

Standard Code Library

FOREIGNERS

East China Normal University

October 8102

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一切的开始

宏定义

- 需要 C++11

```
1 #include <bits/stdc++.h>
2 using namespace std;
3 using LL = long long;
4 #define FOR(i, x, y) for (decay<decltype(y)>::type i = (x), _##i = (y); i < _##i; ++i)
5 #define FORD(i, x, y) for (decay<decltype(x)>::type i = (x), _##i = (y); i > _##i; --i)
6 #ifdef zero1
7 #define dbg(x...) do { cout << "\033[32;1m" << #x << " -> "; err(x); } while (0)
8 #else
9 #define dbg(...)
10 #endif
11 void err() { cout << "\033[39;0m" << endl; }
12 template<typename...> class T, typename t, typename... A>
13 void err(T<t> a, A... x) { for (auto v: a) cout << v << ' '; err(x...); }
14 template<typename T, typename... A>
15 void err(T a, A... x) { cout << a << ' '; err(x...); }
16 // -----
```

- 更多配色:

- 33 黄色
- 34 蓝色
- 31 橙色

- POJ/BZOJ version

```
1 #include <cstdio>
2 #include <iostream>
3 #include <algorithm>
4 #include <cmath>
5 #include <string>
6 #include <vector>
7 #include <set>
8 #include <queue>
9 #include <cstring>
10 #include <cassert>
11 using namespace std;
12 typedef long long LL;
13 #define FOR(i, x, y) for (LL i = (x), _##i = (y); i < _##i; ++i)
14 #define FORD(i, x, y) for (LL i = (x), _##i = (y); i > _##i; --i)
15 #ifdef zero1
16 #define dbg(args...) do { cout << "\033[32;1m" << #args << " -> "; err(args); } while (0)
17 #else
18 #define dbg(...)
19 #endif
20 void err() { cout << "\033[39;0m" << endl; }
21 template<typename T, typename... Args>
22 void err(T a, Args... args) {
23     cout << a << ' '; err(args...);
24 }
25 // -----
```

- HDU Assert Patch

```
1 #ifdef ONLINE_JUDGE
2 #define assert(condition) if (!(condition)) { int x = 1, y = 0; cout << x / y << endl; }
3 #endif
```

快速读

```
1 inline char nc() {
2     static char buf[100000], *p1 = buf, *p2 = buf;
3     return p1 == p2 && (p2 = (p1 = buf) + fread(buf, 1, 100000, stdin), p1 == p2) ? EOF : *p1++;
4 }
5 template <typename T>
```

```

6  bool rn(T& v) {
7      static char ch;
8      while (ch != EOF && !isdigit(ch)) ch = nc();
9      if (ch == EOF) return false;
10     for (v = 0; isdigit(ch); ch = nc())
11         v = v * 10 + ch - '0';
12     return true;
13 }
14
15 template <typename T>
16 void o(T p) {
17     static int stk[70], tp;
18     if (p == 0) { putchar('0'); return; }
19     if (p < 0) { p = -p; putchar('-'); }
20     while (p) stk[++tp] = p % 10, p /= 10;
21     while (tp) putchar(stk[tp--] + '0');
22 }

```

- 需要初始化
- 需要一次读入
- 不支持负数

```

1  const int MAXS = 100 * 1024 * 1024;
2  char buf[MAXS];
3  template<typename T>
4  inline bool read(T& x) {
5      static char* p = buf;
6      x = 0;
7      while (*p && !isdigit(*p)) ++p;
8      if (!*p) return false;
9      while (isdigit(*p)) x = x * 10 + *p++ - 48;
10     return true;
11 }
12
13 fread(buf, 1, MAXS, stdin);

```

对拍

```

1  #!/usr/bin/env bash
2  g++ -o r main.cpp -O2 -std=c++11
3  g++ -o std std.cpp -O2 -std=c++11
4  while true; do
5      python gen.py > in
6      ./std < in > stdout
7      ./r < in > out
8      if test $? -ne 0; then
9          exit 0
10     fi
11     if diff stdout out; then
12         printf "AC\n"
13     else
14         printf "GG\n"
15         exit 0
16     fi
17 done

```

为什么 C++ 不自带这个？

```

1  LL bin(LL x, LL n, LL MOD) {
2      LL ret = MOD - 1;
3      for (x %= MOD; n; n >>= 1, x = x * x % MOD)
4          if (n & 1) ret = ret * x % MOD;
5      return ret;
6  }
7  inline LL get_inv(LL x, LL p) { return bin(x, p - 2, p); }

```

数据结构

ST 表

- 二维

```
1 int f[maxn][maxn][10][10];
2 inline int highbit(int x) { return 31 - __builtin_clz(x); }
3 inline int calc(int x, int y, int xx, int yy, int p, int q) {
4     return max(
5         max(f[x][y][p][q], f[xx - (1 << p) + 1][yy - (1 << q) + 1][p][q]),
6         max(f[xx - (1 << p) + 1][y][p][q], f[x][yy - (1 << q) + 1][p][q])
7     );
8 }
9 void init() {
10     FOR (x, 0, highbit(n) + 1)
11     FOR (y, 0, highbit(m) + 1)
12         FOR (i, 0, n - (1 << x) + 1)
13             FOR (j, 0, m - (1 << y) + 1) {
14                 if (!x && !y) { f[i][j][x][y] = a[i][j]; continue; }
15                 f[i][j][x][y] = calc(
16                     i, j,
17                     i + (1 << x) - 1, j + (1 << y) - 1,
18                     max(x - 1, 0), max(y - 1, 0)
19                 );
20             }
21 }
22 inline int get_max(int x, int y, int xx, int yy) {
23     return calc(x, y, xx, yy, highbit(xx - x + 1), highbit(yy - y + 1));
24 }
```

- 一维

```
1 struct RMQ {
2     int f[22][M];
3     inline int highbit(int x) { return 31 - __builtin_clz(x); }
4     void init(int* v, int n) {
5         FOR (i, 0, n) f[0][i] = v[i];
6         FOR (x, 1, highbit(n) + 1)
7             FOR (i, 0, n - (1 << x) + 1)
8                 f[x][i] = min(f[x - 1][i], f[x - 1][i + (1 << (x - 1))]);
9     }
10     int get_min(int l, int r) {
11         assert(l <= r);
12         int t = highbit(r - l + 1);
13         return min(f[t][l], f[t][r - (1 << t) + 1]);
14     }
15 } rmq;
```

线段树

- 普适

```
1 namespace sg {
2     struct Q {
3         LL setv;
4         explicit Q(LL setv = -1): setv(setv) {}
5         void operator += (const Q& q) { if (q.setv != -1) setv = q.setv; }
6     };
7     struct P {
8         LL min;
9         explicit P(LL min = INF): min(min) {}
10        void up(Q& q) { if (q.setv != -1) min = q.setv; }
11    };
12    template<typename T>
13    P operator & (T&& a, T&& b) {
14        return P(min(a.min, b.min));
15    }
16 }
```

```

16     P p[maxn << 2];
17     Q q[maxn << 2];
18     #define lson o * 2, l, (l + r) / 2
19     #define rson o * 2 + 1, (l + r) / 2 + 1, r
20     void up(int o, int l, int r) {
21         if (l == r) p[o] = P();
22         else p[o] = p[o * 2] & p[o * 2 + 1];
23         p[o].up(q[o]);
24     }
25     void down(int o, int l, int r) {
26         q[o * 2] += q[o]; q[o * 2 + 1] += q[o];
27         q[o] = Q();
28         up(lson); up(rson);
29     }
30     template<typename T>
31     void build(T&& f, int o = 1, int l = 1, int r = n) {
32         if (l == r) q[o] = f(l);
33         else { build(f, lson); build(f, rson); q[o] = Q(); }
34         up(o, l, r);
35     }
36     P query(int ql, int qr, int o = 1, int l = 1, int r = n) {
37         if (ql > r || l > qr) return P();
38         if (ql <= l && r <= qr) return p[o];
39         down(o, l, r);
40         return query(ql, qr, lson) & query(ql, qr, rson);
41     }
42     void update(int ql, int qr, const Q& v, int o = 1, int l = 1, int r = n) {
43         if (ql > r || l > qr) return;
44         if (ql <= l && r <= qr) q[o] += v;
45         else {
46             down(o, l, r);
47             update(ql, qr, v, lson); update(ql, qr, v, rson);
48         }
49         up(o, l, r);
50     }
51 }

```

• SET+ADD

```

1     struct IntervalTree {
2         #define ls o * 2, l, m
3         #define rs o * 2 + 1, m + 1, r
4         static const LL M = maxn * 4, RS = 1E18 - 1;
5         LL addv[M], setv[M], minv[M], maxv[M], sumv[M];
6         void init() {
7             memset(addv, 0, sizeof addv);
8             fill(setv, setv + M, RS);
9             memset(minv, 0, sizeof minv);
10            memset(maxv, 0, sizeof maxv);
11            memset(sumv, 0, sizeof sumv);
12        }
13        void maintain(LL o, LL l, LL r) {
14            if (l < r) {
15                LL lc = o * 2, rc = o * 2 + 1;
16                sumv[o] = sumv[lc] + sumv[rc];
17                minv[o] = min(minv[lc], minv[rc]);
18                maxv[o] = max(maxv[lc], maxv[rc]);
19            } else sumv[o] = minv[o] = maxv[o] = 0;
20            if (setv[o] != RS) { minv[o] = maxv[o] = setv[o]; sumv[o] = setv[o] * (r - l + 1); }
21            if (addv[o]) { minv[o] += addv[o]; maxv[o] += addv[o]; sumv[o] += addv[o] * (r - l + 1); }
22        }
23        void build(LL o, LL l, LL r) {
24            if (l == r) addv[o] = a[l];
25            else {
26                LL m = (l + r) / 2;
27                build(ls); build(rs);
28            }
29            maintain(o, l, r);
30        }
31        void pushdown(LL o) {
32            LL lc = o * 2, rc = o * 2 + 1;
33            if (setv[o] != RS) {

```



```

34         setv[lc] = setv[rc] = setv[o];
35         addv[lc] = addv[rc] = 0;
36         setv[o] = RS;
37     }
38     if (addv[o]) {
39         addv[lc] += addv[o]; addv[rc] += addv[o];
40         addv[o] = 0;
41     }
42 }
43 void update(LL p, LL q, LL o, LL l, LL r, LL v, LL op) {
44     if (p <= r && l <= q)
45     if (p <= l && r <= q) {
46         if (op == 2) { setv[o] = v; addv[o] = 0; }
47         else addv[o] += v;
48     } else {
49         pushdown(o);
50         LL m = (l + r) / 2;
51         update(p, q, ls, v, op); update(p, q, rs, v, op);
52     }
53     maintain(o, l, r);
54 }
55 void query(LL p, LL q, LL o, LL l, LL r, LL add, LL& ssum, LL& smin, LL& smax) {
56     if (p > r || l > q) return;
57     if (setv[o] != RS) {
58         LL v = setv[o] + add + addv[o];
59         ssum += v * (min(r, q) - max(l, p) + 1);
60         smin = min(smin, v);
61         smax = max(smax, v);
62     } else if (p <= l && r <= q) {
63         ssum += sumv[o] + add * (r - l + 1);
64         smin = min(smin, minv[o] + add);
65         smax = max(smax, maxv[o] + add);
66     } else {
67         LL m = (l + r) / 2;
68         query(p, q, ls, add + addv[o], ssum, smin, smax);
69         query(p, q, rs, add + addv[o], ssum, smin, smax);
70     }
71 }
72 } IT;

```

均摊复杂度线段树

- 区间取 max, 区间求和。

```

1  namespace R {
2      #define lson o * 2, l, (l + r) / 2
3      #define rson o * 2 + 1, (l + r) / 2 + 1, r
4      int m1[N], m2[N], cm1[N];
5      LL sum[N];
6      void up(int o) {
7          int lc = o * 2, rc = lc + 1;
8          m1[o] = max(m1[lc], m1[rc]);
9          sum[o] = sum[lc] + sum[rc];
10         if (m1[lc] == m1[rc]) {
11             cm1[o] = cm1[lc] + cm1[rc];
12             m2[o] = max(m2[lc], m2[rc]);
13         } else {
14             cm1[o] = m1[lc] > m1[rc] ? cm1[lc] : cm1[rc];
15             m2[o] = max(min(m1[lc], m1[rc]), max(m2[lc], m2[rc]));
16         }
17     }
18     void mod(int o, int x) {
19         if (x >= m1[o]) return;
20         assert(x > m2[o]);
21         sum[o] -= 1LL * (m1[o] - x) * cm1[o];
22         m1[o] = x;
23     }
24     void down(int o) {
25         int lc = o * 2, rc = lc + 1;

```

```

26     mod(lc, m1[o]); mod(rc, m1[o]);
27 }
28 void build(int o, int l, int r) {
29     if (l == r) { int t; read(t); sum[o] = m1[o] = t; m2[o] = -1; cm1[o] = 1; }
30     else { build(lson); build(rson); up(o); }
31 }
32 void update(int ql, int qr, int x, int o, int l, int r) {
33     if (r < ql || qr < l || m1[o] <= x) return;
34     if (ql <= l && r <= qr && m2[o] < x) { mod(o, x); return; }
35     down(o);
36     update(ql, qr, x, lson); update(ql, qr, x, rson);
37     up(o);
38 }
39 int qmax(int ql, int qr, int o, int l, int r) {
40     if (r < ql || qr < l) return -INF;
41     if (ql <= l && r <= qr) return m1[o];
42     down(o);
43     return max(qmax(ql, qr, lson), qmax(ql, qr, rson));
44 }
45 LL qsum(int ql, int qr, int o, int l, int r) {
46     if (r < ql || qr < l) return 0;
47     if (ql <= l && r <= qr) return sum[o];
48     down(o);
49     return qsum(ql, qr, lson) + qsum(ql, qr, rson);
50 }
51 }

```

持久化线段树

• ADD

```

1 namespace tree {
2     #define mid ((l + r) >> 1)
3     #define lson ql, qr, l, mid
4     #define rson ql, qr, mid + 1, r
5     struct P {
6         LL add, sum;
7         int ls, rs;
8     } tr[maxn * 45 * 2];
9     int sz = 1;
10    int N(LL add, int l, int r, int ls, int rs) {
11        tr[sz] = {add, tr[ls].sum + tr[rs].sum + add * (len[r] - len[l - 1]), ls, rs};
12        return sz++;
13    }
14    int update(int o, int ql, int qr, int l, int r, LL add) {
15        if (ql > r || l > qr) return o;
16        const P& t = tr[o];
17        if (ql <= l && r <= qr) return N(add + t.add, l, r, t.ls, t.rs);
18        return N(t.add, l, r, update(t.ls, lson, add), update(t.rs, rson, add));
19    }
20    LL query(int o, int ql, int qr, int l, int r, LL add = 0) {
21        if (ql > r || l > qr) return 0;
22        const P& t = tr[o];
23        if (ql <= l && r <= qr) return add * (len[r] - len[l - 1]) + t.sum;
24        return query(t.ls, lson, add + t.add) + query(t.rs, rson, add + t.add);
25    }
26 }

```

K-D Tree

最优化问题一定要用全局变量大力剪枝，而且左右儿子先递归潜力大的

- 维护信息
- 带重构（适合在线）
- 插入时左右儿子要标记为 null。

```

1 namespace kd {
2     const int K = 2, inf = 1E9, M = N;
3     const double lim = 0.7;
4     struct P {
5         int d[K], l[K], r[K], sz, val;
6         LL sum;
7         P *ls, *rs;
8         P* up() {
9             sz = ls->sz + rs->sz + 1;
10            sum = ls->sum + rs->sum + val;
11            FOR (i, 0, K) {
12                l[i] = min(d[i], min(ls->l[i], rs->l[i]));
13                r[i] = max(d[i], max(ls->r[i], rs->r[i]));
14            }
15            return this;
16        }
17    } pool[M], *null = new P, *pit = pool;
18    static P *tmp[M], **pt;
19    void init() {
20        null->ls = null->rs = null;
21        FOR (i, 0, K) null->l[i] = inf, null->r[i] = -inf;
22        null->sum = null->val = 0;
23        null->sz = 0;
24    }
25
26    P* build(P** l, P** r, int d = 0) { // [l, r)
27        if (d == K) d = 0;
28        if (l >= r) return null;
29        P** m = l + (r - l) / 2; assert(l <= m && m < r);
30        nth_element(l, m, r, [&](const P* a, const P* b){
31            return a->d[d] < b->d[d];
32        });
33        P* o = *m;
34        o->ls = build(l, m, d + 1); o->rs = build(m + 1, r, d + 1);
35        return o->up();
36    }
37    P* Build() {
38        pt = tmp; FOR (it, pool, pit) *pt++ = it;
39        return build(tmp, pt);
40    }
41    inline bool inside(int p[], int q[], int l[], int r[]) {
42        FOR (i, 0, K) if (r[i] < q[i] || p[i] < l[i]) return false;
43        return true;
44    }
45    LL query(P* o, int l[], int r[]) {
46        if (o == null) return 0;
47        FOR (i, 0, K) if (o->r[i] < l[i] || r[i] < o->l[i]) return 0;
48        if (inside(o->l, o->r, l, r)) return o->sum;
49        return query(o->ls, l, r) + query(o->rs, l, r) +
50            (inside(o->d, o->d, l, r) ? o->val : 0);
51    }
52    void dfs(P* o) {
53        if (o == null) return;
54        *pt++ = o; dfs(o->ls); dfs(o->rs);
55    }
56    P* ins(P* o, P* x, int d = 0) {
57        if (d == K) d = 0;
58        if (o == null) return x->up();
59        P*& oo = x->d[d] <= o->d[d] ? o->ls : o->rs;
60        if (oo->sz > o->sz * lim) {
61            pt = tmp; dfs(o); *pt++ = x;
62            return build(tmp, pt, d);
63        }
64        oo = ins(oo, x, d + 1);
65        return o->up();
66    }
67 }

```

- 维护信息
- 带修改 (适合离线)

```

1 namespace kd {
2     const int K = 3, inf = 1E9, M = N << 3;
3     extern struct P* null;
4     struct P {
5         int d[K], l[K], r[K], val;
6         int Max;
7         P *ls, *rs, *fa;
8         P* up() {
9             Max = max(val, max(ls->Max, rs->Max));
10            FOR (i, 0, K) {
11                l[i] = min(d[i], min(ls->l[i], rs->l[i]));
12                r[i] = max(d[i], max(ls->r[i], rs->r[i]));
13            }
14            return ls->fa = rs->fa = this;
15        }
16    } pool[M], *null = new P, *pit = pool;
17    void upd(P* o, int val) {
18        o->val = val;
19        for (; o != null; o = o->fa)
20            o->Max = max(o->Max, val);
21    }
22    static P *tmp[M], **pt;
23    void init() {
24        null->ls = null->rs = null;
25        FOR (i, 0, K) null->l[i] = inf, null->r[i] = -inf;
26        null->Max = null->val = 0;
27    }
28    P* build(P** l, P** r, int d = 0) { // [l, r)
29        if (d == K) d = 0;
30        if (l >= r) return null;
31        P** m = l + (r - l) / 2; assert(l <= m && m < r);
32        nth_element(l, m, r, [&](const P* a, const P* b){
33            return a->d[d] < b->d[d];
34        });
35        P* o = *m;
36        o->ls = build(l, m, d + 1); o->rs = build(m + 1, r, d + 1);
37        return o->up();
38    }
39    P* Build() {
40        pt = tmp; FOR (it, pool, pit) *pt++ = it;
41        P* ret = build(tmp, pt); ret->fa = null;
42        return ret;
43    }
44    inline bool inside(int p[], int q[], int l[], int r[]) {
45        FOR (i, 0, K) if (r[i] < q[i] || p[i] < l[i]) return false;
46        return true;
47    }
48    int query(P* o, int l[], int r[]) {
49        if (o == null) return 0;
50        FOR (i, 0, K) if (o->r[i] < l[i] || r[i] < o->l[i]) return 0;
51        if (inside(o->l, o->r, l, r)) return o->Max;
52        int ret = 0;
53        if (o->val > ret && inside(o->d, o->d, l, r)) ret = max(ret, o->val);
54        if (o->ls->Max > ret) ret = max(ret, query(o->ls, l, r));
55        if (o->rs->Max > ret) ret = max(ret, query(o->rs, l, r));
56        return ret;
57    }
58 }

```

- 最近点对
- 要用全局变量大力剪枝

```

1 namespace kd {
2     const int K = 3;
3     const int M = N;
4     const int inf = 1E9 + 100;
5     struct P {
6         int d[K];
7         int l[K], r[K];
8         P *ls, *rs;
9         P* up() {

```

```

10         FOR (i, 0, K) {
11             l[i] = min(d[i], min(ls->l[i], rs->l[i]));
12             r[i] = max(d[i], max(ls->r[i], rs->r[i]));
13         }
14         return this;
15     }
16 } pool[M], *null = new P, *pit = pool;
17 static P *tmp[M], **pt;
18 void init() {
19     null->ls = null->rs = null;
20     FOR (i, 0, K) null->l[i] = inf, null->r[i] = -inf;
21 }
22 P* build(P** l, P** r, int d = 0) { // [l, r]
23     if (d == K) d = 0;
24     if (l >= r) return null;
25     P** m = l + (r - l) / 2;
26     nth_element(l, m, r, [&](const P* a, const P* b){
27         return a->d[d] < b->d[d];
28     });
29     P* o = *m;
30     o->ls = build(l, m, d + 1); o->rs = build(m + 1, r, d + 1);
31     return o->up();
32 }
33 LL eval(P* o, int d[]) {
34     // ...
35 }
36 LL dist(int d1[], int d2[]) {
37     // ...
38 }
39 LL S;
40 LL query(P* o, int d[]) {
41     if (o == null) return 0;
42     S = max(S, dist(o->d, d));
43     LL mdl = eval(o->ls, d), mdr = eval(o->rs, d);
44     if (mdl < mdr) {
45         if (S > mdl) S = max(S, query(o->ls, d));
46         if (S > mdr) S = max(S, query(o->rs, d));
47     } else {
48         if (S > mdr) S = max(S, query(o->rs, d));
49         if (S > mdl) S = max(S, query(o->ls, d));
50     }
51     return S;
52 }
53 P* Build() {
54     pt = tmp; FOR (it, pool, pit) *pt++ = it;
55     return build(tmp, pt);
56 }
57 }

```

树状数组

- 注意: 0 是无效下标

```

1 namespace bit {
2     LL c[M];
3     inline int lowbit(int x) { return x & -x; }
4     void add(int x, LL v) {
5         for (; x < M; x += lowbit(x))
6             c[x] += v;
7     }
8     LL sum(int x) {
9         LL ret = 0;
10        for (; x > 0; x -= lowbit(x))
11            ret += c[x];
12        return ret;
13    }
14    int kth(LL k) {
15        int ret = 0;
16        LL cnt = 0;

```

```

17     FORD (i, 20, -1) {
18         ret += 1 << i;
19         if (ret >= M || cnt + c[ret] >= k)
20             ret -= 1 << i;
21         else cnt += c[ret];
22     }
23     return ret + 1;
24 }
25 }

```

- 区间修改 & 区间查询 (单点修改, 查询前缀和的前缀和)

```

1 namespace bit {
2     int c[maxn], cc[maxn];
3     inline int lowbit(int x) { return x & -x; }
4     void add(int x, int v) {
5         for (int i = x; i <= n; i += lowbit(i)) {
6             c[i] += v; cc[i] += x * v;
7         }
8     }
9     void add(int l, int r, int v) { add(l, v); add(r + 1, -v); }
10    int sum(int x) {
11        int ret = 0;
12        for (int i = x; i > 0; i -= lowbit(i))
13            ret += (x + 1) * c[i] - cc[i];
14        return ret;
15    }
16    int sum(int l, int r) { return sum(r) - sum(l - 1); }
17 }

```

- 单点修改, 查询前缀和的前缀和的前缀和 (有用才怪)

```

1 namespace bit {
2     LL c[N], cc[N], ccc[N];
3     inline LL lowbit(LL x) { return x & -x; }
4     void add(LL x, LL v) {
5         for (LL i = x; i < N; i += lowbit(i)) {
6             c[i] = (c[i] + v) % MOD;
7             cc[i] = (cc[i] + x * v) % MOD;
8             ccc[i] = (ccc[i] + x * x % MOD * v) % MOD;
9         }
10    }
11    void add(LL l, LL r, LL v) { add(l, v); add(r + 1, -v); }
12    LL sum(LL x) {
13        static LL INV2 = (MOD + 1) / 2;
14        LL ret = 0;
15        for (LL i = x; i > 0; i -= lowbit(i))
16            ret += (x + 1) * (x + 2) % MOD * c[i] % MOD
17                - (2 * x + 3) * cc[i] % MOD
18                + ccc[i];
19        return ret % MOD * INV2 % MOD;
20    }
21    LL sum(LL l, LL r) { return sum(r) - sum(l - 1); }
22 }

```

- 三维

```

1 inline int lowbit(int x) { return x & -x; }
2 void update(int x, int y, int z, int d) {
3     for (int i = x; i <= n; i += lowbit(i))
4         for (int j = y; j <= n; j += lowbit(j))
5             for (int k = z; k <= n; k += lowbit(k))
6                 c[i][j][k] += d;
7 }
8 LL query(int x, int y, int z) {
9     LL ret = 0;
10    for (int i = x; i > 0; i -= lowbit(i))
11        for (int j = y; j > 0; j -= lowbit(j))
12            for (int k = z; k > 0; k -= lowbit(k))
13                ret += c[i][j][k];
14    return ret;
15 }
16 LL solve(int x, int y, int z, int xx, int yy, int zz) {

```

```

17     return    query(xx, yy, zz)
18             - query(xx, yy, z - 1)
19             - query(xx, y - 1, zz)
20             - query(x - 1, yy, zz)
21             + query(xx, y - 1, z - 1)
22             + query(x - 1, yy, z - 1)
23             + query(x - 1, y - 1, zz)
24             - query(x - 1, y - 1, z - 1);

```

主席树

● 正常主席树

```

1  namespace tree {
2      #define mid ((l + r) >> 1)
3      #define lson l, mid
4      #define rson mid + 1, r
5      const int MAGIC = M * 30;
6      struct P {
7          int sum, ls, rs;
8      } tr[MAGIC] = {{0, 0, 0}};
9      int sz = 1;
10     int N(int sum, int ls, int rs) {
11         if (sz == MAGIC) assert(0);
12         tr[sz] = {sum, ls, rs};
13         return sz++;
14     }
15     int ins(int o, int x, int v, int l = 1, int r = ls) {
16         if (x < l || x > r) return o;
17         const P& t = tr[o];
18         if (l == r) return N(t.sum + v, 0, 0);
19         return N(t.sum + v, ins(t.ls, x, v, lson), ins(t.rs, x, v, rson));
20     }
21     int query(int o, int ql, int qr, int l = 1, int r = ls) {
22         if (ql > r || l > qr) return 0;
23         const P& t = tr[o];
24         if (ql <= l && r <= qr) return t.sum;
25         return query(t.ls, ql, qr, lson) + query(t.rs, ql, qr, rson);
26     }
27 }

