

Contents

| | |
|-----------------|----|
| 一切的开始 | 3 |
| 宏定义 | 3 |
| 快速读 | 3 |
| 对拍 | 4 |
| 为什么 C++ 不自带这个? | 5 |
| 数据结构 | 5 |
| ST 表 | 5 |
| 线段树 | 5 |
| 均摊复杂度线段树 | 7 |
| 持久化线段树 | 8 |
| K-D Tree | 9 |
| 树状数组 | 11 |
| 主席树 | 13 |
| 左偏树 | 15 |
| Treap | 16 |
| Treap-序列 | 18 |
| 可回滚并查集 | 21 |
| 舞蹈链 | 22 |
| CDQ 分治 | 24 |
| 哈希表 | 25 |
| 笛卡尔树 | 26 |
| Trie | 26 |
| pb_ds | 27 |
| Link-Cut Tree | 28 |
| 莫队 | 29 |
| 数学 | 31 |
| 矩阵运算 | 31 |
| 筛 | 31 |
| 亚线性筛 | 32 |
| min_25 | 32 |
| 杜教筛 | 33 |
| 素数测试 | 34 |
| 线性递推 | 34 |
| 扩展欧几里得 | 35 |
| 类欧几里得 | 36 |
| 逆元 | 36 |
| 组合数 | 36 |
| 第二类斯特灵数 | 37 |
| FFT & NTT & FWT | 37 |
| simpson 自适应积分 | 38 |
| 快速乘 | 39 |
| 快速幂 | 39 |
| 高斯消元 | 40 |
| 质因数分解 | 41 |
| 原根 | 42 |
| 公式 | 42 |
| 二次剩余 | 43 |
| 中国剩余定理 | 44 |
| 伯努利数和等幂求和 | 45 |
| 单纯形 | 45 |
| 图论 | 46 |
| LCA | 46 |
| 最短路 | 47 |

| | |
|--------------------------|-----------|
| 网络流 | 47 |
| 树上路径交 | 50 |
| 树上点分治 | 50 |
| 树链剖分 | 53 |
| 二分图匹配 | 57 |
| 虚树 | 59 |
| 欧拉路径 | 59 |
| 强连通分量与 2-SAT | 59 |
| 拓补排序 | 60 |
| 一般图匹配 (没有测试过) | 61 |
| 计算几何 | 62 |
| 圆的反演 | 62 |
| 二维 | 63 |
| 旋转卡壳 | 66 |
| 没有测试过的 | 67 |
| 半平面交 | 68 |
| 字符串 | 69 |
| 后缀自动机 | 69 |
| 回文自动机 | 72 |
| manacher | 72 |
| 哈希 | 73 |
| 后缀数组 | 74 |
| KMP 自动机 | 77 |
| Trie | 79 |
| AC 自动机 | 79 |
| 杂项 | 80 |
| STL | 80 |
| 日期 | 81 |
| 子集枚举 | 81 |
| 权值最大上升子序列 | 82 |
| 数位 DP | 82 |
| 土制 bitset | 82 |
| 随机 | 83 |
| 伪随机数 | 83 |
| Java | 84 |
| Regex | 84 |
| Decimal Format | 84 |
| Sort | 86 |
| 扩栈 (本地使用) | 86 |
| 心态崩了 | 86 |

一切的开始

宏定义

- 需要 C++11

```
1 #include <bits/stdc++.h>
2 using namespace std;
3 using LL = long long;
4 #define FOR(i, x, y) for (decay<decltype(y)>::type i = (x), _##i = (y); i < _##i; ++i)
5 #define FORD(i, x, y) for (decay<decltype(x)>::type i = (x), _##i = (y); i > _##i; --i)
6 #ifdef zero1
7 #define dbg(args...) do { cout << "\033[32;1m" << #args << " -> "; err(args); } while (0)
8 #else
9 #define dbg(...)
10 #endif
11 void err() { cout << "\033[39;0m" << endl; }
12 template<typename...> class T, typename t, typename... Args>
13 void err(T<t> a, Args... args) { for (auto x: a) cout << x << ' '; err(args...); }
14 template<typename T, typename... Args>
15 void err(T a, Args... args) { cout << a << ' '; err(args...); }
16 // -----
```

- 更多配色:

- 33 黄色
- 34 蓝色
- 31 橙色

- POJ/BZOJ version

```
1 #include <cstdio>
2 #include <iostream>
3 #include <algorithm>
4 #include <cmath>
5 #include <string>
6 #include <vector>
7 #include <set>
8 #include <queue>
9 #include <cstring>
10 #include <cassert>
11 using namespace std;
12 typedef long long LL;
13 #define FOR(i, x, y) for (LL i = (x), _##i = (y); i < _##i; ++i)
14 #define FORD(i, x, y) for (LL i = (x), _##i = (y); i > _##i; --i)
15 #ifdef zero1
16 #define dbg(args...) do { cout << "\033[32;1m" << #args << " -> "; err(args); } while (0)
17 #else
18 #define dbg(...)
19 #endif
20 void err() { cout << "\033[39;0m" << endl; }
21 template<typename T, typename... Args>
22 void err(T a, Args... args) {
23     cout << a << ' '; err(args...);
24 }
25 // -----
```

- HDU Assert Patch

```
1 #ifdef ONLINE_JUDGE
2 #define assert(condition) if (!(condition)) { int x = 1, y = 0; cout << x / y << endl; }
3 #endif
```

快速读

```
1 inline char next_char() {
2     static char buf[100000], *p1 = buf, *p2 = buf;
3     return p1 == p2 && (p2 = (p1 = buf) + fread(buf, 1, 100000, stdin), p1 == p2) ? EOF : *p1++;
4 }
5 inline bool maybe_digit(char c) {
```

```

6     return c >= '0' && c <= '9';
7 }
8 template <typename T>
9 void rn(T& _v) {
10     static char ch;
11     static bool negative = false;
12     _v = 0;
13     while (!maybe_digit(ch)) {
14         negative = ch == '-';
15         ch = next_char();
16     }
17     do _v = (_v << 1) + (_v << 3) + ch - '0';
18     while (maybe_digit(ch = next_char()));
19     if (negative) _v = -_v;
20 }
21
22 template <typename T>
23 void o(T p) {
24     static int stk[70], tp;
25     if (p == 0) {
26         putchar('0');
27         return;
28     }
29     if (p < 0) { p = -p; putchar('-'); }
30     while (p) stk[++tp] = p % 10, p /= 10;
31     while (tp) putchar(stk[tp--] + '0');
32 }

```

- 需要初始化
- 需要一次读入
- 不支持负数

```

1 const int MAXS = 100 * 1024 * 1024;
2 char buf[MAXS];
3 template<typename T>
4 inline bool read(T& x) {
5     static char* p = buf;
6     x = 0;
7     while (*p && !isdigit(*p)) ++p;
8     if (!*p) return false;
9     while (isdigit(*p)) x = x * 10 + *p++ - 48;
10    return true;
11 }
12
13 fread(buf, 1, MAXS, stdin);

```

对拍

```

1 #!/usr/bin/env bash
2 g++ -o r main.cpp -O2 -std=c++11
3 g++ -o std std.cpp -O2 -std=c++11
4 while true; do
5     python gen.py > in
6     ./std < in > stdout
7     ./r < in > out
8     if test $? -ne 0; then
9         exit 0
10    fi
11    if diff stdout out; then
12        printf "AC\n"
13    else
14        printf "GG\n"
15        exit 0
16    fi
17 done

```

为什么 C++ 不自带这个？

```
1 LL bin(LL x, LL n, LL MOD) {
2     LL ret = MOD != 1;
3     for (x %= MOD; n; n >>= 1, x = x * x % MOD)
4         if (n & 1) ret = ret * x % MOD;
5     return ret;
6 }
7 inline LL get_inv(LL x, LL p) { return bin(x, p - 2, p); }
```

数据结构

ST 表

● 二维

```
1 int f[maxn][maxn][10][10];
2 inline int highbit(int x) { return 31 - __builtin_clz(x); }
3 inline int calc(int x, int y, int xx, int yy, int p, int q) {
4     return max(
5         max(f[x][y][p][q], f[xx - (1 << p) + 1][yy - (1 << q) + 1][p][q]),
6         max(f[xx - (1 << p) + 1][y][p][q], f[x][yy - (1 << q) + 1][p][q])
7     );
8 }
9 void init() {
10     FOR (x, 0, highbit(n) + 1)
11     FOR (y, 0, highbit(m) + 1)
12     FOR (i, 0, n - (1 << x) + 1)
13     FOR (j, 0, m - (1 << y) + 1) {
14         if (!x && !y) { f[i][j][x][y] = a[i][j]; continue; }
15         f[i][j][x][y] = calc(
16             i, j,
17             i + (1 << x) - 1, j + (1 << y) - 1,
18             max(x - 1, 0), max(y - 1, 0)
19         );
20     }
21 }
22 inline int get_max(int x, int y, int xx, int yy) {
23     return calc(x, y, xx, yy, highbit(xx - x + 1), highbit(yy - y + 1));
24 }
```

● 一维

```
1 struct RMQ {
2     int f[22][M];
3     inline int highbit(int x) { return 31 - __builtin_clz(x); }
4     void init(int* v, int n) {
5         FOR (i, 0, n) f[0][i] = v[i];
6         FOR (x, 1, highbit(n) + 1)
7             FOR (i, 0, n - (1 << x) + 1)
8                 f[x][i] = min(f[x - 1][i], f[x - 1][i + (1 << (x - 1))]);
9     }
10    int get_min(int l, int r) {
11        assert(l <= r);
12        int t = highbit(r - l + 1);
13        return min(f[t][l], f[t][r - (1 << t) + 1]);
14    }
15 } rmq;
```

线段树

● 普适

```
1 namespace sg {
2     struct Q {
3         LL setv;
```

```

4         explicit Q(LL setv = -1): setv(setv) {}
5         void operator += (const Q& q) { if (q.setv != -1) setv = q.setv; }
6     };
7     struct P {
8         LL min;
9         explicit P(LL min = INF): min(min) {}
10        void up(Q& q) { if (q.setv != -1) min = q.setv; }
11    };
12    template<typename T>
13    P operator & (T&& a, T&& b) {
14        return P(min(a.min, b.min));
15    }
16    P p[maxn << 2];
17    Q q[maxn << 2];
18    #define lson o * 2, l, (l + r) / 2
19    #define rson o * 2 + 1, (l + r) / 2 + 1, r
20    void up(int o, int l, int r) {
21        if (l == r) p[o] = P();
22        else p[o] = p[o * 2] & p[o * 2 + 1];
23        p[o].up(q[o]);
24    }
25    void down(int o, int l, int r) {
26        q[o * 2] += q[o]; q[o * 2 + 1] += q[o];
27        q[o] = Q();
28        up(lson); up(rson);
29    }
30    template<typename T>
31    void build(T&& f, int o = 1, int l = 1, int r = n) {
32        if (l == r) q[o] = f(l);
33        else { build(f, lson); build(f, rson); q[o] = Q(); }
34        up(o, l, r);
35    }
36    P query(int ql, int qr, int o = 1, int l = 1, int r = n) {
37        if (ql > r || l > qr) return P();
38        if (ql <= l && r <= qr) return p[o];
39        down(o, l, r);
40        return query(ql, qr, lson) & query(ql, qr, rson);
41    }
42    void update(int ql, int qr, const Q& v, int o = 1, int l = 1, int r = n) {
43        if (ql > r || l > qr) return;
44        if (ql <= l && r <= qr) q[o] += v;
45        else {
46            down(o, l, r);
47            update(ql, qr, v, lson); update(ql, qr, v, rson);
48        }
49        up(o, l, r);
50    }
51 }

```

• SET+ADD

```

1     struct IntervalTree {
2         #define ls o * 2, l, m
3         #define rs o * 2 + 1, m + 1, r
4         static const LL M = maxn * 4, RS = 1E18 - 1;
5         LL addv[M], setv[M], minv[M], maxv[M], sumv[M];
6         void init() {
7             memset(addv, 0, sizeof addv);
8             fill(setv, setv + M, RS);
9             memset(minv, 0, sizeof minv);
10            memset(maxv, 0, sizeof maxv);
11            memset(sumv, 0, sizeof sumv);
12        }
13        void maintain(LL o, LL l, LL r) {
14            if (l < r) {
15                LL lc = o * 2, rc = o * 2 + 1;
16                sumv[o] = sumv[lc] + sumv[rc];
17                minv[o] = min(minv[lc], minv[rc]);
18                maxv[o] = max(maxv[lc], maxv[rc]);
19            } else sumv[o] = minv[o] = maxv[o] = 0;
20            if (setv[o] != RS) { minv[o] = maxv[o] = setv[o]; sumv[o] = setv[o] * (r - l + 1); }
21            if (addv[o]) { minv[o] += addv[o]; maxv[o] += addv[o]; sumv[o] += addv[o] * (r - l + 1); }

```

```

22     }
23     void build(LL o, LL l, LL r) {
24         if (l == r) addv[o] = a[l];
25         else {
26             LL m = (l + r) / 2;
27             build(ls); build(rs);
28         }
29         maintain(o, l, r);
30     }
31     void pushdown(LL o) {
32         LL lc = o * 2, rc = o * 2 + 1;
33         if (setv[o] != RS) {
34             setv[lc] = setv[rc] = setv[o];
35             addv[lc] = addv[rc] = 0;
36             setv[o] = RS;
37         }
38         if (addv[o]) {
39             addv[lc] += addv[o]; addv[rc] += addv[o];
40             addv[o] = 0;
41         }
42     }
43     void update(LL p, LL q, LL o, LL l, LL r, LL v, LL op) {
44         if (p <= r && l <= q)
45             if (p <= l && r <= q) {
46                 if (op == 2) { setv[o] = v; addv[o] = 0; }
47                 else addv[o] += v;
48             } else {
49                 pushdown(o);
50                 LL m = (l + r) / 2;
51                 update(p, q, ls, v, op); update(p, q, rs, v, op);
52             }
53         maintain(o, l, r);
54     }
55     void query(LL p, LL q, LL o, LL l, LL r, LL add, LL& ssum, LL& smin, LL& smax) {
56         if (p > r || l > q) return;
57         if (setv[o] != RS) {
58             LL v = setv[o] + add + addv[o];
59             ssum += v * (min(r, q) - max(l, p) + 1);
60             smin = min(smin, v);
61             smax = max(smax, v);
62         } else if (p <= l && r <= q) {
63             ssum += sumv[o] + add * (r - l + 1);
64             smin = min(smin, minv[o] + add);
65             smax = max(smax, maxv[o] + add);
66         } else {
67             LL m = (l + r) / 2;
68             query(p, q, ls, add + addv[o], ssum, smin, smax);
69             query(p, q, rs, add + addv[o], ssum, smin, smax);
70         }
71     }
72 } IT;

```

均摊复杂度线段树

- 区间取 max, 区间求和。

```

1 namespace R {
2     #define lson o * 2, l, (l + r) / 2
3     #define rson o * 2 + 1, (l + r) / 2 + 1, r
4     int m1[N], m2[N], cm1[N];
5     LL sum[N];
6     void up(int o) {
7         int lc = o * 2, rc = lc + 1;
8         m1[o] = max(m1[lc], m1[rc]);
9         sum[o] = sum[lc] + sum[rc];
10        if (m1[lc] == m1[rc]) {
11            cm1[o] = cm1[lc] + cm1[rc];
12            m2[o] = max(m2[lc], m2[rc]);
13        } else {

```

```

14         cm1[o] = m1[lc] > m1[rc] ? cm1[lc] : cm1[rc];
15         m2[o] = max(min(m1[lc], m1[rc]), max(m2[lc], m2[rc]));
16     }
17 }
18 void mod(int o, int x) {
19     if (x >= m1[o]) return;
20     assert(x > m2[o]);
21     sum[o] -= 1LL * (m1[o] - x) * cm1[o];
22     m1[o] = x;
23 }
24 void down(int o) {
25     int lc = o * 2, rc = lc + 1;
26     mod(lc, m1[o]); mod(rc, m1[o]);
27 }
28 void build(int o, int l, int r) {
29     if (l == r) { int t; read(t); sum[o] = m1[o] = t; m2[o] = -1; cm1[o] = 1; }
30     else { build(lson); build(rson); up(o); }
31 }
32 void update(int ql, int qr, int x, int o, int l, int r) {
33     if (r < ql || qr < l || m1[o] <= x) return;
34     if (ql <= l && r <= qr && m2[o] < x) { mod(o, x); return; }
35     down(o);
36     update(ql, qr, x, lson); update(ql, qr, x, rson);
37     up(o);
38 }
39 int qmax(int ql, int qr, int o, int l, int r) {
40     if (r < ql || qr < l) return -INF;
41     if (ql <= l && r <= qr) return m1[o];
42     down(o);
43     return max(qmax(ql, qr, lson), qmax(ql, qr, rson));
44 }
45 LL qsum(int ql, int qr, int o, int l, int r) {
46     if (r < ql || qr < l) return 0;
47     if (ql <= l && r <= qr) return sum[o];
48     down(o);
49     return qsum(ql, qr, lson) + qsum(ql, qr, rson);
50 }
51 }

```

持久化线段树

• ADD

```

1 namespace tree {
2     #define mid ((l + r) >> 1)
3     #define lson ql, qr, l, mid
4     #define rson ql, qr, mid + 1, r
5     struct P {
6         LL add, sum;
7         int ls, rs;
8     } tr[maxn * 45 * 2];
9     int sz = 1;
10    int N(LL add, int l, int r, int ls, int rs) {
11        tr[sz] = {add, tr[ls].sum + tr[rs].sum + add * (len[r] - len[l - 1]), ls, rs};
12        return sz++;
13    }
14    int update(int o, int ql, int qr, int l, int r, LL add) {
15        if (ql > r || l > qr) return o;
16        const P& t = tr[o];
17        if (ql <= l && r <= qr) return N(add + t.add, l, r, t.ls, t.rs);
18        return N(t.add, l, r, update(t.ls, lson, add), update(t.rs, rson, add));
19    }
20    LL query(int o, int ql, int qr, int l, int r, LL add = 0) {
21        if (ql > r || l > qr) return 0;
22        const P& t = tr[o];
23        if (ql <= l && r <= qr) return add * (len[r] - len[l - 1]) + t.sum;
24        return query(t.ls, lson, add + t.add) + query(t.rs, rson, add + t.add);
25    }
26 }

```


K-D Tree

最优化问题一定要用全局变量大力剪枝，而且左右儿子先递归潜力大的

- 维护信息
- 带重构（适合在线）
- 插入时左右儿子要标记为 null。

```
1 namespace kd {
2     const int K = 2, inf = 1E9, M = N;
3     const double lim = 0.7;
4     struct P {
5         int d[K], l[K], r[K], sz, val;
6         LL sum;
7         P *ls, *rs;
8         P* up() {
9             sz = ls->sz + rs->sz + 1;
10            sum = ls->sum + rs->sum + val;
11            FOR (i, 0, K) {
12                l[i] = min(d[i], min(ls->l[i], rs->l[i]));
13                r[i] = max(d[i], max(ls->r[i], rs->r[i]));
14            }
15            return this;
16        }
17    } pool[M], *null = new P, *pit = pool;
18    static P *tmp[M], **pt;
19    void init() {
20        null->ls = null->rs = null;
21        FOR (i, 0, K) null->l[i] = inf, null->r[i] = -inf;
22        null->sum = null->val = 0;
23        null->sz = 0;
24    }
25
26    P* build(P** l, P** r, int d = 0) { // [l, r)
27        if (d == K) d = 0;
28        if (l >= r) return null;
29        P** m = l + (r - l) / 2; assert(l <= m && m < r);
30        nth_element(l, m, r, [&](const P* a, const P* b){
31            return a->d[d] < b->d[d];
32        });
33        P* o = *m;
34        o->ls = build(l, m, d + 1); o->rs = build(m + 1, r, d + 1);
35        return o->up();
36    }
37    P* Build() {
38        pt = tmp; FOR (it, pool, pit) *pt++ = it;
39        return build(tmp, pt);
40    }
41    inline bool inside(int p[], int q[], int l[], int r[]) {
42        FOR (i, 0, K) if (r[i] < q[i] || p[i] < l[i]) return false;
43        return true;
44    }
45    LL query(P* o, int l[], int r[]) {
46        if (o == null) return 0;
47        FOR (i, 0, K) if (o->r[i] < l[i] || r[i] < o->l[i]) return 0;
48        if (inside(o->l, o->r, l, r)) return o->sum;
49        return query(o->ls, l, r) + query(o->rs, l, r) +
50            (inside(o->d, o->d, l, r) ? o->val : 0);
51    }
52    void dfs(P* o) {
53        if (o == null) return;
54        *pt++ = o; dfs(o->ls); dfs(o->rs);
55    }
56    P* ins(P* o, P* x, int d = 0) {
57        if (d == K) d = 0;
58        if (o == null) return o->up();
59        P*& oo = x->d[d] <= o->d[d] ? o->ls : o->rs;
60        if (oo->sz > o->sz * lim) {
61            pt = tmp; dfs(o); *pt++ = x;
62            return build(tmp, pt, d);
63        }
64    }
```

```

64         oo = ins(oo, x, d + 1);
65         return o->up();
66     }
67 }

```

- 维护信息
- 带修改 (适合离线)

```

1  namespace kd {
2      const int K = 3, inf = 1E9, M = N << 3;
3      extern struct P* null;
4      struct P {
5          int d[K], l[K], r[K], val;
6          int Max;
7          P *ls, *rs, *fa;
8          P* up() {
9              Max = max(val, max(ls->Max, rs->Max));
10             FOR (i, 0, K) {
11                 l[i] = min(d[i], min(ls->l[i], rs->l[i]));
12                 r[i] = max(d[i], max(ls->r[i], rs->r[i]));
13             }
14             return ls->fa = rs->fa = this;
15         }
16     } pool[M], *null = new P, *pit = pool;
17     void upd(P* o, int val) {
18         o->val = val;
19         for (; o != null; o = o->fa)
20             o->Max = max(o->Max, val);
21     }
22     static P *tmp[M], **pt;
23     void init() {
24         null->ls = null->rs = null;
25         FOR (i, 0, K) null->l[i] = inf, null->r[i] = -inf;
26         null->Max = null->val = 0;
27     }
28     P* build(P** l, P** r, int d = 0) { // [l, r)
29         if (d == K) d = 0;
30         if (l >= r) return null;
31         P** m = l + (r - l) / 2; assert(l <= m && m < r);
32         nth_element(l, m, r, [&](const P* a, const P* b){
33             return a->d[d] < b->d[d];
34         });
35         P* o = *m;
36         o->ls = build(l, m, d + 1); o->rs = build(m + 1, r, d + 1);
37         return o->up();
38     }
39     P* Build() {
40         pt = tmp; FOR (it, pool, pit) *pt++ = it;
41         P* ret = build(tmp, pt); ret->fa = null;
42         return ret;
43     }
44     inline bool inside(int p[], int q[], int l[], int r[]) {
45         FOR (i, 0, K) if (r[i] < q[i] || p[i] < l[i]) return false;
46         return true;
47     }
48     int query(P* o, int l[], int r[]) {
49         if (o == null) return 0;
50         FOR (i, 0, K) if (o->r[i] < l[i] || r[i] < o->l[i]) return 0;
51         if (inside(o->l, o->r, l, r)) return o->Max;
52         int ret = 0;
53         if (o->val > ret && inside(o->d, o->d, l, r)) ret = max(ret, o->val);
54         if (o->ls->Max > ret) ret = max(ret, query(o->ls, l, r));
55         if (o->rs->Max > ret) ret = max(ret, query(o->rs, l, r));
56         return ret;
57     }
58 }

```

- 最近点对
- 要用全局变量大力剪枝

```

1  namespace kd {
2      const int K = 3;

```

```

3  const int M = N;
4  const int inf = 1E9 + 100;
5  struct P {
6      int d[K];
7      int l[K], r[K];
8      P *ls, *rs;
9      P* up() {
10         FOR (i, 0, K) {
11             l[i] = min(d[i], min(ls->l[i], rs->l[i]));
12             r[i] = max(d[i], max(ls->r[i], rs->r[i]));
13         }
14         return this;
15     }
16 } pool[M], *null = new P, *pit = pool;
17 static P *tmp[M], **pt;
18 void init() {
19     null->ls = null->rs = null;
20     FOR (i, 0, K) null->l[i] = inf, null->r[i] = -inf;
21 }
22 P* build(P** l, P** r, int d = 0) { // [l, r]
23     if (d == K) d = 0;
24     if (l >= r) return null;
25     P** m = l + (r - l) / 2;
26     nth_element(l, m, r, [&](const P* a, const P* b){
27         return a->d[d] < b->d[d];
28     });
29     P* o = *m;
30     o->ls = build(l, m, d + 1); o->rs = build(m + 1, r, d + 1);
31     return o->up();
32 }
33 LL eval(P* o, int d[]) {
34     // ...
35 }
36 LL dist(int d1[], int d2[]) {
37     // ...
38 }
39 LL S;
40 LL query(P* o, int d[]) {
41     if (o == null) return 0;
42     S = max(S, dist(o->d, d));
43     LL mdl = eval(o->ls, d), mdr = eval(o->rs, d);
44     if (mdl < mdr) {
45         if (S > mdl) S = max(S, query(o->ls, d));
46         if (S > mdr) S = max(S, query(o->rs, d));
47     } else {
48         if (S > mdr) S = max(S, query(o->rs, d));
49         if (S > mdl) S = max(S, query(o->ls, d));
50     }
51     return S;
52 }
53 P* Build() {
54     pt = tmp; FOR (it, pool, pit) *pt++ = it;
55     return build(tmp, pt);
56 }
57 }

