac{1}{\sqrt{2}} \bigg[|0\rangle + \exp\bigg(\frac{2\pi i}{2^n} x\bigg) |1\rangle \bigg] \otimes \frac{1}{\sqrt{2}} \bigg[|0\rangle + \exp\bigg(\frac{2\pi i}{2^n} x\bigg) |1\rangle \bigg] \otimes \frac{1}{\sqrt{2}} \bigg[|0\rangle + \exp\bigg(\frac{2\pi i}{2^n} x\bigg) |1\rangle \bigg] \otimes \frac{1}{\sqrt{2}} \bigg[|0\rangle + \exp\bigg(\frac{2\pi i}{2^n} x\bigg) |1\rangle \bigg] \otimes \frac{1}{\sqrt{2}} \bigg[|0\rangle + \exp\bigg(\frac{2\pi i}{2^n} x\bigg) |1\rangle \bigg] \otimes \frac{1}{\sqrt{2}} \bigg[|0\rangle + \exp\bigg(\frac{2\pi i}{2^n} x\bigg) |1\rangle \bigg] \otimes \frac{1}{\sqrt{2}} \bigg[|0\rangle + \exp\bigg(\frac{2\pi i}{2^n} x\bigg) |1\rangle \bigg] \otimes \frac{1}{\sqrt{2}} \bigg[|0\rangle + \exp\bigg(\frac{2\pi i}{2^n} x\bigg) |1\rangle \bigg] \otimes \frac{1}{\sqrt{2}} \bigg[|0\rangle + \exp\bigg(\frac{2\pi i}{2^n} x\bigg) |1\rangle \bigg] \otimes \frac{1}{\sqrt{2}} \bigg[|0\rangle + \exp\bigg(\frac{2\pi i}{2^n} x\bigg) |1\rangle \bigg] \otimes \frac{1}{\sqrt{2}} \bigg[|0\rangle + \exp\bigg(\frac{2\pi i}{2^n} x\bigg) |1\rangle \bigg] \otimes \frac{1}{\sqrt{2}} \bigg[|0\rangle + \exp\bigg(\frac{2\pi i}{2^n} x\bigg) |1\rangle \bigg] \otimes \frac{1}{\sqrt{2}} \bigg[|0\rangle + \exp\bigg(\frac{2\pi i}{2^n} x\bigg) |1\rangle \bigg] \otimes \frac{1}{\sqrt{2}} \bigg[|0\rangle + \exp\bigg(\frac{2\pi i}{2^n} x\bigg) |1\rangle \bigg] \otimes \frac{1}{\sqrt{2}} \bigg[|0\rangle + \exp\bigg(\frac{2\pi i}{2^n} x\bigg) |1\rangle \bigg] \otimes \frac{1}{\sqrt{2}} \bigg[|0\rangle + \exp\bigg(\frac{2\pi i}{2^n} x\bigg) |1\rangle \bigg] \otimes \frac{1}{\sqrt{2}} \bigg[|0\rangle + \exp\bigg(\frac{2\pi i}{2^n} x\bigg) |1\rangle \bigg] \otimes \frac{1}{\sqrt{2}} \bigg[|0\rangle + \exp\bigg(\frac{2\pi i}{2^n} x\bigg) |1\rangle \bigg] \otimes \frac{1}{\sqrt{2}} \bigg[|0\rangle + \exp\bigg(\frac{2\pi i}{2^n} x\bigg) |1\rangle \bigg] \otimes \frac{1}{\sqrt{2}} \bigg[|0\rangle + \exp\bigg(\frac{2\pi i}{2^n} x\bigg) \bigg[|0\rangle + \exp\bigg(\frac{2\pi i}{2^n} x\bigg) \bigg[|0\rangle + \exp\bigg(\frac{2\pi i}{2^n} x\bigg] \bigg[|0\rangle + \exp\bigg(\frac{2\pi i}{2^n} x\bigg) \bigg[|0\rangle + \exp\bigg(\frac{2\pi i}{2^n} x\bigg] \bigg[|0\rangle + \exp\bigg(\frac{2\pi i}{2^n} x\bigg) \bigg[|0\rangle + \exp\bigg(\frac{2\pi i}{2^n} x\bigg] \bigg[|0\rangle + \exp\bigg(\frac{2\pi i}{2^n}$$

$$= \frac{1}{\sqrt{N}} \bigotimes_{k=1}^{n} (|0\rangle + e^{\frac{2\pi i x_{b10}}{2^k}} |1\rangle)$$



The QFT Formula: An Example

Plugging in x=6 provides the following answer:

$$\frac{1}{\sqrt{8}} \bigotimes_{k=1}^{n} (|0\rangle + e^{\frac{2\pi i6}{2^k}} |1\rangle)$$

$$= \frac{1}{\sqrt{8}} [(|0\rangle + e^{\frac{2\pi i6}{2^1}} |1\rangle) \otimes (|0\rangle + e^{\frac{2\pi i6}{2^2}} |1\rangle) \otimes (|0\rangle + e^{\frac{2\pi i6}{2^3}} |1\rangle)]$$



Binary and Decimal

 The notation for a string of binary and its decimal form can be shown as follows (b is short for base):

$$y_{b2} = y_1 y_2 \dots y_n$$
$$y_{b10} = \sum_{k=1}^{n} 2^{n-k} * y_k$$



Quantum Phase Estimation (QPE)