```

● 第k大

```

1  struct TREE {
2      #define mid ((l + r) >> 1)
3      #define lson l, mid
4      #define rson mid + 1, r
5      struct P {
6          int w, ls, rs;
7      } tr[maxn * 20];
8      int sz = 1;
9      TREE() { tr[0] = {0, 0, 0}; }
10     int N(int w, int ls, int rs) {
11         tr[sz] = {w, ls, rs};
12         return sz++;
13     }
14     int ins(int tt, int l, int r, int x) {
15         if (x < l || r < x) return tt;
16         const P& t = tr[tt];
17         if (l == r) return N(t.w + 1, 0, 0);
18         return N(t.w + 1, ins(t.ls, lson, x), ins(t.rs, rson, x));
19     }
20     int query(int pp, int qq, int l, int r, int k) { // [pp, qq]
21         if (l == r) return l;
22         const P &p = tr[pp], &q = tr[qq];
23         int w = tr[q.ls].w - tr[p.ls].w;
24         if (k <= w) return query(p.ls, q.ls, lson, k);
25         else return query(p.rs, q.rs, rson, k - w);
26     }
27 } tree;

```

- 树状数组套主席树

```

1  typedef vector<int> VI;
2  struct TREE {
3      #define mid ((l + r) >> 1)
4      #define lson l, mid
5      #define rson mid + 1, r
6      struct P {
7          int w, ls, rs;
8      } tr[maxn * 20 * 20];
9      int sz = 1;
10     TREE() { tr[0] = {0, 0, 0}; }
11     int N(int w, int ls, int rs) {
12         tr[sz] = {w, ls, rs};
13         return sz++;
14     }
15     int add(int tt, int l, int r, int x, int d) {
16         if (x < l || r < x) return tt;
17         const P& t = tr[tt];
18         if (l == r) return N(t.w + d, 0, 0);
19         return N(t.w + d, add(t.ls, lson, x, d), add(t.rs, rson, x, d));
20     }
21     int ls_sum(const VI& rt) {
22         int ret = 0;
23         FOR (i, 0, rt.size())
24             ret += tr[rt[i]].ls.w;
25         return ret;
26     }
27     inline void ls(VI& rt) { transform(rt.begin(), rt.end(), rt.begin(), [&](int x)->int{ return tr[x].ls; }); }
28     inline void rs(VI& rt) { transform(rt.begin(), rt.end(), rt.begin(), [&](int x)->int{ return tr[x].rs; }); }
29     int query(VI& p, VI& q, int l, int r, int k) {
30         if (l == r) return l;
31         int w = ls_sum(q) - ls_sum(p);
32         if (k <= w) {
33             ls(p); ls(q);
34             return query(p, q, lson, k);
35         }
36         else {
37             rs(p); rs(q);
38             return query(p, q, rson, k - w);
39         }
40     }
41 } tree;
42 struct BIT {
43     int root[maxn];
44     void init() { memset(root, 0, sizeof root); }
45     inline int lowbit(int x) { return x & -x; }
46     void update(int p, int x, int d) {
47         for (int i = p; i <= m; i += lowbit(i))
48             root[i] = tree.add(root[i], 1, m, x, d);
49     }
50     int query(int l, int r, int k) {
51         VI p, q;
52         for (int i = l - 1; i > 0; i -= lowbit(i)) p.push_back(root[i]);
53         for (int i = r; i > 0; i -= lowbit(i)) q.push_back(root[i]);
54         return tree.query(p, q, 1, m, k);
55     }
56 } bit;
57
58 void init() {
59     m = 10000;
60     tree.sz = 1;
61     bit.init();
62     FOR (i, 1, m + 1)
63         bit.update(i, a[i], 1);
64 }

```


左偏树

```
1 namespace LTree {
2     extern struct P* null, *pit;
3     queue<P*> trash;
4     const int M = 1E5 + 100;
5     struct P {
6         P *ls, *rs;
7         LL v;
8         int d;
9         void operator delete (void* ptr) {
10             trash.push((P*)ptr);
11         }
12         void* operator new(size_t size) {
13             if (trash.empty()) return pit++;
14             void* ret = trash.front(); trash.pop(); return ret;
15         }
16
17         void prt() {
18             if (this == null) return;
19             cout << v << ' ';
20             ls->prt(); rs->prt();
21         }
22     } pool[M], *pit = pool, *null = new P{0, 0, -1, -1};
23     P* N(LL v) {
24         return new P{null, null, v, 0};
25     }
26     P* merge(P* a, P* b) {
27         if (a == null) return b;
28         if (b == null) return a;
29         if (a->v > b->v) swap(a, b);
30         a->rs = merge(a->rs, b);
31         if (a->ls->d < a->rs->d) swap(a->ls, a->rs);
32         a->d = a->rs->d + 1;
33         return a;
34     }
35
36     LL pop(P*& o) {
37         LL ret = o->v;
38         P* t = o;
39         o = merge(o->ls, o->rs);
40         delete t;
41         return ret;
42     }
43 }
```

可持久化

```
1 namespace LTree {
2     extern struct P* null, *pit;
3     queue<P*> trash;
4     const int M = 1E6 + 100;
5     struct P {
6         P *ls, *rs;
7         LL v;
8         int d;
9         void operator delete (void* ptr) {
10             trash.push((P*)ptr);
11         }
12         void* operator new(size_t size) {
13             if (trash.empty()) return pit++;
14             void* ret = trash.front(); trash.pop(); return ret;
15         }
16     } pool[M], *pit = pool, *null = new P{0, 0, -1, -1};
17     P* N(LL v, P* ls = null, P* rs = null) {
18         if (ls->d < rs->d) swap(ls, rs);
19         return new P{ls, rs, v, rs->d + 1};
20     }
21     P* merge(P* a, P* b) {
22         if (a == null) return b;
23         if (b == null) return a;
```

```

24         if (a->v < b->v)
25             return N(a->v, a->ls, merge(a->rs, b));
26         else
27             return N(b->v, b->ls, merge(b->rs, a));
28     }
29
30 LL pop(P*& o) {
31     LL ret = o->v;
32     o = merge(o->ls, o->rs);
33     return ret;
34 }
35 }

```

Treap

- 非旋 Treap
- v 小根堆
- 模板题 bzoj 3224
- lower 第一个大于等于的是第几个 (0-based)
- upper 第一个大于的是第几个 (0-based)
- split 左侧分割出 rk 个元素
- 树套树略

```

1 namespace treap {
2     const int M = maxn * 17;
3     extern struct P* const null;
4     struct P {
5         P *ls, *rs;
6         int v, sz;
7         unsigned rd;
8         P(int v): ls(null), rs(null), v(v), sz(1), rd(rnd()) {}
9         P(): sz(0) {}
10
11         P* up() { sz = ls->sz + rs->sz + 1; return this; }
12         int lower(int v) {
13             if (this == null) return 0;
14             return this->v >= v ? ls->lower(v) : rs->lower(v) + ls->sz + 1;
15         }
16         int upper(int v) {
17             if (this == null) return 0;
18             return this->v > v ? ls->upper(v) : rs->upper(v) + ls->sz + 1;
19         }
20     } *const null = new P, pool[M], *pit = pool;
21
22     P* merge(P* l, P* r) {
23         if (l == null) return r; if (r == null) return l;
24         if (l->rd < r->rd) { l->rs = merge(l->rs, r); return l->up(); }
25         else { r->ls = merge(l, r->ls); return r->up(); }
26     }
27
28     void split(P* o, int rk, P*& l, P*& r) {
29         if (o == null) { l = r = null; return; }
30         if (o->ls->sz >= rk) { split(o->ls, rk, l, o->ls); r = o->up(); }
31         else { split(o->rs, rk - o->ls->sz - 1, o->rs, r); l = o->up(); }
32     }
33 }

```

- 持久化 Treap

```

1 namespace treap {
2     const int M = maxn * 17 * 12;
3     extern struct P* const null, *pit;
4     struct P {
5         P *ls, *rs;
6         int v, sz;
7         LL sum;
8         P(P* ls, P* rs, int v): ls(ls), rs(rs), v(v), sz(ls->sz + rs->sz + 1),
9             sum(ls->sum + rs->sum + v) {}

```

```

10     P() {}
11
12     void* operator new(size_t _) { return pit++; }
13     template<typename T>
14     int rk(int v, T&& cmp) {
15         if (this == null) return 0;
16         return cmp(this->v, v) ? ls->rk(v, cmp) : rs->rk(v, cmp) + ls->sz + 1;
17     }
18     int lower(int v) { return rk(v, greater_equal<int>()); }
19     int upper(int v) { return rk(v, greater<int>()); }
20 } pool[M], *pit = pool, *const null = new P;
21 P* merge(P* l, P* r) {
22     if (l == null) return r; if (r == null) return l;
23     if (rnd() % (l->sz + r->sz) < l->sz) return new P{l->ls, merge(l->rs, r), l->v};
24     else return new P{merge(l, r->ls), r->rs, r->v};
25 }
26 void split(P* o, int rk, P*& l, P*& r) {
27     if (o == null) { l = r = null; return; }
28     if (o->ls->sz >= rk) { split(o->ls, rk, l, r); r = new P{r, o->rs, o->v}; }
29     else { split(o->rs, rk - o->ls->sz - 1, l, r); l = new P{o->ls, l, o->v}; }
30 }
31 }

```

- 帶 pushdown 的持久化 Treap
- 注意任何修改操作前一定要 FIX

```

1  int now;
2  namespace Treap {
3      const int M = 100000000;
4      extern struct P* const null, *pit;
5      struct P {
6          P *ls, *rs;
7          int sz, time;
8          LL cnt, sc, pos, add;
9          bool rev;
10
11          P* up() { sz = ls->sz + rs->sz + 1; sc = ls->sc + rs->sc + cnt; return this; } // MOD
12          P* check() {
13              if (time == now) return this;
14              P* t = new(pit++) P; *t = *this; t->time = now; return t;
15          };
16          P* _do_rev() { rev ^= 1; add *= -1; pos *= -1; swap(ls, rs); return this; } // MOD
17          P* _do_add(LL v) { add += v; pos += v; return this; } // MOD
18          P* do_rev() { if (this == null) return this; return check()->_do_rev(); } // FIX & MOD
19          P* do_add(LL v) { if (this == null) return this; return check()->_do_add(v); } // FIX & MOD
20          P* _down() { // MOD
21              if (rev) { ls = ls->do_rev(); rs = rs->do_rev(); rev = 0; }
22              if (add) { ls = ls->do_add(add); rs = rs->do_add(add); add = 0; }
23              return this;
24          }
25          P* down() { return check()->_down(); } // FIX & MOD
26          void _split(LL p, P*& l, P*& r) { // MOD
27              if (pos >= p) { ls->split(p, l, r); ls = r; r = up(); }
28              else { rs->split(p, l, r); rs = l; l = up(); }
29          }
30          void split(LL p, P*& l, P*& r) { // FIX & MOD
31              if (this == null) l = r = null;
32              else down()->_split(p, l, r);
33          }
34      } pool[M], *pit = pool, *const null = new P;
35      P* merge(P* a, P* b) {
36          if (a == null) return b; if (b == null) return a;
37          if (rand() % (a->sz + b->sz) < a->sz) { a = a->down(); a->rs = merge(a->rs, b); return a->up(); }
38          else { b = b->down(); b->ls = merge(a, b->ls); return b->up(); }
39      }
40  }

```

Treap-序列

- 区间 ADD, SUM

```
1 namespace treap {
2     const int M = 8E5 + 100;
3     extern struct P* const null;
4     struct P {
5         P *ls, *rs;
6         int sz, val, add, sum;
7         P(int v, P* ls = null, P* rs = null): ls(ls), rs(rs), sz(1), val(v), add(0), sum(v) {}
8         P(): sz(0), val(0), add(0), sum(0) {}
9
10        P* up() {
11            assert(this != null);
12            sz = ls->sz + rs->sz + 1;
13            sum = ls->sum + rs->sum + val + add * sz;
14            return this;
15        }
16        void upd(int v) {
17            if (this == null) return;
18            add += v;
19            sum += sz * v;
20        }
21        P* down() {
22            if (add) {
23                ls->upd(add); rs->upd(add);
24                val += add;
25                add = 0;
26            }
27            return this;
28        }
29
30        P* select(int rk) {
31            if (rk == ls->sz + 1) return this;
32            return ls->sz >= rk ? ls->select(rk) : rs->select(rk - ls->sz - 1);
33        }
34    } pool[M], *pit = pool, *const null = new P, *rt = null;
35
36    P* merge(P* a, P* b) {
37        if (a == null) return b->up();
38        if (b == null) return a->up();
39        if (rand() % (a->sz + b->sz) < a->sz) {
40            a->down()->rs = merge(a->rs, b);
41            return a->up();
42        } else {
43            b->down()->ls = merge(a, b->ls);
44            return b->up();
45        }
46    }
47
48    void split(P* o, int rk, P*& l, P*& r) {
49        if (o == null) { l = r = null; return; }
50        o->down();
51        if (o->ls->sz >= rk) {
52            split(o->ls, rk, l, o->ls);
53            r = o->up();
54        } else {
55            split(o->rs, rk - o->ls->sz - 1, o->rs, r);
56            l = o->up();
57        }
58    }
59
60    inline void insert(int k, int v) {
61        P *l, *r;
62        split(rt, k - 1, l, r);
63        rt = merge(merge(l, new (pit++) P(v)), r);
64    }
65
66    inline void erase(int k) {
67        P *l, *r, *_ , *t;
```

```

68     split(rt, k - 1, l, t);
69     split(t, l, _, r);
70     rt = merge(l, r);
71 }
72
73 P* build(int l, int r, int* a) {
74     if (l > r) return null;
75     if (l == r) return new(pit++) P(a[l]);
76     int m = (l + r) / 2;
77     return (new(pit++) P(a[m], build(l, m - 1, a), build(m + 1, r, a)))->up();
78 }
79 };

```

● 区间 REVERSE, ADD, MIN

```

1  namespace treap {
2      extern struct P*const null;
3      struct P {
4          P *ls, *rs;
5          int sz, v, add, m;
6          bool flip;
7          P(int v, P* ls = null, P* rs = null): ls(ls), rs(rs), sz(1), v(v), add(0), m(v), flip(0) {}
8          P(): sz(0), v(INF), m(INF) {}
9
10         void upd(int v) {
11             if (this == null) return;
12             add += v; m += v;
13         }
14         void rev() {
15             if (this == null) return;
16             swap(ls, rs);
17             flip ^= 1;
18         }
19         P* up() {
20             assert(this != null);
21             sz = ls->sz + rs->sz + 1;
22             m = min(min(ls->m, rs->m), v) + add;
23             return this;
24         }
25         P* down() {
26             if (add) {
27                 ls->upd(add); rs->upd(add);
28                 v += add;
29                 add = 0;
30             }
31             if (flip) {
32                 ls->rev(); rs->rev();
33                 flip = 0;
34             }
35             return this;
36         }
37
38         P* select(int k) {
39             if (ls->sz + 1 == k) return this;
40             if (ls->sz >= k) return ls->select(k);
41             return rs->select(k - ls->sz - 1);
42         }
43
44     } pool[M], *const null = new P, *pit = pool, *rt = null;
45
46     P* merge(P* a, P* b) {
47         if (a == null) return b;
48         if (b == null) return a;
49         if (rnd() % (a->sz + b->sz) < a->sz) {
50             a->down()->rs = merge(a->rs, b);
51             return a->up();
52         } else {
53             b->down()->ls = merge(a, b->ls);
54             return b->up();
55         }
56     }
57 }

```

```

58 void split(P* o, int k, P*& l, P*& r) {
59     if (o == null) { l = r = null; return; }
60     o->down();
61     if (o->ls->sz >= k) {
62         split(o->ls, k, l, o->ls);
63         r = o->up();
64     } else {
65         split(o->rs, k - o->ls->sz - 1, o->rs, r);
66         l = o->up();
67     }
68 }
69
70 P* build(int l, int r, int* v) {
71     if (l > r) return null;
72     int m = (l + r) >> 1;
73     return (new (pit++) P(v[m], build(l, m - 1, v), build(m + 1, r, v)))->up();
74 }
75
76 void go(int x, int y, void f(P*&)) {
77     P *l, *m, *r;
78     split(rt, y, l, r);
79     split(l, x - 1, l, m);
80     f(m);
81     rt = merge(merge(l, m), r);
82 }
83
84 using namespace treap;
85 int a[maxn], n, x, y, Q, v, k, d;
86 char s[100];
87
88 int main() {
89     cin >> n;
90     FOR (i, 1, n + 1) scanf("%d", &a[i]);
91     rt = build(1, n, a);
92     cin >> Q;
93     while (Q--) {
94         scanf("%s", s);
95         if (s[0] == 'A') {
96             scanf("%d%d%d", &x, &y, &v);
97             go(x, y, [](P*& o){ o->upd(v); });
98         } else if (s[0] == 'R' && s[3] == 'E') {
99             scanf("%d%d", &x, &y);
100            go(x, y, [](P*& o){ o->rev(); });
101        } else if (s[0] == 'R' && s[3] == 'O') {
102            scanf("%d%d%d", &x, &y, &d);
103            d %= y - x + 1;
104            go(x, y, [](P*& o){
105                P *l, *r;
106                split(o, o->sz - d, l, r);
107                o = merge(r, l);
108            });
109        } else if (s[0] == 'I') {
110            scanf("%d%d", &k, &v);
111            go(k + 1, k, [](P*& o){ o = new (pit++) P(v); });
112        } else if (s[0] == 'D') {
113            scanf("%d", &k);
114            go(k, k, [](P*& o){ o = null; });
115        } else if (s[0] == 'M') {
116            scanf("%d%d", &x, &y);
117            go(x, y, [](P*& o) {
118                printf("%d\n", o->m);
119            });
120        }
121    }
122 }

```

● 持久化

```

1 namespace treap {
2     struct P;
3     extern P*const null;
4     P* N(P* ls, P* rs, LL v, bool fill);

```

```

5 struct P {
6     P *const ls, *const rs;
7     const int sz, v;
8     const LL sum;
9     bool fill;
10    int cnt;
11
12    void split(int k, P*& l, P*& r) {
13        if (this == null) { l = r = null; return; }
14        if (ls->sz >= k) {
15            ls->split(k, l, r);
16            r = N(r, rs, v, fill);
17        } else {
18            rs->split(k - ls->sz - fill, l, r);
19            l = N(ls, l, v, fill);
20        }
21    }
22
23
24    } *const null = new P{0, 0, 0, 0, 0, 1};
25
26    P* N(P* ls, P* rs, LL v, bool fill) {
27        ls->cnt++; rs->cnt++;
28        return new P{ls, rs, ls->sz + rs->sz + fill, v, ls->sum + rs->sum + v, fill, 1};
29    }
30
31    P* merge(P* a, P* b) {
32        if (a == null) return b;
33        if (b == null) return a;
34        if (rand() % (a->sz + b->sz) < a->sz)
35            return N(a->ls, merge(a->rs, b), a->v, a->fill);
36        else
37            return N(merge(a, b->ls), b->rs, b->v, b->fill);
38    }
39
40    void go(P* o, int x, int y, P*& l, P*& m, P*& r) {
41        o->split(y, l, r);
42        l->split(x - 1, l, m);
43    }
44 }

```

可回滚并查集

- 注意这个不是可持久化并查集
- 查找时不进行路径压缩
- 复杂度靠按秩合并解决

```

1 namespace uf {
2     int fa[maxn], sz[maxn];
3     int undo[maxn], top;
4     void init() { memset(fa, -1, sizeof fa); memset(sz, 0, sizeof sz); top = 0; }
5     int findset(int x) { while (fa[x] != -1) x = fa[x]; return x; }
6     bool join(int x, int y) {
7         x = findset(x); y = findset(y);
8         if (x == y) return false;
9         if (sz[x] > sz[y]) swap(x, y);
10        undo[top++] = x;
11        fa[x] = y;
12        sz[y] += sz[x] + 1;
13        return true;
14    }
15    inline int checkpoint() { return top; }
16    void rewind(int t) {
17        while (top > t) {
18            int x = undo[--top];
19            sz[fa[x]] -= sz[x] + 1;
20            fa[x] = -1;
21        }
22    }
23 }

```

```

22     }
23 }

```

舞蹈链

- 注意 link 的 y 的范围是 [1, n]
- 注意在某些情况下替换掉 memset
- 精确覆盖

```

1  struct P {
2      P *L, *R, *U, *D;
3      int x, y;
4  };
5
6  const int INF = 1E9;
7
8  struct DLX {
9      #define TR(i, D, s) for (P* i = s->D; i != s; i = i->D)
10     static const int M = 2E5;
11     P pool[M], *h[M], *r[M], *pit;
12     int sz[M];
13     bool solved;
14     stack<int> ans;
15     void init(int n) {
16         pit = pool;
17         ++n;
18         solved = false;
19         while (!ans.empty()) ans.pop();
20         memset(r, 0, sizeof r);
21         memset(sz, 0, sizeof sz);
22         FOR (i, 0, n)
23             h[i] = new (pit++) P;
24         FOR (i, 0, n) {
25             h[i]->L = h[(i + n - 1) % n];
26             h[i]->R = h[(i + 1) % n];
27             h[i]->U = h[i]->D = h[i];
28             h[i]->y = i;
29         }
30     }
31
32     void link(int x, int y) {
33         sz[y]++;
34         auto p = new (pit++) P;
35         p->x = x; p->y = y;
36         p->U = h[y]->U; p->D = h[y];
37         p->D->U = p->U->D = p;
38         if (!r[x]) r[x] = p->L = p->R = p;
39         else {
40             p->L = r[x]; p->R = r[x]->R;
41             p->L->R = p->R->L = p;
42         }
43     }
44
45     void remove(P* p) {
46         p->L->R = p->R; p->R->L = p->L;
47         TR (i, D, p)
48             TR (j, R, i) {
49                 j->D->U = j->U; j->U->D = j->D;
50                 sz[j->y]--;
51             }
52     }
53
54     void recall(P* p) {
55         p->L->R = p->R->L = p;
56         TR (i, U, p)
57             TR (j, L, i) {
58                 j->D->U = j->U->D = j;
59                 sz[j->y]++;

```



```

60     }
61 }
62
63 bool dfs(int d) {
64     if (solved) return true;
65     if (h[0]->R == h[0]) return solved = true;
66     int m = INF;
67     P* c;
68     TR (i, R, h[0])
69         if (sz[i->y] < m) { m = sz[i->y]; c = i; }
70     remove(c);
71     TR (i, D, c) {
72         ans.push(i->x);
73         TR (j, R, i) remove(h[j->y]);
74         if (dfs(d + 1)) return true;
75         TR (j, L, i) recall(h[j->y]);
76         ans.pop();
77     }
78     recall(c);
79     return false;
80 }
81 } dlx;

```

- 可重复覆盖

```

1  struct P {
2      P *L, *R, *U, *D;
3      int x, y;
4  };
5
6  const int INF = 1E9;
7
8  struct DLX {
9      #define TR(i, D, s) for (P* i = s->D; i != s; i = i->D)
10     static const int M = 2E5;
11     P pool[M], *h[M], *r[M], *pit;
12     int sz[M], vis[M], ans, clk;
13     void init(int n) {
14         clk = 0;
15         ans = INF;
16         pit = pool;
17         ++n;
18         memset(r, 0, sizeof r);
19         memset(sz, 0, sizeof sz);
20         memset(vis, -1, sizeof vis);
21         FOR (i, 0, n)
22             h[i] = new (pit++) P;
23         FOR (i, 0, n) {
24             h[i]->L = h[(i + n - 1) % n];
25             h[i]->R = h[(i + 1) % n];
26             h[i]->U = h[i]->D = h[i];
27             h[i]->y = i;
28         }
29     }
30
31     void link(int x, int y) {
32         sz[y]++;
33         auto p = new (pit++) P;
34         p->x = x; p->y = y;
35         p->U = h[y]->U; p->D = h[y];
36         p->D->U = p->U->D = p;
37         if (!r[x]) r[x] = p->L = p->R = p;
38         else {
39             p->L = r[x]; p->R = r[x]->R;
40             p->L->R = p->R->L = p;
41         }
42     }
43
44     void remove(P* p) {
45         TR (i, D, p) {
46             i->L->R = i->R;
47             i->R->L = i->L;

```

```

48     }
49 }
50
51 void recall(P* p) {
52     TR (i, U, p)
53     i->L->R = i->R->L = i;
54 }
55
56 int eval() {
57     ++clk;
58     int ret = 0;
59     TR (i, R, h[0])
60         if (vis[i->y] != clk) {
61             ++ret;
62             vis[i->y] = clk;
63             TR (j, D, i)
64                 TR (k, R, j)
65                     vis[k->y] = clk;
66         }
67     return ret;
68 }
69
70 void dfs(int d) {
71     if (h[0]->R == h[0]) { ans = min(ans, d); return; }
72     if (eval() + d >= ans) return;
73     P* c;
74     int m = INF;
75     TR (i, R, h[0])
76         if (sz[i->y] < m) { m = sz[i->y]; c = i; }
77     TR (i, D, c) {
78         remove(i);
79         TR (j, R, i) remove(j);
80         dfs(d + 1);
81         TR (j, L, i) recall(j);
82         recall(i);
83     }
84 }
85 } d1x;

```

CDQ 分治

```

1  const int maxn = 2E5 + 100;
2  struct P {
3      int x, y;
4      int* f;
5      bool d1, d2;
6  } a[maxn], b[maxn], c[maxn];
7  int f[maxn];
8
9  void go2(int l, int r) {
10     if (l + 1 == r) return;
11     int m = (l + r) >> 1;
12     go2(l, m); go2(m, r);
13     FOR (i, l, m) b[i].d2 = 0;
14     FOR (i, m, r) b[i].d2 = 1;
15     merge(b + l, b + m, b + m, b + r, c + l, [](const P& a, const P& b)->bool {
16         if (a.y != b.y) return a.y < b.y;
17         return a.d2 > b.d2;
18     });
19     int mx = -1;
20     FOR (i, l, r) {
21         if (c[i].d1 && c[i].d2) *c[i].f = max(*c[i].f, mx + 1);
22         if (!c[i].d1 && !c[i].d2) mx = max(mx, *c[i].f);
23     }
24     FOR (i, l, r) b[i] = c[i];
25 }
26
27 void go1(int l, int r) { // [l, r)
28     if (l + 1 == r) return;

