

Pumping lemma for CFL's - Examples

Pumping lemma (Theorem 2.34): If A is a CFL, then there is a number p (pumping length) where if s is any string in A , $|s| \geq p$, then there exists a partition $s = uvxyz$, satisfying

1. For each $i \geq 0$, $uv^i x y^i z \in A$.
2. $|v| > 0$ (either $v \neq \epsilon$ or $y \neq \epsilon$)
3. $|vxy| \leq p$.

uxz
 $uvxyz$
 $uvvxyz$
 $uvvvxyz$
 \vdots

$\{\epsilon, abc, uabbbcc, a^3b^3c^3, \dots\}$

Example 2.36: $B = \{a^n b^n c^n \mid n \geq 0\} \subseteq a^* b^* c^*$

Show that B is not context-free.

Assume that B is context-free. Then by pumping lemma, there exists a pumping length $p \geq 0$. Let $s = a^p b^p c^p$. Suppose s can be written as $s = uvxyz$.

$aa \dots abbb \dots bbbcc \dots cc$

$\leftarrow v \rightarrow$

$aaa \underline{abb} \underline{abb} bb \dots$

$$\underline{uv^2xy^2z}$$

There are two cases.

1) Either v or y contains two types of symbols.
Then uv^2xy^2z is not of the form $a^*b^*c^*$.

2) If v and y each contain only one type of symbol, then uv^2xy^2z does not contain the same no. of a 's, b 's and c 's.

Hence B is not a CFL.

Example 2.37: $C = \{a^i b^j c^k \mid 0 \leq i \leq j \leq k\} \subseteq a^*b^*c^*$

Like before, we assume that C is context-free.

Pumping lemma implies the existence of a pumping length p . Let $s = a^p b^p c^p$.

$aa \dots a \quad ab \quad b \dots b \quad bcc \dots cc$

$\mid \leftarrow v \rightarrow \mid$

$$uv^2xy^2z = aa \dots a \quad bbaabb \dots$$

1. If v or y contains two different types of symbols, then uv^2xy^2z is not of the type $a^*b^*c^*$.

So v and y each contain only one type of symbol.

2. If vy avoids a , then consider $uv^0x y^0z$.

Since $|vy| \neq 0$, then the count of b 's or c 's must go down.

If vy avoids b , but contains a , then $uv^2x y^2z$ has more a 's than b 's.

If vy avoids b , but contains c , then $uv^0x y^0z$ has more b 's than c 's.

If vy avoids c , then $uv^2x y^2z$ contains more a 's than c 's or more b 's than c 's.

Hence C is not context-free.

Exercise : $D = \{w\#w \mid w \in \{0,1\}^*\}$. $\begin{matrix} 01\#01 \\ 111\#111 \end{matrix}$

Show that D is not context-free.

(Similar to Example 2.38: $\{ww \mid w \in \{0,1\}^*\}$.)