• Imagine we run a qubit through a circuit. This circuit is defined as the following:

$$U|\psi\rangle = e^{i\theta} |\psi\rangle$$



In other words, we rotate the qubit by e^{iθ}.

$$U|\psi\rangle = e^{i\theta} |\psi\rangle$$



 Of course, I did not define θ. That is because in this example, we do not know the value of it.

$$U|\psi\rangle = e^{i\theta} |\psi\rangle$$
$$\theta = ?$$



• Is there a way to make a circuit that gets us this mystery θ-value?

$$U|\psi\rangle = e^{i\theta} |\psi\rangle$$
$$\theta = ?$$



Implementing a 1-bit QPE

 What we want is the exponent of our Unitary gate to be 'saved' in our circuit somehow.

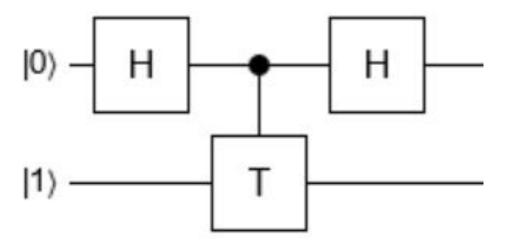
$$U|\psi\rangle = e^{i\theta} |\psi\rangle$$
$$\theta = ?$$



 Going back to our previous lectures, does anyone remember how we were able to manipulate qubits outside of our gate, without modifying the qubits inside the main gate?



 We used phase kickback. Remembering this, we apply this concept to our naïve 1-qubit gate:



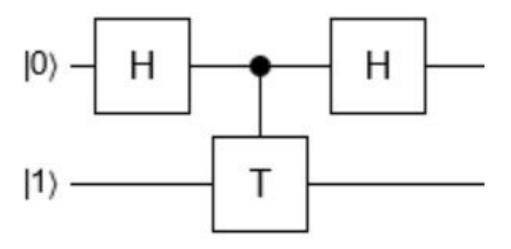
Reminder, our "mystery" T-gate is equal to:

$$T = egin{pmatrix} 1 & 0 \ 0 & \exp\left(rac{i\pi}{4}
ight) \end{pmatrix}$$

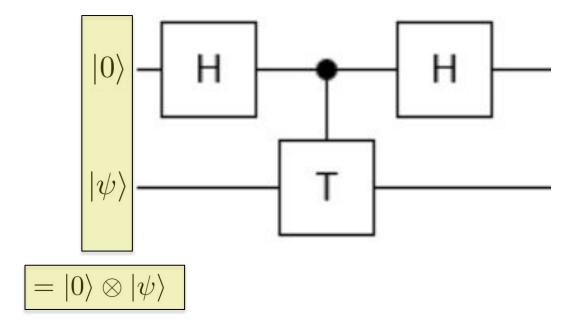
• It rotates our qubit about 45 degrees on the Bloch sphere.



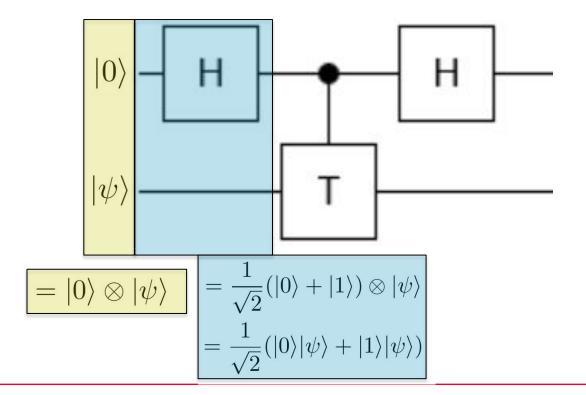
• Back to our circuit:



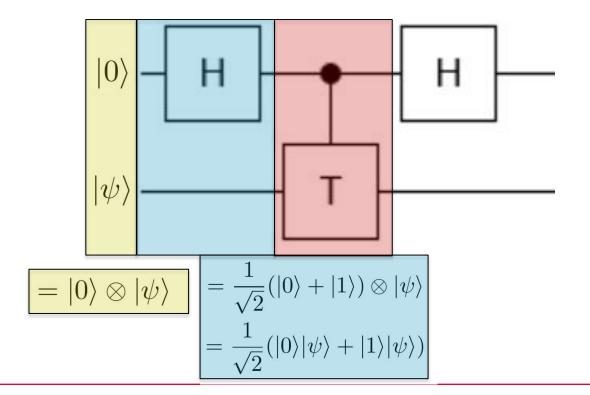
• Step 0:



• Step 1:



• Step 2:



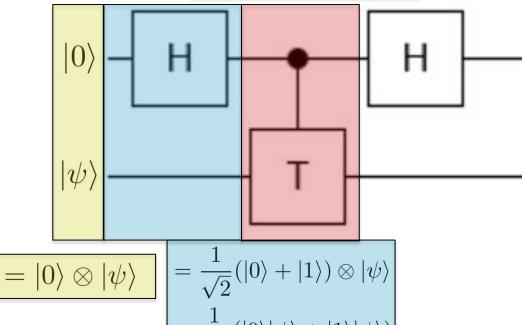


Applying phase kickback in step 2 gives us:

$$\begin{split} &= T[c,\psi] \frac{1}{\sqrt{2}} (|0\rangle|\psi\rangle + |1\rangle|\psi\rangle) \\ &= \frac{1}{\sqrt{2}} (|0\rangle|\psi\rangle + T|1\rangle|\psi\rangle) \\ &= \frac{1}{\sqrt{2}} (|0\rangle|\psi\rangle + e^{\frac{i\pi}{4}} |1\rangle|\psi\rangle) \end{split}$$

• Step 2:

$$= T[c, \psi] \frac{1}{\sqrt{2}} (|0\rangle|\psi\rangle + |1\rangle|\psi\rangle)$$
$$= \frac{1}{\sqrt{2}} (|0\rangle|\psi\rangle + e^{\frac{i\pi}{4}} |1\rangle|\psi\rangle)$$

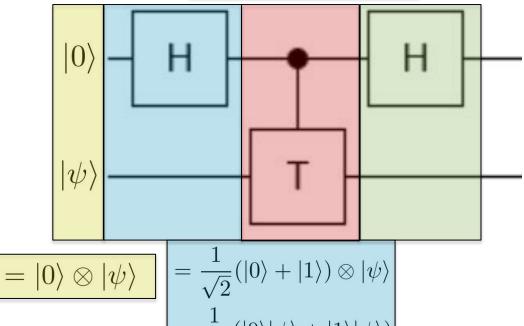


$$= \frac{1}{\sqrt{2}} (|0\rangle + |1\rangle) \otimes |\psi\rangle$$

$$= \frac{1}{\sqrt{2}} (|0\rangle |\psi\rangle + |1\rangle |\psi\rangle)$$

• Step 3:

$$= T[c, \psi] \frac{1}{\sqrt{2}} (|0\rangle|\psi\rangle + |1\rangle|\psi\rangle)$$
$$= \frac{1}{\sqrt{2}} (|0\rangle|\psi\rangle + e^{\frac{i\pi}{4}} |1\rangle|\psi\rangle)$$



$$= \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle) \otimes |\psi\rangle$$
$$= \frac{1}{\sqrt{2}}(|0\rangle|\psi\rangle + |1\rangle|\psi\rangle)$$



Finishing the Hadamard gives us the following:

Wiffig:
$$H(\frac{1}{\sqrt{2}}(|0\rangle|\psi\rangle + e^{\frac{i\pi}{4}}|1\rangle|\psi\rangle))$$

$$= \frac{1}{\sqrt{2}}(H|0\rangle|\psi\rangle + e^{\frac{i\pi}{4}}H|1\rangle|\psi\rangle)$$

$$= \frac{1}{\sqrt{2}}(\frac{1}{\sqrt{2}}(|0\rangle|\psi\rangle + |1\rangle|\psi\rangle) + e^{\frac{i\pi}{4}}\frac{1}{\sqrt{2}}(|0\rangle|\psi\rangle - |1\rangle|\psi\rangle))$$

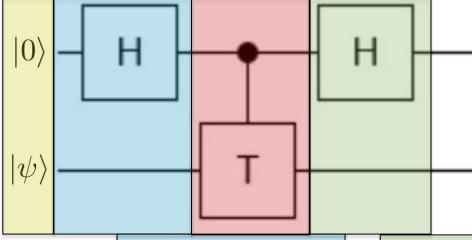
$$= \frac{1}{2}((|0\rangle|\psi\rangle + |1\rangle|\psi\rangle) + e^{\frac{i\pi}{4}}(|0\rangle|\psi\rangle - |1\rangle|\psi\rangle))$$

$$= \frac{1}{2}((|0\rangle + |1\rangle) + e^{\frac{i\pi}{4}}(|0\rangle - |1\rangle))|\psi\rangle$$

$$= \frac{1}{2}(|0\rangle(1 + e^{\frac{i\pi}{4}}) + |1\rangle(1 - e^{\frac{i\pi}{4}}))|\psi\rangle$$

• Step 3:

$$= T[c, \psi] \frac{1}{\sqrt{2}} (|0\rangle|\psi\rangle + |1\rangle|\psi\rangle)$$
$$= \frac{1}{\sqrt{2}} (|0\rangle|\psi\rangle + e^{\frac{i\pi}{4}} |1\rangle|\psi\rangle)$$



$$H(\frac{1}{\sqrt{2}}(|0\rangle|\psi\rangle + e^{\frac{i\pi}{4}}|1\rangle|\psi\rangle))$$

$$= \frac{1}{2}(|0\rangle(1 + e^{\frac{i\pi}{4}}) + |1\rangle(1 - e^{\frac{i\pi}{4}}))|\psi\rangle$$



• Let's look a little closer at the output. Our mystery function (The T-gate) rotates our qubits by $e^{\frac{\pi i}{4}}$.

$$= \frac{1}{2} (|0\rangle (1 + e^{\frac{i\pi}{4}}) + |1\rangle (1 - e^{\frac{i\pi}{4}})) |\psi\rangle$$



 We can generalize this with any gate rotation using a variable θ. In the T-gate's case, θ=1/8.