```

```

29     int m = (l + r) >> 1;
30     go1(l, m);
31     FOR (i, l, m) a[i].d1 = 0;
32     FOR (i, m, r) a[i].d1 = 1;
33     copy(a + l, a + r, b + l);
34     sort(b + l, b + r, [] (const P& a, const P& b) -> bool {
35         if (a.x != b.x) return a.x < b.x;
36         return a.d1 > b.d1;
37     });
38     go2(l, r);
39     go1(m, r);
40 }

```

● k 维 LIS

```

1  struct P {
2      int v[K];
3      LL f;
4      bool d[K];
5  } o[N << 10];
6  P* a[K][N << 10];
7  int k;
8  void go(int now, int l, int r) {
9      if (now == 0) {
10         if (l + 1 == r) return;
11         int m = (l + r) / 2;
12         go(now, l, m);
13         FOR (i, l, m) a[now][i] -> d[now] = 0;
14         FOR (i, m, r) a[now][i] -> d[now] = 1;
15         copy(a[now] + l, a[now] + r, a[now + 1] + l);
16         sort(a[now + 1] + l, a[now + 1] + r, [now] (const P* a, const P* b) {
17             if (a -> v[now] != b -> v[now]) return a -> v[now] < b -> v[now];
18             return a -> d[now] > b -> d[now];
19         });
20         go(now + 1, l, r);
21         go(now, m, r);
22     } else {
23         if (l + 1 == r) return;
24         int m = (l + r) / 2;
25         go(now, l, m); go(now, m, r);
26         FOR (i, l, m) a[now][i] -> d[now] = 0;
27         FOR (i, m, r) a[now][i] -> d[now] = 1;
28         merge(a[now] + l, a[now] + m, a[now] + m, a[now] + r, a[now + 1] + l, [now] (const P* a, const P* b) {
29             if (a -> v[now] != b -> v[now]) return a -> v[now] < b -> v[now];
30             return a -> d[now] > b -> d[now];
31         });
32         copy(a[now + 1] + l, a[now + 1] + r, a[now] + l);
33         if (now < k - 2) {
34             go(now + 1, l, r);
35         } else {
36             LL sum = 0;
37             FOR (i, l, r) {
38                 dbg(a[now][i] -> v[0], a[now][i] -> v[1], a[now][i] -> f,
39                     a[now][i] -> d[0], a[now][i] -> d[1]);
40
41                 int cnt = 0;
42                 FOR (j, 0, now + 1) cnt += a[now][i] -> d[j];
43                 if (cnt == 0) {
44                     sum += a[now][i] -> f;
45                 } else if (cnt == now + 1) {
46                     a[now][i] -> f = (a[now][i] -> f + sum) % MOD;
47                 }
48             }
49         }
50     }
}

```

笛卡尔树

```

1  void build(const vector<int>& a) {
2      static P *stack[M], *x, *last;

```

```

3     int p = 0;
4     FOR (i, 0, a.size()) {
5         x = new P(i + 1, a[i]);
6         last = null;
7         while (p && stack[p - 1]->v > x->v) {
8             stack[p - 1]->maintain();
9             last = stack[--p];
10        }
11        if (p) stack[p - 1]->rs = x;
12        x->ls = last;
13        stack[p++] = x;
14    }
15    while (p)
16        stack[--p]->maintain();
17    rt = stack[0];
18 }

1 void build() {
2     static int s[N], last;
3     int p = 0;
4     FOR (x, 1, n + 1) {
5         last = 0;
6         while (p && val[s[p - 1]] > val[x]) last = s[--p];
7         if (p) G[s[p - 1]][1] = x;
8         if (last) G[x][0] = last;
9         s[p++] = x;
10    }
11    rt = s[0];
12 }

```

Trie

- Trie 二进制版
- M 为二进制的位数
- 使用前必须初始化

```

1 struct Trie2 {
2     int ch[N * M][2], sz;
3     void init() {
4         memset(ch, 0, sizeof ch);
5         sz = 1;
6     }
7     void insert(LL x) {
8         int u = 0;
9         FOR (i, M, -1) {
10            bool b = x & (1LL << i);
11            if (!ch[u][b])
12                ch[u][b] = sz++;
13            u = ch[u][b];
14        }
15    }
16 } trie;

```

- 持久化二进制 Trie
- sz=1

```

1 struct P { int w, ls, rs; };
2 P tr[M] = {{0, 0, 0}};
3 int sz;
4
5 int _new(int w, int ls, int rs) { tr[sz] = {w, ls, rs}; return sz++; }
6 int ins(int oo, int v, int d = 30) {
7     P& o = tr[oo];
8     if (d == -1) return _new(o.w + 1, 0, 0);
9     bool u = v & (1 << d);
10    return _new(o.w + 1, u == 0 ? ins(o.ls, v, d - 1) : o.ls, u == 1 ? ins(o.rs, v, d - 1) : o.rs);
11 }
12 int query(int pp, int qq, int v, int d = 30) {
13     if (d == -1) return 0;

```

```

14     bool u = v & (1 << d);
15     P &p = tr[pp], &q = tr[qq];
16     int lw = tr[q.ls].w - tr[p.ls].w;
17     int rw = tr[q.rs].w - tr[p.rs].w;
18
19     int ret = 0;
20     if (u == 0) {
21         if (rw) { ret += 1 << d; ret += query(p.rs, q.rs, v, d - 1); }
22         else ret += query(p.ls, q.ls, v, d - 1);
23     } else {
24         if (lw) { ret += 1 << d; ret += query(p.ls, q.ls, v, d - 1); }
25         else ret += query(p.rs, q.rs, v, d - 1);
26     }
27     return ret;
28 }

```

pb_ds

- 优先队列
- binary_heap_tag
- pairing_heap_tag 支持修改
- thin_heap_tag 如果修改只有 increase 则较快，不支持 join

```

1 #include<ext/pb_ds/priority_queue.hpp>
2 template<typename _Tv,
3         typename Cmp_Fn = std::less<_Tv>,
4         typename Tag = pairing_heap_tag,
5         typename _Alloc = std::allocator<char> >
6 class priority_queue;
7
1 #include<ext/pb_ds/priority_queue.hpp>
2 using namespace __gnu_pbds;
3
4 typedef __gnu_pbds::priority_queue<LL, less<LL>, pairing_heap_tag> PQ;
5 __gnu_pbds::priority_queue<int, cmp, pairing_heap_tag>::point_iterator it;
6 PQ pq, pq2;
7
8 int main() {
9     auto it = pq.push(2);
10    pq.push(3);
11    assert(pq.top() == 3);
12    pq.modify(it, 4);
13    assert(pq.top() == 4);
14    pq2.push(5);
15    pq.join(pq2);
16    assert(pq.top() == 5);
17 }

```

- 树
- ov_tree_tag
- rb_tree_tag
- splay_tree_tag
- mapped: null_type 或 null_mapped_type (旧版本) 为空
- Node_Update 为 tree_order_statistics_node_update 时才可以 find_by_order & order_of_key
- find_by_order 找 order + 1 小的元素 (其实都是从 0 开始计数)，或者有 order 个元素比它小的 key
- order_of_key 有多少个比 r_key 小的元素
- join & split

```

1 template<typename Key, typename Mapped, typename Cmp_Fn = std::less<Key>,
2         typename Tag = rb_tree_tag,

```

```

3     template<typename Node_Citr, typename Node_Itr,
4             typename Cmp_Fn_, typename _Alloc_>
5     class Node_Update = null_node_update,
6     typename _Alloc = std::allocator<char> >
7     class tree
8
9     #include <ext/pb_ds/assoc_container.hpp>
10    using namespace __gnu_pbds;
11    using Tree = tree<int, null_type, less<int>, rb_tree_tag, tree_order_statistics_node_update>;
12    Tree t;

```

- hash table

```

1    #include<ext/pb_ds/assoc_container.hpp>
2    #include<ext/pb_ds/hash_policy.hpp>
3    using namespace __gnu_pbds;
4
5    gp_hash_table<int, int> mp;
6    cc_hash_table<int, int> mp;

```

Link-Cut Tree

- 图中相邻的结点在伸展树中不一定是父子关系
- 遇事不决 make_root
- 跑左右儿子的时候不要忘记 down

```

1    namespace lct {
2        extern struct P *const null;
3        const int M = N;
4        struct P {
5            P *fa, *ls, *rs;
6            int v, maxv;
7            bool rev;
8
9            bool has_fa() { return fa->ls == this || fa->rs == this; }
10           bool d() { return fa->ls == this; }
11           P& c(bool x) { return x ? ls : rs; }
12           void do_rev() {
13               if (this == null) return;
14               rev ^= 1;
15               swap(ls, rs);
16           }
17           P* up() {
18               maxv = max(v, max(ls->maxv, rs->maxv));
19               return this;
20           }
21           void down() {
22               if (rev) {
23                   rev = 0;
24                   ls->do_rev(); rs->do_rev();
25               }
26           }
27           void all_down() { if (has_fa()) fa->all_down(); down(); }
28       } *const null = new P{0, 0, 0, 0, 0, 0}, pool[M], *pit = pool;
29
30       void rot(P* o) {
31           bool dd = o->d();
32           P *f = o->fa, *t = o->c(!dd);
33           if (f->has_fa()) f->fa->c(f->d()) = o; o->fa = f->fa;
34           if (t != null) t->fa = f; f->c(dd) = t;
35           o->c(!dd) = f->up(); f->fa = o;
36       }
37       void splay(P* o) {
38           o->all_down();
39           while (o->has_fa()) {
40               if (o->fa->has_fa())
41                   rot(o->d() ^ o->fa->d() ? o : o->fa);
42               rot(o);
43           }
44       }

```

```

44     o->up();
45 }
46 void access(P* u, P* v = null) {
47     if (u == null) return;
48     splay(u); u->rs = v;
49     access(u->up()->fa, u);
50 }
51 void make_root(P* o) {
52     access(o); splay(o); o->do_rev();
53 }
54 void split(P* o, P* u) {
55     make_root(o); access(u); splay(u);
56 }
57 void link(P* u, P* v) {
58     make_root(u); u->fa = v;
59 }
60 void cut(P* u, P* v) {
61     split(u, v);
62     u->fa = v->ls = null; v->up();
63 }
64 bool adj(P* u, P* v) {
65     split(u, v);
66     return v->ls == u && u->ls == null && u->rs == null;
67 }
68 bool linked(P* u, P* v) {
69     split(u, v);
70     return u == v || u->fa != null;
71 }
72 P* findrt(P* o) {
73     access(o); splay(o);
74     while (o->ls != null) o = o->ls;
75     return o;
76 }
77 P* findfa(P* rt, P* u) {
78     split(rt, u);
79     u = u->ls;
80     while (u->rs != null) {
81         u = u->rs;
82         u->down();
83     }
84     return u;
85 }
86 }

```

● 维护子树大小

```

1 P* up() {
2     sz = ls->sz + rs->sz + _sz + 1;
3     return this;
4 }
5 void access(P* u, P* v = null) {
6     if (u == null) return;
7     splay(u);
8     u->_sz += u->rs->sz - v->sz;
9     u->rs = v;
10    access(u->up()->fa, u);
11 }
12 void link(P* u, P* v) {
13     split(u, v);
14     u->fa = v; v->_sz += u->sz;
15     v->up();
16 }

```

莫队

● [l, r)

```

1 while (l > q.l) mv(--l, 1);
2 while (r < q.r) mv(r++, 1);

```

```

3 while (l < q.l) mv(l++, -1);
4 while (r > q.r) mv(--r, -1);

    • 树上莫队
    • 注意初始状态 u = v = 1, flip(1)

1 struct Q {
2     int u, v, idx;
3     bool operator < (const Q& b) const {
4         const Q& a = *this;
5         return blk[a.u] < blk[b.u] || (blk[a.u] == blk[b.u] && in[a.v] < in[b.v]);
6     }
7 };
8
9 void dfs(int u = 1, int d = 0) {
10     static int S[maxn], sz = 0, blk_cnt = 0, clk = 0;
11     in[u] = clk++;
12     dep[u] = d;
13     int btm = sz;
14     for (int v: G[u]) {
15         if (v == fa[u]) continue;
16         fa[v] = u;
17         dfs(v, d + 1);
18         if (sz - btm >= B) {
19             while (sz > btm) blk[S[--sz]] = blk_cnt;
20             ++blk_cnt;
21         }
22     }
23     S[sz++] = u;
24     if (u == 1) while (sz) blk[S[--sz]] = blk_cnt - 1;
25 }
26
27 void flip(int k) {
28     dbg(k);
29     if (vis[k]) {
30         // ...
31     } else {
32         // ...
33     }
34     vis[k] ^= 1;
35 }
36
37 void go(int& k) {
38     if (bug == -1) {
39         if (vis[k] && !vis[fa[k]]) bug = k;
40         if (!vis[k] && vis[fa[k]]) bug = fa[k];
41     }
42     flip(k);
43     k = fa[k];
44 }
45
46 void mv(int a, int b) {
47     bug = -1;
48     if (vis[b]) bug = b;
49     if (dep[a] < dep[b]) swap(a, b);
50     while (dep[a] > dep[b]) go(a);
51     while (a != b) {
52         go(a); go(b);
53     }
54     go(a); go(bug);
55 }
56
57 for (Q& q: query) {
58     mv(u, q.u); u = q.u;
59     mv(v, q.v); v = q.v;
60     ans[q.idx] = Ans;
61 }

```


数学

矩阵运算

```
1 struct Mat {
2     static const LL M = 2;
3     LL v[M][M];
4     Mat() { memset(v, 0, sizeof v); }
5     void eye() { FOR (i, 0, M) v[i][i] = 1; }
6     LL* operator [] (LL x) { return v[x]; }
7     const LL* operator [] (LL x) const { return v[x]; }
8     Mat operator * (const Mat& B) {
9         const Mat& A = *this;
10        Mat ret;
11        FOR (i, 0, M)
12            FOR (j, 0, M)
13                FOR (k, 0, M)
14                    ret[i][j] = (ret[i][j] + A[i][k] * B[k][j]) % MOD;
15        return ret;
16    }
17    Mat pow(LL n) const {
18        Mat A = *this, ret; ret.eye();
19        for (; n; n >>= 1, A = A * A)
20            if (n & 1) ret = ret * A;
21        return ret;
22    }
23    Mat operator + (const Mat& B) {
24        const Mat& A = *this;
25        Mat ret;
26        FOR (i, 0, M)
27            FOR (j, 0, M)
28                ret[i][j] = (A[i][j] + B[i][j]) % MOD;
29        return ret;
30    }
31    void prt() const {
32        FOR (i, 0, M)
33            FOR (j, 0, M)
34                printf("%lld%c", (*this)[i][j], j == M - 1 ? '\n' : ' ');
35    }
36 };
```

筛

● 线性筛

```
1 const LL p_max = 1E6 + 100;
2 LL pr[p_max], p_sz;
3 void get_prime() {
4     static bool vis[p_max];
5     FOR (i, 2, p_max) {
6         if (!vis[i]) pr[p_sz++] = i;
7         FOR (j, 0, p_sz) {
8             if (pr[j] * i >= p_max) break;
9             vis[pr[j] * i] = 1;
10            if (i % pr[j] == 0) break;
11        }
12    }
13 }
```

● 线性筛 + 欧拉函数

```
1 const LL p_max = 1E5 + 100;
2 LL phi[p_max] = {-1, 1};
3 void get_phi() {
4     static bool vis[p_max];
5     static LL prime[p_max], p_sz, d;
6     FOR (i, 2, p_max) {
7         if (!vis[i]) {
```

```

8         prime[p_sz++] = i;
9         phi[i] = i - 1;
10    }
11    for (LL j = 0; j < p_sz && (d = i * prime[j]) < p_max; ++j) {
12        vis[d] = 1;
13        if (i % prime[j] == 0) {
14            phi[d] = phi[i] * prime[j];
15            break;
16        }
17        else phi[d] = phi[i] * (prime[j] - 1);
18    }
19 }
20 }

```

● 线性筛 + 莫比乌斯函数

```

1  const LL p_max = 1E5 + 100;
2  LL mu[p_max] = {-1, 1};
3  void get_mu() {
4      static bool vis[p_max];
5      static LL prime[p_max], p_sz, d;
6      mu[1] = 1;
7      FOR (i, 2, p_max) {
8          if (!vis[i]) {
9              prime[p_sz++] = i;
10             mu[i] = -1;
11         }
12         for (LL j = 0; j < p_sz && (d = i * prime[j]) < p_max; ++j) {
13             vis[d] = 1;
14             if (i % prime[j] == 0) {
15                 mu[d] = 0;
16                 break;
17             }
18             else mu[d] = -mu[i];
19         }
20     }
21 }

```

亚线性筛

min_25

```

1  namespace min25 {
2      const int M = 1E6 + 100;
3      LL B, N;
4
5      // g(x)
6      inline LL pg(LL x) { return 1; }
7      inline LL ph(LL x) { return x % MOD; }
8      // Sum[g(i), {x, 2, x}]
9      inline LL psg(LL x) { return x % MOD - 1; }
10     inline LL psh(LL x) {
11         static LL inv2 = (MOD + 1) / 2;
12         x = x % MOD;
13         return x * (x + 1) % MOD * inv2 % MOD - 1;
14     }
15     // f(pp=p^k)
16     inline LL fpk(LL p, LL e, LL pp) { return (pp - pp / p) % MOD; }
17     // f(p) = fgh(g(p), h(p))
18     inline LL fgh(LL g, LL h) { return h - g; }
19
20     LL pr[M], pc, sg[M], sh[M];
21     void get_prime(LL n) {
22         static bool vis[M]; pc = 0;
23         FOR (i, 2, n + 1) {
24             if (!vis[i]) {
25                 pr[pc++] = i;
26                 sg[pc] = (sg[pc - 1] + pg(i)) % MOD;
27                 sh[pc] = (sh[pc - 1] + ph(i)) % MOD;

```

```

28     }
29     FOR (j, 0, pc) {
30         if (pr[j] * i > n) break;
31         vis[pr[j] * i] = 1;
32         if (i % pr[j] == 0) break;
33     }
34 }
35 }
36
37 LL w[M];
38 LL id1[M], id2[M], h[M], g[M];
39 inline LL id(LL x) { return x <= B ? id1[x] : id2[N / x]; }
40
41 LL go(LL x, LL k) {
42     if (x <= 1 || (k >= 0 && pr[k] > x)) return 0;
43     LL t = id(x);
44     LL ans = fgh((g[t] - sg[k + 1]), (h[t] - sh[k + 1]));
45     FOR (i, k + 1, pc) {
46         LL p = pr[i];
47         if (p * p > x) break;
48         ans -= fgh(pg(p), ph(p));
49         for (LL pp = p, e = 1; pp <= x; ++e, pp = pp * p)
50             ans += fpk(p, e, pp) * (1 + go(x / pp, i)) % MOD;
51     }
52     return ans % MOD;
53 }
54
55 LL solve(LL _N) {
56     N = _N;
57     B = sqrt(N + 0.5);
58     get_prime(B);
59     int sz = 0;
60     for (LL l = 1, v, r; l <= N; l = r + 1) {
61         v = N / l; r = N / v;
62         w[sz] = v; g[sz] = psg(v); h[sz] = psh(v);
63         if (v <= B) id1[v] = sz; else id2[r] = sz;
64         sz++;
65     }
66     FOR (k, 0, pc) {
67         LL p = pr[k];
68         FOR (i, 0, sz) {
69             LL v = w[i]; if (p * p > v) break;
70             LL t = id(v / p);
71             g[i] = (g[i] - (g[t] - sg[k]) * pg(p)) % MOD;
72             h[i] = (h[i] - (h[t] - sh[k]) * ph(p)) % MOD;
73         }
74     }
75     return (go(N, -1) % MOD + MOD + 1) % MOD;
76 }
77 }

```

杜教筛

求 $S(n) = \sum_{i=1}^n f(i)$, 其中 f 是一个积性函数。

构造一个积性函数 g , 那么由 $(f * g)(n) = \sum_{d|n} f(d)g(\frac{n}{d})$, 得到 $f(n) = (f * g)(n) - \sum_{d|n, d < n} f(d)g(\frac{n}{d})$ 。

$$g(1)S(n) = \sum_{i=1}^n (f * g)(i) - \sum_{i=1}^n \sum_{d|i, d < i} f(d)g(\frac{n}{d}) \quad (1)$$

$$\stackrel{t=\frac{i}{d}}{=} \sum_{i=1}^n (f * g)(i) - \sum_{t=2}^n g(t)S(\lfloor \frac{n}{t} \rfloor) \quad (2)$$

当然, 要能够由此计算 $S(n)$, 会对 f, g 提出一些要求:

- $f * g$ 要能够快速求前缀和。

- g 要能够快速求分段和（前缀和）。
- 对于正常的积性函数 $g(1) = 1$ ，所以不会有什么问题。

在预处理 $S(n)$ 前 $n^{\frac{2}{3}}$ 项的情况下复杂度是 $O(n^{\frac{2}{3}})$ 。

```

1 namespace dujiao {
2     const int M = 5E6;
3     LL f[M] = {0, 1};
4     void init() {
5         static bool vis[M];
6         static LL pr[M], p_sz, d;
7         FOR (i, 2, M) {
8             if (!vis[i]) { pr[p_sz++] = i; f[i] = -1; }
9             FOR (j, 0, p_sz) {
10                 if ((d = pr[j] * i) >= M) break;
11                 vis[d] = 1;
12                 if (i % pr[j] == 0) {
13                     f[d] = 0;
14                     break;
15                 } else f[d] = -f[i];
16             }
17         }
18         FOR (i, 2, M) f[i] += f[i - 1];
19     }
20     inline LL s_fg(LL n) { return 1; }
21     inline LL s_g(LL n) { return n; }
22
23     LL N, rd[M];
24     bool vis[M];
25     LL go(LL n) {
26         if (n < M) return f[n];
27         LL id = N / n;
28         if (vis[id]) return rd[id];
29         vis[id] = true;
30         LL& ret = rd[id] = s_fg(n);
31         for (LL l = 2, v, r; l <= n; l = r + 1) {
32             v = n / l; r = n / v;
33             ret -= (s_g(r) - s_g(l - 1)) * go(v);
34         }
35         return ret;
36     }
37     LL solve(LL n) {
38         N = n;
39         memset(vis, 0, sizeof vis);
40         return go(n);
41     }
42 }

```

素数测试

- 前置：快速乘、快速幂
- int 范围内只需检查 2, 7, 61
- long long 范围 2, 325, 9375, 28178, 450775, 9780504, 1795265022
- 3E15 内 2, 2570940, 880937, 610386380, 4130785767
- 4E13 内 2, 2570940, 211991001, 3749873356
- <http://miller-rabin.appspot.com/>

```

1 bool checkQ(LL a, LL n) {
2     if (n == 2 || a >= n) return 1;
3     if (n == 1 || !(n & 1)) return 0;
4     LL d = n - 1;
5     while (!(d & 1)) d >>= 1;
6     LL t = bin(a, d, n); // 不一定需要快速乘
7     while (d != n - 1 && t != 1 && t != n - 1) {
8         t = mul(t, t, n);
9         d <<= 1;
10    }
11    return t == n - 1 || d & 1;

```

```

12 }
13
14 bool primeQ(LL n) {
15     static vector<LL> t = {2, 325, 9375, 28178, 450775, 9780504, 1795265022};
16     if (n <= 1) return false;
17     for (LL k: t) if (!checkQ(k, n)) return false;
18     return true;
19 }

```

线性递推

```

1 // k 为 m 最高次数 且 a[m] == 1
2 namespace BerlekampMassey {
3     inline void up(LL& a, LL b) { (a += b) %= MOD; }
4
5     V mul(const V& a, const V& b, const V& m, int k) {
6         V r; r.resize(2 * k - 1);
7         FOR (i, 0, k)
8             FOR (j, 0, k)
9                 up(r[i + j], a[i] * b[j]);
10        FORD (i, k - 2, -1) {
11            FOR (j, 0, k)
12                up(r[i + j], r[i + k] * m[j]);
13            r.pop_back();
14        }
15        return r;
16    }
17
18    V pow(LL n, const V& m) {
19        int k = (int)m.size() - 1; assert(m[k] == -1 || m[k] == MOD - 1);
20        V r(k), x(k); r[0] = x[1] = 1;
21        for (; n >>= 1, x = mul(x, x, m, k))
22            if (n & 1) r = mul(x, r, m, k);
23        return r;
24    }
25
26    LL go(const V& a, const V& x, LL n) {
27        // a: (-1, a1, a2, ..., ak).reverse
28        // x: x1, x2, ..., xk
29        // x[n] = sum[a[i]*x[n-i], {i,1,k}]
30        int k = (int)a.size() - 1;
31        if (n <= k) return x[n - 1];
32        V r = pow(n - 1, a);
33        LL ans = 0;
34        FOR (i, 0, k)
35            up(ans, r[i] * x[i]);
36        return ans;
37    }
38
39    V BM(const V& x) {
40        V a = {-1}, b = {233};
41        FOR (i, 1, x.size()) {
42            b.push_back(0);
43            LL d = 0, la = a.size(), lb = b.size();
44            FOR (j, 0, la) up(d, a[j] * x[i - la + 1 + j]);
45            if (d == 0) continue;
46            V t; for (auto& v: b) t.push_back(d * v % MOD);
47            FOR (j, 0, a.size()) up(t[lb - 1 - j], a[la - 1 - j]);
48            if (lb > la) {
49                b = a;
50                LL inv = -get_inv(d, MOD);
51                for (auto& v: b) v = v * inv % MOD;
52            }
53            a.swap(t);
54        }
55        for (auto& v: a) up(v, MOD);
56        return a;
57    }
58 }

```

扩展欧几里得

- 求 $ax + by = \gcd(a, b)$ 的一组解
- 如果 a 和 b 互素, 那么 x 是 a 在模 b 下的逆元
- 注意 x 和 y 可能是负数

```
1 LL ex_gcd(LL a, LL b, LL &x, LL &y) {
2     if (b == 0) { x = 1; y = 0; return a; }
3     LL ret = ex_gcd(b, a % b, y, x);
4     y -= a / b * x;
5     return ret;
6 }
```