```

树状数组

- 注意: 0 是无效下标

```

1  namespace bit {
2      LL c[M];
3      inline int lowbit(int x) { return x & -x; }
4      void add(int x, LL v) {
5          for (; x < M; x += lowbit(x))
6              c[x] += v;
7      }
8      LL sum(int x) {
9          LL ret = 0;

```

```

10     for (; x > 0; x -= lowbit(x))
11         ret += c[x];
12     return ret;
13 }
14 int kth(LL k) {
15     int ret = 0;
16     LL cnt = 0;
17     FORD (i, 20, -1) {
18         ret += 1 << i;
19         if (ret >= M || cnt + c[ret] >= k)
20             ret -= 1 << i;
21         else cnt += c[ret];
22     }
23     return ret + 1;
24 }
25 }

```

- 区间修改 & 区间查询 (单点修改, 查询前缀和的前缀和)

```

1 namespace bit {
2     int c[maxn], cc[maxn];
3     inline int lowbit(int x) { return x & -x; }
4     void add(int x, int v) {
5         for (int i = x; i <= n; i += lowbit(i)) {
6             c[i] += v; cc[i] += x * v;
7         }
8     }
9     void add(int l, int r, int v) { add(l, v); add(r + 1, -v); }
10    int sum(int x) {
11        int ret = 0;
12        for (int i = x; i > 0; i -= lowbit(i))
13            ret += (x + 1) * c[i] - cc[i];
14        return ret;
15    }
16    int sum(int l, int r) { return sum(r) - sum(l - 1); }
17 }

```

- 单点修改, 查询前缀和的前缀和的前缀和 (有用才怪)

```

1 namespace bit {
2     LL c[N], cc[N], ccc[N];
3     inline LL lowbit(LL x) { return x & -x; }
4     void add(LL x, LL v) {
5         for (LL i = x; i < N; i += lowbit(i)) {
6             c[i] = (c[i] + v) % MOD;
7             cc[i] = (cc[i] + x * v) % MOD;
8             ccc[i] = (ccc[i] + x * x % MOD * v) % MOD;
9         }
10    }
11    void add(LL l, LL r, LL v) { add(l, v); add(r + 1, -v); }
12    LL sum(LL x) {
13        static LL INV2 = (MOD + 1) / 2;
14        LL ret = 0;
15        for (LL i = x; i > 0; i -= lowbit(i))
16            ret += (x + 1) * (x + 2) % MOD * c[i] % MOD
17                - (2 * x + 3) * cc[i] % MOD
18                + ccc[i];
19        return ret % MOD * INV2 % MOD;
20    }
21    LL sum(LL l, LL r) { return sum(r) - sum(l - 1); }
22 }

```

- 三维

```

1 inline int lowbit(int x) { return x & -x; }
2 void update(int x, int y, int z, int d) {
3     for (int i = x; i <= n; i += lowbit(i))
4         for (int j = y; j <= n; j += lowbit(j))
5             for (int k = z; k <= n; k += lowbit(k))
6                 c[i][j][k] += d;
7 }
8 LL query(int x, int y, int z) {
9     LL ret = 0;

```

```

10     for (int i = x; i > 0; i -= lowbit(i))
11         for (int j = y; j > 0; j -= lowbit(j))
12             for (int k = z; k > 0; k -= lowbit(k))
13                 ret += c[i][j][k];
14     return ret;
15 }
16 LL solve(int x, int y, int z, int xx, int yy, int zz) {
17     return
18         query(xx, yy, zz)
19         - query(xx, yy, z - 1)
20         - query(xx, y - 1, zz)
21         - query(x - 1, yy, zz)
22         + query(xx, y - 1, z - 1)
23         + query(x - 1, yy, z - 1)
24         + query(x - 1, y - 1, zz)
25         - query(x - 1, y - 1, z - 1);

```

主席树

● 正常主席树

```

1 namespace tree {
2     #define mid ((l + r) >> 1)
3     #define lson l, mid
4     #define rson mid + 1, r
5     const int MAGIC = M * 30;
6     struct P {
7         int sum, ls, rs;
8     } tr[MAGIC] = {{0, 0, 0}};
9     int sz = 1;
10    int N(int sum, int ls, int rs) {
11        if (sz == MAGIC) assert(0);
12        tr[sz] = {sum, ls, rs};
13        return sz++;
14    }
15    int ins(int o, int x, int v, int l = 1, int r = ls) {
16        if (x < l || x > r) return o;
17        const P& t = tr[o];
18        if (l == r) return N(t.sum + v, 0, 0);
19        return N(t.sum + v, ins(t.ls, x, v, lson), ins(t.rs, x, v, rson));
20    }
21    int query(int o, int ql, int qr, int l = 1, int r = ls) {
22        if (ql > r || l > qr) return 0;
23        const P& t = tr[o];
24        if (ql <= l && r <= qr) return t.sum;
25        return query(t.ls, ql, qr, lson) + query(t.rs, ql, qr, rson);
26    }
27 }

```

● 第k大

```

1 struct TREE {
2     #define mid ((l + r) >> 1)
3     #define lson l, mid
4     #define rson mid + 1, r
5     struct P {
6         int w, ls, rs;
7     } tr[maxn * 20];
8     int sz = 1;
9     TREE() { tr[0] = {0, 0, 0}; }
10    int N(int w, int ls, int rs) {
11        tr[sz] = {w, ls, rs};
12        return sz++;
13    }
14    int ins(int tt, int l, int r, int x) {
15        if (x < l || r < x) return tt;
16        const P& t = tr[tt];
17        if (l == r) return N(t.w + 1, 0, 0);
18        return N(t.w + 1, ins(t.ls, lson, x), ins(t.rs, rson, x));
19    }
20    int query(int pp, int qq, int l, int r, int k) { // (pp, qq]

```

```

21     if (l == r) return l;
22     const P &p = tr[pp], &q = tr[qq];
23     int w = tr[q.ls].w - tr[p.ls].w;
24     if (k <= w) return query(p.ls, q.ls, lson, k);
25     else return query(p.rs, q.rs, rson, k - w);
26 }
27 } tree;

```

● 树状数组套主席树

```

1  typedef vector<int> VI;
2  struct TREE {
3      #define mid ((l + r) >> 1)
4      #define lson l, mid
5      #define rson mid + 1, r
6      struct P {
7          int w, ls, rs;
8      } tr[maxn * 20 * 20];
9      int sz = 1;
10     TREE() { tr[0] = {0, 0, 0}; }
11     int N(int w, int ls, int rs) {
12         tr[sz] = {w, ls, rs};
13         return sz++;
14     }
15     int add(int tt, int l, int r, int x, int d) {
16         if (x < l || r < x) return tt;
17         const P& t = tr[tt];
18         if (l == r) return N(t.w + d, 0, 0);
19         return N(t.w + d, add(t.ls, lson, x, d), add(t.rs, rson, x, d));
20     }
21     int ls_sum(const VI& rt) {
22         int ret = 0;
23         FOR (i, 0, rt.size())
24             ret += tr[rt[i]].ls.w;
25         return ret;
26     }
27     inline void ls(VI& rt) { transform(rt.begin(), rt.end(), rt.begin(), [&](int x)->int{ return tr[x].ls; }); }
28     inline void rs(VI& rt) { transform(rt.begin(), rt.end(), rt.begin(), [&](int x)->int{ return tr[x].rs; }); }
29     int query(VI& p, VI& q, int l, int r, int k) {
30         if (l == r) return l;
31         int w = ls_sum(q) - ls_sum(p);
32         if (k <= w) {
33             ls(p); ls(q);
34             return query(p, q, lson, k);
35         }
36         else {
37             rs(p); rs(q);
38             return query(p, q, rson, k - w);
39         }
40     }
41 } tree;
42 struct BIT {
43     int root[maxn];
44     void init() { memset(root, 0, sizeof root); }
45     inline int lowbit(int x) { return x & -x; }
46     void update(int p, int x, int d) {
47         for (int i = p; i <= m; i += lowbit(i))
48             root[i] = tree.add(root[i], 1, m, x, d);
49     }
50     int query(int l, int r, int k) {
51         VI p, q;
52         for (int i = l - 1; i > 0; i -= lowbit(i)) p.push_back(root[i]);
53         for (int i = r; i > 0; i -= lowbit(i)) q.push_back(root[i]);
54         return tree.query(p, q, 1, m, k);
55     }
56 } bit;
57
58 void init() {
59     m = 10000;
60     tree.sz = 1;
61     bit.init();
62     FOR (i, 1, m + 1)

```

```

63         bit.update(i, a[i], 1);
64     }

```

左偏树

```

1  namespace LTree {
2      extern struct P* null, *pit;
3      queue<P*> trash;
4      const int M = 1E5 + 100;
5      struct P {
6          P *ls, *rs;
7          LL v;
8          int d;
9          void operator delete (void* ptr) {
10             trash.push((P*)ptr);
11         }
12         void* operator new(size_t size) {
13             if (trash.empty()) return pit++;
14             void* ret = trash.front(); trash.pop(); return ret;
15         }
16
17         void prt() {
18             if (this == null) return;
19             cout << v << ' ';
20             ls->prt(); rs->prt();
21         }
22     } pool[M], *pit = pool, *null = new P{0, 0, -1, -1};
23     P* N(LL v) {
24         return new P{null, null, v, 0};
25     }
26     P* merge(P* a, P* b) {
27         if (a == null) return b;
28         if (b == null) return a;
29         if (a->v > b->v) swap(a, b);
30         a->rs = merge(a->rs, b);
31         if (a->ls->d < a->rs->d) swap(a->ls, a->rs);
32         a->d = a->rs->d + 1;
33         return a;
34     }
35
36     LL pop(P*& o) {
37         LL ret = o->v;
38         P* t = o;
39         o = merge(o->ls, o->rs);
40         delete t;
41         return ret;
42     }
43 }

```

可持久化

```

1  namespace LTree {
2      extern struct P* null, *pit;
3      queue<P*> trash;
4      const int M = 1E6 + 100;
5      struct P {
6          P *ls, *rs;
7          LL v;
8          int d;
9          void operator delete (void* ptr) {
10             trash.push((P*)ptr);
11         }
12         void* operator new(size_t size) {
13             if (trash.empty()) return pit++;
14             void* ret = trash.front(); trash.pop(); return ret;
15         }
16     } pool[M], *pit = pool, *null = new P{0, 0, -1, -1};
17     P* N(LL v, P* ls = null, P* rs = null) {
18         if (ls->d < rs->d) swap(ls, rs);
19         return new P{ls, rs, v, rs->d + 1};

```

```

20     }
21     P* merge(P* a, P* b) {
22         if (a == null) return b;
23         if (b == null) return a;
24         if (a->v < b->v)
25             return N(a->v, a->ls, merge(a->rs, b));
26         else
27             return N(b->v, b->ls, merge(b->rs, a));
28     }
29
30     LL pop(P*& o) {
31         LL ret = o->v;
32         o = merge(o->ls, o->rs);
33         return ret;
34     }
35 }

```

Treap

- 非旋 Treap
- v 小根堆
- 模板题 bzoj 3224
- lower 第一个大于等于的是第几个 (0-based)
- upper 第一个大于的是第几个 (0-based)
- split 左侧分割出 rk 个元素
- 树套树略

```

1 namespace treap {
2     const int M = maxn * 17;
3     extern struct P* const null;
4     struct P {
5         P *ls, *rs;
6         int v, sz;
7         unsigned rd;
8         P(int v): ls(null), rs(null), v(v), sz(1), rd(rnd()) {}
9         P(): sz(0) {}
10
11         P* up() { sz = ls->sz + rs->sz + 1; return this; }
12         int lower(int v) {
13             if (this == null) return 0;
14             return this->v >= v ? ls->lower(v) : rs->lower(v) + ls->sz + 1;
15         }
16         int upper(int v) {
17             if (this == null) return 0;
18             return this->v > v ? ls->upper(v) : rs->upper(v) + ls->sz + 1;
19         }
20     } *const null = new P, pool[M], *pit = pool;
21
22     P* merge(P* l, P* r) {
23         if (l == null) return r; if (r == null) return l;
24         if (l->rd < r->rd) { l->rs = merge(l->rs, r); return l->up(); }
25         else { r->ls = merge(l, r->ls); return r->up(); }
26     }
27
28     void split(P* o, int rk, P*& l, P*& r) {
29         if (o == null) { l = r = null; return; }
30         if (o->ls->sz >= rk) { split(o->ls, rk, l, o->ls); r = o->up(); }
31         else { split(o->rs, rk - o->ls->sz - 1, o->rs, r); l = o->up(); }
32     }
33 }

```

- 持久化 Treap

```

1 namespace treap {
2     const int M = maxn * 17 * 12;
3     extern struct P* const null, *pit;
4     struct P {
5         P *ls, *rs;

```



```

6      int v, sz;
7      LL sum;
8      P(P* ls, P* rs, int v): ls(ls), rs(rs), v(v), sz(ls->sz + rs->sz + 1),
9                                     sum(ls->sum + rs->sum + v) {}
10
11      P() {}
12
13      void* operator new(size_t _) { return pit++; }
14      template<typename T>
15      int rk(int v, T&& cmp) {
16          if (this == null) return 0;
17          return cmp(this->v, v) ? ls->rk(v, cmp) : rs->rk(v, cmp) + ls->sz + 1;
18      }
19      int lower(int v) { return rk(v, greater_equal<int>()); }
20      int upper(int v) { return rk(v, greater<int>()); }
21 } pool[M], *pit = pool, *const null = new P;
22 P* merge(P* l, P* r) {
23     if (l == null) return r; if (r == null) return l;
24     if (rnd() % (l->sz + r->sz) < l->sz) return new P{l->ls, merge(l->rs, r), l->v};
25     else return new P{merge(l, r->ls), r->rs, r->v};
26 }
27 void split(P* o, int rk, P*& l, P*& r) {
28     if (o == null) { l = r = null; return; }
29     if (o->ls->sz >= rk) { split(o->ls, rk, l, r); r = new P{r, o->rs, o->v}; }
30     else { split(o->rs, rk - o->ls->sz - 1, l, r); l = new P{o->ls, l, o->v}; }
31 }

```

- 带 pushdown 的持久化 Treap
- 注意任何修改操作前一定要 FIX

```

1  int now;
2  namespace Treap {
3      const int M = 10000000;
4      extern struct P* const null, *pit;
5      struct P {
6          P *ls, *rs;
7          int sz, time;
8          LL cnt, sc, pos, add;
9          bool rev;
10
11          P* up() { sz = ls->sz + rs->sz + 1; sc = ls->sc + rs->sc + cnt; return this; } // MOD
12          P* check() {
13              if (time == now) return this;
14              P* t = new(pit++) P; *t = *this; t->time = now; return t;
15          };
16          P* _do_rev() { rev ^= 1; add *= -1; pos *= -1; swap(ls, rs); return this; } // MOD
17          P* _do_add(LL v) { add += v; pos += v; return this; } // MOD
18          P* do_rev() { if (this == null) return this; return check()->_do_rev(); } // FIX & MOD
19          P* do_add(LL v) { if (this == null) return this; return check()->_do_add(v); } // FIX & MOD
20          P* _down() { // MOD
21              if (rev) { ls = ls->do_rev(); rs = rs->do_rev(); rev = 0; }
22              if (add) { ls = ls->do_add(add); rs = rs->do_add(add); add = 0; }
23              return this;
24          }
25          P* down() { return check()->_down(); } // FIX & MOD
26          void _split(LL p, P*& l, P*& r) { // MOD
27              if (pos >= p) { ls->split(p, l, r); ls = r; r = up(); }
28              else { rs->split(p, l, r); rs = l; l = up(); }
29          }
30          void split(LL p, P*& l, P*& r) { // FIX & MOD
31              if (this == null) l = r = null;
32              else down()->_split(p, l, r);
33          }
34      } pool[M], *pit = pool, *const null = new P;
35      P* merge(P* a, P* b) {
36          if (a == null) return b; if (b == null) return a;
37          if (rand() % (a->sz + b->sz) < a->sz) { a = a->down(); a->rs = merge(a->rs, b); return a->up(); }
38          else { b = b->down(); b->ls = merge(a, b->ls); return b->up(); }
39      }
40  }

```

Treap-序列

- 区间 ADD, SUM

```
1 namespace treap {
2     const int M = 8E5 + 100;
3     extern struct P* const null;
4     struct P {
5         P *ls, *rs;
6         int sz, val, add, sum;
7         P(int v, P* ls = null, P* rs = null): ls(ls), rs(rs), sz(1), val(v), add(0), sum(v) {}
8         P(): sz(0), val(0), add(0), sum(0) {}
9
10        P* up() {
11            assert(this != null);
12            sz = ls->sz + rs->sz + 1;
13            sum = ls->sum + rs->sum + val + add * sz;
14            return this;
15        }
16        void upd(int v) {
17            if (this == null) return;
18            add += v;
19            sum += sz * v;
20        }
21        P* down() {
22            if (add) {
23                ls->upd(add); rs->upd(add);
24                val += add;
25                add = 0;
26            }
27            return this;
28        }
29
30        P* select(int rk) {
31            if (rk == ls->sz + 1) return this;
32            return ls->sz >= rk ? ls->select(rk) : rs->select(rk - ls->sz - 1);
33        }
34    } pool[M], *pit = pool, *const null = new P, *rt = null;
35
36    P* merge(P* a, P* b) {
37        if (a == null) return b->up();
38        if (b == null) return a->up();
39        if (rand() % (a->sz + b->sz) < a->sz) {
40            a->down()->rs = merge(a->rs, b);
41            return a->up();
42        } else {
43            b->down()->ls = merge(a, b->ls);
44            return b->up();
45        }
46    }
47
48    void split(P* o, int rk, P*& l, P*& r) {
49        if (o == null) { l = r = null; return; }
50        o->down();
51        if (o->ls->sz >= rk) {
52            split(o->ls, rk, l, o->ls);
53            r = o->up();
54        } else {
55            split(o->rs, rk - o->ls->sz - 1, o->rs, r);
56            l = o->up();
57        }
58    }
59
60    inline void insert(int k, int v) {
61        P *l, *r;
62        split(rt, k - 1, l, r);
63        rt = merge(merge(l, new (pit++) P(v)), r);
64    }
65
66    inline void erase(int k) {
67        P *l, *r, *_ , *t;
```

```

68     split(rt, k - 1, l, t);
69     split(t, 1, _, r);
70     rt = merge(l, r);
71 }
72
73 P* build(int l, int r, int* a) {
74     if (l > r) return null;
75     if (l == r) return new(pit++) P(a[l]);
76     int m = (l + r) / 2;
77     return (new(pit++) P(a[m], build(l, m - 1, a), build(m + 1, r, a)))->up();
78 }
79 };

```

● 区间 REVERSE, ADD, MIN

```

1  namespace treap {
2      extern struct P*const null;
3      struct P {
4          P *ls, *rs;
5          int sz, v, add, m;
6          bool flip;
7          P(int v, P* ls = null, P* rs = null): ls(ls), rs(rs), sz(1), v(v), add(0), m(v), flip(0) {}
8          P(): sz(0), v(INF), m(INF) {}
9
10         void upd(int v) {
11             if (this == null) return;
12             add += v; m += v;
13         }
14         void rev() {
15             if (this == null) return;
16             swap(ls, rs);
17             flip ^= 1;
18         }
19         P* up() {
20             assert(this != null);
21             sz = ls->sz + rs->sz + 1;
22             m = min(min(ls->m, rs->m), v) + add;
23             return this;
24         }
25         P* down() {
26             if (add) {
27                 ls->upd(add); rs->upd(add);
28                 v += add;
29                 add = 0;
30             }
31             if (flip) {
32                 ls->rev(); rs->rev();
33                 flip = 0;
34             }
35             return this;
36         }
37
38         P* select(int k) {
39             if (ls->sz + 1 == k) return this;
40             if (ls->sz >= k) return ls->select(k);
41             return rs->select(k - ls->sz - 1);
42         }
43
44     } pool[M], *const null = new P, *pit = pool, *rt = null;
45
46     P* merge(P* a, P* b) {
47         if (a == null) return b;
48         if (b == null) return a;
49         if (rnd() % (a->sz + b->sz) < a->sz) {
50             a->down()->rs = merge(a->rs, b);
51             return a->up();
52         } else {
53             b->down()->ls = merge(a, b->ls);
54             return b->up();
55         }
56     }
57 }

```

```

58 void split(P* o, int k, P*& l, P*& r) {
59     if (o == null) { l = r = null; return; }
60     o->down();
61     if (o->ls->sz >= k) {
62         split(o->ls, k, l, o->ls);
63         r = o->up();
64     } else {
65         split(o->rs, k - o->ls->sz - 1, o->rs, r);
66         l = o->up();
67     }
68 }
69
70 P* build(int l, int r, int* v) {
71     if (l > r) return null;
72     int m = (l + r) >> 1;
73     return (new (pit++) P(v[m], build(l, m - 1, v), build(m + 1, r, v)))->up();
74 }
75
76 void go(int x, int y, void f(P*&)) {
77     P *l, *m, *r;
78     split(rt, y, l, r);
79     split(l, x - 1, l, m);
80     f(m);
81     rt = merge(merge(l, m), r);
82 }
83
84 using namespace treap;
85 int a[maxn], n, x, y, Q, v, k, d;
86 char s[100];
87
88 int main() {
89     cin >> n;
90     FOR (i, 1, n + 1) scanf("%d", &a[i]);
91     rt = build(1, n, a);
92     cin >> Q;
93     while (Q--) {
94         scanf("%s", s);
95         if (s[0] == 'A') {
96             scanf("%d%d", &x, &y, &v);
97             go(x, y, [](P*& o){ o->upd(v); });
98         } else if (s[0] == 'R' && s[3] == 'E') {
99             scanf("%d", &x, &y);
100            go(x, y, [](P*& o){ o->rev(); });
101        } else if (s[0] == 'R' && s[3] == 'O') {
102            scanf("%d%d", &x, &y, &d);
103            d %= y - x + 1;
104            go(x, y, [](P*& o){
105                P *l, *r;
106                split(o, o->sz - d, l, r);
107                o = merge(r, l);
108            });
109        } else if (s[0] == 'I') {
110            scanf("%d", &k, &v);
111            go(k + 1, k, [](P*& o){ o = new (pit++) P(v); });
112        } else if (s[0] == 'D') {
113            scanf("%d", &k);
114            go(k, k, [](P*& o){ o = null; });
115        } else if (s[0] == 'M') {
116            scanf("%d", &x, &y);
117            go(x, y, [](P*& o) {
118                printf("%d\n", o->m);
119            });
120        }
121    }
122 }

```

● 持久化

```

1 namespace treap {
2     struct P;
3     extern P*const null;
4     P* N(P* ls, P* rs, LL v, bool fill);

```

```

5 struct P {
6     P *const ls, *const rs;
7     const int sz, v;
8     const LL sum;
9     bool fill;
10    int cnt;
11
12    void split(int k, P*& l, P*& r) {
13        if (this == null) { l = r = null; return; }
14        if (ls->sz >= k) {
15            ls->split(k, l, r);
16            r = N(r, rs, v, fill);
17        } else {
18            rs->split(k - ls->sz - fill, l, r);
19            l = N(ls, l, v, fill);
20        }
21    }
22
23
24    } *const null = new P{0, 0, 0, 0, 0, 0, 1};
25
26    P* N(P* ls, P* rs, LL v, bool fill) {
27        ls->cnt++; rs->cnt++;
28        return new P{ls, rs, ls->sz + rs->sz + fill, v, ls->sum + rs->sum + v, fill, 1};
29    }
30
31    P* merge(P* a, P* b) {
32        if (a == null) return b;
33        if (b == null) return a;
34        if (rand() % (a->sz + b->sz) < a->sz)
35            return N(a->ls, merge(a->rs, b), a->v, a->fill);
36        else
37            return N(merge(a, b->ls), b->rs, b->v, b->fill);
38    }
39
40    void go(P* o, int x, int y, P*& l, P*& m, P*& r) {
41        o->split(y, l, r);
42        l->split(x - 1, l, m);
43    }
44 }

```

可回滚并查集

- 注意这个不是可持久化并查集
- 查找时不进行路径压缩
- 复杂度靠按秩合并解决

```

1 namespace uf {
2     int fa[maxn], sz[maxn];
3     int undo[maxn], top;
4     void init() { memset(fa, -1, sizeof fa); memset(sz, 0, sizeof sz); top = 0; }
5     int findset(int x) { while (fa[x] != -1) x = fa[x]; return x; }
6     bool join(int x, int y) {
7         x = findset(x); y = findset(y);
8         if (x == y) return false;
9         if (sz[x] > sz[y]) swap(x, y);
10        undo[top++] = x;
11        fa[x] = y;
12        sz[y] += sz[x] + 1;
13        return true;
14    }
15    inline int checkpoint() { return top; }
16    void rewind(int t) {
17        while (top > t) {
18            int x = undo[--top];
19            sz[fa[x]] -= sz[x] + 1;
20            fa[x] = -1;
21        }
22    }
23 }