$$= \frac{1}{2} (|0\rangle (1 + e^{\frac{i\pi}{4}}) + |1\rangle (1 - e^{\frac{i\pi}{4}})) |\psi\rangle$$
$$= \frac{1}{2} (|0\rangle (1 + e^{2\pi i\theta}) + |1\rangle (1 - e^{2\pi i\theta})) |\psi\rangle$$



 As you can see here, we have skewed the probability of measuring 1 or 0 a certain way depending on our angle.

$$Prob.|0\rangle = \|\frac{1}{2}(|0\rangle(1 + e^{2\pi i\theta})\|$$

 $Prob.|1\rangle = \|\frac{1}{2}(|1\rangle(1 - e^{2\pi i\theta})\|$



Obviously, we are looking for θ, not for 1 or 0.
 Whatever we are doing, 1 qubit of calculations is not precise enough.

$$Prob.|0\rangle = \|\frac{1}{2}(|0\rangle(1 + e^{2\pi i\theta})\|$$

 $Prob.|1\rangle = \|\frac{1}{2}(|1\rangle(1 - e^{2\pi i\theta})\|$



 We can further skew our output by repeating this phase kickback. Will this give us θ?

$$Prob.|0\rangle = \|\frac{1}{2}(|0\rangle(1 + e^{2\pi i\theta})\|$$

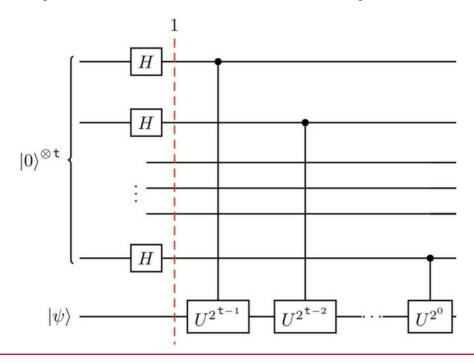
 $Prob.|1\rangle = \|\frac{1}{2}(|0\rangle(1 - e^{2\pi i\theta})\|$



Implementing a N-bit QPE

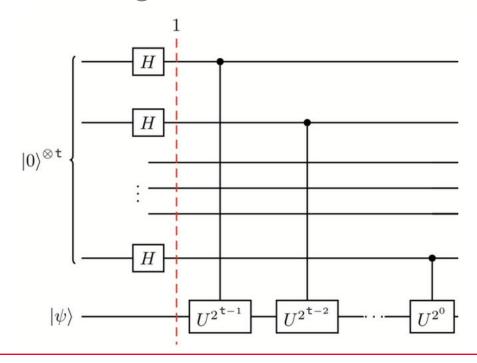


 Let's generalize our previous circuit to repeat over n-qubits instead of 1-qubit.

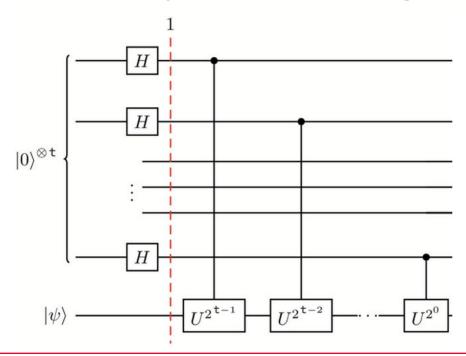




 The pattern seems simple enough, we simply repeat the U-gate for each control bit.

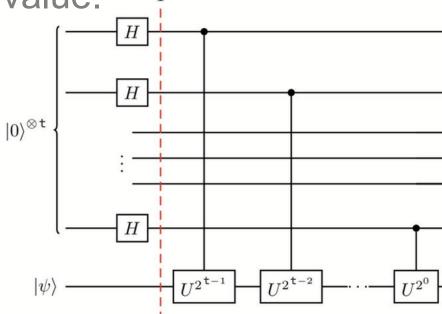


• However, what are these 2^{t-1}, 2^{t-2}, etc. values and how do they effect our U-gate?



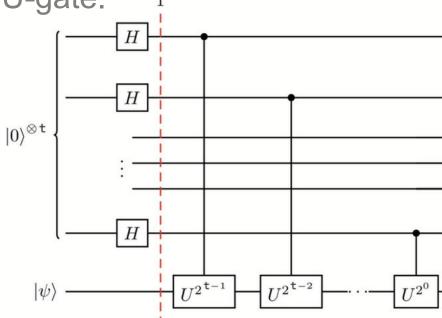


 For this notation, t is the length of our Hadamard-string. The larger the t, the more accurate our value.

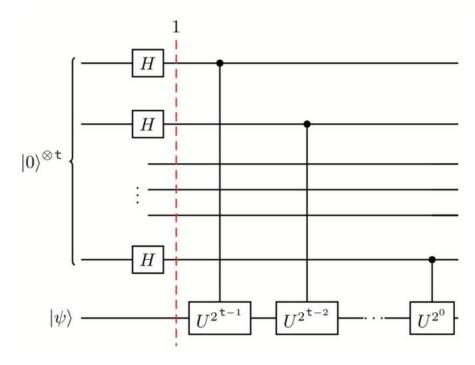


- U^{2^(t-1)} means U*U*U...U repeated 2^(t-1)-times.
- U^{2^(t-2)} is 1 less repetitions of U.

U^{2^0} is a single U-gate.