类欧几里得

- $m = \lfloor \frac{an+b}{c} \rfloor$.
- $f(a, b, c, n) = \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor$: 当 $a \geq c$ or $b \geq c$ 时, $f(a, b, c, n) = (\frac{a}{c})n(n+1)/2 + (\frac{b}{c})(n+1) + f(a \bmod c, b \bmod c, c, n)$; 否则 $f(a, b, c, n) = nm - f(c, c-b-1, a, m-1)$ 。
- $g(a, b, c, n) = \sum_{i=0}^n i \lfloor \frac{ai+b}{c} \rfloor$: 当 $a \geq c$ or $b \geq c$ 时, $g(a, b, c, n) = (\frac{a}{c})n(n+1)(2n+1)/6 + (\frac{b}{c})n(n+1)/2 + g(a \bmod c, b \bmod c, c, n)$; 否则 $g(a, b, c, n) = \frac{1}{2}(n(n+1)m - f(c, c-b-1, a, m-1) - h(c, c-b-1, a, m-1))$ 。
- $h(a, b, c, n) = \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor^2$: 当 $a \geq c$ or $b \geq c$ 时, $h(a, b, c, n) = (\frac{a}{c})^2 n(n+1)(2n+1)/6 + (\frac{b}{c})^2 (n+1) + (\frac{a}{c})(\frac{b}{c})n(n+1) + h(a \bmod c, b \bmod c, c, n) + 2(\frac{a}{c})g(a \bmod c, b \bmod c, c, n) + 2(\frac{b}{c})f(a \bmod c, b \bmod c, c, n)$; 否则 $h(a, b, c, n) = nm(m+1) - 2g(c, c-b-1, a, m-1) - 2f(c, c-b-1, a, m-1) - f(a, b, c, n)$ 。

逆元

- 如果 p 不是素数, 使用拓展欧几里得
- 前置模板: 快速幂 / 扩展欧几里得

```
1 inline LL get_inv(LL x, LL p) { return bin(x, p-2, p); }
2 LL get_inv(LL a, LL M) {
3     static LL x, y;
4     assert(exgcd(a, M, x, y) == 1);
5     return (x % M + M) % M;
6 }
```

- 预处理 $1 \sim n$ 的逆元

```
1 LL inv[N] = {-1, 1};
2 void inv_init(LL n, LL p) {
3     inv[1] = 1;
4     FOR (i, 2, n)
5         inv[i] = (p - p / i) * inv[p % i] % p;
6 }
```

- 预处理阶乘及其逆元

```
1 LL invf[M], fac[M] = {1};
2 void fac_inv_init(LL n, LL p) {
3     FOR (i, 1, n)
4         fac[i] = i * fac[i-1] % p;
5     invf[n-1] = bin(fac[n-1], p-2, p);
6     FORD (i, n-2, -1)
7         invf[i] = invf[i+1] * (i+1) % p;
8 }
```

组合数

- 如果数较小, 模较大时使用逆元
- 前置模板: 逆元-预处理阶乘及其逆元

```

1 inline LL C(LL n, LL m) { // n >= m >= 0
2     return n < m || m < 0 ? 0 : fac[n] * invf[m] % MOD * invf[n - m] % MOD;
3 }

```

- 如果模数较小，数字较大，使用 Lucas 定理
- 前置模板可选 1: 求组合数 (如果使用阶乘逆元，需 fac_inv_init(MOD, MOD);)
- 前置模板可选 2: 模数不固定下使用，无法单独使用。

```

1 LL C(LL n, LL m) { // m >= n >= 0
2     if (m - n < n) n = m - n;
3     if (n < 0) return 0;
4     LL ret = 1;
5     FOR (i, 1, n + 1)
6         ret = ret * (m - n + i) % MOD * bin(i, MOD - 2, MOD) % MOD;
7     return ret;
8 }

1 LL Lucas(LL n, LL m) { // m >= n >= 0
2     return m ? C(n % MOD, m % MOD) * Lucas(n / MOD, m / MOD) % MOD : 1;
3 }

```

- 组合数预处理

```

1 LL C[M][M];
2 void init_C(int n) {
3     FOR (i, 0, n) {
4         C[i][0] = C[i][i] = 1;
5         FOR (j, 1, i)
6             C[i][j] = (C[i - 1][j] + C[i - 1][j - 1]) % MOD;
7     }
8 }

```

斯特灵数

第一类斯特灵数

- 绝对值是 n 个元素划分为 k 个环排列的方案数。
- $s(n, k) = s(n - 1, k - 1) + (n - 1)s(n - 1, k)$

第二类斯特灵数

- n 个元素划分为 k 个等价类的方案数
- $S(n, k) = S(n - 1, k - 1) + kS(n - 1, k)$

```

1 S[0][0] = 1;
2 FOR (i, 1, N)
3     FOR (j, 1, i + 1) S[i][j] = (S[i - 1][j - 1] + j * S[i - 1][j]) % MOD;

```

FFT & NTT & FWT

- NTT
- 前置: 快速幂

```

1 LL wn[N << 2], rev[N << 2];
2 int NTT_init(int n_) {
3     int step = 0; int n = 1;
4     for (; n < n_; n <= 1) ++step;
5     FOR (i, 1, n)
6         rev[i] = (rev[i >> 1] >> 1) | ((i & 1) << (step - 1));
7     int g = bin(G, (MOD - 1) / n, MOD);
8     wn[0] = 1;
9     for (int i = 1; i <= n; ++i)
10         wn[i] = wn[i - 1] * g % MOD;
11     return n;
12 }

```

```

13
14 void NTT(LL a[], int n, int f) {
15     FOR (i, 0, n) if (i < rev[i])
16         std::swap(a[i], a[rev[i]]);
17     for (int k = 1; k < n; k <= 1) {
18         for (int i = 0; i < n; i += (k <= 1)) {
19             int t = n / (k <= 1);
20             FOR (j, 0, k) {
21                 LL w = f == 1 ? wn[t * j] : wn[n - t * j];
22                 LL x = a[i + j];
23                 LL y = a[i + j + k] * w % MOD;
24                 a[i + j] = (x + y) % MOD;
25                 a[i + j + k] = (x - y + MOD) % MOD;
26             }
27         }
28     }
29     if (f == -1) {
30         LL ninv = get_inv(n, MOD);
31         FOR (i, 0, n)
32             a[i] = a[i] * ninv % MOD;
33     }
34 }

```

- FFT
- n 需补成 2 的幂 (n 必须超过 a 和 b 的最高指数之和)

```

1 typedef double LD;
2 const LD PI = acos(-1);
3 struct C {
4     LD r, i;
5     C(LD r = 0, LD i = 0): r(r), i(i) {}
6 };
7 C operator + (const C& a, const C& b) {
8     return C(a.r + b.r, a.i + b.i);
9 }
10 C operator - (const C& a, const C& b) {
11     return C(a.r - b.r, a.i - b.i);
12 }
13 C operator * (const C& a, const C& b) {
14     return C(a.r * b.r - a.i * b.i, a.r * b.i + a.i * b.r);
15 }
16
17 void FFT(C x[], int n, int p) {
18     for (int i = 0, t = 0; i < n; ++i) {
19         if (i > t) swap(x[i], x[t]);
20         for (int j = n >> 1; (t ^= j) < j; j >= 1);
21     }
22     for (int h = 2; h <= n; h <= 1) {
23         C wn(cos(p * 2 * PI / h), sin(p * 2 * PI / h));
24         for (int i = 0; i < n; i += h) {
25             C w(1, 0), u;
26             for (int j = i, k = h >> 1; j < i + k; ++j) {
27                 u = x[j + k] * w;
28                 x[j + k] = x[j] - u;
29                 x[j] = x[j] + u;
30                 w = w * wn;
31             }
32         }
33     }
34     if (p == -1)
35         FOR (i, 0, n)
36             x[i].r /= n;
37 }
38
39 void conv(C a[], C b[], int n) {
40     FFT(a, n, 1);
41     FFT(b, n, 1);
42     FOR (i, 0, n)
43         a[i] = a[i] * b[i];
44     FFT(a, n, -1);
45 }

```


- FWT

```

1  template<typename T>
2  void fwt(LL a[], int n, T f) {
3      for (int d = 1; d < n; d *= 2)
4          for (int i = 0, t = d * 2; i < n; i += t)
5              FOR (j, 0, d)
6                  f(a[i + j], a[i + j + d]);
7  }
8
9  void AND(LL& a, LL& b) { a += b; }
10 void OR(LL& a, LL& b) { b += a; }
11 void XOR (LL& a, LL& b) {
12     LL x = a, y = b;
13     a = (x + y) % MOD;
14     b = (x - y + MOD) % MOD;
15 }

```

simpson 自适应积分

```

1  LD simpson(LD l, LD r) {
2      LD c = (l + r) / 2;
3      return (f(l) + 4 * f(c) + f(r)) * (r - l) / 6;
4  }
5
6  LD asr(LD l, LD r, LD eps, LD S) {
7      LD m = (l + r) / 2;
8      LD L = simpson(l, m), R = simpson(m, r);
9      if (fabs(L + R - S) < 15 * eps) return L + R + (L + R - S) / 15;
10     return asr(l, m, eps / 2, L) + asr(m, r, eps / 2, R);
11 }
12
13 LD asr(LD l, LD r, LD eps) { return asr(l, r, eps, simpson(l, r)); }

```

- FWT

```

1  template<typename T>
2  void fwt(LL a[], int n, T f) {
3      for (int d = 1; d < n; d *= 2)
4          for (int i = 0, t = d * 2; i < n; i += t)
5              FOR (j, 0, d)
6                  f(a[i + j], a[i + j + d]);
7  }
8
9  auto f = [] (LL& a, LL& b) { // xor
10     LL x = a, y = b;
11     a = (x + y) % MOD;
12     b = (x - y + MOD) % MOD;
13 };

```

快速乘

```

1  LL mul(LL a, LL b, LL m) {
2      LL ret = 0;
3      while (b) {
4          if (b & 1) {
5              ret += a;
6              if (ret >= m) ret -= m;
7          }
8          a += a;
9          if (a >= m) a -= m;
10         b >>= 1;
11     }
12     return ret;
13 }

```

- $O(1)$

```

1 LL mul(LL u, LL v, LL p) {
2     return (u * v - LL((long double) u * v / p) * p + p) % p;
3 }

```

快速幂

- 如果模数是素数，则可在函数体内加上 $n \% = \text{MOD} - 1$ ；（费马小定理）。

```

1 LL bin(LL x, LL n, LL MOD) {
2     LL ret = MOD != 1;
3     for (x %= MOD; n; n >>= 1, x = x * x % MOD)
4         if (n & 1) ret = ret * x % MOD;
5     return ret;
6 }

```

- 防爆 LL
- 前置模板：快速乘

```

1 LL bin(LL x, LL n, LL MOD) {
2     LL ret = MOD != 1;
3     for (x %= MOD; n; n >>= 1, x = mul(x, x, MOD))
4         if (n & 1) ret = mul(ret, x, MOD);
5     return ret;
6 }

```

高斯消元

- n - 方程个数， m - 变量个数， a 是 $n * (m + 1)$ 的增广矩阵，free 是否为自由变量
- 返回自由变量个数，-1 无解，-2 无整数解
- 浮点数版本

```

1 typedef double LD;
2 const LD eps = 1E-10;
3 const int maxn = 2000 + 10;
4
5 int n, m;
6 LD a[maxn][maxn], x[maxn];
7 bool free_x[maxn];
8
9 inline int sgn(LD x) { return (x > eps) - (x < -eps); }
10
11
12 int gauss(LD a[maxn][maxn], int n, int m) {
13     memset(free_x, 1, sizeof free_x); memset(x, 0, sizeof x);
14     int r = 0, c = 0;
15     while (r < n && c < m) {
16         int m_r = r;
17         FOR (i, r + 1, n)
18             if (fabs(a[i][c]) > fabs(a[m_r][c])) m_r = i;
19         if (m_r != r)
20             FOR (j, c, m + 1)
21                 swap(a[r][j], a[m_r][j]);
22         if (!sgn(a[r][c])) {
23             a[r][c] = 0;
24             ++c;
25             continue;
26         }
27         FOR (i, r + 1, n)
28             if (a[i][c]) {
29                 LD t = a[i][c] / a[r][c];
30                 FOR (j, c, m + 1) a[i][j] -= a[r][j] * t;
31             }
32         ++r; ++c;
33         // FOR (i, 0, n)
34         //     FOR (j, 0, m + 1)
35         //         printf("%.2f%c", a[i][j], j == _j - 1 ? '\n' : ' '); puts("");

```

```

36     }
37     FOR (i, r, n)
38         if (sgn(a[i][m])) return -1;
39     if (r < m) {
40         FORD (i, r - 1, -1) {
41             int f_cnt = 0, k = -1;
42             FOR (j, 0, m)
43                 if (sgn(a[i][j]) && free_x[j]) {
44                     ++f_cnt;
45                     k = j;
46                 }
47             if(f_cnt > 0) continue;
48             LD s = a[i][m];
49             FOR (j, 0, m)
50                 if (j != k) s -= a[i][j] * x[j];
51             x[k] = s / a[i][k];
52             free_x[k] = 0;
53         }
54         return m - r;
55     }
56     FORD (i, m - 1, -1) {
57         LD s = a[i][m];
58         FOR (j, i + 1, m)
59             s -= a[i][j] * x[j];
60         x[i] = s / a[i][i];
61     }
62     return 0;
63 }

```

- 数据

```

3 4
1 1 -2 2
2 -3 5 1
4 -1 1 5
5 0 -1 7
// many

3 4
1 1 -2 2
2 -3 5 1
4 -1 -1 5
5 0 -1 0 2
// no

3 4
1 1 -2 2
2 -3 5 1
4 -1 1 5
5 0 1 0 7
// one

```

质因数分解

- 前置模板：素数筛
- 带指数

```

1 LL factor[30], f_sz, factor_exp[30];
2 void get_factor(LL x) {
3     f_sz = 0;
4     LL t = sqrt(x + 0.5);
5     for (LL i = 0; pr[i] <= t; ++i)
6         if (x % pr[i] == 0) {
7             factor_exp[f_sz] = 0;

```

```

8         while (x % pr[i] == 0) {
9             x /= pr[i];
10            ++factor_exp[f_sz];
11        }
12        factor[f_sz++] = pr[i];
13    }
14    if (x > 1) {
15        factor_exp[f_sz] = 1;
16        factor[f_sz++] = x;
17    }
18 }

```

- 不带指数

```

1 LL factor[30], f_sz;
2 void get_factor(LL x) {
3     f_sz = 0;
4     LL t = sqrt(x + 0.5);
5     for (LL i = 0; pr[i] <= t; ++i)
6         if (x % pr[i] == 0) {
7             factor[f_sz++] = pr[i];
8             while (x % pr[i] == 0) x /= pr[i];
9         }
10    if (x > 1) factor[f_sz++] = x;
11 }

```

原根

- 前置模板：素数筛，快速幂，分解质因数
- 要求 p 为质数

```

1 LL find_smallest_primitive_root(LL p) {
2     get_factor(p - 1);
3     FOR (i, 2, p) {
4         bool flag = true;
5         FOR (j, 0, f_sz)
6             if (bin(i, (p - 1) / factor[j], p) == 1) {
7                 flag = false;
8                 break;
9             }
10        if (flag) return i;
11    }
12    assert(0); return -1;
13 }

```

公式

一些数论公式

- 当 $x \geq \phi(p)$ 时有 $a^x \equiv a^{x \bmod \phi(p) + \phi(p)} \pmod{p}$
- $\mu^2(n) = \sum_{d^2|n} \mu(d)$
- $\sum_{d|n} \varphi(d) = n$
- $\sum_{d|n} 2^{\omega(d)} = \sigma_0(n^2)$
- $\sum_{d|n} \mu^2(d) = 2^{\omega(n)}$

一些数论函数求和的例子

- $\sum_{i=1}^n i[\gcd(i, n) = 1] = \frac{n\varphi(n) + [n=1]}{2}$
- $\sum_{i=1}^n \sum_{j=1}^m [\gcd(i, j) = x] = \sum_d \mu(d) \lfloor \frac{n}{dx} \rfloor \lfloor \frac{m}{dx} \rfloor$
- $\sum_{i=1}^n \sum_{j=1}^m \gcd(i, j) = \sum_{i=1}^n \sum_{j=1}^m \sum_{d|\gcd(i, j)} \varphi(d) = \sum_d \varphi(d) \lfloor \frac{n}{d} \rfloor \lfloor \frac{m}{d} \rfloor$

- $S(n) = \sum_{i=1}^n \mu(i) = 1 - \sum_{i=1}^n \sum_{d|i, d < i} \mu(d) \stackrel{t=\frac{i}{d}}{=} 1 - \sum_{t=2}^n S(\lfloor \frac{n}{t} \rfloor)$
- 利用 $[n=1] = \sum_{d|n} \mu(d)$
- $S(n) = \sum_{i=1}^n \varphi(i) = \sum_{i=1}^n i - \sum_{i=1}^n \sum_{d|i, d < i} \varphi(i) \stackrel{t=\frac{i}{d}}{=} \frac{i(i+1)}{2} - \sum_{t=2}^n S(\frac{n}{t})$
- 利用 $n = \sum_{d|n} \varphi(d)$
- $\sum_{i=1}^n \mu^2(i) = \sum_{i=1}^n \sum_{d^2|i} \mu(d) = \sum_{d=1}^{\lfloor \sqrt{n} \rfloor} \mu(d) \lfloor \frac{n}{d^2} \rfloor$
- $\sum_{i=1}^n \sum_{j=1}^n gcd^2(i, j) = \sum_d d^2 \sum_t \mu(t) \lfloor \frac{n}{dt} \rfloor^2$
 $\stackrel{x=dt}{=} \sum_x \lfloor \frac{n}{x} \rfloor^2 \sum_{d|x} d^2 \mu(\frac{x}{d})$

斐波那契数列性质

- $F_{a+b} = F_{a-1} \cdot F_b + F_a \cdot F_{b+1}$
- $F_1 + F_3 + \dots + F_{2n-1} = F_{2n}, F_2 + F_4 + \dots + F_{2n} = F_{2n+1} - 1$
- $\sum_{i=1}^n F_i = F_{n+2} - 1$
- $\sum_{i=1}^n F_i^2 = F_n \cdot F_{n+1}$
- $F_n^2 = (-1)^{n-1} + F_{n-1} \cdot F_{n+1}$
- $gcd(F_a, F_b) = F_{gcd(a,b)}$
- 模 n 周期 (皮萨诺周期)
 - $\pi(p^k) = p^{k-1} \pi(p)$
 - $\pi(nm) = lcm(\pi(n), \pi(m)), \forall n \perp m$
 - $\pi(2) = 3, \pi(5) = 20$
 - $\forall p \equiv \pm 1 \pmod{10}, \pi(p) | p-1$
 - $\forall p \equiv \pm 2 \pmod{5}, \pi(p) | 2p+2$

常见生成函数

- $(1+ax)^n = \sum_{k=0}^n \binom{n}{k} a^k x^k$
- $\frac{1-x^{r+1}}{1-x} = \sum_{k=0}^n x^k$
- $\frac{1}{1-ax} = \sum_{k=0}^{\infty} a^k x^k$
- $\frac{1}{(1-x)^2} = \sum_{k=0}^{\infty} (k+1)x^k$
- $\frac{1}{(1-x)^n} = \sum_{k=0}^{\infty} \binom{n+k-1}{k} x^k$
- $e^x = \sum_{k=0}^{\infty} \frac{x^k}{k!}$
- $\ln(1+x) = \sum_{k=0}^{\infty} \frac{(-1)^{k+1}}{k} x^k$

佩尔方程

若一个丢番图方程具有以下形式: $x^2 - ny^2 = 1$ 。且 n 为正整数, 则称此二元二次不定方程为**佩尔方程**。

若 n 是完全平方数, 则这个方程式只有平凡解 $(\pm 1, 0)$ (实际上对任意的 n , $(\pm 1, 0)$ 都是解)。对于其余情况, 拉格朗日证明了佩尔方程总有非平凡解。而这些解可由 \sqrt{n} 的连分数求出。

$$x = [a_0; a_1, a_2, a_3] = x = a_0 + \frac{1}{a_1 + \frac{1}{a_2 + \frac{1}{a_3 + \ddots}}}$$

设 $\frac{p_i}{q_i}$ 是 \sqrt{n} 的连分数表示: $[a_0; a_1, a_2, a_3, \dots]$ 的渐近分数列, 由连分数理论知存在 i 使得 (p_i, q_i) 为佩尔方程的解。取其中最小的 i , 将对应的 (p_i, q_i) 称为佩尔方程的基本解, 或最小解, 记作 (x_1, y_1) , 则所有的解 (x_i, y_i) 可表示成如下形式: $x_i + y_i \sqrt{n} = (x_1 + y_1 \sqrt{n})^i$ 。或者由以下的递回关系式得到:

$$x_{i+1} = x_1 x_i + n y_1 y_i, y_{i+1} = x_1 y_i + y_1 x_i$$

但是: 佩尔方程千万不要去推 (虽然推起来很有趣, 但结果不一定好看, 会是两个式子)。记住佩尔方程结果的形式通常是 $a_n = k a_{n-1} - a_{n-2}$ (a_{n-2} 前的系数通常是 -1)。暴力 / 凑出两个基础解之后加上一个 0, 容易解出 k 并验证。

Burnside & Polya

- $|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|$

注: X^g 是 g 下的不动点数量, 也就是说有多少种东西用 g 作用之后可以保持不变。

- $|Y^X/G| = \frac{1}{|G|} \sum_{g \in G} m^{c(g)}$

注: 用 m 种颜色染色, 然后对于某一种置换 g , 有 $c(g)$ 个置换环, 为了保证置换后颜色仍然相同, 每个置换环必须染成同色。

皮克定理

$$2S = 2a + b - 2$$

- S 多边形面积
- a 多边形内部点数
- b 多边形边上点数

莫比乌斯反演

- $g(n) = \sum_{d|n} f(d) \Leftrightarrow f(n) = \sum_{d|n} \mu(d) g(\frac{n}{d})$
- $f(n) = \sum_{n|d} g(d) \Leftrightarrow g(n) = \sum_{n|d} \mu(\frac{d}{n}) f(d)$

低阶等幂求和

- $\sum_{i=1}^n i^1 = \frac{n(n+1)}{2} = \frac{1}{2}n^2 + \frac{1}{2}n$
- $\sum_{i=1}^n i^2 = \frac{n(n+1)(2n+1)}{6} = \frac{1}{3}n^3 + \frac{1}{2}n^2 + \frac{1}{6}n$
- $\sum_{i=1}^n i^3 = \left[\frac{n(n+1)}{2} \right]^2 = \frac{1}{4}n^4 + \frac{1}{2}n^3 + \frac{1}{4}n^2$
- $\sum_{i=1}^n i^4 = \frac{n(n+1)(2n+1)(3n^2+3n-1)}{30} = \frac{1}{5}n^5 + \frac{1}{2}n^4 + \frac{1}{3}n^3 - \frac{1}{30}n$
- $\sum_{i=1}^n i^5 = \frac{n^2(n+1)^2(2n^2+2n-1)}{12} = \frac{1}{6}n^6 + \frac{1}{2}n^5 + \frac{5}{12}n^4 - \frac{1}{12}n^2$

二次剩余

URAL 1132

```

1 LL q1, q2, w;
2 struct P { // x + y * sqrt(w)
3     LL x, y;
4 };
5
6 P pmul(const P& a, const P& b, LL p) {
7     P res;
8     res.x = (a.x * b.x + a.y * b.y % p * w) % p;
9     res.y = (a.x * b.y + a.y * b.x) % p;
10    return res;
11 }
12
13 P bin(P x, LL n, LL MOD) {

```

```

14     P ret = {1, 0};
15     for (; n; n >>= 1, x = pmul(x, x, MOD))
16         if (n & 1) ret = pmul(ret, x, MOD);
17     return ret;
18 }
19 LL Legendre(LL a, LL p) { return bin(a, (p - 1) >> 1, p); }
20
21 LL equation_solve(LL b, LL p) {
22     if (p == 2) return 1;
23     if ((Legendre(b, p) + 1) % p == 0)
24         return -1;
25     LL a;
26     while (true) {
27         a = rand() % p;
28         w = ((a * a - b) % p + p) % p;
29         if ((Legendre(w, p) + 1) % p == 0)
30             break;
31     }
32     return bin({a, 1}, (p + 1) >> 1, p).x;
33 }
34
35 int main() {
36     int T; cin >> T;
37     while (T--) {
38         LL a, p; cin >> a >> p;
39         a = a % p;
40         LL x = equation_solve(a, p);
41         if (x == -1) {
42             puts("No root");
43         } else {
44             LL y = p - x;
45             if (x == y) cout << x << endl;
46             else cout << min(x, y) << " " << max(x, y) << endl;
47         }
48     }
49 }

```

中国剩余定理

- 无解返回 -1
- 前置模板：拓展欧几里得

```

1 LL CRT(LL *m, LL *r, LL n) {
2     if (!n) return 0;
3     LL M = m[0], R = r[0], x, y, d;
4     FOR (i, 1, n) {
5         d = ex_gcd(M, m[i], x, y);
6         if ((r[i] - R) % d) return -1;
7         x = (r[i] - R) / d * x % (m[i] / d);
8         R += x * M;
9         M = M / d * m[i];
10        R %= M;
11    }
12    return R >= 0 ? R : R + M;
13 }

```

伯努利数和等幂求和

- 预处理逆元
- 预处理组合数
- $\sum_{i=0}^n i^k = \frac{1}{k+1} \sum_{i=0}^k \binom{k+1}{i} B_{k+1-i} (n+1)^i$.
- 也可以 $\sum_{i=0}^n i^k = \frac{1}{k+1} \sum_{i=0}^k \binom{k+1}{i} B_{k+1-i}^+ n^i$ 。区别在于 $B_1^+ = 1/2$ 。(心态崩了)

```

1 namespace Bernoulli {
2     const int M = 100;
3     LL inv[M] = {-1, 1};

```

```

4 void inv_init(LL n, LL p) {
5     FOR (i, 2, n)
6         inv[i] = (p - p / i) * inv[p % i] % p;
7 }
8
9 LL C[M][M];
10 void init_C(int n) {
11     FOR (i, 0, n) {
12         C[i][0] = C[i][i] = 1;
13         FOR (j, 1, i)
14             C[i][j] = (C[i - 1][j] + C[i - 1][j - 1]) % MOD;
15     }
16 }
17
18 LL B[M] = {1};
19 void init() {
20     inv_init(M, MOD);
21     init_C(M);
22     FOR (i, 1, M - 1) {
23         LL& s = B[i] = 0;
24         FOR (j, 0, i)
25             s += C[i + 1][j] * B[j] % MOD;
26         s = (s % MOD * -inv[i + 1] % MOD + MOD) % MOD;
27     }
28 }
29
30 LL p[M] = {1};
31 LL go(LL n, LL k) {
32     n %= MOD;
33     if (k == 0) return n;
34     FOR (i, 1, k + 2)
35         p[i] = p[i - 1] * (n + 1) % MOD;
36     LL ret = 0;
37     FOR (i, 1, k + 2)
38         ret += C[k + 1][i] * B[k + 1 - i] % MOD * p[i] % MOD;
39     ret = ret % MOD * inv[k + 1] % MOD;
40     return ret;
41 }
42 }