```

```

22     }
23 }

```

舞蹈链

- 注意 link 的 y 的范围是 [1, n]
- 注意在某些情况下替换掉 memset
- 精确覆盖

```

1  struct P {
2      P *L, *R, *U, *D;
3      int x, y;
4  };
5
6  const int INF = 1E9;
7
8  struct DLX {
9      #define TR(i, D, s) for (P* i = s->D; i != s; i = i->D)
10     static const int M = 2E5;
11     P pool[M], *h[M], *r[M], *pit;
12     int sz[M];
13     bool solved;
14     stack<int> ans;
15     void init(int n) {
16         pit = pool;
17         ++n;
18         solved = false;
19         while (!ans.empty()) ans.pop();
20         memset(r, 0, sizeof r);
21         memset(sz, 0, sizeof sz);
22         FOR (i, 0, n)
23             h[i] = new (pit++) P;
24         FOR (i, 0, n) {
25             h[i]->L = h[(i + n - 1) % n];
26             h[i]->R = h[(i + 1) % n];
27             h[i]->U = h[i]->D = h[i];
28             h[i]->y = i;
29         }
30     }
31
32     void link(int x, int y) {
33         sz[y]++;
34         auto p = new (pit++) P;
35         p->x = x; p->y = y;
36         p->U = h[y]->U; p->D = h[y];
37         p->D->U = p->U->D = p;
38         if (!r[x]) r[x] = p->L = p->R = p;
39         else {
40             p->L = r[x]; p->R = r[x]->R;
41             p->L->R = p->R->L = p;
42         }
43     }
44
45     void remove(P* p) {
46         p->L->R = p->R; p->R->L = p->L;
47         TR (i, D, p)
48             TR (j, R, i) {
49                 j->D->U = j->U; j->U->D = j->D;
50                 sz[j->y]--;
51             }
52     }
53
54     void recall(P* p) {
55         p->L->R = p->R->L = p;
56         TR (i, U, p)
57             TR (j, L, i) {
58                 j->D->U = j->U->D = j;
59                 sz[j->y]++;

```

```

60     }
61 }
62
63 bool dfs(int d) {
64     if (solved) return true;
65     if (h[0]->R == h[0]) return solved = true;
66     int m = INF;
67     P* c;
68     TR (i, R, h[0])
69         if (sz[i->y] < m) { m = sz[i->y]; c = i; }
70     remove(c);
71     TR (i, D, c) {
72         ans.push(i->x);
73         TR (j, R, i) remove(h[j->y]);
74         if (dfs(d + 1)) return true;
75         TR (j, L, i) recall(h[j->y]);
76         ans.pop();
77     }
78     recall(c);
79     return false;
80 }
81 } dlx;

```

- 可重复覆盖

```

1  struct P {
2      P *L, *R, *U, *D;
3      int x, y;
4  };
5
6  const int INF = 1E9;
7
8  struct DLX {
9      #define TR(i, D, s) for (P* i = s->D; i != s; i = i->D)
10     static const int M = 2E5;
11     P pool[M], *h[M], *r[M], *pit;
12     int sz[M], vis[M], ans, clk;
13     void init(int n) {
14         clk = 0;
15         ans = INF;
16         pit = pool;
17         ++n;
18         memset(r, 0, sizeof r);
19         memset(sz, 0, sizeof sz);
20         memset(vis, -1, sizeof vis);
21         FOR (i, 0, n)
22             h[i] = new (pit++) P;
23         FOR (i, 0, n) {
24             h[i]->L = h[(i + n - 1) % n];
25             h[i]->R = h[(i + 1) % n];
26             h[i]->U = h[i]->D = h[i];
27             h[i]->y = i;
28         }
29     }
30
31     void link(int x, int y) {
32         sz[y]++;
33         auto p = new (pit++) P;
34         p->x = x; p->y = y;
35         p->U = h[y]->U; p->D = h[y];
36         p->D->U = p->U->D = p;
37         if (!r[x]) r[x] = p->L = p->R = p;
38         else {
39             p->L = r[x]; p->R = r[x]->R;
40             p->L->R = p->R->L = p;
41         }
42     }
43
44     void remove(P* p) {
45         TR (i, D, p) {
46             i->L->R = i->R;
47             i->R->L = i->L;

```

```

48     }
49 }
50
51 void recall(P* p) {
52     TR (i, U, p)
53     i->L->R = i->R->L = i;
54 }
55
56 int eval() {
57     ++clk;
58     int ret = 0;
59     TR (i, R, h[0])
60         if (vis[i->y] != clk) {
61             ++ret;
62             vis[i->y] = clk;
63             TR (j, D, i)
64                 TR (k, R, j)
65                     vis[k->y] = clk;
66         }
67     return ret;
68 }
69
70 void dfs(int d) {
71     if (h[0]->R == h[0]) { ans = min(ans, d); return; }
72     if (eval() + d >= ans) return;
73     P* c;
74     int m = INF;
75     TR (i, R, h[0])
76         if (sz[i->y] < m) { m = sz[i->y]; c = i; }
77     TR (i, D, c) {
78         remove(i);
79         TR (j, R, i) remove(j);
80         dfs(d + 1);
81         TR (j, L, i) recall(j);
82         recall(i);
83     }
84 }
85 } dlx;

```

CDQ 分治

```

1  const int maxn = 2E5 + 100;
2  struct P {
3      int x, y;
4      int* f;
5      bool d1, d2;
6  } a[maxn], b[maxn], c[maxn];
7  int f[maxn];
8
9  void go2(int l, int r) {
10     if (l + 1 == r) return;
11     int m = (l + r) >> 1;
12     go2(l, m); go2(m, r);
13     FOR (i, l, m) b[i].d2 = 0;
14     FOR (i, m, r) b[i].d2 = 1;
15     merge(b + l, b + m, b + m, b + r, c + l, [](const P& a, const P& b)->bool {
16         if (a.y != b.y) return a.y < b.y;
17         return a.d2 > b.d2;
18     });
19     int mx = -1;
20     FOR (i, l, r) {
21         if (c[i].d1 && c[i].d2) *c[i].f = max(*c[i].f, mx + 1);
22         if (!c[i].d1 && !c[i].d2) mx = max(mx, *c[i].f);
23     }
24     FOR (i, l, r) b[i] = c[i];
25 }
26
27 void go1(int l, int r) { // [l, r)
28     if (l + 1 == r) return;

```



```

29     int m = (l + r) >> 1;
30     go1(l, m);
31     FOR (i, l, m) a[i].d1 = 0;
32     FOR (i, m, r) a[i].d1 = 1;
33     copy(a + l, a + r, b + l);
34     sort(b + l, b + r, [] (const P& a, const P& b) -> bool {
35         if (a.x != b.x) return a.x < b.x;
36         return a.d1 > b.d1;
37     });
38     go2(l, r);
39     go1(m, r);
40 }

```

● k 维 LIS

```

1  struct P {
2      int v[K];
3      LL f;
4      bool d[K];
5  } o[N << 10];
6  P* a[K][N << 10];
7  int k;
8  void go(int now, int l, int r) {
9      if (now == 0) {
10         if (l + 1 == r) return;
11         int m = (l + r) / 2;
12         go(now, l, m);
13         FOR (i, l, m) a[now][i] -> d[now] = 0;
14         FOR (i, m, r) a[now][i] -> d[now] = 1;
15         copy(a[now] + l, a[now] + r, a[now + 1] + l);
16         sort(a[now + 1] + l, a[now + 1] + r, [now] (const P* a, const P* b) {
17             if (a -> v[now] != b -> v[now]) return a -> v[now] < b -> v[now];
18             return a -> d[now] > b -> d[now];
19         });
20         go(now + 1, l, r);
21         go(now, m, r);
22     } else {
23         if (l + 1 == r) return;
24         int m = (l + r) / 2;
25         go(now, l, m); go(now, m, r);
26         FOR (i, l, m) a[now][i] -> d[now] = 0;
27         FOR (i, m, r) a[now][i] -> d[now] = 1;
28         merge(a[now] + l, a[now] + m, a[now] + m, a[now] + r, a[now + 1] + l, [now] (const P* a, const P* b) {
29             if (a -> v[now] != b -> v[now]) return a -> v[now] < b -> v[now];
30             return a -> d[now] > b -> d[now];
31         });
32         copy(a[now + 1] + l, a[now + 1] + r, a[now] + l);
33         if (now < k - 2) {
34             go(now + 1, l, r);
35         } else {
36             LL sum = 0;
37             FOR (i, l, r) {
38                 dbg(a[now][i] -> v[0], a[now][i] -> v[1], a[now][i] -> f,
39                     a[now][i] -> d[0], a[now][i] -> d[1]);
40
41                 int cnt = 0;
42                 FOR (j, 0, now + 1) cnt += a[now][i] -> d[j];
43                 if (cnt == 0) {
44                     sum += a[now][i] -> f;
45                 } else if (cnt == now + 1) {
46                     a[now][i] -> f = (a[now][i] -> f + sum) % MOD;
47                 }
48             }
49         }
50     }
}

```

哈希表

● 必须初始化

- 备选素数 1572869, 3145739, 6291469, 12582917, 25165843, 50331653

```

1  const LL HASH_MOD=1572869;
2  LL key[HASH_MOD], val[HASH_MOD];
3  int head[HASH_MOD], next[HASH_MOD];
4  struct Hash {
5      int sz;
6      void init() {
7          memset(head, -1, sizeof head);
8          sz = 0;
9      }
10     LL insert(LL x, LL y) {
11         int k = x % HASH_MOD;
12         key[sz] = x;
13         val[sz] = y;
14         next[sz] = head[k];
15         head[k] = sz++;
16     }
17     LL find(LL x) {
18         int k = x % HASH_MOD;
19         for (int i = head[k]; i != -1; i = next[i])
20             if (key[i] == x)
21                 return val[i];
22         return -1;
23     }
24 };

```

笛卡尔树

```

1  void build(const vector<int>& a) {
2      static P *stack[M], *x, *last;
3      int p = 0;
4      FOR (i, 0, a.size()) {
5          x = new P(i + 1, a[i]);
6          last = null;
7          while (p && stack[p - 1]->v > x->v) {
8              stack[p - 1]->maintain();
9              last = stack[--p];
10         }
11         if (p) stack[p - 1]->rs = x;
12         x->ls = last;
13         stack[p++] = x;
14     }
15     while (p)
16         stack[--p]->maintain();
17     rt = stack[0];
18 }

1  void build() {
2      static int s[N], last;
3      int p = 0;
4      FOR (x, 1, n + 1) {
5          last = 0;
6          while (p && val[s[p - 1]] > val[x]) last = s[--p];
7          if (p) G[s[p - 1]][1] = x;
8          if (last) G[x][0] = last;
9          s[p++] = x;
10     }
11     rt = s[0];
12 }

```

Trie

- Trie 二进制版
- M 为二进制的位数
- 使用前必须初始化

```

1 struct Trie2 {
2     int ch[N * M][2], sz;
3     void init() {
4         memset(ch, 0, sizeof ch);
5         sz = 1;
6     }
7     void insert(LL x) {
8         int u = 0;
9         FOR (i, M, -1) {
10             bool b = x & (1LL << i);
11             if (!ch[u][b])
12                 ch[u][b] = sz++;
13             u = ch[u][b];
14         }
15     }
16 } trie;

```

pb_ds

- 优先队列
- binary_heap_tag
- pairing_heap_tag 支持修改
- thin_heap_tag 如果修改只有 increase 则较快, 不支持 join

```

1 #include<ext/pb_ds/priority_queue.hpp>
2 template<typename _Tv,
3         typename Cmp_Fn = std::less<_Tv>,
4         typename Tag = pairing_heap_tag,
5         typename _Alloc = std::allocator<char> >
6 class priority_queue;

1 #include<ext/pb_ds/priority_queue.hpp>
2 using namespace __gnu_pbds;
3
4 typedef __gnu_pbds::priority_queue<LL, less<LL>, pairing_heap_tag> PQ;
5 __gnu_pbds::priority_queue<int, cmp, pairing_heap_tag>::point_iterator it;
6 PQ pq, pq2;
7
8 int main() {
9     auto it = pq.push(2);
10    pq.push(3);
11    assert(pq.top() == 3);
12    pq.modify(it, 4);
13    assert(pq.top() == 4);
14    pq2.push(5);
15    pq.join(pq2);
16    assert(pq.top() == 5);
17 }

```

- 树
- ov_tree_tag
- rb_tree_tag
- splay_tree_tag
- mapped: null_type 或 null_mapped_type (旧版本) 为空
- Node_Update 为 tree_order_statistics_node_update 时才可以 find_by_order & order_of_key
- find_by_order 找 order + 1 小的元素 (其实都是从 0 开始计数), 或者有 order 个元素比它小的 key
- order_of_key 有多少个比 r_key 小的元素
- join & split

```

1  template<typename Key, typename Mapped, typename Cmp_Fn = std::less<Key>,
2      typename Tag = rb_tree_tag,
3      template<typename Node_Citr, typename Node_Itr,
4          typename Cmp_Fn_, typename _Alloc_>
5          class Node_Update = null_node_update,
6      typename _Alloc = std::allocator<char> >
7  class tree
8
9  #include <ext/pb_ds/assoc_container.hpp>
10 using namespace __gnu_pbds;
11 using Tree = tree<int, null_type, less<int>, rb_tree_tag, tree_order_statistics_node_update>;
12 Tree t;

```

- hash table

```

1  #include<ext/pb_ds/assoc_container.hpp>
2  #include<ext/pb_ds/hash_policy.hpp>
3  using namespace __gnu_pbds;
4
5  gp_hash_table<int, int> mp;
6  cc_hash_table<int, int> mp;

```

Link-Cut Tree

- 图中相邻的结点在伸展树中不一定是父子关系
- 遇事不决 make_root
- 跑左右儿子的时候不要忘记 down

```

1  namespace lct {
2      extern struct P *const null;
3      const int M = N;
4      struct P {
5          P *fa, *ls, *rs;
6          int v, maxv;
7          bool rev;
8
9          bool has_fa() { return fa->ls == this || fa->rs == this; }
10         bool d() { return fa->ls == this; }
11         P& c(bool x) { return x ? ls : rs; }
12         void do_rev() {
13             if (this == null) return;
14             rev ^= 1;
15             swap(ls, rs);
16         }
17         P* up() {
18             maxv = max(v, max(ls->maxv, rs->maxv));
19             return this;
20         }
21         void down() {
22             if (rev) {
23                 rev = 0;
24                 ls->do_rev(); rs->do_rev();
25             }
26         }
27         void all_down() { if (has_fa()) fa->all_down(); down(); }
28     } *const null = new P{0, 0, 0, 0, 0, 0}, pool[M], *pit = pool;
29
30     void rot(P* o) {
31         bool dd = o->d();
32         P *f = o->fa, *t = o->c(!dd);
33         if (f->has_fa()) f->fa->c(f->d()) = o; o->fa = f->fa;
34         if (t != null) t->fa = f; f->c(dd) = t;
35         o->c(!dd) = f->up(); f->fa = o;
36     }
37     void splay(P* o) {
38         o->all_down();
39         while (o->has_fa()) {
40             if (o->fa->has_fa())
41                 rot(o->d() ^ o->fa->d() ? o : o->fa);

```

```

42         rot(o);
43     }
44     o->up();
45 }
46 void access(P* u, P* v = null) {
47     if (u == null) return;
48     splay(u); u->rs = v;
49     access(u->up()->fa, u);
50 }
51 void make_root(P* o) {
52     access(o); splay(o); o->do_rev();
53 }
54 void split(P* o, P* u) {
55     make_root(o); access(u); splay(u);
56 }
57 void link(P* u, P* v) {
58     make_root(u); u->fa = v;
59 }
60 void cut(P* u, P* v) {
61     split(u, v);
62     u->fa = v->ls = null; v->up();
63 }
64 bool adj(P* u, P* v) {
65     split(u, v);
66     return v->ls == u && u->ls == null && u->rs == null;
67 }
68 bool linked(P* u, P* v) {
69     split(u, v);
70     return u == v || u->fa != null;
71 }
72 P* findrt(P* o) {
73     access(o); splay(o);
74     while (o->ls != null) o = o->ls;
75     return o;
76 }
77 P* findfa(P* rt, P* u) {
78     split(rt, u);
79     u = u->ls;
80     while (u->rs != null) {
81         u = u->rs;
82         u->down();
83     }
84     return u;
85 }
86 }

```

● 维护子树大小

```

1 P* up() {
2     sz = ls->sz + rs->sz + _sz + 1;
3     return this;
4 }
5 void access(P* u, P* v = null) {
6     if (u == null) return;
7     splay(u);
8     u->_sz += u->rs->sz - v->sz;
9     u->rs = v;
10    access(u->up()->fa, u);
11 }
12 void link(P* u, P* v) {
13     split(u, v);
14     u->fa = v; v->_sz += u->sz;
15     v->up();
16 }

```

莫队

● $[l, r)$

```

1 while (l > q.l) mv(--l, 1);

```

```

2 while (r < q.r) mv(r++, 1);
3 while (l < q.l) mv(l++, -1);
4 while (r > q.r) mv(--r, -1);

    • 树上莫队
    • 注意初始状态 u = v = 1, flip(1)

1 struct Q {
2     int u, v, idx;
3     bool operator < (const Q& b) const {
4         const Q& a = *this;
5         return blk[a.u] < blk[b.u] || (blk[a.u] == blk[b.u] && in[a.v] < in[b.v]);
6     }
7 };
8
9 void dfs(int u = 1, int d = 0) {
10     static int S[maxn], sz = 0, blk_cnt = 0, clk = 0;
11     in[u] = clk++;
12     dep[u] = d;
13     int btm = sz;
14     for (int v: G[u]) {
15         if (v == fa[u]) continue;
16         fa[v] = u;
17         dfs(v, d + 1);
18         if (sz - btm >= B) {
19             while (sz > btm) blk[S[--sz]] = blk_cnt;
20             ++blk_cnt;
21         }
22     }
23     S[sz++] = u;
24     if (u == 1) while (sz) blk[S[--sz]] = blk_cnt - 1;
25 }
26
27 void flip(int k) {
28     dbg(k);
29     if (vis[k]) {
30         // ...
31     } else {
32         // ...
33     }
34     vis[k] ^= 1;
35 }
36
37 void go(int& k) {
38     if (bug == -1) {
39         if (vis[k] && !vis[fa[k]]) bug = k;
40         if (!vis[k] && vis[fa[k]]) bug = fa[k];
41     }
42     flip(k);
43     k = fa[k];
44 }
45
46 void mv(int a, int b) {
47     bug = -1;
48     if (vis[b]) bug = b;
49     if (dep[a] < dep[b]) swap(a, b);
50     while (dep[a] > dep[b]) go(a);
51     while (a != b) {
52         go(a); go(b);
53     }
54     go(a); go(bug);
55 }
56
57 for (Q& q: query) {
58     mv(u, q.u); u = q.u;
59     mv(v, q.v); v = q.v;
60     ans[q.idx] = Ans;
61 }

```

数学

矩阵运算

```
1 struct Mat {
2     static const LL M = 2;
3     LL v[M][M];
4     Mat() { memset(v, 0, sizeof v); }
5     void eye() { FOR (i, 0, M) v[i][i] = 1; }
6     LL* operator [] (LL x) { return v[x]; }
7     const LL* operator [] (LL x) const { return v[x]; }
8     Mat operator * (const Mat& B) {
9         const Mat& A = *this;
10        Mat ret;
11        FOR (i, 0, M)
12            FOR (j, 0, M)
13                FOR (k, 0, M)
14                    ret[i][j] = (ret[i][j] + A[i][k] * B[k][j]) % MOD;
15        return ret;
16    }
17    Mat pow(LL n) const {
18        Mat A = *this, ret; ret.eye();
19        for (; n; n >>= 1, A = A * A)
20            if (n & 1) ret = ret * A;
21        return ret;
22    }
23    Mat operator + (const Mat& B) {
24        const Mat& A = *this;
25        Mat ret;
26        FOR (i, 0, M)
27            FOR (j, 0, M)
28                ret[i][j] = (A[i][j] + B[i][j]) % MOD;
29        return ret;
30    }
31    void prt() const {
32        FOR (i, 0, M)
33            FOR (j, 0, M)
34                printf("%lld%c", (*this)[i][j], j == M - 1 ? '\n' : ' ');
35    }
36 };
```

筛

● 线性筛

```
1 const LL p_max = 1E6 + 100;
2 LL pr[p_max], p_sz;
3 void get_prime() {
4     static bool vis[p_max];
5     FOR (i, 2, p_max) {
6         if (!vis[i]) pr[p_sz++] = i;
7         FOR (j, 0, p_sz) {
8             if (pr[j] * i >= p_max) break;
9             vis[pr[j] * i] = 1;
10            if (i % pr[j] == 0) break;
11        }
12    }
13 }
```

● 线性筛 + 欧拉函数

```
1 const LL p_max = 1E5 + 100;
2 LL phi[p_max] = {-1, 1};
3 void get_phi() {
4     static bool vis[p_max];
5     static LL prime[p_max], p_sz, d;
6     FOR (i, 2, p_max) {
7         if (!vis[i]) {
```

```

8         prime[p_sz++] = i;
9         phi[i] = i - 1;
10    }
11    for (LL j = 0; j < p_sz && (d = i * prime[j]) < p_max; ++j) {
12        vis[d] = 1;
13        if (i % prime[j] == 0) {
14            phi[d] = phi[i] * prime[j];
15            break;
16        }
17        else phi[d] = phi[i] * (prime[j] - 1);
18    }
19 }
20 }

```

● 线性筛 + 莫比乌斯函数

```

1  const LL p_max = 1E5 + 100;
2  LL mu[p_max] = {-1, 1};
3  void get_mu() {
4      static bool vis[p_max];
5      static LL prime[p_max], p_sz, d;
6      mu[1] = 1;
7      FOR (i, 2, p_max) {
8          if (!vis[i]) {
9              prime[p_sz++] = i;
10             mu[i] = -1;
11         }
12         for (LL j = 0; j < p_sz && (d = i * prime[j]) < p_max; ++j) {
13             vis[d] = 1;
14             if (i % prime[j] == 0) {
15                 mu[d] = 0;
16                 break;
17             }
18             else mu[d] = -mu[i];
19         }
20     }
21 }

```

亚线性筛

min_25

```

1  namespace min25 {
2      const int M = 1E6 + 100;
3      LL B, N;
4
5      // g(x)
6      inline LL pg(LL x) { return 1; }
7      inline LL ph(LL x) { return x % MOD; }
8      // Sum[g(i), {x, 2, x}]
9      inline LL psg(LL x) { return x % MOD - 1; }
10     inline LL psh(LL x) {
11         static LL inv2 = (MOD + 1) / 2;
12         x = x % MOD;
13         return x * (x + 1) % MOD * inv2 % MOD - 1;
14     }
15     // f(pp=p^k)
16     inline LL fpk(LL p, LL e, LL pp) { return (pp - pp / p) % MOD; }
17     // f(p) = fgh(g(p), h(p))
18     inline LL fgh(LL g, LL h) { return h - g; }
19
20     LL pr[M], pc, sg[M], sh[M];
21     void get_prime(LL n) {
22         static bool vis[M]; pc = 0;
23         FOR (i, 2, n + 1) {
24             if (!vis[i]) {
25                 pr[pc++] = i;
26                 sg[pc] = (sg[pc - 1] + pg(i)) % MOD;
27                 sh[pc] = (sh[pc - 1] + ph(i)) % MOD;

```



```

28     }
29     FOR (j, 0, pc) {
30         if (pr[j] * i > n) break;
31         vis[pr[j] * i] = 1;
32         if (i % pr[j] == 0) break;
33     }
34 }
35 }
36
37 LL w[M];
38 LL id1[M], id2[M], h[M], g[M];
39 inline LL id(LL x) { return x <= B ? id1[x] : id2[N / x]; }
40
41 LL go(LL x, LL k) {
42     if (x <= 1 || (k >= 0 && pr[k] > x)) return 0;
43     LL t = id(x);
44     LL ans = fgh((g[t] - sg[k + 1]), (h[t] - sh[k + 1]));
45     FOR (i, k + 1, pc) {
46         LL p = pr[i];
47         if (p * p > x) break;
48         ans -= fgh(pg(p), ph(p));
49         for (LL pp = p, e = 1; pp <= x; ++e, pp = pp * p)
50             ans += fpk(p, e, pp) * (1 + go(x / pp, i)) % MOD;
51     }
52     return ans % MOD;
53 }
54
55 LL solve(LL _N) {
56     N = _N;
57     B = sqrt(N + 0.5);
58     get_prime(B);
59     int sz = 0;
60     for (LL l = 1, v, r; l <= N; l = r + 1) {
61         v = N / l; r = N / v;
62         w[sz] = v; g[sz] = psg(v); h[sz] = psh(v);
63         if (v <= B) id1[v] = sz; else id2[r] = sz;
64         sz++;
65     }
66     FOR (k, 0, pc) {
67         LL p = pr[k];
68         FOR (i, 0, sz) {
69             LL v = w[i]; if (p * p > v) break;
70             LL t = id(v / p);
71             g[i] = (g[i] - (g[t] - sg[k]) * pg(p)) % MOD;
72             h[i] = (h[i] - (h[t] - sh[k]) * ph(p)) % MOD;
73         }
74     }
75     return (go(N, -1) % MOD + MOD + 1) % MOD;
76 }
77 }

```

杜教筛

```

1 namespace dujiao {
2     const int M = 5E6;
3     LL f[M] = {0, 1};
4     void init() {
5         static bool vis[M];
6         static LL pr[M], p_sz, d;
7         FOR (i, 2, M) {
8             if (!vis[i]) { pr[p_sz++] = i; f[i] = -1; }
9             FOR (j, 0, p_sz) {
10                 if ((d = pr[j] * i) >= M) break;
11                 vis[d] = 1;
12                 if (i % pr[j] == 0) {
13                     f[d] = 0;
14                     break;
15                 } else f[d] = -f[i];
16             }
17         }
18     }
19 }

```

```

18     FOR (i, 2, M) f[i] += f[i - 1];
19 }
20 inline LL s_fg(LL n) { return 1; }
21 inline LL s_g(LL n) { return n; }
22
23 LL N, rd[M];
24 bool vis[M];
25 LL go(LL n) {
26     if (n < M) return f[n];
27     LL id = N / n;
28     if (vis[id]) return rd[id];
29     vis[id] = true;
30     LL& ret = rd[id] = s_fg(n);
31     for (LL l = 2, v, r; l <= n; l = r + 1) {
32         v = n / l; r = n / v;
33         ret -= (s_g(r) - s_g(l - 1)) * go(v);
34     }
35     return ret;
36 }
37 LL solve(LL n) {
38     N = n;
39     memset(vis, 0, sizeof vis);
40     return go(n);
41 }
42 }