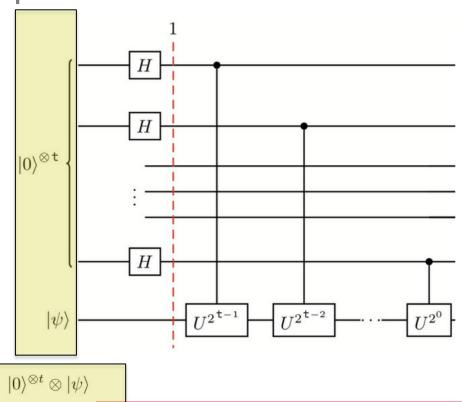


So, what does this do?

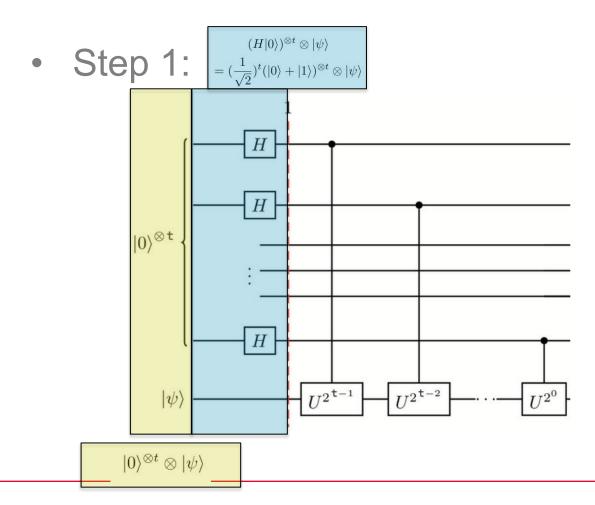




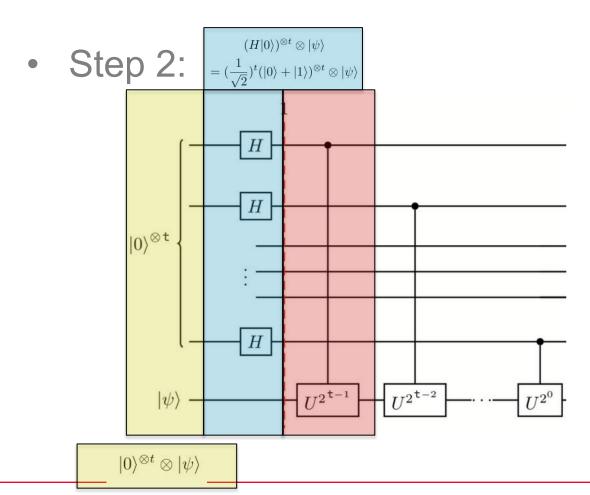
• Step 0:













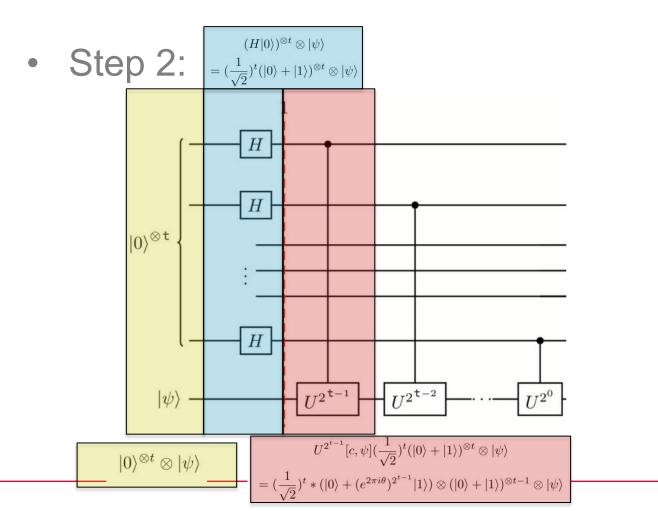
Applying phase kickback in step 2 gives us:

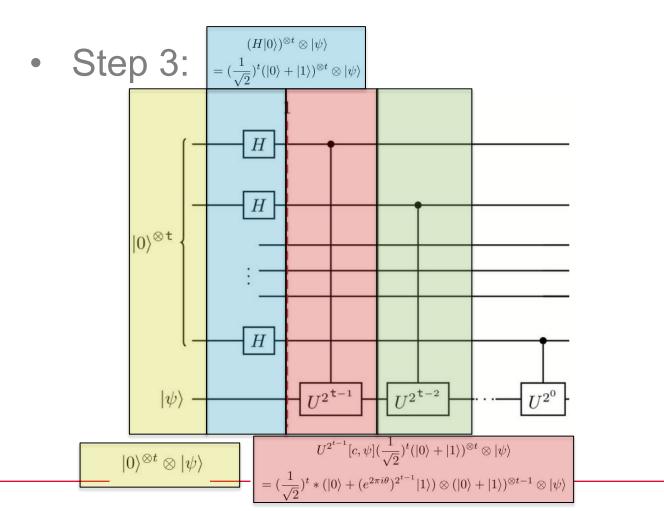
$$U^{2^{t-1}}[c,\psi](\frac{1}{\sqrt{2}})^t(|0\rangle + |1\rangle)^{\otimes t} \otimes |\psi\rangle$$

$$= (\frac{1}{\sqrt{2}})^t * (|0\rangle + U^{2^{t-1}}|1\rangle) \otimes (|0\rangle + |1\rangle)^{\otimes t-1} \otimes |\psi\rangle$$

