# HO CHI MINH UNIVERSITY OF TECHNOLOGY FACULTY OF COMPUTER SCIENCE AND ENGINEERING



# SIMPLE OPERATING SYSTEM IMPLEMENTATION WITH PROCESS SCHEDULING AND MEMORY MANAGEMENT SIMULATION

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# I. WORKLOAD

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#### II. QUESTION

• QUESTION 1: What is the advantage of using priority queue in comparison with other scheduling algorithms you have learned?

The relative importance of each process can be precisely defined beforehand allowing the cpu to choose processes based on their importance and it is easier to manage the queue while processing.

• QUESTION 2: In this simple OS, we implement a design of multiple memory segments or memory areas in source code declaration. What is the advantage of the proposed design of multiple segments?

The multiple memory segments design allows many processes to share data easily, different areas can be used for stack, code, heap or data, easier to manage the size of each area like stack or heap as we can manage how many segments will be used in each area. It extends the address ability of a processor, for example, giving an addressing capability of 1 MB, with segmentation it only requires 16 bit registers, otherwise it needs 20 bit registers. It also allows processes to work independently of each other.

• QUESTION 3: What will happen if we divide the address to more than 2-levels in the paging memory management system?

This approach is known as multi-level paging or hierarchical paging, where the page number is divided into multiple levels, and each level contains a smaller number of entries. Each level of the page table points to the next level of the page table until the final level points to the physical page frame, reducing the size of the page table.

 QUESTION 4: What is the advantage and disadvantage of segmentation with paging?

Advantage: It does not cause external fragmentation. It's also easier to manage memory allocation and the algorithm is simpler

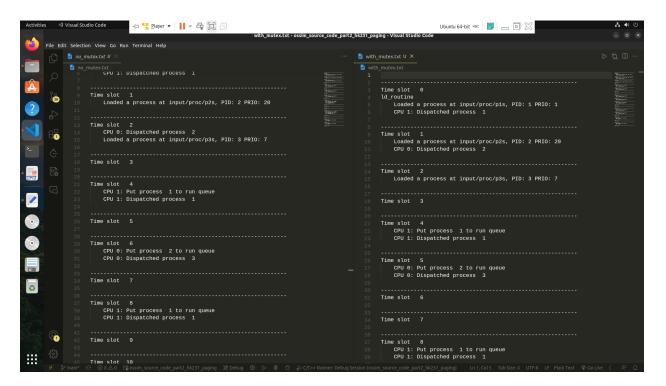
Disadvantage: It can cause internal fragmentation, require more hardware resources, may increase memory access time, leading to increasing complexity. It also require memory to store the page table

• QUESTION 5: What will happen if the synchronization is not handled in your simple OS? Illustrate by example the problem of your simple OS if you have any.

The major problem is race condition. Normally, without synchronization, when both cpus are running at the same time, a single variable which stores the remaining time left of

each process will be accessed by both cpus, leading to race condition. Thankfully, in our code we store each remaining time of different processes in another array outside of the process, avoiding race condition entirely, so even if the synchronization is not handled in our simple OS it would still work.

Running input "sched" with 2 cpus, without mutex and with mutex, respectively:



### III. RESULTS

# 1. Scheduling:

Take input os\_sched\_0 as an example, we have the following time slot table:



## 2. Memory management:

At the end of each time slot, we add a line title RAM to show the status of the RAM after each time slot, only meaningful value will be shown (value different from 0).

```
Time slot 0

ld_routine

Loaded a process at input/proc/p0s, PID: 1 PRIO: 0

RAM:

CPU 0: Dispatched process 1

-
Time slot 1
```

```
Time slot 2
   Loaded a process at input/proc/p1s, PID: 2 PRIO: 15
RAM:
Time slot 3
   CPU 1: Dispatched process 2
   Loaded a process at input/proc/p1s, PID: 3 PRIO: 0
RAM:
Time slot 4
   Loaded a process at input/proc/p1s, PID: 4 PRIO: 0
RAM:
Time slot 5
write region=1 offset=20 value=100
print_pgtbl: 0 - 1280
00000000: 80000001
00000004: 80000001
00000008: 80000003
00000012: 80000003
00000016: 80000004
Time slot 6
   CPU 0: Put process 1 to run queue
   CPU 0: Dispatched process 3
Time slot 7
Time slot 8
```

```
Time slot 9
   CPU 1: Put process 2 to run queue
   CPU 1: Dispatched process 4
Time slot 10
Time slot 11
Time slot 12
   CPU 0: Put process 3 to run queue
   CPU 0: Dispatched process 1
read region=1 offset=20 value=100
print_pgtbl: 0 - 1280
00000000: 80000001
00000004: 80000001
00000012: 80000003
00000016: 80000004
RAM content: 100
Time slot 13
write region=2 offset=20 value=102
print_pgtbl: 0 - 1280
00000000: 80000001
00000004: 80000001
00000008: 80000003
00000012: 80000003
00000016: 80000004
RAM content: 100
Time slot 14
read region=2 offset=20 value=102
print_pgtbl: 0 - 1280
00000000: 80000001
```

```
00000004: 80000001
00000008: 80000003
00000012: 80000003
RAM content: 102
Time slot 15
   CPU 1: Put process 4 to run queue
   CPU 1: Dispatched process 3
write region=3 offset=20 value=103
print_pgtbl: 0 - 1280
00000000: 80000001
00000004: 80000001
00000008: 80000003
00000012: 80000003
00000016: 80000004
RAM content: 102
Time slot 16
   CPU 0: Processed 1 has finished
   CPU 0: Dispatched process 4
Time slot 17
Time slot 18
Time slot 19
   CPU 1: Processed 3 has finished
   CPU 1: Dispatched process 2
Time slot 20
   CPU 0: Processed 4 has finished
  CPU 0 stopped
```

```
Time slot 21

Time slot 22

Time slot 23

CPU 1: Processed 2 has finished
CPU 1 stopped

Time slot 24
```