```

素数测试

- 前置：快速乘、快速幂
- int 范围内只需检查 2, 7, 61
- long long 范围 2, 325, 9375, 28178, 450775, 9780504, 1795265022
- 3E15 内 2, 2570940, 880937, 610386380, 4130785767
- 4E13 内 2, 2570940, 211991001, 3749873356
- <http://miller-rabin.appspot.com/>

```

1 bool checkQ(LL a, LL n) {
2     if (n == 2 || a >= n) return 1;
3     if (n == 1 || !(n & 1)) return 0;
4     LL d = n - 1;
5     while (!(d & 1)) d >>= 1;
6     LL t = bin(a, d, n); // 不一定需要快速乘
7     while (d != n - 1 && t != 1 && t != n - 1) {
8         t = mul(t, t, n);
9         d <<= 1;
10    }
11    return t == n - 1 || d & 1;
12 }
13
14 bool primeQ(LL n) {
15     static vector<LL> t = {2, 325, 9375, 28178, 450775, 9780504, 1795265022};
16     if (n <= 1) return false;
17     for (LL k: t) if (!checkQ(k, n)) return false;
18     return true;
19 }

```

线性递推

```

1 // k 为 m 最高次数 且 a[m] == 1
2 namespace BerlekampMassey {
3     inline void up(LL& a, LL b) { (a += b) %= MOD; }
4
5     V mul(const V& a, const V& b, const V& m, int k) {
6         V r; r.resize(2 * k - 1);
7         FOR (i, 0, k)
8             FOR (j, 0, k)
9                 up(r[i + j], a[i] * b[j]);

```

```

10     FORD (i, k - 2, - 1) {
11         FOR (j, 0, k)
12             up(r[i + j], r[i + k] * m[j]);
13         r.pop_back();
14     }
15     return r;
16 }
17
18 V pow(LL n, const V& m) {
19     int k = (int)m.size() - 1; assert(m[k] == -1 || m[k] == MOD - 1);
20     V r(k), x(k); r[0] = x[1] = 1;
21     for (; n; n >>= 1, x = mul(x, x, m, k))
22         if (n & 1) r = mul(x, r, m, k);
23     return r;
24 }
25
26 LL go(const V& a, const V& x, LL n) {
27     // a: (-1, a1, a2, ..., ak).reverse
28     // x: x1, x2, ..., xk
29     // x[n] = sum[a[i]*x[n-i], {i, 1, k}]
30     int k = (int)a.size() - 1;
31     if (n <= k) return x[n - 1];
32     V r = pow(n - 1, a);
33     LL ans = 0;
34     FOR (i, 0, k)
35         up(ans, r[i] * x[i]);
36     return ans;
37 }
38
39 V BM(const V& x) {
40     V a = {-1}, b = {233};
41     FOR (i, 1, x.size()) {
42         b.push_back(0);
43         LL d = 0, la = a.size(), lb = b.size();
44         FOR (j, 0, la) up(d, a[j] * x[i - la + 1 + j]);
45         if (d == 0) continue;
46         V t; for (auto& v: b) t.push_back(d * v % MOD);
47         FOR (j, 0, a.size()) up(t[lb - 1 - j], a[la - 1 - j]);
48         if (lb > la) {
49             b = a;
50             LL inv = -get_inv(d, MOD);
51             for (auto& v: b) v = v * inv % MOD;
52         }
53         a.swap(t);
54     }
55     for (auto& v: a) up(v, MOD);
56     return a;
57 }
58 }

```

扩展欧几里得

- 求 $ax + by = \gcd(a, b)$ 的一组解
- 如果 a 和 b 互素, 那么 x 是 a 在模 b 下的逆元
- 注意 x 和 y 可能是负数

```

1 LL ex_gcd(LL a, LL b, LL &x, LL &y) {
2     if (b == 0) {
3         x = 1;
4         y = 0;
5         return a;
6     }
7     LL ret = ex_gcd(b, a % b, y, x);
8     y -= a / b * x;
9     return ret;
10 }

```

类欧几里得

- $m = \lfloor \frac{an+b}{c} \rfloor$.
- $f(a, b, c, n) = \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor$: 当 $a \geq c$ or $b \geq c$ 时, $f(a, b, c, n) = (\frac{a}{c})n(n+1)/2 + (\frac{b}{c})(n+1) + f(a \bmod c, b \bmod c, c, n)$; 否则 $f(a, b, c, n) = nm - f(c, c-b-1, a, m-1)$ 。
- $g(a, b, c, n) = \sum_{i=0}^n i \lfloor \frac{ai+b}{c} \rfloor$: 当 $a \geq c$ or $b \geq c$ 时, $g(a, b, c, n) = (\frac{a}{c})n(n+1)(2n+1)/6 + (\frac{b}{c})n(n+1)/2 + g(a \bmod c, b \bmod c, c, n)$; 否则 $g(a, b, c, n) = \frac{1}{2}(n(n+1)m - f(c, c-b-1, a, m-1) - h(c, c-b-1, a, m-1))$ 。
- $h(a, b, c, n) = \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor^2$: 当 $a \geq c$ or $b \geq c$ 时, $h(a, b, c, n) = (\frac{a}{c})^2 n(n+1)(2n+1)/6 + (\frac{b}{c})^2 (n+1) + (\frac{a}{c})(\frac{b}{c})n(n+1) + h(a \bmod c, b \bmod c, c, n) + 2(\frac{a}{c})g(a \bmod c, b \bmod c, c, n) + 2(\frac{b}{c})f(a \bmod c, b \bmod c, c, n)$; 否则 $h(a, b, c, n) = nm(m+1) - 2g(c, c-b-1, a, m-1) - 2f(c, c-b-1, a, m-1) - f(a, b, c, n)$ 。

逆元

- $ax \equiv 1 \pmod{p}$
- 如果 p 不是素数, 使用拓展欧几里得
- 模数是素数, 求一个数的逆元
- 前置模板: 快速幂

```
1 inline LL get_inv(LL x, LL p) { return bin(x, p - 2, p); }
```

- 预处理

$$1 - n$$

的逆元

```
1 LL inv[N] = {-1, 1};
2 void inv_init(LL n, LL p) {
3     inv[1] = 1;
4     FOR (i, 2, n)
5         inv[i] = (p - p / i) * inv[p % i] % p;
6 }
```

- 预处理阶乘及其逆元

```
1 LL invf[M], fac[M] = {1};
2 void fac_inv_init(LL n, LL p) {
3     FOR (i, 1, n)
4         fac[i] = i * fac[i - 1] % p;
5     invf[n - 1] = bin(fac[n - 1], p - 2, p);
6     FORD (i, n - 2, -1)
7         invf[i] = invf[i + 1] * (i + 1) % p;
8 }
```

组合数

- 如果数较小, 模较大时使用逆元
- 前置模板: 逆元-预处理阶乘及其逆元

```
1 inline LL C(LL n, LL m) { // n >= m >= 0
2     return n < m || m < 0 ? 0 : fac[n] * invf[m] % MOD * invf[n - m] % MOD;
3 }
```

- 如果模数较小, 数字较大, 使用 Lucas 定理
- 前置模板可选 1: 求组合数 (如果使用阶乘逆元, 需 `fac_inv_init(MOD, MOD);`)
- 前置模板可选 2: 模数不固定下使用, 无法单独使用。

```
1 LL C(LL n, LL m) { // m >= n >= 0
2     if (m - n < n) n = m - n;
3     if (n < 0) return 0;
4     LL ret = 1;
5     FOR (i, 1, n + 1)
6         ret = ret * (m - n + i) % MOD * bin(i, MOD - 2, MOD) % MOD;
```

```

7     return ret;
8 }

1 LL Lucas(LL n, LL m) { // m >= n >= 0
2     return m ? C(n % MOD, m % MOD) * Lucas(n / MOD, m / MOD) % MOD : 1;
3 }

```

- 组合数预处理

```

1 LL C[M][M];
2 void init_C(int n) {
3     FOR (i, 0, n) {
4         C[i][0] = C[i][i] = 1;
5         FOR (j, 1, i)
6             C[i][j] = (C[i - 1][j] + C[i - 1][j - 1]) % MOD;
7     }
8 }

```

第二类斯特灵数

```

1 S[0][0] = 1;
2 FOR (i, 1, N)
3     FOR (j, 1, i + 1) S[i][j] = (S[i - 1][j - 1] + j * S[i - 1][j]) % MOD;

```

FFT & NTT & FWT

- NTT
- 前置：快速幂

```

1 LL wn[N << 2], rev[N << 2];
2 int NTT_init(int n_) {
3     int step = 0; int n = 1;
4     for (; n < n_; n <= 1) ++step;
5     FOR (i, 1, n)
6         rev[i] = (rev[i >> 1] >> 1) | ((i & 1) << (step - 1));
7     int g = bin(G, (MOD - 1) / n, MOD);
8     wn[0] = 1;
9     for (int i = 1; i <= n; ++i)
10         wn[i] = wn[i - 1] * g % MOD;
11     return n;
12 }
13
14 void NTT(LL a[], int n, int f) {
15     FOR (i, 0, n) if (i < rev[i])
16         std::swap(a[i], a[rev[i]]);
17     for (int k = 1; k < n; k <= 1) {
18         for (int i = 0; i < n; i += (k <= 1)) {
19             int t = n / (k <= 1);
20             FOR (j, 0, k) {
21                 LL w = f == 1 ? wn[t * j] : wn[n - t * j];
22                 LL x = a[i + j];
23                 LL y = a[i + j + k] * w % MOD;
24                 a[i + j] = (x + y) % MOD;
25                 a[i + j + k] = (x - y + MOD) % MOD;
26             }
27         }
28     }
29     if (f == -1) {
30         LL ninv = get_inv(n, MOD);
31         FOR (i, 0, n)
32             a[i] = a[i] * ninv % MOD;
33     }
34 }

```

- FFT
- n 需补成 2 的幂 (n 必须超过 a 和 b 的最高指数之和)

```

1  typedef double LD;
2  const LD PI = acos(-1);
3  struct C {
4      LD r, i;
5      C(LD r = 0, LD i = 0): r(r), i(i) {}
6  };
7  C operator + (const C& a, const C& b) {
8      return C(a.r + b.r, a.i + b.i);
9  }
10 C operator - (const C& a, const C& b) {
11     return C(a.r - b.r, a.i - b.i);
12 }
13 C operator * (const C& a, const C& b) {
14     return C(a.r * b.r - a.i * b.i, a.r * b.i + a.i * b.r);
15 }
16
17 void FFT(C x[], int n, int p) {
18     for (int i = 0, t = 0; i < n; ++i) {
19         if (i > t) swap(x[i], x[t]);
20         for (int j = n >> 1; (t ^= j) < j; j >>= 1);
21     }
22     for (int h = 2; h <= n; h <= 1) {
23         C wn(cos(p * 2 * PI / h), sin(p * 2 * PI / h));
24         for (int i = 0; i < n; i += h) {
25             C w(1, 0), u;
26             for (int j = i, k = h >> 1; j < i + k; ++j) {
27                 u = x[j + k] * w;
28                 x[j + k] = x[j] - u;
29                 x[j] = x[j] + u;
30                 w = w * wn;
31             }
32         }
33     }
34     if (p == -1)
35         FOR (i, 0, n)
36             x[i].r /= n;
37 }
38
39 void conv(C a[], C b[], int n) {
40     FFT(a, n, 1);
41     FFT(b, n, 1);
42     FOR (i, 0, n)
43         a[i] = a[i] * b[i];
44     FFT(a, n, -1);
45 }

```

• FWT

```

1  template<typename T>
2  void fwt(LL a[], int n, T f) {
3      for (int d = 1; d < n; d *= 2)
4          for (int i = 0, t = d * 2; i < n; i += t)
5              FOR (j, 0, d)
6                  f(a[i + j], a[i + j + d]);
7  }
8
9  void AND(LL& a, LL& b) { a += b; }
10 void OR(LL& a, LL& b) { b += a; }
11 void XOR (LL& a, LL& b) {
12     LL x = a, y = b;
13     a = (x + y) % MOD;
14     b = (x - y + MOD) % MOD;
15 }

```

simpson 自适应积分

```

1  LD simpson(LD l, LD r) {
2      LD c = (l + r) / 2;
3      return (f(l) + 4 * f(c) + f(r)) * (r - l) / 6;
4  }

```

```

5
6 LD asr(LD l, LD r, LD eps, LD S) {
7     LD m = (l + r) / 2;
8     LD L = simpson(l, m), R = simpson(m, r);
9     if (fabs(L + R - S) < 15 * eps) return L + R + (L + R - S) / 15;
10    return asr(l, m, eps / 2, L) + asr(m, r, eps / 2, R);
11 }
12
13 LD asr(LD l, LD r, LD eps) { return asr(l, r, eps, simpson(l, r)); }

```

● FWT

```

1 template<typename T>
2 void fwt(LL a[], int n, T f) {
3     for (int d = 1; d < n; d *= 2)
4         for (int i = 0, t = d * 2; i < n; i += t)
5             FOR (j, 0, d)
6                 f(a[i + j], a[i + j + d]);
7 }
8
9 auto f = [](LL& a, LL& b) { // xor
10     LL x = a, y = b;
11     a = (x + y) % MOD;
12     b = (x - y + MOD) % MOD;
13 };

```

快速乘

```

1 LL mul(LL a, LL b, LL m) {
2     LL ret = 0;
3     while (b) {
4         if (b & 1) {
5             ret += a;
6             if (ret >= m) ret -= m;
7         }
8         a += a;
9         if (a >= m) a -= m;
10        b >>= 1;
11    }
12    return ret;
13 }

```

● $O(1)$

```

1 LL mul(LL u, LL v, LL p) {
2     return (u * v - LL((long double) u * v / p) * p + p) % p;
3 }

```

快速幂

- 如果模数是素数，则可在函数体内加上 $n \% = \text{MOD} - 1$ ；（费马小定理）。

```

1 LL bin(LL x, LL n, LL MOD) {
2     LL ret = MOD != 1;
3     for (x %= MOD; n; n >>= 1, x = x * x % MOD)
4         if (n & 1) ret = ret * x % MOD;
5     return ret;
6 }

```

● 防爆 LL

- 前置模板：快速乘

```

1 LL bin(LL x, LL n, LL MOD) {
2     LL ret = MOD != 1;
3     for (x %= MOD; n; n >>= 1, x = mul(x, x, MOD))
4         if (n & 1) ret = mul(ret, x, MOD);
5     return ret;
6 }

```

高斯消元

- n - 方程个数, m - 变量个数, a 是 $n * (m + 1)$ 的增广矩阵, $free$ 是否为自由变量
- 返回自由变量个数, -1 无解, -2 无整数解
- 浮点数版本

```
1  typedef double LD;
2  const LD eps = 1E-10;
3  const int maxn = 2000 + 10;
4
5  int n, m;
6  LD a[maxn][maxn], x[maxn];
7  bool free_x[maxn];
8
9  inline int sgn(LD x) { return (x > eps) - (x < -eps); }
10
11
12 int guass(LD a[maxn][maxn], int n, int m) {
13     memset(free_x, 1, sizeof free_x); memset(x, 0, sizeof x);
14     int r = 0, c = 0;
15     while (r < n && c < m) {
16         int m_r = r;
17         FOR (i, r + 1, n)
18             if (fabs(a[i][c]) > fabs(a[m_r][c])) m_r = i;
19         if (m_r != r)
20             FOR (j, c, m + 1)
21                 swap(a[r][j], a[m_r][j]);
22         if (!sgn(a[r][c])) {
23             a[r][c] = 0;
24             ++c;
25             continue;
26         }
27         FOR (i, r + 1, n)
28             if (a[i][c]) {
29                 LD t = a[i][c] / a[r][c];
30                 FOR (j, c, m + 1) a[i][j] -= a[r][j] * t;
31             }
32         ++r; ++c;
33         // FOR (i, 0, n)
34         //     FOR (j, 0, m + 1)
35         //         printf("%.2f%c", a[i][j], j == _j - 1 ? '\n' : ' '); puts("");
36     }
37     FOR (i, r, n)
38         if (sgn(a[i][m])) return -1;
39     if (r < m) {
40         FORD (i, r - 1, -1) {
41             int f_cnt = 0, k = -1;
42             FOR (j, 0, m)
43                 if (sgn(a[i][j]) && free_x[j]) {
44                     ++f_cnt;
45                     k = j;
46                 }
47             if (f_cnt > 0) continue;
48             LD s = a[i][m];
49             FOR (j, 0, m)
50                 if (j != k) s -= a[i][j] * x[j];
51             x[k] = s / a[i][k];
52             free_x[k] = 0;
53         }
54         return m - r;
55     }
56     FORD (i, m - 1, -1) {
57         LD s = a[i][m];
58         FOR (j, i + 1, m)
59             s -= a[i][j] * x[j];
60         x[i] = s / a[i][i];
61     }
62     return 0;
63 }
```


- 数据

```

3 4
1 1 -2 2
2 -3 5 1
4 -1 1 5
5 0 -1 7
// many

3 4
1 1 -2 2
2 -3 5 1
4 -1 -1 5
5 0 -1 0 2
// no

3 4
1 1 -2 2
2 -3 5 1
4 -1 1 5
5 0 1 0 7
// one

```

质因数分解

- 前置模板：素数筛

- 带指数

```

1 LL factor[30], f_sz, factor_exp[30];
2 void get_factor(LL x) {
3     f_sz = 0;
4     LL t = sqrt(x + 0.5);
5     for (LL i = 0; pr[i] <= t; ++i)
6         if (x % pr[i] == 0) {
7             factor_exp[f_sz] = 0;
8             while (x % pr[i] == 0) {
9                 x /= pr[i];
10                ++factor_exp[f_sz];
11            }
12            factor[f_sz++] = pr[i];
13        }
14    if (x > 1) {
15        factor_exp[f_sz] = 1;
16        factor[f_sz++] = x;
17    }
18 }

```

- 不带指数

```

1 LL factor[30], f_sz;
2 void get_factor(LL x) {
3     f_sz = 0;
4     LL t = sqrt(x + 0.5);
5     for (LL i = 0; pr[i] <= t; ++i)
6         if (x % pr[i] == 0) {
7             factor[f_sz++] = pr[i];
8             while (x % pr[i] == 0) x /= pr[i];
9         }
10    if (x > 1) factor[f_sz++] = x;
11 }

```

原根

- 前置模板：素数筛，快速幂，分解质因数
- 要求 p 为质数

```
1 LL find_smallest_primitive_root(LL p) {
2     get_factor(p - 1);
3     FOR (i, 2, p) {
4         bool flag = true;
5         FOR (j, 0, f_sz)
6             if (bin(i, (p - 1) / factor[j], p) == 1) {
7                 flag = false;
8                 break;
9             }
10        if (flag) return i;
11    }
12    assert(0); return -1;
13 }
```

公式

- 当 $x \geq \phi(p)$ 时有 $a^x \equiv a^{x \bmod \phi(p) + \phi(p)} \pmod{p}$

斐波那契数列性质

- $F_{a+b} = F_{a-1} \cdot F_b + F_a \cdot F_{b+1}$

常见生成函数

- $(1 + ax)^n = \sum_{k=0}^n \binom{n}{k} a^k x^k$
- $\frac{1 - x^{r+1}}{1 - x} = \sum_{k=0}^n x^k$
- $\frac{1}{1 - ax} = \sum_{k=0}^{\infty} a^k x^k$
- $\frac{1}{(1 - x)^2} = \sum_{k=0}^{\infty} (k + 1) x^k$
- $\frac{1}{(1 - x)^n} = \sum_{k=0}^{\infty} \binom{n+k-1}{k} x^k$
- $e^x = \sum_{k=0}^{\infty} \frac{x^k}{k!}$
- $\ln(1 + x) = \sum_{k=0}^{\infty} \frac{(-1)^{k+1}}{k} x^k$

佩尔方程

若一个丢番图方程具有以下的形式： $x^2 - ny^2 = 1$ 。且 n 为正整数，则称此二元二次不定方程为**佩尔方程**。

若 n 是完全平方数，则这个方程式只有平凡解 $(\pm 1, 0)$ （实际上对任意的 n ， $(\pm 1, 0)$ 都是解）。对于其余情况，拉格朗日证明了佩尔方程总有非平凡解。而这些解可由 \sqrt{n} 的连分数求出。

$$x = [a_0; a_1, a_2, a_3] = x = a_0 + \frac{1}{a_1 + \frac{1}{a_2 + \frac{1}{a_3 + \ddots}}}$$

设 $\frac{p_i}{q_i}$ 是 \sqrt{n} 的连分数表示： $[a_0; a_1, a_2, a_3, \dots]$ 的渐近分数列，由连分数理论知存在 i 使得 (p_i, q_i) 为佩尔方程的解。取其中最小的 i ，将对应的 (p_i, q_i) 称为佩尔方程的基本解，或最小解，记作 (x_1, y_1) ，则所有的解 (x_i, y_i) 可表示成如下形式： $x_i + y_i \sqrt{n} = (x_1 + y_1 \sqrt{n})^i$ 。或者由以下的递回关系式得到：

$$x_{i+1} = x_1 x_i + n y_1 y_i, y_{i+1} = x_1 y_i + y_1 x_i.$$

但是：佩尔方程千万不要去推（虽然推起来很有趣，但结果不一定好看，会是两个式子）。记住佩尔方程结果的形式通常是 $a_n = k a_{n-1} - a_{n-2}$ (a_{n-2} 前的系数通常是 -1)。暴力 / 凑出两个基础解之后加上一个 0，容易解出 k 并验证。一般的话还可以找规律得到，但本题由于是两个序列并在一起，规律并不容易发现。

Burnside & Polya

- $|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|$

注： X^g 是 g 下的不动点数量，也就是说有多少种东西用 g 作用之后可以保持不变。

- $|Y^X/G| = \frac{1}{|G|} \sum_{g \in G} m^{c(g)}$

注：用 m 种颜色染色，然后对于某一种置换 g ，有 $c(g)$ 个置换环，为了保证置换后颜色仍然相同，每个置换环必须染成同色。

二次剩余

```

1  LL q1, q2;
2
3  LL w;
4  struct Point { //x + y*sqrt(w)
5      LL x;
6      LL y;
7  };
8
9  LL mod(LL a, LL p) {
10     a %= p;
11     if (a < 0) {
12         a += p;
13     }
14     return a;
15 }
16
17 LL quick_mul(LL x, LL y, LL p) {
18     x = mod(x, p);
19     y = mod(y, p);
20     LL ans = 0;
21     while (y > 0) {
22         if (y & 1) {
23             ans = ans + x;
24             if (ans >= p) ans -= p;
25         }
26         x = x + x;
27         if (x >= p) x -= p;
28         y >>= 1;
29     }
30     return ans;
31 }
32
33 Point point_mul(Point a, Point b, LL p) {
34     Point res;
35     res.x = quick_mul(a.x, b.x, p);
36     res.x += quick_mul(w, quick_mul(a.y, b.y, p), p);
37     res.x = mod(res.x, p);
38     res.y = quick_mul(a.x, b.y, p);
39     res.y += quick_mul(a.y, b.x, p);
40     res.y = mod(res.y, p);
41     return res;
42 }
43
44 Point power(Point a, LL b, LL p) {
45     Point res;
46     res.x = 1;
47     res.y = 0;
48     while (b) {
49         if (b & 1) {
50             res = point_mul(res, a, p);

```

```

51     }
52     a = point_mul(a, a, p);
53     b = b >> 1;
54 }
55
56 return res;
57 }
58
59 LL quick_power(LL a, LL b, LL p) //(a^b)%p
60 {
61     LL res = 1;
62     while (b) {
63         if (b & 1) {
64             res = quick_mul(res, a, p);
65         }
66
67         a = quick_mul(a, a, p);
68         b = b >> 1;
69     }
70
71     return res;
72 }
73
74
75 LL Legendre(LL a, LL p) { // a^((p-1)/2)
76     return quick_power(a, (p - 1) >> 1, p);
77 }
78
79 LL equation_solve(LL b, LL p) { //求解 x^2=b(%p) 方程解
80     if ((Legendre(b, p) + 1) % p == 0) {
81         return -1; //表示没有解
82     }
83
84     LL a, t;
85     while (true) {
86         a = rand() % p;
87         t = quick_mul(a, a, p) - b;
88         t = mod(t, p);
89         if ((Legendre(t, p) + 1) % p == 0) {
90             break;
91         }
92     }
93
94     w = t;
95     Point temp, res;
96     temp.x = a;
97     temp.y = 1;
98     res = power(temp, (p + 1) >> 1, p);
99     return res.x;
100 }
101
102 bool getInv(LL p) {
103     LL x = equation_solve(p - 3, p);
104     if (x == -1) return false;
105     LL inv2 = (p + 1) >> 1;
106     q1 = quick_mul(p - 1 - x, inv2, p), q2 = quick_mul(x - 1, inv2, p);
107     return true;
108 }

```

中国剩余定理

- 无解返回 -1
- 前置模板：拓展欧几里得

```

1 LL CRT(LL *m, LL *r, LL n) {
2     if (!n) return 0;
3     LL M = m[0], R = r[0], x, y, d;
4     FOR (i, 1, n) {
5         d = ex_gcd(M, m[i], x, y);

```

```

6         if ((r[i] - R) % d) return -1;
7         x = (r[i] - R) / d * x % (m[i] / d);
8         R += x * M;
9         M = M / d * m[i];
10        R %= M;
11    }
12    return R >= 0 ? R : R + M;
13 }