• Which if we substitute U with $e^{2\pi i\theta}$ gives us:

$$= \left(\frac{1}{\sqrt{2}}\right)^t * (|0\rangle + (e^{2\pi i\theta})^{2^{t-1}}|1\rangle) \otimes (|0\rangle + |1\rangle)^{\otimes t-1} \otimes |\psi\rangle$$

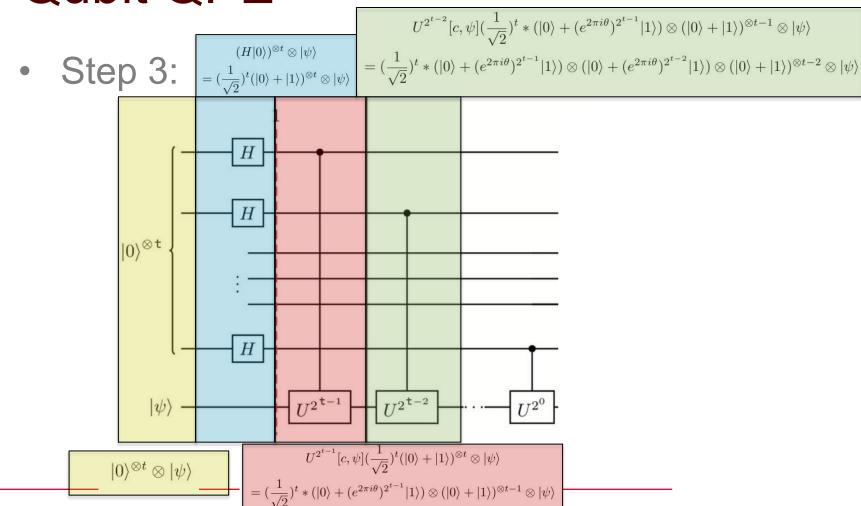


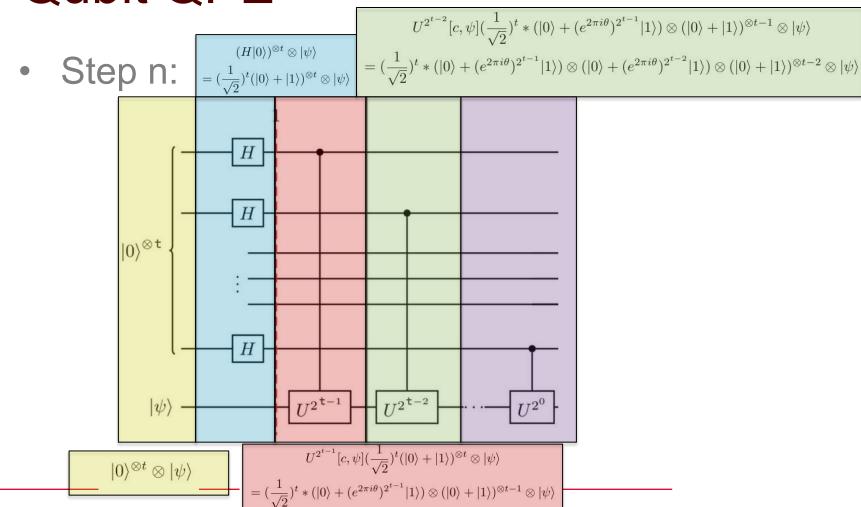




 Because we are just repeating step 2 but with the second qubit and one less U gate, the logic done before can be applied again:

$$\begin{split} U^{2^{t-2}}[c,\psi](\frac{1}{\sqrt{2}})^t * (|0\rangle + (e^{2\pi i\theta})^{2^{t-1}}|1\rangle) \otimes (|0\rangle + |1\rangle)^{\otimes t-1} \otimes |\psi\rangle \\ &= (\frac{1}{\sqrt{2}})^t * (|0\rangle + U^{2^{t-1}}|1\rangle) \otimes (|0\rangle + U^{2^{t-2}}|1\rangle) \otimes (|0\rangle + |1\rangle)^{\otimes t-2} \otimes |\psi\rangle \\ &= (\frac{1}{\sqrt{2}})^t * (|0\rangle + (e^{2\pi i\theta})^{2^{t-1}}|1\rangle) \otimes (|0\rangle + (e^{2\pi i\theta})^{2^{t-2}}|1\rangle) \otimes (|0\rangle + |1\rangle)^{\otimes t-2} \otimes |\psi\rangle \end{split}$$