```

单纯形

- 要求有基本解，也就是 x 为零向量可行
- v 要初始化为 0, n 表示向量长度, m 表示约束个数

```

1 // min{ b x } / max { c x }
2 // A x >= c / A x <= b
3 // x >= 0
4 namespace lp {
5     int n, m;
6     double a[M][N], b[M], c[N], v;
7
8     void pivot(int l, int e) {
9         b[l] /= a[l][e];
10        FOR (j, 0, n) if (j != e) a[l][j] /= a[l][e];
11        a[l][e] = 1 / a[l][e];
12
13        FOR (i, 0, m)
14            if (i != l && fabs(a[i][e]) > 0) {
15                b[i] -= a[i][e] * b[l];
16                FOR (j, 0, n)
17                    if (j != e) a[i][j] -= a[i][e] * a[l][j];
18                a[i][e] = -a[i][e] * a[l][e];
19            }
20        v += c[e] * b[l];
21        FOR (j, 0, n) if (j != e) c[j] -= c[e] * a[l][j];
22        c[e] = -c[e] * a[l][e];
23    }
24    double simplex() {

```



```

25     while (1) {
26         v = 0;
27         int e = -1, l = -1;
28         FOR (i, 0, n) if (c[i] > eps) { e = i; break; }
29         if (e == -1) return v;
30         double t = INF;
31         FOR (i, 0, m)
32             if (a[i][e] > eps && t > b[i] / a[i][e]) {
33                 t = b[i] / a[i][e];
34                 l = i;
35             }
36         if (l == -1) return INF;
37         pivot(l, e);
38     }
39 }
40 }

```

图论

LCA

● 倍增

```

1 void dfs(int u, int fa) {
2     pa[u][0] = fa; dep[u] = dep[fa] + 1;
3     FOR (i, 1, SP) pa[u][i] = pa[pa[u][i-1]][i-1];
4     for (int& v: G[u]) {
5         if (v == fa) continue;
6         dfs(v, u);
7     }
8 }
9
10 int lca(int u, int v) {
11     if (dep[u] < dep[v]) swap(u, v);
12     int t = dep[u] - dep[v];
13     FOR (i, 0, SP) if (t & (1 << i)) u = pa[u][i];
14     FORD (i, SP-1, -1) {
15         int uu = pa[u][i], vv = pa[v][i];
16         if (uu != vv) { u = uu; v = vv; }
17     }
18     return u == v ? u : pa[u][0];
19 }

```

最短路

```

1 bool BF() {
2     queue<int> q;
3     FOR (i, 1, n) d[i] = INF;
4     d[0] = 0; inq[0] = true; q.push(0);
5     while (!q.empty()) {
6         int u = q.front(); q.pop();
7         inq[u] = false;
8         for (E& e: G[u]) {
9             int v = e.to;
10            if (d[u] < INF && d[v] > d[u] + e.d) {
11                d[v] = d[u] + e.d;
12                if (!inq[v]) {
13                    q.push(v); inq[v] = true;
14                    if (++cnt[v] > n) return false;
15                }
16            }
17        }
18    }
19    return true;
20 }

```

网络流

- 最大流

```
1 struct E {
2     int to, cp;
3     E(int to, int cp): to(to), cp(cp) {}
4 };
5
6 struct Dinic {
7     static const int M = 1E5 * 5;
8     int m, s, t;
9     vector<E> edges;
10    vector<int> G[M];
11    int d[M];
12    int cur[M];
13
14    void init(int n, int s, int t) {
15        this->s = s; this->t = t;
16        for (int i = 0; i <= n; i++) G[i].clear();
17        edges.clear(); m = 0;
18    }
19
20    void addedge(int u, int v, int cap) {
21        edges.emplace_back(v, cap);
22        edges.emplace_back(u, 0);
23        G[u].push_back(m++);
24        G[v].push_back(m++);
25    }
26
27    bool BFS() {
28        memset(d, 0, sizeof d);
29        queue<int> Q;
30        Q.push(s); d[s] = 1;
31        while (!Q.empty()) {
32            int x = Q.front(); Q.pop();
33            for (int& i: G[x]) {
34                E &e = edges[i];
35                if (!d[e.to] && e.cp > 0) {
36                    d[e.to] = d[x] + 1;
37                    Q.push(e.to);
38                }
39            }
40        }
41        return d[t];
42    }
43
44    int DFS(int u, int cp) {
45        if (u == t || !cp) return cp;
46        int tmp = cp, f;
47        for (int& i = cur[u]; i < G[u].size(); i++) {
48            E& e = edges[G[u][i]];
49            if (d[u] + 1 == d[e.to]) {
50                f = DFS(e.to, min(cp, e.cp));
51                e.cp -= f;
52                edges[G[u][i] ^ 1].cp += f;
53                cp -= f;
54                if (!cp) break;
55            }
56        }
57        return tmp - cp;
58    }
59
60    int go() {
61        int flow = 0;
62        while (BFS()) {
63            memset(cur, 0, sizeof cur);
64            flow += DFS(s, INF);
65        }
66        return flow;
67    }
```

```
68 } DC;
```

● 费用流

```
1 struct E {
2     int from, to, cp, v;
3     E() {}
4     E(int f, int t, int cp, int v) : from(f), to(t), cp(cp), v(v) {}
5 };
6
7 struct MCMF {
8     int n, m, s, t;
9     vector<E> edges;
10    vector<int> G[maxn];
11    bool inq[maxn]; //是否在队列
12    int d[maxn]; //Bellman_ford 单源最短路径
13    int p[maxn]; //p[i] 表从 s 到 i 的最小费用路径上的最后一条弧编号
14    int a[maxn]; //a[i] 表示从 s 到 i 的最小残量
15
16    void init(int _n, int _s, int _t) {
17        n = _n; s = _s; t = _t;
18        FOR (i, 0, n + 1) G[i].clear();
19        edges.clear(); m = 0;
20    }
21
22    void addedge(int from, int to, int cap, int cost) {
23        edges.emplace_back(from, to, cap, cost);
24        edges.emplace_back(to, from, 0, -cost);
25        G[from].push_back(m++);
26        G[to].push_back(m++);
27    }
28
29    bool BellmanFord(int &flow, int &cost) {
30        FOR (i, 0, n + 1) d[i] = INF;
31        memset(inq, 0, sizeof inq);
32        d[s] = 0, a[s] = INF, inq[s] = true;
33        queue<int> Q; Q.push(s);
34        while (!Q.empty()) {
35            int u = Q.front(); Q.pop();
36            inq[u] = false;
37            for (int& idx: G[u]) {
38                E &e = edges[idx];
39                if (e.cp && d[e.to] > d[u] + e.v) {
40                    d[e.to] = d[u] + e.v;
41                    p[e.to] = idx;
42                    a[e.to] = min(a[u], e.cp);
43                    if (!inq[e.to]) {
44                        Q.push(e.to);
45                        inq[e.to] = true;
46                    }
47                }
48            }
49        }
50        if (d[t] == INF) return false;
51        flow += a[t];
52        cost += a[t] * d[t];
53        int u = t;
54        while (u != s) {
55            edges[p[u]].cp -= a[t];
56            edges[p[u] ^ 1].cp += a[t];
57            u = edges[p[u]].from;
58        }
59        return true;
60    }
61
62    int go() {
63        int flow = 0, cost = 0;
64        while (BellmanFord(flow, cost));
65        return cost;
66    }
67 } MM;
```

- zkw 费用流（代码长度没有优势）
- 不允许有负权边

```

1  struct E {
2      int to, cp, v;
3      E() {}
4      E(int to, int cp, int v): to(to), cp(cp), v(v) {}
5  };
6
7  struct MCMF {
8      int n, m, s, t, cost, D;
9      vector<E> edges;
10     vector<int> G[maxn];
11     bool vis[maxn];
12
13     void init(int _n, int _s, int _t) {
14         n = _n; s = _s; t = _t;
15         FOR (i, 0, n + 1) G[i].clear();
16         edges.clear(); m = 0;
17     }
18
19     void addedge(int from, int to, int cap, int cost) {
20         edges.emplace_back(to, cap, cost);
21         edges.emplace_back(from, 0, -cost);
22         G[from].push_back(m++);
23         G[to].push_back(m++);
24     }
25
26     int aug(int u, int cp) {
27         if (u == t) {
28             cost += D * cp;
29             return cp;
30         }
31         vis[u] = true;
32         int tmp = cp;
33         for (int idx: G[u]) {
34             E& e = edges[idx];
35             if (e.cp && !e.v && !vis[e.to]) {
36                 int f = aug(e.to, min(cp, e.cp));
37                 e.cp -= f;
38                 edges[idx ^ 1].cp += f;
39                 cp -= f;
40                 if (!cp) break;
41             }
42         }
43         return tmp - cp;
44     }
45
46     bool modlabel() {
47         int d = INF;
48         FOR (u, 0, n + 1)
49             if (vis[u])
50                 for (int& idx: G[u]) {
51                     E& e = edges[idx];
52                     if (e.cp && !vis[e.to]) d = min(d, e.v);
53                 }
54         if (d == INF) return false;
55         FOR (u, 0, n + 1)
56             if (vis[u])
57                 for (int& idx: G[u]) {
58                     edges[idx].v -= d;
59                     edges[idx ^ 1].v += d;
60                 }
61         D += d;
62         return true;
63     }
64
65     int go(int k) {
66         cost = D = 0;
67         int flow = 0;
68         while (true) {

```

```

69         memset(vis, 0, sizeof vis);
70         int t = aug(s, INF);
71         if (!t && !modlabel()) break;
72         flow += t;
73     }
74     return cost;
75 }
76 } MM;

```

● 带下界网络流:

- 无源汇: $u \rightarrow v$ 边容量为 $[l, r]$, 连容量 $r - l$, 虚拟源点到 v 连 l , u 到虚拟汇点连 l 。
- 有源汇: 为了让流能循环使用, 连 $T \rightarrow S$, 容量 ∞ 。
- 最大流: 跑完可行流后, 加 $S' \rightarrow S$, $T \rightarrow T'$, 最大流就是答案 ($T \rightarrow S$ 的流量自动退回去了, 这一部分就是下界部分的流量)。
- 最小流: T 到 S 的那条边的实际流量, 减去删掉那条边后 T 到 S 的最大流。
- 网上说可能会减成负的, 还要有限地供应 S 之后, 再跑一遍 S 到 T 的。
- 费用流: 必要的部分 (下界以下的) 不要钱, 剩下的按照最大流。

树上路径交

```

1  int intersection(int x, int y, int xx, int yy) {
2      int t[4] = {lca(x, xx), lca(x, yy), lca(y, xx), lca(y, yy)};
3      sort(t, t + 4);
4      int r = lca(x, y), rr = lca(xx, yy);
5      if (dep[t[0]] < min(dep[r], dep[rr]) || dep[t[2]] < max(dep[r], dep[rr]))
6          return 0;
7      int tt = lca(t[2], t[3]);
8      int ret = 1 + dep[t[2]] + dep[t[3]] - dep[tt] * 2;
9      return ret;
10 }

```

树上点分治

```

1  int get_sz(int u, int fa) {
2      int& s = sz[u] = 1;
3      for (E& e: G[u]) {
4          int v = e.to;
5          if (vis[v] || v == fa) continue;
6          s += get_sz(v, u);
7      }
8      return s;
9  }
10
11 void get_rt(int u, int fa, int s, int& m, int& rt) {
12     int t = s - sz[u];
13     for (E& e: G[u]) {
14         int v = e.to;
15         if (vis[v] || v == fa) continue;
16         get_rt(v, u, s, m, rt);
17         t = max(t, sz[v]);
18     }
19     if (t < m) { m = t; rt = u; }
20 }
21
22 void dfs(int u) {
23     int tmp = INF; get_rt(u, -1, get_sz(u, -1), tmp, u);
24     vis[u] = true;
25     get_dep(u, -1, 0);
26     // ...
27     for (E& e: G[u]) {
28         int v = e.to;
29         if (vis[v]) continue;
30         // ...
31         dfs(v);
32     }
33 }

```

- 动态点分治

```

1  const int maxn = 15E4 + 100, INF = 1E9;
2  struct E {
3      int to, d;
4  };
5  vector<E> G[maxn];
6  int n, Q, w[maxn];
7  LL A, ans;
8
9  bool vis[maxn];
10 int sz[maxn];
11
12 int get_rt(int u) {
13     // dbg(u);
14     static int q[N], fa[N], sz[N], mx[N];
15     int p = 0, cur = -1;
16     q[p++] = u; fa[u] = -1;
17     while (++cur < p) {
18         u = q[cur]; mx[u] = 0; sz[u] = 1;
19         for (int& v: G[u])
20             if (!vis[v] && v != fa[u]) fa[q[p++]] = v;
21     }
22     FORD (i, p - 1, -1) {
23         u = q[i];
24         mx[u] = max(mx[u], p - sz[u]);
25         if (mx[u] * 2 <= p) return u;
26         sz[fa[u]] += sz[u];
27         mx[fa[u]] = max(mx[fa[u]], sz[u]);
28     }
29     assert(0);
30 }
31
32 int get_sz(int u, int fa) {
33     int& s = sz[u] = 1;
34     for (E& e: G[u]) {
35         int v = e.to;
36         if (vis[v] || v == fa) continue;
37         s += get_sz(v, u);
38     }
39     return s;
40 }
41
42 void get_rt(int u, int fa, int s, int& m, int& rt) {
43     int t = s - sz[u];
44     for (E& e: G[u]) {
45         int v = e.to;
46         if (vis[v] || v == fa) continue;
47         get_rt(v, u, s, m, rt);
48         t = max(t, sz[v]);
49     }
50     if (t < m) { m = t; rt = u; }
51 }
52
53 int dep[maxn], md[maxn];
54 void get_dep(int u, int fa, int d) {
55     dep[u] = d; md[u] = 0;
56     for (E& e: G[u]) {
57         int v = e.to;
58         if (vis[v] || v == fa) continue;
59         get_dep(v, u, d + e.d);
60         md[u] = max(md[u], md[v] + 1);
61     }
62 }
63
64 struct P {
65     int w;
66     LL s;
67 };
68 using VP = vector<P>;
69 struct R {
70     VP *rt, *rt2;

```

```

71     int dep;
72 };
73 VP pool[maxn << 1], *pit = pool;
74 vector<R> tr[maxn];
75
76 void go(int u, int fa, VP* rt, VP* rt2) {
77     tr[u].push_back({rt, rt2, dep[u]});
78     for (E& e: G[u]) {
79         int v = e.to;
80         if (v == fa || vis[v]) continue;
81         go(v, u, rt, rt2);
82     }
83 }
84
85 void dfs(int u) {
86     int tmp = INF; get_rt(u, -1, get_sz(u, -1), tmp, u);
87     vis[u] = true;
88     get_dep(u, -1, 0);
89     VP* rt = pit++; tr[u].push_back({rt, nullptr, 0});
90     for (E& e: G[u]) {
91         int v = e.to;
92         if (vis[v]) continue;
93         go(v, u, rt, pit++);
94         dfs(v);
95     }
96 }
97
98 bool cmp(const P& a, const P& b) { return a.w < b.w; }
99
100 LL query(VP& p, int d, int l, int r) {
101     l = lower_bound(p.begin(), p.end(), P{l, -1}, cmp) - p.begin();
102     r = upper_bound(p.begin(), p.end(), P{r, -1}, cmp) - p.begin() - 1;
103     return p[r].s - p[l - 1].s + 1LL * (r - l + 1) * d;
104 }
105
106 int main() {
107     cin >> n >> Q >> A;
108     FOR (i, 1, n + 1) scanf("%d", &w[i]);
109     FOR (_, 1, n) {
110         int u, v, d; scanf("%d%d%d", &u, &v, &d);
111         G[u].push_back({v, d}); G[v].push_back({u, d});
112     }
113     dfs(1);
114     FOR (i, 1, n + 1)
115         for (R& x: tr[i]) {
116             x.rt->push_back({w[i], x.dep});
117             if (x.rt2) x.rt2->push_back({w[i], x.dep});
118         }
119     FOR (it, pool, pit) {
120         it->push_back({-INF, 0});
121         sort(it->begin(), it->end(), cmp);
122         FOR (i, 1, it->size())
123             (*it)[i].s += (*it)[i - 1].s;
124     }
125     while (Q--) {
126         int u; LL a, b; scanf("%d%lld%lld", &u, &a, &b);
127         a = (a + ans) % A; b = (b + ans) % A;
128         int l = min(a, b), r = max(a, b);
129         ans = 0;
130         for (R& x: tr[u]) {
131             ans += query(*x.rt, x.dep, l, r);
132             if (x.rt2) ans -= query(*x.rt2, x.dep, l, r);
133         }
134         printf("%lld\n", ans);
135     }
136 }

```

树链剖分

- 初始化需要清空 clk
- 使用 `hld::predfs(1, 1); hld::dfs(1, 1);`

```
1  int fa[N], dep[N], idx[N], out[N], ridx[N];
2  namespace hld {
3      int sz[N], son[N], top[N], clk;
4      void predfs(int u, int d) {
5          dep[u] = d; sz[u] = 1;
6          int& maxs = son[u] = -1;
7          for (int& v: G[u]) {
8              if (v == fa[u]) continue;
9              fa[v] = u;
10             predfs(v, d + 1);
11             sz[u] += sz[v];
12             if (maxs == -1 || sz[v] > sz[maxs]) maxs = v;
13         }
14     }
15     void dfs(int u, int tp) {
16         top[u] = tp; idx[u] = ++clk; ridx[clk] = u;
17         if (son[u] != -1) dfs(son[u], tp);
18         for (int& v: G[u])
19             if (v != fa[u] && v != son[u]) dfs(v, v);
20         out[u] = clk;
21     }
22     template<typename T>
23     int go(int u, int v, T&& f = [] (int, int) {}) {
24         int uu = top[u], vv = top[v];
25         while (uu != vv) {
26             if (dep[uu] < dep[vv]) { swap(uu, vv); swap(u, v); }
27             f(idx[uu], idx[vv]);
28             u = fa[uu]; uu = top[u];
29         }
30         if (dep[u] < dep[v]) swap(u, v);
31         // choose one
32         // f(idx[v], idx[u]);
33         // if (u != v) f(idx[v] + 1, idx[u]);
34         return v;
35     }
36     int up(int u, int d) {
37         while (d) {
38             if (dep[u] - dep[top[u]] < d) {
39                 d -= dep[u] - dep[top[u]];
40                 u = top[u];
41             } else return ridx[idx[u] - d];
42             u = fa[u]; --d;
43         }
44         return u;
45     }
46 }
```

二分图匹配

- 最小覆盖数 = 最大匹配数
- 最大独立集 = 顶点数 - 二分图匹配数
- DAG 最小路径覆盖数 = 结点数 - 拆点后二分图最大匹配数

```
1  struct MaxMatch {
2      int n;
3      vector<int> G[maxn];
4      int vis[maxn], left[maxn], clk;
5
6      void init(int n) {
7          this->n = n;
8          FOR (i, 0, n + 1) G[i].clear();
9          memset(left, -1, sizeof left);
10         memset(vis, -1, sizeof vis);
11     }
12 }
```



```

11     }
12
13     bool dfs(int u) {
14         for (int v: G[u])
15             if (vis[v] != clk) {
16                 vis[v] = clk;
17                 if (left[v] == -1 || dfs(left[v])) {
18                     left[v] = u;
19                     return true;
20                 }
21             }
22         return false;
23     }
24
25     int match() {
26         int ret = 0;
27         for (clk = 0; clk <= n; ++clk)
28             if (dfs(clk)) ++ret;
29         return ret;
30     }
31 } MM;

```

● 二分图最大权完美匹配 KM

```

1 namespace R {
2     const int maxn = 300 + 10;
3     int n, m;
4     int left[maxn], L[maxn], R[maxn];
5     int w[maxn][maxn], slack[maxn];
6     bool visL[maxn], visR[maxn];
7
8     bool dfs(int u) {
9         visL[u] = true;
10        FOR (v, 0, m) {
11            if (visR[v]) continue;
12            int t = L[u] + R[v] - w[u][v];
13            if (t == 0) {
14                visR[v] = true;
15                if (left[v] == -1 || dfs(left[v])) {
16                    left[v] = u;
17                    return true;
18                }
19            } else slack[v] = min(slack[v], t);
20        }
21        return false;
22    }
23
24    int go() {
25        memset(left, -1, sizeof left);
26        memset(R, 0, sizeof R);
27        memset(L, 0, sizeof L);
28        FOR (i, 0, n)
29            FOR (j, 0, m)
30                L[i] = max(L[i], w[i][j]);
31
32        FOR (i, 0, n) {
33            memset(slack, 0x3f, sizeof slack);
34            while (1) {
35                memset(visL, 0, sizeof visL); memset(visR, 0, sizeof visR);
36                if (dfs(i)) break;
37                int d = 0x3f3f3f3f;
38                FOR (j, 0, m) if (!visR[j]) d = min(d, slack[j]);
39                FOR (j, 0, n) if (visL[j]) L[j] -= d;
40                FOR (j, 0, m) if (visR[j]) R[j] += d; else slack[j] -= d;
41            }
42        }
43        int ret = 0;
44        FOR (i, 0, m) if (left[i] != -1) ret += w[left[i]][i];
45        return ret;
46    }
47 }

```

虚树

```
1 void go(vector<int>& V, int& k) {
2     int u = V[k]; f[u] = 0;
3     dbg(u, k);
4     for (auto& e: G[u]) {
5         int v = e.to;
6         if (v == pa[u][0]) continue;
7         while (k + 1 < V.size()) {
8             int to = V[k + 1];
9             if (in[to] <= out[v]) {
10                go(V, ++k);
11                if (key[to]) f[u] += w[to];
12                else f[u] += min(f[to], (LL)w[to]);
13            } else break;
14        }
15    }
16    dbg(u, f[u]);
17 }
18 inline bool cmp(int a, int b) { return in[a] < in[b]; }
19 LL solve(vector<int>& V) {
20     static vector<int> a; a.clear();
21     for (int& x: V) a.push_back(x);
22     sort(a.begin(), a.end(), cmp);
23     FOR (i, 1, a.size())
24         a.push_back(lca(a[i], a[i - 1]));
25     a.push_back(1);
26     sort(a.begin(), a.end(), cmp);
27     a.erase(unique(a.begin(), a.end()), a.end());
28     dbg(a);
29     int tmp; go(a, tmp = 0);
30     return f[1];
31 }
```

欧拉路径

```
1 int S[N << 1], top;
2 Edge edges[N << 1];
3 set<int> G[N];
4
5 void DFS(int u) {
6     S[top++] = u;
7     for (int eid: G[u]) {
8         int v = edges[eid].get_other(u);
9         G[u].erase(eid);
10        G[v].erase(eid);
11        DFS(v);
12        return;
13    }
14 }
15
16 void fleury(int start) {
17     int u = start;
18     top = 0; path.clear();
19     S[top++] = u;
20     while (top) {
21         u = S[--top];
22         if (!G[u].empty())
23             DFS(u);
24         else path.push_back(u);
25     }
26 }
```

强连通分量与 2-SAT

```
1 int n, m;
2 vector<int> G[N], rG[N], vs;
```

```

3  int used[N], cmp[N];
4
5  void add_edge(int from, int to) {
6      G[from].push_back(to);
7      rG[to].push_back(from);
8  }
9
10 void dfs(int v) {
11     used[v] = true;
12     for (int u: G[v]) {
13         if (!used[u])
14             dfs(u);
15     }
16     vs.push_back(v);
17 }
18
19 void rdfs(int v, int k) {
20     used[v] = true;
21     cmp[v] = k;
22     for (int u: rG[v])
23         if (!used[u])
24             rdfs(u, k);
25 }
26
27 int scc() {
28     memset(used, 0, sizeof(used));
29     vs.clear();
30     for (int v = 0; v < n; ++v)
31         if (!used[v]) dfs(v);
32     memset(used, 0, sizeof(used));
33     int k = 0;
34     for (int i = (int) vs.size() - 1; i >= 0; --i)
35         if (!used[vs[i]]) rdfs(vs[i], k++);
36     return k;
37 }
38
39 int main() {
40     cin >> n >> m;
41     n *= 2;
42     for (int i = 0; i < m; ++i) {
43         int a, b; cin >> a >> b;
44         add_edge(a - 1, (b - 1) ^ 1);
45         add_edge(b - 1, (a - 1) ^ 1);
46     }
47     scc();
48     for (int i = 0; i < n; i += 2) {
49         if (cmp[i] == cmp[i + 1]) {
50             puts("NIE");
51             return 0;
52         }
53     }
54     for (int i = 0; i < n; i += 2) {
55         if (cmp[i] > cmp[i + 1]) printf("%d\n", i + 1);
56         else printf("%d\n", i + 2);
57     }
58 }

```

拓补排序

```

1  vector<int> toporder(int n) {
2      vector<int> orders;
3      queue<int> q;
4      for (int i = 0; i < n; i++)
5          if (!deg[i]) {
6              q.push(i);
7              orders.push_back(i);
8          }
9      while (!q.empty()) {
10         int u = q.front(); q.pop();

```

```

11         for (int v: G[u])
12             if (!--deg[v]) {
13                 q.push(v);
14                 orders.push_back(v);
15             }
16     }
17     return orders;
18 }

```

一般图匹配

带花树。复杂度 $O(n^3)$ 。

```

1  int n;
2  vector<int> G[N];
3  int fa[N], mt[N], pre[N], mk[N];
4  int lca_clk, lca_mk[N];
5  pair<int, int> ce[N];
6
7  void connect(int u, int v) {
8      mt[u] = v;
9      mt[v] = u;
10 }
11 int find(int x) { return x == fa[x] ? x : fa[x] = find(fa[x]); }
12
13 void flip(int s, int u) {
14     if (s == u) return;
15     if (mk[u] == 2) {
16         int v1 = ce[u].first, v2 = ce[u].second;
17         flip(mt[u], v1);
18         flip(s, v2);
19         connect(v1, v2);
20     } else {
21         flip(s, pre[mt[u]]);
22         connect(pre[mt[u]], mt[u]);
23     }
24 }
25
26 int get_lca(int u, int v) {
27     lca_clk++;
28     for (u = find(u), v = find(v); ; u = find(pre[u]), v = find(pre[v])) {
29         if (u && lca_mk[u] == lca_clk) return u;
30         lca_mk[u] = lca_clk;
31         if (v && lca_mk[v] == lca_clk) return v;
32         lca_mk[v] = lca_clk;
33     }
34 }
35
36 void access(int u, int p, const pair<int, int>& c, vector<int>& q) {
37     for (u = find(u); u != p; u = find(pre[u])) {
38         if (mk[u] == 2) {
39             ce[u] = c;
40             q.push_back(u);
41         }
42         fa[find(u)] = find(p);
43     }
44 }
45
46 bool aug(int s) {
47     fill(mk, mk + n + 1, 0);
48     fill(pre, pre + n + 1, 0);
49     iota(fa, fa + n + 1, 0);
50     vector<int> q = {s};
51     mk[s] = 1;
52     int t = 0;
53     for (int t = 0; t < (int) q.size(); ++t) {
54         // q size can be changed
55         int u = q[t];
56         for (int &v: G[u]) {