```

伯努利数和等幂求和

- 预处理逆元
- 预处理组合数
- $\sum_{i=0}^n i^k = \frac{1}{k+1} \sum_{i=0}^k \binom{k+1}{i} B_{k+1-i} (n+1)^i$.
- 也可以 $\sum_{i=0}^n i^k = \frac{1}{k+1} \sum_{i=0}^k \binom{k+1}{i} B_{k+1-i}^+ n^i$ 。区别在于 $B_1^+ = 1/2$ 。(心态崩了)

```

1 namespace Bernoulli {
2     const int M = 100;
3     LL inv[M] = {-1, 1};
4     void inv_init(LL n, LL p) {
5         FOR (i, 2, n)
6             inv[i] = (p - p / i) * inv[p % i] % p;
7     }
8
9     LL C[M][M];
10    void init_C(int n) {
11        FOR (i, 0, n) {
12            C[i][0] = C[i][i] = 1;
13            FOR (j, 1, i)
14                C[i][j] = (C[i-1][j] + C[i-1][j-1]) % MOD;
15        }
16    }
17
18    LL B[M] = {1};
19    void init() {
20        inv_init(M, MOD);
21        init_C(M);
22        FOR (i, 1, M-1) {
23            LL& s = B[i] = 0;
24            FOR (j, 0, i)
25                s += C[i+1][j] * B[j] % MOD;
26            s = (s % MOD * -inv[i+1] % MOD + MOD) % MOD;
27        }
28    }
29
30    LL p[M] = {1};
31    LL go(LL n, LL k) {
32        n %= MOD;
33        if (k == 0) return n;
34        FOR (i, 1, k+2)
35            p[i] = p[i-1] * (n+1) % MOD;
36        LL ret = 0;
37        FOR (i, 1, k+2)
38            ret += C[k+1][i] * B[k+1-i] % MOD * p[i] % MOD;
39        ret = ret % MOD * inv[k+1] % MOD;
40        return ret;
41    }
42 }

```

单纯形

- 要求有基本解，也就是 x 为零向量可行
- v 要初始化为 0, n 表示向量长度, m 表示约束个数

```

1 // min{ b x } / max{ c x }
2 // A x >= c / A x <= b

```

```

3 // x >= 0
4 namespace lp {
5     int n, m;
6     double a[M][N], b[M], c[N], v;
7
8     void pivot(int l, int e) {
9         b[l] /= a[l][e];
10        FOR (j, 0, n) if (j != e) a[l][j] /= a[l][e];
11        a[l][e] = 1 / a[l][e];
12
13        FOR (i, 0, m)
14            if (i != l && fabs(a[i][e]) > 0) {
15                b[i] -= a[i][e] * b[l];
16                FOR (j, 0, n)
17                    if (j != e) a[i][j] -= a[i][e] * a[l][j];
18                a[i][e] = -a[i][e] * a[l][e];
19            }
20        v += c[e] * b[l];
21        FOR (j, 0, n) if (j != e) c[j] -= c[e] * a[l][j];
22        c[e] = -c[e] * a[l][e];
23    }
24    double simplex() {
25        while (1) {
26            v = 0;
27            int e = -1, l = -1;
28            FOR (i, 0, n) if (c[i] > eps) { e = i; break; }
29            if (e == -1) return v;
30            double t = INF;
31            FOR (i, 0, m)
32                if (a[i][e] > eps && t > b[i] / a[i][e]) {
33                    t = b[i] / a[i][e];
34                    l = i;
35                }
36            if (l == -1) return INF;
37            pivot(l, e);
38        }
39    }
40 }

```

图论

LCA

- 倍增

```

1 void dfs(int u, int fa) {
2     pa[u][0] = fa; dep[u] = dep[fa] + 1;
3     FOR (i, 1, SP) pa[u][i] = pa[pa[u][i - 1]][i - 1];
4     for (int& v: G[u]) {
5         if (v == fa) continue;
6         dfs(v, u);
7     }
8 }
9
10 int lca(int u, int v) {
11     if (dep[u] < dep[v]) swap(u, v);
12     int t = dep[u] - dep[v];
13     FOR (i, 0, SP) if (t & (1 << i)) u = pa[u][i];
14     FORD (i, SP - 1, -1) {
15         int uu = pa[u][i], vv = pa[v][i];
16         if (uu != vv) { u = uu; v = vv; }
17     }
18     return u == v ? u : pa[u][0];
19 }

```

最短路

```
1  bool BF() {
2      queue<int> q;
3      FOR (i, 1, n) d[i] = INF;
4      d[0] = 0; inq[0] = true; q.push(0);
5      while (!q.empty()) {
6          int u = q.front(); q.pop();
7          inq[u] = false;
8          for (E& e: G[u]) {
9              int v = e.to;
10             if (d[u] < INF && d[v] > d[u] + e.d) {
11                 d[v] = d[u] + e.d;
12                 if (!inq[v]) {
13                     q.push(v); inq[v] = true;
14                     if (++cnt[v] > n) return false;
15                 }
16             }
17         }
18     }
19     return true;
20 }
```

网络流

• 最大流

```
1  struct E {
2      int to, cp;
3      E(int to, int cp): to(to), cp(cp) {}
4  };
5
6  struct Dinic {
7      static const int M = 1E5 * 5;
8      int m, s, t;
9      vector<E> edges;
10     vector<int> G[M];
11     int d[M];
12     int cur[M];
13
14     void init(int n, int s, int t) {
15         this->s = s; this->t = t;
16         for (int i = 0; i <= n; i++) G[i].clear();
17         edges.clear(); m = 0;
18     }
19
20     void addedge(int u, int v, int cap) {
21         edges.emplace_back(v, cap);
22         edges.emplace_back(u, 0);
23         G[u].push_back(m++);
24         G[v].push_back(m++);
25     }
26
27     bool BFS() {
28         memset(d, 0, sizeof d);
29         queue<int> Q;
30         Q.push(s); d[s] = 1;
31         while (!Q.empty()) {
32             int x = Q.front(); Q.pop();
33             for (int& i: G[x]) {
34                 E &e = edges[i];
35                 if (!d[e.to] && e.cp > 0) {
36                     d[e.to] = d[x] + 1;
37                     Q.push(e.to);
38                 }
39             }
40         }
41         return d[t];
42     }
```

```

43
44 int DFS(int u, int cp) {
45     if (u == t || !cp) return cp;
46     int tmp = cp, f;
47     for (int& i = cur[u]; i < G[u].size(); i++) {
48         E& e = edges[G[u][i]];
49         if (d[u] + 1 == d[e.to]) {
50             f = DFS(e.to, min(cp, e.cp));
51             e.cp -= f;
52             edges[G[u][i] ^ 1].cp += f;
53             cp -= f;
54             if (!cp) break;
55         }
56     }
57     return tmp - cp;
58 }
59
60 int go() {
61     int flow = 0;
62     while (BFS()) {
63         memset(cur, 0, sizeof cur);
64         flow += DFS(s, INF);
65     }
66     return flow;
67 }
68 } DC;

```

● 费用流

```

1 struct E {
2     int from, to, cp, v;
3     E() {}
4     E(int f, int t, int cp, int v) : from(f), to(t), cp(cp), v(v) {}
5 };
6
7 struct MCMF {
8     int n, m, s, t;
9     vector<E> edges;
10    vector<int> G[maxn];
11    bool inq[maxn]; //是否在队列
12    int d[maxn]; //Bellman_ford 单源最短路径
13    int p[maxn]; //p[i] 表从 s 到 i 的最小费用路径上的最后一条弧编号
14    int a[maxn]; //a[i] 表示从 s 到 i 的最小残量
15
16    void init(int _n, int _s, int _t) {
17        n = _n; s = _s; t = _t;
18        FOR (i, 0, n + 1) G[i].clear();
19        edges.clear(); m = 0;
20    }
21
22    void addedge(int from, int to, int cap, int cost) {
23        edges.emplace_back(from, to, cap, cost);
24        edges.emplace_back(to, from, 0, -cost);
25        G[from].push_back(m++);
26        G[to].push_back(m++);
27    }
28
29    bool BellmanFord(int &flow, int &cost) {
30        FOR (i, 0, n + 1) d[i] = INF;
31        memset(inq, 0, sizeof inq);
32        d[s] = 0, a[s] = INF, inq[s] = true;
33        queue<int> Q; Q.push(s);
34        while (!Q.empty()) {
35            int u = Q.front(); Q.pop();
36            inq[u] = false;
37            for (int& idx: G[u]) {
38                E &e = edges[idx];
39                if (e.cp && d[e.to] > d[u] + e.v) {
40                    d[e.to] = d[u] + e.v;
41                    p[e.to] = idx;
42                    a[e.to] = min(a[u], e.cp);
43                    if (!inq[e.to]) {

```



```

44         Q.push(e.to);
45         inq[e.to] = true;
46     }
47 }
48 }
49 }
50 if (d[t] == INF) return false;
51 flow += a[t];
52 cost += a[t] * d[t];
53 int u = t;
54 while (u != s) {
55     edges[p[u]].cp -= a[t];
56     edges[p[u] ^ 1].cp += a[t];
57     u = edges[p[u]].from;
58 }
59 return true;
60 }
61
62 int go() {
63     int flow = 0, cost = 0;
64     while (BellmanFord(flow, cost));
65     return cost;
66 }
67 } MM;

```

- zkw 费用流（代码长度没有优势）
- 不允许有负权边

```

1 struct E {
2     int to, cp, v;
3     E() {}
4     E(int to, int cp, int v): to(to), cp(cp), v(v) {}
5 };
6
7 struct MCMF {
8     int n, m, s, t, cost, D;
9     vector<E> edges;
10    vector<int> G[maxn];
11    bool vis[maxn];
12
13    void init(int _n, int _s, int _t) {
14        n = _n; s = _s; t = _t;
15        FOR (i, 0, n + 1) G[i].clear();
16        edges.clear(); m = 0;
17    }
18
19    void addedge(int from, int to, int cap, int cost) {
20        edges.emplace_back(to, cap, cost);
21        edges.emplace_back(from, 0, -cost);
22        G[from].push_back(m++);
23        G[to].push_back(m++);
24    }
25
26    int aug(int u, int cp) {
27        if (u == t) {
28            cost += D * cp;
29            return cp;
30        }
31        vis[u] = true;
32        int tmp = cp;
33        for (int idx: G[u]) {
34            E& e = edges[idx];
35            if (e.cp && !e.v && !vis[e.to]) {
36                int f = aug(e.to, min(cp, e.cp));
37                e.cp -= f;
38                edges[idx ^ 1].cp += f;
39                cp -= f;
40                if (!cp) break;
41            }
42        }
43        return tmp - cp;

```

```

44     }
45
46     bool modlabel() {
47         int d = INF;
48         FOR (u, 0, n + 1)
49             if (vis[u])
50                 for (int& idx: G[u]) {
51                     E& e = edges[idx];
52                     if (e.cp && !vis[e.to]) d = min(d, e.v);
53                 }
54         if (d == INF) return false;
55         FOR (u, 0, n + 1)
56             if (vis[u])
57                 for (int& idx: G[u]) {
58                     edges[idx].v -= d;
59                     edges[idx ^ 1].v += d;
60                 }
61         D += d;
62         return true;
63     }
64
65     int go(int k) {
66         cost = D = 0;
67         int flow = 0;
68         while (true) {
69             memset(vis, 0, sizeof vis);
70             int t = aug(s, INF);
71             if (!t && !modlabel()) break;
72             flow += t;
73         }
74         return cost;
75     }
76 } MM;

```

树上路径交

```

1  int intersection(int x, int y, int xx, int yy) {
2      int t[4] = {lca(x, xx), lca(x, yy), lca(y, xx), lca(y, yy)};
3      sort(t, t + 4);
4      int r = lca(x, y), rr = lca(xx, yy);
5      if (dep[t[0]] < min(dep[r], dep[rr]) || dep[t[2]] < max(dep[r], dep[rr]))
6          return 0;
7      int tt = lca(t[2], t[3]);
8      int ret = 1 + dep[t[2]] + dep[t[3]] - dep[tt] * 2;
9      return ret;
10 }

```

树上点分治

```

1  int get_sz(int u, int fa) {
2      int& s = sz[u] = 1;
3      for (E& e: G[u]) {
4          int v = e.to;
5          if (vis[v] || v == fa) continue;
6          s += get_sz(v, u);
7      }
8      return s;
9  }
10
11 void get_rt(int u, int fa, int s, int& m, int& rt) {
12     int t = s - sz[u];
13     for (E& e: G[u]) {
14         int v = e.to;
15         if (vis[v] || v == fa) continue;
16         get_rt(v, u, s, m, rt);
17         t = max(t, sz[v]);
18     }

```

```

19     if (t < m) { m = t; rt = u; }
20 }
21
22 void dfs(int u) {
23     int tmp = INF; get_rt(u, -1, get_sz(u, -1), tmp, u);
24     vis[u] = true;
25     get_dep(u, -1, 0);
26     // ...
27     for (E& e: G[u]) {
28         int v = e.to;
29         if (vis[v]) continue;
30         // ...
31         dfs(v);
32     }
33 }

```

● 动态点分治

```

1  const int maxn = 15E4 + 100, INF = 1E9;
2  struct E {
3      int to, d;
4  };
5  vector<E> G[maxn];
6  int n, Q, w[maxn];
7  LL A, ans;
8
9  bool vis[maxn];
10 int sz[maxn];
11
12 int get_rt(int u) {
13     // dbg(u);
14     static int q[N], fa[N], sz[N], mx[N];
15     int p = 0, cur = -1;
16     q[p++] = u; fa[u] = -1;
17     while (++cur < p) {
18         u = q[cur]; mx[u] = 0; sz[u] = 1;
19         for (int& v: G[u])
20             if (!vis[v] && v != fa[u]) fa[q[p++]] = v; u;
21     }
22     FORD (i, p - 1, -1) {
23         u = q[i];
24         mx[u] = max(mx[u], p - sz[u]);
25         if (mx[u] * 2 <= p) return u;
26         sz[fa[u]] += sz[u];
27         mx[fa[u]] = max(mx[fa[u]], sz[u]);
28     }
29     assert(0);
30 }
31
32 int get_sz(int u, int fa) {
33     int& s = sz[u] = 1;
34     for (E& e: G[u]) {
35         int v = e.to;
36         if (vis[v] || v == fa) continue;
37         s += get_sz(v, u);
38     }
39     return s;
40 }
41
42 void get_rt(int u, int fa, int s, int& m, int& rt) {
43     int t = s - sz[u];
44     for (E& e: G[u]) {
45         int v = e.to;
46         if (vis[v] || v == fa) continue;
47         get_rt(v, u, s, m, rt);
48         t = max(t, sz[v]);
49     }
50     if (t < m) { m = t; rt = u; }
51 }
52
53 int dep[maxn], md[maxn];
54 void get_dep(int u, int fa, int d) {

```

```

55     dep[u] = d; md[u] = 0;
56     for (E& e: G[u]) {
57         int v = e.to;
58         if (vis[v] || v == fa) continue;
59         get_dep(v, u, d + e.d);
60         md[u] = max(md[u], md[v] + 1);
61     }
62 }
63
64 struct P {
65     int w;
66     LL s;
67 };
68 using VP = vector<P>;
69 struct R {
70     VP *rt, *rt2;
71     int dep;
72 };
73 VP pool[maxn << 1], *pit = pool;
74 vector<R> tr[maxn];
75
76 void go(int u, int fa, VP* rt, VP* rt2) {
77     tr[u].push_back({rt, rt2, dep[u]});
78     for (E& e: G[u]) {
79         int v = e.to;
80         if (v == fa || vis[v]) continue;
81         go(v, u, rt, rt2);
82     }
83 }
84
85 void dfs(int u) {
86     int tmp = INF; get_rt(u, -1, get_sz(u, -1), tmp, u);
87     vis[u] = true;
88     get_dep(u, -1, 0);
89     VP* rt = pit++; tr[u].push_back({rt, nullptr, 0});
90     for (E& e: G[u]) {
91         int v = e.to;
92         if (vis[v]) continue;
93         go(v, u, rt, pit++);
94         dfs(v);
95     }
96 }
97
98 bool cmp(const P& a, const P& b) { return a.w < b.w; }
99
100 LL query(VP& p, int d, int l, int r) {
101     l = lower_bound(p.begin(), p.end(), P{l, -1}, cmp) - p.begin();
102     r = upper_bound(p.begin(), p.end(), P{r, -1}, cmp) - p.begin() - 1;
103     return p[r].s - p[l - 1].s + 1LL * (r - l + 1) * d;
104 }
105
106 int main() {
107     cin >> n >> Q >> A;
108     FOR (i, 1, n + 1) scanf("%d", &w[i]);
109     FOR (_, 1, n) {
110         int u, v, d; scanf("%d%d%d", &u, &v, &d);
111         G[u].push_back({v, d}); G[v].push_back({u, d});
112     }
113     dfs(1);
114     FOR (i, 1, n + 1)
115         for (R& x: tr[i]) {
116             x.rt->push_back({w[i], x.dep});
117             if (x.rt2) x.rt2->push_back({w[i], x.dep});
118         }
119     FOR (it, pool, pit) {
120         it->push_back({-INF, 0});
121         sort(it->begin(), it->end(), cmp);
122         FOR (i, 1, it->size())
123             (*it)[i].s += (*it)[i - 1].s;
124     }
125     while (Q--) {

```

```

126     int u; LL a, b; scanf("%d%lld%lld", &u, &a, &b);
127     a = (a + ans) % A; b = (b + ans) % A;
128     int l = min(a, b), r = max(a, b);
129     ans = 0;
130     for (R& x: tr[u]) {
131         ans += query(*(x.rt), x.dep, l, r);
132         if (x.rt2) ans -= query(*(x.rt2), x.dep, l, r);
133     }
134     printf("%lld\n", ans);
135 }
136 }

```

树链剖分

```

1  int fa[N], dep[N], idx[N], out[N], ridx[N];
2  namespace hld {
3      int sz[N], son[N], top[N], clk;
4      void predfs(int u, int d) {
5          dep[u] = d; sz[u] = 1;
6          int& maxs = son[u] = -1;
7          for (int& v: G[u]) {
8              if (v == fa[u]) continue;
9              fa[v] = u;
10             predfs(v, d + 1);
11             sz[u] += sz[v];
12             if (maxs == -1 || sz[v] > sz[maxs]) maxs = v;
13         }
14     }
15     void dfs(int u, int tp) {
16         top[u] = tp; idx[u] = ++clk; ridx[clk] = u;
17         if (son[u] != -1) dfs(son[u], tp);
18         for (int& v: G[u])
19             if (v != fa[u] && v != son[u]) dfs(v, v);
20         out[u] = clk;
21     }
22     template<typename T>
23     int go(int u, int v, T&& f = [] (int, int) {}) {
24         int uu = top[u], vv = top[v];
25         while (uu != vv) {
26             if (dep[uu] < dep[vv]) { swap(uu, vv); swap(u, v); }
27             f(idx[uu], idx[u]);
28             u = fa[uu]; uu = top[u];
29         }
30         if (dep[u] < dep[v]) swap(u, v);
31         // f(idx[v], idx[u]);
32         // if (u != v) f(idx[v] + 1, idx[u]);
33         return v;
34     }
35     int up(int u, int d) {
36         while (d) {
37             if (dep[u] - dep[top[u]] < d) {
38                 d -= dep[u] - dep[top[u]];
39                 u = top[u];
40             } else return ridx[idx[u] - d];
41             u = fa[u]; --d;
42         }
43         return u;
44     }
45 }

```

• HDU 3966

```

1  #include <bits/stdc++.h>
2  using namespace std;
3  typedef long long LL;
4  #define FOR(i, x, y) for (decay<decltype(y)>::type i = (x), _##i = (y); i < _##i; ++i)
5  #define FORD(i, x, y) for (decay<decltype(x)>::type i = (x), _##i = (y); i > _##i; --i)
6  #ifdef zero1
7  #define dbg(args...) do { cout << "\033[32;1m" << #args << " -> "; err(args); } while (0)
8  #else

```

```

9  #define dbg(...)
10 #endif
11 void err() { cout << "\033[39;0m" << endl; }
12 template<typename T, typename... Args>
13 void err(T a, Args... args) { cout << ' '; err(args...); }
14 // -----
15 const int maxn = 5E4 + 100;
16 vector<int> G[maxn];
17 int dep[maxn], sz[maxn], son[maxn], fa[maxn], idx[maxn], top[maxn];
18 int clk, n, Q;
19
20 struct IntervalTree {
21     #define ls o * 2, l, (l + r) >> 1
22     #define rs o * 2 + 1, ((l + r) >> 1) + 1, r
23     static const int M = maxn << 2;
24     int addv[M];
25     void init() { memset(addv, 0, sizeof addv); }
26     int query(int k, int o, int l, int r, int add = 0) {
27         if (k < l || r < k) return 0;
28         if (l == r) return add + addv[o];
29         return query(k, ls, add + addv[o]) + query(k, rs, add + addv[o]);
30     }
31     void update(int p, int q, int o, int l, int r, int add) {
32         assert(l <= r && r <= n);
33         if (q < l || r < p) return;
34         if (p <= l && r <= q) addv[o] += add;
35         else { update(p, q, ls, add); update(p, q, rs, add); }
36     }
37 } IT;
38
39 void predfs(int u, int d) {
40     dep[u] = d;
41     sz[u] = 1;
42     int& maxs = son[u] = -1;
43     for (int v : G[u])
44         if (v != fa[u]) {
45             fa[v] = u;
46             predfs(v, d + 1);
47             sz[u] += sz[v];
48             if (maxs == -1 || sz[v] > sz[maxs])
49                 maxs = v;
50         }
51 }
52
53 void dfs(int u, int tp) {
54     top[u] = tp;
55     idx[u] = ++clk;
56     if (son[u] != -1) dfs(son[u], tp);
57     for (int v : G[u])
58         if (v != son[u] && v != fa[u])
59             dfs(v, v);
60 }
61
62 void update(int u, int v, int add) {
63     int uu = top[u], vv = top[v];
64     while (uu != vv) {
65         if (dep[uu] < dep[vv]) { swap(uu, vv); swap(u, v); }
66         IT.update(idx[uu], idx[u], 1, 1, n, add);
67         u = fa[uu];
68         uu = top[u];
69     }
70     if (dep[u] < dep[v]) swap(u, v);
71     dbg(u, v, idx[u], idx[v]);
72     IT.update(idx[v], idx[u], 1, 1, n, add);
73 }
74
75 int a[maxn];
76 void init();
77 int main() {
78     int u, v, l, r, k, d;
79     char s[100];

```

```

80 while (cin >> n >> Q >> Q) {
81     init();
82     FOR (i, 1, n + 1) scanf("%d", &a[i]);
83     FOR (i, 1, n) {
84         scanf("%d%d", &u, &v);
85         G[u].push_back(v);
86         G[v].push_back(u);
87     }
88     predfs(1, 1);
89     dfs(1, 1);
90     while (Q--) {
91         scanf("%s", s);
92         if (s[0] == 'I') {
93             scanf("%d%d%d", &l, &r, &d);
94             update(l, r, d);
95         } else if (s[0] == 'D') {
96             scanf("%d%d%d", &l, &r, &d);
97             update(l, r, -d);
98         } else {
99             scanf("%d", &k);
100             printf("%d\n", a[k] + IT.query(idx[k], 1, 1, n));
101         }
102     }
103 }
104 }
105
106 void init() {
107     clk = 0;
108     fa[1] = 0;
109     IT.init();
110     FOR (i, 0, n + 1) G[i].clear();
111 }

```

• SPOJ QTREE

```

1  #include <bits/stdc++.h>
2  using namespace std;
3  typedef long long LL;
4  #define FOR(i, x, y) for (decay<decltype(y)>::type i = (x), _##i = (y); i < _##i; ++i)
5  #define FORR(i, x, y) for (decay<decltype(x)>::type i = (x), _##i = (y); i > _##i; --i)
6  #ifdef zero1
7  #define dbg(args...) do { cout << "\033[32;1m" << #args << " -> "; err(args); } while (0)
8  #else
9  #define dbg(...)
10 #endif
11 void err() { cout << "\033[39m" << endl; }
12 template<typename T, typename... Args>
13 void err(T a, Args... args) {
14     cout << a << ' ';
15     err(args...);
16 }
17 // -----
18 const int maxn = 10000 * 2 * 4 + 100;
19 struct Edge {
20     int from, to, c;
21     Edge(int u, int v, int c): from(u), to(v), c(c) {}
22 };
23 vector<Edge> edge;
24 vector<int> G[maxn];
25 int fa[maxn], dep[maxn], sz[maxn], son[maxn], top[maxn], idx[maxn], w[maxn], val[maxn];
26 LL sum[maxn];
27 int n, clk, len;
28
29 struct IntervalTree {
30     #define lson p, q, o * 2, l, m
31     #define rson p, q, o * 2 + 1, m + 1, r
32     int maxv[maxn];
33     void init() { memset(maxv, 0, sizeof maxv); }
34     int query(int p, int q, int o, int l, int r) {
35         // dbg(p, q);
36         assert(p <= q);
37         if (p > r || q < l) return 0;