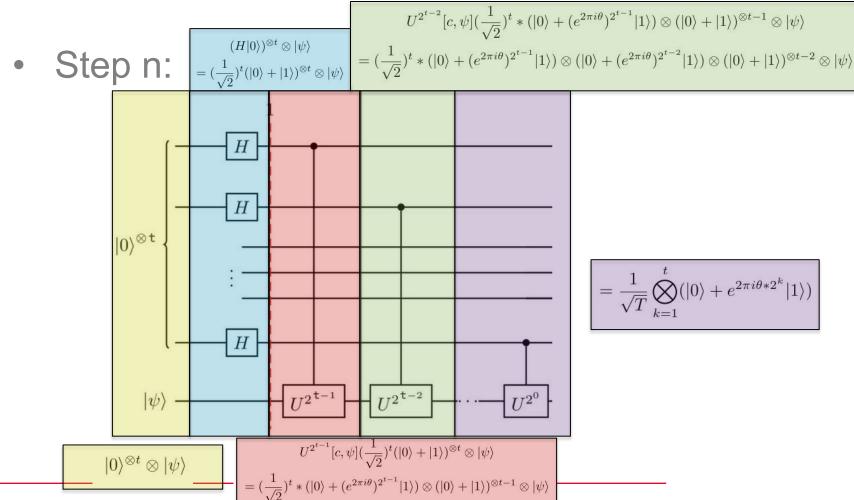


 Repeating our steps over and over gives us the following formula:

$$\implies \left(\frac{1}{\sqrt{2}}\right)^t * (|0\rangle + (e^{2\pi i\theta})^{2^{t-1}}|1\rangle) \otimes (|0\rangle + (e^{2\pi i\theta})^{2^{t-2}}|1\rangle) \otimes \dots \otimes (|0\rangle + (e^{2\pi i\theta})^{2^0}|1\rangle)$$

$$= \frac{1}{\sqrt{T}} \bigotimes_{k=1}^t (|0\rangle + e^{2\pi i\theta * 2^k}|1\rangle)$$

• The first control qubit is modified by the (t-1) gate, the second qubit is modified by the (t-2) gate, etc.







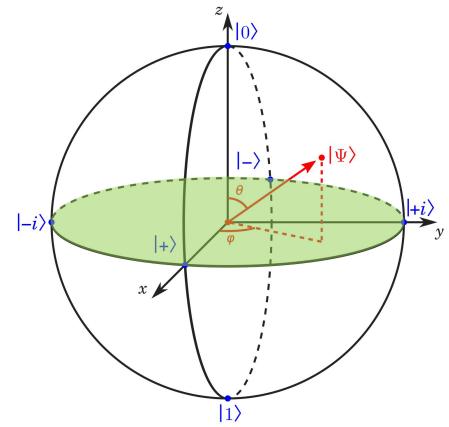
 We simply repeated our naïve 1-qubit algorithm iteratively until we reached our nth qubit.



- As such, we shifted our probabilities of measuring 1 and 0 depending on our angle θ.
- This shifted is more and more precise the more qubits we use.

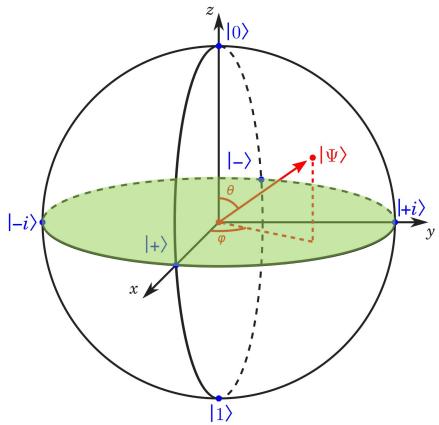


- Because of the precision, each value of θ should have a unique probability.
- To reword this, each value of θ should have a unique coordinate on the Bloch Sphere's x/y axis.





Sound Familiar?



Turns out, the formula is almost the exact same as the QFT, with the difference being 1/2ⁿ.

$$\frac{1}{\sqrt{T}} \bigotimes_{k=1}^{t} (|0\rangle + e^{2\pi i \theta * 2^k} |1\rangle)$$
$$\frac{1}{\sqrt{N}} \bigotimes_{k=1}^{n} (|0\rangle + e^{\frac{2\pi i x_{b10}}{2^k}} |1\rangle)$$



$$= \frac{1}{\sqrt{T}} \bigotimes_{k=1}^{t} (|0\rangle + e^{2\pi i\theta * 2^{k}} |1\rangle)$$

$$\Longrightarrow \frac{1}{\sqrt{T}} \bigotimes_{k=1}^{t} (|0\rangle + e^{\frac{1}{2^{n}}} * e^{2\pi i\theta * 2^{k}} |1\rangle)$$

$$= \frac{1}{\sqrt{T}} \bigotimes_{k=1}^{t} (|0\rangle + e^{\frac{1}{2^{n}}} * e^{2\pi i\theta * 2^{k}} |1\rangle)$$

$$= \frac{1}{\sqrt{T}} \bigotimes_{k=1}^{t} (|0\rangle + e^{\frac{1}{2^{n}} * 2^{k}} |1\rangle)$$

$$= \frac{1}{\sqrt{T}} \bigotimes_{k=1}^{t} (|0\rangle + e^{\frac{2\pi i\theta * 2^{k}}{2^{n}}} |1\rangle)$$

$$= \frac{1}{\sqrt{T}} \bigotimes_{k=1}^{t} (|0\rangle + e^{\frac{2\pi i\theta * 2^{k}}{2^{k}}} |1\rangle)$$

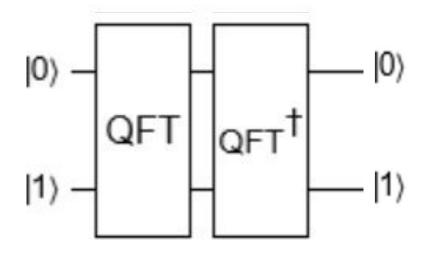


When you multiply out the constant, it is the same formula, with some variables renamed.