```

```

57         if (find(v) == find(u)) continue;
58         if (!mk[v] && !mt[v]) {
59             flip(s, u);
60             connect(u, v);
61             return true;
62         } else if (!mk[v]) {
63             int w = mt[v];
64             mk[v] = 2; mk[w] = 1;
65             pre[w] = v; pre[v] = u;
66             q.push_back(w);
67         } else if (mk[find(v)] == 1) {
68             int p = get_lca(u, v);
69             access(u, p, {u, v}, q);
70             access(v, p, {v, u}, q);
71         }
72     }
73 }
74 return false;
75 }
76
77 int match() {
78     fill(mt + 1, mt + n + 1, 0);
79     lca_clk = 0;
80     int ans = 0;
81     FOR (i, 1, n + 1)
82         if (!mt[i]) ans += aug(i);
83     return ans;
84 }
85
86 int main() {
87     int m; cin >> n >> m;
88     while (m--) {
89         int u, v; scanf("%d%d", &u, &v);
90         G[u].push_back(v); G[v].push_back(u);
91     }
92     printf("%d\n", match());
93     FOR (i, 1, n + 1) printf("%d%c", mt[i], i == _i - 1 ? '\n' : ' ');
94     return 0;
95 }

```

Tarjan

割点

- 判断割点
- 注意原图可能不连通

```

1  int dfn[N], low[N], clk;
2  void init() { clk = 0; memset(dfn, 0, sizeof dfn); }
3  void tarjan(int u, int fa) {
4      low[u] = dfn[u] = ++clk;
5      int cc = fa != -1;
6      for (int& v: G[u]) {
7          if (v == fa) continue;
8          if (!dfn[v]) {
9              tarjan(v, u);
10             low[u] = min(low[u], low[v]);
11             cc += low[v] >= dfn[u];
12         } else low[u] = min(low[u], dfn[v]);
13     }
14     if (cc > 1) // ...
15 }

```

桥

- 注意原图不连通和重边

```

1  int dfn[N], low[N], clk;
2  void init() { memset(dfn, 0, sizeof dfn); clk = 0; }
3  void tarjan(int u, int fa) {
4      low[u] = dfn[u] = ++clk;
5      int _fst = 0;
6      for (E& e: G[u]) {
7          int v = e.to; if (v == fa && ++_fst == 1) continue;
8          if (!dfn[v]) {
9              tarjan(v, u);
10             if (low[v] > dfn[u]) // ...
11                 low[u] = min(low[u], low[v]);
12             } else low[u] = min(low[u], dfn[v]);
13     }
14 }

```

强连通分量缩点

```

1  int low[N], dfn[N], clk, B, belong[N];
2  vector<int> Bc[N];
3  void init() { B = clk = 0; memset(dfn, 0, sizeof dfn); }
4  void tarjan(int u) {
5      static int st[N], p;
6      static bool in[N];
7      dfn[u] = low[u] = ++clk;
8      st[p++] = u; in[u] = true;
9      for (int& v: G[u]) {
10         if (!dfn[v]) {
11             tarjan(v);
12             low[u] = min(low[u], low[v]);
13         } else if (in[v]) low[u] = min(low[u], dfn[v]);
14     }
15     if (dfn[u] == low[u]) {
16         while (1) {
17             int x = st[--p]; in[x] = false;
18             belong[x] = B; Bc[B].push_back(x);
19             if (x == u) break;
20         }
21         ++B;
22     }
23 }

```

- 点双连通分量
- 数组开两倍大，按照边数而不是点数开

点双连通分量 / 广义圆方树

- 数组开两倍
- 一条边也被计入点双了（适合拿来建圆方树），可以用点数 \leq 边数过滤

```

1  struct E { int to, nxt; } e[N];
2  int hd[N], ecnt;
3  void addedge(int u, int v) {
4      e[ecnt] = {v, hd[u]};
5      hd[u] = ecnt++;
6  }
7  int low[N], dfn[N], clk, B, bno[N];
8  vector<int> bc[N], be[N];
9  bool vise[N];
10 void init() {
11     memset(vise, 0, sizeof vise);
12     memset(hd, -1, sizeof hd);
13     memset(dfn, 0, sizeof dfn);
14     memset(bno, -1, sizeof bno);
15     B = clk = 0;
16 }
17
18 void tarjan(int u, int feid) {
19     static int st[N], p;

```

```

20 static auto add = [&](int x) {
21     if (bno[x] != B) { bno[x] = B; bc[B].push_back(x); }
22 };
23 low[u] = dfn[u] = ++clk;
24 for (int i = hd[u]; ~i; i = e[i].nxt) {
25     if ((feid ^ i) == 1) continue;
26     if (!vise[i]) { st[p++] = i; vise[i] = vise[i ^ 1] = true; }
27     int v = e[i].to;
28     if (!dfn[v]) {
29         tarjan(v, i);
30         low[u] = min(low[u], low[v]);
31         if (low[v] >= dfn[u]) {
32             bc[B].clear(); be[B].clear();
33             while (1) {
34                 int eid = st[--p];
35                 add(e[eid].to); add(e[eid ^ 1].to);
36                 be[B].push_back(eid);
37                 if ((eid ^ i) <= 1) break;
38             }
39             ++B;
40         }
41     } else low[u] = min(low[u], dfn[v]);
42 }
43 }

```

圆方树

- 从仙人掌建圆方树
- N 至少边数 $\times 2$

```

1 vector<int> G[N];
2 int nn;
3
4 struct E { int to, nxt; };
5 namespace C {
6     E e[N * 2];
7     int hd[N], ecnt;
8     void addedge(int u, int v) {
9         e[ecnt] = {v, hd[u]};
10        hd[u] = ecnt++;
11    }
12    int idx[N], clk, fa[N];
13    bool ring[N];
14    void init() { ecnt = 0; memset(hd, -1, sizeof hd); clk = 0; }
15    void dfs(int u, int feid) {
16        idx[u] = ++clk;
17        for (int i = hd[u]; ~i; i = e[i].nxt) {
18            if ((i ^ feid) == 1) continue;
19            int v = e[i].to;
20            if (!idx[v]) {
21                fa[v] = u; ring[u] = false;
22                dfs(v, i);
23                if (!ring[u]) { G[u].push_back(v); G[v].push_back(u); }
24            } else if (idx[v] < idx[u]) {
25                ++nn;
26                G[nn].push_back(v); G[v].push_back(nn); // 强行把环的根放在最前面
27                for (int x = u; x != v; x = fa[x]) {
28                    ring[x] = true;
29                    G[nn].push_back(x); G[x].push_back(nn);
30                }
31                ring[v] = true;
32            }
33        }
34    }
35 }

```

计算几何

圆的反演

```
1 typedef double LD;
2 const LD PI = 3.14159265358979323846;
3 const LD eps = 1E-10;
4 const LD R2 = 1.0;
5 int sgn(LD x) { return fabs(x) < eps ? 0 : (x > 0 ? 1 : -1); }
6 struct P {
7     LD x, y;
8     P(LD x = 0, LD y = 0): x(x), y(y) {}
9     P operator * (LD k) { return P(x * k, y * k); }
10    P operator / (LD k) { return P(x / k, y / k); }
11    string prt() const {
12        char s[100];
13        sprintf(s, "(%.2f, %.2f)", x, y);
14        return string(s);
15    }
16 };
17 typedef P V;
18 P operator - (const P& a, const P& b) { return P(a.x - b.x, a.y - b.y); }
19 P operator + (const P& a, const P& b) { return P(a.x + b.x, a.y + b.y); }
20 struct C {
21     P p;
22     LD r;
23     C(LD x = 0, LD y = 0, LD r = 0): p(x, y), r(r) {}
24 };
25 LD dist(V v) { return sqrt(v.x * v.x + v.y * v.y); }
26
27 C inv(C c, const P& o) {
28     LD d = dist(c.p - o);
29     assert(sgn(d) != 0);
30     LD a = 1 / (d - c.r);
31     LD b = 1 / (d + c.r);
32     c.p = (a - b) / 2 * R2;
33     c.p = o + (c.p - o) * ((a + b) * R2 / 2 / d);
34     return c;
35 }
```

二维

- nxt 宏要求多边形变量名为 s
- L 可隐式转换为 V(P)
- 可以自定义结构体 PP, 可隐式转换为 P

```
1 #define y1 yy1
2 #define nxt(i) ((i + 1) % s.size())
3 typedef double LD;
4 const LD PI = 3.14159265358979323846;
5 const LD eps = 1E-10;
6 int sgn(LD x) { return fabs(x) < eps ? 0 : (x > 0 ? 1 : -1); }
7 struct L;
8 struct P;
9 //struct PP;
10 typedef P V;
11 struct P {
12     LD x, y;
13     explicit P(LD x = 0, LD y = 0): x(x), y(y) {}
14     P(const L& l);
15     // P(const PP& pp);
16 };
17 struct L {
18     P s, t;
19     L() {}
20     L(P s, P t): s(s), t(t) {}
21 };
```



```

22
23 P operator + (const P& a, const P& b) { return P(a.x + b.x, a.y + b.y); }
24 P operator - (const P& a, const P& b) { return P(a.x - b.x, a.y - b.y); }
25 P operator * (const P& a, LD k) { return P(a.x * k, a.y * k); }
26 P operator / (const P& a, LD k) { return P(a.x / k, a.y / k); }
27 inline int operator < (const P& a, const P& b) {
28     return sgn(a.x - b.x) < 0 || (sgn(a.x - b.x) == 0 && sgn(a.y - b.y) < 0);
29 }
30 bool operator == (const P& a, const P& b) { return !sgn(a.x - b.x) && !sgn(a.y - b.y); }
31 P::P(const L& l) { *this = l.t - l.s; }
32 ostream &operator << (ostream &os, const P &p) {
33     return (os << "(" << p.x << ", " << p.y << ")");
34 }
35 istream &operator >> (istream &is, P &p) {
36     return (is >> p.x >> p.y);
37 }
38
39 // -----
40
41 //struct PP {
42 //    P p;
43 //    LD v, l;
44 //};
45 //P::P(const PP& pp) { *this = pp.p; }
46 typedef P PP;
47
48 typedef vector<PP> S;
49
50 // -----
51 LD dist(const P& p) { return sqrt(p.x * p.x + p.y * p.y); }
52 LD dot(const V& a, const V& b) { return a.x * b.x + a.y * b.y; }
53 LD det(const V& a, const V& b) { return a.x * b.y - a.y * b.x; }
54 LD cross(const P& s, const P& t, const P& o = P()) { return det(s - o, t - o); }
55
56 // 如需支持 unique, 需要加 eps
57 bool cmp_xy(const P& a, const P& b) { return a.x < b.x || a.x == b.x && a.y < b.y; }
58
59 // 象限
60 int quad(P p) {
61     int x = sgn(p.x), y = sgn(p.y);
62     if (x > 0 && y >= 0) return 1;
63     if (x <= 0 && y > 0) return 2;
64     if (x < 0 && y <= 0) return 3;
65     if (x >= 0 && y < 0) return 4;
66     assert(0);
67 }
68
69 // 仅适用于参照点在所有点一侧的情况
70 struct cmp_angle {
71     P p;
72     bool operator () (const P& a, const P& b) {
73         // int qa = quad(a), qb = quad(b);
74         // if (qa != qb) return qa < qb;
75         int d = sgn(cross(a, b, p));
76         if (d) return d > 0;
77         return dist(a - p) < dist(b - p);
78     }
79 };
80
81 // -----线-----
82
83 // 是否平行
84 bool parallel(const L& a, const L& b) {
85     return !sgn(det(a, b));
86 }
87 // 直线是否相等
88 bool l_eq(const L& a, const L& b) {
89     return parallel(a, b) && parallel(L(a.s, b.t), L(b.s, a.t));
90 }
91 // 逆时针旋转 r 弧度
92 P rotation(const P& p, const LD& r) { return P(p.x * cos(r) - p.y * sin(r), p.x * sin(r) + p.y * cos(r)); }

```

```

93 P RotateCCW90(const P& p) { return P(-p.y, p.x); }
94 P RotateCW90(const P& p) { return P(p.y, -p.x); }
95 // 单位法向量
96 V normal(const V& v) { return V(-v.y, v.x) / dist(v); }
97
98
99 // -----点和线-----
100
101 // 点在线段上 <= 0 包含端点 < 0 则不包含
102 bool p_on_seg(const P& p, const L& seg) {
103     P a = seg.s, b = seg.t;
104     return !sgn(det(p - a, b - a)) && sgn(dot(p - a, p - b)) <= 0;
105 }
106 // 点到直线距离
107 LD dist_to_line(const P& p, const L& l) {
108     return fabs(cross(l.s, l.t, p)) / dist(l);
109 }
110 // 点到线段距离
111 LD dist_to_seg(const P& p, const L& l) {
112     if (l.s == l.t) return dist(p - l);
113     V vs = p - l.s, vt = p - l.t;
114     if (sgn(dot(l, vs)) < 0) return dist(vs);
115     else if (sgn(dot(l, vt)) > 0) return dist(vt);
116     else return dist_to_line(p, l);
117 }
118
119
120 // -----线和线-----
121
122 // 求直线交 需要事先保证有界
123 P l_intersection(const L& a, const L& b) {
124     LD s1 = det(a, b.s - a.s), s2 = det(a, b.t - a.s);
125     return (b.s * s2 - b.t * s1) / (s2 - s1);
126 }
127 // 向量夹角的弧度
128 LD angle(const V& a, const V& b) {
129     LD r = asin(fabs(det(a, b)) / dist(a) / dist(b));
130     if (sgn(dot(a, b)) < 0) r = PI - r;
131     return r;
132 }
133 // 线段和直线是否有交 1 = 规范, 2 = 不规范
134 int s_l_cross(const L& seg, const L& line) {
135     int d1 = sgn(cross(line.s, line.t, seg.s));
136     int d2 = sgn(cross(line.s, line.t, seg.t));
137     if ((d1 ^ d2) == -2) return 1; // proper
138     if (d1 == 0 || d2 == 0) return 2;
139     return 0;
140 }
141 // 线段的交 1 = 规范, 2 = 不规范
142 int s_cross(const L& a, const L& b, P& p) {
143     int d1 = sgn(cross(a.t, b.s, a.s)), d2 = sgn(cross(a.t, b.t, a.s));
144     int d3 = sgn(cross(b.t, a.s, b.s)), d4 = sgn(cross(b.t, a.t, b.s));
145     if ((d1 ^ d2) == -2 && (d3 ^ d4) == -2) { p = l_intersection(a, b); return 1; }
146     if (!d1 && p_on_seg(b.s, a)) { p = b.s; return 2; }
147     if (!d2 && p_on_seg(b.t, a)) { p = b.t; return 2; }
148     if (!d3 && p_on_seg(a.s, b)) { p = a.s; return 2; }
149     if (!d4 && p_on_seg(a.t, b)) { p = a.t; return 2; }
150     return 0;
151 }
152
153
154 // -----多边形-----
155
156 // 点是否在多边形中 0 = 在外部 1 = 在内部 -1 = 在边界上
157 int inside(const S& s, const P& p) {
158     int cnt = 0;
159     FOR (i, 0, s.size()) {
160         P a = s[i], b = s[nxt(i)];
161         if (p_on_seg(p, L(a, b))) return -1;
162         if (sgn(a.y - b.y) <= 0) swap(a, b);
163         if (sgn(p.y - a.y) > 0) continue;

```

```

164         if (sgn(p.y - b.y) <= 0) continue;
165         cnt += sgn(cross(b, a, p)) > 0;
166     }
167     return bool(cnt & 1);
168 }
169 // 多边形面积, 有向面积可能为负
170 LD polygon_area(const S& s) {
171     LD ret = 0;
172     FOR (i, 1, (LL)s.size() - 1)
173         ret += cross(s[i], s[i + 1], s[0]);
174     return ret / 2;
175 }
176 // 构建凸包 点不可以重复 < 0 边上可以有点, <= 0 则不能
177 // 会改变输入点的顺序
178 const int MAX_N = 1000;
179 S convex_hull(S& s) {
180     // assert(s.size() >= 3);
181     sort(s.begin(), s.end(), cmp_xy);
182     S ret(MAX_N * 2);
183     int sz = 0;
184     FOR (i, 0, s.size()) {
185         while (sz > 1 && sgn(cross(ret[sz - 1], s[i], ret[sz - 2])) < 0) --sz;
186         ret[sz++] = s[i];
187     }
188     int k = sz;
189     FOR (i, (LL)s.size() - 2, -1) {
190         while (sz > k && sgn(cross(ret[sz - 1], s[i], ret[sz - 2])) < 0) --sz;
191         ret[sz++] = s[i];
192     }
193     ret.resize(sz - (s.size() > 1));
194     return ret;
195 }
196
197 P ComputeCentroid(const vector<P> &p) {
198     P c(0, 0);
199     LD scale = 6.0 * polygon_area(p);
200     for (unsigned i = 0; i < p.size(); i++) {
201         unsigned j = (i + 1) % p.size();
202         c = c + (p[i] + p[j]) * (p[i].x * p[j].y - p[j].x * p[i].y);
203     }
204     return c / scale;
205 }
206
207 // -----圖-----
208
209 P ComputeCircleCenter(P a, P b, P c) {
210     b = (a + b) / 2;
211     c = (a + c) / 2;
212     return l_intersection({b, b + RotateCW90(a - b)}, {c, c + RotateCW90(a - c)});
213 }
214 vector<P> CircleLineIntersection(P a, P b, P c, LD r) {
215     vector<P> ret;
216     b = b - a;
217     a = a - c;
218     LD A = dot(b, b), B = dot(a, b), C = dot(a, a) - r * r;
219     LD D = B * B - A * C;
220     if (sgn(D) < 0) return ret;
221     ret.push_back(c + a + b * (-B + sqrt(D + eps)) / A);
222     if (sgn(D) > 0) ret.push_back(c + a + b * (-B - sqrt(D)) / A);
223     return ret;
224 }
225 vector<P> CircleCircleIntersection(P a, P b, LD r, LD R) {
226     vector<P> ret;
227     LD d = dist(a - b);
228     if (sgn(d) == 0 || sgn(d - (r + R)) > 0 || sgn(d + min(r, R) - max(r, R)) < 0) return ret;
229     LD x = (d * d - R * R + r * r) / (2 * d);
230     LD y = sqrt(r * r - x * x);
231     P v = (b - a) / d;
232     ret.push_back(a + v * x + RotateCCW90(v) * y);
233     if (sgn(y) > 0) ret.push_back(a + v * x - RotateCCW90(v) * y);
234     return ret;

```

```

235 }
236
237 // -----模板结束-----

```

旋转卡壳

```

1  LD rotatingCalipers(S& qs) {
2      int n = qs.size();
3      if (n == 2)
4          return dist(qs[0] - qs[1]);
5      int i = 0, j = 0;
6      FOR (k, 0, n) {
7          if (!(qs[i] < qs[k])) i = k;
8          if (qs[j] < qs[k]) j = k;
9      }
10     LD res = 0;
11     int si = i, sj = j;
12     while (i != sj || j != si) {
13         res = max(res, dist(qs[i] - qs[j]));
14         if (sgn(cross(qs[(i+1)%n] - qs[i], qs[(j+1)%n] - qs[j])) < 0)
15             i = (i + 1) % n;
16         else j = (j + 1) % n;
17     }
18     return res;
19 }
20
21 int main() {
22     int n;
23     while (cin >> n) {
24         S v(n);
25         FOR (i, 0, n) cin >> v[i].x >> v[i].y;
26         convex_hull(v);
27         printf("%.0f\n", rotatingCalipers(v));
28     }
29 }

```

没有测试过的

```

1  int relation(Point p, Circle a) { //点和圆的关系
2      //0: 圆外 1: 圆上 2: 圆内
3      double d = dis(p, a.p);
4      if (dcmp(d - a.r) == 0) return 1;
5      return (dcmp(d - a.r) < 0 ? 2 : 0);
6  }
7
8  int relation(Line a, Circle b) { //直线和圆的关系
9      //0: 相离 1: 相切 2: 相交
10     double p = point_to_line(b.p, a);
11     if (dcmp(p - b.r) == 0) return 1;
12     return (dcmp(p - b.r) < 0 ? 2 : 0);
13 }
14
15 int relation(Circle a, Circle v) { //两圆的位置关系
16     //1: 内含 2: 内切 3: 相交 4: 外切 5: 相离
17     double d = dis(a.p, v.p);
18     if (dcmp(d - a.r - v.r) > 0) return 5;
19     if (dcmp(d - a.r - v.r) == 0) return 4;
20     double l = fabs(a.r - v.r);
21     if (dcmp(d - a.r - v.r) < 0 && dcmp(d - l) > 0) return 3;
22     if (dcmp(d - l) == 0) return 2;
23     if (dcmp(d - l) < 0) return 1;
24     assert (0);
25 }
26
27 double circle_triangle_area(Point a, Point b, Circle c) { //圆心三角形的面积
28     //a.output (), b.output (), c.output ();
29     Point p = c.p;

```

```

30     double r = c.r; //cout << cross (p-a, p-b) << endl;
31     if (dcmp(cross(p - a, p - b)) == 0) return 0;
32     Point q[5];
33     int len = 0;
34     q[len++] = a;
35     Line l(a, b);
36     Point p1, p2;
37     if (line_circle_intersection(l, c, q[1], q[2]) == 2) {
38         if (dcmp(dot(a - q[1], b - q[1])) < 0) q[len++] = q[1];
39         if (dcmp(dot(a - q[2], b - q[2])) < 0) q[len++] = q[2];
40     }
41     q[len++] = b;
42     if (len == 4 && dcmp(dot(q[0] - q[1], q[2] - q[1])) > 0)
43         swap(q[1], q[2]);
44     double res = 0;
45     for (int i = 0; i < len - 1; i++) {
46         if (relation(q[i], c) == 0 || relation(q[i + 1], c) == 0) {
47             double arg = rad(q[i] - p, q[i + 1] - p);
48             res += r * r * arg / 2.0;
49         } else {
50             res += fabs(cross(q[i] - p, q[i + 1] - p)) / 2;
51         }
52     }
53     return res;
54 }

```

半平面交

```

1  struct Line {
2      PT p, v;
3      double ang;
4
5      Line() {}
6      Line(PT from, PT to) : p(from), v(to - from) { ang = atan2(v.y, v.x); }
7      friend bool operator<(Line a, Line b) {
8          return a.ang < b.ang;
9      }
10 };
11
12 bool OnLeft(Line L, PT p) {
13     return dcmp(cross(L.v, p - L.p)) >= 0;
14 }
15
16 PT GetIntersection(Line a, Line b) {
17     PT u = a.p - b.p;
18     ld t = cross(b.v, u) / cross(a.v, b.v);
19     return a.p + a.v * t;
20 }
21
22 vector<PT> HalfplaneIntersection(vector<Line>& L) {
23     int n = L.size();
24     sort(L.begin(), L.end());
25
26     int first, last;
27     vector<PT> p(n);
28     vector<Line> q(n);
29     q[first = last = 0] = L[0];
30     for (int i = 1; i < n; i++) {
31         while (first < last && !OnLeft(L[i], p[last - 1])) last--;
32         while (first < last && !OnLeft(L[i], p[first])) first++;
33         q[++last] = L[i];
34         if (dcmp(cross(q[last].v, q[last - 1].v)) == 0) {
35             last--;
36             if (OnLeft(q[last], L[i].p)) q[last] = L[i];
37         }
38         if (first < last) p[last - 1] = GetIntersection(q[last - 1], q[last]);
39     }
40     while (first < last && !OnLeft(q[first], p[last - 1])) last--;
41     if (last - first <= 1) return vector<PT>();

```

```

42     p[last] = GetIntersection(q[last], q[first]);
43
44     return vector<PT>(p.begin() + first, p.begin() + last + 1);
45 }
46
47 vector<PT> convexIntersection(const vector<PT> &v1, const vector<PT> &v2) {
48     vector<Line> h;
49     int n = v1.size(), m = v2.size();
50     for (int i = 0; i < n; ++i)
51         h.push_back(Line(v1[i], v1[(i + 1) % n]));
52     for (int i = 0; i < m; ++i)
53         h.push_back(Line(v2[i], v2[(i + 1) % m]));
54     return HalfplaneIntersection(h);
55 }
56
57 ld ComputeSignedArea(const vector<PT> &p) {
58     ld area = 0;
59     for (unsigned i = 0; i < p.size(); i++) {
60         unsigned j = (i + 1) % p.size();
61         area += p[i].x * p[j].y - p[j].x * p[i].y;
62     }
63     return area / 2.0;
64 }
65
66 ld ComputeArea(const vector<PT> &p) {
67     return fabs(ComputeSignedArea(p));
68 }

```

三维计算几何

```

1  struct P;
2  struct L;
3  typedef P V;
4
5  struct P {
6      LD x, y, z;
7      explicit P(LD x = 0, LD y = 0, LD z = 0): x(x), y(y), z(z) {}
8      explicit P(const L& l);
9  };
10
11 struct L {
12     P s, t;
13     L() {}
14     L(P s, P t): s(s), t(t) {}
15 };
16
17 P operator + (const P& a, const P& b) { return P(a.x + b.x, a.y + b.y, a.z + b.z); }
18 P operator - (const P& a, const P& b) { return P(a.x - b.x, a.y - b.y, a.z - b.z); }
19 P operator * (const P& a, LD k) { return P(a.x * k, a.y * k, a.z * k); }
20 P operator / (const P& a, LD k) { return P(a.x / k, a.y / k, a.z / k); }
21 inline int operator < (const P& a, const P& b) {
22     return sgn(a.x - b.x) < 0 || (sgn(a.x - b.x) == 0 && (sgn(a.y - b.y) < 0 ||
23         (sgn(a.y - b.y) == 0 && sgn(a.z - b.z) < 0)));
24 }
25 bool operator == (const P& a, const P& b) { return !sgn(a.x - b.x) && !sgn(a.y - b.y) && !sgn(a.z - b.z); }
26 P::P(const L& l) { *this = l.t - l.s; }
27 ostream &operator << (ostream &os, const P &p) {
28     return (os << "(" << p.x << "," << p.y << "," << p.z << ")");
29 }
30 istream &operator >> (istream &is, P &p) {
31     return (is >> p.x >> p.y >> p.z);
32 }
33
34 // -----
35 LD dist2(const P& p) { return p.x * p.x + p.y * p.y + p.z * p.z; }
36 LD dist(const P& p) { return sqrt(dist2(p)); }
37 LD dot(const V& a, const V& b) { return a.x * b.x + a.y * b.y + a.z * b.z; }
38 P cross(const P& v, const P& w) {
39     return P(v.y * w.z - v.z * w.y, v.z * w.x - v.x * w.z, v.x * w.y - v.y * w.x);