```

```

38         if (p <= l && r <= q) return maxv[o];
39         int m = (l + r) / 2;
40         return max(query(lson), query(rson));
41     }
42     void maintain(int o, int l, int r) {
43         if (l < r)
44             maxv[o] = max(maxv[o * 2], maxv[o * 2 + 1]);
45     }
46     void update(int p, int q, int o, int l, int r, int v) {
47         // dbg(p, q, o, l, r, v);
48         assert(p <= q);
49         if (p > r || q < l) return;
50         if (p <= l && r <= q) maxv[o] = v;
51         else {
52             int m = (l + r) / 2;
53             update(lson, v); update(rson, v);
54             maintain(o, l, r);
55         }
56     }
57 } IT;
58
59 void dfs1(int u, int d) {
60     dep[u] = d;
61     sz[u] = 1;
62     FOR (i, 0, G[u].size()) {
63         Edge& e = edge[G[u][i]];
64         int v = e.to;
65         if (v == fa[u]) continue;
66         val[v] = e.c;
67         // dbg(v, e.from, e.to, e.c);
68         fa[v] = u;
69         dfs1(v, d + 1);
70         sz[u] += sz[v];
71         if (son[u] == -1 || sz[v] > sz[son[u]])
72             son[u] = v;
73     }
74 }
75
76 void dfs2(int u, int tp) {
77     top[u] = tp;
78     idx[u] = ++clk;
79     w[idx[u]] = tp;
80     if (son[u] == -1) return;
81     dfs2(son[u], tp);
82     FOR (i, 0, G[u].size()) {
83         int v = edge[G[u][i]].to;
84         if (v != son[u] && v != fa[u])
85             dfs2(v, v);
86     }
87 }
88
89 int query(int u, int v) {
90     dbg(u, v);
91     int uu = top[u], vv = top[v], ret = 0;
92     while (uu != vv) {
93         if (dep[uu] < dep[vv]) { swap(u, v); swap(uu, vv); }
94         // dbg(u, v, uu, vv, dep[uu], dep[vv], idx[uu], idx[u]);
95         ret = max(ret, IT.query(idx[uu], idx[u], 1, 1, len));
96         u = fa[uu];
97         uu = top[u];
98     }
99     if (dep[u] < dep[v]) swap(u, v);
100    // dbg(idx[v], idx[u]);
101    if (u != v) ret = max(ret, IT.query(idx[v] + 1, idx[u], 1, 1, len));
102    return ret;
103 }
104
105 void init();
106 void add_edge(int u, int v, int c);
107
108 int main() {

```



```

109  #ifndef zerol
110      freopen("in", "r", stdin);
111  #endif
112      int T, u, v, c;
113      char s[100];
114      cin >> T;
115      while (T--) {
116          cin >> n;
117          for (len = 1; len < n; len *= 2);
118          init();
119          FOR (i, 1, n) {
120              scanf("%d%d%d", &u, &v, &c);
121              add_edge(u, v, c);
122              add_edge(v, u, c);
123          }
124          dfs1(1, 0);
125          dfs2(1, 1);
126          // FOR (i, 1, n + 1) dbg(idx[i], w[i]);
127          FOR (i, 2, n + 1)
128              IT.update(idx[i], idx[i], 1, 1, len, val[i]);
129          while (scanf("%s", s) && s[0] != 'D') {
130              scanf("%d%d", &u, &v);
131              if (s[0] == 'C') {
132                  Edge& e = edge[u * 2 - 1];
133                  dbg(u, e.from, e.to);
134                  int t = max(idx[e.from], idx[e.to]);
135                  IT.update(t, t, 1, 1, len, v);
136                  dbg("upd", t, v);
137              }
138              if (s[0] == 'Q') printf("%d\n", query(u, v));
139          }
140          FOR (i, 1, n + 1) if (idx[i] == 2) dbg(i, idx[i]);
141          dbg(IT.query(idx[2], idx[2], 1, 1, len));
142          dbg(IT.query(idx[6], idx[6], 1, 1, len));
143      }
144  }
145
146  void init() {
147      edge.clear();
148      memset(son, -1, sizeof son);
149      memset(sum, 0, sizeof sum);
150      IT.init();
151      FOR (i, 0, n + 1) G[i].clear();
152      clk = 0;
153      fa[1] = 0;
154      sum[0] = sum[1] = 0;
155  }
156
157  void add_edge(int u, int v, int c) {
158      edge.emplace_back(u, v, c);
159      G[u].push_back(edge.size() - 1);
160  }

```

二分图匹配

- 最小覆盖数 = 最大匹配数
- 最大独立集 = 顶点数 - 二分图匹配数
- DAG 最小路径覆盖数 = 结点数 - 拆点后二分图最大匹配数

```

1  struct MaxMatch {
2      int n;
3      vector<int> G[maxn];
4      int vis[maxn], left[maxn], clk;
5
6      void init(int n) {
7          this->n = n;
8          FOR (i, 0, n + 1) G[i].clear();
9          memset(left, -1, sizeof left);
10         memset(vis, -1, sizeof vis);

```

```

11     }
12
13     bool dfs(int u) {
14         for (int v: G[u])
15             if (vis[v] != clk) {
16                 vis[v] = clk;
17                 if (left[v] == -1 || dfs(left[v])) {
18                     left[v] = u;
19                     return true;
20                 }
21             }
22         return false;
23     }
24
25     int match() {
26         int ret = 0;
27         for (clk = 0; clk <= n; ++clk)
28             if (dfs(clk)) ++ret;
29         return ret;
30     }
31 } MM;

```

● 二分图最大权完美匹配 KM

```

1 namespace R {
2     const int maxn = 300 + 10;
3     int n, m;
4     int left[maxn], L[maxn], R[maxn];
5     int w[maxn][maxn], slack[maxn];
6     bool visL[maxn], visR[maxn];
7
8     bool dfs(int u) {
9         visL[u] = true;
10        FOR (v, 0, m) {
11            if (visR[v]) continue;
12            int t = L[u] + R[v] - w[u][v];
13            if (t == 0) {
14                visR[v] = true;
15                if (left[v] == -1 || dfs(left[v])) {
16                    left[v] = u;
17                    return true;
18                }
19            } else slack[v] = min(slack[v], t);
20        }
21        return false;
22    }
23
24    int go() {
25        memset(left, -1, sizeof left);
26        memset(R, 0, sizeof R);
27        memset(L, 0, sizeof L);
28        FOR (i, 0, n)
29            FOR (j, 0, m)
30                L[i] = max(L[i], w[i][j]);
31
32        FOR (i, 0, n) {
33            memset(slack, 0x3f, sizeof slack);
34            while (1) {
35                memset(visL, 0, sizeof visL); memset(visR, 0, sizeof visR);
36                if (dfs(i)) break;
37                int d = 0x3f3f3f3f;
38                FOR (j, 0, m) if (!visR[j]) d = min(d, slack[j]);
39                FOR (j, 0, n) if (visL[j]) L[j] -= d;
40                FOR (j, 0, m) if (visR[j]) R[j] += d; else slack[j] -= d;
41            }
42        }
43        int ret = 0;
44        FOR (i, 0, m) if (left[i] != -1) ret += w[left[i]][i];
45        return ret;
46    }
47 }

```

虚树

```
1 void go(vector<int>& V, int& k) {
2     int u = V[k]; f[u] = 0;
3     dbg(u, k);
4     for (auto& e: G[u]) {
5         int v = e.to;
6         if (v == pa[u][0]) continue;
7         while (k + 1 < V.size()) {
8             int to = V[k + 1];
9             if (in[to] <= out[v]) {
10                go(V, ++k);
11                if (key[to]) f[u] += w[to];
12                else f[u] += min(f[to], (LL)w[to]);
13            } else break;
14        }
15    }
16    dbg(u, f[u]);
17 }
18 inline bool cmp(int a, int b) { return in[a] < in[b]; }
19 LL solve(vector<int>& V) {
20     static vector<int> a; a.clear();
21     for (int& x: V) a.push_back(x);
22     sort(a.begin(), a.end(), cmp);
23     FOR (i, 1, a.size())
24         a.push_back(lca(a[i], a[i - 1]));
25     a.push_back(1);
26     sort(a.begin(), a.end(), cmp);
27     a.erase(unique(a.begin(), a.end()), a.end());
28     dbg(a);
29     int tmp; go(a, tmp = 0);
30     return f[1];
31 }
```

欧拉路径

```
1 int S[N << 1], top;
2 Edge edges[N << 1];
3 set<int> G[N];
4
5 void DFS(int u) {
6     S[top++] = u;
7     for (int eid: G[u]) {
8         int v = edges[eid].get_other(u);
9         G[u].erase(eid);
10        G[v].erase(eid);
11        DFS(v);
12        return;
13    }
14 }
15
16 void fleury(int start) {
17     int u = start;
18     top = 0; path.clear();
19     S[top++] = u;
20     while (top) {
21         u = S[--top];
22         if (!G[u].empty())
23             DFS(u);
24         else path.push_back(u);
25     }
26 }
```

强连通分量与 2-SAT

```
1 int n, m;
2 vector<int> G[N], rG[N], vs;
```

```

3  int used[N], cmp[N];
4
5  void add_edge(int from, int to) {
6      G[from].push_back(to);
7      rG[to].push_back(from);
8  }
9
10 void dfs(int v) {
11     used[v] = true;
12     for (int u: G[v]) {
13         if (!used[u])
14             dfs(u);
15     }
16     vs.push_back(v);
17 }
18
19 void rdfs(int v, int k) {
20     used[v] = true;
21     cmp[v] = k;
22     for (int u: rG[v])
23         if (!used[u])
24             rdfs(u, k);
25 }
26
27 int scc() {
28     memset(used, 0, sizeof(used));
29     vs.clear();
30     for (int v = 0; v < n; ++v)
31         if (!used[v]) dfs(v);
32     memset(used, 0, sizeof(used));
33     int k = 0;
34     for (int i = (int) vs.size() - 1; i >= 0; --i)
35         if (!used[vs[i]]) rdfs(vs[i], k++);
36     return k;
37 }
38
39 int main() {
40     cin >> n >> m;
41     n *= 2;
42     for (int i = 0; i < m; ++i) {
43         int a, b; cin >> a >> b;
44         add_edge(a - 1, (b - 1) ^ 1);
45         add_edge(b - 1, (a - 1) ^ 1);
46     }
47     scc();
48     for (int i = 0; i < n; i += 2) {
49         if (cmp[i] == cmp[i + 1]) {
50             puts("NIE");
51             return 0;
52         }
53     }
54     for (int i = 0; i < n; i += 2) {
55         if (cmp[i] > cmp[i + 1]) printf("%d\n", i + 1);
56         else printf("%d\n", i + 2);
57     }
58 }

```

拓补排序

```

1  vector<int> toporder(int n) {
2      vector<int> orders;
3      queue<int> q;
4      for (int i = 0; i < n; i++)
5          if (!deg[i]) {
6              q.push(i);
7              orders.push_back(i);
8          }
9      while (!q.empty()) {
10         int u = q.front(); q.pop();

```

```

11         for (int v: G[u])
12             if (!--deg[v]) {
13                 q.push(v);
14                 orders.push_back(v);
15             }
16     }
17     return orders;
18 }

```

一般图匹配 (没有测试过)

```

1 // maximum matching (graph not necessarily bipartite)
2 // whatever code uses this needs to set N in main()
3 // author claims N^4 running time
4
5 #define PB push_back
6 #define SZ(x) ((int)(x).size())
7 #define REP(i, n) for(int i=0; i<n; ++i)
8 #define FOR(i, b, e) for(typeof(e) i=b; i!=e; ++i)
9
10 int N; // the number of vertices in the graph
11
12 typedef vector<int> VI;
13 typedef vector<vector<int>> VVI;
14
15 VI match;
16 VI vis;
17
18 void couple(int n, int m) {
19     match[n] = m;
20     match[m] = n;
21 }
22
23 // returns true if something interesting has been found, thus a
24 // augmenting path or a blossom (if blossom is non-empty).
25 // the dfs returns true from the moment the stem of the flower is
26 // reached and thus the base of the blossom is an unmatched node.
27 // blossom should be empty when dfs is called and
28 // contains the nodes of the blossom when a blossom is found.
29 bool dfs(int n, VVI &conn, VI &blossom) {
30     vis[n] = 0;
31     REP(i, N) if (conn[n][i]) {
32         if (vis[i] == -1) {
33             vis[i] = 1;
34             if (match[i] == -1 || dfs(match[i], conn, blossom)) {
35                 couple(n, i);
36                 return true;
37             }
38         }
39         if (vis[i] == 0 || SZ(blossom)) { // found flower
40             blossom.PB(i);
41             blossom.PB(n);
42             if (n == blossom[0]) {
43                 match[n] = -1;
44                 return true;
45             }
46             return false;
47         }
48     }
49     return false;
50 }
51
52 // search for an augmenting path.
53 // if a blossom is found build a new graph (newconn) where the
54 // (free) blossom is shrunk to a single node and recurse.
55 // if a augmenting path is found it has already been augmented
56 // except if the augmented path ended on the shrunk blossom.
57 // in this case the matching should be updated along the appropriate
58 // direction of the blossom.

```

```

59 bool augment(VVI & conn) {
60     REP(m, N) if (match[m] == -1) {
61         VI blossom;
62         vis = VI(N, -1);
63         if (!dfs(m, conn, blossom)) continue;
64         if (SZ(blossom) == 0) return true; // augmenting path found
65
66         // blossom is found so build shrunken graph
67         int base = blossom[0], S = SZ(blossom);
68         VVI newconn = conn;
69         FOR(i, 1, S - 1) REP(j, N)
70             newconn[base][j] = newconn[j][base] |= conn[blossom[i]][j];
71         FOR(i, 1, S - 1) REP(j, N)
72             newconn[blossom[i]][j] = newconn[j][blossom[i]] = 0;
73         newconn[base][base] = 0; // is now the new graph
74         if (!augment(newconn)) return false;
75         int n = match[base];
76
77         // if n!=-1 the augmenting path ended on this blossom
78         if (n != -1)
79             REP(i, S) if (conn[blossom[i]][n]) {
80                 couple(blossom[i], n);
81                 if (i & 1) for (int j = i + 1; j < S; j += 2)
82                     couple(blossom[j], blossom[j + 1]);
83                 else for (int j = 0; j < i; j += 2)
84                     couple(blossom[j], blossom[j + 1]);
85                 break;
86             }
87         return true;
88     }
89     return false;
90 }
91
92 // conn should have N VI's, each of length N. conn[i][j] = 1 if (i,j) is an
93 // edge, and conn[i][j] = 0 otherwise. returns the number of edges in a max
94 // matching.
95 int edmonds(VVI & conn) {
96     int res = 0;
97     match = VI(N, -1);
98     while (augment(conn)) res++;
99     return res;
100 }

```

计算几何

圆的反演

```

1  typedef double LD;
2  const LD PI = 3.14159265358979323846;
3  const LD eps = 1E-10;
4  const LD R2 = 1.0;
5  int sgn(LD x) { return fabs(x) < eps ? 0 : (x > 0 ? 1 : -1); }
6  struct P {
7      LD x, y;
8      P(LD x = 0, LD y = 0): x(x), y(y) {}
9      P operator * (LD k) { return P(x * k, y * k); }
10     P operator / (LD k) { return P(x / k, y / k); }
11     string prt() const {
12         char s[100];
13         sprintf(s, "(%.2f, %.2f)", x, y);
14         return string(s);
15     }
16 };
17 typedef P V;
18 P operator - (const P& a, const P& b) { return P(a.x - b.x, a.y - b.y); }
19 P operator + (const P& a, const P& b) { return P(a.x + b.x, a.y + b.y); }
20 struct C {
21     P p;

```

```

22     LD r;
23     C(LD x = 0, LD y = 0, LD r = 0): p(x, y), r(r) {}
24 };
25 LD dist(V v) { return sqrt(v.x * v.x + v.y * v.y); }
26
27 C inv(C c, const P& o) {
28     LD d = dist(c.p - o);
29     assert(sgn(d) != 0);
30     LD a = 1 / (d - c.r);
31     LD b = 1 / (d + c.r);
32     c.r = (a - b) / 2 * R2;
33     c.p = o + (c.p - o) * ((a + b) * R2 / 2 / d);
34     return c;
35 }

```

二维

- nxt 宏要求多边形变量名为 s
- L 可隐式转换为 V(P)
- 可以自定义结构体 PP, 可隐式转换为 P

```

1  #define y1 yy1
2  #define nxt(i) ((i + 1) % s.size())
3  typedef double LD;
4  const LD PI = 3.14159265358979323846;
5  const LD eps = 1E-10;
6  int sgn(LD x) { return fabs(x) < eps ? 0 : (x > 0 ? 1 : -1); }
7  struct L;
8  struct P;
9  //struct PP;
10 typedef P V;
11 struct P {
12     LD x, y;
13     explicit P(LD x = 0, LD y = 0): x(x), y(y) {}
14     P(const L& l);
15     // P(const PP& pp);
16 };
17 struct L {
18     P s, t;
19     L() {}
20     L(P s, P t): s(s), t(t) {}
21 };
22
23 P operator + (const P& a, const P& b) { return P(a.x + b.x, a.y + b.y); }
24 P operator - (const P& a, const P& b) { return P(a.x - b.x, a.y - b.y); }
25 P operator * (const P& a, LD k) { return P(a.x * k, a.y * k); }
26 P operator / (const P& a, LD k) { return P(a.x / k, a.y / k); }
27 inline int operator < (const P& a, const P& b) {
28     return sgn(a.x - b.x) < 0 || (sgn(a.x - b.x) == 0 && sgn(a.y - b.y) < 0);
29 }
30 bool operator == (const P& a, const P& b) { return !sgn(a.x - b.x) && !sgn(a.y - b.y); }
31 P::P(const L& l) { *this = l.t - l.s; }
32 ostream &operator << (ostream &os, const P &p) {
33     return (os << "(" << p.x << ", " << p.y << ")");
34 }
35 istream &operator >> (istream &is, P &p) {
36     return (is >> p.x >> p.y);
37 }
38
39 // -----
40
41 //struct PP {
42 //    P p;
43 //    LD v, l;
44 //};
45 //P::P(const PP& pp) { *this = pp.p; }
46 typedef P PP;
47
48 typedef vector<PP> S;

```

```

49 // -----
50 LD dist(const P& p) { return sqrt(p.x * p.x + p.y * p.y); }
51 LD dot(const V& a, const V& b) { return a.x * b.x + a.y * b.y; }
52 LD det(const V& a, const V& b) { return a.x * b.y - a.y * b.x; }
53 LD cross(const P& s, const P& t, const P& o = P()) { return det(s - o, t - o); }
54
55 // 如需支持 unique, 需要加 eps
56 bool cmp_xy(const P& a, const P& b) { return a.x < b.x || a.x == b.x && a.y < b.y; }
57
58 // 象限
59 int quad(P p) {
60     int x = sgn(p.x), y = sgn(p.y);
61     if (x > 0 && y >= 0) return 1;
62     if (x <= 0 && y > 0) return 2;
63     if (x < 0 && y <= 0) return 3;
64     if (x >= 0 && y < 0) return 4;
65     assert(0);
66 }
67
68 // 仅适用于参照点在所有点一侧的情况
69 struct cmp_angle {
70     P p;
71     bool operator () (const P& a, const P& b) {
72         // int qa = quad(a), qb = quad(b);
73         // if (qa != qb) return qa < qb;
74         int d = sgn(cross(a, b, p));
75         if (d) return d > 0;
76         return dist(a - p) < dist(b - p);
77     }
78 };
79
80 // -----线-----
81
82 // 是否平行
83 bool parallel(const L& a, const L& b) {
84     return !sgn(det(a, b));
85 }
86
87 // 直线是否相等
88 bool l_eq(const L& a, const L& b) {
89     return parallel(a, b) && parallel(L(a.s, b.t), L(b.s, a.t));
90 }
91
92 // 逆时针旋转 r 弧度
93 P rotation(const P& p, const LD& r) { return P(p.x * cos(r) - p.y * sin(r), p.x * sin(r) + p.y * cos(r)); }
94 P RotateCCW90(const P& p) { return P(-p.y, p.x); }
95 P RotateCW90(const P& p) { return P(p.y, -p.x); }
96
97 // 单位法向量
98 V normal(const V& v) { return V(-v.y, v.x) / dist(v); }
99
100 // -----点和线-----
101
102 // 点在线段上 <= 0 包含端点 < 0 则不包含
103 bool p_on_seg(const P& p, const L& seg) {
104     P a = seg.s, b = seg.t;
105     return !sgn(det(p - a, b - a)) && sgn(dot(p - a, p - b)) <= 0;
106 }
107
108 // 点到直线距离
109 LD dist_to_line(const P& p, const L& l) {
110     return fabs(cross(l.s, l.t, p)) / dist(l);
111 }
112
113 // 点到线段距离
114 LD dist_to_seg(const P& p, const L& l) {
115     if (l.s == l.t) return dist(p - l);
116     V vs = p - l.s, vt = p - l.t;
117     if (sgn(dot(l, vs)) < 0) return dist(vs);
118     else if (sgn(dot(l, vt)) > 0) return dist(vt);
119     else return dist_to_line(p, l);
120 }

```



```

120 // -----线和线-----
121
122 // 求直线交 需要事先保证有界
123 P l_intersection(const L& a, const L& b) {
124     LD s1 = det(a, b.s - a.s), s2 = det(a, b.t - a.s);
125     return (b.s * s2 - b.t * s1) / (s2 - s1);
126 }
127 // 向量夹角的弧度
128 LD angle(const V& a, const V& b) {
129     LD r = asin(fabs(det(a, b)) / dist(a) / dist(b));
130     if (sgn(dot(a, b)) < 0) r = PI - r;
131     return r;
132 }
133 // 线段和直线是否有交 1 = 规范, 2 = 不规范
134 int s_l_cross(const L& seg, const L& line) {
135     int d1 = sgn(cross(line.s, line.t, seg.s));
136     int d2 = sgn(cross(line.s, line.t, seg.t));
137     if ((d1 ^ d2) == -2) return 1; // proper
138     if (d1 == 0 || d2 == 0) return 2;
139     return 0;
140 }
141 // 线段的交 1 = 规范, 2 = 不规范
142 int s_cross(const L& a, const L& b, P& p) {
143     int d1 = sgn(cross(a.t, b.s, a.s)), d2 = sgn(cross(a.t, b.t, a.s));
144     int d3 = sgn(cross(b.t, a.s, b.s)), d4 = sgn(cross(b.t, a.t, b.s));
145     if ((d1 ^ d2) == -2 && (d3 ^ d4) == -2) { p = l_intersection(a, b); return 1; }
146     if (!d1 && p_on_seg(b.s, a)) { p = b.s; return 2; }
147     if (!d2 && p_on_seg(b.t, a)) { p = b.t; return 2; }
148     if (!d3 && p_on_seg(a.s, b)) { p = a.s; return 2; }
149     if (!d4 && p_on_seg(a.t, b)) { p = a.t; return 2; }
150     return 0;
151 }
152
153 // -----多边形-----
154
155 // 点是否在多边形中 0 = 在外部 1 = 在内部 -1 = 在边界上
156 int inside(const S& s, const P& p) {
157     int cnt = 0;
158     FOR (i, 0, s.size()) {
159         P a = s[i], b = s[nxt(i)];
160         if (p_on_seg(p, L(a, b))) return -1;
161         if (sgn(a.y - b.y) <= 0) swap(a, b);
162         if (sgn(p.y - a.y) > 0) continue;
163         if (sgn(p.y - b.y) <= 0) continue;
164         cnt += sgn(cross(b, a, p)) > 0;
165     }
166     return bool(cnt & 1);
167 }
168 // 多边形面积, 有向面积可能为负
169 LD polygon_area(const S& s) {
170     LD ret = 0;
171     FOR (i, 1, (LL)s.size() - 1)
172         ret += cross(s[i], s[i + 1], s[0]);
173     return ret / 2;
174 }
175
176 // 构建凸包 点不可以重复 < 0 边上可以有点, <= 0 则不能
177 // 会改变输入点的顺序
178 const int MAX_N = 1000;
179 S convex_hull(S& s) {
180     // assert(s.size() >= 3);
181     sort(s.begin(), s.end(), cmp_xy);
182     S ret(MAX_N * 2);
183     int sz = 0;
184     FOR (i, 0, s.size()) {
185         while (sz > 1 && sgn(cross(ret[sz - 1], s[i], ret[sz - 2])) < 0) --sz;
186         ret[sz++] = s[i];
187     }
188     int k = sz;
189     FOR (i, (LL)s.size() - 2, -1) {
190         while (sz > k && sgn(cross(ret[sz - 1], s[i], ret[sz - 2])) < 0) --sz;

```

```

191     ret[sz++] = s[i];
192 }
193 ret.resize(sz - (s.size() > 1));
194 return ret;
195 }
196
197 P ComputeCentroid(const vector<P> &p) {
198     P c(0, 0);
199     LD scale = 6.0 * polygon_area(p);
200     for (unsigned i = 0; i < p.size(); i++) {
201         unsigned j = (i + 1) % p.size();
202         c = c + (p[i] + p[j]) * (p[i].x * p[j].y - p[j].x * p[i].y);
203     }
204     return c / scale;
205 }
206
207 // -----圆-----
208
209 P ComputeCircleCenter(P a, P b, P c) {
210     b = (a + b) / 2;
211     c = (a + c) / 2;
212     return l_intersection({b, b + RotateCW90(a - b)}, {c, c + RotateCW90(a - c)});
213 }
214 vector<P> CircleLineIntersection(P a, P b, P c, LD r) {
215     vector<P> ret;
216     b = b - a;
217     a = a - c;
218     LD A = dot(b, b), B = dot(a, b), C = dot(a, a) - r * r;
219     LD D = B * B - A * C;
220     if (sgn(D) < 0) return ret;
221     ret.push_back(c + a + b * (-B + sqrt(D + eps)) / A);
222     if (sgn(D) > 0) ret.push_back(c + a + b * (-B - sqrt(D)) / A);
223     return ret;
224 }
225 vector<P> CircleCircleIntersection(P a, P b, LD r, LD R) {
226     vector<P> ret;
227     LD d = dist(a - b);
228     if (sgn(d) == 0 || sgn(d - (r + R)) > 0 || sgn(d + min(r, R) - max(r, R)) < 0) return ret;
229     LD x = (d * d - R * R + r * r) / (2 * d);
230     LD y = sqrt(r * r - x * x);
231     P v = (b - a) / d;
232     ret.push_back(a + v * x + RotateCCW90(v) * y);
233     if (sgn(y) > 0) ret.push_back(a + v * x - RotateCCW90(v) * y);
234     return ret;
235 }
236
237 // -----模板结束-----