#### 3. Result

Scheduling functionality test by with input "sched\_0"

```
Time slot 0

ld_routine

Loaded a process at input/proc/s0, PID: 1 PRIO: 1

CPU 1: Dispatched process 1

Time slot 1

Time slot 2
```

```
CPU 1: Put process 1 to run queue
     CPU 1: Dispatched process 1
Time slot 3
     Loaded a process at input/proc/s1, PID: 2 PRIO: 3
     CPU 0: Dispatched process 2
     CPU 1: Put process 1 to run queue
     CPU 1: Dispatched process 1
Time slot 5
Time slot 6
     CPU 0: Put process 2 to run queue
     CPU 0: Dispatched process 2
     CPU 1: Put process 1 to run queue
     CPU 1: Dispatched process 1
Time slot 7
Time slot 8
     CPU 0: Put process 2 to run queue
     CPU 0: Dispatched process 2
     CPU 1: Put process 1 to run queue
     CPU 1: Dispatched process 1
Time slot 9
Time slot 10
     CPU 0: Put process 2 to run queue
     CPU 1: Put process 1 to run queue
     CPU 1: Dispatched process 1
     CPU 0: Dispatched process 2
```

```
Time slot 11

CPU 0: Processed 2 has finished
CPU 0 stopped

Time slot 12

CPU 1: Put process 1 to run queue
CPU 1: Dispatched process 1

Time slot 13

Time slot 14

CPU 1: Put process 1 to run queue
CPU 1: Dispatched process 1

Time slot 15

CPU 1: Processed 1 has finished
CPU 1 stopped

Time slot 16
```

#### Paging

Paging functionality test result with input "os\_0\_mlq\_paging"

```
Time slot 2
     Loaded a process at input/proc/p1s, PID: 2 PRIO: 15
Time slot 3
     CPU 1: Dispatched process 2
     Loaded a process at input/proc/p1s, PID: 3 PRIO: 0
Time slot 4
     Loaded a process at input/proc/p1s, PID: 4 PRIO: 0
Time slot 5
write region=1 offset=20 value=100
print_pgtbl: 0 - 1280
00000000: 80000001
00000008: 80000003
00000012: 80000003
00000016: 80000004
Time slot 6
     CPU 0: Put process 1 to run queue
     CPU 0: Dispatched process 3
Time slot 7
Time slot 8
```

```
Time slot 9
     CPU 1: Put process 2 to run queue
     CPU 1: Dispatched process 4
Time slot 10
Time slot 11
Time slot 12
     CPU 0: Put process 3 to run queue
     CPU 0: Dispatched process 1
read region=1 offset=20 value=100
print_pgtbl: 0 - 1280
00000000: 80000001
00000004: 80000001
00000008: 80000003
00000012: 80000003
00000016: 80000004
Time slot 13
write region=2 offset=20 value=102
print_pgtbl: 0 - 1280
00000000: 80000001
00000004: 80000001
00000012: 80000003
00000016: 80000004
Time slot 14
read region=2 offset=20 value=102
print_pgtbl: 0 - 1280
00000000: 80000001
00000004: 80000001
```

```
00000008: 80000003
00000012: 80000003
00000016: 80000004
Time slot 15
write region=3 offset=20 value=103
print pgtbl: 0 - 1280
00000000: 80000001
00000004: 80000001
00000008: 80000003
00000012: 80000003
00000016: 80000004
     CPU 1: Put process 4 to run queue
     CPU 1: Dispatched process 3
Time slot 16
     CPU 0: Processed 1 has finished
     CPU 0: Dispatched process 4
Time slot 17
Time slot 18
Time slot 19
     CPU 1: Processed 3 has finished
     CPU 1: Dispatched process 2
Time slot 20
     CPU 0: Processed 4 has finished
     CPU 0 stopped
```

```
Time slot 21

Time slot 22

Time slot 23

CPU 1: Processed 2 has finished
CPU 1 stopped

Time slot 24
```

#### IV. REFERENCES

- 1. "Operating System Concepts and Design" by Milan Milenkovic
- 2. "operating-system-concepts-10th-Copy" by Peter Baer Galvin