```

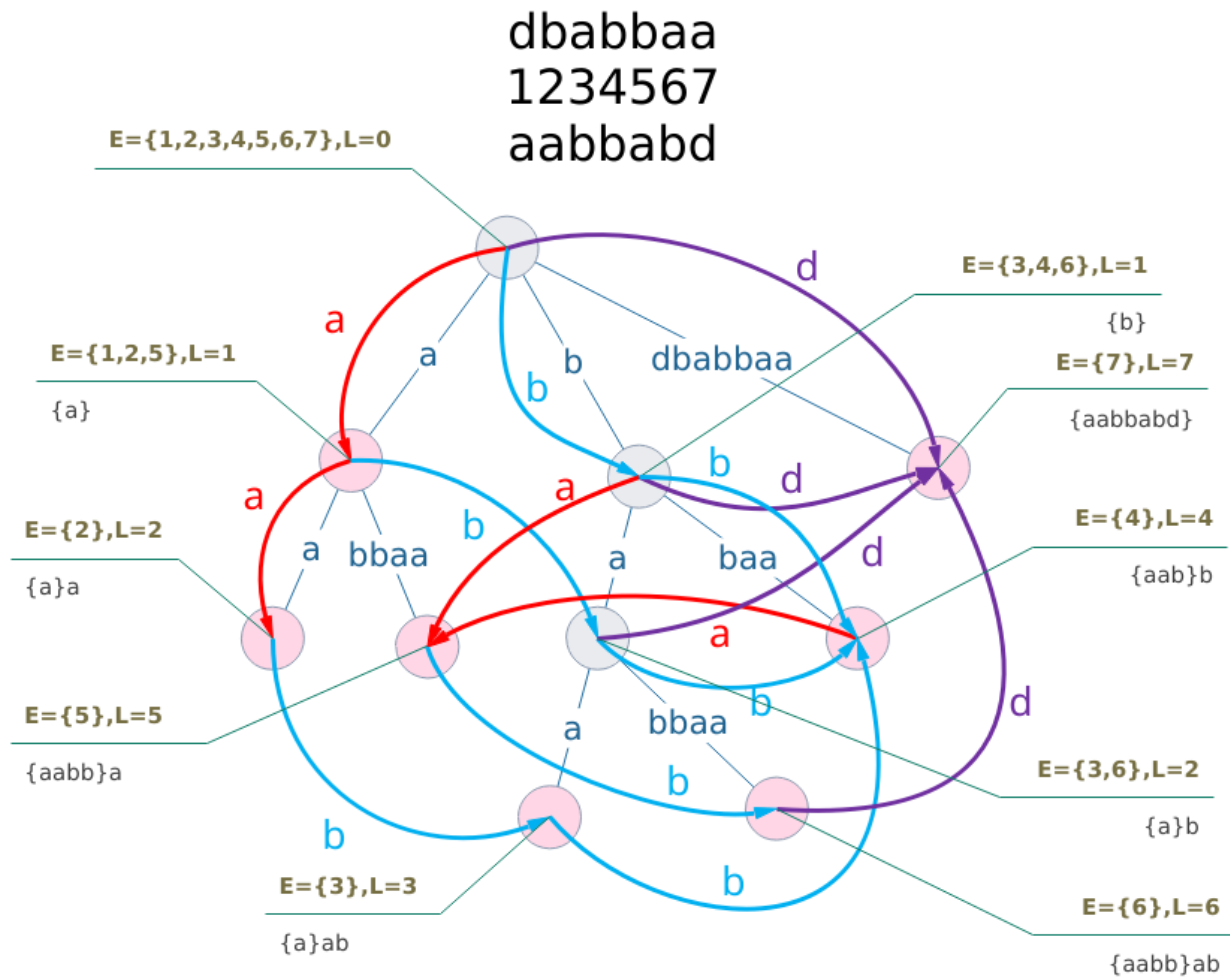
```

40 }
41
42 // -----线-----
43
44 // 逆时针旋转 r 弧度
45 // axis = 0 绕 x 轴
46 // axis = 1 绕 y 轴
47 // axis = 2 绕 z 轴
48 P rotation(const P& p, const LD& r, int axis = 0) {
49     if (axis == 0)
50         return P(p.x, p.y * cos(r) - p.z * sin(r), p.y * sin(r) + p.z * cos(r));
51     else if (axis == 1)
52         return P(p.z * cos(r) - p.x * sin(r), p.y, p.z * sin(r) + p.x * cos(r));
53     else if (axis == 2)
54         return P(p.x * cos(r) - p.y * sin(r), p.x * sin(r) + p.y * cos(r), p.z);
55 }
56 // n 是单位向量 表示旋转轴
57 // 模板是顺时针的
58 P rotation(const P& p, const LD& r, const P& n) {
59     LD c = cos(r), s = sin(r), x = n.x, y = n.y, z = n.z;
60     // dbg(c, s);
61     return P((x * x * (1 - c) + c) * p.x + (x * y * (1 - c) + z * s) * p.y + (x * z * (1 - c) - y * s) * p.z,
62             (x * y * (1 - c) - z * s) * p.x + (y * y * (1 - c) + c) * p.y + (y * z * (1 - c) + x * s) * p.z,
63             (x * z * (1 - c) + y * s) * p.x + (y * z * (1 - c) - x * s) * p.y + (z * z * (1 - c) + c) * p.z);
64 }
65
66 // -----点和线-----
67
68 // 点在线段上 <= 0 包含端点 < 0 则不包含
69 bool p_on_seg(const P& p, const L& seg) {
70     P a = seg.s, b = seg.t;
71     return !sgn(dist2(cross(p - a, b - a))) && sgn(dot(p - a, p - b)) <= 0;
72 }
73 // 点到直线距离
74 LD dist_to_line(const P& p, const L& l) {
75     return dist(cross(l.s - p, l.t - p)) / dist(l);
76 }
77 // 点到线段距离
78 LD dist_to_seg(const P& p, const L& l) {
79     if (l.s == l.t) return dist(p - l.s);
80     V vs = p - l.s, vt = p - l.t;
81     if (sgn(dot(l, vs)) < 0) return dist(vs);
82     else if (sgn(dot(l, vt)) > 0) return dist(vt);
83     else return dist_to_line(p, l);
84 }

```

字符串

后缀自动机



- 广义后缀自动机如果直接使用以下代码的话会产生一些冗余状态（置 last 为 1），所以要用拓扑排序。用 len 基数排序不能。
- 字符集大的话要使用 map。
- 树上 dp 时注意边界（root 和 null）。
- rsort 需要初始化

```
1 namespace sam {
2     const int M = N << 1;
3     int t[M][26], len[M] = {-1}, fa[M], sz = 2, last = 1;
4     void init() { sz = 2; last = 1; memset(t, 0, (sz + 10) * sizeof t[0]); }
5     void ins(int ch) {
6         int p = last, np = last = sz++;
7         len[np] = len[p] + 1;
8         for (; p && !t[p][ch]; p = fa[p]) t[p][ch] = np;
9         if (!p) { fa[np] = 1; return; }
10        int q = t[p][ch];
11        if (len[p] + 1 == len[q]) fa[np] = q;
12        else {
13            int nq = sz++; len[nq] = len[p] + 1;
14            memcpy(t[nq], t[q], sizeof t[0]);
15            fa[nq] = fa[q];
16            fa[np] = fa[q] = nq;
17            for (; t[p][ch] == q; p = fa[p]) t[p][ch] = nq;
18        }
19    }
```



```

20
21     int c[M] = {1}, a[M];
22     void rsort() {
23         FOR (i, 1, sz) c[i] = 0;
24         FOR (i, 1, sz) c[len[i]]++;
25         FOR (i, 1, sz) c[i] += c[i - 1];
26         FOR (i, 1, sz) a[--c[len[i]]] = i;
27     }
28 }

```

- 真·广义后缀自动机

```

1     int t[M][26], len[M] = {-1}, fa[M], sz = 2, last = 1;
2     LL cnt[M][2];
3     void ins(int ch, int id) {
4         int p = last, np = 0, nq = 0, q = -1;
5         if (!t[p][ch]) {
6             np = sz++;
7             len[np] = len[p] + 1;
8             for (; p && !t[p][ch]; p = fa[p]) t[p][ch] = np;
9         }
10        if (!p) fa[np] = 1;
11        else {
12            q = t[p][ch];
13            if (len[p] + 1 == len[q]) fa[np] = q;
14            else {
15                nq = sz++; len[nq] = len[p] + 1;
16                memcpy(t[nq], t[q], sizeof t[0]);
17                fa[nq] = fa[q];
18                fa[np] = fa[q] = nq;
19                for (; t[p][ch] == q; p = fa[p]) t[p][ch] = nq;
20            }
21        }
22        last = np ? np : nq ? nq : q;
23        cnt[last][id] = 1;
24    }

```

- 按字典序建立后缀树注意逆序插入

```

1     void ins(int ch, int pp) {
2         int p = last, np = last = sz++;
3         len[np] = len[p] + 1; one[np] = pos[np] = pp;
4         for (; p && !t[p][ch]; p = fa[p]) t[p][ch] = np;
5         if (!p) { fa[np] = 1; return; }
6         int q = t[p][ch];
7         if (len[q] == len[p] + 1) fa[np] = q;
8         else {
9             int nq = sz++; len[nq] = len[p] + 1; one[nq] = one[q];
10            memcpy(t[nq], t[q], sizeof t[0]);
11            fa[nq] = fa[q];
12            fa[q] = fa[np] = nq;
13            for (; p && t[p][ch] == q; p = fa[p]) t[p][ch] = nq;
14        }
15    }
16
17    int up[M], c[256] = {2}, a[M];
18    void rsort2() {
19        FOR (i, 1, 256) c[i] = 0;
20        FOR (i, 2, sz) up[i] = s[one[i] + len[fa[i]]];
21        FOR (i, 2, sz) c[up[i]]++;
22        FOR (i, 1, 256) c[i] += c[i - 1];
23        FOR (i, 2, sz) a[--c[up[i]]] = i;
24        FOR (i, 2, sz) G[fa[a[i]]].push_back(a[i]);
25    }

```

- 广义后缀自动机建后缀树，必须反向插入

```

1     int t[M][26], len[M] = {0}, fa[M], sz = 2, last = 1;
2     char* one[M];
3     void ins(int ch, char* pp) {
4         int p = last, np = 0, nq = 0, q = -1;
5         if (!t[p][ch]) {
6             np = sz++; one[np] = pp;

```

```

7         len[np] = len[p] + 1;
8         for (; p && !t[p][ch]; p = fa[p]) t[p][ch] = np;
9     }
10    if (!p) fa[np] = 1;
11    else {
12        q = t[p][ch];
13        if (len[p] + 1 == len[q]) fa[np] = q;
14        else {
15            nq = sz++; len[nq] = len[p] + 1; one[nq] = one[q];
16            memcpy(t[nq], t[q], sizeof t[0]);
17            fa[nq] = fa[q];
18            fa[np] = fa[q] = nq;
19            for (; t[p][ch] == q; p = fa[p]) t[p][ch] = nq;
20        }
21    }
22    last = np ? np : nq ? nq : q;
23 }
24 int up[M], c[256] = {2}, aa[M];
25 vector<int> G[M];
26 void rsort() {
27     FOR (i, 1, 256) c[i] = 0;
28     FOR (i, 2, sz) up[i] = *(one[i] + len[fa[i]]);
29     FOR (i, 2, sz) c[up[i]]++;
30     FOR (i, 1, 256) c[i] += c[i - 1];
31     FOR (i, 2, sz) aa[--c[up[i]]] = i;
32     FOR (i, 2, sz) G[fa[aa[i]]].push_back(aa[i]);
33 }

```

● 匹配

```

1 int u = 1, l = 0;
2 FOR (i, 0, strlen(s)) {
3     int ch = s[i] - 'a';
4     while (u && !t[u][ch]) { u = fa[u]; l = len[u]; }
5     ++l; u = t[u][ch];
6     if (!u) u = 1;
7     if (l) // do something...
8 }

```

● 获取子串状态

```

1 int get_state(int l, int r) {
2     int u = rpos[r], s = r - l + 1;
3     FOR (i, SP - 1, -1) if (len[pa[u][i]] >= s) u = pa[u][i];
4     return u;
5 }

```

● 配合 LCT

```

1 namespace lct_sam {
2     extern struct P *const null;
3     const int M = N;
4     struct P {
5         P *fa, *ls, *rs;
6         int last;
7
8         bool has_fa() { return fa->ls == this || fa->rs == this; }
9         bool d() { return fa->ls == this; }
10        P& c(bool x) { return x ? ls : rs; }
11        P* up() { return this; }
12        void down() {
13            if (ls != null) ls->last = last;
14            if (rs != null) rs->last = last;
15        }
16        void all_down() { if (has_fa()) fa->all_down(); down(); }
17    } *const null = new P[0, 0, 0, 0], pool[M], *pit = pool;
18    P* G[N];
19    int t[M][26], len[M] = {-1}, fa[M], sz = 2, last = 1;
20
21    void rot(P* o) {
22        bool dd = o->d();
23        P *f = o->fa, *t = o->c(!dd);
24        if (f->has_fa()) f->fa->c(f->d()) = o; o->fa = f->fa;

```

```

25     if (t != null) t->fa = f; f->c(dd) = t;
26     o->c(!dd) = f->up(); f->fa = o;
27 }
28 void splay(P* o) {
29     o->all_down();
30     while (o->has_fa()) {
31         if (o->fa->has_fa())
32             rot(o->d() ^ o->fa->d() ? 0 : o->fa);
33         rot(o);
34     }
35     o->up();
36 }
37 void access(int last, P* u, P* v = null) {
38     if (u == null) { v->last = last; return; }
39     splay(u);
40     P *t = u;
41     while (t->ls != null) t = t->ls;
42     int L = len[fa[t - pool]] + 1, R = len[u - pool];
43
44     if (u->last) bit::add(u->last - R + 2, u->last - L + 2, 1);
45     else bit::add(1, 1, R - L + 1);
46     bit::add(last - R + 2, last - L + 2, -1);
47
48     u->rs = v;
49     access(last, u->up()->fa, u);
50 }
51 void insert(P* u, P* v, P* t) {
52     if (v != null) { splay(v); v->rs = null; }
53     splay(u);
54     u->fa = t; t->fa = v;
55 }
56
57 void ins(int ch, int pp) {
58     int p = last, np = last = sz++;
59     len[np] = len[p] + 1;
60     for (; p && !t[p][ch]; p = fa[p]) t[p][ch] = np;
61     if (!p) fa[np] = 1;
62     else {
63         int q = t[p][ch];
64         if (len[p] + 1 == len[q]) { fa[np] = q; G[np]->fa = G[q]; }
65         else {
66             int nq = sz++; len[nq] = len[p] + 1;
67             memcpy(t[nq], t[q], sizeof t[0]);
68             insert(G[q], G[fa[q]], G[nq]);
69             G[nq]->last = G[q]->last;
70             fa[nq] = fa[q];
71             fa[np] = fa[q] = nq;
72             G[np]->fa = G[nq];
73             for (; t[p][ch] == q; p = fa[p]) t[p][ch] = nq;
74         }
75     }
76     access(pp + 1, G[np]);
77 }
78
79 void init() {
80     ++pit;
81     FOR (i, 1, N) {
82         G[i] = pit++;
83         G[i]->ls = G[i]->rs = G[i]->fa = null;
84     }
85     G[1] = null;
86 }
87 }

```

回文自动机

```

1 namespace pam {
2     int t[N][26], fa[N], len[N], rs[N], cnt[N], num[N];
3     int sz, n, last;

```

```

4     int _new(int l) {
5         memset(t[sz], 0, sizeof t[0]);
6         len[sz] = l; cnt[sz] = num[sz] = 0;
7         return sz++;
8     }
9     void init() {
10        rs[n = sz = 0] = -1;
11        last = _new(0);
12        fa[last] = _new(-1);
13    }
14    int get_fa(int x) {
15        while (rs[n - 1 - len[x]] != rs[n]) x = fa[x];
16        return x;
17    }
18    void ins(int ch) {
19        rs[++n] = ch;
20        int p = get_fa(last);
21        if (!t[p][ch]) {
22            int np = _new(len[p] + 2);
23            num[np] = num[fa[np]] = t[get_fa(fa[p])][ch] + 1;
24            t[p][ch] = np;
25        }
26        ++cnt[last = t[p][ch]];
27    }
28 }

```

manacher

```

1  int RL[N];
2  void manacher(int* a, int n) { // "abc" => "#a#b#a#"
3      int r = 0, p = 0;
4      FOR (i, 0, n) {
5          if (i < r) RL[i] = min(RL[2 * p - i], r - i);
6          else RL[i] = 1;
7          while (i - RL[i] >= 0 && i + RL[i] < n && a[i - RL[i]] == a[i + RL[i]])
8              RL[i]++;
9          if (RL[i] + i - 1 > r) { r = RL[i] + i - 1; p = i; }
10     }
11     FOR (i, 0, n) --RL[i];
12 }

```

哈希

内置了自动双哈希开关 (小心 TLE)。

```

1  #include <bits/stdc++.h>
2  using namespace std;
3
4  #define ENABLE_DOUBLE_HASH
5
6  typedef long long LL;
7  typedef unsigned long long ULL;
8
9  const int x = 135;
10 const int N = 4e5 + 10;
11 const int p1 = 1e9 + 7, p2 = 1e9 + 9;
12 ULL xp1[N], xp2[N], xp[N];
13
14 void init_xp() {
15     xp1[0] = xp2[0] = xp[0] = 1;
16     for (int i = 1; i < N; ++i) {
17         xp1[i] = xp1[i - 1] * x % p1;
18         xp2[i] = xp2[i - 1] * x % p2;
19         xp[i] = xp[i - 1] * x;
20     }
21 }
22
23 struct String {

```

```

24     char s[N];
25     int length, subsize;
26     bool sorted;
27     ULL h[N], hl[N];
28
29     ULL hash() {
30         length = strlen(s);
31         ULL res1 = 0, res2 = 0;
32         h[length] = 0; // ATTENTION!
33         for (int j = length - 1; j >= 0; --j) {
34             #ifdef ENABLE_DOUBLE_HASH
35                 res1 = (res1 * x + s[j]) % p1;
36                 res2 = (res2 * x + s[j]) % p2;
37                 h[j] = (res1 << 32) | res2;
38             #else
39                 res1 = res1 * x + s[j];
40                 h[j] = res1;
41             #endif
42             // printf("%llu\n", h[j]);
43         }
44         return h[0];
45     }
46
47     // 获取子串哈希, 左闭右开区间
48     ULL get_substring_hash(int left, int right) const {
49         int len = right - left;
50         #ifdef ENABLE_DOUBLE_HASH
51             // get hash of s[left...right-1]
52             unsigned int mask32 = ~(0u);
53             ULL left1 = h[left] >> 32, right1 = h[right] >> 32;
54             ULL left2 = h[left] & mask32, right2 = h[right] & mask32;
55             return (((left1 - right1 * xp1[len] % p1 + p1) % p1) << 32) |
56                 (((left2 - right2 * xp2[len] % p2 + p2) % p2));
57         #else
58             return h[left] - h[right] * xp[len];
59         #endif
60     }
61
62     void get_all_subs_hash(int sublen) {
63         subsize = length - sublen + 1;
64         for (int i = 0; i < subsize; ++i)
65             hl[i] = get_substring_hash(i, i + sublen);
66         sorted = 0;
67     }
68
69     void sort_substring_hash() {
70         sort(hl, hl + subsize);
71         sorted = 1;
72     }
73
74     bool match(ULL key) const {
75         if (!sorted) assert (0);
76         if (!subsize) return false;
77         return binary_search(hl, hl + subsize, key);
78     }
79
80     void init(const char *t) {
81         length = strlen(t);
82         strcpy(s, t);
83     }
84 };
85
86 int LCP(const String &a, const String &b, int ai, int bi) {
87     // Find LCP of a[ai...] and b[bi...]
88     int l = 0, r = min(a.length - ai, b.length - bi);
89     while (l < r) {
90         int mid = (l + r + 1) / 2;
91         if (a.get_substring_hash(ai, ai + mid) == b.get_substring_hash(bi, bi + mid))
92             l = mid;
93         else r = mid - 1;
94     }

```

```

95     return l;
96 }
97
98 int check(int ans) {
99     if (T.length < ans) return 1;
100     T.get_all_subs_hash(ans); T.sort_substring_hash();
101     for (int i = 0; i < S.length - ans + 1; ++i)
102         if (!T.match(S.get_substring_hash(i, i + ans)))
103             return 1;
104     return 0;
105 }
106
107 int main() {
108     init_xp(); // DON'T FORGET TO DO THIS!
109
110     for (int tt = 1; tt <= kases; ++tt) {
111         scanf("%d", &n); scanf("%s", str);
112         S.init(str);
113         S.hash(); T.hash();
114     }
115 }

```

后缀数组

构造时间: $O(L \log L)$; 查询时间 $O(\log L)$ 。suffix 数组是排好序的后缀下标, suffix 的反数组是后缀数组。

```

1  #include <bits/stdc++.h>
2  using namespace std;
3
4  const int N = 2e5 + 10;
5  const int Nlog = 18;
6
7  struct SuffixArray {
8      const int L;
9      vector<vector<int>> > P;
10     vector<pair<pair<int, int>, int>> > M;
11     int s[N], sa[N], rank[N], height[N];
12     // s: raw string
13     // sa[i]=k: s[k...L-1] ranks i (0 based)
14     // rank[i]=k: the rank of s[i...L-1] is k (0 based)
15     // height[i] = lcp(sa[i-1], sa[i])
16
17     SuffixArray(const string &raw_s) : L(raw_s.length()), P(1, vector<int>(L, 0)), M(L) {
18         for (int i = 0; i < L; i++)
19             P[0][i] = this->s[i] = int(raw_s[i]);
20         for (int skip = 1, level = 1; skip < L; skip *= 2, level++) {
21             P.push_back(vector<int>(L, 0));
22             for (int i = 0; i < L; i++)
23                 M[i] = make_pair(make_pair(P[level - 1][i], i + skip < L ? P[level - 1][i + skip] : -1000), i);
24             sort(M.begin(), M.end());
25             for (int i = 0; i < L; i++)
26                 P[level][M[i].second] = (i > 0 && M[i].first == M[i - 1].first) ? P[level][M[i - 1].second] : i;
27         }
28         for (unsigned i = 0; i < P.back().size(); ++i) {
29             rank[i] = P.back()[i];
30             sa[rank[i]] = i;
31         }
32     }
33
34     // This is a traditional way to calculate LCP
35     void getHeight() {
36         memset(height, 0, sizeof height);
37         int k = 0;
38         for (int i = 0; i < L; ++i) {
39             if (rank[i] == 0) continue;
40             if (k) k--;
41             int j = sa[rank[i] - 1];
42             while (i + k < L && j + k < L && s[i + k] == s[j + k]) ++k;
43             height[rank[i]] = k;

```

```

44     }
45     rmq_init(height, L);
46 }
47
48 int f[N][Nlog];
49 inline int highbit(int x) {
50     return 31 - __builtin_clz(x);
51 }
52
53 int rmq_query(int x, int y) {
54     int p = highbit(y - x + 1);
55     return min(f[x][p], f[y - (1 << p) + 1][p]);
56 }
57
58 // arr has to be 0 based
59 void rmq_init(int *arr, int length) {
60     for (int x = 0; x <= highbit(length); ++x)
61         for (int i = 0; i <= length - (1 << x); ++i) {
62             if (!x) f[i][x] = arr[i];
63             else f[i][x] = min(f[i][x - 1], f[i + (1 << (x - 1))][x - 1]);
64         }
65 }
66
67 #ifdef NEW
68 // returns the length of the longest common prefix of s[i...L-1] and s[j...L-1]
69 int LongestCommonPrefix(int i, int j) {
70     int len = 0;
71     if (i == j) return L - i;
72     for (int k = (int) P.size() - 1; k >= 0 && i < L && j < L; k--) {
73         if (P[k][i] == P[k][j]) {
74             i += 1 << k;
75             j += 1 << k;
76             len += 1 << k;
77         }
78     }
79     return len;
80 }
81 #else
82 int LongestCommonPrefix(int i, int j) {
83     // getHeight() must be called first
84     if (i == j) return L - i;
85     if (i > j) swap(i, j);
86     return rmq_query(i + 1, j);
87 }
88 #endif
89
90 int checkNonOverlappingSubstring(int K) {
91     // check if there is two non-overlapping identical substring of length K
92     int minsa = 0, maxsa = 0;
93     for (int i = 0; i < L; ++i) {
94         if (height[i] < K) {
95             minsa = sa[i]; maxsa = sa[i];
96         } else {
97             minsa = min(minsa, sa[i]);
98             maxsa = max(maxsa, sa[i]);
99             if (maxsa - minsa >= K) return 1;
100         }
101     }
102     return 0;
103 }
104
105 int checkBelongToDifferentSubstring(int K, int split) {
106     int minsa = 0, maxsa = 0;
107     for (int i = 0; i < L; ++i) {
108         if (height[i] < K) {
109             minsa = sa[i]; maxsa = sa[i];
110         } else {
111             minsa = min(minsa, sa[i]);
112             maxsa = max(maxsa, sa[i]);
113             if (maxsa > split && minsa < split) return 1;
114         }
115     }

```

```

115     }
116     return 0;
117 }
118
119 } *S;
120
121 int main() {
122     string s, t;
123     cin >> s >> t;
124     int sp = s.length();
125     s += "*" + t;
126     S = new SuffixArray(s);
127     S->getHeight();
128     int left = 0, right = sp;
129     while (left < right) {
130         int mid = (left + right + 1) / 2;
131         if (S->checkBelongToDifferentSubstring(mid, sp))
132             left = mid;
133         else right = mid - 1;
134     }
135     printf("%d\n", left);
136 }