```

旋转卡壳

```

1 LD rotatingCalipers(S& qs) {
2     int n = qs.size();
3     if (n == 2)
4         return dist(qs[0] - qs[1]);
5     int i = 0, j = 0;
6     FOR (k, 0, n) {
7         if (!(qs[i] < qs[k])) i = k;
8         if (qs[j] < qs[k]) j = k;
9     }
10    LD res = 0;
11    int si = i, sj = j;
12    while (i != sj || j != si) {
13        res = max(res, dist(qs[i] - qs[j]));
14        if (sgn(cross(qs[(i+1)%n] - qs[i], qs[(j+1)%n] - qs[j])) < 0)
15            i = (i + 1) % n;
16        else j = (j + 1) % n;
17    }
18    return res;
19 }

```

```

20
21 int main() {
22     int n;
23     while (cin >> n) {
24         S v(n);
25         FOR (i, 0, n) cin >> v[i].x >> v[i].y;
26         convex_hull(v);
27         printf("%.0f\n", rotatingCalipers(v));
28     }
29 }

```

没有测试过的

```

1  int relation(Point p, Circle a) { //点和圆的关系
2      //0: 圆外 1: 圆上 2: 圆内
3      double d = dis(p, a.p);
4      if (dcmp(d - a.r) == 0) return 1;
5      return (dcmp(d - a.r) < 0 ? 2 : 0);
6  }
7
8  int relation(Line a, Circle b) { //直线和圆的关系
9      //0: 相离 1: 相切 2: 相交
10     double p = point_to_line(b.p, a);
11     if (dcmp(p - b.r) == 0) return 1;
12     return (dcmp(p - b.r) < 0 ? 2 : 0);
13 }
14
15 int relation(Circle a, Circle v) { //两圆的位置关系
16     //1: 内含 2: 内切 3: 相交 4: 外切 5: 相离
17     double d = dis(a.p, v.p);
18     if (dcmp(d - a.r - v.r) > 0) return 5;
19     if (dcmp(d - a.r - v.r) == 0) return 4;
20     double l = fabs(a.r - v.r);
21     if (dcmp(d - a.r - v.r) < 0 && dcmp(d - l) > 0) return 3;
22     if (dcmp(d - l) == 0) return 2;
23     if (dcmp(d - l) < 0) return 1;
24     assert (0);
25 }
26
27 double circle_triangle_area(Point a, Point b, Circle c) { //圆心三角形的面积
28     //a.output (), b.output (), c.output ();
29     Point p = c.p;
30     double r = c.r; //cout << cross (p-a, p-b) << endl;
31     if (dcmp(cross(p - a, p - b)) == 0) return 0;
32     Point q[5];
33     int len = 0;
34     q[len++] = a;
35     Line l(a, b);
36     Point p1, p2;
37     if (line_circle_intersection(l, c, q[1], q[2]) == 2) {
38         if (dcmp(dot(a - q[1], b - q[1])) < 0) q[len++] = q[1];
39         if (dcmp(dot(a - q[2], b - q[2])) < 0) q[len++] = q[2];
40     }
41     q[len++] = b;
42     if (len == 4 && dcmp(dot(q[0] - q[1], q[2] - q[1])) > 0)
43         swap(q[1], q[2]);
44     double res = 0;
45     for (int i = 0; i < len - 1; i++) {
46         if (relation(q[i], c) == 0 || relation(q[i + 1], c) == 0) {
47             double arg = rad(q[i] - p, q[i + 1] - p);
48             res += r * r * arg / 2.0;
49         } else {
50             res += fabs(cross(q[i] - p, q[i + 1] - p)) / 2;
51         }
52     }
53     return res;
54 }

```

半平面交

```
1 struct Line {
2     PT p, v;
3     double ang;
4
5     Line() {}
6     Line(PT from, PT to) : p(from), v(to - from) { ang = atan2(v.y, v.x); }
7     friend bool operator<(Line a, Line b) {
8         return a.ang < b.ang;
9     }
10 };
11
12 bool OnLeft(Line L, PT p) {
13     return dcmp(cross(L.v, p - L.p)) >= 0;
14 }
15
16 PT GetIntersection(Line a, Line b) {
17     PT u = a.p - b.p;
18     ld t = cross(b.v, u) / cross(a.v, b.v);
19     return a.p + a.v * t;
20 }
21
22 vector<PT> HalfplaneIntersection(vector<Line>& L) {
23     int n = L.size();
24     sort(L.begin(), L.end());
25
26     int first, last;
27     vector<PT> p(n);
28     vector<Line> q(n);
29     q[first = last = 0] = L[0];
30     for (int i = 1; i < n; i++) {
31         while (first < last && !OnLeft(L[i], p[last - 1])) last--;
32         while (first < last && !OnLeft(L[i], p[first])) first++;
33         q[++last] = L[i];
34         if (dcmp(cross(q[last].v, q[last - 1].v)) == 0) {
35             last--;
36             if (OnLeft(q[last], L[i].p)) q[last] = L[i];
37         }
38         if (first < last) p[last - 1] = GetIntersection(q[last - 1], q[last]);
39     }
40     while (first < last && !OnLeft(q[first], p[last - 1])) last--;
41     if (last - first <= 1) return vector<PT>();
42     p[last] = GetIntersection(q[last], q[first]);
43
44     return vector<PT>(p.begin() + first, p.begin() + last + 1);
45 }
46
47 vector<PT> convexIntersection(const vector<PT> &v1, const vector<PT> &v2) {
48     vector<Line> h;
49     int n = v1.size(), m = v2.size();
50     for (int i = 0; i < n; ++i)
51         h.push_back(Line(v1[i], v1[(i + 1) % n]));
52     for (int i = 0; i < m; ++i)
53         h.push_back(Line(v2[i], v2[(i + 1) % m]));
54     return HalfplaneIntersection(h);
55 }
56
57 ld ComputeSignedArea(const vector<PT> &p) {
58     ld area = 0;
59     for (unsigned i = 0; i < p.size(); i++) {
60         unsigned j = (i + 1) % p.size();
61         area += p[i].x * p[j].y - p[j].x * p[i].y;
62     }
63     return area / 2.0;
64 }
65
66 ld ComputeArea(const vector<PT> &p) {
67     return fabs(ComputeSignedArea(p));
68 }
```

字符串

后缀自动机

- 广义后缀自动机如果直接使用以下代码的话会产生一些冗余状态（置 last 为 1），所以要用拓扑排序。用 len 基数排序不能。
- 字符集大的话要使用 map。
- 树上 dp 时注意边界（root 和 null）。
- rsort 需要初始化

```
1 namespace sam {
2     const int M = N << 1;
3     int t[M][26], len[M] = {-1}, fa[M], sz = 2, last = 1;
4     void ins(int ch) {
5         int p = last, np = last = sz++;
6         len[np] = len[p] + 1;
7         for (; p && !t[p][ch]; p = fa[p]) t[p][ch] = np;
8         if (!p) { fa[np] = 1; return; }
9         int q = t[p][ch];
10        if (len[p] + 1 == len[q]) fa[np] = q;
11        else {
12            int nq = sz++; len[nq] = len[p] + 1;
13            memcpy(t[nq], t[q], sizeof t[0]);
14            fa[nq] = fa[q];
15            fa[np] = fa[q] = nq;
16            for (; t[p][ch] == q; p = fa[p]) t[p][ch] = nq;
17        }
18    }
19
20    int c[M] = {1}, a[M];
21    void rsort() {
22        FOR (i, 1, sz) c[i] = 0;
23        FOR (i, 1, sz) c[len[i]]++;
24        FOR (i, 1, sz) c[i] += c[i - 1];
25        FOR (i, 1, sz) a[--c[len[i]]] = i;
26    }
27 }
```

- 真·广义后缀自动机

```
1 int t[M][26], len[M] = {-1}, fa[M], sz = 2, last = 1;
2 LL cnt[M][2];
3 void ins(int ch, int id) {
4     int p = last, np = 0, nq = 0, q = -1;
5     if (!t[p][ch]) {
6         np = sz++;
7         len[np] = len[p] + 1;
8         for (; p && !t[p][ch]; p = fa[p]) t[p][ch] = np;
9     }
10    if (!p) fa[np] = 1;
11    else {
12        q = t[p][ch];
13        if (len[p] + 1 == len[q]) fa[np] = q;
14        else {
15            nq = sz++; len[nq] = len[p] + 1;
16            memcpy(t[nq], t[q], sizeof t[0]);
17            fa[nq] = fa[q];
18            fa[np] = fa[q] = nq;
19            for (; t[p][ch] == q; p = fa[p]) t[p][ch] = nq;
20        }
21    }
22    last = np ? np : nq ? nq : q;
23    cnt[last][id] = 1;
24 }
```

- 按字典序建立后缀树注意逆序插入

```
1 void ins(int ch, int pp) {
2     int p = last, np = last = sz++;
3     len[np] = len[p] + 1; one[np] = pos[np] = pp;
4     for (; p && !t[p][ch]; p = fa[p]) t[p][ch] = np;
```

```

5     if (!p) { fa[np] = 1; return; }
6     int q = t[p][ch];
7     if (len[q] == len[p] + 1) fa[np] = q;
8     else {
9         int nq = sz++; len[nq] = len[p] + 1; one[nq] = one[q];
10        memcpy(t[nq], t[q], sizeof t[0]);
11        fa[nq] = fa[q];
12        fa[q] = fa[np] = nq;
13        for (; p && t[p][ch] == q; p = fa[p]) t[p][ch] = nq;
14    }
15 }
16
17 int up[M], c[256] = {2}, a[M];
18 void rsort2() {
19     FOR (i, 1, 256) c[i] = 0;
20     FOR (i, 2, sz) up[i] = s[one[i] + len[fa[i]]];
21     FOR (i, 2, sz) c[up[i]]++;
22     FOR (i, 1, 256) c[i] += c[i - 1];
23     FOR (i, 2, sz) a[--c[up[i]]] = i;
24     FOR (i, 2, sz) G[fa[a[i]]].push_back(a[i]);
25 }

```

- 广义后缀自动机建后缀树，必须反向插入

```

1  int t[M][26], len[M] = {0}, fa[M], sz = 2, last = 1;
2  char* one[M];
3  void ins(int ch, char* pp) {
4      int p = last, np = 0, nq = 0, q = -1;
5      if (!t[p][ch]) {
6          np = sz++; one[np] = pp;
7          len[np] = len[p] + 1;
8          for (; p && !t[p][ch]; p = fa[p]) t[p][ch] = np;
9      }
10     if (!p) fa[np] = 1;
11     else {
12         q = t[p][ch];
13         if (len[p] + 1 == len[q]) fa[np] = q;
14         else {
15             nq = sz++; len[nq] = len[p] + 1; one[nq] = one[q];
16             memcpy(t[nq], t[q], sizeof t[0]);
17             fa[nq] = fa[q];
18             fa[np] = fa[q] = nq;
19             for (; t[p][ch] == q; p = fa[p]) t[p][ch] = nq;
20         }
21     }
22     last = np ? np : nq ? nq : q;
23 }
24 int up[M], c[256] = {2}, aa[M];
25 vector<int> G[M];
26 void rsort() {
27     FOR (i, 1, 256) c[i] = 0;
28     FOR (i, 2, sz) up[i] = *(one[i] + len[fa[i]]);
29     FOR (i, 2, sz) c[up[i]]++;
30     FOR (i, 1, 256) c[i] += c[i - 1];
31     FOR (i, 2, sz) aa[--c[up[i]]] = i;
32     FOR (i, 2, sz) G[fa[aa[i]]].push_back(aa[i]);
33 }

```

- 匹配

```

1  int u = 1, l = 0;
2  FOR (i, 0, strlen(s)) {
3      int ch = s[i] - 'a';
4      while (u && !t[u][ch]) { u = fa[u]; l = len[u]; }
5      ++l; u = t[u][ch];
6      if (!u) u = 1;
7      // do something...
8  }

```

- 获取子串状态

```

1  int get_state(int l, int r) {
2      int u = rpos[r], s = r - l + 1;

```

```

3     FORD (i, SP - 1, -1) if (len[pa[u][i]] >= s) u = pa[u][i];
4     return u;
5 }

```

● 配合 LCT

```

1 namespace lct_sam {
2     extern struct P *const null;
3     const int M = N;
4     struct P {
5         P *fa, *ls, *rs;
6         int last;
7
8         bool has_fa() { return fa->ls == this || fa->rs == this; }
9         bool d() { return fa->ls == this; }
10        P*& c(bool x) { return x ? ls : rs; }
11        P* up() { return this; }
12        void down() {
13            if (ls != null) ls->last = last;
14            if (rs != null) rs->last = last;
15        }
16        void all_down() { if (has_fa()) fa->all_down(); down(); }
17    } *const null = new P{0, 0, 0, 0}, pool[M], *pit = pool;
18    P* G[N];
19    int t[M][26], len[M] = {-1}, fa[M], sz = 2, last = 1;
20
21    void rot(P* o) {
22        bool dd = o->d();
23        P *f = o->fa, *t = o->c(!dd);
24        if (f->has_fa()) f->fa->c(f->d()) = o; o->fa = f->fa;
25        if (t != null) t->fa = f; f->c(dd) = t;
26        o->c(!dd) = f->up(); f->fa = o;
27    }
28    void splay(P* o) {
29        o->all_down();
30        while (o->has_fa()) {
31            if (o->fa->has_fa())
32                rot(o->d() ^ o->fa->d() ? o : o->fa);
33            rot(o);
34        }
35        o->up();
36    }
37    void access(int last, P* u, P* v = null) {
38        if (u == null) { v->last = last; return; }
39        splay(u);
40        P *t = u;
41        while (t->ls != null) t = t->ls;
42        int L = len[fa[t - pool]] + 1, R = len[u - pool];
43
44        if (u->last) bit::add(u->last - R + 2, u->last - L + 2, 1);
45        else bit::add(1, 1, R - L + 1);
46        bit::add(last - R + 2, last - L + 2, -1);
47
48        u->rs = v;
49        access(last, u->up()->fa, u);
50    }
51    void insert(P* u, P* v, P* t) {
52        if (v != null) { splay(v); v->rs = null; }
53        splay(u);
54        u->fa = t; t->fa = v;
55    }
56
57    void ins(int ch, int pp) {
58        int p = last, np = last = sz++;
59        len[np] = len[p] + 1;
60        for (; p && !t[p][ch]; p = fa[p]) t[p][ch] = np;
61        if (!p) fa[np] = 1;
62        else {
63            int q = t[p][ch];
64            if (len[p] + 1 == len[q]) { fa[np] = q; G[np]->fa = G[q]; }
65            else {
66                int nq = sz++; len[nq] = len[p] + 1;

```

```

67         memcpy(t[nq], t[q], sizeof t[0]);
68         insert(G[q], G[fa[q]], G[nq]);
69         G[nq]->last = G[q]->last;
70         fa[nq] = fa[q];
71         fa[np] = fa[q] = nq;
72         G[np]->fa = G[nq];
73         for (; t[p][ch] == q; p = fa[p]) t[p][ch] = nq;
74     }
75 }
76 access(pp + 1, G[np]);
77 }
78
79 void init() {
80     ++pit;
81     FOR (i, 1, N) {
82         G[i] = pit++;
83         G[i]->ls = G[i]->rs = G[i]->fa = null;
84     }
85     G[1] = null;
86 }
87 }

```

回文自动机

```

1  namespace pam {
2      int t[N][26], fa[N], len[N], rs[N], cnt[N];
3      int sz, n, last;
4      int _new(int l) {
5          memset(t[sz], 0, sizeof t[0]);
6          len[sz] = l; cnt[sz] = 0;
7          return sz++;
8      }
9      void init() {
10         rs[n = sz = 0] = -1;
11         last = _new(0);
12         fa[last] = _new(-1);
13     }
14     int get_fa(int x) {
15         while (rs[n - 1 - len[x]] != rs[n]) x = fa[x];
16         return x;
17     }
18     void ins(int ch) {
19         rs[++n] = ch;
20         int p = get_fa(last);
21         if (!t[p][ch]) {
22             int np = _new(len[p] + 2);
23             fa[np] = t[get_fa(fa[p])][ch];
24             t[p][ch] = np;
25         }
26         ++cnt[last = t[p][ch]];
27     }
28 }

```

manacher

```

1  int RL[N];
2  void manacher(int* a, int n) { // "abc" => "#a#b#a#"
3      int r = 0, p = 0;
4      FOR (i, 0, n) {
5          if (i < r) RL[i] = min(RL[2 * p - i], r - i);
6          else RL[i] = 1;
7          while (i - RL[i] >= 0 && i + RL[i] < n && a[i - RL[i]] == a[i + RL[i]])
8              RL[i]++;
9          if (RL[i] + i - 1 > r) { r = RL[i] + i - 1; p = i; }
10     }
11     FOR (i, 0, n) --RL[i];
12 }

```


哈希

内置了自动双哈希开关（小心 TLE）。

```
1  #include <bits/stdc++.h>
2  using namespace std;
3
4  #define ENABLE_DOUBLE_HASH
5
6  typedef long long LL;
7  typedef unsigned long long ULL;
8
9  const int x = 135;
10 const int N = 4e5 + 10;
11 const int p1 = 1e9 + 7, p2 = 1e9 + 9;
12 ULL xp1[N], xp2[N], xp[N];
13
14 void init_xp() {
15     xp1[0] = xp2[0] = xp[0] = 1;
16     for (int i = 1; i < N; ++i) {
17         xp1[i] = xp1[i - 1] * x % p1;
18         xp2[i] = xp2[i - 1] * x % p2;
19         xp[i] = xp[i - 1] * x;
20     }
21 }
22
23 struct String {
24     char s[N];
25     int length, subsize;
26     bool sorted;
27     ULL h[N], hl[N];
28
29     ULL hash() {
30         length = strlen(s);
31         ULL res1 = 0, res2 = 0;
32         h[length] = 0; // ATTENTION!
33         for (int j = length - 1; j >= 0; --j) {
34             #ifdef ENABLE_DOUBLE_HASH
35                 res1 = (res1 * x + s[j]) % p1;
36                 res2 = (res2 * x + s[j]) % p2;
37                 h[j] = (res1 << 32) | res2;
38             #else
39                 res1 = res1 * x + s[j];
40                 h[j] = res1;
41             #endif
42             // printf("%llu\n", h[j]);
43         }
44         return h[0];
45     }
46
47     // 获取子串哈希，左闭右开区间
48     ULL get_substring_hash(int left, int right) const {
49         int len = right - left;
50         #ifdef ENABLE_DOUBLE_HASH
51             // get hash of s[left...right-1]
52             unsigned int mask32 = ~(0u);
53             ULL left1 = h[left] >> 32, right1 = h[right] >> 32;
54             ULL left2 = h[left] & mask32, right2 = h[right] & mask32;
55             return (((left1 - right1 * xp1[len] % p1 + p1) % p1) << 32) |
56                 (((left2 - right2 * xp2[len] % p2 + p2) % p2));
57         #else
58             return h[left] - h[right] * xp[len];
59         #endif
60     }
61
62     void get_all_subs_hash(int sublen) {
63         subsize = length - sublen + 1;
64         for (int i = 0; i < subsize; ++i)
65             hl[i] = get_substring_hash(i, i + sublen);
66         sorted = 0;
67     }
68 }
```

```

68
69 void sort_substring_hash() {
70     sort(hl, hl + subsize);
71     sorted = 1;
72 }
73
74 bool match(ULL key) const {
75     if (!sorted) assert (0);
76     if (!subsize) return false;
77     return binary_search(hl, hl + subsize, key);
78 }
79
80 void init(const char *t) {
81     length = strlen(t);
82     strcpy(s, t);
83 }
84 };
85
86 int LCP(const String &a, const String &b, int ai, int bi) {
87     // Find LCP of a[ai...] and b[bi...]
88     int l = 0, r = min(a.length - ai, b.length - bi);
89     while (l < r) {
90         int mid = (l + r + 1) / 2;
91         if (a.get_substring_hash(ai, ai + mid) == b.get_substring_hash(bi, bi + mid))
92             l = mid;
93         else r = mid - 1;
94     }
95     return l;
96 }
97
98 int check(int ans) {
99     if (T.length < ans) return 1;
100    T.get_all_subs_hash(ans); T.sort_substring_hash();
101    for (int i = 0; i < S.length - ans + 1; ++i)
102        if (!T.match(S.get_substring_hash(i, i + ans)))
103            return 1;
104    return 0;
105 }
106
107 int main() {
108     init_xp(); // DON'T FORGET TO DO THIS!
109
110     for (int tt = 1; tt <= kases; ++tt) {
111         scanf("%d", &n); scanf("%s", str);
112         S.init(str);
113         S.hash(); T.hash();
114     }
115 }

```

后缀数组

构造时间: $O(L \log L)$; 查询时间 $O(\log L)$ 。suffix 数组是排好序的后缀下标, suffix 的反数组是后缀数组。

```

1  #include <bits/stdc++.h>
2  using namespace std;
3
4  const int N = 2e5 + 10;
5  const int Nlog = 18;
6
7  struct SuffixArray {
8      const int L;
9      vector<vector<int>> > P;
10     vector<pair<pair<int, int>, int> > M;
11     int s[N], sa[N], rank[N], height[N];
12     // s: raw string
13     // sa[i]=k: s[k...L-1] ranks i (0 based)
14     // rank[i]=k: the rank of s[i...L-1] is k (0 based)
15     // height[i] = lcp(sa[i-1], sa[i])
16

```

```

17 SuffixArray(const string &raw_s) : L(raw_s.length()), P(1, vector<int>(L, 0)), M(L) {
18     for (int i = 0; i < L; i++)
19         P[0][i] = this->s[i] = int(raw_s[i]);
20     for (int skip = 1, level = 1; skip < L; skip *= 2, level++) {
21         P.push_back(vector<int>(L, 0));
22         for (int i = 0; i < L; i++)
23             M[i] = make_pair(make_pair(P[level - 1][i], i + skip < L ? P[level - 1][i + skip] : -1000), i);
24         sort(M.begin(), M.end());
25         for (int i = 0; i < L; i++)
26             P[level][M[i].second] = (i > 0 && M[i].first == M[i - 1].first) ? P[level][M[i - 1].second] : i;
27     }
28     for (unsigned i = 0; i < P.back().size(); ++i) {
29         rank[i] = P.back()[i];
30         sa[rank[i]] = i;
31     }
32 }
33
34 // This is a traditional way to calculate LCP
35 void getHeight() {
36     memset(height, 0, sizeof height);
37     int k = 0;
38     for (int i = 0; i < L; ++i) {
39         if (rank[i] == 0) continue;
40         if (k) k--;
41         int j = sa[rank[i] - 1];
42         while (i + k < L && j + k < L && s[i + k] == s[j + k]) ++k;
43         height[rank[i]] = k;
44     }
45     rmq_init(height, L);
46 }
47
48 int f[N][Nlog];
49 inline int highbit(int x) {
50     return 31 - __builtin_clz(x);
51 }
52
53 int rmq_query(int x, int y) {
54     int p = highbit(y - x + 1);
55     return min(f[x][p], f[y - (1 << p) + 1][p]);
56 }
57
58 // arr has to be 0 based
59 void rmq_init(int *arr, int length) {
60     for (int x = 0; x <= highbit(length); ++x)
61         for (int i = 0; i <= length - (1 << x); ++i) {
62             if (!x) f[i][x] = arr[i];
63             else f[i][x] = min(f[i][x - 1], f[i + (1 << (x - 1))][x - 1]);
64         }
65 }
66
67 #ifdef NEW
68 // returns the length of the longest common prefix of s[i...L-1] and s[j...L-1]
69 int LongestCommonPrefix(int i, int j) {
70     int len = 0;
71     if (i == j) return L - i;
72     for (int k = (int) P.size() - 1; k >= 0 && i < L && j < L; k--) {
73         if (P[k][i] == P[k][j]) {
74             i += 1 << k;
75             j += 1 << k;
76             len += 1 << k;
77         }
78     }
79     return len;
80 }
81 #else
82 int LongestCommonPrefix(int i, int j) {
83     // getHeight() must be called first
84     if (i == j) return L - i;
85     if (i > j) swap(i, j);
86     return rmq_query(i + 1, j);
87 }

```

```

88     #endif
89
90     int checkNonOverlappingSubstring(int K) {
91         // check if there is two non-overlapping identical substring of length K
92         int minsa = 0, maxsa = 0;
93         for (int i = 0; i < L; ++i) {
94             if (height[i] < K) {
95                 minsa = sa[i]; maxsa = sa[i];
96             } else {
97                 minsa = min(minsa, sa[i]);
98                 maxsa = max(maxsa, sa[i]);
99                 if (maxsa - minsa >= K) return 1;
100             }
101         }
102         return 0;
103     }
104
105     int checkBelongToDifferentSubstring(int K, int split) {
106         int minsa = 0, maxsa = 0;
107         for (int i = 0; i < L; ++i) {
108             if (height[i] < K) {
109                 minsa = sa[i]; maxsa = sa[i];
110             } else {
111                 minsa = min(minsa, sa[i]);
112                 maxsa = max(maxsa, sa[i]);
113                 if (maxsa > split && minsa < split) return 1;
114             }
115         }
116         return 0;
117     }
118 } *S;
119
120
121 int main() {
122     string s, t;
123     cin >> s >> t;
124     int sp = s.length();
125     s += "*" + t;
126     S = new SuffixArray(s);
127     S->getHeight();
128     int left = 0, right = sp;
129     while (left < right) {
130         int mid = (left + right + 1) / 2;
131         if (S->checkBelongToDifferentSubstring(mid, sp))
132             left = mid;
133         else right = mid - 1;
134     }
135     printf("%d\n", left);
136 }