$$= \frac{1}{\sqrt{T}} \bigotimes_{k=1}^{t} (|0\rangle + e^{2\pi i \theta * 2^{k}} |1\rangle)$$
$$= \frac{1}{\sqrt{N}} \bigotimes_{k=1}^{n} (|0\rangle + e^{\frac{2\pi i x_{b10}}{2^{k}}} |1\rangle)$$

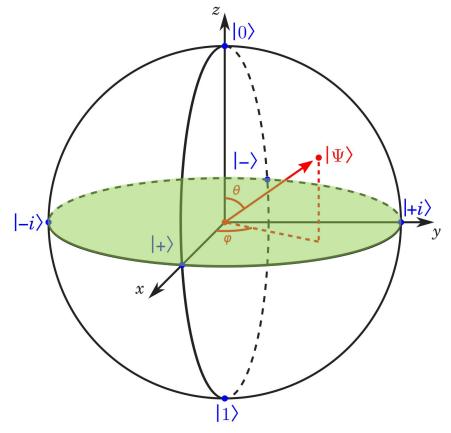


- Finally, we can remember that every function is unitary in Quantum Computing.
- This means that for every function, there is an inverse function.

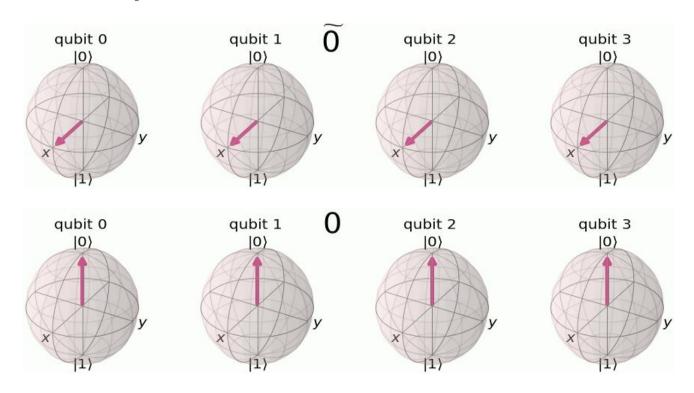




- If a QFT takes a number and turns it into a rotation angle along the Bloch Sphere...
- Then the Inverse QFT takes that rotation and gets us our number.



The input of a QFT is our unknown θ!





 ...Except our QPE is equal to the QFT function, multiplied by a constant 2ⁿ. which means we must divide by 2ⁿ to get our true value for θ.

$$QPE = \frac{1}{\sqrt{T}} \bigotimes_{k=1}^{t} (|0\rangle + e^{2\pi i \theta * 2^k} |1\rangle)$$

$$QFT = \frac{1}{\sqrt{N}} \bigotimes_{k=1}^{n} (|0\rangle + e^{\frac{2\pi i x_{b10}}{2^k}} |1\rangle)$$

• In the base case of 1-qubit, we only skewed our algorithm slightly. This resulted in an output with low precision and only gave us 0 or $\frac{1}{2}$ as our estimation for θ . With value closest to the true answer having the highest probability.

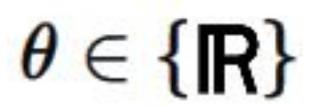
$$egin{aligned} heta \in egin{cases} 0 & a=0 \ rac{1}{2} & a=1. \end{cases}$$

 In the case of a 2-qubit QPE, we skewed our algorithm slightly more. This resulted in an output with higher precision and gave us more potential answers. With value closest to the true answer having the highest probability.

$$\theta \in \{0, 1/4, 1/2, 3/4\}$$



 Naturally, as n increases, so does our decimal range. With enough precision, we can get our desired angle with any number of decimal places.

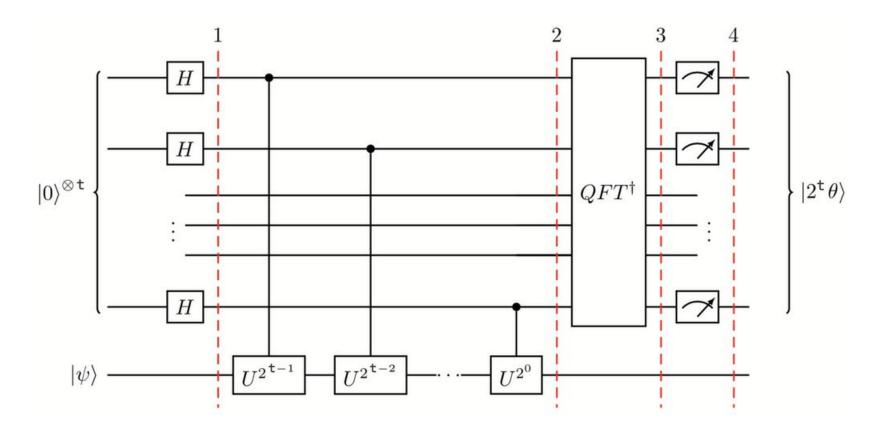




The Final Circuit



The Final Circuit





Lets Get Started