```

- SA-IS
- 仅在后缀自动机被卡内存或者卡常且需要 $O(1)$ LCA 的情况下使用（比赛中敲这个我觉得不行）
- UOJ 35

```

1 // rk [0..n-1] -> [1..n], sa/ht [1..n]
2 // s[i] > 0 && s[n] = 0
3 // b: normally as bucket
4 // c: normally as bucket1
5 // d: normally as bucket2
6 // f: normally as cntbuf
7
8 template<size_t size>
9 struct SuffixArray {
10     bool t[size << 1];
11     int b[size], c[size];
12     int sa[size], rk[size], ht[size];
13     inline bool isLMS(const int i, const bool *t) { return i > 0 && t[i] && !t[i - 1]; }
14     template<class T>
15     inline void inducedSort(T s, int *sa, const int n, const int M, const int bs,
16                             bool *t, int *b, int *f, int *p) {
17         fill(b, b + M, 0); fill(sa, sa + n, -1);
18         FOR (i, 0, n) b[s[i]]++;
19         f[0] = b[0];
20         FOR (i, 1, M) f[i] = f[i - 1] + b[i];
21         FORD (i, bs - 1, -1) sa[--f[s[p[i]]]] = p[i];
22         FOR (i, 1, M) f[i] = f[i - 1] + b[i - 1];
23         FOR (i, 0, n) if (sa[i] > 0 && !t[sa[i] - 1]) sa[f[s[sa[i] - 1]]++] = sa[i] - 1;
24         f[0] = b[0];
25         FOR (i, 1, M) f[i] = f[i - 1] + b[i];
26         FORD (i, n - 1, -1) if (sa[i] > 0 && t[sa[i] - 1]) sa[--f[s[sa[i] - 1]]] = sa[i] - 1;
27     }
28     template<class T>
29     inline void sais(T s, int *sa, int n, bool *t, int *b, int *c, int M) {
30         int i, j, bs = 0, cnt = 0, p = -1, x, *r = b + M;
31         t[n - 1] = 1;
32         FORD (i, n - 2, -1) t[i] = s[i] < s[i + 1] || (s[i] == s[i + 1] && t[i + 1]);
33         FOR (i, 1, n) if (t[i] && !t[i - 1]) c[bs++] = i;
34         inducedSort(s, sa, n, M, bs, t, b, r, c);
35         for (i = bs = 0; i < n; i++) if (isLMS(sa[i], t)) sa[bs++] = sa[i];
36         FOR (i, bs, n) sa[i] = -1;
37         FOR (i, 0, bs) {
38             x = sa[i];
39             for (j = 0; j < n; j++) {
40                 if (p == -1 || s[x + j] != s[p + j] || t[x + j] != t[p + j]) { cnt++, p = x; break; }
41                 else if (j > 0 && (isLMS(x + j, t) || isLMS(p + j, t))) break;
42             }
43             x = (~x & 1 ? x >> 1 : x - 1 >> 1), sa[bs + x] = cnt - 1;
44         }
45     }
46 }

```



```

45     for (i = j = n - 1; i >= bs; i--) if (sa[i] >= 0) sa[j--] = sa[i];
46     int *s1 = sa + n - bs, *d = c + bs;
47     if (cnt < bs) sais(s1, sa, bs, t + n, b, c + bs, cnt);
48     else FOR (i, 0, bs) sa[s1[i]] = i;
49     FOR (i, 0, bs) d[i] = c[sa[i]];
50     inducedSort(s, sa, n, M, bs, t, b, r, d);
51 }
52 template<typename T>
53 inline void getHeight(T s, const int n, const int *sa) {
54     for (int i = 0, k = 0; i < n; i++) {
55         if (rk[i] == 0) k = 0;
56         else {
57             if (k > 0) k--;
58             int j = sa[rk[i] - 1];
59             while (i + k < n && j + k < n && s[i + k] == s[j + k]) k++;
60         }
61         ht[rk[i]] = k;
62     }
63 }
64 template<class T>
65 inline void init(T s, int n, int M) {
66     sais(s, sa, ++n, t, b, c, M);
67     for (int i = 1; i < n; i++) rk[sa[i]] = i;
68     getHeight(s, n, sa);
69 }
70 };
71
72 const int N = 2E5 + 100;
73 SuffixArray<N> sa;
74
75 int main() {
76     string s; cin >> s; int n = s.length();
77     sa.init(s, n, 128);
78     FOR (i, 1, n + 1) printf("%d%c", sa.sa[i] + 1, i == _i - 1 ? '\n' : ' ');
79     FOR (i, 2, n + 1) printf("%d%c", sa.ht[i], i == _i - 1 ? '\n' : ' ');
80 }

```

KMP 自动机

```

1  int m; int pat[N];
2  namespace kmp {
3      int f[N]; // f[i] 表示已匹配成功 i 个, 失配要去哪里
4
5      template<typename T>
6      int go(int stat, T c, bool& acc) {
7          // stat 是当前态 (表示已经匹配了 stat 个字符), c 是要走的边
8          while (stat && c != pat[stat]) stat = f[stat];
9          if (c == pat[stat]) stat++;
10         if (stat == m) acc = true;
11         return stat;
12     }
13
14     void getFail() {
15         static int f2[N];
16         f[0] = f[1] = 0;
17         f2[0] = f2[1] = 0;
18         FOR (i, 1, m) {
19             int j = f2[i];
20             while (j && pat[i] != pat[j]) j = f2[j];
21             f2[i+1] = f[j+1] = (pat[i] == pat[j]) ? j+1 : 0;
22             if (f[i+1] == j+1 && pat[i+1] == pat[j+1]) f[i+1] = f[j+1];
23         }
24         FOR (i, 0, m) dbg(i, f[i]);
25     }
26 }

```

● 拓展 KMP

```

1  #include <bits/stdc++.h>
2  using namespace std;

```

```

3
4  /*
5   Define template S, pattern T, len(S)=n, len(T)=m
6   Find the longest common prefix of T and every suffix of S
7   ex[i]: the LCP between T and S[i..n-1]
8   */
9
10 const int maxn = 1e6 + 10;
11 int nt[maxn], ex[maxn];
12 char s[maxn], t[maxn];
13
14 void get_next(char *str) {
15     int i = 0, j, po, len = strlen(str);
16     nt[0] = len;
17     while (str[i] == str[i + 1] && i + 1 < len)
18         i++;
19     nt[1] = i;
20     po = 1;
21     for (i = 2; i < len; i++) {
22         if (nt[i - po] + i < nt[po] + po)
23             nt[i] = nt[i - po];
24         else {
25             j = nt[po] + po - i;
26             if (j < 0) j = 0;
27             while (i + j < len && str[j] == str[j + i])
28                 j++;
29             nt[i] = j;
30             po = i;
31         }
32     }
33 }
34
35 void exkmp(char *s1, char *s2) {
36     int i = 0, j, po, len = strlen(s1), l2 = strlen(s2);
37     get_next(s2);
38     while (s1[i] == s2[i] && i < l2 && i < len)
39         i++;
40     ex[0] = i;
41     po = 0;
42     for (i = 1; i < len; i++) {
43         if (nt[i - po] + i < ex[po] + po)
44             ex[i] = nt[i - po];
45         else {
46             j = ex[po] + po - i;
47             if (j < 0) j = 0;
48             while (i + j < len && j < l2 && s1[j + i] == s2[j])
49                 j++;
50             ex[i] = j;
51             po = i;
52         }
53     }
54 }
55
56 int main() {
57     const int modn = 1e9 + 7;
58     int T; scanf("%d", &T);
59     while (T--) {
60         memset(nt, 0, sizeof nt);
61         memset(ex, 0, sizeof ex);
62         scanf("%s", s); scanf("%s", t);
63         int slen = strlen(s), tlen = strlen(t);
64         reverse(s, s + slen);
65         reverse(t, t + tlen);
66         exkmp(s, t);
67         int ans = 0;
68         for (int i = 0; i < slen; ++i)
69             ans = (ans + 1LL * ex[i] * (ex[i] + 1) / 2) % modn;
70         printf("%d\n", ans);
71     }
72 }

```

Trie

```
1 namespace trie {
2     int t[N][26], sz, ed[N];
3     void init() { sz = 2; memset(ed, 0, sizeof ed); }
4     int _new() { memset(t[sz], 0, sizeof t[sz]); return sz++; }
5     void ins(char* s, int p) {
6         int u = 1;
7         FOR (i, 0, strlen(s)) {
8             int c = s[i] - 'a';
9             if (!t[u][c]) t[u][c] = _new();
10            u = t[u][c];
11        }
12        ed[u] = p;
13    }
14 }
```

AC 自动机

```
1 const int N = 1e6 + 100, M = 26;
2
3 int mp(char ch) { return ch - 'a'; }
4
5 struct ACA {
6     int ch[N][M], danger[N], fail[N];
7     int sz;
8     void init() {
9         sz = 1;
10        memset(ch[0], 0, sizeof ch[0]);
11        memset(danger, 0, sizeof danger);
12    }
13    void insert(const string &s, int m) {
14        int n = s.size(); int u = 0, c;
15        FOR (i, 0, n) {
16            c = mp(s[i]);
17            if (!ch[u][c]) {
18                memset(ch[sz], 0, sizeof ch[sz]);
19                danger[sz] = 0; ch[u][c] = sz++;
20            }
21            u = ch[u][c];
22        }
23        danger[u] |= 1 << m;
24    }
25    void build() {
26        queue<int> Q;
27        fail[0] = 0;
28        for (int c = 0, u; c < M; c++) {
29            u = ch[0][c];
30            if (u) { Q.push(u); fail[u] = 0; }
31        }
32        while (!Q.empty()) {
33            int r = Q.front(); Q.pop();
34            danger[r] |= danger[fail[r]];
35            for (int c = 0, u; c < M; c++) {
36                u = ch[r][c];
37                if (!u) {
38                    ch[r][c] = ch[fail[r]][c];
39                    continue;
40                }
41                fail[u] = ch[fail[r]][c];
42                Q.push(u);
43            }
44        }
45    }
46 } ac;
47
48 char s[N];
49
50 int main() {
```

```

51     int n; scanf("%d", &n);
52     ac.init();
53     while (n--) {
54         scanf("%s", s);
55         ac.insert(s, 0);
56     }
57     ac.build();
58
59     scanf("%s", s);
60     int u = 0; n = strlen(s);
61     FOR (i, 0, n) {
62         u = ac.ch[u][mp(s[i])];
63         if (ac.danger[u]) {
64             puts("YES");
65             return 0;
66         }
67     }
68     puts("NO");
69     return 0;
70 }

```

杂项

STL

- copy

```

1  template <class InputIterator, class OutputIterator>
2  OutputIterator copy (InputIterator first, InputIterator last, OutputIterator result);

```

- merge (如果相等, 第一个优先)

```

1  template <class InputIterator1, class InputIterator2,
2           class OutputIterator, class Compare>
3  OutputIterator merge (InputIterator1 first1, InputIterator1 last1,
4                       InputIterator2 first2, InputIterator2 last2,
5                       OutputIterator result, Compare comp);

```

- for_each

```

1  template <class InputIterator, class Function>
2  Function for_each (InputIterator first, InputIterator last, Function fn);

```

- transform

```

1  template <class InputIterator, class OutputIterator, class UnaryOperation>
2  OutputIterator transform (InputIterator first1, InputIterator last1,
3                           OutputIterator result, UnaryOperation op);

```

- numeric_limits

```

1  template <class T> numeric_limits;

```

- iota

```

1  template< class ForwardIterator, class T >
2  void iota( ForwardIterator first, ForwardIterator last, T value );

```

日期

```

1  // Routines for performing computations on dates. In these routines,
2  // months are expressed as integers from 1 to 12, days are expressed
3  // as integers from 1 to 31, and years are expressed as 4-digit
4  // integers.
5
6  string dayOfWeek[] = {"Mo", "Tu", "We", "Th", "Fr", "Sa", "Su"};
7
8  // converts Gregorian date to integer (Julian day number)

```

```

9
10 int DateToInt (int m, int d, int y){
11     return
12         1461 * (y + 4800 + (m - 14) / 12) / 4 +
13         367 * (m - 2 - (m - 14) / 12 * 12) / 12 -
14         3 * ((y + 4900 + (m - 14) / 12) / 100) / 4 +
15         d - 32075;
16 }
17
18 // converts integer (Julian day number) to Gregorian date: month/day/year
19
20 void IntToDate (int jd, int &m, int &d, int &y){
21     int x, n, i, j;
22
23     x = jd + 68569;
24     n = 4 * x / 146097;
25     x -= (146097 * n + 3) / 4;
26     i = (4000 * (x + 1)) / 1461001;
27     x -= 1461 * i / 4 - 31;
28     j = 80 * x / 2447;
29     d = x - 2447 * j / 80;
30     x = j / 11;
31     m = j + 2 - 12 * x;
32     y = 100 * (n - 49) + i + x;
33 }
34
35 // converts integer (Julian day number) to day of week
36
37 string IntToDay (int jd){
38     return dayOfWeek[jd % 7];
39 }

```

子集枚举

- 枚举真子集

```

1 for (int s = (S - 1) & S; s; s = (s - 1) & S)

```

- 枚举大小为 k 的子集

```

1 template<typename T>
2 void subset(int k, int n, T&& f) {
3     int t = (1 << k) - 1;
4     while (t < 1 << n) {
5         f(t);
6         int x = t & -t, y = t + x;
7         t = ((t & ~y) / x >> 1) | y;
8     }
9 }

```

权值最大上升子序列

```

1 const LL maxn = 1E5 + 10;
2 const LL INF = 1E10;
3 struct P {
4     LL k, v;
5     bool operator < (const P& rhs) const {
6         return k < rhs.k || (k == rhs.k && v < rhs.v);
7     }
8 };
9 LL k[maxn], v[maxn], n, T;
10 set<P> s;
11
12 int main() {
13     cin >> T;
14     while (T--) {
15         s.clear();
16         s.insert({-INF, 0});

```

```

17     cin >> n;
18     FOR (i, 0, n) scanf("%lld", &k[i]);
19     FOR (i, 0, n) scanf("%lld", &v[i]);
20     FOR (i, 0, n) {
21         auto it = s.lower_bound({k[i], INF});
22         LL vv = (--it)->v + v[i];
23         ++it;
24         while (it != s.end() && it->v <= vv)
25             it = s.erase(it);
26         if (it == s.end() || it->k != k[i]) s.insert({k[i], vv});
27     }
28     cout << s.rbegin()->v << endl;
29 }
30 }

```

数位 DP

```

1  LL dfs(LL base, LL pos, LL len, LL s, bool limit) {
2      if (pos == -1) return s ? base : 1;
3      if (!limit && dp[base][pos][len][s] != -1) return dp[base][pos][len][s];
4      LL ret = 0;
5      LL ed = limit ? a[pos] : base - 1;
6      FOR (i, 0, ed + 1) {
7          tmp[pos] = i;
8          if (len == pos)
9              ret += dfs(base, pos - 1, len - (i == 0), s, limit && i == a[pos]);
10         else if (s && pos < (len + 1) / 2)
11             ret += dfs(base, pos - 1, len, tmp[len - pos] == i, limit && i == a[pos]);
12         else
13             ret += dfs(base, pos - 1, len, s, limit && i == a[pos]);
14     }
15     if (!limit) dp[base][pos][len][s] = ret;
16     return ret;
17 }
18
19 LL solve(LL x, LL base) {
20     LL sz = 0;
21     while (x) {
22         a[sz++] = x % base;
23         x /= base;
24     }
25     return dfs(base, sz - 1, sz - 1, 1, true);
26 }

```

土制 bitset

- 可以用 `auto p = reinterpret_cast<unsigned*>(&x);` (`p[0]` 的最低位就是 bitset 的最低位)

```

1  // M 要开大至少 1 个 64
2  const int M = (1E4 + 200) / 64;
3  typedef unsigned long long ULL;
4  const ULL ONE = 1;
5
6  struct Bitset {
7      ULL a[M];
8      void go(int x) {
9          int offset = x / 64; x %= 64;
10         for (int i = offset, j = 0; i + 1 < M; ++i, ++j) {
11             a[j] |= a[i] >> x;
12             if (x) a[j] |= a[i + 1] << (64 - x); // 不能左移 64 位
13         }
14     }
15     void init() { memset(a, 0, sizeof a); }
16     void set(int x) {
17         int offset = x / 64; x %= 64;
18         a[offset] |= (ONE << x);
19     }
20     void prt() {

```

```

21     FOR (i, 0, M) FOR (j, 0, 64) putchar((a[i] & (ONE << j)) ? '1' : '0');
22     puts("");
23 }
24 int lowbit() {
25     FOR (i, 0, M) if (a[i]) return i * 64 + __builtin_ctzll(a[i]);
26     assert (0);
27 }
28 int highbit(int x) {
29     // [0,x) 的最高位
30     int offset = x / 64; x %= 64;
31     FOR (i, offset, -1) {
32         if (!a[i]) continue;
33         if (i == offset) {
34             FOR (j, x - 1, -1) if ((ONE << j) & a[i]) { return i * 64 + j; }
35         } else return i * 64 + 63 - __builtin_clzll(a[i]);
36     }
37     assert (0);
38 }
39 };

```

随机

- 不要使用 rand()。
- chrono::steady_clock::now().time_since_epoch().count() 可用于计时。
- 64 位可以使用 mt19937_64。

```

1 int main() {
2     mt19937 rng(chrono::steady_clock::now().time_since_epoch().count());
3     vector<int> permutation(N);
4
5     for (int i = 0; i < N; i++)
6         permutation[i] = i;
7     shuffle(permutation.begin(), permutation.end(), rng);
8
9     for (int i = 0; i < N; i++)
10         permutation[i] = i;
11     for (int i = 1; i < N; i++)
12         swap(permutation[i], permutation[uniform_int_distribution<int>(0, i)(rng)]);

```

伪随机数

```

1 unsigned rnd() {
2     static unsigned A = 1 << 16 | 3, B = 33333331, C = 2341;
3     return C = A * C + B;
4 }

```

随机素数表

42737, 46411, 50101, 52627, 54577, 191677, 194869, 210407, 221831, 241337, 578603, 625409, 713569, 788813, 862481, 2174729, 2326673, 2688877, 2779417, 3133583, 4489747, 6697841, 6791471, 6878533, 7883129, 9124553, 10415371, 11134633, 12214801, 15589333, 17148757, 17997457, 20278487, 27256133, 28678757, 38206199, 41337119, 47422547, 48543479, 52834961, 76993291, 85852231, 95217823, 108755593, 132972461, 171863609, 173629837, 176939899, 207808351, 227218703, 306112619, 311809637, 322711981, 330806107, 345593317, 345887293, 362838523, 373523729, 394207349, 409580177, 437359931, 483577261, 490845269, 512059357, 534387017, 698987533, 764016151, 906097321, 914067307, 954169327

1572869, 3145739, 6291469, 12582917, 25165843, 50331653 (适合哈希的素数)

```

1 from random import randint
2
3 def is_prime(num, test_count):
4     if num == 1:
5         return False
6     if test_count >= num:
7         test_count = num - 1

```

```

8     for x in range(test_count):
9         val = randint(1, num - 1)
10        if pow(val, num-1, num) != 1:
11            return False
12    return True
13
14    def generate_big_prime(n):
15        found_prime = False
16        while not found_prime:
17            p = randint(2**(n-1), 2**n)
18            if is_prime(p, 1000):
19                return p

```

Java

Regex

```

1  // Code which demonstrates the use of Java's regular expression libraries.
2  // This is a solution for
3  //
4  // Loglan: a logical language
5  // http://acm.uva.es/p/v1/134.html
6
7  import java.util.*;
8  import java.util.regex.*;
9
10 public class LogLan {
11
12     public static void main(String args[]) {
13
14         String regex = BuildRegex();
15         Pattern pattern = Pattern.compile(regex);
16
17         Scanner s = new Scanner(System.in);
18         while (true) {
19
20             // In this problem, each sentence consists of multiple lines, where the last
21             // line is terminated by a period. The code below reads lines until
22             // encountering a line whose final character is a '.'. Note the use of
23             //
24             // s.length() to get length of string
25             // s.charAt() to extract characters from a Java string
26             // s.trim() to remove whitespace from the beginning and end of Java string
27             //
28             // Other useful String manipulation methods include
29             //
30             // s.compareTo(t) < 0 if s < t, lexicographically
31             // s.indexOf("apple") returns index of first occurrence of "apple" in s
32             // s.lastIndexOf("apple") returns index of last occurrence of "apple" in s
33             // s.replace(c,d) replaces occurrences of character c with d
34             // s.startsWith("apple") returns (s.indexOf("apple") == 0)
35             // s.toLowerCase() / s.toUpperCase() returns a new lower/uppercased string
36             //
37             // Integer.parseInt(s) converts s to an integer (32-bit)
38             // Long.parseLong(s) converts s to a long (64-bit)
39             // Double.parseDouble(s) converts s to a double
40
41             String sentence = "";
42             while (true) {
43                 sentence = (sentence + " " + s.nextLine()).trim();
44                 if (sentence.equals("#")) return;
45                 if (sentence.charAt(sentence.length() - 1) == '.') break;
46             }
47
48             // now, we remove the period, and match the regular expression
49
50             String removed_period = sentence.substring(0, sentence.length() - 1).trim();
51             if (pattern.matcher(removed_period).find()) {
52                 System.out.println("Good");

```



```

53         } else {
54             System.out.println("Bad!");
55         }
56     }
57 }
58 }

```

Decimal Format

```

1  // examples for printing floating point numbers
2
3  import java.util.*;
4  import java.io.*;
5  import java.text.DecimalFormat;
6
7  public class DecFormat {
8      public static void main(String[] args) {
9          DecimalFormat fmt;
10
11         // round to at most 2 digits, leave of digits if not needed
12         fmt = new DecimalFormat("#.##");
13         System.out.println(fmt.format(12345.6789)); // produces 12345.68
14         System.out.println(fmt.format(12345.0)); // produces 12345
15         System.out.println(fmt.format(0.0)); // produces 0
16         System.out.println(fmt.format(0.01)); // produces .1
17
18         // round to precisely 2 digits
19         fmt = new DecimalFormat("#.00");
20         System.out.println(fmt.format(12345.6789)); // produces 12345.68
21         System.out.println(fmt.format(12345.0)); // produces 12345.00
22         System.out.println(fmt.format(0.0)); // produces .00
23
24         // round to precisely 2 digits, force leading zero
25         fmt = new DecimalFormat("0.00");
26         System.out.println(fmt.format(12345.6789)); // produces 12345.68
27         System.out.println(fmt.format(12345.0)); // produces 12345.00
28         System.out.println(fmt.format(0.0)); // produces 0.00
29
30         // round to precisely 2 digits, force leading zeros
31         fmt = new DecimalFormat("000000000.00");
32         System.out.println(fmt.format(12345.6789)); // produces 000012345.68
33         System.out.println(fmt.format(12345.0)); // produces 000012345.00
34         System.out.println(fmt.format(0.0)); // produces 000000000.00
35
36         // force leading '+'
37         fmt = new DecimalFormat("+0;-0");
38         System.out.println(fmt.format(12345.6789)); // produces +12346
39         System.out.println(fmt.format(-12345.6789)); // produces -12346
40         System.out.println(fmt.format(0)); // produces +0
41
42         // force leading positive/negative, pad to 2
43         fmt = new DecimalFormat("positive 00;negative 0");
44         System.out.println(fmt.format(1)); // produces "positive 01"
45         System.out.println(fmt.format(-1)); // produces "negative 01"
46
47         // quote special chars (#)
48         fmt = new DecimalFormat("text with '#' followed by #");
49         System.out.println(fmt.format(12.34)); // produces "text with # followed by 12"
50
51         // always show "."
52         fmt = new DecimalFormat("#.##");
53         fmt.setDecimalSeparatorAlwaysShown(true);
54         System.out.println(fmt.format(12.34)); // produces "12.3"
55         System.out.println(fmt.format(12)); // produces "12."
56         System.out.println(fmt.format(0.34)); // produces "0.3"
57
58         // different grouping distances:
59         fmt = new DecimalFormat("#,###.###");
60         System.out.println(fmt.format(123456789.123)); // produces "1,2345,6789.123"
61

```

```

62 // scientific:
63 fmt = new DecimalFormat("0.000E00");
64 System.out.println(fmt.format(123456789.123)); // produces "1.235E08"
65 System.out.println(fmt.format(-0.000234)); // produces "-2.34E-04"
66
67 // using variable number of digits:
68 fmt = new DecimalFormat("0");
69 System.out.println(fmt.format(123.123)); // produces "123"
70 fmt.setMinimumFractionDigits(8);
71 System.out.println(fmt.format(123.123)); // produces "123.12300000"
72 fmt.setMaximumFractionDigits(0);
73 System.out.println(fmt.format(123.123)); // produces "123"
74
75 // note: to pad with spaces, you need to do it yourself:
76 // String out = fmt.format(...)
77 // while (out.length() < targlength) out = " "+out;
78 }
79 }

```

Sort

```

1 import java.util.ArrayList;
2 import java.util.Collections;
3 import java.util.List;
4
5 public class Employee implements Comparable<Employee> {
6     private int id;
7     private String name;
8     private int age;
9
10    public Employee(int id, String name, int age) {
11        this.id = id;
12        this.name = name;
13        this.age = age;
14    }
15
16    @Override
17    public int compareTo(Employee o) {
18        if (id > o.id) {
19            return 1;
20        } else if (id < o.id) {
21            return -1;
22        }
23        return 0;
24    }
25
26    public static void main(String[] args) {
27        List<Employee> list = new ArrayList<Employee>();
28        list.add(new Employee(2, "Java", 20));
29        list.add(new Employee(1, "C", 30));
30        list.add(new Employee(3, "C#", 10));
31        Collections.sort(list);
32    }
33 }

```

扩栈（本地使用）

```

1 #include <sys/resource.h>
2 void init_stack(){
3     const rlim_t kStackSize = 512 * 1024 * 1024;
4     struct rlimit rl;
5     int result;
6     result = getrlimit(RLIMIT_STACK, &rl);
7     if (result == 0) {
8         if (rl.rlim_cur < kStackSize) {
9             rl.rlim_cur = kStackSize;
10            result = setrlimit(RLIMIT_STACK, &rl);
11            if (result != 0) {

```

```

12         fprintf(stderr, "setrlimit returned result = %d\n", result);
13     }
14 }
15 }
16 }

```

心态崩了

- `(int)v.size()`
- `1LL << k`
- 递归函数用全局或者 `static` 变量要小心
- 预处理组合数注意上限
- 想清楚到底是要 `multiset` 还是 `set`
- 提交之前看一下数据范围，测一下边界
- 数据结构注意数组大小（2 倍，4 倍）
- 字符串注意字符集
- 如果函数中使用了默认参数的话，注意调用时的参数个数。
- 注意要读完
- 构造参数无法使用自己
- 树链剖分/dfs 序，初始化或者询问不要忘记 `idx, ridx`
- 排序时注意结构体的所有属性是不是考虑了
- 不要把 `while` 写成 `if`
- 不要把 `int` 开成 `char`
- 清零的时候全部用 `0~n+1`。
- 模意义下不要用除法
- 哈希不要自然溢出
- 最短路不要 SPFA，乖乖写 Dijkstra
- 上取整以及 GCD 小心负数
- `mid` 用 `l + (r - l) / 2` 可以避免溢出和负数的问题
- 小心模板自带的意料之外的隐式类型转换