```

- SA-IS
- 仅在后缀自动机被卡内存或者卡常且需要 $O(1)$ LCA 的情况下使用（比赛中敲这个我觉得不行）

```

1 // rk [0..len-1] -> [1..len], sa/ht [1..len]
2 // s[i] > 0 && s[len] = 0
3 template<size_t size>
4 struct SuffixArray {
5     bool type[size << 1];
6     int bucket[size], bucketl[size];
7     int sa[size], rk[size], ht[size];
8     inline bool isLMS(const int i, const bool *type) { return i > 0 && type[i] && !type[i - 1]; }
9     template<class T>
10     inline void inducedSort(T s, int *sa, const int len, const int sigma, const int bucketSize, bool *type, int
11         ↪ *bucket, int *cntbuf, int *p) {
12         memset(bucket, 0, sizeof(int) * sigma);
13         memset(sa, -1, sizeof(int) * len);
14         for (int i = 0; i < len; i++) bucket[s[i]]++;
15         cntbuf[0] = bucket[0];
16         for (int i = 1; i < sigma; i++) cntbuf[i] = cntbuf[i - 1] + bucket[i];
17         for (int i = bucketSize - 1; i >= 0; i--) sa[--cntbuf[s[p[i]]]] = p[i];
18         for (int i = 1; i < sigma; i++) cntbuf[i] = cntbuf[i - 1] + bucket[i - 1];

```

```

18     for (int i = 0; i < len; i++) if (sa[i] > 0 && !type[sa[i] - 1]) sa[cntbuf[s[sa[i] - 1]]++] = sa[i] - 1;
19     cntbuf[0] = bucket[0];
20     for (int i = 1; i < sigma; i++) cntbuf[i] = cntbuf[i - 1] + bucket[i];
21     for (int i = len - 1; i >= 0; i--) if (sa[i] > 0 && type[sa[i] - 1]) sa[--cntbuf[s[sa[i] - 1]]] = sa[i] - 1;
22 }
23 template<typename T>
24 inline void sais(T s, int *sa, int len, bool *type, int *bucket, int *bucket1, int sigma) {
25     int i, j, bucketSize = 0, cnt = 0, p = -1, x, *cntbuf = bucket + sigma;
26     type[len - 1] = 1;
27     for (i = len - 2; i >= 0; i--) type[i] = s[i] < s[i + 1] || (s[i] == s[i + 1] && type[i + 1]);
28     for (i = 1; i < len; i++) if (type[i] && !type[i - 1]) bucket1[bucketSize++] = i;
29     inducedSort(s, sa, len, sigma, bucketSize, type, bucket, cntbuf, bucket1);
30     for (i = bucketSize = 0; i < len; i++) if (isLMS(sa[i], type)) sa[bucketSize++] = sa[i];
31     for (i = bucketSize; i < len; i++) sa[i] = -1;
32     for (i = 0; i < bucketSize; i++) {
33         x = sa[i];
34         for (j = 0; j < len; j++) {
35             if (p == -1 || s[x + j] != s[p + j] || type[x + j] != type[p + j]) { cnt++, p = x; break; }
36             else if (j > 0 && (isLMS(x + j, type) || isLMS(p + j, type))) break;
37         }
38         x = (~x & 1 ? x >> 1 : x - 1 >> 1), sa[bucketSize + x] = cnt - 1;
39     }
40     for (i = j = len - 1; i >= bucketSize; i--) if (sa[i] >= 0) sa[j--] = sa[i];
41     int *s1 = sa + len - bucketSize, *bucket2 = bucket1 + bucketSize;
42     if (cnt < bucketSize) sais(s1, sa, bucketSize, type + len, bucket, bucket1 + bucketSize, cnt);
43     else for (i = 0; i < bucketSize; i++) sa[s1[i]] = i;
44     for (i = 0; i < bucketSize; i++) bucket2[i] = bucket1[sa[i]];
45     inducedSort(s, sa, len, sigma, bucketSize, type, bucket, cntbuf, bucket2);
46 }
47 template<typename T>
48 inline void getHeight(T s, const int len, const int *sa) {
49     for (int i = 0, k = 0; i < len; i++) {
50         if (rk[i] == 0) k = 0;
51         else {
52             if (k > 0) k--;
53             int j = sa[rk[i] - 1];
54             while (i + k < len && j + k < len && s[i + k] == s[j + k]) k++;
55         }
56         ht[rk[i]] = k;
57     }
58 }
59 template<class T>
60 inline void init(T s, int len, int sigma) {
61     sais(s, sa, ++len, type, bucket, bucket1, sigma);
62     for (int i = 1; i < len; i++) rk[sa[i]] = i;
63     getHeight(s, len, sa);
64 }
65 };

```

KMP 自动机

```

1  int m; int pat[N];
2  namespace kmp {
3      int f[N]; // f[i] 表示已匹配成功 i 个, 失配要去哪里
4
5      template<typename T>
6      int go(int stat, T c, bool& acc) {
7          // stat 是当前态 (表示已经匹配了 stat 个字符), c 是要走的边
8          while (stat && c != pat[stat]) stat = f[stat];
9          if (c == pat[stat]) stat++;
10         if (stat == m) acc = true;
11         return stat;
12     }
13
14     void getFail() {
15         static int f2[N];
16         f[0] = f[1] = 0;
17         f2[0] = f2[1] = 0;
18         FOR (i, 1, m) {

```

```

19         int j = f2[i];
20         while (j && pat[i] != pat[j]) j = f2[j];
21         f2[i+1] = f[i+1] == (pat[i] == pat[j]) ? j+1 : 0;
22         if (f[i+1] == j+1 && pat[i+1] == pat[j+1]) f[i+1] = f[j+1];
23     }
24     FOR (i, 0, m) dbg(i, f[i]);
25 }
26 }

```

● 拓展 KMP

```

1  #include <bits/stdc++.h>
2  using namespace std;
3
4  /*
5   Define template S, pattern T, len(S)=n, len(T)=m
6   Find the longest common prefix of T and every suffix of S
7   ex[i]: the LCP between T and S[i..n-1]
8   */
9
10 const int maxn = 1e6 + 10;
11 int nt[maxn], ex[maxn];
12 char s[maxn], t[maxn];
13
14 void get_next(char *str) {
15     int i = 0, j, po, len = strlen(str);
16     nt[0] = len;
17     while (str[i] == str[i + 1] && i + 1 < len)
18         i++;
19     nt[1] = i;
20     po = 1;
21     for (i = 2; i < len; i++) {
22         if (nt[i - po] + i < nt[po] + po)
23             nt[i] = nt[i - po];
24         else {
25             j = nt[po] + po - i;
26             if (j < 0) j = 0;
27             while (i + j < len && str[j] == str[j + i])
28                 j++;
29             nt[i] = j;
30             po = i;
31         }
32     }
33 }
34
35 void exkmp(char *s1, char *s2) {
36     int i = 0, j, po, len = strlen(s1), l2 = strlen(s2);
37     get_next(s2);
38     while (s1[i] == s2[i] && i < l2 && i < len)
39         i++;
40     ex[0] = i;
41     po = 0;
42     for (i = 1; i < len; i++) {
43         if (nt[i - po] + i < ex[po] + po)
44             ex[i] = nt[i - po];
45         else {
46             j = ex[po] + po - i;
47             if (j < 0) j = 0;
48             while (i + j < len && j < l2 && s1[j + i] == s2[j])
49                 j++;
50             ex[i] = j;
51             po = i;
52         }
53     }
54 }
55
56 int main() {
57     const int modn = 1e9 + 7;
58     int T; scanf("%d", &T);
59     while (T--) {
60         memset(nt, 0, sizeof nt);
61         memset(ex, 0, sizeof ex);

```

```

62     scanf("%s", s); scanf("%s", t);
63     int slen = strlen(s), tlen = strlen(t);
64     reverse(s, s + slen);
65     reverse(t, t + tlen);
66     exkmp(s, t);
67     int ans = 0;
68     for (int i = 0; i < slen; ++i)
69         ans = (ans + 1LL * ex[i] * (ex[i] + 1) / 2) % modn;
70     printf("%d\n", ans);
71 }
72 }

```

Trie

```

1  namespace trie {
2      int t[N][26], sz, ed[N];
3      void init() { sz = 2; memset(ed, 0, sizeof ed); }
4      int _new() { memset(t[sz], 0, sizeof t[sz]); return sz++; }
5      void ins(char* s, int p) {
6          int u = 1;
7          FOR (i, 0, strlen(s)) {
8              int c = s[i] - 'a';
9              if (!t[u][c]) t[u][c] = _new();
10             u = t[u][c];
11         }
12         ed[u] = p;
13     }
14 }

```

AC 自动机

```

1  const int N = 1e6 + 100, M = 26;
2
3  int mp(char ch) { return ch - 'a'; }
4
5  struct ACA {
6      int ch[N][M], danger[N], fail[N];
7      int sz;
8      void init() {
9          sz = 1;
10         memset(ch[0], 0, sizeof ch[0]);
11         memset(danger, 0, sizeof danger);
12     }
13     void insert(const string &s, int m) {
14         int n = s.size(); int u = 0, c;
15         FOR (i, 0, n) {
16             c = mp(s[i]);
17             if (!ch[u][c]) {
18                 memset(ch[sz], 0, sizeof ch[sz]);
19                 danger[sz] = 0; ch[u][c] = sz++;
20             }
21             u = ch[u][c];
22         }
23         danger[u] |= 1 << m;
24     }
25     void build() {
26         queue<int> Q;
27         fail[0] = 0;
28         for (int c = 0, u; c < M; c++) {
29             u = ch[0][c];
30             if (u) { Q.push(u); fail[u] = 0; }
31         }
32         while (!Q.empty()) {
33             int r = Q.front(); Q.pop();
34             danger[r] |= danger[fail[r]];
35             for (int c = 0, u; c < M; c++) {
36                 u = ch[r][c];

```

```

37         if (!u) {
38             ch[r][c] = ch[fail[r]][c];
39             continue;
40         }
41         fail[u] = ch[fail[r]][c];
42         Q.push(u);
43     }
44 }
45 }
46 } ac;
47
48 char s[N];
49
50 int main() {
51     int n; scanf("%d", &n);
52     ac.init();
53     while (n--) {
54         scanf("%s", s);
55         ac.insert(s, 0);
56     }
57     ac.build();
58
59     scanf("%s", s);
60     int u = 0; n = strlen(s);
61     FOR (i, 0, n) {
62         u = ac.ch[u][mp(s[i])];
63         if (ac.danger[u]) {
64             puts("YES");
65             return 0;
66         }
67     }
68     puts("NO");
69     return 0;
70 }

```

杂项

STL

- copy

```

1 template <class InputIterator, class OutputIterator>
2 OutputIterator copy (InputIterator first, InputIterator last, OutputIterator result);

```

- merge (如果相等，第一个优先)

```

1 template <class InputIterator1, class InputIterator2,
2           class OutputIterator, class Compare>
3 OutputIterator merge (InputIterator1 first1, InputIterator1 last1,
4                       InputIterator2 first2, InputIterator2 last2,
5                       OutputIterator result, Compare comp);

```

- for_each

```

1 template <class InputIterator, class Function>
2 Function for_each (InputIterator first, InputIterator last, Function fn);

```

- transform

```

1 template <class InputIterator, class OutputIterator, class UnaryOperation>
2 OutputIterator transform (InputIterator first1, InputIterator last1,
3                           OutputIterator result, UnaryOperation op);

```

- numeric_limits

```

1 template <class T> numeric_limits;

```

- iota


```

1  template< class ForwardIterator, class T >
2  void iota( ForwardIterator first, ForwardIterator last, T value );

```

日期

```

1  // Routines for performing computations on dates. In these routines,
2  // months are expressed as integers from 1 to 12, days are expressed
3  // as integers from 1 to 31, and years are expressed as 4-digit
4  // integers.
5
6  string dayOfWeek[] = {"Mo", "Tu", "We", "Th", "Fr", "Sa", "Su"};
7
8  // converts Gregorian date to integer (Julian day number)
9
10 int DateToInt (int m, int d, int y){
11     return
12         1461 * (y + 4800 + (m - 14) / 12) / 4 +
13         367 * (m - 2 - (m - 14) / 12 * 12) / 12 -
14         3 * ((y + 4900 + (m - 14) / 12) / 100) / 4 +
15         d - 32075;
16 }
17
18 // converts integer (Julian day number) to Gregorian date: month/day/year
19
20 void IntToDate (int jd, int &m, int &d, int &y){
21     int x, n, i, j;
22
23     x = jd + 68569;
24     n = 4 * x / 146097;
25     x -= (146097 * n + 3) / 4;
26     i = (4000 * (x + 1)) / 1461001;
27     x -= 1461 * i / 4 - 31;
28     j = 80 * x / 2447;
29     d = x - 2447 * j / 80;
30     x = j / 11;
31     m = j + 2 - 12 * x;
32     y = 100 * (n - 49) + i + x;
33 }
34
35 // converts integer (Julian day number) to day of week
36
37 string IntToDay (int jd){
38     return dayOfWeek[jd % 7];
39 }

```

子集枚举

- 枚举真子集

```

1  for (int s = (S - 1) & S; s; s = (s - 1) & S)

```

- 枚举大小为 k 的子集

```

1  template<typename T>
2  void subset(int k, int n, T&& f) {
3      int t = (1 << k) - 1;
4      while (t < 1 << n) {
5          f(t);
6          int x = t & -t, y = t + x;
7          t = ((t & ~y) / x >> 1) | y;
8      }
9  }

```

权值最大上升子序列

```
1  const LL maxn = 1E5 + 10;
2  const LL INF = 1E10;
3  struct P {
4      LL k, v;
5      bool operator < (const P& rhs) const {
6          return k < rhs.k || (k == rhs.k && v < rhs.v);
7      }
8  };
9  LL k[maxn], v[maxn], n, T;
10 set<P> s;
11
12 int main() {
13     cin >> T;
14     while (T--) {
15         s.clear();
16         s.insert({-INF, 0});
17         cin >> n;
18         FOR (i, 0, n) scanf("%lld", &k[i]);
19         FOR (i, 0, n) scanf("%lld", &v[i]);
20         FOR (i, 0, n) {
21             auto it = s.lower_bound({k[i], INF});
22             LL vv = (--it)->v + v[i];
23             ++it;
24             while (it != s.end() && it->v <= vv)
25                 it = s.erase(it);
26             if (it == s.end() || it->k != k[i]) s.insert({k[i], vv});
27         }
28         cout << s.rbegin()->v << endl;
29     }
30 }
```

数位 DP

```
1  LL dfs(LL base, LL pos, LL len, LL s, bool limit) {
2      if (pos == -1) return s ? base : 1;
3      if (!limit && dp[base][pos][len][s] != -1) return dp[base][pos][len][s];
4      LL ret = 0;
5      LL ed = limit ? a[pos] : base - 1;
6      FOR (i, 0, ed + 1) {
7          tmp[pos] = i;
8          if (len == pos)
9              ret += dfs(base, pos - 1, len - (i == 0), s, limit && i == a[pos]);
10             else if (s && pos < (len + 1) / 2)
11                 ret += dfs(base, pos - 1, len, tmp[len - pos] == i, limit && i == a[pos]);
12             else
13                 ret += dfs(base, pos - 1, len, s, limit && i == a[pos]);
14         }
15         if (!limit) dp[base][pos][len][s] = ret;
16         return ret;
17     }
18
19     LL solve(LL x, LL base) {
20         LL sz = 0;
21         while (x) {
22             a[sz++] = x % base;
23             x /= base;
24         }
25         return dfs(base, sz - 1, sz - 1, 1, true);
26     }
```

土制 bitset

```
1  // M 要开大至少 1 个 64
2  const int M = (1E4 + 200) / 64;
3  typedef unsigned long long ULL;
```

```

4  const ULL ONE = 1;
5
6  struct Bitset {
7      ULL a[M];
8      void go(int x) {
9          int offset = x / 64; x %= 64;
10         for (int i = offset, j = 0; i + 1 < M; ++i, ++j) {
11             a[j] |= a[i] >> x;
12             if (x) a[j] |= a[i + 1] << (64 - x); // 不能左移 64 位
13         }
14     }
15     void init() { memset(a, 0, sizeof a); }
16     void set(int x) {
17         int offset = x / 64; x %= 64;
18         a[offset] |= (ONE << x);
19     }
20     void prt() {
21         FOR (i, 0, M) FOR (j, 0, 64) putchar((a[i] & (ONE << j)) ? '1' : '0');
22         puts("");
23     }
24     int lowbit() {
25         FOR (i, 0, M) if (a[i]) return i * 64 + __builtin_ctzll(a[i]);
26         assert (0);
27     }
28     int highbit(int x) {
29         // [0,x) 的最高位
30         int offset = x / 64; x %= 64;
31         FOR (i, offset, -1) {
32             if (!a[i]) continue;
33             if (i == offset) {
34                 FOR (j, x - 1, -1) if ((ONE << j) & a[i]) { return i * 64 + j; }
35             } else return i * 64 + 63 - __builtin_clzll(a[i]);
36         }
37         assert (0);
38     }
39 };

```

随机

- 不要使用 rand()。
- chrono::steady_clock::now().time_since_epoch().count() 可用于计时。
- 64 位可以使用 mt19937_64。

```

1  int main() {
2      mt19937 rng(chrono::steady_clock::now().time_since_epoch().count());
3      vector<int> permutation(N);
4
5      for (int i = 0; i < N; i++)
6          permutation[i] = i;
7      shuffle(permutation.begin(), permutation.end(), rng);
8
9      for (int i = 0; i < N; i++)
10         permutation[i] = i;
11     for (int i = 1; i < N; i++)
12         swap(permutation[i], permutation[uniform_int_distribution<int>(0, i)(rng)]);

```

伪随机数

```

1  unsigned rnd() {
2      static unsigned A = 1 << 16 | 3, B = 33333331, C = 2341;
3      return C = A * C + B;
4  }

```

Java

Regex

```
1 // Code which demonstrates the use of Java's regular expression libraries.
2 // This is a solution for
3 //
4 // Loglan: a logical language
5 // http://acm.uva.es/p/v1/134.html
6
7 import java.util.*;
8 import java.util.regex.*;
9
10 public class LogLan {
11
12     public static void main(String args[]) {
13
14         String regex = BuildRegex();
15         Pattern pattern = Pattern.compile(regex);
16
17         Scanner s = new Scanner(System.in);
18         while (true) {
19
20             // In this problem, each sentence consists of multiple lines, where the last
21             // line is terminated by a period. The code below reads lines until
22             // encountering a line whose final character is a '.'. Note the use of
23             //
24             // s.length() to get length of string
25             // s.charAt() to extract characters from a Java string
26             // s.trim() to remove whitespace from the beginning and end of Java string
27             //
28             // Other useful String manipulation methods include
29             //
30             // s.compareTo(t) < 0 if s < t, lexicographically
31             // s.indexOf("apple") returns index of first occurrence of "apple" in s
32             // s.lastIndexOf("apple") returns index of last occurrence of "apple" in s
33             // s.replace(c,d) replaces occurrences of character c with d
34             // s.startsWith("apple") returns (s.indexOf("apple") == 0)
35             // s.toLowerCase() / s.toUpperCase() returns a new lower/uppercased string
36             //
37             // Integer.parseInt(s) converts s to an integer (32-bit)
38             // Long.parseLong(s) converts s to a long (64-bit)
39             // Double.parseDouble(s) converts s to a double
40
41             String sentence = "";
42             while (true) {
43                 sentence = (sentence + " " + s.nextLine()).trim();
44                 if (sentence.equals("#")) return;
45                 if (sentence.charAt(sentence.length() - 1) == '.') break;
46             }
47
48             // now, we remove the period, and match the regular expression
49
50             String removed_period = sentence.substring(0, sentence.length() - 1).trim();
51             if (pattern.matcher(removed_period).find()) {
52                 System.out.println("Good");
53             } else {
54                 System.out.println("Bad!");
55             }
56         }
57     }
58 }
```

Decimal Format

```
1 // examples for printing floating point numbers
2
3 import java.util.*;
4 import java.io.*;
```

```

5  import java.text.DecimalFormat;
6
7  public class DecFormat {
8      public static void main(String[] args) {
9          DecimalFormat fmt;
10
11          // round to at most 2 digits, leave of digits if not needed
12          fmt = new DecimalFormat("#.##");
13          System.out.println(fmt.format(12345.6789)); // produces 12345.68
14          System.out.println(fmt.format(12345.0)); // produces 12345
15          System.out.println(fmt.format(0.0)); // produces 0
16          System.out.println(fmt.format(0.01)); // produces .1
17
18          // round to precisely 2 digits
19          fmt = new DecimalFormat("#.00");
20          System.out.println(fmt.format(12345.6789)); // produces 12345.68
21          System.out.println(fmt.format(12345.0)); // produces 12345.00
22          System.out.println(fmt.format(0.0)); // produces .00
23
24          // round to precisely 2 digits, force leading zero
25          fmt = new DecimalFormat("0.00");
26          System.out.println(fmt.format(12345.6789)); // produces 12345.68
27          System.out.println(fmt.format(12345.0)); // produces 12345.00
28          System.out.println(fmt.format(0.0)); // produces 0.00
29
30          // round to precisely 2 digits, force leading zeros
31          fmt = new DecimalFormat("000000000.00");
32          System.out.println(fmt.format(12345.6789)); // produces 000012345.68
33          System.out.println(fmt.format(12345.0)); // produces 000012345.00
34          System.out.println(fmt.format(0.0)); // produces 000000000.00
35
36          // force leading '+'
37          fmt = new DecimalFormat("+0;-0");
38          System.out.println(fmt.format(12345.6789)); // produces +12346
39          System.out.println(fmt.format(-12345.6789)); // produces -12346
40          System.out.println(fmt.format(0)); // produces +0
41
42          // force leading positive/negative, pad to 2
43          fmt = new DecimalFormat("positive 00;negative 0");
44          System.out.println(fmt.format(1)); // produces "positive 01"
45          System.out.println(fmt.format(-1)); // produces "negative 01"
46
47          // quote special chars (#)
48          fmt = new DecimalFormat("text with '#' followed by #");
49          System.out.println(fmt.format(12.34)); // produces "text with # followed by 12"
50
51          // always show "."
52          fmt = new DecimalFormat("#.##");
53          fmt.setDecimalSeparatorAlwaysShown(true);
54          System.out.println(fmt.format(12.34)); // produces "12.3"
55          System.out.println(fmt.format(12)); // produces "12."
56          System.out.println(fmt.format(0.34)); // produces "0.3"
57
58          // different grouping distances:
59          fmt = new DecimalFormat("#,###.###");
60          System.out.println(fmt.format(123456789.123)); // produces "1,2345,6789.123"
61
62          // scientific:
63          fmt = new DecimalFormat("0.000E00");
64          System.out.println(fmt.format(123456789.123)); // produces "1.235E08"
65          System.out.println(fmt.format(-0.000234)); // produces "-2.34E-04"
66
67          // using variable number of digits:
68          fmt = new DecimalFormat("0");
69          System.out.println(fmt.format(123.123)); // produces "123"
70          fmt.setMinimumFractionDigits(8);
71          System.out.println(fmt.format(123.123)); // produces "123.12300000"
72          fmt.setMaximumFractionDigits(0);
73          System.out.println(fmt.format(123.123)); // produces "123"
74
75          // note: to pad with spaces, you need to do it yourself:

```

```

76         // String out = fmt.format(...)
77         // while (out.length() < targlength) out = " "+out;
78     }
79 }

```

Sort

```

1  import java.util.ArrayList;
2  import java.util.Collections;
3  import java.util.List;
4
5  public class Employee implements Comparable<Employee> {
6      private int id;
7      private String name;
8      private int age;
9
10     public Employee(int id, String name, int age) {
11         this.id = id;
12         this.name = name;
13         this.age = age;
14     }
15
16     @Override
17     public int compareTo(Employee o) {
18         if (id > o.id) {
19             return 1;
20         } else if (id < o.id) {
21             return -1;
22         }
23         return 0;
24     }
25
26     public static void main(String[] args) {
27         List<Employee> list = new ArrayList<Employee>();
28         list.add(new Employee(2, "Java", 20));
29         list.add(new Employee(1, "C", 30));
30         list.add(new Employee(3, "C#", 10));
31         Collections.sort(list);
32     }
33 }

```

扩栈（本地使用）

```

1  #include <sys/resource.h>
2  void init_stack(){
3      const rlim_t kStackSize = 512 * 1024 * 1024;
4      struct rlimit rl;
5      int result;
6      result = getrlimit(RLIMIT_STACK, &rl);
7      if (result == 0) {
8          if (rl.rlim_cur < kStackSize) {
9              rl.rlim_cur = kStackSize;
10             result = setrlimit(RLIMIT_STACK, &rl);
11             if (result != 0) {
12                 fprintf(stderr, "setrlimit returned result = %d\n", result);
13             }
14         }
15     }
16 }

```

心态崩了

- (int)v.size()
- 1LL << k
- 递归函数用全局或者 static 变量要小心

- 预处理组合数注意上限
- 想清楚到底是要 `multiset` 还是 `set`
- 提交之前看一下数据范围，测一下边界
- 数据结构注意数组大小（2 倍，4 倍）
- 字符串注意数据集
- 如果函数中使用了默认参数的话，注意调用时的参数个数。
- 注意要读完
- 构造参数无法使用自己
- 树链剖分/dfs 序，初始化或者询问不要忘记 `idx`, `ridx`
- 排序时注意结构体的所有属性是不是考虑了
- 不要把 `while` 写成 `if`
- 不要把 `int` 开成 `char`
- 清零的时候全部用 `0~n+1`。
- 模意义下不要用除法
- 哈希不要自然溢出
- 最短路不要 SPFA，乖乖写 Dijkstra
- 上取整以及 GCD 小心负数
- `mid` 用 $\lfloor \frac{l + (r - l)}{2} \rfloor$ 可以避免溢出和负数的问题